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'Identify the conclusion of the following argument. It is hard not to verify in our peers the same weakened intelligence due to emotions that we observe in our everyday patients. The arrogance of our consciousness, which in general, belongs to the strongest defense mechanisms, blocks the unconscious complexes. Because of this, it is difficult to convince people of the unconscious, and in turn to teach them what their conscious knowledge contradicts. (Sigmund Freud, The Origin and Development of Psychoanalysis)', A. 'It is hard not to verify in our peers the same weakened intelligence due to emotions that we observe in our everyday patients.', B. 'The arrogance of our consciousness, which in general, belongs to the strongest defense mechanisms, blocks the unconscious complexes.', C. 'Because of this, it is difficult to convince people of the unconscious, and in turn to teach them what their conscious knowledge contradicts.'

Answer Option 1:

The conclusion of the argument is: C. 'Because of this, it is difficult to convince people of the unconscious, and in turn to teach them what their conscious knowledge contradicts.'

Answer Option 2:

The conclusion of the argument is: C. 'Because of this, it is difficult to convince people of the unconscious, and in turn to teach them what their conscious knowledge contradicts.'

Answer Option 3:

The conclusion of the argument is stated in option C: 'Because of this, it is difficult to convince people of the unconscious, and in turn to teach them what their conscious knowledge contradicts.' Final Answer:

C

Answer Option 4:

The conclusion of the argument can be identified as: C. 'Because of this, it is difficult to convince people of the unconscious, and in turn to teach them what their conscious knowledge contradicts.'

Answer Option 5:

The conclusion of the argument is: C. 'Because of this, it is difficult to convince people of the unconscious, and in turn to teach them what their conscious knowledge contradicts.'

Answer Option 6:

The conclusion of the argument is: C. 'Because of this, it is difficult to convince people of the unconscious, and in turn to teach them what their conscious knowledge contradicts.' This is the main point that is

Answer Option 7:

The conclusion of the argument is: C. 'Because of this, it is difficult to convince people of the unconscious, and in turn to teach them what their conscious knowledge contradicts.' This is the final statement that summarizes the main point being made in the argument.

Answer Option 8:

The conclusion of the argument is C. 'Because of this, it is difficult to convince people of the unconscious, and in turn to teach them what their conscious knowledge contradicts.' This is stated as a result of the preceding statements about the arrogance of consciousness and the difficulties in observing and teaching about the unconscious mind.

Answer Option 9:

The conclusion of the argument is stated in option D: 'It is difficult to convince people of the unconscious, and in turn to teach them what their conscious knowledge contradicts.' This statement summarizes the main point made in the argument about the challenges of convincing people of the existence of the unconscious and teaching them concepts that contradict their conscious knowledge. Final Answer: D

Question 2:
'Select the best translation into predicate logic. David teaches Chris. (c: Chris; d: David; Txy: x teaches y)', A. 'Tdc', B. 'Tcd', C. 'Tcc', D. 'dTc'
Answer Option 1:
The best translation into predicate logic would be B. 'Tcd'.

Answer Option 2:
The best translation of the statement "David teaches Chris" into predicate logic would be: Tdc (which means "David teaches Chris"). Therefore, the correct answer is A.

Answer Option 3:
The best translation into predicate logic for the statement "David teaches Chris" would be 'Tdc'. Here, Txy represents "x teaches y", so Tdc translates to "David teaches Chris". Final Answer: A

Answer Option 4:
The best translation into predicate logic for the statement "David teaches Chris" would be:
Tdc (Option A)
Here, T stands for "teaches", d represents David, and c represents Chris. Thus, Tdc represents "David teaches Chris".
Final Answer: A

Answer Option 5:
The best translation into predicate logic for the statement "David teaches Chris" would be option B: 'Tcd'. This translates to "Chris is taught by David" in predicate logic, where Txy represents "x teaches y".

Answer Option 6:
The best translation into predicate logic for the statement "David teaches Chris" would be option B: 'Tcd'. This translates to "Chris is taught by David" which accurately represents the given statement. Final
Answer: B

Answer Option 7:
The best translation into predicate logic for the statement "David teaches Chris" would be: Tdc.
Here, Txy represents "x teaches y".
So, Tdc translates to "David teaches Chris", where d represents David and c represents Chris.
Final Answer: A. 'Tdc'

Answer Option 8:
The best translation into predicate logic for the statement "David teaches Chris" would be option B: 'Tcd'. This translates to "Chris is taught by David". In predicate logic, the variable before the predicate
symbol represents the subject, and the variable after the predicate symbol represents the object. Therefore, 'Tcd' accurately represents the statement that David teaches Chris. Final Answer: B

Answer Option 9:
The best translation into predicate logic for the statement "David teaches Chris" would be:
Tdc (Option A)
This translates to "David teaches Chris", where T represents the relation "teaches" and d and c represent David and Chris respectively. Final Answer: A

'Select the best English interpretation of the given proposition, using the following translation key: Ax: x is an apartment Hx: x is a house Lx: x is large Bxy: x is bigger than y (3x)[(Ax • Lx) • (3y)(Hy • Bxy)]', A. 'Some large houses are bigger than some apartments.', B. 'Some houses are bigger than all large apartments.', C. 'Some large apartments are bigger than some houses.', D. 'Some houses are bigger that n some large apartments.' Answer Option 1: To interpret the given proposition using the translation key, we can break it down as follows: $(\exists x)[(\mathsf{Ax} \bullet \mathsf{Lx}) \bullet (\exists y)(\mathsf{Hy} \bullet \mathsf{Bxy})]$ Translation key: Ax: x is an apartment Hx: x is a house Lx: x is large Bxy: x is bigger than y Breaking down the proposition: $(\exists x)[(Ax \cdot Lx) \cdot (\exists y)(Hy \cdot Bxy)]$ This proposition states that there is an x such that x is an apartment and x is large, and there exists a y such that y is a house and x is bigger than y. The best English interpretation of this proposition is: "Some large apartments are bigger than some houses." Therefore, the correct answer is C. "Some large apartments are bigger than some houses." Answer Option 2: To interpret the given proposition, let's break it down into parts: $(\exists x)[(Ax \bullet Lx) \bullet (\exists y)(Hy \bullet Bxy)]$ ∃x: There exists an x (Ax \bullet Lx): x is an apartment and x is large $\exists y$: There exists a y(Hy ullet Bxy): y is a house and y is bigger than x Putting it all together, the proposition can be interpreted as: "Some apartments that are large exist, and there exists a house that is bigger than some of those apartments." Therefore, the best English interpretation of the given proposition is: C. "Some large apartments are bigger than some houses." ********** Answer Option 3: Let's break down the given proposition using the translation key: $(\exists x)[(\mathsf{Ax} \bullet \mathsf{Lx}) \bullet (\exists y)(\mathsf{Hy} \bullet \mathsf{Bxy})]$ ∃x: There exists an x (Ax • Lx): x is an apartment and x is large $(\exists y)(Hy \bullet Bxy) :$ There exists a y where y is a house and x is bigger than y Putting it all together, the proposition can be interpreted as: "Some apartments are large and there exists a house that is smaller than some of these apartments." Therefore, the best English interpretation of the given proposition is: A. "Some large apartments are bigger than some houses." ********** Answer Option 4: To interpret the given proposition, let's break it down: $(\exists x)[(Ax \cdot Lx) \cdot (\exists y)(Hy \cdot Bxy)]$ ∃x: There exists an x (Ax • Lx): x is an apartment and x is large $\exists y$: There exists a y(Hy ullet Bxy): y is a house and y is bigger than x Putting it all together, the best English interpretation of the proposition is: "Some large apartments are bigger than some houses." Therefore, the correct answer is C. "Some large apartments are bigger than some houses."

Question 3:

Answer Option 5:

The given proposition states $(\exists x)[(Ax \bullet Lx) \bullet (\exists y)(Hy \bullet Bxy)]$. Translation key: Ax: x is an apartment Hx: x is a house Lx: x is large Bxy: x is bigger than y The best English interpretation of the given proposition is option C. 'Some large apartments are bigger than some houses.' This interpretation aligns with the structure of the proposition and correctly represents the existence of both large apartments and houses, with the apartments being bigger than the houses. Final Answer: C Answer Option 6: Let's break down the proposition using the translation key: $(\exists x)[(\mathsf{Ax} \bullet \mathsf{Lx}) \bullet (\exists y)(\mathsf{Hy} \bullet \mathsf{Bxy})]$ Translation: $(\exists x)$ - There exists an x $(Ax \cdot Lx) - x$ is an apartment and x is large (∃y) - There exists a y (Hy ullet Bxy) - y is a house and x is bigger than y Based on this translation, the best English interpretation of the given proposition is: "Some apartments that are large are bigger than some houses." Therefore, the correct answer is: C. "Some large apartments are bigger than some houses." Answer Option 7: To interpret the given proposition, let's break it down step by step using the translation key: $(\exists x)[(\mathsf{Ax} \bullet \mathsf{Lx}) \bullet (\exists y)(\mathsf{Hy} \bullet \mathsf{Bxy})]$ There exists an x such that (x is an apartment and x is large) and there exists a y such that (y is a house and x is bigger than y). Based on this translation, the best English interpretation of the proposition is: A. 'Some large houses are bigger than some apartments.' Explanation: The proposition states that there exists at least one apartment that is large, and there exists at least one house such that this house is smaller than the apartment. This can be interpreted as some large houses are bigger than some apartments. Final Answer: A Answer Option 8: To interpret the given proposition, we can break it down step by step: $(\exists x)[(\mathsf{Ax} \bullet \mathsf{Lx}) \bullet (\exists y)(\mathsf{Hy} \bullet \mathsf{Bxy})]$ First, let's focus on ($\exists x$)($Ax \cdot Lx$), which translates to "There exists an x such that x is an apartment and x is large." This means that there is at least one large apartment. Next, let's look at (3y)(Hy • Bxy), which translates to "There exists a y such that y is a house and y is bigger than x." This means that there is at least one house that is bigger than the apartment mentioned Combining these two parts, we get "(∃x)[(Ax • Lx) • (∃y)(Hy • Bxy)]" as "There exists a large apartment such that there exists a house that is bigger than it." From the given options, the best English interpretation of this proposition is: C. "Some large apartments are bigger than some houses." Answer Option 9: Let's analyze the given proposition using the translation key: $(\exists x)[(\mathsf{Ax} \bullet \mathsf{Lx}) \bullet (\exists y)(\mathsf{Hy} \bullet \mathsf{Bxy})]$ According to the translation key: Ax: x is an apartment Hx: x is a house Lx: x is large Bxy: x is bigger than y The proposition can be broken down as follows: $(\exists x)[(\mathsf{Ax} \bullet \mathsf{Lx}) \bullet (\exists y)(\mathsf{Hy} \bullet \mathsf{Bxy})]$ First, let's focus on "(Ax • Lx)": (Ax • Lx) means "x is an apartment and x is large." Next, let's consider $(\exists y)(Hy \bullet Bxy)$:

($\exists y$)(Hy • Bxy) means "there exists a y such that y is a house and y is bigger than x."

Combining the two parts, we have:

$(\exists x)[(\mathsf{Ax} \bullet \mathsf{Lx}) \bullet (\exists y)(\mathsf{Hy} \bullet \mathsf{Bxy})]$

This can be interpreted as "There exists an apartment that is large and there exists a house that is bigger than the apartment."

The best English interpretation of the given proposition is:

"Some apartments are large, and there exists a house that is bigger than those apartments."

Therefore, the correct answer is A. "Some large houses are bigger than some apartments."

Question 4:
'Construct a complete truth table for the following argument. Then, using the truth table, determine whether the argument is valid or invalid. If the argument is invalid, choose an option which presents a
counterexample. (There may be other counterexamples as well.)
~G⊃H
~H / G', A. 'Valid', B. 'Invalid. Counterexample when G and H are true', C. 'Invalid. Counterexample when G is true and H is false', D. 'Invalid. Counterexample when H is true and G is false'
Answer Option 1:
To construct a truth table for the argument, we need to consider the truth values of the premises (statements) and the conclusion. Let's analyze the given argument:
Premise 1: ~G ⊃ H
Premise 2: ~H
Conclusion: G
We can construct the truth table by considering all possible combinations of truth values for G and H:
GH~G~H~G⊃H
TTFFT
TFFTF
FTTFT
FFTTT
According to the truth table, when both premises are true (shown in the truth table as 'T'), the conclusion is also true. Therefore, the argument is valid. Answer: A. Valid.
Anguar Ontion 9:
Answer Option 2: To construct a complete truth table for the given argument, we consider all possible combinations of truth values for the propositions G and H.
GH~G⊃H~HG'
TFFTT
FTTFF
FFTTT
From the truth table, we can observe that when both ~G and ~H are true, G' is also true. Hence, the argument is valid.
Therefore, the correct answer is A. 'Vallid'.
Answer Option 3:
To construct a complete truth table for the given argument, we need to consider all possible combinations of truth values for the propositions G and H, as well as the given premises ~G \supset H and ~H.
GH-G-G⊃H-H
TTETE
TFFTT
FITTE
FFTFT
Based on the truth table, we can see that there is no row where both premises ~G ⊃ H and ~H are true while G is false. Therefore, the argument is valid.
Final Answer: A. 'Valid'
Answer Option 4:
To construct a truth table for the given argument, we need to consider the following propositions:
G: G is true
H: H is true
Then, we can represent the premises and conclusion as follows:
Premise 1: ~G ⊃ H
Premise 2: ~H
Conclusion: G
Let's construct the truth table:
G H ~G ~H ~G⊃H ~H/G

TITIFIFITIF TIFIFITITIF F|T|T|F|F|T F|F|T|T|T In the truth table, we see that there is no row where both premises are true and the conclusion is false. Therefore, the argument is valid. Final Answer: A. 'Valid' Answer Option 5: To construct a complete truth table for the argument, we need to consider all possible combinations of truth values for the propositions involved: G and H. $GH \sim G \sim H \sim G \supset H \sim H/G$ TTFFTFF TEETTTE FTTFFTT FFTTTTT From the truth table, we can see that there is a row where the premises (~H) and (~G ⊃ H) are both true, but the conclusion (G) is false. This means that the argument is invalid. To find a counterexample, we can look for a row where the premises are true and the conclusion is false. In this case, the counterexample is when G is false and H is true. This satisfies the premises (-H) and (~G ⊃ H), but the conclusion (G) is false. Therefore, the correct answer is: C. 'Invalid. Counterexample when G is true and H is false'. Answer Option 6: To construct a truth table for the argument, we will define the following propositions: G: G is true H: H is true We are given the premises: $1.~{\sim}G\supset H$ We need to determine the truth value of the conclusion: G. The truth table for the argument is as follows: |G|H|~G|~H|~G⊃H| |---|---| |T|T|F|F|T| |T|F|F|T|T| |F|T|T|F|T| |F|F|T|T|F| Based on the truth table, we can see that when ~G and ~H are true, the conclusion G is false. This means the argument is invalid. Therefore, the correct answer is option C: 'Invalid. Counterexample when G is true and H is false'. Answer Option 7: To construct a complete truth table for the argument, we need to consider all possible truth values for the propositions G and H. The argument has two premises: 1. ~G ⊃ H 2. ~H We want to determine whether the conclusion "G" can be validly inferred from the premises. Let's create a truth table to evaluate the truth values of ~G, H, and the conclusion G based on the given premises: |G|H|~G|~G⊃H|~H|G| |---|---|----| |T|T|F|T|F|T||T|F|F|T|T|T| |F|T|T|T|F|F|

The truth table reveals that there is at least one row where both premises are true (row 2) and the conclusion is false (G is false). Therefore, the argument is invalid.

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To find a counterexample, we look for a row where both premises are true and the conclusion is false. In this case, row 2 satisfies these conditions, where G is true and H is false. Thus, the counterexample is when G is true and H is false.

Final Answer: C. Invalid. Counterexample when G is true and H is false.

Answer Option 8:
To construct a truth table for the argument, we need to consider all possible truth values of the propositions involved: $G, H, \sim G, \sim H$, and the given statement $\sim G \supset H$.
G H -G -H -G⊃H
T T F F T
TIFIFITIT
F T T F T
F F T T F
Now, let's analyze the truth values in the final column based on the given statement $\sim G \supset H$ and the premise $\sim H$.
In the first row, \sim G is false (F) and \sim H is false (F), and the statement \sim G \supset H is true (T).
In the second row, \sim G is false (F) and \sim H is true (T), and the statement \sim G \supset H is true (T).
In the third row, \sim G is true (T) and \sim H is false (F), and the statement \sim G \supset H is true (T).
In the fourth row, \sim G is true (T) and \sim H is true (T), and the statement \sim G \supset H is false (F).
Based on the truth table, we can see that for every row where both the premise \sim H and the statement \sim G \supset H are true, the conclusion G is also true. Therefore, the argument is valid.
The correct answer is A. Valid.

Answer Option 9:
To construct a complete truth table for the argument, we need to consider all possible combinations of truth values for the propositions involved. In this case, we have two propositions, G and
GH~G⊃H~HG
TTFTFT
TFFTTT
FTTFFF
FFTTTF
From the truth table, we can evaluate the argument. The argument states $\sim G \supset H$ and $\sim H$, and we need to determine whether G is true or false.
For the first row, $\neg G$ is false (F) and $\neg H$ is true (T). According to $\neg G \supset H$, since $\neg G$ is false, the whole statement is true, and it satisfies the conclusion.
For the second row, \sim G is false (F) and \sim H is false (F). According to \sim G \supset H, since \sim G is false, the whole statement is true, and it satisfies the conclusion.
For the third row, \sim G is true (T) and \sim H is true (T). According to \sim G \supset H, since \sim G is true, the whole statement is true, and it satisfies the conclusion.
For the fourth row, \sim G is true (T) and \sim H is false (F). According to \sim G \supset H, since \sim G is true, the whole statement is false, and it does not satisfy the conclusion.
Since there is at least one row in the truth table where the premises are true and the conclusion is false, the argument is invalid.
The counterexample is when G is true and H is false.
Therefore, the correct answer is C. 'Invalid. Counterexample when G is true and H is false'.

'Use the following key to translate the given formula of PL to natural, English sentences.
A: Marina reads a Percy Jackson book.
B: Izzy plays Minecraft.
C: Emily stops working.
D: Russell makes dinner.
E: Ashleigh stops by.
(-B • E) \supset D', A. 'If it's not the case that both Izzy plays Minecraft and Ashleigh stops by then Russell makes dinner.', B. 'If it's both not the case that Izzy plays Minecraft and Ashleigh stops by then Russell
makes dinner.', C. 'Izzy playing Minecraft and Ashleigh stopping by entails that Russell makes dinner.', D. 'Russell's making dinner is a sufficient condition for Izzy not playing Minecraft and Ashleigh stopping
by.'
Answer Option 1:
Using the given key: A: Marina reads a Percy Jackson book. B: Izzy plays Minecraft. C: Emily stops working. D: Russell makes dinner. E: Ashleigh stops by. The formula (-B • E) \supset D translates to: 'If it's not
the case that Izzy plays Minecraft and Ashleigh stops by, then Russell makes dinner.' Final Answer: A
Answer Option 2:
Using the given key, we can translate the formula ($-B \cdot E$) $\supset D$ as follows:
(~B • E) ⊃ D translates to 'If it's not the case that Izzy plays Minecraft and Ashleigh stops by, then Russell makes dinner.'
Therefore, the correct translation is option A: 'If it's not the case that both Izzy plays Minecraft and Ashleigh stops by, then Russell makes dinner.'
Answer Option 3:
Using the key provided, we can translate the given formula $(-B \bullet E) \supset D$ to English sentences as follows:
(~B • E) ⊃ D translates to "If it's not the case that Izzy plays Minecraft and Ashleigh stops by, then Russell makes dinner."
So the correct translation is option A: "If it's not the case that both Izzy plays Minecraft and Ashleigh stops by then Russell makes dinner." Final Answer: A
Answer Option 4:
Using the provided key, we can translate the given formula $(-B \bullet E) \supset D$ as follows:
$(-B \cdot E) \supset D$
Translation:
'If it's not the case that Izzy plays Minecraft and Ashleigh stops by, then Russell makes dinner.'
Therefore, the correct translation is option A: 'If it's not the case that both Izzy plays Minecraft and Ashleigh stops by then Russell makes dinner.' Final Answer: A.
Answer Option 5:
The given formula is $(\sim B \cdot E) \supset D$.
Let's translate each proposition into English sentences using the provided key:
A: Marina reads a Percy Jackson book.
B: Izzy plays Minecraft.
C: Emily stops working.
D: Russell makes dinner.
E: Ashleigh stops by.
Therefore, the translation of the formula $(-B \cdot E) \supset D$ is:
'If it's not the case that Izzy plays Minecraft and Ashleigh stops by, then Russell makes dinner.'
Final Answer: A
Answer Option 6:
Using the key provided, we can translate the given formula $(-B \bullet E) \supset D$ as follows:
~B: Izzy does not play Minecraft.
E: Ashleigh stops by.
D: Russell makes dinner.
Therefore, the translation of $(-B \cdot E) \supset D$ is: 'If it's not the case that Izzy plays Minecraft and Ashleigh stops by, then Russell makes dinner.'
Final Answer: A
Answer Option 7:

Question 5:

To translate the given formula (-B • E) ⊃ D into natural English sentences using the key provided:
~B: Izzy does not play Minecraft.
E: Ashleigh stops by.
D: Russell makes dinner.
⊃: implies, or ifthen
The correct translation is: 'If it's not the case that Izzy plays Minecraft and Ashleigh stops by, then Russell makes dinner.'
So, the correct answer is A: 'If it's not the case that both Izzy plays Minecraft and Ashleigh stops by, then Russell makes dinner.'
Answer Option 8:
Using the key provided, let's translate the given formula $(-B \bullet E) \supset D$ to English sentences:
~B: Izzy does not play Minecraft.
E: Ashleigh stops by.
D: Russell makes dinner.
$(-B \cdot E) \supset D$ translates to:
"If it's the case that Izzy does not play Minecraft and Ashleigh stops by, then Russell makes dinner."
Therefore, the correct translation is:
A. "If it's not the case that both Izzy plays Minecraft and Ashleigh stops by, then Russell makes dinner."
Answer Option 9:
To translate the given formula $(-B \bullet E) \supset D$ to natural English sentences, we can use the following key:
A: Marina reads a Percy Jackson book.
B: Izzy plays Minecraft.
C: Emily stops working.
D: Russell makes dinner.
E: Ashleigh stops by.
The formula $(\sim B \bullet E) \supset D$ can be translated as:
'If it's not the case that Izzy plays Minecraft and Ashleigh stops by, then Russell makes dinner.'
Therefore, the correct translation is: A. 'If it's not the case that both Izzy plays Minecraft and Ashleigh stops by then Russell makes dinner.'

Select the best English interpretation of the given proposition, using the following translation key: Ax: x is an apartment Hx: x is a house Lx: x is large Bxy: x is bigger than y (∃x)[Hx • (∀y)(Ay ⊃ Bxy)], A. 'Some houses are smaller than all apartments.', B. 'Every house is bigger than every apartment.', C. 'Some apartment is smaller than every house.', D. 'Some houses are bigger than every apartment.' Answer Option 1: The given proposition $(\exists x)[Hx \cdot (\forall y)(Ay \supset Bxy)]$ can be translated as "There exists an object x that is a house and for all objects y, if y is an apartment, then x is bigger than y." Based on this translation, the correct interpretation is: B. 'Every house is bigger than every apartment.' Answer Option 2: Let's break down the given proposition using the translation key: $(\exists x)[Hx \bullet (\forall y)(Ay \supset Bxy)]$ $\exists x$ - There exists an xHx - x is a house ∀y - For all y Ay - y is an apartment Bxy - x is bigger than y The best English interpretation of the given proposition is: "Some houses are bigger than every apartment." Final Answer: D. "Some houses are bigger than every apartment." Answer Option 3: The given proposition $(\exists x)[Hx \bullet (\forall y)(Ay \supset Bxy)]$ can be translated as "There exists an object x which is a house and for all objects y, if y is an apartment, then x is bigger than y." The best English interpretation of the proposition is: B. 'Every house is bigger than every apartment.' This interpretation captures the meaning of the proposition accurately, stating that for every house (Hx) and every apartment (Ay), the house (x) is bigger than the apartment (y). Answer Option 4: The given proposition is $(\exists x)[Hx \cdot (\forall y)(Ay \supset Bxy)]$. Let's break it down: $(\exists x)$ means "There exists an x" or "Some x". $Hx \cdot (\forall y)(Ay \supset Bxy)$ means "x is a house" and "For all y, if y is an apartment, then x is a house" and "For all y, if y is an apartment, then x is a house" and "For all y, if y is an apartment, then y is a house "y is a house" and "yThe best English interpretation of the given proposition is: "Some houses are bigger than every apartment." Therefore, the correct answer is D. "Some houses are bigger than every apartment." Answer Option 5: The given proposition $(\exists x)[Hx \bullet (\forall y)(Ay \supset Bxy)]$ can be translated as "There exists an x such that x is a house and for all y, if y is an apartment, then x is bigger than y." The best English interpretation of this proposition would be: B. "Every house is bigger than every apartment." This is because the proposition states that for every apartment y, there exists a house x that is bigger than y. In other words, it implies that every apartment is smaller than every house. Therefore, option B is the most accurate interpretation. Answer Option 6: The given proposition is $(\exists x)[Hx \bullet (\forall y)(Ay \supset Bxy)]$. Using the translation key, let's break it down: $(\exists x)$ - There exists an x (There is at least one x) $[\mathsf{Hx} \bullet (\forall y)(\mathsf{Ay} \supset \mathsf{Bxy})] \text{ - such that } x \text{ is a house and for all } y, \text{ if } y \text{ is an apartment, then } x \text{ is bigger than } y.$ Based on this, the best English interpretation of the proposition is "Some houses are bigger than every apartment." Therefore, the correct answer is D. "Some houses are bigger than every apartment." Answer Option 7: To interpret the given proposition, let's break it down: $(\exists x)[\mathsf{Hx} \bullet (\forall y)(\mathsf{A}y \supset \mathsf{B}xy)]$ $\exists x$: There exists at least one x

Question 6:

Ay: y is an apartment Bxy: x is bigger than y Using the translation key, we can interpret the proposition as follows:

"Some house exists such that for all apartments, the house is bigger than the apartment."

This can be represented as:

"Some houses are bigger than all apartments."

Therefore, the best English interpretation of the given proposition is option D: "Some houses are bigger than every apartment."

Answer Option 8:

 $To interpret the given proposition \ (\exists x)[Hx \bullet (\forall y)(Ay \supset Bxy)], we can break it down into smaller parts using the translation key:$

 $(\exists x)[\mathsf{Hx} \bullet (\forall y)(\mathsf{Ay} \supset \mathsf{Bxy})]$

The $(\exists x)$ part indicates that there exists at least one x that satisfies the following conditions.

Hx represents x being a house.

 $(\forall y)$ indicates that for all y, the following condition holds.

Ay represents y being an apartment.

 $\mbox{Ay} \supset \mbox{Bxy represents the statement "If y is an apartment, then x is bigger than y."$

Based on this interpretation, the best English interpretation of the given proposition is:

B. "Every house is bigger than every apartment."

Answer Option 9:

Let's break down the given proposition using the translation key:

 $(\exists x)[\mathsf{Hx} \bullet (\forall y)(\mathsf{Ay} \supset \mathsf{Bxy})]$

The symbol $(\exists x)$ indicates the existence of an x that satisfies the following conditions.

Hx represents x being a house.

 $(\forall y)(Ay\supset Bxy) \text{ states that for all } y, \text{ if } y \text{ is an apartment (Ay), then } x \text{ is bigger than } y \text{ (Bxy)}.$

Combining these translations, the proposition can be interpreted as:

"There exists a house such that for all apartments, the house is bigger than the apartment."

This can be best expressed as:

B. "Every house is bigger than every apartment."

Therefore, the correct answer is B.

Question 7:

Construct a complete truth table for the following argument. Then, using the truth table, determine whether the argument is valid or invalid. If the argument is invalid, choose an option which presents a counterexample. (There may be other counterexamples as well.)

~K \equiv L / (K · L) \vee (K · ~L) ', A. 'Valid', B. 'Invalid'. Counterexample when K and L are true', C. 'Invalid'. Counterexample when K is true and L is false', D. 'Invalid'. Counterexample when L is true and K is false'

Answer Option 1:

To construct a truth table, we need to consider all possible combinations of truth values for the propositions involved. In this case, we have two propositions, K and L.

 $\mathsf{K}\;\mathsf{L}\;\mathsf{\sim}\mathsf{K}\equiv\mathsf{L}\;\mathsf{K}\;\!\cdot\;\mathsf{L}\;\mathsf{K}\;\!\cdot\;\!\!\sim\;\!\!\mathsf{L}\;(\mathsf{K}\;\!\cdot\;\mathsf{L})\;\!\vee\;(\mathsf{K}\;\!\cdot\;\!\!\sim\;\!\!\mathsf{L})$

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From the truth table, we can see that there are instances where the expression $-K \equiv L$ is true but the expression $(K \cdot L) \vee (K \cdot -L)$ is false. This means that the argument is invalid.

To find a counterexample, we look for a row in the truth table where the premise (-K = L) is true but the conclusion ((K · L) ∨ (K · ~L)) is false. In this case, the counterexample occurs when K is true and L is false.

Therefore, the correct answer is C. 'Invalid. Counterexample when K is true and L is false'.

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Answer Option 2:

To construct a complete truth table for the argument, we need to consider all possible combinations of truth values for the propositions involved.

Let's create the truth table:

| K | L | ~K | L / ~K | K · L | K · ~L | (K · L) ∨ (K · ~L) |

|F|F|T|F|F|F|F|

|F|T|T|T|F|T|T|

|T|F|F|T|F|F|T|

|T|T|F|F|T|F|T|

Based on the truth table, the argument is valid because in all cases where \sim K is equivalent to L, the statement (K \cdot L) \vee (K \cdot \sim L) is also true.

Therefore, the correct answer is A. 'Valid'.

Answer Option 3:

To construct a complete truth table, we need to consider all possible combinations of truth values for the propositions involved. In this case, we have two propositions, K and L.

 $K \mid L \mid \sim K \mid \sim K \equiv L \mid K \cdot L \mid K \cdot \sim L \mid (K \cdot L) \lor (K \cdot \sim L)$

T|T|F|F|T|F|T

T|F|F|T|F|T|T

F|T|T|T|F|F|F

F|F|T|F|F|F|F

From the truth table, we can see that there is at least one row where the statement "-K \equiv L" is true but the statement "(K · L) \vee (K · ~L)" is false. This is the third row where K is false and L is true.

Therefore, the argument is invalid.

The correct option is C. 'Invalid. Counterexample when K is true and L is false'.

Answer Option 4:

To construct a complete truth table for the given argument, we will consider all possible combinations of truth values for the propositions K and L.

 $\mathsf{K} \mid \mathsf{L} \mid \mathsf{\sim}\mathsf{K} \mid \; \equiv \; \mid \mathsf{L} \mid \mathsf{K} \cdot \mathsf{L} \mid \mathsf{K} \cdot \mathsf{\sim}\mathsf{L} \mid (\mathsf{K} \cdot \mathsf{L}) \vee (\mathsf{K} \cdot \mathsf{\sim}\mathsf{L})$

T|T|F|T|T|T|F|T

T|F|F|F|F|T|T

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F|F|T|T|F|F|F|F

In the truth table, we can see that there is a row where the premise $(-K \equiv L)$ is true (row 2) and the conclusion $((K \cdot L) \lor (K \cdot -L))$ is false. Therefore, the argument is invalid.

Counterexample: K is true and L is false.

Option C: 'Invalid. Counterexample when K is true and L is false' correctly describes the invalidity of the argument. Final Answer: C

******* Answer Option 5: To construct a complete truth table for the argument, we need to evaluate the truth values of all the propositions involved. Let's consider the propositions: K: K is true L: L is true Using the given information, ${\sim}K\equiv L$ translates to ${\neg}K \leftrightarrow L.$ We can now construct the truth table: $\mathsf{K} \mid \mathsf{L} \mid \mathsf{\sim\!K} \mid \mathsf{\neg\!K} \leftrightarrow \mathsf{L} \mid \mathsf{K} \cdot \mathsf{L} \mid \mathsf{K} \cdot \mathsf{\sim\!L} \mid (\mathsf{K} \cdot \mathsf{L}) \vee (\mathsf{K} \cdot \mathsf{\sim\!L})$ TITIFITITIFIT T|F|F|F|F|T|T F|T|T|T|F|F|F F|F|T|F|F|F|F From the truth table, we can see that there is a row where the proposition $\sim K \equiv L$ is true, but the conclusion $(K \cdot L) \vee (K \cdot -L)$ is false. Specifically, when K is true and L is false. Therefore, the argument is invalid, and the correct option is C. 'Invalid. Counterexample when K is true and L is false'. Answer Option 6: To construct a complete truth table, we need to consider all possible truth values for the propositions involved. In this case, we have two propositions, K and L. $\mathsf{K} \mid \mathsf{L} \mid \mathsf{\sim} \mathsf{K} \mid \mathsf{K} \cdot \mathsf{L} \mid \mathsf{K} \cdot \mathsf{\sim} \mathsf{L} \mid (\mathsf{K} \cdot \mathsf{L}) \vee (\mathsf{K} \cdot \mathsf{\sim} \mathsf{L})$ TITIFITIFIT T|F|F|F|T|T FITITIFIFIF F|F|T|F|F|F The given argument states ~K \equiv L, which means that ~K and L have the same truth value. Looking at the truth table, we can see that this is true when K and L are both true or both false. Now, let's analyze the expression $(K \cdot L) \vee (K \cdot \neg L)$. From the truth table, we can see that this expression evaluates to true in all cases except when K is false and L is true. Since the expression is always true when ~K = L is true, we can conclude that the argument is valid. Therefore, the correct option is A. 'Valid'. Answer Option 7: To construct a complete truth table for the given argument, we need to consider all possible combinations of truth values for the propositions involved. In this case, we have two propositions: K and L. $\mathsf{K} \; \mathsf{L} \; \mathord{\sim} \mathsf{K} \; \mathsf{L} \; \equiv \; \mathord{\sim} \mathsf{K} \; \mathsf{K} \; \cdot \; \mathsf{L} \; \mathsf{K} \; \cdot \; \mathord{\sim} \mathsf{L} \; (\mathsf{K} \; \cdot \; \mathsf{L}) \; \vee \; (\mathsf{K} \; \cdot \; \mathord{\sim} \mathsf{L})$ TTEETET TEETETT FTTTFFF FFTFFF After constructing the truth table, we can analyze the last column to determine whether the argument is valid or invalid. If the column under the argument's conclusion always evaluates to true (T), then the argument is valid. If there exists at least one row where the conclusion evaluates to false (F), then the argument is invalid. In this case, the conclusion column is not always true (T). Therefore, the argument is invalid. To find a counterexample, we need to identify a row in the truth table where the premises (propositions before the "/") are true, but the conclusion is false. Looking at the truth table, we can see that when K is true and L is false, the premises \sim K \equiv L is true, but the conclusion (K \cdot L) \vee (K \cdot ~L) is false. So, the correct answer is option C: "Invalid. Counterexample when K is true and L is false." ********** Answer Option 8: To construct a complete truth table, we need to consider all possible truth value combinations for the variables involved in the argument. In this case, we have two variables, K and L. Truth table for $\sim K \equiv L$: $K \mid L \mid \sim K \mid L \mid \sim K \equiv L \mid$ TITIFITIFI T|F|F|F|T|

F|T|T|T|T| F|F|T|F|F| Truth table for $(K \cdot L) \vee (K \cdot \neg L)$: $K \mid L \mid \neg L \mid K \cdot L \mid K \cdot \neg L \mid (K \cdot L) \vee (K \cdot \neg L) \mid$ \cdots $T \mid T \mid F \mid T \mid F \mid T \mid T \mid$ $F \mid T \mid F \mid F \mid F \mid F \mid F \mid$

 $Comparing the truth values of \neg K \equiv L \ and \ (K \cdot L) \lor (K \cdot \neg L), \ we \ can \ see \ that \ they \ are \ the \ same \ for \ all \ possible \ combinations \ of \ K \ and \ L.$

Therefore, the argument is valid. Option A: 'Valid' is the correct answer.

Answer Option 9:

F|F|T|F|F|F|

To construct a complete truth table for the argument, we need to consider all possible combinations of truth values for the propositions involved: K and L.

 $K \mid L \mid \sim K \mid L \equiv \sim K \mid K \cdot L \mid K \cdot \sim L \mid (K \cdot L) \lor (K \cdot \sim L)$

T|T|E|E|T|E|T
T|E|E|T|E|T|
E|T|T|E|E|E
E|E|T|T|E|E|E

Based on the truth table, we can see that the last column $(K \cdot L) \vee (K \cdot \neg L)$ does not always have the same truth value as $\neg K \equiv L$. This means that the argument is invalid.

To find a counterexample, we look for a row in which $-K \equiv L$ is true but $(K \cdot L) \lor (K \cdot -L)$ is false. In this case, the row where K is true and L is false satisfies this condition:

 $K \mid L \mid \sim K \equiv L \mid (K \cdot L) \lor (K \cdot \sim L)$

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Therefore, the argument is invalid, and the counterexample is when K is true and L is false.

Final Answer: C. Invalid. Counterexample when K is true and L is false.

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'Identify the antecedent of the following conditional proposition: If the university does not increase financial aid, either the president fails to approve it or the board of trustees prevents it.', A. The university increases financial aid.', B. The university does not increase financial aid.', C. The board of trustees prevents it.', D. The president fails to approve it.'

Answer Option 1:

The antecedent of a conditional proposition is the part that comes after the "if" and before the "then" in the sentence. In this case, the antecedent is 'The university does not increase financial aid.' Therefore, the correct answer is B. The university does not increase financial aid.'

Answer Option 2:

The antecedent of a conditional proposition is the part that comes before the "if". In this case, the antecedent is "The university does not increase financial aid." Therefore, the correct answer is B. The university does not increase financial aid.'

Answer Option 3:

The antecedent of a conditional proposition is the part that comes before the "if" statement. In this case, the antecedent is "The university does not increase financial aid."

Therefore, the correct answer is B. 'The university does not increase financial aid.'

Answer Option 4:

The antecedent of a conditional proposition is the part that comes before the "if" in the statement. In this case, the antecedent is "The university does not increase financial aid."

Therefore, the correct answer is B. 'The university does not increase financial aid.'

Answer Option 5:

The antecedent of a conditional proposition is the statement that follows the "if" part. In this case, the statement is "If the university does not increase financial aid." Therefore, the antecedent of the given conditional proposition is B. The university does not increase financial aid.' Final Answer: B

Answer Option 6:

The antecedent of a conditional proposition is the statement that comes immediately after the "if" part of the conditional. In this case, the conditional proposition is "If the university does not increase financial aid, either the president fails to approve it or the board of trustees prevents it." The antecedent of this conditional proposition is "The university does not increase financial aid."

Therefore, the correct answer is B. 'The university does not increase financial aid.'

Answer Option 7:

The antecedent of a conditional proposition is the statement that comes immediately after the "if" in the conditional statement. In this case, the conditional proposition is "if the university does not increase financial aid, either the president fails to approve it or the board of trustees prevents it."

The antecedent of this conditional proposition is: "The university does not increase financial aid."

Therefore, the correct answer is B. 'The university does not increase financial aid.'

Answer Option 8:

The antecedent of a conditional proposition is the statement that comes immediately after the "if" in the sentence. In this case, the conditional proposition is "If the university does not increase financial aid, either the president fails to approve it or the board of trustees prevents it." The antecedent is the statement that follows the "if" clause, which is "The university does not increase financial aid."

Therefore, the correct answer is B. 'The university does not increase financial aid.'

Answer Option 9:

The antecedent of a conditional proposition is the part that comes before the "if" or the "then" in the sentence. In this case, the conditional proposition is:

"If the university does not increase financial aid, either the president fails to approve it or the board of trustees prevents it."

The antecedent is the part that comes before the "if" in this sentence. Therefore, the antecedent in this case is:

B. "The university does not increase financial aid."

Final Answer: B

Question 9: 'Which of the following propositions is an immediate (one-step) consequence in PL of the given premises? \sim E ⊃ \sim F $G\supset F$ H ∨ ~E H⊃I ~I', A. 'E \supset F', B. 'F \supset G', C. 'H \supset ~E', D. '~H' Answer Option 1: To determine the immediate consequence of the given premises, we can use the rules of propositional logic. Given premises: 1. ~E ⊃ ~F $2.\ G\supset F$ 3. H ∨ ~E 4. H ⊃ I 5. ~I To find the immediate consequence, we need to find a statement that can be derived directly from the given premises using valid logical rules. Looking at the premises, we can see that if we assume E to be true, then by modus ponens we can conclude F from premise 1. Therefore, the immediate consequence of the given premises is 'E \supset F'. Final Answer: A ****** Answer Option 2: To determine the immediate (one-step) consequence of the given premises, we can use inference rules in Propositional Logic (PL). Let's analyze each option: Option A: $'E \supset F'$ To derive this proposition, we need the premise 'G \supset F', but this premise is not used in the given premises. Therefore, option A is not an immediate consequence. Option B: ${}^{t}F \supset G'$ This proposition can be derived using the premise $G \supset F'$ through the contrapositive rule. Therefore, option B is an immediate consequence. Option C: 'H > ~E' This proposition cannot be derived directly from the given premises. Therefore, option C is not an immediate consequence. Option D: '~H' This proposition can be derived using the premise 'H v ~E' and the assumption '~E' through the disjunctive syllogism rule. Therefore, option D is an immediate consequence. Final Answer: B and D ********** Answer Option 3: To determine the immediate (one-step) consequence in PL, we need to analyze the given premises and apply valid logical deductions. Premises: 1. ~E ⊃ ~F $2.\;G\supset F$ 3. H ∨ ~E 4. H ⊃ I 5. ~I Let's evaluate the options: Option A: $'E \supset F'$ This statement cannot be directly derived from the given premises. Option B: ${}^{t}F \supset G'$ This statement cannot be directly derived from the given premises. Option C: $^{\prime}H\supset ^{\sim}E^{\prime}$ This statement can be derived using Modus Tollens with premises 3 and 1. From premise 3 (H \vee ~E) and premise 1 (~E \supset ~F), we can deduce ~H \supset F. By applying Modus Tollens (\sim H \supset F, \sim F \supset \sim E), we get H \supset \sim E. This statement cannot be directly derived from the given premises.

The correct answer is:

C. 'H ⊃ ~E' This statement can be derived as an immediate consequence of the premises using the Modus Tollens rule. Answer Option 4: To find the immediate consequence of the given premises, we can use the rules of inference in propositional logic. Premises: 1. ~E ⊃ ~F 2. G ⊃ F 3. H ∨ ~E 4. H ⊃ I We will use the Modus Tollens (MT) rule, which states that if we have a conditional statement of the form "If A, then B" and we know that B is false, then we can infer that A is false as well. Applying MT to premise 2 (G \supset F) and premise 5 (~I), we can infer ~G. Next, we will use the Disjunctive Syllogism (DS) rule, which states that if we have a disjunction of the form "A or B" and we know that A is false, then we can infer B. Applying DS to premise 3 (H v ~E) and the derived ~E, we can infer H. Finally, we will use the Simplification (SIMP) rule, which allows us to extract a single component from a conjunction. Applying SIMP to premise 1 (\sim E $\supset \sim$ E), we can infer E $\supset \sim$ E. Therefore, the immediate consequence that can be derived from the given premises is 'E $\supset \sim F'$. The correct answer is A. Answer Option 5: To find the immediate consequence in propositional logic (PL) of the given premises, we can use the method of proof by contradiction. We assume the negation of the proposition we want to prove and derive a contradiction from the premises. Let's consider the proposition 'E \supset F'. Assume its negation: ~(E ⊃ F) Using the rule of implication, we can rewrite $\mathord{\sim}(E\supset F)$ as: Now let's examine the given premises: 1) ~F ⊃ ~F 2) G ⊃ F 3) H v ~E 4) H ⊃ I 5) ~I From premise 1, we can infer ~F (modus ponens). From premise 2, $G \supset F$, and $\sim F$ (inferred from premise 1), we can infer G (modus ponens). From premise 3, $H \lor \sim E$, and $\sim E$ (inferred from premise 5 $\sim I$ and premise 4 $H \supset I$), we can infer H (disjunctive syllogism). Now, we have the following information: ~F, G, H. But we cannot infer E from the given premises. Therefore, we cannot derive the proposition 'E \supset F' as an immediate consequence. The correct answer is: A. $\mbox{'E}\supset\mbox{F'}$ ********** Answer Option 6: Let's analyze the given premises and try to deduce an immediate consequence. Premise 1: ~E ⊃ ~F This premise states that if E is false, then F must also be false. This premise states that if G is true, then F must also be true. Premise 3: H v ~E This premise states that either H is true or E is false. Premise 4: H ⊃ I This premise states that if H is true, then I must also be true. Premise 5: ~I This premise states that I is false. Now, let's consider the options:

Option A: ${}^{{}^{{}}}\!\!E\supset F'$

There is no direct relation between E and F in the given premises. Therefore, this option is not an immediate consequence. Option B: ${}^{t}F\supset G'$ There is no direct relation between F and G in the given premises. Therefore, this option is not an immediate consequence. Option C: 'H ⊃ ~E' This option matches the structure of the premises. From premise 3 (H v ~E), we can deduce ~E \supset H. By contrapositive, this can be written as H \supset ~E. Therefore, this option is an immediate consequence. Option D: '~H' There is no direct relation between ~H and the given premises. Therefore, this option is not an immediate consequence. Final Answer: C. 'H ⊃ ~E' ********** Answer Option 7: To determine the immediate consequence of the given premises, we can use the method of logical deduction. Premises: 1. ~E ⊃ ~F 2. G ⊃ F 3. H ∨ ~E 4. H ⊃ I 5. ~I To find an immediate consequence, we try to derive a statement directly from the premises without any additional steps or assumptions. From premise 2 (G \supset F) and premise 5 (\sim I), we can use modus tollens to infer \sim G. (If G \supset F and \sim I, then \sim F, and from there \sim G.) From premise 3 (H \vee ~E) and premise 4 (H \supset I), we can use disjunctive syllogism to infer ~E \supset I. (If H \vee ~E and H \supset I, then ~E \supset I.) Now, from premise 1 (-E \supset -F) and -E \supset I (derived from premises 3 and 4), we can use hypothetical syllogism to infer -E \supset -F. (If -E \supset -F and -E \supset I, then -E \supset -F.) Therefore, the immediate consequence in PL of the given premises is $\sim E \supset \sim F$. Final Answer: \sim E $\supset \sim$ F (Option D) ******** Answer Option 8: To determine the immediate (one-step) consequence of the given premises, we can use the rules of propositional logic. Premise 1: ~E ⊃ ~F Premise 2: G ⊃ F Premise 3: H ∨ ~E Premise 4: H ⊃ I Premise 5: ~I Let's analyze the answer options: Option A: $'E \supset F'$ There is no direct implication or inference between 'E' and 'F' based on the given premises. Therefore, option A is not an immediate consequence. Option B: ${}^{t}F \supset G'$ There is no direct implication or inference between 'F' and 'G' based on the given premises. Therefore, option B is not an immediate consequence. Option C: 'H ⊃ ~E' We can see that premise 3 states 'H \lor ~E', and premise 4 states 'H \supset I'. By combining these two premises, we can infer that '-E \supset ~H'. However, option C states 'H \supset ~E', which is not the immediate consequence. Therefore, option C is not an immediate consequence. Option D: '~H' There is no direct implication or inference that leads to '~H' based on the given premises. Therefore, option D is not an immediate consequence. After analyzing all the options, we can conclude that none of the given options is an immediate consequence of the premises. Thus, the correct answer is None (or there is no immediate consequence). Answer Option 9: To find an immediate consequence (one-step) in propositional logic (PL) of the given premises, we can use the method of resolution. Let's analyze the premises: 1. ~E ⊃ ~F 2. G ⊃ F 3. H ∨ ~E 4. H ⊃ I 5. ~I We can start by assuming the negation of the proposition we want to derive and see if we can derive a contradiction. Let's assume the negation of the proposition 'E > F', which is '-(E > F)'.

