



Sudoku Solver with Backtracking and CSP Techniques

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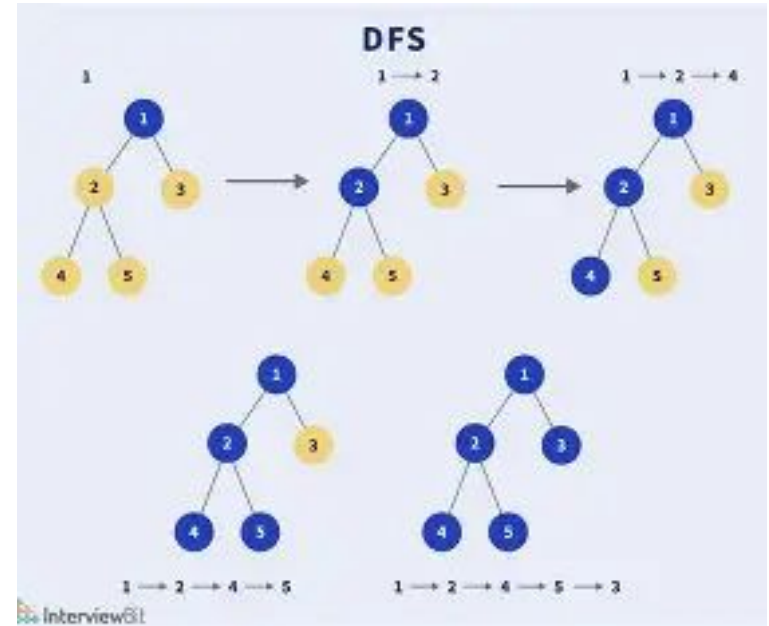
What is Sudoku

- Puzzle where you need to fill in a 9x9 board with numbers 1-9
- Each cell must have a unique number within its column, row, and in the 3x3 area it sits in
- Boards start pre-filled, usually with enough cells filled so that there is only 1 possible solution
- Solving the board is an exercise in logic- making educated guesses and ruling out possibilities until the remaining cells only have one option left

4		1	2	9			7	5
2			3			8		
	7			8				6
			1		3		6	2
1		5				4		3
7	3		6		8			
6				2			3	
		7			1			4
8	9			6	5	1		7

What is DFS

- DFS is a graph search which makes use of a stack data structure in order to process nodes for exploration (LIFO behavior)
- Once it is done with a path it will backtrack to the next path and repeat the process until it either explores all nodes or finds a goal node
- This is how we implement backtracking for Sudoku



Constraint Satisfaction Problems (CSP)

A mathematical problem where variables must be assigned a value in their domain such that it meets all given constraints

- Variables: These are what we need to find values for (each territory)
- Domains: The set of possible values for a given variable (Colors Red, Green, Blue)
- Constraints: The rules that dictate what values are allowed for a given variable (If two territories touch then they can't have the same color)

We will model Sudoku as a CSP in order to improve backtracking





Basic Backtracking vs CSP Backtracking

Hypothesis

- The basic backtracking should explore more states since it is doing no pruning on the initial domains of all the variables
- Using forward checking in the backtracking will result in fewer possible states since when exploring a branch the domains of neighboring states will be pruned down so there will be less to explore



Methods

- Solving Sudoku with DFS
 - Try every possible board using DFS
- Solving Sudoku with DFS and simple check
 - Before placing a value check if it is safe to place
- Solving Sudoku with DFS enhanced with forward checking
 - Treat the puzzle as a CSP and use forward checking in order to process less branches



Sudoku with DFS

State space: The starting board and all board states made from filling in the blank spaces in the start board

- Successor function: Assign a value 1-9 to an empty space
- Start state: The starting board, which has one or more spaces filled in
- Goal test: The current board is completely filled in, and each space is a different number than those in its column, row, and group

A sample starting board

0s are empty spaces

0 0 9	0 4 0	0 0 0
0 0 0	0 0 5	3 1 0
0 6 1	0 0 8	0 5 0
0 0 5	4 0 0	2 0 3
0 1 0	0 0 7	0 0 8
0 8 0	0 0 0	7 6 0
3 0 6	0 1 9	4 0 0
7 0 0	0 0 0	0 0 0
0 0 4	0 5 0	6 2 7

One possible solution

5 3 9	1 4 6	8 7 2
8 4 7	9 2 5	3 1 6
2 6 1	3 7 8	9 5 4
6 7 5	4 8 1	2 9 3
9 1 2	6 3 7	5 4 8
4 8 3	5 9 2	7 6 1
3 2 6	7 1 9	4 8 5
7 5 8	2 6 4	1 3 9
1 9 4	8 5 3	6 2 7

Solving Steps

0	0	9	0	4	0	0	0	0
0	0	0	0	0	5	3	1	0
0	6	1	0	0	8	0	5	0
0	0	5	4	0	0	2	0	3
0	1	0	0	0	7	0	0	8
0	8	0	0	0	0	7	6	0
3	0	6	0	1	9	4	0	0
7	0	0	0	0	0	0	0	0
0	0	4	0	5	0	6	2	7

1	0
0	

2	0
0	

3	0
0	

4	0
0	0

...

