



Computing the
Modular Inverse
of a Polynomial
Function over
 $GF(2^P)$
Using Bit Wise
Operation

Seshagiri Prabhu
N

2 Introduction

Language and
platform
Qt Framework

3 Collage Creator
Tool

Key features

4 Demo

Computing the Modular Inverse of a Polynomial Function over $GF(2^P)$ Using Bit Wise Operation

Seshagiri Prabhu N

M.Tech, Cyber Security and Networks (First Year)
Amrita School of Engineering, Amritapuri Campus

April 27, 2013





Outline

Computing the
Modular Inverse
of a Polynomial
Function over
 $GF(2^P)$
Using Bit Wise
Operation

Seshagiri Prabhu
N

2 Introduction

Language and
platform
Qt Framework

3 Collage Creator
Tool

Key features

4 Demo

- 1 Introduction
- 2 Collage Creator Tool
- 3 Demo



Introduction

Computing the Modular Inverse of a Polynomial Function over $GF(2^P)$ Using Bit Wise Operation

Seshagiri Prabhu
N

2 Introduction

Language and
platform
Qt Framework

3 Collage Creator Tool

Key features

4 Demo

- 1 Relevance of Modulo arithmetic in public key crypto system
- 2 The use of **Extended Euclidean Algorithm (EEA)** to evaluate the multiplicative inverse



Contribution of this paper

Computing the
Modular Inverse
of a Polynomial
Function over
 $GF(2^P)$
Using Bit Wise
Operation

Seshagiri Prabhu
N

2 Introduction

Language and
platform
Qt Framework

3 Collage Creator
Tool

Key features

4 Demo

Contribution of this paper

Computerized algorithm for the determination of the multiplicative inverse of a polynomial over $GF(2^P)$ using simple bit wise shift and XOR operations.



Problem Description

Computing the
Modular Inverse
of a Polynomial
Function over
 $GF(2^P)$
Using Bit Wise
Operation

Seshagiri Prabhu
N

2 Introduction
Language and
platform
Qt Framework

3 Collage Creator
Tool
Key features

4 Demo

EEA

Let $A(x)$ and $B(x)$ be polynomials.

EEA gives U and V such that

$$\gcd(A, B) = U * A + V * B$$

Note

If A is irreducible, then its gcd is 1 , and we are only interested in V , which is the inverse of $B[mod A]$



Problem Description

Computing the
Modular Inverse
of a Polynomial
Function over
 $GF(2^p)$
Using Bit Wise
Operation

Seshagiri Prabhu
N

2 Introduction

Language and
platform
Qt Framework

3 Collage Creator
Tool
Key features

4 Demo

Polynomial representation

The finite field is a representative of a polynomial function with respect to one variable x :

$$GF(2^p) = x^{p-1} + x^{p-2} + \dots + x^2 + x^1$$

Example

$$\text{Finite field } GF(2^8) = x^8 + x^4 + x^3 + x + 1$$

$$53_{10} \rightarrow 1010011_2 \rightarrow (x^6 + x^4 + x + 1)$$

$$\text{The EEA of } 53 \text{ on } GF(2^8) \text{ is } x^7 + x^6 + x^3 + x$$



Proposed Algorithm

Computing the
Modular Inverse
of a Polynomial
Function over
 $GF(2^P)$
Using Bit Wise
Operation

Seshagiri Prabhu
N

2 Introduction

Language and
platform
Qt Framework

3 Collage Creator
Tool

Key features

4 Demo

procedure MULTIPLICATIVE INVERSE($A_3[]$, $B_3[]$)

$C_1 = A_2 = B_2 = 0$;

while ($B_3 > 1$) **do**

▷ Step 1 do

$Q = 0$;

$Temp = B_3$;

while ($A_3 > Temp \parallel BitSize(C) \geq BitSize(Temp)$) **do**

$Q_1 = 1$;

while ($A_{3MSB} == B_{3MSB}$) **do**

$B_3 = B_3 \ll LinearLeftShift$;

