Why I Like Functional Programming

Adelbert Chang Box, Inc.

Let's go back..



http://gauss.cs.ucsb.edu/home/images/UCSB-from-air.jpg



Took first two college physics courses. Considered changing to physics. 2011

Started learning Haskell.
Took graduate static
analysis course.
2013

Presenting at an awesome FP conference.
2015





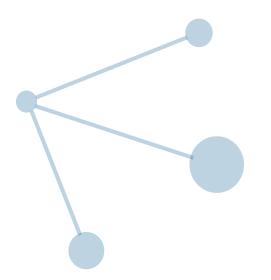
2010 Entered UCSB as CS major.

2012

Took upper-division math courses. Considered minoring in mathematics. Started learning Scala.

2014

Graduated UCSB as a CS major. Working at Box in a functional codebase.



Why CS

- Computers are pretty cool
- Applied problems to solve
- Fairly interactive
- Rooted in logic
- Can be done pretty much anywhere

2011

- Languages learned/used: Python, C, C++
- Took first two college physics courses
 - Considered switching major to physics

```
swapped = True
while swapped:
    swapped = False
    for i in range(len(xs) - 1):
        if xs[i] > xs[i + 1]:
            xs[i], xs[i + 1] = xs[i + 1], xs[i]
        swapped = True
```

Physics and Math

- Problem solving activity
- Can be done pretty much anywhere
- Simplicity and consistency across problems
- Example: Deriving law of motion from first principles

•
$$x = x_0 + v_0 + \frac{1}{2}a + \frac{2}{2}$$

Physics and Math

Deriving $x = x_0 + v_0 t + \frac{1}{2} a t^2$

$$v_{avg} = \frac{1}{2}(v_f + v_0)$$

$$\frac{1}{2}(v_0 + a\Delta t + v_0)$$

$$\frac{1}{2}(2v_0 + a\Delta t)$$

$$v_{avg} = v_0 + \frac{1}{2}a\Delta t$$

$$a = \Delta v / \Delta t$$

$$a = (v_f - v_0) / \Delta t$$

$$v_f - v_0 = a \Delta t$$

$$v_f = v_0 + a \Delta t$$

$$v_{avg} = \Delta x/\Delta t$$

$$\Delta x = v_{avg} * \Delta t$$

$$(x - x_0) = v_{avg} * \Delta t$$

$$x = x_0 + (v_{avg} * \Delta t)$$

$$x = x_0 + ((v_0 + \frac{1}{2}a\Delta t) * \Delta t)$$

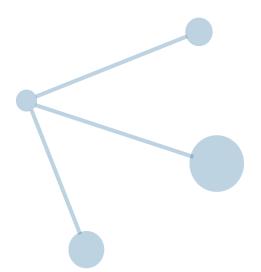
$$x_0 + v_0 t + \frac{1}{2}at^2$$

2012

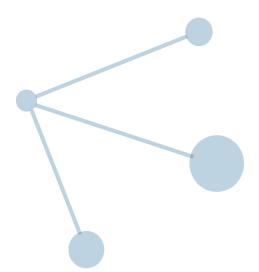
```
void insert(Node* &tree, int value) {
    if (!tree) return;
    Node *trav = tree, *parent;
    while (trav) {
        parent = trav;
        if (value < trav->data) trav = trav->left;
        else if (value > trav->data) trav = trav->right;
        else break;
    if (value < parent->data)
        parent->left = new Node(value);
    else
        parent->right = new Node(value);
```

2012

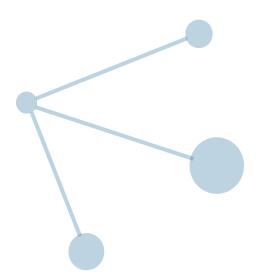
```
sealed abstract class Tree
case class Branch(data: Int, left: Tree, right: Tree)
  extends Tree
case class Leaf() extends Tree
def insert(tree: Tree, value: Int): Tree =
  tree match {
    case Leaf() => Branch(value, Leaf(), Leaf())
    case b@Branch(d, l, r) =>
      if (value < d) Branch(d, insert(l, value), r)</pre>
      else if (value > d) Branch(d, l, insert(r, value)
      else b
```



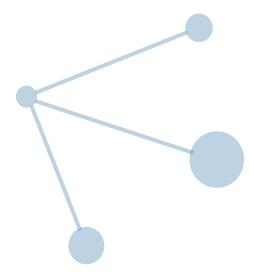
Hmm.. functional programming is pretty cool.



