

CONFIDENTIAL

# C Programming Basic – week 8

*Gdb – Make*

*Tree*

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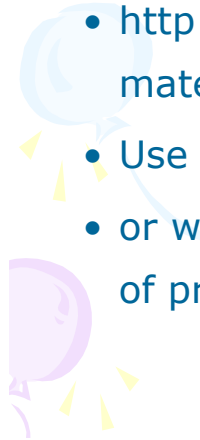
## Topics of this week

- How to use debugger tool(gdb)
- Tree data structure
  - Binary Tree
  - Binary Search Tree
- Recursive processing on Tree

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## **gdb** for debugging (1)



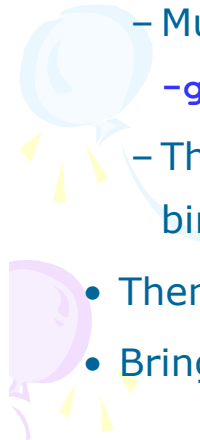
- **gdb**: the Gnu DeBugger
- <http://www.cs.caltech.edu/courses/cs11/material/c/mike/misc/gdb.html>
- Use when program core dumps
- or when want to walk through execution of program line-by-line



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## **gdb** for debugging (2)



- Before using **gdb**:
  - Must compile C code with additional flag:  
**-g**
  - This puts all the source code into the binary executable
- Then can execute as: **gdb myprogram**
- Brings up an interpreted environment

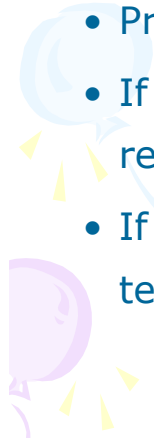


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## **gdb** for debugging (3)

**gdb**> run



- Program runs...
- If all is well, program exits successfully, returning you to prompt
- If there is (e.g.) a core dump, **gdb** will tell you and abort the program

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## **gdb** – basic commands (1)



- Stack backtrace ("**where**")
  - Your program core dumps
  - Where was the last line in the program that was executed before the core dump?
  - That's what the **where** command tells you

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## gdb – basic commands (2)

`gdb> where`

last call

last call in your code

```
#0 0x4006cb26 in free () from /lib/libc.so.6
#1 0x4006ca0d in free () from /lib/libc.so.6
#2 0x8048951 in board_updater (array=0x8049bd0,
ncells=2) at 1dCA2.c:148
#3 0x80486be in main (argc=3, argv=0xbffff7b4) at
1dCA2.c:44
#4 0x40035a52 in __libc_start_main () from
/lib/libc.so.6
```

stack backtrace

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## gdb – basic commands (3)

- Look for topmost location in stack backtrace that corresponds to your code
- Watch out for
  - freeing memory you didn't allocate
  - accessing arrays beyond their maximum elements
  - dereferencing pointers that don't point to part of a `malloc()`ed block

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## **gdb – basic commands (4)**

- **break**, **continue**, **next**, **step** commands
- **break** causes execution to stop on a given line  

```
gdb> break foo.c: 100
```

 (setting a breakpoint)
- **continue** resumes execution from that point
- **next** executes the next line, then stops
- **step** executes the next statement
  - goes into functions if necessary (**next** doesn't)

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## **gdb – basic commands (5)**

- **print** and **display** commands
- **print** prints the value of any program expression  

```
gdb> print i
```

```
$1 = 100
```
- **display** prints a particular value every time execution stops  

```
gdb> display i
```

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## gdb – printing arrays (1)

- `print` will print arrays as well

```
int arr[] = { 1, 2, 3 };
```



```
gdb> print arr
```

```
$1 = {1, 2, 3}
```

- N.B. the `$1` is just a name for the result



```
print $1
```

```
$2 = {1, 2, 3}
```

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## gdb – printing arrays (2)

- `print` has problems with dynamically-allocated arrays



```
int *arr;
```

```
arr = (int *)malloc(3 * sizeof(int));
```

```
arr[0] = 1; arr[1] = 2; arr[2] = 3;
```



```
gdb> print arr
```

```
$1 = (int *) 0x8094610
```

- Not very useful...

