

TSPLIB 95

Gerhard Reinelt
Universität Heidelberg
Institut für Angewandte Mathematik
Im Neuenheimer Feld 294
D-69120 Heidelberg
Gerhard.Reinelt@IWR.Uni-Heidelberg.DE

TSPLIB is a library of sample instances for the TSP (and related problems) from various sources and of various types. Instances of the following problem classes are available.

Symmetric traveling salesman problem (TSP)

Given a set of n nodes and distances for each pair of nodes, find a roundtrip of minimal total length visiting each node exactly once. The distance from node i to node j is the same as from node j to node i .

Hamiltonian cycle problem (HCP)

Given a graph, test if the graph contains a Hamiltonian cycle or not.

Asymmetric traveling salesman problem (ATSP)

Given a set of n nodes and distances for each pair of nodes, find a roundtrip of minimal total length visiting each node exactly once. In this case, the distance from node i to node j and the distance from node j to node i may be different.

Sequential ordering problem (SOP)

This problem is an asymmetric traveling salesman problem with additional constraints. Given a set of n nodes and distances for each pair of nodes, find a Hamiltonian path from node 1 to node n of minimal length which takes given precedence constraints into account. Each precedence constraint requires that some node i has to be visited before some other node j .

Capacitated vehicle routing problem (CVRP)

We are given $n - 1$ nodes, one depot and distances from the nodes to the depot, as well as between nodes. All nodes have demands which can be satisfied by the depot. For delivery to the nodes, trucks with identical capacities are available. The problem is to find tours for the trucks of minimal total length that satisfy the node demands without violating truck capacity constraint. The number of trucks is not specified. Each tour visits a subset of the nodes and starts and terminates at the depot. (Remark: In some data files a collection of alternate depots is given. A CVRP is then given by selecting one of these depots.)

Except, for the Hamiltonian cycle problems, all problems are defined on a complete graph and, at present, all distances are integer numbers. There is a possibility to require that certain edges appear in the solution of a problem.

1. The file format

Each file consists of a **specification part** and of a **data part**. The specification part contains information on the file format and on its contents. The data part contains explicit data.

1.1 The specification part

All entries in this section are of the form $\langle keyword \rangle : \langle value \rangle$, where $\langle keyword \rangle$ denotes an alphanumerical keyword and $\langle value \rangle$ denotes alphanumerical or numerical data. The terms $\langle string \rangle$, $\langle integer \rangle$ and $\langle real \rangle$ denote character string, integer or real data, respectively. The order of specification of the keywords in the data file is arbitrary (in principle), but must be consistent, i.e., whenever a keyword is specified, all necessary information for the correct interpretation of the keyword has to be known. Below we give a list of all available keywords.

1.1.1 NAME : $\langle string \rangle$

Identifies the data file.

1.1.2 TYPE : $\langle string \rangle$

Specifies the type of the data. Possible types are

TSP	Data for a symmetric traveling salesman problem
ATSP	Data for an asymmetric traveling salesman problem
SOP	Data for a sequential ordering problem
HCP	Hamiltonian cycle problem data
CVRP	Capacitated vehicle routing problem data
TOUR	A collection of tours

1.1.3 COMMENT : $\langle string \rangle$

Additional comments (usually the name of the contributor or creator of the problem instance is given here).

1.1.4 DIMENSION : $\langle integer \rangle$

For a TSP or ATSP, the dimension is the number of its nodes. For a CVRP, it is the total number of nodes and depots. For a TOUR file it is the dimension of the corresponding problem.

1.1.5 CAPACITY : $\langle integer \rangle$

Specifies the truck capacity in a CVRP.

