



Percolation

Jonathan Marriott

Supervised by Dr Edward Crane
Level 6
20 Credit Points

29th April 2022

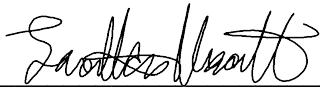
Acknowledgement of Sources

Acknowledgement of Sources

For all ideas taken from other sources (books, articles, internet), the source of the ideas is mentioned in the main text and fully referenced at the end of the report.

All material which is quoted essentially word-for-word from other sources is given in quotation marks and referenced.

Pictures and diagrams copied from the internet or other sources are labelled with a reference to the web page or book, article etc.

Signed 

Date _____

Percolation

Jonathan Marriott

Contents

| | | |
|----------|---|-----------|
| 1 | Introduction | 3 |
| 2 | Introducing the Percolation Model | 4 |
| 2.1 | Cubic Lattices | 5 |
| 2.2 | Probability Space | 5 |
| 2.3 | The Percolation Model | 6 |
| 3 | Critical values | 9 |
| 3.1 | Existence of a critical value on \mathbb{Z} | 9 |
| 3.2 | Existence of a critical value on \mathbb{Z}^2 | 9 |
| 3.3 | Existence of a critical value on \mathbb{Z}^d | 11 |
| 4 | Harris-Kesten Theorem | 12 |
| 4.1 | Some preliminary results | 12 |
| 4.2 | Crossing rectangles | 16 |
| 4.3 | Harris' Theorem | 19 |
| 4.4 | Kesten's Theorem | 23 |
| 5 | Simulation | 29 |
| 5.1 | Percolation simulation design | 29 |
| 5.2 | Simulation Implementation | 34 |
| 6 | References | 36 |
| 7 | Appendix | 38 |

1 Introduction

Originally the percolation model was devised to capture the physical properties of liquids flowing through porous media. The model introduced in 1957 by Broadbent and Hammersley [6] has produced a rich mathematical theory with wide-ranging applications in areas including disease spread [15], network robustness [7], ferromagnetism [1] and forest fire spread [17]. In this report we first introduce the percolation model, with no expectation

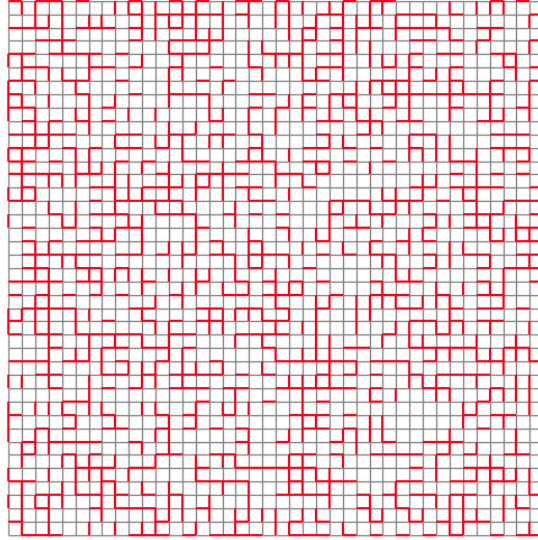


Figure 1: An example bond percolation configuration on a 40 by 40 box of the square lattice when $p = 0.4$. Where red edges are open and grey edges are closed.

on the reader to have a strong understanding of measure theory nor probability theory as the necessary definitions are included here. Subsequently, our main aim is to prove a key result in percolation theory - the Harris-Kesten Theorem. This theorem tells us the point at which percolation on the square lattice undergoes its phase transition, notably percolation is one of the simplest mathematical models to exhibit such behaviour. Finally, we look at percolation through the lens of computer simulations, to see how one might investigate the critical behaviour in unknown graph structures.

2 Introducing the Percolation Model

Broadly percolation theory can be split into 2 main categories: bond percolation and site percolation. Given some initial graph, in site percolation instead we assign the vertices of the graph to be open or closed. Whereas in bond percolation we assign edges to be open or closed, see Figure 1. Due to historic reasons the terminology used in percolation theory differs from that used in modern graph theory. Where graph edges and vertices are often referred to as bonds and sites respectively. Throughout this report we will be largely focusing on bond percolation, although it is interesting to note that all bond percolation problems can be reformulated as site percolation problem but not always the other way around [11].

2.1 Cubic Lattices

We start by introducing some basic notation and definitions for cubic lattices. In general, we match the notation and definitions given in [18]. The measure and probability theory definitions given largely follow [2], but we also use [4] for percolation specific details.

Definition 2.1 (Integers). $\mathbb{Z} = \{\dots, -2, -1, 0, 1, 2, \dots\}$ and $\mathbb{Z}^d = \{(x_1, x_2, \dots, x_d) : x_i \in \mathbb{Z}\}$

Definition 2.2 (1-norm distance). For $x, y \in \mathbb{Z}^d$, define the distance from x to y , denoted $\delta(x, y)$, by

$$\delta(x, y) := \sum_{i=1}^d |x_i - y_i|$$

Definition 2.3 (d-dimensional cubic lattice). We construct the d-dimensional cubic lattice with vertex set \mathbb{Z}^d and edges between vertices at distance one.

$$E(\mathbb{Z}^d) = \{\{u, v\} : u, v \in V(\mathbb{Z}^d), \delta(u, v) = 1\}$$

We will often abuse notation and refer to d-dimensional cubic lattice by the vertex set \mathbb{Z}^d without specifying the edge set. We will also denote the origin by 0.

2.2 Probability Space

We now introduce the measure theory basics required to define the probability measure and subsequently the probability space for our percolation model.

Definition 2.4 (σ -algebra). For some set X , we call $\mathcal{A} \subseteq \mathcal{P}(X)$, a subset of the power set of X , a σ -algebra of X if:

1. $\emptyset, X \in \mathcal{A}$
2. $A \in \mathcal{A} \implies A^c := X \setminus A \in \mathcal{A}$ (Closed under complement)
3. For $A_i \in \mathcal{A}, i \in \mathbb{N}$ we have that $\bigcup_{i=1}^{\infty} A_i \in \mathcal{A}$ (Closed under countable unions)

We call the pair (X, \mathcal{A}) a measurable space, and elements of \mathcal{A} measurable sets.

We note by De Morgan's Laws that a σ -algebra is also closed under countable intersections.

Definition 2.5 (Measure). A measure μ on a measurable space (X, \mathcal{A}) is a function $\mu : \mathcal{A} \rightarrow [0, \infty]$ where $\mu(\emptyset) = 0$ and for disjoint $A_i \in \mathcal{A}$, $i \in \mathbb{N}$ we have that

$$\mu\left(\bigcup_{i=1}^{\infty} A_i\right) = \sum_{i=1}^{\infty} \mu(A_i)$$

Then we call the triple (X, \mathcal{A}, μ) a measure space.

In the context above if we think about X being our set of outcomes, and the sigma algebra of X representing the set of events we wish to assign a probability to, it seems sensible to give these events a probability via measure. However, we clearly need to restrict our measure which assigns these probabilities such that its domain is the interval $[0, 1]$

Definition 2.6 (Probability Measure). Let Ω be the set of all outcomes (the sample space) and \mathcal{F} be a σ -algebra of Ω where the elements are events we wish to consider (the event space). Then a measure \mathbb{P} on (Ω, \mathcal{F}) is a probability measure if:

1. $\mathbb{P} : \mathcal{F} \rightarrow [0, 1]$
2. $\mathbb{P}(\Omega) = 1$

Then we call the triple $(\Omega, \mathcal{F}, \mathbb{P})$ a probability space.

2.3 The Percolation Model

We now take some parameter $p \in [0, 1]$, this will specify the probability any given edge is open. Setting $q = 1 - p$ we say each edge is independently open with probability p and closed with probability q . We can think of the open and closed edges defining a random subgraph of \mathbb{Z}^d where only edges set to open are retained.

Now we define our sample space

$$\Omega = \prod_{e \in E(\mathbb{Z}^d)} \{0, 1\}$$

where an edge in state 1 represents it is open, and 0 that it is closed.

We may refer to Ω as the set of configurations, where a configuration $\omega \in \Omega$ is a function assigning edges to be open or closed, meaning

$$\begin{aligned} \omega : E &\rightarrow \{0, 1\} \\ e &\mapsto \omega_e \end{aligned}$$

Then we must define \mathcal{F} , some σ -algebra of Ω , the set events we wish to assign a probability to. Clearly we cannot just take $\mathcal{F} = \mathcal{P}(\Omega)$, since we have uncountably many configurations in Ω . This can be easily seen via a

diagonalization argument. It turns out we can generate the σ -algebra we want by the cylinder sets of Ω . We base our definition on the one given by [4].

Definition 2.7 (Cylinder Set). We say a subset $S \subseteq \Omega$ is a cylinder set if and only if there exists a finite subset $F \subseteq E$ and $\sigma \in \{0, 1\}^F$ such that

$$S(F, \sigma) = \{\omega \in \Omega : \omega_f = \sigma_f \text{ for } f \in F\}$$

In essence the cylinder set is the set of configurations ω which map all the edges in F to the same state that σ does.

Then we define \mathcal{F} to be the set of all unions of cylinder sets of Ω , then \mathcal{F} is said to be generated by the cylinder sets of Ω . In this case \mathcal{F} is a σ -algebra of Ω . Intuitively \mathcal{F} is the set of events which only depend on a finite number of edges.

Next we define our probability measure on (Ω, \mathcal{F}) by

$$\mathbb{P}_p(S(F, \sigma)) = \prod_{f \in F} (p(\sigma_f) + q(1 - \sigma_f))$$

Where again $q = 1 - p$, and we use the subscript p in \mathbb{P}_p to emphasize that p is the parameter in our model, such that later if we fix p to a specific value it will be clear. Notice that we have defined \mathbb{P}_p for cylinder sets rather than elements of the σ -algebra, however it turns out by Carathéodory's extension theorem the probability measure above has a unique extension to the whole σ -algebra. For the details of this theorem see Theorem 1.3.10 in [2]. Then $(\Omega, \mathcal{F}, \mathbb{P}_p)$ is our probability space in which we examine the percolation model.

We now lay out some definitions specific for percolation.

Definition 2.8. Let $C(x)$ denote the open cluster (component) containing x , which is the set of vertices in \mathbb{Z}^d which are connected to x by a path of open edges. We abbreviate the open cluster containing the origin $C(0)$ by C

Definition 2.9 (Increasing Event). An event $A \subseteq \Omega$ is increasing if when $\omega \in A$ and

$$\forall e \in E(\mathbb{Z}^d), \omega_e = 1 \implies \omega'_e = 1$$

Then $\omega' \in A$.

We notice that the event $\{|C| = \infty\}$ is clearly an increasing event, since adding open edges to a configuration with an infinite connected cluster containing the origin will not remove the infinite cluster.

Definition 2.10 (Percolation function). We define the percolation function $\theta(p)$ as follows

$$\theta(p) = \mathbb{P}_p(|C| = \infty)$$

In words the percolation function is simply the probability that we can reach an infinite number of vertices from the origin by open edges. Furthermore, we also note that this is the same as asking what is the probability of having an infinite length self-avoiding path of open edges starting at the origin.

We intend to show that this percolation function is non-decreasing, but first we show a more general result for increasing events.

Lemma 2.11. *If $A \subseteq \Omega$ is an increasing event then $\mathbb{P}_p(A)$ is non-decreasing in p .*

Proof. We use the coupling of percolation processes to show that $\mathbb{P}_p(A)$ is non-decreasing. We use the definition of couplings given by [16].

