Community detection with spectral methods

Laurent Massoulié

MSR-Inria Joint Centre

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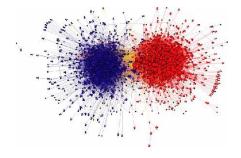


Community detection

Definition: Identification of groups of similar items, based on observed interactions between items

Observation: a graph (e.g. represented by its adjacency matrix); edges could be directed or not, labeled or not.

 \rightarrow Same as clustering, specialized to "graphical" observations



Example: From observations of citation graph between blog posts during US presidential campaign, partitioning into democrats / republicans communities

From "friendship" graph of facebook infer communities of "similar" users to guide recommendations of potential new contacts



An example of assortative communities: stronger connectivity within than across communities

From "friendship" graph of facebook infer communities of "similar" users to guide recommendations of potential new contacts



An example of *assortative* communities: stronger connectivity within than across communities

Variation: NSA's "co-traveler programme": spot group of suspect persons meeting regularly in unusual places

From matrix of user ratings of items, infer communities of "similar" items to guide item recommendations ("users who liked this also liked that")

User / Movie	f_1	f_2	 f_m
u_1	?	**	***
u_2	***	?	?
u_n	****	**	**

 \rightarrow e.g. Netflix "Cinematch" movie recommendation engine (see Netflix prize)



Lists of proteins involved in chemical reactions of cell biology

- \rightarrow Graph: co-involvement in some reaction
- → Infer groups of proteins with same functional role (see "Functional cartography of complex metabolic networks", R. Guimera, L. Nunes Amaral, Nature 2005 http://www.ncbi.nlm.nih.gov/pmc/articles/PMC2175124/)

Potentially *disassortative* communities (connections within community may be rarer than across)



Knowledge graphs

A generic representation of data:

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A1 has with B1 interaction of type C1
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A2 has with B2 interaction of type C2

A3 has with B3 interaction of type C3

...

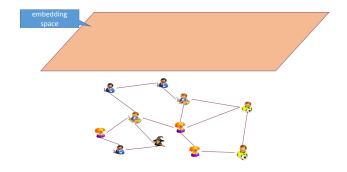
ightarrow Popularized by companies like Google as a universal format

Outline

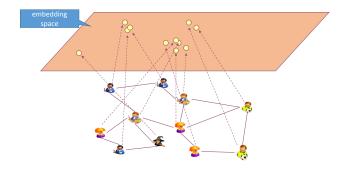
Outline

- Spectral methods
- The Stochastic Block Model (SBM)
- Consistent community detection in a "high signal" regime
- Tools: linear algebra and random matrices

Embedding



Embedding



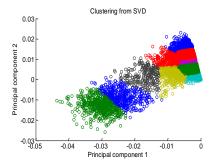
- Extract top two (more generally top k) eigenvalues λ_1, λ_2 of graph's adjacency matrix $A \in \mathbb{R}^{n \times n}$ (ordered by absolute value: $|\lambda_1| \ge \lambda_2 \ge \cdots$)
- Let $x_1, x_2 \in \mathbb{R}^n$: corresponding normalized eigenvectors
- Embed vertex $k \in [n]$ into \mathbb{R}^2 by letting $z_k := \sqrt{n}(x_1(k), x_2(k))$

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- → Also known as Principal Component Analysis (PCA)

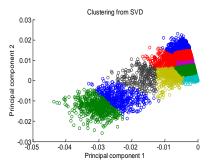
Due to Karl Pearson "On Lines and Planes of Closest Fit to Systems of Points in Space", 1901



Example: 2D-spectral embedding of Netflix prize data



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From embedding to clustering: K-means algorithm

- **1** Choose K cluster centers $y_1, \ldots, y_K \in \mathbb{R}^2$
- ② Form clusters $C_k := \{i \in [n] : ||y_k z_i|| = \min_{\ell \in [K]} \{||y_\ell z_i||\} \}$
- **3** Reset y_k to cluster average $\frac{1}{|C_k|} \sum_{i \in C_k} z_i$
- Go back to step 2



