Modeling Choice in Co-Design

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Contents

1   A Brief Introduction to Co-Design

Co-Design investigates *feasibility relations* between *resources* and *functionalities*. We describe these in turn.

1.1   Resource and Functionality Preorders

We model the resources and functionalities using preordered sets. The order relation $a \to b$ signifies that we can freely transform a into b. Equivalently, $a \to b$ can also mean that a is implicitly already b.

Example 1.1. (Money as Resources) The set $\mathbb{N} := \{0, 1, 2, \ldots\}$ with the order relation $a \to b \Leftrightarrow a \geq b$ can be used to describe money (of a single currency, say dollars) as a resource. The element a represents having a dollars, while $a \to b$ means that if you have a dollars, you also implicitly have b dollars. For example, $3 \to 1$ signifies that if we have 3 dollars, we also possess the purchasing power 1 dollar. Observe that being able to freely transform $3 \to 1$ presupposes people are willing to accept 3 dollars in lieu of 1 dollar. This hypothesis is system relative. For example, a vending machine might only accept bills of low denomination. So in this case one might well not have $50 \to 1$.

Returning to the general theory, we observe that the preorder axioms precisely describe the relation of free transformations.

Definition 1.1. (Preorder) A *preorder* is a set with a binary relation $R \subseteq A \times A$ satisfying reflexivity (r) and transitivity (t):

(r) $(a, a) \in R$ for all $a \in A$,

(t) if $(a, b) \in R$ and $(b, c) \in R$, then $(a, c) \in R$.

Usually we write the pair (a, b) using an infix symbol. In our case, $a \to b$ means (a, b) is part of the relation R.

Using our infix notation, the two axioms become

(r) $a \to a$ for all a,

(t) $a \to b$ and $b \to c$ imply $a \to c$.

In other words, each resource is freely transformable into itself (do nothing). Further, if we can freely transform a into b, and b into c, then we can freely transform a into c be doing one transformation after the other.

Everything that has been said here applies equally to functionalities. In fact, what we view as resources or functionalities is relative to the process we are considering. When I purchase a candy bar from a vending machine, the dollar I spend is the resource, and the candy bar the functionality. When I eat the candy bar to gain energy, the candy bar has become the resource.

1.2 Feasibility Relations

Given a preorder of resources $\mathcal{R}$ and a preorder of functionalities $\mathcal{F}$, a feasibility relation $\Phi : \mathcal{R} \nrightarrow \mathcal{F}$ between them expresses which functionalities I am able to obtain from which resources. We write $\Phi(r, f) = \texttt{true}$ to indicate that f is obtainable from r.

Example 1.2. Consider buying groceries. Our resources $\mathcal{R}$ will be money as in Example 1.1. The possible groceries we can buy are the functionalities $\mathcal{F}$. The feasibility relation $\Phi : \mathcal{R} \nrightarrow \mathcal{F}$ will describe what groceries we can purchase given a specific amount of cash. For example, a `carrot` might cost 1\$. Hence, $\Phi(1\$, \texttt{carrot}) = \texttt{true}$. On the other hand, $\Phi(0\$, \texttt{carrot}) = \texttt{false}$, since we can't get a `carrot` without paying anything. Suppose now that a `bag of carrots` costs 5\$. This means $\Phi(5\$, \texttt{bag of carrots}) = \texttt{true}$. Observe that also $\Phi(5\$, \texttt{carrot}) = \texttt{true}$ should hold. Either we can pick up a single carrot in the store and pay for it with 5\$, or we could buy the whole bag for 5\$ and this would also give us a single `carrot`.

This shows that feasibility interacts with the free transformations we have for the resources and functionalities. In our example 5\$ $\to$ 1\$ and $\Phi(1\$, \texttt{carrot}) = \texttt{true}$ implied $\Phi(5\$, \texttt{carrot}) = \texttt{true}$. Similarly, `bag of carrots` $\to$ `carrot` and $\Phi(5\$, \texttt{bag of carrots}) = \texttt{true}$ implied $\Phi(5\$, \texttt{carrot}) = \texttt{true}$.

The interaction we observed between free transformations and feasibility justifies the formal definition of a feasibility relation.

Definition 1.2 (Feasibility Relation)**.** Given preorders $\mathcal{R}$ and $\mathcal{F}$, a *feasibility relation* $\Phi : \mathcal{R} \nrightarrow \mathcal{F}$ is a relation $\Phi \subseteq \mathcal{R} \times \mathcal{F}$ such that

Figure 1: Internal diagram of a feasibility relation.

(i) if $r \to s$ in $\mathcal{R}$ and $(s, f) \in \Phi$, then $(r, f) \in \Phi$,

(ii) if $f \to g$ in $\mathcal{F}$ and $(r, f) \in \Phi$, then $(r, g) \in \Phi$.

If $(r, f) \in \Phi$ we write $\Phi(r, f) = \mathtt{true}$ of simply say that $\Phi(r, f)$ holds.

In other words, a feasibility relation is a relation from $\mathcal{R}$ to $\mathcal{F}$ that is downward closed with respect to free transformations in $\mathcal{R}$ and upward closed with respect to free transformations in $\mathcal{F}$. We will often draw internal diagrams of feasibility relations as in Figure 1.2. We put an arrow between an $r \in \mathcal{R}$ and $f \in \mathcal{F}$ iff $\Phi(r, f) = \mathtt{true}$. The arrows coming from Φ can be thought of the transformations that Φ enables. Observe that the closure conditions in Definition 1.2 mean that the arrows coming from Φ are transitively closed with respect to the internal arrows in $\mathcal{R}$ and $\mathcal{F}$: If there is a path from an $r \in \mathcal{R}$ to an $f \in \mathcal{F}$ following any type of arrow, then we also have a direct arrow between r and f coming from Φ.

We could have also defined feasibility relations as special types of monotone maps.

Lemma 1.1. *There is a one-to-one correspondence between feasibility relations* $\Phi : \mathcal{R} \nrightarrow \mathcal{F}$ *and monotone maps* $\mathcal{R}^{op} \times \mathcal{F} \to \mathsf{Bool}$, *where* Bool *denotes the preorder generated by* $\{\mathtt{false} \to \mathtt{true}\}$.

2 Choice between Resources

We will now extend what we've seen so far to encompass choice and uncertainty. What is meant by these terms is best illustrated with an example.

Example 2.1 (Menu Options). In a restaurant we might see the following menu options.

- Coffee or tea (customer choice)

- Asparagus or beetroot soup (depending on availability)

In the first case, we as customers are given the choice between coffee or tea. In the second, it is chosen for us which soup we will get, depending on what

the cooks have available. From our perspective as customers, this represents an uncertainty. In both cases, however, we will end up with exactly one of the options.

We might refer to the two modes of choice presented in Example 2.1 as *internal choice* and *external choice*, respectively. Alternatively, we could call them *choice* and *uncertainty*, or *free* and *forced* choice. Regardless of the names we choose, the terms are relative to which perspective we are taking. If we look through the eyes of the restaurant, the customer choice of coffee or tea is external, while the choice of soup is internal.

