143A: Operating Systems

Lecture 3: Stack, function invocations, and calling conventions

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Last two lectures: x86 instruction set

- Three main groups
 - Data movement (from memory and between registers)
 - Arithmetic operations (addition, subtraction, etc.)
 - Control flow (jumps, function calls)

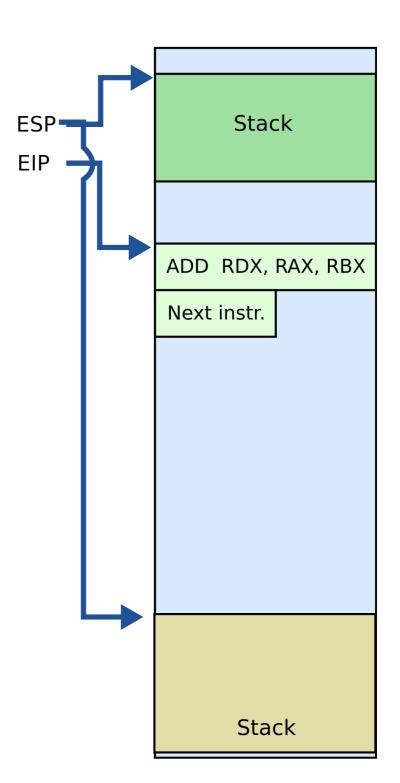
This is enough to write all programs you can think of

Stack and procedure calls

What is stack?

Stack

- It's just a region of memory
 - Pointed by a special register ESP
- You can change ESP
 - Get a new stack



Why do we need stack?

Calling functions

```
// some code...
foo();
// more code..
```

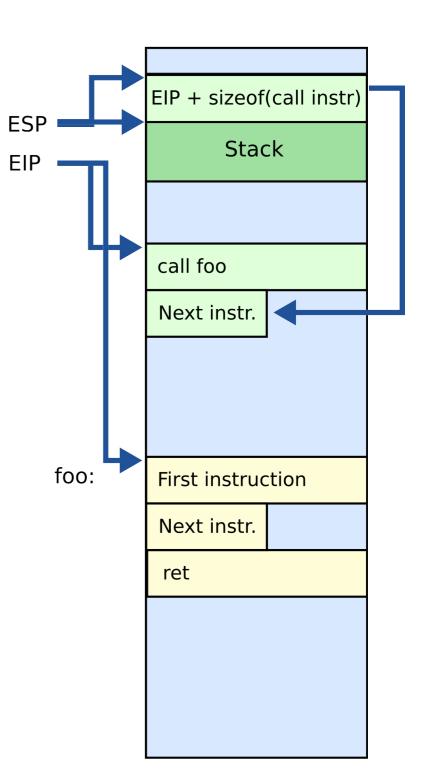
- Stack contains information for how to return from a subroutine
 - i.e., from foo()

 Functions can be called from different places in the program

```
if (a == 0) {
    foo();
    ...
} else {
    foo();
    ...
}
```

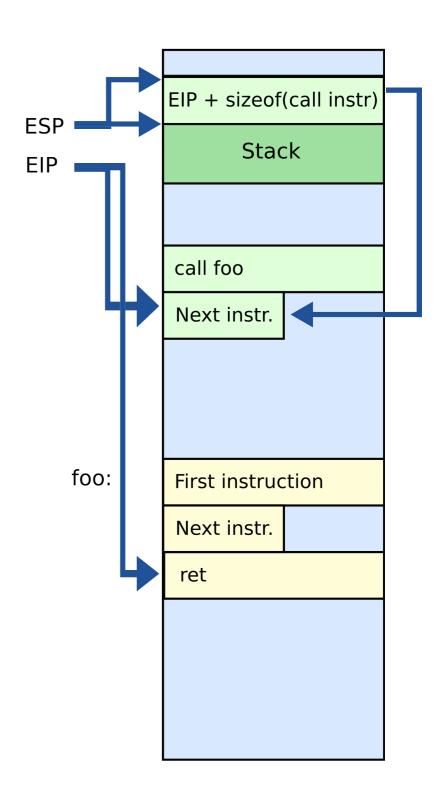
Stack

- Main purpose:
 - Store the return address for the current procedure
 - Caller pushes return address on the stack
 - Callee pops it and jumps



Stack

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 - Store the return address for the current procedure
 - Caller pushes return address on the stack
 - Callee pops it and jumps