Definition 2.12 (Coupling). Let \mathbb{P} and \mathbb{P}' be probability measures on the same measurable space (Ω, \mathcal{F}) . A coupling of \mathbb{P} and \mathbb{P}' is a probability measure \mathbf{P} on $(\Omega \times \Omega, \mathcal{F} \times \mathcal{F})$ such that the marginals of \mathbf{P} coincide with \mathbb{P} and \mathbb{P}' . Meaning

$$\mathbf{P}(A \times \Omega) = \mathbb{P}(A) \quad \text{and} \quad \mathbf{P}(\Omega \times A) = \mathbb{P}'(A), \quad \forall A \in \mathcal{F}$$

We say (X, Y) is coupling of \mathbb{P} and \mathbb{P}' for random variables X, Y if the law of (X, Y) is a coupling of \mathbb{P} and \mathbb{P}' as defined above.

We construct a pair of configurations $\omega, \omega' \in \Omega$ using a family of uniform random variables $(U_e)_{e \in E}$ which are i.i.d. on $[0, 1]$. Then fixing a pair $p, p' \in [0, 1]$ such that $p < p'$, for each $e \in E$ we define its state in ω by:

$$\omega_e = \begin{cases} 1 & \text{if } U_e \leq p \\ 0 & \text{otherwise} \end{cases}$$

So each edge in ω is independently open with probability p . Specifically the law of ω is \mathbb{P}_p . Similarly define ω' by:

$$\omega'_e = \begin{cases} 1 & \text{if } U_e \leq p' \\ 0 & \text{otherwise} \end{cases}$$

Then the law of ω'_e is $\mathbb{P}_{p'}$. Then since $p \leq p'$ it is clear that

$$\forall e \in E, \omega_e \leq \omega'_e$$

Which we denote by the shorthand $\omega \leq \omega'$.

Then (ω, ω') is a coupling of \mathbb{P}_p and $\mathbb{P}_{p'}$ with the property that $\mathbf{P}(\omega \leq \omega') = 1$. Let some increasing event $A \subseteq \Omega$, by definition if $\omega \in A$, then $\omega' \in A$ so:

$$\mathbb{P}_p(A) = \mathbf{P}(\omega \in A) \leq \mathbf{P}(\omega' \in A) = \mathbb{P}_{p'}(A)$$

So we conclude $\mathbb{P}_p(A)$ is non-decreasing in p . □

Corollary 2.13. *The percolation function $\theta(p)$ is non-decreasing in p .*

Proof. Since we already saw that the event $\{|C| = \infty\}$ is an increasing event, then by Lemma 2.11 we know that $\theta(p) = \mathbb{P}_p(|C| = \infty)$ is non-decreasing in p . □

An interesting property of our percolation model is that it exhibits a phase transition. Where at some point the behaviour of the model changes notably. For the percolation model we see there is some value of our parameter p , before which we know almost surely there is no infinitely connected cluster, and after which there may be one. We call this value of p the critical value or percolation threshold.

Definition 2.14 (Critical Value). We define the critical value $p_c(d)$ formally as follows

$$p_c(d) = \sup\{p \in [0, 1] : \theta(p) = 0\}$$

where d denotes the dimensionality of our graph \mathbb{Z}^d , we may sometimes drop the d and just refer to p_c when it is clear what the model dimension is.

3 Critical values

3.1 Existence of a critical value on \mathbb{Z}

Trivially the critical value is $p = 1$. Consider the event $X_n = \{\text{There is an open self-avoiding path of length } n \text{ starting at the origin}\}$ Then $X_n \supseteq X_{n+1}$ and so

$$\lim_{n \rightarrow \infty} \mathbb{P}_p(X_n) = \theta(p)$$

And since $\mathbb{P}_p(X_n) = 2p^n$, as the path can go left or right from the origin. We have for all $p < 1$, $\theta(p) = 0$. Thus, $\theta(p) > 0$ if and only if $p = 1$.

3.2 Existence of a critical value on \mathbb{Z}^2

We show the existence of the critical value in this case by bounding it from above and below. We follow the proofs given in [18].

Theorem 3.1. *If $p < 1/3$, $\theta(p) = 0$.*

Proof. Let $X_n = \{\text{There is an open self-avoiding path of length } n \text{ starting at the origin}\}$ as in Section 3.1. Then the probability for a path of length n to be open on every edge is p^n . The number of paths of length n from the origin is at most $4(3^{n-1})$ since there are 4 edges to choose from at the

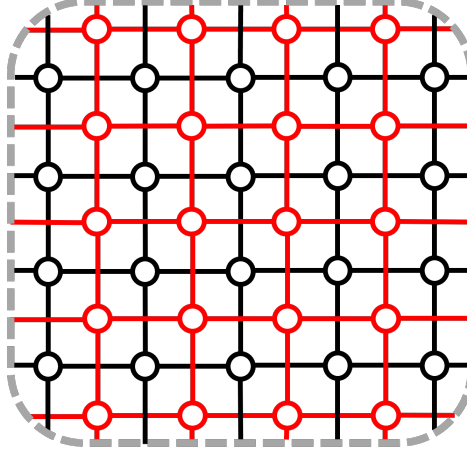


Figure 2: An example snippet of the primal graph \mathbb{Z}^2 shown in black and the dual graph $(\mathbb{Z}^2)^*$ in red.

origin, then for each next step in the path there are at most 3 edges we can pick as the path is self-avoiding. Hence, we have

$$\mathbb{P}_p(X_n) \leq 4(3^{n-1})p^n$$

Then we take the limit since $\lim_{n \rightarrow \infty} \mathbb{P}_p(X_n) = \theta(p)$. So,

$$\begin{aligned} \lim_{n \rightarrow \infty} \mathbb{P}_p(X_n) &\leq \lim_{n \rightarrow \infty} 4(3^{n-1})p^n \\ &\leq 4 \cdot 3^{-1} \lim_{n \rightarrow \infty} (3p)^n \end{aligned}$$

Since $p < 1/3$ clearly the limit on the final line is zero. Thus, we have $\theta(p) = \lim_{n \rightarrow \infty} \mathbb{P}_p(X_n) = 0$. \square

Theorem 3.2. *For p close to 1, we have $\theta(p) > 0$*

Proof. We introduce the dual graph $(\mathbb{Z}^2)^*$ which has vertices in $(\mathbb{Z}^2 + (\frac{1}{2}, \frac{1}{2}))$, and edges as you would expect between vertices at distance 1. Then we can see there is a clear correspondence between the edges of \mathbb{Z}^2 and its dual, since each edge in the dual intersects a unique edge in the original graph, see Figure 2.

Thus, we can create a mapping from the open and closed edges of \mathbb{Z}^2 to the dual graph, where the edge in the dual is open if and only if the intersecting edge in \mathbb{Z}^2 is closed, and vice versa.

Then we notice that if there exists a cycle of open edges in the dual graph enclosing the origin then the size of the open cluster at the origin is finite.

Lemma 3.3. *$|C| < \infty$ if and only if there exists a cycle of open edges in $(\mathbb{Z}^2)^*$ enclosing the origin*

Proof. Thinking visually since there is a ring of open edges in the dual, the open edges from the origin in the original graph cannot extend beyond this ring. A more formal pure graph theory proof can be made but is omitted here. \square

Let $X_n = \{\text{There is a length } n \text{ cycle of open edges in } (\mathbb{Z}^2)^* \text{ which surrounds the origin}\}$ Then using Lemma 3.3 we see

$$\mathbb{P}_p(|C| < \infty) = \mathbb{P}_p\left(\bigcup_{n=4}^{\infty} \mathbb{P}_p(X_n)\right) \leq \sum_{n=4}^{\infty} \mathbb{P}_p(X_n) \leq \sum_{n=4}^{\infty} n \cdot 4(3^{n-1})q^n$$

Where in the last summation we get $4(3^{n-1})q^n$ as before by counting the number of possible edges a self avoiding walk can take then multiplying by the probability this path is open in the dual graph. We multiply this term by n since we know for every cycle around the origin in the dual there is some smallest $i \in \mathbb{N}$ for which the cycle crosses the edge $\{i, i-1\}$, i.e. the positive x-axis. Then we can translate this cycle left and right along the 'x-axis', this gives up to n cycles enclosing the origin in the dual graph.

This sum is finite when $q < 1/3$, which is when $p > 2/3$. We can make the sum arbitrarily small when $p \rightarrow 1$, when the sum is smaller than 1 this implies $\theta(p) > 0$. Taking the value $p = 0.9$ for example, we can calculate that the sum converges to

$$\sum_{n=4}^{\infty} n \cdot 4(3^{n-1})(1 - 0.9)^n = 0.0683265$$

Thus $\theta(0.9) > 0$. \square

Hence, the critical value $p_c(2) \in (\frac{1}{3}, \frac{9}{10})$

3.3 Existence of a critical value on \mathbb{Z}^d

We can easily derive the result that $p_c(d) \in (0, 1)$ from the previous section. We notice that for all $d \geq 2$ we can extract \mathbb{Z}^d from \mathbb{Z}^{d+1} by simply ignoring the extra edges and their states. Thus, if there exists an infinite open cluster in \mathbb{Z}^d this would also imply there is one in \mathbb{Z}^{d+1} . Since we have just shown $p_c(2) \in (0, 1)$ then we must also have $p_c(d) \in (0, 1)$ for all d . In fact by this argument it is clear we have for $d \geq 1$ that,

$$p_c(d) \geq p_c(d+1)$$

Since, if we couple the configurations of each lattice as mentioned by embedding \mathbb{Z}^d into \mathbb{Z}^{d+1} at the origin. Then as the value of p increases the origin lies in an infinite open cluster of \mathbb{Z}^{d+1} at worst at the same value of p for which the origin lies in an infinite open cluster of \mathbb{Z}^d .

4 Harris-Kesten Theorem

An important result in percolation theory is that the critical value for the square lattice is $\frac{1}{2}$. This result came in two parts, firstly Harris [12] showed that $p_c(2) \geq \frac{1}{2}$ in 1960. Then 20 years later Kesten [13] proved $p_c(2) \leq \frac{1}{2}$, these proofs combined resulted in the following theorem.

Theorem 4.1 (Harris-Kesten theorem). *The critical value on \mathbb{Z}^2 , $p_c(2) = \frac{1}{2}$.*

The aim of this section is to construct a proof of this theorem. We largely follow the proofs and methods of [5], which in turn is largely based on the proofs by Russo, Seymour and Welsh. To do this we first introduce some more foundational results.

4.1 Some preliminary results

Recall from earlier the definition of an increasing event.

Definition 4.2 (Increasing Event). An event $A \subseteq \Omega$ is increasing if when $\omega \in A$ and

$$\forall e \in E(\mathbb{Z}^d), \omega_e = 1 \implies \omega'_e = 1$$

Then $\omega' \in A$.

We now introduce an important result, known as Harris' Inequality, which tells us that the probability of increasing events are positively correlated, we follow the proof for Harris' Inequality given by [8].

Theorem 4.3 (Harris' Inequality). *If $A, B \subseteq \Omega$ are increasing events then*

$$\mathbb{P}(A \cap B) \geq \mathbb{P}(A)\mathbb{P}(B)$$

This is a specific case of a more general inequality, the FKG inequality. If f and g are two bounded increasing functions then

$$\mathbb{E}_p(fg) \geq \mathbb{E}_p(f)\mathbb{E}_p(g)$$

Proof. We show that the FKG inequality holds, since we can easily derive Harris' inequality by setting f and g to be indicator functions of the increasing events A and B . We first show the result holds for increasing functions which depend on a finite number of edges, then we see how this restriction can be removed using the Martingale convergence theorem.

Working by induction on the number of edges that the functions f and g depend on, we label the edges in $E = \{e_i : i \geq 1, e_i \in E\}$. We also shorten our notation for the state of a given edge $\omega_{e_i} = \omega_i$ for readability during this proof.

