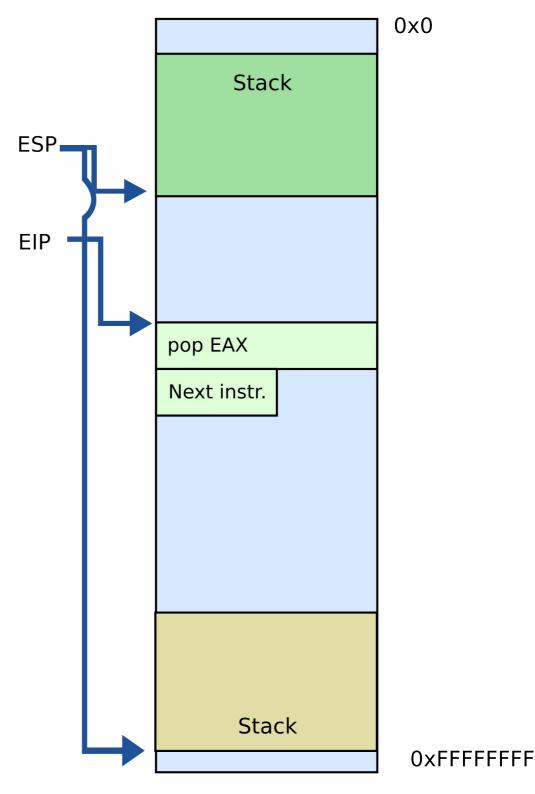
# Function invocations, and calling conventions

**Anton Burtsev** 

## Recap: stack

#### Stack

- It's just a region of memory
  - Pointed by a special register ESP
- You can change ESP
  - Get a new stack





## Calling functions

```
// some code...
foo();
// more code..
```

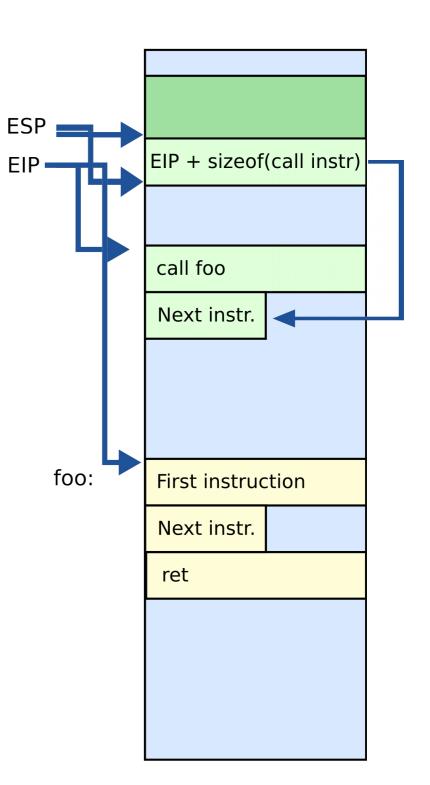
- Stack contains information for how to return from a subroutine
  - i.e., from foo()

 Functions can be called from different places in the program

```
if (a == 0) {
    foo();
    ...
} else {
    foo();
    ...
}
```

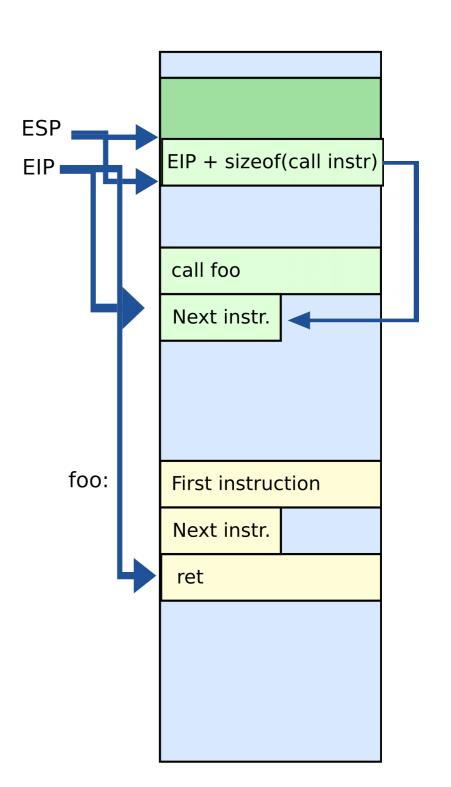
#### Stack

- Main purpose:
  - Store the return address for the current procedure
  - Caller pushes return address on the stack
  - Callee pops it and jumps