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## **gdb** – printing arrays (3)

- Can print this array by using `@` (**gdb** special syntax)

```
int *arr;  
arr = (int *)malloc(3 * sizeof(int));  
arr[0] = 1; arr[1] = 2; arr[2] = 3;  
  
gdb> print *arr@3  
$2 = {1, 2, 3}
```

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## **gdb** – abbreviations

- Common **gdb** commands have abbreviations

**p** (same as **print**)

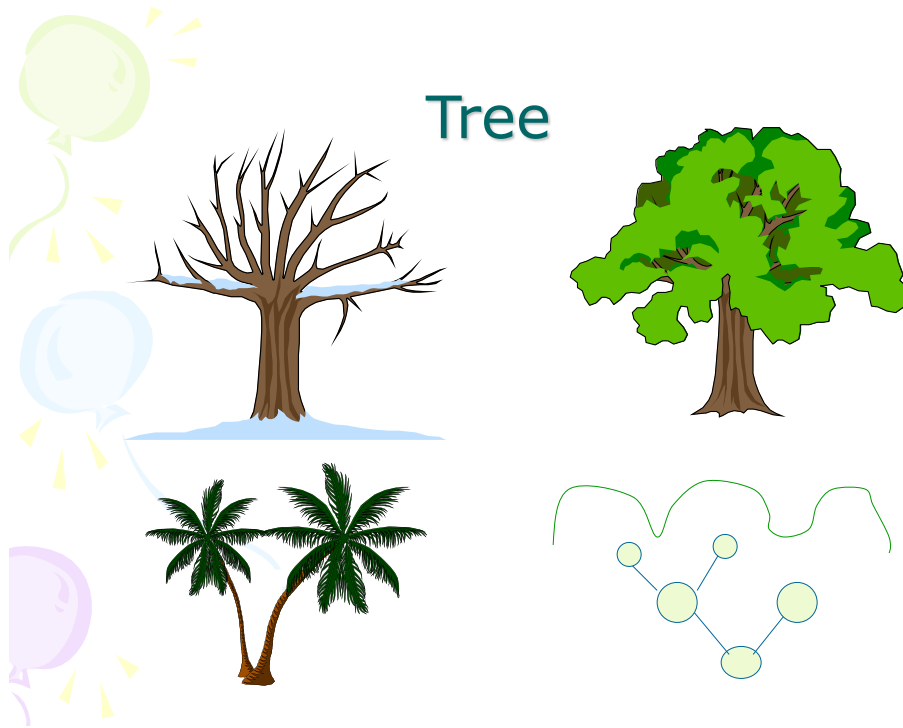
**c** (same as **continue**)

**n** (same as **next**)

**s** (same as **step**)

- More convenient to use when interactively debugging

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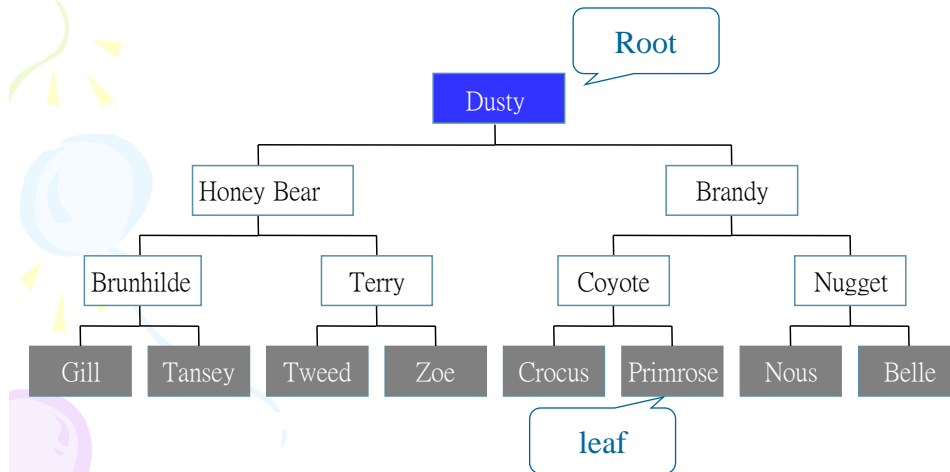
## Trees, Binary Trees, and Binary Search Trees

- Linked lists are **linear structures** and it is difficult to use them to organize a **hierarchical** representation of objects.
- Although stacks and queues reflect some hierarchy, they are limited to only **one dimension**.
- To overcome this limitation, we create a new data type called a **tree** that consists of **nodes** and **arcs**. Unlike natural trees, these trees are **depicted upside down** with the **root** at the top and the **leaves** at the bottom.