2.1 Choice Connectives

To formalize the two modes of choice described above, we introduce two connectives, denoted $\sqcup$ and $\sqcap$, respectively. Because the distinction between what is internal and external is relative to what perspective we are taking, we will hold off on generally fixing which symbol denotes which mode of choice. Instead, we will always indicate what the interpretation should be in each specific case.

For the moment, fix $\sqcap$ as free choice and $\sqcup$ as forced choice. Hence, to denote a free choice from a set A we write $\bigsqcap_{a \in A} a$, while for a forced choice from a set B we write $\bigsqcup_{b \in B} b$. We will now see how such choices should interact with the free transformations we have in a resource preorder, as well as each other.

2.1.1 Interaction of Choice with Free Transformations

Let $\mathcal{R}$ be a preorder of resources. If we have a free choice of any $a \in A$, then we can freely transform that choice into any specific $a \in A$ by making the corresponding choice. For example, if I have a choice between coffee or tea, I implicitly have coffee, and I implicitly have tea. In other words, I can downgrade any choice I get to make to the certainty of having one of the items I get to choose from. Formally, we have

$$\bigsqcap_{a \in A} a \to a, \quad \forall a \in A. \tag{1}$$

Now suppose we can freely obtain all a in a set A from a fixed resource t. In that case, having t means implicitly having a free choice between any

$a \in A$. For any a you end up choosing, there is a free way to get that a from t. Formally,

$$\forall a \in A : t \to a \quad \Rightarrow \quad t \to \prod_{a \in A} a. \tag{2}$$

Dually, if I have any $b \in B$ for certain, I can freely transform it into a forced choice among all $b \in B$. For example, having asparagus soup also implicitly means you have asparagus or beetroot soup (depending on availability). In other words, we can downgrade the certainty of having a specific $b \in B$ to the uncertainty of having some $b \in B$. Formally, this becomes

$$b \to \bigsqcup_{b \in B} b, \quad \forall b \in B. \tag{3}$$

Finally, suppose that for every $b \in B$ there is a free transformation to a fixed resource t. In that case, if someone chooses an arbitrary $b \in B$ for you, you can still freely obtain t using the corresponding free transformation. Hence,

$$\forall b \in B : b \to t \quad \Rightarrow \quad \bigsqcup_{b \in B} b \to t. \tag{4}$$

Mathematically, (1) and (2) mean that $\sqcap$ is the *meet* over $a \in A$, while (3) and (4) show that $\sqcup$ is the *join* over $b \in B$.

2.1.2 Mutual Interaction of the Two Modes of Choice

We now look at how the two connectives should interact with each other. First we consider the finite case. For simplicity we write

$$\prod_{x \in \{a,b\}} x =: a \sqcap b,$$

$$\bigsqcup_{x \in \{a,b\}} x =: a \sqcup b.$$

Suppose we have a series of choices given by $a \sqcap (b \sqcup c)$. This means we get to freely choose between getting a for certain, or between the uncertainty of getting b, or c. In this scenario, we can always guarantee a if we want, but we are uncertain whether we will get b or c, if we don't choose a. Compare this with the series of choices given by $(a \sqcap b) \sqcup (a \sqcap c)$. Here too we can always guarantee the a, but are uncertain whether we will get b or c, if we

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don't choose a. In fact, in terms of which resources we can guarantee, the two formulations are equivalent from our perspective. Hence, we can assume the following distributive law

$$a \sqcap (b \sqcup c) \cong (a \sqcap b) \sqcup (a \sqcap c).$$

Dually, we can compare $a \sqcup (b \sqcap c)$ with $(a \sqcup b) \sqcap (a \sqcup c)$. In the first expression we are uncertain whether we will get a, or a free choice between b and c. This means that if we don't get a, we can guarantee either b or c. This is also the case in the second expression. Hence, we may assume the distributive law

$$a \sqcup (b \sqcap c) \cong (a \sqcup b) \sqcap (a \sqcup c).$$

More generally, these equalities should also hold for the infinite case. The statement of the infinite distributive law is more complicated: For any doubly indexed family of elements $\{b_{j,k} : j \in J, k \in K_j\}$ of $\mathcal{R}$ we have

$$\prod_{j \in J} \left(\bigsqcup_{k \in K_j} b_{j,k} \right) \cong \bigsqcup_{f \in F} \left(\prod_{j \in J} b_{j,f(j)} \right), \tag{5}$$

where F is the set of all choice functions choosing for each index $j \in J$ some index $f(j) \in K_j$.

Suppose we are presented with a choice of the form $\prod_{j \in J} \left(\bigsqcup_{k \in K_j} b_{j,k} \right)$. Then we can guarantee $\bigsqcup_{k \in K_j} b_{j,k}$ for any $j \in J$ we choose. That is we can guarantee uncertainty only between the $b_{j,k}$ for one specific j of our choice. Now consider $\bigsqcup_{f \in F} \left(\prod_{j \in J} b_{j,f(j)} \right)$. Here, some $f \in F$ is chosen for us. We then get to chose freely between the $b_{j,f(j)}$ for that fixed f. Hence we still have uncertainty among a set of $b_{j,k}$ for a j of our choice. Moreover, f picks precisely one $k = f(j)$ among all $k \in K_j$, so the set of $b_{j,k}$ for which there is uncertainty remains the same in both cases.

We can also formulate the dual distributive law

$$\bigsqcup_{j \in J} \left(\prod_{k \in K_j} b_{j,k} \right) \cong \prod_{f \in F} \left(\bigsqcup_{j \in J} b_{j,f(j)} \right), \tag{6}$$

which happens to be equivalent to (5) provided all both $\prod_{a \in A}$ and $\bigsqcup_{a \in A}$ exist for any subset $A \subseteq \mathcal{R}$.

2.1.3   Laws for the Choice Connectives

We have seen that the choice connectives $\sqcap$ and $\sqcup$ on a resource preorder should behave like meets and joins over arbitrary subsets of $\mathcal{R}$. Moreover, they should distribute over one another, even in the infinite case. Hence, any structure that we use to model choice should be a completely distributive lattice. Conversely, such a model shouldn't satisfy any other laws that are not implied by our desired laws. This leads us to consider the free distributive lattice on $\mathcal{R}$ to model choice.

There are several notions of free distributive lattice one could use, depending on what set of maps one is interested in. We will follow [Mor04] which shows that the upper sets of lower sets of $\mathcal{R}$, denoted $\mathsf{U}(\mathsf{L}(\mathcal{R}))$ satisfy the following universal property:

> There is an order embedding $i : \mathcal{R} \hookrightarrow \mathsf{U}(\mathsf{L}(\mathcal{R}))$, such that for any monotone function $f : \mathcal{R} \to M$ into a completely distributive lattice M, there is a unique complete homomorphisms $u : \mathsf{U}(\mathsf{L}(\mathcal{R})) \to M$ (preserving arbitrary meets and joins) satisfying $u \circ i = f$.

In particular, if we put any completely distributive choice structure M on $\mathcal{R}$, there will be a unique map from $u : \mathsf{U}(\mathsf{L}(\mathcal{R})) \to M$, preserving choices.