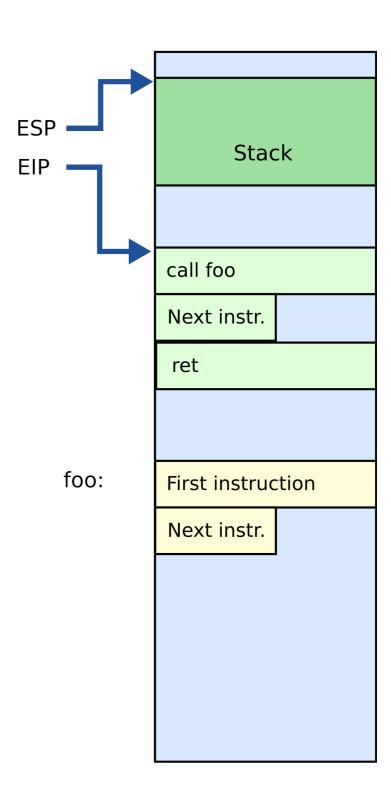
Call/return

- CALL instruction
 - Makes an unconditional jump to a subprogram and pushes the address of the next instruction on the stack

- RET instruction
 - Pops off an address and jumps to that address

Stack

- Other uses:
 - Local data storage
 - Parameter passing
 - Evaluation stack
 - Register spill

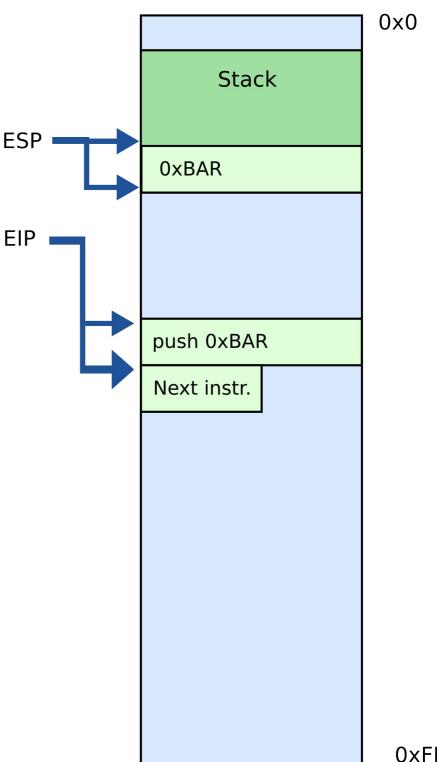


Manipulating stack

- ESP register
 - Contains the memory address of the topmost element in the stack
- PUSH instruction

push OxBAR

- Subtract 4 from ESP
- Insert data on the stack



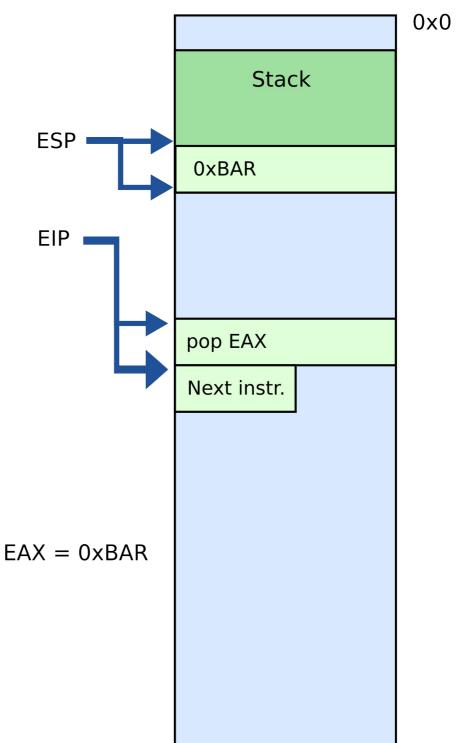
OxFFFFFFF

Manipulating stack

POP instruction

pop EAX

- Removes data from the stack
- Saves in register or memory
- Adds 4 to ESP



0xFFFFFFF

Calling conventions

Calling conventions

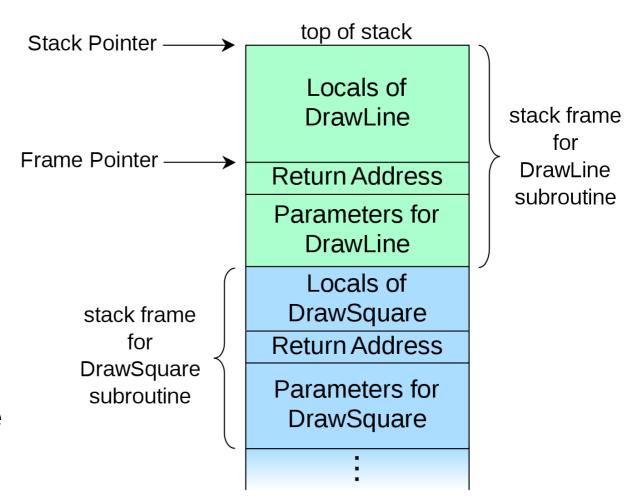
- Goal: re-entrant programs
 - How to pass arguments
 - On the stack?
 - In registers?
 - How to return values
 - On the stack?
 - In registers?
- Conventions differ from compiler, optimizations, etc.