Using the implication rule of inference, we can rewrite '~(E \supset F)' as 'E \wedge ~F'.

Now, let's try to prove 'E $_{\wedge}$ ~F' using the premises:

- From premise 1, '~E \supset ~F', and by contrapositive, we can infer 'F \supset E'.
- From premise 2, 'G \supset F'.
- By applying the transitive property, we have 'G \supset E'.
- From premise 3, 'H v ~E'.
- We consider two cases:
- 1. If 'H' is true, then we can infer 'H \supset E' by premise 4.
- 2. If '~E' is true, then we have 'E \wedge ~F' as desired.

Since we can derive 'E $_{\wedge}$ ~F' using the premises, the immediate consequence is the proposition 'E $_{\supset}$ F'.

Final Answer: A. ${}^{{}^{{}}}\!\!E\supset F'$

Question 10:

'Which of the given formulas of PL is the best symbolization of the following sentence?

Everybody loves Raymond, or not.', A. 'L', B. '~L', C. 'L • ~L', D. 'L \vee ~L'

Answer Option 1:

The best symbolization of the sentence "Everybody loves Raymond, or not." would be D. 'L v ~L'. This captures the meaning that either 'L' (Everybody loves Raymond) is true or '~L' (Not everybody loves Raymond) is true.

Answer Option 2:

The best symbolization of the sentence "Everybody loves Raymond, or not" would be "L v ~L". This formula states that either "L" (Everybody loves Raymond) is true, or "~L" (Not everybody loves Raymond) is true. Final Answer: D

Answer Option 3:

The best symbolization of the given sentence, "Everybody loves Raymond, or not" would be: D. 'L \lor ~L'

This symbolization represents the statement that either everybody loves Raymond (L), or it is not the case that everybody loves Raymond (~L), which covers all possible scenarios. Final Answer: D

Answer Option 4:

The best symbolization of the given sentence "Everybody loves Raymond, or not" would be option D. 'L v ~L'.

This formula represents the logical statement "L or not L", which means either somebody loves Raymond (L) or somebody does not love Raymond (-L). This covers all possible scenarios and accurately represents the given sentence. Therefore, the correct answer is option D.

Answer Option 5

The best symbolization of the sentence "Everybody loves Raymond, or not" would be option D: "L v ~L". This formula represents the logical statement that either everybody (represented by L) loves Raymond, or it is not the case that everybody loves Raymond. This captures the idea that either everyone loves Raymond or there is at least one person who does not love Raymond. Final Answer: D

Answer Option 6:

The best symbolization for the sentence "Everybody loves Raymond, or not" is option D: 'L \vee ~L'.

This is because the sentence states that either everybody loves Raymond (L) or not (-L). The symbol v represents the logical OR operation, which indicates that at least one of the two options is true.

Therefore, option D correctly symbolizes the given sentence.

Final Answer: D

Answer Option 7:

The best symbolization of the sentence "Everybody loves Raymond, or not" would be option D, "L \vee ~L".

This formula represents the logical statement that either "L" (Everybody loves Raymond) is true or "~L" (not everybody loves Raymond) is true.

Option A, "L", represents the statement that everybody loves Raymond, but it does not consider the possibility that not everybody loves Raymond.

Option B, "~L", represents the statement that nobody loves Raymond, which is not equivalent to the given sentence.

Option C, "L - -L", represents the contradictory statement that everybody loves Raymond and nobody loves Raymond at the same time, which does not accurately represent the original sentence.

Therefore, option D, "L \vee ~L", is the best symbolization of the given sentence. Final Answer: D

Answer Option 8:

The best symbolization of the sentence "Everybody loves Raymond, or not" is option D, "L v ~L". This formula represents the logical statement that either everyone loves Raymond (L) or it is not the case that everyone loves Raymond (-L), covering all possibilities. Option A, "L", only represents the statement that everyone loves Raymond but does not account for the possibility that it might not be true. Option B, "-L", represents the statement that it is not the case that everyone loves Raymond, but it does not cover the possibility that everyone does love Raymond. Option C, "L • -L", represents a contradiction, as it asserts both that everyone loves Raymond and does not love Raymond simultaneously, which is not what the original sentence conveys. Therefore, option D, "L v ~L", is the best symbolization. Final Answer: D

Answer Option 9:

The best symbolization of the sentence "Everybody loves Raymond, or not" is D. 'L \vee ~L'.

This symbolization captures the meaning of the sentence accurately. It states that either L (stands for "Everybody loves Raymond") is true or ~L (stands for "not everybody loves Raymond") is true. Since it covers both possibilities, it represents the given sentence correctly.

Option A, 'L', only represents the "Everybody loves Raymond" part of the sentence but does not consider the possibility of not loving Raymond.

Option B, '-L', only represents the "not everybody loves Raymond" part of the sentence but does not consider the possibility of loving Raymond.

Option C, 'L • ~L', represents the contradictory statement that "Everybody loves Raymond" and "not everybody loves Raymond" are both true, which is not consistent with the original sentence.

Therefore, the best symbolization of the sentence is D. 'L \vee ~L'.

Question 11:

'Select the best translation into predicate logic. All kings are luckier than all paupers. (Kx: x is a king; Px: x is a pauper; Lxy: x is luckier than y)', A. ' $(\forall x)[Kx \supset (\forall y)(Py \supset Lxy)]'$, B. ' $(\forall x)[Kx \supset (\forall y)(Lxy \supset Py)]'$, C. ' $(\forall x)[(Kx \bullet Py) \supset Lyx]'$, D. ' $(\forall x)[(Kx \bullet (\forall y)(Py)) \supset Lxy]'$

Answer Option 1:

The best translation into predicate logic for the given statement "All kings are luckier than all paupers" is: A. $(\forall x)[Kx \cup (\forall y)(Py \cup Lxy)]$ " In this translation, $(\forall x)$ represents "for all x", Kx represents "x is a king", $(\forall y)$ represents "for all y", Py represents "y is a pauper", and Lxy represents "x is luckier than y". The translation states that for all x (kings), if x is a king (Kx), then for all y (paupers), if y is a pauper (Py), then y is luckier than y (Lxy). This accurately represents the given statement. Final Answer: A

Answer Option 2:

The best translation into predicate logic for the statement "All kings are luckier than all paupers" is:

 $(\forall x)[Kx \to (\forall y)(Py \to Lxy)]$

This translation correctly expresses that for all x (any individual), if x is a king (Kx), then for all y (any individual), if y is a pauper (Py), then x is luckier than y (Lxy).

Hence, the correct answer is A. $(\forall x)[Kx \rightarrow (\forall y)(Py \rightarrow Lxy)].$

Answer Option 3:

To translate the statement "All kings are luckier than all paupers" into predicate logic, we can use the following symbols:

Kx: x is a king

Px: x is a pauper

Lxy: x is luckier than y

The correct translation would be:

 $(\forall x)[Kx\supset (\forall y)(Py\supset Lxy)]$

Let's break down the translation:

(∀x) - For all x

 $[Kx\supset (\forall y)(Py\supset Lxy)] \text{ - If } x \text{ is a king, then for all } y, \text{ if } y \text{ is a pauper, then } x \text{ is luckier than } y.$

This translation accurately represents the statement "All kings are luckier than all paupers."

The correct answer is A. $(\forall x)[Kx\supset (\forall y)(Py\supset Lxy)]$

Answer Option 4:

 $The best translation into predicate logic for the statement "All kings are luckier than all paupers" is option A: '(\forall x)[Kx \supset (\forall y)(Py \supset Lxy)]'.$

Let's break down the translation:

- Kx: x is a king
- Px: x is a pauper
- Lxy: x is luckier than y

The statement "All kings are luckier than all paupers" can be translated as "For all x, if x is a king, then for all y, if y is a pauper, then x is luckier than y."

This translates to '($\forall x$)[Kx \supset ($\forall y$)(Py \supset Lxy)]'.

Hence, the correct answer is A.

Answer Option 5:

The best translation into predicate logic for the statement "All kings are luckier than all paupers" would be:

 $(\forall x)[\mathsf{Kx} \supset (\forall y)(\mathsf{Py} \supset \mathsf{Lxy})]$

Explanation:

- Kx: x is a king
- Px: x is a pauper
- Lxy: x is luckier than y

The translation "(∀x)[Kx > (∀y)[Py > Lxy)]" states that for all x, if x is a king, then for all y, if y is a pauper, then x is luckier than y. This translation accurately represents the given statement.

 $\label{eq:control} \text{Option A: '}(\forall x)[Kx\supset (\forall y)(Py\supset Lxy)]' \text{ is the correct translation and represents the statement accurately. Final Answer: A }$

Answer Option 6:

The best translation into predicate logic for the given statement is: (A) $(\forall x)[Kx\supset (\forall y)(Py\supset Lxy)]$

Explanation:

Let's break down the original statement:

All kings are luckier than all paupers.

To represent "x is a king" using predicate logic, we can use Kx. To represent "x is a pauper" using predicate logic, we can use Px. To represent "x is luckier than y" using predicate logic, we can use Lxy. Now, let's translate the statement using predicate logic: "All kings are luckier than all paupers." For every x, if x is a king (Kx), then for every y, if y is a pauper (Py), then x is luckier than y (Lxy). This can be written as: $(\forall x)[\mathsf{K} x \supset (\forall y)(\mathsf{P} y \supset \mathsf{L} xy)]$ Therefore, the best translation into predicate logic is option A. Answer Option 7: The best translation into predicate logic for the given statement "All kings are luckier than all paupers" would be: $(\forall x)[Kx \supset (\forall y)(Py \supset Lxy)]$ - (∀x) denotes "for all x" - Kx represents "x is a king" - Px represents "x is a pauper" - Lxy represents "x is luckier than y" $The translation states that for every \ x, if \ x is a king \ (Kx), then for every \ y, if \ y is a pauper \ (Py), then \ x is luckier than \ y \ (Lxy).$ Option A: $(\forall x)[Kx\supset (\forall y)(Py\supset Lxy)]$ - This is the correct translation. Option B: $(\forall x)[Kx\supset (\forall y)(Lxy\supset Py)]$ - This does not accurately capture the original statement. Option C: $(\forall x)(\forall y)[(Kx \cdot Py) \supset Lyx]$ - This does not accurately capture the original statement. $Option \ D: \ (\forall x)[(Kx \bullet (\forall y)(Py)) \supset Lxy] \ - \ This \ does \ not \ accurately \ capture \ the \ original \ statement.$ Final Answer: A Answer Option 8: The best translation into predicate logic for the statement "All kings are luckier than all paupers" would be: $(\forall x)[\mathsf{K}\mathsf{x}\supset (\forall y)(\mathsf{P}\mathsf{y}\supset \mathsf{L}\mathsf{x}\mathsf{y})]$ This translates to "For all x, if x is a king, then for all y, if y is a pauper, then x is luckier than y." Option A: $(\forall x)[Kx \supset (\forall y)(Py \supset Lxy)]$ - Correct translation. Option B: $(\forall x)[Kx\supset (\forall y)(Lxy\supset Py)]$ - Incorrect translation. This would mean "For all x, if x is a king, then for all y, if x is luckier than y, then y is a pauper." This does not capture the meaning of the original Option C: $(\forall x)(\forall y)[(Kx \cdot Py) = Lyx]$ - Incorrect translation. This would mean "For all x and y, if x is a king and y is a pauper, then y is luckier than x." This reverses the intended meaning of the statement. Option D: $(\forall x)[(Kx \cdot (\forall y)(Py)) \supset Lxy]$ - Incorrect translation. This would mean "For all x, if x is a king and for all y, y is a pauper, then x is luckier than y." This incorrectly assumes that all paupers are luckier than kings. Final Answer: A Answer Option 9: The best translation into predicate logic for the given statement "All kings are luckier than all paupers" is: $(\forall x)[Kx\supset (\forall y)(Py\supset Lxy)]$ Option A: $(\forall x)[Kx \supset (\forall y)(Py \supset Lxy)]$ - This represents the correct translation, stating that for all x (which are kings), if x implies that for all y (which are paupers), then Py implies Lxy. This accurately captures the statement that "All kings are luckier than all paupers." Option B: $(\forall x)[Kx\supset (\forall y)(Lxy\supset Py)]$ - This translation swaps the order of Lxy and Py, which changes the meaning of the statement. It states that for all x (which are kings), if x implies that for all y, Lxy implies Py. This is not an accurate representation of the given statement. Option C: $(\forall x)(\forall y)[(Kx \cdot Py) - Lyx]$ - This translation introduces an unnecessary conjunction $(Kx \cdot Py)$ that is not present in the original statement. It states that for all x and y, if Kx and Py are both true, then Lyx is true. This does not capture the intended meaning. $Option \ D: \ (\forall x)[(Kx \bullet (\forall y)(Py)) \supset Lxy] - This translation incorrectly places the universal quantifier \ (\forall y)(Py) within the parentheses with Kx. It states that for all x, if both Kx and \ (\forall y)(Py) are true, then Lxy is true. \\$ This does not accurately represent the original statement. Final Answer: A

Question 12:

Construct a complete truth table for the following pairs of propositions. Then, using the truth tables, determine whether the statements are logically equivalent or contradictory. If neither, determine whether they are consistent or inconsistent. Justify your answers.

 $(-M\supset -N)\lor (O\equiv N) \text{ and } (-M\cdot N)\cdot [(-O\lor -N)]", \text{ A. `Logically equivalent', B. `Contradictory', C. `Neither logically equivalent nor contradictory, but consistent', D. 'Inconsistent', D

Answer Option 1:

To construct a complete truth table for the given pair of propositions, we need to consider all possible combinations of truth values for the variables M, N, and O.

Let's break down the given propositions:

1. $(\sim M \supset \sim N) \lor (O \equiv N)$

2. $(\sim M \cdot N) \cdot [(\sim O \vee \sim N) \cdot (O \vee N)]$

Complete Truth Table:

 $M\mid N\mid O\mid \sim M\mid \sim N\mid (\sim M\supset \sim N)\vee (O\equiv N)\mid (\sim M\cdot N)\cdot [(\sim O\vee \sim N)\cdot (O\vee N)]$

T|T|T|F|F|T|F

T|T|F|F|F|F|F

.

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T|F|F|F|T|T|F

F|T|T|T|F|T|F

F|T|F|T|F|T|F

F|F|T|T|T|F

F|F|F|T|T|F

Based on the truth table, we can observe that the two propositions have different truth values for different combinations of truth values of M, N, and O. Thus, the statements are not logically equivalent.

Since they do not contradict each other for any row in the truth table, we can conclude that they are consistent. Therefore, the answer is C. 'Neither logically equivalent nor contradictory, but consistent'.

.....

Answer Option 2:

To construct a complete truth table for the given pairs of propositions, we need to consider the truth values of the individual propositions M, N, and O.

Inputs: M, N, C

Outputs: (~M \supset ~N) \vee (O \equiv N) and (~M \cdot N) \cdot [(~O \vee ~N) \cdot (O \vee N)]

Using the truth tables, we can evaluate the truth values of the propositions for all possible combinations of inputs:

 $M\mid N\mid O\mid\mid (\sim\!M\supset \sim\!N)\vee (O\equiv N)\mid (\sim\!M\cdot N)\cdot [(\sim\!O\vee \sim\!N)\cdot (O\vee N)]$

TITITIFIE

TITIFIFIF

T|F|T||T|F

TIFIFIITIF

.

F|T|T||T|F

F|T|F||F|F

F|F|T||T|F F|F|F||T|F

From the truth table, we observe that the two propositions have different outputs for at least one input combination (M = T, N = F, O = T). Therefore, the statements are not logically equivalent.

Additionally, there is no row in the truth table where both propositions have the same truth value. Hence, the statements are not consistent.

Final Answer: D. 'Inconsistent'

Answer Option 3:

To construct a complete truth table for the given propositions, we have the following variables:

M, N, O

We will evaluate the truth values of both propositions:

1. $(\sim M \supset \sim N) \lor (O \equiv N)$

2. $(\sim M \cdot N) \cdot [(\sim O \lor \sim N) \cdot (O \lor N)]$

Constructing the truth table:

 $M \ N \ O \ \sim M \ \sim N \ \sim N \ O \ \equiv N \ (\sim M \ \supset \ \sim N) \ \lor \ (O \ \equiv N) \ \sim M \ \cdot N \ \sim O \ \lor \ \sim N \ O \ \lor \ N \ [(\sim O \ \lor \ \sim N)) \ \cdot \ (O \ \lor \ N)] \ (\sim M \ \cdot \ N) \ \cdot \ [(\sim O \ \lor \ \sim N) \ \cdot \ (O \ \lor \ N)] \ (\sim M \ \cdot \ N) \ (\sim M \ \cdot \$

TTTFFTTTTTTT

TTFFFTFTTTFFF

Looking at the final column of the truth table, we can see that both propositions have the same truth values for all possible combinations of truth values for M, N, and O.

Since the truth values of both propositions are the same for all possible combinations of truth values, the statements are logically equivalent.

Therefore, the answer is A. 'Logically equivalent'.

Answer Option 4:

To construct the complete truth table for the given propositions, we need to consider all possible combinations of truth values for the propositions M, N, and O.

Let's start by assigning truth values to M, N, and O:

 $M \mid N \mid O \mid \neg M \mid \neg N \mid \neg M \supset \neg N \mid O \equiv N \mid (\neg M \supset \neg N) \lor (O \equiv N) \mid \neg M \cdot N \mid (\neg O \lor \neg N) \cdot (O \lor N) \mid (\neg M \cdot N) \cdot [(\neg O \lor \neg N) \cdot (O \lor N)]$

Based on the truth table, we can see that the values of the propositions $(-M \supset -N) \lor (O \equiv N)$ and $(-M \lor N) \cdot [(-O \lor -N) \cdot (O \lor N)]$ are the same for all possible truth value combinations. Therefore, the statements are logically equivalent.

Final Answer: A. Logically equivalent

Answer Option 5:

To construct a complete truth table for the given pairs of propositions, we need to consider all possible combinations of truth values for the variables M, N, and O.

Let's represent the propositions as follows:

(~M ⊃ ~N) ∨ (O ≡ N) as P

 $({\sim}M\cdot N)\cdot[({\sim}O\vee{\sim}N)\cdot(O\vee N)]$ as Q

Using the truth tables, we can evaluate the truth values for P and Q as follows:

 $M\mid N\mid O\mid \sim M\mid \sim N\mid \sim M\supset \sim N\mid O\equiv N\mid P$

```
F|F|T|T|T|F|T|T|T|F|F
  F|F|F|T|T|T|F|F|F|F
  By comparing the truth values of P and Q in the truth table, we can see that they are not logically equivalent, as there are rows where the truth values of P and Q differ.
 However, they are not contradictory either, as there are rows where both {\sf P} and {\sf Q} are true.
  Therefore, the given statements (\neg M \supset \neg N) \lor (O \equiv N) and (\neg M \cdot N) \cdot [(\neg O \lor \neg N)] are neither logically equivalent nor contradictory, but consistent.
  Final Answer: C. 'Neither logically equivalent nor contradictory, but consistent'
  Answer Option 6:
  To construct a complete truth table, we need to consider all possible combinations of truth values for the propositions M, N, and O.
  Let's start by listing out all the possible truth values for M, N, and O:
 MNO
  ттт
  TTF
  TFT
  TFF
  FTT
  FTF
  FFT
  FFF
  Next, we evaluate the truth values for the given propositions using the truth values from the table:
  (\sim M \supset \sim N) \lor (O \equiv N)
\mid M\mid N\mid O\mid \sim M\mid \sim N\mid (\sim M\supset \sim N)\mid (O\equiv N)\mid (\sim M\supset \sim N)\vee (O\equiv N)\mid
 |\mathsf{T}|\mathsf{T}|\mathsf{T}|\mathsf{F}|\mathsf{F}|\mathsf{T}|\mathsf{T}|\mathsf{T}|
 |T|T|F|F|F|T|F|T|
|T|F|T|F|T|T|F|T|
|T|F|F|F|T|T|T|T|
 |F|T|T|T|F|F|T|T|
|F|T|F|T|F|F|F|F|
|F|F|T|T|T|T|F|T|
|F|F|F|T|T|T|T|T|
  (\sim M \cdot N) \cdot [(\sim O \lor \sim N) \cdot (O \lor N)]
\mid M \mid N \mid O \mid \neg M \mid \neg O \mid \neg N \mid (\neg M \cdot N) \mid (\neg O \vee \neg N) \mid (O \vee N) \mid [(\neg O \vee \neg N) \cdot (O \vee N)] \mid (\neg M \cdot N) \cdot [(\neg O \vee \neg N) \cdot (O \vee N)] \mid (O \vee N) \mid (O
 |\mathsf{T}|\mathsf{T}|\mathsf{T}|\mathsf{F}|\mathsf{F}|\mathsf{F}|\mathsf{F}|\mathsf{T}|\mathsf{T}|\mathsf{T}|\mathsf{F}|
 |T|T|F|F|T|F|F|T|F|F|F|
|T|F|T|F|F|T|F|T|T|F|
|T|F|F|F|T|T|F|T|F|F|F
```

Based on the truth tables, we can see that the truth values for the two propositions are not the same for all possible combinations of truth values. Therefore, the statements are not logically equivalent.

Since there are some combinations of truth values where the statements are both true, they are not contradictory either. Therefore, the statements are consistent.

Final Answer: C. 'Neither logically equivalent nor contradictory, but consistent'

Answer Option 7:

To construct a complete truth table, we need to consider all possible combinations of truth values for the propositions involved. Let's break down the given expressions and construct the truth table for each:

```
(\text{~M}\supset\text{~N})\vee(O\equiv N):
```

```
M N O ~M ~N (~M \supset ~N) (O \equiv N) [(~M \supset ~N) \lor (O \equiv N)]
```

TTTFFTTT

TTFFFTFF

```
TFTFTTFT
  TEFETTTT
  FTTTFFFT
 FTFTFFTT
  FFTTTTT
 FFFTTTTF
 (\sim M \cdot N) \cdot [(\sim O \vee \sim N) \cdot (O \vee N)]:
 M\ N\ O\ \sim M\ \sim N\ (\sim M\ \cdot\ N)\ (\sim O\ \vee\ \sim N)\ (O\ \vee\ N)\ [(\sim O\ \vee\ \sim N)\ \cdot\ (O\ \vee\ N)]\ (\sim M\ \cdot\ N)\ \cdot\ [(\sim O\ \vee\ \sim N)\ \cdot\ (O\ \vee\ N)]
  TTTFFTFTT
 TTFFFFFFFF
  TFTFTFTTTFF
 TFFFTFFTFF
  FTTTFFTFFF
  FTFTFFFTFF
  FFTTTFTTFFF
  FFFTTFTFFF
 Now, let's \ compare \ the \ truth \ tables \ to \ determine \ whether \ the \ statements \ are \ logically \ equivalent, \ contradictory, \ or \ neither:
  The truth table for (-M \supset -N) \lor (O \equiv N) and (-M \lor N) \lor [(-O \lor -N) \lor (O \lor N)] shows that the two expressions have the same truth values for all possible combinations of truth values. Therefore, the statements
 are logically equivalent. Answer: A. 'Logically equivalent'
  Answer Option 8:
  To determine whether the statements are logically equivalent, contradictory, consistent, or inconsistent, we need to construct a complete truth table for both statements and compare the values of their
  respective columns.
  Statement 1: (\sim M \supset \sim N) \lor (O \equiv N)
 Statement 2: (~M \cdot N) \cdot [(~O \vee ~N) \cdot (O \vee N)]
 Let's construct the truth table for both statements:
\mid M \mid N \mid O \mid \neg M \supset \neg N \mid O \equiv N \mid (\neg M \supset \neg N) \lor (O \equiv N) \mid \neg M \cdot N \mid (\neg O \lor \neg N) \cdot (O \lor N) \mid (\neg M \cdot N) \cdot [(\neg O \lor \neg N) \cdot (O \lor N)] \mid (\neg M \cdot N) \cdot [(\neg O \lor \neg N) \cdot (O \lor N)] \mid (\neg M \lor N) \cdot [(\neg O \lor \neg N) \cdot (O \lor N)] \mid (\neg M \lor N) \cdot [(\neg O \lor \neg N) \cdot (O \lor N)] \mid (\neg M \lor N) \cdot [(\neg O \lor \neg N) \cdot (O \lor N)] \mid (\neg M \lor N) \cdot [(\neg O \lor \neg N) \cdot (O \lor N)] \mid (\neg M \lor N) \cdot [(\neg O \lor \neg N) \cdot (O \lor N)] \mid (\neg M \lor N) \cdot [(\neg O \lor \neg N) \cdot (O \lor N)] \mid (\neg M \lor N) \cdot [(\neg O \lor \neg N) \cdot (O \lor N)] \mid (\neg M \lor N) \cdot [(\neg O \lor \neg N) \cdot (O \lor N)] \mid (\neg M \lor N) \cdot [(\neg O \lor \neg N) \cdot (O \lor N)] \mid (\neg M \lor N) \cdot [(\neg O \lor \neg N) \cdot (O \lor N)] \mid (\neg M \lor N) \cdot [(\neg O \lor \neg N) \cdot (O \lor N)] \mid (\neg M \lor N) \cdot [(\neg O \lor \neg N) \cdot (O \lor N)] \mid (\neg M \lor N) \cdot [(\neg O \lor \neg N) \cdot (O \lor N)] \mid (\neg M \lor N) \cdot [(\neg O \lor \neg N) \cdot (O \lor N)] \mid (\neg M \lor N) \cdot [(\neg O \lor \neg N) \cdot (O \lor N)] \mid (\neg M \lor N) \cdot [(\neg O \lor \neg N) \cdot (O \lor N)] \mid (\neg M \lor N) \cdot [(\neg O \lor \neg N) \cdot (O \lor N)] \mid (\neg M \lor N) \cdot [(\neg O \lor \neg N) \cdot (O \lor N)] \mid (\neg M \lor N) \cdot [(\neg O \lor \neg N) \cdot (O \lor N)] \mid (\neg M \lor N) \cdot [(\neg O \lor \neg N) \cdot (O \lor N)] \mid (\neg M \lor N) \cdot [(\neg O \lor \neg N) \cdot (O \lor N)] \mid (\neg M \lor N) \cdot [(\neg O \lor \neg N) \cdot (O \lor N)] \mid (\neg M \lor N) \cdot [(\neg O \lor \neg N) \cdot (O \lor N)] \mid (\neg M \lor N) \cdot [(\neg O \lor \neg N) \cdot (O \lor N)] \mid (\neg M \lor N) \cdot [(\neg O \lor \neg N) \cdot (O \lor N)] \mid (\neg M \lor N) \cdot [(\neg O \lor \neg N) \cdot (O \lor N)] \mid (\neg M \lor N) \cdot [(\neg O \lor \neg N) \cdot (O \lor N)] \mid (\neg M \lor N) \cdot [(\neg O \lor \neg N) \cdot (O \lor N)] \mid (\neg M \lor N) \cdot [(\neg O \lor \neg N) \cdot (O \lor N)] \mid (\neg M \lor N) \cdot [(\neg O \lor \neg N) \cdot (O \lor \neg N)] \mid (\neg M \lor N) \cdot [(\neg O \lor \neg N) \cdot (O \lor \neg N)] \mid (\neg M \lor N) \cdot [(\neg O \lor \neg N) \cdot (O \lor \neg N)] \mid (\neg M \lor N) \cdot [(\neg O \lor \neg N) \cdot (O \lor \neg N)] \mid (\neg M \lor \neg N) \cdot [(\neg O \lor \neg N) \cdot (O \lor \neg N)] \mid (\neg M \lor \neg N) \cdot [(\neg O \lor \neg N) \cdot (O \lor \neg N)] \mid (\neg M \lor \neg N) \cdot [(\neg O \lor \neg N) \cdot (O \lor \neg N)] \mid (\neg M \lor \neg N) \cdot [(\neg O \lor \neg N) \cdot (O \lor \neg N)] \mid (\neg M \lor \neg N) \cdot [(\neg O \lor \neg N) \cdot (O \lor \neg N)] \mid (\neg M \lor \neg N) \cdot [(\neg O \lor \neg N) \cdot (O \lor \neg N)] \mid (\neg M \lor \neg N) \cdot [(\neg O \lor \neg N) \cdot (O \lor \neg N)] \mid (\neg M \lor \neg N) \cdot [(\neg O \lor \neg N) \cdot (O \lor \neg N)] \mid (\neg M \lor \neg N) \cdot [(\neg O \lor \neg N) \cdot (O \lor \neg N)] \mid (\neg M \lor \neg N) \cdot [(\neg O \lor \neg N) \cdot (O \lor \neg N)] \mid (\neg O \lor \neg N) \cdot [(\neg O \lor \neg N) \cdot (O \lor \neg N)] \mid (\neg O \lor \neg N) \cdot [(\neg O \lor \neg N) \cdot (O \lor \neg N)] \mid (\neg O \lor \neg N) \cdot [(\neg O \lor \neg N) \cdot (O \lor \neg N)] \mid (\neg O \lor \neg N) \cdot [(\neg O \lor \neg N) \cdot (O \lor \neg N)] \mid (\neg O \lor \neg N) \cdot [(\neg O \lor \neg N)
 |---|---|--------|-------|------
 |T|T|T|F|T|T|T|T|T
```

From the truth table, we can see that the columns for both statements have the same values. Therefore, the statements are logically equivalent.

Final Answer: A. 'Logically equivalent'

Answer Option 9:

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To determine whether the statements are logically equivalent, contradictory, consistent, or inconsistent, we need to construct a complete truth table for both propositions.

Proposition 1: $(\sim M \supset \sim N) \lor (O \equiv N)$

Proposition 2: $(\sim M \cdot N) \cdot [(\sim O \vee \sim N) \cdot (O \vee N)]$

Let's construct the truth table for Proposition 1:

M N O ~M ~N (~M \supset ~N) (O \equiv N) (~M \supset ~N) \vee (O \equiv N)

TTTFFTTT

TTFFFTTT

TFFFTTTT

TFFFTTTT

FTTTFFFF

FTFTFFFTT

FFFTTTTT

Now, let's construct the truth table for Proposition 2:

 $\text{M N O } \sim \text{M } \sim \text{N } \sim \text{O } (\sim \text{O } \vee \sim \text{N}) \text{ O } \vee \text{N } (\sim \text{M} \cdot \text{N}) \left[(\sim \text{O } \vee \sim \text{N}) \cdot (\text{O } \vee \text{N}) \right] (\sim \text{M} \cdot \text{N}) \cdot \left[(\sim \text{O } \vee \sim \text{N}) \cdot (\text{O } \vee \text{N}) \right]$

TTTFFFTTTTT

TTFFFTTTTT

TETETETTEE

TFFFTTTTFFF

FTTTFFFTFF

FTFTFTTTFFF

FFTTTFTFFF

FFFTTTTFFFF

By comparing the truth tables, we can see that the two propositions have the same truth values for all combinations of M, N, and O. Therefore, the statements are logically equivalent.

Final Answer: A. 'Logically equivalent'

```
'Which of the following propositions is an immediate (one-step) consequence in PL of the given premises?
E⊃~F
\simF ⊃ G
~G', A. 'E \supset F', B. 'F \supset G', C. 'G \supset E', D. '~~F'
To find the immediate (one-step) consequence in propositional logic (PL) of the given premises, we can use the method of modus ponens.
The given premises are:
1. E ⊃ ~F
2. ~F ⊃ G
3. ~G
Modus ponens states that if we have a conditional statement of the form ^{\prime}A \supset B^{\prime} and we also have ^{\prime}A^{\prime}, then we can infer ^{\prime}B^{\prime}.
From the premises, we can see that we have \simF as a premise and also \simF \supset G as a premise. By applying modus ponens, we can infer G.
Therefore, the immediate consequence is 'G'.
Final Answer: C. 'G \supset E'
*******
Answer Option 2:
To determine the immediate (one-step) consequence in PL of the given premises, we can use the process of Modus Ponens.
The given premises are:
1. E ⊃ ~F
2. ~F ⊃ G
3. ~G
From premise 1 (E \supset ~F) and premise 3 (~G), we can infer ~F using Modus Ponens.
From premise 2 ({}^{\sim}F\supset G) and the derived {}^{\sim}F, we can infer G using Modus Ponens.
Therefore, the immediate (one-step) consequence in PL of the given premises is 'G'.
Final Answer: C. 'G \supset E'
**********
To determine the immediate (one-step) consequence in propositional logic (PL) of the given premises, we need to apply the modus ponens rule of inference.
The modus ponens rule states that if we have two premises of the form A \to B and A, we can conclude B.
Given premises:
1. E ⊃ ~F
2. ~F ⊃ G
3. ~G
We can apply modus ponens as follows:
Using premise 1 and premise 3, we have ~F (from premise 1) and ~G (from premise 3).
Using premise 2, since we have ~F (premise 1) and ~F \supset G (premise 2), we can conclude G.
Therefore, the immediate consequence in PL of the given premises is 'G'.
Final Answer: None of the options provided (A, B, C, D)
Answer Option 4:
To find the immediate consequence of the given premises, we can use the method of hypothetical syllogism.
Premise 1: E ⊃ ~F
Premise 2: ~F ⊃ G
Premise 3: ~G
From Premise 1 and Premise 2, using the hypothetical syllogism, we can conclude:
However, none of the given answer options match this result. Therefore, there is no immediate consequence in PL of the given premises among the options provided.
Final Answer: None (No option is correct)
Answer Option 5:
```

Question 13:

To determine the immediate (one-step) consequence of the given premises, we can use the process of Modus Tollens, which states that if we have a conditional statement of the form $P \rightarrow Q$ and the negation
of the consequent, ~Q, then we can conclude the negation of the antecedent, ~P.
Given premises:
1.E⊃~F
2.~F⊃G
3.~G
Using Modus Tollens on premise 2, we can conclude ~~F, which simplifies to F.
Therefore, the immediate consequence in PL of the given premises is 'F'.
Final Answer: D. '~~F'
Answer Option 6:
To determine the immediate consequence of the given premises, we can use the modus ponens rule of inference.
The modus ponens rule states that if we have a conditional statement and its antecedent is true, then we can conclude that the consequent is also true.
In this case, the premises are:
1.E⊃~F
2. ~F ⊃ G
3. ~G (which is equivalent to ~~G)
We can apply the modus ponens rule to the premises 1 and 3, which gives us:
From 1. E \supset ~F and 3. ~G, we can conclude the immediate consequence: E \supset ~G.
Therefore, the correct option is not provided among the given options.

Answer Option 7:
To determine the immediate (one-step) consequence of the given premises in Propositional Logic (PL), we can use the process of Modus Ponens.
The premises are:
1.E⊃~F
2.~F⊃G
3.~G
Modus Ponens states that if we have a premise in the form of "A \supset B" and another premise that is "A," then we can infer "B."
In this case, from premise 1 "E ⊃ ~F" and premise 3 "~G," we can infer "~F" using Modus Ponens.
Therefore, the immediate consequence in PL is "-F."
Among the given options:
A. 'E⊃F'
B. 'F⊃G'
C. 'G ⊃ E'
D. '~~F'
None of the options match the immediate consequence of "~F."
Final Answer: None of the given options

Answer Option 8:
To determine the immediate consequence of the given premises, we can use the rules of implication in propositional logic.
Premises:
1.E⊃~F
2. ~F⊃G
3.~G
Based on these premises, we can apply the rule of Modus Ponens to derive the immediate consequence. Modus Ponens states that if we have a conditional statement P \rightarrow Q and we also have P, then we
can infer Q.
Applying Modus Ponens to premises 1 and 3, we can conclude ~F.
Now, let's examine the answer choices:
A. 'E ⊃ F': This is not a valid immediate consequence based on the given premises. P. 'E ⊃ C': This is not a valid immediate consequence based on the given premises.
B. 'F ⊃ G': This is not a valid immediate consequence based on the given premises.
C. 'G ⊃ E': This is not a valid immediate consequence based on the given premises.
D. '~~F': This is not a valid immediate consequence based on the given premises.
Therefore, none of the answer choices provided correspond to the immediate consequence of the given premises.

No Answer.
Answer Option 9:
To determine the immediate (one-step) consequence in PL of the given premises, we can use the method of Modus Ponens.
The premises are:
1. E⊃~F
2F ⊃ G
3G
We can use Modus Ponens to derive new statements. Modus Ponens states that if we have a conditional statement $P \rightarrow Q$ and we also have P , then we can conclude Q .
Applying Modus Ponens to the premises:
From premise 1 (E \supset ~F) and premise 3 (~G), we can infer ~F.
From premise 2 ($\neg F \supset G$) and the derived $\neg F$, we can infer G .
Therefore, the immediate consequence in PL of the given premises is 'G'.
None of the answer options provided ('E \supset F', 'F \supset G', 'G \supset E', ' \sim -F') match the immediate consequence that can be derived from the premises.
Final Answer: None of the above.

Question 14:

"Use indirect truth tables to determine whether the following argument is valid. If the argument is invalid, choose an option which presents a counterexample. (There may be other counterexamples as well.)

F ⊃ (F ∨ ~G)

F \supset (\sim E \sim \sim G) / \sim (E = F)', A. 'Valid', B. 'Invalid. Counterexample when E, F, and G are true', C. 'Invalid. Counterexample when E and F are true and G is false', D. 'Invalid. Counterexample when E, F, and G are false'

Answer Option 1:

We can use an indirect truth table to determine the validity of the argument.

 $\mathsf{E}\supset (\mathsf{F}\vee \mathsf{\sim}\mathsf{G})\;\mathsf{F}\supset (\mathsf{\sim}\mathsf{E}\vee \mathsf{\sim}\mathsf{G})\;\mathsf{\sim}(\mathsf{E}\equiv\mathsf{F})$

TTTFT

TTFFF

TETTE

TFFTT

FTTTF

FTFFT

FFTTF

FFFTF

In the truth table, we can see that there is no row where both premises (E \supset (F \lor \sim G) and F \supset (\sim E \lor \sim G)) are true and the conclusion (\sim (E = F)) is false. Therefore, there is no counterexample, and the argument is valid.

Final Answer: A. Valid

Answer Option 2:

To determine the validity of the argument, we can construct a truth table for all the propositions involved.

 $\mathsf{E}\;\mathsf{F}\;\mathsf{G}\;\mathsf{E}\supset (\mathsf{F}\vee \mathsf{\sim}\mathsf{G})\;\mathsf{F}\supset (\mathsf{\sim}\mathsf{E}\vee \mathsf{\sim}\mathsf{G})\;\mathsf{\sim}(\mathsf{E}\equiv\mathsf{F})$

TTTTFF

TTFTTFF

TFTFTTT

TFFFTTT

FTTTFTT

FTFTFTT

FFTTTT

FFFTTTT

The argument is valid if and only if the conclusion is always true whenever all the premises are true. In this case, the conclusion is $\neg (E \equiv F)$, which means E is not logically equivalent to F.

From the truth table, we can see that when E and F are both true and G is false (row 3), the premises are true, but the conclusion is false. Therefore, this is a counterexample to the argument.

The correct answer is C. 'Invalid. Counterexample when E and F are true and G is false.'

Answer Option 3:

To determine the validity of the argument, we can use an indirect truth table. Let's assign truth values to the propositions E, F, and G:

EIFIG

TITIT

TITIF

T|F|T

T|F|F

F|T|T

 $F \mid T \mid F$

F|F|T F|F|F

Now, let's evaluate the premises and the conclusion of the argument:

1. $E \supset (F \lor \sim G)$

 $2. \ F \supset (\text{~E} \lor \text{~G})$

~(E ≡ F)

We will look for a row in which both premises are true and the conclusion is false. In the given truth table, we can see that when E is true, F is true, and G is false, both premises are true, but the conclusion is false. Therefore, the argument is invalid. Counterexample: E and F are true, and G is false. Option C: 'Invalid. Counterexample when E and F are true, and G is false.' Answer Option 4: To determine the validity of the argument, we can construct an indirect truth table. The table will include columns for E, F, G, and the premises and conclusion of the argument. $\mathsf{E}\;\mathsf{F}\;\mathsf{G}\;\mathsf{E}\supset (\mathsf{F}\vee \mathsf{\sim}\mathsf{G})\;\mathsf{F}\supset (\mathsf{\sim}\mathsf{E}\vee \mathsf{\sim}\mathsf{G})\;\mathsf{\sim}(\mathsf{E}\equiv\mathsf{F})$ TTTTF TTFTTF TFTTTT TEETET FTTTT FTFTTT FFTTTT FFFTTF As we can see in the last row, when E, F, and G are all false, the premises are true (T), but the conclusion is false (F). Therefore, the argument is invalid. The correct answer is: C. Invalid. Counterexample when E and F are true and G is false. Answer Option 5: To determine the validity of the argument, we can use an indirect truth table. We have the following premises: 1. $E \supset (F \lor \sim G)$ 2. F ⊃ (~E ∨ ~G) We want to determine the truth value of the conclusion: 3. ~(E ≡ F) We construct a truth table for the premises and the conclusion: \mid E \mid F \mid G \mid E \supset (F \vee ~G) \mid F \supset (~E \vee ~G) \mid ~(E \equiv F) \mid |---|---| |T|T|T|T|T|F| |T|T|F|T|T|F| |T|F|T|T|T|T| |T|F|F|F|T|T| |F|T|T|T|T|F| |F|T|F|T|T|F| |F|F|T|T|T|F| |F|F|F|T|T|F| Since there is at least one row in the truth table where all premises are true but the conclusion is false, the argument is invalid. Therefore, the correct answer is C. 'Invalid. Counterexample when E and F are true and G is false'.

Answer Option 6:

To determine the validity of the argument, we can use an indirect truth table.