1.1.6 EDGE_WEIGHT_TYPE : $\langle string \rangle$

Specifies how the edge weights (or distances) are given. The values are

EXPLICIT	Weights are listed explicitly in the corresponding section
EUC_2D	Weights are Euclidean distances in 2-D
EUC_3D	Weights are Euclidean distances in 3-D

MAX_2D	Weights are maximum distances in 2-D
MAX_3D	Weights are maximum distances in 3-D
MAN_2D	Weights are Manhattan distances in 2-D
MAN_3D	Weights are Manhattan distances in 3-D
CEIL_2D	Weights are Euclidean distances in 2-D rounded up
GEO	Weights are geographical distances
ATT	Special distance function for problems att48 and att532
XRAY1	Special distance function for crystallography problems (Version 1)
XRAY2	Special distance function for crystallography problems (Version 2)
SPECIAL	There is a special distance function documented elsewhere

1.1.7 EDGE_WEIGHT_FORMAT : <string>

Describes the format of the edge weights if they are given explicitly. The values are

FUNCTION	Weights are given by a function (see above)
FULL_MATRIX	Weights are given by a full matrix
UPPER_ROW	Upper triangular matrix (row-wise without diagonal entries)
LOWER_ROW	Lower triangular matrix (row-wise without diagonal entries)
UPPER_DIAG_ROW	Upper triangular matrix (row-wise including diagonal entries)
LOWER_DIAG_ROW	Lower triangular matrix (row-wise including diagonal entries)
UPPER_COL	Upper triangular matrix (column-wise without diagonal entries)
LOWER_COL	Lower triangular matrix (column-wise without diagonal entries)
UPPER_DIAG_COL	Upper triangular matrix (column-wise including diagonal entries)
LOWER_DIAG_COL	Lower triangular matrix (column-wise including diagonal entries)

1.1.7 EDGE_DATA_FORMAT : <string>

Describes the format in which the edges of a graph are given, if the graph is not complete. The values are

EDGE_LIST	The graph is given by an edge list
ADJ_LIST	The graph is given as an adjacency list

1.1.9 NODE_COORD_TYPE : <string>

Specifies whether coordinates are associated with each node (which, for example may be used for either graphical display or distance computations). The values are

TWOD_COORDS	Nodes are specified by coordinates in 2-D
THREED_COORDS	Nodes are specified by coordinates in 3-D
NO_COORDS	The nodes do not have associated coordinates

The default value is NO_COORDS.

1.1.10 DISPLAY_DATA_TYPE : <string>

Specifies how a graphical display of the nodes can be obtained. The values are

COORD_DISPLAY	Display is generated from the node coordinates
TWOD_DISPLAY	Explicit coordinates in 2-D are given
NO_DISPLAY	No graphical display is possible

The default value is COORD_DISPLAY if node coordinates are specified and NO_DISPLAY otherwise.

1.1.11 EOF :

Terminates the input data. This entry is optional.

1.2 The data part

Depending on the choice of specifications some additional data may be required. These data are given in corresponding data sections following the specification part. Each data section begins with the corresponding keyword. The length of the section is either implicitly known from the format specification, or the section is terminated by an appropriate end-of-section identifier.

1.2.1 NODE_COORD_SECTION :

Node coordinates are given in this section. Each line is of the form

<integer> <real> <real>

if NODE_COORD_TYPE is TWOD_COORDS, or

<integer> <real> <real> <real>

if NODE_COORD_TYPE is THREED_COORDS. The integers give the number of the respective nodes. The real numbers give the associated coordinates.

1.2.2 DEPOT_SECTION :

Contains a list of possible alternate depot nodes. This list is terminated by a **-1**.

1.2.3 DEMAND_SECTION :

The demands of all nodes of a CVRP are given in the form (per line)

<integer> <integer>

The first integer specifies a node number, the second its demand. The depot nodes must also occur in this section. Their demands are 0.

1.2.4 EDGE_DATA_SECTION :

Edges of a graph are specified in either of the two formats allowed in the EDGE_DATA_FORMAT entry. If the type is EDGE_LIST, then the edges are given as a sequence of lines of the form

<integer> <integer>

each entry giving the terminal nodes of some edge. The list is terminated by a **-1**.

If the type is ADJ_LIST, the section consists of a list of adjacency lists for nodes. The adjacency list of a node x is specified as

<integer> <integer> ... <integer> -1

where the first integer gives the number of node x and the following integers (terminated by **-1**) the numbers of nodes adjacent to x . The list of adjacency lists is terminated by an additional **-1**.

1.2.5 FIXED_EDGES_SECTION :

In this section, edges are listed that are required to appear in each solution to the problem. The edges to be fixed are given in the form (per line)

<integer> <integer>

meaning that the edge (arc) from the first node to the second node has to be contained in a solution. This section is terminated by a **-1**.

