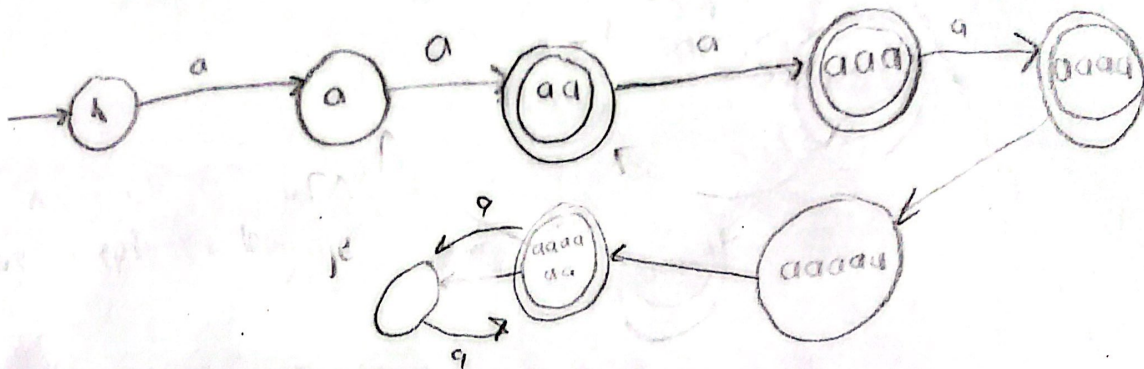


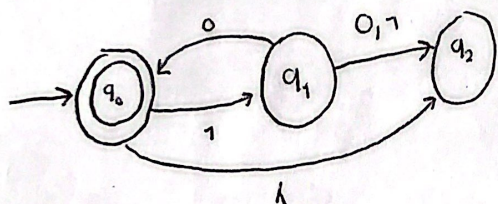
# ACEPTADORES FINITOS NO DETERMINISTAS

② Find a dfa that accepts the language defined by the nfa in figure 2.8

Language

$$\{aaa, (aa)^n, n \geq 0\}$$


④ In Figure 2.9, find  $s^*(90, 1011)$  and  $s^{\pm}(9, 01)$



a)  $(q_0, 1011) = \{q_0, q_1, q_2\}$

$$\begin{array}{ccccc}
 q_0 & \xrightarrow{1} & q_1 & \xrightarrow{0} & q_2 \rightarrow \emptyset \\
 \vdots & & \vdots & & \vdots \\
 q_2 \rightarrow \emptyset & & q_0 & \xrightarrow{1} & q_1 \xrightarrow{1} q_2 \\
 & & \vdots & & \vdots \\
 & & q_2 \rightarrow \emptyset & & 
 \end{array}$$