In the base case $n = 1$ we suppose f and g depend solely on ω_1 . Note that we can just show the inequality holds when we assume $f(0) = g(0) = 0$ since additional constants in f and g will not affect the inequality. Since f and g are increasing we know that $f(1) \geq 0$ and $g(1) \geq 0$. So we find by the definition of expectation

$$\mathbb{E}_p(fg) - \mathbb{E}_p(f)\mathbb{E}_p(g) = (p \cdot f(1)g(1) + q \cdot 0) - (p \cdot f(1) + q \cdot 0)(p \cdot g(1) + q \cdot 0)$$

Then simplifying by removing terms multiplied by zero

$$\mathbb{E}_p(fg) - \mathbb{E}_p(f)\mathbb{E}_p(g) = pf(1)g(1) - p^2f(1)g(1) \geq 0$$

Since $p \geq 0$, $f(1) \geq 0$ and $g(1) \geq 0$. So the result holds for $n = 1$.

Now assuming as our inductive hypothesis that the inequality holds for functions dependent on $n = k$ edges. We consider when $n = k + 1$. We fix the state of the first k edges, $\omega_1, \dots, \omega_k$, and condition on them since we aim to use the tower law for expectation later. Then,

$$\mathbb{E}_p(fg|\omega_1, \dots, \omega_k) = pf(\omega_1, \dots, \omega_k, 1)g(\omega_1, \dots, \omega_k, 1) + qf(\omega_1, \dots, \omega_k, 0)g(\omega_1, \dots, \omega_k, 0)$$

Where the functions $f(\omega_1, \dots, \omega_k, 1)$ and $g(\omega_1, \dots, \omega_k, 1)$ now depend only on the edge ω_{k+1} . Then letting $\mathbb{P}_{\omega_{k+1}}$ be the law (distribution) of ω_{k+1} and similarly $\mathbb{E}_{\omega_{k+1}}$. We see that

$$\begin{aligned} \mathbb{E}_p(fg|\omega_1, \dots, \omega_k) &= pf(\omega_1, \dots, \omega_k, 1)g(\omega_1, \dots, \omega_k, 1) \\ &\quad + qf(\omega_1, \dots, \omega_k, 0)g(\omega_1, \dots, \omega_k, 0) \\ &= \mathbb{E}_{\omega_{k+1}}(f(\omega_1, \dots, \omega_k, \cdot)g(\omega_1, \dots, \omega_k, \cdot)) \end{aligned}$$

Then applying our inductive hypothesis.

$$\mathbb{E}_{\omega_{k+1}}(f(\omega_1, \dots, \omega_k, \cdot)g(\omega_1, \dots, \omega_k, \cdot)) \geq \mathbb{E}_{\omega_{k+1}}(f(\omega_1, \dots, \omega_k, \cdot))\mathbb{E}_{\omega_{k+1}}(g(\omega_1, \dots, \omega_k, \cdot))$$

Which means

$$\mathbb{E}_p(fg|\omega_1, \dots, \omega_k) \geq \mathbb{E}_p(f|\omega_1, \dots, \omega_k)\mathbb{E}_p(g|\omega_1, \dots, \omega_k)$$

So applying the tower law of expectation

$$\begin{aligned} \mathbb{E}_p(fg) &= \mathbb{E}_p(\mathbb{E}_p(fg|\omega_1, \dots, \omega_k)) \\ &\geq \mathbb{E}_p(\mathbb{E}_p(f|\omega_1, \dots, \omega_k)\mathbb{E}_p(g|\omega_1, \dots, \omega_k)) \\ &\geq \mathbb{E}_p(\mathbb{E}_p(f|\omega_1, \dots, \omega_k))\mathbb{E}_p(\mathbb{E}_p(g|\omega_1, \dots, \omega_k)) \\ &= \mathbb{E}_p(f)\mathbb{E}_p(g) \end{aligned}$$

So we have shown the result for all increasing functions which depend on a finite number of edges. Now to extend this to arbitrary bounded increasing

functions we consider the martingale convergence theorem. For our bounded increasing functions f and g we define:

$$\begin{aligned} f_n &:= \mathbb{E}_p(f|\omega_1, \dots, \omega_n) \\ g_n &:= \mathbb{E}_p(g|\omega_1, \dots, \omega_n) \end{aligned}$$

We see that (f_n) and (g_n) form martingales due to the tower law of expectation. Then since we assumed f and g are bounded we can apply the martingale convergence theorem to see that f_n converges to f as $n \rightarrow \infty$ and similarly g_n converges to g . We do not give a rigorous proof of this here, but the reader may wish to seek further information on Doob's martingale convergence theorem from [21]. □

We also note a useful rearrangement of Harris' Inequality which captures the idea that the events A and B are positively correlated.

Corollary 4.4. *If $A, B \subseteq \Omega$ are increasing events then*

$$\mathbb{P}_p(A|B) \geq \mathbb{P}_p(A)$$

Intuitively this makes sense, the probability A occurs increases when we know more edges are open, and knowing B has occurred tells us about open edges.

Furthermore, we have an immediate analogous result for decreasing events.

Corollary 4.5. *If $A, B \subseteq \Omega$ are decreasing events then*

$$\mathbb{P}(A \cap B) \geq \mathbb{P}(A)\mathbb{P}(B)$$

This means decreasing events are also positively correlated.

Proof. Simply set f and g to be the negative of the indicator functions of A and B and apply the FKG inequality. □

We now prove an interesting inequality for events that occur disjointly. We do not need to use this inequality in our proof of the Harris-Kesten Theorem, but we include it simply for the readers interest since it was used in the original proof of Kesten's theorem. First we define what we mean for events to occur disjointly.

Definition 4.6 (Disjoint realisation of events). Let A and B be events. For a configuration $\omega \in A$, there is some witness set $I \subset E$ for ω such that for any other configuration ω' which satisfies:

$$\forall e \in I, \omega_e = \omega'_e$$

Then $\omega' \in A$, in essence the states of edges in I tell us the event A has occurred in ω .

We say the events A and B occur disjointly in ω , denoted $A \circ B$, if there exists disjoint witnesses I and J for A and B respectively. I.e. $I \cap J = \emptyset$

This concept is different to the intersection of events, since intuitively the edges on which each event occurs cannot overlap. For example given some $n \times n$ box of \mathbb{Z}^2 , the event there is an open crossing from top to bottom and the event there is an open crossing left to right cannot occur disjointly. Since the open paths will cross at some point. We now state the inequality we aim to prove.

Theorem 4.7 (BK inequality). *If $A, B \subseteq \Omega$ are increasing events then*

$$\mathbb{P}_p(A \circ B) \leq \mathbb{P}_p(A)\mathbb{P}_p(B)$$

Proof. We follow the proof given by [4] which is easier than the original. We work by induction on the size of the edge set E which we denote by n , remembering the sample space $\Omega = \{0, 1\}^E$. When $n = 0$ the inequality holds trivially. Supposing as our inductive hypothesis that the inequality holds when the edge set has size less than n . We show it also holds when E has size n . First we define for any subset $D \subseteq \{0, 1\}^E$:

$$\begin{aligned} D_0 &= \{(\omega_1, \dots, \omega_{n-1}) : (\omega_1, \dots, \omega_{n-1}, 0) \in D\} \\ D_1 &= \{(\omega_1, \dots, \omega_{n-1}) : (\omega_1, \dots, \omega_{n-1}, 1) \in D\} \end{aligned}$$

Then we can decompose the probability of D since the states of edges are independent.

$$\mathbb{P}_p(D) = p\mathbb{P}_p(D_1) + q\mathbb{P}_p(D_0)$$

Now for our increasing $A, B \subseteq \{0, 1\}^E$, define $C = A \circ B$, we now see that

$$C_0 = A_0 \circ B_0$$

$$C_1 = (A_1 \circ B_0) \cup (A_0 \circ B_1)$$

Applying our assumption that A and B are increasing its clear that $A_0 \subseteq A_1$ and $B_0 \subseteq B_1$. Thus we find,

$$C_0 = A_0 \circ B_0 \subseteq (A_1 \circ B_0) \cap (A_0 \circ B_1)$$

$$C_1 = (A_1 \circ B_0) \cup (A_0 \circ B_1) \subseteq A_1 \circ B_1$$

Using this decomposition in conjunction with our inductive hypothesis we find

$$\mathbb{P}_p(C_0) = \mathbb{P}_p(A_0 \circ B_0) \leq \mathbb{P}_p(A_0)\mathbb{P}_p(B_0)$$

$$\mathbb{P}_p(C_1) \leq \mathbb{P}_p(A_1 \circ B_1) = \mathbb{P}_p(A_1)\mathbb{P}_p(B_1)$$

Now we can also get by adding the probabilities of C_0 and C_1 that

$$\begin{aligned} \mathbb{P}_p(C_0) + \mathbb{P}_p(C_1) &\leq \mathbb{P}_p((A_1 \circ B_0) \cap (A_0 \circ B_1)) + \mathbb{P}_p((A_1 \circ B_0) \cup (A_0 \circ B_1)) \\ &= \mathbb{P}_p(A_0 \circ B_1) + \mathbb{P}_p(A_1 \circ B_0) \\ &\leq \mathbb{P}_p(A_0)\mathbb{P}_p(B_1) + \mathbb{P}_p(A_1)\mathbb{P}_p(B_0) \end{aligned}$$

The next step is to combine these by finding

$$(1-p)^2\mathbb{P}_p(C_0) + p^2\mathbb{P}_p(C_1) + p(1-p)(\mathbb{P}_p(C_0) + \mathbb{P}_p(C_1))$$

$$= ((1-p)^2 + p(1-p))\mathbb{P}_p(C_0) + (p^2 + p(1-p))\mathbb{P}_p(C_1) = (1-p)\mathbb{P}_p(C_0) + p\mathbb{P}_p(C_1)$$

Then looking at the right-hand sides

$$(1-p)\mathbb{P}_p(C_0) + p\mathbb{P}_p(C_1) \leq (1-p)^2[\mathbb{P}_p(A_0)\mathbb{P}_p(B_0)]$$

$$+ p^2[\mathbb{P}_p(A_1)\mathbb{P}_p(B_1)]$$

$$+ p(1-p)[\mathbb{P}_p(A_0)\mathbb{P}_p(B_1) + \mathbb{P}_p(A_1)\mathbb{P}_p(B_0)]$$

Which simplifies to

$$(1-p)\mathbb{P}_p(C_0) + p\mathbb{P}_p(C_1) \leq [p\mathbb{P}_p(A_1) + (1-p)\mathbb{P}_p(A_0)][p\mathbb{P}_p(B_1) + (1-p)\mathbb{P}_p(B_0)]$$

Then using $\mathbb{P}_p(D) = p\mathbb{P}_p(D_1) + (1-p)\mathbb{P}_p(D_0)$ we see that

$$\mathbb{P}_p(A \circ B) = \mathbb{P}_p(C) \leq \mathbb{P}_p(A)\mathbb{P}_p(B)$$

□

4.2 Crossing rectangles

We now introduce some notation to allow us to discuss rectangles and their crossings in a succinct manner. Again we follow the definitions and proofs given in [5].

Definition 4.8. Let $R = [a, b] \times [c, d]$ for integers $a < b$, $c < d$ be a rectangle, we identify this rectangle with a subgraph of \mathbb{Z}^2 which includes all vertices and edges on the boundary and inside. Let $k = b - a$ and $l = d - c$, then we say that R is a k by l rectangle.

We now let $H(R)$ denote the event that there is a horizontal open crossing of R , i.e. a path from the left boundary of R to the right only using open edges. Similarly, we define $V(R)$ to be the event that there is a vertical open crossing of R .