2.2   The Upper-Lower Model

In this section we introduce and interpret the choice structure on a resource preorder $\mathcal{R}$ given by $\mathsf{U}(\mathsf{L}(\mathcal{R}))$. First, we will show how to use the upper-lower construction as a purely formal tool that realizes the desired properties we derived in Sections 2.1.1 and 2.1.2. Then, we will reconstruct the upper-lower model by giving an interpretation to the sets appearing therein. This will include deriving the corresponding laws based on that interpretation. The fact that both approaches agree suggests that we have chosen a good model.

2.2.1   Formally using the Upper-Lower Model

We begin by describing $\mathsf{U}(\mathsf{L}(\mathcal{R}))$ in greater detail. We again fix $\sqcap$ as free choice and $\sqcup$ as forced choice. First, we define the lattice of lower sets.

Definition 2.1 (Lower Sets). Let $(P, \to)$ be a preorder. A subset $L \subseteq P$ is a lower set, iff for any $l \in L$, $x \to l$ implies $x \in L$. The set of lower sets of P, ordered by inclusion is denoted by $\mathsf{L}(P)$.

Lemma 2.1. *For any preorder P, $\mathsf{L}(P)$ is a completely distributive lattice with meets given by set intersection and joins by set union.*

Definition 2.2 (Lower Closure). Given a subset $A \subseteq P$, its lower closure is the lower set given by

$$\downarrow A := \{x \in P : x \to a \text{ for some a } \in A\}.$$

Lemma 2.2. *For any preorder P, we can embed $i : P \hookrightarrow \mathsf{L}(P)$ by taking lower closures $p \mapsto \downarrow \{p\}$. This map is injective and satisfies $p \to q \Leftrightarrow i(p) \subseteq i(q)$.*

Dually, we can make the same definitions for upper sets.

Definition 2.3 (Upper Sets). Let $(P, \to)$ be a preorder. A subset $U \subseteq P$ is an upper set, iff for any $u \in U$, $u \to x$ implies $x \in U$. The set of upper sets of P, ordered by containment is denoted by $\mathsf{U}(P)$.

Lemma 2.3. *For any preorder P, $\mathsf{U}(P)$ is a completely distributive lattice with meets given by set union and joins by set intersection.*

Definition 2.4 (Lower Closure). Given a subset $A \subseteq P$, its lower closure is the lower set given by

$$\uparrow A := \{x \in P : a \to x \text{ for some a } \in A\}.$$

Lemma 2.4. *For any preorder P, we can embed $i : P \hookrightarrow \mathsf{U}(P)$ by taking upper closures $p \mapsto \uparrow \{p\}$. This map is injective and satisfies $p \to q \Leftrightarrow i(p) \supseteq i(q)$.*

In particular, for any preorder of resources $\mathcal{R}$, $\mathsf{U}(\mathsf{L}(\mathcal{R}))$ is completely distributive with meets given by set union and joins by set intersection. Moreover, we have an injective map $i : \mathcal{R} \hookrightarrow \mathsf{U}(\mathsf{L}(\mathcal{R}))$ given by $r \mapsto \uparrow \downarrow \{r\}$ that satisfies
$$r \to s \Leftrightarrow i(r) \supseteq i(s).$$

We implement our connectives $\sqcap$ and $\sqcup$ in $\mathsf{U}(\mathsf{L}(\mathcal{R}))$ by setting $\sqcap$ to be the meet and $\sqcup$ to be the join in $\mathsf{U}(\mathsf{L}(\mathcal{R}))$: For a subset $A \subseteq \mathsf{U}(\mathsf{L}(\mathcal{R}))$,

$$\prod_{a \in A} a := \bigcup_{a \in A} a,$$

$$\bigsqcup_{a \in A} a := \bigcap_{a \in A} a.$$

In particular, if we are interested choice between elements of $\mathcal{R}$, we first embed those elements into $\mathsf{U}(\mathsf{L}(\mathcal{R}))$ and then apply the corresponding operation. That is if $B \subseteq \mathcal{R}$,

$$\prod_{r \in B} a := \bigcup_{r \in B} i(r),$$

$$\bigsqcup_{r \in B} a := \bigcap_{r \in B} i(r).$$

Example 2.2 (Menu 2). In a restaurant we are allowed to choose between a vegetarian and a meat option. If we choose the vegetarian option, the restaurant will provide us a curry dish, or a casserole, depending on availability. If we choose the meat option we will get chicken or beef, depending on availability. Using our connectives, we can describe this as

$$(\texttt{curry} \sqcup \texttt{casserole}) \sqcap (\texttt{chicken} \sqcup \texttt{beef}).$$

We model our options as the discrete preorder $\{\texttt{curry}, \texttt{casserole}, \texttt{chicken}, \texttt{beef}\}$. Using the rules above for describing this choice in the upper-lower model, we represent

$$\begin{aligned}
\texttt{curry} \sqcup \texttt{casserole} &= i(\texttt{curry}) \cap i(\texttt{casserole}) \\
&= \uparrow \{\{\texttt{curry}\}\} \cap \uparrow \{\{\texttt{casserole}\}\} \\
&= \uparrow \{\{\texttt{curry}, \texttt{casserole}\}\}.
\end{aligned}$$

A similar calculation shows $\texttt{chicken} \sqcup \texttt{beef} = \uparrow \{\{\texttt{chicken}, \texttt{beef}\}\}$. Hence,

$$\begin{aligned}
(\texttt{curry} \sqcup \texttt{casserole}) \sqcap (\texttt{chicken} \sqcup \texttt{beef}) &= \uparrow \{\{\texttt{curry}, \texttt{casserole}\}\} \cup \uparrow \{\{\texttt{chicken}, \texttt{beef}\}\} \\
&= \uparrow \{\{\texttt{curry}, \texttt{casserole}\}, \{\texttt{chicken}, \texttt{beef}\}\}.
\end{aligned}$$

Example 2.3 (Menu 3). A restaurant offers two menus, depending on the day. As customers we can freely choose any option on the current menu. Suppose the menu items the first menu are $\{\texttt{curry}, \texttt{beef}\}$ and those on the second are $\{\texttt{casserole}, \texttt{chicken}\}$. In terms of our connectives, this choice becomes

$$(\texttt{curry} \sqcap \texttt{beef}) \sqcup (\texttt{casserole} \sqcap \texttt{chicken}).$$

In the upper-lower model we represent

$$\texttt{curry} \sqcap \texttt{beef} = i(\texttt{curry}) \cup i(\texttt{beef})$$
$$= \uparrow \{\{\texttt{curry}\}\} \cup \uparrow \{\{\texttt{beef}\}\}$$
$$= \uparrow \{\{\texttt{curry}\}, \{\texttt{beef}\}\}.$$

Similarly, $\texttt{casserole} \sqcap \texttt{chicken} = \uparrow \{\{\texttt{casserole}\}, \{\texttt{chicken}\}\}$. Hence,

$$(\texttt{curry} \sqcap \texttt{beef}) \sqcup (\texttt{casserole} \sqcap \texttt{chicken}) = \uparrow \{\{\texttt{curry}\}, \{\texttt{beef}\}\} \cap \uparrow \{\{\texttt{casserole}\}, \{\texttt{chicken}\}\}$$
$$= \uparrow \{\{\texttt{curry}, \texttt{casserole}\}, \{\texttt{curry}, \texttt{chicken}\},$$
$$\{\texttt{beef}, \texttt{casserole}\}, \{\texttt{beef}, \texttt{chicken}\}\}.$$

The examples above show the need for some rules for calculating unions and intersections in $\mathsf{U}(\mathsf{L}(\mathcal{R}))$.