But is it as cool as physics and math?

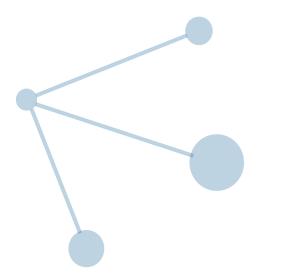


Let's see..



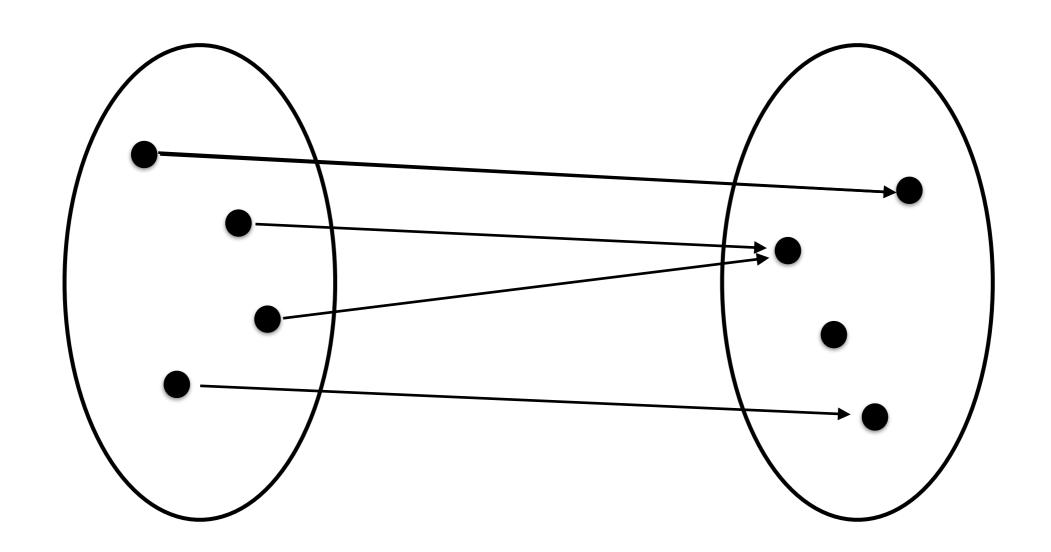
Functional Programing

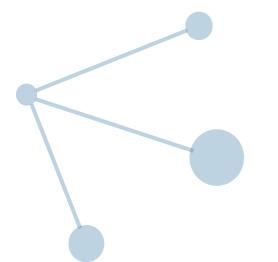
programming with functions



Functional Programming

programming with pure functions





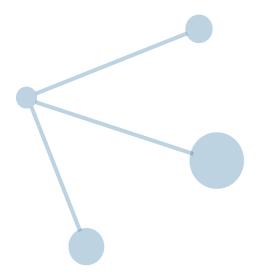
Functions

```
def parseIntPartial(s: String): Int =
    s.toInt // can throw NumberFormatException

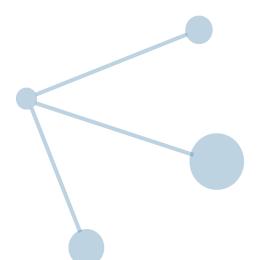
def parseIntTotal(s: String): Option[Int] =
    try {
        Some(s.toInt)
    } catch {
        case nfe: NumberFormatException => None
    }
```

Functions

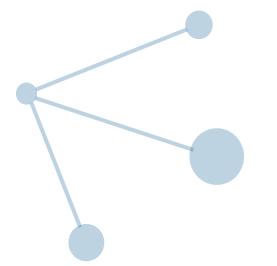
```
class Rng(var seed: Long) {
  def nextInt(): Int = {
    val int = getInt(seed)
    mutate(seed)
    int
case class Rng(seed: Long) {
  def nextInt: (Rng, Int) = {
    val int = getInt(seed)
    val newSeed = f(seed)
    (Rng(newSeed), int)
```



