



جامعة أم القرى  
UMM AL-QURA UNIVERSITY

*Algorithms Programming*

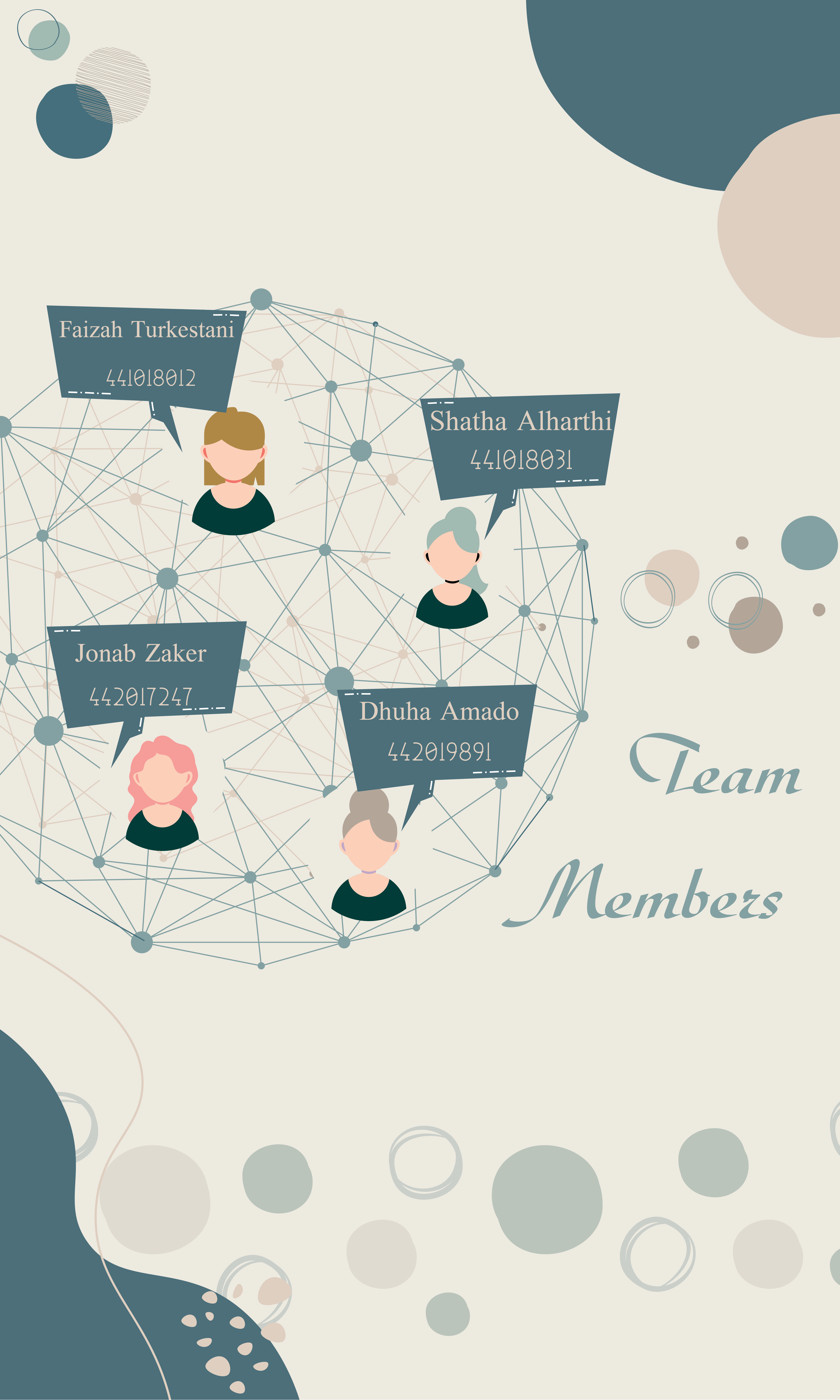
# *Network Flow*

Project Report

*Instructor : Areej Al-sini*

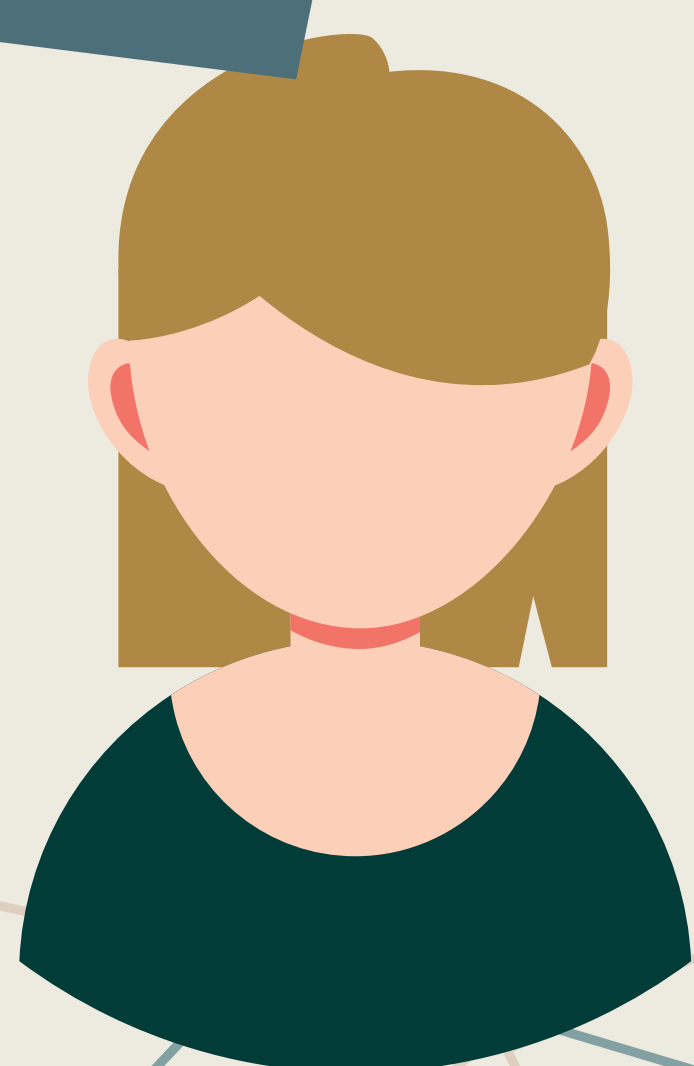
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Faizah Turkestani

441018012