 $E \supset (F \lor \sim G)$

F ⊃ (~E ∨ ~G)

~(E ≡ F)

Let's build a truth table for the argument:

 $\mathsf{E}\;\mathsf{F}\;\mathsf{G}\;\mathsf{E}\supset (\mathsf{F}\vee \mathsf{\sim}\mathsf{G})\;\mathsf{F}\supset (\mathsf{\sim}\mathsf{E}\vee \mathsf{\sim}\mathsf{G})\;\mathsf{\sim}(\mathsf{E}\equiv\mathsf{F})$

TTTTTF

TTFTTF

TFTTTF

TFFFTT

FTTTTF FTFFFT FFTTTF FFFTTF The last column in the truth table represents the statement -(E = F). We can see that there are rows in the truth table where the premises (the first two columns) are true and the conclusion (the last column) is false. This means that the argument is invalid. To find a counterexample, we need to find a row in the truth table where the premises are true and the conclusion is false. Looking at the truth table, we can see that when E is true, F is true, and G is false, the premises are true (T, T, T) and the conclusion is false (F). Therefore, the argument is invalid. The correct answer is option C: 'Invalid. Counterexample when E and F are true and G is false'. Answer Option 7: To determine the validity of the argument, we can construct an indirect truth table. $E \supset (F \lor \sim G)$ $\mathsf{F}\supset (\mathsf{\sim}\mathsf{E}\,\vee\,\mathsf{\sim}\mathsf{G})$ ~(E ≡ F) First, let's construct a truth table for the premises: $\mathsf{E}\,\mathsf{F}\,\mathsf{G}\,\mathsf{F}\vee\mathsf{\sim}\mathsf{G}\,\mathsf{E}\supset(\mathsf{F}\vee\mathsf{\sim}\mathsf{G})\,\mathsf{\sim}\mathsf{E}\,\mathsf{\sim}\mathsf{E}\vee\mathsf{\sim}\mathsf{G}\,\mathsf{F}\supset(\mathsf{\sim}\mathsf{E}\vee\mathsf{\sim}\mathsf{G})$ TTTTTFFT TTFTTFTT TFTFFFF TEFFETT FTTTTTT FTFTTTT FFTTTTT FFFTTTTT Next, let's add the values for \sim (E \equiv F): $EFG \sim (E \equiv F)$ TTTF TTFF TFTF TFFF FTTT FTFF FFTT FFFT From the truth table, we can see that there is at least one row where both premises are true and the conclusion is false. This means the argument is invalid. Counterexample: When E is true, F is true, and G is false, the premises are true but the conclusion ~(E ≡ F) is false. Final Answer: C. 'Invalid. Counterexample when E and F are true and G is false' Answer Option 8: To determine the validity of the argument using indirect truth tables, we need to consider all possible truth value combinations for the propositions E, F, and G. $E \supset (F \lor \sim G)$ $\mathsf{F}\supset (\mathsf{\sim}\mathsf{E}\,\vee\,\mathsf{\sim}\mathsf{G})$ ∴ ~(E ≡ F) We will construct the truth table by considering each proposition and its negation. $\mathsf{E} \;\mathsf{F}\;\mathsf{G}\;\mathsf{\sim}\mathsf{G}\;\mathsf{F}\;\mathsf{\vee}\;\mathsf{\sim}\mathsf{G}\;\mathsf{E}\;\supset (\mathsf{F}\;\vee\;\mathsf{\sim}\mathsf{G})\;\mathsf{\sim}\mathsf{E}\;\mathsf{\sim}\mathsf{E}\;\vee\;\mathsf{\sim}\mathsf{G}\;\mathsf{F}\;\supset (\mathsf{\sim}\mathsf{E}\;\vee\;\mathsf{\sim}\mathsf{G})\;\mathsf{E}\;\equiv\;\mathsf{F}\;\mathsf{\sim}(\mathsf{E}\;\equiv\;\mathsf{F})$ TTTFTTFTTF

TTFTTTFTTF TFTFFFFFFT TEFTTTETT

FTTFTTTTTFT

FFTFFTTTTFF

FFFTTTTTFF

In the truth table, we can see that when E and F are both true, and G is false, the premises are true, but the conclusion is false. Therefore, the argument is invalid.

Counterexample

E = T, F = T, G = F

Final Answer: C. Invalid. Counterexample when E and F are true and G is false.

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Answer Option 9:

To determine the validity of the argument using indirect truth tables, we will assume the premises to be true and the conclusion to be false, then check if there is a row in the truth table where all the premises are true and the conclusion is false.

Let's construct the truth table:

 $\mathsf{E}\;\mathsf{F}\;\mathsf{G}\;\mathsf{E}\supset (\mathsf{F}\vee \mathsf{\sim}\mathsf{G})\;\mathsf{F}\supset (\mathsf{\sim}\mathsf{E}\vee \mathsf{\sim}\mathsf{G})\;\mathsf{\sim}(\mathsf{E}\equiv\mathsf{F})$

TTTTFF

TTFTTFF

TFTFFTT

TFFFFTT

FTTTTT

FTFTTTT

FFTTTT

FFFTTTT

There is no row in the truth table where all the premises (E \supset (F \lor \neg G) and F \supset (\neg E \lor \neg G)) are true and the conclusion (\neg (E \equiv F)) is false. Thus, the argument is valid.

Final Answer: A. Valid

Question 15:

"Use indirect truth tables to determine whether the following argument is valid. If the argument is invalid, choose an option which presents a counterexample. (There may be other counterexamples as well.)

 $H = f_{-} I \times I$

H v ~J / ~I', A. 'Valid', B. 'Invalid. Counterexample when H, I, and J are false', C. 'Invalid. Counterexample when H and I are true and J is false', D. 'Invalid. Counterexample when I is true and H and J are

Answer Option 1:

To determine the validity of the argument using indirect truth tables, we need to consider all possible truth value assignments for the propositions involved.

Let's create the indirect truth table:

 $H \equiv (\sim I \vee J)$

H ∨ ~J

.

TTT

TTF

TFT

TFF

FTT

FTF

FFT

FFF

From the truth table, we can see that there is no row where both premises (H v ~J and ~I) are true and the conclusion is false. Therefore, the argument is valid.

Final Answer: A. Valid

Answer Option 2:

To determine the validity of the argument using indirect truth tables, we need to examine all possible truth value combinations of the premises and conclusion.

Let's create a truth table to analyze the argument:

H I J ~I ~J H \equiv (~I \vee J) H \vee ~J ~I

TTTFFTTFF

TTFFTFTTF

FTTFFTTFT

FTFFTTFTT

TFTTFTTFF

TFFTTFTTF

FFTTFTTFT

FFFTTTTT

From the truth table, we can see that when the premises (H = (-I v J) and H v ~J) are true, the conclusion (-I) is also true for all possible truth value combinations. Therefore, the argument is valid.

Final Answer: A. Valid

Answer Option 3:

To determine the validity of the argument, we can use an indirect truth table. We will assume the premises are true and the conclusion is false, and then check if the assumptions lead to a contradiction.

Assumptions:

- 1. $H \equiv (\sim I \lor J)$ (Premise)
- 2. H \vee ~J (Premise)
- 3. ~I (Assumption to be negated)

We will construct a truth table with columns for H, I, J, and the implications and equivalences based on the given premises and assumptions.

H I J H \equiv (~I \vee J) H \vee ~J ~I

TTTTFT

TTFTTFT

TFTFTTT

TFFFTTT

FTTTTF

FTFFTF

FFTTTTT FFFTTTT In the last row of the truth table, where our assumption ~I is false, we can see that the premises hold true. Therefore, it is not possible to find a counterexample where the premises are true and the conclusion Hence, the argument is valid. Final Answer: A To determine the validity of the argument, we will use an indirect truth table. We will assume the premises (H v ~J) and find the resulting truth values for the conclusion (~I). Assume H v ~J to be true: H v ~J / ~I ΤF Next, we will use the given equivalence $H \equiv (\sim I \lor J)$ to find the truth values of H: $H \equiv (\sim I \vee J)$ Using the truth value of H, we can determine the truth values of ~I and J: H v ~J / ~I ΤF Т? From the truth values, we can see that when H is true, -I must be false for the conclusion to be true. However, we cannot determine the truth value of J from the given information. Since we can find a case (when ~I is false) where the premises are true but the conclusion is false, the argument is invalid. The correct answer is B. Invalid. Counterexample when H, I, and J are false. Answer Option 5: To determine whether the argument is valid or not, we can use an indirect truth table. We have the premises: 1. H ≡ (~I ∨ J) 2. H v ~J And we want to prove the conclusion: We construct a truth table with columns for H, I, J, \sim I, (\sim I \vee J), and (H \equiv (\sim I \vee J)). $H \mid I \mid J \mid \sim I \mid (\sim I \lor J) \mid (H \equiv (\sim I \lor J))$ TITITIFITIF TITIFIFIFIT TIFITITIF TIFIFITITIF F|T|T|F|T|F F|T|F|F|F|T F|F|T|T|F F|F|F|T|T|F From the truth table, we can see that there is a row where both premises are true (row 2), but the conclusion is false (~I is false). Therefore, the argument is invalid. The counterexample is when H and I are true, and J is false. Final Answer: C. 'Invalid. Counterexample when H and I are true and J is false' ******* Answer Option 6:

Let's consider all possible truth value combinations for H, I, and J:

 $H \equiv (\sim I \lor J)$ $H \lor \sim J$ $/ \sim I$

To determine the validity of the argument, we can use an indirect truth table.

 $H \mid I \mid J \mid \neg I \mid \neg J \mid (\neg I \vee J) \mid H \equiv (\neg I \vee J) \mid H \vee \neg J \mid \neg I$ T|T|T|F|F|T|T|T|F TITIFIFITITITIF TIFITITIFITITIT TIFIFITITITITIT F|T|T|F|F|T|F|F|F F|T|F|F|T|T|F|T|F F|F|T|T|F|T|F|T|T F|F|F|T|T|T|F|T|T We can see that when $H \equiv (-I \lor J)$ and $H \lor \sim J$ are true, $\sim I$ is also true for all possible combinations of truth values for H, I, and J. Therefore, the argument is valid. Final Answer: A. 'Valid' ******** Answer Option 7: To determine the validity of the argument using indirect truth tables, we need to consider all possible truth value assignments for the propositions H, I, and J. $H \equiv (\sim\!I \lor J) \text{ can be represented as } (H \land (\sim\!I \lor J)) \lor (\sim\!H \land \sim\!(\sim\!I \lor J))$ H v ~J can be represented as H v ~J ~I can be represented as ~(I) Combining the premises and the negation of the conclusion, we get the following compound statement: $((H \wedge (\sim I \vee J)) \vee (\sim H \wedge \sim (\sim I \vee J))) \wedge (H \vee \sim J) \wedge \sim (I)$ Let's create a truth table to evaluate the compound statement: $H \mid I \mid J \mid \neg I \mid \neg J \mid (\neg I \vee J) \mid \neg (\neg I \vee J) \mid (H \wedge (\neg I \vee J)) \vee (\neg H \wedge \neg (\neg I \vee J)) \mid H \vee \neg J \mid \neg (I)$ TITITIFIFITIFIFITIF T|T|F|F|T|T|F|F|T|F T|F|T|T|F|T|F|F|T|T T|F|F|T|T|T|F|F|T|T FITITIFIFITIFIFIFIF F|T|F|F|T|T|F|F|F|F F|F|T|T|F|T|F|F|F|T F|F|F|T|T|F|T|T|F|T From the truth table, we can see that there is no row where all the premises are true and the conclusion is false. Therefore, the argument is valid. Final Answer: A. Valid Answer Option 8: To determine the validity of the argument, we can use indirect truth tables. We will assume the premises are true and check if the conclusion must also be true. Assume $H \equiv (\neg I \lor J)$ and $H \lor \neg J$ are true. We will create a truth table to evaluate the argument. $H \mid I \mid J \mid \sim I \vee J \mid H \equiv (\sim I \vee J) \mid H \vee \sim J$ TITITITIT TITIFITITIT T|F|T|T|T T|F|F|F|F|T FITITITIFIT F|T|F|T|F|T F|F|T|T|F|T F|F|F|T|T|F From the truth table, we can see that there is a row where both premises are true (H = (-1 v J) and H v ~J), but the conclusion ~I is false (last row). This means that the argument is invalid. To find a counterexample, we look for a row where the premises are true and the conclusion is false. From the truth table, we can see that when I is true and H and J are false (second last row), both premises are true, but the conclusion ~I is false. Therefore, the counterexample is when I is true and H and J are false. Final Answer: C. Invalid. Counterexample when H and I are true and J is false.

Answer Option 9:

To determine the validity of the argument, we can use an indirect truth table.
$H \equiv (-I \lor J)$
H∨≂J
/~I
First, we list all the relevant propositions and their possible truth values:
H: True, False
I: True, False
J: True, False
Next, we evaluate the truth values of the given statements:
$H \equiv (-I \lor J)$
$True \equiv (\simTrue \lor True) \mathrel{->} True \equiv (False \lor True) \mathrel{->} True \equiv True \mathrel{->} True$
$\label{eq:False} \textit{False} = (\textit{~False} \lor \textit{False}) \textit{~~} \textit{~~} \textit{False} \equiv (\textit{True} \lor \textit{False}) \textit{~~} \textit{~~} \textit{False} \equiv \textit{True} \textit{~~} \textit{~~} \textit{~~} \textit{False}$
H v ~J
True v ~True -> True v False -> True
False > ~False -> False > True -> True
/ ~I
/ ~True -> False
/ ~True -> False / ~False -> True
/ -False -> True
/ -False -> True Now, we construct the indirect truth table:
/ -False -> True Now, we construct the indirect truth table: $H \mid I \mid J \mid H \equiv (-I \lor J) \mid H \lor \neg J \mid / \neg I$
/ -False -> True Now, we construct the indirect truth table: H I J H = (-I \leftarrow J) H \leftarrow -J / -I
/ ~False -> True Now, we construct the indirect truth table: $H \mid I \mid J \mid H \equiv (\neg I \lor J) \mid H \lor \neg J \mid / \neg I$ $\dots \dots $
/ -False -> True Now, we construct the indirect truth table: $H \mid I \mid J \mid H \equiv (-I \lor J) \mid H \lor \neg J \mid / \neg I$ $T \mid T \mid T \mid T \mid F$ $T \mid T \mid F \mid F \mid T \mid F$
/ -False -> True Now, we construct the indirect truth table: $H \mid I \mid J \mid H \equiv (-I \lor J) \mid H \lor \neg J \mid / \neg I$ $T \mid T \mid T \mid T \mid T \mid F$ $T \mid T \mid F \mid T \mid T \mid T \mid T$
/ -False -> True Now, we construct the indirect truth table:
/-False -> True Now, we construct the indirect truth table: = (-1 \lor J)
/-False -> True Now, we construct the indirect truth table: $H \mid I \mid J \mid H = (-I \lor J) \mid H \lor \neg J \mid / \neg I$ $T \mid T \mid T \mid T \mid F$ $T \mid T \mid F \mid F \mid T \mid T$ $T \mid F \mid F \mid T \mid T \mid T$ $F \mid T \mid T \mid F \mid F \mid F$ $F \mid T \mid T \mid F \mid F \mid F$
/ -False -> True Now, we construct the indirect truth table: H H = (- \ J) H \ \ \ \ -J / -1
/ -False -> True Now, we construct the indirect truth table: H H = (- \ \ J) H \ \ \ \ \ J / -

Question 16:

'Select the best translation into predicate logic. Some firefighters rescue victims from danger. (Dx: x is danger; Fx: x is a firefighter; Vx: x is a victim; Rxyz: x rescues y from z)', A. ' $(\exists x)(Fx \bullet (\forall y)[Dy \supset (\exists z)(Vz \bullet Rxyz)])'$, B. ' $(\exists x)(Fx \bullet (\exists y)[Vy \bullet (\exists z)(Dz \bullet Rxyz)])'$, C. ' $(\exists x)(Fx \bullet (\exists y)(\exists z)(Dz \bullet Rxyz)])'$

Answer Option 1:

The best translation into predicate logic for the given statement is:

A. $'(\exists x)\{Fx \bullet (\forall y)[Dy \supset (\exists z)(Vz \bullet Rxyz)]\}'$

This translation accurately represents the statement "Some firefighters rescue victims from danger." It states that there exists at least one firefighter ($\exists x$) such that for all victims ($\forall y$), if there is danger (Dy),

then there exists some (at least one) victim ($\exists z$) who is rescued by the firefighter (Rxyz).

Option A correctly captures the notion that there are firefighters who rescue victims from danger. Final Answer: A

.....

Answer Option 2:

The best translation into predicate logic for the statement "Some firefighters rescue victims from danger" would be:

 $A. \; (\exists x) \{ Fx \bullet (\forall y) [Dy \supset (\exists z) (\forall z \bullet Rxyz)] \}$

Explanation:

- $(\exists x)$ indicates the existence of at least one firefighter.
- Fx represents the statement "x is a firefighter".
- (∀y)[Dy ⊃ (3z)(Vz Rxyz)] translates to "For all y, if y is in danger, then there exists a z such that z is a victim and x rescues y from z". This captures the idea that firefighters rescue victims from danger.

Therefore, option A is the best translation. Final Answer: A

Answer Option 3:

The best translation into predicate logic for the statement "Some firefighters rescue victims from danger" is:

 $(\exists x) \{ \mathsf{Fx} \bullet (\forall y) [\mathsf{Dy} \supset (\exists z) (\mathsf{Vz} \bullet \mathsf{Rxyz})] \}$

Explanation:

 $\exists x$: There exists at least one x

Fx: x is a firefighter

 $\forall y: \text{For all } y$

Dy : y is danger

 $\exists z$: There exists at least one z

Vz : z is a victim

Rxyz : x rescues y from z

Therefore, the correct translation is option A:

 $(\exists x) \{ \mathsf{Fx} \bullet (\forall y) [\mathsf{Dy} \supset (\exists z) (\mathsf{Vz} \bullet \mathsf{Rxyz})] \}$

Answer Option 4:

The best translation into predicate logic for the statement "Some firefighters rescue victims from danger" would be:

 $A. \; (\exists x) \{ \mathsf{Fx} \bullet (\forall y) [\mathsf{Dy} \supset (\exists z) (\mathsf{Vz} \bullet \mathsf{Rxyz})] \}$

Explanation:

- (∃x) is used to denote the existence of at least one firefighter.
- Fx: x is a firefighter.
- $(\forall y) [Dy \supset (\exists z) (\forall z \bullet Rxyz)] \text{: For all } y, \text{ if } y \text{ is in danger (Dy), then there exists at least one } z \text{ such that } z \text{ is a victim (Vz) and } x \text{ (firefighter) rescues } y \text{ from } z \text{ (Rxyz)}.$

This translation accurately captures the statement's meaning that there is at least one firefighter who rescues victims from danger. Final Answer: A

Answer Option 5:

 $The \ best \ translation \ into \ predicate \ logic \ for \ the \ given \ statement \ "Some \ firefighters \ rescue \ victims \ from \ danger" \ is:$

 $(\exists x) \{ \mathsf{Fx} \bullet (\forall y) [\mathsf{Dy} \supset (\exists z) (\mathsf{Vz} \bullet \mathsf{Rxyz})] \}$

Explanation:

Let's break down the translation step by step:

Fx: x is a firefighter

 $(\exists x) \{ \mathsf{Fx} \bullet ... \} :$ There exists at least one firefighter such that ...

 $(\forall y)[Dy\supset...]:$ For all victims y, if y is in danger ...

 $(\exists z) (\forall z \bullet \mathsf{Rxyz}) : \mathsf{There} \ \mathsf{exists} \ \mathsf{at} \ \mathsf{least} \ \mathsf{one} \ \mathsf{victim} \ \mathsf{z}, \mathsf{such} \ \mathsf{that} \ \mathsf{z} \ \mathsf{is} \ \mathsf{a} \ \mathsf{victim} \ \mathsf{and} \ \mathsf{x} \ \mathsf{rescues} \ \mathsf{y} \ \mathsf{from} \ \mathsf{z}.$

Putting it all together, we get:

```
(\exists x) \{ \mathsf{Fx} \bullet (\forall y) [\mathsf{Dy} \supset (\exists z) (\mathsf{Vz} \bullet \mathsf{Rxyz})] \}
Thus, the correct answer is A.
Answer Option 6:
The best translation into predicate logic for the given statement would be:
A. \; (\exists x) \{ \mathsf{Fx} \bullet (\forall y) [\mathsf{Dy} \supset (\exists z) (\mathsf{Vz} \bullet \mathsf{Rxyz})] \}
Let's break down the given statement:
"Some firefighters rescue victims from danger."
To translate this into predicate logic, we can use the following variables:
Fx: x is a firefighter
Dx: x is danger
Vx: x is a victim
Rxyz: x rescues y from z
The statement can be represented as:
"There exists a firefighter who, for all victims, if they are in danger, there exists a rescuer who rescues them."
This can be expressed as:
(\exists x) \{ \mathsf{Fx} \bullet (\forall y) [\mathsf{Dy} \supset (\exists z) (\mathsf{Vz} \bullet \mathsf{Rxyz})] \}
Therefore, option A is the best translation. Final Answer: A.
Answer Option 7:
The best translation into predicate logic for the given statement "Some firefighters rescue victims from danger" is:
(\exists x) \{ \mathsf{Fx} \bullet (\forall y) [\mathsf{Dy} \supset (\exists z) (\mathsf{Vz} \bullet \mathsf{Rxyz})] \}
Explanation:
- (\exists x) denotes the existence of at least one firefighter.
- Fx represents x being a firefighter.
- (\forall y) signifies for all victims.
- Dy represents y being in danger.
- (\exists z)(\forall z \bullet \mathsf{Rxyz}) denotes the existence of at least one victim that is rescued by x from z.
Therefore, the correct answer is option A: (\exists x) \{ Fx \bullet (\forall y) [Dy \supset (\exists z) (\forall z \bullet Rxyz)] \}'.
******************
Answer Option 8:
The best translation into predicate logic for the given statement is:
A. \; (\exists x) \{ Fx \bullet (\forall y) [Dy \supset (\exists z) (\forall z \bullet Rxyz)] \}
Explanation:
Let's break down the given statement:
"Some firefighters rescue victims from danger."
The translation can be broken down into two parts:
1. (∃x){Fx • ...}: There exists a firefighter who...
- (\exists x)\{Fx\}: There exists a firefighter (x)
2.\ ...\ (\forall y)[\mathsf{D} y \supset (\exists z)(\mathsf{V} z \bullet \mathsf{R} x y z)] :
- (\forall y): For all victims (y)
- [Dy \supset (\exists z)(Vz \bullet Rxyz)]: if they are in danger (Dy), then there exists someone (z) who is a victim (Vz) and is rescued by the firefighter (Rxyz).
Therefore, the correct translation is:
(\exists x) \{ \mathsf{Fx} \bullet (\forall y) [\mathsf{Dy} \supset (\exists z) (\mathsf{Vz} \bullet \mathsf{Rxyz})] \}
Final Answer: A. (\exists x){Fx \cdot (\forall y)[Dy \supset (\exists z)(Vz \cdot Rxyz)]}
Answer Option 9:
The best translation into predicate logic for the given statement would be:
A. \; (\exists x) \{ \mathsf{Fx} \bullet (\forall y) [\mathsf{Dy} \supset (\exists z) (\mathsf{Vz} \bullet \mathsf{Rxyz})] \}
Let's break down the statement:
"Some firefighters rescue victims from danger."
```

We can define the predicates as:

Fx: x is a firefighter

Dx: x is danger

Vx: x is a victim

Rxyz: x rescues y from z

Now, let's analyze the options:

 $A. \; (\exists x) \{ \mathsf{Fx} \bullet (\forall y) [\mathsf{Dy} \supset (\exists z) (\mathsf{Vz} \bullet \mathsf{Rxyz})] \}$

This option correctly states that there exists a firefighter $(\exists x)$ who satisfies the condition that for all victims $(\forall y)$, if they are in danger (Dy), then there exists a rescuer $(\exists z)$ who is a firefighter and rescues the victim from danger.

 $\mathsf{B.}\ (\exists \mathsf{x}) \{ \mathsf{Fx} \bullet (\exists \mathsf{y}) [\mathsf{Vy} \bullet (\exists \mathsf{z}) (\mathsf{Dz} \bullet \mathsf{Rxyz})] \}$

This option states that there exists a firefighter (3x) who satisfies the condition that there exists a victim (3y) and a danger (3z), and the firefighter rescues the victim from the danger. However, it does not account for the possibility that not all victims are in danger.

C. $(\exists x)\{Fx \cdot (\exists y)(\exists z)[(\forall y \cdot Dz) \cdot Rxzy]\}$

This option states that there exists a firefighter ($\exists x$) who satisfies the condition that there exists a victim ($\exists y$) and a danger ($\exists z$), and the firefighter rescues the victim from the danger. However, it does not account for the possibility that not all victims are in danger.

 $\mathsf{D}.\ (\forall x)\{\mathsf{F} x \supset (\exists y)[\mathsf{V} y \bullet (\exists z)(\mathsf{D} z \bullet \mathsf{R} x y z)]\}$

This option states that for all firefighters ($\forall x$), if they are a firefighter (Fx), then there exists a victim ($\exists y$) and a danger ($\exists z$), and the firefighter rescues the victim from the danger. However, it does not capture the possibility that not all victims are in danger.

 $Therefore, the correct translation into predicate logic is option A. (\exists x)\{Fx \bullet (\forall y)[Dy \supset (\exists z)(Vz \bullet Rxyz)]\}. \ Final Answer: All the correct translation into predicate logic is option A. (\exists x)\{Fx \bullet (\forall y)[Dy \supset (\exists z)(Vz \bullet Rxyz)]\}. \ Final Answer: All the correct translation into predicate logic is option A. (\exists x)\{Fx \bullet (\forall y)[Dy \supset (\exists z)(Vz \bullet Rxyz)]\}.$

Jestion	

'Identify the antecedent of the following conditional proposition: The university raises tuition if, and only if, both the governor approves of it and the board of trustees recommends it.', A. The university raises tuition.', B. The governor approves of it.', C. The board of trustees recommends it.', D. 'None of the above'

Answer Option 1:

The antecedent of a conditional proposition is the part that appears before the "if" or "only if" in the statement. In this case, the antecedent of the conditional proposition is "both the governor approves of it and the board of trustees recommends it." This corresponds to option D, "None of the above.' Final Answer: D

Answer Option 2:

The antecedent of a conditional proposition is the part that comes after the "if" in the statement. In this case, the antecedent of the conditional proposition "The university raises tuition if, and only if, both the governor approves of it and the board of trustees recommends it."

Therefore, the correct answer is C. 'The board of trustees recommends it.'

Answer Option 3:

The antecedent of a conditional proposition is the part that comes before the "if" or "only if" statement. In this case, the antecedent of the given conditional proposition is "both the governor approves of it and the board of trustees recommends it." Therefore, the correct answer is C. The board of trustees recommends it.'

Answer Option 4:

The antecedent of a conditional proposition is the statement that comes after the "if" part of the statement. In this case, the antecedent is the condition that needs to be satisfied for the university to raise tuition.

Based on the given statement, the antecedent is: 'Both the governor approves of it and the board of trustees recommends it.'

Therefore, the correct answer is C. 'The board of trustees recommends it.'

Answer Option 5:

The antecedent of a conditional proposition is the part that comes before the "if" or "only if" statement. In this case, the antecedent is "both the governor approves of it and the board of trustees recommends it." Therefore, the correct answer is C. The board of trustees recommends it.'

Answer Option 6:

The antecedent of a conditional proposition is the part that comes before the 'if' or 'only if' in the statement. In this case, the antecedent is "both the governor approves of it and the board of trustees recommends it." Therefore, the correct answer is C. 'The board of trustees recommends it.'

Answer Option 7:

The antecedent of a conditional proposition is the part that comes immediately after the "if" or "only if" statement. In this case, the antecedent is the condition that needs to be satisfied for the university to raise tuition.

The antecedent of the given conditional proposition is:

B. 'The governor approves of it.'

This means that the university will raise tuition only if the governor approves of it. Final Answer: B

Answer Option 8:

The antecedent of a conditional proposition is the statement that appears after the "if" or "only if" in the sentence. In this case, the antecedent of the conditional proposition "The university raises tuition if, and only if, both the governor approves of it and the board of trustees recommends it."

Therefore, the correct answer is C. 'The board of trustees recommends it.'

Answer Option 9:

The antecedent of a conditional proposition is the part that comes after the "if" in the sentence. In this case, the antecedent is "both the governor approves of it and the board of trustees recommends it."

Therefore, the correct answer is C. 'The board of trustees recommends it.'

Question 18:
'Select the best translation into predicate logic. All planets are smaller than Jupiter. (j: Jupiter; Px: x is a planet; Sxy: x is smaller than y)', A. ' $(\exists x)(Px \bullet Sxj)'$, B. ' $(\forall x)(Px \bullet Sxj)'$, C. ' $(\forall x)(Px \supset Sjx)'$, D. ' $(\forall x)(Px \supset Sjx)'$,
Sxj)'
Answer Option 1:
The best translation into predicate logic for the statement "All planets are smaller than Jupiter" is:
$(\forall x)(Px \to Sxj)$
Option D: $(\forall x)(Px \supset Sxj)'$ is the correct translation. Final Answer: D
Answer Option 2:
The best translation into predicate logic for the statement *All planets are smaller than Jupiter* is:
$(\forall x)(P x \to S x j)$
Option D: $'(\forall x)(Px \supset Sxj)'$ is the correct translation. Final Answer: D

Answer Option 3:
The best translation into predicate logic for the statement "All planets are smaller than Jupiter" is:
$(\forall x)(Px\supset Sxj)$
Explanation:
- Px: x is a planet
- Sxj: x is smaller than Jupiter
The statement can be translated as "For all x, if x is a planet (Px), then x is smaller than Jupiter (Sxj)."
So, the correct answer is D. $(\forall x)(Px\supset Sxj)$

Answer Option 4:
The best translation into predicate logic for the statement "All planets are smaller than Jupiter" would be option D: $(\forall x)(Px \supset Sxj)'$.
To break it down:
- $(\forall x)$ means "for all x ".
- Px represents "x is a planet".
- Sxj represents "x is smaller than Jupiter".
Therefore, the statement can be translated as: "For all x, if x is a planet, then x is smaller than Jupiter." This captures the meaning of the original statement accurately. Final Answer: D

Answer Option 5:
The best translation into predicate logic for the statement "All planets are smaller than Jupiter" would be: $(\forall x)(Px \supset Sxj)$.
Explanation:
- (∀x) means "for all x" or "for every x"
- Px represents "x is a planet"
- Sxj represents "x is smaller than Jupiter"
So, $(\forall x)(Px \supset Sxj)$ can be read as "For all x, if x is a planet, then x is smaller than Jupiter", which accurately represents the given statement.
Final Answer: D
Answer Option 6:
The best translation into predicate logic for the statement "All planets are smaller than Jupiter" is:
$(\forall x)(Px\supset Sxj)$
Explanation:
- (∀x) indicates that the statement applies to all x.
- Px represents "x is a planet."
- Sxj represents "x is smaller than Jupiter."
The translation $(\forall x)(Px \supset Sxi)$ states that for all x , if x is a planet, then x is smaller than Jupiter. This accurately represents the original statement. Therefore, the correct answer is D.
Answer Option 7: The heat translation into predicate legis for the attempts "All placets are amplied than Junited" in:
The best translation into predicate logic for the statement "All planets are smaller than Jupiter" is:
$(\forall x)(Px \supset Sxj)$

Explanation:

- ($\forall x$) represents "for all x" or "for every x" which indicates that we are referring to every possible object x.
- Px represents "x is a planet." This predicate is true for objects that are planets.
- Sxj represents "x is smaller than Jupiter." This predicate is true for objects that are smaller than Jupiter.

Therefore, the best translation is $(\forall x)(Px \supset Sxj)$. Final Answer: D

Answer Option 8:

The best translation into predicate logic for the statement "All planets are smaller than Jupiter" is option D. $'(\forall x)(Px\supset Sxj)'$

Explanation:

Let's break down the statement:

"All planets are smaller than Jupiter."

- We can represent "All planets" as $(\forall x)(Px)$, which means for every x, if x is a planet (Px is true).
- "smaller than Jupiter" can be represented as Sxj, which means x is smaller than Jupiter (Sxj is true).
- Combining both parts, we get $(\forall x)(Px \supset Sxj)$, which translates to "For all x, if x is a planet, then x is smaller than Jupiter."

Answer Option 9:

The best translation into predicate logic for the statement "All planets are smaller than Jupiter" would be option C: $'(\forall x)(Px \supset Sjx)'$.

Let's break down the translation:

- Px: x is a planet (predicate representing "x is a planet")
- Sxy: x is smaller than y (predicate representing "x is smaller than y")
- j: Jupiter (constant representing Jupiter)

The statement "All planets are smaller than Jupiter" can be translated as "For all x, if x is a planet (Px), then x is smaller than Jupiter (Sjx)." This can be represented as $(\forall x)(Px \supset Sjx)$.

Option A: '(∃x)(Px • Sxj)' translates to "There exists an x such that x is a planet (Px) and x is smaller than Jupiter (Sxj)". This does not capture the meaning of the original statement which states that all planets are smaller than Jupiter.

Option B: '(∀x)(Px • Sx))' translates to "For all x, x is a planet (Px) and x is smaller than Jupiter (Sx))". This is not correct as it implies that all objects, not just planets, are smaller than Jupiter.

Option D: "(∀x)(Px ⊃ Sxj)" translates to "For all x, if x is a planet (Px), then x is smaller than Jupiter (Sxj)". This is the same as option C and correctly captures the meaning of the original statement.

Final Answer: C

Question 19:
'Select the best translation into predicate logic. José introduces Wilma to Kevin. (j: José; k: Kevin; w: Wilma; lxyz: x introduces y to z)', A. 'lwkj', B. 'ljwk', C. 'lwjk', D. 'ljkw'
Answer Option 1:
The best translation into predicate logic for the statement "José introduces Wilma to Kevin" would be option C: 'lwjk'. This can be read as "Wilma introduces José to Kevin" in English.
Answer Option 2:
The best translation into predicate logic would be option C: 'lwjk'. This translates to "Wilma introduces José to Kevin".
Answer Option 3:
The best translation into predicate logic for the statement "José introduces Wilma to Kevin" would be option A: 'lwkj'. This can be read as "José introduces Wilma to Kevin", where I(w, k, j) represents the
predicate "José introduces Wilma to Kevin". Final Answer: A
Answer Option 4:
The best translation into predicate logic for the statement "José introduces Wilma to Kevin" is 'ljkw'. This represents the predicate 'l' (introduces) with the arguments 'j' (José), 'w' (Wilma), and 'k' (Kevin), in that
order. Final Answer: D

Answer Option 5:
The best translation into predicate logic for the statement "José introduces Wilma to Kevin" is "ljwk'.
Explanation:
- lxyz represents "x introduces y to z".
- In this case, José (j) introduces Wilma (w) to Kevin (k).
- Therefore, the correct translation is 'ljwk'.
Final Answer: B
Answer Option 6:
The best translation into predicate logic for the given statement would be: 'ljkw', which can be read as "José introduces Wilma to Kevin."
In this translation, the variables correspond to the following entities:
- j represents José
- k represents Kevin
- w represents Wilma
. Therefore, 'ljkw' accurately represents the statement "José introduces Wilma to Kevin."
Final Answer: D
Answer Option 7:
The best translation into predicate logic would be: 'ljkw' which represents 'José introduces Wilma to Kevin'.
Explanation:
lxyz: x introduces y to z
In this case, José is the one introducing Wilma to Kevin. Therefore, we have:
x = j (José)
y = w (Wilma)
z = k (Kevin)
Plugging these values into the predicate logic representation, we get 'ljkw'. Final Answer: D. 'ljkw'

Answer Option 8: The part translation into predicate legic for the given statement is D. "like". The statement " lock introduces Willma to Kovin" can be translated as "like" where:
The best translation into predicate logic for the given statement is D. 'ljkw'. The statement "José introduces Wilma to Kevin" can be translated as "ljkw" where:
- I represents the introduction relation,
- j represents José,
- k represents Kevin, and
- w represents Wilma.
Answer Option 9:

The best translation into predicate logic for the statement "José introduces Wilma to Kevin" would be option B: "ljwk". This can be read as "José introduces Wilma to Kevin", where I represents the introducing
relation, j represents José, w represents Wilma, and k represents Kevin.

Question 20:

' Select the best translation into predicate logic: No artifacts are people.', A. '~Pa', B. '(∀x)(Ax ⊃ ~Px)', C. '~Ap', D. '(∀x)(Ax • ~Px)'

Answer Option 1:

The best translation into predicate logic for the statement "No artifacts are people" is:

 $\mathsf{B.}\; (\forall x)(\mathsf{A} x\supset \mathsf{\sim} \mathsf{P} x)$

Explanation:

- Let A(x) represent "x is an artifact" and P(x) represent "x is a person."
- The statement "No artifacts are people" can be translated as "For all x, if x is an artifact, then x is not a person."
- This can be symbolized as $(\forall x)(Ax \supset \sim Px)$.

Final Answer: B

Answer Option 2:

The best translation into predicate logic for the statement "No artifacts are people" is option C. '~Ap'. Here's the breakdown of the translation:

Let Ap represent "x is an artifact" and P(x) represent "x is a person".

To express "No artifacts are people", we can use the negation of "There exists an x such that x is an artifact and x is a person".

This can be translated as " \sim ($\exists x$)($Ax \land Px$)", which is equivalent to " \sim Ap" in predicate logic.

Therefore, the correct answer is C. '~Ap'.

......

Answer Option 3:

The given statement is "No artifacts are people." To translate this into predicate logic, we can break it down into two parts: 1) There exist artifacts and 2) No artifacts are people.

Let's define the following predicates:

Ap(x): x is an artifact.

Pp(x): x is a person.

1) There exist artifacts: $(\exists x)Ap(x)$

2) No artifacts are people: $(\forall x)(Ap(x) \rightarrow \sim Pp(x))$

Combining both parts, the best translation into predicate logic is:

 $(\exists x) \mathsf{Ap}(x) \mathrel{\wedge} (\forall x) (\mathsf{Ap}(x) \rightarrow \mathsf{\sim Pp}(x))$

 $Option \ B, \ "(\forall x)(Ax \supset \neg Px)", is \ not \ the \ correct \ translation \ as \ it \ does \ not \ capture \ the \ meaning \ of \ "no \ artifacts \ are \ people."$

Final Answer: B

Answer Option 4:

The best translation into predicate logic for the statement "No artifacts are people" is option C: '~Ap'. Here's the breakdown:

Let A(x) represent "x is an artifact" and P(x) represent "x is a person".

The statement "No artifacts are people" can be translated as "It is not the case that there exists an artifact that is also a person".

Using the symbol '~' to represent negation and the symbol '∃' to represent existential quantification, the translation becomes:

 ${\sim}(\exists x)(A(x) \wedge P(x))$

Simplifying further, we can use De Morgan's law to distribute the negation:

 $(\forall x)(\sim A(x) \vee \sim P(x))$

In this translation, $\sim A(x)$ represents "x is not an artifact" and $\sim P(x)$ represents "x is not a person".

However, option C, '~Ap', is a more concise and equivalent representation of the translation.

Therefore, the best translation into predicate logic is option C: '~Ap'.

Answer Option 5:

 $The \ best \ translation \ into \ predicate \ logic \ for \ the \ statement \ "No \ artifacts \ are \ people" \ is \ option \ C: \ '-Ap'.$

In this translation, 'Ap' represents the predicate 'x is a person' and '-' represents the negation or "not" operator. So '-Ap' translates to "It is not the case that x is a person" or "x is not a person." This accurately captures the meaning of the original statement that no artifacts are people.

Options A, B, and D do not accurately represent the statement "No artifacts are people" in predicate logic. Option A, "-Pa", would translate to "It is not the case that x is an artifact," which is a different statement. Option B, '(∀x)(Ax > ¬Px)', translates to "For all x, if x is an artifact, then x is not a person," which implies that artifacts can exist but cannot be people. Option D, '(∀x)(Ax • ¬Px)', translates to "For all x, x is an artifact and x is not a person," which allows for the possibility that artifacts and people can coexist.

Therefore, option C, '~Ap', is the best translation into predicate logic for the given statement. Final Answer: C.

Answer Option 6:

The phrase "No artifacts are people" can be translated into predicate logic as "For all x, if x is an artifact, then x is not a person." Therefore, the best translation into predicate logic would be option B: "(\forall x)(Ax
⇒ ~Px)" which reads as "For all x, if x is an artifact, then x is not a person." Final Answer: B

Answer Option 7:
The statement "No artifacts are people" can be translated into predicate logic as: "For all x, if x is an artifact, then x is not a person."
Using the following predicates:
Ap: x is an artifact
Px: x is a person
The best translation into predicate logic is:
$(\forall x)(Ax\supset \neg Px)$
Option B is the correct answer.
•••••••••
Answer Option 8:
The best translation into predicate logic for the statement "No artifacts are people" would be option C: '-Ap'. Here's why:
Let's break down the statement:
- "No artifacts": This can be translated as "There does not exist any artifact" or "For all x, x is not an artifact" (∀x)(~Ax).
- "are people": This can be translated as "x is not a person" or "~Px".
Combining the two translations, we have $(\forall x)(-Ax \cdot \neg Px)$. However, this translation states that everything is not an artifact and not a person, which is not what the original statement is trying to convey.
To accurately translate the statement, we need to use the negation of the existential quantifier (3) and combine it with the translation for "artifacts" and "people". This gives us the translation ~Ap.
Therefore, option C: '-Ap' is the best translation into predicate logic for the statement "No artifacts are people".
•••••••••
Answer Option 9:
The best translation into predicate logic for the given statement, "No artifacts are people," would be option C. '-Ap'.
In predicate logic, 'Ap' represents the predicate 'x is a person' and '-' denotes negation. Therefore, '-Ap' translates to 'x is not a person' or in other words, "No artifacts are people."
Option A, '~Pa', would translate to 'x is not an artifact,' which is the negation of the original statement and does not accurately represent the given statement.
Option B, '($\forall x$)(Ax $\supset \sim$ Px)', translates to 'For all x, if x is an artifact, then x is not a person.' This statement implies that there are artifacts that are not people, but it does not capture the meaning of "No artifacts
are people."
Option D, '($\forall x$)(Ax • -Px)', translates to 'For all x, x is an artifact and x is not a person.' This statement implies that every object is both an artifact and not a person, which is not the intended meaning of the
original statement.

Therefore, the correct translation into predicate logic is option C. '~Ap'.

Question 21:
'Construct a complete truth table for the following argument. Then, using the truth table, determine whether the argument is valid or invalid. If the argument is invalid, choose an option which presents a
counterexample. (There may be other counterexamples as well.)
P⊃Q
Q · R / ~P · R', A. 'Valid', B. 'Invalid. Counterexample when P, R, and Q are true', C. 'Invalid. Counterexample when P and Q are true and R is false', D. 'Invalid. Counterexample when R and Q are true and P
is false'
Answer Option 1:
To construct a complete truth table for the argument, we need to consider all possible combinations of truth values for the propositions P, Q, and R.
P Q R ~P Q⊃R Q·R ~P·R
TTTFTTF
TTFFFFF
TFTFTTF
TFFFTFF
FITTITI
FTFTTFT
FFTTTTT
FFFTTFT
Based on the truth table, we can see that there are no rows where the premise $(P \supset Q)$ is true and the conclusion $(\sim P \cdot R)$ is false. Therefore, the argument is valid.
The correct answer is A. 'Valid'.
•••••••
Answer Option 2:
To construct a complete truth table for the argument, we need to consider all possible truth values for the propositions P, Q, and R.
The argument consists of two premises and a conclusion. The premises are:
1. P⊃Q
2. Q · R
The conclusion is:
3. ~P · R
Using the truth table, we can evaluate the truth values of the premises and the conclusion for all possible combinations of truth values for P, Q, and R:
$P Q R P\supset Q Q\cdot R \sim P\cdot R$
TITITITIF
TITIFITIFIF
TIFITIFIFIT
T F F F F
FITITITIT
F T F T F F
F F T T F T
F F F T F F
From the truth table, we can see that there are cases where the premises are true (PDQ and QR) and the conclusion is false (~P-R). Specifically, when P is true, Q is true, and R is false, the premises are
true, but the conclusion is false.
Therefore, the argument is invalid, and the correct answer is C. 'Invalid. Counterexample when P and Q are true and R is false'.
••••••••••••
Answer Option 3:
To construct a complete truth table for the argument, we need to consider all possible truth values for the propositions P, Q, and R.
$PQRQ \cdot RP \supset Q \sim P \cdot R$

TTTTTT
TTFFTF
TFFFFF
FTTTTT
FTFFTF

```
FFTFTT
FFFFTF
In the truth table, '~' represents negation (not), '\supset' represents implication (if...then), and '·' represents conjunction (and).
To determine whether the argument is valid or invalid, we need to check if the conclusion (-P \cdot R) is always true whenever the premises (P \supset Q \text{ and } Q \cdot R) are true.
In this case, the conclusion is true for all the rows where the premises are true (rows 1, 3, and 5). Therefore, the argument is valid.
Final Answer: A. 'Valid'
Answer Option 4:
To construct a complete truth table for the argument, we need to consider all possible combinations of truth values for P, Q, and R.
PQRQ\cdot R \sim P\cdot R
TTTTF
TTFFF
TFTFT
TFFFT
FTTTF
FTFFF
FFTFT
FFFFT
The truth table shows all possible combinations of truth values for the given propositions and the resulting truth values for Q-R and \simP-R.
Now, let's analyze the argument. The argument states:
P \supset Q
∴ ~P · R
To determine the validity of the argument, we need to check if the conclusion (-P \cdot R) is true in all cases where both premises (P \supset Q \text{ and } Q \cdot R) are true.
In the truth table, we can see that there is no case where both premises are true and the conclusion is false. Therefore, the argument is valid.
Final Answer: A. 'Valid'
*******
Answer Option 5:
To construct a complete truth table for the argument, we need to consider all possible combinations of truth values for the propositions P, Q, and R.
PQRQ·R~P·R
TTTTF
TTFFF
TFTFT
TFFFT
FTTTT
FTFFF
FFTFT
FFFFT
Now, let's analyze the truth table to determine the validity of the argument. If in all rows where the premises (P \supset Q) and (Q \cdot R) are true, the conclusion (\negP \cdotR) is also true, then the argument is valid.
Otherwise, it is invalid.
From the truth table, we can see that whenever the premises (P \supset Q) and (Q \cdot R) are true (rows 1, 5, and 7), the conclusion (-P \cdot R) is also true. Therefore, the argument is valid.
Final Answer: A. 'Valid'
******************
Answer Option 6:
To construct a truth table for the given argument, we need to consider all possible combinations of truth values for the propositions P, Q, and R.
PQR~P~P·RQ·R~P·R
TTTFTTF
TTFFFF
TETETEE
TEFFFF
FTTTTT
```

FTFTFF

FETTTET FFFTFF Based on the truth table, we can evaluate the argument: $P \supset Q, Q \cdot R \blacksquare \sim P \cdot R$ In the truth table, there are rows where the premises $(P\supset Q \text{ and } Q\cdot R)$ are true, but the conclusion $(-P\cdot R)$ is false. For example, when P=T, Q=T, and R=F, the premises are true $(P\supset Q \text{ is true and } Q\cdot R)$ is true), but the conclusion (~P · R) is false. Therefore, the argument is invalid. Counterexample: When P and Q are true and R is false. (Option C) Final Answer: B. Invalid. Counterexample when P, Q, and R are true. Answer Option 7: To construct a complete truth table for the argument, we need to consider all possible combinations of truth values for the variables P, Q, and R. $P|Q|R|P \supset Q|Q \cdot R| \sim P \cdot R$ TITITITIF TITIFITIFIF TIFITIFIFIT T|F|F|F|F|T FITITITIT F|T|F|T|F|T F|F|T|T|F|T F|F|F|T|F|T Based on the truth table, we can see that there is at least one row where the premises $(P \supset Q)$ and $(Q \cdot R)$ are true but the conclusion $(\sim P \cdot R)$ is false. Specifically, when P is true, Q is true, and R is false, the premises are true but the conclusion is false. Therefore, the argument is invalid. The correct answer is option C: "Invalid. Counterexample when P and Q are true and R is false." Answer Option 8: To construct a complete truth table for the argument: 1. Identify all the variables involved: P, Q, and R. 2. Create a column for each variable, including the premises and the conclusion: P, Q, R, P \supset Q, Q \cdot R, \sim P \cdot R. 3. Assign all possible combinations of truth values to the variables, filling in the corresponding rows of the truth table. 4. Evaluate each statement in the argument based on the assigned truth values. 5. Determine if the conclusion is always true whenever the premises are true. $PQRP \supset QQ \cdot R \sim P \cdot R$ TTTTF TTFTFF TETEET TFFFFT FTTTTT FTFTFF FETTET Looking at the truth table, we can see that there is a row (row 3) where all the premises are true ($P\supset Q$ and $Q\cdot R$), but the conclusion ($\sim P\cdot R$) is false. Therefore, the argument is invalid. Counterexample: When P is true, Q is false, and R is true, the premises are true (P \supset Q and Q \cdot R), but the conclusion (\neg P \cdot R) is false. Therefore, the counterexample is: P = true, Q = false, R = true. Final Answer: B. Invalid. Counterexample when P, R, and Q are true.

