

1. What is Address Binding?
  - a) going to an address in memory
  - b) locating an address with the help of another address
  - c) binding two addresses together to form a new address in a different memory space
  - d) **a mapping from one address space to another**
2. Binding of instructions and data to memory addresses can be done at \_\_\_\_\_
  - a) Compile time
  - b) Load time
  - c) Execution time
  - d) **All of the mentioned**
3. If the process can be moved during its execution from one memory segment to another, then binding must be \_\_\_\_\_
  - a) **delayed until run time**
  - b) preponed to compile time
  - c) preponed to load time
  - d) none of the mentioned
4. What is Dynamic loading?
  - a) loading multiple routines dynamically
  - b) **loading a routine only when it is called**
  - c) loading multiple routines randomly
  - d) none of the mentioned
5. The idea of overlays is to \_\_\_\_\_
  - a) data that are needed at any given time
  - b) enable a process to be larger than the amount of memory allocated to it
  - c) keep in memory only those instructions
  - d) **all of the mentioned**
6. The \_\_\_\_\_ swaps processes in and out of the memory.
  - a) **Memory manager**
  - b) CPU
  - c) CPU manager
  - d) User
7. If a higher priority process arrives and wants service, the memory manager can swap out the lower priority process to execute the higher priority process. When the higher priority process finishes, the lower priority process

is swapped back in and continues execution. This variant of swapping is sometimes called?

- a) priority swapping
- b) pull out, push in
- c) roll out, roll in**
- d) none of the mentioned

8. If binding is done at assembly or load time, then the process \_\_\_\_ be moved to different locations after being swapped out and in again.

- a) can
- b) must
- c) can never**
- d) may

9. In a system that does not support swapping \_\_\_\_\_

- a) the compiler normally binds symbolic addresses (variables) to relocatable addresses**
- b) the compiler normally binds symbolic addresses to physical addresses
- c) the loader binds relocatable addresses to physical addresses
- d) binding of symbolic addresses to physical addresses normally takes place during execution

10. The address generated by the CPU is referred to as \_\_\_\_\_

- a) Physical address
- b) Logical address**
- c) Neither physical nor logical
- d) None of the mentioned

11. The address loaded into the memory address register of the memory is referred to as \_\_\_\_\_

- a) Physical address**
- b) Logical address
- c) Neither physical nor logical
- d) None of the mentioned

12. The run time mapping from virtual to physical addresses is done by a hardware device called the \_\_\_\_\_

- a) Virtual to physical mapper
- b) Memory management unit**
- c) Memory mapping unit
- d) None of the mentioned

13. The size of a process is limited to the size of \_\_\_\_\_

- a) physical memory**
- b) external storage
- c) secondary storage
- d) none of the mentioned

14. If execution time binding is being used, then a process \_\_\_\_\_ be swapped to a different memory space.
- a) has to be
  - b) can never
  - c) must
  - d) may**
15. Swapping requires a \_\_\_\_\_
- a) motherboard
  - b) keyboard
  - c) monitor
  - d) backing store**
16. The backing store is generally a \_\_\_\_\_
- a) fast disk
  - b) disk large enough to accommodate copies of all memory images for all users
  - c) disk to provide direct access to the memory images
  - d) all of the mentioned**
17. The \_\_\_\_\_ consists of all processes whose memory images are in the backing store or in memory and are ready to run.
- a) wait queue
  - b) ready queue**
  - c) cpu
  - d) secondary storage
18. The \_\_\_\_\_ time in a swap out of a running process and swap in of a new process into the memory is very high.
- a) context – switch**
  - b) waiting
  - c) execution
  - d) all of the mentioned
19. The major part of swap time is \_\_\_\_\_ time.
- a) waiting
  - b) transfer**
  - c) execution
  - d) none of the mentioned
20. Swapping \_\_\_\_\_ be done when a process has pending I/O, or has to execute I/O operations only into operating system buffers.
- a) must
  - b) can
  - c) must never**
  - d) maybe
21. Which one of the following is the address generated by CPU?