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# Family Tree



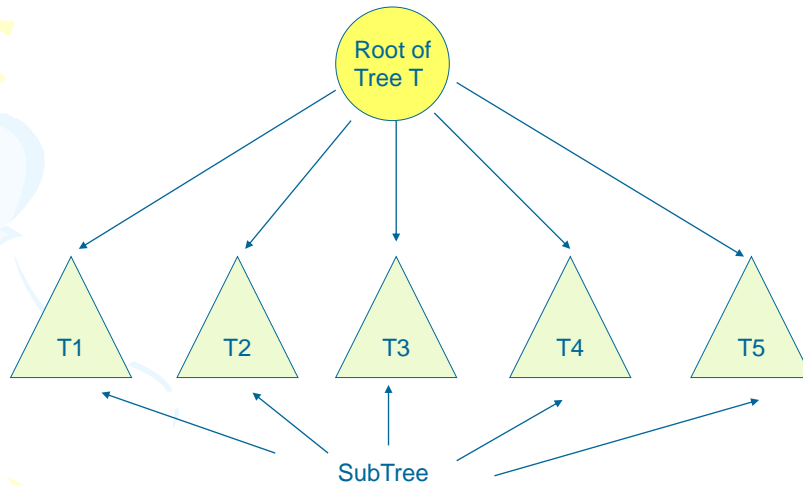
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## Definition of tree

- A tree is a finite set of one or more nodes such that:
- There is a specially designated node called the root.
- The remaining nodes are partitioned into  $n \geq 0$  disjoint sets  $T_1, \dots, T_n$ , where each of these sets is a tree.
- We call  $T_1, \dots, T_n$  the subtrees of the root.

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## Recursive definition



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## Binary Tree

- A binary tree is a tree in which no node can have more than two children.
- Each node has 0, 1, or 2 children

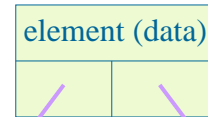
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## Linked Representation

- Each tree node is represented as an object whose data type is
- The space required by an  $n$  node binary tree is  $n \times$  (space required by one node)

```
typedef ... elmType;  
//whatever type of element  
typedef struct nodeType {  
    elmType element;  
    struct nodeType *left, *right;  
};  
typedef struct nodeType *treetype;
```



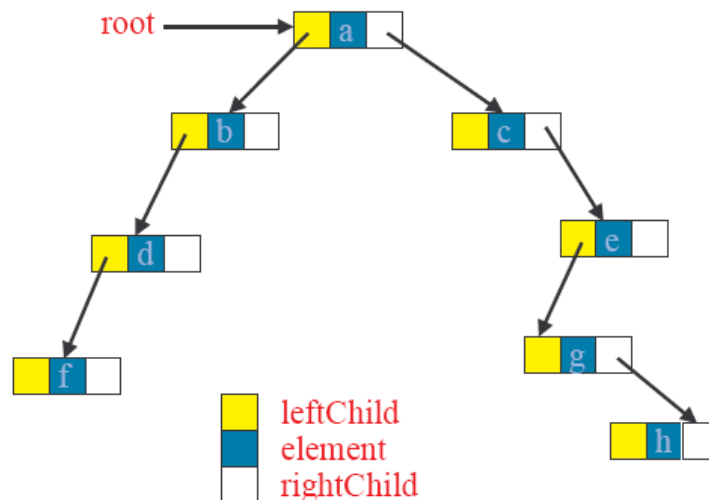
left child   right child



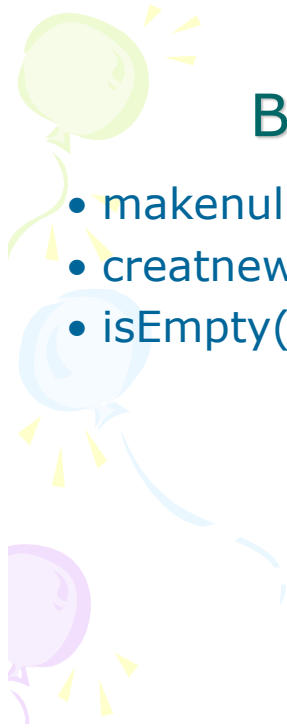
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## A linked binary tree



