Natural Computing, Assignment 1

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1

Due to no eliteness, we can treat the best member as any other member in the pool.

(a) There are 100 individuals, with a mean of 76. That means that the total size of the roulette wheel will be 7600. The current best member has a fitness of 157, so it will occupy $\frac{157}{7600}$th of the roulette wheel. This means that it has roughly a 2% chance of being selected. If we select 100 members, we expect to select the best member $\frac{157}{7600} * 100 \approx 2$ times.

(b) The chance that we don't select the fittest member is $1 - \frac{157}{7600}$. So the chance that we never select it is $\left(1 - \frac{157}{7600}\right)^{100} \approx 0.124$.

2

The fitness function is

$$f(x) = x^2.$$

The members of the pool are $x = 3, x = 5, \text{and} x = 7$. So the total fitness is $3^2+5^2+7^2 = 83$. The chance to select each individual is:

$x = 3 :\ 9/83 \approx 0.108,$

$x = 5 :\ 25/83 \approx 0.301,$

$x = 5 :\ 49/83 \approx 0.590.$

When using the alternative selection function

$$f_1(x) = x^2 + 8,$$

the total fitness is $83 + 24 = 107$ and we obtain the following results:

$x = 3 :\ 17/107 \approx 0.159,$

$x = 5 :\ 33/107 \approx 0.308,$

$x = 5 :\ 57/107 \approx 0.533.$

We can see that the second function yields a lower selection pressure. This is because the relative boost of 8 is much higher for the less fit individuals. It almost doubles the fitness of the $x = 3$ individual, and only amounts to a 0.16-th increase in fitness of the $x = 7$ individual.

3

(a) When running the algorithm with $n = 100$ for 1500 iterations, we can observe the result found in figure 1. We can see that the Monte-Carlo search as specified in the assignment did not find the optimal solution.

(b) Running the algorithm ten times, the algorithm finds the optimal solution 0 times.

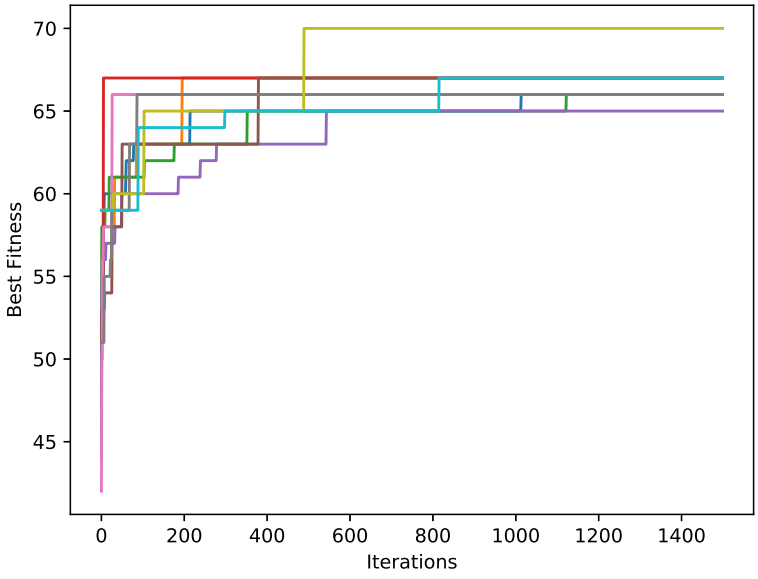

Figure 1: Monte Carlo Search for the Counting Ones problem, ran 10 times.

4

(a) When running the (1+1)-GA algorithm for the Counting Ones problem, we found the results from figure 2.

(b) As we can observe from our results, the algorithm found the optimum 10 times out of 10 runs.

(c) If we compare these results to those of the Monte-Carlo algorithm, we can easily see that the Genetic Algorithm is a significant improvement, as the Monte-Carlo algorithm did not find the optimum once and the GA algorithm every time.

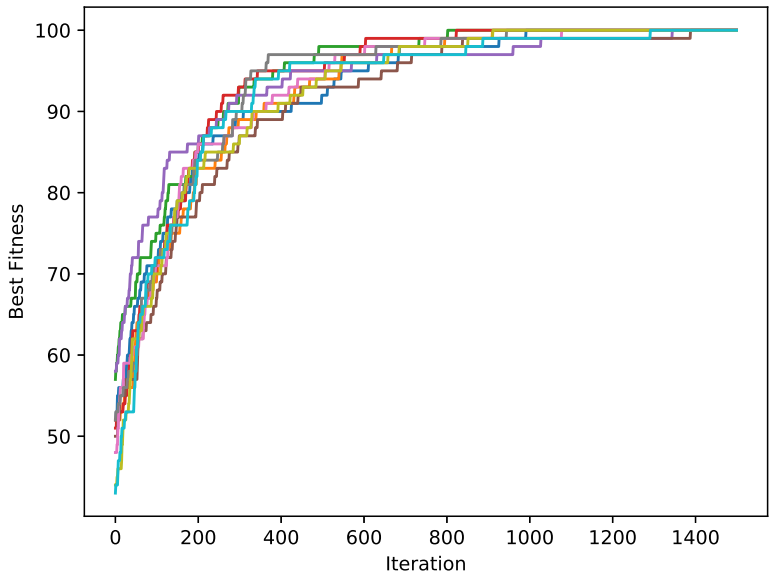

Figure 2: (1+1)-GA for the Counting Ones problem, ran 10 times.

