

ICS143A: Principles of Operating Systems

Lecture 15: Locking

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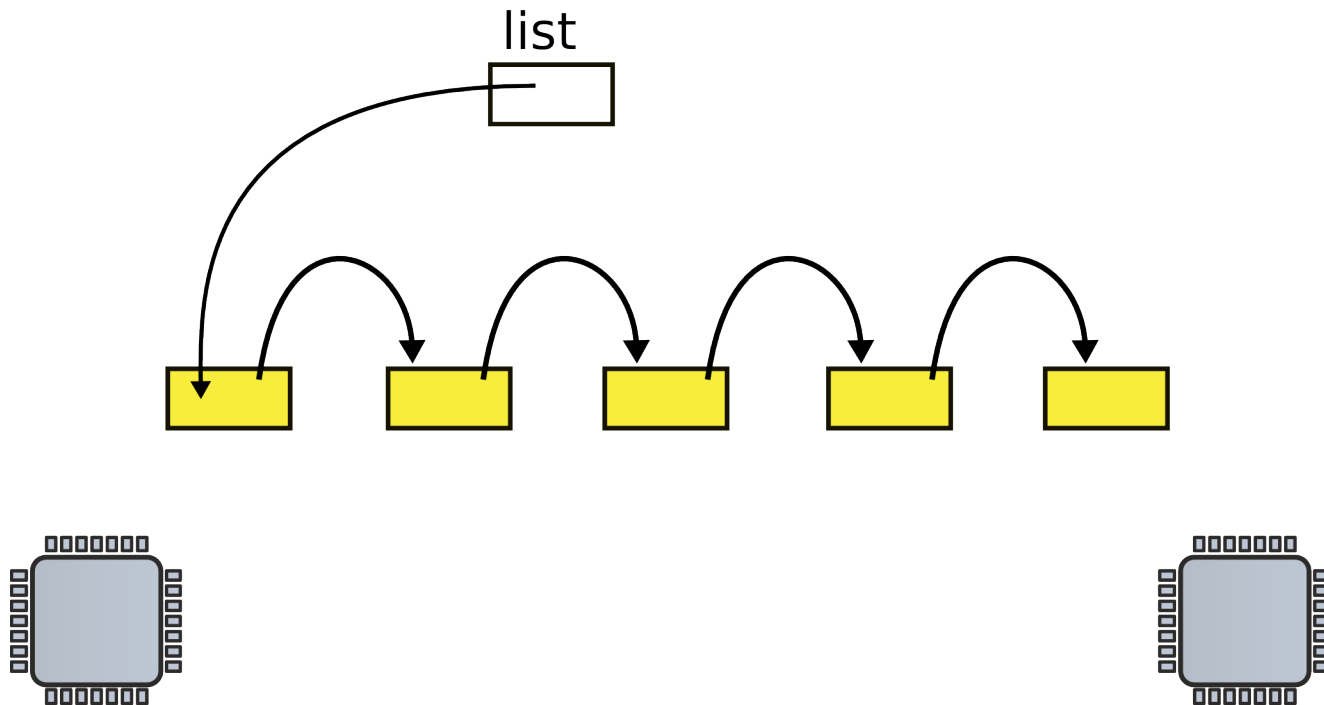
Race conditions

- Disk driver maintains a list of outstanding requests
- Each process can add requests to the list