$Q_1 = Q_1 * 2$;

end while

$Q = Q + Q_1$;

$A_3 = A_3[] \oplus B_3[]$;

$B_3 = Temp$;

end while

$A_2 = B_2$; $B_3 = A_3$; $A_3 = Temp$;

$N = BitSize(Q)$;

▷ Binary Bit Size of Q

$Temp = B_2$; $C_2 = 0$;

▷ Step2

while ($N \geq 1$) **do**

$C_2 = 0_d$;

if ($Q_N == 1$) **then**

▷ Testing if Nth bit of Q is 1

$C_1 = B_2 \ll N - 1$;

▷ Linear left shift by N-1 times

$C_2 = C_2 \oplus C_1$;

end if

$N = N - 1$;

end while

$B_2 = C_2$; $A_2 = Temp$; $B_2 = B_2 \oplus A_2$;

end while

▷ Multiplicative Inverse



Implementation of Algorithm

Let's apply EEA to $A = 283$ and $B = 42$

Computing the
Modular Inverse
of a Polynomial
Function over
 $GF(2^P)$
Using Bit Wise
Operation

Seshagiri Prabhu
N

2 Introduction

Language and
platform
Qt Framework

3 Collage Creator
Tool

Key features

4 Demo

i	Operation	Binary	U	V
0	A	100011011	1	0
1	B	000101010	0	1
2	$3 \ll B$	101010000	0	1000
	$A \leftarrow A \oplus (3 \ll B)$	001001011	1	1000
3	$1 \ll B$	001010100	0	0010
	$A \leftarrow A \oplus (1 \ll B)$	000011111	1	1010
4	$A < B \ A \rightleftharpoons B$			
	A	000101010	00	00001
	B	000011111	01	01010
	$1 \ll B$	000111110	10	10100
5	$A \leftarrow A \oplus (1 \ll B)$	000010100	10	10101
	$A < B \ A \rightleftharpoons B$			
6	A	000011111	01	01010
	B	000010100	10	10101
	$A \leftarrow A \oplus B$	000001011	11	11111
	$A < B \ A \rightleftharpoons B$			
7	A	000010100	010	010101
	B	000001011	011	011111
	$1 \ll B$	000010110	110	111110
	$A \leftarrow A \oplus (1 \ll B)$	000000010	100	101011
8	$A < B \ A \rightleftharpoons B$			
	A	000001011	00011	00011111
	B	000000010	00100	00101011
	$2 \ll B$	000001000	10000	10101100
9	$A \leftarrow A \oplus (1 \ll B)$	000000011	10011	10110011
	$A \leftarrow A \oplus B$	000000001	10111	10011000



Conclusion

Computing the Modular Inverse of a Polynomial Function over $GF(2^P)$ Using Bit Wise Operation

Seshagiri Prabhu
N

2 Introduction

Language and
platform
Qt Framework

3 Collage Creator Tool

Key features

4 Demo

- 1 This algorithm can be easily extended for determining the elements of the S-Box used in *AES*.
- 2 This algorithm is efficient for determining the multiplicative inverse of polynomial over $GF(2^P)$



Computing the
Modular Inverse
of a Polynomial
Function over
 $GF(2^P)$
Using Bit Wise
Operation

Seshagiri Prabhu
N

2 Introduction

Language and
platform
Qt Framework

3 Collage Creator
Tool

Key features

4 Demo

Possible future works

- 1 Optimize the algorithm
- 2 Comparative study with many existing algorithm
- 3 Implementation in hardware for real time applications



Computing the
Modular Inverse
of a Polynomial
Function over
 $GF(2^P)$
Using Bit Wise
Operation

Seshagiri Prabhu
N

Questions ?

2 Introduction

Language and
platform
Qt Framework

3 Collage Creator Tool

Key features

4 Demo

Seshagiri Prabhu
seshagiriprabhu@gmail.com