We note from this definition that if R and R' are distinct rectangles, which includes their boundaries being distinct. Then the events $H(R)$, $V(R)$ are independent of $H(R')$, $V(R')$.

It is also useful to specify how the subgraph identified with a rectangle is related to the dual graph, $(\mathbb{Z}^2)^*$.

Definition 4.9. The (horizontal) dual of a rectangle $R = [a, b] \times [c, d]$ in \mathbb{Z}^2 is the rectangle $R^h = [a + 1/2, b - 1/2] \times [c - 1/2, d + 1/2]$ in $(\mathbb{Z}^2)^*$. As a reminder an edge in the dual is open if and only if the edges it intersects in the primal lattice is closed. Let $V^*(R^h)$ be the event there is a vertical crossing of R^h in the dual graph by open edges.

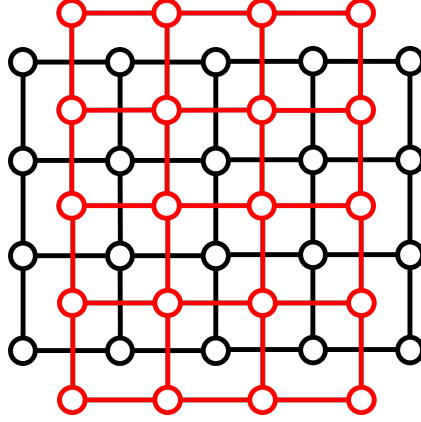


Figure 3: An example 5 by 4 rectangle R shown in black and the dual R^h in red.

We now prove an important result similar to that of Lemma 3.3, for rectangles. Relating the horizontal crossings in primal rectangles to vertical crossings in their duals.

Lemma 4.10. *Let R be a rectangle in \mathbb{Z}^2 . For any configuration ω of \mathbb{Z}^2 , exactly one of the events $H(R)$ and $V^*(R^h)$ holds.*

Proof. Without loss of generality we scale the lattice to avoid working with fractions. Let the lattice $L = (0, 2) + 4\mathbb{Z}^2$, and hence its dual $L^* = (2, 0) + 4\mathbb{Z}^2$. Then let R be a k by l rectangle on L with vertex set,

$$\{(4i, 4j + 2) : 0 \leq i \leq k, 0 \leq j \leq l - 1\}$$

Thus we have that R^h has vertex set

$$\{(4i + 2, 4j) : 0 \leq i \leq k - 1, 0 \leq j \leq l\}$$

Suppose some percolation configuration of \mathbb{Z}^2 and let G be the graph induced by the open edges of R , similarly let the open edges of R^h induce the graph G^h . Now it suffices to show that either G contains a left-right path or G^h contains a top-bottom path, but not both. If we wish to show that G has a left-right path, i.e. a path from $(0, 4m + 2)$ to $(4k, 4n + 2)$ for some $m, n \in \mathbb{N}$, then we can assume without loss of generality that G contains all $2(l - 1)$ edges on its left and right most sides. We make a similar assumption that G^h has all $2(k - 1)$ on its top and bottom. This means we only have to consider one vertex on each of the sides of the rectangles.

We now construct a third graph, M , which lies between G and G^h . Let R_{odd} be a $2k \times 2l$ rectangle with vertex set,

$$\{(2i + 1, 2j + 1) : 0 \leq i \leq 2k - 1, 0 \leq j \leq 2l - 1\}$$

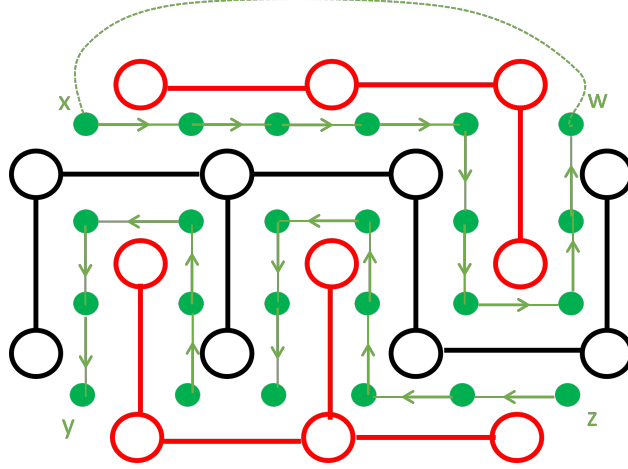


Figure 4: An example with G in black, G^h in red and M in green. The dashed green line shows the Jordan Curve construction.

Then let M be the subgraph of R_{odd} which just contains edges that do not intersect edges of G nor G^h . We direct all the edges of M such that edges of G lie on the right and G^h lie on the left. It is evident that this results in all vertices of M having one edge in and one edge out, except for 4 vertices of degree one which lie in the “corners” of the rectangle, label them counter-clockwise x, y, z, w as in Figure 4.

Clearly due to the way we previously oriented the edges of M , x and z have one out edge, y and w have one in edge. Hence, the component of M which x lies in must be a path, say P , which ends at either y or w . In the case where P ends at y , then we have that the subgraph of G^h to the left of P must be connected from the top to the bottom. Thus, by basic graph theory this subgraph must contain a top-bottom path. If we are in the case where P ends at w then we must have the subgraph of G to the right of P is connected from left to right, thus meaning it admits a left-right path. Clearly then at least one of $H(R)$ and $V^*(R^h)$ occurs.

To show that exactly one $H(R)$ and $V^*(R^h)$ occurs we utilise the Jordan Curve Theorem. A summary of this result is, given a non-self-intersecting continuous closed loop in the plane, J , called a Jordan Curve, then J divides the plane into two distinct regions, an inside and an outside whose boundary is J . We refer the reader to [14] for a proof of the Jordan Curve Theorem.

If the path P runs from x to w , then P can be extended to a Jordan curve J by simply adding a curve from w to x over the top of all the vertices of R^h . This results in all the vertices on the top of R^h lying inside the Jordan curve and the vertices at the bottom of R^h lie on the outside of the curve. Hence, any top-bottom path in G^h must cross J which means it crosses an edge of M , but this contradicts the definition of the edges of M .

So only $H(R)$ occurs. Similarly, in the case that P ends at y , we form a Jordan curve by connecting x to y on the left of the vertices on the left side of G . Then by an analogous argument to before, we cannot have a left-right path in G since the left and right sides lie in the interior and exterior of the Jordan curve respectively. Hence, only $V^*(R^h)$ occurs in this case. \square

We note that the edges of L^* above are open with probability $1 - p$, and so we can form some useful results for rectangles in \mathbb{Z}^2 directly from this lemma.

Corollary 4.11. *If R is a k by $l - 1$ rectangle in \mathbb{Z}^2 and R' is a $k - 1$ by l rectangle. Then,*

$$\mathbb{P}_p(H(R)) + \mathbb{P}_{1-p}(V(R')) = 1$$

Further to this, if $k = l = n + 1$ for some n , then R' is simply a 90-degree rotation of R . Hence, $\mathbb{P}_p(H(R)) = \mathbb{P}_{1-p}(V(R'))$. From this we find the key statement of self-duality of \mathbb{Z}^2 we discussed in Section 3.2.

Corollary 4.12. *If R is an $n+1$ by n rectangle then $\mathbb{P}_{1/2}(H(R)) = 1/2$. So for n sided square S , which is just defined as an n by n rectangle. We must have,*

$$\mathbb{P}_{1/2}(H(S)) = \mathbb{P}_{1/2}(V(S)) \geq 1/2$$

While it might now feel intuitively obvious that $p_c(2) \geq 1/2$, and certainly Corollary 4.12 is essential in the reasoning for this, we are yet to prove it rigorously.

4.3 Harris' Theorem

In order to make this idea rigorous, we now aim to show that $\theta(1/2) = 0$. This is Harris' Theorem, although the method we use will be different to Harris' original as previously mentioned, we follow the ideas of Russo, Seymour and Welsh as outlined in [5].

Lemma 4.13. *Let $R = [0, m] \times [0, 2n]$ be an m by $2n$ rectangle with $m \geq n$. Let $X(R)$ be the event that there are paths P_1 and P_2 of open edges, where P_1 crosses the square $S = [0, n] \times [0, n]$ from top to bottom, and P_2 lies inside R and joins some vertex on P_1 to some vertex on the right side of R . Then,*

$$\mathbb{P}_p(X(R)) \geq \frac{1}{2} \mathbb{P}_p(H(R)) \mathbb{P}_p(V(S))$$

Proof. We first suppose that $V(S)$ holds, so there is a path P of open edges crossing the square top to bottom. This path separates S into two regions, by the Jordan Curve Theorem. In fact, we choose the left most path P , call it $LV(S)$, and we note that the event $\{LV(S) = P_1\}$ does not depend on any edges to the right of P_1 . We can do this due to our proof of Lemma 4.10, if

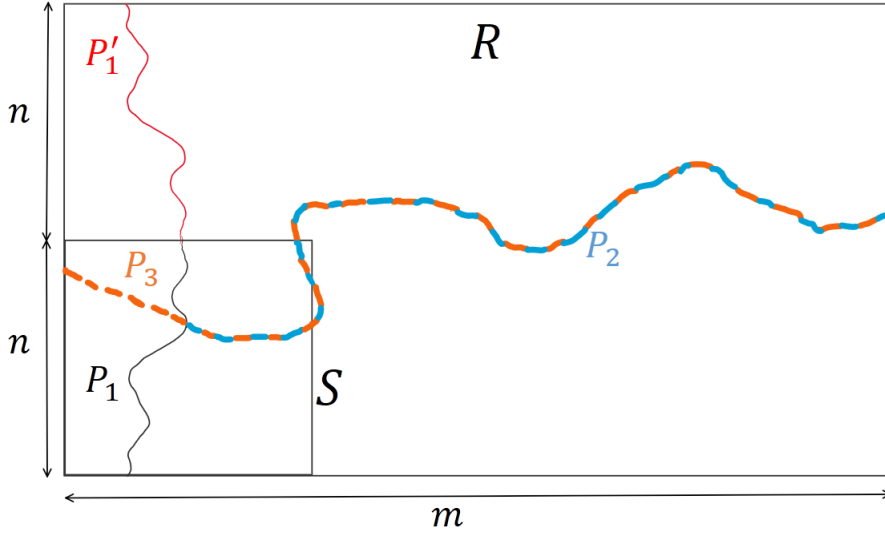


Figure 5: An example m by $2n$ rectangle R , where P_3 crosses P_1 and thus the event $X(R)$ has occurred.

$V(S)$ holds we have a path from x to y in M , and $LV(S)$ is directly adjacent on the left of the path (when viewed from the edge direction) from x to y , in doing so we only needed to look at the edges to the left of $LV(S)$.

Now let P be the path constructed by joining P_1 to its reflection P'_1 in the line $y = n$, note P may not necessarily be open. Then P crosses the left $[0, n] \times [0, 2n]$ section of R top to bottom. By definition, we know with probability $\mathbb{P}_p(H(R))$ there is an open path P_3 crossing R left to right. We know this path P_3 must cross P at some vertex. Due to symmetry by which we constructed P , we know the probability that P_3 crosses P at a vertex in P_1 (the bottom half of P) is at least $\mathbb{P}_p(H(R))/2$. Now let $Y(P_1)$ be the event there is an open path P_2 from the right-hand side of R to some vertex of P_1 , then clearly $Y(P_1)$ must occur with probability at least $\mathbb{P}_p(H(R))/2$ too. We note that the event $Y(P_1)$ only depends on the edges to the right of P_1 . So since the event $\{LV(S) = P_1\}$ does not depend on any edges to the right of P_1 as previously mentioned, $Y(P_1)$ and $\{LV(S) = P_1\}$ are independent. Hence, we have:

$$\mathbb{P}_p(Y(P_1) \mid LV(S) = P_1) = \mathbb{P}_p(Y(P_1)) \geq \mathbb{P}_p(H(R))/2$$

Furthermore we can see that $Y(P_1)$ and $\{LV(S) = P_1\}$ occurring implies that $X(R)$ has occurred. So, it is clear that:

$$\mathbb{P}_p(X(R) \mid LV(S) = P_1) \geq \mathbb{P}_p(H(R))/2$$

Finally we get that the event there is a vertical crossing of S , denoted $V(S)$, is the disjoint union of the events $LV(S) = P_1$ for all P_1 in S . Subsequently,

we get that:

$$\mathbb{P}_p(X(R) | V(S)) \geq \mathbb{P}_p(H(R))/2$$

Then by simple rearrangement we get the required result. \square

We now prove an important corollary of the above lemma which bounds below the probability of a horizontal crossing in specific shapes of rectangles. We follow the method used in [5] by Bollobás and Riordan which differs from the original methods of Russo, Seymour and Welsh, which can be found in [11].

