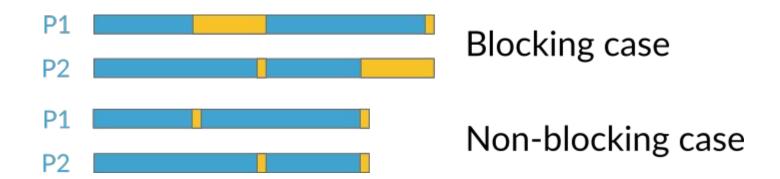


Outline

- 1. MPI Overview
- 2. Basic Structure of a MPI program
- 3. Messages and Point-to-Point Communication
- 4. Non-blocking Communication
- 5. Derived Data Types
- 6. Collective Communication



Why nonblocking communication?



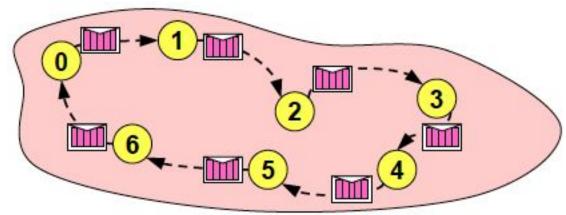
Blocking Communication requires processes to stop working to send and receive messages

Non Blocking Communication permits processes to keep working while checking once in a while if a message has been sent

Deadlock

•Code in each MPI process:

```
MPI_Ssend(..., right_rank,...)
MPI_Recv (..., left_rank,...)
```



will block and never return because MPI_Recv cannot be called in the right-hand MPI process

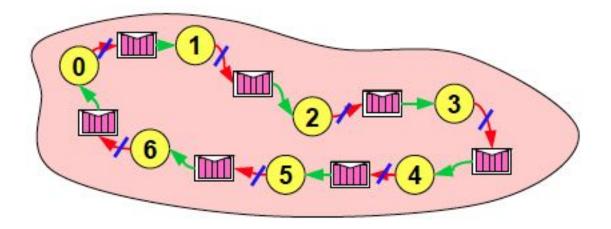
Same problem with standard send mode (MPI_Send), if MPI implementation chooses synchronous protocol

Separate communication intro three phases:

- •Initiate nonblocking communication that returns Immediately (routine name starting with MPI_I...)
 In the example: Initiate nonblocking send to the right neighbor (□)
- Do some work

In the example: Receiving the message from left neighbor (\square)

•Wait for nonblocking communication to complete (/)





Nonblocking Send

```
MPI_Isend(void *buf, int count, MPI_Datatype datatype, int
dest, int tag, MPI_Comm comm, MPI_Request *request)

MPI_Wait(MPI_Request *request, MPI_Status *status)
```

- •buf must not be modified between Isend and Wait
- •<u>Isend</u> + <u>Wait</u> directly after *Isend* is equivalent to blocking call (*Ssend*)
- •<u>status</u> is not used.

Nonblocking Receive

```
MPI_Irecv (void *buf, int count, MPI_Datatype datatype, int
source, int tag, MPI_Comm comm, MPI_Request *request)

MPI_Wait(MPI_Request *request, MPI_Status *status)
```

- •<u>buf</u> must not be modified between *Irecv* and *Wait*
- Message <u>status</u> is returned in Wait

•Request handle:

- Must be stored in local variables: MPI_Request
- Is generated by a nonblocking communication routine
- Is used (and freed) in the MPI_WAIT routine



Blocking and Non-Blocking

- Send and receive can be blocking or nonblocking
- •A blocking send can be used with a nonblocking receive, and vice-versa
- Nonblocking sends can use any mode

– standard: MPI ISEND

– synchronous: MPI ISSEND

– buffered: MPI_IBSEND

– ready: MPI_IRSEND

•Synchronous mode affects completion, i.e. MPI_Wait / MPI_Test,

Completion

```
MPI_Wait(MPI_Request *request, MPI_Status *status)
MPI_Test(MPI_Request *request,int *flag, MPI_Status *status)
```

Completion can be checked by:

- Wait: It blocks the task (does not go back of the routine) until the communication has finalized.
- Test: It gives back a value TRUE or FALSE depending if the communication has finalized or not.

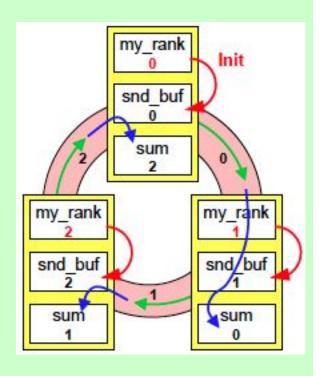
Example

```
/* Blocking*/
MPI RECV(x,N,MPI Datatype,...,&status)
/* Non Blocking with wait */
MPI IRECV(x,N,MPI Datatype,...,&request)
... Realise unrelated work with the vector X
MPI WAIT(&request, &status)
... Realise work related with the vector X
/* Non Blocking with test */
MPI IRECV(x,N,MPI Datatype,...,&request)
MPI TEST(&request, &flag, &status)
while (flag==FALSE) {
     ... Realise unrelated work with the vector X
     MPI TEST(&request, &flag, &status)
... Realise work related with the vector X
```



Activity 4: Rotating information around a ring

- 1.A set of processes are arranged in a ring.
- 2.Each process stores its rank into an integer variable *snd_buf*.
- 3. Each process passes this on to its neighbour on the right.
- 4. Each processor calculates the sum of the values.
- 5.Repeat steps 2-5 with "size" iterations (size = number of processes)
- 6.Implement the program using blocking operations and verify the correctness.
- 7. Use nonblocking MPI_Issend and verify the correctness.





Activity4: Rotating information around a ring