Maintain stack as frames

 Each function has a new frame

```
void DrawSquare(...)
{
     ...
     DrawLine(x, y, z);
}
```

- Use dedicated register
 EBP (frame pointer)
 - Points to the base of the frame

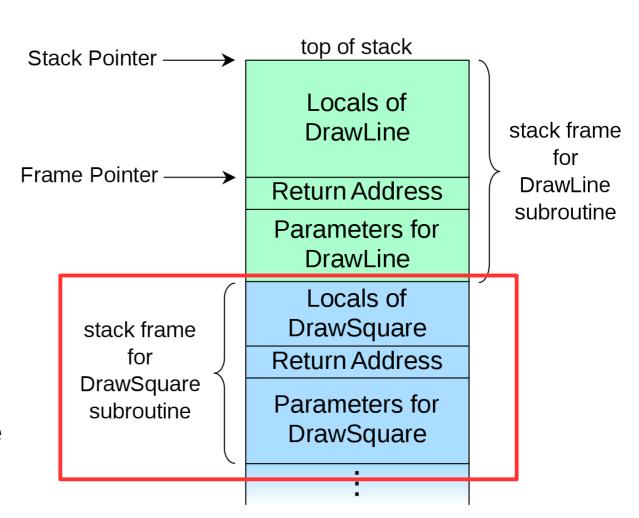


Maintain stack as frames

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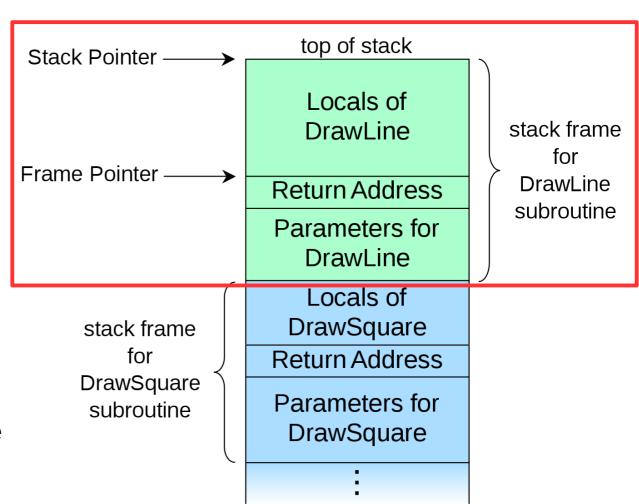


Stack consists of frames

 Each function has a new frame

```
void DrawSquare(...)
{
    ...
    DrawLine(x, y, z);
}
```

- Use dedicated register
 EBP (frame pointer)
 - Points to the base of the frame



Prologue/epilogue

- Each function maintains the frame
 - A dedicated register EBP is used to keep the frame pointer
 - Each function uses prologue code (blue), and epilogue (yellow) to maintain the frame

Local variables

What types of variables do you know?

Or where these variables are in memory?

What types of variables do you know?

- Global variables
 - Initialized → data section
 - Uninitalized → BSS
- Dynamic variables
 - Heap
- Local variables
 - Stack

Global variables

```
1. #include <stdio.h>
2.
3. char hello[] = "Hello";
4. int main(int ac, char **av)
5. {
       static char world[] = "world!";
6.
       printf("%s %s\n", hello, world);
7.
8.
      return 0;
9. }
```

Global variables

```
1. #include <stdio.h>
2.
3. char hello[] = "Hello";
4. int main(int ac, char **av)
5. {
6. static char world[] = "world!";
7. printf("%s %s\n", hello, world);
8. return 0;
9. }
```

- Allocated in the data section
 - It is split in initialized (non-zero), and non-initialized (zero)
 - As well as read/write, and read only data section

Dynamic variables (heap)

```
1. #include <stdio.h>
2. #include <string.h>
3. #include <stdlib.h>
4.
5. char hello[] = "Hello";
6. int main(int ac, char **av)
7. {
8.
     char world[] = "world!";
9.
      char *str = malloc(64);
      memcpy(str, "beautiful", 64);
10.
11. printf("%s %s %s\n", hello, str, world);
12.
      return 0;
13.}
```