#### Stack

- Main purpose:
  - Store the return address for the current procedure
  - Caller pushes return address on the stack
  - Callee pops it and jumps



```
foo(int a) {
    if (a == 0)
        return;
    a--;
    foo(a);
    return;
}
```

## Calling conventions

## Calling conventions

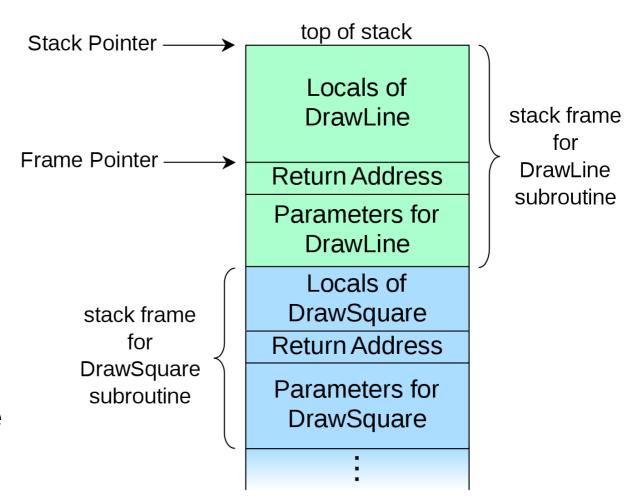
- Goal: re-entrant programs
  - How to pass arguments
    - On the stack?
    - In registers?
  - How to return values
    - On the stack?
    - In registers?
- Conventions differ from compiler, optimizations, etc.

#### Maintain stack as frames

 Each function has a new frame

```
void DrawSquare(...)
{
     ...
     DrawLine(x, y, z);
}
```

- Use dedicated register
   EBP (frame pointer)
  - Points to the base of the frame

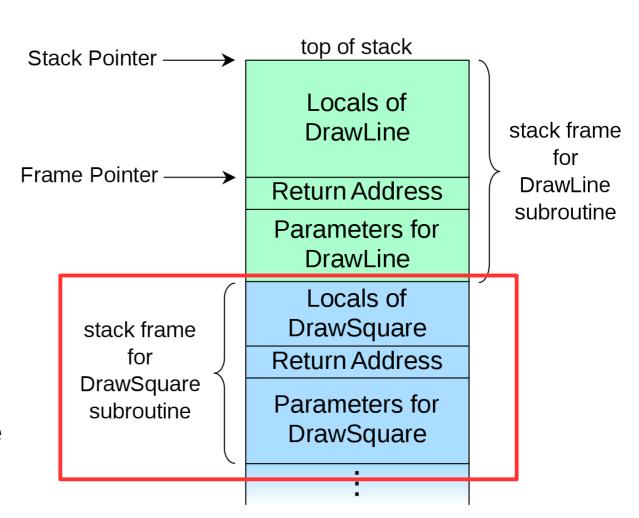


#### Maintain stack as frames

 Each function has a new frame

```
void DrawSquare(...)
{
     ...
     DrawLine(x, y, z);
}
```

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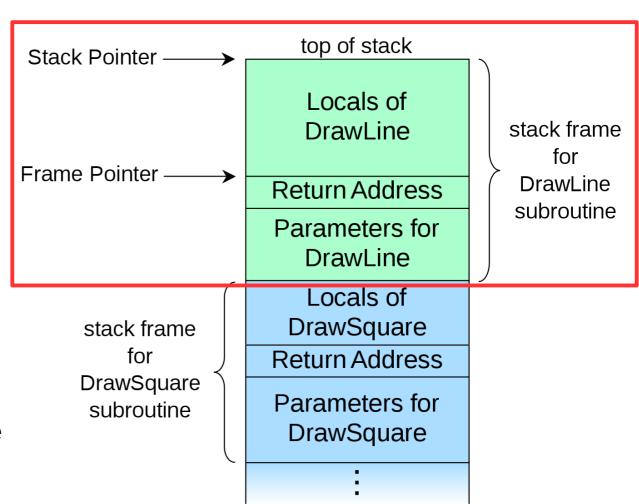


#### Stack consists of frames

 Each function has a new frame

```
void DrawSquare(...)
{
    ...
    DrawLine(x, y, z);
}
```

- Use dedicated register
   EBP (frame pointer)
  - Points to the base of the frame



## Prologue/epilogue

- Each function maintains the frame
  - A dedicated register EBP is used to keep the frame pointer
  - Each function uses prologue code (blue), and epilogue (yellow) to maintain the frame

## Local variables

## What types of variables do you know?

 Or where these variables are allocated in memory?

