# ICS143A: Principles of Operating Systems

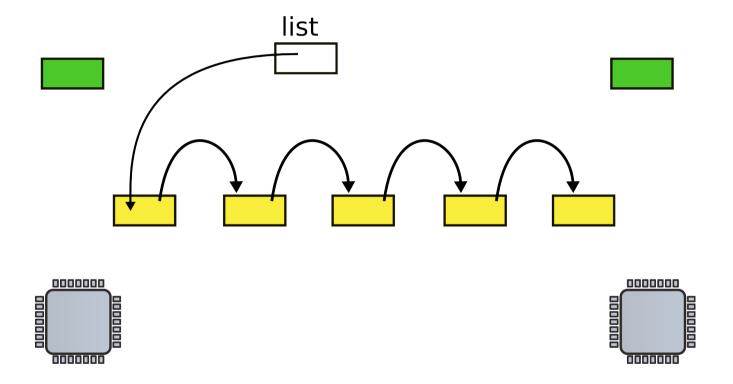
Lecture 16: Locking (continued)

Anton Burtsev February, 2017

#### Recap: Race conditions

- Disk driver maintains a list of outstanding requests
- Each process can add requests to the list

### Request queue (e.g. incoming network packets)



 Linked list, list is pointer to the first element

### List implementation with locks

```
9 insert(int data)
10 {
11 struct list *1;
13 \quad l = malloc(size of *l);
     acquire(&listlock);
14
     1->data = data;
15 l \rightarrow next = list;

    Critical section

16
     list = 1;
     release(&listlock);
17 }
```

#### Xchg instruction

- Swap a word in memory with a new value
  - Atomic!
  - Return old value

## Correct implementation

```
1573 void
1574 acquire(struct spinlock *lk)
1575 {
1580 // The xchg is atomic.
while(xchg(&lk->locked, 1) != 0)
1582
1592 }
```

#### One last detail...

```
9 insert(int data)
10 {
11 struct list *1;
13  l = malloc(sizeof *1);
    acquire(&listlock);
14
    1->data = data;
15 l->next = list;
16
    list = 1;
    release(&listlock);
17 }
```

## Correct implementation

```
1573 void
1574 acquire(struct spinlock *lk)
1575 {
. . .
1580
      // The xchg is atomic.
1581
       while(xchg(&lk->locked, 1) != 0)
1582
1584
      // Tell the C compiler and the processor to not move loads or
          stores
1585
       // past this point, to ensure that the critical section's memory
1586
      // references happen after the lock is acquired.
1587
      __sync_synchronize();
. . .
1592 }
```

#### Deadlocks

#### Deadlocks

```
acquire(A)

acquire(B)

acquire(B) {
    while(xchg(&B->locked, 1) != 0)
}
acquire(A) {
    while(xchg(&A->locked, 1) != 0)
}
```





#### Lock ordering

Locks need to be acquired in the same order

#### Locks and interrupts

```
Network
                        interrupt
                                     network_packet(){
network_packet(){
                                        insert() {
  insert() {
                                        _ acquire(A)
     acquire(A)
                     0000000
```

#### Locks and interrupts

Never hold a lock with interrupts enabled

```
1573 void
1574 acquire(struct spinlock *lk)
1575 {
      pushcli(); // disable interrupts to avoid deadlock.
1576
1577
      if(holding(lk))
        panic("acquire");
1578
1580 // The xchg is atomic.
      while(xchg(&lk->locked, 1) != 0)
1581
1582
. . .
1587    __sync_synchronize();
                           Disabling interrupts
1592 }
```

#### Simple disable/enable is not enough

- If two locks are acquired
- Interrupts should be re-enabled only after the second lock is released

Pushcli() uses a counter

```
1655 pushcli(void)
1656 {
1657
       int eflags;
1658
1659
       eflags = readeflags();
1660
       cli();
1661
       if(cpu->ncli == 0)
1662
      cpu->intena = eflags & FL IF;
1663
       cpu->ncli += 1;
1664 }
. . .
1667 popcli(void)
1668 {
1669
       if(readeflags()&FL_IF)
1670
         panic("popcli - interruptible");
1671
       if(--cpu->ncli < 0)</pre>
1672
         panic("popcli");
       if(cpu->ncli == 0 && cpu->intena)
1673
1674
         sti();
1675 }
```