- a) physical address
  - b) absolute address
  - c) logical address**
  - d) none of the mentioned
22. Run time mapping from virtual to physical address is done by \_\_\_\_\_
- a) **Memory management unit**
  - b) CPU
  - c) PCI
  - d) None of the mentioned
23. Memory management technique in which system stores and retrieves data from secondary storage for use in main memory is called?
- a) fragmentation
  - b) paging**
  - c) mapping
  - d) none of the mentioned
24. The address of a page table in memory is pointed by \_\_\_\_\_
- a) stack pointer
  - b) page table base register**
  - c) page register
  - d) program counter
25. Program always deals with \_\_\_\_\_
- a) logical address**
  - b) absolute address
  - c) physical address
  - d) relative address
26. The page table contains \_\_\_\_\_
- a) base address of each page in physical memory**
  - b) page offset
  - c) page size
  - d) none of the mentioned
27. What is compaction?
- a) a technique for overcoming internal fragmentation
  - b) a paging technique
  - c) a technique for overcoming external fragmentation**
  - d) a technique for overcoming fatal error
28. Operating System maintains the page table for \_\_\_\_\_
- a) each process**
  - b) each thread
  - c) each instruction
  - d) each address
29. In contiguous memory allocation \_\_\_\_\_

- a) **each process is contained in a single contiguous section of memory**
  - b) all processes are contained in a single contiguous section of memory
  - c) the memory space is contiguous
  - d) none of the mentioned
30. The relocation register helps in \_\_\_\_\_
- a) providing more address space to processes
  - b) a different address space to processes
  - c) to **protect the address spaces of processes**
  - d) none of the mentioned
31. The operating system and the other processes are protected from being modified by an already running process because \_\_\_\_\_
- a) they are in different memory spaces
  - b) they are in different logical addresses
  - c) they have a protection algorithm
  - d) **every address generated by the CPU is being checked against the relocation and limit registers**
32. When memory is divided into several fixed sized partitions, each partition may contain \_\_\_\_\_
- a) **exactly one process**
  - b) at least one process
  - c) multiple processes at once
  - d) none of the mentioned
33. In fixed size partition, the degree of multiprogramming is bounded by \_\_\_\_\_
- a) **the number of partitions**
  - b) the CPU utilization
  - c) the memory size
  - d) all of the mentioned
34. The first fit, best fit and worst fit are strategies to select a \_\_\_\_\_
- a) process from a queue to put in memory
  - b) processor to run the next process
  - c) **free hole from a set of available holes**
  - d) all of the mentioned
35. In internal fragmentation, memory is internal to a partition and \_\_\_\_\_
- a) is being used
  - b) **is not being used**
  - c) is always used
  - d) none of the mentioned
36. A solution to the problem of external fragmentation is \_\_\_\_\_
- a) **compaction**
  - b) larger memory space

- c) smaller memory space
  - d) none of the mentioned
37. The disadvantage of moving all process to one end of memory and all holes to the other direction, producing one large hole of available memory is \_\_\_\_\_
- a) the cost incurred**
  - b) the memory used
  - c) the CPU used
  - d) all of the mentioned
38. \_\_\_\_\_ is generally faster than \_\_\_\_\_ and \_\_\_\_\_
- a) first fit, best fit, worst fit**
  - b) best fit, first fit, worst fit
  - c) worst fit, best fit, first fit
  - d) none of the mentioned
39. External fragmentation exists when?
- a) enough total memory exists to satisfy a request but it is not contiguous**
  - b) the total memory is insufficient to satisfy a request
  - c) a request cannot be satisfied even when the total memory is free
  - d) none of the mentioned
40. When the memory allocated to a process is slightly larger than the process, then \_\_\_\_\_
- a) internal fragmentation occurs**
  - b) external fragmentation occurs
  - c) both internal and external fragmentation occurs
  - d) neither internal nor external fragmentation occurs
41. Physical memory is broken into fixed-sized blocks called \_\_\_\_\_
- a) frames**
  - b) pages
  - c) backing store
  - d) none of the mentioned
42. Logical memory is broken into blocks of the same size called \_\_\_\_\_
- a) frames
  - b) pages**
  - c) backing store
  - d) none of the mentioned
43. very address generated by the CPU is divided into two parts. They are \_\_\_\_\_
- a) frame bit & page number
  - b) page number & page offset**
  - c) page offset & frame bit