5

Suitable fonction : $F = \{\wedge, \rightarrow, \vee, \leftrightarrow\}$
Terminal set : $T = \{x, y, z, true\}$
S-expression : $(\rightarrow (\wedge\ x\ true)(\vee(\vee\ x\ y)(\leftrightarrow\ z(\wedge\ x\ y))))$

6

Before we can implement the search algorithm, we have to make some changes to the primitive function set. First, $\log 0 = -\infty$, we change this and set the value to 0, the *exp* function grows exponentially, in order to prevent overflows, we limit it to a value maximum value and finally *div*, since division by zero is not defined. For this we created protected versions of the functions.

For the implementation of this assignment we used the python module `DEAP`. For the tree structure we used a `genHalfAndHalf` tree, since we think it has more flexibility than using a *grown* or *full* tree.

In figure 3, we see the best of generation fitness per generation. What we can observe from this figure and the output of the algorithm is that the best solution has a *-sum of absolute errors* of -0.02, which is pretty good. It is not perfect, we think that a perfect solution is hard to achieve with the current settings.

Do note that the algorithm could not find a solution within 1000 generations, but it kept (marginally) improving. So giving it more time would lead to a better solution. We hypothesize that the function is not actually generated by a possible function, so we are instead approximating the best approximation. A solution to speed up the search may be

to allow mutations, in that way we search a bigger space. We did a preliminary test, which converged on an error of -0.0008, so we can see that this speeds up the search.

In figure 4, we see the best of generation size per generation. From this figure, we may observe that the best solution of the algorithm does not necessarily have to grow in size to improve. Although we can see a tendency to grow.

After repeating this experiment a 100 times, we believe to have found the solution. The algorithm managed to find an individual whose error was $-8.68353e-16$, which is near 0. This validates our idea that we could not search the search space fast enough. Running the algorithm for longer would probably lead to the same results.

The results may be found in figure 5-6. The latter may be more evidence to the hypothesis about the best of generation size. The function that achieved this error is the following:

```
mul(add(add(x, mul(x, x)), protected_div(sub(x, x), mul(x, mul(x, x)))), add(mul(x, x),
cos(protected_log(protected_div(mul(protected_exp(sub(add(x, mul(x, x)), sin(mul(add(x,
mul(x, x)), sin(sin(sub(x, x)))))))), sin(sin(protected_log(protected_div(mul(x, x),
sub(mul(x, x), x)))))), sub(sub(x, x), x))))))
```

Although the following function also produces a near zero score. Interestingly, after the best solution is found, equivalent but longer solutions are found too. They are then added since there is no penalty for size, and the Hall of Fame adds equally fit individuals as well.

```
add(mul(add(x, mul(add(mul(x, x), x), x)), x), x)
```

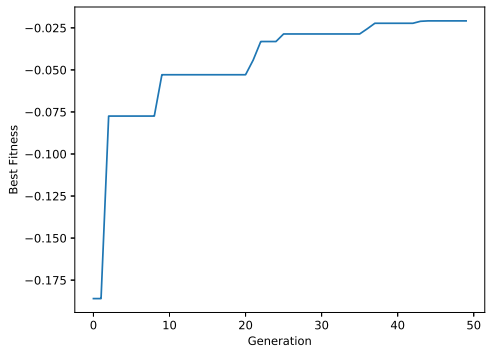

Figure 3: Best of generation fitness per generation.

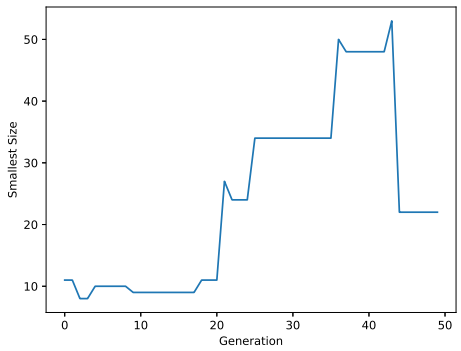

Figure 4: Best of generation size per generation.

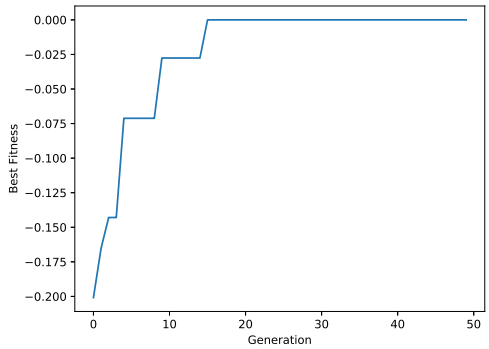

Figure 5: Best of generation fitness per generation.

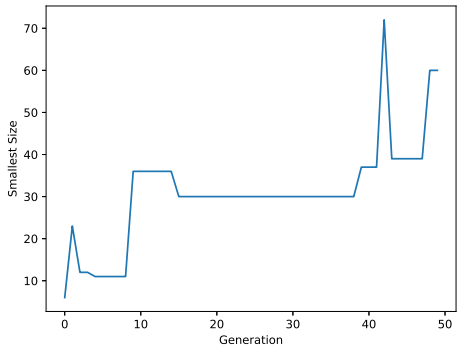

Figure 6: Best of generation size per generation.