

# Julia Sets

## Theory and Algorithms

Josh Lipschultz & Ricky LeVan  
*MATH/CAAM 435*



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Quick refresher from last time:

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*An orbit is bounded if there exists a  $K$  such that  $|Q_c^{on}(z)| < K$  for all  $n$ . Otherwise the orbit is unbounded.*

## Remark

*The points of  $S^1$  are supersensitive under  $Q_0$ . That is, any open ball  $B$  around  $z \in S^1$  has the property that  $\bigcup_{n=0}^{\infty} Q_0^{on}(B) = \mathbb{C} \setminus \{p\}$  for at most one point  $p$ .*

## Definition

*The Julia Set  $J_c$  is the boundary of the filled Julia set  $K_c$ . (The filled Julia set is the set of bounded points of  $Q_c$ .) We could alternatively define  $J_c$  as the closure of the set of repelling points of  $Q_c$ .*

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## Remark

*For the quadratic map  $Q_0(z) = z^2$  we saw chaotic behavior only on  $S^1$  by angle doubling ( $\theta \rightarrow 2\theta$ ). We also saw that  $|Q_0(z)| \rightarrow \infty$  for all  $|z| > 1$  and  $|Q_0(z)| \rightarrow 0$  for all  $|z| < 1$ .*

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## Remark

*Finally, we learned that for  $Q_{-2}(z) = z^2 - 2$ , we have  $J_2 = K_2 = [-2, 2]$ , so any  $z \in \mathbb{C} \setminus [-2, 2] \rightarrow \infty$  as we compose  $Q_{-2}(z)$  infinitely many times.*



## Definition

*A set  $\Lambda \subseteq \mathbb{C}$  is a Cantor set if it is a closed, totally disconnected, and perfect subset of  $\mathbb{C}$ . (Adapted from the version for an interval of the real line.)*

- Goal: discuss and prove parts of:

## Theorem

*If  $|c|$  is sufficiently large,  $\Lambda$ , the set of points whose entire forward orbits lie within the circle  $|z| = |c|$ , is a Cantor set on which  $Q_c$  is topologically conjugate to the shift map on two symbols. All points in  $\mathbb{C} - \Lambda$  tend to  $\infty$  under iteration of  $Q_c$ . Hence,  $J_c = K_c$ .*

## Theorem (The Escape Criterion)

*Suppose  $2 < |c| \leq |z|$ . Then we have that  $|Q_c^n(z)| \rightarrow \infty$  as  $n \rightarrow \infty$ .*

### Proof.

We use the triangle inequality to get the following estimate:

$$|Q_c(z)| \geq |z|^2 - |c| \geq |z|^2 - |z| = |z|(|z| - 1) \quad (1)$$

Since  $|z| > 2$ , we know that  $|z| - 1 > 1$ , so there exists  $\lambda > 0$  such that

$$|Q_c^n(z)| \geq (1 + \lambda)^n |z| \quad (2)$$

Since  $|z|$  is fixed and  $(1 + \lambda)^n$  grows arbitrarily large,  $|Q_c^n(z)|$  also grows arbitrarily large, as desired.  $\square$

## Theorem

Let  $D$  be the closed disk (i.e.  $\{z : |z| \leq |c|\}$ ), with  $|c| > 2$ . Then the filled Julia set of  $Q_c$  is given by

$$K_c = \bigcap_{n \geq 0} Q_c^{\circ - n}(D)$$

where  $Q_c^{\circ - n}(D)$  denotes the preimage of  $D$  under  $Q_c^{\circ n}$

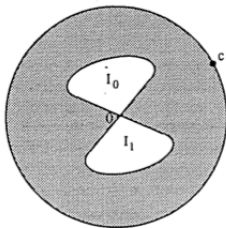
## Proof.

Consider the case where  $z \notin \bigcap_{n \geq 0} Q_c^{\circ - n}(D)$ . Then there exists some  $k \in \mathbb{N}$  such that  $z \notin Q_c^{\circ - k}(D)$ , so we have that  $Q_c^{\circ k}(z) \notin D$ . Thus by the escape criterion,  $z$  escapes to infinity under iteration of  $Q_c$ , so  $z \notin K_c$ .

Otherwise, if  $z \in \bigcap_{n \geq 0} Q_c^{\circ - n}(D)$ , then  $Q_c^{\circ n}(z) \in D$  for all  $n \in \mathbb{N}$ . Thus  $z$  is bounded under iteration of  $Q_c$ , so  $z \in K_c$ . □

Let's return to the theorem we wish to prove.