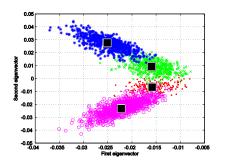
The stochastic block model

A multi-type version of the Erdős-Rényi random graph

- *n* vertices partitioned into *K* communities
- Type (community) of node $i : \sigma(i) \in [K]$
- For each $k \in [K]$, number $\sum_{i} \mathbb{I}_{\sigma(i)=k}$ of type-k-nodes: $\sim \alpha_k n$ for fixed $\alpha_k > 0$
- For each pair $i, j \in [n]$: edge (i, j) present with probability $B_{\sigma(i),\sigma(j)}\frac{d}{n}$, where $B \in \mathbb{R}_+^{K \times K}$: fixed matrix, and d: may increase as $n \to \infty$



Example: spectral embedding for SBM



A case with K = 4 communities Spectral embedding seems to reflect community structure \rightarrow Why / when do spectral methods work?

Theorem

Assume communities are **distinguishable**, i.e. for each $k \neq \ell \in [K]$, there exists $m \in [K]$ such that $B_{km} \neq B_{\ell m}$. Assume $d \sim n^{\delta}$ for some $\delta \in]0,1[$. Let R: rank of matrix B. Then with high probability:

- (i) the spectrum of A consists of R eigenvalues of order $\Theta(d)$ and n-R eigenvalues of order O(d).
- (ii) R-dimensional spectral embedding reveals underlying communities: except for vanishing fraction of nodes $i \in [n]$,

$$||z_i - z_j|| =$$

$$\begin{cases} o(1) & \text{if } \sigma(i) = \sigma(j), \\ \Omega(1) & \text{if } \sigma(i) \neq \sigma(j) \end{cases}$$

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Corollary

Under these conditions any sensible clustering scheme (eg K-means properly initialized) correctly classifies all but vanishing fraction of nodes.



Proof strategy



- \bar{A} : block matrix (useful "signal")
- Write adjacency matrix as $A = \overline{A} + W$ with $\overline{A}_{ij} = \frac{d}{n} B_{\sigma(i),\sigma(j)}$
- R leading eigen-elements of \overline{A} capture community structure
- Control perturbation of eigen-elements of a symmetric matrix \overline{A} by addition of symmetric matrix W in terms of **spectral** radius $\rho(W)$ of noise matrix
- Prove bound on $\rho(W)$ for random noise matrix W



Eigenstructure of \overline{A}

Block structure of $\overline{A} \Rightarrow \overline{A}x$ constant on each block \Rightarrow eigenvectors associated to non-zero eigenvalue are block-constant.

For
$$t \in \mathbb{R}^K$$
 define $x := \phi(t) = (t_{\sigma(i)})_{i \in [n]} \in \mathbb{R}^n$.

Then $\overline{A}\phi(t) = d\phi(Mt)$, where $M_{uv} := B_{uv}\alpha_v$.

Lemma

Spectrum of \overline{A} :

R eigen-pairs $(\lambda_u = d\mu_u, \overline{x}_u = \phi(t_u))$ where (μ_u, t_u) : eigen-pairs of M with $\mu_u \neq 0$;

0: eigenvalue with multiplicity n - R

Eigenstructure of \overline{A} (continued)

Lemma

Under distinguishability hypothesis there exists $\epsilon > 0$ function of B, α such that for any choice of normalized leading eigenvectors $\overline{x}_1, \ldots, \overline{x}_R, \overline{z}_i = \sqrt{n}(\overline{x}_1(i), \ldots, \overline{x}_R(i))^T$ verify

$$\sigma(i) \neq \sigma(j) \Rightarrow ||\overline{z}_i - \overline{z}_j|| \geq \epsilon > 0$$

Proof: Let $t_u \in \mathbb{R}^K$ be such that $\sqrt{n}\overline{x}_u = \phi(t_u)$, and $\sqrt{\alpha} = \text{Diag}(\sqrt{\alpha_u})$.