Calculation Rules

Lemma 2.5. *For arbitrary upper closures of $A, B \subseteq \mathsf{L}(\mathcal{R})$ we have:*

(i) $\uparrow A \cup \uparrow B = \uparrow \{x : x \in A \text{ or } x \in B\} = \uparrow (A \cup B)$

(ii) $\bigcup_{i \in I}(\uparrow A_i) = \uparrow (\bigcup_{i \in I} A_i)$

(iii) $\uparrow A \cap \uparrow B = \uparrow \{a \cup b : a \in A, b \in B\} = \{a \cup b : a \in \uparrow A, b \in \uparrow B\}$

(iv) $\bigcap_{i \in I}(\uparrow A_i) = \uparrow\{\bigcup a_i : a_i \in A_i \text{ for all } i\} = \{\bigcup a_i : a_i \in \uparrow A_i \text{ for all } i\}$

Proof. It suffices to prove (ii) and (iv).

(ii) Suppose $a \in \bigcup_{i \in I}(\uparrow A_i)$. Then $a \in \uparrow A_i$ for some i. In particular, $a_i \subseteq a$ for $a_i \in A_i$. Hence $a_i \subseteq a$ for $a_i \in \bigcup_{i \in I} A_i$, showing $a \in \uparrow (\bigcup_{i \in I} A_i)$.
Conversely, if $a \in \uparrow (\bigcup_{i \in I} A_i)$, then $a_i \subseteq a$ for some $a_i \in \bigcup_{i \in I} A_i$, that is for some $a_i \in A_i$. Thus $a \in \uparrow A_i$, meaning $a \in \bigcup_{i \in I}(\uparrow A_i)$.

(iv) Suppose $a \in \bigcap_{i \in I}(\uparrow A_i)$. Then for all i there exists $a_i \in A_i$ such that $a_i \subseteq a$. Thus $\bigcup a_i \subseteq a$, meaning $a \in \uparrow\{\bigcup a_i : a_i \in A_i \text{ for all } i\}$.
Conversely, if $a \in \uparrow\{\bigcup a_i : a_i \in A_i \text{ for all } i\}$, there are i such that $\bigcup a_i \subseteq a$, where $a_i \in A_i$. Hence for all i, $a_i \subseteq \bigcup a_i \subseteq a$, meaning $a \in \bigcap_{i \in I}(\uparrow A_i)$.
The final equality assume $a \in \uparrow\{\bigcup a_i : a_i \in A_i \text{ for all } i\}$. Then there are i such that $\bigcup a_i \subseteq a$, where $a_i \in A_i$. This implies $a_i \subseteq a$ for all i, whence $a \in \uparrow A_i$ for all i. Hence $a = \bigcup a \in \{\bigcup a_i : a_i \in \uparrow A_i \text{ for all } i\}$.
If $a \in \{\bigcup a_i : a_i \in \uparrow A_i \text{ for all } i\}$, then there are i such that $\bigcup a_i = a$, where $a_i \in \uparrow A_i$. For each i, we have $x_i \subseteq a_i$ for $x_i \in A_i$. Hence $\bigcup x_i \subseteq \bigcup a_i = a$, showing $a \in \uparrow \{\bigcup a_i : a_i \in A_i \text{ for all } i\}$. ⌐

Lemma 2.6. *The following equalities hold for $a, b, \in \mathcal{R}$:*

(i) $i(a) \cup i(b) = \uparrow \{\downarrow \{a\}, \downarrow \{b\}\}$

(ii) $i(a) \cap i(b) = \uparrow \{\downarrow \{a\} \cup \downarrow \{b\}\} = \uparrow \{\downarrow \{a, b\}\}$

Proof. Use Lemma 2.5 and the fact that $\downarrow \{a\} \cup \downarrow \{b\} = \downarrow \{a, b\}$. ⌐

2.2.2 Interpretation-Based Reconstruction

In this section we will reconstruct $\mathsf{U}(\mathsf{L}(\mathcal{R}))$ based on interpretation. For this we fix $\sqcap$ as free choice and $\sqcup$ as forced choice. Because of distributivity, we can always represent any choice as a free choice over a set of forced choices.

We start by considering a forced choice of resources from a set $A \subseteq \mathcal{R}$. We will argue that $\bigsqcup_{a \in A} a \cong \bigsqcup_{a \in \downarrow A} a$, in the sense that we have the same information on the outcome of the choice in both cases. Observe that we have the following monotonicity law

$$a \to b \quad \Leftrightarrow \quad a \sqcup b \cong b.$$

If $a \to b$, then since $b \to b$, we have $a \sqcup b \to b$ by law (4). By law (3), $b \to a \sqcup b$, so $a \sqcup b \cong b$. Conversely, if $a \sqcup b \cong b$, then $a \to a \sqcup b \cong b$.

More generally,

$$\bigsqcup_{a \in A} a \cong b \quad \Leftrightarrow \quad a \to b, \ \forall a \in A.$$

In particular, for every $a \in A$, we have $\bigsqcup_{x \in \downarrow \{a\}} x \cong a$. Now $\downarrow A = \bigcup_{a \in A} \downarrow a$. Therefore,

$$\bigsqcup_{a \in \downarrow A} a = \bigsqcup_{x \in \bigcup_{a \in A} \downarrow \{a\}} x = \bigsqcup_{a \in A} \bigsqcup_{x \in \downarrow \{a\}} x \cong \bigsqcup_{a \in A} a.$$

This shows that for each forced choice, we can choose a representative of that forced choice, where the set that is being chosen from is a lower set. Hence, $\mathsf{L}(\mathcal{R})$ can be seen as the set of equivalence classes of forced choices between elements of $\mathcal{R}$.

By duality we have

$$\bigsqcap_{a \in \uparrow A} a \cong \bigsqcap_{a \in A} a.$$

This implies that we can represent each free choice (up to isomorphism) by an upper set. Hence, $\mathsf{U}(P)$ is the set of equivalence classes of free choices between elements of any preorder P. In particular, if we set $P := \mathsf{L}(\mathcal{R})$ we see that we can represent the set of free choices among forced choices among elements of $\mathcal{R}$ by $\mathsf{U}(\mathsf{L}(\mathcal{R}))$.

It remains to show that the natural meets and joins on $\mathsf{U}(\mathsf{L}(\mathcal{R}))$ align with the operations given by our interpretation. Let $A \in \mathsf{U}(\mathsf{L}(\mathcal{R}))$ represent a free choice over some forced choices a_i. That is, $A = \{a_i\}_{i \in I}$ represents

$\prod_{i \in I} a_i$. Similarly, suppose $B \in \mathsf{U}(\mathsf{L}(\mathcal{R}))$ is given by $B = \{b_j\}_{j \in J}$, representing $\prod_{j \in J} b_j$. Because

$$\prod_{i \in I} a_i \sqcap \prod_{j \in J} b_j = \prod_{a \in A} a \sqcap \prod_{b \in B} b \cong \prod_{x \in A \cup B} x,$$

we see that $A \cup B$ represents the free choice between A and B.