Corollary 4.14. *Let $\rho > 1$ be a fixed integer. Then there is a constant $c_2(\rho) > 0$ which is only dependent on ρ such that, for any $2\rho n$ by $2n$ rectangle R we have $\mathbb{P}_{1/2}(H(R)) \geq c_2(\rho)$*

Proof. Let $h_{m,n} = \mathbb{P}_{1/2}(H(R_{m,n}))$ where $R_{m,n}$ is an m by n rectangle. We aim to show that for $m > n$ we have

$$h_{2m-n,2n} \geq \frac{h_{m,2n}^2}{2^5} \quad (\star)$$

Since then we can iteratively apply the inequality using $h_{2n,2n} \geq 1/2$ from Corollary 4.12 to begin, repeating until we reach $h_{2\rho n,2n}$. For example the first step would be:

$$h_{3n,2n} = h_{2(2n)-n,2n} \geq \frac{h_{2n,2n}^2}{2^5} \geq \frac{1/2}{2^5} = 2^{-6}$$

So in order to show (\star) holds we consider intersecting rectangles

$$R = [0, m] \times [0, 2n]$$

and,

$$R' = [n - m, n] \times [0, 2n]$$

In addition to the lower square $S = [0, n] \times [0, n]$ inside both. Let $E_1 = X(R)$ then let E_2 be the corresponding event reflected in the line $y = n/2$ for R' . Let $E_3 = H(S)$ be the event there is a horizontal crossing of S . Then if E_1, E_2, E_3 all hold we have that $H(R \cup R')$ holds, by the fact that the horizontal crossing of S must cross the two vertical crossing of S which occur due to E_1, E_2 . This is shown in Figure 6. Then we note that E_1, E_2, E_3 are all increasing events, so we can use Harris' Inequality which gives us:

$$\begin{aligned} \mathbb{P}_{1/2}(H(R \cup R')) &\geq \mathbb{P}_{1/2}(E_1 \cap E_2 \cap E_3) \\ &\geq \mathbb{P}_{1/2}(E_1) \mathbb{P}_{1/2}(E_2) \mathbb{P}_{1/2}(E_3) \\ &\geq \mathbb{P}_{1/2}(X(R))^2 \mathbb{P}_{1/2}(H(S)) \end{aligned}$$

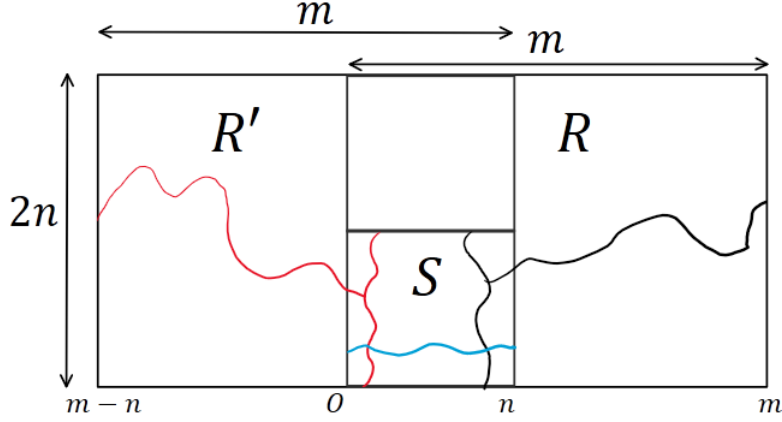


Figure 6: An example of an overlapping R and R' , with the edges of the event E_1 in black, E_2 in red and E_3 in blue.

Where the last line is a result of the symmetry of E_1 and E_2 . From Corollary 4.12 we know $\mathbb{P}_{1/2}(V(S)) = \mathbb{P}_{1/2}(H(S)) \geq 1/2$, and using the result Lemma 4.13 we can plug into the above to get.

$$\begin{aligned}
\mathbb{P}_{1/2}(H(R \cup R')) &\geq \mathbb{P}_{1/2}(X(R))^2 \mathbb{P}_{1/2}(H(S)) \\
&\geq \{2^{-1} \mathbb{P}_{1/2}(H(R)) \mathbb{P}_{1/2}(V(S))\}^2 * 1/2 \\
&\geq (1/2)^2 \mathbb{P}_{1/2}(H(R))^2 (1/2)^2 * 1/2 \\
&= \mathbb{P}_{1/2}(H(R))^2 / 2^5
\end{aligned}$$

This proves (\star) and hence the Corollary holds. \square

We now need one more result before we can prove Harris' Theorem, the second Borel-Cantelli Lemma which we give without proof.

Lemma 4.15 (Second Borel-Cantelli Lemma). *Let $(E_n)_{n=1}^\infty$ be an infinite sequence of events. If $(E_n)_{n=1}^\infty$ are all independent and $\sum_{n=1}^\infty \mathbb{P}(E_n) = \infty$ Then $\mathbb{P}(\text{infinitely many } E_n \text{ occur}) = 1$.*

For a proof we suggest the reader looks at one given in [3]. We are now in a position to prove Harris' theorem.

Theorem 4.16 (Harris' theorem). *For Bernoulli bond percolation on \mathbb{Z}^2 we have that $\theta(1/2) = 0$.*

Proof. We work by constructing a sequence of square annuli in the dual graph. The first annulus is formed by arranging four $6n$ by $2n$ rectangles around the origin such that the $2n$ by $2n$ ends of each rectangle overlap. See Figure 7 for the arrangement, we note that the chosen direction to offset the annulus in the dual does not matter so long as it encloses the origin of the

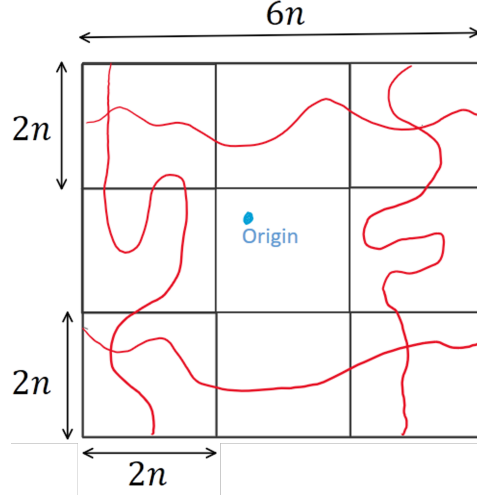


Figure 7: An example of the first square annulus in the dual graph with the four rectangle crossings shown in red which enclose the origin.

primal graph. Then by Corollary 4.14 each rectangle has an open crossing in the long direction with probability at least $c_2(3)$, and thus all of them do with probability at least $c_2(3)^4$. When all four rectangles have an open crossing they form an open path in the dual which encloses the origin. Let A_k be the square annulus in the dual graph formed by overlapping four $(2k+4)n$ by $2n$ rectangles as above. Then let E_k be the event that the annulus contains an open cycle enclosing the origin. Then $\mathbb{P}_{1/2}(E_k) \geq c_2(k+2)$ by Corollary 4.14. Now since the annuli (A_k) lie on disjoint regions, we have that the (E_k) are all independent. Since $\sum_{k=1}^{\infty} \mathbb{P}_{1/2}(E_k) \geq \sum_{k=1}^{\infty} c_2(k+2)^4 = \infty$ we can apply the second Borel-Cantelli Lemma.

$$\mathbb{P}_{1/2}(\text{infinitely many } E_k \text{ occur}) = 1$$

Take the first E_k which occurs, then the annulus A_k in the dual graph contains an open cycle enclosing the origin. Thus, there cannot be an infinite cluster containing the origin in the primal graph. \square

4.4 Kesten's Theorem

We now move onto proving Kesten's theorem, specifically that the critical value for the square lattice is less than or equal to a half. In this section we again follow the definitions and methods detailed in [5]. Firstly we introduce a powerful result due to Friedgut and Kalai [9]. We assume the reader has some basic group theory knowledge, although it is not critical to understanding.

Definition 4.17 (Symmetric event). An event $\mathcal{A} \subseteq \mathcal{P}(X)$ is symmetric if

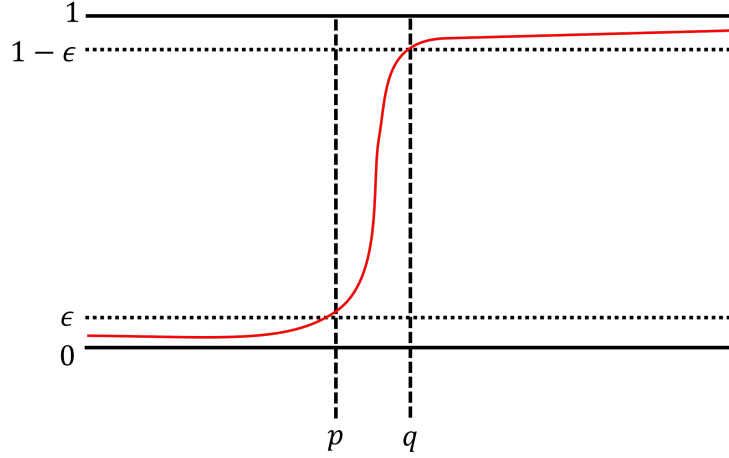


Figure 8: An example of a sharp threshold, plotted in red is the probability of some symmetric increasing event as the percolation parameter increases.

there is a permutation group G acting transitively on $\mathcal{P}(X)$ which preserves \mathcal{A}

To explain briefly, the action of a group G on a set X is transitive if for all $x, y \in X$, there exists $g \in G$ such that $g \cdot x = y$. This means G shuffles the elements of X , in essence given some $x \in X$ you can obtain any $y \in X$ by some $g \in G$. Then an event \mathcal{A} being preserved under G just means that if $A \in \mathcal{A}$ then there exists $g \in G$ such that $g \cdot A \in \mathcal{A}$.

Theorem 4.18 (Friedgut-Kalai sharp threshold theorem). *There is an absolute constant c_1 such that if $\mathcal{A} \subset \mathcal{P}(X)$ is a symmetric increasing event and if $\mathbb{P}_p(\mathcal{A}) > \epsilon$ for some $0 < \epsilon < 1/2$ then,*

$$\mathbb{P}_q(\mathcal{A}) > 1 - \epsilon$$

whenever we have

$$q - p \geq c_1 \frac{\log(1/(2\epsilon))}{\log|X|}$$

We do not give the proof since it is too long to include within the scope of this project, but the reader is recommended to view [9] for details. This theorem tells us all symmetric increasing events have 'sharp threshold', meaning a small increase of order $1/\log|X|$ to the parameter p is enough to take the probability of the event from a value close 0 to nearly 1, see Figure 8.

Using the Friedgut-Kalai sharp threshold theorem we will now work to prove that when $p > 1/2$ the probability of crossing rectangles larger than a certain size is close to 1.