Dynamic variables (heap)

```
1. #include <stdio.h>
2. #include <string.h>
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      memcpy(str, "beautiful", 64);
      printf("%s %s %s\n", hello, str, world);
11.
12.
      return 0:
13.}
```

- Allocated on the heap
 - Special area of memory provided by the OS from where malloc() can allocate memory

Local variables

Local variables

```
1. #include <stdio.h>
2.
3. char hello[] = "Hello";
4. int main(int ac, char **av)
5. {
6.
      //static char world[] = "world!";
      char world[] = "world!";
7.
      printf("%s %s\n", hello, world);
8.
9.
      return 0;
10.}
```

Local variables...

Each function has private instances of local variables

```
foo(int x) {
    int a, b, c;
    ...
    return;
}
```

Function can be called recursively

```
foo(int x) {
    int a, b, c;
    a = x + 1;
    if ( a < 100 )
        foo(a);
    return;
}</pre>
```

How to allocate local variables?

```
void my_function()
{
    int a, b, c;
    ...
}
```

How to allocate local variables?

```
void my_function()
{
    int a, b, c;
    ...
}
```

On the stack!

Allocating local variables

- Stored right after the saved EBP value in the stack
- Allocated by subtracting the number of bytes required from ESP

```
void my_function() {
   int a, b, c;
...
```

```
mov [ebp - 4], 10 ; location of variable a
mov [ebp - 8], 5 ; location of b
mov [ebp - 12], 2 ; location of c
```

```
void my_function() {
   int a, b, c;
   ...

_my_function:

  push ebp   ; save the value of ebp
  mov ebp, esp ; ebp = esp, set ebp to be top of the stack (esp)
  sub esp, 12 ; move esp down to allocate space for the
      ; local variables on the stack
```

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mov [ebp - 4], 10 ; location of variable a
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```

How to pass arguments?

- Possible options:
 - In registers
 - On the stack

How to pass arguments?

- x86 32 bit
 - Pass arguments on the stack
 - Return value is in EAX and EDX
- x86 64 bit more registers!
 - Pass first 6 arguments in registers
 - RDI, RSI, RDX, RCX, R8, and R9
 - The rest on the stack
 - Return value is in RAX and RDX

x86_32: passing arguments on the stack

Example function

```
void my_function(int x, int y, int z)
{     ... }
```

Example invocation

```
my_function(2, 5, 10);
```

Generated code

```
push 10
push 5
push 2
call _my_function
```

```
10 | [ebp + 16] (3rd function argument)
 5 | [ebp + 12] (2nd argument)
 2 | [ebp + 8] (1st argument)
RA | [ebp + 4] (return address)
FP | [ebp] (old ebp value) ← EBP points here
   | [ebp - 4] (1st local variable)
   | [ebp - X] (esp - the current stack pointer)
```

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```

```
int callee(int, int, int);
int caller(void)
{
   int ret;

   ret = callee(1, 2, 3);
   ret += 5;
   return ret;
}
```

```
caller:
  ; manage own stack frame
 push
         ebp
      ebp, esp
 mov
  ; push call arguments
         3
 push
 push
 push
  : call subroutine 'callee'
 call callee
  ; remove arguments from frame
 add
     esp, 12
  : use subroutine result
 add eax, 5
  : restore old call frame
 pop ebp
  ; return
 ret
```

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 ret
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        ebp
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         3
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  : call subroutine 'callee'
 call callee
  ; remove arguments from frame
     esp, 12
 add
  : use subroutine result
 add eax, 5
  ; restore old call frame
         ebp
  : return
 ret
```

Wait, where is "return ret;"?

```
int callee(int, int, int);
int caller(void)
{
   int ret;

   ret = callee(1, 2, 3);
   ret += 5;
   return ret;
}
```

```
caller:
 ; manage own stack frame
 push
        ebp
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 ; push call arguments
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 : call subroutine 'callee'
 call callee
 ; remove arguments from frame
 add esp, 12
 : use subroutine result
 add eax, 5
 : restore old call frame
     ebp
 pop
 ; return
```

```
void my_function(int x, int y, int z)
                          int a, b, c;
                          return;
_my_function:
 push ebp
 mov ebp, esp
  sub esp, 12; allocate local varaibles
               ; sizeof(a) + sizeof(b) + sizeof(c)
  ; x = [ebp + 8], y = [ebp + 12], z = [ebp + 16]
  ; a = [ebp-4] = [esp+8],
  ; b=[ebp-8]=[esp+4], c=[ebp-12]=[esp]
  mov esp, ebp; deallocate local variables
 pop ebp
  ret
```

side code

_my_function:

push ebp

pop ebp

ret

mov ebp, esp