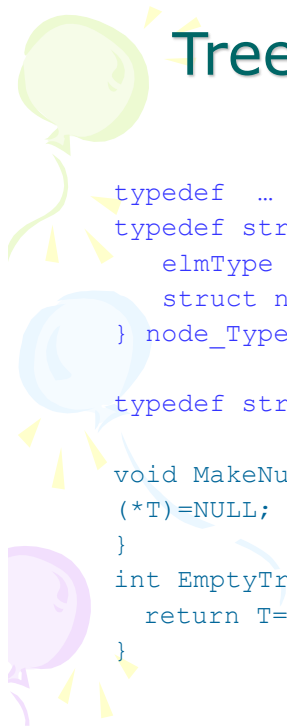
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## Binary Tree ADT

- `makenullTree(treetype *t)`
- `creatnewNode()`
- `isEmpty()`

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## Tree initialization and verification

```
typedef ... elmType;
typedef struct nodeType {
    elmType element;
    struct nodeType *left, *right;
} node_Type;

typedef struct nodeType *treetype;

void MakeNullTree(treetype *T) {
    (*T)=NULL;
}

int EmptyTree(treetype T) {
    return T==NULL;
}
```

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## Access left and right child

```
treetype LeftChild(treetype n)
{
    if (n!=NULL) return n->left;
    else return NULL;
}

treetype RightChild(treetype n)
{
    if (n!=NULL) return n->right;
    else return NULL;
}
```

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## create a new node

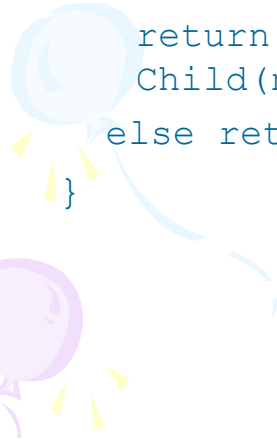
```
node_type *create_node(elmttype NewData)
{
    N=(node_type*)malloc(sizeof(node_type));
    if (N != NULL)
    {
        N->left = NULL;
        N->right = NULL;
        N->element = NewData;
    }
    return N;
}
```

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## check if a node is a leaf

```
int IsLeaf(treetype n) {  
    if (n!=NULL)  
        return (LeftChild(n)==NULL) && (Right  
            Child(n)==NULL);  
    else return -1;  
}
```



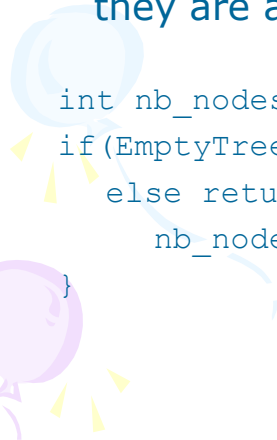
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## Recursive processing: Number of nodes

- As tree is a recursive data structure, recursive algorithms are useful when they are applied on tree.

```
int nb_nodes(treetype T) {  
    if (EmptyTree(T)) return 0;  
    else return 1+nb_nodes(LeftChild(T))+  
        nb_nodes(RightChild(T));  
}
```



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## Create a tree from two sub-trees

```
treetype createfrom2(elmtype v,
treetype l, treetype r){
    treetype N;
    N=(node_type*)malloc(sizeof(node_type));
    N->element=v;
    N->left=l;
    N->right=r;
    return N;
}
```

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## Adding a new node to the left most position

```
treetype Add_Left(treetype *Tree, elmtype NewData){
    node_type *NewNode = Create_Node(NewData);
    if (NewNode == NULL) return (NewNode);
    if (*Tree == NULL)
        *Tree = NewNode;
    else{
        node_type *Lnode = *Tree;
        while (Lnode->left != NULL)
            Lnode = Lnode->left;
        Lnode->left = NewNode;
    }
    return (NewNode);
}
```

