

# Hammering Away A User's Guide to Sledgehammer for Isabelle/HOL

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# Contents

1	Introduction	2	
2	Installation	3	
3	First Steps		
4	4 Hints 4.1 Presimplify the goal		
	4.1 Trestmptify the goal	7	
	1.2 Mand sure B, SI 1188, Varieties, and 28 are recally instance	•	

	4.3	Familiarize yourself with the main options	7
5	Free	quently Asked Questions	8
	5.1	Which facts are passed to the automatic provers?	8
	5.2	Why does Metis fail to reconstruct the proof?	9
	5.3	How can I tell whether a suggested proof is sound?	10
	5.4	What are the full_types, no_types, and mono_tags arguments	
		to Metis?	10
	5.5	And what are the lifting and hide_lams arguments to Metis? .	11
	5.6	Are generated proofs minimal?	11
	5.7	A strange error occurred—what should I do?	12
	5.8	Auto can solve it—why not Sledgehammer?	12
	5.9	Why are there so many options?	12
	0.5	wing are mere so many opiions:	14
6	Cor	nmand Syntax	13
	6.1	·	13
	6.2	Metis	15
	• -		
7	Opt	tion Reference	15
	$7.\bar{1}$	Mode of Operation	16
	7.2		21
	7.3	Problem Encoding	22
	7.4	Output Format	$\frac{-}{25}$
	7.5	Authentication	26
	7.6	Timeouts	26

# 1 Introduction

Sledgehammer is a tool that applies automatic theorem provers (ATPs) and satisfiability-modulo-theories (SMT) solvers on the current goal.<sup>1</sup> The supported ATPs are agsyHOL [12], Alt-Ergo [3], E [14], E-SInE [9], E-ToFoF [16], iProver [10], iProver-Eq [11], LEO-II [2], Satallax [6], SNARK [15], SPASS [18], Vampire [13], and Waldmeister [8]. The ATPs are run either locally or remotely via the SystemOnTPTP web service [17]. In addition to the ATPs, a selection of the SMT solvers CVC3 [1], Yices [7], and Z3 [19]

<sup>&</sup>lt;sup>1</sup>The distinction between ATPs and SMT solvers is convenient but mostly historical. The two communities are converging, with more and more ATPs supporting typical SMT features such as arithmetic and sorts, and a few SMT solvers parsing ATP syntaxes. There is also a strong technological connection between instantiation-based ATPs (such as iProver and iProver-Eq) and SMT solvers.

are run by default; these are run either locally or (for CVC3 and Z3) on a server at the TU München.

The problem passed to the automatic provers consists of your current goal together with a heuristic selection of hundreds of facts (theorems) from the current theory context, filtered by relevance.

The result of a successful proof search is some source text that usually (but not always) reconstructs the proof within Isabelle. For ATPs, the reconstructed proof relies on the general-purpose *metis* proof method, which integrates the Metis ATP in Isabelle/HOL with explicit inferences going through the kernel. Thus its results are correct by construction.

For Isabelle/jEdit users, Sledgehammer provides an automatic mode that can be enabled via the "Auto Sledgehammer" option under "Plugins > Plugin Options > Isabelle > General." In this mode, Sledgehammer is run on every newly entered theorem.

To run Sledgehammer, you must make sure that the theory *Sledgehammer* is imported—this is rarely a problem in practice since it is part of *Main*. Examples of Sledgehammer use can be found in Isabelle's src/HOL/Metis\_Examples directory. Comments and bug reports concerning Sledgehammer or this manual should be directed to the author at blanchette@in.tum.de.

# 2 Installation

Sledgehammer is part of Isabelle, so you do not need to install it. However, it relies on third-party automatic provers (ATPs and SMT solvers).

Among the ATPs, agsyHOL, Alt-Ergo, E, LEO-II, Satallax, SPASS, and Vampire can be run locally; in addition, agsyHOL, E, E-SInE, E-ToFoF, iProver, iProver-Eq, LEO-II, Satallax, SNARK, Vampire, and Waldmeister are available remotely via SystemOnTPTP [17]. If you want better performance, you should at least install E and SPASS locally.

The SMT solvers CVC3, Yices, and Z3 can be run locally, and CVC3 and Z3 can be run remotely on a TU München server. If you want better performance and get the ability to replay proofs that rely on the *smt* proof method without an Internet connection, you should at least have Z3 locally installed.

There are three main ways to install automatic provers on your machine:

- If you installed an official Isabelle package, it should already include properly setup executables for CVC3, E, SPASS, and Z3, ready to use.<sup>2</sup> For Z3, you must additionally set the variable Z3\_NON\_COMMERCIAL to "yes" to confirm that you are a noncommercial user, either in the environment in which Isabelle is launched or in your \$ISABELLE\_HOME\_USER/etc/settings file.
- Alternatively, you can download the Isabelle-aware CVC3, E, SPASS, and Z3 binary packages from http://isabelle.in.tum.de/components/. Extract the archives, then add a line to your \$ISABELLE\_HOME\_USER/etc/components³ file with the absolute path to CVC3, E, SPASS, or Z3. For example, if the components file does not exist yet and you extracted SPASS to /usr/local/spass-3.8ds, create it with the single line

/usr/local/spass-3.8ds

in it.

• If you prefer to build agsyHOL, Alt-Ergo, E, LEO-II, Satallax, or SPASS manually, or found a Vampire executable somewhere (e.g., http://www.vprover.org/), set the environment variable AGSYHOL\_HOME, E\_HOME, LEO2\_HOME, SATALLAX\_HOME, SPASS\_HOME, or VAMPIRE\_HOME to the directory that contains the agsyHOL, eprover (and/or eproof or eproof\_ram), leo, satallax, SPASS, or vampire executable; for Alt-Ergo, set the environment variable WHY3\_HOME to the directory that contains the why3 executable. Sledgehammer has been tested with agsyHOL 1.0, Alt-Ergo 0.95.1, E 1.0 to 1.8, LEO-II 1.3.4, Satallax 2.2 to 2.7, SPASS 3.8ds, and Vampire 0.6 to 3.0.4Since the ATPs' output formats are neither documented nor stable, other versions might not work well with Sledgehammer. Ideally, you should also set E\_VERSION, LEO2\_VERSION, SATALLAX\_VERSION, SPASS\_VERSION, or VAMPIRE\_VERSION to the prover's version number (e.g., "1.8").