An expression **e** is referentially transparent if for all programs **p**every occurrence of **e** in **p**can be replaced with the result of evaluating **e** without changing the result of evaluating **p**.*



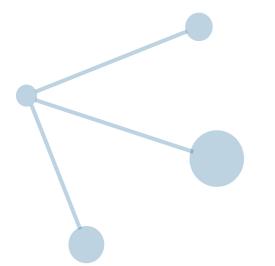
$$x^{4} - 12x^{2} + 36 = 0$$
 $ax^{2} + bx + c = 0$
 $y = x^{2}$ $(-b \pm \sqrt{b^{2} - 4ac}) / 2a$
 $y^{2} - 12y + 36 = 0$ $a = 1$
 $b = -12$
 $y = 6$ $c = 36$
 $x^{2} = 6$ $(-(-12) \pm \sqrt{(-12)^{2} - 4(1)(36)}) / 2(1)$
 $x = \pm \sqrt{6}$ $12 / 2$



```
def userInfo(id: UserId): Future[UserData]
val userId = . . .

val fetchData = userInfo(userId)

fetchData.retry {
   case StatusCode(429) => fetchData
}
```

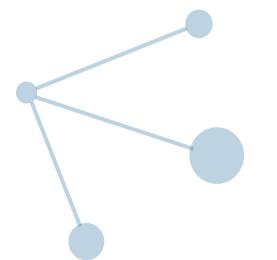


```
def userInfo(id: UserId): Future[UserData]

val userId = . . .

val fetchData = userInfo(userId)

fetchData.retry {
   case StatusCode(429) => userInfo(userId)
}
```

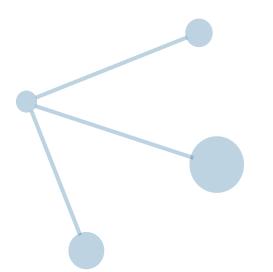


```
def userInfo(id: UserId): Task[UserData]

val userId = . . .

val fetchData = userInfo(userId)

fetchData.retry {
   case StatusCode(429) => fetchData
}
```

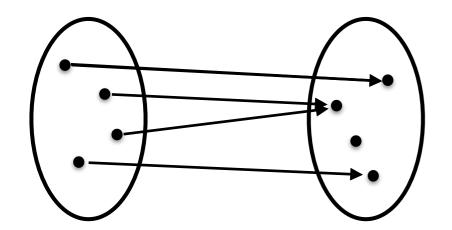


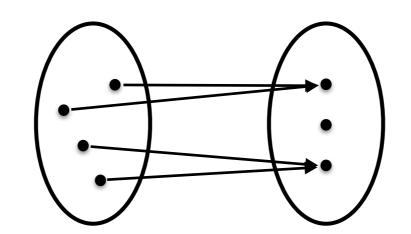
Functions

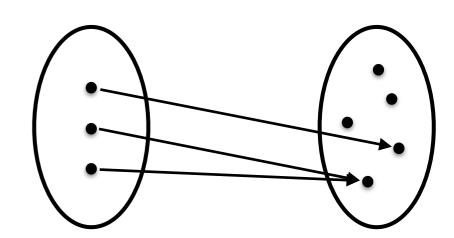
foo: A => B

bar: B => C

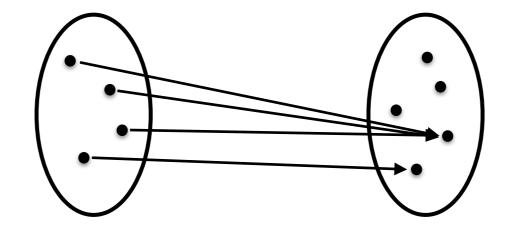
baz: $C \Rightarrow D$







baz.compose(bar).compose(foo): A => D



- Interesting programs will have effects
 - Optional values, error handling, asynchronous computation, input/output
 - Usual means aren't quite nice

```
def lookup(map: Map[Foo, Bar], foo: Foo): Bar =
  if (map.contains(foo)) map(foo) else null

val bar = lookup(someMap, someFoo)
  if (bar != null) {
        . . .
} else {
        . . .
}
```

- Interesting programs will have effects
 - Optional values, error handling, asynchronous computation, input/output
 - Usual means aren't quite nice

```
def lookup(map: Map[Foo, Bar], foo: Foo): Bar =
   map(foo)

try {
  val bar = lookup(map, someFoo)
} catch {
  case e: NoSuchElementException => . . .
}
```