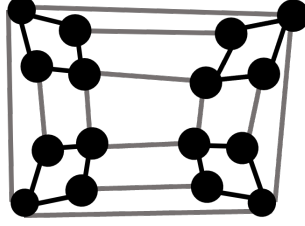


Figure 9: \mathbb{T}_4 , the 4 by 4 discrete torus.

Lemma 4.19. *Let $p > 1/2$ and fix some integer $\rho > 1$. Then there exists constants $\delta = \delta(p) > 0$ and $n_0 = n_0(p, \rho)$ such that when $n > n_0$ and R_n is a $4\rho n$ by $4n$ rectangle on the square lattice then,*

$$\mathbb{P}_p(H(R_n)) \geq 1 - n^{-\delta}$$

Proof. For the sake of simplicity, we prove the case of $\rho = 3$ since this implies the result for $\rho = 2$ and these are the only cases we will later need. The other cases can be proved by an analogous argument. Let $p > 1/2$, then by Corollary 4.14 with $\rho = 7$ we know there exists a constant $0 < c_2 < 1/2$ such that for any $14n$ by $2n$ rectangle R , the probability of a horizontal crossing is bounded below by c_2 , $\mathbb{P}_{1/2}(H(R)) \geq c_2$. Now we cannot directly apply the Friedgut-Kalai sharp threshold theorem since $H(R)$ is not a symmetric event, since X in this case is the set of all edges in \mathbb{Z}^2 . In order to get around this we introduce new symmetric events on the discrete torus.

Definition 4.20 (Discrete Torus). For $n \geq 3$, the n by n discrete torus, \mathbb{T}_n , is the Cartesian product of two length n cycles. Equivalently \mathbb{T}_n can be formed by identifying all pairs of vertices on \mathbb{Z}^2 whose coordinates are congruent modulo n . We note that \mathbb{T}_n has n^2 vertices and $2n^2$ edges.

Now for $1 \leq k, l \leq n - 2$, we define a k by l rectangle R in \mathbb{T}_n to be an induced subgraph in \mathbb{T}_n which has an associated k by l rectangle $R' = [a, a + k] \times [b, b + l]$ in \mathbb{Z}^2 , for some $a, b \in \mathbb{Z}$. Due to the restriction in size of k and l , our rectangles in the torus never wrap around on themselves and thus always have an associated planar rectangle in \mathbb{Z}^2 .

We define Bernoulli bond percolation on \mathbb{T}_n as you would expect, where each edge is open independently with probability p . We will write $\mathbb{P}_p^{\mathbb{T}_n} = \mathbb{P}_p^{\mathbb{T}}$ for the corresponding probability measure, dropping the n when it is clear from the context for readability.

Then define E_n to be the event that \mathbb{T}_{16n} admits a $14n$ by $2n$ rectangle with an open horizontal crossing, or it admits a $2n$ by $14n$ rectangle with an open vertical crossing. Now E_n is a symmetric event on $\mathcal{P}(X)$ where X is the edge set of the $16n$ by $16n$ torus, we note $|X| = 512n^2$. Suppose a fixed $14n$ by $2n$ rectangle R in \mathbb{T}_{16n} and its corresponding rectangle R' in

\mathbb{Z}^2 , we know

$$\mathbb{P}_{1/2}^{\mathbb{T}}(E_n) \geq \mathbb{P}_{1/2}^{\mathbb{T}}(H(R)) = \mathbb{P}_{1/2}(H(R')) \geq c_2$$

Now we aim to apply Theorem 4.18, firstly setting

$$\delta = \delta(p) = \frac{p - 1/2}{64c_1}$$

$$\epsilon = n^{-128\delta}$$

where c_1 is the absolute constant from the sharp threshold theorem. Then since δ is only dependent on p , clearly there exists an $n_0 = n_0(p)$ such that $\epsilon < c_2 < 1/2$ for all $n \geq n_0$. So by some rearrangement,

$$p - 1/2 = 64c_1\delta = c_1 \frac{-(-128\delta)\log(n)}{2\log(n)} = c_1 \frac{\log(1/\epsilon)}{\log(n^2)} > c_1 \frac{\log(1/(2\epsilon))}{\log(512n^2)}$$

Now since $\mathbb{P}_{1/2}^{\mathbb{T}}(E_n) \geq c_2 > \epsilon$, by Theorem 4.18, we obtain the bound

$$\mathbb{P}_p^{\mathbb{T}}(E_n) \geq 1 - \epsilon = 1 - n^{-128\delta}$$

for all $n > n_0$

Now if we consider covering \mathbb{T}_{16n} with $64 \cdot 12n$ by $4n$ rectangles labelled R_1, \dots, R_{64} , positioning a corner of each R_i on a multiple of $2n$. Then any $14n$ by $2n$ rectangle R on \mathbb{T}_{16n} crosses one of the R_i the long way and the width of R is inside the R_i , see Figure 10. With an equivalent argument, we can place $4n$ by $12n$ rectangles R_{65}, \dots, R_{128} to cover \mathbb{T}_{16n} such that any $2n$ by $14n$ rectangle in \mathbb{T}_{16n} crosses one the long way. Let $E_{n,i}$, $i = 1, \dots, 128$ be the event that rectangle R_i has an open crossing the long way. Clearly if E_n occurs then so does one of the $E_{n,i}$, as illustrated by the purple path in Figure 10. Hence, E_n^c , the complement of E_n , must contain the intersection of all $E_{n,i}^c$.

Then for each i we can use the corollary of Harris' Lemma (Corollary 4.5) to get that the decreasing events $E_{n,i}^c$ and $\cap_{j < i} E_{n,j}^c$ are positively correlated. Thus we see,

$$\mathbb{P}_p^{\mathbb{T}}(E_n^c) \geq \mathbb{P}_p^{\mathbb{T}}\left(\bigcap_{i=1}^{128} E_{n,i}^c\right) \geq \prod_{i=1}^{128} \mathbb{P}_p^{\mathbb{T}}(E_{n,i}^c) = \mathbb{P}_p^{\mathbb{T}}(E_{n,1}^c)^{128}$$

Where we have repeatedly applied Corollary 4.5 to get the second inequality. Now combining this with our earlier result that $\mathbb{P}_p^{\mathbb{T}}(E_n) \geq 1 - n^{-128\delta}$, we find that for $n > n_0$:

$$\mathbb{P}_p^{\mathbb{T}}(E_{n,1}^c) \leq \mathbb{P}_p^{\mathbb{T}}(E_n^c)^{1/128} \leq n^{-\delta}$$

Hence,

$$\mathbb{P}_p^{\mathbb{T}}(E_{n,1}) \geq 1 - n^{-\delta}$$

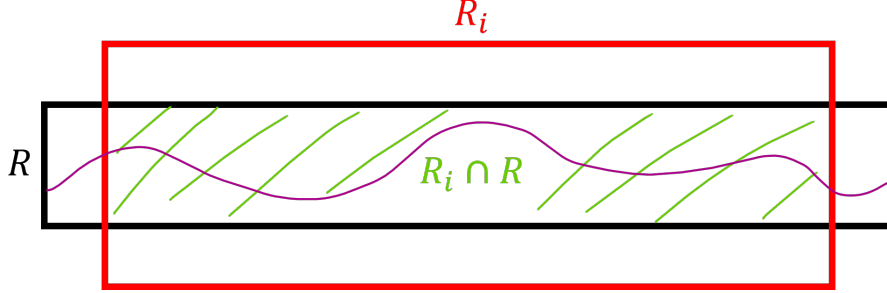


Figure 10: An example $14n$ by $2n$ rectangle in black, with overlapping $12n$ by $4n$ rectangle R_i in red. The $12n$ by $2n$ rectangle induced by their intersection is shaded in green. The purple open horizontal crossing in R clearly implies an open horizontal crossing in R_i .

Since $E_{n,1}$ is the event there is an open horizontal crossing of a $12n$ by $4n$ rectangle, say R , in \mathbb{T}_{16n} . Then by our previous reasoning R has a corresponding $12n$ by $4n$ rectangle R' in \mathbb{Z}^2 . Hence, we get our final result for $n > n_0$ that:

$$\mathbb{P}_p(H(R')) = \mathbb{P}_p^{\mathbb{T}}(H(R)) \geq 1 - n^{-\delta}$$

□

Using this lemma we are in a position to prove Kesten's Theorem, for the proof we follow the method outlined in [4].

Theorem 4.21 (Kesten's Theorem). *For Bernoulli bond percolation on \mathbb{Z}^2 , if $p > 1/2$ then $\theta(p) > 0$*

Proof. Let $p > 1/2$, and define $\delta = \delta(p)$ and $n_0 = n_0(p, 2)$ as in Lemma 4.19. Initially suppose some integer n such that $n > n_0$, we will put another bound on n later. Define R_k for $k = 0, 1, 2, \ell$ to be the rectangle with side lengths $2^k n$ and $2^{k+1} n$, with the longer side vertical if k is even and horizontal otherwise. We position the bottom left corner of all R_k rectangles at the origin so that R_i encloses R_{i-1} , see Figure 11.

Let E_k be the event that R_k admits an open crossing of edges the long way. Clearly if E_k and E_{k+1} hold then the open crossings of R_k and R_{k+1} must meet, as depicted in Figure 11. Then using Lemma 4.19 and the formula for the sum of an infinite geometric series we see that,

$$\sum_{k=0}^{\infty} \mathbb{P}_p(E_k^c) \leq \sum_{k=0}^{\infty} (2^k n)^{-\delta} = \frac{n^{-\delta}}{1 - 2^{-\delta}}$$

Clearly we impose a bound on n to be large enough such that,

$$\sum_{k=0}^{\infty} \mathbb{P}_p(E_k^c) \leq \frac{n^{-\delta}}{1 - 2^{-\delta}} < 1$$

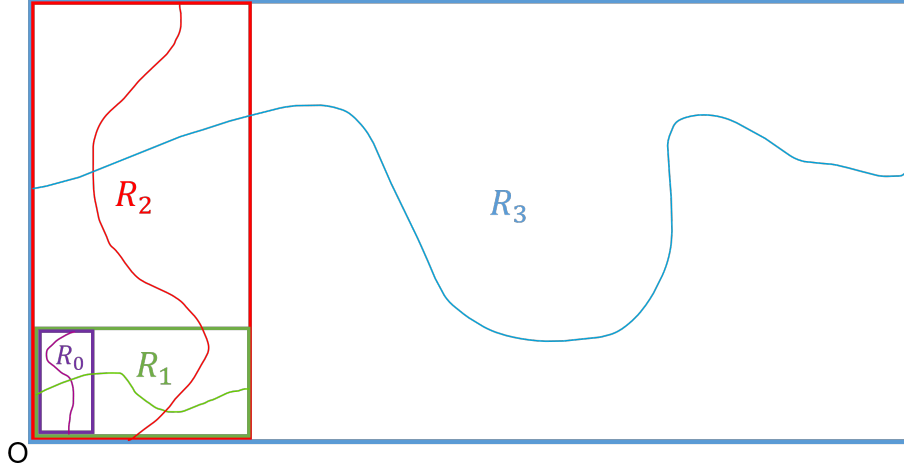


Figure 11: Rectangles R_k for $k = 0, 1, 2, 3$ with their open crossings corresponding to all E_k holding. R_0 in purple, R_1 in green, R_2 in red and R_3 in blue.

Which implies that

$$\mathbb{P}_p \left(\bigcap_{k=0}^{\infty} E_k \right) > 0$$

Now if E_k holds for all k then the event, $E_{\infty} = |C(x)| = \infty : x \in \mathbb{Z}^2$, that there is an infinite open cluster must also hold. Thus,

$$\mathbb{P}_p(E_{\infty}) \geq \mathbb{P}_p \left(\bigcap_{k=0}^{\infty} E_k \right) > 0$$

Attentive readers may note that while E_{∞} implies there is an infinite open path, this path may not necessarily start at the origin. However clearly the probability of an open path from the origin along the bottom edge of R_0 is positive, thus when $p > 1/2$ we have $\theta(p) > 0$.