1.2.6 DISPLAY_DATA_SECTION :

If **DISPLAY_TYPE** is **TWOD_DISPLAY**, the 2-dimensional coordinates from which a display can be generated are given in the form (per line)

<integer> <real> <real>

The integers specify the respective nodes and the real numbers give the associated coordinates.

1.2.7 TOUR_SECTION :

A collection of tours is specified in this section. Each tour is given by a list of integers giving the sequence in which the nodes are visited in this tour. Every such tour is terminated by a **-1**. An additional **-1** terminates this section.

1.2.8 EDGE_WEIGHT_SECTION :

The edge weights are given in the format specified by the **EDGE_WEIGHT_FORMAT** entry. At present, all explicit data is integral and is given in one of the (self-explanatory) matrix formats. with implicitly known lengths.

2. The distance functions

For the various choices of EGDE_WEIGHT_TYPE, we now describe the computations of the repective distances. In each case we give a (simplified) C-implementation for computing the distances from the input coordinates. All computations involving floating-point numbers are carried out in double precision arithmetic. The integers are assumed to be represented in 32-bit words. Since distances are required to be integral, we round to the nearest integer (in most cases). Below we have used the rounding function “nint”.

2.1 Euclidean distance (L_2 -metric)

For edge weight type EUC_2D and EUC_3D, floating point coordinates must be specified for each node. Let $x[i]$, $y[i]$, and $z[i]$ be the coordinates of node i .

In the 2-dimensional case the distance between two points i and j is computed as follows:

```
xd = x[i] - x[j];
yd = y[i] - y[j];
dij = nint( sqrt( xd*xd + yd*yd ) );
```

In the 3-dimensional case we have:

```
xd = x[i] - x[j];
yd = y[i] - y[j];
zd = z[i] - z[j];
dij = nint( sqrt( xd*xd + yd*yd + zd*zd ) );
```

where `sqrt` is the C square root function.

2.2 Manhattan distance (L_1 -metric)

Distances are given as Manhattan distances if the edge weight type is MAN_2D or MAN_3D. They are computed as follows.

2-dimensional case:

```
xd = abs( x[i] - x[j] );
yd = abs( y[i] - y[j] );
dij = nint( xd + yd );
```

3-dimensional case:

```
xd = abs( x[i] - x[j] );
yd = abs( y[i] - y[j] );
zd = abs( z[i] - z[j] );
dij = nint( xd + yd + zd );
```

2.3 Maximum distance (L_∞ -metric)

Maximum distances are computed if the edge weight type is MAX_2D or MAX_3D.

2-dimensional case:

```
xd = abs( x[i] - x[j] );
yd = abs( y[i] - y[j] );
dij = max( nint( xd ), nint( yd ) );
```

3-dimensional case:

```
xd = abs( x[i] - x[j] );
yd = abs( y[i] - y[j] );
zd = abs( z[i] - z[j] );
dij = max( nint( xd ), nint( yd ), nint( zd ) );
```

2.4 Geographical distance

If the traveling salesman problem is a geographical problem, then the nodes correspond to points on the earth and the distance between two points is their distance on the idealized sphere with radius 6378.388 kilometers. The node coordinates give the geographical latitude and longitude of the corresponding point on the earth. Latitude and longitude are given in the form DDD.MM where DDD are the degrees and MM the minutes. A positive latitude is assumed to be “North”, negative latitude means “South”. Positive longitude means “East”, negative longitude is assumed to be “West”. For example, the input coordinates for Augsburg are 48.23 and 10.53, meaning $48^{\circ}23'$ North and $10^{\circ}53'$ East.

Let $x[i]$ and $y[i]$ be coordinates for city i in the above format. First the input is converted to geographical latitude and longitude given in radians.

```
PI = 3.141592;
deg = nint( x[i] );
min = x[i] - deg;
latitude[i] = PI * (deg + 5.0 * min / 3.0) / 180.0;
deg = nint( y[i] );
min = y[i] - deg;
longitude[i] = PI * (deg + 5.0 * min / 3.0) / 180.0;
```

The distance between two different nodes i and j in kilometers is then computed as follows:

```
RRR = 6378.388;
q1 = cos( longitude[i] - longitude[j] );
q2 = cos( latitude[i] - latitude[j] );
q3 = cos( latitude[i] + latitude[j] );
dij = (int) ( RRR * acos( 0.5*((1.0+q1)*q2 - (1.0-q1)*q3) ) + 1.0 );
```

The function “acos” is the inverse of the cosine function.