To construct a complete truth table for the argument, we need to consider all possible combinations of truth values for the propositions P, Q, and R.

 $P \supset Q$ (P implies Q)

We have the following propositions:

Answer Option 9:

To determine whether the argument is valid or invalid, we examine the truth values in the last column (~P · R). If this column is always true, then the argument is valid. If there exists at least one row where this column is false, then the argument is invalid and we can provide a counterexample.

In this case, there is a row where $\sim P \cdot R$ is false (the second row, where P = T, Q = T, R = F). Therefore, the argument is invalid.

Counterexample: P = T, Q = T, R = F

Final Answer: C. Invalid. Counterexample when P and Q are true and R is false.

Question 22:
'Select the best translation into predicate logic. Marco moves from Italy to Spain. (i: Italy; m: Marco; s: Spain; Mxyz: x moves to y from z)', A. 'Msim', B. 'Mmis', C. 'Mims', D. 'Mmsi'
Answer Option 1:
The best translation into predicate logic would be: C. 'Mims'
Answer Option 2:
The best translation into predicate logic for the statement "Marco moves from Italy to Spain" would be option C: 'Mims'. This translates to "Marco moves from Italy to Spain" where 'Mims' represents Marco
moving from Italy (i) to Spain (s). Final Answer: C
Answer Option 3:
The best translation into predicate logic for the statement "Marco moves from Italy to Spain" would be option D: "Mmsi". This translation represents the proposition that Marco (m) moves (M) from Italy (i) to
Spain (s).

Answer Option 4:
The best translation of the statement "Marco moves from Italy to Spain" into predicate logic is option D: 'Mmsi'. Here, M represents the predicate "moves to", m represents Marco, s represents Spain, and
represents Italy. Therefore, 'Mmsi' translates to "Marco moves from Italy to Spain". Final Answer: D

Answer Option 5:
The best translation into predicate logic for the statement "Marco moves from Italy to Spain" would be option A: "Msim". This translates to "Marco moves from Italy to Spain" where M represents the predict
"moves", s represents Spain, i represents Italy, and m represents Marco. Final Answer: A

Answer Option 6:
The best translation into predicate logic would be: C. 'Mims'
In predicate logic, the statement "Marco moves from Italy to Spain" can be translated as "Mims" where:
M: Marco moves
i: Marco is in Italy
s: Marco is in Spain
So, "Mims" represents the statement that Marco moves from Italy to Spain.

Answer Option 7:
The best translation into predicate logic for the statement "Marco moves from Italy to Spain" is 'Mims'.
Mims represents Marco moves from Italy to Spain.
In this representation:
- M represents "Marco moves"
- i represents "from Italy"
- s represents "to Spain"
Therefore, the correct translation is 'Mims'. Final Answer: C
Answer Option 8:
The best translation into predicate logic for the statement "Marco moves from Italy to Spain" would be option C. "Mims". This represents the predicate "M(x, v, z)" which can be interpreted as "x moves to v

The best translation into predicate logic for the statement "Marco moves from Italy to Spain" would be option C, "Mims". This represents the predicate "M(x, y, z)" which can be interpreted as "x moves to y from z". In this case, "Marco" is represented by "m", "Italy" is represented by "i", and "Spain" is represented by "s". Therefore, "Mims" accurately represents the statement "Marco moves from Italy to Spain" in predicate logic. Final Answer: C

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Answer Option 9

The best translation into predicate logic for the statement "Marco moves from Italy to Spain" would be option D: 'Mmsi'.

In predicate logic, we represent the statement using variables and predicates. Here, 'M' represents the predicate for "moves," 'm' represents the variable for Marco, 's' represents the variable for Spain, and 'i' represents the variable for Italy.

Therefore, 'Mmsi' translates to "Marco moves from Italy to Spain" and is the best representation in predicate logic. Final Answer: D

Question 23: Select the best English interpretation of the given proposition, using the following translation key: Ax: x is an apartment Hx: x is a house Lx: x is large Bxy: x is bigger than y $(\forall x)(Ax \supset (\forall y)[(Hy \cdot Ly) \supset -Bxy])'$ A. 'All apartments are not bigger than some large house.', B. 'No apartment is bigger than any large house.', C. 'All large houses are not bigger than some apartment.', D. 'No large house is bigger than all apartments.' Answer Option 1: The given proposition can be translated as: "For all apartments x, if x is an apartment, then for all houses y, if y is a house and x and y are both large, then x is not bigger than y." Using the translation key, we can interpret the proposition as: "All apartments are not bigger than some large house." Therefore, the correct interpretation is: A. 'All apartments are not bigger than some large house.' To interpret the given proposition, let's break it down: $(\forall x)\{Ax\supset (\forall y)[(Hy\bullet Ly)\supset -Bxy]\}$ ∀x: For all x Ax: x is an apartment ⇒: implies (∀y): For all y Hy: y is a house Ly: y is large ~Bxy: x is not bigger than y Based on this translation key, we can interpret the proposition as: All apartments are not bigger than some large house. Therefore, the correct interpretation is: A. 'All apartments are not bigger than some large house.' Answer Option 3: To interpret the given proposition $(\forall x)\{Ax\supset (\forall y)[(Hy\bullet Ly)\supset \neg Bxy]\}$, we can break it down as follows: $\forall x$: For all x, Ax: x is an apartment. ⇒: implies. $(\forall y)$: For all y, (Hy • Ly): y is a house and y is large. ⇒: implies. -Bxy: x is not bigger than y. A correct interpretation using the translation key would be: B. 'No apartment is bigger than any large house.' This interpretation captures the meaning of the proposition by stating that for every apartment x, there does not exist a large house y such that x is bigger than y. In other words, no apartment is bigger than any large house. Answer Option 4: Let's analyze the given proposition using the translation key: $(\forall x)\{Ax\supset (\forall y)[(Hy \bullet Ly)\supset \neg Bxy]\}$ The proposition starts with $(\forall x)$, which means "for all x". This indicates that the statement applies to every x in the domain. The next part is $Ax \supset (\forall y)[(Hy \bullet Ly) \supset \neg Bxy].$ Breaking it down further: - Ax represents "x is an apartment". $\text{- } (\forall y) \text{[(Hy \bullet Ly)} \supset \neg \text{Bxy] represents "for all y, if y is a house and y is large, then x is not bigger than y".}$

Answer Option 5:

To interpret the given proposition using the translation key, we can break it down as follows:

Therefore, the correct answer is A. "All apartments are not bigger than some large house."

Putting it all together, the English interpretation of the given proposition is:

"All apartments are not bigger than some large house."

```
(\forall x)\{Ax\supset (\forall y)[(Hy\bullet Ly)\supset -Bxy]\}
Translation key:
Ax: x is an apartment
Hx: x is a house
Lx: x is large
Bxy: x is bigger than y
Breaking down the proposition:
(\forall x)\{Ax\supset (\forall y)[(Hy\bullet Ly)\supset -Bxy]\}
For all x, if x is an apartment (Ax), then for all y, if y is a house (Hy) and y is large (Ly), then it is not the case that x is bigger than y (-Bxy).
Interpreting the proposition in English:
The best English interpretation of the given proposition is:
"All apartments are not bigger than some large house."
Option A: "All apartments are not bigger than some large house." matches the interpretation of the given proposition. Final Answer: A
Answer Option 6:
The given proposition (\forall x)\{Ax\supset (\forall y)[(Hy \bullet Ly)\supset \neg Bxy]\} can be interpreted as follows:
"For every x, if x is an apartment, then for every y, if y is a house and y is large, it is not the case that x is bigger than y."
To find the best English interpretation, we can break down the proposition:
1. (∀x) - For every x
2. Ax - x is an apartment
3. \supset \text{- implies}
4. (∀y) - For every y
5. Hy - y is a house
6. Ly - y is large
7. \supset - implies
8. -Bxy - x is not bigger than y
Based on this breakdown, the best English interpretation of the given proposition would be:
B. "No apartment is bigger than any large house."
This interpretation captures the meaning of the given proposition by stating that there does not exist any apartment that is bigger than any large house. Final Answer: B
Answer Option 7:
To interpret the given proposition, let's break it down step by step:
(\forall x)\{Ax \supset (\forall y)[(Hy \cdot Ly) \supset -Bxy]\}
The quantifier (\forall x) indicates that the proposition is true for all x.
Ax: x is an apartment
Hy: y is a house
Ly: x is large
Bxy: x is bigger than y
The proposition can be translated as:
"For every apartment, if it is the case that every house is large, then it is not the case that the apartment is bigger than any house."
Now, let's analyze the answer choices:
A. 'All apartments are not bigger than some large house.' This interpretation does not match the given proposition because it uses the phrase "some large house" instead of "every house."
B. 'No apartment is bigger than any large house.' This interpretation matches the given proposition exactly as it states that no apartment is bigger than any large house.
C. 'All large houses are not bigger than some apartment.' This interpretation does not match the given proposition because it talks about "some apartment" instead of "every apartment."
D. No large house is bigger than all apartments. This interpretation does not match the given proposition because it switches the roles of houses and apartments and states that no large house is bigger than
all apartments.
Therefore, the best English interpretation of the given proposition is:
B. 'No apartment is bigger than any large house.'
*******
Answer Option 8:
Let's break down the given proposition using the translation key: (\forall x)\{Ax\supset (\forall y)[(Hy\bullet Ly)\supset -Bxy]\}
∀x: For all x
Ax: x is an apartment
```

□: implies
\forall y: For all y
Hy: y is a house
Ly: y is large
~: not
Bxy: x is bigger than y
Now, let's translate the proposition into English:
"For all apartments, if there exists a house that is both large and a house, then the apartment is not bigger than that house."
The best English interpretation of the given proposition is:
'No apartment is bigger than any large house.'
Therefore, the correct answer is B.

Answer Option 9:
Using the translation key, let's break down the given proposition: $(\forall x)\{Ax \supset (\forall y)[(Hy \bullet Ly) \supset -Bxy]\}$
∀x: For all x
Ax: x is an apartment
⊐: implies
∀y: For all y
Hy: y is a house
Ly: y is large
~Bxy: x is not bigger than y
Rearranging the proposition to fit a more natural language structure, we get: "For all apartments, if for all houses, if the house is large, then the apartment is not bigger than the house."
Now let's analyze the answer choices:
Option A: 'All apartments are not bigger than some large house.'
This interpretation does not capture the "for all houses" part. It only states that there is at least one large house that an apartment is not bigger than. This is not an accurate interpretation.
Option B: 'No apartment is bigger than any large house.'
This interpretation accurately captures the meaning of the proposition. It states that for all apartments and all large houses, no apartment is bigger than any large house. This aligns with the provided
translation.
Option C: 'All large houses are not bigger than some apartment.'
This interpretation reverses the roles of apartments and houses. It states that for all large houses, there is at least one apartment that the large house is not bigger than. This is not an accurate interpretation
Option D: 'No large house is bigger than all apartments.'
This interpretation does not capture the "for all houses" part. It only states that there is no large house that is bigger than all apartments. This is not an accurate interpretation.
Based on the above analysis, the best English interpretation of the given proposition is: B. 'No apartment is bigger than any large house.'

Question 24:

*Construct a complete truth table for the following argument. Then, using the truth table, determine whether the argument is valid or invalid. If the argument is invalid, choose an option which presents a counterexample. (There may be other counterexamples as well.)

 $\mathsf{J} \vee \mathsf{L}$

~K/L', A. "Valid', B. 'Invalid. Counterexample when J and K are true and L is false', C. 'Invalid. Counterexample when J is true and K and L are false', D. 'Invalid. Counterexample when K is true and J and L

Answer Option 1:

To construct a complete truth table, we need to consider all possible combinations of truth values for the propositions involved.

Let's create the truth table for the argument:

 $| \ J \ | \ K \ | \ L \ | \ \sim K \ | \ J \equiv \sim K \ | \ J \lor L \ | \ \sim K \ / \ L \ |$ |---|---|------|------| |T|T|T|F|F|T|T| |T|T|F|F|F|T|F| |T|F|T|T|T|T|T||T|F|F|T|T|T|T| |F|T|T|F|T|T|T| |F|T|F|F|T|F|F| |F|F|T|T|F|T|T| |F|F|F|T|T|F|F|

Now, let's analyze the truth table to determine the validity of the argument.

We can see that there are no rows where $J \equiv \kappa K$ is true and $\kappa K / L$ is false. Therefore, the argument is valid.

Final Answer: A

Answer Option 2:

To construct a truth table for the argument, we need to consider all possible combinations of truth values for the propositions J, K, and L.

- J = ~K represents the equivalence between J and the negation of K.
- $J \vee L$ represents the disjunction (OR) between J and L.
- ~K represents the negation of K.

L represents the proposition L.

Using these propositions, we can construct the following truth table:

 $J | K | L | \sim K | J = \sim K | J \vee L | \sim K / L$

TITITIFIFITIT TITIFIFIFITIT TIFITITITIT T|F|F|T|T|T FITITIFITIT F|T|F|F|T|F|F

F|F|T|T|F|T|T

F|F|F|T|T|F|T

Now, let's analyze the truth values of the last column, ${\sim}K\,/\,L.$

In all cases where ~K is true (F in the truth table), ~K / L is also true (T in the truth table).

Therefore, the argument is valid.

Final Answer: A. 'Valid'

Answer Option 3:

To construct a complete truth table, we need to consider all possible combinations of truth values for the variables J, K, and L.

- J = ~K means that J is true when K is false, and J is false when K is true.
- $J \lor L$ represents the logical OR operation between J and L, where the result is true if either J or L (or both) is true.
- ~K represents the negation of K, which is true when K is false, and false when K is true.

Finally, we have to check whether the premises (J \vee L and \sim K) logically imply the conclusion L.

Here is the complete truth table:	
J K L -K J = -K J v L -K / L	
TITITIFIFITIF	
TITIFIFIFITIT	
TIFITITITIT	
TIFIFITITIF	
FITITIFITIF	
FITIFIFIT	
F F T T F T T	
F F F T F F T	
Based on the truth table, we can see that there are cases where the premises (J v L and ~K) are true, but the conclusion L is false. For example, when J is true, K is true, and L is false:	
J=T,K=T,L=F	
Therefore, the argument is invalid. The correct answer is:	
B. Invalid. Counterexample when J and K are true and L is false	
••••••••••	
Answer Option 4:	
To construct a complete truth table for the argument, we need to consider all the possible combinations of truth values for the propositions J, K, and L.	
$J = -K \ J \lor L \ -K \ L$	
TFFF	
TTTT	
FTFT	
FTTF	
Now, let's analyze the argument based on the truth table:	
1. When J is true (T) and K is false (F), J = ~K is True (T) because ~K is also True (T). J ∨ L is True (T) because the disjunction is True (T) when at least one of the propositions is True (T). ~K is False (F).	
Therefore, the argument is valid in this case.	
2. When J is true (T) and K is true (T), J = ~K is False (F) because ~K is False (F). J ∨ L is True (T) because the disjunction is True (T) when at least one of the propositions is True (T). ~K is True (T).	
Therefore, the argument is invalid in this case.	
3. When J is false (F) and K is false (F), J = -K is True (T) because -K is also True (T). J v L is True (T) because the disjunction is True (T) when at least one of the propositions is True (T)K is False (F).	
Therefore, the argument is valid in this case.	
4. When J is false (F) and K is true (T), $J = K$ is False (F) because K is False (F). $J \lor L$ is True (T) because the disjunction is True (T) when at least one of the propositions is True (T).	
Therefore, the argument is invalid in this case.	
Since the argument is invalid for some truth value combinations, we need to find a counterexample. Let's consider Option B as a counterexample:	
Counterexample: When J and K are true (T) and L is false (F). In this case, $J = K$ is False (F), $J \lor L$ is True (T), and K is True (T), which contradicts the conclusion L being false (F).	
Therefore, the correct answer is B. "Invalid. Counterexample when J and K are true and L is false."	
•••••••	
Answer Option 5:	
To construct a complete truth table for the argument, we need to consider all possible combinations of truth values for the propositions J, K, and L.	
J = ~K	
JvL	
-K	
/L	
Let's start by listing all possible combinations of truth values for J, K, and L:	
JKL	
TTT	
TTF	
TFT	
TFF	
FTT	
FTF	
FFT	
FFF	

Now, let's evaluate each statement in the argument for each row of the truth table: For J ≡ ~K: $J \equiv {\sim} K \; (J \; is \; equivalent \; to \; not \; K)$ T ≡ F -> F T ≡ F -> F $T\equiv T -> T$ $T\equiv T -\!\!\!\!> T$ $F \equiv F -\!\!\!> T$ F ≡ F -> T F ≡ T -> F $F\equiv T -\!\!\!\!> F$ For J v L: $\mathsf{J} \vee \mathsf{L} \, (\mathsf{J} \, \, \mathsf{or} \, \, \mathsf{L})$ $\mathsf{T} \vee \mathsf{T} -\!\!\!\!> \mathsf{T}$ T v F -> T $T \lor T \rightarrow T$ T v F -> T $F \lor T \rightarrow T$ F v T -> T F ∨ T -> T $F \vee F \rightarrow F$ For ~K: ~K (not K) F -> T F -> T T -> F T -> F F -> T F -> T T -> F T -> F For / L: L (not L) F -> T T -> F T -> F F -> T F -> T T -> F Now, let's evaluate the argument by comparing the last statement ("/ L") with the previous statements: ~K (not K) and L should have the same truth value. From the truth table, we can see that there is at least one row where ~K and L have different truth values (row 3: ~K is true and L is false). Therefore, the argument is invalid. The correct answer is B. 'Invalid. Counterexample when J and K are true and L is false'.

 $J \equiv \sim K$ means that J and $\sim K$ have the same truth value. Thus, we can write $J \equiv \sim K$ as $J = \sim K$.

To construct a complete truth table, we need to consider all possible combinations of truth values for the given propositions: J, K, and L.

Using this information, we can construct the truth table as follows: |J|K|L|~K|J∨L| |---|---|----| |T|T|T|F|T| |T|T|F|F|T| |T|F|T|T|T| |T|F|F|T|T| |F|T|T|F|T| |F|T|F|F|F| |F|F|T|T|T| |F|F|F|T|F| Now, let's analyze the argument based on the truth table. The argument states the following premises: 1. J ≡ ~K $2.\;J\vee L$ We need to determine whether the conclusion (L) follows from these premises in every row of the truth table. If L is true in every row where all the premises are true, then the argument is valid. Otherwise, if there is at least one row where all the premises are true but the conclusion is false, then the argument is invalid. In the truth table, we can see that there is no row where all the premises $(J \equiv \sim K, \ J \lor L, \sim K)$ are true and the conclusion (L) is false. Therefore, the argument is valid. Final Answer: A. Valid ******* Answer Option 7: To construct a complete truth table for the argument, we need to consider all possible combinations of truth values for the propositions J, K, and L. $J \equiv \sim K$ can be expressed as $(J \land \sim K) \lor (\sim J \land K)$. Using this information, we can construct the truth table as follows: $JKL \sim KJ \equiv \sim KJ \vee L \sim K/L$ TTTFFTT TTEFETT TFTTTT TFFTTTF FTTFTTT FTFFTFT FFTTFTT FFFTFFT Now, let's analyze the truth table to determine whether the argument is valid or invalid. In this case, we can see that whenever the premises $(J \equiv \neg K, J \lor L, and \neg K)$ are true, the conclusion (L) is also true. Therefore, the argument is valid. So, the correct answer is A. 'Valid'. Answer Option 8: $To \ construct \ a \ truth \ table, \ we \ need \ to \ consider \ all \ possible \ combinations \ of \ truth \ values \ for \ the \ propositions \ J, \ K, \ and \ L.$ $J \equiv \sim K$ means that J is equivalent to the negation of K. The argument is as follows: 1. J v L 2. ~K 3. Therefore, L Let's construct the truth table: $J | K | L | \sim K | J = \sim K | J \vee L | \sim K / L$ TITITIFIFITIT

T|T|F|F|F|T|F

T|F|T|T|T|T|TTIFIFITITIF F|T|T|F|T|T F|T|F|F|T|F|T F|F|T|T|F|T|T F|F|F|T|F|F|T Looking at the truth table, we see that whenever ~K is true (T), L is also true (T). Therefore, the argument is valid. Final Answer: A. 'Valid' ****************** Answer Option 9: To construct a complete truth table, we need to consider all possible truth value combinations for the variables involved in the argument: J, K, and L. $J \equiv {\sim} K$ represents the equivalence between J and the negation of K. $J \lor L$ represents the disjunction (logical OR) between J and L. ~K represents the negation of K. L represents the proposition L. We can construct a truth table as follows: | J | K | L | ~K | J ≡ ~K | J ∨ L | |---|---| |T|T|T|F|F|T| |T|T|F|F|F|T| |T|F|T|T|T|T| |T|F|F|T|T|T| |F|T|T|F|T|T||F|T|F|F|T|F|

Now let's evaluate the argument:

Premise 1: $J \equiv \sim K$

|F|F|T|T|F|T| |F|F|F|T|F|F|

Premise 2: J v L

Premise 3: ~K

Conclusion: L

To determine the validity of the argument, we need to check if all rows where the premises are true also have the conclusion true.

In our truth table, there are two rows (row 3 and row 7) where all the premises are true, but the conclusion is false (L is false). Therefore, the argument is invalid.

The correct answer is: C. 'Invalid. Counterexample when J is true and K and L are false'.

Question 25:

'Which of the given formulas of PL is the best symbolization of the following sentence?

Pablo will go to the store if, and only if, his brother drives him and pays for gas.', A. '(P \equiv D) \equiv G', B. '(D \bullet G) \supset P', C. 'P \supset (D \bullet G)', D. 'P \equiv (D \bullet G)'

Answer Option 1:

The best symbolization of the given sentence is D. $P = (D \cdot G)$ '. Explanation: The sentence states that Pablo will go to the store if, and only if, his brother drives him and pays for gas. This can be represented by the logical equivalence (if and only if) operator =. P represents 'Pablo will go to the store', D represents 'his brother drives him', and G represents 'his brother pays for gas'. Therefore, the correct symbolization is $P = (D \cdot G)$. Final Answer: D

Answer Option 2:

The best symbolization of the given sentence would be option D: 'P ≡ (D • G)'. This symbolizes that 'Pablo will go to the store if, and only if, his brother drives him and pays for gas.' The ≡ (equivalence) represents "if and only if," and the (D • G) represents "his brother drives him and pays for gas." The P on the left side of the equivalence symbol represents "Pablo will go to the store." Therefore, option D is the best symbolization. Final Answer: D

Answer Option 3:

The best symbolization of the given sentence would be option D. 'P \equiv (D \bullet G)'.

Explanation:

Pablo will go to the store if, and only if, his brother drives him and pays for gas.

Symbolization:

P: Pablo will go to the store.

D: His brother drives him.

G: His brother pays for gas.

In the given sentence, we have the condition "if, and only if", which is represented by the biconditional symbol (a). The statement "his brother drives him and pays for gas" can be represented by the conjunction (b) of D and G.

Therefore, the best symbolization of the sentence is 'P \equiv (D \bullet G)'. Final Answer: D.

Answer Option 4:

The statement "Pablo will go to the store if, and only if, his brother drives him and pays for gas" can be symbolized as "P = (D • G)", where P represents "Pablo will go to the store", D represents "his brother drives him", and G represents "his brother pays for gas".

Option A: $'(P \equiv D) \equiv G'$ is not the best symbolization as it does not capture the "and" condition between D and G.

 $\text{Option B: } '(D \bullet G) \supset P' \text{ is also not the best symbolization as it does not capture the "if and only if" condition. }$

Option C: 'P \supset (D \bullet G)' is not the best symbolization as it only captures the condition of P being true if D and G are true, but not the reverse condition.

Option D: 'P = (D • G)' is the best symbolization as it properly captures the "if and only if" condition and represents the statement accurately.

Therefore, the correct answer is D.

Answer Option 5:

The given sentence can be symbolized as follows:

P: Pablo will go to the store.

D: His brother drives him.

G: His brother pays for gas.

The sentence states that Pablo will go to the store if, and only if, his brother drives him and pays for gas. This can be symbolized as $P \equiv (D \cdot G)$.

Option A: '(P = D) = G' - This formula does not accurately represent the statement because it implies that Pablo going to the store is equivalent to his brother driving him, which is not the intended meaning.

Option B: '(D • G) \supset P' - This formula states that if his brother drives him and pays for gas, then Pablo will go to the store. However, it does not capture the "if and only if" relationship.

Option C: 'P \supset (D • G)' - This formula states that if Pablo goes to the store, then his brother drives him and pays for gas. It does not capture the "if and only if" relationship.

Option D: 'P = (D • G)' - This formula accurately symbolizes the given sentence, stating that Pablo will go to the store if, and only if, his brother drives him and pays for gas.

Final Answer: D

Answer Option 6:

The best symbolization of the sentence "Pablo will go to the store if, and only if, his brother drives him and pays for gas" is option D. $P \equiv (D \cdot G)$

Explanation:

P: Pablo will go to the store

D: His brother drives him

G: His brother pays for gas

The sentence can be understood as "P is true if and only if (D is true and G is true)". This can be symbolized as 'P \equiv (D \cdot G)'.

Option A: '(P = D) = G' - This symbolization implies that P is true if and only if D is true, and that is equivalent to G. It does not capture the requirement that both D and G need to be true for P to be true.

Option B: '(D • G) \supset P' - This symbolization states that if D and G are both true, then P is true. However, it does not capture the requirement that P is true if and only if D and G are both true.

Option C: "P \supset (D \bullet G)" - This symbolization states that if P is true, then D and G are both true. However, it does not capture the requirement that P is true if and only if D and G are both true.

Therefore, the best symbolization is option D. 'P \equiv (D \bullet G)'.

Answer Option 7:

To symbolize the sentence "Pablo will go to the store if, and only if, his brother drives him and pays for gas", we can break it down into two parts:

1. Pablo will go to the store if his brother drives him and pays for gas.

This can be symbolized as: $P \rightarrow (D \wedge G)$

Explanation: If Pablo's brother drives him and pays for gas (D \wedge G), then Pablo will go to the store (P).

2. Pablo will go to the store only if his brother drives him and pays for gas.

This can be symbolized as: $(D \land G) \rightarrow P$

Explanation: If Pablo's brother drives him and pays for gas (D \wedge G), then and only then Pablo will go to the store (P).

Now, to symbolize the entire sentence, we can combine the two parts using the biconditional operator (=) since it represents "if and only if":

 $(P \to (D \mathrel{\wedge} G)) \equiv ((D \mathrel{\wedge} G) \to P)$

Comparing this with the given options:

A. '(P = D) = G': This does not properly symbolize the sentence, as it implies that Pablo going to the store being equivalent to his brother driving him (P = D), and that being equivalent to G (G). This is not the intended meaning.

B. '(D • G) \supset P': This does not capture the "if and only if" condition, as it only represents the implication from his brother driving him and paying for gas to Pablo going to the store.

C. "P \supset (D \bullet G)": This only represents the implication from Pablo going to the store to his brother driving him and paying for gas, but it does not capture the "if and only if" condition.

D. $P = (D \cdot G)$: This is the best symbolization as it correctly symbolizes the entire sentence, capturing both the implication from Pablo going to the store to his brother driving him and paying for gas, and the implication from his brother driving him and paying for gas to Pablo going to the store.

Therefore, the correct answer is D. $P \equiv (D \cdot G)'$.

Answer Option 8:

To symbolize the sentence "Pablo will go to the store if, and only if, his brother drives him and pays for gas", we can break it down into the following components:

P: Pablo will go to the store

D: His brother drives him

G: His brother pays for gas

The sentence can be rephrased as "Pablo will go to the store if, and only if, his brother drives him and pays for gas." This means that Pablo will go to the store if his brother drives him and pays for gas, and he will not go to the store if his brother does not drive him or does not pay for gas.

Using these components, let's analyze the given options:

Option A: $'(P \equiv D) \equiv G'$

This option states that Pablo going to the store if and only if his brother drives him is equivalent to his brother paying for gas. However, it does not account for the condition that both his brother driving him and paying for gas are required.

Option B: $'(D \bullet G) \supset P'$

This option states that if his brother drives him and pays for gas, then Pablo will go to the store. However, it does not account for the condition that this is the only scenario in which Pablo will go to the store.

Option C: 'P ⊃ (D • G)'

This option states that if Pablo goes to the store, then his brother drives him and pays for gas. This does not capture the "if and only if" condition of the original sentence.

Option D: $P \equiv (D \cdot G)'$

This option states that Pablo will go to the store if and only if his brother drives him and pays for gas. This accurately represents the original sentence, capturing the "if and only if" condition.

Therefore, the best symbolization of the sentence is Option D: $^{\prime}P\equiv (D\bullet G)^{\prime}.$

Answer Option 9:

The best symbolization of the given sentence is: 'Pablo will go to the store if, and only if, his brother drives him and pays for gas.'

Let's break down the sentence:

- "Pablo will go to the store" can be represented by the proposition P.
- "his brother drives him" can be represented by the proposition D.
- "pays for gas" can be represented by the proposition G.

The sentence states that Pablo will go to the store if, and only if, his brother drives him and pays for gas. This can be symbolized as "P if and only if (D and G)".

Option A: $'(P \equiv D) \equiv G'$

This option represents that P is logically equivalent to D, and this equivalence is further logically equivalent to G. However, this does not represent the required condition that D and G are both necessary conditions for P. Therefore, this option is not the best symbolization.

Option B: $'(D \bullet G) \supset P'$

This option represents that if both D and G are true, then P must be true. However, it does not capture the condition that P is true only if D and G are both true. Therefore, this option is not the best symbolization.

Option C: $P \supset (D \cdot G)'$

This option represents that if P is true, then both D and G must be true. However, it does not capture the condition that D and G are both necessary conditions for P to be true. Therefore, this option is not the best symbolization.

Option D: $P \equiv (D \cdot G)'$

This option represents that P is true if and only if both D and G are true. This accurately captures the given condition that Pablo will go to the store if, and only if, his brother drives him and pays for gas.

Therefore, this option is the best symbolization.

Final Answer: D

' Use indirect truth tables to determine whether each set of propositions is consistent. If the set is consistent, choose an option with a consistent valuation. (There may be other consistent valuations.) $D \equiv E$ $F \lor D$ ~E · ~F', A. 'Inconsistent', B. 'Consistent. Consistent valuation when D, E, and F are true', C. 'Consistent. Consistent valuation when D and E are true and F is false', D. 'Consistent. Consistent valuation when D and F are true and E is false! Answer Option 1: To determine the consistency of the set of propositions, we can create a truth table and examine if there exists at least one row where all propositions are true. $D \equiv E$ $F \lor D$ ~E · ~F Creating a truth table for the given propositions: $D E F D \equiv E F \vee D \sim E \cdot \sim F$ TTTTF TTFTTF TETETE TFFTTF FTTFTF FTFFFT FFTFFF FFFTTT From the truth table, we can see that there is no row where all propositions are true. Therefore, the set is inconsistent. Final Answer: A. 'Inconsistent' Answer Option 2: To determine the consistency of the given set of propositions, we can create an indirect truth table: $D \equiv E$ $F \lor D$ ~E · ~F We will create a truth table and check for any row where all propositions are true. $D E F D \equiv E F \lor D \sim E \cdot \sim F$ TTTTF TTETTE TETETE TFFTTT FTTFTF FTFFTF FFTFTF FFFTFT From the truth table, we can see that there is no row where all the propositions are true. Therefore, the set of propositions is inconsistent. Final Answer: A. 'Inconsistent' ******* Answer Option 3: To determine whether the set of propositions is consistent, we can use an indirect truth table. $D \equiv E$ $\mathsf{F} \vee \mathsf{D}$ ~E · ~F To construct the indirect truth table, we assign truth values to D, E, and F and evaluate the given propositions. DEFD = EF v D ~ E · ~ F TTTTTF TTFTTF

Question 26:

TFTFTF

TEFETT FTTFTF FTFFFT FFTFFT FFFTTT From the truth table, we can see that there is at least one row where all the propositions are true. Therefore, the set of propositions is consistent. To find a consistent valuation, we look for a row where all the propositions are true. In the fourth row, D is true, E is false, and F is false. This is a consistent valuation. Therefore, the correct answer is C. "Consistent. Consistent valuation when D and E are true and F is false". Answer Option 4: To determine the consistency of the set of propositions, we need to create an indirect truth table. We need to consider all possible combinations of truth values for D, E, and F and check if there is any row where all the propositions in the set are true. $D E F D \equiv E F \vee D \sim E \cdot \sim F$ TTTTF TTFTTF TETETE TFFFTT FTTFTF FTFFFT FFTTFF After constructing the indirect truth table, we can see that there is no row where all the propositions in the set are true. Therefore, the set of propositions is inconsistent. Final Answer: A. 'Inconsistent' Answer Option 5: To determine whether the set of propositions is consistent, we can create an indirect truth table to evaluate the truth values of the propositions. D≡E $F \lor D$ ~E · ~F Let's create the truth table: $D E F D \equiv E F \lor D \sim E \cdot \sim F$ TTTTF TTETTE TFTFTF TFFFTT FTTFTF FTFFF FFTTTF FFFTTT Based on the truth table, we can see that there is at least one row where all the propositions are true. Therefore, the set of propositions is consistent. However, there are multiple consistent valuations: - Consistent valuation when D, E, and F are true - Consistent valuation when D and E are true and F is false - Consistent valuation when D and F are true and E is false Therefore, the correct answer is C. 'Consistent. Consistent valuation when D and E are true and F is false'. Answer Option 6:

To construct the truth table, we assign truth values to each proposition and evaluate the resulting expression:

To determine whether the set of propositions is consistent, we can use indirect truth tables:

D≡E F∨D ~E·~F

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DEFD = EF v D ~ E · ~ F
TTTTF
TTFTTF
TFTFTT
TFFFTF
FTTFTF
FTFFFT
FETEET
FFFTTT
From the truth table, we can see that there is at least one row where the entire set of propositions is true (T). Therefore, the set of propositions is consistent.
Now, let's determine a consistent valuation:
From the truth table, we can see that for a consistent valuation:
- When D is true, E is also true (D \equiv E)
- When D and F are true, the entire set of propositions is true (F \vee D, ~E \cdot ~F)
Therefore, the consistent valuation is when D and F are true and E is false.
Final Answer: C. Consistent. Consistent valuation when D and E are true and F is false.
Answer Option 7:
To determine whether the set of propositions is consistent, we can construct an indirect truth table by assuming each proposition to be true and evaluating the truth value of the other propositions accordingly.
Assuming D \equiv E to be true, we can replace E with D in the given set of propositions:
D \equiv D
F v D
~D · ~F
Now, let's evaluate the truth value of each proposition based on the assumed truth values:
1. D \equiv D: This is a tautology since any proposition is logically equivalent to itself. It is always true regardless of the truth value of D.
2. F v D: If D is true, then F must also be true for this proposition to be true. If D is false, then F can take any truth value. Therefore, this proposition is consistent when D and F are both true.
3. ~D · ~F: If D is true, then ~D is false and therefore ~D · ~F is false. If D is false, then ~D is true and ~D · ~F is true regardless of the truth value of F. Therefore, this proposition is consistent when D is false.
Based on the above evaluations, the set of propositions is consistent. The consistent valuation options are:
B. Consistent. Consistent valuation when D, E, and F are true
C. Consistent. Consistent valuation when D and E are true and F is false
D. Consistent. Consistent valuation when D and F are true and E is false
Final Answer: B, C, D
******************
Answer Option 8:
To determine the consistency of the set of propositions, we can create an indirect truth table.
D \equiv E \text{ means that } D \text{ and } E \text{ have the same truth value. We can represent this as } D \leftrightarrow E.
Let's create the indirect truth table:
D \leftrightarrow E \mid F \lor D \mid \sim E \cdot \sim F \mid Consistency
T | T | F | Inconsistent
T | T | T | Consistent
T | F | T | Consistent
T | F | F | Consistent
F | T | T | Consistent
F | T | F | Consistent
F | F | F | Consistent
F | F | T | Consistent
From the truth table, we can see that there are consistent valuations for each combination of D, E, and F.
Option C: 'Consistent. Consistent valuation when D and E are true and F is false' is a valid option. (D = T, E = T, F = F)
Option B: 'Consistent, Consistent valuation when D. E. and F are true' is also a valid option. (D = T, E = T, E = T)
Option D: 'Consistent. Consistent valuation when D and F are true and E is false' is another valid option. (D = T, E = F, F = T)
Hence, the set of propositions is consistent. Final Answer: B, C, or D.
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Answer Option 9: To determine the consistency of the set of propositions, we can create an indirect truth table: $\mathsf{D} \equiv \mathsf{E}$ $\mathsf{F} \vee \mathsf{D}$ ~E · ~F First, we assign truth values to each proposition: D|E|F TITIT TITIE T|F|T T|F|F FITIT F|T|F F|F|T F|F|F Next, we evaluate the given propositions based on the truth values assigned: $\mathsf{D}\equiv\mathsf{E}$ Т F F $\mathsf{F} \vee \mathsf{D}$ т т F ~E · ~F F F

Comparing the truth values of all three propositions, we see that there is no row where all three propositions are true. Therefore, the set of propositions is inconsistent.

Final Answer: A. 'Inconsistent'

Question 27:

'Which of the given formulas of PL is the best symbolization of the following sentence?

The restaurant serves pie or cheesecake or ice cream for dessert.', A. 'P \vee (C \vee I)', B. '(P \vee C) \vee I', C. 'P \vee C \vee I', D. 'Options A and B'

Answer Option 1:

The best symbolization of the given sentence would be "P v C v I" which corresponds to option C. This formula accurately represents the statement that the restaurant serves pie or cheesecake or ice cream

Answer Option 2:

The correct symbolization of the given sentence would be 'P v C v I', where P represents pie, C represents cheesecake, and I represents ice cream. This formula accurately represents the statement that the restaurant serves pie or cheesecake or ice cream for dessert. Option C is the correct answer. Final Answer: C

Answer Option 3:

The correct symbolization of the given sentence is 'P v C v I', where P represents pie, C represents cheesecake, and I represents ice cream. This statement accurately represents that the restaurant serves pie or cheesecake or ice cream for dessert. Option C is the correct symbolization. Final Answer: C

Answer Option 4:

The best symbolization of the sentence "The restaurant serves pie or cheesecake or ice cream for dessert" would be:

C. 'P v C v I'

This is because the sentence states that the restaurant serves pie OR cheesecake OR ice cream for dessert, without any specific grouping or ordering. Therefore, the best representation is a simple disjunction (OR) of the three options. Final Answer: C

Answer Option 5:

The best symbolization of the given sentence is 'P v C v I' (Option C). This is because the sentence states that the restaurant serves pie or cheesecake or ice cream for dessert, without any specific grouping or order. The formula 'P v C v I' represents the logical disjunction (OR) of the propositions P (pie), C (cheesecake), and I (ice cream), indicating that any or all of these options may be chosen. Final Answer: C

Answer Option 6:

The best symbolization of the sentence "The restaurant serves pie or cheesecake or ice cream for dessert" would be option C: 'P v C v I'. This formula represents that the restaurant serves pie (P), cheesecake (C), or ice cream (I) for dessert. Option A, 'P v (C v I)', suggests that the restaurant serves either pie or both cheesecake and ice cream. Option B, '(P v C) v I', suggests that the restaurant serves either pie and/or cheesecake, or ice cream. So, option C is the most accurate representation of the given sentence. Final Answer: C

Answer Option 7:

The best symbolization of the given sentence, "The restaurant serves pie or cheesecake or ice cream for dessert." would be: C. 'P v C v I'.

This is because the sentence states that the restaurant serves either pie or cheesecake or ice cream for dessert, without any specific grouping or combination required. Therefore, it can be symbolized as a simple disjunction (OR) of the three options: P (pie), C (cheesecake), and I (ice cream).

 $Option\ A,\ 'P\lor (C\lor I)',\ and\ Option\ B,\ '(P\lor C)\lor I',\ both\ introduce\ unnecessary\ grouping\ and\ do\ not\ accurately\ represent\ the\ given\ sentence.$

So, the correct answer is C. 'P \vee C \vee I'.

Answer Option 8:

The best symbolization of the given sentence would be: C. 'P \vee C \vee I'

This is because the sentence states that the restaurant serves pie OR cheesecake OR ice cream for dessert. The logical OR operator (v) is used to represent the concept of *or* in symbolic logic.

Option A: $'P \lor (C \lor I)'$ represents the restaurant serving pie OR (cheesecake OR ice cream), which is not the intended meaning of the sentence.

Option B: '(P v C) v l' represents the restaurant serving (pie OR cheesecake) OR ice cream, which is not the intended meaning of the sentence.

 $Option \ C \ correctly \ symbolizes \ the \ sentence \ by \ stating \ that \ the \ restaurant \ serves \ pie \ OR \ cheese \ cheese \ OR \ ice \ cream \ for \ dessert.$

Option D is incorrect as it includes both option A and option B, which are not the best symbolizations of the given sentence. Final Answer: C. 'P \vee C \vee I'

Answer Option 9:

The best symbolization of the sentence "The restaurant serves pie or cheesecake or ice cream for dessert" would be the logical formula that accurately represents the options provided.

The sentence states that the restaurant serves either pie or cheesecake or ice cream. This can be symbolized using the logical operator OR (v) to connect the individual propositions representing each dessert option.

Looking at the options:

A. 'P v (C v I)' - This option represents serving pie, then either cheesecake or ice cream. It does not accurately represent the given sentence.

B. '(P v C) v l' - This option represents serving either pie or cheesecake, and then also serving ice cream. It does not accurately represent the given sentence.

 $C. \ 'P \lor C \lor I' - This \ option \ represents \ serving \ pie, \ cheese cake, \ and \ ice \ cream \ separately. \ It \ accurately \ represents \ the \ given \ sentence.$

D. 'Options A and B' - This option is not a valid choice as it combines both options A and B, which we have already determined to be incorrect.

Therefore, the best symbolization of the sentence is option C: 'P \vee C \vee I'. Final Answer: C

Question 28:

'Which of the given formulas of PL is the best symbolization of the following sentence?

If Beth rides her bike, she gets to work earlier.', A. 'B \supset W', B. 'W \supset B', C. 'B \equiv W', D. 'Both A and B'

Answer Option 1:

The best symbolization of the given sentence, "if Beth rides her bike, she gets to work earlier," is option A: "B \supset W. This symbolizes the conditional statement correctly, where B represents "Beth rides her bike" and W represents "she gets to work earlier." The arrow represents the implication or "if...then" relationship between the two propositions. Thus, option A is the best symbolization for the given sentence. Final Answer: A

Answer Option 2:

The best symbolization of the given sentence would be 'B \supset W'. This symbolization represents the logical statement "If Beth rides her bike, she gets to work earlier." Option A: 'B \supset W' represents the correct symbolization. Option B: 'W \supset B' represents the incorrect symbolization and translates to "If she gets to work earlier, then Beth rides her bike." Option C: 'B = W' represents the incorrect symbolization and translates to "Beth rides her bike if and only if she gets to work earlier." Option D: None of the given options are correct. Final Answer: A

Answer Option 3:

The best symbolization of the given sentence would be A. 'B \supset W', which translates to "If Beth rides her bike, then she gets to work earlier." This captures the conditional relationship between Beth riding her bike (B) and her getting to work earlier (W). Option B, 'W \supset B', would mean "If Beth gets to work earlier, then she rides her bike," which is the opposite of the intended meaning. Option C, 'B = W', would mean "Beth rides her bike if and only if she gets to work earlier," suggesting a bidirectional relationship that is not implied by the original sentence. Therefore, the correct answer is A.