- d) frame offset & page offset
44. The \_\_\_\_\_ is used as an index into the page table.
- a) frame bit
  - b) page number**
  - c) page offset
  - d) frame offset
45. The \_\_\_\_ table contains the base address of each page in physical memory.
- a) process
  - b) memory
  - c) page**
  - d) frame
46. The size of a page is typically \_\_\_\_\_
- a) varied
  - b) power of 2**
  - c) power of 4
  - d) none of the mentioned
47. If the size of logical address space is  $2^m$ , and a page size is  $2^n$  addressing units, then the high order \_\_\_\_ bits of a logical address designate the page number, and the \_\_\_\_ low order bits designate the page offset.
- a) m, n
  - b) n, m
  - c) m - n, m
  - d) m - n, n**
48. With paging there is no \_\_\_\_\_ fragmentation.
- a) internal
  - b) external**
  - c) either type of
  - d) none of the mentioned
49. The operating system maintains a \_\_\_\_\_ table that keeps track of how many frames have been allocated, how many are there, and how many are available.
- a) page
  - b) mapping
  - c) frame**
  - d) memory
50. Paging increases the \_\_\_\_\_ time.
- a) waiting
  - b) execution
  - c) context - switch**
  - d) all of the mentioned

51. Smaller page tables are implemented as a set of \_\_\_\_\_  
a) queues  
b) stacks  
c) counters  
**d) registers**
52. The page table registers should be built with \_\_\_\_\_  
a) very low speed logic  
**b) very high speed logic**  
c) a large memory space  
d) none of the mentioned
53. For larger page tables, they are kept in main memory and a \_\_\_\_\_ points to the page table.  
**a) page table base register**  
b) page table base pointer  
c) page table register pointer  
d) page table base
54. Time taken in memory access through PTBR is \_\_\_\_\_  
a) extended by a factor of 3  
b) extended by a factor of 2  
c) slowed by a factor of 3  
**d) slowed by a factor of 2**
55. Each entry in a translation lookaside buffer (TLB) consists of \_\_\_\_\_  
**a) key**  
b) value  
c) bit value  
d) constant
56. If a page number is not found in the TLB, then it is known as a \_\_\_\_\_  
**a) TLB miss**  
b) Buffer miss  
c) TLB hit  
d) All of the mentioned
57. An \_\_\_\_\_ uniquely identifies processes and is used to provide address space protection for that process.  
a) address space locator  
**b) address space identifier**  
c) address process identifier  
d) none of the mentioned
58. The percentage of times a page number is found in the TLB is known as \_\_\_\_\_  
a) miss ratio  
**b) hit ratio**