1	1
0	

1	2
0	

1	3
0	

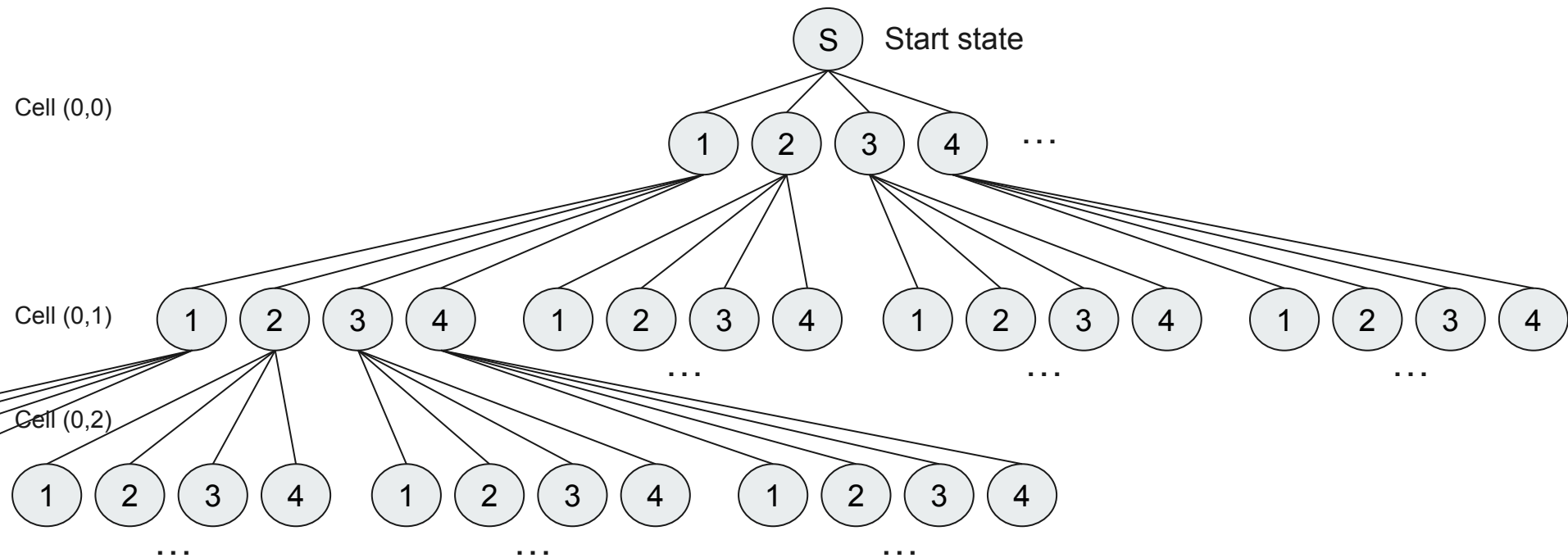
1	4
0	0

...

As a Search Tree

Given this initial row 0:

