Virtual address

(page #, ott set)

Physical address

(trame #, ottset)

loge table

Page #	frame #
Ô	
1	
2	
:	

Suppose we have a 16-6/t oddress space pager = 1KB = framer

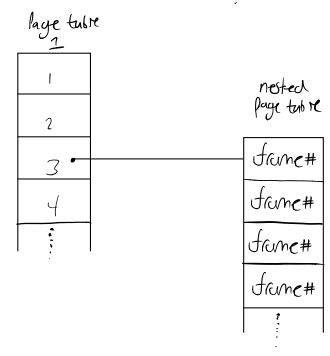
Virtual adars 4322 <u>4096</u> 226 <u>128</u> 98 0001000011100010 <u>54</u> 38 32 2 0ff set .00010000011100010 paye 4 he tulked about trues befor this.

"internal flagmentation - to much money given to poceer

32 bit address face 4 KB pager

° you need a paye take for each process running.

## 



Page tables need to be start.

· new to be in conore

· or a dedicated hardware

i.e. translation look asion buffer - TLB

Multi level page tales us noted money taker

Stame physical memory OS-test 2 wechseday  $24^{th}$ Uppare we have a 32 bit oddron space 4kB pager =  $2^{12}$  by tor  $2^{20}$  pager  $2^{20}$  pager

Contiguous space not required.

Inverted page tables we have 32 mB Ram  $2^{5} \cdot 2^{20} \text{ by ter}$   $\frac{2^{5}}{2^{12}} = 2^{13} \text{ formes}$ 

1 table for machine each entry tells ownlon process o' Vintual page

Virtual address (page #, offset)

trame#
0 ProcessID Page#

Segmentation

- Sagrents do not new to be the sure size

- Not transport to a programmer

- Compiler organizer data into segments
- global Utriables

Ubroutin-cr

- procedure call Stack

- local variable for subscortine

- growing data Structurer

- Thursd data.

- Virtual address ( Jeyment #, offset)
Jeyment table used to Goode addresses.

Jegment table

Seyment #	address	Jize Gytu
0	1000	120
1	10000	250
Z	6000	400

Virtual address (1, 46)
Physical address (10000 + 46)

butter overflow attack

Test Synchrono sa tran
Devolocil
Monary Managnent,

try and read chaper 7 + problem 7.2, 7.6, 7.12
7.13, 7.14