2.5 Pseudo-Euclidean distance

The edge weight type ATT corresponds to a special “pseudo-Euclidean” distance function. Let $x[i]$ and $y[i]$ be the coordinates of node i . The distance between two points i and j is computed as follows:

```
xd = x[i] - x[j];
yd = y[i] - y[j];
rij = sqrt( (xd*xd + yd*yd) / 10.0 );
tij = nint( rij );
if (tij<rij) dij = tij + 1;
else dij = tij;
```

2.6 Ceiling of the Euclidean distance

The edge weight type `CEIL_2D` requires that the 2-dimensional Euclidean distances is rounded up to the next integer.

2.7 Distance for crystallography problems

We have included into TSPLIB the crystallography problems as described in [1]. These problems are not explicitly given but subroutines are provided to generate the 12 problems mentioned in this reference and subproblems thereof (see section 3.2).

To compute distances for these problems the movement of three motors has to be taken into consideration. There are two types of distance functions: one that assumes equal speed of the motors (`XRAY1`) and one that uses different speeds (`XRAY2`). The corresponding distance functions are given as FORTRAN implementations (files `deq.f`, resp. `duneq.f`) in the distribution file.

For obtaining integer distances, we propose to multiply the distances computed by the original subroutines by 100.0 and round to the nearest integer.

We list our modified distance function for the case of equal motor speeds in the FORTRAN version below.

```
INTEGER FUNCTION ICOST(V,W)
INTEGER V,W
DOUBLE PRECISION DMIN1,DMAX1,DABS
DOUBLE PRECISION DISTP,DISTC,DISTT,COST
DISTP=DMIN1(DABS(PHI(V)-PHI(W)),DABS(DABS(PHI(V)-PHI(W))-360.0E+0))
DISTC=DABS(CHI(V)-CHI(W))
DISTT=DABS(TWOTH(V)-TWOTH(W))
C COST=DMAX1(DISTP/1.00E+0,DISTC/1.0E+0,DISTT/1.00E+0)
C *** Make integral distances ***
ICOST=AINT(100.0E+0*COST+0.5E+0)
RETURN
END
```

The numbers `PHI()`, `CHI()`, and `TWOTH()` are the respective x -, y -, and z -coordinates of the points in the generated traveling salesman problems. Note, that TSPLIB95 contains only the original distance computation without the above modification.

2.7 Verification

To verify correctness of the distance function implementations we give the length of some “canonical” tours $1, 2, 3, \dots, n$.

The canonical tours for `pcb442`, `gr666`, and `att532` have lengths 221 440, 423 710, and 309 636, respectively.

The canonical tour for the problem `xray14012` (the 8th problem considered in [21]) with distance `XRAY1` has length 15 429 219. With distance `XRAY2` it has the length 12 943 294.

3. Description of the library files

In this section we give a list of all problem instances that are currently available together with information on the length of optimal tours or lower and upper bounds for this length (if available).

3.1 Symmetric traveling salesman problems

The TSP instances are contained in directory `tsp`. Table 1 gives the problem names along with number of cities, problem type, and known lower and upper bounds for the optimal tour length (a single number indicating that the optimal length is known). The entry `MATRIX` indicates that the data is given in one of the matrix formats of 1.1.7. The names of the corresponding data files are obtained by appending the suffix “`.tsp`” to the problem name. Some optimal tours are also provided. The corresponding files have names with suffix “`.opt.tour`”.