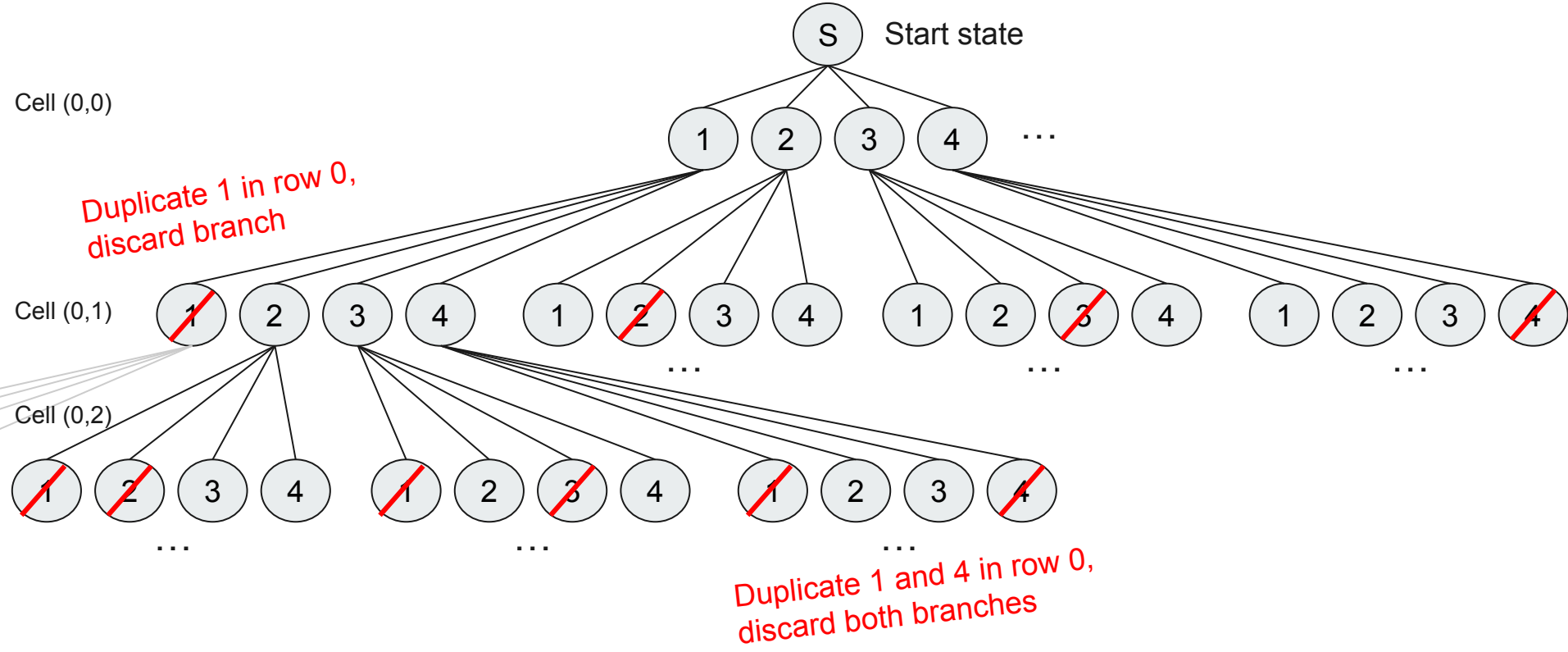
0	0	9	0	7	0	0	0	0
---	---	---	---	---	---	---	---	---



With CSP Backtracking

Given this initial row 0:

0	0	9	0	7	0	0	0	0
---	---	---	---	---	---	---	---	---





Backtracking in Sudoku

At each step, check that the assigned value will be valid in the grid

Saves time evaluating results for a grid that will never be possible (e.g. grid where we assign 1s to every single empty space)

Test:

Do any column, row, or group neighbors have the same value? If so, abandon this branch of the search tree



Backtracking Cont.

No need to deduplicate nodes- each DFS node is a unique state

No inherent advantage to using BFS- solution depth is constant

Using DFS optimizes memory usage

Pruning is performed on branches that will never produce valid board states-
backtracking solution is the same as the naive solution

Sudoku as a CSP

We framed Sudoku as a Constraint Satisfaction Problem

- Identification problem
- High-order constraints (81 variables, 1-9 is initial domain)
- Start with a partial assignment
- Build complete assignments
- Eventually settle on a solution, a complete assignment that meets constraints

8	3		1			6		5	8	3	7	1	9	4	6	2	5
							8		5	4	9	6	2	3	7	8	1
			7			9			6	2	1	7	8	5	9	3	4
	5			1	7				2	5	6	8	1	7	4	9	3
		3				2			4	1	3	5	6	9	2	7	8
			3	4			1		9	7	8	3	4	2	5	1	6
		4			8				1	6	4	2	7	8	3	5	9
	9								7	9	5	4	3	1	8	6	2
3		2			6		4	7	3	8	2	9	5	6	1	4	7



Forward checking in Sudoku

For every empty cell:

- Assign each empty cell a domain list, with the numbers 1-9

- Assign the pre-placed value to the domain of cells that are pre-filled

Every recursion:

- Update domain list by removing numbers used by neighboring cells from the domain list

- Recurse on every possible cell value for the domain

Results

- Plain DFS
 - Never finished (57 0's = $6.3 \cdot 10^{15}$ branches)
- DFS with backtracking
 - Solve time: 32.5s on average
 - Total branches explored: 535051
- DFS with backtracking and forward checking
 - Solve time: 17s on average
 - Total branches explored: 436329
- Optimization effects with forward checking:
 - 15.5s reduction in search time
 - 18.5% reduction in branch exploration