$$b) (a_1, 01) = \{a_1\}$$

$$q_0 \xrightarrow{0} a_0 \xrightarrow{1} q_1 \quad \checkmark$$

⑤ In Figure 2.10, find  $\delta^+(a, a)$  and  $\delta^+(a, k)$

a)  $f^*(q_0, a) = \{q_0, q_1, q_2\}$

$$q_0 \xrightarrow{a} q_1$$

$$\quad \quad \quad \downarrow$$

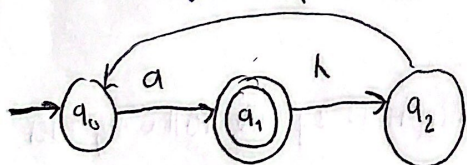
$$\quad \quad \quad q_2$$

$$\quad \quad \quad \downarrow$$

$$\quad \quad \quad q_3$$

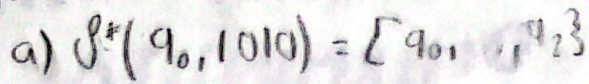
$$b)(q_1, k) = \{q_0, q_1, q_2\}$$

$$\begin{array}{c} a_1 \\ \vdots \\ a_2 \\ \vdots \\ a_0 \end{array}$$





110 111 112



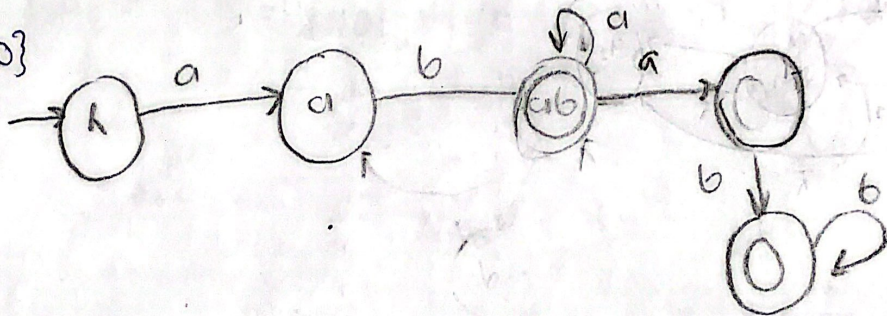
$$\begin{array}{c}
 \begin{array}{c}
 \vdots \\
 q_2 \xrightarrow{1} q_1 \xrightarrow{0} q_0
 \end{array}
 \end{array}$$

$$q_1 \xrightarrow{0} q_2 \xrightarrow{0} \emptyset$$

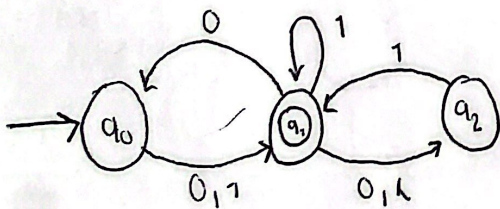
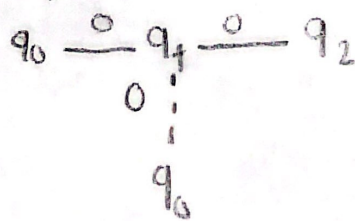
$$q_2 \xrightarrow{0} 0$$