Reify effects as values

```
// A value that might exist is either..
sealed abstract class Option[A]

// ...there
final case class Some[A](a: A) extends Option[A]

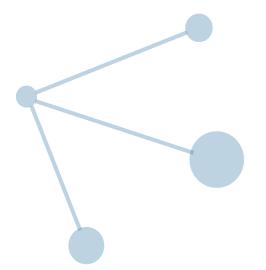
// ... or not there
final case class None[A]() extends Option[A]
```

Reify effects as values

```
def lookup(map: Map[Foo, Bar], foo: Foo): Option[Bar] =
  if (map.contains(foo)) Some(map(foo))
  else None
```

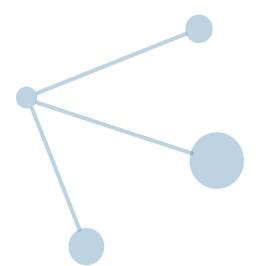
Missing values

```
sealed abstract class Option[A]
final case class Some[A](a: A) extends Option[A]
final case class None[A]() extends Option[A]
```



Errors

```
sealed abstract class Either[+E, +A]
final case class Success[+A](a: A) extends Either[Nothing, A]
final case class Failure[+E](e: E) extends Either[E, Nothing]
```



Ex. Errors

Manipulating Effects

 The type signature of our functions reflect the effects involved

```
def tokenFor(uid: UserId): IO[EncryptedToken]

def decrypt(token: EncryptedToken): Token

/** Client side */
val encrypted: IO[EncryptedToken] =
   tokenFor(. . .)

// Want EncryptedToken, have IO[EncryptedToken]
val decrypted = decrypt(???)
```

Manipulating Effects

```
/** Apply a pure function to an effectful value */
trait Functor[F[_]] {
  def map[A, B](fa: F[A])(f: A => B): F[B]
new Functor[I0] {
  def map[A, B](fa: IO[A])(f: A => B): IO[B] =
    new IO[B] {
      def unsafePerformIO(): B =
        f(fa.unsafePerformIO())
```

Manipulating Effects

```
def tokenFor(uid: UserId): IO[EncryptedToken]

def decrypt(token: EncryptedToken): Token

/** Client side */
val encrypted: IO[EncryptedToken] =
   tokenFor(. . .)

// Want EncryptedToken, have IO[EncryptedToken]
val decrypted: IO[Token] =
   encrypted.map(et => decrypt(et))
```

Similar mechanisms for manipulating multiple effects

- Often want getters/setters when working with data
- If getters/setters are per-object, cannot compose
- As with effects, reify as data
 - Getter: get an A field in an object S
 - Setter: change an A field in an object S

```
trait Lens[S, A] {
  def get(s: S): A

  def set(s: S, a: A): S
}
```

```
trait Lens[S, A] { outer =>
 def get(s: S): A
 def set(s: S, a: A): S
  def modify(s: S, f: A \Rightarrow A): S =
    set(s, f(get(s)))
  def compose[B](other: Lens[A, B]): Lens[S, B] =
    new Lens[S, B] {
      def get(s: S): B = other.get(get(s))
      def set(s: S, b: B): S =
        set(s, other.set(outer.get(s), b))
```

```
case class Employee(position: Position)
object Employee {
  val position: Lens[Employee, Position] = . . .
case class Team(manager: Employee, . . .)
object Team {
  val manager: Lens[Team, Employee] = . . .
/** Client side */
val 1: Lens[Team, Position] =
  Team.manager.compose(Employee.position)
l.set(someTeam, somePosition)
```

- Effect-ful modifications can be useful
 - Possibly failing: A => Option[A]
 - Many possible values: A => List[A]
 - Async: A => Task[A]

```
trait Lens[S, A] {
  def modifyOption(s: S, f: A => Option[A]): Option[S]

def modifyList(s: S, f: A => List[A]): List[S]

def modifyTask(s: S, f: A => Task[A]): Task[S]
}
```

Lens

```
trait Lens[S, A] {
  def modifyOption(s: S, f: A => Option[A]): Option[S] =
    f(get(s)).map(a => set(s, a))

def modifyList(s: S, f: A => List[A]): List[S] =
    f(get(s)).map(a => set(s, a))

def modifyTask(s: S, f: A => Task[A]): Task[S] =
    f(get(s)).map(a => set(s, a))
}
```

Lens

```
trait Lens[S, A] {
  def modifyF[F[_] : Functor](s: S, f: A => F[A]): F[S] =
    f(get(s)).map(a => set(s, a))
}
```

- Example functors:
 - Option, Either, Future, List, IO
- What if we want to just modify without any effect?