- c) miss percent
  - d) none of the mentioned
59. Memory protection in a paged environment is accomplished by \_\_\_\_\_
- a) protection algorithm with each page
  - b) restricted access rights to users
  - c) restriction on page visibility
  - d) protection bit with each page**
60. When the valid – invalid bit is set to valid, it means that the associated page \_\_\_\_\_
- a) is in the TLB
  - b) has data in it
  - c) is in the process's logical address space**
  - d) is the system's physical address space
61. In a paged memory, the page hit ratio is 0.35. The required to access a page in secondary memory is equal to 100 ns. The time required to access a page in primary memory is 10 ns. The average time required to access a page is?
- a) 3.0 ns
  - b) 68.0 ns
  - c) 68.5 ns**
  - d) 78.5 ns
62. To obtain better memory utilization, dynamic loading is used. With dynamic loading, a routine is not loaded until it is called. For implementing dynamic loading \_\_\_\_\_
- a) special support from hardware is required
  - b) special support from operating system is essential
  - c) special support from both hardware and operating system is essential
  - d) user programs can implement dynamic loading without any special support from hardware or operating system**
63. In segmentation, each address is specified by \_\_\_\_\_
- a) a segment number & offset**
  - b) an offset & value
  - c) a value & segment number
  - d) a key & value
64. In paging the user provides only \_\_\_\_\_ which is partitioned by the hardware into \_\_\_\_\_ and \_\_\_\_\_
- a) one address, page number, offset**
  - b) one offset, page number, address
  - c) page number, offset, address
  - d) none of the mentioned
65. Each entry in a segment table has a \_\_\_\_\_
- a) segment base**

- b) segment peak
  - c) segment value
  - d) none of the mentioned
66. The segment limit contains the \_\_\_\_\_
- a) starting logical address of the process
  - b) starting physical address of the segment in memory
  - c) segment length**
  - d) none of the mentioned
67. The offset 'd' of the logical address must be \_\_\_\_\_
- a) greater than segment limit
  - b) between 0 and segment limit**
  - c) between 0 and the segment number
  - d) greater than the segment number
68. When the entries in the segment tables of two different processes point to the same physical location \_\_\_\_\_
- a) the segments are invalid
  - b) the processes get blocked
  - c) segments are shared**
  - d) all of the mentioned
69. The protection bit is 0/1 based on \_\_\_\_\_
- a) write only
  - b) read only
  - c) read - write**
  - d) none of the mentioned
70. If there are 32 segments, each of size 1Kb, then the logical address should have \_\_\_\_\_
- a) 13 bits**
  - b) 14 bits
  - c) 15 bits
  - d) 16 bits
71. Consider a computer with 8 Mbytes of main memory and a 128K cache. The cache block size is 4 K. It uses a direct mapping scheme for cache management. How many different main memory blocks can map onto a given physical cache block?
- a) 2048
  - b) 256
  - c) 64**
  - d) 8
72. A multilevel page table is preferred in comparison to a single level page table for translating virtual address to physical address because \_\_\_\_\_
- a) it reduces the memory access time to read or write a memory location

**b) it helps to reduce the size of page table needed to implement the virtual address space of a process**

c) it is required by the translation lookaside buffer

d) it helps to reduce the number of page faults in page replacement algorithms

**73.** \_\_\_\_ is the concept in which a process is copied into the main memory from the secondary memory according to the requirement.

a) Paging

**b) Demand paging**

c) Segmentation

d) Swapping

**74.** The pager concerns with the \_\_\_\_\_

**a) individual page of a process**

b) entire process

c) entire thread

d) first page of a process

**75.** Swap space exists in \_\_\_\_\_

a) primary memory

**b) secondary memory**

c) cpu

d) none of the mentioned

**76.** When a program tries to access a page that is mapped in address space but not loaded in physical memory, then \_\_\_\_\_

a) segmentation fault occurs

b) fatal error occurs

**c) page fault occurs**

d) no error occurs

**77.** In FIFO page replacement algorithm, when a page must be replaced \_\_\_\_\_

**a) oldest page is chosen**

b) newest page is chosen

c) random page is chosen

d) none of the mentioned

**78.** Which algorithm chooses the page that has not been used for the longest period of time whenever the page required to be replaced?