Name	#cities	Type	Bounds
a280	280	EUC_2D	2579
ali535	535	GEO	202310
att48	48	ATT	10628
att532	532	ATT	27686
bayg29	29	GEO	1610
bays29	29	GEO	2020
berlin52	52	EUC_2D	7542
bier127	127	EUC_2D	118282
brazil58	58	MATRIX	25395
brd14051	14051	EUC_2D	[468942, 469445]
brg180	180	MATRIX	1950
burma14	14	GEO	3323
ch130	130	EUC_2D	6110
ch150	150	EUC_2D	6528
d198	198	EUC_2D	15780
d493	493	EUC_2D	35002
d657	657	EUC_2D	48912
d1291	1291	EUC_2D	50801
d1655	1655	EUC_2D	62128
d2103	2103	EUC_2D	[79952, 80450]
d15112	15112	EUC_2D	[1564590, 1573152]
d18512	18512	EUC_2D	[644650, 645488]
dantzig42	42	MATRIX	699
dsj1000	1000	CEIL_2D	18659688
eil51	51	EUC_2D	426
eil76	76	EUC_2D	538
eil101	101	EUC_2D	629

Table 1 Symmetric traveling salesman problems (Part I)

Name	#cities	Type	Bounds
f1417	417	EUC_2D	11861
f11400	1400	EUC_2D	20127
f11577	1577	EUC_2D	[22204,22249]
f13795	3795	EUC_2D	[28723,28772]
fnl4461	4461	EUC_2D	182566
fri26	26	MATRIX	937
gil262	262	EUC_2D	2378
gr17	17	MATRIX	2085
gr21	21	MATRIX	2707
gr24	24	MATRIX	1272
gr48	48	MATRIX	5046
gr96	96	GEO	55209
gr120	120	MATRIX	6942
gr137	137	GEO	69853
gr202	202	GEO	40160
gr229	229	GEO	134602
gr431	431	GEO	171414
gr666	666	GEO	294358
hk48	48	MATRIX	11461
kroA100	100	EUC_2D	21282
kroB100	100	EUC_2D	22141
kroC100	100	EUC_2D	20749
kroD100	100	EUC_2D	21294
kroE100	100	EUC_2D	22068
kroA150	150	EUC_2D	26524
kroB150	150	EUC_2D	26130
kroA200	200	EUC_2D	29368
kroB200	200	EUC_2D	29437
lin105	105	EUC_2D	14379
lin318	318	EUC_2D	42029
linhp318	318	EUC_2D	41345
nrw1379	1379	EUC_2D	56638
p654	654	EUC_2D	34643
pa561	561	MATRIX	2763
pcb442	442	EUC_2D	50778
pcb1173	1173	EUC_2D	56892
pcb3038	3038	EUC_2D	137694
pla7397	7397	CEIL_2D	23260728
pla33810	33810	CEIL_2D	[65913275,66116530]
pla85900	85900	CEIL_2D	[141904862,142487006]

Table 1 Symmetric traveling salesman problems (Part II)

Name	#cities	Type	Bounds
pr76	76	EUC_2D	108159
pr107	107	EUC_2D	44303
pr124	124	EUC_2D	59030
pr136	136	EUC_2D	96772
pr144	144	EUC_2D	58537
pr152	152	EUC_2D	73682
pr226	226	EUC_2D	80369
pr264	264	EUC_2D	49135
pr299	299	EUC_2D	48191
pr439	439	EUC_2D	107217
pr1002	1002	EUC_2D	259045
pr2392	2392	EUC_2D	378032
rat99	99	EUC_2D	1211
rat195	195	EUC_2D	2323
rat575	575	EUC_2D	6773
rat783	783	EUC_2D	8806
rd100	100	EUC_2D	7910
rd400	400	EUC_2D	15281
rl1304	1304	EUC_2D	252948
rl1323	1323	EUC_2D	270199
rl1889	1889	EUC_2D	316536
rl5915	5915	EUC_2D	[565040,565530]
rl5934	5934	EUC_2D	[554070,556045]
rl11849	11849	EUC_2D	[920847,923368]
si175	175	MATRIX	21407
si535	535	MATRIX	48450
si1032	1032	MATRIX	92650
st70	70	EUC_2D	675
swiss42	42	MATRIX	1273
ts225	225	EUC_2D	126643
tsp225	225	EUC_2D	3919
u159	159	EUC_2D	42080
u574	574	EUC_2D	36905
u724	724	EUC_2D	41910
u1060	1060	EUC_2D	224094
u1432	1432	EUC_2D	152970
u1817	1817	EUC_2D	57201
u2152	2152	EUC_2D	64253
u2319	2319	EUC_2D	234256
ulysses16	16	GEO	6859
ulysses22	22	GEO	7013
usa13509	13509	EUC_2D	[19947008,19982889]
vm1084	1084	EUC_2D	239297
vm1748	1748	EUC_2D	336556