To see how forced choice is represented, we take $A, B \in \mathsf{U}(\mathsf{L}(\mathcal{R}))$. Then $A = \{a_i\}_{i \in I}$, where $a_i = \{r_{i,k}\}_{k \in K_i}$. Similarly, $B = \{b_l\}_{l \in L}$, where $b_l = \{r_{l,m}\}_{m \in M_l}$. Hence A represents $\prod_{i \in I} \bigsqcup_{k \in K_i} r_{i,k}$ and B represents $\prod_{l \in L} \bigsqcup_{m \in M_l} r_{l,m}$, where we assume that all index sets are disjoint.

Using distributivity

$$\prod_{i \in I} \bigsqcup_{k \in K_i} r_{i,k} \sqcup \prod_{l \in L} \bigsqcup_{m \in M_l} r_{l,m} \cong \prod_{i \in I} \prod_{l \in L} \left(\bigsqcup_{k \in K_i} r_{i,k} \sqcup \bigsqcup_{m \in M_l} r_{l,m} \right)$$

$$\cong \prod_{(i,l) \in I \times L} \bigsqcup_{j \in K_i \cup M_l} r_{h(j),j},$$

where $h(j) = \begin{cases} i, & \text{if } j \in K_i \\ l, & \text{if } j \in M_l \end{cases}$.

Therefore, the set representing the forced choice between A and B is $\{a \cup b : a \in A, b \in B\}$, which be Lemma 2.5 (iii) is precisely equal to $A \cap B$.

In summary, we have shown that $\mathsf{U}(\mathsf{L}(\mathcal{R}))$ corresponds exactly to the set of equivalence classes of free choices among forced choices among elements of $\mathcal{R}$. Moreover, the meet $A \cup B$ in $\mathsf{U}(\mathsf{L}(\mathcal{R}))$ represents the free choice between the choices represented by A and B. The join $A \cap B$ represents the forced choice between the choices represented by A and B. Analogous arguments can be made for arbitrary meets and joins.

3   Feasibility between Choices

Having constructed a model for choices between the resources in a preorder $\mathcal{R}$, we turn to the task of extending feasibility relations to the model. The first task will be to lift a feasibility relation $\Phi : \mathcal{R} \nrightarrow \mathcal{F}$ to the upper-lower models. We will see that such lifted feasibility relations display a certain kind of optimality in comparison to arbitrary feasibility relations between the upper-lower models.

3.1 Lifting Feasibility to Upper-Lower Models

We have seen in Section 2.2.2 that the choice of model is essentially determined by the laws we fixed in Sections 2.1.1 and 2.1.2. In a similar spirit, the way we are going to naturally lift feasibility relations is already determined by the choices we have made so far.

Let $\Phi : \mathcal{R} \nrightarrow \mathcal{F}$ be a feasibility relation between preorders $\mathcal{R}$ and $\mathcal{F}$. Lift this to a feasibility relation $\mathsf{L}(\Phi) : \mathsf{L}(\mathcal{R}) \nrightarrow \mathsf{L}(\mathcal{F})$ by setting

$$\mathsf{L}(\Phi)(A, B) = \mathtt{true} \iff \forall a \in A \, \exists b \in B : \Phi(a, b) = \mathtt{true}.$$

This is a feasibility relation since 'for all' is preserved by taking subsets and 'there exists' by taking supersets. Note that the choice of quantifiers is completely determined by the fact that we are ordering lower sets by inclusion. This choice was in turn fixed by the fact that forced choice is represented by lower sets, and forced choices over set can be weakened to forced choices over a superset. Hence, our lifting is fixed by the fact that forced choice behaves like a join.

Moving up a level, let $\Psi : \mathsf{L}(\mathcal{R}) \nrightarrow \mathsf{L}(\mathcal{F})$ be a feasibility relation. Lift this to $\mathsf{U}(\Psi) : \mathsf{U}(\mathsf{L}(\mathcal{R})) \nrightarrow \mathsf{U}(\mathsf{L}(\mathcal{F}))$ by setting

$$\mathsf{U}(\Psi)(A, B) = \mathtt{true} \iff \forall b \in B \, \exists a \in A : \Psi(a, b) = \mathtt{true}.$$

This is again a feasibility relation because the choice of quantifiers is compatible with the fact that we order the upper sets by containment. Dual to the above, the way we lift is fixed by the fact that we want free choice to act like a meet.

Putting both lifts together, given $\Phi : \mathcal{R} \nrightarrow \mathcal{F}$, we can lift it to $\mathsf{U}(\mathsf{L}(\Phi)) : \mathsf{U}(\mathsf{L}(\mathcal{R})) \nrightarrow \mathsf{U}(\mathsf{L}(\mathcal{F}))$ by iterating the two lifts. Explicitly,

$$\mathsf{U}(\mathsf{L}(\Phi))(A, B) = \mathtt{true} \iff \forall b \in B \, \exists a \in A : \forall x \in a \, \exists y \in b : \Phi(a, b).$$

3.2 Characterizing Lifted Feasibility Relations

3.2.1 Lower Set Level

Suppose we have a feasibility relation $\Phi : \mathcal{R} \nrightarrow \mathcal{F}$ which we lift to $\mathsf{L}(\Phi) : \mathsf{L}(\mathcal{R}) \nrightarrow \mathsf{L}(\mathcal{F})$ by setting

$$\mathsf{L}(\Phi)(A, B) = \mathtt{true} \iff \forall a \in A \, \exists b \in B : \Phi(a, b) = \mathtt{true}.$$

Decomposition Property   If $\mathsf{L}(\Phi)(A, B) = \mathtt{true}$, then

$$\forall a \in A \ \exists b \in B : \Phi(a, b) = \mathtt{true}.$$

In particular, for every $a \in A$ there is $b \in B$ such that

$$\forall x \in\ \downarrow \{a\} \ \exists y \in\ \downarrow \{b\} : \Phi(x, y) = \mathtt{true}.$$

This is because for any $x \in\ \downarrow \{a\}$ we have $\Phi(x, b)$ by monotonicity, and $b \in\ \downarrow \{b\}$. Hence, for every $a \in A$ there is $b \in B$ such that $\mathsf{L}(\Phi)(\downarrow \{a\}, \downarrow \{b\})$.

Union Property   Assume $\mathsf{L}(\Phi)(X_i, B) = \mathtt{true}$ for each X_i in a collection $\{X_i\}_{i \in I}$. Then for each X_i we have that

$$\forall x \in X_i \ \exists b \in B : \Phi(x, b) = \mathtt{true}.$$

This implies that

$$\forall x \in \bigcup_{i \in I} X_i \ \exists b \in B : \Phi(x, b) = \mathtt{true},$$

whence $\mathsf{L}(\Phi)(\bigcup_{i \in I} X_i, B) = \mathtt{true}$.

