

Dependently Typed Languages in Statix

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My Thesis

- Followed Compiler Construction A+B from Eelco
- Open question from Eelco: **How powerful is Statix?**

Background: What are Dependent Types?

- Types may depend on values!

Example

```
concat : (A: Set) -> (n m : Nat) -> Vec A n -> Vec A m  
        -> Vec A (n + m)
```

- Proof assistants
- For example: Agda, Coq, Lean, ...

Research Question

How suitable is Statix for defining a dependently-typed language?

- Will it be easier than doing it in Haskell?
- Defining system F was a challenge¹, will it even be possible?

¹Hendrik van Antwerpen et al. Scopes as types.

Why is this important?

Spoofax perspective

Developing a language with a complex type system tests the boundaries of what Spoofax can do.

Dependent Types perspective

Using a language workbench helps with rapid prototyping.

Primary Contribution: Calculus of Constructions in Statix

A lambda calculus with dependent types.

Example 1

```
(\v: Type. v) Type
```

Example 2

```
let f = \T: Type. \x: T. x;  
f (T: Type -> Type) (\y: Type. y)
```

Type Checking

Type checking relation

`typeOfExpr : scope * Expr -> Expr`

How do we use scopes?

A scope is used as a replacement for an environment and a context. One edge p. One relation `name → NameEntry`, `NameEntry` is either:

- `NType`: Stores a type (Corresponds to a context)
- `NSubst`: Stores a substitution (Corresponds to an environment)

Type Checking: Requires Evaluation

Example 1

```
let T = if false then Int else Bool end;  
let b: T = true;
```

Evaluation relation

```
betaHeadReduce : scope * Expr -> scope * Expr  
betaReduce : scope * Expr -> Expr  
exectBetaEq : (scope * Expr) * (scope * Expr)
```


Type Checking: From inference rules to Statix code

$$\frac{\langle s \mid e \rangle : t_e \quad \langle \text{sPutSubst}(s, x, (s, e)) \mid b \rangle : t_b}{\langle s \mid \text{Let}(x, e, b) \rangle : t_b}$$

Equivalent Statix code

```
typeOfExpr (s, Let(x, e, b)) = bt :-  
  typeOfExpr (s, e) == et,  
  typeOfExpr (sPutSubst (s, x, (s, e)), b) == bt
```

Extra contributions

Features

- ① Implemented Inference
- ② Implemented Inductive Data Types
- ③ Implemented Universes
- ④ Interpreter
- ⑤ Compiler to Clojure

Evaluation

- ① Comparison with implementation in Haskell
- ② Comparison with implementation in LambdaPi
- ③ Evaluation of Spoofax

Conclusions

Spoofax is a great tool for developing dependently typed languages!

- We can use scopes to represent environments and contexts
- Statix can still use improvements