CS143A: Principles of Operating Systems

Lecture 12: File systems

Anton Burtsev November, 2018

The role of file systems

The role of file systems

- Sharing
 - Sharing of data across users and applications
- Persistent storage
 - Data is available after reboot

Architecture

- On-disk and in-memory data structures that represent
 - The tree of named files and directories
 - Record identities of disk blocks which hold data for each file
 - Record which areas of the disk are free

Crash recovery

- File systems must support crash recovery
 - A power loss may interrupt a sequence of updates
 - And leave the file system in an inconsistent state
 - E.g., a block both marked free and used

Speed

- Access to a block device is several orders of magnitude slower
 - Memory: 200 cycles
 - Disk: 20 000 000 cycles
- A file system must maintain a cache of disk blocks in memory

Block layer

System calls	File descriptors
Pathnames	Recursive lookup
Directories	Directory inodes
Files	Inodes and block allocator
Transactions	Logging
Blocks	Buffer cache

- Read and write data
 - From a block device
 - Into a buffer cache
- Synchronize across multiple readers and writers

Transactions

System calls	File descriptors
Pathnames	Recursive lookup
Directories	Directory inodes
Files	Inodes and block allocator
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Transactions	Logging

Group multiple writes into an atomic transaction

Files

File descriptors
Recursive lookup
Directory inodes
Inodes and block allocator
Logging
Buffer cache

Unnamed files

- Represented as inodes
- Sequence of blocks holding file's data

Directories

System calls	File descriptors
Pathnames	Recursive lookup
Directories	Directory inodes
Files	Inodes and block allocator
Transactions	Logging
Blocks	Buffer cache

- Special kind of inode
 - Sequence of directory entries
 - Each contains name and a pointer to an unnamed inode

Pathnames

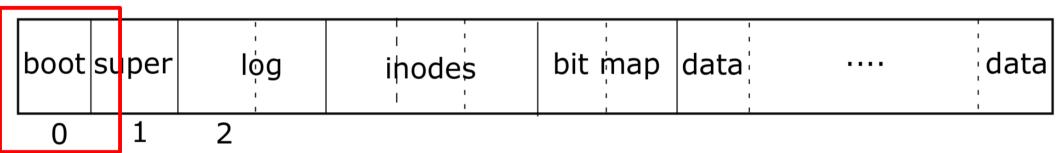
System calls	File descriptors
Pathnames	Recursive lookup
Directories	Directory inodes
Files	Inodes and block allocator
Transactions	Logging
Blocks	Buffer cache

- Hierarchical path names
 - /usr/bin/sh
 - Recursive lookup

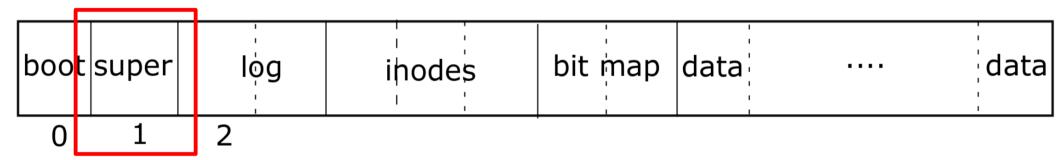
System call

System calls	File descriptors
Pathnames	Recursive lookup
Directories	Directory inodes
Files	Inodes and block allocator
Transactions	Logging
Blocks	Buffer cache
	•

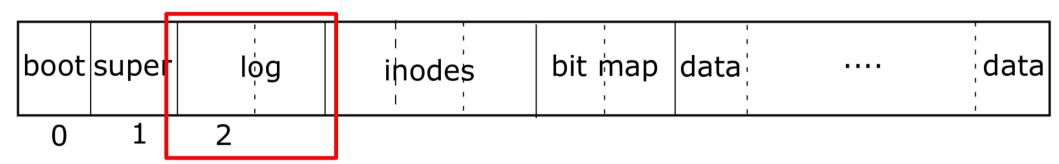
- Abstract UNIX resources as files
 - Files, sockets, devices, pipes, etc.
- Unified programming interface



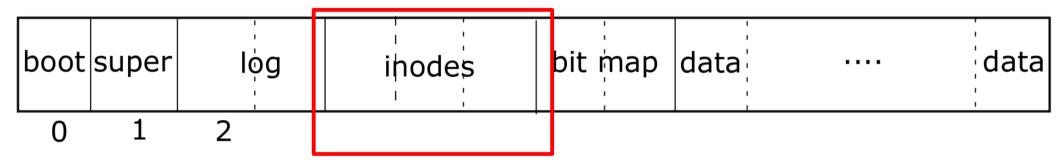
Block #0: Boot code



- Block #0: Boot code
- Block #1: (superblock) Metadata about the file system
 - Size (number of blocks)
 - Number of data blocks
 - Number of inodes
 - Number of blocks in log

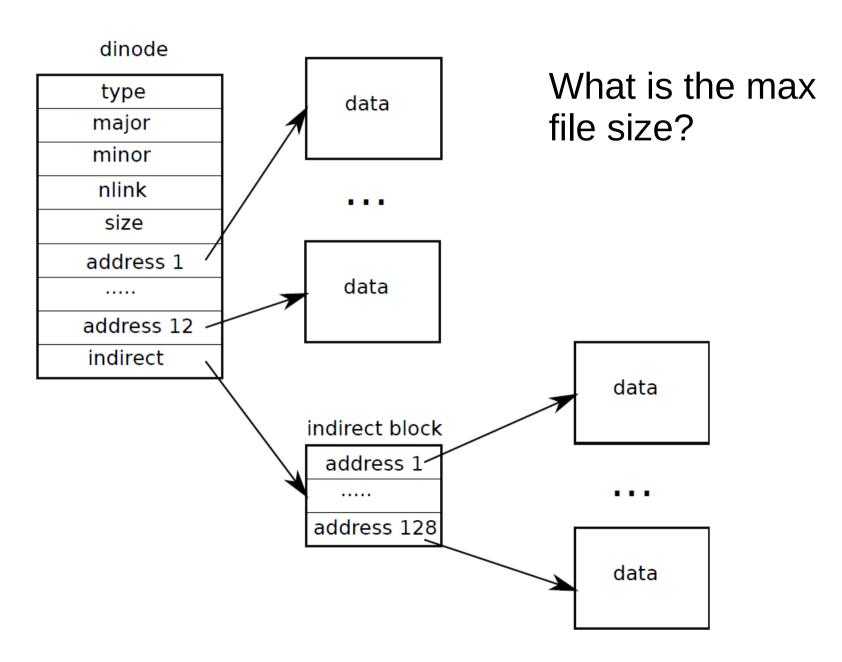


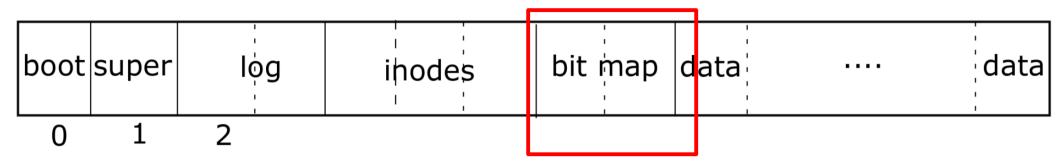
 Block #2: Log area: maintaining consistency in case of a power outage or system crash



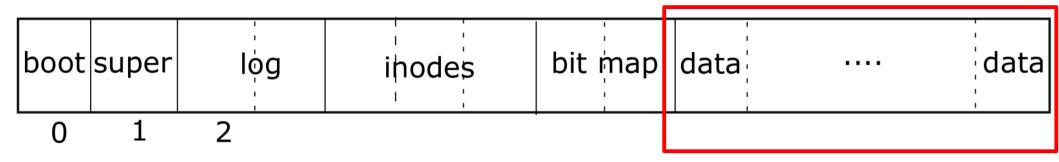
- Inode area
 - Unnamed files

Representing files on disk





• Bit map area: track which blocks are in use



Data area: actual file data

Buffer cache layer

Buffer cache layer

Two goals:

- Synchronization:
 - Only one copy of a data block exist in the kernel
 - Only one writer updates this copy at a time
- Caching
 - Frequently used copies are cached for efficient reads and writes

```
3750 struct buf {
                                                   Buffer cache
3751
      int flags;
3752
     uint dev;
3753
     uint blockno;
3754 struct buf *prev; // LRU cache list
3755
    struct buf *next;
3756 struct buf *qnext; // disk queue
3757 uchar data[BSIZE];
3758 };
3759 #define B BUSY 0x1 // buffer is locked by some process
3760 #define B VALID 0x2 // buffer has been read from disk
3761 #define B DIRTY 0x4 // buffer needs to be written to disk
```

```
4329 struct {
4330    struct spinlock lock;
4331    struct buf buf[NBUF];
4332
4333    // Linked list of all buffers, through prev/next.
4334    // head.next is most recently used.
4335    struct buf head;
4336 } bcache;
```

```
struct buf {
3750
                                                     Buffer cache
3751
      int flags;
3752
      uint dev;
3753
      uint blockno;
3754
      struct buf *prev; // LRU cache list
      struct buf *next;
3755

    Array of buffers

3756
      struct buf *qnext; // disk queue
3757
      uchar data[BSIZE];
3758 }:
3759 #define B BUSY 0x1 // buffer is locked by some process
3760 #define B VALID 0x2 // buffer has been read from disk
3761 #define B DIRTY 0x4 // buffer needs to be written to disk
4329 struct {
      struct spinlock lock;
4330
4331
      struct buf buf[NBUF];
4332
4333
     // Linked list of all buffers, through prev/next.
4334
      // head.next is most recently used.
4335
      struct buf head;
4336 } bcache;
```

```
3750 struct buf {
                                                         Cached data
3751
      int flags;
3752
     uint dev;
                                                          • 512 bytes
3753 uint blockno;
3754 struct buf *prev; // LRU cache list
3755 struct buf *next;
3756
     struct buf *qnext; // disk queue
3757
     uchar data[BSIZE];
3758 };
3759 #define B BUSY 0x1 // buffer is locked by some process
3760 #define B VALID 0x2 // buffer has been read from disk
3761 #define B DIRTY 0x4 // buffer needs to be written to disk
4329 struct {
4330
      struct spinlock lock;
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     struct buf buf [NBUF];
4332
4333
     // Linked list of all buffers, through prev/next.
4334
      // head.next is most recently used.
4335
      struct buf head;
4336 } bcache;
```

```
3750 struct buf {
                                                          Flags
3751
      int flags;
3752
     uint dev;
3753
     uint blockno;
3754 struct buf *prev; // LRU cache list
3755
     struct buf *next;
3756
     struct buf *qnext; // disk gaeue
     uchar data[BSIZE];
3757
3758 };
3759 #define B BUSY 0x1 // buffer is locked by some process
3760 #define B VALID 0x2 // buffer has been read from disk
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4329 struct {
4330
      struct spinlock lock;
4331
      struct buf buf [NBUF];
4332
4333
     // Linked list of all buffers, through prev/next.
4334
      // head.next is most recently used.
4335
      struct buf head;
4336 } bcache;
```

```
3750 struct buf {
                                                         Device
3751
      int flags;
3752
     uint dev;

    We might have

3753
     uint blockno;
                                                             multiple disks
3754 struct buf *prev; // LRU cache list
3755 struct buf *next;
3756 struct buf *qnext; // disk queue
3757
    uchar data[BSIZE];
3758 };
3759 #define B BUSY 0x1 // buffer is locked by some process
3760 #define B VALID 0x2 // buffer has been read from disk
3761 #define B DIRTY 0x4 // buffer needs to be written to disk
4329 struct {
4330
      struct spinlock lock;
4331
     struct buf buf [NBUF];
4332
4333
     // Linked list of all buffers, through prev/next.
4334
      // head.next is most recently used.
4335
      struct buf head;
4336 } bcache;
```

```
3750 struct buf {
                                                         Block number on disk
3751
      int flags;
3752
     uint dev;
     uint blockno:
3753
3754
     struct buf *prev; // LRU cache list
3755
     struct buf *next;
3756
     struct buf *qnext; // disk queue
3757
     uchar data[BSIZE];
3758 };
3759 #define B BUSY 0x1 // buffer is locked by some process
3760 #define B VALID 0x2 // buffer has been read from disk
3761 #define B DIRTY 0x4 // buffer needs to be written to disk
4329 struct {
4330
      struct spinlock lock;
4331
      struct buf buf [NBUF];
4332
4333
     // Linked list of all buffers, through prev/next.
4334
      // head.next is most recently used.
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      struct buf head;
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```

