Cubical Abstract Machine

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We will write $M[\mathcal{E}]$ for substituting the free dimensions and variables of a term M using an environment \mathcal{E} .

Stack frames and machine states A stack frame (K, \mathcal{E}, Ψ) represents a continuation whose whole binds Ψ dimensions. For a stack π , let $\vec{\Psi}$ be the aggregation of all the dimension bindings Ψ mentioned in π ; then the extension of a stack $(K, \mathcal{E}, \Phi) :: \pi$ is wellformed when \mathcal{E} is an environment for $\mathbf{FD}(K) \setminus \vec{\Psi}$; moreover, K is allowed to mention the dimensions in $\vec{\Psi}$.

A machine state $\langle M \triangleleft \mathcal{E} \parallel \pi \rangle$, where $\vec{\Psi}$ is the aggregation of dimension bindings in π , is wellformed when \mathcal{E} is an environment for $\mathbf{FD}(M) \setminus \vec{\Psi}$; moreover, M is allowed to mention the dimensions in $\vec{\Psi}$.

Selected transition rules We define a transition judgment $C \succ_{\mu} C'$ with μ ranging over the following possible modes:

- 1. , denoting a cubically stable transition
- 2. •, denoting a non-cubically stable transition

Moreover, if $\mathcal{C} \succ_{\square} \mathcal{C}'$ then $\mathcal{C} \succ_{\bullet} \mathcal{C}'$.

$$\langle a \triangleleft (\gamma, \psi) \parallel \pi \rangle \succ_{\text{fil}} \langle \gamma(a) \triangleleft (\gamma, \psi) \parallel \pi \rangle \tag{1}$$

$$\langle M(N) \triangleleft \mathcal{E} \parallel \pi \rangle \succ_{\text{fil}} \langle M \triangleleft \mathcal{E} \parallel (\square(N), \mathcal{E}, []) :: \pi \rangle \tag{2}$$

$$\langle \lambda a. M \triangleleft \mathcal{E} \parallel (\square(N), \mathcal{E}', []) :: \pi \rangle \succ_{\text{fil}} \langle M \triangleleft [\mathcal{E}, a \hookrightarrow N \triangleleft \mathcal{E}] \parallel \pi \rangle \tag{3}$$

$$\langle \mathsf{fst}(M) \triangleleft \mathcal{E} \parallel \pi \rangle \succ_{\square} \langle M \triangleleft \mathcal{E} \parallel (\mathsf{fst}(\square), \mathcal{E}, []) :: \pi \rangle \tag{4}$$

$$\langle \operatorname{snd}(M) \triangleleft \mathcal{E} \parallel \pi \rangle \succ_{\square} \langle M \triangleleft \mathcal{E} \parallel (\operatorname{snd}(\square), \mathcal{E}, []) :: \pi \rangle \tag{5}$$

$$\langle \langle M, N \rangle \triangleleft \mathcal{E} \parallel (\mathsf{fst}(\square), \mathcal{E}', []) :: \pi \rangle \succ_{\square} \langle M \triangleleft \mathcal{E} \parallel \pi \rangle \tag{6}$$

$$\langle \langle M, N \rangle \triangleleft \mathcal{E} \parallel (\operatorname{snd}(\square), \mathcal{E}', []) :: \pi \rangle \succ_{\square} \langle N \triangleleft \mathcal{E} \parallel \pi \rangle \tag{7}$$

$$\langle \mathsf{loop}_r \triangleleft (\gamma, \psi) \parallel \pi \rangle \succ_{\square} \langle \mathsf{base} \triangleleft (\gamma, \psi) \parallel \pi \rangle \text{ when } r\psi = \epsilon \tag{8}$$

(9)

$$\langle \mathsf{coe}_{x,A}^{r \leadsto r'}(M) \triangleleft \mathcal{E} \parallel \pi \rangle \succ_{\square} \langle A \triangleleft \mathcal{E} \parallel (\mathsf{coe}_{\square}^{r \leadsto r'}(M), \mathcal{E}', [x]) :: \pi \rangle \tag{10}$$

$$\langle \mathsf{bool} \triangleleft \mathcal{E} \parallel (\mathsf{coe}_{\square}^{r \leadsto r'}(M), \mathcal{E}', [x]) :: \pi \rangle \succ_{\mathsf{fil}} \langle M \triangleleft \mathcal{E}' \parallel \pi \rangle \tag{11}$$

Unloading the machine We can unload the machine at any time; this is useful if we are computing an open term and hit a variable.

$$\frac{\langle K[\Psi.M[\mathcal{E}]] \triangleleft \mathcal{E}' \parallel \pi \rangle \Longrightarrow N}{\langle M \triangleleft \mathcal{E} \parallel \cdot \rangle \Longrightarrow M[\mathcal{E}]}$$

$$\frac{\langle K[\Psi.M[\mathcal{E}]] \triangleleft \mathcal{E}' \parallel \pi \rangle \Longrightarrow N}{\langle M \triangleleft \mathcal{E} \parallel (K[\square], \mathcal{E}', \Psi) :: \pi \rangle \Longrightarrow N}$$