#### Pushcli()/popcli()



#### Send/receive queue

```
100 struct q {
                                      112 void*
101 void *ptr;
                                      113 recv(struct q *q)
102 };
                                      114 {
103
                                      115 void *p;
104 void*
                                      116
105 send(struct q *q, void *p)
                                            while((p = q->ptr) == 0)
                                      117
106 {
                                      118 ;
107 while (q->ptr != 0)
                                      119 q->ptr = 0;
108
109 q \rightarrow ptr = p;
                                      120
                                            return p;
110 }
                                      121 }
```

#### Send/receive queue

```
100 struct q {
                                      112 void*
101 void *ptr;
                                      113 recv(struct q *q)
102 };
                                      114 {
103
                                      115 void *p;
104 void*
                                      116
105 send(struct q *q, void *p)
                                            while((p = q->ptr) == 0)
                                      117
106 {
                                      118 ;
107 while (q->ptr != 0)
                                      119 q \rightarrow ptr = 0;
108
109 q - ptr = p;
                                      120
                                            return p;
110 }
                                      121 }
```

- Expensive if communication is rare
  - Receiver wastes CPU cycles

#### Sleep and wakeup syscalls

- sleep(channel)
  - Put calling process to sleep
  - Release CPU for other work
- wakeup(channel)
  - Wakes all processes sleeping on a channel
    - If any
  - i.e., causes sleep() calls to return

#### Send/receive queue

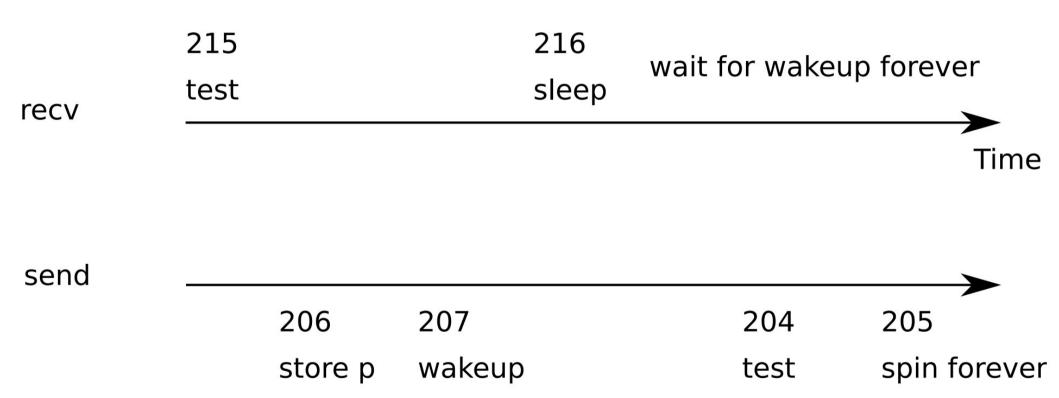
```
210 void*
201 void*
                                 211 recv(struct q *q)
202 send(struct q *q, void *p)
                                 212 {
203 {
                                 213 void *p;
    while(q->ptr != 0)
204
                                 214
205
                                       while((p = q->ptr) == 0)
                                 215
206 	 q->ptr = p;
                                 216
                                         sleep(q);
      wakeup(q); /*wake recv*/
207
                                 217 	 q->ptr = 0;
208 }
                                 218
                                       return p;
                                 219 }
```

#### Send/receive queue

```
210 void*
201 void*
                                 211 recv(struct q *q)
202 send(struct q *q, void *p)
                                  212 {
203 {
                                  213 void *p;
    while(q->ptr != 0)
204
                                  214
205
                                       while((p = q->ptr) == 0)
                                  215
206 	 q->ptr = p;
                                  216 sleep(q);
      wakeup(q); /*wake recv*/
207
                                  217 	 q->ptr = 0;
208 }
                                  218
                                       return p;
                                  219 }
```

- recv() gives up the CPU to other processes
  - But there is a problem...

#### Lost wakeup problem



### Thank you!