```
3750 struct buf {
                                                          LRU list
3751
      int flags;
3752
     uint dev;
3753
     uint blockno;
3754 struct buf *prev; // LRU cache list
3755
     struct buf *next;
3756 struct buf *qnext; // disk queue
3757 uchar data[BSIZE];
3758 };
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4329 struct {
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      struct buf buf[NBUF];
4331
4332
4333
     // Linked list of all buffers, through prev/next.
4334
      // head.next is most recently used.
4335
      struct buf head;
4336 } bcache;
```

Buffer cache layer: interface

- bread() and bwrite() obtain a copy for reading or writing
 - Owned until brelse()
 - Locking with a flag (B_BUSY)
- Other threads will be blocked and wait until brelse()

```
4401 struct buf*
4402 bread(uint dev, uint sector)
4403 {
4404 struct buf *b;
4405
4406 b = bget(dev, sector);
4407 if(!(b->flags & B_VALID)) {
        iderw(b);
4408
4409 }
4410 return b;
4411 }
```

bread(): block read

- bget() obtain a copy of the block in the buffer caceh
 - Owned until brelse()
 - Locked with a flag (B_BUSY)

```
bget(): Getting a
4365 static struct buf*
4366 bget(uint dev, uint blockno)
                                          block from the
4367 {
4368
      struct buf *b;
                                            buffer cache
4369
4370
      acquire(&bcache.lock);
                                                 (part 1)
4371
4372
     loop:
4373
      // Is the block already cached?
4374
      for(b = bcache.head.next; b != &bcache.head; b = b->next){
4375
        if(b->dev == dev && b->blockno == blockno){
4376
          if(!(b->flags & B BUSY)){
4377
           b->flags |= B BUSY;
4378
            release(&bcache.lock);
4379
            return b:
4380
4381
          sleep(b, &bcache.lock);
4382
          goto loop;
4383
        }
4384
```

```
4385
      // Not cached; recycle some non-busy and clean buffer.
4386
      // "clean" because B_DIRTY and !B_BUSY means log.c
4387
4388
      // hasn't yet committed the changes to the buffer.
      for(b = bcache.head.prev; b != &bcache.head; b = b->prev){
4389
        if((b->flags & B_BUSY)== 0 && (b->flags & B_DIRTY)== 0){
4390
4391
          b->dev = dev:
4392
          b->blockno = blockno;
4393
          b->flags = B_BUSY;
          release(&bcache.lock);
4394
4395
          return b;
4396
                                  bget(): Getting a
      }
4397
4398
      panic("bget: no buffers");
                                    block from the
4399 }
                                      buffer cache
                                          (part 2)
```

```
4401 struct buf*
4402 bread(uint dev, uint sector)
4403 {
4404 struct buf *b;
4405
4406 b = bget(dev, sector);
4407 if(!(b->flags & B_VALID)) {
4408
        iderw(b);
4409 }
4410 return b;
4411 }
```

bread(): block read

- bget() reserved an entry in the buffer cache
- If the data is not there read it from disk

- Block data was updated
 - Mark it as dirty in the buffer cache
 - Write data to disk

```
4423 // Release a B BUSY buffer.
4424 // Move to the head of the MRU list.
4425 void
4426 brelse(struct buf *b)
4427 {
       if((b->flags & B BUSY) == 0)
4428
4429
     panic("brelse");
4430
4431
      acquire(&bcache.lock);
4432
4433
      b->next->prev = b->prev;
4434
      b->prev->next = b->next;
4435
      b->next = bcache.head.next;
4436
      b->prev = &bcache.head;
4437
      bcache.head.next->prev = b;
4438
      bcache.head.next = b;
4439
4440
      b->flags &= ~B BUSY;
4441
      wakeup(b);
4442
4443
      release(&bcache.lock);
4444 }
```

brelse(): Release buffer

- Maintain least recently used list
 - Move to the head

Common block usage pattern

```
bread()
bwrite()
brelse()
```

- Read
- Write
- Release

Logging layer

Logging layer

- Consistency
 - File system operations involve multiple writes to disk
 - During the crash, subset of writes might leave the file system in an inconsistent state
 - E.g. if crash happens during file delete operation it can leave the file system with:
 - Ex #1: Directory entry pointing to a free inode
 - Ex #2: Allocated but unlinked inode

Logging

- Writes don't directly go to disk
 - Instead they are logged in a journal
 - Once all writes are logged, the system writes a special commit record
 - Indicating that log contains a complete operation
- At this point file system copies writes to the ondisk data structures
 - After copy completes, log record is erased

Recovery

- After reboot, copy the log
 - For operations marked as complete
 - Copy blocks to disk
 - For operations partially complete
 - Discard all writes
 - Information might be lost (output consistency, e.g. can launch the rocket twice)

```
begin_op();
...
bp = bread(...);
bp->data[...] = ...;
log_write(bp);
...
end_op();
```

```
4532 struct logheader {
4533 int n;
                                 Log (in memory)
4534 int block[LOGSIZE];
4535 };
4536
4537 struct log {
4538 struct spinlock lock;
4539 int start;
4540 int size;
4541 int outstanding; // how many FS sys calls are
                         executing.
4542 int committing; // in commit(), please wait.
4543 int dev;
4544 struct logheader lh;
4545 };
```

```
begin_op();
...
bp = bread(...);
bp->data[...] = ...;
log_write(bp);
...
end_op();
```

```
4626 // called at the start of each FS system call.
4627 void
4628 begin op(void)
4629 {
                                            begin op()
4630
      acquire(&log.lock);
4631
      while(1){
4632
        if(log.committing){
4633
          sleep(&log, &log.lock);
4634
        } else if(log.lh.n + (log.outstanding+1)*MAXOPBLOCKS > LOGSIZE){
4635
          // this op might exhaust log space; wait for commit.
4636
          sleep(&log, &log.lock);
4637
        } else {

    Case #1

          log.outstanding += 1;
4638
4639
          release(&log.lock);
                                           Log is being
4640
          break;
                                            committed
4641
4642 }

    Sleep

4643 }
```

```
4626 // called at the start of each FS system call.
4627 void
4628 begin op(void)
4629 {
                                           begin op()
      acquire(&log.lock);
4630
4631
      while(1){
4632
        if(log.committing){
4633
          sleep(&log, &log.lock);
4634
        } else if(log.lh.n + (log.outstanding+1)*MAXOPBLOCKS > LOGSIZE){
4635
          // this op might exhaust log space; wait for commit.
4636
          sleep(&log, &log.lock);
4637
        } else {

    Case #2

          log.outstanding += 1;
4638
4639
          release(&log.lock);
                                          Log doesn't have
4640
          break;
                                          enough space for the
4641
                                           new transaction
4642 }
4643 }
```

```
4626 // called at the start of each FS system call.
4627 void
4628 begin op(void)
4629 {
                                            begin op()
4630
      acquire(&log.lock);
4631
      while(1){
4632
        if(log.committing){
4633
          sleep(&log, &log.lock);
4634
        } else if(log.lh.n + (log.outstanding+1)*MAXOPBLOCKS > LOGSIZE){
4635
          // this op might exhaust log space; wait for commit.
4636
          sleep(&log, &log.lock);
4637
        } else {

    Case #3

          log.outstanding += 1;
4638
4639
          release(&log.lock);

    All ok, reserve space

4640
          break;
                                           in the log for the new
4641
                                           transaction
4642 }
4643 }
```

```
begin_op();
...
bp = bread(...);
bp->data[...] = ...;
log_write(bp);
...
end_op();
```

log_write() replacesbwrite(); brelse()

```
4722 log_write(struct buf *b)
                                              log write
4723 {
4724
      int i:
4725
4726
       if (log.lh.n >= LOGSIZE || log.lh.n >= log.size - 1)
4727
         panic("too big a transaction");
4728
       if (log.outstanding < 1)
4729
         panic("log_write outside of trans");
4730
4731
      acquire(&log.lock);
4732
       for (i = 0; i < log.lh.n; i++) {
4733
         if (log.lh.block[i] == b->blockno) // log absorbtion
4734
           break:
4735
       }

    Check if already

4736
       log.lh.block[i] = b->blockno;
4737
       if (i == log.lh.n)
                                                 in log
4738
         log.lh.n++;
       b->flags |= B_DIRTY; // prevent eviction
4739
       release(&log.lock);
4740
4741 }
```

```
4722 log_write(struct buf *b)
                                              log write
4723 {
4724
      int i:
4725
4726
       if (log.lh.n >= LOGSIZE || log.lh.n >= log.size - 1)
4727
         panic("too big a transaction");
4728
       if (log.outstanding < 1)
4729
         panic("log_write outside of trans");
4730
4731
      acquire(&log.lock);
       for (i = 0; i < log.lh.n; i++) {
4732
4733
         if (log.lh.block[i] == b->blockno) // log absorbtion
4734
           break:
4735
      }

    Add to the log

4736
       log.lh.block[i] = b->blockno;
4737
       if (i == log.lh.n)

    Prevent eviction

4738
         log.lh.n++;
4739
       b->flags |= B DIRTY; // prevent eviction
      release(&log.lock);
4740
4741 }
```

```
begin_op();
...
bp = bread(...);
bp->data[...] = ...;
log_write(bp);
...
end_op();
```

```
4653 end op(void)
4654 {
                                                          end op()
      int do commit = 0;
4655
4656
4657
      acquire(&log.lock);
4658
      log.outstanding -= 1;
      if(log.outstanding == 0){
4661
4662
        do commit = 1;
4663
        log.committing = 1;
4664
      } else {
4665
        // begin_op() may be waiting for log space.
4666
        wakeup(&log);
4667
4668
      release(&log.lock);
4669
4670
      if(do commit){
4671
     // call commit w/o holding locks, since not allowed
4672
        // to sleep with locks.
4673
     commit();
4674
     acquire(&log.lock);
        log.committing = 0;
4675
        wakeup(&log);
4676
        release(&log.lock);
4677
4678
     }
4679 }
```

```
4653 end op(void)
4654 {
                                                          end op()
      int do commit = 0;
4655
4656
4657
      acquire(&log.lock);
4658
      log.outstanding -= 1;
4661
      if(log.outstanding == 0){
4662
      do commit = 1;
4663
     log.committing = 1;
      } else {
4664
4665
        // begin_op() may be waiting for log space.
4666
        wakeup(&log);
4667
4668
      release(&log.lock);
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      if(do commit){
4670
4671
        // call commit w/o holding locks, since not allowed
4672
        // to sleep with locks.
4673
        commit();
4674
     acquire(&log.lock);
        log.committing = 0;
4675
        wakeup(&log);
4676
        release(&log.lock);
4677
4678
     }
4679 }
```

