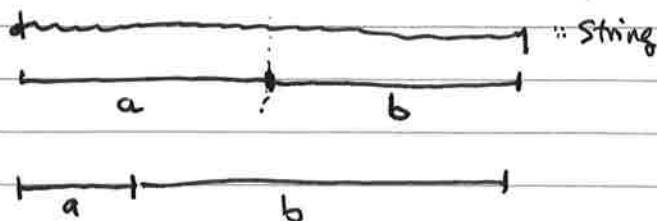


```
newtype Parser a :: Parser (String → [(String, a)])
```



```
(>>=) :: Parser a → (a → Parser b) → Parser b
```

$px \gg= f$ is a parser that first uses px to parse a value $x :: a$. The value x is used as a parameter to the function f , such that $fx :: \text{Parser } b$. We apply fx to the remaining tokens that were not consumed when parsing x .

NB. There may have been more than one result (x, ts) from applying the parser px , so we may have different fx combinations and final outputs.

$\text{oneOf} :: [\text{Char}] \rightarrow \text{Parser Char}$
 $\text{oneOf } xs = \text{satisfy } (\lambda x \rightarrow x \text{'elem' } xs)$

$\alpha :: \text{Parser Char}$
 $\alpha = \text{satisfy isAlpha}$

\uparrow
 $\text{isAlpha} :: \text{Char} \rightarrow \text{Bool}$
 from the prelude

or

$\alpha :: \text{Parser Char}$
 $\alpha = \text{oneOf} ([\text{'a'.. 'z'}] \# [\text{'A'.. 'Z'}])$

$\text{nat} :: \text{Parser Integer}$

$\text{nat} = \text{read} \langle \$ \rangle \text{ some } (\text{satisfy isDigit})$

\uparrow
 $:: \text{String} \rightarrow \text{Integer}$

\uparrow
 mysteries
 so far.

\uparrow
 we understand

from prelude



The parser $\text{some} :: \text{Parser } a \rightarrow \text{Parser } [a]$ is equivalent to the $+$ operation in BNF, where $\text{some } p$ is the same as p^+ . It is related to the parser $\text{many} :: \text{Parser } a \rightarrow \text{Parser } [a]$, which is where $\text{many } p$ corresponds to p^* in BNF.

$\text{some} :: \text{Alternative } f \Rightarrow f\ a \rightarrow f\ [a]$
which specializes to:

$\text{some} :: \text{Parser } a \rightarrow \text{Parser } [a]$

$\text{some } v = (\underbrace{(:)}_{a \rightarrow [a] \rightarrow [a]}) \underbrace{\langle \$ \rangle}_{\text{Parser } a} v \underbrace{\langle * \rangle}_{\text{Parser } [a]} \underbrace{\text{many } v}_{\text{Parser } [a]}$

$\text{many} :: \text{Alternative } \Rightarrow f\ a \rightarrow f\ [a]$
 $\text{many } v = \underbrace{\text{some } v}_{\text{Parser } [a]} \underbrace{\langle | \rangle}_{\text{Parser } [a]} \underbrace{\text{produce } []}_{\text{Parser } [a]}$

To understand some and many , we need to understand: $\langle * \rangle$, $\langle | \rangle$, and $\langle \$ \rangle$.

The simplest is: $\langle 1 \rangle$.

$p \langle 1 \rangle q$. This is an example of Alternative:

class Alternative f where

empty $:: f a$

$\langle 1 \rangle :: f a \rightarrow f a \rightarrow f a$

This has to satisfy some laws:

$$\text{empty} \langle 1 \rangle q = q$$

$$p \langle 1 \rangle \text{empty} = p$$

We won't force associativity i.e. $(p \langle 1 \rangle q) \langle 1 \rangle r = p \langle 1 \rangle (q \langle 1 \rangle r)$

Instances of this exist in Mayhead [3].

instance Alternative Maybe where
empty = Nothing

$$\text{Just } x \langle 1 \rangle q = \text{Just } x$$

$$\text{Nothing} \langle 1 \rangle q = q$$

instance Alternative [] where

empty = []

$xs <|> ys = xs \# ys$

The parser instance is a bit more complex:

instance Alternative Parser where

empty = failure

$(<|>) = \text{orElse}$

$\text{orElse} :: \text{Parser } a \rightarrow \text{Parser } a \rightarrow \text{Parser } a$
 $\text{orElse (Parser } px) (\text{Parser } py) = \text{Parser } (\lambda ts \rightarrow$

case px ts of

[] \rightarrow py ts

xs \rightarrow xs)

↑

$:: [(String, a)]$

The result of $\text{orElse } px \ py$ is a parser that first tries to parse px but if it fails, it tries py instead. If it succeeds then it is px .

The $\langle * \rangle$ is a member of the Applicative class:

class Applicative f where
 pure :: ~~f~~ a \rightarrow f a
 $\langle * \rangle$:: f (a \rightarrow b) \rightarrow f a \rightarrow f b.
 ↑
 pronounced 'ap' for apply.

instance Applicative [] where
 -- pure :: a \rightarrow [a]
 pure x = [x] -- $\langle * \rangle$:: [a \rightarrow b] \rightarrow [a] \rightarrow [b]
 fs $\langle * \rangle$ xs = [f x | f \leftarrow fs, x \leftarrow xs]
 ↑
 this applies each function in fs to each value in xs.

instance Applicative Maybe where
 -- pure :: a \rightarrow Maybe a
 pure x = Just x
 -- $\langle * \rangle$:: Maybe (a \rightarrow b) \rightarrow Maybe a \rightarrow Maybe b
~~Nothing~~ $\langle * \rangle$ ~~mx~~
 Just f $\langle * \rangle$ Just x = Just (f x)
 — $\langle * \rangle$ — = Nothing.