Julia Sets Theory and Algorithms

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Remark

The points of S^1 are supersensitive under Q_0 . That is, any open ball B around $z \in S^1$ has the property that $\bigcup_{n=0}^{\infty} Q_0^{\circ n}(B) = \mathbb{C} \setminus \{p\}$ for at most one point p.

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For the quadratic map $Q_0(z)=z^2$ we saw chaotic behavior only on S^1 by angle doubling $(\theta \to 2\theta)$. We also saw that $|Q_0(z)| \to \infty$ for all |z|>1 and $|Q_0(z)|\to 0$ for all |z|<1.

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Remark

Finally, we learned that for $Q_{-2}(z)=z^2-2$, we have $J_2=K_2=[-2,2]$, so any $z\in\mathbb{C}\setminus[-2,2]\to\infty$ as we compose $Q_{-2}(z)$ infinitely many times.

► Goal: discuss and prove parts of:

Theorem

If |c| is sufficiently large, Λ , the set of points whose entire forward orbits lie within the circle |z| = |c|, is a Cantor set on which Q_c is topologically conjugate to the shift map on two symbols. All points in $\mathbb{C} - \Lambda$ tend to ∞ under iteration of Q_c . Hence, $J_c = K_c$.

How large is 'sufficiently large'? We will discuss the case where |c|>2.

Theorem (The Escape Criterion)

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Proof.

We use the triangle inequality to get the following estimate:

$$|Q_c(z)| \ge |z|^2 - |c| \ge |z|^2 - |z| = |z|(|z| - 1)$$
 (1)

Since |z|>2, we know that |z|-1>1, so there exists $\lambda>0$ such that

$$|Q_c^n(z)| \ge (1+\lambda)^n |z| \tag{2}$$

Since |z| is fixed and $(1 + \lambda)^n$ grows arbitrarily large, $|Q_c^n(z)|$ also grows arbitrarily large, as desired.

Theorem

Let D be the closed disk (i.e. $\{z:|z|\leq |c|\}$), with |c|>2. Then the filled Julia set of Q_c is given by

$$K_c = \bigcap_{n \ge 0} Q_c^{\circ - n}(D)$$

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Let D be the closed disk (i.e. $\{z: |z| \le |c|\}$), with |c| > 2. Then the filled Julia set of Q_c is given by

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Proof.

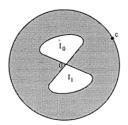
Consider the case where $z \notin \bigcap_{n \geq 0} Q_c^{\circ - n}(D)$. Then there exists some $k \in \mathbb{N}$ such that $z \notin Q_c^{\circ - k}(D)$, so we have that $Q_c^{\circ k}(z) \notin D$. Thus by the escape criterion, z escapes to infinity under iteration of Q_c , so $z \notin K_c$.

Otherwise, if $z \in \bigcap_{n \geq 0} Q_c^{\circ -n}(D)$, then $Q_c^{\circ n}(z) \in D$ for all $n \in \mathbb{N}$. Thus z is bounded under iteration of Q_c , so $z \in K_c$.

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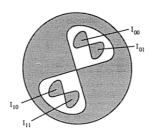
Take C to be the circle around the origin of radius |c|. What does $Q_c^{\circ-1}(C)$ look like? Recall from last week that $\pm \sqrt{C-c}$ is the following Lemniscate-like shape (by the Boundary Mapping Principle).



Where each 'lobe' is mapped diffeomorphically (continuously bijected) to C.

Repeating this process yields nested lobes, each a diffeomorphism to its 'parent' lobe from the previous iteration. Below is a visualization of this process for two backwards-iterations.

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$$I_{s_0s_1\dots s_n} = \{z \in D \mid z \in I_{s_0}, Q_c(z) \in I_{s_1}, \dots, Q_c^{on}(z) \in I_{s_n}\}$$

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As you may have noticed, this is identical to our previous use of symbolic dynamics (in relation to the Cantor Middle-Thirds set). By similar logic to our discussion before, the following intersection is a nonempty set.

$$\bigcap_{n>0} I_{s_0s_1...s_n}$$

This gives us that, for any $z\in \cap_{n\geq 0}I_{s_0s_1...s_n}$, $Q_c^{\circ k}(z)\in D$ for all k. Hence, $z\in K_c$.

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Conversely, any $z \in K_c$ must lie in exactly one of these components, since any infinite string of 1s and 0 maps to any z by $S(z) = s_0 s_1 \dots s_n$ provided $z \in \cap_{n \geq 0} I_{s_0 s_1 \dots s_n}$.

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Similar arguments to those we used for the Cantor set yield that [1] this correspondence is continuous and [2] that each of these infinite intersections is exactly a single point in the complex plane. The mechanics behind this exceed the scope of this lecture (and class), so we will not include them here.

Remark

 Q_c is supersensitive on J_c . Roughly, (as handwavy as Devaney), this is because we can, for any point $z \in I_{s_0s_1...s_k}$, find a small enough disk that contains z such that for an arbitrarily large k we have $Q_c^{\circ k} = D$. Further iterations of Q_c will eventually diffeomorphically map that disk to all of $\mathbb C$ (except for at most 1 point).

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When |c| > 2, the orbit of 0 tents to ∞ . This will be important on Thursday.

This website gives an amazing visual intuition for the process we just described:

► Acko - How to Fold a Julia Fractal

Algorithm 1: Algorithm to plot K_c

Input: grid, a list of evenly-spaced complex numbers in a rectangular region of the complex plane

```
1 for j \leftarrow 1 to N do
2 | for z in grid do
3 | z \leftarrow Q_c(z)
```

Preliminaries

// Shade all points which didn't escape

```
4 for z in grid do

5 | if |z| \le max(|c|, 2) then

6 | paintBlack(z)
```

Algorithm 2: Algorithm to plot J_c

Input: MaxIter, the maximum number of iterations to perform

```
1 Z ← randomComplexNumber()
2 for i \leftarrow 1 to MaxIter do
     binaryRand \leftarrow randomBit()
3
                         // Pick a random backwards iteration
4
     if binaryRand = 0 then
      z \leftarrow \sqrt{(z-c)}
5
     else
6
      z \leftarrow -\sqrt{(z-c)}
                                      // Don't plot stray points
     if i > 100 then
8
         paintBlack(z)
9
```

Explanation for an alternative plotting algorithm

The above algorithm works well for K_c but as well for J_c . To remedy this, recall that Q_c is supersensitive on J_c . This means that any neighborhood B of some point $z \in J_c$ has the property that $\bigcup_{n=0}^{\infty} Q_0^{\circ n}(B) = \mathbb{C} \setminus \{p\}$ for at most one point p.

So supersensitive points serve as ``attractors'' for the backward iteration of Q_c , in the sense that for each supersensitive point, all of $\mathbb C$ except at most one point must come arbitrarily close to it on reverse iteration.

- ► A demo of Julia-Mandelbrot correspondence
- ► Some pictures of generalized Julia sets