```
4701 commit()
                                  commit()
4702 {
4703 if (log.lh.n > 0) {
4704
         write_log(); // Write modified blocks
                          from cache to log
4705
         write head(); // Write header to disk --
                          the real commit
4706
         install_trans(); // Now install writes
                             to home locations
4707
         log.lh.n = 0;
         write_head(); // Erase the transaction
4708
                          from the log
4709 }
4710 }
```

```
4681 // Copy modified blocks from cache to log.
4682 static void
4683 write_log(void)
                                        write log()
4684 {
4685 int tail;
4686
4687 for (tail = 0; tail < log.lh.n; tail++) {
        struct buf *to = bread(log.dev,
4688
                          log.start+tail+1); // log block
        struct buf *from = bread(log.dev,
4689
                          log.lh.block[tail]); // cache block
        memmove(to->data, from->data, BSIZE);
4690
        bwrite(to); // write the log
4691
4692
        brelse(from);

    Loop through the entire log

4693
        brelse(to);
4694 }
4695 }
```

```
4681 // Copy modified blocks from cache to log.
4682 static void
4683 write_log(void)
                                        write log()
4684 {
4685 int tail;
4686
4687 for (tail = 0; tail < log.lh.n; tail++) {
        struct buf *to = bread(log.dev,
4688
                          log.start+tail+1); // log block
        struct buf *from = bread(log.dev,
4689
                          log.lh.block[tail]); // cache block
        memmove(to->data, from->data, BSIZE);
4690
        bwrite(to); // write the log
4691
4692
        brelse(from);

    Read the log block

4693
        brelse(to);
4694 }

    Log goes to

4695 }
                             log.start+tail+1
```

```
4681 // Copy modified blocks from cache to log.
4682 static void
4683 write_log(void)
                                         write log()
4684 {
4685 int tail;
4686
4687 for (tail = 0; tail < log.lh.n; tail++) {
         struct buf *to = bread(log.dev,
4688
                           log.start+tail+1); // log block
         struct buf *from = bread(log.dev,
4689
                           log.lh.block[tail]); // cache block
         memmove(to->data, from->data, BSIZE);
4690
         bwrite(to); // write the log
4691
4692
         brelse(from);

    Read the actual block

4693
         brelse(to);
4694 }

    It's in the buffer cache

4695 }

    Block number is in

                              log.lh.block[tail]
```

```
4681 // Copy modified blocks from cache to log.
4682 static void
4683 write_log(void)
                                        write log()
4684 {
4685 int tail;
4686
4687 for (tail = 0; tail < log.lh.n; tail++) {
        struct buf *to = bread(log.dev,
4688
                          log.start+tail+1); // log block
        struct buf *from = bread(log.dev,
4689
                          log.lh.block[tail]); // cache block
        memmove(to->data, from->data, BSIZE);
4690
        bwrite(to); // write the log
4691
4692
        brelse(from);

    Copy block data into the log

4693
        brelse(to);
4694 }
4695 }
```

```
4681 // Copy modified blocks from cache to log.
4682 static void
4683 write_log(void)
                                        write log()
4684 {
4685 int tail;
4686
4687 for (tail = 0; tail < log.lh.n; tail++) {
        struct buf *to = bread(log.dev,
4688
                           log.start+tail+1); // log block
        struct buf *from = bread(log.dev,
4689
                          log.lh.block[tail]); // cache block
        memmove(to->data, from->data, BSIZE);
4690
        bwrite(to); // write the log
4691
4692
        brelse(from);

    Write the log block (to)

        brelse(to);
4693
4694 }

    Release both blocks

4695 }
```

```
4701 commit()
                                  commit()
4702 {
4703 if (log.lh.n > 0) {
4704
         write_log(); // Write modified blocks
                          from cache to log
4705
         write head(); // Write header to disk --
                          the real commit
4706
         install_trans(); // Now install writes
                             to home locations
4707
         log.lh.n = 0;
         write_head(); // Erase the transaction
4708
                          from the log
4709 }
4710 }
```

```
4600 // Write in-memory log header to disk.
4601 // This is the true point at which the
4602 // current transaction commits.
4603 static void
                                    write head()
4604 write_head(void)
4605 {
4606 struct buf *buf = bread(log.dev, log.start);
      struct logheader *hb = (struct logheader *)
4607
                                         (buf->data);
     int i;
4608
4609
      hb->n = log.lh.n;
4610
      for (i = 0; i < log.lh.n; i++) {
4611
        hb->block[i] = log.lh.block[i];
4612

    Read the log header block

      bwrite(buf);
4613
4614 brelse(buf);
                          • It's in log.start
4615 }
```

```
4600 // Write in-memory log header to disk.
4601 // This is the true point at which the
4602 // current transaction commits.
4603 static void
                                    write head()
4604 write_head(void)
4605 {
4606 struct buf *buf = bread(log.dev, log.start);
      struct logheader *hb = (struct logheader *)
4607
                                         (buf->data);
     int i;
4608
4609
      hb->n = log.lh.n;
4610
      for (i = 0; i < log.lh.n; i++) {
4611
        hb->block[i] = log.lh.block[i];
4612

    Interpret buf->data as log

      bwrite(buf);
4613
                         header
4614 brelse(buf);
4615 }

    See how type casts work in C
```

```
4600 // Write in-memory log header to disk.
4601 // This is the true point at which the
4602 // current transaction commits.
4603 static void
                                    write head()
4604 write_head(void)
4605 {
4606 struct buf *buf = bread(log.dev, log.start);
      struct logheader *hb = (struct logheader *)
4607
                                         (buf->data);
      int i;
4608
4609
      hb->n = log.lh.n;
4610
      for (i = 0; i < log.lh.n; i++) {
4611
        hb->block[i] = log.lh.block[i];
4612

    Write log size (log.lh.n)

4613
      bwrite(buf);
                         into block of the logheader
      brelse(buf);
4614
4615 }
```

```
4600 // Write in-memory log header to disk.
4601 // This is the true point at which the
4602 // current transaction commits.
4603 static void
                                    write head()
4604 write_head(void)
4605 {
4606 struct buf *buf = bread(log.dev, log.start);
      struct logheader *hb = (struct logheader *)
4607
                                         (buf->data);
     int i;
4608
4609
      hb->n = log.lh.n;
4610
      for (i = 0; i < log.lh.n; i++) {
4611
        hb->block[i] = log.lh.block[i];
4612

    Write the entire log

      bwrite(buf);
4613
                         (numbers of blocks in the
4614 brelse(buf);
                         log) into log header
4615 }
```

```
4600 // Write in-memory log header to disk.
4601 // This is the true point at which the
4602 // current transaction commits.
4603 static void
                                    write head()
4604 write_head(void)
4605 {
4606 struct buf *buf = bread(log.dev, log.start);
4607 struct logheader *hb = (struct logheader *)
                                         (buf->data);
     int i;
4608
4609
      hb->n = log.lh.n;
4610
      for (i = 0; i < log.lh.n; i++) {
4611
        hb->block[i] = log.lh.block[i];
4612

    Write block to disk

      bwrite(buf);
4613

    Release

      brelse(buf);
4614
4615 }
```

```
4701 commit()
                                  commit()
4702 {
4703 if (log.lh.n > 0) {
4704
         write_log(); // Write modified blocks
                          from cache to log
4705
         write head(); // Write header to disk --
                          the real commit
         install_trans(); // Now install writes
4706
                             to home locations
4707
         log.lh.n = 0;
         write_head(); // Erase the transaction
4708
                          from the log
4709 }
4710 }
```

```
4570 // Copy committed blocks from log to their home location
4571 static void
4572 install_trans(void)
                                       install trans()
4573 {
4574 int tail;
4575
      for (tail = 0; tail < log.lh.n; tail++) {</pre>
4576
         struct buf *lbuf = bread(log.dev,
4577
                       log.start+tail+1); // read log block
        struct buf *dbuf = bread(log.dev,
4578
                       log.lh.block[tail]); // read dst
        memmove(dbuf->data, lbuf->data, BSIZE); // copy block
4579
                                            // to dst
        bwrite(dbuf); // write dst to disk
4580
4581
        brelse(lbuf);

    Read the block from the log

4582
        brelse(dbuf);
                           area (log.start+tail+1)
4583
      }
4584 }
```

```
4570 // Copy committed blocks from log to their home location
4571 static void
4572 install_trans(void)
                                       install trans()
4573 {
4574 int tail;
4575
4576 for (tail = 0; tail < log.lh.n; tail++) {
         struct buf *lbuf = bread(log.dev,
4577
                       log.start+tail+1); // read log block
         struct buf *dbuf = bread(log.dev,
4578
                       log.lh.block[tail]); // read dst
        memmove(dbuf->data, lbuf->data, BSIZE); // copy block
4579
                                             // to dst
4580
         bwrite(dbuf); // write dst to disk
4581
        brelse(lbuf);

    Read the block where data

4582
        brelse(dbuf);
                           should go on disk
4583
      }
4584 }

    It's a block number in

                              log.lh.block[tail]
```

```
4570 // Copy committed blocks from log to their home location
4571 static void
4572 install_trans(void)
                                       install trans()
4573 {
4574 int tail;
4575
4576 for (tail = 0; tail < log.lh.n; tail++) {
4577
         struct buf *lbuf = bread(log.dev,
                       log.start+tail+1); // read log block
4578
         struct buf *dbuf = bread(log.dev,
                       log.lh.block[tail]); // read dst
        memmove(dbuf->data, lbuf->data, BSIZE); // copy block
4579
                                             // to dst
4580
         bwrite(dbuf); // write dst to disk
4581
        brelse(lbuf);

    Copy data

4582
        brelse(dbuf);
4583 }
4584 }
```

```
4570 // Copy committed blocks from log to their home location
4571 static void
4572 install_trans(void)
                                       install trans()
4573 {
4574 int tail;
4575
     for (tail = 0; tail < log.lh.n; tail++) {</pre>
4576
         struct buf *lbuf = bread(log.dev,
4577
                        log.start+tail+1); // read log block
4578
         struct buf *dbuf = bread(log.dev,
                        log.lh.block[tail]); // read dst
        memmove(dbuf->data, lbuf->data, BSIZE); // copy block
4579
                                             // to dst
4580
         bwrite(dbuf); // write dst to disk
4581
         brelse(lbuf);