Answer Option 4:

The best symbolization of the given sentence would be option A: $^{\prime}B\supset W'$.

This translates to "If Beth rides her bike, then she gets to work earlier."

This captures the conditional relationship between Beth riding her bike (B) and her getting to work earlier (W).

Option B ("W > B") would mean "If she gets to work earlier, then Beth rides her bike," which is the opposite of the intended meaning.

Option C ('B = W') would mean "Beth rides her bike if and only if she gets to work earlier," implying a bidirectional relationship, which is not stated in the original sentence.

Therefore, the correct answer is A: 'B \supset W'.

Answer Option 5:

The given sentence can be symbolized as 'B \supset W', where B represents "Beth rides her bike" and W represents "she gets to work earlier". This accurately captures the conditional relationship between the two statements. Option A: 'B \supset W' represents the correct symbolization of the sentence. Option B: 'W \supset B' represents the converse of the sentence, which is not the same as the original sentence. Option C: 'B \equiv W' represents a biconditional relationship between the two statements, which is not necessary in this case. Option D: 'Both A and B' is not a valid option as only option A is the correct symbolization. Final Answer: A

Allowel. A

Answer Option 6:

The best symbolization for the given sentence 'lf Beth rides her bike, she gets to work earlier' is option A: 'B \supset W'. This translates to "Beth rides her bike implies she gets to work earlier". Option B: 'W \supset B' translates to "She gets to work earlier implies Beth rides her bike", which is the opposite of the original sentence. Option C: 'B \equiv W' translates to "Beth rides her bike if and only if she gets to work earlier", which implies a bidirectional relationship between riding her bike and getting to work earlier, whereas the original sentence implies a unidirectional relationship. So, option C is not the best symbolization. Therefore, the correct answer is option A: 'B \supset W'. Final Answer: A

Answer Option 7:

The best symbolization of the sentence "If Beth rides her bike, she gets to work earlier" is option A: "B \to W", which translates to "If Beth rides her bike (B), then she gets to work earlier (W)." This symbolization accurately captures the conditional relationship between Beth riding her bike and getting to work earlier. Option B: "W \to B" would mean "If she gets to work earlier (W), then Beth rides her bike (B)," which does not accurately represent the original sentence. Option C: "B \to W" would mean "Beth rides her bike (B) if and only if (\to) she gets to work earlier (W)," which suggests a bi-conditional relationship rather than a conditional relationship. Option D: "Both A and B" is incorrect as it does not accurately represent the sentence. Final Answer: A

Answer Option 8:

The best symbolization of the given sentence would be option A: 'B \supset W'.

In this symbolization, B represents the proposition 'Beth rides her bike' and W represents the proposition 'she gets to work earlier'. The arrow (\supset) denotes the implication, which is appropriate for representing the conditional relationship between the two propositions in the sentence.

Option B, 'W \supset B', would represent the opposite relationship, implying that if she gets to work earlier, then she rides her bike, which is not the intended meaning of the sentence.

Option C, 'B = W', represents a biconditional relationship, implying that Beth riding her bike is equivalent to her getting to work earlier. This is not accurately capturing the meaning of the original sentence, which is a conditional statement.

Therefore, the best symbolization is option A: 'B \supset W'. Final Answer: A.

Answer Option 9:

The best symbolization of the given sentence would be option A: 'B \supset W'.

This symbolization represents the logical implication that if Beth rides her bike (B), then she gets to work earlier (W). The arrow (\supset) denotes the conditional relationship between the two propositions.

Option B: 'W \supset B' represents the opposite implication, stating that if she gets to work earlier (W), then she rides her bike (B). This does not accurately represent the original sentence.

Option C: 'B = W' represents a biconditional relationship, stating that Beth rides her bike if and only if she gets to work earlier. This goes beyond the original sentence, as it implies that these two events are always connected in both directions.

Option D: 'Both A and B' is not a valid symbolization and does not represent the original sentence.

Therefore, the best symbolization is option A: 'B \supset W'. Final Answer: A

Question 29:

'Select the best translation into predicate logic: Either Taylor or Eric is a latte boy.', A. 'Lt \lor Le', B. 'Lx \lor Ly', C. 'tL \lor eL', D. 'TI \supset El'

Answer Option 1:

The best translation into predicate logic for the statement "Either Taylor or Eric is a latte boy" would be option A: "Lt v Le". This translates to "Taylor is a latte boy or Eric is a latte boy." Final Answer: A.

Answer Option 2:

The best translation into predicate logic for the given statement, "Either Taylor or Eric is a latte boy", is option B: $'Lx \lor Ly'$.

In this translation:

- Lx represents "x is a latte boy", where x can be any person.
- Ly represents "y is a latte boy", where y can be any person.

By using Lx and Ly, we cover the possibility that either Taylor (represented by x) or Eric (represented by y) is a latte boy.

Final Answer: B

Answer Option 3:

The best translation into predicate logic for the statement "Either Taylor or Eric is a latte boy" is option B: 'Lx \vee Ly'.

In this translation, Lx represents the statement "x is a latte boy" and Ly represents the statement "y is a latte boy". The v symbol denotes logical disjunction, which means "either/or".

So, $'Lx \lor Ly'$ correctly captures the meaning that either Taylor or Eric (represented by the variables x and y) is a latte boy.

Therefore, the answer is B.

Answer Option 4:

The best translation into predicate logic for the statement "Either Taylor or Eric is a latte boy" would be option B: 'Lx v Ly'.

In this translation, "Lx' represents the statement "x is a latte boy", and "Ly' represents the statement "y is a latte boy". By using variables 'x' and 'y', we allow for the possibility of different individuals, in this case, Taylor and Eric, to be latte boys.

So, the correct translation is 'Lx \vee Ly'. Final Answer: B

Answer Option 5:

The best translation into predicate logic for the given statement "Either Taylor or Eric is a latte boy" is option B: $Lx \vee Ly'$.

Explanation

In predicate logic, we represent individuals using variables, and we use predicates to represent properties or relationships. Here, we can represent Taylor as 'x' and Eric as 'y'. The predicate 'L' can represent the property of being a latte boy.

So, the statement "Either Taylor or Eric is a latte boy" can be translated as 'Lx v Ly', which means "Taylor is a latte boy or Eric is a latte boy".

Therefore, the correct answer is option B: 'Lx \vee Ly'.

Answer Option 6:

The best translation into predicate logic for the statement "Either Taylor or Eric is a latte boy" would be option B: $\ Lx \lor Ly'$.

In this translation, "Lx' represents the statement "x is a latte boy," where 'x' can be replaced by any individual. Similarly, "Ly' represents the statement "y is a latte boy," where 'y' can be replaced by any individual.

By using 'Lx' and 'Ly' in the expression 'Lx v Ly', we are saying that either 'x' or 'y' (which can represent Taylor and Eric) is a latte boy. Therefore, option B is the correct translation. Final Answer: B

Answer Option 7:

The best translation into predicate logic for the statement "Either Taylor or Eric is a latte boy" is option B: $^{\prime}$ Lx $^{\prime}$ Ly'.

In this translation, "Lx' represents the statement "x is a latte boy" where 'x' can be replaced by any individual. Similarly, "Ly' represents the statement "y is a latte boy" where 'y' can be replaced by any individual.

By using 'x' to represent Taylor and 'y' to represent Eric, the statement 'Lx \times Ly' correctly represents the statement "Either Taylor or Eric is a latte boy."

Therefore, the correct answer is option $\ensuremath{\mathsf{B}}.$

Answer Option 8:

The best translation into predicate logic for the statement "Either Taylor or Eric is a latte boy" is: B. $Lx \vee Ly'$

Here, Lx represents "x is a latte boy" and Ly represents "y is a latte boy". By using the logical connective "<" (OR), we are expressing the disjunction between the two individuals, Taylor and Eric, being latte

Option A, 'Lt \lor Le', does not properly represent the idea of individuals being latte boys.

Option C, 'tL \vee eL', does not follow the correct order and format for predicate logic.

Option D, TI \supset EI', represents the implication that if Taylor is a latte boy, then Eric is a latte boy, which is not the intended meaning of the original statement.

Answer Option 9:

The statement "Either Taylor or Eric is a latte boy" can be translated into predicate logic as: Lt \vee Le.

Option A: 'Lt v Le' is the correct translation because it represents the logical disjunction (OR) between the predicates Lt (Taylor is a latte boy) and Le (Eric is a latte boy).

Option B: $^{\prime}$ Lx $^{\prime}$ Ly $^{\prime}$ does not specify the individuals Taylor and Eric, and uses variables x and y instead.

Option C: 'tL v eL' is not a correct translation as it does not represent the correct relationship between the individuals and the predicate.

Option D: TI \supset El' represents the implication (IF-THEN) that if Taylor is a latte boy, then Eric is a latte boy. This does not accurately capture the meaning of the original statement.

Therefore, the best translation into predicate logic is A. 'Lt \vee Le'.

Question 30:

"Use indirect truth tables to determine whether the following argument is valid. If the argument is invalid, choose an option which presents a counterexample. (There may be other counterexamples as well.)

 $E\supset F$

G > ~F / ~G v ~E', A. "Valid', B. 'Invalid. Counterexample when E, F, and G are true', C. 'Invalid. Counterexample when E and G are true and F is false', D. 'Invalid. Counterexample when F and G are true and

E is false'

Answer Option 1:

To determine the validity of the argument, we can use an indirect truth table.

Assume E is true and F is false and G is true.

E ⊃ F is false (since E is true and F is false)

 $G \supset {}^{\sim}F$ is true (since G is true and ${}^{\sim}F$ is true)

Therefore, both premises are true, but the conclusion \sim G \vee \sim E is false (since \sim G is false and \sim E is false).

Since we have found a case where the premises are true and the conclusion is false, we have found a counterexample to the argument.

So, the correct answer is:

B. Invalid. Counterexample when E, F, and G are true

Answer Option 2:

To determine the validity of the argument, we will construct an indirect truth table.

E ¬ F

G ⊃ ~F

/ ~G v ~E

We will consider all possible truth value combinations for E, F, and G.

EFG~G∨~E

TTTF

TTFT

TFTT

TFFT

FTTF

FTFT FFTT

FFFT

In all the rows where both premises (E \supset F and G \supset \sim F) are true, the conclusion (\sim G \vee \sim E) is also true. There is no row where both premises are true and the conclusion is false.

Therefore, the argument is valid as there is no counterexample.

Final Answer: A. Valid

Answer Option 3:

To determine the validity of the argument using indirect truth tables, we create a truth table with columns for each premise and the conclusion. We will assume that the premises are true and check if the conclusion is always true as well.

E ⊃ F | G ⊃ ~F | ~G ∨ ~E

TITIT

T|F|T

FITIT

ELTLE

From the truth table, we can see that there is a row where both premises are true (row 1), but the conclusion is false. Therefore, the argument is invalid.

Counterexample: When E is true, F is true, and G is true, the premises are true, but the conclusion (\sim G \vee \sim E) is false.

Final Answer: C. Invalid. Counterexample when E and G are true and F is false.

Answer Option 4:

To determine the validity of the argument using indirect truth tables, we need to consider all possible truth value combinations for the premises and the conclusion.

Let's construct the truth table:

 $\mathsf{E}\,\mathsf{F}\,\mathsf{G}\,\mathsf{\sim}\!\mathsf{G}\,\vee\,\mathsf{\sim}\!\mathsf{E}\,\mathsf{E}\supset\mathsf{F}\,\mathsf{G}\supset\mathsf{\sim}\!\mathsf{F}$

TTTFTF

TTETTT TETEET TFFTFT FTTFTF FTFTTT FFTFTT FFFTTT For the conclusion \sim G \vee \sim E to be true, either \sim G or \sim E (or both) must be true. Looking at the truth table, we can see that in all cases where \sim G \vee \sim E is true, either \sim G or \sim E holds true. Therefore, the conclusion is satisfied in all rows of the truth table. Since the conclusion is always true when the premises are true, the argument is valid. Final Answer: A. Valid Answer Option 5: To determine the validity of the argument using an indirect truth table, we construct a truth table with all the propositions involved. $E \supset F$ is equivalent to $\sim E \lor F$. $G\supset {\sim} F \text{ is equivalent to } {\sim} G \vee {\sim} F.$ We need to check if the conclusion $\sim G \vee \sim E$ is always true when the premises $E \supset F$ and $G \supset \sim F$ are true. Here is the truth table: E F G ~E ~F E ⊃ F G ⊃ ~F ~G ∨ ~E TTTFFTFTT TTFFTTTTT TFTFFFTTT TEFETTTE FTTTFTTFT FTFTTTTFT FFTTFTTFT FFFTTTTFT In the truth table, whenever the premises $E \supset F$ and $G \supset F$ are true, the conclusion $F \subseteq F$ is also true. Therefore, the argument is valid. Answer Option 6: To determine the validity of the argument using indirect truth tables, we will assume the premises are true and check if the conclusion must also be true. Assume: 1. E ⊃ F 2. G ⊃ ¬F We will build the truth table for the premises and the conclusion: $\mathsf{E}\,\mathsf{F}\,\mathsf{G}\,\neg\mathsf{E}\,\neg\mathsf{G}\,\mathsf{E}\supset\mathsf{F}\,\mathsf{G}\supset\neg\mathsf{F}\,\neg\mathsf{G}\vee\neg\mathsf{E}$ TTTFFTFTF TTFFTTTTT TFTFFFTTT TEFETETTT FTTTFTTFT FTFTTTTFT FFTTFTTFF FFFTTTTFT From the truth table, we can see that there is a row where both premises are true (row 2), but the conclusion is false (¬G ∨ ¬E is false). This means that the argument is invalid. When E is true, F is true, and G is false, the premises are true (E \supset F and G $\supset \neg$ F), but the conclusion (\neg G $\lor \neg$ E) is false. Therefore, the answer is: B. Invalid. Counterexample when E, F, and G are true. *******

To determine the validity of the argument using an indirect truth table, we will assume the premises to be true and check if the conclusion is always true.

Answer Option 7:

Let's construct the truth table:
EFG-G-FEDFGD-F-GV-E
TTTFFTFTF
TTFTTTTTT
TETETTEET
TEFTITITI
FTTFFFTTT
FTFTTTFFT
FFTFTTTTT
FFFTTTTTF
From the truth table, we can see that when E is true, F is true, and G is true, the premises are true. However, the conclusion ~G v ~E is false in this case. Therefore, the argument is invalid.
Option C presents a counterexample when E is true, G is true, and F is false. In this case, the premises are true, but the conclusion ~G v ~E is false, which contradicts the argument.
Final Answer: C. 'Invalid. Counterexample when E and G are true and F is false'
Answer Option 8:
To determine the validity of the argument using indirect truth tables, we will construct a truth table to check if the premises imply the conclusion.
E ⊃ F (Premise 1)
$G \supset F$ (Premise 2)
~G ∨ ~E (Conclusion)
We have 3 atomic propositions: E, F, and G. Therefore, we will need 2 ⁴ 3 = 8 rows in our truth table.
$E F G (E \supset F) (G \supset \neg F) (\neg G \lor \neg E)$
TTTTFF
TIFTIT
TETEFF
TEFETT
FTTTFF
FTFTTT
FFTTFF
FFFTTT
From the truth table, we can see that there is at least one row where both premises are true (row 2: T, T, F) but the conclusion is false (-G v -E is false). Therefore, the argument is invalid.
Counterexample: When E is true, F is true, and G is false, the premises are true but the conclusion is false. Therefore, the correct answer is B. 'Invalid. Counterexample when E, F, and G are true'.
Answer Option 9:
To determine the validity of the argument using indirect truth tables, we need to consider all possible truth value combinations for the premises and the conclusion.
Let's create a truth table considering the propositions E, F, G, and the premises E ⊃ F and G ⊃ ¬F, as well as the conclusion ¬G ∨ ¬E.
E F G E>F G>~F ~Gv~E
TITITITIFIF
TITIFITIT
TIFITIFIF
TIFIFIFITIT
FITITIFIT
FITIFITIT
FIFITIFIF
FIFIFITIT
In the last row of the truth table, we see that when E = F = G = F = False, the premises E ⊃ F and G ⊃ ~F are both true, but the conclusion ~G ∨ ~E is false.
Therefore, the argument is invalid.
The correct answer is C. 'Invalid. Counterexample when E and G are true and F is false'.

Question 31: ' Use indirect truth tables to determine whether each set of propositions is consistent. If the set is consistent, choose an option with a consistent valuation. (There may be other consistent valuations.) $\mathsf{A} \supset (\mathsf{\sim}\mathsf{B} \vee \mathsf{C})$ $\mathsf{C}\supset\mathsf{D}$ A ∨ (B · ~D) A = (B \supset D)', A. 'Inconsistent', B. 'Consistent. Consistent valuation when A and D are true and B and C are false', C. 'Consistent. Consistent valuation when A, B, and D are true and C is false', D. 'Consistent. Consistent valuation when C and D are true and A and B are false' Answer Option 1: To determine the consistency of the set of propositions using indirect truth tables, we need to consider all possible truth value combinations for the propositions and check if there is any row where all propositions in the set are true. Using indirect truth tables: $A\supset (\sim B\lor C)\ C\supset D\ A\lor (B\cdot \sim D)\ A\equiv (B\supset D)$ TTTTTTTTTTT TTFTTTFFTFTT TFTTFFTTTTFF TEFTETETTET FTTTTFTTF FTFTTTFFTFTT FFTTFTTTFTFF FFFTFTTTFTFT Based on the truth table, there is no row where all propositions in the set are true. Therefore, the set of propositions is inconsistent. Final Answer: A. 'Inconsistent' ****************** Answer Option 2: To determine the consistency of the given set of propositions, we can construct an indirect truth table. $\mathsf{A}\supset (\mathsf{\sim}\mathsf{B}\vee\mathsf{C})$ $\mathsf{C}\supset\mathsf{D}$ $\mathsf{A} \vee (\mathsf{B} \cdot \mathsf{\sim} \mathsf{D})$ $A \equiv (B \supset D)$ We will consider all possible combinations of truth values for A, B, C, and D and evaluate the truth value of each proposition. $A \ B \ C \ D \ \sim\! B \ \sim\! D \ \sim\! B \ \vee \ C \ A \supset (\sim\! B \ \vee \ C) \ C \supset D \ B \ \cdot \ \sim\! D \ A \ \vee \ (B \ \cdot \ \sim\! D) \ A \equiv (B \supset D)$ TTTTFFTTTFTT TTTFFTTTTFTF TTFTFFTTFFTF TTFFFTTTFTTF TFTTTFTFTTFT TFTFTTTFTFTF TEETTETETE TEFETTTETTTE FTTTFFTFTT FTTFFTTFTF FTFTFFTFTT FTFFFTTFTT FFTTTFTFTTTT FFTFTTTFTTTF FFFTTFTTTT

After evaluating the truth values for each proposition, we can see that there is no row where all propositions are true. Therefore, the set of propositions is inconsistent.

Final Answer: A. 'Inconsistent'

FFFFTTTFTTT

Answer Option 3:

To determine whether the set of propositions is consistent, we can create an indirect truth table by assigning truth values to each proposition and checking if there is a row where all propositions are true.

Let's create the truth table:

FTFTTT

 $A \ B \ C \ D \ \sim B \ \sim D \ \sim B \ \sim C \ A \supset (\sim B \lor C) \ C \supset D \ B \ \sim \sim D \ A \lor (B \lor \sim D) \ A \equiv (B \supset D) \ Consistent?$

TTTTFFTTTFTTNo ${\tt TTTFFTTTFNo}$ TTFTFFFFTFTFNo TTFFTFFTFNo TFTTTFTTFTFNo TFTFTTTTFTTFNo TFFTTFFFTFTNo TFFTTFFTFNo FTTTFFTTTFTFNo FTTFFTTTFNo ${\tt FTFTFFFTTFNo}$ FTFFTFTTTNo ${\tt FFTTTTTTTFTNo}$ FFTFTTTTTFFTNo FFFTTTFTTTFTNo FFFFTTFTTTFTNo None of the rows have all propositions being true. Therefore, the set of propositions is inconsistent. Final Answer: A. 'Inconsistent' Answer Option 4: Let's create an indirect truth table to determine the consistency of the given set of propositions: $\mathsf{A}\supset (\mathsf{\sim}\mathsf{B}\vee\mathsf{C})$ $\mathsf{C}\supset\mathsf{D}$ $\mathsf{A} \vee (\mathsf{B} \cdot \mathsf{\sim} \mathsf{D})$ $\mathsf{A}\equiv (\mathsf{B}\supset \mathsf{D})$ We will assign truth values to each proposition and check for any contradictions. For A \supset (~B \vee C): $\mathsf{A} \sim \!\! \mathsf{B} \; \mathsf{C} \; (\sim \!\! \mathsf{B} \vee \mathsf{C}) \; \mathsf{A} \supset (\sim \!\! \mathsf{B} \vee \mathsf{C})$ TTTTT TTFTT TFTTT TFFFF FTTTT FTFTT FFTTTFFFTT For $C \supset D$: $C\;D\;C\supset D$ $\mathsf{T}\,\mathsf{T}\,\mathsf{T}$ TFF FTT FFT For A \vee (B \cdot ~D): A B D ~D (B · ~D) A \vee (B · ~D) TTTFFT TTFTTT TETEET TFFTFT FTTFFF

```
FFTFFF
FFFTFF
For A \equiv (B \supset D):
A \ B \ D \ (B \supset D) \ A \equiv (B \supset D)
\mathsf{T}\,\mathsf{T}\,\mathsf{T}\,\mathsf{T}\,\mathsf{T}\,\mathsf{T}
TTEFE
TFTTT
TFFTT
 FTTTF
 FTFTF
 FFTFT
 FFFTT
From the truth table, we can see that there are no contradictions. Therefore, the set of propositions is consistent.
 Among the given options, the consistent valuation is when A, B, and D are true and C is false. Therefore, the correct option is C. 'Consistent. Consistent valuation when A, B, and D are true and C is false'.
Final Answer: C
 *******
 Answer Option 5:
 Let's create a truth table to determine the consistency of the given set of propositions:
\mid A \mid B \mid C \mid D \mid A \supset (\sim B \lor C) \mid C \supset D \mid A \lor (B \cdot \sim D) \mid A \equiv (B \supset D) \mid
|T|T|T|F|T|F|T|F|
|T|T|F|T|F|T|T|T|
|T|T|F|F|F|F|T|F|
|T|F|T|F|T|F|T|F|
|T|F|F|T|F|T|T|T|
|T|F|F|F|F|T|T|F|
|F|T|T|T|T|T|F|F|
|F|T|T|F|T|F|T|F|
|F|T|F|T|T|T|F|F|
|F|T|F|F|T|T|F|F|
|F|F|T|T|T|T|F|F|
|F|F|T|F|T|F|F|T|
|F|F|F|T|T|T|F|F|
|F|F|F|F|T|T|F|T|
From the truth table, we can see that there are consistent valuations where the propositions are true. Specifically, when A and D are true and B and C are false. Thus, the correct option is B: 'Consistent.
 Consistent valuation when A and D are true and B and C are false'. Final Answer: B
Answer Option 6:
To determine the consistency of the given set of propositions, we can construct an indirect truth table.
 \mathsf{A} \supset (\mathsf{\sim}\mathsf{B} \lor \mathsf{C})
\mathsf{C}\supset\mathsf{D}
A \vee (B \cdot \sim D)
A\equiv (B\supset D)
We can assign truth values to each proposition and evaluate the truth values of the compound propositions.
 Let's create a truth table considering the truth values of A, B, C, and D:
A \; B \; C \; D \; | \; A \supset (\sim\!B \lor C) \; | \; C \supset D \; | \; A \lor (B \cdot \sim\!D) \; | \; A \equiv (B \supset D)
TTTT|T|T|T
 TTTF|T|F|T|F
TTFT|T|T|T
```

TTFF|T|T|T

```
TFTT|T|T|T
TETELTIFITIT
\mathsf{TFFT} | \mathsf{T} | \mathsf{T} | \mathsf{T} | \mathsf{T} | \mathsf{T}
TEFFITITIT
FTTT|T|T|T
FTTF|T|F|F|F
FTFT|T|T|F|F
FTFF|T|T|F|F
FFTT|T|T|F|F
FFTF|T|F|F|F
FFFT|T|T|F|F
FFFF|T|T|F|F
From the truth table, we can see that there are consistent valuations where A is true, B is false, C is false, and D is true.
Therefore, the set of propositions is consistent.
The correct option would be B. 'Consistent. Consistent valuation when A and D are true and B and C are false'.
Answer Option 7:
To determine whether the set of propositions is consistent, we can construct an indirect truth table by considering all possible truth value combinations for the propositions A, B, C, and D.
\mathsf{A}\supset (\mathsf{\sim}\mathsf{B}\vee\mathsf{C})
\mathsf{C}\supset\mathsf{D}
\mathsf{A} \vee (\mathsf{B} \cdot \mathsf{\sim} \mathsf{D})
\mathsf{A}\equiv (\mathsf{B}\supset \mathsf{D})
We will analyze the set of propositions by considering all possible truth value combinations for A, B, C, and D.
A \ B \ C \ D \ \sim\! B \ \sim\! B \ \vee \ C \ A \supset (\sim\! B \ \vee \ C) \ B \cdot \sim\! D \ A \ \vee \ (B \cdot \sim\! D) \ B \supset D \ A \equiv (B \supset D)
TTTTFTTFTT
TTTFFTTFTF
TTFTFFFFF
TTFFFFFFF
TFTTTTTTT
TETETTTTTE
TFFTTFFFF
TFFFTFFFT
FTTTFTTFTF
FTTFFTTFTT
FTFTFFTF
FTFFFFTFTT
FFTTTTTFTT
FFTFTTTFTT
FFFTTFTT
FFFFTFTT
From the truth table, we can see that there is at least one row where all propositions are true. Therefore, the set of propositions is consistent.
Option C: 'Consistent. Consistent valuation when A, B, and D are true and C is false' is a valid answer as it satisfies the consistency condition.
Final Answer: C. 'Consistent. Consistent valuation when A, B, and D are true and C is false'
To determine the consistency of the set of propositions using indirect truth tables, we will create a table and evaluate the propositions under all possible truth value combinations.
We have 4 propositions: A, B, C, D.
1. A \supset (\sim B \lor C)
2.\;C\supset D
3. A ∨ (B · ~D)
4. A ≡ (B ⊃ D)
Let's start by creating the truth table:
```

 $\mid A\mid B\mid C\mid D\mid \sim B\vee C\mid A\supset (\sim B\vee C)\mid C\supset D\mid B\cdot \sim D\mid A\vee (B\cdot \sim D)\mid B\supset D\mid A\equiv (B\supset D)\mid A=(B)$

|--|--|--|---|------| |T|T|T|T|T|T|F|T|T|T $|\mathsf{T}|\mathsf{T}|\mathsf{T}|\mathsf{F}|\mathsf{T}|\mathsf{T}|\mathsf{F}|\mathsf{F}|\mathsf{T}$ |T|T|F|T|T|T|F|T|T|T |T|T|F|F|F|F|T|T|T|T| |T|F|T|T|T|T|F|T|T|T |T|F|T|F|T|T|F|T|T|F|F||T|F|F|T|F|F|T|F|F|T|F||T|F|F|F|F|F|T|T|T|T| |F|T|T|T|T|T|F|T|T|F||F|T|T|F|T|T|F|T|T|F|F| |F|T|F|T|T|T|F|F|T|F||F|T|F|F|F|T|T|T|T|T| |F|F|T|T|T|T|F|T|T|F| |F|F|T|F|T|T|F|T|T|F|F| |F|F|F|T|T|T|F|F|T|F| |F|F|F|F|T|T|T|T|T|T| Now, let's analyze the table to determine the consistency:

From the table, we can see that there is no row where all propositions are false. Therefore, the set of propositions is consistent.

Among the given options, the consistent valuation when A and D are true and B and C are false matches the truth table.

Final Answer: B. 'Consistent. Consistent valuation when A and D are true and B and C are false'

Answer Option 9:

To determine the consistency of the given set of propositions, we can use an indirect truth table.

We will consider all possible truth values for the propositions A, B, C, and D and check if there is any row where all the propositions in the set are true.

 $\mathsf{A}\supset (\mathsf{\sim}\mathsf{B}\vee\mathsf{C})\;\mathsf{C}\supset\mathsf{D}\;\mathsf{A}\vee(\mathsf{B}\cdot\mathsf{\sim}\mathsf{D})\;\mathsf{A}\equiv(\mathsf{B}\supset\mathsf{D})$

 ${\tt TTTTTTTTTT}$

TTTTTTFTFFT

TTTTTFTTFTF

TTTTTFTFTTFT

TFFTTTTTFTTF

TFFTTTTFFFFT

TFFTTFTTFFTF

TFFTTFTFFTFT

FTTTTTTTTF

FTTTTTTFTFF

FTTTTFTTT

FTTTTFTFTTFF

FFTTTTTTTF

FFTTTTTFTFF

FFTTTFTTT

FFTTTFTFTTFF

TTTTFTTTTTT

TTTTFTTFTFFT

TTTTFFTTTFTF TTTTFFTFTTFT

TFFTFTTTFTTF

TEETETTEEET

TEETEETTEETE

TEETEETEET

FTTTFTTTTTF

FTTTFTTFFF

FTTTFFTTTFTT

FTTTFFTTTFF
FFTTFFTTTFFF
FFTTFFTFFTFFF

After examining all the possible truth value combinations, we find that there is no row where all the propositions in the set are true. Therefore, the set of propositions is inconsistent.

Final Answer: A. 'Inconsistent'

Question 32:
'Construct a complete truth table for the following argument. Then, using the truth table, determine whether the argument is valid or invalid. If the argument is invalid, choose an option which presents a
counterexample. (There may be other counterexamples as well.)
$M \times N$
~M · O / N', A. 'Valid', B. 'Invalid. Counterexample when M and O are true and N is false', C. 'Invalid. Counterexample when M is true and O and N are false', D. 'Invalid. Counterexample when O is true and M
and N are false'
Answer Option 1:
To construct a complete truth table for the argument, we need to consider all possible combinations of truth values for the propositions M, N, and O.
$MNO \sim MM \vee N \sim M \cdot O/N$
TITEITT
TTFFTFF
TFTFTTF
TFFFTFF
FTTTTTT
FTFTTFT
FFTTFTF
FFFTFFT
From the truth table, we can see that there are no rows where both premises (M v N) and (-M · O) are true and the conclusion N is false. Thus, the argument is valid.
The correct answer is A. 'Valid'.
Answer Option 2:
To construct a complete truth table for the argument, we need to consider all possible combinations of truth values for the propositional variables M, N, and O.
$MNOM \lor N \sim M \cdot O/N$
TTTTFT
TTFTFT
TETTET
TFFTFT
FTTTTT
FTFTTT
FFTFTT
FFFFTF
From the truth table, we can see that there is no row where both premises, M v N and ~M · O, are true and the conclusion, N, is false. Therefore, the argument is valid.
Final Answer: A. Valid
Answer Option 3:
To construct a complete truth table for the argument, we need to consider all possible combinations of truth values for the propositions M, N, and O.
$MNOM \lor N \sim M \cdot O/N$
TTTTFTT
TTFTFTT
TFTTFFF
TFFTFFF
FTTTTTT
FTFTTTT
FFTFTFT
FFFFTFT
From the truth table, we can see that there is no row where both premises M v N and ~M · O are true while the conclusion N is false. Therefore, the argument is valid.
Final Answer: A. 'Valid'

To construct a complete truth table for the argument, we need to consider all possible combinations of truth values for the variables M, N, and O.

M N O M \vee N \sim M \cdot O / N

Answer Option 4:

TITIFT
TTETET
TFTTTF
TFFTTF
FTTTTT
FTFTFT
FFTFTF
FFFFFT
Based on the truth table, we can see that there are no rows where both the premises (M ∨ N) and (-M · O) are true and the conclusion N is false.
Therefore, the argument is valid.
Final Answer: A. Valid
Answer Option 5:
To construct a complete truth table for the argument, we need to consider all possible combinations of truth values for the propositions M, N, and O.
M N O ~M M ∨ N ~M · O N'
TITETIT
TTFFTFT
TETETEE
TFFFTFF
FITITI
FTFTTFT
FFTTFTF
FFFTFTF
The argument is invalid because there is a row in the truth table where the premises (M v N and ~M · O) are true, but the conclusion (N') is false. In particular, in the third row, both M v N and ~M · O are true,
but N' is false.
Therefore, the correct answer is B. "Invalid. Counterexample when M and O are true and N is false."

Answer Option 6:
Answer Option 6: To construct a truth table for the argument, we need to consider all possible combinations of truth values for the propositions M, N, and O.
To construct a truth table for the argument, we need to consider all possible combinations of truth values for the propositions M, N, and O.
To construct a truth table for the argument, we need to consider all possible combinations of truth values for the propositions M, N, and O. M N O M × N -M · O / N
To construct a truth table for the argument, we need to consider all possible combinations of truth values for the propositions M, N, and O. M N O M v N -M · O / N T T T T F T
To construct a truth table for the argument, we need to consider all possible combinations of truth values for the propositions M, N, and O. M N O M ∨ N -M · O / N T T T T F T T T F T F T
To construct a truth table for the argument, we need to consider all possible combinations of truth values for the propositions M, N, and O. M N O M V N -M · O / N TTTTFT TFTTFF
To construct a truth table for the argument, we need to consider all possible combinations of truth values for the propositions M, N, and O. M N O M V N -M · O / N TTTFFT TFFTFF TFFTFF
To construct a truth table for the argument, we need to consider all possible combinations of truth values for the propositions M, N, and O. M N O M V N -M · O / N TTTTFT TTFTFF TFTTFF TFTTFF TFTTFF TFTTTTT
To construct a truth table for the argument, we need to consider all possible combinations of truth values for the propositions M, N, and O. M N O M × N -M · O / N TITTET TETTEE TETTEE TETTEE TETTEE TETTET
To construct a truth table for the argument, we need to consider all possible combinations of truth values for the propositions M, N, and O. M N O M × N -M · O / N TTTTFT TTFTFF TFFTFF FTFTFF FTFTFF FTFTFF FTFFFF
To construct a truth table for the argument, we need to consider all possible combinations of truth values for the propositions M, N, and O. M N O M V N - M · O / N TITTET TIFTET TFTTEF TFTTEF FFTTEF FFTFFF FFFFFF FFFFFF FFFFFFF
To construct a truth table for the argument, we need to consider all possible combinations of truth values for the propositions M, N, and O. M N O M × N - M · O / N TITTFT TIFTFT TFTFFF FFTFFF FFTFFF FFFFFF FFFFFF
To construct a truth table for the argument, we need to consider all possible combinations of truth values for the propositions M, N, and O. M N O M × N -M · O / N TTTFT TTFTFT TFTFFF TFTTFF FFTFFF FFTFFF FFFFFF
To construct a truth table for the argument, we need to consider all possible combinations of truth values for the propositions M, N, and O. M N O M > N - M · O / N TITTFT TIFTFT TFTTFF TFTTFF FFTFFF FFTFFF FFFFFF
To construct a truth table for the argument, we need to consider all possible combinations of truth values for the propositions M, N, and O. M N O M \ N - M \ O / N TTTTFT TTFTFT TFTFF FFTFFF FFTFFF FFTFFF Looking at the truth table, we can see that there are cases where both premises are true (M \ N \ N \ and \ -M \ O), but the conclusion (N) is false. Specifically, when M is true, O is true, and N is false. Therefore, the argument is invalid. The counterexample is when M is true, O is true, and N is false. Hence, the correct answer is B. 'Invalid. Counterexample when M and O are true and N is false'.
To construct a truth table for the argument, we need to consider all possible combinations of truth values for the propositions M, N, and O. M N O M V N - M · O / N TTTTFT TTFTFT TFTTFF TFTTFF FFTFFF FFTFFF FFFFFF
To construct a truth table for the argument, we need to consider all possible combinations of truth values for the propositions M, N, and O. M N O M > N - M · O / N TITTET TETTET TETTEF TETTEF FTTTTT FTFFF FTFFF Looking at the truth table, we can see that there are cases where both premises are true (M > N and -M · O), but the conclusion (N) is false. Specifically, when M is true, O is true, and N is false. Therefore, the argument is invalid. The counterexample is when M is true, O is true, and N is false. Hence, the correct answer is B. 'Invalid. Counterexample when M and O are true and N is false'.
To construct a truth table for the argument, we need to consider all possible combinations of truth values for the propositions M, N, and O. M N O M V N - M · O / N TITIFF TIFIFF TIFIFF FITTIT FIFTF FFFFF Looking at the truth table, we can see that there are cases where both premises are true (M V N and -M · O), but the conclusion (N) is false. Specifically, when M is true, O is true, and N is false. Therefore, the argument is invalid. The counterexample is when M is true, O is true, and N is false. Hence, the correct answer is B. 'Invalid. Counterexample when M and O are true and N is false'.
To construct a truth table for the argument, we need to consider all possible combinations of truth values for the propositions M, N, and O. M N O M v N - M · O / N TTTTFT TFTTFF TFTTFF FFTFFF FFTFFF FFTFFF Looking at the truth table, we can see that there are cases where both premises are true (M v N and - M · O), but the conclusion (N) is false. Specifically, when M is true, O is true, and N is false. Therefore, the argument is invalid. The counterexample is when M is true, O is true, and N is false. Hence, the correct answer is B. 'Invalid. Counterexample when M and O are true and N is false'.
To construct a truth table for the argument, we need to consider all possible combinations of truth values for the propositions M, N, and O. M N O M V N - M · O / N TTTTFF TTFTFF TFTTFF TFTTFF FFTFFF FFTFFF FFTFFF Locking at the truth table, we can see that there are cases where both premises are true (M v N and -M · O), but the conclusion (N) is false. Specifically, when M is true, O is true, and N is false. The counterexample is when M is true, O is true, and N is false. Hence, the correct answer is B. 'Invalid'. Counterexample when M and O are true and N is false'.
To construct a truth table for the argument, we need to consider all possible combinations of truth values for the propositions M, N, and O. M N O M V N -M · O / N TTTTFT TTFTFT TFTTFF TFTTFF TFTTFF FFTFFF FFTFFF FFTFFF Looking at the truth table, we can see that there are cases where both premises are true (M ∨ N and -M · O), but the conclusion (N) is false. Specifically, when M is true, O is true, and N is false. Therefore, the argument is invalid. The counterexample is when M is true, O is true, and N is false. Hence, the correct answer is B. 'Invalid. Counterexample when M and O are true and N is false'.
To construct a truth table for the argument, we need to consider all possible combinations of truth values for the propositions M, N, and O. M N O M × N -M · O / N T T T T F T F T F T T F F T F T F F F T T T T
To construct a truth table for the argument, we need to consider all possible combinations of truth values for the propositions M, N, and O. M N O M × N -M · O / N T T T T F T T F T T F F T F T T F F F T F T
To construct a truth table for the argument, we need to consider all possible combinations of truth values for the propositions M, N, and O. M N O M × N -M · O / N T T T T F T F T F T T F F T F T F F F T T T T

TITITITIFIT T|T|F|T|F|T TIFITITIF T|F|F|T|T|F FITITITIT FITIFITIT F|F|T|F|T|F F|F|F|F|T|F From the truth table, we can see that whenever the premises are true, the conclusion is also true. Therefore, the argument is valid. Final Answer: A. 'Valid' Answer Option 8: To construct a complete truth table for the argument, we need to consider all possible combinations of truth values for the propositional variables M, N, and O. Let's represent the given premises and conclusion as propositions: Premise 1: $M \vee N$ Premise 2: ~M · O We will evaluate the truth values of each proposition for all possible combinations of truth values for M, N, and O. M N O M \vee N \sim M \cdot O N TTTTFT TTETET TFTTFF TFFTFF FTTTTT FTFTFT FFTFTF FFFFTF Based on the truth table, we can see that there is no row where both premises are true and the conclusion is false. Therefore, the argument is valid. Final Answer: A. Valid Answer Option 9: To construct a complete truth table for the argument, we need to consider all possible combinations of truth values for the propositions M, N, and O. $M \vee N$ represents the disjunction (logical OR) of M and N. ${\sim}M\cdot O$ represents the conjunction (logical AND) of the negation of M and O. We can construct the truth table as follows: | M | N | O | M ∨ N | ~M · O | N' | |---|---|-----| |T|T|T|T|T|F||T|T|F|T|F|F| |T|F|T|T|T|T| |T|F|F|T|F|T| |F|T|T|T|T|F| |F|T|F|T|F|F| |F|F|T|F|T|T| |F|F|F|F|F|T| Based on the truth table, we can determine whether the argument is valid or invalid. In this case, we find that there is no row where both premises (M ∨ N and ~M · O) are true but the conclusion (N') is false. This means that the argument is valid.

Therefore, the correct answer is A. 'Valid'.

Question 33: 'Which of the following propositions is an immediate (one-step) consequence in PL of the given premises? $\mathsf{U}\supset\mathsf{W}$ $W\supset (X\equiv \sim Y)$ $(X \equiv \sim Y) \supset Z$ ~Z', A. 'U \supset Z', B. 'U \vee W', C. '~W', D. '~(X \equiv ~Y)' Answer Option 1: We can use the method of logical inference to determine the immediate (one-step) consequence of the given premises. Given premises: $1.\;U\supset W$ 2. W \supset (X \equiv \sim Y) 3. (X ≡ ~Y) ⊃ Z 4. ~Z To find the immediate consequence, we can apply modus tollens to premise 3 and premise 4: From premise 3: $(X \equiv \sim Y) \supset Z$ From premise 4: ~Z Using modus tollens: \sim (X = \sim Y) Therefore, the immediate consequence in PL of the given premises is ' \sim (X \equiv \sim Y)'. Therefore, the correct answer is D. Answer Option 2: To determine the immediate consequence, we need to simplify the premises using logical equivalences and inference rules. Premises: $1.\;U\supset W$ 2. W \supset (X \equiv ~Y) 3. (X ≡ ~Y) ⊃ Z 4. ~Z Using the premise 1 (U \supset W) and modus ponens, we can infer W. Using premise 2 (W \supset (X \equiv ~Y)) and modus ponens, we can infer (X \equiv ~Y). Using premise 3 ((X \equiv ~Y) \supset Z) and modus ponens, we can infer Z. Now, we have W and Z as inferred propositions. Since we have \sim Z as premise 4, we can use modus tollens to infer \sim ((X \equiv \sim Y) \supset Z). This is equivalent to \sim (X \equiv \sim Y) \supset \sim Z. The immediate consequence in PL of the given premises is therefore D. ' \sim (X \equiv \sim Y)'. Answer Option 3: To determine the immediate (one-step) consequence in propositional logic (PL) of the given premises, we need to apply the valid inference rules. Premises: $1.\;U\supset W$ 2. W ⊃ (X ≡ ~Y) 3. $(X \equiv \sim Y) \supset Z$ 4. ~Z By applying the modus ponens rule (if P implies Q and P is true, then Q must be true), we can conclude the following: From premise 1 (U \supset W) and premise 4 (~Z), we can infer ~W using modus tollens. From premise 3 ((X \equiv ~Y) \supset Z) and premise 4 (~Z), we can infer ~(X \equiv ~Y) using modus tollens. None of the given options corresponds to the immediate consequence inferred from the premises. Hence, none of the options A, B, C, or D is correct. No answer can be chosen from the given options. ********** Answer Option 4: To determine the immediate consequence of the given premises, we can simplify the premises step by step using logical equivalences. Given premises:

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4. ~Z
First, let's rewrite the implication in premise 1 using the logical equivalence "p \rightarrow q \equiv ~p \vee q":
1. ~U ∨ W
Next, we can apply the transitive property of implications to premises 1 and 2:
\simU \vee (X \equiv \simY)
Then, using the logical equivalence "p \rightarrow q \equiv ~p \vee q" on premise 3:
Now, using the logical equivalence "p \rightarrow q = ~p \vee q" on premise 4:
Z \to \text{False}
Finally, we can apply the contrapositive of an implication to the previous step:
{\sim}\mathsf{False} \to {\sim}\mathsf{Z}
True \to {\sim} Z
~Z
Therefore, the immediate consequence in PL of the given premises is '~Z'. Final Answer: D
Answer Option 5:
To determine the immediate consequence in PL, we can apply the Modus Ponens rule, which states that if we have premises of the form P and P \rightarrow Q, we can infer Q.
1. U \supset W
2. W \supset (X \equiv ~Y)
3. (X \equiv \sim Y) \supset Z
We can use Modus Ponens with premises 1 and 2 to infer (X \equiv ~Y):
5. X ≡ ~Y (from 1 and 2)
Now, we can use Modus Ponens with premises 3 and 5 to infer Z:
6. Z (from 3 and 5)
Finally, since we have ~Z as a premise, we can use Modus Tollens to infer ~X \equiv Y:
7. \sim X \equiv Y \text{ (from 4 and 6)}
Therefore, the immediate consequence in PL of the given premises is '~(X \equiv ~Y)'.
Final Answer: D. '\sim(X \equiv \simY)'
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Answer Option 6:
We can use the rules of implication and equivalence to determine the immediate consequences of the given premises.
Premise 1: U \supset W
Premise 2: W \supset (X \equiv \sim Y)
Premise 3: (X \equiv \sim Y) \supset Z
Premise 4: ~Z
To find an immediate consequence, we need to find a proposition that can be derived directly from the given premises without any additional steps.
From Premise 1 (U \supset W) and Premise 3 ((X \equiv \simY) \supset Z), we can use the rule of transitivity to conclude U \supset Z. This means that if U is true, then Z must also be true. Therefore, the immediate consequence is 'U
Option A: U \supset Z' is the correct immediate consequence.
Option B: ^{\prime}U\vee W^{\prime} cannot be derived directly from the given premises.
Option C: '~W' cannot be derived directly from the given premises.
Option D: '\sim(X\equiv\sim Y)' cannot be derived directly from the given premises.
Final Answer: A
********
Answer Option 7:
To find the immediate consequence in propositional logic (PL), we can apply the rules of implication and equivalence to the given premises.
Given premises:
1.\;U\supset W
2. W ⊃ (X ≡ ~Y)
3. (X \equiv \sim Y) \supset Z
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4. ~Z

Let's analyze the premises:

From premise 4, ~Z, we can infer that Z is false.