```
Example: callee void my_function(int x, int y, int z)
                           int a, b, c;
                           return;
    sub esp, 12; allocate local varaibles
                ; sizeof(a) + sizeof(b) + sizeof(c)
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```

```
void my_function(int x, int y, int z)
                           int a, b, c;
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_my_function:
 push ebp
 mov ebp, esp; ebp = esp
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 mov esp, ebp ;deallocate local variables (esp = ebp)
 pop ebp
 ret
```

```
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                           int a, b, c;
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    mov esp, ebp ; deallocate local variables (esp = ebp)
    pop ebp
    ret.
```

leave instruction

```
void my_function(int x, int y, int z)
                           int a, b, c;
                           return;
_my_function:
  push ebp
  mov ebp, esp; ebp = esp
  sub esp, 12; allocate local varaibles
               : sizeof(a) + sizeof(b) + sizeof(c)
  ; x = [ebp + 8], y = [ebp + 12], z = [ebp + 16]
  ; a = [ebp-4] = [esp+8],
  ; b=[ebp-8]=[esp+4], c=[ebp-12] = [esp]
 mov esp, ebp

    x86 has a special instruction

  pop ebp
                           for this
  ret
```

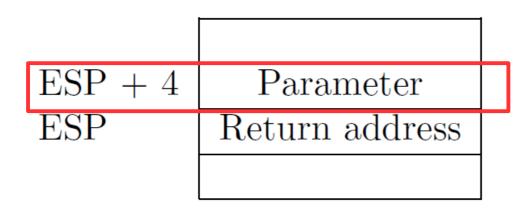
leave

Back to stack frames, so why do we need them?

- ... They are not strictly required
- GCC compiler option -fomit-frame-pointer can disable them

Don't keep the frame pointer in a register for functions that don't need one. This avoids the instructions to save, set up and restore frame pointers; it also makes an extra register available in many functions. It also makes debugging impossible on some machines.

Referencing args without frames



ESP + 8	Parameter	
ESP + 4	Return address	
ESP	subprogram data	

Initially parameter is

• [ESP + 4]

Later as the function pushes things on the stack it changes, e.g.

• [ESP + 8]

- Debugging becomes hard
 - As ESP changes one has to manually keep track where local variables are relative to ESP (ESP + 4 or +8)
 - Compiler can easily do this and generate correct code!
 - But it's hard for a human
 - It's hard to unwind the stack in case of a crash
 - To print out a backtrace

And you only save...

- A couple instructions required to maintain the stack frame
- And 1 register (EBP)
 - x32 has 8 registers (and one is ESP)
 - So taking another one is 12.5% of register space
 - Sometimes its worse it!
 - x64 has 16 registers, so it doesn't really matter
- That said, GCC sets -fomit-frame-pointer to "on"
 - At -O, -O1, -O2 ...
 - Don't get surprised

Relevant part of the GCC manual

3.10 Options That Control Optimization

https://gcc.gnu.org/onlinedocs/gcc/Optimize-Options.html

-O

-01

With -O, the compiler tries to reduce code size and execution time, without performing any optimizations that take a great deal of compilation time.

- -O turns on the following optimization flags:
- -fauto-inc-dec
- -fbranch-count-reg

. . .

- -fomit-frame-pointer
- -freorder-blocks

Saving and restoring registers

Saving register state across invocations

- Processor doesn't save registers
 - General purpose, segment, flags
- Again, a calling convention is needed
 - Agreement on what gets saved by the callee and the caller

Saving register state across invocations

- Registers EAX, ECX, and EDX are caller-saved
 - The function is free to use them
- ... the rest are callee-saved
 - If the function uses them it has to restore them to the original values

- In general there multiple calling conventions
 - We described cdecl
 - Make sure you know what you're doing
 - https://en.wikipedia.org/wiki/X86_calling_convention s#List_of_x86_calling_conventions
 - It's easy as long as you know how to read the table

Questions?

References

- https://en.wikibooks.org/wiki/X86_Disassembly/ Functions_and_Stack_Frames
- https://en.wikipedia.org/wiki/Calling_convention
- https://en.wikipedia.org/wiki/ X86_calling_conventions
- http://stackoverflow.com/questions/14666665/ trying-to-understand-gcc-option-fomit-framepointer