    Write the block to disk

4582
         brelse(dbuf);
4583
      }

    Release both blocks

4584 }
```

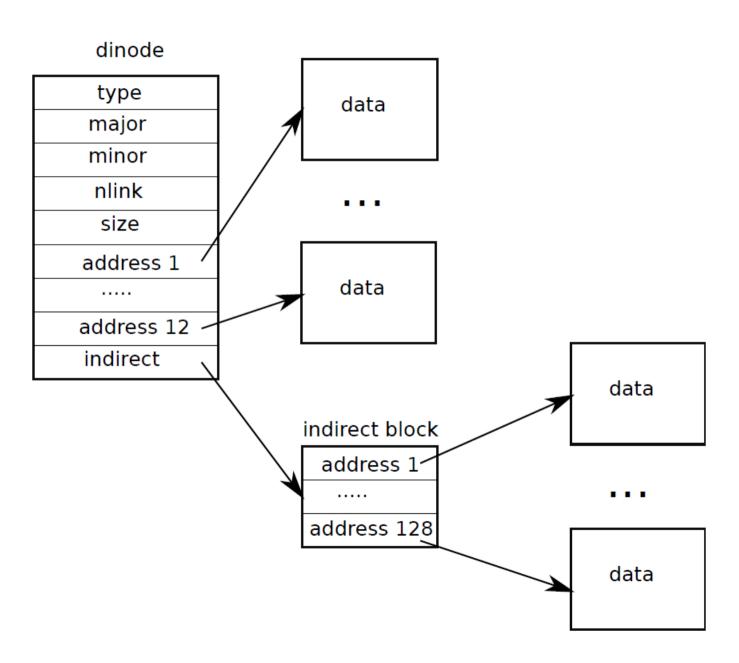
```
4701 commit()
                                  commit()
4702 {
4703 if (log.lh.n > 0) {
4704
         write_log(); // Write modified blocks
                          from cache to log
4705
         write head(); // Write header to disk --
                          the real commit
4706
         install_trans(); // Now install writes
                             to home locations
4707
         log.lh.n = 0;
4708
         write head(); // Erase the transaction
                          from the log
4709 }
4710 }
```

Inode layer

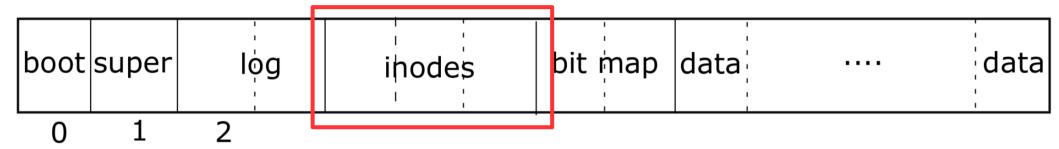
Inode

- Describes a single unnamed file
- The inode on disk holds metadata
 - File type, size, # of links referring to it, list of blocks with data
 - In memory
 - A copy of an on-disk inode + some additional kernel information
 - Reference counter (ip->ref)
 - Synchronization flags (ip->flags)

Representing files on disk



File system layout on disk

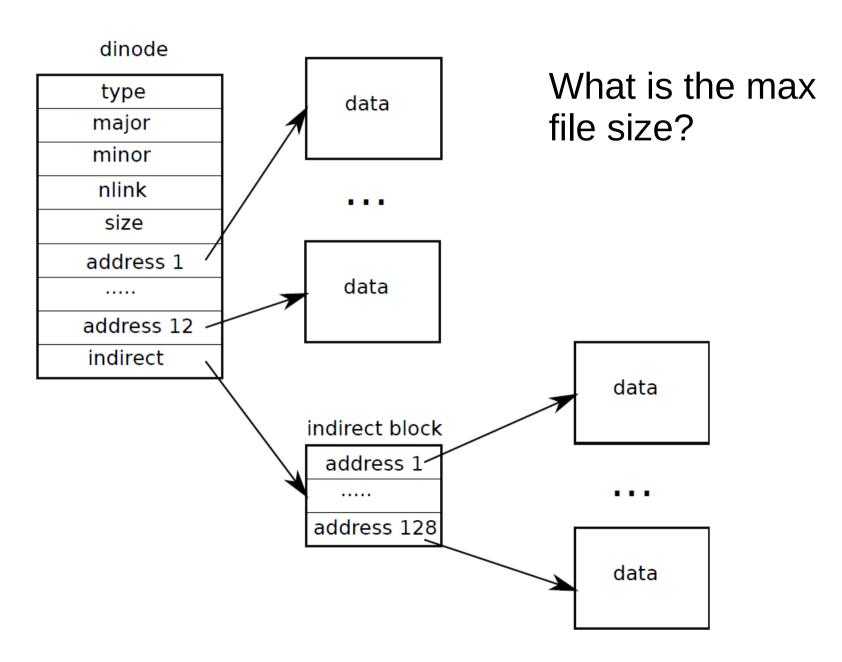


- Inodes are stored as an array on disk
 - sb.startinode
- Each inode has a number (indicating its position on disk)
- The kernel keeps a cache of inodes in memory
 - Synchronization

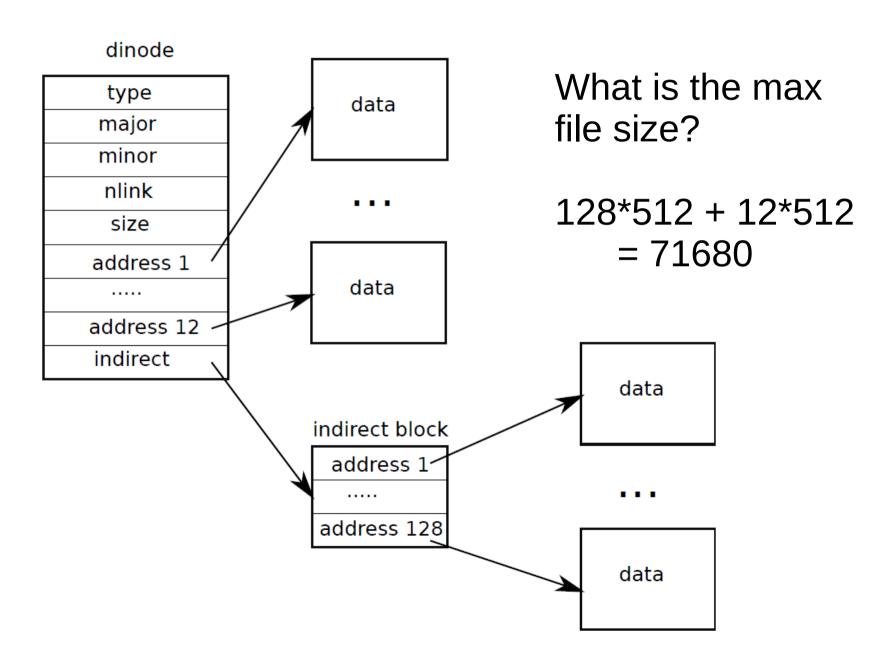
Inode on disk

```
3927 // On-disk inode structure
3928 struct dinode {
      short type; // File type
3929
3930 short major; // Major device number (T_DEV
                       only)
3931 short minor; // Minor device number (T DEV
                       only)
3932 short nlink; // Number of links to inode in
                       file system
3933 uint size; // Size of file (bytes)
3934 uint addrs[NDIRECT+1]; // Data block addresses
3935 };
```

Representing files on disk



Representing files on disk



Inode in memory

```
4011 // in-memory copy of an inode
4012 struct inode {
4013
      uint dev; // Device number
4014
      uint inum; // Inode number
       int ref; // Reference count
4015
4016
       int flags; // I_BUSY, I_VALID
4017
4018
       short type; // copy of disk inode
4019
       short major;
4020
       short minor;
4021
       short nlink;
4022 uint size;
4023 uint addrs[NDIRECT+1];
4024 };
```

In-memory cache of inodes

```
4912 struct {
4913    struct spinlock lock;
4914    struct inode inode[NINODE];
4915 } icache;
```

Lifecycle of inode

- Allocation (on disk)
 - ialloc()
 - iput() -- deallocates
- Referencing in cache
 - ip->ref tracks the number of active pointers to an inode in memory
 - iget()/iput()

Accessing inodes

```
4894 // Thus a typical sequence is:

4895 // ip = iget(dev, inum)

4896 // ilock(ip)

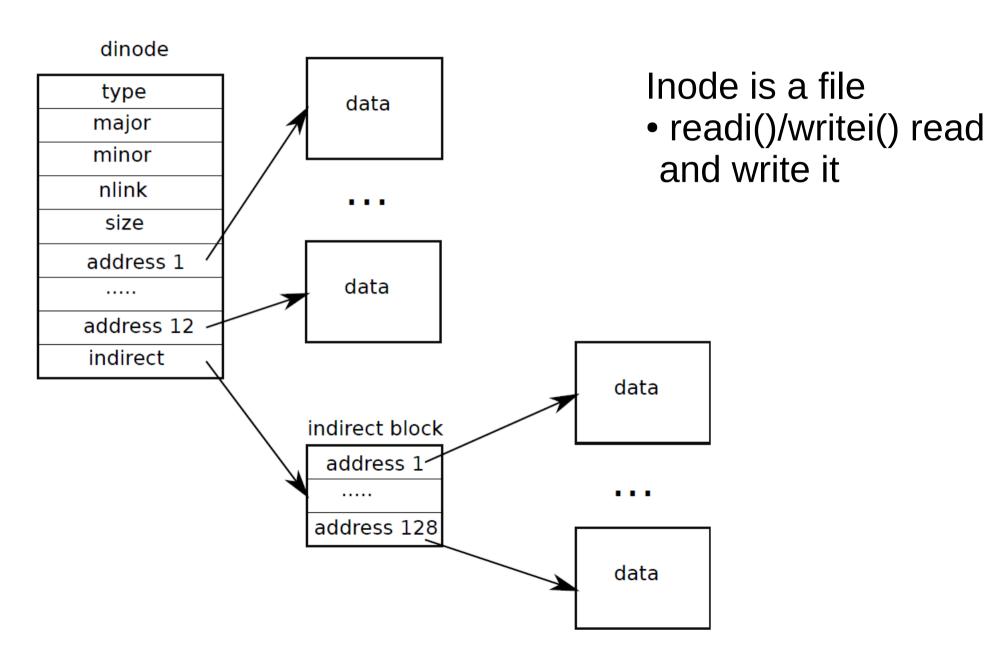
4897 // ... examine and modify ip->xxx ...

4898 // iunlock(ip)

4899 // iput(ip)
```

```
5004 iget(uint dev, uint inum) {
                                                     iget()
. . .
5008
       acquire(&icache.lock);
5010
       // Is the inode already cached?
5011
       emptv = 0;
5012
       for(ip = &icache.inode[0]; ip < &icache.inode[NINODE]; ip++){</pre>
5013
         if(ip->ref > 0 && ip->dev == dev && ip->inum == inum){
5014
           ip->ref++;
           release(&icache.lock);
5015
5016
           return ip;
5017
         }
5018
         if(empty == 0 && ip->ref == 0) // Remember empty slot.
5019
           empty = ip;
5020
. . .
5029 ip->ref = 1;
. . .
5031
       release(&icache.lock);
5033
       return ip;
5034 }
```

Reading and writing inodes



```
5864 int
                                                Example:
5865 sys_read(void)
5866 {
                                              sys_read()
5867 struct file *f;
5868 int n;
5869
    char *p;
5870
      if(argfd(0, 0, &f) < 0 \mid | argint(2, &n) < 0 \mid | argptr(1, &p, n) < 0)
5871
5872
        return -1:
5873
      return fileread(f, p, n);
5874 }
```

- Question:
 - Where does f come from?