List implementation no locks

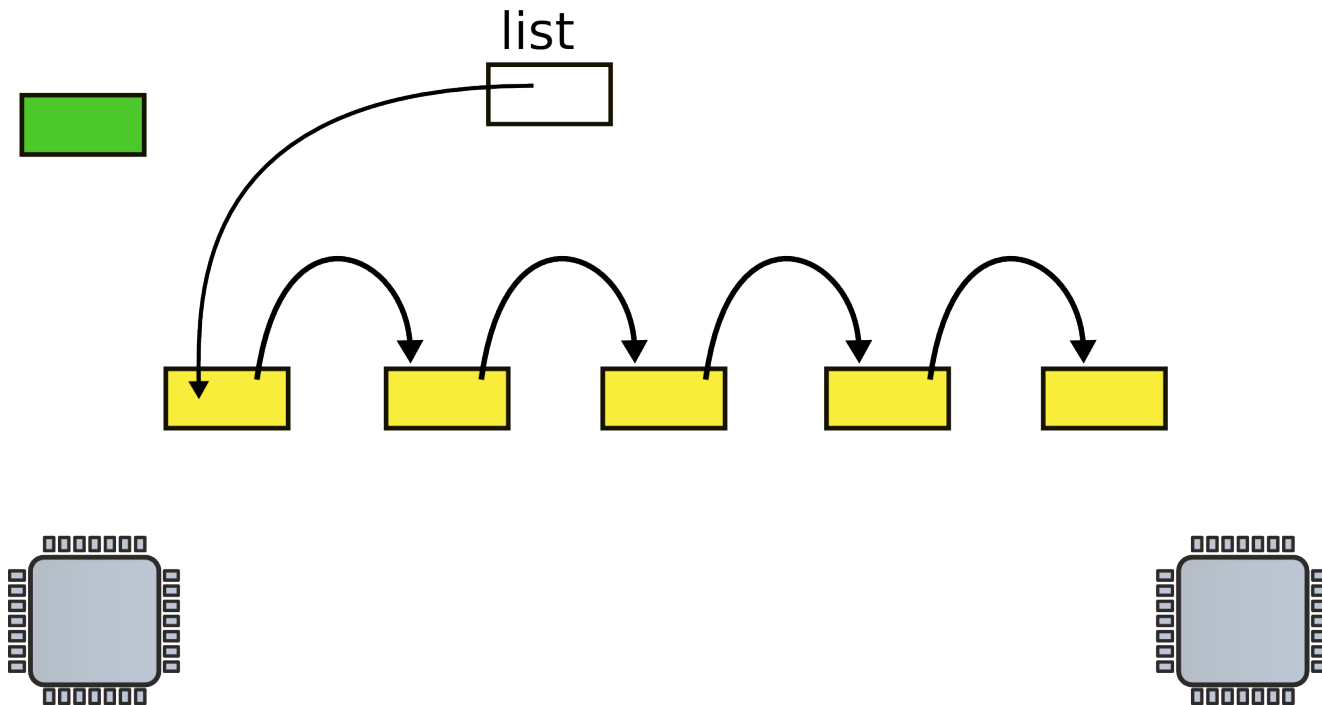
```
1 struct list {  
2     int data;  
3     struct list *next;  
4 };  
  
...  
6 struct list *list = 0;  
  
...  
9 insert(int data)  
10 {  
11     struct list *l;  
12  
13     l = malloc(sizeof *l);  
14     l->data = data;  
15     l->next = list;  
16     list = l;  
17 }
```

Request queue (e.g. incoming network packets)