Table 1 Symmetric traveling salesman problems (Part III)

Crystallography problems

In the file `xray.problems` in directory `tsp` we distribute the routines written by Bland and Shallcross and the necessary data to generate the crystallography problems discussed in [1]. The file `xray.problems` is one file into which the single files mentioned in the sequel have been merged. These single files have to be extracted from `xray.problems` using an editor. The following original files are provided

read.me	deq.f	duneq.f	daux.f	gentsp.f
a.data	b.data	d.data	e.data	f.data

In addition we have included specially prepared data files to generate the 12 problems mentioned in [1]. The files have the names `xray1.data` through `xray12.data`.

Using these data files 12 symmetric TSPs can be generated using the program `gentsp.f`. We propose to name the respective problem instances `xray4472`, `xray2950`, `xray7008`, `xray2762`, `xray6922`, `xray9070`, `xray5888`, `xray14012`, `xray5520`, `xray13804`, `xray14464`, and `xray13590`.

To verify the correct use of the generating routines we list part of the file `xray14012.tsp`.

```
NAME : xray14012
COMMENT : Crystallography problem 8 (Bland/Shallcross)
TYPE : TSP
DIMENSION : 14012
EDGE_WEIGHT_TYPE : XRAY2
NODE_COORD_SECTION

    1    -91.802854544029    -6.4097888697337    176.39830490027
    2    -87.715643397938    -6.4659384343446    165.56800324542
    3    -83.587211962870    -6.4895404648110    163.53828545043
    4    -79.460007412434    -6.4797580053949    165.86438271158
    :
    :
14009    100.539992581837    6.4797580053949    165.86438271158
14010    96.412788031401    6.4895404648110    163.53828545043
14011    92.284356596333    6.4659384343446    165.56800324542
14012    88.197145450242    6.4097888697337    176.39830490027
```

3.2 Hamiltonian cycle problems

Instances of the Hamiltonian cycle problem are contained in the directory `hcp`. At present, we have the data files

alb1000.hcp	alb3000b.hcp	alb3000e.hcp	alb2000.hcp	alb3000c.hcp
alb4000.hcp	alb3000a.hcp	alb3000d.hcp	alb5000.hcp	

Every instance contains a Hamiltonian cycle which is given in the corresponding `.opt.tour` file. In problem instance `alb4000` two edges are fixed.

In addition to these files, the directory contains the C-program `tspleap.c` by M. Jünger and G. Rinaldi. This program can be used to generate TSP instances (in TSPLIB format)

originating from the problem of deciding whether an (r,s) -leaper on a mxn chess board can start at some square of the board, visit each square exactly once, and return to its starting square. A detailed documentation is given in the file `tspleap.c`.

3.3 Asymmetric traveling salesman problems

Table 2 lists the ATSP instances (in directory `atsp`) together with their optimal solution values. The names of the corresponding data files are obtained by appending the suffix “`.atsp`” to the problem name. The data files for problems `ftv90`, `ftv100`, `ftv110`, `ftv120`, `ftv130`, `ftv140`, `ftv150`, and `ftv160` are not present. These instances are obtained from `ftv170`. E.g., `ftv120` is the subproblem of `ftv170` defined by the first 121 nodes, `ftv130` is defined by the first 131 nodes, etc.