□

It may interest the reader to know that the final step in the proof is not strictly necessary, the event that there is an infinite open cluster containing the origin is translation invariant. This means that probability of an infinite cluster is not affected by a shift in edge configurations. We not do prove this property here, but the reader may wish to see [8] for details.

Furthermore, the event that there is an infinite open cluster in the lattice is an example of a tail event, since clearly it is independent of the configuration of any finite number of edges. Thus, Kolmogorov's 0-1 law can be applied to this event. We take the definition from [4].

Theorem 4.22 (Kolmogorov's 0-1 law). *Let $X = (X_1, X_2, \dots)$ be a sequence of independent random variables, and let A be an event generated by the σ -algebra generated by X . Suppose that for all n , the event A is independent of X_1, \dots, X_n . Then,*

$$\mathbb{P}(A) \in \{0, 1\}.$$

This theorem immediately tells us when the square lattice admits an infinite open cluster.

Corollary 4.23. *By Theorem 4.22, we have*

$$\mathbb{P}_p(E_\infty) = \begin{cases} 0 & \text{if } p \leq 1/2 \\ 1 & \text{if } p > 1/2 \end{cases}$$

Finally, since we have shown Harris' Theorem that for $p = 1/2$ we have that $\theta(p) = 0$ in conjunction with Kesten's Theorem that for $p > 1/2$ we have that $\theta(p) > 0$. It must be that $p_c(2) = 1/2$ by the definition of the critical value, which completes our proof of the Harris-Kesten Theorem.

5 Simulation

In this section we discuss techniques and algorithms which can be used to explore percolation via simulations. We aim to outline a general method to investigate the critical crossing probability of finite boxes in such a way that it can be applied to any graph with minimal modification. We use the analytical proof for the critical value on the square lattice presented in Section 4 to verify correctness of our method.

5.1 Percolation simulation design

As we have seen previously the task of determining the percolation threshold for a given lattice structure is highly non-trivial and so for many lattices the task of finding exact values is still an open problem [19]. Subsequently, a great deal of work has been carried out to determine approximate values for a variety of lattices and distributions of open sites or bonds by way of computer simulations. In this section we explore an algorithm I used to estimate the bond percolation threshold for boxes of the square lattice, since we can compare the results to the exact value of $1/2$ we proved earlier. The algorithm described makes no assumptions specific to bond percolation on the square lattice and thus will have applicability to other percolation models.

While we know that a configuration having an open crossing from one side to another of a single finite box in \mathbb{Z}^2 does not affect the probability of there being an infinite open cluster since it is a tail event. It is clear that we would expect the probability of such a crossing in a very large box to

tend towards the percolation threshold for the graph. Hence, we can exploit this behaviour to find an estimate for the percolation threshold by creating some large square grid, assign some configuration of open and closed edges with probability p , then check for a left-right crossing. We alter the value of p repeatedly until we find the lowest value at which a crossing occurs.

There are a number of factors to consider when setting up this simulation such that it can run sufficiently fast, firstly how do we assign each edge to be open or closed so that we can traverse the graph with different values of p .

The simplest approach would be to assign a Bernoulli random variable with probability p to each edge, however this has the drawback of requiring us to reevaluate the random variable on each edge when we change the value of p , and to mirror the non-decreasing nature of the percolation function $\theta(p)$, we would have to only do this for previously closed edges. To avoid these issues I chose to implement an idea similar to one we used to prove Lemma 2.11, where we coupled configurations using uniform random variables. To recap this, by assigning an edge with a value from independent uniform random on $[0, 1]$ we can then set an edge to be open if this value is less than or equal to p and closed otherwise. In this way every edge is open with probability p , and increasing p has the desired property that edges which were previously open remain so. Additionally, after the values between 0 and 1 have been assigned to each edge, we gain the benefit of only needing to query the state of edges as we come to them by comparing the assigned value to p , rather than needing to assign the open or closed states to every edge each time we alter p .

We could remove the initial pass of assigning the uniform random variable values to each edge initially. Instead, we could define some hash function mapping the existing attributes of each edge to the range $[0, 1]$, allowing us to compute the states of edges on the fly. However, I did not find this speed up necessary.

Secondly how do we check quickly for a given box if there is a left-right crossing. The first step in simplifying this problem to add two auxiliary vertices to the graph, one connected to every vertex on the left-hand side of the box, the other to the right-hand side vertices. Then our problem reduces to whether there is a path between these two auxiliary vertices via open edges. The fastest way to go about this, since we have an unweighted graph, is to use a version of breadth first search. Taking the subgraph induced by the open edges, note we always set edges incident to the auxiliary edges to be open, we follow a simple algorithm.

The algorithm uses a queue, which is a simple data structure that works in the same way as a queue at a shop. The first item added to the queue is the first to then be removed. Algorithm 1 runs in $O(|V| + |E|)$ worst-case time, where V and E are the vertex and edge sets of the graph G respectively. With a little thought it is clear no algorithm can do better for

Algorithm 1 Breadth first search from vertex u to v

```
1: procedure BFSPATH( $G, u, v$ )
2:   Let  $Q$  be a queue
3:   Let  $explored$  be a list of booleans
4:   Let  $parent$  be a list of nodes
5:    $explored[u] \leftarrow \text{True}$ 
6:    $parent[u] \leftarrow \text{None}$ 
7:    $Q.enqueue(u)$ 
8:   while  $Q$  is not empty do
9:      $x \leftarrow Q.dequeue()$ 
10:    if  $x$  is  $v$  then ▷ Retrace parents to find path
11:       $path \leftarrow [x]$ 
12:      while  $parent[x]$  is not None do
13:         $path.prepend(parent[x])$ 
14:         $x \leftarrow parent[x]$ 
15:      end while
16:      return  $path$ 
17:    end if
18:    for every edge  $w$  in  $G.adjacent(x)$  do
19:      if  $explored[w]$  is False then
20:         $explored[w] \leftarrow \text{True}$ 
21:         $parent[w] \leftarrow x$ 
22:         $Q.enqueue(w)$ 
23:      end if
24:    end for
25:  end while
26:  if  $explored[v]$  is False then return None
27:  end if
28: end procedure
```

this problem.

We also note that Algorithm 1 returns the list of nodes on the path between the two vertices. Depending on the use case this can be easily modified to just return true if a path exists.

Now we can determine if there is a path from one side of a box to another for a given percolation parameter p , we want to determine the lowest threshold for which a crossing path exists for a given configuration of uniform random variables. For the case of the planar square lattice we could use some heuristic based on our prior knowledge that the threshold should lie around $1/2$. However, to enable the general applicability of the algorithm described to any graph, we do not make any specific assumptions about the threshold distribution.

While a variety of possible approaches exist, a rather efficient one is an adapted form of binary search on the parameter. While usually a binary search can be used to identity sorted items in a list, in some sense our parameter values are also sorted. All the values of p which admit a left right crossing appear after those which do have a crossing, this is only true due to our coupling implementation with uniform random variables assigned to each edge. In order to check if a given value of p is the crossing threshold, we must check a value very close to p to see if it is different. Say we have some precision constant ϵ then we observe:

- If p admits a box crossing, then p is the crossing threshold if and only if $p - \epsilon$ does not.
- If p does not admit a box crossing, then $p + \epsilon$ is the crossing threshold if and only if $p + \epsilon$ does admit a crossing.

Now we can follow an adapted binary search to centre in on the threshold, where we let $\text{isPath}(G, u, v)$ be the aforementioned change to Algorithm 1 so that it returns a boolean value.

You can inspect Algorithm 2 to see how whenever the current value of p is not the threshold we can safely throw out one half of the current search range. This technique, commonly referred to as divide and conquer, gives good asymptotic search performance. We would say in general a binary search finds the target in $O(n \log(n))$ worst-case time where n is the number of elements to be searched. In our case the effective number of “items” depends on the precision parameter ϵ , where the number of potential threshold values to check is $\frac{1}{\epsilon}$. Thus, Algorithm 2 runs in $O(\frac{1}{\epsilon} \log(\frac{1}{\epsilon}))$ worst-case time.

You may notice a possible optimization when both the algorithms are used together. Instead of first constructing the induced subgraph by the open edges then seeing if there is a crossing, we can directly only consider the open edges during our breadth first search by requiring in line 18 of Algorithm 1 that the edges w must have associated uniform random variable

Algorithm 2 Find the crossing threshold for a given graph

```

1: procedure FINDTHRESHOLD( $G, u, v, \epsilon$ )  $\triangleright$   $u$  &  $v$  are the auxiliary nodes
2:    $lo \leftarrow 0$ 
3:    $hi \leftarrow 1$ 
4:   while  $lo \leq hi$  do
5:      $p \leftarrow (lo + hi)/2$ 
6:      $openEdges \leftarrow \{e \in E(G): e \text{ has uniform r.v. value } \leq p\}$ 
7:     let  $OpenG$  be the subgraph of  $G$  induced by  $openEdges$ 
8:     if  $isPath(openG, u, v)$  then
9:        $altOpenEdges \leftarrow \{e \in E(G): unifRV(e) \leq p - \epsilon\}$ 
10:      let  $AltG$  be the subgraph of  $G$  induced by  $altOpenEdges$ 
11:      if  $isPath(AltG, u, v)$  then
12:         $hi \leftarrow p - \epsilon$ 
13:      else
14:        return  $p$ 
15:      end if
16:    else
17:       $altOpenEdges \leftarrow \{e \in E(G): unifRV(e) \leq p + \epsilon\}$ 
18:      let  $AltG$  be the subgraph of  $G$  induced by  $altOpenEdges$ 
19:      if  $isPath(AltG, u, v)$  then
20:        return  $p + \epsilon$ 
21:      else
22:         $lo \leftarrow p + \epsilon$ 
23:      end if
24:    end if
25:  end while
26: end procedure

```

value of less than or equal to p . In essence, we can construct the subgraph induced by the open edges on the fly, saving us one pass over the edge set.

5.2 Simulation Implementation

My initial plan to investigate the percolation threshold for the planar square lattice was to simulate many boxes each of increasing width to see how the variation in the crossing threshold decreases with box width. Using the previously described algorithms to find the threshold for each box.

I implemented the simulation using Python due to its extensive set of community libraries and ease of prototyping. Since Python is an interpreted language it can be less performant than other compiled languages which can be an important consideration for very large simulations. However, alternative interpreters of Python such as PyPy can mitigate this issue [20]. PyPy uses just-in-time compilation and other optimization techniques to speed up native python code. I decided to simulate boxes of edge widths from width 2 (3x3 vertices) to width 1950, with a precision of 0.0001. I then used the matplotlib library to draw a scatter plot of the box width against the computed crossing threshold, which can be seen in Figure 12. The red dashed line lies at the threshold value 0.5 which we know is the true value for an infinite lattice.