0 0 0	0 0 0	0 0 0
0 0 0	7 0 9	0 0 0
0 0 6	0 3 0	9 0 0
0 6 0	2 7 5	0 1 0
0 0 5	9 0 6	3 0 0
0 1 0	3 4 8	0 9 0
0 0 7	0 8 0	5 0 0
0 0 0	4 0 7	0 0 0
0 0 0	0 0 0	0 0 0

Stress test with
mostly-blank board on
a laptop

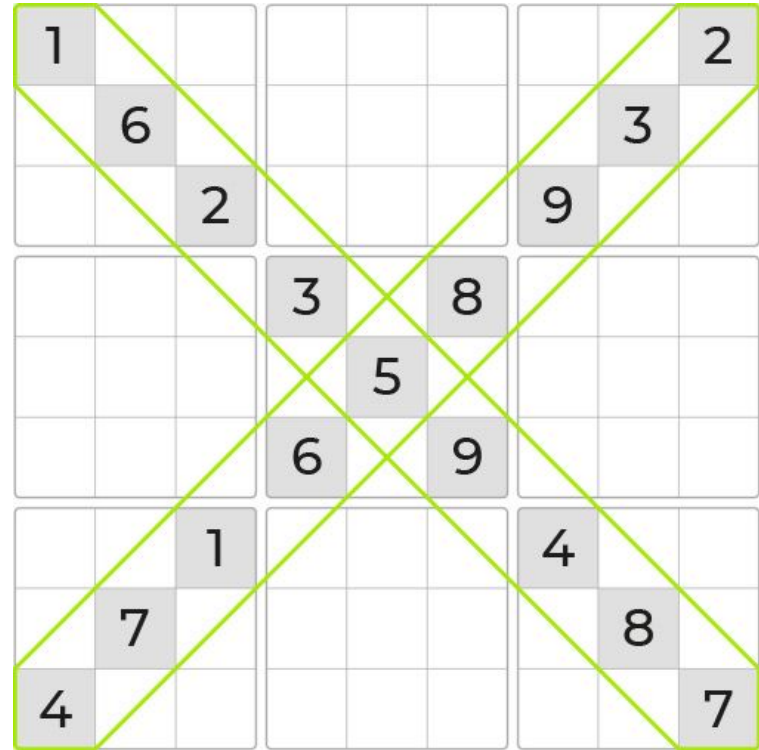


Conclusions

- Sudoku can be feasibly treated as a tree search problem
- The solver can take advantage of Constraint Satisfaction Problem solving techniques for further optimization
- More optimizations to consider:
 - Could cache neighbor list to reduce list access (memory vs cpu)
 - Would be orders of magnitude faster in a compiled language
 - Domain updates can be cached and made deterministic based on only earlier cell updates

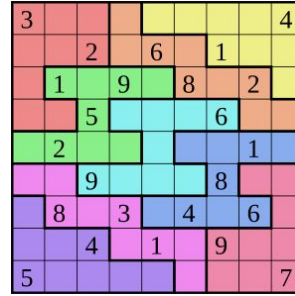
Extension: Diagonals

- Along with the base rules for Sudoku we implemented another custom rule for the diagonals
- This is where along both main diagonals you can only have exactly one of each number 1 to N
- Usually referred to as X Sudoku



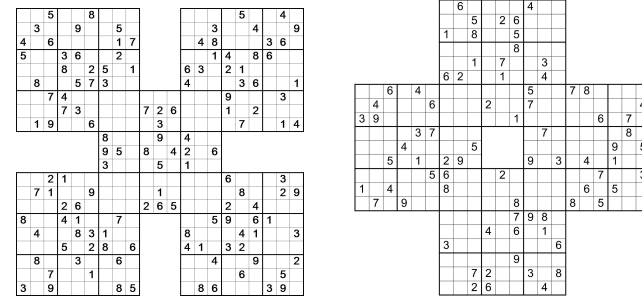
Further Research

- Solving for more advanced board shapes
 - Would require a more complex structure for storing the board state and figuring out which areas are connected
 - CSP solution methodologies would remain largely unchanged
- More complicated solution criteria
 - Requires more complicated solution tests
 - Backtracking and forward checking needs reimagining



Jigsaw Sudoku

Giant Sudoku boards



Killer Sudoku, before and after solving



References

- Course materials on CSP
- <https://en.wikipedia.org/wiki/Sudoku>
- <https://prntbl.concejomunicipaldechinu.gov.co/sudoku-printables-6-per-page/>
- <https://www.geeksforgeeks.org/artificial-intelligence/constraint-satisfaction-problems-csp-in-artificial-intelligence/>