Id

```
type Id[A] = A

new Functor[Id] {
  def map[A, B](fa: Id[A])(f: A => B): Id[B] =
}
```

ld

```
type Id[A] = A

new Functor[Id] {
  def map[A, B](fa: A)(f: A => B): B =
}
```

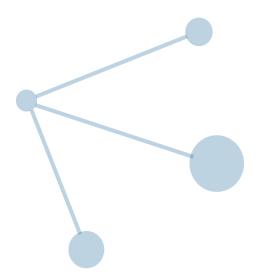
ld

```
type Id[A] = A

new Functor[Id] {
  def map[A, B](fa: A)(f: A => B): B = f(fa)
}
```

```
trait Lens[S, A] {
  def modifyF[F[_] : Functor](s: S, f: A => F[A]): F[S] =
    f(get(s)).map(a => set(s, a))

  def modify(s: S, f: A => A): S = modifyF[Id](s, f)
}
```



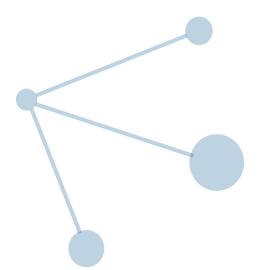
Setting

```
trait Lens[S, A] {
  def modifyF[F[_] : Functor](s: S, f: A => F[A]): F[S] =
    f(get(s)).map(a => set(s, a))

  def modify(s: S, f: A => A): S = modifyF[Id](s, f)

  def set(s: S, a: A): S = modify(s, const(a))
}
```

def const[A, B](a: A)(b: B): A = a



Setting

```
trait Lens[S, A] {
  def modifyF[F[_] : Functor](s: S, f: A => F[A]): F[S]

  def get(s: S): A

  def modify(s: S, f: A => A): S = modifyF[Id](s, f)

  def set(s: S, a: A): S = modify(s, const(a))
}
```

def const[A, B](a: A)(b: B): A = a

```
trait Lens[S, A] {
  def modifyF[F[_] : Functor](s: S, f: A => F[A]): F[S]
  def get(s: S): A
}
```

- Can we define `get` in terms of `modify`?
 - modifyF` gives us some F[S].. but we want an A
 - We need some way of "ignoring" the S parameter and still get an A back
 - Type-level constant function

```
def const[A, B](a: A)(b: B): A = a
```

```
type Const[Z, A] = Z

new Functor[Const[Z 2]] {
  def map[A, B](fa: Const[Z, A])(f: A => B): Const[Z, B]
}
```

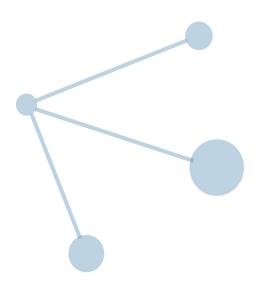
```
type Const[Z, A] = Z

new Functor[Const[Z, ?]] {
  def map[A, B](fa: Z)(f: A => B): Z =
}
```

```
type Const[Z, A] = Z
new Functor[Const[Z, ?]] {
  def map[A, B](fa: \mathbb{Z})(f: A \Rightarrow B): \mathbb{Z} = fa
trait Lens[S, Z] {
  def modifyF[F[_] : Functor](s: S, f: Z => F[Z]): F[S]
  def get(s: S): Z = {
    val const: Const[Z, S] = modifyF[Const[Z, ?]](s, z \Rightarrow z)
    const
```

Interested?

- Functional Programming
 - Cats, Scalaz
 - Argonaut, Atto, Doobie, HTTP4S, Monocle,
 Remotely, Scalaz-stream, Shapeless, Typelevel.org
- Algebra
 - Algebird, Algebra, Breeze, Spire
- Big Data
 - Scalding + Algebird, Scoobi, Spark



EOF