```
5816 // Fetch the nth word-sized system call argument as a file descriptor
5817 // and return both the descriptor and the corresponding struct file.
5818 static int
5819 argfd(int n, int *pfd, struct file **pf)
5820 {
                                                  argfd()
5821 int fd;
5822 struct file *f;
5823
5824
      if(argint(n, &fd) < 0)
5825
        return -1;
5826
      if(fd < 0 || fd >= NOFILE || (f=proc->ofile[fd]) == 0)
5827
        return -1;
      if(pfd)
5828

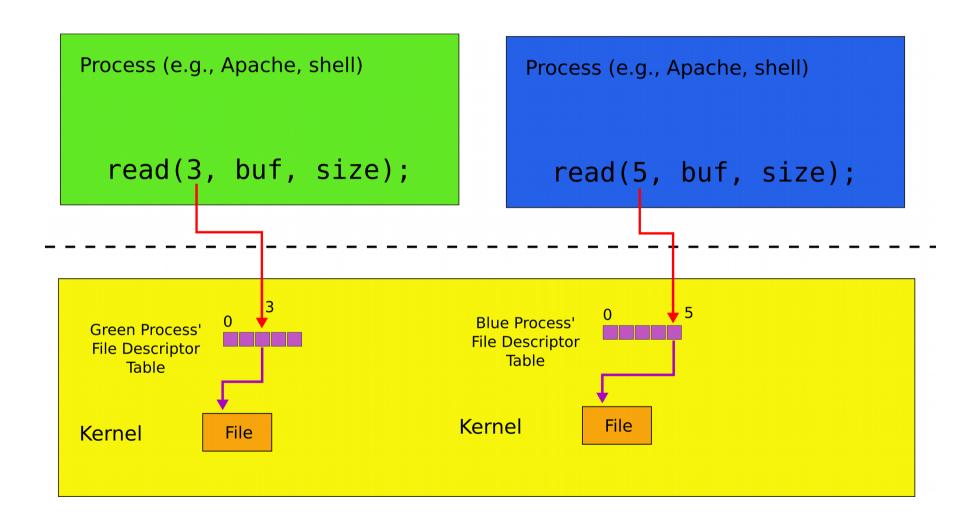
    Remember file descriptors?

        *pfd = fd;
5829
5830
      if(pf)

    Each process has a table

        *pf = f;
5831
5832
      return 0;
                             proc->ofile[]
5833 }
```

File descriptors: two processes



```
2353 struct proc {
2354 uint sz;
                                // Size of process memory (bytes)
2355 pde_t* pgdir;
                                // Page table
2356 char *kstack;
                                // Bottom of kernel stack for this
process
2357 enum procstate state; // Process state
2358
      int pid;
                                // Process ID
      struct proc *parent; // Parent process
2359
2360
      struct trapframe *tf; // Trap frame for current syscall
2361
      struct context *context; // swtch() here to run process
2362 void *chan;
                                // If non-zero, sleeping on chan
2363
      int killed;
                                // If non-zero, have been killed
      struct file *ofile[NOFILE]; // Open files
2364
2365 struct inode *cwd; // Current directory
2366 char name[16];
                    // Process name (debugging)
2367 };
```

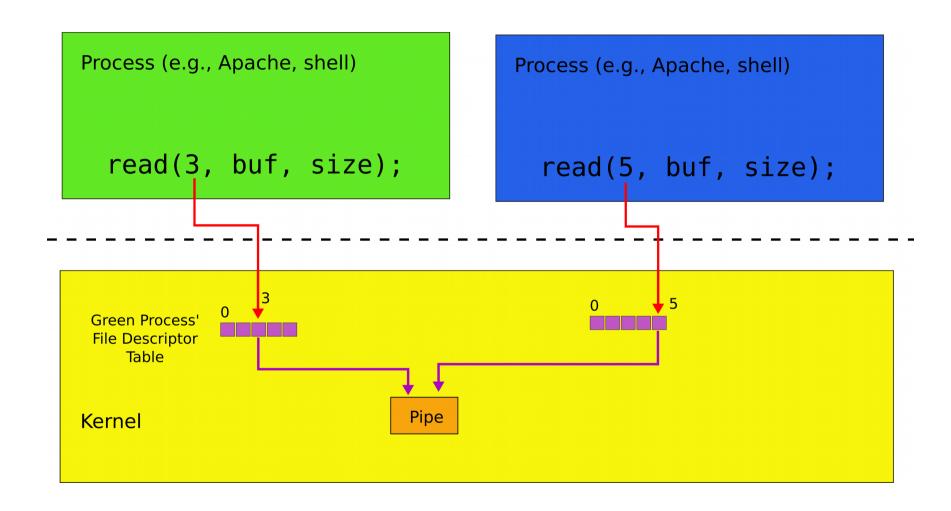
- struct proc has an array of struct file pointers
 - Each element is a "file descriptor"

```
4000 struct file {
4001   enum { FD_NONE, FD_PIPE, FD_INODE } type;
4002   int ref; // reference count
4003   char readable;
4004   char writable;
4005   struct pipe *pipe;
4006   struct inode *ip;
4007   uint off;
4008 };

Struct file
```

- A file can be a pipe or an inode
 - It can be readable and/or writable
 - Each file has current offset (off)

Two file descriptors pointing to a pipe



```
5714 int
5715 fileread(struct file *f, char *addr, int n)
5716 {
                                           readi()
5717 int r;
5718
if(f->readable == 0)
        return -1;
5720
5721
      if(f->type == FD PIPE)
        return piperead(f->pipe, addr, n);
5722
5723
      if(f->type == FD_INODE){
5724
        ilock(f->ip);
        if((r = readi(f->ip, addr, f->off, n)) > 0)
5725
5726
          f \rightarrow off += r;
        iunlock(f->ip);
5727
5728
        return r;
5729 }
5730 panic("fileread");
5731 }
```

```
5714 int
5715 fileread(struct file *f, char *addr, int n)
5716 {
                                           readi()
5717 int r;
5718
if(f->readable == 0)
        return -1;
5720
      if(f->type == FD PIPE)
5721
        return piperead(f->pipe, addr, n);
5722
      if(f->type == FD_INODE){
5723
5724
        ilock(f->ip);
        if((r = readi(f->ip, addr, f->off, n)) > 0)
5725
          f \rightarrow off += r;
5726
        iunlock(f->ip);
5727
5728
        return r;
5729 }
5730 panic("fileread");
5731 }
```

```
5714 int
5715 fileread(struct file *f, char *addr, int n)
5716 {
                                          readi()
5717 int r;
5718
if(f->readable == 0)
        return -1;
5720
if(f->type == FD PIPE)
        return piperead(f->pipe, addr, n);
5722
5723
      if(f->type == FD_INODE){
5724
        ilock(f->ip);
        if((r = readi(f->ip, addr, f->off, n)) > 0)
5725
          f \rightarrow off += r;
5726
        iunlock(f->ip);
5727
5728
        return r;
5729 }
5730 panic("fileread");
5731 }
```

```
5714 int.
5715 fileread(struct file *f, char *addr, int n)
5716 {
                                             readi()
5717
     int r;
5718
if(f->readable == 0)
5720
         return -1;
       if(f->type == FD PIPE)
5721
         return piperead(f->pipe, addr, n);
5722
5723
       if(f->type == FD_INODE){
5724
         ilock(f->ip);
         if((r = readi(f->ip, addr, f->off, n)) > 0)
5725
           f \rightarrow off += r;
5726
         iunlock(f->ip);
5727

    Note

5728
         return r;
5729 }

    Read starts with the

5730 panic("fileread");
                                 current offset (f->off)
5731 }
```

```
5252 readi(struct inode *ip, char *dst, uint off, uint n)
5253 {
5254 uint tot, m;
                                            readi()
5255 struct buf *bp;
5256
       if(off > ip->size || off + n < off)</pre>
5263
         return -1;
5264
5265
       if(off + n > ip->size)
5266
         n = ip \rightarrow size - off;
5267
5268
       for(tot=0; tot<n; tot+=m, off+=m, dst+=m){
         bp = bread(ip->dev, bmap(ip, off/BSIZE));
5269
         m = min(n - tot, BSIZE - off%BSIZE);
5270
         memmove(dst, bp->data + off%BSIZE, m);
5271
         brelse(bp);
5272

    What is this check for?

5273 }
5274
       return n;
5275 }
```

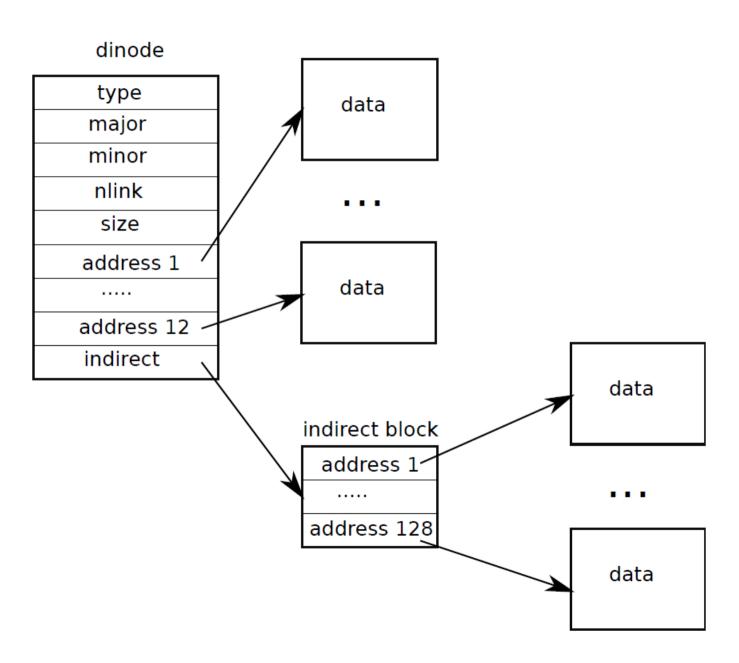
```
5252 readi(struct inode *ip, char *dst, uint off, uint n)
5253 {
5254 uint tot, m;
                                             readi()
5255 struct buf *bp;
5256
. . .
       if(off > ip->size || off + n < off)</pre>
5263
         return -1;
5264
5265
       if(off + n > ip->size)
5266
         n = ip \rightarrow size - off;
5267
5268
       for(tot=0; tot<n; tot+=m, off+=m, dst+=m){</pre>
         bp = bread(ip->dev, bmap(ip, off/BSIZE));
5269
         m = min(n - tot, BSIZE - off%BSIZE);
5270
         memmove(dst, bp->data + off%BSIZE, m);
5271
         brelse(bp);
5272
5273 }
5274 return n;
5275 }
```

```
5252 readi(struct inode *ip, char *dst, uint off, uint n)
5253 {
5254 uint tot, m;
                                            readi()
5255 struct buf *bp;
5256
       if(off > ip->size || off + n < off)</pre>
5263
         return -1;
5264
       if(off + n > ip->size)
5265
5266
         n = ip \rightarrow size - off;
5267
5268
       for(tot=0; tot<n; tot+=m, off+=m, dst+=m){
         bp = bread(ip->dev, bmap(ip, off/BSIZE));
5269
         m = min(n - tot, BSIZE - off%BSIZE);
5270
         memmove(dst, bp->data + off%BSIZE, m);
5271
         brelse(bp);
5272

    What is this bmap()

5273 }
                               function?
5274 return n;
5275 }
```

Representing files on disk



```
5159 static uint
5160 bmap(struct inode *ip, uint bn)
                                                        bmap()
5161 {
. . .
5165
      if(bn < NDIRECT){
5166
        if((addr = ip->addrs[bn]) == 0)
5167
          ip->addrs[bn] = addr = balloc(ip->dev);
5168
        return addr;
5169
5170
     bn -= NDIRECT:
5171
5172
      if(bn < NINDIRECT){</pre>
5173
        // Load indirect block, allocating if necessary.
        if((addr = ip->addrs[NDIRECT]) == 0)
5174
5175
          ip->addrs[NDIRECT] = addr = balloc(ip->dev);
5176
        bp = bread(ip->dev, addr);
        a = (uint*)bp->data;
5177
        if((addr = a[bn]) == 0){
5178
          a[bn] = addr = balloc(ip->dev);
5179
5180
          log write(bp);