Similarly, if you want to build CVC3, or found a Yices or Z3 executable somewhere (e.g., http://yices.csl.sri.com/download.shtml or http://research.microsoft.com/en-us/um/redmond/projects/

<sup>&</sup>lt;sup>2</sup>Vampire's and Yices's licenses prevent us from doing the same for these otherwise remarkable tools.

<sup>&</sup>lt;sup>3</sup>The variable \$ISABELLE\_HOME\_USER is set by Isabelle at startup. Its value can be retrieved by executing isabelle getenv ISABELLE\_HOME\_USER on the command line.

<sup>&</sup>lt;sup>4</sup>Following the rewrite of Vampire, the counter for version numbers was reset to 0; hence the (new) Vampire versions 0.6, 1.0, 1.8, 2.6, and 3.0 are more recent than 9.0 or 11.5.

z3/download.html), set the environment variable CVC3\_SOLVER, YICES\_SOLVER, or Z3\_SOLVER to the complete path of the executable, including the file name. Sledgehammer has been tested with CVC3 2.2 and 2.4.1, Yices 1.0.28 and 1.0.33, and Z3 3.0 to 4.0. Since the SMT solvers' output formats are somewhat unstable, other versions of the solvers might not work well with Sledgehammer. Ideally, also set CVC3\_VERSION, YICES\_VERSION, or Z3\_VERSION to the solver's version number (e.g., "4.0").

To check whether E, SPASS, Vampire, and/or Z3 are successfully installed, try out the example in §3. If the remote versions of any of these provers is used (identified by the prefix "remote\_"), or if the local versions fail to solve the easy goal presented there, something must be wrong with the installation.

Remote prover invocation requires Perl with the World Wide Web Library (libwww-perl) installed. If you must use a proxy server to access the Internet, set the http\_proxy environment variable to the proxy, either in the environment in which Isabelle is launched or in your \$ISABELLE\_HOME\_USER/etc/settings file. Here are a few examples:

```
http_proxy=http://proxy.example.org
http_proxy=http://proxy.example.org:8080
http_proxy=http://joeblow:pAsSwRd@proxy.example.org
```

# 3 First Steps

To illustrate Sledgehammer in context, let us start a theory file and attempt to prove a simple lemma:

```
theory Scratch
imports Main
begin
lemma "[a] = [b] \implies a = b"
sledgehammer
```

Instead of issuing the **sledgehammer** command, you can also use the Sledgehammer panel in Isabelle/jEdit. Sledgehammer produces the following output after a few seconds:

```
Sledgehammer: "e" on goal [a] = [b] \implies a = b
```

```
Try this: by (metis last_ConsL) (64 ms).
```

$$[a] = [b] \implies a = b$$

Try this: **by** (metis list.inject) (20 ms).

$$[a] = [b] \implies a = b$$

Try this: **by** (metis hd.simps) (14 ms).

Sledgehammer: "spass" on goal

$$[a] = [b] \implies a = b$$

Try this: **by** (metis list.inject) (17 ms).

Sledgehammer: "remote\_e\_sine" on goal

$$[a] = [b] \implies a = b$$

Try this: by (metis hd.simps) (18 ms).

Sledgehammer ran E, E-SInE, SPASS, Vampire, and Z3 in parallel. Depending on which provers are installed and how many processor cores are available, some of the provers might be missing or present with a *remote\_* prefix. Wald-meister is run only for unit equational problems, where the goal's conclusion is a (universally quantified) equation.

For each successful prover, Sledgehammer gives a one-liner *metis* or *smt* method call. Rough timings are shown in parentheses, indicating how fast the call is. You can click the proof to insert it into the theory text.

In addition, you can ask Sledgehammer for an Isar text proof by enabling the *isar\_proofs* option (§7.4):

#### sledgehammer [isar\_proofs]

When Isar proof construction is successful, it can yield proofs that are more readable and also faster than the *metis* or *smt* one-liners. This feature is experimental and is only available for ATPs.

#### 4 Hints

This section presents a few hints that should help you get the most out of Sledgehammer. Frequently asked questions are answered in §5.

# 4.1 Presimplify the goal

For best results, first simplify your problem by calling *auto* or at least *safe* followed by  $simp\_all$ . The SMT solvers provide arithmetic decision procedures, but the ATPs typically do not (or if they do, Sledgehammer does not use it yet). Apart from Waldmeister, they are not particularly good at heavy rewriting, but because they regard equations as undirected, they often prove theorems that require the reverse orientation of a *simp* rule. Higher-order problems can be tackled, but the success rate is better for first-order problems. Hence, you may get better results if you first simplify the problem to remove higher-order features.

# 4.2 Make sure E, SPASS, Vampire, and Z3 are locally installed

Locally installed provers are faster and more reliable than those running on servers. See §2 for details on how to install them.

# 4.3 Familiarize yourself with the main options

Sledgehammer's options are fully documented in §6. Many of the options are very specialized, but serious users of the tool should at least familiarize themselves with the following options:

- provers (§7.1) specifies the automatic provers (ATPs and SMT solvers) that should be run whenever Sledgehammer is invoked (e.g., "provers = e spass remote\_vampire"). For convenience, you can omit "provers =" and simply write the prover names as a space-separated list (e.g., "e spass remote\_vampire").
- max\_facts (§7.2) specifies the maximum number of facts that should be passed to the provers. By default, the value is prover-dependent but varies between about 50 and 1000. If the provers time out, you can try lowering this value to, say, 25 or 50 and see if that helps.
- *isar\_proofs* (§7.4) specifies that Isar proofs should be generated, in addition to one-liner *metis* or *smt* proofs. The length of the Isar proofs can be controlled by setting *isar\_compress* (§7.4).