a) first in first out algorithm

b) additional reference bit algorithm

**c) least recently used algorithm**

d) counting based page replacement algorithm

**79.** A process is thrashing if \_\_\_\_\_

**a) it is spending more time paging than executing**

- b) it is spending less time paging than executing
  - c) page fault occurs
  - d) swapping can not take place
80. In virtual memory, the programmer \_\_\_\_\_ of overlays.
- a) has to take care
  - b) does not have to take care
  - c) all of the mentioned**
  - d) none of the mentioned
81. Virtual memory is normally implemented by \_\_\_\_\_
- a) demand paging**
  - b) buses
  - c) virtualization
  - d) all of the mentioned
82. A swapper manipulates \_\_\_\_\_ whereas the pager is concerned with individual \_\_\_\_\_ of a process.
- a) the entire process, parts
  - b) all the pages of a process, segments
  - c) the entire process, pages**
  - d) none of the mentioned
83. The valid – invalid bit, in this case, when valid indicates?
- a) the page is not legal
  - b) the page is illegal
  - c) the page is in memory**
  - d) the page is not in memory
84. page fault occurs when?
- a) a page gives inconsistent data
  - b) a page cannot be accessed due to its absence from memory**
  - c) a page is invisible
  - d) all of the mentioned
85. When a page fault occurs, the state of the interrupted process is \_\_\_\_\_
- a) disrupted
  - b) invalid
  - c) saved**
  - d) none of the mentioned
86. When a process begins execution with no pages in memory?
- a) process execution becomes impossible
  - b) a page fault occurs for every page brought into memory**
  - c) process causes system crash
  - d) none of the mentioned

87. If the memory access time is denoted by 'ma' and 'p' is the probability of a page fault ( $0 \leq p \leq 1$ ). Then the effective access time for a demand paged memory is \_\_\_\_\_
- a)  $p \times ma + (1-p) \times \text{page fault time}$
  - b)  $ma + \text{page fault time}$
  - c)  **$(1-p) \times ma + p \times \text{page fault time}$**
  - d) none of the mentioned
88. Locality of reference implies that the page reference being made by a process \_\_\_\_\_
- a) will always be to the page used in the previous page reference
  - b) **is likely to be one of the pages used in the last few page references**
  - c) will always be one of the pages existing in memory
  - d) will always lead to page faults
89. Which of the following page replacement algorithms suffers from Belady's Anomaly?
- a) Optimal replacement
  - b) LRU
  - c) **FIFO**
  - d) Both optimal replacement and FIFO
90. A process refers to 5 pages, A, B, C, D, E in the order : A, B, C, D, A, B, E, A, B, C, D, E. If the page replacement algorithm is FIFO, the number of page transfers with an empty internal store of 3 frames is?
- a) 8
  - b) 10
  - c) **9**
  - d) 7
91. A process refers to 5 pages, A, B, C, D, E in the order : A, B, C, D, A, B, E, A, B, C, D, E. If the page replacement algorithm is FIFO, the number of page transfers with an empty internal store of 3 frames is?
- a) 8
  - b) **10**
  - c) 9
  - d) 7
92. A virtual memory system uses First In First Out (FIFO) page replacement policy and allocates a fixed number of frames to a process. Consider the following statements.
- P : Increasing the number of page frames allocated to a process sometimes increases the page fault rate
- Q : Some programs do not exhibit locality of reference
- Which of the following is TRUE?
- a) Both P and Q are true, and Q is the reason for P

b) Both P and Q are true, but Q is not the reason for P

**c) P is false but Q is true**

d) Both P and Q are false

**93.** What is the Optimal page – replacement algorithm?

a) Replace the page that has not been used for a long time

b) Replace the page that has been used for a long time

**c) Replace the page that will not be used for a long time**

d) None of the mentioned

**94.** Optimal page – replacement algorithm is difficult to implement, because

a) it requires a lot of information

**b) it requires future knowledge of the reference string**

c) it is too complex

d) it is extremely expensive

**95.** For 3 page frames, the following is the reference string:

7 0 1 2 0 3 0 4 2 3 0 3 2 1 2 0 1 7 0 1

How many page faults does the LRU page replacement algorithm produce?

a) 10

b) 15

c) 11

**d) 12**

**96.** Applying the LRU page replacement to the following reference string.

1 2 4 5 2 1 2 4

The main memory can accommodate 3 pages and it already has pages 1 and

2. Page 1 came in before page 2.

How many page faults will occur?