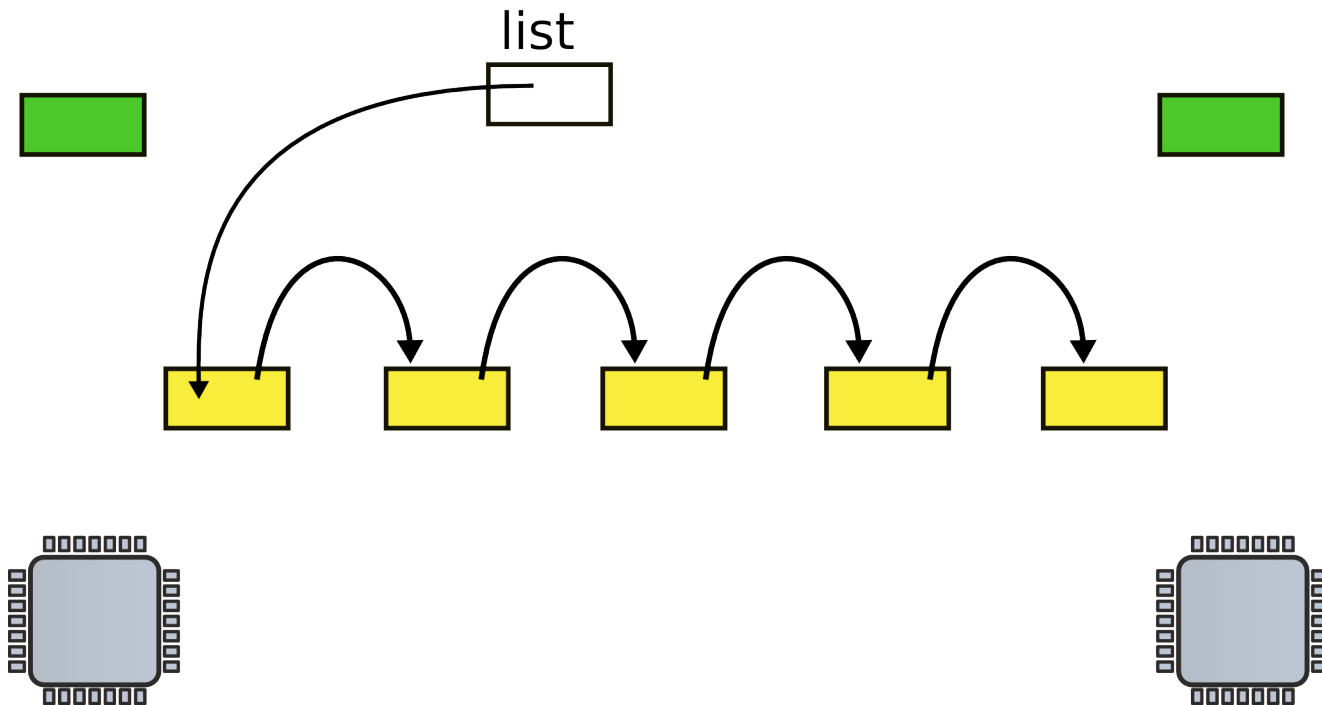


- Linked list, list is pointer to the first element

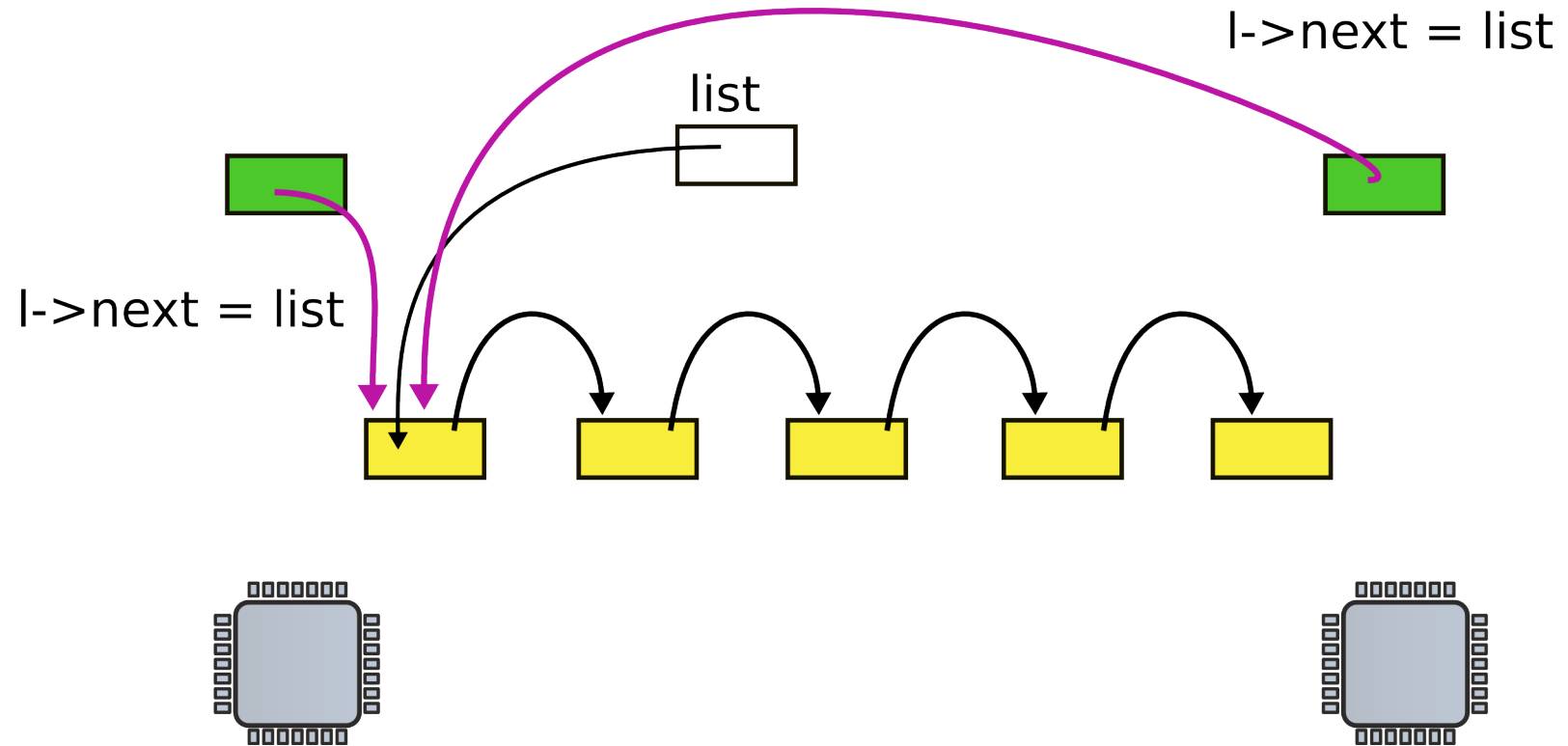
CPU1 allocates new request



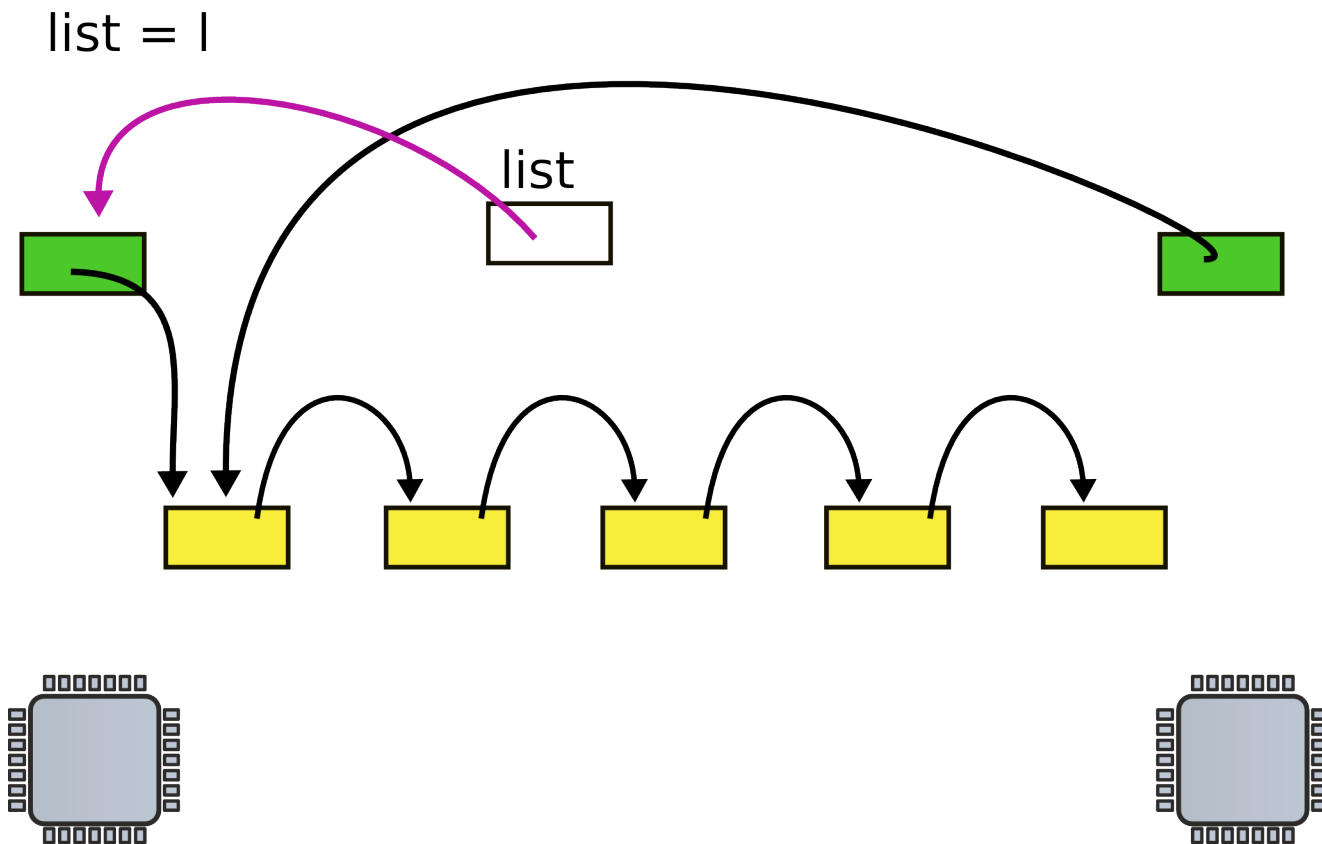
CPU2 allocates new request



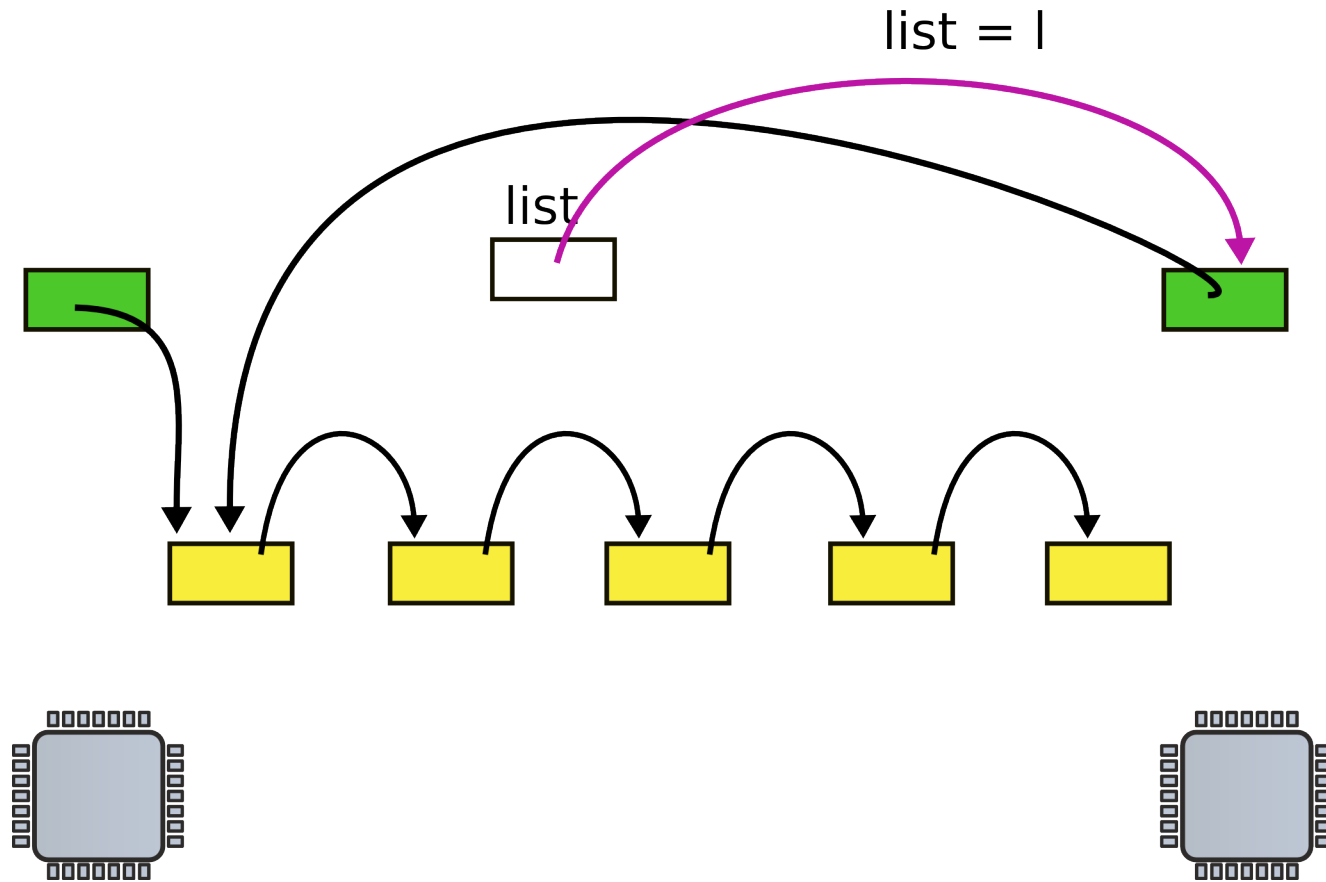
CPU 1 and 2 update next pointer



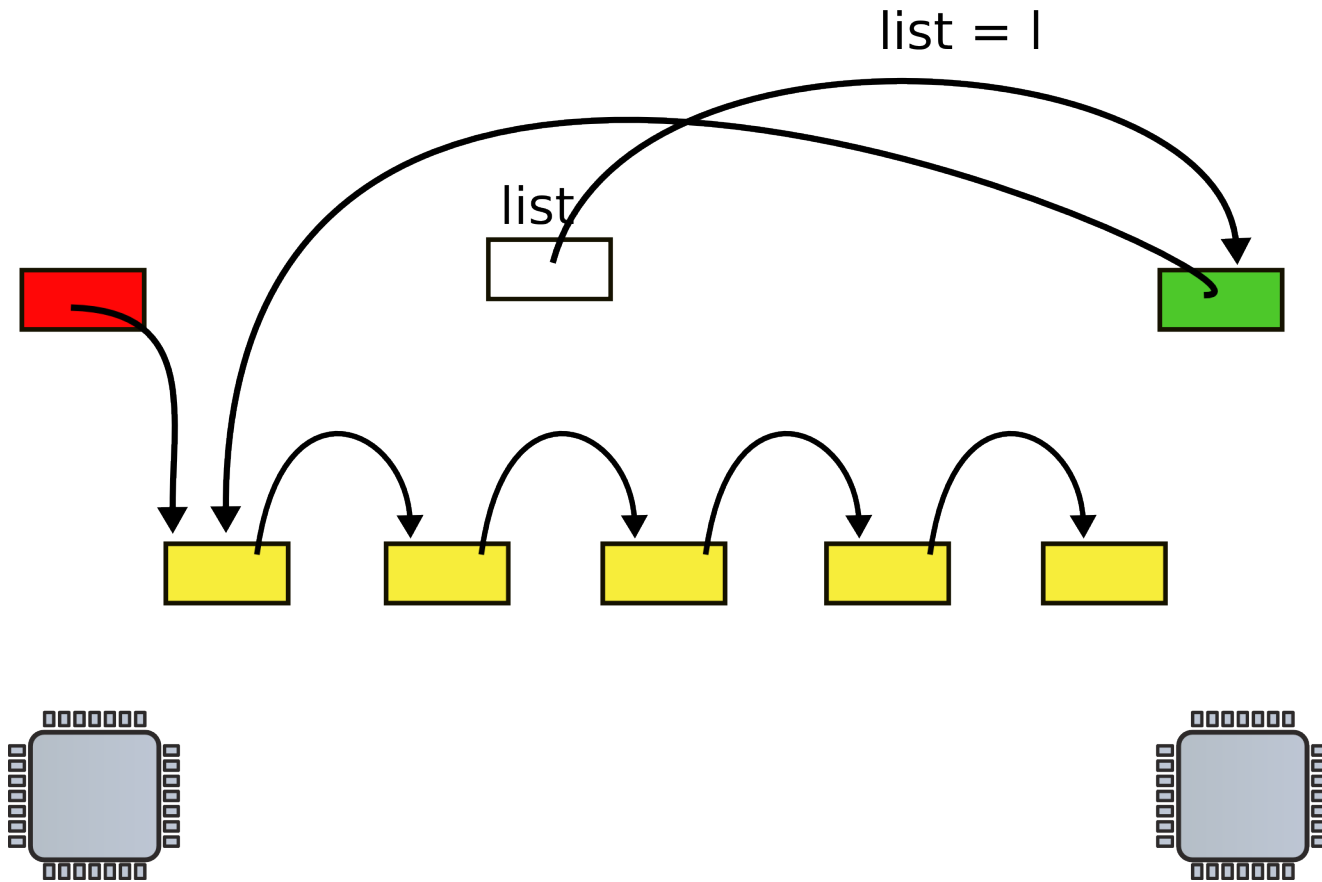
CPU1 updates head pointer



CPU2 updates head pointer



State after the race



Mutual exclusion

- Only one CPU can update list at a time

List implementation with locks