Using premise 3, $(X = \neg Y) \supset Z$, and the fact that Z is false, we can conclude that $(X = \neg Y)$ must also be false. This means that X and $\neg Y$ cannot have the same truth value.

Now, using premise 2, $W \supset (X \equiv \neg Y)$, and the fact that $(X \equiv \neg Y)$ is false, we can deduce that W must also be false. This follows from the rule that a true statement implies anything, but a false statement implies nothing

Finally, using premise 1, U \supset W, and the fact that W is false, we can conclude that U must also be false. Again, this follows from the rule that a true statement implies anything, but a false statement implies nothing.

Therefore, the immediate consequence in PL of the given premises is '~U' or 'U is false'.

The correct option is D. ' \sim (X \equiv \sim Y)'.

Answer Option 8:

To determine which proposition is an immediate consequence of the given premises, we can analyze each option based on the premises.

Option A: $U \supset Z'$

This proposition cannot be derived directly from the given premises because there is no direct implication involving U and Z.

Option B: 'U v W'

This proposition also cannot be derived directly from the given premises. The premises do not provide any information about the relationship between U and W.

Option C: '~W'

This proposition can be derived directly from the given premises. From the first premise 'U \supset W' and the second premise 'W \supset (X = \sim Y)', we can infer that 'U \supset (X = \sim Y)' using the transitive property of implication. Then, using modus tollens with the third premise '(X = \sim Y) \supset Z' and the fourth premise ' \sim Z', we can conclude ' \sim W'.

Option D: $'\sim(X \equiv \sim Y)$

This proposition cannot be derived directly from the given premises. The premises only provide information about the implications involving U, W, X, Y, and Z, but not about the equivalence between X and

Therefore, the proposition that is an immediate consequence of the given premises is C. '~W'. Final Answer: C

Answer Option 9:

To find the immediate consequence in propositional logic (PL) of the given premises, we can use the process of Modus Ponens. Modus Ponens states that if we have a premise of the form 'A \rightarrow B' and another premise of 'A', then we can infer 'B' as the immediate consequence. Let's analyze the given premises:

Premise 1: $U \supset W$

Premise 2: $W \supset (X \equiv \sim Y)$

Premise 3: $(X \equiv \sim Y) \supset Z$

Premise 4: ~Z

From Premise 1 (U \supset W) and Premise 2 (W \supset (X \equiv ~Y)), we can apply Modus Ponens to infer (X \equiv ~Y):

(X ≡ ~Y) [from Premise 2]

W [from Premise 1]

Therefore, $(X \equiv \sim Y) \supset Z$ becomes:

Z [from Modus Ponens]

Now, we have Z as a premise. From Premise 3 ((X \equiv ~Y) \supset Z) and Premise 4 (~Z), we can once again apply Modus Ponens to obtain ~(X \equiv ~Y):

 \sim (X \equiv \sim Y) [from Premise 3]

~Z [from Premise 4]

Therefore, the immediate consequence in PL of the given premises is D. ' \sim (X \equiv \sim Y)'.

Question 3	34
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'Identify the antecedent of the following conditional proposition: If the Bees don't win their first game, then both the Aardvarks and the Chipmunks win their first games.', A. 'The Aardvarks win their first games.',

B. The Bees don't win their first game.', C. The Chipmunks win their first game.', D. 'Both the Aardvarks and the Chipmunks win their first games.'

Answer Option 1:

The antecedent of a conditional proposition is the statement that comes after the "if" part. In this case, the antecedent is "The Bees don't win their first game." Therefore, the correct answer is B. The Bees don't win their first game."

Answer Option 2:

The antecedent of a conditional proposition is the part that comes immediately after the "if" statement. In this case, the antecedent is "The Bees don't win their first game." Therefore, the correct answer is B.

The Bees don't win their first game.'

Answer Option 3:

The antecedent of a conditional proposition is the statement that follows the "if" part. In this case, the antecedent is "The Bees don't win their first game.' Therefore, the correct answer is B. The Bees don't win their first game.'

Answer Option 4:

The antecedent of a conditional proposition is the part that comes before the "if" statement. In this case, the antecedent is "The Bees don't win their first game." Therefore, the correct answer is B. The Bees don't win their first game.'

Answer Option 5:

The antecedent of a conditional proposition is the statement that comes before the "if" in the conditional statement. In this case, the antecedent is "The Bees don't win their first game." Therefore, the correct answer is B. The Bees don't win their first game.'

Answer Option 6:

The antecedent of a conditional proposition is the statement that comes immediately after the "if" in the sentence. In this case, the antecedent is 'The Bees don't win their first game.' Therefore, the correct answer is B. 'The Bees don't win their first game.'

Answer Option 7:

The antecedent of a conditional proposition is the statement that appears immediately after the "If" keyword. In this case, the antecedent is "The Bees don't win their first game." Therefore, the correct answer is B. The Bees don't win their first game.'

Answer Option 8:

The antecedent of a conditional proposition is the part of the statement that comes after "if." In this case, the antecedent is "The Bees don't win their first game." Therefore, the correct answer is B. The Bees don't win their first game.'

Answer Option 9:

The antecedent of a conditional proposition is the part that comes before the "then" statement. In this case, the conditional proposition is "If the Bees don't win their first game, then both the Aardvarks and the Chipmunks win their first games."

The antecedent is "The Bees don't win their first game."

Therefore, the correct answer is B. 'The Bees don't win their first game.'

Question 35:
'Select the best translation into predicate logic. Holly loves Frances; (f: Frances; h: Holly; Lxy: x loves y)', A. 'hLf', B. 'fLh', C. 'Lfh', D. 'Lhf'
Answer Option 1:
The correct translation into predicate logic would be: C. 'Lfh'
Answer Option 2:
The best translation into predicate logic for the statement "Holly loves Frances" is "Lhf". This means that Holly loves Frances, where "Lxy" represents x loves y. Final Answer: D
Answer Option 3:
The best translation of the statement "Holly loves Frances" into predicate logic is: D. 'Lhf'. This translates to "There exists an individual h (Holly) who loves an individual f (Frances)."
Answer Option 4:
The best translation into predicate logic for the statement "Holly loves Frances" is: D. 'Lhf'
Here, Lxy represents "x loves y", so 'Lhf' translates to "Holly loves Frances".
Answer Option 5:
The best translation into predicate logic for the statement "Holly loves Frances" would be:
Lhf
This translates to "There exists an x such that Holly loves Frances."
Therefore, the correct option is D.
Answer Option 6:
The best translation into predicate logic for the statement "Holly loves Frances" is:
Lhf
Here, Lxy represents "x loves y", so Lhf represents "Holly loves Frances".
Final Answer: D
Answer Option 7:
The best translation into predicate logic for the statement "Holly loves Frances" using the given symbols would be: "Lhf". This represents the predicate "L" (loves) with "h' representing Holly and "f" representing
Frances. Final Answer: D.
Answer Option 8:
The best translation into predicate logic for the statement "Holly loves Frances" can be represented as 'Lhf', where L represents the predicate "loves", h represents Holly, and f represents Frances. This
translates to "Holly loves Frances" in predicate logic. Final Answer: D. 'Lhf'
Answer Option 9:
The best translation into predicate logic for the statement "Holly loves Frances" is:
Lhf
In this translation, L represents the "loves" relationship, h represents Holly, and f represents Frances. The order in which the variables appear indicates the subject (Holly) and object (Frances) of the
sentence. Therefore, the correct translation is option D: 'Lhf'.

Question 36: "Use indirect truth tables to determine whether the following argument is valid. If the argument is invalid, choose an option which presents a counterexample. (There may be other counterexamples as well.) $P \supset (Q \vee R)$ $(Q \cdot S) \supset \sim P$ ~(~Q v R) / ~P', A. "Valid', B. 'Invalid. Counterexample when P, Q, R, and S are true', C. 'Invalid. Counterexample when P and Q are true and R and S are false', D. 'Invalid. Counterexample when P, R, and S are true and Q is false! Answer Option 1: To determine the validity of the argument, we can construct an indirect truth table. $PQRS(Q \cdot S) \supset \sim P \sim (\sim Q \vee R) \sim P$ TTTTTFFF TTTFFFTF TETTTEFF TETETETE TFFTTFFF TFFFTTTF FTTTTFFF FTTFFFTF FFTTTFFT FFTFTFTT FFFTTFFT FFFFTTTT In the given argument, the premises are P \supset (Q \vee R), (Q \cdot S) \supset ~P and ~(~Q \vee R), and the conclusion is ~P. From the truth table, we can see that whenever the premises are true, the conclusion is also true. There is no row in the truth table where all the premises are true and the conclusion is false. Therefore, the argument is valid. Final Answer: A. Valid Answer Option 2: To determine the validity of the argument, we will construct an indirect truth table. $P \supset (Q \vee R)$ $(\mathsf{Q}\cdot\mathsf{S})\supset \mathsf{\sim}\mathsf{P}$ ~(~Q ∨ R) / ~P Truth Table: PQRSQ \vee RQ \cdot S \sim P \sim (\sim Q \vee R) \sim P TTTTTFFF TTTFTFFF TTFTTTFFF TTFFTFFF TETTTTEFF TFTFTFFF TEETEEFE TFFFFFFF

From the indirect truth table, we can see that there is no row where all the premises are true and the conclusion is false. Therefore, the argument is valid.

Final Answer: A. 'Valid'

Answer Option 3:

To determine the validity of the argument using indirect truth tables, we will construct a truth table with all the relevant propositions and check if there is any row where all the premises are true and the conclusion is false

 $P\mid Q\mid R\mid S\mid (P\supset (Q\vee R))\mid ((Q\cdot S)\supset \sim P)\mid \sim (\sim Q\vee R)\mid \sim P$

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Based on the truth table, we can see that there is no row where all the premises are true and the conclusion is false. Therefore, the argument is valid.

Final Answer: A. 'Valid'

Answer Option 4:

We will construct an indirect truth table to determine the validity of the argument.

 $P \supset (Q \vee R)$

 $(\mathsf{Q}\cdot\mathsf{S})\supset \mathsf{\sim}\mathsf{P}$

~(~Q ∨ R) / ~P

We will assume the opposite of the conclusion, $\sim\!\!P$, is true. Let's denote the assumptions as follows:

1. $P \supset (Q \lor R)$

2. (Q · S) ⊃ ~P

3. ~(~Q ∨ R)

4. ~P (Assumption)

Next, we will use the assumptions and apply logical deductions to get the truth values of the premises and conclusion.

1. $P \supset (Q \lor R)$ (Premise)

2. (Q · S) \supset ~P (Premise)

3. ~(~Q v R) (Premise)

4. ~P (Assumption)

5. ~Q (Modus Tollens: 2, 4)

6. ~Q v R (Addition: 5)

7. ~~(~Q v R) (Double Negation: 6)

8. ~P \supset (~Q \vee R) (Implication: 7)

9. \sim P \supset (Q \vee R) (Commutation: 8)

10. Q \vee R (Modus Ponens: 1, 9)

11. ~Q ∧ ~R (De Morgan's Law: 10)

12. ~Q (Simplification: 11)

13. ~~Q (Double Negation: 12)

14. Q (Double Negation: 13)

15. Q ^ ~Q (Conjunction: 14, 12)

At this point, we have reached a contradiction, as Q \land ~Q cannot be true. Therefore, our assumption, ~P, must be false, and thus P must be true.

Since we have found that the assumption ~P leads to a contradiction and must be false, the argument is valid. Therefore, the correct answer is A. 'Valid'.

Answer Option 5:

To determine the validity of the argument, we can use an indirect truth table to evaluate the premises and the conclusion.

 $P \supset (Q \vee R)$

 $(Q \cdot S) \supset \sim P$

~(~Q ∨ R) / ~P

We need to consider all possible truth value combinations for the propositions P, Q, R, and S and check if the conclusion is always true when all the premises are true.

 $P \mathrel{Q} R \mathrel{S} \mathrel{Q} \vee R \mathrel{P} \supset (Q \vee R) \; (Q \cdot S) \supset {\sim} P \; {\sim} ({\sim} Q \vee R) \; {\sim} P$

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From the truth table, we can see that there are cases where all the premises are true, but the conclusion is false. Specifically, if P is false, Q is false, R is false, and S is false, then all the premises are true but

the conclusion ~P is false.

Therefore, the argument is invalid. The correct answer is option D: 'Invalid. Counterexample when P, R, and S are true and Q is false'.

Answer Option 6:

To determine the validity of the argument, we can create an indirect truth table to analyze all possible truth value combinations.

P ⊃ (Q ∨ R)

 $(\mathsf{Q}\cdot\mathsf{S})\supset \mathsf{\sim}\mathsf{P}$

~(~Q ∨ R) / ~P

We will list all the premises and the conclusion, and then create a truth table to check if there is any row where all the premises are true and the conclusion is false.

 $P\mid Q\mid R\mid S\mid Q\vee R\mid P\supset (Q\vee R)\mid (Q\cdot S)\supset {\sim}P\mid {\sim}({\sim}Q\vee R)\mid {\sim}P$

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F|F|F|F|F|T|T|T

From the truth table, we can see that there is no row where all the premises are true and the conclusion is false. Therefore, the argument is valid.

Final Answer: A. 'Valid'

Answer Option 7:

To determine the validity of the argument using indirect truth tables, we need to consider all possible truth value assignments for the premises and the conclusion.

Let's create a truth table to analyze the argument:

 $P\mid Q\mid R\mid S\mid P\supset (Q\vee R)\mid (Q\cdot S)\supset \sim P\mid \sim (\sim Q\vee R)\mid \sim P$

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F|T|F|F|T|T|F|T

F|F|T|T|T|T|T

F|F|T|F|T|T|T

F|F|F|T|T|T|T

F|F|F|F|T|T|T

From the truth table, we can see that there is at least one row where all the premises are true (T) and the conclusion is false (F). This indicates that the argument is invalid.

To find a counterexample, we can look for a row where all the premises are true, but the conclusion is false. In this case, we can see that when P is false (F), Q is true (T), R is false (F), and S is true (T), all

the premises are true, but the conclusion is false.

 $Therefore, the argument is invalid. \ The correct option is B. 'Invalid. \ Counterexample \ when \ P, \ Q, \ R, \ and \ S \ are \ true'.$

Answer Option 8:

To determine the validity of the argument using indirect truth tables, we start by assuming that the conclusion is false and then determine if it leads to a contradiction.

Assuming ~P is false, we can conclude that P is true.

Using the premises:

 $1. \ P \supset (Q \vee R)$

2. (Q · S) ⊃ ~P

3. ~(~Q ∨ R)

We can construct the truth table as follows:

 $P \mid Q \mid R \mid S \mid \neg P \mid (Q \vee R) \mid P \supset (Q \vee R) \mid (Q \cdot S) \mid \neg (\neg Q \vee R)$

T|T|T|F|T|T|F

T|T|T|F|F|T|T|F|F

T|T|F|T|F|T|T|F

TITIFIFIFITITIFIF

T|F|T|T|F|T|T|F

TIFITIFIFITITIFIF

TIFIFITIFIFIFITIF

TIFIFIFIFIFIFIF

FITITITITITITIT

F|T|T|F|T|T|T|F|T

F|T|F|T|T|T|T|T

F|T|F|F|T|T|T|F|T

F|F|T|T|T|T|T

F|F|T|F|T|T|F|T

F|F|F|T|T|F|T|T

F|F|F|F|T|F|T|F|T

From the truth table, we can see that when ~P is false, the premises 1, 2, and 3 are all true. Therefore, we have not found a counterexample.

Since there is no counterexample, we can conclude that the argument is valid.

Final Answer: A. Valid

Answer Option 9:

To determine the validity of the argument using an indirect truth table, we need to consider all possible truth value combinations for the premises and conclusion. Let's create a truth table:

 $P \mid Q \mid R \mid S \mid Q \vee R \mid P \supset (Q \vee R) \mid (Q \cdot S) \supset {\sim}P \mid {\sim}({\sim}Q \vee R) \mid {\sim}P$

We can see that there are cases where the premises are true $(-(-Q \lor R)$ is true) but the conclusion (-P) is false. Specifically, when P is false, Q is true, R is true, and S is false.

Therefore, the argument is Invalid. A counterexample is when P is false, Q is true, R is true, and S is false.

Final Answer: C. Invalid. Counterexample when P and Q are true and R and S are false.

 Question 37: *Construct a complete truth table for the following argument. Then, using the truth table, determine whether the argument is valid or invalid. If the argument is invalid, choose an option which presents a counterexample. (There may be other counterexamples as well.) W · T / U', A. 'Valid', B. 'Invalid. Counterexample when T and W are true and U is false', C. 'Invalid. Counterexample when T is true and W and U are false', D. 'Invalid. Counterexample when W is true and T and U are false Answer Option 1: To construct a complete truth table for the argument, we need to consider all possible combinations of truth values for the variables T, U, and W. $\mathsf{T}\,\mathsf{U}\,\mathsf{W}\,\mathsf{T}\,\mathsf{\vee}\,\mathsf{U}\,\mathsf{W}\cdot\mathsf{T}/\,\mathsf{U}$ TTTTT TTFTF TFTTT TFFTF FTTTT FTFFT FFTFT FFFFT From the truth table, we can see that there is no row where T \vee U is true and W \cdot T / U is false. Therefore, the argument is valid. Final Answer: A. 'Valid' Answer Option 2: To construct a complete truth table, we need to consider all possible combinations of truth values for the variables T, U, and W. TUWTVUW·T/U TTTTT TTFTFF TETTTT TFFTFT FTTTTF FTFTFF FFTFTT From the truth table, we can observe that there are no rows in which the premise $(T \vee U)$ is true and the conclusion $(W \cdot T / U)$ is false. Therefore, the argument is valid. The correct answer is A. 'Valid'. Answer Option 3: To construct a complete truth table for the argument, we need to consider all possible combinations of truth values for the variables T, U, and W. TUWTVUW·T/U TTTTT TTFTF TFTTT TEETE

FTTTT

FTFFF

FFTFT

FFFFT

From the truth table, we can see that there are no rows where the premise, $T \lor U$, is true and the conclusion, $W \lor T / U$, is false. Therefore, the argument is valid.

The correct answer is A. 'Valid'.

Answer Option 4:

To construct a complete truth table for the argument, we need to consider all possible combinations of truth values for the variables T, U, and W.

 $T~U~W~T~\vee~U~W~\cdot~T~/~U$

00000
00101
01010
01110
10010
10110
11010
11111
In the truth table, we evaluate the compound statement $T \vee U$ and $W \cdot T / U$.
For the argument to be valid, the compound statement T v U must always imply W · T / U.
Looking at the truth table, we can see that whenever T v U is true, W · T / U is also true. Therefore, the argument is valid.
Final Answer: A. "Valid"
Answer Option 5:
To construct a truth table for the argument, we need to consider all possible combinations of truth values for the variables T, U, and W.
T U W T√U W-T/U
TITIT
TIFIF
TETTT
TFFTF
FITTI
FTFTF
FFTFF
FFFFF
Looking at the truth table, we can see that for every row where the statement T v U is true, the statement W · T / U is also true. There are no counterexamples where T and W are true, and U is false.
Therefore, the argument is valid.
Final Answer: A. Valid
Answer Option 6:
To construct the truth table for the argument, we need to consider all possible combinations of truth values for the propositions T, U, and W.
TUWT~UW-T/U
TIFIT
TFTTF
TFFTF
FITTI
FTFTF
FFTFF
FFFFF
From the truth table, we can see that there are cases where the statement 'T v U' is true (when T is true or U is true) and cases where the statement 'W · T / U' is false (when W is false and either T or U
false).
Therefore, the argument is invalid.
Option B presents a counterexample when T and W are true and U is false, which matches one of the cases in the truth table.
Final Answer: B. Invalid. Counterexample when T and W are true and U is false.

Answer Option 7:
To construct a complete truth table for the argument, we need to consider all possible combinations of truth values for the variables T, U, and W.
TUWT~UW·T/U
00000
00101
01010

01110 10010 10111 11010 11111 Looking at the truth table, we can see that there are no rows where the premise (T v U) is true and the conclusion (W · T / U) is false. In other words, there are no counterexamples where the premise is true but the conclusion is false. Therefore, the argument is valid. Final Answer: A. 'Valid' Answer Option 8: Truth table for the given argument: T|U|W|T \ U|W \ T / U TITITIT T|T|F|T|TTIFITITIF TIFIFITIF FITITIF FITIFITIF F|F|T|F|F F|F|F|F|F From the truth table, we can see that there are instances where the premises $(T \lor U)$ and $(W \cdot T / U)$ are true, but the conclusion $(W \cdot T / U)$ is false. Specifically, when T is true, W is true, and U is false. Therefore, the argument is invalid. The correct answer is B. "Invalid. Counterexample when T and W are true and U is false". Answer Option 9: To construct a complete truth table for the argument, we need to consider all possible combinations of truth values for the variables T, U, and W. First, let's create a truth table for T \vee U: $\mathsf{T} \mid \mathsf{U} \mid \mathsf{T} \vee \mathsf{U}$ TITIT TIFIT FITIT F|F|F Next, let's create a truth table for W \cdot T / U: $\mathsf{T} \, | \, \mathsf{U} \, | \, \mathsf{W} \, | \, \mathsf{T} \, | \, \mathsf{W} \cdot \mathsf{T} \, / \, \mathsf{U}$ TITITIT TITITIFIF TITIFITIF T|T|F|F|F TIFITITIT T|F|T|F|F TIFIFITIT T|F|F|F|T FITITIT FITITIFIF

F|T|F|T|F F|T|F|F|F F|F|T|T|F
F|F|T|F
F|F|F|F|F

Now, let's compare the truth values of T \vee U and W \cdot T / U. The argument states that the two statements are equivalent.

We can see that for every row in the truth table, the truth values of T v U and W · T / U are always the same. Therefore, the argument is valid. The correct answer is A. 'Valid'.

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'Identify the antecedent of the following conditional proposition: The university raises tuition if both the governor approves of it and the board of trustees recommends it.', A. 'The university raises tuition.', B.

The governor approves of it.', C. The board of trustees recommends it.', D. The governor approves of it and the board of trustees recommends it.'

Answer Option 1:

The antecedent of a conditional proposition is the part that comes before the "if" or "when" in the statement. In this case, the antecedent is "both the governor approves of it and the board of trustees recommends it." Therefore, the correct answer is D. The governor approves of it and the board of trustees recommends it."

Answer Option 2:

The antecedent of a conditional proposition is the part that comes before the "if" statement. In this case, the antecedent is "The governor approves of it and the board of trustees recommends it." Therefore,

the correct answer is D. The governor approves of it and the board of trustees recommends it.

Answer Option 3:

The antecedent of a conditional proposition is the part that comes before the "if" statement. In this case, the antecedent is "both the governor approves of it and the board of trustees recommends it."

Therefore, the correct answer is D. 'The governor approves of it and the board of trustees recommends it.'

Answer Option 4:

The antecedent of a conditional proposition is the part that comes before the "if" keyword. In this case, the antecedent is "both the governor approves of it and the board of trustees recommends it."

Therefore, the correct answer is D. 'The governor approves of it and the board of trustees recommends it.'

Answer Option 5:

The antecedent of a conditional proposition is the statement that comes after the "if" in the conditional statement. In this case, the antecedent is "both the governor approves of it and the board of trustees recommends it."

So, the correct answer is D. 'The governor approves of it and the board of trustees recommends it.'

.....

Answer Option 6:

The antecedent of a conditional proposition is the part of the statement that comes before the "if" or "when" clause. In this case, the antecedent is "both the governor approves of it and the board of trustees recommends it." Therefore, the correct answer is D. The governor approves of it and the board of trustees recommends it."

Answer Option 7:

The antecedent of a conditional proposition is the part that comes immediately after the "if" in the statement. In this case, the antecedent is "The governor approves of it and the board of trustees recommends

it.' Therefore, the correct answer is D. 'The governor approves of it and the board of trustees recommends it.'

Answer Option 8:

The antecedent of the conditional proposition is the part of the statement that comes before the "if" keyword. In this case, the antecedent is the condition that needs to be satisfied for the consequent (the part

The antecedent of the given conditional proposition, "The university raises tuition if both the governor approves of it and the board of trustees recommends it," is:

D. 'The governor approves of it and the board of trustees recommends it.'

Answer Option 9:

The antecedent of a conditional proposition is the part that comes before the arrow (\rightarrow) symbol. In this case, the antecedent is the condition that needs to be satisfied in order for the university to raise tuition.

The antecedent of the given conditional proposition is: 'The governor approves of it and the board of trustees recommends it.'

Therefore, the correct answer is D. 'The governor approves of it and the board of trustees recommends it.'

Question 39:

'Identify the antecedent of the following conditional proposition: The restaurant does not have peas, so instead they serve corn and potatoes.', A. 'The restaurant has peas.', B. The restaurant does not have peas.', C. The restaurant serves corn.', D. The restaurant serves potatoes.'

Answer Option 1:

The antecedent of a conditional proposition is the part that comes before the "so" or "if" in the statement. In this case, the antecedent is "The restaurant does not have peas."

Therefore, the correct answer is B. 'The restaurant does not have peas.'

Answer Option 2:

The antecedent of a conditional proposition is the statement that comes before the 'if' in the conditional statement. In this case, the conditional proposition is "The restaurant does not have peas, so instead they serve corn and potatoes."

The antecedent is "The restaurant does not have peas."

Therefore, the correct answer is B. 'The restaurant does not have peas.'

Answer Option 3:

The antecedent of a conditional proposition is the statement that comes after the "if" or "so" part of the conditional statement. In this case, the sentence "The restaurant does not have peas" is the antecedent of the conditional proposition.

Therefore, the correct answer is B. 'The restaurant does not have peas.'

.....

Answer Option 4:

The antecedent of a conditional proposition is the part that comes before the "if" or "so" in the statement. In this case, the antecedent is the condition that is being stated. The conditional proposition is "The restaurant does not have peas, so instead they serve corn and potatoes." The antecedent is the part that comes before the "so" in the statement, which is "The restaurant does not have peas." Therefore, the correct answer is B. The restaurant does not have peas.'

Answer Option 5:

The antecedent of a conditional proposition is the statement that comes before the "if" or the "so" in the sentence. In this case, the antecedent is 'The restaurant does not have peas', as it is the statement that is followed by the word 'so'.

Therefore, the correct answer is B. 'The restaurant does not have peas.'

Answer Option 6:

The antecedent of a conditional proposition is the part that comes before the "if." In this case, the conditional proposition is "The restaurant does not have peas, so instead they serve corn and potatoes." The part that comes before the "if" is "The restaurant does not have peas." Option B is the correct answer. Final Answer: B

Answer Option 7:

The antecedent of a conditional proposition is the statement that comes before the "if" or "so" in the conditional statement. In this case, the conditional proposition is "The restaurant does not have peas, so instead they serve corn and potatoes." The statement that comes before the "so" is "The restaurant does not have peas." Therefore, the antecedent in this case is "The restaurant does not have peas."

Final Answer: B. The restaurant does not have peas.'

Answer Option 8:

The antecedent of a conditional proposition is the part that comes before the "if" or "so" in the sentence. In this case, the antecedent is the statement that is being negated or denied. The given conditional proposition is "The restaurant does not have peas, so instead they serve corn and potatoes." The antecedent in this case is the statement that is being denied or negated, which is "The restaurant has peas."

Therefore, the correct answer is A. 'The restaurant has peas.'

Answer Option 9:

The antecedent of a conditional proposition is the part that comes before the "if" or "so" in the statement. In this case, the antecedent is the condition that leads to the consequence.

The statement is: "The restaurant does not have peas, so instead they serve corn and potatoes."

The antecedent is the condition that leads to the consequence. In this case, it is "The restaurant does not have peas."

Therefore, the antecedent of the given conditional proposition is:

B. 'The restaurant does not have peas.'

Question 40: Select the best translation into predicate logic. All jellybeans are sweeter than all fruit. (Fx: x is a fruit; Jx: x is a jellybean; Sxy: x is sweeter than y)', A. '(∀x)[Fx ⊃ (∀y)](Jy ⊃ Sxy)]', B. '(∀x)[Fx • Jy) • Sxy]', $C. \ '(\forall x)(\forall y)[Sxy \supset (Fy \bullet Jx)]', \ D. \ '(\forall x)[Jx \supset (\forall y)(Fy \supset Sxy)]'$ Answer Option 1: The best translation into predicate logic for the statement "All jellybeans are sweeter than all fruit" is: $(\forall x)[Jx\supset (\forall y)(Fy\supset Sxy)]$ This translates to: For all x, if x is a jellybean, then for all y, if y is a fruit, then x is sweeter than y. Therefore, the correct answer is option D: '($\forall x$)[Jx \supset ($\forall y$)(Fy \supset Sxy)]'. Answer Option 2: The statement "All jellybeans are sweeter than all fruit" can be translated into predicate logic as follows: Fx: x is a fruit Jx: x is a jellybean Sxy: x is sweeter than y The best translation into predicate logic is option D, which is: $(\forall x)[Jx \supset (\forall y)(Fy \supset Sxy)]$ This translates to "For all x, if x is a jellybean, then for all y, if y is a fruit, x is sweeter than y." Final Answer: D ****************** Answer Option 3: The best translation into predicate logic for the statement "All jellybeans are sweeter than all fruit" would be: $(\forall x)[Jx \to (\forall y)(Fy \to Sxy)]$ Explanation: - Fx: x is a fruit - Jx: x is a jellybean - Sxy: x is sweeter than y The translation captures the idea that for all objects x that are jellybeans, if there is any object y that is a fruit, then x is sweeter than y. Therefore, the correct answer is option D: $(\forall x)[Jx \to (\forall y)(Fy \to Sxy)]$ Answer Option 4: To translate the statement "All jellybeans are sweeter than all fruit" into predicate logic, we can use the following translations: Fx: x is a fruit Jx: x is a jellybean Sxy: x is sweeter than y The statement can be translated as: $(\forall x)[Jx \rightarrow (\forall y)(Fy \rightarrow Sxy)]$ Option D: $(\forall x)[Jx \rightarrow (\forall y)(Fy \rightarrow Sxy)]$ This option correctly captures the meaning of the original statement, stating that for all x, if x is a jellybean, then for all y, if y is a fruit, x is sweeter than y. Answer Option 5: The best translation into predicate logic for the statement "All jellybeans are sweeter than all fruit" is: $(\forall x)[Jx \supset (\forall y)(Fy \supset Sxy)]$ Explanation: Jx: x is a jellybean Fx: x is a fruit Sxy: x is sweeter than y The statement can be broken down as follows:

For all x, if x is a jellybean (Jx), then for all y, if y is a fruit (Fy), then x is sweeter than y (Sxy).

Therefore, the correct translation is: $(\forall x)[Jx\supset (\forall y)(Fy\supset Sxy)], \text{ which is option D}.$

Fx: x is a fruit Jx: x is a jellybean Sxv: x is sweeter than v The statement "All jellybeans are sweeter than all fruit" can be broken down into the following logical structure: For every x, if x is a jellybean (Jx), then for every y, if y is a fruit (Fy), then x is sweeter than y (Sxy). Translated into predicate logic, this becomes: $(\forall x)[Jx\supset (\forall y)(Fy\supset Sxy)]$ Option D, $(\forall x)[Jx \supset (\forall y)(Fy \supset Sxy)]$, is the correct translation. Final Answer: D Answer Option 7: The best translation into predicate logic for the given statement is option A: $(\forall x)[Fx\supset (\forall y)(Jy\supset Sxy)]^t.$ Let's break down the translation: Fx: x is a fruit. Jx: x is a jellybean. Sxy: x is sweeter than y. The given statement states that "All jellybeans are sweeter than all fruit." In predicate logic, we can represent this as $(\forall x)[(Jx\supset (y)(Fx\supset Sxy))].$ This translates to "For all x, if x is a jellybean, then for all y, if y is a fruit, then x is sweeter than y." Simplifying the statement, we get $(\forall x)[Fx \supset (\forall y)(Jy \supset Sxy)]$. This represents the given statement accurately and is the best translation into predicate logic. Final Answer: A. Answer Option 8: The given statement "All jellybeans are sweeter than all fruit" can be translated into predicate logic as: $(\forall x)(\forall y)[(Jx \land Fy) \rightarrow Sxy]$. A. $(\forall x)[Fx \supset (\forall y)(Jy \supset Sxy)]$ This option states that if something is a fruit, then everything is a jellybean that is sweeter than it. This is not the correct translation. B. $(\forall x)(\forall y)[(Fx \cdot Jy) \cdot Sxy]$ This option states that everything is both a fruit and a jellybean, and everything is sweeter than everything else. This is not the correct translation. C. $(\forall x)(\forall y)[Sxy \supset (Fy \bullet Jx)]$ This option states that if something is sweeter than something else, then the other thing is both a fruit and a jellybean. This is not the correct translation. D. $(\forall x)[Jx \supset (\forall y)(Fy \supset Sxy)]$ This option states that if something is a jellybean, then everything is a fruit that is sweeter than it. This is not the correct translation. Therefore, the correct translation is option B: $(\forall x)(\forall y)[(Jx \land Fy) \rightarrow Sxy]$. Answer Option 9: The best translation into predicate logic for the statement "All jellybeans are sweeter than all fruit" is option D: $(\forall x)[Jx \supset (\forall y)(Fy \supset Sxy)]^T$. Let's break down the options to determine the correct translation: Option A: "(∀x)[Fx ⊃ (∀y)(Jy ⊃ Sxy)]" translates to "For all x, if x is a fruit then for all y, if y is a jellybean then x is sweeter than y." This does not capture the statement that all jellybeans are sweeter than all Option B: '(∀x)(∀y)[(Fx • Jy) • Sxy]' translates to "For all x and for all y, if x is a fruit and y is a jellybean, then x is sweeter than y." This does not capture the statement that all jellybeans are sweeter than all Option C: $(\forall x)(\forall y)[Sxy \supset (Fy \bullet Jx)]^*$ translates to "For all x and for all y, if x is sweeter than y, then y is a fruit and x is a jellybean." This does not capture the statement that all jellybeans are sweeter than all Option D: '(∀x)[Jx \supset (∀y)(Fy \supset Sxy)]' translates to "For all x, if x is a jellybean then for all y, if y is a fruit then x is sweeter than x." This captures the statement that all jellybeans are sweeter than all fruit.

Answer Option 6:

Final Answer: D

The best translation into predicate logic for the statement "All jellybeans are sweeter than all fruit" is: $(\forall x)[Jx \supset (\forall y)(Fy \supset Sxy)]$

Question 41:
'Identify the conclusion of the following argument. That the world is my idea is a truth valid for every living creature, though only man can contemplate it. In doing so, he attains philosophical wisdom. No truth
is more absolutely certain than that all that exists for knowledge and therefore this world is only object in relation to subject, perception of a perceiver. The world is an idea. (Arthur Schopenhauer, The World
as Will and Idea)', A. 'The world is an idea.', B. 'That the world is my idea is a truth valid for every living creature, though only man can contemplate it.', C. 'In doing so, he attains philosophical wisdom.', D. 'No
truth is more absolutely certain than that all that exists for knowledge and therefore this world is only object in relation to subject, perception of a perceiver.'
Answer Option 1:
The conclusion of the argument is: A. 'The world is an idea.'

Answer Option 2:
The conclusion of the argument is: A. 'The world is an idea.'

Answer Option 3:
The conclusion of the argument is stated in option A: 'The world is an idea.' Final Answer: A

Answer Option 4:
The conclusion of the argument is stated in the last sentence: "The world is an idea." Therefore, the correct answer is A. 'The world is an idea.'

Answer Option 5:

The conclusion of the argument is A. 'The world is an idea.' This is stated explicitly in the last sentence of the argument: "The world is an idea." Final Answer: A

Answer Option 6:

The conclusion of the argument is: A. 'The world is an idea.' The statement "The world is an idea" is explicitly mentioned as the conclusion of the argument in the last sentence of the provided passage. Final Answer: A

Answer Option 7:
The conclusion of the argument is: A. 'The world is an idea.' The argument leads to the statement that the world is an idea, which is the central claim that is being made. Final Answer: A

Answer Option 8:

The conclusion of the argument is: A. 'The world is an idea.' This is the main point that is being asserted and supported throughout the argument. The other statements provide additional context and explanations, but the ultimate conclusion is that the world is an idea. Final Answer: A

Answer Option 9:
The conclusion of the argument is: A. 'The world is an idea.'

This is stated directly in the last sentence of the argument: "The world is an idea." Therefore, option A is the correct choice.

Question 42:
'Select the best translation into predicate logic. Leo is taller than Cathy. (c: Cathy; l: Leo; Txy: x is taller than y)', A. 'Tlc', B. 'Tcl', C. 'ITc', D. 'cTl'
Answer Option 1:
The best translation into predicate logic for the statement "Leo is taller than Cathy" would be "Tlc". This statement can be read as "There exists an x (Leo) and a y (Cathy) such that x is taller than y." Final
Answer: A
Answer Option 2:
The best translation into predicate logic for the statement "Leo is taller than Cathy" is: Tlc.
In this translation, Txy represents the predicate "x is taller than y". I' represents Leo and 'c' represents Cathy. Therefore, Tlc translates to "Leo is taller than Cathy".
Final Answer: A
Answer Option 3:
The best translation into predicate logic would be A. 'Tlc', where Txy represents x is taller than y. In this case, Leo (I) is taller than Cathy (c), so the correct translation would be Tlc. Final Answer: A
Answer Option 4:
The best translation into predicate logic for the statement "Leo is taller than Cathy" would be option A: 'Tlc'. This translates to the predicate T(I, c), which asserts that Leo (I) is taller than Cathy (c). Final
Answer: A
Answer Option 5:
The best translation into predicate logic for the statement "Leo is taller than Cathy" would be:
Tic
Here, Txy represents "x is taller than y".
In this case, x represents Leo (I) and y represents Cathy (c).
Therefore, "Leo is taller than Cathy" can be translated as Tlc.
Final Answer: A

Answer Option 6:
The best translation into predicate logic for the statement "Leo is taller than Cathy" would be option A: "Tic'. This translates to "Leo is taller than Cathy" or "Leo (I) is taller (T) than Cathy (c)". Final Answer: A
Answer Option 7:
The best translation into predicate logic for the statement "Leo is taller than Cathy" is:
Tic
This represents the predicate T(x,y) where x is Leo and y is Cathy. The statement Tlc asserts that Leo (I) is taller than Cathy (c). Therefore, the correct answer is A. Final Answer: A
Answer Option 8:
The best translation into predicate logic for the statement "Leo is taller than Cathy" would be option A: "Tic'. In this notation, T represents the "is taller than" relation, I represents Leo, and c represents Cathy
Therefore, 'Tlc' translates to "Leo is taller than Cathy". Final Answer: A

Answer Option 9:
The best translation into predicate logic for the statement "Leo is taller than Cathy" would be:
Tic (Leo is taller than Cathy).
Option A: 'Tlc' - This option correctly represents the statement "Leo is taller than Cathy".
Option B: 'Tcl' - This option incorrectly represents the statement as "Cathy is taller than Leo".
Option C: "ITc' - This option incorrectly represents the statement as "Leo is taller than Cathy" but in reverse order.
Option D: 'cTl' - This option incorrectly represents the statement as "Cathy is taller than Leo".
Therefore, the correct translation into predicate logic is A. 'Tlc'.

Question 43:

Construct a complete truth table for the following pairs of propositions. Then, using the truth tables, determine whether the statements are logically equivalent or contradictory. If neither, determine whether they are consistent or inconsistent. Justify your answers.

 $C\supset (D\lor C) \text{ and } C\cdot \neg D', \text{ A. 'Logically equivalent', B. 'Contradictory', C. 'Neither logically equivalent nor contradictory, but consistent', D. 'Inconsistent', D. 'Inconsistent$

Answer Option 1:

To construct a complete truth table for the given propositions, we consider all possible combinations of truth values for the variables C and D.

First proposition: $C \supset (D \vee C)$

 $C \ D \ D \lor C \ C \supset (D \lor C)$

TTTT

TFTT

FTTT

FFFT

Second proposition: C \cdot ~D

C D ~D C · ~D

TTFF

TFTT

FTFF

FFTF

Comparing the two truth tables, we can see that the propositions are not logically equivalent since they have different truth values in some rows. Specifically, the first proposition is true when both C and D are true, whereas the second proposition is false in this case. Therefore, the statements are neither logically equivalent nor contradictory. They are consistent because there exist combinations of truth values where both statements are true.

Final Answer: C. 'Neither logically equivalent nor contradictory, but consistent'

Answer Option 2:

To construct a complete truth table, we need to consider all possible combinations of truth values for the propositions involved.

Let's consider the propositions C and D, and their negations ~C and ~D:

C D ~C ~D D \vee C C \supset (D \vee C) C \cdot ~D

TTEETTE

TFFTTTT

FTTFTTF

FFTTFTF

From the truth table, we can see that both $C\supset (D\lor C)$ and $C\cdot \neg D$ have the same truth values in all rows. Therefore, the statements are logically equivalent.

So, the answer is A. 'Logically equivalent'.

Answer Option 3:

To construct a complete truth table, we need to consider all possible combinations of truth values for the propositions involved.

Let's start with the propositions C and D:

 $C \ D \ \text{$^{\circ}$D} \ V \ C \ C \ \supset \ (D \ \lor \ C) \ C \ \text{$^{\circ}$D}$

TTFTTF

TFTTTT

FTFTTF

FFTFTF

From the truth table, we can see that the propositions C \supset (D \vee C) and C \cdot ~D have different truth values for some rows. Therefore, they are not logically equivalent.

However, we can also see that they are not contradictory since there are rows where both propositions are true.

Therefore, the statements are neither logically equivalent nor contradictory, but consistent. The correct answer is C. 'Neither logically equivalent nor contradictory, but consistent'.

Answer Option 4:

To construct a complete truth table for the given pairs of propositions:

Let C, D be the propositions

 $C \supset (\mathsf{D} \vee \mathsf{C})$ and $\mathsf{C} \cdot \mathsf{\sim} \mathsf{D}$

We have 4 possible combinations for the truth values of C and D: $|C|D|D \lor C|C \supset (D \lor C)| \sim D|C \cdot \sim D|$ |---|---|------| |T|T|T|T|F|F| |T|F|T|T|T|T| |F|T|T|T|F|F| |F|F|F|T|T|F| From the truth table, we can see that the truth values of both propositions, C \supset (D \vee C) and C \cdot \sim D, are the same for all combinations of truth values of C and D. Therefore, the statements are logically Final Answer: A. 'Logically equivalent' ****************** Answer Option 5: To construct a complete truth table for the given propositions, we need to consider all possible combinations of truth values for the propositions C and D. $C \supset (D \vee C) \text{ can be represented as } (\neg C \vee (D \vee C)) \text{ since } C \supset (D \vee C) \text{ is equivalent to } \neg C \vee (D \vee C).$ C · ~D can be represented as C ∧ ¬D. The truth table for the given propositions is as follows: $|C|D|\neg C|D\lor C|\neg C\lor (D\lor C)|C\land \neg D|$ |---|---|------| |T|T|F|T|T|F| |T|F|F|T|T|T| |F|T|T|T|T|F| |F|F|T|F|F|F| By comparing the truth values of $\neg C \lor (D \lor C)$ and $C \land \neg D$, we can see that the two propositions have the same truth values for each combination of truth values for C and D. Therefore, the statements $C \supset (D \lor C)$ ∨ C) and C · ~D are logically equivalent. The correct answer is: A. 'Logically equivalent' Answer Option 6: To construct a complete truth table for the given propositions: $C \supset (D \lor C)$ and $C \cdot \sim D$ We have the following atomic propositions: C and D. We first construct the truth table for $C \supset (D \vee C)$: $C \mid D \mid D \lor C \mid C \supset (D \lor C)$ $T \mid T \mid T \mid T$ TIFITIT FITITIT F|F|F|T Next, we construct the truth table for C \cdot ~D: C | ~D | C · ~D T|F|F TITIT F|F|F To determine whether the statements are logically equivalent or contradictory, we compare the truth tables for the two statements. For the combination of inputs (C = T, D = T), the two statements have different truth values. Therefore, the statements are not logically equivalent. For the combination of inputs (C = T, D = F), the two statements have the same truth value. For the combination of inputs (C = F, D = T), the two statements have the same truth value. For the combination of inputs (C = F, D = F), the two statements have the same truth value. Since there are at least some combinations of inputs where the two statements have different truth values, they are not logically equivalent. Thus, the answer is C. 'Neither logically equivalent nor contradictory, but consistent'.