• *timeout* (§7.6) controls the provers' time limit. It is set to 30 seconds, but since Sledgehammer runs asynchronously you should not hesitate to raise this limit to 60 or 120 seconds if you are the kind of user who can think clearly while ATPs are active.

Options can be set globally using **sledgehammer\_params** (§6). The command also prints the list of all available options with their current value. Fact selection can be influenced by specifying " $(add: my\_facts)$ " after the **sledgehammer** call to ensure that certain facts are included, or simply " $(my\_facts)$ " to force Sledgehammer to run only with  $my\_facts$ .

# 5 Frequently Asked Questions

This sections answers frequently (and infrequently) asked questions about Sledgehammer. It is a good idea to skim over it now even if you do not have any questions at this stage. And if you have any further questions not listed here, send them to the author at blanchette@in.tum.de.

# 5.1 Which facts are passed to the automatic provers?

Sledgehammer heuristically selects a few hundred relevant lemmas from the currently loaded libraries. The component that performs this selection is called *relevance filter*.

- The traditional relevance filter, called MePo (Meng-Paulson), assigns a score to every available fact (lemma, theorem, definition, or axiom) based upon how many constants that fact shares with the conjecture. This process iterates to include facts relevant to those just accepted. The constants are weighted to give unusual ones greater significance. MePo copes best when the conjecture contains some unusual constants; if all the constants are common, it is unable to discriminate among the hundreds of facts that are picked up. The filter is also memoryless: It has no information about how many times a particular fact has been used in a proof, and it cannot learn.
- An experimental alternative to MePo is *MaSh* (<u>Machine Learner for Sledgehammer</u>). It relies on an external Python tool that applies machine learning to the problem of finding relevant facts.
- The MeSh filter combines MePo and MaSh.

The default is either MePo or MeSh, depending on whether the environment variable MASH is set and what class of provers the target prover belongs to (§7.2).

The number of facts included in a problem varies from prover to prover, since some provers get overwhelmed more easily than others. You can show the number of facts given using the *verbose* option ( $\S7.4$ ) and the actual facts using debug ( $\S7.4$ ).

Sledgehammer is good at finding short proofs combining a handful of existing lemmas. If you are looking for longer proofs, you must typically restrict the number of facts, by setting the max-facts option (§7.2) to, say, 25 or 50.

You can also influence which facts are actually selected in a number of ways. If you simply want to ensure that a fact is included, you can specify it using the "(add: my\_facts)" syntax. For example:

```
sledgehammer (add: hd.simps tl.simps)
```

The specified facts then replace the least relevant facts that would otherwise be included; the other selected facts remain the same. If you want to direct the selection in a particular direction, you can specify the facts via **using**:

```
using hd.simps tl.simps
sledgehammer
```

The facts are then more likely to be selected than otherwise, and if they are selected at iteration j they also influence which facts are selected at iterations j + 1, j + 2, etc. To give them even more weight, try

```
using hd.simps tl.simps
apply –
sledgehammer
```

# 5.2 Why does Metis fail to reconstruct the proof?

There are many reasons. If Metis runs seemingly forever, that is a sign that the proof is too difficult for it. Metis's search is complete, so it should eventually find it, but that's little consolation. There are several possible solutions:

• Try the *isar\_proofs* option (§7.4) to obtain a step-by-step Isar proof where each step is justified by *metis*. Since the steps are fairly small, *metis* is more likely to be able to replay them.

- Try the *smt* proof method instead of *metis*. It is usually stronger, but you need to either have Z3 available to replay the proofs, trust the SMT solver, or use certificates. See the documentation in the *SMT* theory (\$ISABELLE\_HOME/src/HOL/SMT.thy) for details.
- Try the *blast* or *auto* proof methods, passing the necessary facts via **unfolding**, **using**, *intro*:, *elim*:, *dest*:, or *simp*:, as appropriate.

In some rare cases, *metis* fails fairly quickly, and you get the error message

One-line proof reconstruction failed.

This message indicates that Sledgehammer determined that the goal is provable, but the proof is, for technical reasons, beyond metis's power. You can then try again with the strict option (§7.3).

If the goal is actually unprovable and you did not specify an unsound encoding using  $type\_enc$  (§7.3), this is a bug, and you are strongly encouraged to report this to the author at blanchette@in.tum.de.

# 5.3 How can I tell whether a suggested proof is sound?

Earlier versions of Sledgehammer often suggested unsound proofs—either proofs of nontheorems or simply proofs that rely on type-unsound inferences. This is a thing of the past, unless you explicitly specify an unsound encoding using type\_enc (§7.3). Officially, the only form of "unsoundness" that lurks in the sound encodings is related to missing characteristic theorems of datatypes. For example,

```
lemma "\exists xs. \ xs \neq []" sledgehammer ()
```

suggests an argumentless metis call that fails. However, the conjecture does actually hold, and the metis call can be repaired by adding list.distinct. We hope to address this problem in a future version of Isabelle. In the meantime, you can avoid it by passing the strict option (§7.3).

# 5.4 What are the full\_types, no\_types, and mono\_tags arguments to Metis?

The metis (full\_types) proof method and its cousin metis (mono\_tags) are fully-typed versions of Metis. It is somewhat slower than metis, but the

proof search is fully typed, and it also includes more powerful rules such as the axiom " $x = \mathsf{True} \lor x = \mathsf{False}$ " for reasoning in higher-order places (e.g., in set comprehensions). The method kicks in automatically as a fallback when *metis* fails, and it is sometimes generated by Sledgehammer instead of *metis* if the proof obviously requires type information or if *metis* failed when Sledgehammer preplayed the proof. (By default, Sledgehammer tries to run *metis* with various options for up to 3 seconds each time to ensure that the generated one-line proofs actually work and to display timing information. This can be configured using the *preplay\_timeout* and *dont\_preplay* options (§7.6).) At the other end of the soundness spectrum, *metis* (*no\_types*) uses no type information at all during the proof search, which is more efficient but often fails. Calls to *metis* (*no\_types*) are occasionally generated by Sledgehammer. See the *type\_enc* option (§7.3) for details.