    Each inode has some

5181
5182
        brelse(bp);
                                     number (NDIRECT) of
5183
        return addr;
                                     direct pointers
5184 }
```

```
5159 static uint
5160 bmap(struct inode *ip, uint bn)
                                                          bmap()
5161 {
. . .
5165
      if(bn < NDIRECT){
5166
         if((addr = ip->addrs[bn]) == 0)
5167
           ip->addrs[bn] = addr = balloc(ip->dev);
5168
         return addr;
5169
5170
      bn -= NDIRECT:
5171
5172
       if(bn < NINDIRECT){</pre>
5173
         // Load indirect block, allocating if necessary.
         if((addr = ip->addrs[NDIRECT]) == 0)
5174
5175
           ip->addrs[NDIRECT] = addr = balloc(ip->dev);
5176
         bp = bread(ip->dev, addr);
         a = (uint*)bp->data;
5177
         if((addr = a[bn]) == 0){
5178
           a[bn] = addr = balloc(ip->dev);
5179
5180
           log write(bp);

    No it's beyond NDIRECT

5181
5182
         brelse(bp);
5183
         return addr;
5184
```

```
5159 static uint
5160 bmap(struct inode *ip, uint bn)
                                                           bmap()
5161 {
. . .
5165
       if(bn < NDIRECT){</pre>
5166
         if((addr = ip->addrs[bn]) == 0)
5167
           ip->addrs[bn] = addr = balloc(ip->dev);
5168
         return addr;
5169
5170
      bn -= NDIRECT:
5171
5172
       if(bn < NINDIRECT){</pre>
5173
         // Load indirect block, allocating if necessary.
         if((addr = ip->addrs[NDIRECT]) == 0)
5174
5175
           ip->addrs[NDIRECT] = addr = balloc(ip->dev);
5176
         bp = bread(ip->dev, addr);
         a = (uint*)bp->data;
5177
         if((addr = a[bn]) == 0){
5178
           a[bn] = addr = balloc(ip->dev);
5179
           log write(bp);
5180

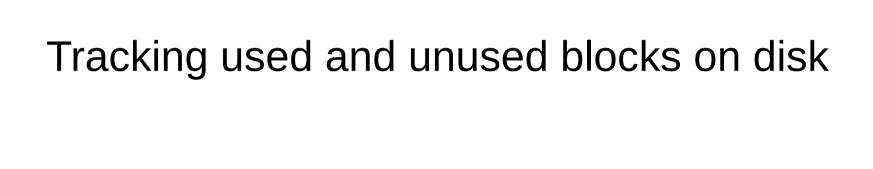
    Read an indirect block

5181
5182
         brelse(bp);
5183
         return addr;
5184
```

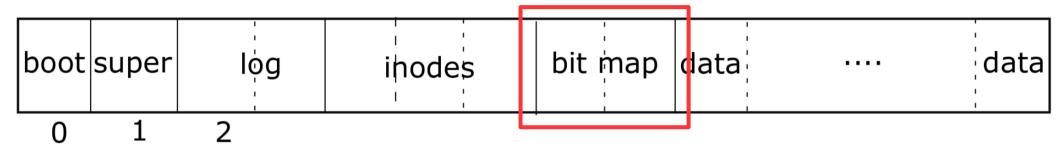
```
5159 static uint
5160 bmap(struct inode *ip, uint bn)
                                                         bmap()
5161 {
. . .
5165
      if(bn < NDIRECT){
5166
         if((addr = ip->addrs[bn]) == 0)
5167
           ip->addrs[bn] = addr = balloc(ip->dev);
5168
        return addr;
5169
5170
      bn -= NDIRECT:
5171
5172
      if(bn < NINDIRECT){</pre>
5173
        // Load indirect block, allocating if necessary.
         if((addr = ip->addrs[NDIRECT]) == 0)
5174
5175
           ip->addrs[NDIRECT] = addr = balloc(ip->dev);
5176
        bp = bread(ip->dev, addr);
        a = (uint*)bp->data;
5177
         if((addr = a[bn]) == 0){
5178
          a[bn] = addr = balloc(ip->dev);
5179
          log write(bp);
5180

    Check if a pointer in the

5181
5182
        brelse(bp);
                                      indirect block is already
5183
        return addr;
                                      allocated
5184
```



Block allocator



- Bitmap of free blocks
 - balloc()/bfree()
- Read the bitmap block by block
 - Scan for a "free" bit
- Access to the bitmap is synchronized with bread()/bwrite()/brelse() operations

```
4802 // Allocate a zeroed disk block.
4803 static uint
4804 balloc(uint dev)
4805 {
4806
       int b, bi, m;
4807
     struct buf *bp;
4808
4809
     bp = 0;
4810
      for(b = 0; b < sb.size; b += BPB) {
4811
         bp = bread(dev, BBLOCK(b, sb));
4812
         for(bi = 0; bi < BPB && b + bi < sb.size; bi++){
4813
           m = 1 << (bi % 8):
           if((bp->data[bi/8] \& m) == 0){ // Is block free?}
4814
4815
             bp->data[bi/8] |= m; // Mark block in use.
4816
             log write(bp);
             brelse(bp);
4817
4818
             bzero(dev, b + bi);
4819
             return b + bi;
4820
4821
         brelse(bp);
4822
4823
4824
       panic("balloc: out of blocks");
4825 }
```

```
4802 // Allocate a zeroed disk block.
4803 static uint
4804 balloc(uint dev)
4805 {
4806
       int b, bi, m;
4807
      struct buf *bp;
4808
4809
     bp = 0;
4810
      for(b = 0; b < sb.size; b += BPB){
4811
         bp = bread(dev, BBLOCK(b, sb));
4812
         for(bi = 0; bi < BPB && b + bi < sb.size; bi++){
4813
           m = 1 << (bi % 8):
4814
           if((bp->data[bi/8] \& m) == 0){ // Is block free?}
4815
             bp->data[bi/8] |= m; // Mark block in use.
4816
             log write(bp);
4817
             brelse(bp);
4818
             bzero(dev, b + bi);
4819
             return b + bi;
4820
4821
         brelse(bp);
4822
4823
4824
       panic("balloc: out of blocks");
4825 }
```

- Check every bit (bi) of a block
 - BPB bits per block

```
4802 // Allocate a zeroed disk block.
4803 static uint
4804 balloc(uint dev)
4805 {
4806
       int b, bi, m;
4807
     struct buf *bp;
4808
4809
     bp = 0;
4810
      for(b = 0; b < sb.size; b += BPB){
4811
      bp = bread(dev, BBLOCK(b, sb));
4812
         for(bi = 0; bi < BPB && b + bi < sb.size; bi++){
4813
           m = 1 << (bi % 8):
           if((bp->data[bi/8] \& m) == 0){ // Is block free?}
4814
4815
             bp->data[bi/8] |= m; // Mark block in use.
4816
             log write(bp);
             brelse(bp);
4817
4818
             bzero(dev, b + bi);
4819
             return b + bi;
4820
4821
         brelse(bp);
4822
4823
4824
       panic("balloc: out of blocks");
4825 }
```

```
4802 // Allocate a zeroed disk block.
4803 static uint
4804 balloc(uint dev)
4805 {
4806
       int b, bi, m;
4807
     struct buf *bp;
4808
4809
     bp = 0;
4810
     for(b = 0; b < sb.size; b += BPB){
4811
         bp = bread(dev, BBLOCK(b, sb));
4812
         for(bi = 0; bi < BPB && b + bi < sb.size; bi++){
4813
           m = 1 << (bi % 8):
           if((bp->data[bi/8] \& m) == 0){ // Is block free?}
4814
4815
             bp->data[bi/8] |= m; // Mark block in use.
4816
             log write(bp);
             brelse(bp);
4817
4818
             bzero(dev, b + bi);
4819
             return b + bi;
4820
4821
         brelse(bp);
4822
4823
4824
       panic("balloc: out of blocks");
4825 }
```

 Why do we need log_write() instead of bwrite()?

Directory layer

Directory inodes

- A directory inode is a sequence of directory entries and inode numbers
 - Each name is max of 14 characters
 - Has a special inode type T_DIR
- dirlookup() searches for a directory with a given name
- dirlink() adds new file to a directory

Directory entry

```
3965 struct dirent {
3966 ushort inum;
3967 char name[DIRSIZ];
3968 };
```

```
5360 struct inode*
5361 dirlookup(struct inode *dp, char *name, uint *poff)
5362 {
. . .
                                              dirlookup()
       if(dp->type != T_DIR)
5366
5367
         panic("dirlookup not DIR");
5368
5369
       for(off = 0; off < dp->size; off += sizeof(de)){
5370
         if(readi(dp, (char*)&de, off, sizeof(de)) != sizeof(de))
           panic("dirlink read");
5371
5372
         if(de.inum == 0)
5373
           continue:
         if(namecmp(name, de.name) == 0){
5374
5375
           // entry matches path element

    Inode is a directory

5376
           if(poff)
             *poff = off;
5377
5378
             inum = de.inum;
5379
             return iget(dp->dev, inum);
           }
5380
5381
      }
5382
5383
       return 0;
5384 }
```

```
5360 struct inode*
5361 dirlookup(struct inode *dp, char *name, uint *poff)
5362 {
. . .
                                              dirlookup()
5366
       if(dp->type != T DIR)
5367
        panic("dirlookup not DIR");
5368
5369
       for(off = 0; off < dp->size; off += sizeof(de)){
5370
         if(readi(dp, (char*)&de, off, sizeof(de)) != sizeof(de))
           panic("dirlink read");
5371
5372
         if(de.inum == 0)
5373
           continue:
5374
         if(namecmp(name, de.name) == 0){
5375
           // entry matches path element

    Iterate through all

5376
           if(poff)
             *poff = off;
5377
                                            entries?
5378
             inum = de.inum;
5379
             return iget(dp->dev, inum);
           }
5380
5381
      }
5382
5383
       return 0;
5384 }
```

```
5360 struct inode*
5361 dirlookup(struct inode *dp, char *name, uint *poff)
5362 {
. . .
                                             dirlookup()
5366
      if(dp->type != T DIR)
5367
        panic("dirlookup not DIR");
5368
5369
      for(off = 0; off < dp->size; off += sizeof(de)){
5370
         if(readi(dp, (char*)&de, off, sizeof(de)) != sizeof(de))
          panic("dirlink read");
5371
5372
         if(de.inum == 0)
5373
          continue:
         if(namecmp(name, de.name) == 0){
5374
5375
          // entry matches path element

    Read the inode

5376
          if(poff)
            *poff = off;
5377
                                           Compare names
5378
            inum = de.inum;
5379
            return iget(dp->dev, inum);
          }
5380
5381
      }
5382
5383
      return 0;
5384 }
```

Path names layer

- Series of directory lookups to resolve a path
 - E.g. /usr/bin/sh
- Namei() resolves a path into an inode
 - If path starts with "/" evaluation starts at the root
 - Otherwise current directory

namei()

```
5539 struct inode*
5540 namei(char *path)
5541 {
5542    char name[DIRSIZ];
5543    return namex(path, 0, name);
5544 }
```