Answer Option 7:

First, let's construct the truth table for $C \supset (D \lor C)$: $C\mid D\mid (D\vee C)\mid C\supset (D\vee C)$ TITITIT TIFITIT FITITIT F|F|F|T Next, let's construct the truth table for C \cdot ~D: C | D | ~D | C · ~D TITIFIF TIFITIT FITIFIF F|F|T|F Now, let's compare the truth values for the two propositions: $C \supset (D \vee C) \mid C \cdot {\scriptscriptstyle \sim} D$ TIF ТІТ TIF Based on the truth table, we can see that the truth values for the two propositions are not the same for any combination of truth values of C and D. Therefore, the statements C \supset (D \vee C) and C \cdot \neg D are not logically equivalent. However, the statements are also not contradictory since there are some rows in the truth table where both statements are true (e.g., the second row). Therefore, the statements $C \supset (D \lor C)$ and $C \cdot \neg D$ are consistent. Final Answer: C. 'Neither logically equivalent nor contradictory, but consistent' ******* To construct a complete truth table for the given propositions, we have two propositions $C \supset (D \lor C)$ and $C \cdot \sim D$. We have three variables, C, D, and the compound proposition. The truth table for $C \supset (D \lor C)$ will have four rows (2^2) and the truth table for $C \cdot \sim D$ will have eight rows (2^3). Let's construct the truth table: $\mid C \mid D \mid D \vee C \mid C \supset (D \vee C) \mid \neg D \mid C \cdot \neg D \mid$ |---|---|------| |T|T|T|T|F|F| |T|F|T|T|T|T| |F|T|T|T|F|F| |F|F|F|T|T|F| Now, let's analyze the truth table to determine whether the statements are logically equivalent, contradictory, or neither but consistent. The compound proposition $C \supset (D \lor C)$ is true for all possible combinations of truth values of C and D. It means that the proposition is always true, regardless of the truth values of C and D. On the other hand, the compound proposition $C \cdot \sim D$ is not always true. There are rows in the truth table where the proposition is false. Since the two propositions do not have the same truth values for all rows in the truth table, they are not logically equivalent. However, the two propositions are not contradictory, as there are rows in the truth table where both propositions are true. Therefore, the statements are neither logically equivalent nor contradictory but consistent. The correct answer is C. 'Neither logically equivalent nor contradictory, but consistent'. ********** Answer Option 9: To construct a complete truth table, we need to consider all possible combinations of truth values for the propositions C and D. Let's fill in the truth table for the first statement, $C \supset (D \lor C)$: $C \mid D \mid (D \vee C) \mid C \supset (D \vee C)$

TITITIT

To construct the truth table for the propositions $C \supset (D \lor C)$ and $C \cdot \sim D$, we need to consider all possible combinations of truth values for the propositions C and D.

T F T T
F T T T
F F F T
Next, let's fill in the truth table for the second statement, C \cdot ~D:
C D ~D C · ~D
T T F F
T F T T
F T F F
FIFITIF

Now, let's compare the two truth tables to determine whether the statements are logically equivalent, contradictory, consistent, or inconsistent.

For the given statements to be logically equivalent, the truth values of the statements should be the same for all possible combinations of truth values for C and D. Looking at the truth tables, we can see that the statements have different truth values for different combinations of truth values for C and D. Therefore, the statements are not logically equivalent.

For the given statements to be contradictory, the truth values of the statements should be opposite (one true and the other false) for all possible combinations of truth values for C and D. Looking at the truth tables, we can see that the statements have the same truth value (both true) for some combinations of truth values for C and D. Therefore, the statements are not contradictory.

Since the statements are neither logically equivalent nor contradictory, we need to check whether they are consistent or inconsistent. If the statements are consistent, it means there exists at least one combination of truth values for C and D where both statements are true. If the statements are inconsistent, it means there is no combination of truth values for C and D where both statements are true. Looking at the truth tables, we can see that there exists at least one combination of truth values (C = T, D = F) where both statements are true. Therefore, the statements are consistent.

Final Answer: C. 'Neither logically equivalent nor contradictory, but consistent'

Question 44:

Construct a complete truth table for the following pairs of propositions. Then, using the truth tables, determine whether the statements are logically equivalent or contradictory. If neither, determine whether they are consistent or inconsistent. Justify your answers.

Answer Option 1:

To construct a complete truth table for the given pair of propositions, we need to consider all possible truth values for the propositions involved.

Let's break down the given propositions:

1. \sim 0 \supset P (Proposition A)

2. O v P (Proposition B)

We have two atomic propositions involved: O and P. Each atomic proposition can take on two truth values, either True (T) or False (F). Therefore, we will have 2² = 4 rows in our truth table.

Using the truth tables, we can evaluate the truth values of the given propositions for all possible combinations of truth values for O and P:

Now, let's compare the truth values of Proposition A (\sim 0 \supset P) and Proposition B (0 \vee P) in the truth table.

We can see that in all four rows, the truth values of Proposition A and Proposition B are exactly the same. Therefore, we can conclude that the statements are logically equivalent.

Final Answer: A. 'Logically equivalent'

Answer Option 2:

Let's construct the truth table for ${\sim}O \supset P$ and $O \lor P$:

O P ~O ~O ⊃ P O ∨ P

TTFTT

TFFFT

FTTTT

FFTTF

From the truth table, we can see that $\neg O \supset P$ and $O \lor P$ have different truth values for some rows. In row 2, $\neg O \supset P$ is false while $O \lor P$ is true. Therefore, $\neg O \supset P$ and $O \lor P$ are not logically equivalent.

However, \sim 0 \supset P and 0 \vee P are also not contradictory since they are both true in row 1 and false in row 3.

Therefore, the statements $\sim\!0\supset P$ and $0\lor P$ are neither logically equivalent nor contradictory but consistent.

The answer is C. 'Neither logically equivalent nor contradictory, but consistent'.

Answer Option 3:

To construct a complete truth table for the given pair of propositions, let's consider the propositions individually:

~O ⊃ P:

O P ~O ⊃ P

TTT

TFF

FTT

FFT

0 v P:

OPOvP

TTT

TFT

FTT FFF

Now, let's compare the truth values of the two propositions:

 $\text{~O} \supset P \text{ O} \vee P$

ттт

```
TFF
FTT
\mathsf{FTF}
From the truth table, we can see that the two propositions are not logically equivalent since there are different truth values for some rows.
However, the two propositions are not contradictory either since there are some rows where both propositions can still be true simultaneously.
Therefore, the statements are neither logically equivalent nor contradictory, but consistent.
The correct answer is C. 'Neither logically equivalent nor contradictory, but consistent'
Answer Option 4:
To construct a complete truth table for the given propositions, we consider all possible combinations of truth values for the atomic propositions O and P.
The first proposition is \sim 0 \supset P, which can also be written as \neg 0 \to P.
Truth Table for \sim0 \supset P:
O ~O P \negO \negO \rightarrow P
TFTFT
TFFFT
FTTTT
FTFTF
The second proposition is O \vee P.
Truth Table for O \vee P:
0 P 0 v P
TTT
FTT
Comparing the truth tables, we can see that the two propositions are not logically equivalent. In rows 2 and 4 of the truth table, the values for the propositions differ.
Therefore, the statements \sim 0 \sim P and 0 \sim P are neither logically equivalent nor contradictory. However, they are consistent since there are rows in the truth table where both propositions are true.
Final Answer: C. 'Neither logically equivalent nor contradictory, but consistent'
*******
To construct a complete truth table for the given proposition pairs, we need to consider all possible combinations of truth values for the propositions involved.
Let's analyze the propositions:
1. \sim0 \supset P
2. O v P
Constructing a truth table for \sim\!0 \supset P:
\mathsf{O} \, \mathord{\sim} \mathsf{O} \, \mathord{\sim} \mathsf{O} \supset \mathsf{P}
TFT
TFT
FTT
FTF
Constructing a truth table for O \vee P:
0 P O v P
TTT
TFT
FTT
FFF
```

Now, let's compare the truth values of the two propositions:

~0>P0 v P
......

TT

TT

FF
....

Based on the truth table, we can see that the propositions $\neg O \supset P$ and $O \lor P$ have the same truth values for all combinations of truth values for O and P. Therefore, the statements are logically equivalent.

The correct answer is: A. 'Logically equivalent'

Answer Option 6:

To construct a complete truth table for the propositions \sim O \supset P and O \vee P, we need to consider all possible combinations of truth values for O and P.

Let's denote ~O as the negation of O.

 $O | P | \sim O | \sim O \supset P | O \lor P$

T|T|F|T|T

FITITIT

F|F|T|F|F

Based on the truth table, we can see that the propositions \sim O \supset P and O \vee P have different truth values for some combinations of O and P. Therefore, they are not logically equivalent.

However, they are also not contradictory because there are some combinations of O and P where both ~O \supset P and O \lor P are true (e.g., when O is false and P is true). Therefore, the statements are consistent

The correct answer is C. 'Neither logically equivalent nor contradictory, but consistent'.

Answer Option 7:

To construct a complete truth table for the propositions \sim O \supset P and O \vee P, we need to consider all possible combinations of truth values for the propositions O and P.

Here is the truth table for the given propositions:

 $|\; 0\; |\; P\; |\; \text{$^{\sim}$} 0\; |\; \text{$^{\sim}$} 0\; \Rightarrow P\; |\; 0\; \vee P\; |\;$

|---|---|

|T|T|F|T|T|

|T|F|F|T|T|

|F|T|T|T|T|

|F|F|T|F|F|

From the truth table, we can see that the truth values of $\neg O \supset P$ and $O \lor P$ are not the same for all rows. Therefore, the statements $\neg O \supset P$ and $O \lor P$ are not logically equivalent.

However, there is no row in the truth table where both $\neg O \supset P$ and $O \lor P$ are false. Therefore, the statements $\neg O \supset P$ and $O \lor P$ are not contradictory.

Based on this information, the statements \sim O \supset P and O \vee P are neither logically equivalent nor contradictory, but they are consistent.

Therefore, the correct answer is C. 'Neither logically equivalent nor contradictory, but consistent'.

Answer Option 8:

To construct a truth table for the given pair of propositions ~O \supset P and O \vee P, we need to consider all possible combinations of truth values for the variables O and P.

Here is the truth table:

| O | P | ~O | ~O ⊃ P | O ∨ P |

|---|---|-----|

|T|T|F|T|T|

|T|F|F|T|T|

|F|T|T|T|T|

|F|F|T|F|F|

Now, let's analyze the truth table.

In the column for $-O \supset P$, we see that it is true for all rows except for the last row where O is false and P is false. In this case, the implication $-O \supset P$ is false because the antecedent -O is true and the consequent P is false.

In the column for $O \lor P$, we see that it is true for all rows except for the last row where both O and P are false. In this case, the disjunction $O \lor P$ is false because neither O nor P is true.

Based on the truth table, we can see that the statements $-O \supset P$ and $O \lor P$ are not logically equivalent because they have different truth values in the last row. They are also not contradictory because they have the same truth values in all other rows except the last one.

Answer Option 9: To construct a complete truth table for the given propositions, we need to consider all possible truth values for the propositions involved. Let's denote ~O as proposition A and P as proposition B. Here is the truth table for \sim 0 \supset P and 0 \vee P: $| A | B | \sim A | \sim A \supset B | O | P | O \lor P |$ |---|---| |T|T|F|T|T|T|T|T|F|F|T|T|F|T| |F|T|T|F|F|T|T| |F|F|T|T|F|F|F| Now, let's analyze the truth table: For the proposition \sim O \supset P, the column \sim A \supset B evaluates to T in all rows except for the third row where A is F and B is T. For the proposition $O \lor P$, the column $O \lor P$ evaluates to T in the first three rows where either O or P (or both) are T. $Comparing the columns ~~A \supset B \text{ and } O \lor P, we can see that they have the same truth values in all rows except for the third row.$ Therefore, the statements ~O ⊃ P and O ∨ P are NOT logically equivalent, as they do not have the same truth values in all rows of the truth table. However, they are also NOT contradictory because there are some rows where both propositions are true (rows 1 and 2). Thus, the statements $\sim\!0\supset\!P$ and $0\vee\!P$ are consistent, as there is at least one row where both propositions are true.

Therefore, the correct answer is C. 'Neither logically equivalent nor contradictory, but consistent'.

The answer is C. 'Neither logically equivalent nor contradictory, but consistent'.

Question 45: 'Use indirect truth tables to determine whether the following argument is valid. If the argument is invalid, choose an option which presents a counterexample. (There may be other counterexamples as well.) $(T\cdot {\scriptscriptstyle \sim} U) \supset {\scriptscriptstyle \sim} W$ \sim W \supset X ${\sim} Z \supset W$ ~(X · Z) / T ⊃ U', A. 'Valid', B. 'Invalid. Counterexample when T, X, and Z are true and U and W are false', C. 'Invalid. Counterexample when T and X are true and U, W, and Z are false', D. 'Invalid. Counterexample when T, W, and Z are true and U and X are false' Answer Option 1: To determine the validity of the argument using indirect truth tables, we list all the propositions involved in the argument and their negations. U: U is true W: W is true X: X is true Z: Z is true ¬T: T is false ¬U: U is false ¬W: W is false ¬X: X is false ¬Z: Z is false We construct a truth table to evaluate the argument: $T \ U \ W \ X \ Z \ T \cdot \neg U \ \neg W \ \neg (X \cdot Z) \ T \supset U$ TTTTTTFFT TTTTFTFFT TTTFTTFFT TTTFFTFFT TFFTTFTTF TEETEETTE TFFFTFTTF TEFFETTE FTTTTFFTT FTTTFFFTT FTTFTFFTT FTTFFFFTT FFFTTFTTT FFFTFFTTT FFFFTTTT FFFFFTTT From the truth table, we can see that the conclusion $T\supset U$ is true for all rows where the premises $(T \cdot \neg U) \supset \neg W$, $\neg W \supset X$, and $\neg (X \cdot Z)$ are true. Therefore, the argument is valid. Final Answer: A. Valid Answer Option 2: To determine the validity of the argument using indirect truth tables, we will create a truth table with the premises and the negation of the conclusion and check if there is any row where all the premises are true and the conclusion is false. $(T\cdot {\scriptscriptstyle \sim} U) \supset {\scriptscriptstyle \sim} W$ \sim W \supset X ${\sim} Z \supset W$ ~(X · Z) $/T\supset U$ Truth Table: $T\;U\;W\;Z\;X\;(T\;\cdot\;\sim U)\supset \sim W\;\sim W\supset X\;\sim Z\supset W\;\sim (X\;\cdot\;Z)\;\neg (T\supset U)$ TTTTTFTTTFF

TTTTFFTTTFF

```
TTTFFFTTTFF
TFFTTFTTTT
TEETEETTTTT
TFFFTTTTTT
TEFFETTTTTT
FTTTTTTTT
FTTTFTTTTT
FTTFTTTTTT
FTTFFTTTTT
FFFTTTTTTT
FFFTFTTTTT
FFFFTTTTTT
FFFFFTTTTT
From the truth table, we can see that in every row where all the premises are true, the conclusion is also true. Therefore, the argument is valid.
Final Answer: A. Valid
*******
Answer Option 3:
To determine the validity of the argument, we can use an indirect truth table. We will assume the premises are true and check if the conclusion must also be true.
The argument is as follows:
1. (T · ~U) ⊃ ~W
2. ~W ⊃ X
3.~ \text{~~} Z \supset W
4. ~(X · Z)
\therefore \, T \supset U
Creating the truth table:
T~U~W~X~Z~(T\cdot {\sim} U)\supset {\sim} W~{\sim} W \supset X~{\sim} Z\supset W~{\sim} (X\cdot Z)~T\supset U
TTTTTFTTTFT
TTTTFFTTTFT
TTTFTFTTTTT
TTTFFFTTTFT
TTFTTFFTFFT
TTFTFFFTFFT
TTFFTFFTTFT
TTFFFFFTTFT
TFTTTTTTFFF
TFTTFFTTTFFF
TFTFTTTTTFF
TETEFETTTTEE
TFFTTTFTFFF
TEETEEFEE
TEFETTETTEE
TEFFFFFTTFFF
Based on the truth table, we can see that when T is true and U is false, the conclusion T \supset U is false. Therefore, the argument is invalid.
The correct answer is: B. Invalid. Counterexample when T, X, and Z are true and U and W are false.
*********
Answer Option 4:
To determine the validity of the argument, we can construct an indirect truth table.
T~U~W~X~Z~\sim U~\sim (T~\sim U)~(T~\sim U) \supset \sim W~\sim W~\supset X~\sim Z~\supset W~\sim (X~\sim Z)~T~\supset U~Valid?
```

TTTFTFTTTFF

TFFTTTTTFTFTT TEFTETETETT TEFETTTTETETT TEFETTTETT FTTTTFTTFTTF FTTTFFTTFTTF FTTFTFTTTF FTTFFFTTFTTTF FFFTTTTTTTF FFFTFTTTTTF FFFFTTTTTTF FFFFFTTTTTTF From the truth table, we can see that there is at least one row where all the premises are true (T · ~U) ¬ ~W, ~W ¬ X, ~Z ¬ W, ~(X · Z), and the conclusion T ¬ U is false. Therefore, the argument is invalid. The counterexample is when T and X are true, and U, W, and Z are false. Final Answer: C. 'Invalid. Counterexample when T and X are true and U, W, and Z are false' ******* Answer Option 5: To determine the validity of the argument, let's construct an indirect truth table: $(T\cdot {\sim} U)\supset {\sim} W\mid {\sim} W\supset X\mid {\sim} Z\supset W\mid {\sim} (X\cdot Z)\mid T\supset U$ T|T|T|T|U T|T|T|F|U $\mathsf{T} \, | \, \mathsf{T} \, | \, \mathsf{F} \, | \, \mathsf{T} \, | \, \mathsf{U}$ T|T|F|F|U $\mathsf{T} | \mathsf{F} | \mathsf{T} | \mathsf{T} | \mathsf{U}$ T|F|T|F|U T|F|F|T|U T|F|F|F|U F|T|T|T|U F|T|T|F|U F|T|F|T|U F|T|F|F|U F|F|T|T|U F|F|T|F|U F|F|F|T|U F|F|F|U We can see that in all cases where the premises are true, the conclusion is also true. Therefore, the argument is valid. Final Answer: A. Valid Answer Option 6: To determine the validity of the argument using indirect truth tables, we need to consider all possible truth value combinations for the propositions involved. Let's create a truth table to analyze the argument: $T~U~W~X~Z~(T\cdot {\sim} U)\supset {\sim} W~{\sim} W\supset X~{\sim} Z\supset W~{\sim} (X\cdot Z)~T\supset U$ FFFFFTTTT FFFFTTTTT FFFTFTTTT FFFTTTTTT FTTFFTTTTT

FTTFTTTTT

FTTTFTTTTT

FTTFTFFTF

TFTFTTFFTF

TETTETTETE TETTTTTETE TTTFFFFTTT TTTFTFFTTT TTTTFFTFTT TTTTTFTFTT Based on the truth table, we can see that there is no row where all the premises are true and the conclusion is false. Therefore, the argument is valid. Final Answer: A. 'Valid' ****************** Answer Option 7: To determine the validity of the argument, we can create an indirect truth table. We need to assign truth values to the premises and check if the conclusion is always true. Premises: 1. $(T \cdot \sim U) \supset \sim W$ $2.~\text{~}\text{W} \supset X$ $3.~ {\sim} Z \supset W$ 4. ~(X · Z) Conclusion: $5.\ T\supset U$ We will consider all possible truth value combinations for the variables T, U, W, X, and Z. $T~U~W~X~Z~(T\cdot {\sim} U)\supset {\sim} W~{\sim} W\supset X~{\sim} Z\supset W~{\sim} (X\cdot Z)~T\supset U$ TTTTTTTT TTTTFTTTT TTTFTTFTTT TTTFFTFTTT TEETTTTTTE TFFTFTTTF TEFETTETTE TFFFFTFTTF FTTTTTTTT FTTTFTTTT FTTFTTFTTT FTTFFTFTTT FFTTTTTTT FFTTFTTTT FFTFTTFTTT FFTFFTFTT From the truth table, we can see that for all combinations of truth values, the conclusion $T \supset U$ is true. Hence, the argument is valid. Final Answer: A. Valid ******* Answer Option 8: To determine the validity of the argument using indirect truth tables, we need to consider all possible truth value combinations for the premises and check if the conclusion is true in all cases. The argument is as follows: 1. $(T \cdot \sim U) \supset \sim W$ $2. \ \text{~} W \supset X$ 3. ~Z ⊃ W 4. \sim (X · Z) (premises) 5. T \supset U (conclusion) We will build a truth table considering the truth values of T, U, W, X, and Z. $T~U~W~X~Z~(T\cdot {\sim} U)\supset {\sim} W~{\sim} W\supset X~{\sim} Z\supset W~{\sim} (X\cdot Z)~T\supset U$ TTTTTFTFTF

TTTTFFTFTF

TTTFTFTFF
TTTEFFTFTF
TTFTTFFTTT
TTFTFFFTTT
TTFFTFFTTT
TTFFFFFTTT
TFTTTFTFTT
TFTTFFTFTT
TETETETTT
TFTFFFTTTT
TFFTTFTTTT
TFFTFFTTTT
TEFETETITT
TFFFFFTTTT
FTTTTFTTFT
FTTTFFTTFT
FTTFTFTTFT
FTTFFFTTFT
FTFTTFTTFT
FTFTFFTTFT
FTFFTFTTFT
FTFFFFTTFT
FFTTTTTTT
FFTTFTTTTT
FFTFTTTTT
FFTFFTTTT
FFFTTTTTT
FFFTFTTTTT
FFFFTTTTT
FFFFFTTTTT
In all cases, we observe that the conclusion T ⊃ U is true whenever the premises are true. Therefore, the argument is valid.
Final Answer: A. Valid
Answer Option 9:
To determine the validity of the argument using indirect truth tables, we will create a truth table with all the premises and the negation of the conclusion, and check if there is any row in which all the premises
are true and the conclusion is false.
Let's represent the premises and the conclusion using propositional variables:
P1: (T ~ U) ⊃ ~W
P2: -W⊃ X
P3: -Z ⊃ W
P4: ~(X · Z)
C:ToU
We'll create a truth table with columns for P1, P2, P3, P4, and ~C (negation of C).
T U W Z X P1 P2 P3 P4 ~C
Now, let's fill in the truth table by considering all possible combinations of truth values for T, U, W, Z, and X:
T U W Z X P1 P2 P3 P4 ~C
TITITITITIFITITIFIFI
TITITIFITIFITIFIFI

T|T|T|F|F|F|T|T|F|

Now, let's check if there is any row where all the premises are true and the conclusion is false. In the given truth table, we can see that there is no such row.

Therefore, the argument is valid. Final Answer: A

Question 46:
$"Select the best translation into predicate logic: Cecilia breaks my heart if, and only if, Maria is either a problem or a girl.', A. "Bc = (Pm \lor Gm)", B. "Bc = (Pm \bullet Gm)", C. "Bc \supset (Pm \lor Gm)", D. "Cb \supset (Pm \bullet Gm$
Answer Option 1:
The best translation into predicate logic for the given statement would be:
$Bc \equiv (Pm \vee Gm)$
Explanation:
- Bc represents the statement "Cecilia breaks my heart."
- Pm represents the statement "Maria is a problem."
- Gm represents the statement "Maria is a girl."
The symbol ≡ denotes "if and only if" in predicate logic.
Therefore, the correct answer is A.

Answer Option 2:
The best translation into predicate logic would be: A. 'Bc ≡ (Pm ∨ Gm)'
Explanation:
Let's break down the given statement:
- Cecilia breaks my heart: Bc
- Maria is either a problem or a girl: (Pm v Gm)
The phrase "if, and only if" in logic is represented by the biconditional operator "≡".
Therefore, the best translation into predicate logic is: Bc ≡ (Pm ∨ Gm)
Final Answer: A
Answer Option 3:
The given statement can be translated into predicate logic as follows:
Bc: Cecilia breaks my heart
Pm: Maria is a problem
Gm: Maria is a girl
The statement "Cecilia breaks my heart if, and only if, Maria is either a problem or a girl" can be expressed as "Bc if and only if (Pm or Gm)".
The logical equivalence can be represented by the biconditional operator (=).
Therefore, the best translation into predicate logic is:
Bc ≡ (Pm ∨ Gm)
Final Answer: A

Answer Option 4:
The sentence "Cecilia breaks my heart if, and only if, Maria is either a problem or a girl" can be translated into predicate logic as follows:
Let Bc represent "Cecilia breaks my heart", Pm represent "Maria is a problem", and Gm represent "Maria is a girl".
The phrase "if, and only if" is equivalent to the biconditional operator (=).
Therefore, the best translation into predicate logic is:
$Bc \equiv (Pm \vee Gm)$
Final Answer: A
Answer Option 5:
The given statement can be translated into predicate logic as follows:
Let Bc be the statement "Cecilia breaks my heart."
Let Pm be the statement "Maria is a problem."
Let Gm be the statement "Maria is a girl."
The given statement "Cecilia breaks my heart if, and only if, Maria is either a problem or a girl" can be translated as: Bc if and only if (Pm or Gm).
Therefore, the best translation into predicate logic is: Bc ≡ (Pm ∨ Gm).
Final Answer: A
Answer Option 6:

The best translation into predicate logic would be option A: 'Bc \equiv (Pm \vee Gm)'.

Explanation:
- Bc represents "Cecilia breaks my heart".
- Pm represents "Maria is a problem".
- Gm represents "Maria is a girl".
- The statement "if, and only if" can be represented by the logical equivalence symbol "≡".
- The statement "Maria is either a problem or a girl" can be represented by the disjunction symbol "\".
Putting it all together, the translation becomes: "Cecilia breaks my heart if, and only if, Maria is either a problem or a girl" is represented as 'Bc = (Pm v Gm)'.
Final Answer: A
Answer Option 7:
The best translation into predicate logic for the given statement is: $Bc \equiv (Pm \vee Gm)$
Explanation:
Let Bc represent "Cecilia breaks my heart", Pm represent "Maria is a problem", and Gm represent "Maria is a girl".
The statement "Cecilia breaks my heart if, and only if, Maria is either a problem or a girl" can be translated as "Cecilia breaks my heart if Maria is a problem or Maria is a girl, and Cecilia breaks my heart only if
Maria is a problem or Maria is a girl."
Symbolically, we can represent this as $Bc = (Pm \vee Gm)$.
Option A (Bc \equiv (Pm \vee Gm)) is the correct translation into predicate logic. Final Answer: A
Answer Option 8:
The given statement can be translated into predicate logic as follows:
Bc: Cecilia breaks my heart
Pm: Maria is a problem
Gm: Maria is a girl
The best translation into predicate logic is:
$Bc \equiv (Pm \vee Gm)$
Explanation: The statement "Cecilia breaks my heart if, and only if, Maria is either a problem or a girl" can be represented as "Cecilia breaks my heart if Maria is a problem or Maria is a girl, and Cecilia breaks
my heart only if Maria is a problem or Maria is a girl." This can be further simplified as "Cecilia breaks my heart if and only if Maria is either a problem or a girl." In predicate logic, the "if and only if" condition is
represented by the equivalence operator (\equiv). Hence, the best translation is Bc \equiv (Pm \vee Gm).
Final Answer: A

Answer Option 9:
The statement "Cecilia breaks my heart if, and only if, Maria is either a problem or a girl" can be translated as follows:
Let Bc be the predicate "Cecilia breaks my heart".
Let Pm be the predicate "Maria is a problem".
Let Combe the products "Mario in a girl"

Let Gm be the predicate "Maria is a girl".

The translation of the given statement into predicate logic is:

 $Bc \equiv (Pm \vee Gm)$

Option A: $Bc \equiv (Pm \vee Gm)'$ - This is the correct translation that captures the meaning of the given statement. Cecilia breaking my heart is equivalent to Maria being either a problem or a girl.

Option B: $"Bc = (Pm \bullet Gm)"$ - This translation uses the conjunction operator $"\bullet"$, which implies that both Pm and Gm must be true for Bc to be true. However, the given statement does not require both Pm and Gm to be true simultaneously.

Option C: 'Bc \supset (Pm \lor Gm)' - This translation uses the implication operator " \supset ", which implies that if Bc is true, then either Pm or Gm must be true. However, the given statement requires an "if and only if" condition, meaning that both sides of the statement are dependent on each other.

Option D: 'Cb \supset (Pm \bullet Gm)' - This translation uses the wrong predicates ("Cb" instead of "Bc") and the conjunction operator " \bullet " instead of the disjunction operator " \lor ".

Final Answer: A

Question 47: Select the best translation into predicate logic: If Delilah is in New York City, then Juana is in Bogota or Layla is in San Juan.', A. 'Nd \supset (Bj \bullet SI)', B. 'Dn \equiv (Bx \bullet Ly)', C. 'Nd \equiv (By \bullet SI)', D. 'Nd \supset (Bj \bullet SI)', D. 'Nd \supset (Bj \bullet SI)' Answer Option 1: The best translation into predicate logic for the given statement would be: $Dn \rightarrow (Bi \vee SI)$ This translation accurately represents the statement "If Delilah is in New York City, then Juana is in Bogota or Layla is in San Juan." Final Answer: D ****************** Answer Option 2: Let's break down the given sentence into its components: Delilah is in New York City - Dn (D is in New York City) Juana is in Bogota - Bj (J is in Bogota) Layla is in San Juan - SI (L is in San Juan) The sentence can be translated as: $Dn \rightarrow (Bj \vee Sl)$ Therefore, the best translation into predicate logic is: D. 'Nd \supset (Bj \vee Sl)' Answer Option 3: The best translation into predicate logic would be: D. 'Nd \supset (Bj \vee SI)' Explanation: Let's break down the given statement into logical components: - Delilah is in New York City: Nd - Juana is in Bogota: Bj - Layla is in San Juan: SI - If Delilah is in New York City, then Juana is in Bogota or Layla is in San Juan: Nd \supset (Bj \vee SI) Therefore, option D correctly represents the given statement in predicate logic. Final Answer: D Answer Option 4: The best translation into predicate logic for the given statement would be: D. 'Nd \supset (Bj \vee Sl)' Let's assign the following predicates: Nd: Delilah is in New York City Bj: Juana is in Bogota SI: Layla is in San Juan The given statement "If Delilah is in New York City, then Juana is in Bogota or Layla is in San Juan" can be represented as $Nd \supset (Bj \vee SI)$. This translation accurately captures the conditional relationship between Delilah being in New York City and Juana being in Bogota or Layla being in San Juan. Therefore, the correct choice is D. 'Nd \supset (Bj \vee SI)'. Answer Option 5: The best translation into predicate logic for the given statement is: D. 'Nd \supset (Bj \vee SI)' Let's break down the statement: "If Delilah is in New York City, then Juana is in Bogota or Layla is in San Juan." Let's define the predicate symbols: Nd: Delilah is in New York City Bj: Juana is in Bogota SI: Layla is in San Juan The statement can be translated as: Nd implies (Bj or SI). In predicate logic, implication is represented by the symbol ' \supset ', and 'or' is represented by the symbol ' \lor '. Thus, the best translation into predicate logic is: 'Nd ⊃ (Bj ∨ SI)'.

Answer Option 6:

Let's break down the given sentence into propositional logic:
P: Delilah is in New York City
Q: Juana is in Bogota
R: Layla is in San Juan
The given sentence can be translated as "if P, then (Q or R)".
In predicate logic, we use variables to represent different objects or individuals. Let's use "d" to represent Delilah, "n" to represent New York City, "j" to represent Bogota, and "s" to represent San Juan.
The best translation into predicate logic would be:
P: d is in n
Q: Juana is in j
R: Layla is in s
Therefore, the correct translation would be:
$Nd \supset (Bj \vee SI)$
Final Answer: D

Answer Option 7:
The best translation into predicate logic for the given statement would be:
D: 'Nd \supset (Bj \vee SI)'
Explanation:
- 'Nd' represents "Delilah is in New York City."
- 'Bj' represents "Juana is in Bogota."
- 'SI' represents "Layla is in San Juan."
- '¬' represents "ifthen" or implication.
- 'v' represents "or."
The statement "If Delilah is in New York City, then Juana is in Bogota or Layla is in San Juan" can be translated as "Nd implies (Bj or SI)".
Hence, the best translation into predicate logic is option D: 'Nd \supset (Bj \vee SI)'.

Answer Option 8:
The best translation into predicate logic for the given statement, "If Delilah is in New York City, then Juana is in Bogota or Layla is in San Juan" is:
Option D: 'Nd \supset (Bj \lor SI)'
Explanation:
Nd represents "Deliiah is in New York City"
Bj represents "Juana is in Bogota"
SI represents "Layla is in San Juan"
The symbol ⊃ represents the implication "ifthen"
The symbol ∨ represents the logical OR
So the translated statement would be "If Delilah is in New York City, then Juana is in Bogota or Layla is in San Juan" which can be represented as 'Nd \supset (Bj \vee SI)'
Final Answer: D

Answer Option 9:
The best translation into predicate logic for the given statement is: D. $^{1}Nd \supset (B_{j} \vee S_{j})^{1}$.
Translation breakdown:
- "Delilah is in New York City" is represented by the predicate Nd.
- "Juana is in Bogota" is represented by the predicate Bj.
- "Layla is in San Juan" is represented by the predicate Sl.
- "If Delilah is in New York City, then Juana is in Bogota or Layla is in San Juan" can be translated as $Nd \supset (Bj \lor Sl)$.
This translation accurately captures the conditional relationship, stating that if Delilah is in New York City (Nd), then either Juana is in Bogota (Bj) or Layla is in San Juan (Sl). Final Answer: D. 'Nd \supset (Bj \vee Sl)'.

'Use the following key to translate the given formula of PL to natural, English sentences. A: Marina reads a Percy Jackson book. B: Izzy plays Minecraft. C: Emily stops working. D: Russell makes dinner. E: Ashleigh stops by. ~(A \supset B) • (B \supset ~E)', A. 'It's not the case that Marina's reading a Percy Jackson book entails that Izzy plays Minecraft, but Izzy's playing Minecraft does entail that Ashleigh doesn't stop by.', B. 'If Marina doesn't read a Percy Jackson book, then Izzy plays Minecraft, which entails that Ashleigh doesn't stop by.', C. 'Marina's reading a Percy Jackson book does not entail that Izzy plays Minecraft, but Izzy plays Minecraft provided that Ashleigh doesn't stop by.', D. 'It's not true that Marina reads a Percy Jackson book only when Izzy plays Minecraft, but Izzy plays Minecraft only when Ashleigh stops by.' Answer Option 1: ~(A \supset B) • (B \supset ~E) can be translated as: "It's not the case that Marina's reading a Percy Jackson book entails that Izzy plays Minecraft, but Izzy's playing Minecraft does entail that Ashleigh doesn't stop by." Final Answer: A ******* Answer Option 2: ~(A ⊃ B) • (B ⊃ ~E) can be translated as 'It's not the case that Marina's reading a Percy Jackson book entails that Izzy plays Minecraft, but Izzy's playing Minecraft does entail that Ashleigh doesn't stop by.' Therefore, the correct option is A. 'It's not the case that Marina's reading a Percy Jackson book entails that Izzy plays Minecraft, but Izzy's playing Minecraft does entail that Ashleigh doesn't stop by.' Answer Option 3: Using the given key, we can translate the given formula as follows: $\mathord{\sim} (\mathsf{A} \supset \mathsf{B}) \bullet (\mathsf{B} \supset \mathord{\sim} \mathsf{E})$ ~(Marina reads a Percy Jackson book entails that Izzy plays Minecraft) and (Izzy plays Minecraft entails that Ashleigh doesn't stop by) By rearranging and simplifying the sentence, we get: "It's not the case that Marina's reading a Percy Jackson book entails that Izzy plays Minecraft, but Izzy's playing Minecraft does entail that Ashleigh doesn't stop by." Therefore, the correct translation is option A. Answer Option 4: Using the given key: A: Marina reads a Percy Jackson book. B: Izzy plays Minecraft. C: Emily stops working. D: Russell makes dinner. E: Ashleigh stops by. The formula ${\scriptstyle \sim}(A\supset B)$ • $(B\supset {\scriptstyle \sim} E)$ can be translated as follows: "It's not the case that Marina's reading a Percy Jackson book entails that Izzy plays Minecraft, but Izzy's playing Minecraft does entail that Ashleigh doesn't stop by." Therefore, the correct translation is option A. Answer Option 5: The given formula is ${\sim}(A\supset B)$ • $(B\supset{\sim}E)$ Let's break it down and translate it step by step: $\sim \! (A \supset B) \text{ translates to "It's not the case that Marina's reading a Percy Jackson book entails that Izzy plays Minecraft."}$ (B ⊃ ~E) translates to "Izzy's playing Minecraft does entail that Ashleigh doesn't stop by." Combining these two translations, the overall translation becomes: "It's not the case that Marina's reading a Percy Jackson book entails that Izzy plays Minecraft, but Izzy's playing Minecraft does entail that Ashleigh doesn't stop by." Therefore, the correct answer is A. Answer Option 6: Using the given key, we can translate the given formula to natural English sentences. $\mathord{\sim} (\mathsf{A} \supset \mathsf{B}) \bullet (\mathsf{B} \supset \mathord{\sim} \mathsf{E})$ ~(Marina reads a Percy Jackson book entails that Izzy plays Minecraft) • (Izzy plays Minecraft entails that Ashleigh doesn't stop by) Translating this to natural English sentences, we can rewrite it as:

Question 48:

'It's not the case that Marina's reading a Percy Jackson book entails that Izzy plays Minecraft, but Izzy's playing Minecraft does entail that Ashleigh doesn't stop by.'
Therefore, the correct option is A.

Answer Option 7:
Using the given key:
A: Marina reads a Percy Jackson book.
B: Izzy plays Minecraft.
C: Emily stops working.
D: Russell makes dinner.
E: Ashleigh stops by.
$\sim (A \supset B) \bullet (B \supset \sim E)$
~(Marina reads a Percy Jackson book entails that Izzy plays Minecraft) • (Izzy plays Minecraft entails that Ashleigh doesn't stop by)
The correct translation of the given formula to natural English sentences is:
A. 'It's not the case that Marina's reading a Percy Jackson book entails that Izzy plays Minecraft, but Izzy's playing Minecraft does entail that Ashleigh doesn't stop by.'

Answer Option 8:
The given formula is \sim (A \supset B) • (B \supset \sim E).
Translation:
A: Marina reads a Percy Jackson book.
B: Izzy plays Minecraft.
C: Emily stops working.
D: Russell makes dinner.
E: Ashleigh stops by.
\sim (A \supset B) • (B \supset \sim E) can be translated as:
"It's not the case that Marina's reading a Percy Jackson book entails that Izzy plays Minecraft, but Izzy's playing Minecraft does entail that Ashleigh doesn't stop by."
Therefore, the correct translation is option A: "It's not the case that Marina's reading a Percy Jackson book entails that Izzy plays Minecraft, but Izzy's playing Minecraft does entail that Ashleigh doesn't stop
by."

Answer Option 9:
Let's break down the given formula and translate it using the given key:
$\sim (A \supset B) \bullet (B \supset \sim E)$
~(A \supset B) translates to "It's not the case that if Marina reads a Percy Jackson book, then Izzy plays Minecraft."
$(B \supset \sim E)$ translates to "Izzy playing Minecraft entails that Ashleigh doesn't stop by."
Combining the translations, we get:
"It's not the case that if Marina reads a Percy Jackson book, then Izzy plays Minecraft, but Izzy playing Minecraft entails that Ashleigh doesn't stop by."
Therefore, the correct translation is A. "It's not the case that Marina's reading a Percy Jackson book entails that Izzy plays Minecraft, but Izzy's playing Minecraft does entail that Ashleigh doesn't stop by."

Question 49: 'Use indirect truth tables to determine whether the following argument is valid. If the argument is invalid, choose an option which presents a counterexample. (There may be other counterexamples as well.) $(\mathsf{A} \vee \mathsf{B}) \supset \mathsf{C}$ C $\supset \sim$ D / D \supset A', A. 'Valid', B. 'Invalid', B. 'Invalid', Counterexample when B, C, and D are true and A is false', C. 'Invalid', Counterexample when B and D are true and A and C are false', D. 'Invalid', Counterexample when A, B, and C are false and D is true' To determine the validity of the argument, we will construct an indirect truth table. Let's consider the following propositions: A: A is true B: B is true C: C is true D: D is true Based on the premises and conclusion, we have: Premise 1: (A \vee B) \supset C Premise 2: $C \supset \sim D$ Conclusion: D ⊃ A Using the indirect truth table method, we will assume that the conclusion is false (D \supset A is false) and derive the truth values of the premises. Assume D is true and A is false: $\mathsf{D}\supset\mathsf{A}\;\mathsf{is}\;\mathsf{False}$ A v B is True C is True $C \supset {\scriptscriptstyle \sim} D$ is True The premises hold true in this case, which contradicts the assumption that the conclusion is false. Therefore, the argument is valid. Final Answer: A. Valid Answer Option 2: To determine the validity of the argument, we can use an indirect truth table. $A \ B \ C \ D \ (A \lor B) \supset C \ C \supset \neg D \ D \supset A \ D$ TTTTTTT TTTFTFTT TTFTFTTT TTFFTFTT TFTTTTT TETETETT TFFTFTTT TFFFTFFT FTTTTTF FTTFTFTF FTFTFTTF FTFFTFFT FFTTTTF FFTFTFTF FFFTFTTF FFFFFFF The truth table shows that there is no row where all the premises are true and the conclusion is false. Therefore, the argument is valid. Final Answer: A Answer Option 3: To determine the validity of the argument, we will use indirect truth tables to evaluate all possible combinations of truth values for the premises and the conclusion. $A \ B \ C \ D \ (A \lor B) \supset C \ C \supset {\sim} D \ D \supset A$ TTTTTTT

TTTFTFTT

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TTFTFTTT
TTEETETT
TFTTTTT
TETETETT
TFFTFTTT
TFFFTFTT
FTTTTTF
FTTFTF
FTFTFTTF
FTFFTFTF
FFTTFTTF
FFTFTF
FFFTTTT
FFFFTTTT
After evaluating all possible combinations, we can see that there are no rows where all the premises are true and the conclusion is false. Therefore, the argument is valid.
Final Answer: A. Valid
Answer Option 4:
To determine the validity of the argument, we can construct an indirect truth table. We will assume the premises are true and check if the conclusion is always true.
A \mathrel{B} C \mathrel{D} (A \vee B) \supset C \mathrel{C} \supset \neg D \mathrel{D} \supset A
TTTTTTT
TTTFTTFT
TTFTTFTT
TTFFTFTT
TFTTTTT
TETETTET
TFFTTTTT
TFFFTTFT
FTTTTTF
FTTFTTFF
FTFTTFTF
FTFFTFTF
FFTTTTF
FFTFTTFF
FFFTTTF
FFFFTTF
In all the rows where the premises are true (T), the conclusion D \supset A is also true (T). Therefore, the argument is valid.
Final Answer: A. Valid
*******
Answer Option 5:
To determine the validity of the argument, we will construct an indirect truth table.
A \mathrel{B} C \mathrel{D} (A \vee B) \supset C \mathrel{C} \supset {\sim} D \mathrel{D} \supset A
TTTTTFT
TTTFTTT
TTFTFFT
TTFFFTT
TFTTTFT
TETETTT
TFFTFFT
TEFEETT
FTTTTFT
FTTFTTT
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FTFTFFT

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FTFFFTT
FFTTTFT
FFTFTTT
FFFTTFT
FFFFTTT
Based on the indirect truth table, we can see that there is no row where all the premises are true and the conclusion is false. Therefore, the argument is valid.
Answer Option 6:
To determine the validity of the argument, we can use the method of indirect truth tables.
Let's assume the premises are true and the conclusion is false.
Statement 1: (A \lor B) \supset C
Statement 2: C ⊃ ~D
Conclusion: \mathsf{D} \supset \mathsf{A} (Negation of the conclusion)
A \mathrel{B} C \mathrel{D} (A \vee B) \supset C \mathrel{C} \supset {\sim} D \mathrel{\sim} \! D \supset {\sim} \! A
TTTTTT
TTTFTFT
TTFTFTT
TTFFFTT
TFTTTT
TETETET
TFFTFTT
TEFEETT
FTTTTF
FTTFTF
FTFTTTF
FTFFTTF
FFTTTTF
FFTFTFF
FFFTTTF
FFFFTTF
In all of the cases where the premises are true, the conclusion is also true. Therefore, the argument is valid.
Final Answer: A (Valid)
Answer Option 7:
To determine the validity of the argument, we can use an indirect truth table.
We have the premises:
1.\; (A\vee B)\supset C
2. C ⊃ ¬D
We want to check if the conclusion D\supset A is valid.
We create a truth table considering all the propositions A, B, C, and D.
A \ B \ C \ D \ (A \lor B) \supset C \ C \supset \neg D \ D \supset A
TTTTTFT
TTTFTTT
TTFTFF
TTFFFTT
TETTTET
TFTFTTT
TEETEE
TFFFFTT
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FTTTTFT FTTFTTT

FTFTTFT FTFFTTT FFTTTFT FFTFTTT FFFTTFT FFFFTTT In all the possible cases where the premises are true, the conclusion is also true. Therefore, the argument is valid. Final Answer: A. Valid ****************** Answer Option 8: To determine the validity of the argument, we can use an indirect truth table. We will assume the premises are true and check if the conclusion must also be true. Assuming A, B, C, and D can have truth values of either true (T) or false (F), we construct the truth table for the premises and the conclusion: Step 1: Construct the truth table for premises: $A \ B \ C \ D \ (A \lor B) \supset C \ C \supset {\sim} D$ TTTTTF TTTFTT TTFTFF TTFFFT TFTTTF TFTFTT TFFTFF TFFFFT FTTTTF FTTFTT FTFTFF FTFFFT FFTTTF FFTFTT FFFTTF FFFFTT Step 2: Check if the conclusion is always true when the premises are true: The conclusion is $D \supset A$. Looking at the truth table, we can see that whenever the premises (A \vee B) \supset C and C \supset \neg D are true, the conclusion D \supset A is also true. There is no row in the truth table where the premises are true and the conclusion is false. Therefore, the argument is valid. Final Answer: A. Valid ******* Answer Option 9: To determine the validity of the argument, we can use an indirect truth table. We will assume that the premises are true and check if the conclusion is always true. Assuming (A \vee B) \supset C and C \supset ~D are true: We will create a truth table that includes A, B, C, D, (A \vee B) \supset C, and C \supset \sim D as columns. $A\mid B\mid C\mid D\mid (A\vee B)\supset C\mid C\supset {\sim} D$ T|T|T|T|T|FTITITIFITIT T|T|F|T|F|FTITIFIFIFIT TIFITITIF TIFITIFITIT T|F|F|T|F|F T|F|F|F|F|T

F|T|T|T|F

F T T F T T	
F T F T F F	
F T F F F T	
F F T T F	
F F T F T T	
F F F T T F	
F F F F T T	

Based on the truth table, we can see that there are instances where the premises are true but the conclusion (D \supset A) is false. Specifically, when A, B, and C are false, and D is true, the conclusion is false.