Name	#cities	Type	Optimum
br17	17	MATRIX	39
ft53	53	MATRIX	6905
ft70	70	MATRIX	38673
ftv33	34	MATRIX	1286
ftv35	36	MATRIX	1473
ftv38	39	MATRIX	1530
ftv44	45	MATRIX	1613
ftv47	48	MATRIX	1776
ftv55	56	MATRIX	1608
ftv64	65	MATRIX	1839
ftv70	71	MATRIX	1950
ftv90	91	MATRIX	1579
ftv100	101	MATRIX	1788
ftv110	111	MATRIX	1958
ftv120	121	MATRIX	2166
ftv130	131	MATRIX	2307
ftv140	141	MATRIX	2420
ftv150	151	MATRIX	2611
ftv160	161	MATRIX	2683
ftv170	171	MATRIX	2755
kro124	100	MATRIX	36230
p43	43	MATRIX	5620
rbg323	323	MATRIX	1326
rbg358	358	MATRIX	1163
rbg403	403	MATRIX	2465
rbg443	443	MATRIX	2720
ry48p	48	MATRIX	14422

Table 2 Asymmetric traveling salesman problems

3.4 Sequential ordering problems

Every instance of a sequential ordering problem is given by a full matrix C of the following kind. If node i has to precede node j , then C_{ji} is set to -1 . C is assumed to be transitively closed with respect to precedences, i.e., if i has to precede j and j has to precede k , then

it is implied that i has to precede k and, therefore, also C_{ki} has to be set to -1 . Because we require, that node 1 is the first node and node n is the last node in each feasible path, a SOP problem instance always has $C_{i1} = -1$, for all $i = 2, \dots, n$, and $C_{nj} = -1$, for all $j = 1, \dots, n - 1$. The entry C_{1n} is set to infinity. All other entries of C are nonnegative integer values.

Name	#nodes	#prec.	Type	Bounds
ESC07	9	6	MATRIX	2125
ESC11	13	3	MATRIX	2075
ESC12	14	7	MATRIX	1675
ESC25	27	9	MATRIX	1681
ESC47	49	10	MATRIX	1288
ESC63	65	95	MATRIX	62
ESC78	80	77	MATRIX	18230
br17.10	17	10	MATRIX	55
br17.12	17	12	MATRIX	55
ft53.1	54	12	MATRIX	[7438, 7570]
ft53.2	54	25	MATRIX	[7630, 8335]
ft53.3	54	48	MATRIX	[9473, 10935]
ft53.4	54	63	MATRIX	14425
ft70.1	71	17	MATRIX	39313
ft70.2	71	35	MATRIX	[39739, 41778]
ft70.3	71	68	MATRIX	[41305, 44732]
ft70.4	71	86	MATRIX	[52269, 53882]
kro124p.1	101	25	MATRIX	[37722, 42845]
kro124p.2	101	49	MATRIX	[38534, 45848]
kro124p.3	101	97	MATRIX	[40967, 55649]
kro124p.4	101	131	MATRIX	[64858, 80753]
p43.1	44	9	MATRIX	27990
p43.2	44	20	MATRIX	[28175, 28330]
p43.3	44	37	MATRIX	[28366, 28680]
p43.4	44	50	MATRIX	[69569, 82960]
prob42	42	10	MATRIX	243
prob100	100	41	MATRIX	[1024, 1385]
rbg048a	50	192	MATRIX	351
rbg050c	52	256	MATRIX	467
rbg109a	111	622	MATRIX	1038
rbg150a	152	952	MATRIX	[1748, 1750]
rbg174a	176	1113	MATRIX	2053
rbg253a	255	1721	MATRIX	[2928, 2987]
rbg323a	325	2412	MATRIX	[3136, 3221]
rbg341a	343	2542	MATRIX	[2543, 2854]
rbg358a	360	3239	MATRIX	[2518, 2758]
rbg378a	380	3069	MATRIX	[2761, 3142]
ry48p.1	49	11	MATRIX	[15220, 15935]
ry48p.2	49	23	MATRIX	[15524, 17071]
ry48p.3	49	42	MATRIX	[18156, 20051]
ry48p.4	49	58	MATRIX	[29967, 31446]

Table 3 Sequential ordering problems

Table 3 lists the SOP instances (in directory `sop`) together with their known lower and upper bounds for the optimal path length. The names of the corresponding data files are obtained by appending the suffix “`.sop`” to the problem name.