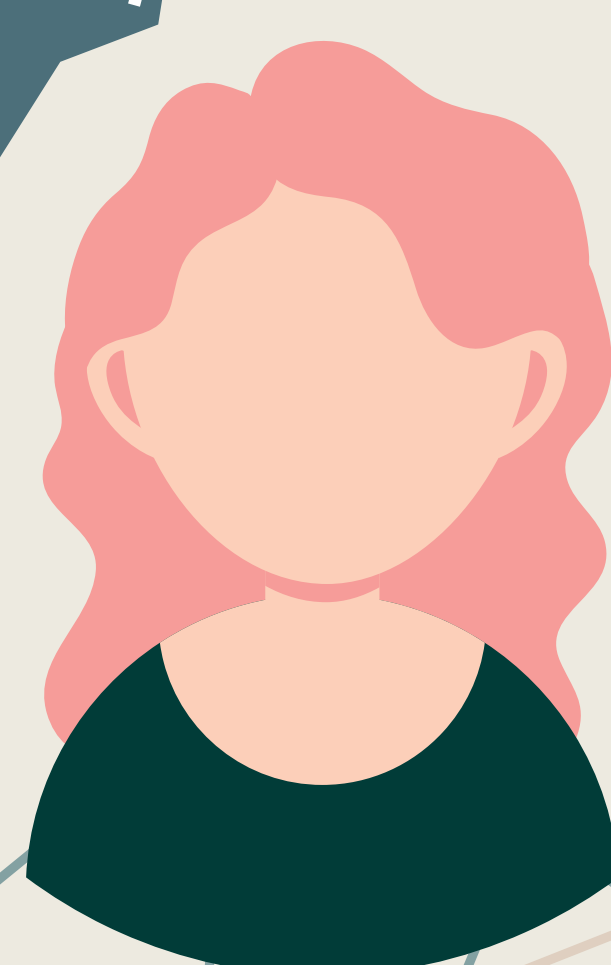
Shatha Alharthi

441018031



Jonab Zaker

442017247



Dhuha Amado

442019891



*Team  
Members*



# Network Flow

is a directed graph where each edge has a capacity and each edge receives a flow. The amount of flow on an edge cannot exceed the capacity of the edge.