```
5505 namex(char *path, int nameiparent, char *name)
5506 {
. . .
       if(*path == ',')
5509
         ip = iget(ROOTDEV, ROOTINO);
5510
5511
       else
5512
         ip = idup(proc->cwd);
       // skipelem("a/bb/c", name) = "bb/c", setting name = "a"
5513
       while((path = skipelem(path, name)) != 0){
5514
5515
         ilock(ip);
         if(ip->type != T_DIR){
5516
5517
           iunlockput(ip);
5518
           return 0;
5519
. . .
5525
         if((next = dirlookup(ip, name, 0)) == 0){
5526
           iunlockput(ip);
5527
           return 0;
5528
5529
         iunlockput(ip);
5530
         ip = next;
5531
5532
       if(nameiparent){
5533
         iput(ip);
5534
         return 0;
5535
5536
       return ip;
5537 }
```

namex()

 If path == "/" start with the inode number of the root

```
5505 namex(char *path, int nameiparent, char *name)
5506 {
                                                           namex()
. . .
      if(*path == ',')
5509
5510
        ip = iget(ROOTDEV, ROOTINO);
5511
      else
5512
        ip = idup(proc->cwd);
      // skipelem("a/bb/c", name) = "bb/c", setting name = "a"
5513
      while((path = skipelem(path, name)) != 0){
5514
5515
        ilock(ip);
5516
        if(ip->type != T DIR){
5517
          iunlockput(ip);
5518
          return 0;
5519
. . .
5525
        if((next = dirlookup(ip, name, 0)) == 0){
5526
          iunlockput(ip);

    Get the prefix

5527
          return 0;
5528
5529
        iunlockput(ip);
                                                  "a/bb/c"
5530
        ip = next;
5531

    name = a

5532
      if(nameiparent){
5533
        iput(ip);
                                                      - path = "bb/c"
        return 0;
5534
5535
5536
      return ip;
5537 }
```

```
5505 namex(char *path, int nameiparent, char *name)
5506 {
                                                           namex()
. . .
      if(*path == ',')
5509
5510
        ip = iget(ROOTDEV, ROOTINO);
5511
      else
5512
        ip = idup(proc->cwd);
      // skipelem("a/bb/c", name) = "bb/c", setting name = "a"
5513
      while((path = skipelem(path, name)) != 0){
5514
5515
        ilock(ip);
        if(ip->type != T_DIR){
5516
5517
          iunlockput(ip);
5518
          return 0;
5519
. . .
5525
        if((next = dirlookup(ip, name, 0)) == 0){
5526
           iunlockput(ip);
5527
          return 0;

    Lookup that name in

5528
                                                the directory
5529
        iunlockput(ip);
5530
        ip = next;
5531
5532
      if(nameiparent){
5533
        iput(ip);
5534
        return 0;
5535
5536
      return ip;
5537 }
```

```
6101 sys_open(void)
6102 {
. . .
       if(argstr(0, &path) < 0 || argint(1, &omode) < 0)</pre>
6108
6109
         return -1;
6110
6111
       begin op();
6112
. . .
6120
         if((ip = namei(path)) == 0){
6121
           end op();
6122
           return -1;
6123
         }
. . .
       if((f = filealloc()) == 0 \mid | (fd = fdalloc(f)) < 0){
6132
6133
         if(f)
6134
         fileclose(f);
         iunlockput(ip);
6135
6136
         end op();
         return -1;
6137
6138
6139
       iunlock(ip);
6140
       end_op();
6141
6142
       f->type = FD_INODE;
6143
       f \rightarrow ip = ip;
. . .
6147
       return fd;
6148 }
```

Eaxmple: sys_open

File descriptor layer

File descriptors

- Uniform access to
 - Files
 - Devices, e.g., console
 - Pipes

```
4000 struct file {
4001    enum { FD_NONE, FD_PIPE, FD_INODE } type;
4002    int ref; // reference count
4003    char readable;
4004    char writable;
4005    struct pipe *pipe;
4006    struct inode *ip;
4007    uint off;
4008 };
```

```
6101 sys open(void)
6102 {
. . .
       if(argstr(0, &path) < 0 || argint(1, &omode) < 0)</pre>
6108
6109
         return -1:
6110
6111
       begin op();
6112
. . .
6120
         if((ip = namei(path)) == 0){
            end op();
6121
6122
            return -1;
6123
          }
. . .
       if(f = filealloc()) == 0 \mid \mid (fd = fdalloc(f)) < 0)
6132
6133
          if(f)
           fileclose(f);
6134
         iunlockput(ip);
6135
6136
         end op();
6137
         return -1;
6138
6139
       iunlock(ip);
6140
       end_op();
6141
6142
       f->type = FD_INODE;
6143
       f \rightarrow ip = ip;
. . .
6147
       return fd;
6148 }
```

Eaxmple: sys_open

 Allocate new file data structure

```
5612 struct {
5613 struct spinlock lock;
5614 struct file file[NFILE];
5615 } ftable:
5624 struct file*
5625 filealloc(void)
5626 {
5627
     struct file *f;
5628
5629
      acquire(&ftable.lock);
5630
       for(f = ftable.file; f < ftable.file + NFILE; f++){</pre>
5631
         if(f->ref == 0){
5632
           f \rightarrow ref = 1:
5633
           release(&ftable.lock);
5634
           return f;
5635
5636 }
5637
       release(&ftable.lock);
5638
       return 0;
5639 }
```

Files and filealloc()

- Linear search for an available element of the ftable array
 - f->ref == 0

```
6101 sys open(void)
6102 {
. . .
       if(argstr(0, &path) < 0 || argint(1, &omode) < 0)</pre>
6108
6109
         return -1:
6110
6111
       begin op();
6112
. . .
6120
         if((ip = namei(path)) == 0){
           end op();
6121
6122
           return -1;
6123
. . .
       if((f = filealloc()) == 0 | | (fd = fdalloc(f)) < 0)
6132
6133
         if(f)
          fileclose(f);
6134
6135
         iunlockput(ip);
6136
         end op();
6137
         return -1;
6138
6139
       iunlock(ip);
6140
       end_op();
6141
6142
       f->type = FD_INODE;
       f \rightarrow ip = ip;
6143
. . .
6147
       return fd;
6148 }
```

Eaxmple: sys_open

Allocate a new file descriptor

```
5835 // Allocate a file descriptor for the given file.
5836 // Takes over file reference from caller on
success.
                               File descriptors
5837 static int
                                 and fdalloc()
5838 fdalloc(struct file *f)
5839 {
5840 int fd;
5841
       for(fd = 0; fd < NOFILE; fd++){</pre>
5842
         if(proc->ofile[fd] == 0){
5843
           proc->ofile[fd] = f;
5844
5845
           return fd;

    Allocate a file

5846
                              descriptor
5847 }
5848
      return -1;
5849 }
```

Pipes

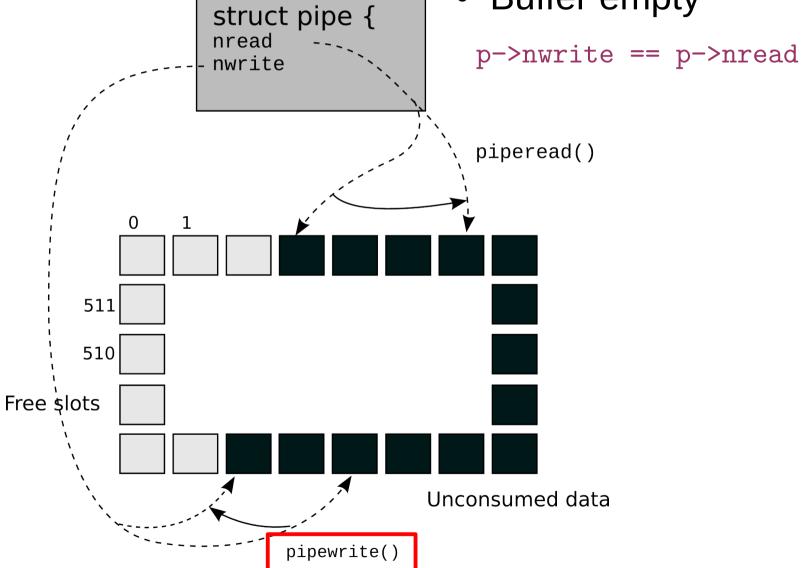
```
Pipe
6459 #define PIPESIZE 512
6460
6461 struct pipe {
6462
      struct spinlock lock;
6463 char data[PIPESIZE];
6464 uint nread; // number of bytes read
6465 uint nwrite; // number of bytes written
6466
      int readopen; // read fd is still open
6467
      int writeopen; // write fd is still open
6468 };
```

```
Pipe
6459 #define PIPESIZE 512
6460
6461 struct pipe {
      struct spinlock lock;
6462
6463 char data[PIPESIZE];
6464 uint nread; // number of bytes read
6465 uint nwrite; // number of bytes written
6466
      int readopen; // read fd is still open
6467
      int writeopen; // write fd is still open
6468 };
```

Buffer full

p->nwrite == p->nread + PIPESIZE

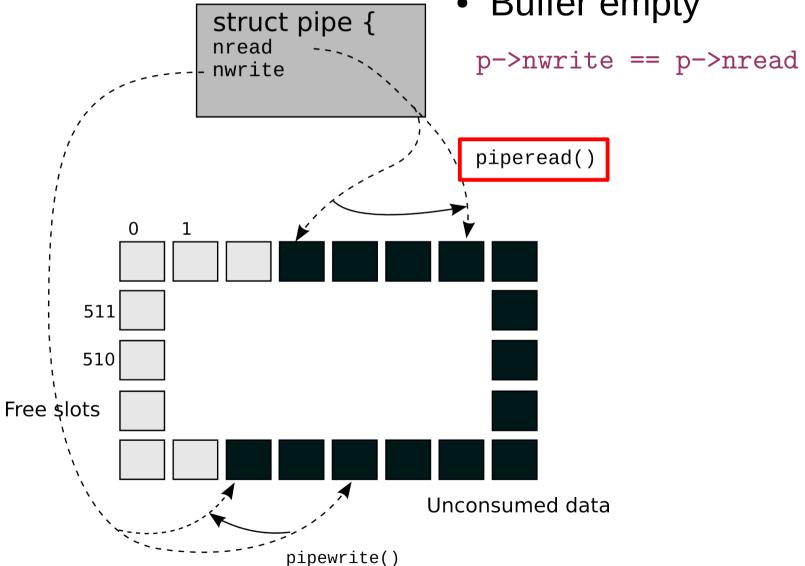
Buffer empty



Buffer full

p->nwrite == p->nread + PIPESIZE

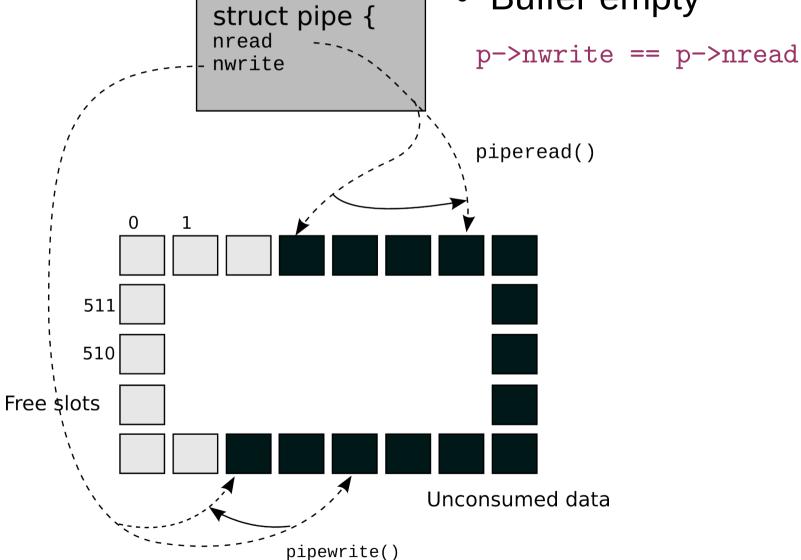




Buffer full

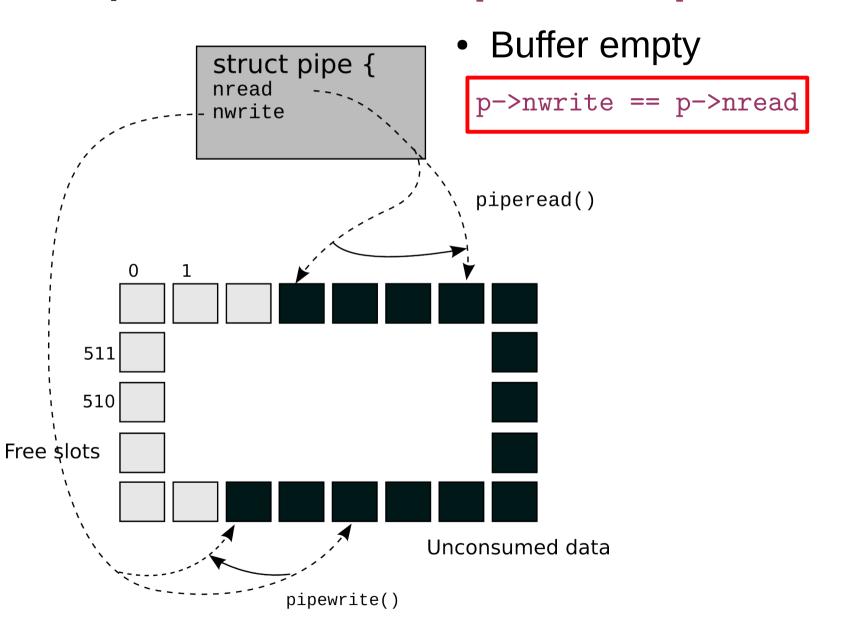
p->nwrite == p->nread + PIPESIZE

Buffer empty



Buffer full

p->nwrite == p->nread + PIPESIZE