## What types of variables do you know?

- Global variables
  - Initialized → data section
  - Uninitalized → BSS
- Dynamic variables
  - Heap
- Local variables
  - Stack

#### Global variables

```
1. #include <stdio.h>
2.
3. char hello[] = "Hello";
4. int main(int ac, char **av)
5. {
       static char world[] = "world!";
6.
       printf("%s %s\n", hello, world);
7.
8.
      return 0;
9. }
```

#### Global variables

```
1. #include <stdio.h>
2.
3. char hello[] = "Hello";
4. int main(int ac, char **av)
5. {
6. static char world[] = "world!";
7. printf("%s %s\n", hello, world);
8. return 0;
9. }
```

- Allocated in the data section
  - It is split in initialized (non-zero), and non-initialized (zero)
  - As well as read/write, and read only data section

#### Global variables

## Dynamic variables (heap)

```
1. #include <stdio.h>
2. #include <string.h>
3. #include <stdlib.h>
4.
5. char hello[] = "Hello";
6. int main(int ac, char **av)
7. {
8.
     char world[] = "world!";
9.
      char *str = malloc(64);
      memcpy(str, "beautiful", 64);
10.
11. printf("%s %s %s\n", hello, str, world);
12.
      return 0;
13.}
```

## Dynamic variables (heap)

```
1. #include <stdio.h>
2. #include <string.h>
3. #include <stdlib.h>
4.
5. char hello[] = "Hello";
6. int main(int ac, char **av)
7. {
8. char world[] = "world!";
9. char *str = malloc(64);
10.
      memcpy(str, "beautiful", 64);
      printf("%s %s %s\n", hello, str, world);
11.
12.
      return 0:
13.}
```

- Allocated on the heap
  - Special area of memory provided by the OS from where malloc() can allocate memory

## Dynamic variables (heap)

#### Local variables

#### Local variables

```
1. #include <stdio.h>
2.
3. char hello[] = "Hello";
4. int main(int ac, char **av)
5. {
6.
      //static char world[] = "world!";
      char world[] = "world!";
7.
      printf("%s %s\n", hello, world);
8.
9.
      return 0;
10.}
```

#### Local variables...

Each function has private instances of local variables

```
foo(int x) {
    int a, b, c;
    ...
    return;
}
```

Function can be called recursively

```
foo(int x) {
    int a, b, c;
    a = x + 1;
    if ( a < 100 )
        foo(a);
    return;
}</pre>
```

#### How to allocate local variables?

```
void my_function()
{
    int a, b, c;
    ...
}
```

#### How to allocate local variables?

```
void my_function()
{
    int a, b, c;
    ...
}
```

On the stack!

### Allocating local variables

- Stored right after the saved EBP value on the stack
- Allocated by subtracting the number of bytes required from ESP

```
void my_function() {
   int a, b, c;
...
```

```
mov [ebp - 4], 10 ; location of variable a
mov [ebp - 8], 5 ; location of b
mov [ebp - 12], 2 ; location of c
```

```
void my_function() {
   int a, b, c;
   ...

_my_function:

  push ebp   ; save the value of ebp
  mov ebp, esp ; ebp = esp, set ebp to be top of the stack (esp)
  sub esp, 12 ; move esp down to allocate space for the
      ; local variables on the stack
```

```
mov [ebp - 4], 10 ; location of variable a
mov [ebp - 8], 5 ; location of b
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mov [ebp - 4], 10 ; location of variable a mov [ebp - 8], 5 ; location of b mov [ebp - 12], 2 ; location of c
```

### How to pass arguments?

- Possible options:
  - In registers
  - On the stack

### How to pass arguments?

- x86 32 bit
  - Pass arguments on the stack
  - Return value is in EAX and EDX
- x86 64 bit more registers!
  - Pass first 6 arguments in registers
    - RDI, RSI, RDX, RCX, R8, and R9
  - The rest on the stack
  - Return value is in RAX and RDX

## x86\_32: passing arguments on the stack

Example function

```
void my_function(int x, int y, int z)
{    ... }
```

Example invocation

```
my_function(2, 5, 10);
```

Generated code

```
push 10
push 5
push 2
call _my_function
```

## Example stack

```
10 | [ebp + 16] (3rd function argument)
 5 | [ebp + 12] (2nd argument)
 2 | [ebp + 8] (1st argument)
RA | [ebp + 4] (return address)
FP | [ebp] (old ebp value) ← EBP points here
   | [ebp - 4] (1st local variable)
   | [ebp - X] (esp - the current stack pointer)
```