Incidentally, if you ever see warnings such as

Metis: Falling back on "metis (full\_types)".

for a successful *metis* proof, you can advantageously pass the *full\_types* option to *metis* directly.

# 5.5 And what are the lifting and hide\_lams arguments to Metis?

Orthogonally to the encoding of types, it is important to choose an appropriate translation of  $\lambda$ -abstractions. Metis supports three translation schemes, in decreasing order of power: Curry combinators (the default),  $\lambda$ -lifting, and a "hiding" scheme that disables all reasoning under  $\lambda$ -abstractions. The more powerful schemes also give the automatic provers more rope to hang themselves. See the  $lam_{-}trans$  option (§7.3) for details.

# 5.6 Are generated proofs minimal?

Automatic provers frequently use many more facts than are necessary. Sledge-hammer inclues a minimization tool that takes a set of facts returned by a given prover and repeatedly calls the same prover, *metis*, or *smt* with subsets of those axioms in order to find a minimal set. Reducing the number of axioms typically improves Metis's speed and success rate, while also removing superfluous clutter from the proof scripts.

In earlier versions of Sledgehammer, generated proofs were systematically accompanied by a suggestion to invoke the minimization tool. This step is now performed implicitly if it can be done in a reasonable amount of time (something that can be guessed from the number of facts in the original proof and the time it took to find or preplay it).

In addition, some provers (e.g., Yices) do not provide proofs or sometimes produce incomplete proofs. The minimizer is then invoked to find out which facts are actually needed from the (large) set of facts that was initially given to the prover. Finally, if a prover returns a proof with lots of facts, the minimizer is invoked automatically since Metis would be unlikely to re-find the proof. Automatic minimization can be forced or disabled using the *minimize* option (§7.1).

# 5.7 A strange error occurred—what should I do?

Sledgehammer tries to give informative error messages. Please report any strange error to the author at blanchette@in.tum.de. This applies doubly if you get the message

The prover derived "False" using "foo", "bar", and "baz". This could be due to inconsistent axioms (including "sorry"s) or to a bug in Sledge-hammer. If the problem persists, please contact the Isabelle developers.

# 5.8 Auto can solve it—why not Sledgehammer?

Problems can be easy for *auto* and difficult for automatic provers, but the reverse is also true, so do not be discouraged if your first attempts fail. Because the system refers to all theorems known to Isabelle, it is particularly suitable when your goal has a short proof from lemmas that you do not know about.

# 5.9 Why are there so many options?

Sledgehammer's philosophy should work out of the box, without user guidance. Many of the options are meant to be used mostly by the Sledgehammer developers for experiments. Of course, feel free to try them out if you are so inclined.

# 6 Command Syntax

# 6.1 Sledgehammer

Sledgehammer can be invoked at any point when there is an open goal by entering the **sledgehammer** command in the theory file. Its general syntax is as follows:

**sledgehammer**  $\langle subcommand \rangle$ ?  $\langle options \rangle$ ?  $\langle facts\_override \rangle$ ?  $\langle num \rangle$ ?

In the general syntax, the  $\langle subcommand \rangle$  may be any of the following:

- run (the default): Runs Sledgehammer on subgoal number  $\langle num \rangle$  (1 by default), with the given options and facts.
- min: Attempts to minimize the facts specified in the  $\langle facts\_override \rangle$  argument to obtain a simpler proof involving fewer facts. The options and goal number are as for run.
- messages: Redisplays recent messages issued by Sledgehammer. This allows you to examine results that might have been lost due to Sledgehammer's asynchronous nature. The  $\langle num \rangle$  argument specifies a limit on the number of messages to display (10 by default).
- supported\_provers: Prints the list of automatic provers supported by Sledgehammer. See §2 and §7.1 for more information on how to install automatic provers.
- running\_provers: Prints information about currently running automatic provers, including elapsed runtime and remaining time until timeout.
- *kill\_all*: Terminates all running threads (automatic provers and machine learners).
- refresh\_tptp: Refreshes the list of remote ATPs available at System-OnTPTP [17].

In addition, the following subcommands provide finer control over machine learning with MaSh:

- unlearn: Resets MaSh, erasing any persistent state.
- *learn\_isar*: Invokes MaSh on the current theory to process all the available facts, learning from their Isabelle/Isar proofs. This happens

automatically at Sledgehammer invocations if the learn option (§7.2) is enabled.

- learn\_prover: Invokes MaSh on the current theory to process all the available facts, learning from proofs generated by automatic provers. The prover to use and its timeout can be set using the prover (§7.1) and timeout (§7.6) options. It is recommended to perform learning using an efficient first-order ATP (such as E, SPASS, and Vampire) as opposed to a higher-order ATP or an SMT solver.
- relearn\_isar: Same as unlearn followed by learn\_isar.
- relearn\_prover: Same as unlearn followed by learn\_prover.
- running\_learners: Prints information about currently running machine learners, including elapsed runtime and remaining time until timeout.

Sledgehammer's behavior can be influenced by various  $\langle options \rangle$ , which can be specified in brackets after the **sledgehammer** command. The  $\langle options \rangle$  are a list of key-value pairs of the form " $[k_1 = v_1, \ldots, k_n = v_n]$ ". For Boolean options, "= true" is optional. For example:

 $sledgehammer [isar\_proofs, timeout = 120]$ 

Default values can be set using **sledgehammer\_params**:

sledgehammer\_params \( \langle options \rangle \)

The supported options are described in §7.

The  $\langle facts\_override \rangle$  argument lets you alter the set of facts that go through the relevance filter. It may be of the form " $(\langle facts \rangle)$ ", where  $\langle facts \rangle$  is a space-separated list of Isabelle facts (theorems, local assumptions, etc.), in which case the relevance filter is bypassed and the given facts are used. It may also be of the form " $(add: \langle facts_1 \rangle)$ ", " $(del: \langle facts_2 \rangle)$ ", or " $(add: \langle facts_1 \rangle del: \langle facts_2 \rangle)$ ", where the relevance filter is instructed to proceed as usual except that it should consider  $\langle facts_1 \rangle$  highly-relevant and  $\langle facts_2 \rangle$  fully irrelevant.