```
6250 int
6251 sys_pipe(void)
6252 {
6253
      int *fd;
6254
      struct file *rf, *wf;
6255
      int fd0, fd1;
6256
      if(argptr(0, (void*)&fd, 2*sizeof(fd[0])) < 0)
6257
6258
       return -1;
       if(pipealloc(&rf, &wf) < 0)</pre>
6259
6260
         return -1:
      fd0 = -1;
6261
      if((fd0 = fdalloc(rf)) < 0 \mid | (fd1 = fdalloc(wf)) < 0){
6262
       if(fd0 >= 0)
6263
      proc->ofile[fd0] = 0;
6264
6265
       fileclose(rf);
       fileclose(wf);
6266
6267
        return -1;
6268
      }
6269
      fd[0] = fd0;
6270
      fd[1] = fd1;
6271
      return 0;
6272 }
```

pipe(): create a new pipe

 Allocate pipe object

```
6470 int
6471 pipealloc(struct file **f0, struct file **f1)
6472 {
       struct pipe *p;
6473
6474
      p = 0;
6475
6476
      *f0 = *f1 = 0:
     if((*f0 = filealloc()) == 0 \mid | (*f1 = filealloc()) == 0)
6477
6478
       goto bad;
6479
       if((p = (struct pipe*)kalloc()) == 0)
6480
        goto bad;
      p->readopen = 1;
6481
6482
      p->writeopen = 1;
      p->nwrite = 0;
6483
6484
      p->nread = 0;
6485
       initlock(&p->lock, "pipe");
       (*f0)->type = FD_PIPE;
6486
       (*f0)->readable = 1;
6487
6488
       (*f0)->writable = 0;
       (*f0)->pipe = p;
6489
       (*f1)->type = FD_PIPE;
6490
6491
       (*f1)->readable = 0;
       (*f1)->writable = 1;
6492
       (*f1)->pipe = p;
6493
6494
      return 0;
. . .
```

pipealloc()

- One end is writable
- One is readable

```
6250 int
6251 sys_pipe(void)
6252 {
6253
      int *fd;
6254
      struct file *rf, *wf;
6255
      int fd0, fd1;
6256
      if(argptr(0, (void*)&fd, 2*sizeof(fd[0])) < 0)
6257
6258
       return -1;
       if(pipealloc(&rf, &wf) < 0)</pre>
6259
6260
       return -1;
      fd0 = -1;
6261
       if((fd0 = fdalloc(rf)) < 0 \mid | (fd1 = fdalloc(wf)) < 0){
6262
      if(fd0 >= 0)
6263
      proc->ofile[fd0] = 0;
6264
6265
      fileclose(rf);
      fileclose(wf);
6266
6267
        return -1;
6268
      }
6269
      fd[0] = fd0;
      fd[1] = fd1;
6270
6271
      return 0;
6272 }
```

pipe(): create a new pipe

- Find unopened file descriptors
 - You need two

```
6551 piperead(struct pipe *p, char *addr, int n)
6552 {
6553
       int i:
6554
6555
       acquire(&p->lock);
6556
       while(p->nread == p->nwrite && p->writeopen){
6557
         if(proc->killed){
6558
           release(&p->lock);
           return -1:
6559
         }
6560
6561
         sleep(&p->nread, &p->lock);
6562
       for(i = 0; i < n; i++){
6563
6564
         if(p->nread == p->nwrite)
6565
           break:
         addr[i] = p->data[p->nread++ % PIPESIZE];
6566
6567
6568
       wakeup(&p->nwrite);
       release(&p->lock);
6569
6570
       return i;
6571 }
```

piperead()

- Acquire pipe lock
 - All pipe
 operations are
 are protected
 with the lock

```
6551 piperead(struct pipe *p, char *addr, int n)
6552 {
6553
       int i:
6554
6555
       acquire(&p->lock);
6556
       while(p->nread == p->nwrite && p->writeopen){
         if(proc->killed){
6557
6558
           release(&p->lock);
           return -1;
6559
         }
6560
6561
         sleep(&p->nread, &p->lock);
6562
       for(i = 0; i < n; i++){
6563
6564
         if(p->nread == p->nwrite)
6565
           break:
         addr[i] = p->data[p->nread++ % PIPESIZE];
6566
6567
6568
       wakeup(&p->nwrite);
       release(&p->lock);
6569
6570
       return i;
6571 }
```

piperead()

- If the buffer is empty && the write end is still open
 - Go to sleep

```
6551 piperead(struct pipe *p, char *addr, int n)
6552 {
6553
       int i:
6554
6555
       acquire(&p->lock);
6556
       while(p->nread == p->nwrite && p->writeopen){
6557
         if(proc->killed){
6558
           release(&p->lock);
           return -1:
6559
         }
6560
6561
         sleep(&p->nread, &p->lock);
6562
       for(i = 0; i < n; i++){
6563
6564
         if(p->nread == p->nwrite)
6565
           break:
         addr[i] = p->data[p->nread++ % PIPESIZE];
6566
6567
6568
       wakeup(&p->nwrite);
       release(&p->lock);
6569
6570
       return i;
6571 }
```

piperead()

- After reading some data from the buffer
 - Wakeup the writer

```
6530 pipewrite(struct pipe *p, char *addr, int n)
6531 {
6532
       int i;
6533
6534
       acquire(&p->lock);
6535
       for(i = 0; i < n; i++){
         while(p->nwrite == p->nread + PIPESIZE){
6536
           if(p->readopen == 0 || proc->killed){
6537
6538
             release(&p->lock);
6539
             return -1;
6540
           wakeup(&p->nread);
6541
           sleep(&p->nwrite, &p->lock);
6542
6543
         }
         p->data[p->nwrite++ % PIPESIZE] = addr[i];
6544
6545
6546
       wakeup(&p->nread);
6547
       release(&p->lock);
6548
       return n;
6549 }
```

pipewrite()

- If the buffer is full
 - Wakeup reader
 - Go to sleep

```
6530 pipewrite(struct pipe *p, char *addr, int n)
6531 {
6532
       int i;
6533
6534
       acquire(&p->lock);
6535
       for(i = 0; i < n; i++){
6536
         while(p->nwrite == p->nread + PIPESIZE){
6537
           if(p->readopen == 0 || proc->killed){
6538
             release(&p->lock);
6539
             return -1;
6540
           }
           wakeup(&p->nread);
6541
6542
           sleep(&p->nwrite, &p->lock);
         }
6543
         p->data[p->nwrite++ % PIPESIZE] = addr[i];
6544
6545
       wakeup(&p->nread);
6546
       release(&p->lock);
6547
6548
       return n;
6549 }
```

pipewrite()

- If the buffer is full
 - Wakeup reader
 - Go to sleep
- However if the read end is closed
 - Return an error
 - (-1)

```
6530 pipewrite(struct pipe *p, char *addr, int n)
6531 {
6532
       int i;
6533
6534
       acquire(&p->lock);
       for(i = 0; i < n; i++){
6535
6536
         while(p->nwrite == p->nread + PIPESIZE){
6537
           if(p->readopen == 0 || proc->killed){
6538
             release(&p->lock);
6539
             return -1;
6540
           wakeup(&p->nread);
6541
6542
           sleep(&p->nwrite, &p->lock);
6543
         p->data[p->nwrite++ % PIPESIZE] = addr[i];
6544
6545
       wakeup(&p->nread);
6546
       release(&p->lock);
6547
6548
       return n;
6549 }
```

pipewrite()

- Otherwise keep writing bytes into the pipe
- When done
 - Wakeup reader

Thank you!

Example: write system call

```
5476 int
                                  Write() syscall
5477 sys_write(void)
5478 {
5479 struct file *f;
5480 int n;
5481 char *p;
5482
       if(argfd(0, 0, &f) < 0
5483
        || \operatorname{argint}(2, \&n) < 0 || \operatorname{argptr}(1, \&p, n) < 0)
5484
          return -1;
5485 return filewrite(f, p, n);
5486 }
```

```
5352 filewrite(struct file *f, char *addr, int n)
5353 {
if(f->type == FD_INODE)
. . .
5368
        int i = 0;
5369 while(i < n){
5373
          begin_trans();
5374
          ilock(f->ip);
5375
          if ((r = writei(f->ip, addr + i, f->off, n1)) > 0)
5376
5377
            f \rightarrow off += r;
          iunlock(f->ip);
5378
          commit_trans();
5379
                                     Write several
5386 }
                                   blocks at a time
5390 }
```

```
4570 // Copy committed blocks from log to their home location
4571 static void
4572 install trans(void)
4573 {
4574
       int tail;
4575
4576
       for (tail = 0; tail < log.lh.n; tail++) {</pre>
4577
         struct buf *lbuf = bread(log.dev, log.start+tail+1); // read log
block
4578
         struct buf *dbuf = bread(log.dev, log.lh.block[tail]); // read dst
4579
         memmove(dbuf->data, lbuf->data, BSIZE); // copy block to dst
4580
         bwrite(dbuf); // write dst to disk
4581
         brelse(lbuf):
4582
         brelse(dbuf);
4583 }
4584 }
```

Example