a) 2

b) 3

**c) 4**

d) 5

**97.** The minimum number of page frames that must be allocated to a running process in a virtual memory environment is determined by \_\_\_\_\_

**a) the instruction set architecture**

b) page size

c) physical memory size

d) number of processes in memory

**98.** Consider the following page reference string.

1 2 3 4 2 1 5 6 2 1 2 3 7 6 3 2 1 2 3 6

For LRU page replacement algorithm with 4 frames, the number of page faults is?

**a) 10**

- b) 14
- c) 8
- d) 11

**99.** Consider the following page reference string.

1 2 3 4 2 1 5 6 2 1 2 3 7 6 3 2 1 2 3 6

For LRU page replacement algorithm with 5 frames, the number of page faults is?

- a) 10
- b) 14
- c) 8**
- d) 11

**100.** Consider the following page reference string.

1 2 3 4 2 1 5 6 2 1 2 3 7 6 3 2 1 2 3 6

For FIFO page replacement algorithms with 3 frames, the number of page faults is?

- a) 16**
- b) 15
- c) 14
- d) 11

**101.** Consider the following page reference string.

1 2 3 4 2 1 5 6 2 1 2 3 7 6 3 2 1 2 3 6

For FIFO page replacement algorithms with 4 frames, the number of page faults is?

- a) 16
- b) 15
- c) 14**
- d) 11

**102.** Consider the following page reference string.

1 2 3 4 2 1 5 6 2 1 2 3 7 6 3 2 1 2 3 6

For Optimal page replacement algorithms with 3 frames, the number of page faults is?

- a) 16
- b) 15
- c) 14
- d) 11**

**103.** \_\_\_\_\_ is a unique tag, usually a number identifies the file within the file system.

- a) File identifier**
- b) File name
- c) File type
- d) None of the mentioned

104. File type can be represented by \_\_\_\_\_
- a) file name
  - b) **file extension**
  - c) file identifier
  - d) none of the mentioned
105. What is the mounting of file system?
- a) creating of a filesystem
  - b) deleting a filesystem
  - c) **attaching portion of the file system into a directory structure**
  - d) removing the portion of the file system into a directory structure
106. Mapping of file is managed by \_\_\_\_\_
- a) **file metadata**
  - b) page table
  - c) virtual memory
  - d) file system
107. Which one of the following explains the sequential file access method?
- a) random access according to the given byte number
  - b) **read bytes one at a time, in order**
  - c) read/write sequentially by record
  - d) read/write randomly by record
108. A file control block contains the information about \_\_\_\_\_
- a) file ownership
  - b) file permissions
  - c) location of file contents
  - d) **all of the mentioned**
109. Which table contains the information about each mounted volume?
- a) mount table
  - b) system-wide open-file table
  - c) per-process open-file table
  - d) **all of the mentioned**
110. What is raw disk?
- a) **disk without file system**
  - b) empty disk
  - c) disk lacking logical file system
  - d) disk having file system
111. The data structure used for file directory is called \_\_\_\_\_
- a) mount table
  - b) **hash table**
  - c) file table
  - d) process table



112. In which type of allocation method each file occupy a set of contiguous block on the disk?
- a) **contiguous allocation**
  - b) dynamic-storage allocation
  - c) linked allocation
  - d) indexed allocation
113. If the block of free-space list is free then bit will \_\_\_\_\_
- a) **1**
  - b) 0
  - c) any of 0 or 1
  - d) none of the mentioned
114. Which protocol establishes the initial logical connection between a server and a client?
- a) transmission control protocol
  - b) user datagram protocol
  - c) **mount protocol**
  - d) datagram congestion control protocol
115. File attributes consist of \_\_\_\_\_
- a) name
  - b) type
  - c) identifier
  - d) **all of the mentioned**
116. The information about all files is kept in \_\_\_\_\_
- a) swap space
  - b) operating system
  - c) **seperate directory structure**
  - d) none of the mentioned
117. The larger the block size, the \_\_\_\_\_ the internal fragmentation.
- a) **greater**
  - b) lesser
  - c) same
  - d) none of the mentioned
118. Sequential access method \_\_\_\_\_ on random access devices.
- a) **works well**
  - b) doesnt work well
  - c) maybe works well and doesnt work well
  - d) none of the mentioned
119. A relative block number is an index relative to \_\_\_\_\_
- a) **the beginning of the file**
  - b) the end of the file
  - c) the last written position in file