Now assume that we have an arbitrary feasibility relation $\Psi : \mathsf{L}(\mathcal{R}) \nrightarrow \mathsf{L}(\mathcal{F})$ satisfying decomposition property and the union property:

- If $\Psi(A, B) = \mathtt{true}$, then for every $a \in A$ there is $b \in B$ such that $\Psi(\downarrow \{a\}, \downarrow \{b\}) = \mathtt{true}$.

- For any collection $\{X_i\}_{i \in I} \subseteq \mathsf{L}(\mathcal{R})$,

$$\Psi(X_i, B) = \mathtt{true} \ \forall i \in I \quad \Rightarrow \quad \Psi(\bigcup_{i \in I} X_i, B) = \mathtt{true}.$$

We can obtain a feasibility relation $\tilde{\Psi} : \mathcal{R} \nrightarrow \mathcal{F}$ by setting

$$\tilde{\Psi}(r, f) := \Psi(\downarrow \{r\}, \downarrow \{f\}).$$

If $s \to r$, then $\downarrow \{s\} \subseteq\ \downarrow \{r\}$, so in this case

$$\tilde{\Psi}(r, f) \Leftrightarrow \Psi(\downarrow \{r\}, \downarrow \{f\}) \Rightarrow \Psi(\downarrow \{s\}, \downarrow \{f\}) \Leftrightarrow \tilde{\Psi}(r, f).$$

On the other hand if $f \to g$, then $\downarrow \{f\} \subseteq \downarrow \{g\}$, hence again

$$\tilde{\Psi}(r, f) \Leftrightarrow \Psi(\downarrow \{r\}, \downarrow \{f\}) \Rightarrow \Psi(\downarrow \{r\}, \downarrow \{g\}) \Leftrightarrow \tilde{\Psi}(r, g).$$

We show that $\mathsf{L}(\tilde{\Psi}) = \Psi$. This requires proving that $\mathsf{L}(\tilde{\Psi})(A, B) = \texttt{true} \Leftrightarrow \Psi(A, B) = \texttt{true}$. Suppose $\mathsf{L}(\tilde{\Psi})(A, B) = \texttt{true}$. This means that

$$\forall a \in A \; \exists b \in B : \Psi(\downarrow \{a\}, \downarrow \{b\}) = \texttt{true}.$$

In particular, by monotonicity,

$$\forall a \in A \; \Psi(\downarrow \{a\}, B) = \texttt{true}.$$

By the union property in the source, $\Psi(\bigcup_{a \in A} \downarrow \{a\}, B) = \texttt{true}$. It remains to observe that $\bigcup_{a \in A} \downarrow \{a\} = \downarrow \bigcup_{a \in A} a = \downarrow A = A$, since A is a lower set. Hence, $\Psi(A, B) = \texttt{true}$.

Conversely, suppose $\Psi(A, B) = \texttt{true}$. Then for each $a \in A$ there is $b \in B$ satisfying $\Psi(\downarrow \{a\}, \downarrow \{b\}) = \texttt{true}$. In symbols,

$$\forall a \in A \; \exists b \in B : \Psi(\downarrow \{a\}, \downarrow \{b\}) = \texttt{true}.$$

But this is precisely the condition showing that $\mathsf{L}(\tilde{\Psi})(A, B) = \texttt{true}$. Therefore Ψ is induced by $\tilde{\Psi}$.

The above proves the following characterization of induced feasibility relations on the lower set level.

Proposition 3.1 (Induction Criterion for Lower Sets). *Let $\Psi : \mathsf{L}(\mathcal{R}) \twoheadrightarrow \mathsf{L}(\mathcal{F})$ be a feasibility relation. Then $\Phi = \mathsf{L}(\tilde{\Phi})$ for some $\tilde{\Phi} : \mathcal{R} \twoheadrightarrow \mathcal{F}$ if and only if Φ satisfies the following:*

- *If $\Psi(A, B) = \texttt{true}$, then for every $a \in A$ there is $b \in B$ such that $\Psi(\downarrow \{a\}, \downarrow \{b\}) = \texttt{true}$.*

- *For any collection $\{X_i\}_{i \in I} \subseteq \mathsf{L}(\mathcal{R})$,*

$$\Psi(X_i, B) = \texttt{true} \; \forall i \in I \quad \Rightarrow \quad \Psi\Big(\bigcup_{i \in I} X_i, B\Big) = \texttt{true}.$$

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3.2.2 Upper Set Level

We want to derive a similar criterion for upper sets. Let $\Psi : \mathsf{U}(\mathcal{R}) \nrightarrow \mathsf{U}(\mathcal{F})$ be a feasibility relation. Observe that $\mathsf{U}(\mathcal{R}) = \mathsf{L}(\mathcal{R}^{\mathrm{op}})^{\mathrm{op}}$. Hence we can view Ψ as being a feasibility relation $\mathsf{L}(\mathcal{R}^{\mathrm{op}})^{\mathrm{op}} \nrightarrow \mathsf{L}(\mathcal{F}^{\mathrm{op}})^{\mathrm{op}}$, which corresponds bijectively to a feasibility relation $\check{\Psi} : \mathsf{L}(\mathcal{F}^{\mathrm{op}}) \nrightarrow \mathsf{L}(\mathcal{R}^{\mathrm{op}})$. By the induction criterion (Proposition 3.1) we know that $\check{\Psi}$ is $\mathsf{L}(\tilde{\Psi})$ for $\tilde{\Psi} : \mathcal{F}^{\mathrm{op}} \nrightarrow \mathcal{R}^{\mathrm{op}}$ iff we satisfy the decomposition and union properties:

- If $\check{\Psi}(A, B) = \mathtt{true}$, then for every $a \in A$ there is $b \in B$ such that $\check{\Psi}(\downarrow \{a\}, \downarrow \{b\}) = \mathtt{true}$.

- For any collection $\{X_i\}_{i \in I}$, where $X_i \in \mathsf{L}(\mathcal{F}^{\mathrm{op}})$,
$$\check{\Psi}(X_i, B) = \mathtt{true} \; \forall i \in I \quad \Rightarrow \quad \check{\Psi}(\bigcup_{i \in I} X_i, B) = \mathtt{true}.$$

Since $\check{\Psi}(A, B) = \Psi(B, A)$, and lower closures become upper closures under taking opposites, this means that $\check{\Psi}$ is induced iff

- If $\Psi(B, A) = \mathtt{true}$, then for every $a \in A$ there is $b \in B$ such that $\Psi(\uparrow \{b\}, \uparrow \{a\}) = \mathtt{true}$.