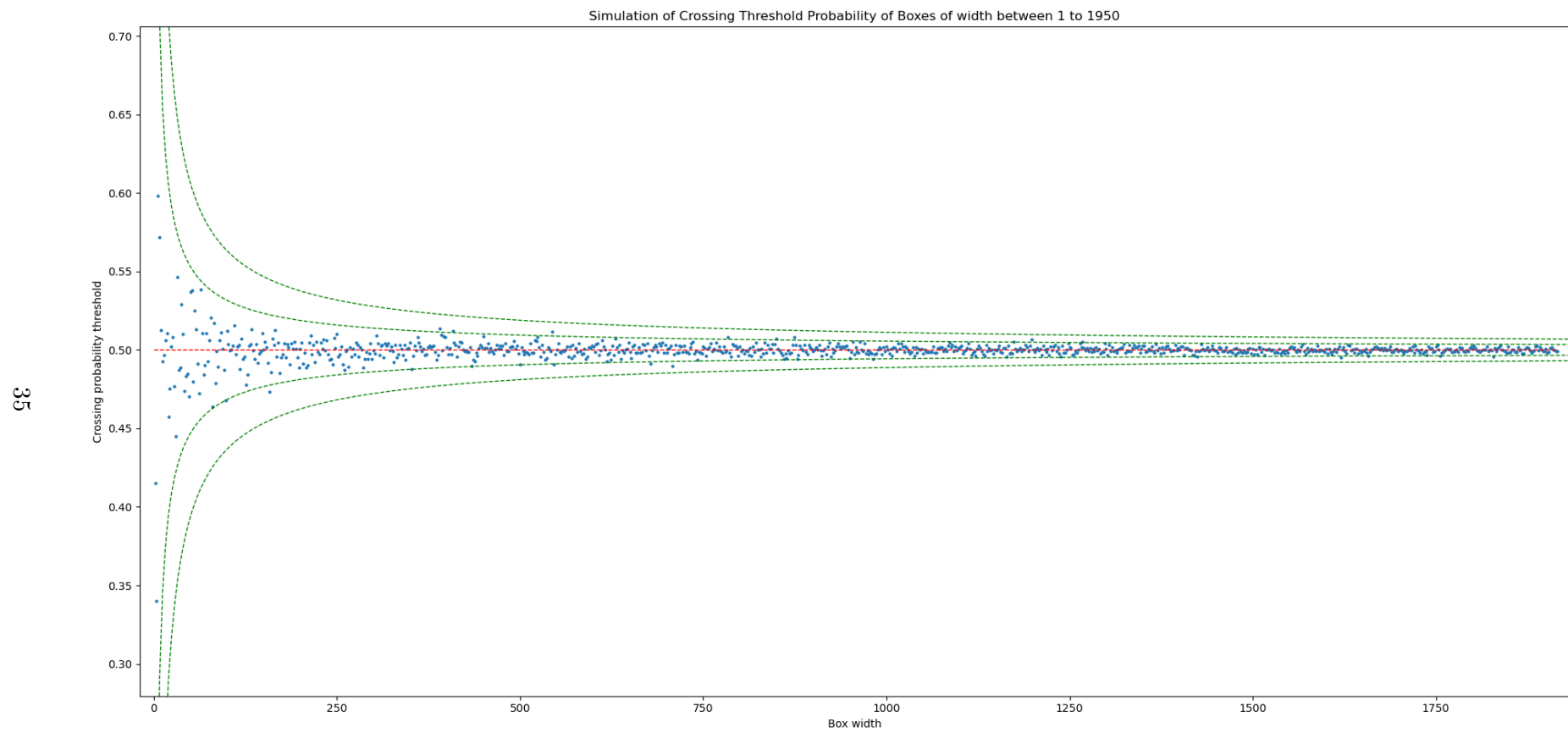


Figure 12: Simulated crossing thresholds for boxes of the square lattice. Expected standard deviations marked in green.

You can clearly see how the variability of crossing threshold decreases as the box width increases, becoming more concentrated around 0.5. Supposing we did not already know the true value of the infinite square lattice, this would provide us with a strong indication that the percolation threshold lies close to 0.5. In fact if we calculate the mean of the crossing thresholds for each box, disregarding the first 50 due to their small size, we would get a value of 0.499979 which when rounded to our precision of 0.0001 is 0.5000.

Another way to check our simulation technique is working correctly is to compare our observed distribution of crossing probabilities as p increases with what is known analytically. Proposition 6.8 of [10] tells how ‘sharp’ the threshold is around the critical value. It tells us that the derivative with respect to p of the probability of a horizontal open path around an n by n torus evaluated at the critical value $p = 1/2$ is

$$n^{3/4+o(1)}$$

We would expect a very similar result for horizontal open crossings of $n+1$ by n boxes. This informs us that the standard deviation of crossing thresholds should tend to $n^{3/4}$ as the box width n tends to infinity. Therefore, I plotted the first and second expected standard deviations in green on Figure 12. We can see almost all points lie within one standard deviation which provides further evidence that the simulation results are correct.

This verifies the method we have used gives a good approximation for the percolation threshold. Since no part of our simulation or algorithm depended on specific structure of the square planar lattice, it would be very easy to use this approach to investigate other planar graph structures. In fact, with some adaptations the Algorithm 1 could be extended to consider an n -dimension hypercube to allow the investigation of Bernoulli percolation on arbitrary dimension lattices. Further to this we can easily assign values from different random distributions to each edge, and explore how this affects the crossing threshold.

6 References

- [1] Michael Aizenman and Geoffrey Grimmett. ‘Strict monotonicity for critical points in percolation and ferromagnetic models’. In: *Journal of Statistical Physics* 63.5 (1991), pp. 817–835.
- [2] Robert B Ash et al. *Probability and measure theory*. Academic Press, 2000.
- [3] Patrick Billingsley. *Probability and measure*. John Wiley & Sons, 2008.
- [4] Bela Bollobás and Oliver Riordan. *Percolation*. Cambridge University Press, 2006.
DOI: 10.1017/CB09781139167383.

- [5] Béla Bollobás and Oliver Riordan.
‘A short proof of the Harris–Kesten theorem’. In: *Bulletin of the London Mathematical Society* 38.3 (2006), pp. 470–484.
- [6] S. R. Broadbent and J. M. Hammersley.
‘Percolation processes: I. Crystals and mazes’.
In: *Mathematical Proceedings of the Cambridge Philosophical Society* 53.3 (1957), pp. 629–641. DOI: 10.1017/S0305004100032680.
- [7] Duncan S Callaway et al.
‘Network robustness and fragility: Percolation on random graphs’.
In: *Physical review letters* 85.25 (2000), p. 5468.
- [8] Hugo Duminil-Copin. ‘Introduction to Bernoulli percolation’.
In: *Lecture notes available on the webpage of the author* (2018).
- [9] Ehud Friedgut and Gil Kalai.
‘Every monotone graph property has a sharp threshold’.
In: *Proceedings of the American mathematical Society* 124.10 (1996), pp. 2993–3002.
- [10] Christophe Garban and Jeffrey E Steif.
Noise sensitivity of Boolean functions and percolation. Vol. 5.
Cambridge University Press, 2014.
- [11] Geoffrey Grimmett. ‘Bond Percolation in Two Dimensions’. In:
Percolation. Springer, Berlin, Heidelberg, 1999, pp. 281–348.
DOI: 10.1007/978-3-662-03981-6_11.
- [12] T. E. Harris. ‘A lower bound for the critical probability in a certain percolation process’. In: *Mathematical Proceedings of the Cambridge Philosophical Society* 56.1 (1960), pp. 13–20.
DOI: 10.1017/S0305004100034241.
- [13] Harry Kesten et al. ‘The critical probability of bond percolation on the square lattice equals $1/2$ ’.
In: *Communications in mathematical physics* 74.1 (1980), pp. 41–59.
- [14] Czes Kosniowski. ‘12A The Jordan curve theorem’. In:
A First Course in Algebraic Topology.
Cambridge University Press, 1980, pp. 100–109.
DOI: 10.1017/CB09780511569296.015.
- [15] Lauren Meyers. ‘Contact network epidemiology: Bond percolation applied to infectious disease prediction and control’. In: *Bulletin of the American Mathematical Society* 44.1 (2007), pp. 63–86.
- [16] S Roch. *Modern discrete probability: An essential toolkit. chapter 4*. 2015.
- [17] Dietrich Stauffer and Ammon Aharony.
Introduction to percolation theory. Taylor & Francis, 2018.

- [18] Jeffrey E Steif. ‘A mini course on percolation theory’.
In: *Jyväsylä lectures in mathematics* 3 (2011).
- [19] Christopher Stover and Eric W. Weisstein.
Percolation Threshold. From MathWorld—A Wolfram Web Resource.
Visited on 28/01/22. URL:
<https://mathworld.wolfram.com/PercolationThreshold.html>.
- [20] The PyPy Team. *PyPy Features*. Visited on 09/02/22. Dec. 2019.
URL: <https://www.pypy.org/features.html>.
- [21] David Williams. *Probability with Martingales*.
Cambridge University Press, 1991.
DOI: 10.1017/CB09780511813658.

7 Appendix

Below is the Python code I used to carry out my simulations:

```
import networkx as nx
import random
import time

def gen2dLattice(n):
    '''Create a 2n by 2n square graph'''
    return nx.grid_2d_graph(range(-n,n+1),range(-n,n+1))

def percGraph(n):
    G = gen2dLattice(n)
    num_edges = nx.number_of_edges(G)
    unifRvs = [random.random() for _ in range(num_edges)]
    # Create Dictionary from edges to their unif rv
    dictUnifRvs = {coord:rv for coord,rv in zip(G.edges,iter(unifRvs))}
    # bind the unif rvs to the edges in the graph
    nx.set_edge_attributes(G,dictUnifRvs,name="unifRv")
    #create aug graph with s and f nodes which connect to the left an right sides resp
    GAug = G.copy()
    GAug.add_node('s')
    GAug.add_node('f')
    for i in range(-n,n+1): # Add the edges to s and f
        GAug.add_edge('s',(-n,i))
        GAug.add_edge((n,i),'f')
    return G , GAug

def find_crossing_path(openEdges,GAug) -> list:
    '''finds left/right crossing by open edges if one exists
```

```

Returns: List of edges on the path'''
    # find the subgraph of open edges plus the aux left right edges
    GOpen = nx.edge_subgraph(GAug, list(GAug.edges('f'))
        + list(GAug.edges('s'))
        + [(e[0], e[1]) for e in openEdges])
    edgepath = []
    try:
        openPaths = nx.algorithms.shortest_paths.bidirectional_shortest_path(GOpen, 's', 'f')
    except:
        #If no path is found
        openPaths = []
    for i in range(len(openPaths)): #Convert from list of nodes to list of edges
        edgepath.append((openPaths[i-1], openPaths[i]))
    openPaths = edgepath
    return openPaths

def find_threshold_binary_search(i, precision):
    '''Finds the threshold for a left right crossing on a 2i x 2i box'''
    G, GAug = percGraph(i)
    lo=0
    hi=1
    p = (lo+hi)/2
    openEdges = [e for e in G.edges(data=True) if e[2]['unifRv']<=p]
    crossing = find_crossing_path(openEdges, GAug)
    while lo <= hi:
        p = (lo+hi)/2
        openEdges = [e for e in G.edges(data=True) if e[2]['unifRv']<=p]
        crossing = len(find_crossing_path(openEdges, GAug))
        if crossing == 0:
            pAlt = p + precision
            openEdgesAlt = [e for e in G.edges(data=True) if e[2]['unifRv']<=pAlt]
            crossingAlt = len(find_crossing_path(openEdgesAlt, GAug))
            if crossingAlt > 0:
                return pAlt
            else:
                lo = pAlt
        else:
            pAlt = p - precision
            openEdgesAlt = [e for e in G.edges(data=True) if e[2]['unifRv']<=pAlt]
            crossingAlt = len(find_crossing_path(openEdgesAlt, GAug))
            if crossingAlt == 0:
                return p
            else:

```

```
        hi = pAlt
    return 'Error'

start_time = time.time()
thresholds = [find_threshold_binary_search(i,0.0001) for i in range(1,1000)]
print(thresholds)
end_time = time.time()
print(f"It took {end_time-start_time:.2f} seconds to compute")
```