Then: $\{\sqrt{\alpha}t_u\}_{u\in[R]}$: orthonormal family by orthonormality of the \overline{x}_u .

 t_u eigenvectors of matrix $M = B\alpha$, hence $\sqrt{\alpha}t_u$: orthonormal family of eigenvectors of matrix $\sqrt{\alpha}B\sqrt{\alpha}$.

Thus
$$\sqrt{\alpha}B\sqrt{\alpha} = \sum_{u \in [R]} \mu_u(\sqrt{\alpha}t_u)(\sqrt{\alpha}t_u)^T$$
.

Equivalently: $B = \sum_{u \in [R]} \mu_u t_u t_u^T$.

Hence minimum of $||\overline{z}_i - \overline{z}_j||$ over $\sigma(i) \neq \sigma(j)$ strictly positive, for otherwise B has two identical rows, i.e. distinguishability fails.

Controlling perturbation of eigenvalues

Lemma

(Weyl's inequality) Order eigenvalues of A (resp. $\overline{A} = A + W$ for symmetric A, W as $\lambda_1 \geq \lambda_2 \geq \cdots$ (respectively, $\overline{\lambda}_1 \geq \overline{\lambda}_2 \geq \cdots$). Then for all $i \in [n]$, $|\lambda_i - \overline{\lambda}_i| \leq \rho(W)$

Proof: by Courant-Fisher theorem,

$$\lambda_i = \sup_{\dim(E)=i} \inf_{x \in E, ||x||=1} x^T A x$$

Apply to $E = \text{Vect}\{\overline{x}_1, \dots, \overline{x}_i\}$ to obtain

$$\begin{array}{ll} \lambda_{i} & \geq \inf_{x \in E, ||x|| = 1} x^{T} A x \\ & \geq \inf_{x \in E, ||x|| = 1} x^{T} \overline{A} x + \inf_{x \in E, ||x|| = 1} x^{T} W x \\ & \geq \overline{\lambda}_{i} - \rho(W). \end{array}$$

By symmetry, $\overline{\lambda}_i \geq \lambda_i - \rho(W)$ hence the result



Controlling perturbation of eigenvectors

Lemma

Let $\Delta := \inf_{\overline{\lambda}_i \neq \overline{\lambda}_j} |\overline{\lambda}_i - \overline{\lambda}_j|$. Assume $\rho(W) < \Delta/2$. Then for any normed eigenvector x_i of A associated with λ_i there exists \overline{x}_i normed eigenvector of \overline{A} associated with $\overline{\lambda}_i$ such that

$$\langle x_i, \overline{x}_i \rangle \geq \sqrt{1 - \left(\frac{\rho(W)}{\Delta - \rho(W)}\right)^2}$$

Proof: Decomposition $x_i = \sum_j \theta_j \overline{x}_j$ yields

$$Ax_i = \lambda_i x_i = \sum_j \theta_j \overline{\lambda}_j \overline{x}_j + Wx_i$$
 hence $Wx_i = \sum_j (\lambda_i - \overline{\lambda}_j) \theta_j \overline{x}_j$
By Weyl's inequality, $|\lambda_i - \overline{\lambda}_i| \ge \Delta - \rho(W)$ if $\overline{\lambda}_i \ne \overline{\lambda}_i$. Thus

$$ho(\mathcal{W}) \geq (\Delta -
ho(\mathcal{W})) \sqrt{1 - \sum_{k: \overline{\lambda}_k = \overline{\lambda}_i} | heta_k|^2}$$

$$\Rightarrow \sum_{k:\overline{\lambda}_k=\overline{\lambda}_i} |\theta_k|^2 \ge 1 - \left(\frac{\rho(W)}{\Delta - \rho(W)}\right)^2$$



Summary of argument

- Matrix \overline{A} of rank R, spectral gaps $|\overline{\lambda}_i \overline{\lambda}_j| = \Omega(d)$, R-dimensional spectral embedding with $\overline{x}_1, \dots, \overline{x}_R$ reveals clusters
- Assuming $\rho = \rho(A \overline{A}) << d$, Weyl's inequality: R eigenvalues λ_i close to $\overline{\lambda}_i = \Omega(d)$, others of order $\rho << d$
- Associated eigenvectors x_i such that $\langle x_i, \overline{x}_i \rangle = 1 O((\rho/d)^2)$