- d) none of the mentioned
120. For large files, when the index itself becomes too large to be kept in memory?
- a) index is called
  - b) **an index is created for the index file**
  - c) secondary index files are created
  - d) all of the mentioned
121. What will happen in the two level directory structure?
- a) **each user has his/her own user file directory**
  - b) the system doesn't its own master file directory
  - c) all of the mentioned
  - d) none of the mentioned
122. What will happen in the two level directory structure?
- a) each user has his/her own user file directory
  - b) the system doesn't its own master file directory
  - c) **all of the mentioned**
  - d) none of the mentioned
123. In the tree structured directories \_\_\_\_\_
- a) the tree has the stem directory
  - b) the tree has the leaf directory
  - c) **the tree has the root directory**
  - d) all of the mentioned
124. The current directory contains, most of the files that are \_\_\_\_\_
- a) **of current interest to the user**
  - b) stored currently in the system
  - c) not used in the system
  - d) not of current interest to the system
125. Which of the following are the types of Path names?
- a) **absolute & relative**
  - b) local & global
  - c) global & relative
  - d) relative & local
126. An absolute path name begins at the \_\_\_\_\_
- a) leaf
  - b) stem
  - c) current directory
  - d) **root**
127. A relative path name begins at the \_\_\_\_\_
- a) leaf
  - b) stem
  - c) **current directory**

- d) root
128. In a tree structure, when deleting a directory that is not empty?
- a) The contents of the directory are safe
  - b) **The contents of the directory are also deleted**
  - c) contents of the directory are not deleted
  - d) none of the mentioned
129. In contiguous allocation \_\_\_\_\_
- a) **each file must occupy a set of contiguous blocks on the disk**
  - b) each file is a linked list of disk blocks
  - c) all the pointers to scattered blocks are placed together in one location
  - d) none of the mentioned
130. In linked allocation \_\_\_\_\_
- a) each file must occupy a set of contiguous blocks on the disk
  - b) **each file is a linked list of disk blocks**
  - c) all the pointers to scattered blocks are placed together in one location
  - d) none of the mentioned
131. In indexed allocation \_\_\_\_\_
- a) each file must occupy a set of contiguous blocks on the disk
  - b) each file is a linked list of disk blocks
  - c) **all the pointers to scattered blocks are placed together in one location**
  - d) none of the mentioned
132. On systems where there are multiple operating system, the decision to load a particular one is done by \_\_\_\_\_
- a) **boot loader**
  - b) bootstrap
  - c) process control block
  - d) file control block
133. \_\_\_\_\_ and \_\_\_\_\_ are the most common strategies used to select a free hole from the set of available holes.
- a) **First fit, Best fit**
  - b) Worst fit, First fit
  - c) Best fit, Worst fit
  - d) None of the mentioned
134. To solve the problem of external fragmentation \_\_\_\_\_ needs to be done periodically.
- a) **compaction**
  - b) check
  - c) formatting
  - d) replacing memory