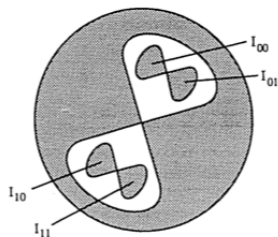
Take  $C$  to be the circle around the origin of radius  $|c|$ . What does  $Q_c^{\circ-1}(C)$  look like? Recall from last week that  $\pm\sqrt{C-c}$  is the following Lemniscate-like shape (by the Boundary Mapping Principle).



Where each 'lobe' is mapped diffeomorphically (continuously bijected) to  $C$ .

Repeating this process yields nested lobes, each a diffeomorphism to its 'parent' lobe from the previous iteration. Below is a visualization of this process for two backwards-iterations.

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Define

$$I_{s_0 s_1 \dots s_n} = \{z \in D \mid z \in I_{s_0}, Q_c(z) \in I_{s_1}, \dots, Q_c^{\circ n}(z) \in I_{s_n}\}$$

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As you may have noticed, this is identical to our previous use of symbolic dynamics (in relation to the Cantor Middle-Thirds set). By similar logic to our discussion before, the following intersection is a nonempty set.

$$\bigcap_{n \geq 0} I_{s_0 s_1 \dots s_n}$$

This gives us that, for any  $z \in \cap_{n \geq 0} I_{s_0 s_1 \dots s_n}$ ,  $Q_c^{ok}(z) \in D$  for all  $k$ .  
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Conversely, any  $z \in K_c$  must lie in exactly one of these components, since any infinite string of 1s and 0 maps to any  $z$  by  $S(z) = s_0 s_1 \dots s_n$  provided  $z \in \bigcap_{n \geq 0} I_{s_0 s_1 \dots s_n}$ .

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Similar arguments to those we used for the Cantor set yield that [1] this correspondence is continuous and [2] that each of these infinite intersections is exactly a single point in the complex plane. The mechanics behind this exceed the scope of this lecture (and class), so we will not include them here.



## Remark

*$Q_c$  is supersensitive on  $J_c$ . (As handwavy as Devaney)*

## Remark

*Since  $K_c$  is a Cantor set (each component of  $K_c$  is a point) when  $|c| > 2$ , the Julia and filled Julia sets are identical ( $K_c = J_c$ ).*

## Remark

*When  $|c| > 2$ , the orbit of 0 tends to  $\infty$ . This will be important on Thursday.*

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## Algorithm 1: Algorithm to plot $K_c$

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Input: *grid*, a list of evenly-spaced complex numbers in a rectangular region of the complex plane

```
1 for  $j \leftarrow 1$  to  $N$  do
2   for  $z$  in grid do
3      $z \leftarrow Q_c(z)$ 
                                     // Shade all points which didn't escape
4 for  $z$  in grid do
5   if  $|z| \leq \max(|c|, 2)$  then
6     paintBlack( $z$ )
```

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## Acko -- How to Fold a Julia Fractal

## Explanation for an alternative plotting algorithm

The above algorithm works well for  $K_c$  but as well for  $J_c$ . To remedy this, recall that  $Q_c$  is supersensitive on  $J_c$ . This means that any neighborhood  $B$  of some point  $z \in J_c$  has the property that  $\bigcup_{n=0}^{\infty} Q_0^{\circ n}(B) = \mathbb{C} \setminus \{p\}$  for at most one point  $p$ .

So supersensitive points serve as "attractors" for the backward iteration of  $Q_c$ , in the sense that for each supersensitive point, all of  $\mathbb{C}$  except at most one point must come arbitrarily close to it on reverse iteration.

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## Algorithm 2: Algorithm to plot $J_c$

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Input: *MaxIter*, the maximum number of iterations to perform

```
1  $z \leftarrow \text{randomComplexNumber}()$ 
2 for  $j \leftarrow 1$  to  $\text{MaxIter}$  do
3    $\text{binaryRand} \leftarrow \text{randomBit}()$ 
4   // Pick a random backwards iteration
5   if  $\text{binaryRand} = 0$  then
6      $z \leftarrow \sqrt{(z - c)}$ 
7   else
8      $z \leftarrow -\sqrt{(z - c)}$ 
9   // Don't plot stray points
10  if  $j > 100$  then
11     $\text{paintBlack}(z)$ 
```

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A demo of Julia-Mandelbrot correspondence

Some pictures of generalized Julia sets