Often in operations research, a directed graph is called a network. A flow must satisfy the restriction that the amount of flow into a node equals the amount of flow out of it, unless it is a source, which has only outgoing flow, or sink, which has only incoming flow. A network can be used to model traffic in a computer network or anything similar in which something travels through a network of nodes.

**DINC'S ALGORITHM**

**EDMONDS' ALGORITHM**

**FORD-FULKERSON ALGORITHM**



# Dinic's Algorithm

is a strongly polynomial algorithm for computing the maximum flow in a flow network, The algorithm runs in  $O(V^2E)$  time . in that it uses .shortest augmenting paths

The Data structures used are :

- GRAPH
- ADJACENCY LIST
- ADJACENCY MATRIX





# Algorithm

## Pseudocode

DINIC(SOURCE, SINK, ADJACENCY\_LIST)

INITIALIZE THE LEVEL GRAPH AND THE RESIDUAL GRAPH

WHILE BFS(SOURCE, SINK)

WHILE THERE IS AN AUGMENTING PATH FROM SOURCE TO SINK IN THE RESIDUAL GRAPH

CALCULATE THE BOTTLENECK CAPACITY OF THE PATH

AUGMENT THE FLOW ALONG THE PATH BY THE BOTTLENECK CAPACITY

RETURN THE FLOW IN THE RESIDUAL GRAPH

BFS(SOURCE, SINK)

