Implementing a library for scoped algebraic effects in Agda

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Introduction

1.1 Goals

Preliminaries

2.1 Agda

Agda is a functional language with dependent types. The current version, originally known as Agda2, was originally developed by Ulf Norell [Nor07]

2.1.1 Dependent Types

```
data Vec (A : Set) : \mathbb{N} \to Set where
\underline{\quad} :: \underline{\quad} : \{n : \mathbb{N}\} \to A \to Vec \ A \ n \to Vec \ A \ (suc \ n)
[] : Vec \ A \ 0
```

2.1.2 Propositions as Types

2.1.3 Strict Positivity

Container

2.1.4 Termination Checking

Well Founded Induction / Sized Types

```
open import Agda.Builtin.Size public
  renaming ( SizeU to SizeUniv ) -- sort SizeUniv
  using (Size
                                                 -- Size : SizeUniv
          ; Size<_
                                                 -- Size<_ : Size → SizeUniv
                                                 -- ↑_ : Size → Size
          ; ↑_
           ; <u>_</u>⊔<sup>s</sup>_
                                                 -- _□s_ : Size → Size → Size
                                                 -- œ : Size
           \infty
data List (A : Set) : Set where
   \underline{\phantom{a}}::\underline{\phantom{a}}:A\to\operatorname{\mathsf{List}} A\to\operatorname{\mathsf{List}} A
  []: List A
\mathsf{map-list}: \{A \ B : \mathsf{Set}\} \to (A \to B) \to (\mathsf{List} \ A \to \mathsf{List} \ B)
\mathsf{map-list}\ f\left(x::xs\right) = f\,x::\mathsf{map-list}\ f\,xs
map-list f
data Rose (A : \mathsf{Set}) : \mathsf{Size} \to \mathsf{Set} where
  rose : \forall \{i\} \rightarrow A \rightarrow \mathsf{List} \; (\mathsf{Rose} \; A \; i) \rightarrow \mathsf{Rose} \; A \; (\uparrow \; i)
\mathsf{map\text{-}rose} : \{A \ B : \mathsf{Set}\} \ \{i : \mathsf{Size}\} \to (A \to B) \to (\mathsf{Rose} \ A \ i \to \mathsf{Rose} \ B \ i)
map-rose f (rose x xs) = rose (f x) (map-list (map-rose f) xs)
```

2.2 Curry and Call-Time-Choice

 $let\ x = coin\ in\ x + x$

Algebraic Effects

- 3.1 Free Monads
- 3.2 Handler
- **3.2.1** Nondet
- **3.2.2** State
- 3.3 Scoped Effects
- 3.3.1 Cut
- 3.4 Call-Time-Choice as Effect

Higher Order

Conclusion

5.1 Summary

Bibliography

[Nor07] Ulf Norell. "Towards a practical programming language based on dependent type theory". PhD thesis. SE-412 96 Göteborg, Sweden: Department of Computer Science and Engineering, Chalmers University of Technology, Sept. 2007.