If you use Isabelle/jEdit, Sledgehammer also provides an automatic mode that can be enabled via the "Auto Sledgehammer" option under "Plugins > Plugin Options > Isabelle > General." For automatic runs, only the first prover set using provers (§7.1) is considered, fewer facts are passed to the prover, slice (§7.1) is disabled, strict (§7.3) is enabled, verbose (§7.4) and debug (§7.4) are disabled, and timeout (§7.6) is superseded by the "Auto Time Limit" option in jEdit. Sledgehammer's output is also more concise.

#### 6.2 Metis

The *metis* proof method has the syntax

```
metis (\langle options \rangle)^? \langle facts \rangle^?
```

where  $\langle facts \rangle$  is a list of arbitrary facts and  $\langle options \rangle$  is a comma-separated list consisting of at most one  $\lambda$  translation scheme specification with the same semantics as Sledgehammer's  $lam\_trans$  option (§7.3) and at most one type encoding specification with the same semantics as Sledgehammer's  $type\_enc$  option (§7.3). The supported  $\lambda$  translation schemes are  $hide\_lams$ , lifting, and combs (the default). All the untyped type encodings listed in §7.3 are supported. For convenience, the following aliases are provided:

• full\_types: Alias for poly\_guards\_query.

• partial\_types: Alias for poly\_args.

• no\_types: Alias for erased.

# 7 Option Reference

Sledgehammer's options are categorized as follows: mode of operation ( $\S7.1$ ), problem encoding ( $\S7.3$ ), relevance filter ( $\S7.2$ ), output format ( $\S7.4$ ), authentication ( $\S7.5$ ), and timeouts ( $\S7.6$ ).

The descriptions below refer to the following syntactic quantities:

- $\langle string \rangle$ : A string.
- $\langle bool \rangle$ : true or false.
- $\langle smart\_bool \rangle$ : true, false, or smart.
- $\langle int \rangle$ : An integer.
- (float\_pair): A pair of floating-point numbers (e.g., 0.6 0.95).
- $\langle smart\_int \rangle$ : An integer or smart.
- $\langle float\_or\_none \rangle$ : A floating-point number (e.g., 60 or 0.5) expressing a number of seconds, or the keyword *none* ( $\infty$  seconds).

Default values are indicated in curly brackets ( $\{\}$ ). Boolean options have a negative counterpart (e.g., blocking vs. non\_blocking). When setting Boolean options or their negative counterparts, "= true" may be omitted.

# 7.1 Mode of Operation

# $[provers =] \langle string \rangle$

Specifies the automatic provers to use as a space-separated list (e.g., "e spass remote\_vampire"). Provers can be run locally or remotely; see §2 for installation instructions.

The following local provers are supported:

- agsyHOL: agsyHOL is an automatic higher-order prover developed by Fredrik Lindblad [12], with support for the TPTP typed higher-order syntax (THF0). To use agsyHOL, set the environment variable AGSYHOL\_HOME to the directory that contains the agsyHOL executable. Sledgehammer has been tested with version 1.0.
- alt\_ergo: Alt-Ergo is a polymorphic ATP developed by Bobot et al. [3]. It supports the TPTP polymorphic typed first-order format (TFF1) via Why3 [4]. To use Alt-Ergo, set the environment variable WHY3\_HOME to the directory that contains the why3 executable. Sledgehammer has been tested with Alt-Ergo 0.95.1 and an unidentified development version of Why3.
- cvc3: CVC3 is an SMT solver developed by Clark Barrett, Cesare Tinelli, and their colleagues [1]. To use CVC3, set the environment variable CVC3\_SOLVER to the complete path of the executable, including the file name, or install the prebuilt CVC3 package from http://isabelle.in.tum.de/components/. Sledgehammer has been tested with version 2.2 and 2.4.1.
- e: E is a first-order resolution prover developed by Stephan Schulz [14]. To use E, set the environment variable E\_HOME to the directory that contains the eproof executable and E\_VERSION to the version number (e.g., "1.8"), or install the prebuilt E package from http://isabelle.in.tum.de/components/. Sledgehammer has been tested with versions 1.0 to 1.8.
- e\_males: E-MaLeS is a metaprover developed by Daniel Kühlwein that implements strategy scheduling on top of E. To use E-MaLeS, set the environment variable E\_MALES\_HOME to the directory that contains the emales.py script. Sledgehammer has been tested with version 1.1.
- *e\_par*: E-Par is a metaprover developed by Josef Urban that implements strategy scheduling on top of E. To use E-Par, set the

environment variable E\_HOME to the directory that contains the runepar.pl script and the eprover and epclextract executables, or use the prebuilt E package from http://isabelle.in.tum.de/components/.

- *iprover*: iProver is a pure instantiation-based prover developed by Konstantin Korovin [10]. To use iProver, set the environment variable IPROVER\_HOME to the directory that contains the iprover and vclausify\_rel executables. Sledgehammer has been tested with version 0.99.
- *iprover\_eq*: iProver-Eq is an instantiation-based prover with native support for equality developed by Konstantin Korovin and Christoph Sticksel [11]. To use iProver-Eq, set the environment variable IPROVER\_EQ\_HOME to the directory that contains the iprover-eq and vclausify\_rel executables. Sledgehammer has been tested with version 0.8.
- leo2: LEO-II is an automatic higher-order prover developed by Christoph Benzmüller et al. [2], with support for the TPTP typed higher-order syntax (THF0). To use LEO-II, set the environment variable LEO2\_HOME to the directory that contains the leo executable. Sledgehammer requires version 1.3.4 or above.
- *metis*: Although it is less powerful than the external provers, Metis itself can be used for proof search.
- satallax: Satallax is an automatic higher-order prover developed by Chad Brown et al. [6], with support for the TPTP typed higher-order syntax (THF0). To use Satallax, set the environment variable SATALLAX\_HOME to the directory that contains the satallax executable. Sledgehammer requires version 2.2 or above.
- smt: The smt proof method with the current settings (usually: Z3 with proof reconstruction) can be used for proof search.
- spass: SPASS is a first-order resolution prover developed by Christoph Weidenbach et al. [18]. To use SPASS, set the environment variable SPASS\_HOME to the directory that contains the SPASS executable and SPASS\_VERSION to the version number (e.g., "3.8ds"), or install the prebuilt SPASS package from http://isabelle.in.tum.de/components/. Sledgehammer requires version 3.8ds or above.
- *vampire*: Vampire is a first-order resolution prover developed by Andrei Voronkov and his colleagues [13]. To use Vampire, set