#### Example stack

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10 | [ebp + 16] (3rd function argument)
 5 | [ebp + 12] (2nd argument)
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     [ebp - 4] (1st local variable)
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#### Example stack

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RA | [ebp + 4] (return address)
FP | [ebp] (old ebp value) ← EBP points here
    [ebp - 4] (1st local variable)
     [ebp - X] (esp - the current stack pointer)
```

```
int callee(int, int, int);
int caller(void)
{
   int ret;

   ret = callee(1, 2, 3);
   ret += 5;
   return ret;
}
```

```
caller:
  ; manage own stack frame
 push
         ebp
      ebp, esp
 mov
  ; push call arguments
         3
 push
 push
 push
  : call subroutine 'callee'
 call callee
  ; remove arguments from frame
 add
     esp, 12
  : use subroutine result
 add eax, 5
  : restore old call frame
 pop ebp
  ; return
 ret
```

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int callee(int, int, int);
int caller(void)
{
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   ret = calle(1, 2, 3);
   ret += 5;
   return ret;
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 pop ebp
  ; return
 ret
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int callee(int, int, int);
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caller:
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 add eax, 5
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 pop ebp
  ; return
 ret
```

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int callee(int, int, int);
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caller:
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 add eax, 5
  : restore old call frame
 pop ebp
  ; return
 ret
```

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int callee(int, int, int);
int caller(void)
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caller:
  ; manage own stack frame
 push
        ebp
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  ; push call arguments
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  : call subroutine 'callee'
 call callee
  ; remove arguments from frame
 add esp, 12
  ; use subroutine result
      eax, 5
 add
  : restore old call frame
 pop ebp
  ; return
 ret
```

```
int callee(int, int, int);
int caller(void)
{
   int ret;

   ret = callee(1, 2, 3);
   ret += 5;
   return ret;
}
```

```
caller:
  ; manage own stack frame
 push
        ebp
 mov ebp, esp
  ; push call arguments
 push
         3
 push
 push
  : call subroutine 'callee'
 call callee
  ; remove arguments from frame
     esp, 12
 add
  : use subroutine result
 add eax, 5
  ; restore old call frame
         ebp
  : return
 ret
```

# Wait, where is "return ret;"?

```
int callee(int, int, int);
int caller(void)
{
   int ret;

   ret = callee(1, 2, 3);
   ret += 5;
   return ret;
}
```

```
caller:
 ; manage own stack frame
 push
        ebp
 mov ebp, esp
 ; push call arguments
 push 3
 push 2
 push
 : call subroutine 'callee'
 call callee
 ; remove arguments from frame
 add esp, 12
 : use subroutine result
 add eax, 5
 : restore old call frame
     ebp
 pop
 ; return
```

```
void my_function(int x, int y, int z)
                          int a, b, c;
                          return;
_my_function:
 push ebp
 mov ebp, esp
  sub esp, 12; allocate local varaibles
               ; sizeof(a) + sizeof(b) + sizeof(c)
  ; x = [ebp + 8], y = [ebp + 12], z = [ebp + 16]
  ; a = [ebp-4] = [esp+8],
  ; b=[ebp-8]=[esp+4], c=[ebp-12]=[esp]
  mov esp, ebp; deallocate local variables
 pop ebp
  ret
```

### side code

\_my\_function:

push ebp

pop ebp

ret

mov ebp, esp

```
Example: callee void my_function(int x, int y, int z)
                           int a, b, c;
                           return;
    sub esp, 12; allocate local varaibles
                ; sizeof(a) + sizeof(b) + sizeof(c)
   ; x = [ebp + 8], y = [ebp + 12], z = [ebp + 16]
    ; a=[ebp-4]=[esp+8],
    ; b=[ebp-8]=[esp+4], c=[ebp-12]=[esp]
    mov esp, ebp; deallocate local variables
```

```
void my_function(int x, int y, int z)
                           int a, b, c;
                           return;
_my_function:
 push ebp
 mov ebp, esp; ebp = esp
  sub esp, 12; allocate local varaibles
              ; sizeof(a) + sizeof(b) + sizeof(c)
  ; x = [ebp + 8], y = [ebp + 12], z = [ebp + 16]
 ; a=[ebp-4]=[esp+8],
  ; b=[ebp-8]=[esp+4], c=[ebp-12]=[esp]
 mov esp, ebp ;deallocate local variables (esp = ebp)
 pop ebp
 ret
```

```
void my_function(int x, int y, int z)
                           int a, b, c;
                           return;
_my_function:
 push ebp
 mov ebp, esp; ebp = esp
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 mov esp, ebp ;deallocate local variables (esp = ebp)
 pop ebp
  ret
```