3.5 Capacitated vehicle routing problems

Data for capacitated vehicle routing problems is contained in the directory `vrp`. Data files have suffix “`.vrp`”. At present, we have the data files

att48.vrp	eil30.vrp	eil7.vrp	eilB76.vrp	eil13.vrp
eil31.vrp	eilA101.vrp	eilC76.vrp	eil22.vrp	eil33.vrp
eilA76.vrp	eilD76.vrp	eil23.vrp	eil51.vrp	eilB101.vrp
gil262.vrp				

Various problems can be defined on these data sets, e.g., depending on whether the number of vehicles is fixed, so we do not list optimal solutions here. Some values are given in the data files themselves.

3.6 Further special files

In addition to the data and solution files, the following special files are contained in the library.

`TSPLIB_VERSION`: Gives the current version of the library

`README`: A short information on TSPLIB

`DOC.PS`: Description of TSPLIB (PostScript)

4. Remarks

1. The problem `lin318` is originally a Hamiltonian path problem. One obtains this problem by adding the additional requirement that the edge from 1 to 214 is contained in the tour. The data is given in `linhp318.tsp`.
2. Some data sets are referred to by different names in the literature. Below we give the corresponding names used in [3] and [2].

TSPLIB	[3]	[2]	TSPLIB	[3]	[2]
att48	ATT048	-	att532	ATT532	-
dantzig42	-	42	eil101	EIL10	-
eil51	EIL08	-	eil76	EIL09	-
gil262	GIL249	-	gr120	-	120
gr137	GH137	137	gr202	GH202	202
gr229	GH229	229	gr431	GH431	431
gr666	GH666	666	gr96	GH096	96
hk48	-	48H	kroA100	KR0124	100A
kroB100	KR0125	100B	kroC100	KR0126	100C
kroD100	KR0127	100D	kroE100	KR0128	100E

kroA150	KR030	-	kroB150	KR031	-
kroA200	KR032	-	kroB200	KR033	-
lin105	LK105	-	lin318	LK318	-
linhp318	LK318P	-	pr1002	TK1002	-
pr107	TK107	-	pr124	TK124	-
pr136	TK136	-	pr144	TK144	-
pr152	TK152	-	pr226	TK226	-
pr2392	TK2392	-	pr264	TK264	-
pr299	TK299	-	pr439	TK439	-
pr76	TK076	-	st70	KR0070	70

3. Some vehicle routing problems are also available in a TSP version. Here the depots are just treated as normal nodes. The problem *gi1262* originally contained two identical nodes, of which one was eliminated.

4. Potential contributors to this library should provide their data files in appropriate format and contact

Gerhard Reinelt

Institut für Angewandte Mathematik, Universität Heidelberg

Im Neuenheimer Feld 294, D-69120 Heidelberg

Germany

Tel (6221) 56 3171

Fax (6221) 56 5634

E-Mail Gerhard.Reinelt@IWR.Uni-Heidelberg.DE

5. Informations on new bounds or optimal solutions for library problems as well as references to computational studies (to be included in the list of references) are also appreciated.

5. Access

TSPLIB is available at various locations.

Interdisziplinäres Zentrum für Wissenschaftliches Rechnen (IWR), Heidelberg

World Wide Web:

<http://www.iwr.uni-heidelberg.de/iwr/comopt/soft/TSPLIB95/TSPLIB.html>

Konrad-Zuse-Zentrum für Informationstechnik (ZIB), Berlin

World Wide Web:

<ftp://elib.zib-berlin.de/pub/mp-testdata/tsp/index.html>

For further information on how to use the electronic library facilities at ZIB send a mail just containing "help" to elib@ZIB-Berlin.de.

Center for Research on Parallel Computation (CRPC), Houston

Gopher:

<gopher://softlib.rice.edu/11/softlib/tsplib>

Anonymous ftp:

<ftp://softlib.rice.edu/pub/tsplib>

References

1. R.E. Bland & D.F. Shallcross (1989). *Large Traveling Salesman Problems Arising from Experiments in X-ray Crystallography: A Preliminary Report on Computation*, Operations Research Letters 8, 125–128.
2. M. Grötschel & O. Holland (1991). *Solution of Large-Scale Symmetric Travelling Salesman Problems*, Mathematical Programming 51, 141–202.
3. M.W. Padberg & G. Rinaldi (1991). *A Branch & Cut Algorithm for the Resolution of Large-scale Symmetric Traveling Salesman Problems*, SIAM Review 33, 60–100.