the environment variable VAMPIRE\_HOME to the directory that contains the vampire executable and VAMPIRE\_VERSION to the version number (e.g., "2.6"). Sledgehammer has been tested with versions 0.6 to 3.0. Versions strictly above 1.8 support the TPTP typed first-order format (TFF0).

- *yices*: Yices is an SMT solver developed at SRI [7]. To use Yices, set the environment variable YICES\_SOLVER to the complete path of the executable, including the file name. Sledgehammer has been tested with version 1.0.28.
- z3: Z3 is an SMT solver developed at Microsoft Research [19]. To use Z3, set the environment variable Z3\_SOLVER to the complete path of the executable, including the file name, and set Z3\_NON\_COMMERCIAL to "yes" to confirm that you are a noncommercial user. Sledgehammer has been tested with versions 3.0, 3.1, 3.2, and 4.0.

The following remote provers are supported:

- remote\_agsyhol: The remote version of agsyHOL runs on Geoff Sutcliffe's Miami servers [17].
- remote\_cvc3: The remote version of CVC3 runs on servers at the TU München (or wherever REMOTE\_SMT\_URL is set to point).
- remote\_e: The remote version of E runs on Geoff Sutcliffe's Miami servers [17].
- remote\_e\_sine: E-SInE is a metaprover developed by Kryštof Hoder [9] based on E. It runs on Geoff Sutcliffe's Miami servers.
- remote\_e\_tofof: E-ToFoF is a metaprover developed by Geoff Sutcliffe [16] based on E running on his Miami servers. This ATP supports the TPTP typed first-order format (TFF0). The remote version of E-ToFoF runs on Geoff Sutcliffe's Miami servers.
- remote\_iprover: The remote version of iProver runs on Geoff Sutcliffe's Miami servers [17].
- remote\_iprover\_eq: The remote version of iProver-Eq runs on Geoff Sutcliffe's Miami servers [17].
- remote\_leo2: The remote version of LEO-II runs on Geoff Sutcliffe's Miami servers [17].
- remote\_satallax: The remote version of Satallax runs on Geoff Sutcliffe's Miami servers [17].

- remote\_snark: SNARK is a first-order resolution prover developed by Stickel et al. [15]. It supports the TPTP typed first-order format (TFF0). The remote version of SNARK runs on Geoff Sutcliffe's Miami servers.
- remote\_vampire: The remote version of Vampire runs on Geoff Sutcliffe's Miami servers.
- remote\_waldmeister: Waldmeister is a unit equality prover developed by Hillenbrand et al. [8]. It can be used to prove universally quantified equations using unconditional equations, corresponding to the TPTP CNF UEQ division. The remote version of Waldmeister runs on Geoff Sutcliffe's Miami servers.
- remote\_z3: The remote version of Z3 runs on servers at the TU München (or wherever REMOTE\_SMT\_URL is set to point).

By default, Sledgehammer runs a selection of CVC3, E, E-SInE, SPASS, Vampire, Yices, and Z3 in parallel—either locally or remotely, depending on the number of processor cores available.

It is generally a good idea to run several provers in parallel. Running E, SPASS, and Vampire for 5 seconds yields a similar success rate to running the most effective of these for 120 seconds [5].

For the *min* subcommand, the default prover is *metis*. If several provers are set, the first one is used.

 $prover = \langle string \rangle$ 

Alias for provers.

$$blocking [= \langle bool \rangle] \{false\}$$
 (neg.:  $non\_blocking$ )

Specifies whether the **sledgehammer** command should operate synchronously. The asynchronous (non-blocking) mode lets the user start proving the putative theorem manually while Sledgehammer looks for a proof, but it can also be more confusing. Irrespective of the value of this option, Sledgehammer is always run synchronously if *debug* (§7.4) is enabled.

$$slice \ [=\langle bool \rangle] \ \{true\}$$
 (neg.:  $dont\_slice$ )

Specifies whether the time allocated to a prover should be sliced into several segments, each of which has its own set of possibly prover-dependent options. For SPASS and Vampire, the first slice tries the fast but incomplete set-of-support (SOS) strategy, whereas the second

slice runs without it. For E, up to three slices are tried, with different weighted search strategies and number of facts. For SMT solvers, several slices are tried with the same options each time but fewer and fewer facts. According to benchmarks with a timeout of 30 seconds, slicing is a valuable optimization, and you should probably leave it enabled unless you are conducting experiments. This option is implicitly disabled for (short) automatic runs.

See also verbose (§7.4).

$$minimize \ [= \langle smart\_bool 
angle] \ \{smart\}$$
 (neg.:  $dont\_minimize$ )

Specifies whether the minimization tool should be invoked automatically after proof search. By default, automatic minimization takes place only if it can be done in a reasonable amount of time (as determined by the number of facts in the original proof and the time it took to find or preplay it) or the proof involves an unreasonably large number of facts.

See also preplay\_timeout (§7.6) and dont\_preplay (§7.6).

Specifies whether Sledgehammer should record statistics in \$ISABELLE\_HOME\_USER/spy\_sledgehammer. These statistics can be useful to the developers of Sledgehammer. If you are willing to have your interactions recorded in the name of science, please enable this feature and send the statistics file every now and then to the author of this manual (blanchette@in.tum.de). To change the default value of this option globally, set the environment variable SLEDGEHAMMER\_SPY to yes.