INITIALIZE THE LEVEL GRAPH WITH -1 FOR ALL NODES EXCEPT THE SOURCE NODE, WHICH HAS LEVEL 0

INITIALIZE A QUEUE AND ADD THE SOURCE NODE TO IT

WHILE THE QUEUE IS NOT EMPTY

DEQUEUE A NODE  $u$  FROM THE QUEUE

FOR EACH ADJACENT NODE  $v$  OF  $u$

IF THE LEVEL OF  $v$  IS -1 AND THERE IS AVAILABLE CAPACITY IN THE RESIDUAL GRAPH FROM  $u$  TO  $v$

SET THE LEVEL OF  $v$  TO THE LEVEL OF  $u + 1$

ADD  $v$  TO THE QUEUE

END IF

END FOR

END WHILE

RETURN THE LEVEL OF SINK NODE IS NOT -1



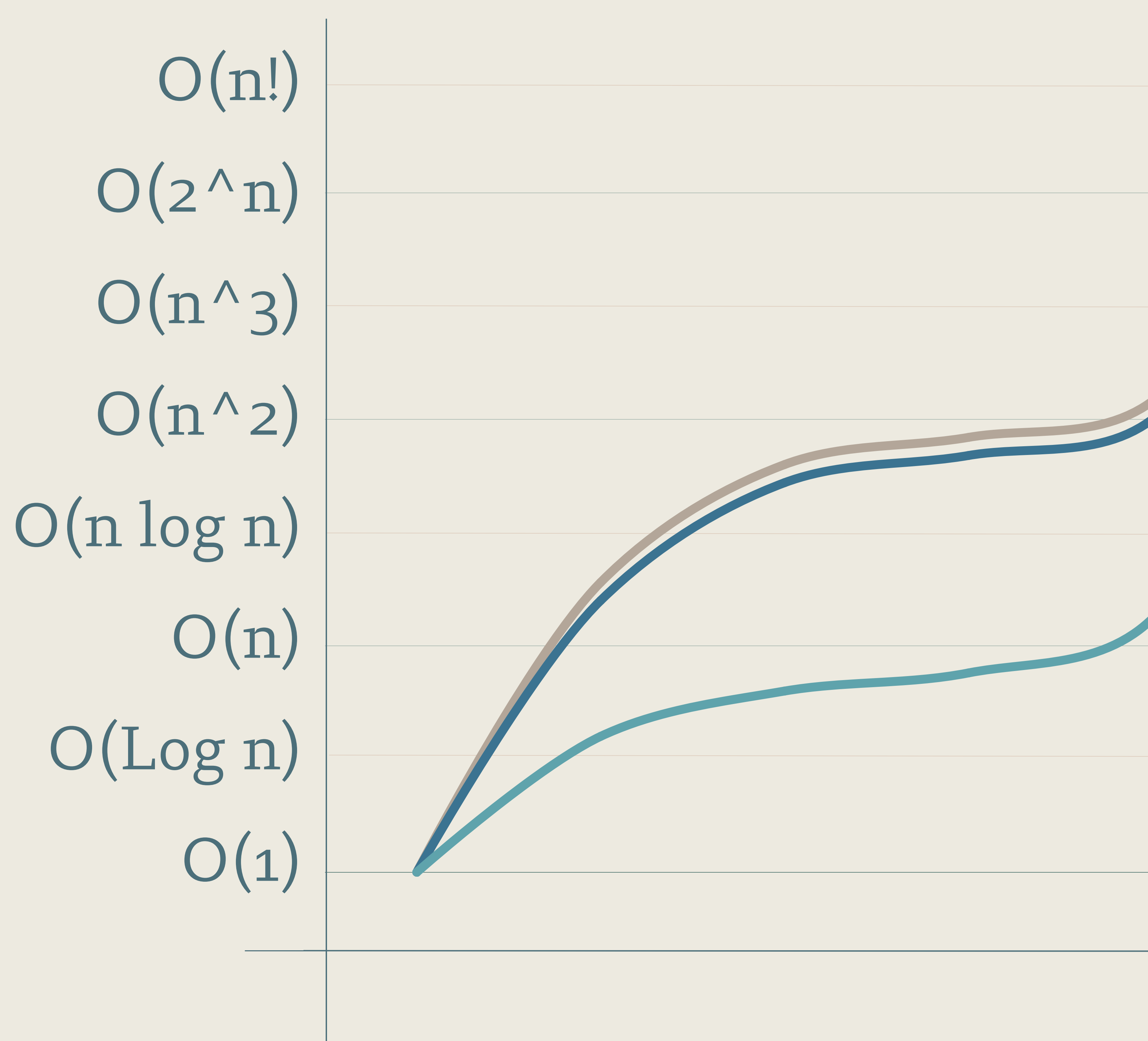


# Algorithm

## Discussion

As we observe that arraylist is the best data structure for this algorithm because it takes  $O(n)$  while array matrix and graph take  $O(n^2)$

### BIG(O) NOTATION



- GRAPH
- ADJACENCY LIST
- ADJACENCY MATRIX





# Edmonds-Karp algorithm

The Edmonds-Karp Algorithm is a specific implementation of the Ford-Fulkerson algorithm. Like Ford-Fulkerson, Edmonds-Karp is also an algorithm that deals with the max-flow min-cut problem.

is an algorithm for finding a spanning arborescence of minimum weight (sometimes called an optimum branching). It is the directed analog of the minimum spanning tree problem.

The Data structures used are :

- **STACK & 2D ARRAY**
- **QUEUE & GRAPH**
- **QUEUE & 2D ARRAY**



# Edmonds-Karp algorithm

## Pseudocode

INPUT

$C[N \times N]$  : CAPACITY MATRIX

$E[N \times N]$  : ADJACENCY MATRIX

$S$  : SOURCE

$T$  : SINK

OUTPUT

$F$  : MAXIMUM FLOW

EDMONDS-KARP:

$F = 0$

$F = [N \times N]$

WHILE TRUE:

$M, P = \text{BREADTH-FIRST-SEARCH}(C, E, S, T, F)$

IF  $M = 0$ :

BREAK

$F = F + M$

$V = T$

WHILE  $V \neq S$ :