```
1 struct list {
2     int data;
3     struct list *next;
4 };
5
6 struct list *list = 0;
7     struct lock listlock;
8
9 insert(int data)
10 {
11     struct list *l;
12
13     l = malloc(sizeof *l);
14     acquire(&listlock);
15     l->data = data;
16     l->next = list;
17     list = l;
18     release(&listlock);
19 }
20 }
```

- Critical section

- How can we implement `acquire()`?

Spinlock

```
21 void
22 acquire(struct spinlock *lk)
23 {
24     for(;;) {
25         if(!lk->locked) {
26             lk->locked = 1;
27             break;
28         }
29     }
30 }
```

- Spin until lock is 0
- Set it to 1

Still incorrect

```
21 void
22 acquire(struct spinlock *lk)
23 {
24     for(;;) {
25         if(!lk->locked) {
26             lk->locked = 1;
27             break;
28         }
29     }
30 }
```

- Two CPUs can reach line #25 at the same time
 - See not locked, and
 - Acquire the lock
- Lines #25 and #26 need to be atomic
 - I.e. indivisible

Compare and swap: xchg

- Swap a word in memory with a new value
 - Return old value

Correct implementation

```
1573 void
1574 acquire(struct spinlock *lk)
1575 {
...
1580     // The xchg is atomic.
1581     while(xchg(&lk->locked, 1) != 0)
1582         ;
...
1592 }
```

xchgl instruction

```
0568 static inline uint
0569 xchg(volatile uint *addr, uint newval)
0570 {
0571     uint result;
0572
0573     // The + in "+m" denotes a read-modify-write
        operand.
0574     asm volatile("lock; xchgl %0, %1" :
0575                 "+m" (*addr), "=a" (result) :
0576                 "1" (newval) :
0577                 "cc");
0578     return result;
0579 }
```

Correct implementation

```
1573 void
1574 acquire(struct spinlock *lk)
1575 {
1576     ...
1580     // The xchg is atomic.
1581     while(xchg(&lk->locked, 1) != 0)
1582         ;
1584     // Tell the C compiler and the processor to not move loads or
1585     // stores
1586     // past this point, to ensure that the critical section's memory
1587     // references happen after the lock is acquired.
1588     __sync_synchronize();
1589     ...
1592 }
```

Thank you!