135. Contiguous allocation has two problems \_\_\_\_\_ and \_\_\_\_\_ that linked allocation solves.
- a) **external – fragmentation & size – declaration**
  - b) internal – fragmentation & external – fragmentation
  - c) size – declaration & internal – fragmentation
  - d) memory – allocation & size – declaration
136. Indexed allocation \_\_\_\_\_ direct access.
- a) **supports**
  - b) does not support
  - c) is not related to
  - d) none of the mentioned
137. Consider a disk where blocks 2, 3, 4, 5, 8, 9, 10, 11, 12, 13, 17, 18, 25, 26 and 27 are free and the rest of the blocks are allocated. Then the free space bitmap would be \_\_\_\_\_
- a) 10000110000001110011111100011111...
  - b) 11000011000000111001111110001111...
  - c) 01111001111110001100000011100000...
  - d) **001111001111110001100000011100000...**
138. The heads of the magnetic disk are attached to a \_\_\_\_\_ that moves all the heads as a unit.
- a) spindle
  - b) **disk arm**
  - c) track
  - d) none of the mentioned
139. The time taken to move the disk arm to the desired cylinder is called the \_\_\_\_\_
- a) positioning time
  - b) random access time
  - c) **seek time**
  - d) rotational latency
140. The time taken for the desired sector to rotate to the disk head is called \_\_\_\_\_
- a) positioning time
  - b) random access time
  - c) seek time
  - d) **rotational latency**
141. What is the host controller?
- a) controller built at the end of each disk
  - b) **controller at the computer end of the bus**
  - c) all of the mentioned
  - d) none of the mentioned

- 142.** Consider a disk queue with requests for I/O to blocks on cylinders.  
98 183 37 122 14 124 65 67  
Considering FCFS (first cum first served) scheduling, the total number of head movements is, if the disk head is initially at 53 is?
- a) 600
  - b) 620
  - c) 630
  - d) 640**
- 143.** Consider a disk queue with requests for I/O to blocks on cylinders.  
98 183 37 122 14 124 65 67  
Considering SSTF (shortest seek time first) scheduling, the total number of head movements is, if the disk head is initially at 53 is?
- a) 224
  - b) 236**
  - c) 245
  - d) 240
- 144.** SSTF algorithm, like SJF \_\_\_\_\_ of some requests.
- a) may cause starvation**
  - b) will cause starvation
  - c) does not cause starvation
  - d) causes aging
- 145.** In the \_\_\_\_\_ algorithm, the disk arm starts at one end of the disk and moves toward the other end, servicing requests till the other end of the disk. At the other end, the direction is reversed and servicing continues.
- a) LOOK
  - b) SCAN**
  - c) C-SCAN
  - d) C-LOOK
- 146.** In the \_\_\_\_\_ algorithm, the disk head moves from one end to the other, servicing requests along the way. When the head reaches the other end, it immediately returns to the beginning of the disk without servicing any requests on the return trip.
- a) LOOK
  - b) SCAN
  - c) C-SCAN**
  - d) C-LOOK
- 147.** In the \_\_\_\_\_ algorithm, the disk arm goes as far as the final request in each direction, then reverses direction immediately without going to the end of the disk.
- a) LOOK**
  - b) SCAN

- c) C-SCAN
- d) C-LOOK

**148.** The \_\_\_\_\_ program initializes all aspects of the system, from CPU registers to device controllers and the contents of main memory, and then starts the operating system.

- a) main
- b) bootloader
- c) bootstrap**
- d) rom

**149.** In SCSI disks used in high end PCs, the controller maintains a list of \_\_\_\_\_ on the disk. The disk is initialized during \_\_\_\_\_ which sets aside spare sectors not visible to the operating system.

- a) destroyed blocks, high level formatting
- b) bad blocks, partitioning
- c) bad blocks, low level formatting**
- d) destroyed blocks, partitioning

**150.** In SCSI disks used in high end PCs, the controller maintains a list of bad blocks on the disk. The disk is initialized during low-level formatting which sets aside spare sectors not visible to the operating system. The scheme used is known as \_\_\_\_\_ or \_\_\_\_\_

- a) sector sparing & forwarding**
- b) forwarding & sector utilization
- c) backwarding & forwarding
- d) sector utilization & backwarding