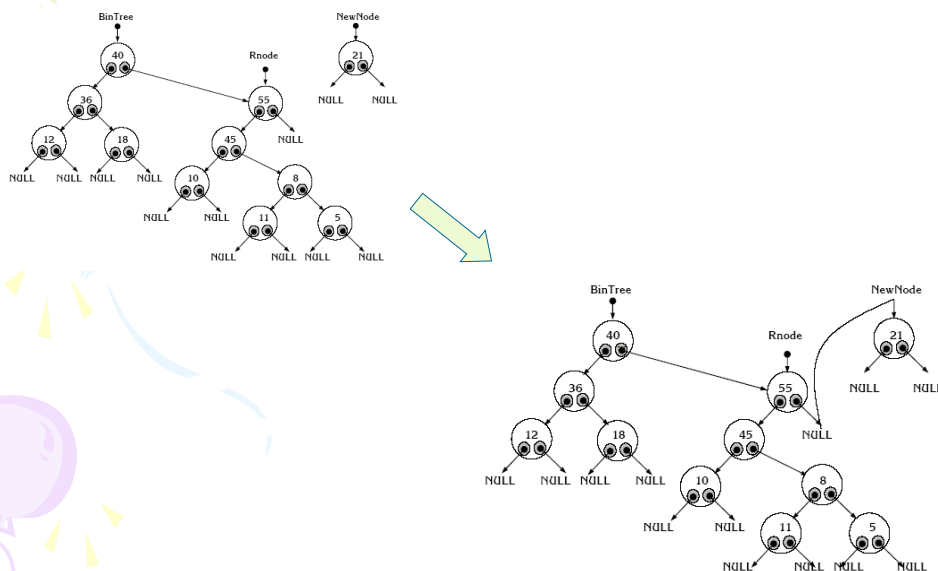
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## Adding a new node to the right most position

```
treetype Add_Left(treetype *Tree, elmttype NewData){
    node_type *NewNode = Create_Node(NewData);
    if (NewNode == NULL) return (NewNode);
    if (*Tree == NULL)
        *Tree = NewNode;
    else{
        node_type *Rnode = *Tree;
        while (Rnode->right != NULL)
            Rnode = Rnode->right;
        Rnode->right = NewNode;
    }
    return (NewNode);
}
```

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## Illustration



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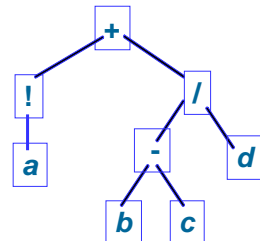
## Exercise

- Develop the following helper functions for a tree:
  - return the height of a binary tree.
  - return the number of leafs
  - return the number of internal nodes
  - count the number of right children.

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## Exercise

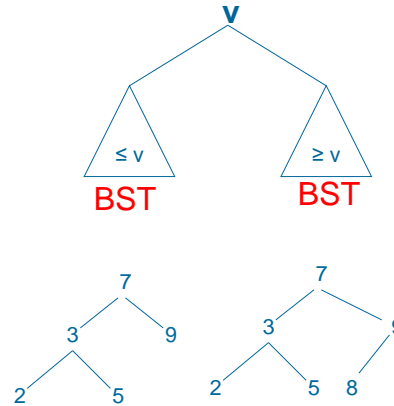
- A binary tree can represent an arithmetic expression: The leaves are operands and the other nodes are operators.
- The left and right subtrees of an operator node represent **subexpressions** that must be evaluated **before** applying the operator at the root of the subtree.
- For example  
 $a + (b - c)/d$
- Write a program create a tree representing this expression



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## Binary Search Tree

- Every element has a unique key.
- The keys in a nonempty left subtree (right subtree) are smaller (larger) than the key in the root of subtree.
- The left and right subtrees are also binary search trees.



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## Binary Search Tree Implementation

```
#include <stdio.h>
#include <stdlib.h>
typedef . . . KeyType; // specify a type for
the data
typedef struct Node{
    KeyType key;
    struct Node* left, right;
} NodeType;
typedef Node* TreeType;
```

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## Search on BST

```
TreeType Search(KeyType x, TreeType Root) {
    if (Root == NULL) return NULL; // not found
    else if (Root->key == x) /* found x */
        return Root;
    else if (Root->key < x)
        //continue searching in the right sub tree
        return Search(x, Root->right);
    else {
        // continue searching in the left sub tree
        return Search(x, Root->left);
    }
}
```

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## Insert a node from a BST

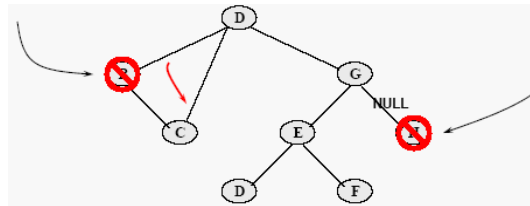
- In a binary, there are not two nodes with the same key.

```
void InsertNode(KeyType x, TreeType *Root ) {
    if (*Root == NULL) {
        /* Create a new node for key x */
        *Root = (NodeType*) malloc(sizeof(NodeType));
        (*Root)->key = x;
        (*Root)->left = NULL;
        (*Root)->right = NULL;
    }
    else if (x < (*Root)->key) InsertNode(x, &(*Root)->left);
    else if (x > (*Root)->key) InsertNode(x, &(*Root)->right);
}
```