See also debug (§7.4).

$$overlord \ igl[=\langle bool 
angleigr] \ \ \{false\}$$
 (neg.:  $no\_overlord)$ 

Specifies whether Sledgehammer should put its temporary files in \$ISA-BELLE\_HOME\_USER, which is useful for debugging Sledgehammer but also unsafe if several instances of the tool are run simultaneously. The files are identified by the prefixes prob\_ and mash\_; you may safely remove them after Sledgehammer has run.

See also debug (§7.4).

#### 7.2 Relevance Filter

#### $fact\_filter = \langle string \rangle \ \{smart\}$

Specifies the relevance filter to use. The following filters are available:

- *mepo*: The traditional memoryless MePo relevance filter.
- mash: The experimental MaSh machine learner. MaSh relies on the external Python program mash.py, which is part of Isabelle. To enable MaSh, set the environment variable MASH to yes. Persistent data is stored in the directory \$ISABELLE\_HOME\_USER/mash.
- *mesh*: The MeSh filter, which combines the rankings from MePo and MaSh.
- *smart*: A mixture of MePo, MaSh, and MeSh if MaSh is enabled; otherwise, MePo.

#### $max\_facts = \langle smart\_int \rangle \ \{smart\}$

Specifies the maximum number of facts that may be returned by the relevance filter. If the option is set to *smart*, it is set to a value that was empirically found to be appropriate for the prover. Typical values range between 50 and 1000.

For the MaSh-related commands  $learn\_isar$ ,  $learn\_prover$ ,  $relearn\_isar$ , and  $relearn\_prover$ , this option specifies the maximum number of facts from the background library that should be learned ( $\infty$  by default).

#### $fact\_thresholds = \langle float\_pair \rangle \ \{0.45 \ 0.85\}$

Specifies the thresholds above which facts are considered relevant by the relevance filter. The first threshold is used for the first iteration of the relevance filter and the second threshold is used for the last iteration (if it is reached). The effective threshold is quadratically interpolated for the other iterations. Each threshold ranges from 0 to 1, where 0 means that all theorems are relevant and 1 only theorems that refer to previously seen constants.

$$learn [= \langle bool \rangle] \{true\}$$
 (neg.:  $dont\_learn$ )

Specifies whether MaSh should be run automatically by Sledgehammer to learn the available theories (and hence provide more accurate results). Learning takes place only if MaSh is enabled.

# $max\_new\_mono\_instances = \langle int \rangle \ \{smart\}$

Specifies the maximum number of monomorphic instances to generate

beyond max\_facts. The higher this limit is, the more monomorphic instances are potentially generated. Whether monomorphization takes place depends on the type encoding used. If the option is set to smart, it is set to a value that was empirically found to be appropriate for the prover. For most provers, this value is 100.

See also  $type\_enc$  (§7.3).

#### $max\_mono\_iters = \langle int \rangle \ \{smart\}$

Specifies the maximum number of iterations for the monomorphization fixpoint construction. The higher this limit is, the more monomorphic instances are potentially generated. Whether monomorphization takes place depends on the type encoding used. If the option is set to *smart*, it is set to a value that was empirically found to be appropriate for the prover. For most provers, this value is 3.

See also  $type\_enc$  (§7.3).

# 7.3 Problem Encoding

# $lam\_trans = \langle string \rangle \ \{smart\}$

Specifies the  $\lambda$  translation scheme to use in ATP problems. The supported translation schemes are listed below:

- $hide\_lams$ : Hide the  $\lambda$ -abstractions by replacing them by unspecified fresh constants, effectively disabling all reasoning under  $\lambda$ -abstractions.
- *lifting*: Introduce a new supercombinator c for each cluster of n  $\lambda$ -abstractions, defined using an equation c  $x_1 \ldots x_n = t$  ( $\lambda$ -lifting).
- combs: Rewrite lambdas to the Curry combinators (I, K, S, B, C). Combinators enable the ATPs to synthesize  $\lambda$ -terms but tend to yield bulkier formulas than  $\lambda$ -lifting: The translation is quadratic in the worst case, and the equational definitions of the combinators are very prolific in the context of resolution.
- $combs\_and\_lifting$ : Introduce a new supercombinator c for each cluster of  $\lambda$ -abstractions and characterize it both using a lifted equation c  $x_1 \ldots x_n = t$  and via Curry combinators.
- $combs\_or\_lifting$ : For each cluster of  $\lambda$ -abstractions, heuristically choose between  $\lambda$ -lifting and Curry combinators.

- $keep\_lams$ : Keep the  $\lambda$ -abstractions in the generated problems. This is available only with provers that support the THF0 syntax.
- *smart*: The actual translation scheme used depends on the ATP and should be the most efficient scheme for that ATP.

For SMT solvers, the  $\lambda$  translation scheme is always *lifting*, irrespective of the value of this option.

Specifies whether fresh function symbols should be generated as aliases for applications of curried functions in ATP problems.

```
type\_enc = \langle string \rangle \ \{smart\}
```

Specifies the type encoding to use in ATP problems. Some of the type encodings are unsound, meaning that they can give rise to spurious proofs (unreconstructible using metis). The type encodings are listed below, with an indication of their soundness in parentheses. An asterisk (\*) indicates that the encoding is slightly incomplete for reconstruction with metis, unless the strict option (described below) is enabled.

- *erased* (unsound): No type information is supplied to the ATP, not even to resolve overloading. Types are simply erased.
- $poly\_guards$  (sound): Types are encoded using a predicate  $g(\tau, t)$  that guards bound variables. Constants are annotated with their types, supplied as extra arguments, to resolve overloading.
- $poly\_tags$  (sound): Each term and subterm is tagged with its type using a function  $t(\tau, t)$ .
- poly\_args (unsound): Like for poly\_guards constants are annotated with their types to resolve overloading, but otherwise no type information is encoded. This is the default encoding used by the metis command.
- raw\_mono\_guards, raw\_mono\_tags (sound); raw\_mono\_args (unsound):

Similar to *poly\_guards*, *poly\_tags*, and *poly\_args*, respectively, but the problem is additionally monomorphized, meaning that type variables are instantiated with heuristically chosen ground types. Monomorphization can simplify reasoning but also leads to larger fact bases, which can slow down the ATPs.