$U = P[V]$

$F[U, V] = F[U, V] - M$

$F[V, U] = F[V, U] + M$

$V = U$

RETURN



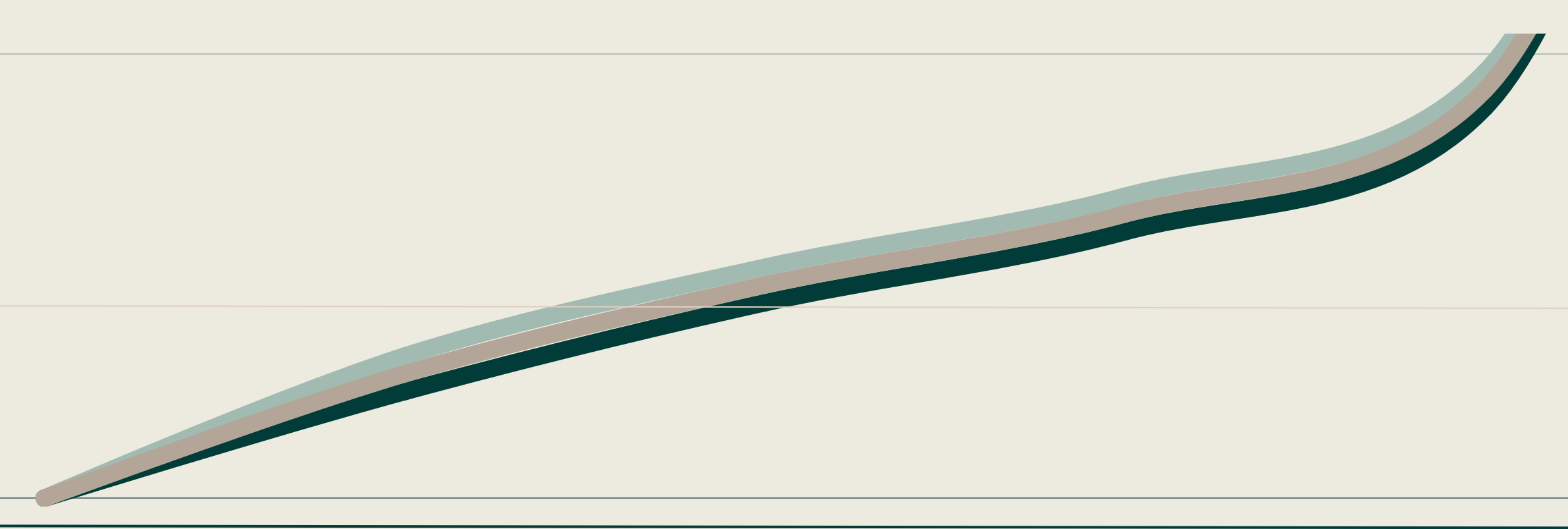
# Edmonds-karp algorithm

## Discussion

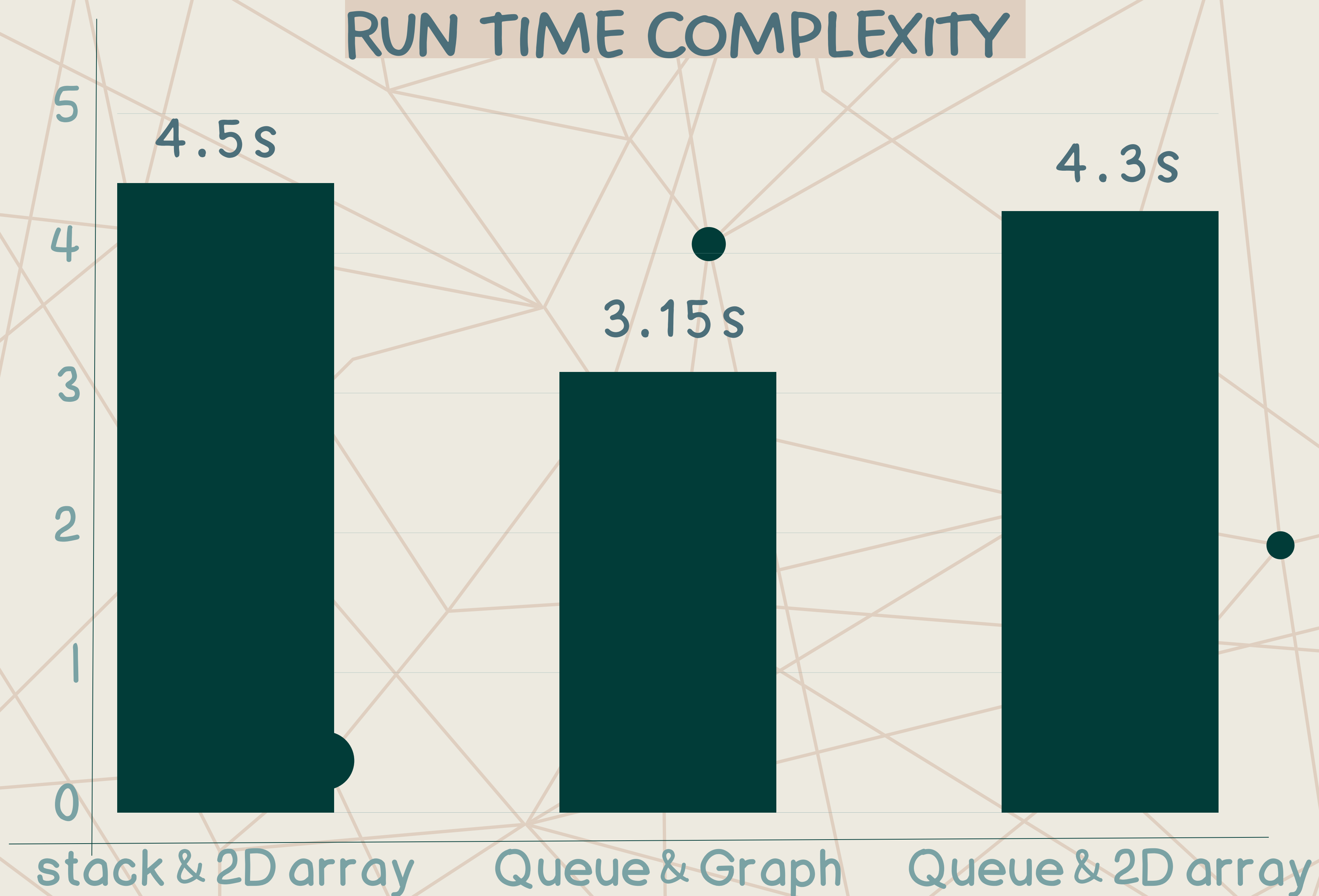
We observed that all data structures take the same Big O, but (Queue with Graph) was the best of them of Runtime Complexity and we concluded that by reading 124675 data.