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## Delete a node from a BST

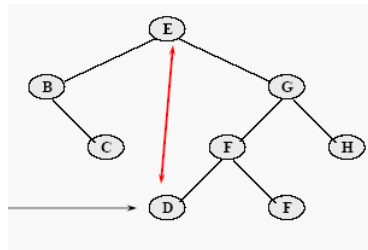
- Removing a leaf node is trivial, just set the relevant child pointer in the parent node to NULL.
- Removing an internal node which has only one subtree is also trivial, just set the relevant child pointer in the parent node to target the root of the subtree.



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## Delete a node from a BST

- Removing an internal node which has two subtrees is more complex
  - Find the left-most node of the right subtree, and then swap data values between it and the targeted node.
  - Delete the swapped value from the right subtree.



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## Find the left-most node of right sub tree

- This function find the leftmost node then delete it.

```
KeyType DeleteMin (TreeType *Root ){
    KeyType k;
    if ((*Root)->left == NULL){
        k=(*Root)->key;
        (*Root) = (*Root)->right;
        return k;
    }
    else return DeleteMin(&(*Root)->left);
}
```

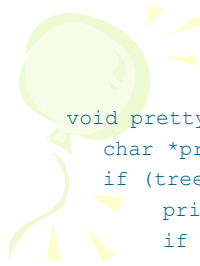
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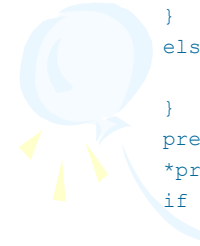
## Delete a node from a BST

```
void DeleteNode(key X,TreeType *Root){
    if (*Root!=NULL)
        if (x < (*Root)->Key) DeleteNode(x, &(*Root)->
        >left)
        else if (x > (*Root)->Key)
            DeleteNode(x, &(*Root)->right)
        else if
            ((*Root)->left==NULL) && ((*Root)->right==NULL)
                *Root=NULL;
        else if ((*Root)->left == NULL)
            *Root = (*Root)->right
        else if ((*Root)->right==NULL)
            *Root = (*Root)->left
        else (*Root)->Key = DeleteMin(&(*Root)->right);
}
```

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## Pretty print a BST



```
void prettyprint(TreeType tree, char *prefix) {
    char *prefixend = prefix + strlen(prefix);
    if (tree != NULL) {
        printf("%04d", tree->key);
        if (tree->left != NULL) if (tree->right == NULL) {
            printf("\304"); strcat(prefix, "    ");
        }
        else {
            printf("\302"); strcat(prefix, "\263    ");
        }
        prettyprint(tree->left, prefix);
        *prefixend = '\0';
        if (tree->right != NULL) if (tree->left != NULL) {
            printf("\n%s", prefix); printf("\300");
        } else printf("\304");
        strcat(prefix, "    ");
        prettyprint(tree->right, prefix);
    }
}
```

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## Exercise

- Write a function to delete all node of a tree. This function must be called before terminating program.

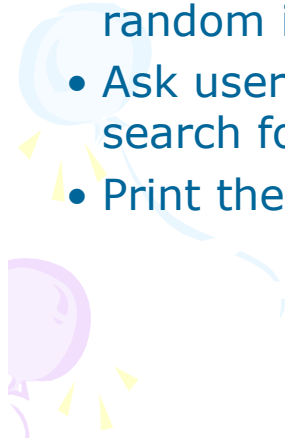


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## Exercise

- Create an binary search tree with 10 nodes. Each node contains an random integer.
- Ask user to input an number and search for it.
- Print the content of the trees.



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## Exercise

- We assume that you make a mobile phone's address book.
- Declare a structure which can store at least "name", "telephone number", "e-mail address".
- Declare a structure for a binary tree which can stores the structure of an address book inside. Read data of about 10 from an input file to this binary tree as the following rules.
  - An address data which is smaller in the dictionary order for the e-mail address is stored to the left side of a node.
  - An address data which is larger in the dictionary order for the e-mail address is stored to the right side of a node.
- (1) Confirm the address data is organized in the binary tree structure with some methods (printing, debugger, etc).
- (2) Find a specified e-mail address in the binary tree and output it to a file if found.
- (3) Output all the data stored in the binary tree in ascending order for the e-mail address. (Reserve it for the next week)



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