- For any collection $\{X_i\}_{i \in I}$, where $X_i \in \mathsf{U}(\mathcal{F})$,
$$\Psi(B, X_i) = \mathtt{true} \; \forall i \in I \quad \Rightarrow \quad \Psi(B, \bigcup_{i \in I} X_i) = \mathtt{true}.$$

Because all the correspondences we used to get from $\Psi : \mathsf{U}(\mathcal{R}) \nrightarrow \mathsf{U}(\mathcal{F})$ to $\check{\Psi} : \mathsf{L}(\mathcal{F}^{\mathrm{op}}) \nrightarrow \mathsf{L}(\mathcal{R}^{\mathrm{op}})$ are invertible, this means that the original Ψ is induced by $\tilde{\Psi} : \mathcal{F}^{\mathrm{op}} \nrightarrow \mathcal{R}^{\mathrm{op}}$ iff the above conditions are fulfilled. We conclude by observing that $\tilde{\Psi} : \mathcal{F}^{\mathrm{op}} \nrightarrow \mathcal{R}^{\mathrm{op}}$ corresponds bijectively to a feasibility relation $\mathcal{R} \nrightarrow \mathcal{F}$. Therefore,

Proposition 3.2 (Induction Criterion for Upper Sets)**.** *Let $\Psi : \mathsf{U}(\mathcal{R}) \nrightarrow \mathsf{U}(\mathcal{F})$ be a feasibility relation. Then $\Phi = \mathsf{L}(\tilde{\Phi})$ for some $\tilde{\Phi} : \mathcal{R} \nrightarrow \mathcal{F}$ if and only if Φ satisfies the following:*

- *If $\Psi(A, B) = \mathtt{true}$, then for every $b \in B$ there is $a \in A$ such that $\Psi(\uparrow \{a\}, \uparrow \{b\}) = \mathtt{true}$.*

- *For any collection $\{X_i\}_{i \in I} \subseteq \mathsf{U}(\mathcal{F})$,*
$$\Psi(A, X_i) = \mathtt{true} \; \forall i \in I \quad \Rightarrow \quad \Psi(A, \bigcup_{i \in I} X_i) = \mathtt{true}.$$

3.2.3 Upper-Lower Induction

We prove the following characterization directly.

Theorem 3.1. *A feasibility relation* $\Psi : \mathsf{U}(\mathsf{L}(\mathcal{R})) \twoheadrightarrow \mathsf{U}(\mathsf{L}(\mathcal{F}))$ *has the form* $\mathsf{U}(\mathsf{L}(\tilde{\Psi}))$ *for* $\tilde{\Psi} : \mathcal{R} \twoheadrightarrow \mathcal{F}$ *if and only if the following three conditions are fulfilled:*

(i) If $\Psi(A, B) = \mathtt{true}$*, then*

$$\forall b \in B \, \exists a \in A : \forall r \in a \, \exists y \in b : \Psi(\uparrow\downarrow\{r\}, \uparrow\downarrow\{f\}) = \mathtt{true}.$$

(ii) For any collection $\{X_i\}_{i \in I} \subseteq \mathsf{U}(\mathcal{F})$*,*

$$\Psi(A, X_i) = \mathtt{true} \; \forall i \in I \quad \Rightarrow \quad \Psi(A, \bigcup_{i \in I} X_i) = \mathtt{true}.$$

(iii) For any collection $\{X_i\}_{i \in I} \subseteq \mathsf{U}(\mathcal{R})$*,*

$$\Psi(X_i, B) = \mathtt{true} \; \forall i \in I \quad \Rightarrow \quad \Psi(\bigcap_{i \in I} X_i, B) = \mathtt{true}.$$

In that case, $\tilde{\Psi}$ *is obtained by setting* $\tilde{\Psi}(r, f) = \Psi(\uparrow\downarrow\{r\}, \uparrow\downarrow\{f\})$*.*

Proof. First, we show that given a feasibility relation $\Phi : \mathcal{R} \twoheadrightarrow \mathcal{F}$, the induced feasibility relation $\mathsf{U}(\mathsf{L}(\Phi)) : \mathsf{U}(\mathsf{L}(\mathcal{R})) \twoheadrightarrow \mathsf{U}(\mathsf{L}(\mathcal{F}))$ satisfies (i)-(iii).

(i)

(ii)

(iii) Next, assume we have a feasibility relation $\Psi : \mathsf{U}(\mathsf{L}(\mathcal{R})) \twoheadrightarrow \mathsf{U}(\mathsf{L}(\mathcal{F}))$ satisfying (i)-(iii). We show that $\Psi = \mathsf{U}(\mathsf{L}(\tilde{\Psi}))$, where $\tilde{\Psi}(r, f) = \Psi(\uparrow\downarrow\{r\}, \uparrow\downarrow\{f\})$.
Suppose that $\Psi(A, B) = \mathtt{true}$. Then by (i) we know that

$$\forall b \in B \, \exists a \in A : \forall r \in a \, \exists y \in b : \Psi(\uparrow\downarrow\{r\}, \uparrow\downarrow\{f\}) = \mathtt{true}.$$

But this implies, by definition, that $\mathsf{U}(\mathsf{L}(\tilde{\Psi}))(A, B) = \mathtt{true}$.
Conversely, suppose $\mathsf{U}(\mathsf{L}(\tilde{\Psi}))(A, B) = \mathtt{true}$. Then we know that

$$\forall b \in B \, \exists a \in A : \forall r \in a \, \exists y \in b : \Psi(\uparrow\downarrow\{r\}, \uparrow\downarrow\{f\}) = \mathtt{true}.$$

$\square$

Proposition 3.3. *Suppose* $\Psi : \mathsf{U}(\mathsf{L}(\mathcal{R})) \twoheadrightarrow \mathsf{U}(\mathsf{L}(\mathcal{F}))$ *is induced as in Theorem 3.1. Then the following equivalences hold for principal upper sets*

(i) $\Psi(\bigcup_{i \in I} X_i, \uparrow \{z\}) \Leftrightarrow \Psi(X_i, \uparrow \{z\})$ *for some* $i \in I$.

(ii) $\Psi(\uparrow \{\downarrow \{s\}\}, \bigcap_{i \in I} X_i) \Leftrightarrow \Psi(\uparrow \{\downarrow \{s\}\}, X_i)$ *for some* $i \in I$.

Proof. Note that for both (i) and (ii) the right-hand sides imply the left by monotonicity.

For (i) suppose $\Psi(\bigcup_{i \in I} X_i, \uparrow \{z\})$ holds. Then we know that

$$\forall b \in \uparrow \{z\} \; \exists a \in \bigcup_{i \in I} X_i : \forall r \in a \; \exists f \in b : \tilde{\Psi}(r, f).$$

In particular, for z we have an $a \in \bigcup_{i \in I} X_i : \forall r \in a \; \exists f \in z : \tilde{\Psi}(r, f)$. This a will lie in some X_i, so we have an $a \in X_i : \forall r \in a \; \exists f \in z : \tilde{\Psi}(r, f)$. It remains to show that this choice of a also works for any $b \in \uparrow \{z\}$. Let $b \in \uparrow \{z\}$. Then $z \subseteq b$. Hence, for the $a \in X_i$ we have that for all $r \in a$ there is $f \in z \subseteq b$ such that $\tilde{\Psi}(r, f)$. This shows that for some X_i,

$$\forall b \in \uparrow \{z\} \; \exists a \in X_i : \forall r \in a \; \exists f \in b : \tilde{\Psi}(r, f),$$

meaning that $\Psi(X_i, \uparrow \{z\})$ for some $i \in I$.