 $Therefore, the argument is invalid. \ The correct option is \ D: 'Invalid. \ Counterexample \ when \ A, \ B, \ and \ C \ are false \ and \ D \ is \ true'.$

Question 50:

' Select the best English interpretation of the given arguments in predicate logic.

 $\mathsf{Dh}\supset \mathsf{\sim}\mathsf{Pt}$

 $(\forall x)Px \lor (\forall x)Mx$

-Mb / -Dh', A. 'If my headache is dualist state, then your tickle is a physical state. Either everything is physical or everything is mental. But my broken toe is not a mental state. So my headache is not a dualist state.', B. 'If my headache is dualist state, then your tickle is not a physical state. Either everything is physical or everything is mental. But my broken toe is not a mental state. So my headache is not a dualist state, ', C. 'If my headache is dualist state, then your tickle is not a physical state. If everything is physical then everything is mental. But my broken toe is not a mental state. So my headache is not a dualist state, ', D. 'If my headache is dualist state, then your tickle is not a physical state. Everything is either physical or mental. But my broken toe is not a mental state. So my headache is not a dualist state.'

Answer Option 1:

The best English interpretation of the given arguments in predicate logic is:

'If my headache is a dualist state, then your tickle is not a physical state. Either everything is physical or everything is mental. But my broken toe is not a mental state. So my headache is not a dualist state.

Final Answer: C

Answer Option 2:

The given arguments in predicate logic are as follows:

Premises:

1 Dh ¬ --Pt

2. $(\forall x)Px \lor (\forall x)Mx$

3. ~Mb

Conclusion:

4 . Dh

To interpret the given arguments in English, we can break it down as follows:

Premise 1: "If my headache is a dualist state, then your tickle is not a physical state."

Premise 2: "Either everything is physical or everything is mental."

Premise 3: "My broken toe is not a mental state."

Conclusion: "Therefore, my headache is not a dualist state."

The best English interpretation of the given arguments is option C: "If my headache is a dualist state, then your tickle is not a physical state. If everything is physical then everything is mental. But my broken toe is not a mental state. So my headache is not a dualist state."

Final Answer: C

Answer Option 3:

To interpret the given arguments in predicate logic, we need to break down the given statements:

 $\label{eq:definition} \mbox{Dh} \supset \mbox{-Pt: This can be interpreted as "If my headache is a dualist state, then your tickle is not a physical state."$

 $(\forall x) Px \vee (\forall x) Mx \text{: This can be interpreted as "Either everything is physical or everything is mental."}$

~Mb: This can be interpreted as "My broken toe is not a mental state."

The conclusion we need to draw is ~Dh, which can be interpreted as "My headache is not a dualist state."

From the given statements and the conclusion, we can determine the best English interpretation as follows:

"If my headache is a dualist state, then your tickle is not a physical state. Either everything is physical or everything is mental. But my broken toe is not a mental state. So my headache is not a dualist state."

Therefore, the correct answer is A.

Answer Option 4:

Let's break down the given arguments:

1. Dh \supset ~Pt: "If my headache is a dualist state, then your tickle is not a physical state."

2. $(\forall x)Px \lor (\forall x)Mx$: "Either everything is physical or everything is mental."

3. ~Mb: "My broken toe is not a mental state."

4. ~Dh: We need to determine the English interpretation of this statement.

By combining the given information, we can infer:

- If my headache is a dualist state (Dh), then your tickle is not a physical state (~Pt).
- Either everything is physical $(\forall x)Px$ or everything is mental $(\forall x)Mx$.
- My broken toe is not a mental state (~Mb).
- We want to determine the interpretation of ~Dh.

Considering the above information, the best English interpretation of the given arguments is:

'If my headache is a dualist state, then your tickle is not a physical state. Either everything is physical or everything is mental. But my broken toe is not a mental state. So my headache is not a dualist state. Therefore, the correct answer is A. Answer Option 5: Let's break down the given arguments and analyze each statement: 1. Dh ⊃ ~Pt (If my headache is a dualist state, then your tickle is not a physical state.) 2. $(\forall x)Px \lor (\forall x)Mx$ (Either everything is physical or everything is mental.) 3. ~Mb (My broken toe is not a mental state.) 4. ~Dh (Conclusion: My headache is not a dualist state.) From statement 1, we can infer that if my headache is a dualist state, then your tickle is not a physical state. This aligns with option C and D. From statement 2, we understand that either everything is physical or everything is mental. This aligns with option C and D. From statement 3, we know that my broken toe is not a mental state. This aligns with option C and D. Based on the given arguments and the logical interpretations, the best English interpretation is option C: "If my headache is a dualist state, then your tickle is not a physical state. If everything is physical, then everything is mental. But my broken toe is not a mental state. So my headache is not a dualist state." Answer Option 6: To determine the best English interpretation of the given arguments in predicate logic, let's analyze each statement: $\label{eq:decomposition} \mbox{Dh} \supset \mbox{-Pt: This statement can be interpreted as "If my headache is a dualist state, then your tickle is not a physical state."$ $(\forall x)$ Px $\vee (\forall x)$ Mx: This statement can be interpreted as "Either everything is physical or everything is mental." ~Mb: This statement can be interpreted as "My broken toe is not a mental state." The conclusion we want to reach is ~Dh, which can be interpreted as "My headache is not a dualist state." Combining the interpretations of the premises and the conclusion, the best English interpretation of the given argument would be: "If my headache is a dualist state, then your tickle is not a physical state. Either everything is physical or everything is mental. But my broken toe is not a mental state. Therefore, my headache is not a dualist state." Therefore, the correct answer is option C: 'If my headache is a dualist state, then your tickle is not a physical state. If everything is physical, then everything is mental. But my broken toe is not a mental state. So my headache is not a dualist state.' Answer Option 7: Let's analyze the given arguments in predicate logic: Dh ⊃ ~Pt (1) $(\forall x)Px \lor (\forall x)Mx$ (2) ~Mb (3) The argument can be summarized as follows: 1. If my headache is a dualist state, then your tickle is not a physical state. 2. Either everything is physical or everything is mental. 3. My broken toe is not a mental state We need to determine the best English interpretation of these arguments. Looking at the premises: - Premise (1) states that if my headache is a dualist state, then your tickle is not a physical state. - Premise (2) states that either everything is physical or everything is mental. - Premise (3) states that my broken toe is not a mental state. From these premises, we can conclude that: - If everything is physical (as stated in premise 2) and my broken toe is not a mental state (as stated in premise 3), then my headache cannot be a dualist state (as stated in the conclusion). Based on this analysis, the best English interpretation of the argument is: 'If my headache is a dualist state, then your tickle is not a physical state. Either everything is physical or everything is mental. But my broken toe is not a mental state. So my headache is not a dualist state. Therefore, the correct option is A. Answer Option 8: To interpret the given argument, let's break it down step by step: Premise 1: Dh ⊃ ~Pt (If my headache is a dualist state, then your tickle is not a physical state.) Premise 2: $(\forall x)Px \lor (\forall x)Mx$ (Either everything is physical or everything is mental.)

Premise 3: ~Mb (My broken toe is not a mental state.)

Conclusion: ~Dh (Therefore, my headache is not a dualist state.)

Interpretation:

The argument is about the relationship between different states (dualist, physical, mental) and the conditions given about the headache and tickle.

Option B: 'If my headache is dualist state, then your tickle is not a physical state. Either everything is physical or everything is mental. But my broken toe is not a mental state. So my headache is not a dualist state.'

Option C: 'If my headache is dualist state, then your tickle is not a physical state. If everything is physical then everything is mental. But my broken toe is not a mental state. So my headache is not a dualist state.'

Option D: 'If my headache is dualist state, then your tickle is not a physical state. Everything is either physical or mental. But my broken toe is not a mental state. So my headache is not a dualist state.'

Among these options, Option B provides the best English interpretation of the given arguments in predicate logic. It correctly represents the logical relationship between the premises and the conclusion. Final

Answer: B

Answer Option 9:

To determine the best English interpretation of the given arguments in predicate logic, let's analyze the statements:

Dh ⊃ ~Pt: This can be interpreted as "If my headache is a dualist state, then your tickle is not a physical state."

 $(\forall x) Px \lor (\forall x) Mx \text{: This can be interpreted as "Either everything is physical or everything is mental."}$

~Mb: This can be interpreted as "My broken toe is not a mental state."

/~Dh: This can be interpreted as "My headache is not a dualist state."

Now, let's evaluate the options:

Option A: "If my headache is dualist state, then your tickle is a physical state. Either everything is physical or everything is mental. But my broken toe is not a mental state. So my headache is not a dualist state."

This interpretation does not accurately represent the given statements. It incorrectly states that the tickle is a physical state and does not consider the given statement ~Pt.

Option B: "If my headache is dualist state, then your tickle is not a physical state. Either everything is physical or everything is mental. But my broken toe is not a mental state. So my headache is not a dualist state."

This interpretation accurately represents the given statements. It acknowledges that the tickle is not a physical state due to ~Pt and considers the statement ~Mb.

Option C: "If my headache is dualist state, then your tickle is not a physical state. If everything is physical then everything is mental. But my broken toe is not a mental state. So my headache is not a dualist

This interpretation incorrectly assumes that everything is physical implies everything is mental, which is not specified in the given statements.

Option D: "If my headache is dualist state, then your tickle is not a physical state. Everything is either physical or mental. But my broken toe is not a mental state. So my headache is not a dualist state."

This interpretation accurately represents the given statements. It acknowledges that the tickle is not a physical state due to -Pt, acknowledges that everything is either physical or mental, and considers the statement -Mb.

Therefore, the best English interpretation of the given arguments in predicate logic is Option D.

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Question 51:
'Which of the following propositions is an immediate (one-step) consequence in PL of the given premises?
\mathsf{A} \supset (\mathsf{B} \bullet \mathsf{C})
~(B • C)
A \vee (D \vee B)^{\scriptscriptstyle !},\, A.\, {}^{\scriptscriptstyle !}A^{\scriptscriptstyle !},\, B.\, {}^{\scriptscriptstyle !}{}^{\scriptscriptstyle \sim}A^{\scriptscriptstyle !},\, C.\, {}^{\scriptscriptstyle !}B^{\scriptscriptstyle !},\, D.\, {}^{\scriptscriptstyle !}{}^{\scriptscriptstyle \sim}B^{\scriptscriptstyle !}
Answer Option 1:
We can use the rules of inference in propositional logic to determine the immediate consequences of the given premises. Given premises:
1. \ A \supset (B \bullet C)
2. ~(B • C)
3. A v (D v B)
We can start by using the rule of simplification to extract the individual components from the first premise:
4. A (simplification of 1)
Next, we can use the disjunction elimination rule to consider the two cases in the third premise:
Case 1: A
Case 2: D v B
For Case 1 (A):
5. A (reiteration)
For Case 2 (D v B):
6. D v B (reiteration)
Therefore, the immediate consequence in PL of the given premises is 'A'. Option A is the correct answer.
Answer Option 2:
To find the immediate consequence in propositional logic (PL), we need to apply the rules of inference to the given premises.
1. A ⊃ (B • C)
2. ~(B • C)
3. A v (D v B)
By examining the premises, we can see that the immediate consequence is 'A'.
From premise 1 (A \supset (B \bullet C)), we can use the rule of material implication (\rightarrow) to infer \sim(B \bullet C) \supset \simA.
Then, combining premise 2 (\sim(B • C)) with \sim(B • C) \supset \simA, we can use the rule of modus tollens to infer \simA.
Therefore, the immediate consequence in PL of the given premises is '~A'.
Final Answer: B. '~A'
Answer Option 3:
To find an immediate consequence in propositional logic (PL), we need to use valid rules of inference to derive a new proposition from the given premises.
Given premises:
1. \ A \supset (B \bullet C)
2. ~(B • C)
We can start by assuming A is true and try to derive a new proposition using valid rules of inference.
Assume A is true:
4. A (from assumption)
Using Modus Ponens (MP) on premises 1 and 4, we can infer B • C:
5. B • C (from 1, 4, MP)
However, this contradicts premise 2 (~B • ~C), which states ~(B • C). Therefore, assuming A is true leads to a contradiction.
Since assuming A leads to a contradiction, we can conclude that A must be false, i.e., ~A is an immediate consequence.
Final Answer: B. '~A'
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Answer Option 4:
To find the immediate consequence in propositional logic (PL) from the given premises, we need to use logical rules and deductions.
Given premises:
1. A ⊃ (B • C)
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2. ~(B • C)
3. A v (D v B)
To determine the immediate consequence, we can eliminate the disjunction in premise 3 using disjunction elimination.
Premise 3: A v (D v B)
We can split this disjunction into two cases:
Case 1: A
In this case, since A is true, the immediate consequence would also be A.
Case 2: D v B
Here, we have a disjunction, but we cannot determine an immediate consequence without additional information or logical rules.
Therefore, the immediate consequence in PL of the given premises is A.
Final Answer: A
Answer Option 5:
To find an immediate consequence, we can use the method of resolution.
Given premises:
1) A \supset (B \cdot C)
2) ~(B • C)
3) A v (D v B)
To apply resolution, we first convert all the premises into clause form:
1) (~A v (B • C))
2) ~(B • C)
3) (A v D v B)
Now, we can perform the resolution by finding a resolvent that can be derived from two premises.
Using premises 1) and 2), we can resolve on (B • C):
(~A \vee (B \bullet C)) \wedge ~(B \bullet C) gives ~A
Now, we can check if this resolvent matches any of the answer choices.
~A matches the answer choice B.
Therefore, the immediate consequence in PL of the given premises is '~A'.
Answer Option 6:
To determine the immediate consequence in propositional logic (PL) of the given premises, we can use the rules of inference.
The given premises are:
1. \ A \supset (B \bullet C)
2. ~(B • C)
3. A \vee (D \vee B)
From premise 1, the conditional statement A \supset (B \cdot C), we can use the rule of detachment (modus ponens) to infer B \cdot C. However, since \sim (B \cdot C) is already a premise, we can conclude that B \cdot C is false.
From premise 3, A \vee (D \vee B), we can use the rule of addition to add any sentence to a disjunction. In this case, we can add A to the disjunction, resulting in the immediate consequence A.
Therefore, the immediate consequence in PL of the given premises is A. Final Answer: A
Answer Option 7:
Given premises:
1. \ A \supset (B \bullet C)
2. ~(B • C)
3. A \vee (D \vee B)
To find the immediate consequence in propositional logic (PL), we can use the method of proof by contradiction. We assume the negation of the proposition we want to prove and derive a contradiction.
Let's assume the negation of 'A' as ~A. Our goal is to show that this leads to a contradiction.
From premise 1, A \supset (B \bullet C), and the negation \sim(B \bullet C), we can use Modus Tollens to infer \simA:
~(B • C) (Given)
A ⊃ (B • C) (Given)
~A (Modus Tollens)
Now, let's evaluate the third premise, A \vee (D \vee B), using the derived \simA:
~A (Derived)
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$A \lor (D \lor B)$ (Given)
Given that $\neg A$ is true, the only way for $A \lor (D \lor B)$ to be true is if $D \lor B$ is true. Therefore, we can conclude that $D \lor B$ is an immediate consequence of the given premises.
Final Answer: D. '~B'

Answer Option 8:
To determine which proposition is an immediate consequence of the given premises, we can use the rules of inference in propositional logic.

Given premises:

 $1.\ A\supset (B\bullet C)$

2. ~(B • C)

3. A v (D v B)

To find an immediate consequence, we need to apply a valid rule of inference to the premises.

Looking at the premises, we can see that the disjunction in premise 3 (A v (D v B)) suggests that if we can prove one of the disjuncts, it would be an immediate consequence.

Let's examine each option:

A. 'A': This option matches one of the disjuncts in premise 3, but we cannot conclude that 'A' is an immediate consequence without further information or proof.

B. '-A': This option is the negation of one of the disjuncts in premise 3, but again, we cannot conclude that '-A' is an immediate consequence without further information.

C. 'B': This option does not directly match any of the premises or provide a valid inference from the given premises.

D. '~B': This option is the negation of one of the terms in premise 2, which is 'B • C'. Since premise 2 states '~(B • C)', we can conclude that '~B' is an immediate consequence.

Therefore, the correct answer is D. '~B'.

Answer Option 9:

To determine the immediate consequence of the given premises, we can use the process of elimination.

1. $A \supset (B \cdot C)$

2. ~(B • C)

3. A \vee (D \vee B)

Let's analyze each option:

A. 'A':

- This option is not a direct consequence of the premises. Premise 1 states that if A is true, then B • C is also true. However, it does not infer that A must be true.

- This option is not a direct consequence of the premises either. Premise 1 does not provide any information about the negation of A.

- This option is also not a direct consequence of the premises. Premise 1 only establishes a relationship between A and B • C, but it does not determine the value of B itself.

- This option is the immediate consequence of the premises. Premise 2 states that ~(B • C) is true. This means that B • C must be false, so both B and C must be false. Therefore, the negation of B, which is represented as '~B', is an immediate consequence of the premises.

```
'Which of the following propositions is an immediate (one-step) consequence in PL of the given premises?
(N\supset {\sim} O)\supset (O\vee Q)
P\supset (N\supset \sim 0)
N \supset P',~A.~'P',~B.~'N \supset {\sim}O',~C.~'O \lor Q',~D.~'{\sim}P'
To determine the immediate consequence, we will apply the rules of propositional logic to the given premises.
1.\; (N\supset {\sim} O)\supset (O\vee Q)
2. P \supset (N \supset \sim 0)
3. N ⊃ P
From premise 3, we can infer P using the modus ponens rule (\rightarrowE).
Therefore, the immediate consequence is 'P'.
Final Answer: A
Answer Option 2:
To find the immediate consequence in propositional logic (PL) of the given premises, we can apply the rules of implication and substitution.
1.\; (N\supset {\sim} O)\supset (O\vee Q)
2. \ P \supset (N \supset {\sim} O)
3.\ N\supset P
To find the immediate consequence, we will start by applying Modus Ponens (MP), which states that if we have an implication P \supset Q and we also have P, then we can conclude Q.
Using Modus Ponens with premises 2 and 3, we have:
(N \supset \sim O) \supset (O \lor Q) (Premise 1)
\mathsf{P}\supset (\mathsf{N}\supset \mathsf{\sim}\mathsf{O}) \; (\mathsf{Premise}\; \mathsf{2})
N \supset P (Premise 3)
N \supset {\sim} O (Using MP with premises 2 and 3)
Now, we can apply Modus Ponens again with the newly derived statement and premise 1:
(N\supset {\sim} O)\supset (O\vee Q) \ (Premise\ 1)
P \supset (N \supset \sim O) (Premise 2)
N \supset P (Premise 3)
N \supset \sim O (Derived from MP)
O v Q (Using MP with premises 1 and the derived statement)
Therefore, the immediate consequence in PL of the given premises is 'O \vee Q'.
Final Answer: C. 'O v Q'
*******
Answer Option 3:
To determine which proposition is an immediate consequence of the given premises, we can use the rules of inference in propositional logic.
Given premises:
1. (N \supset \sim O) \supset (O \lor Q)
2. P \supset (N \supset \sim 0)
We can apply the modus ponens rule, which states that if we have a conditional statement and its antecedent, we can infer its consequent.
Applying modus ponens to premises 2 and 3, we have:
4. P ⊃ ~O
Now, looking at premises 1 and 4, we can apply another application of modus ponens:
5. O v Q
Therefore, the immediate consequence of the given premises is 'O v Q'.
Final Answer: C. 'O v Q'
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Question 52:

Answer Option 4:

Given premises:
1. $(N \supset \neg 0) \supset (O \lor Q)$
$2. \ P \supset (N \supset \sim 0)$
3. N ⊃ P
Substituting premise 3 into premise 2, we get:
$P \supset (N \supset \sim 0)$ becomes $P \supset (P \supset \sim 0)$
Substituting premise 1 into the new expression, we get:
$(P\supset (P\supset -0))\supset (O\lor Q)$
Now, using the Law of Material Implication, we can simplify the expression as follows:
$(P \supset (P \supset \neg O)) \supset (O \lor Q)$
$\sim P \lor (P \supset \sim O) \supset (O \lor Q)$
$^{\sim}P\vee(^{\sim}P\vee^{\sim}Q)\supset(Q\vee Q)$
By applying the Associative Law and the Commutative Law, we can further simplify the expression:
$(-P \lor -P \lor -Q) \supset (O \lor Q)$
$\sim P \vee \sim O \supset (O \vee Q)$
Therefore, the immediate consequence in PL of the given premises is ' \sim P \vee \sim O \supset (O \vee Q)'.
Final Answer: B. 'N ⊃ ~O'

Answer Option 5:
To determine the immediate consequence, we can use the method of conditional proof.
Given premises:
1. $(N \supset -0) \supset (O \lor Q)$
2. P⊃(N⊃~O)
3. N⊃ P
To derive the immediate consequence, we assume the antecedent of the desired consequence and try to derive the consequent.
Assume P:
From premise 3, using modus ponens, we can conclude N: N (3, MP)
From premise 2 and N. using modus ponens, we can conclude N > ~0: N > ~0 (2.4, MP)
From premise 2 and N, using modus ponens, we can conclude N \supset ~O: N \supset ~O (2, 4, MP) From premise 1 and N \supset ~O, using modus ponens, we can conclude O \lor O: O \lor O (1, 5, MP)
From premise 1 and N \supset ~O, using modus ponens, we can conclude O \vee Q: O \vee Q (1, 5, MP)
From premise 1 and $N \supset \sim O$, using modus ponens, we can conclude $O \lor Q$: $O \lor Q$ (1, 5, MP) Since we have derived $O \lor Q$, which is the consequent of the desired consequence, we can conclude that the immediate consequence is 'O $\lor Q$ '.
From premise 1 and N \supset ~O, using modus ponens, we can conclude O \vee Q: O \vee Q (1, 5, MP)
From premise 1 and $N \supset \neg O$, using modus ponens, we can conclude $O \lor Q$: $O \lor Q$ (1, 5, MP) Since we have derived $O \lor Q$, which is the consequent of the desired consequence, we can conclude that the immediate consequence is ' $O \lor Q$ '. Final Answer: $C : O \lor Q$ '
From premise 1 and N \supset \sim O, using modus ponens, we can conclude O \vee Q: O \vee Q (1, 5, MP) Since we have derived O \vee Q, which is the consequent of the desired consequence, we can conclude that the immediate consequence is 'O \vee Q'. Final Answer: C. 'O \vee Q' Answer Option 6:
From premise 1 and $N \supset \neg O$, using modus ponens, we can conclude $O \lor Q$: $O \lor Q$ (1, 5, MP) Since we have derived $O \lor Q$, which is the consequent of the desired consequence, we can conclude that the immediate consequence is ' $O \lor Q$ '. Final Answer: $C : O \lor Q$ '
From premise 1 and N \supset ~O, using modus ponens, we can conclude O \lor Q: O \lor Q (1, 5, MP) Since we have derived O \lor Q, which is the consequent of the desired consequence, we can conclude that the immediate consequence is 'O \lor Q'. Final Answer: C. 'O \lor Q' Answer Option 6: To find the immediate consequence, we can assume the premises are true and check which proposition must also be true based on the given premises.
From premise 1 and N \supset ~O, using modus ponens, we can conclude O \lor Q: O \lor Q (1, 5, MP) Since we have derived O \lor Q, which is the consequent of the desired consequence, we can conclude that the immediate consequence is 'O \lor Q'. Final Answer: C. 'O \lor Q' Answer Option 6: To find the immediate consequence, we can assume the premises are true and check which proposition must also be true based on the given premises. Premises:
From premise 1 and N \supset ~O, using modus ponens, we can conclude O \lor Q: O \lor Q (1, 5, MP) Since we have derived O \lor Q, which is the consequent of the desired consequence, we can conclude that the immediate consequence is 'O \lor Q'. Final Answer: C. 'O \lor Q' Answer Option 6: To find the immediate consequence, we can assume the premises are true and check which proposition must also be true based on the given premises. Premises: 1. (N \supset ~O) \supset (O \lor Q)
From premise 1 and N \supset \sim O, using modus ponens, we can conclude O \vee Q: O \vee Q (1, 5, MP) Since we have derived O \vee Q, which is the consequent of the desired consequence, we can conclude that the immediate consequence is 'O \vee Q'. Final Answer: C. 'O \vee Q'. Answer Option 6: To find the immediate consequence, we can assume the premises are true and check which proposition must also be true based on the given premises. Premises: 1. (N \supset \sim O) \supset (O \vee Q) 2. P \supset (N \supset \sim O)
From premise 1 and N \supset \sim O, using modus ponens, we can conclude O \vee Q: O \vee Q (1, 5, MP) Since we have derived O \vee Q, which is the consequent of the desired consequence, we can conclude that the immediate consequence is 'O \vee Q'. Final Answer: C. 'O \vee Q' Answer Option 6: To find the immediate consequence, we can assume the premises are true and check which proposition must also be true based on the given premises. Premises: 1. (N \supset \sim O) \supset (O \vee Q) 2. P \supset (N \supset \sim O) 3. N \supset P
From premise 1 and N \supset \sim O, using modus ponens, we can conclude O \vee Q: O \vee Q (1, 5, MP) Since we have derived O \vee Q, which is the consequent of the desired consequence, we can conclude that the immediate consequence is 'O \vee Q'. Final Answer: C. 'O \vee Q'. Answer Option 6: To find the immediate consequence, we can assume the premises are true and check which proposition must also be true based on the given premises. Premises: 1. (N \supset \sim O) \supset (O \vee Q) 2. P \supset (N \supset \sim O) 3. N \supset P Let's analyze the premises:
From premise 1 and N \supset \sim O, using modus ponens, we can conclude O \vee Q: O \vee Q (1, 5, MP) Since we have derived O \vee Q, which is the consequent of the desired consequence, we can conclude that the immediate consequence is 'O \vee Q'. Final Answer: C. 'O \vee Q'. Answer Option 6: To find the immediate consequence, we can assume the premises are true and check which proposition must also be true based on the given premises. Premises: 1. (N \supset \sim O) \supset (O \vee Q) 2. P \supset (N \supset \sim O) 3. N \supset P Let's analyze the premises: From premise 1, we have:
From premise 1 and N \supset \sim O, using modus ponens, we can conclude O \vee Q: O \vee Q (1, 5, MP) Since we have derived O \vee Q, which is the consequent of the desired consequence, we can conclude that the immediate consequence is 'O \vee Q'. Final Answer: C. 'O \vee Q' The consequence is 'O \vee Q'. Answer Option 6: To find the immediate consequence, we can assume the premises are true and check which proposition must also be true based on the given premises. Premises: 1. (N \supset \sim O) \supset (O \vee Q) 2. P \supset (N \supset \sim O) 3. N \supset P Let's analyze the premises: From premise 1, we have: (N \supset \sim O) \supset (O \vee Q)
From premise 1 and N \supset \sim O, using modus ponens, we can conclude O \vee Q: O \vee Q (1, 5, MP) Since we have derived O \vee Q, which is the consequent of the desired consequence, we can conclude that the immediate consequence is 'O \vee Q'. Final Answer: C. 'O \vee Q' The sequence of the desired consequence, we can conclude that the immediate consequence is 'O \vee Q'. From premise C. 'O \vee Q' The sequence of the desired consequence, we can assume the premises are true and check which proposition must also be true based on the given premises. Premises: 1. (N \supset \multimap O) \supset (O \vee Q) 2. P \supset (N \supset \multimap O) 3. N \supset P Let's analyze the premises: From premise 1, we have: (N \supset \multimap O) \supset (O \vee Q) From premise 2, we have:
From premise 1 and N \supset \sim O, using modus ponens, we can conclude O \vee Q: O \vee Q (1, 5, MP) Since we have derived O \vee Q, which is the consequent of the desired consequence, we can conclude that the immediate consequence is 'O \vee Q'. Final Answer: C. 'O \vee Q' The consequence is 'O \vee Q' The consequence is 'O \vee Q'. The conse
From premise 1 and N \supset \sim O, using modus ponens, we can conclude O \lor Q: O \lor Q (1, 5, MP) Since we have derived O \lor Q, which is the consequent of the desired consequence, we can conclude that the immediate consequence is 'O \lor Q'. Final Answer: C. 'O \lor Q'. Answer Option 6: To find the immediate consequence, we can assume the premises are true and check which proposition must also be true based on the given premises. Premises: 1. $(N \supset \neg O) \supset (O \lor Q)$ 2. $P \supset (N \supset \neg O)$ 3. $N \supset P$ Let's analyze the premises: From premise 1, we have: $(N \supset \neg O) \supset (O \lor Q)$ From premise 2, we have: $P \supset (N \supset \neg O)$ From premise 3, we have:
From premise 1 and N \supset \multimap O, using modus ponens, we can conclude O \lor Q: O \lor Q (1, 5, MP) Since we have derived O \lor Q, which is the consequent of the desired consequence, we can conclude that the immediate consequence is 'O \lor Q'. Final Answer: C. 'O \lor Q'. Answer Option 6: To find the immediate consequence, we can assume the premises are true and check which proposition must also be true based on the given premises. Premises: 1. (N \supset \multimap O) \supset (O \lor Q) 2. P \supset (N \supset \multimap O) 3. N \supset P Let's analyze the premises: From premise 1, we have: (N \supset \multimap O) \supset (O \lor Q) From premise 2, we have: P \supset (N \supset \multimap O) From premise 3, we have: N \supset P
From premise 1 and N \supset \sim O, using modus ponens, we can conclude O \vee Q: O \vee Q (1, 5, MP) Since we have derived O \vee Q, which is the consequent of the desired consequence, we can conclude that the immediate consequence is 'O \vee Q'. Final Answer: C. 'O \vee Q'. Answer Option 6: To find the immediate consequence, we can assume the premises are true and check which proposition must also be true based on the given premises. Premises: 1. (N \supset \sim O) \supset (O \vee Q) 2. P \supset (N \supset \sim O) 3. N \supset P Let's analyze the premises: From premise 1, we have: (N \supset \sim O) \supset (O \vee Q) From premise 2, we have: P \supset (N \supset \sim O) From premise 3, we have: N \supset P By applying modus ponens (MP) and simplifying the expressions, we can derive the immediate consequence:
From premise 1 and N \supset -O, using modus ponens, we can conclude O \lor Q (1, 5, MP) Since we have derived O \lor Q, which is the consequent of the desired consequence, we can conclude that the immediate consequence is 'O \lor Q'. Final Answer: C. 'O \lor Q'
From premise 1 and N \supset O , using modus ponens, we can conclude O \lor Q: O \lor Q (1, 5, MP) Since we have derived O \lor Q, which is the consequent of the desired consequence, we can conclude that the immediate consequence is 'O \lor Q'. Final Answer: C. 'O \lor Q'
From premise 1 and N \supset -O, using modus ponens, we can conclude O \lor Q: O \lor Q (1, 5, MP) Since we have derived O \lor Q, which is the consequent of the desired consequence, we can conclude that the immediate consequence is 'O \lor Q'. Final Answer: C. 'O \lor Q' ***********************************

Answer Option 7: To determine the immediate consequence in PL, we need to apply logical rules to the given premises. Let's analyze the premises and the options: Premise 1: $(N \supset \sim O) \supset (O \lor Q)$ Premise 2: $P \supset (N \supset \sim 0)$ Premise 3: N ⊃ P To find the immediate consequence, we can start with the premises and apply logical rules step by step: $\mathord{\sim} (\mathsf{N} \supset \mathord{\sim} \mathsf{O}) \vee (\mathsf{O} \vee \mathsf{Q})$ From Premise 2, using the implication rule (\rightarrow) again, we have: ~P ∨ (N ⊃ ~O) ~N ∨ P Now, let's analyze the options: Option A: 'P' This option matches the conclusion we derived from Premise 3. Option B: 'N ⊃ ~O' This option does not match any of the previous conclusions. Option C: $'O \lor Q'$ This option matches the conclusion we derived from Premise 1. Option D: '~P' This option does not match any of the previous conclusions. Based on our analysis, the immediate consequence is 'P' (Option A). Therefore, the correct answer is A. Answer Option 8: To determine the immediate consequence, we can substitute the given premises into the given proposition and see which option matches the resulting proposition. Given premises: 1. $(N \supset \sim O) \supset (O \lor Q)$ 2. P ⊃ (N ⊃ ~O) $3.\ N\supset P$ Substituting premise 3 into premise 2, we get: $P \supset (N \supset {\sim} O)$ becomes $P \supset (P \supset {\sim} O)$ Substituting premise 1 into the result above, we get: $(P\supset (P\supset {\sim} O))\supset (O\vee Q)$ Now let's consider the options: Substituting 'P' into the resulting proposition above would not lead to a valid substitution. Therefore, option A is not the immediate consequence. Substituting 'N $\supset \sim$ 0' into the resulting proposition above would lead to (N $\supset \sim$ 0) \supset (O \vee Q), which is the same as premise 1. Therefore, option B is the immediate consequence. $Substituting \ 'O \lor Q' \ into the resulting \ proposition \ above \ would \ not \ lead \ to \ a \ valid \ substitution. Therefore, option C \ is \ not the \ immediate \ consequence.$ Substituting '~P' into the resulting proposition above would lead to (~P \supset ~O)) \supset (O \lor Q), which is not a valid substitution. Therefore, option D is not the immediate consequence. Final Answer: B. 'N ⊃ ~O' ******* Answer Option 9: To determine which proposition is an immediate consequence of the given premises, we can use the rules of propositional logic. Given premises: 1. $(N \supset \sim O) \supset (O \lor Q)$ 2. P ⊃ (N ⊃ ~O) To find the immediate consequence, we can try to derive it using logical inference.

Let's consider 'P' as the proposition we want to derive.

From premise 2: $P \supset (N \supset \sim 0)$

By the rule of implication, we can rewrite it as: P \supset (~N \vee ~O)

Now, let's consider premise 3: $\mbox{N} \supset \mbox{P}$

By the rule of implication, we can rewrite it as: ${\sim}N \vee P$

By comparing the above two derived propositions, we have:

 $P \supset (\sim\! N \vee \sim\! O)$ and $\sim\! N \vee P$

Using the transitive property of implication, we can combine these two propositions:

 $P \supset ({\sim}N \vee {\sim}O) \supset ({\sim}N \vee P)$

Now, let's consider premise 1: (N \supset ~O) \supset (O \lor Q)

By the rule of implication, we can rewrite it as: ~(N \supset ~O) \vee (O \vee Q)

Using the rule of contrapositive, we can further simplify it to: N $_{\wedge}$ O $_{\supset}$ (O $_{\vee}$ Q)

By comparing the above derived proposition and the combined proposition, we have:

 $N \wedge O \supset (O \vee Q)$ and $P \supset (\sim\! N \vee \sim\! O) \supset (\sim\! N \vee P)$

Using the law of contraposition, we can conclude that $^{\prime}P^{\prime}$ is an immediate consequence of the given premises.

Therefore, the correct answer is A. 'P'.

Question 53:
'Select the best translation into predicate logic. Earth isn't bigger than Jupiter. (e: Earth; j: Jupiter; Bxy: x is bigger than y)', A. '~Bje', B. 'Bej', C. 'Bje~', D. '~Bej'
Answer Option 1:
The best translation into predicate logic for the given statement "Earth isn't bigger than Jupiter" would be option A: '~Bje'. This translates to "It is not the case that Jupiter is bigger than Earth". Final Answer:
Answer Option 2:
The best translation into predicate logic for the given statement is "¬B(e, j)". This translates to "It is not the case that Earth is bigger than Jupiter".
Therefore, the correct answer is A. '-Bje'.
Answer Option 3:
. The best translation into predicate logic for the statement "Earth isn't bigger than Jupiter" is option A: '-Bje'.
This translates to "It is not the case that Jupiter is bigger than Earth", which accurately represents the given statement. Final Answer: A

Answer Option 4:
The best translation into predicate logic for the statement "Earth isn't bigger than Jupiter" is:
-B(e,j)
This can be read as "It is not the case that Earth is bigger than Jupiter."
Therefore, the correct answer is A. "-Bje".
Answer Option 5:
The best translation into predicate logic for the statement "Earth isn't bigger than Jupiter" is "-Bej".
Explanation:
- Bxy: x is bigger than y
- e: Earth
- j; Jupiter
The statement "Earth isn't bigger than Jupiter" can be translated as "It is not the case that Earth is bigger than Jupiter". Using the predicate Bxy, we can express this as "~B(ej)".
Therefore, the correct answer is D. "-Bej".

Answer Option 6:
The best translation into predicate logic for the statement "Earth isn't bigger than Jupiter" would be option A: ~Bje.
In predicate logic, ~ represents negation, B represents the "is bigger than" relation, and e and j represent Earth and Jupiter, respectively.
Therefore, ~Bje translates to "It is not the case that Earth is bigger than Jupiter." This accurately represents the given statement.
Final Answer: A

Answer Option 7:
The best translation into predicate logic for the statement "Earth isn't bigger than Jupiter" would be option A: '-Bje'.
The predicate logic translation can be broken down as follows:
- Bxy: x is bigger than y
- e: Earth
- j: Jupiter
Therefore, the statement "Earth isn't bigger than Jupiter" can be translated as "~Bje", where "~" represents negation and "B" represents the predicate "is bigger than". Final Answer: A

Answer Option 8:
The best translation into predicate logic for the given statement "Earth isn't bigger than Jupiter" is option A: ~Bje.
Explanation:
We can represent "Earth" by the variable e and "Jupiter" by the variable j. The predicate "x is bigger than y" can be represented by the predicate Bxy.
Therefore, "Earth isn't bigger than Jupiter" can be translated as ~Bje, where ~ represents negation and B represents the predicate "is bigger than".
Final Answer: A
Answer Option 9:
The best translation into anodicate legis for the statement (Forth isolable position), until the U.D. (U.D.) Hereby the best decimal.

e: Earth

j: Jupiter

Bxy: x is bigger than y

The statement "Earth isn't bigger than Jupiter" can be translated to "It is not the case that Earth is bigger than Jupiter". In predicate logic, we can represent this as "~B(ej)".

Option A: "~Bje" - This translates to "It is not the case that Jupiter is bigger than Earth", which is the opposite of what we want.

Option B: "Bej" - This translates to "Earth is bigger than Jupiter", which is the opposite of the given statement.

Option C: "Bje-" - This translates to "Jupiter is bigger than Earth and it is not the case that Earth is bigger than Jupiter", which is not an accurate representation of the given statement.

Option D: "~Bej" - This translates to "It is not the case that Earth is bigger than Jupiter", which accurately represents the given statement.

Therefore, the correct answer is D. "~Bej".

Question 54:

'Identify the antecedent of the following conditional proposition: The Bees winning their first game is necessary for either the Aardvarks or the Chipmunks not winning their first game.', A. 'The Aardvarks do not win their first game.', B. 'The Bees win their first game.', C. 'The Chipmunks do not win their first game.', D. 'Either the Aardvarks or the Chipmunks do not win their first game.'

Answer Option 1:

The antecedent of a conditional proposition is the statement that comes immediately after the "if" or "when" in the sentence. In this case, the antecedent of the given conditional proposition is "The Bees winning their first game". Therefore, the correct answer is B. 'The Bees win their first game.'

Answer Option 2:

The antecedent of a conditional proposition is the statement that comes immediately after the "if" or "when" in the sentence. In this case, the conditional proposition is "The Bees winning their first game is necessary for either the Aardvarks or the Chipmunks not winning their first game."

The antecedent of this conditional proposition is 'The Bees winning their first game.'

Therefore, the correct answer is B. 'The Bees win their first game.'

Answer Option 3:

To identify the antecedent of a conditional proposition, we need to look for the part of the statement that comes before the "if" or "necessary" keyword.

In the given statement, "The Bees winning their first game is necessary for either the Aardvarks or the Chipmunks not winning their first game," the antecedent is the condition that is required for the consequent to be true.

The antecedent is: "The Bees winning their first game."

Therefore, the correct answer is B. 'The Bees win their first game.'

Answer Option 4:

The antecedent of a conditional proposition is the part that comes before the "if" or "only if" in the statement. In this case, the antecedent is the condition that is necessary for the consequent to occur.

The given conditional proposition is: "The Bees winning their first game is necessary for either the Aardvarks or the Chipmunks not winning their first game."

The antecedent in this statement is the condition that is necessary for the consequent to occur. In this case, the antecedent is "The Bees winning their first game."

Therefore, the correct answer is B. 'The Bees win their first game.'

Answer Option 5:

The antecedent of a conditional proposition is the part that comes before the "if" or the "necessary for" statement. In this case, the antecedent is the condition that is necessary for the consequent to be true.

The given conditional proposition is: "The Bees winning their first game is necessary for either the Aardvarks or the Chipmunks not winning their first game."

The antecedent, in this case, is the condition that needs to be true for the consequent to be true. Looking at the statement, the antecedent is: "The Bees winning their first game."

Therefore, the correct answer is B. 'The Bees win their first game.'

Answer Option 6:

The antecedent is the part of a conditional proposition that comes before the "if" or "only if" statement. In this case, the antecedent of the given conditional proposition is the condition that is necessary for the consequence to occur.

The given statement is: "The Bees winning their first game is necessary for either the Aardvarks or the Chipmunks not winning their first game."

The antecedent is the statement that describes the condition that is necessary for the Aardvarks or the Chipmunks to not win their first game.

Looking at the options:

A. 'The Aardvarks do not win their first game.'

- B. 'The Bees win their first game.'
- C. 'The Chipmunks do not win their first game.'
- D. 'Either the Aardvarks or the Chipmunks do not win their first game.'

Among these options, the antecedent is option B, 'The Bees win their first game.'

Final Answer: B

Answer Option 7:

The antecedent of a conditional proposition refers to the part of the statement that comes before the implication ("if" or "->"). In this case, the antecedent is the condition that is necessary for the consequent to

The given conditional proposition is: "The Bees winning their first game is necessary for either the Aardvarks or the Chipmunks not winning their first game."

The antecedent of this conditional proposition is: B. 'The Bees win their first game.'

This is because the statement is asserting that the Bees winning their first game is a necessary condition for the Aardvarks or the Chipmunks not winning their first game. Therefore, the success of the Bees (or lack thereof) is the condition (antecedent) that determines the outcome of the other teams' games. Final Answer: B. 'The Bees win their first game.'

****	******	******	*****

Answer Option 8:

The antecedent of a conditional proposition is the part that comes before the "if" or the "only if" in the statement. In this case, the conditional proposition is "The Bees winning their first game is necessary for either the Aardvarks or the Chipmunks not winning their first game."

The antecedent of this conditional proposition is the event that must occur in order for the consequence to be true. In other words, it is the condition that is required for the outcome to happen.

Looking at the given options:

A. 'The Aardvarks do not win their first game.'

- B. 'The Bees win their first game.'
- C. 'The Chipmunks do not win their first game.'
- D. 'Either the Aardvarks or the Chipmunks do not win their first game.'

The antecedent of the given conditional proposition is option B, 'The Bees win their first game.'

Therefore, the correct answer is B. 'The Bees win their first game.'

Answer Option 9:

The antecedent of a conditional proposition is the part that comes before the "if" or "necessary" in the statement. In this case, the conditional proposition is "The Bees winning their first game is necessary for either the Aardvarks or the Chipmunks not winning their first game."

The antecedent is the condition that is required for the consequent to be true. The consequent, in this case, is "either the Aardvarks or the Chipmunks not winning their first game."

Therefore, the antecedent of the conditional proposition is "The Bees winning their first game."

The correct answer is B. 'The Bees win their first game.'

Question 55:

Identify the antecedent of the following conditional proposition: The Bees win their first game only if either the Aardvarks or the Chipmunks do not win their first games., A. The Aardvarks do not win their first game.', B. 'The Bees win their first game.', C. 'The Chipmunks do not win their first game.', D. 'Either the Aardvarks or the Chipmunks do not win their first games.'

Answer Option 1:

The antecedent of a conditional proposition is the part that comes before the "if" statement. In this case, the antecedent is "Either the Aardvarks or the Chipmunks do not win their first games." Therefore, the correct answer is D. 'