FYS3150 oppgavesett 1

Jon Aleksander Prøitz and Marius Torsvoll

Relevant code can be found at: https://github.com/Jonaproitz/Project_1

PROBLEM 1.

Given

$$u(x) = 1 - (1 - e^{-10})x - e^{-10x}$$
(1)

Inserting x = 0

$$u(0) = 1 - (1 - e^{-10}) \cdot 0 - e^{-10 \cdot 0} = 1 - 0 - 1 = 0$$

and x = 1

$$u(1) = 1 - (1 - e^{-10}) \cdot 1 - e^{-10 \cdot 1} = 1 - 1 + e^{-10} - e^{-10} = 0$$

furthermore the one-dimensional possion equation can be written

$$-\frac{d^2u}{dx^2} = -\frac{d^2}{dx^2}\left(1 - (1 - e^{-10})x - e^{-10x}\right) = -\frac{d}{dx}\left((1 - e^{-10}) + 10e^{-10x}\right) = 100e^{-10x} = f(x)$$

Hence equation 1 is an exact solution to our problem.

PROBLEM 2.

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PROBLEM 3.

The one-dimensional poisson equation can be written

$$-\frac{d^2u}{dx^2} = \frac{u(x-h) + 2u(x) - u(x+h)}{h^2} + O(h^2) = f(x)$$

Discretizing x with m values and a given distance h between each distinct value then gives

$$x \to x_0, x_1, x_2, ..., x_{m-1}$$

 $u(x) \to u_0, u_1, u_2, ..., u_{m-1}$
 $f(x) \to f_0, f_1, f_2, ..., f_{m-1}$

with $u_i = v_i$, such that

$$-\frac{d^2v_i}{dx^2} = -v_{i-1} + 2v_i - v_{i+1} = f_i h^2$$

PROBLEM 4.

The set of equations from problem 3 can be written as

$$-v_0 + 2v_1 - v_2 = h^2 f_1$$

$$-v_1 + 2v_2 - v_3 = h^2 f_2$$

$$-v_2 + 2v_3 - v_4 = h^2 f_3$$

$$\vdots$$

$$-v_{m-3} + 2v_{m-2} - v_{m-1} = h^2 f_{m-2}$$

Wich for

$$g_{1} = h^{2} f_{1} + v_{0}$$

$$g_{2} = h^{2} f_{2}$$

$$g_{3} = h^{2} f_{3}$$

$$\vdots$$

$$g_{m-3} = h^{2} f_{m-3}$$

$$q_{m-1} = h^{2} f_{m-2} + v_{m-1}$$

can be written as the matrix equation

$$\begin{pmatrix} 2 & -1 & 0 & 0 & 0 & \dots \\ -1 & 2 & -1 & 0 & 0 & \dots \\ 0 & -1 & 2 & -1 & 0 & \dots \\ \vdots & \vdots & \ddots & \ddots & \ddots & \vdots \\ 0 & 0 & 0 & -1 & 2 & -1 \\ 0 & 0 & 0 & 0 & -1 & 2 \end{pmatrix} \begin{pmatrix} v_1 \\ v_2 \\ v_3 \\ \vdots \\ v_{m-3} \\ v_{m-2} \end{pmatrix} = \begin{pmatrix} g_1 \\ g_2 \\ g_3 \\ \vdots \\ g_{m-3} \\ g_{m-2} \end{pmatrix}$$

on the form $A\vec{v} = \vec{g}$.

PROBLEM 5

a

When finding the matrix, A, in problem 4 it is assumed that v_0 and v_{m-1} are known. Hence theese values are not calculated and

$$n = m - 2$$

b

Since m = n + 2 and \vec{v}^* is of length m the vector thus has to be given by the following:

$$\vec{v}^* = \left[v_0, \vec{V}, v_m\right] \tag{2}$$

This is a result of task 4 as the first and last elements of the equation set is known.

PROBLEM 6

 \mathbf{a}

A general tridiagonal matrix is given by the following: syntax:

Forward substitution gives us the following relation:

$$R_2 \to R_2 - \frac{a_2}{b_1} R_1 = 0 \ b_2 - \frac{a_2}{b_1} c_1 \ 0 \ 0$$

Continuing this substitution with the variables gives us the following results for the variables:

$$\begin{split} \tilde{b}_1 &= b_1 \quad \text{and} \quad \tilde{b}_i = b_i - \frac{a_i}{\tilde{b}_{i-1}} c_{i-1}, \qquad \text{For i} = 2,3,4...., \text{n} \\ \tilde{g}_1 &= g_1 \quad \text{and} \quad \tilde{g}_i = g_i - \frac{a_i}{\tilde{q}_{i-1}} \tilde{g}_{i-1}, \qquad \text{For i} = 2,3,4...., \text{n} \end{split}$$

This then gives us the matrix:

We then see that:

$$R_n = \frac{R_n}{\tilde{b}_n} = 0 \ 0$$

and that:

$$\frac{\tilde{g}_n}{\tilde{b}_n} = v_n$$

Bacwards substitution then gives us the following:

$$v_i = \frac{\tilde{g}_i - c_i v_{i+1}}{\tilde{b}_i}$$
, For $i = n - 1, n - 2, n - 3, \dots 2, 1$

b

This algorithm is thus:

$$\tilde{b}_1 = b_1$$

$$\tilde{b}_i = b_i - \frac{a_i}{b_{i-1}} c_{i-1}, \text{ For } i = 2, 3, 4 \dots n$$

$$\tilde{g}_1 = g_1 \tilde{g}_i = g_i - \frac{a_i}{\tilde{b}_{i-1}} \tilde{g}_{i-1}, \text{ For } i = 2, 3, 4 \dots n$$

$$v_n = \frac{\tilde{g}_n}{\tilde{b}_n}$$

$$v_i = \frac{\tilde{g}_i - c_i v_{i+1}}{\tilde{b}_i}, \text{ For } i = n-1, n-2, n-3, \dots, 1$$

The total amount of FLOPs will then be described as the sum of FLOPs per n. Thus the FLOPs will then be described as:

$$8(n-1)(+1)$$

TASK 7

See the GitHub link

TASK 8

See the GitHub link

TASK 9

 \mathbf{a}

By making the special algorithm we then get the following results.

$$\tilde{b}_1 = 2$$

$$\tilde{b}_i = 2 - \frac{1}{\tilde{b}_{i-1}}$$

$$\tilde{g}_1 = g_1$$

$$\tilde{g}_i = g_i + \frac{1}{\tilde{b}_{i-1}} \tilde{g}_{i-1}$$

$$v_n = \frac{\tilde{g}_n}{\tilde{b}_n}$$

$$v_i = \frac{\tilde{g}_i + v_{i+1}}{\tilde{b}_i}$$

b

The number of flops for this special algorithm will then be described as:

$$6(n-1)(+1)$$

 \mathbf{c}

See the GitHub link

TASK 10

See the GitHub link

TASK 11

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