For (ii) suppose $\Psi(\uparrow \{\downarrow \{s\}\}, \bigcap_{i \in I} X_i) = \mathtt{true}$. Recall from Lemma 2.5 that $\bigcap_{i \in I} X_i = \{\cup x_i : x_i \in X_i\}$. Hence,

$$\forall b \in \{\cup x_i : x_i \in X_i\} \; \exists a \in \uparrow \{\downarrow \{s\}\} : \forall r \in a \; \exists f \in b : \tilde{\Psi}(r, f).$$

Fix some $\cup x_i$. Then there is $a \supseteq \downarrow \{s\}$ such that for all $r \in a$ there is $f \in \cup x_i$ satisfying $\tilde{\Psi}(r, f)$. In particular, the same holds if we restrict a to $\downarrow \{s\}$. Hence,

$$\forall b \in \{\cup x_i : x_i \in X_i\} \; \forall r \in \downarrow \{s\} \; \exists f \in b : \tilde{\Psi}(r, f).$$

Since

Suppose that $\Psi(\uparrow \{\downarrow \{s\}\}, X_i) = \mathtt{false}$ for all i. That is, for all $i \in I$,

$$\exists b_i \in X_i \; \forall a \in \uparrow \{\downarrow \{s\}\} : \exists r \in a \; \forall f \in b_i : \tilde{\Psi}(r, f) = \mathtt{false}.$$

In particular for any such b_i we have that

$$\exists r \in \downarrow \{s\} \; \forall f \in b_i : \tilde{\Psi}(r, f) = \mathtt{false}.$$

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By monotonicity, this means in particular that $\forall f \in b_i : \tilde{\Psi}(s, f) = \texttt{false}$, from which it follows that $\tilde{\Psi}(s, f) = \texttt{false}$ for all $f \in \cup b_i$. Therefore $\cup b_i \in \{\cup x_i : x_i \in X_i\}$ is an element of the intersection for which there exists $r \in \downarrow \{s\}$, (namely s) such that for all $f \in \cup b_i$ we have $\tilde{\Psi}(s, f) = \texttt{false}$. This contradicts our assumption that $\Psi(\uparrow \{\downarrow \{s\}\}, \bigcap_{i \in I} X_i) = \texttt{true}$. $\quad$ ⌐⌐

3.3 Formulas for Queries on Induced Feasibility

Let $\Psi : \mathsf{U}(\mathsf{L}(\mathcal{R})) \twoheadrightarrow \mathsf{U}(\mathsf{L}(\mathcal{F}))$ be an induced feasibility relation, that is $\Psi = \mathsf{U}(\mathsf{L}(\tilde{\Psi}))$ for $\tilde{\Psi}(r, f) = \Psi(\uparrow \downarrow \{r\}, \uparrow \downarrow \{f\})$. We define the following sets called queries.

Definition 3.1 (Querying Resources). For a fixed $A \in \mathsf{U}(\mathsf{L}(\mathcal{R}))$ define

$$r(A) := \{f \in \mathcal{F} : \Psi(A, i(f))\}.$$

The result is an upper set in $\mathcal{F}$.

Definition 3.2 (Querying Functionalities). For a fixed $B \in \mathsf{U}(\mathsf{L}(\mathcal{F}))$ define

$$f(B) := \{r \in \mathcal{R} : \Psi(i(r), B)\}.$$

The result is a lower set in $\mathcal{R}$.

Using the characterization in Theorem 3.1, we can immediately obtain formulas for queries between choices. Recall from Section 2.2 that for $A \subseteq \mathsf{U}(\mathsf{L}(\mathcal{R}))$ we define

$$\prod_{a \in A} a := \bigcup_{a \in A} a,$$

$$\bigsqcup_{a \in A} a := \bigcap_{a \in A} a.$$

Proposition 3.4. *For*

(i) $r(\prod_{a \in A} a) = r(\bigcup_{a \in A} a) = \bigcup_{a \in A} r(a)$

(ii) $r(\bigsqcup_{a \in A} a) = r(\bigcap_{a \in A} a) = \bigcap_{a \in A} r(a)$

(iii) $f(\prod_{b \in B} b) = f(\bigcup_{b \in B} b) = \bigcap_{b \in B} f(b)$

(iv) $f(\bigsqcup_{b \in B} b) = f(\bigcap_{b \in B} b) = \bigcup_{b \in B} f(b)$

Proof. We prove each in turn.

(i) Follows from 3.3 (i)

(ii) By definition $B \in r(\bigcap_{a \in A} a) \Leftrightarrow \Psi(\bigcap_{a \in A} a, B)$. By monotonicity and Theorem 3.1 (iii) we know $\Psi(\bigcap_{a \in A} a, B) \Leftrightarrow \Psi(a, B), \forall a \in A$. In summary, $B \in r(\bigcap_{a \in A} a) \Leftrightarrow B \in r(a), \forall a \in A \Leftrightarrow B \in \bigcap_{a \in A} r(A)$.

(iii) Obeserve that by monotonicity and Theorem 3.1 (ii) we have $\Psi(A, \bigcup_{b \in B}) \Leftrightarrow \Psi(A, b), \forall b \in B$. Hence $A \in f(\bigcup_{b \in B}) \Leftrightarrow \Psi(A, \bigcup_{b \in B}) \Leftrightarrow \Psi(A, b), \forall b \in B, \Leftrightarrow A \in f(b), \forall b \in B \Leftrightarrow A \in \bigcap_{b \in B} f(b)$.

(iv) Follows from 3.3 (ii)

3.4 General Interpretation for Feasibility between Upper-Lower Models

3.4.1 Induced Feasibility

4 Choice between Transformations

5 Towards Modeling Dependencies

A Laws for Choice

Derive some general laws for distributivity ect... Especially those used is section on interpretation based reconstruction

Lemma A.1.
$$\prod_{a \in A} a \sqcap \prod_{b \in B} b \cong \prod_{x \in A \cup B} x$$

and its dual.

Lemma A.2. *The folowing holds for sets $\{r_i\}_{i \in I}$ and $\{r_l\}_{l \in L}$:*

$$\prod_{i \in I} r_i \sqcup \prod_{l \in L} r_l = \prod_{i \in I} \prod_{l \in L} (r_i \sqcup r_l)$$

Proof. Write

$$\prod_{i \in I} r_i \sqcup \prod_{l \in L} r_l = \bigsqcup_{X \in \{I,L\}} \prod_{j \in X} r_{X,j}.$$

Let F be the set of choice functions f from $\{I, L\}$ to $I \cup L$, such that $f(I) \in I$ and $f(L) \in L$. Observe that $F \cong I \times L$. Hence,

$$\begin{aligned}
\bigsqcup_{X \in \{I,L\}} \prod_{j \in X} r_{X,j} &= \prod_{f \in F} \bigsqcup_{X \in \{I,L\}} r_{X,f(X)} \\
&= \prod_{f \in F} r_{I,f(I)} \sqcup r_{L,f(L)} \\
&= \prod_{(i,l) \in I \times L} r_{I,i} \sqcup r_{L,l} \\
&= \prod_{i \in I} \prod_{l \in L} (r_i \sqcup r_l).
\end{aligned}$$

$\qed$

and its dual.

References

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