## BIG(O) NOTATION

$O(n!)$   
 $O(2^n)$   
 $O(n^3)$   
 $O(n^2)$   
 $O(n \log n)$   
 $O(n)$   
 $O(\log n)$   
 $O(1)$



## RUN TIME COMPLEXITY



- STACK & 2D ARRAY
- QUEUE & GRAPH
- QUEUE & 2D ARRAY



# Ford-Fulkerson algorithm

is a greedy algorithm that computes the maximum flow in a flow network.

The idea behind the algorithm is as follows: as long as there is a path from the source (start node) to the sink (end node), with available capacity on all edges in the path, we send flow along one of the paths. Then we find another path, and so on

The Data structures used are :

- ADJACENCY LIST
- ADJACENCY GRAPH
- ADJACENCY ARRAY



# Ford-Fulkerson algorithm

## ● Pseudocode

FORD\_FULKERSON( $G, s, t$ )

FOR EACH EDGE  $(u, v) \in E[G]$  DO

$F[u, v] \leftarrow 0$

$F[v, u] \leftarrow 0$

WHILE THERE EXISTS A PATH  $P$  FROM  $s$  TO  $t$  IN THE RESIDUAL NETWORK  $G_f$  DO

$CF(P) \leftarrow \min \{ CF(u, v) : (u, v) \text{ IS IN } P \}$

    FOR EACH EDGE  $(u, v)$  IN  $P$  DO

$F[u, v] \leftarrow F[u, v] + CF(P)$

$F[v, u] \leftarrow -F[u, v]$

    END FOR

REBUILD  $G$  BASED ON NEW FLOW  $F$

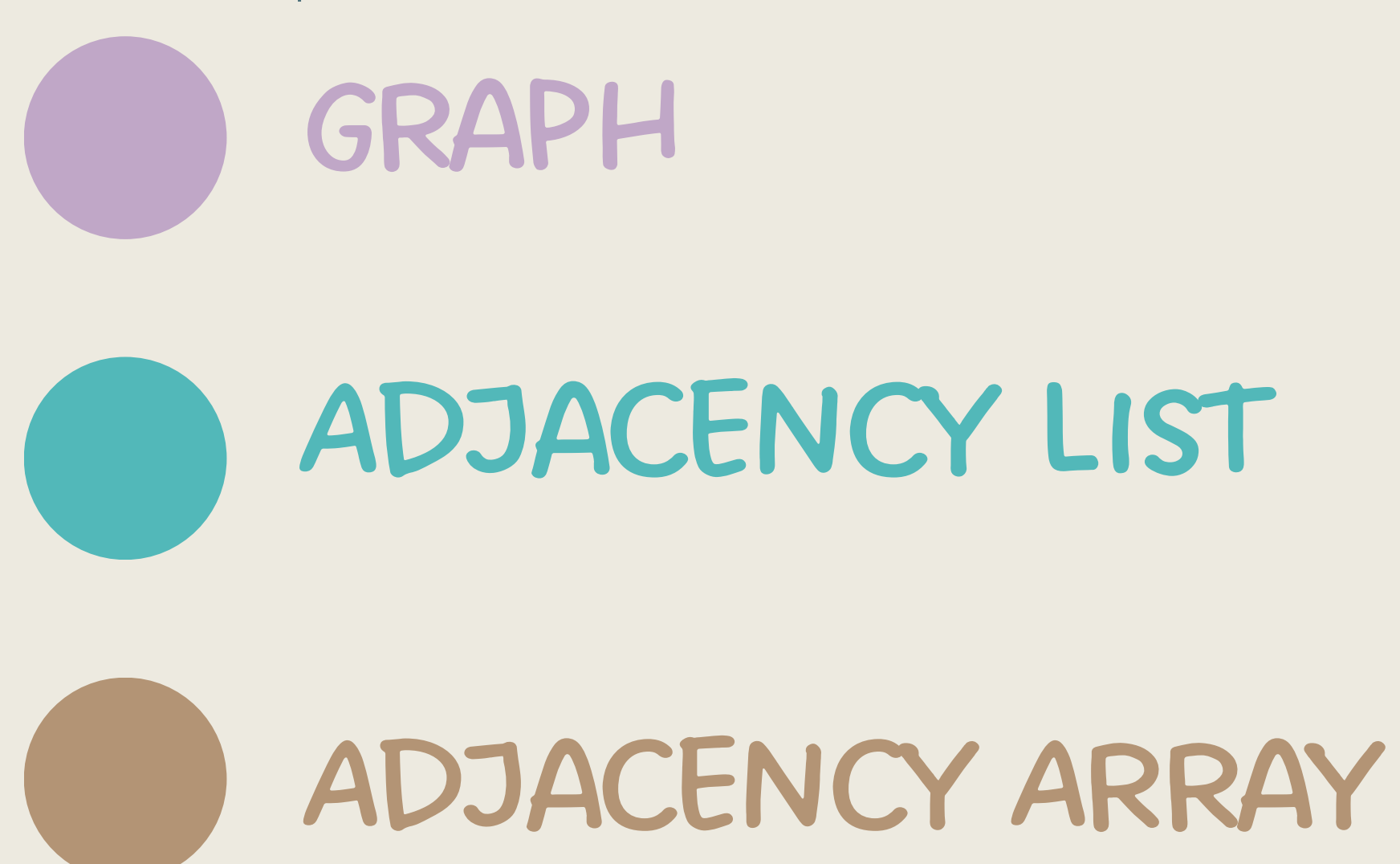
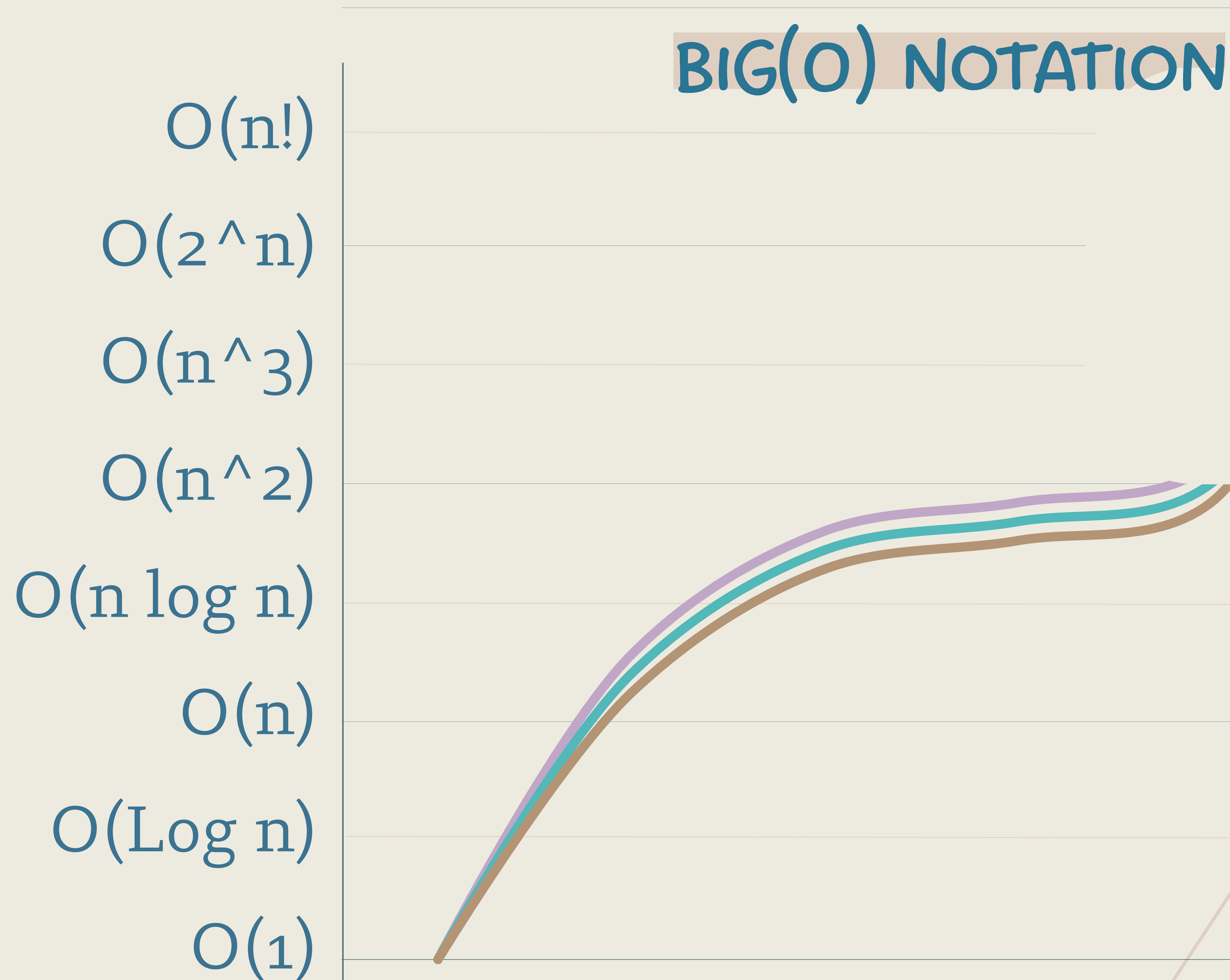
END WHILE