Then
$$\sum_{i \in [n]} ||z_i - \overline{z}_i||^2 = n \sum_{u \in [R]} ||x_u - \overline{x}_u||^2 = n\theta$$
 with $\theta = O((\rho/d)^2) = o(1)$

Hence (Tchebitchev inequality):

$$|\{i: ||z_i - \overline{z}_i|| \ge \theta^{1/3}\}| \le n\theta^{1/3} = o(n)$$

Yields desired conclusion: except for vanishing fraction $\theta^{1/3}$ of nodes, spectral representatives z_i $\theta^{1/3}$ -close of corresponding \overline{z}_i , themselves clustered according to community structure



Bounding $\rho = \rho(W)$

Lemma

Let $W \in \mathbb{R}^{n \times n}$: symmetric matrix with entries independent up to symmetry, bounded by 1, and such that $\mathbb{E}(W_{ij}^2) \leq O(d/n)$. Then for any fixed ϵ , with high probability, $\rho(W) \leq O(\sqrt{d} n^{\epsilon})$.

Corollary

Under assumptions of main result, $d = n^{\delta}$ for some $\delta \in]0,1[$, then with high probability $\rho(W) = o(d)$.

Proof: Take $\epsilon < \delta/2$ to obtain $\rho(W) = O(n^{\delta/2 + \epsilon}) = o(n^{\delta})$.



Bounding $\rho = \rho(W)$: proof of Lemma

For fixed
$$k \in \mathbb{N}$$
, write $\rho^{2k} \leq \sum_{i \in [n]} \lambda_i(W)^{2k} = \text{Trace}(W^{2k})$

Thus $\mathbb{P}(\rho \ge x) \le x^{-2k} \mathbb{E}(\rho^{2k}) \le x^{-2k} \mathbb{E} \operatorname{Trace}(W^{2k})$ Combinatorial view of trace:

Trace
$$(W^{2k}) = \sum_{i_0^{2k} \in [n]^{2k+1}: i_0 = i_{2k}} \prod_{j=1}^{2k} W_{i_{j-1}i_j}$$

Recall W_{ii} : centered and independent

 \rightarrow Only paths contributing non-zero expectation: traverse each edge at least twice

$$\Rightarrow \mathbb{E}\mathsf{Trace}(W^{2k}) \leq \sum_{e=1}^k \sum_{v=2}^{e+1} C(e,v) n^v O((d/n)^e)$$

Yields
$$\mathbb{E}\mathsf{Trace}(W^{2k}) = O(nd^k)$$
.

For
$$x = \sqrt{d} n^{\epsilon}$$
, yields $\mathbb{P}(\rho \ge x) \le O(n^{1-2k\epsilon})$

Result follows by taking $k > 1/(2\epsilon)$

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Stronger bounds on ho= ho(W)

Theorem (Feige and Ofek, 2005)

Let $A \in \mathbb{R}^{n \times n}$: symmetric matrix with entries independent up to symmetry, $A_{ij} \in [0,1]$, and such that $\mathbb{E}(A_{ij}) \leq d/n$, where $d \leq n^{1/5}$.

Then for some (universal) constant $\kappa > 0$, with high probability $\rho(A - \mathbb{E}(A)) \leq \kappa \sqrt{\max(d, \log(n))}$.

Hence result on spectral methods for Stochastic Block Model still valid as long as $d >> \sqrt{\max(d, \log(n))}$, i.e. $d >> \sqrt{\log(n)}$.

Takeaway messages

- Community detection a generic inference problem with many applications
- Basic spectral methods successful in scenarios well described by stochastic block model with strong enough signal and fixed number of large communities
- Variants with improved efficiency under weaker signal and in more difficult scenarios (many communities, small communities, overlapping communities,...) subject of ongoing research