- $mono\_guards$ ,  $mono\_tags$  (sound);  $mono\_args$  (unsound): Similar to  $raw\_mono\_guards$ ,  $raw\_mono\_tags$ , and  $raw\_mono\_args$ , respectively but types are mangled in constant names instead of being supplied as ground term arguments. The binary predicate  $g(\tau, t)$  becomes a unary predicate  $g\_\tau(t)$ , and the binary function  $t(\tau, t)$  becomes a unary function  $t\_\tau(t)$ .
- mono\_native (sound): Exploits native first-order types if the prover supports the TFF0, TFF1, or THF0 syntax; otherwise, falls back on mono\_guards. The problem is monomorphized.
- mono\_native\_higher (sound): Exploits native higher-order types if the prover supports the THF0 syntax; otherwise, falls back on mono\_native or mono\_guards. The problem is monomorphized.
- poly\_native (sound): Exploits native first-order polymorphic types if the prover supports the TFF1 syntax; otherwise, falls back on mono\_native.
- poly\_guards?, poly\_tags?, raw\_mono\_guards?, raw\_mono\_tags?, mono\_guards?, mono\_tags?, mono\_native? (sound\*):

The type encodings poly\_guards, poly\_tags, raw\_mono\_guards, raw\_mono\_tags, mono\_guards, mono\_tags, and mono\_native are fully typed and sound. For each of these, Sledgehammer also provides a lighter variant identified by a question mark ('?') that detects and erases monotonic types, notably infinite types. (For mono\_native, the types are not actually erased but rather replaced by a shared uniform type of individuals.) As argument to the metis proof method, the question mark is replaced by a "\_query" suffix.

• poly\_guards??, poly\_tags??, raw\_mono\_guards??, raw\_mono\_tags??, mono\_guards??, mono\_tags?? (sound\*):

Even lighter versions of the '?' encodings. As argument to the *metis* proof method, the '??' suffix is replaced by "\_query\_query".

• poly\_guards@, poly\_tags@, raw\_mono\_guards@, raw\_mono\_tags@ (sound\*):

Alternative versions of the '??' encodings. As argument to the *metis* proof method, the '@' suffix is replaced by "\_at".

• poly\_args?, raw\_mono\_args? (unsound): Lighter versions of poly\_args and raw\_mono\_args. • *smart*: The actual encoding used depends on the ATP and should be the most efficient sound encoding for that ATP.

For SMT solvers, the type encoding is always *mono\_native*, irrespective of the value of this option.

See also max\_new\_mono\_instances (§7.2) and max\_mono\_iters (§7.2).

$$strict \ [= \langle bool \rangle] \ \{false\}$$
 (neg.:  $non\_strict$ )

Specifies whether Sledgehammer should run in its strict mode. In that mode, sound type encodings marked with an asterisk (\*) above are made complete for reconstruction with *metis*, at the cost of some clutter in the generated problems. This option has no effect if *type\_enc* is deliberately set to an unsound encoding.

# 7.4 Output Format

$$verbose \ [= \langle bool 
angle] \ \ \{false\}$$
 (neg.:  $quiet$ )

Specifies whether the **sledgehammer** command should explain what it does. This option is implicitly disabled for automatic runs.

Specifies whether Sledgehammer should display additional debugging information beyond what verbose already displays. Enabling debug also enables verbose and blocking (§7.1) behind the scenes. The debug option is implicitly disabled for automatic runs.

See also spy (§7.1) and overload (§7.1).

Specifies whether Isar proofs should be output in addition to one-liner *metis* proofs. The construction of Isar proof is still experimental and may sometimes fail; however, when they succeed they are usually faster and more intelligible than *metis* proofs. If the option is set to *smart* (the default), Isar proofs are only generated when no working one-liner *metis* proof is available.

$$isar\_compress = \langle int \rangle \ \{10\}$$

Specifies the granularity of the generated Isar proofs if  $isar\_proofs$  is explicitly enabled. A value of n indicates that each Isar proof step should correspond to a group of up to n consecutive proof steps in the ATP proof.

# $dont\_compress\_isar \ [= true]$

Alias for " $isar\_compress = 0$ ".

$$isar\_try0 \ [= \langle bool \rangle] \ \{true\}$$
 (neg.:  $dont\_try0\_isar$ )

Specifies whether standard proof methods such as auto and blast should be tried as alternatives to metis and smt in Isar proofs. The collection of methods is roughly the same as for the try0 command.

#### 7.5 Authentication

```
expect = \langle string \rangle
```

Specifies the expected outcome, which must be one of the following:

- some: Sledgehammer found a proof.
- none: Sledgehammer found no proof.
- *timeout*: Sledgehammer timed out.
- unknown: Sledgehammer encountered some problem.

Sledgehammer emits an error (if *blocking* is enabled) or a warning (otherwise) if the actual outcome differs from the expected outcome. This option is useful for regression testing.

See also blocking ( $\S7.1$ ) and timeout ( $\S7.6$ ).

#### 7.6 Timeouts

$$timeout = \langle float\_or\_none \rangle$$
 {30}

Specifies the maximum number of seconds that the automatic provers should spend searching for a proof. This excludes problem preparation and is a soft limit. For automatic runs, the "Auto Time Limit" option under "Plugins > Plugin Options > Isabelle > General" is used instead.

$$preplay\_timeout = \langle float\_or\_none \rangle$$
 {3}

Specifies the maximum number of seconds that *metis* or *smt* should spend trying to "preplay" the found proof. If this option is set to 0, no preplaying takes place, and no timing information is displayed next to the suggested *metis* calls.

See also minimize (§7.1).

# $dont\_preplay [= true]$

Alias for " $preplay\_timeout = 0$ ".

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