# Ford-Fulkerson algorithm

## Discussion

We also observation that all data structures take the same Big O, but (Adjacency List) was the best of them of Runtime Complexity and we concluded that by reading 12213 data.







*Run*

*Time Device*

*Specifications*

LENOVO IDEAPAD

SDD: 1TB - RAM: 16GB

AMD RYZEN7 | 2.9GHZ

SPYDER - VERSION 5

PROGRAMMING LANGUAGE:PYTHON



Big(O) Notation

Data Structure	Graph	Adjacency List	Adjacency matrix
Dinic's algorithm	$O(n^2)$	$O(n)$	$O(n^2)$
Data Structure	Stack & 2D Array	Queue & Graph	Queue & 2D Array
Edmonds' - Karp algorithm	$O(\log n)$	$O(\log n)$	$O(\log n)$
Data Structure	Adjacency Array	Graph	Adjacency List
Ford Fulkerson algorithm	$O(n^2)$	$O(n^2)$	$O(n^2)$

Best Algorithm

We observation the best Algorithm for Network Flow is EDMONONDS'-KARP Algorithm because it's take  $O(\log n)$

Run Time Complexity

Data Structure	Graph	Adjacency List	Adjacency matrix
Dinic's algorithm	-	-	-
Data Structure	Stack & 2D Array	Queue & Graph	Queue & 2D Array
Edmonds' algorithm	4.5 s	3.15s	4.3s
Data Structure	Adjacency Array	Graph	Adjacency List
Ford Fulkerson algorithm	52.09s	1.89s	0.06s



# References

## *Dinic's Algorithm*

<https://www.geeksforgeeks.org/dinics-algorithm-maximum-flow/>

<https://iq.opengenus.org/dinics-algorithm/#algorithm>

<https://github.com/williamfiset/Algorithms/blob/master/src/main/java/com/williamfiset/algorithms>

[/graphtheory/networkflow/Dinics.java](#)

<https://www.topcoder.com/thrive/articles/edmonds-karp-and-dinics-algorithms-for-maximum-flow>

## *Edmonds' - Karp Algorithm*

Edmonds Karp Algorithm for maximum flow (opengenus.org)

python - Edmonds—Karp time complexity - Stack Overflow

<https://brilliant.org/wiki/edmonds-karp-algorithm/>

Creating capacity graph for edmonds karp maximum flow algorithm in Python - Stack Overflow

## *Ford- Fulkerson Algorithm*

<https://www.programiz.com/dsa/ford-fulkerson-algorithm>

<https://www.autoscripts.net/pseudocode-of-ford-fulkerson/>



I hope you  
like it and  
satisfy you.