### side code

```
Example: callee void my_function(int x, int y, int z)
                             int a, b, c;
                             return;
  _my_function:
    push ebp
    mov ebp, esp; ebp = esp
    sub esp, 12; allocate local varaibles
                 ; sizeof(a) + sizeof(b) + sizeof(c)
    ; x = [ebp + 8], y = [ebp + 12], z = [ebp + 16]
    ; a = [ebp-4] = [esp+8],
    ; b=[ebp-8]=[esp+4], c=[ebp-12]=[esp]
    mov esp, ebp ; deallocate local variables (esp = ebp)
    pop ebp
    ret.
```

#### leave instruction

```
void my_function(int x, int y, int z)
                           int a, b, c;
                           return;
_my_function:
  push ebp
  mov ebp, esp; ebp = esp
  sub esp, 12; allocate local varaibles
               : sizeof(a) + sizeof(b) + sizeof(c)
  ; x = [ebp + 8], y = [ebp + 12], z = [ebp + 16]
  ; a = [ebp-4] = [esp+8],
  ; b=[ebp-8]=[esp+4], c=[ebp-12] = [esp]
 mov esp, ebp

    x86 has a special instruction

  pop ebp
                           for this
  ret
```

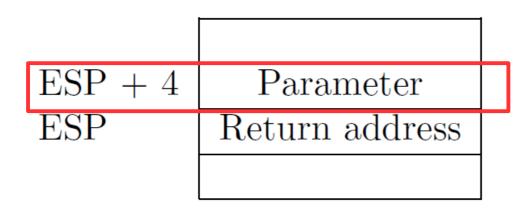
leave

### Back to stack frames, so why do we need them?

- ... They are not strictly required
- GCC compiler option -fomit-frame-pointer can disable them

Don't keep the frame pointer in a register for functions that don't need one. This avoids the instructions to save, set up and restore frame pointers; it also makes an extra register available in many functions. It also makes debugging impossible on some machines.

#### Referencing args without frames



ESP + 8	Parameter	
ESP + 4	Return address	
ESP	subprogram data	

Initially parameter is

• [ESP + 4]

Later as the function pushes things on the stack it changes, e.g.

• [ESP + 8]

- Debugging becomes hard
  - As ESP changes one has to manually keep track where local variables are relative to ESP (ESP + 4 or +8)
    - Compiler can easily do this and generate correct code!
    - But it's hard for a human
  - It's hard to unwind the stack in case of a crash
    - To print out a backtrace

#### And you only save...

- A couple instructions required to maintain the stack frame
- 1 register (EBP)
  - x32 has 8 registers (and one is ESP, so 7 are left)
    - So taking another one is 1/7 or 14.28% of register space
    - Sometimes its worse it!
  - x64 has 16 registers, so it doesn't really matter
- That said, GCC sets -fomit-frame-pointer to "on"
  - At -O, -O1, -O2 ...
  - Don't get surprised

#### Relevant part of the GCC manual

3.10 Options That Control Optimization

https://gcc.gnu.org/onlinedocs/gcc/Optimize-Options.html

**-**O

**-01** 

With -O, the compiler tries to reduce code size and execution time, without performing any optimizations that take a great deal of compilation time.

- -O turns on the following optimization flags:
- -fauto-inc-dec
- -fbranch-count-reg

. . .

- -fomit-frame-pointer
- -freorder-blocks

### Saving and restoring registers

## Saving register state across invocations

- Processor doesn't save registers
  - General purpose, segment, flags
- Again, a calling convention is needed
  - Agreement on what gets saved by the callee and the caller

## Saving register state across invocations

- Registers EAX, ECX, and EDX are caller-saved
  - The function is free to use them
- ... the rest are callee-saved
  - If the function uses them it has to restore them to the original values

- In general there multiple calling conventions
  - We described cdecl
  - Make sure you know what you're doing
  - https://en.wikipedia.org/wiki/X86\_calling\_convention s#cdecl

It's easy as long as you know how to read the table

### Questions?

#### References

- https://en.wikibooks.org/wiki/X86\_Disassembly/ Functions\_and\_Stack\_Frames
- https://en.wikipedia.org/wiki/Calling\_convention
- https://en.wikipedia.org/wiki/ X86\_calling\_conventions
- http://stackoverflow.com/questions/14666665/tr ying-to-understand-gcc-option-fomit-frame-point er