

Homework 1 – Due Tuesday, May 28th, 2024 09:00 pm

- Provide step-by-step explanations, not just answers. Answers without explanations will earn a small fraction of the points.
- Submit your solutions on Gradescope. Don't forget to include information about your collaborators (or say "Collaborators: none").

Problems

0. (0 points) The following steps are required to get you started in the course. Please complete them today.

- (a) Make sure you are signed up on Piazza at <https://piazza.com/bu/summer2024/cascs330s> using your BU email address.
- (b) Sign up on Gradescope using your BU email address and the code **PWDGPG**.
- (c) Read and sign the Collaboration and Honesty Policy and submit it on Gradescope. We will be able to grade your homework only after you complete this step.
- (d) (**Nameplate**) Please print out (or make by hand) a nameplate with your name and bring it to every lecture and lab. A template is available at the bottom of the course web page.
- (e) Check out the course webpage: <https://cs-people.bu.edu/januario/teaching/cs330/su24/index.html>.
- (f) Familiarize yourself with the homework template files at the bottom of the course webpage. Each problem must include a note about collaborators (even if you did the problem by yourself).

1. (Tracing Algorithms, 10 points)

Consider the following algorithm:

// A is an array of integers, indexed from 0 to n-1.

```
mystery( int A[] ){
    int i = 1
    while(i < A.length){
        int j = i
        while(j > 0 && A[j-1] > A[j]) {
            swap(A[j], A[j - 1])
            j = j - 1
        }
        i = i + 1
    }
}
```

- (a) Trace the run of the algorithm on the following inputs. You can draw a table showing the progress of the nested loop (the progress of the counters i, j) and the contents of the array A as the algorithm progresses. If the contents of A do not change, you can leave the corresponding cell blank.

- $A = [1, 2, 3, 4, 5]$
- $A = [4, 1, 5, 3, 2]$
- $A = [5, 4, 3, 2, 1]$

Solution:

$A = [1, 2, 3, 4, 5]$:

i, j	4	3	2	1
1				
2				
3				
4				

$A = [4, 1, 5, 3, 2]$:

i, j	4	3	2	1
1				$[1, 4, 5, 3, 2]$
2				
3		$[1, 4, 3, 5, 2]$	$[1, 3, 4, 5, 2]$	
4	$[1, 3, 4, 2, 5]$	$[1, 3, 2, 4, 5]$	$[1, 2, 3, 4, 5]$	

$A = [5, 4, 3, 2, 1]$:

i, j	4	3	2	1
1				$[4, 5, 3, 2, 1]$
2			$[4, 3, 5, 2, 1]$	$[3, 4, 5, 2, 1]$
3		$[3, 4, 2, 5, 1]$	$[3, 2, 4, 5, 1]$	$[2, 3, 4, 5, 1]$
4	$[2, 3, 4, 1, 5]$	$[2, 3, 1, 4, 5]$	$[2, 1, 3, 4, 5]$	$[1, 2, 3, 4, 5]$

(b) Briefly describe what the algorithm does in general. One word can be enough!

Sort

(c) For an input array A of length n , how many steps does the algorithm need to finish (in the worst case)? For simplicity, count the number of times the variable j is updated.

$$\frac{n(n-1)}{2} \cdot 2 = \boxed{n(n-1)}$$

2. (Proof Techniques: Contradiction and Contraposition, 15 points) In this problem, let n be a positive integer. You will prove the following claim using two different proof techniques.

Claim 1. If $3n + 2$ is odd, then n is odd.

(a) Prove the above claim by **contradiction**.

Assume $3n+2$ is odd and n is even. So $n = 2k$ for some integer k . $3n+2 = 3(2k)+2 = 6k+2$. $6k+2$ must be even as it can be written as $2(3k+1)$.
Therefore $3n+2$ is even, contradiction. \square

(b) Rewrite the contrapositive equivalent of the claim.

if n is even, then $3n+2$ is even

(c) Prove the claim you wrote in part (b). Please do not use a proof by contradiction in this question item because we want you to train on both techniques.

assume n is even. Therefore $n = 2k$ for some k .
if you multiply n by 3 and add 2 you get
 $3n+2 = 6k+2$. $6k+2$ can be written as $2(3k+1)$
and therefore is even. \square

3. (**Proofs by Induction, 15 points**) A *binary tree* is a rooted tree in which each node has at most two children. Let T be a binary tree with n nodes and let L denote the set of leaves and I the set of nodes with exactly two children. Prove by induction that it's always true that $|I| = |L| - 1$.

Hint: Start by writing down the statement you want to prove formally. Start by proving the base case. Then, prove the inductive case by first identifying the inductive hypothesis.

Proof. Base case:

Let T be a binary tree with only one node. Then $|I| = 0$, $|L| = 1$.

So, $|I| = |L| - 1$ is $0 = 1 - 1$ which holds.

Inductive Step: Assume that $|I| = |L| - 1$ holds for all binary trees with n nodes where $n \leq k$ for some integer k .

Inductive Hypothesis:

We must prove this holds for any tree w/ $k+1$ nodes. When adding a new node, there are two resulting cases.

(Case 1: The new node is an only child.)

In this case, $|I_{k+1}| = |I_k| + 0$ because no nodes w/ two children were created.

$|L_{k+1}| = (|L_k| - 1) + 1$ because the new leaf node is added as a child of an existing leaf node.

Therefore $|I_{k+1}| = |L_{k+1}| - 1$ equals $|I_k| = |L_k| - 1$ and by IH holds.

(Case 2: The new node is a sibling.)

In this case, $|I_{k+1}| = |I_k| + 1$ because the added node results in the parent having two children.

$|L_{k+1}| = |L_k| + 1$ because the added node is also a leaf node.

Therefore $|I_{k+1}| = |L_{k+1}| - 1$ equals $|I_k| + 1 = |L_k| + 1 - 1$ which equals $|I_k| = |L_k| - 1$. By IH this holds.

Conclusion:

For any binary tree with n nodes,

$$|I| = |L| - 1 \text{ holds.}$$

□

Hint 1: In the inductive step, consider deleting a leaf node.

Hint 2: What happens if it has a sibling?

When you delete the node, what happens if the only child?

4. (**Programming Assignment, 10 points**) Login to [Vjudge](#) and submit your solution to the programming assignment. You can choose either Pypy 3 or Python 3 as your language. No other programming language will be accepted.