

✓ ¡Felicitaciones! ¡Aprobaste!

Calificación recibida 100 % Para Aprobar 80 % o más

Ir al siguiente elemento

Problem Set #3

Calificación de la entrega más reciente: 100 %

1. Suppose you implement the functionality of a priority queue using a *sorted* array (e.g., from biggest to smallest). What is the worst-case running time of Insert and Extract-Min, respectively? (Assume that you have a large enough array to accommodate the Insertions that you face.)

1 / 1 punto

- ☐ $\Theta(1)$ and $\Theta(n)$
- ☐ $\Theta(n)$ and $\Theta(n)$
- ☒ $\Theta(n)$ and $\Theta(1)$
- ☐ $\Theta(\log n)$ and $\Theta(1)$

✓ Correcto

2. Suppose you implement the functionality of a priority queue using an *unsorted* array. What is the worst-case running time of Insert and Extract-Min, respectively? (Assume that you have a large enough array to accommodate the Insertions that you face.)

1 / 1 punto

- ☒ $\Theta(1)$ and $\Theta(n)$
- ☐ $\Theta(n)$ and $\Theta(1)$
- ☐ $\Theta(n)$ and $\Theta(n)$
- ☐ $\Theta(1)$ and $\Theta(\log n)$

✓ Correcto

3. You are given a heap with n elements that supports Insert and Extract-Min. Which of the following tasks can you achieve in $O(\log n)$ time?

1 / 1 punto

- ☐ Find the median of the elements stored in the heap.
- ☐ Find the largest element stored in the heap.
- ☒ Find the fifth-smallest element stored in the heap.
- ☐ None of these.

✓ Correcto

4. You are given a binary tree (via a pointer to its root) with n nodes. As in lecture, let $\text{size}(x)$ denote the number of nodes in the subtree rooted at the node x . How much time is necessary and sufficient to compute $\text{size}(x)$ for every node x of the tree?

1 / 1 punto

- ☐ $\Theta(n^2)$
- ☐ $\Theta(\text{height})$
- ☐ $\Theta(n \log n)$

☒ $\Theta(n)$

✓ **Correcto**

For the lower bound, note that a linear number of quantities need to be computed. For the upper bound, recursively compute the sizes of the left and right subtrees, and use the formula $\text{size}(x) = 1 + \text{size}(y) + \text{size}(z)$ from lecture.

5. Suppose we relax the third invariant of red-black trees to the property that there are no *three* reds in a row. That is, if a node and its parent are both red, then both of its children must be black. Call these *relaxed* red-black trees. Which of the following statements is *not* true?

1 / 1 punto

- ☒ Every binary search tree can be turned into a relaxed red-black tree (via some coloring of the nodes as black or red).
- ☐ There is a relaxed red-black tree that is not also a red-black tree.
- ☐ The height of every relaxed red-black tree with n nodes is $O(\log n)$.
- ☐ Every red-black tree is also a relaxed red-black tree.

✓ **Correcto**

A chain with four nodes is a counterexample.