$$\{a^i b a^j : i, j \geq 0\} \cup \{a^i b a^j : i \geq 0, j \geq 1\}$$

ab, aba, abaa

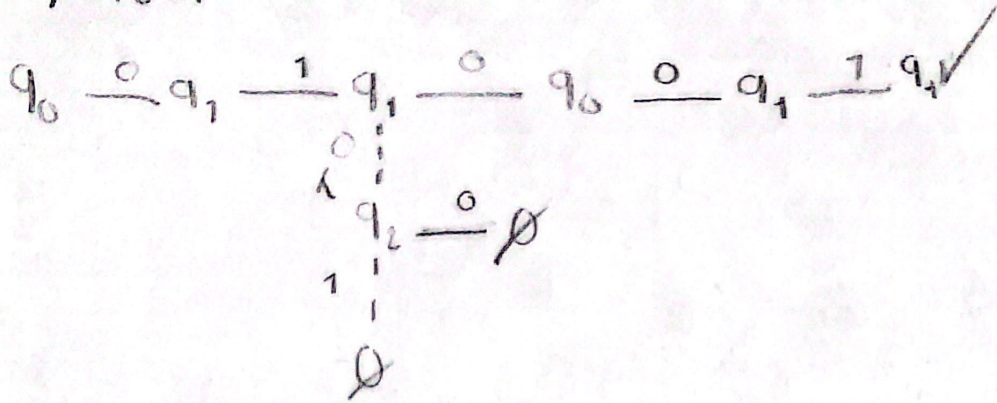

$$\{1, a, aa, \dots\}$$

a) 00

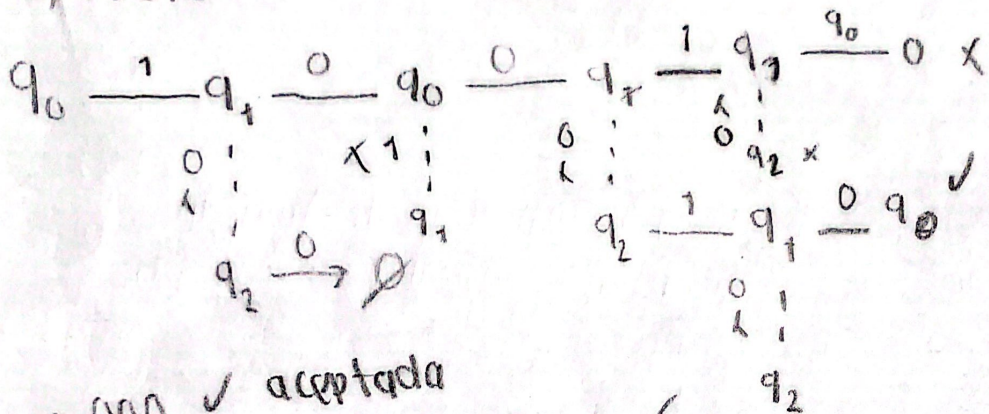




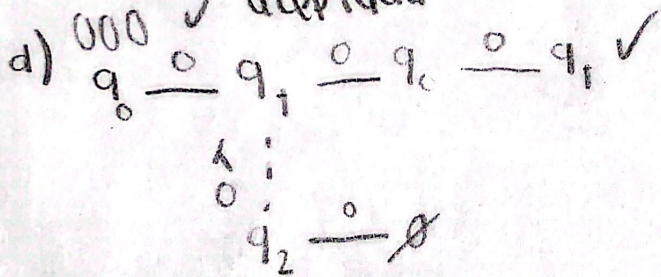
b) 01001 ✓ aceptada



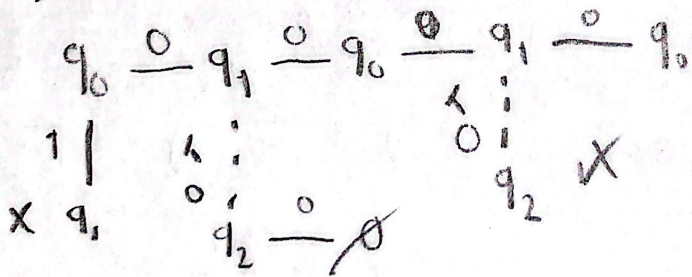
c) 10010 ✓



d) 000 ✓ aceptada  
 $q_0 \xrightarrow{0} q_1 \xrightarrow{0} q_2$



d) 0000 x

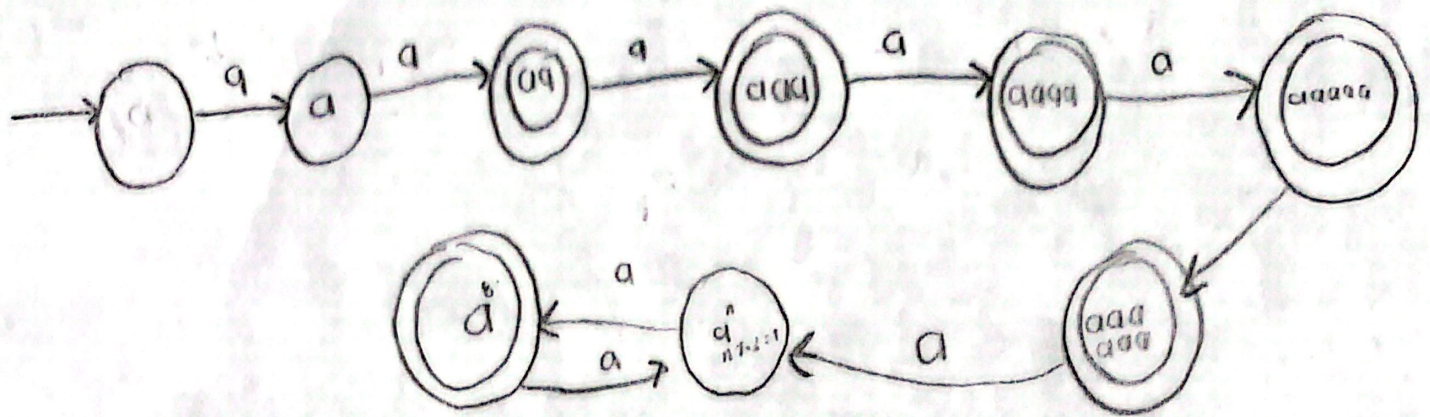




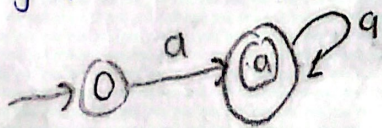
(14) Let  $L$  be the language accepted by the nfa in figure 2.8.

Find an nfa accepting  $L \cup \{a^5\}$

$$L = \{a^n : n > 0, n \text{ is par or } n=3\} \cup \{a^5\}$$



(16) Find an nfa that accepts  $\{a^*\}$  and is such that if in its transition graph a single edge is removed, the resulting automaton accepts  $\{a^3\}$



(27) An nfa in which

a) there's no  $\lambda$ -transition

b) for all  $q \in Q$  and all  $a \in \Sigma$ ,  $\delta(q, a)$  contains at most 1 element is sometimes called an incomplete dfa. This is reasonable since the condition make it such that there is only choice or move.

For  $\Sigma = \{a, b\}$ , convert the incomplete dfa into a standard dfa

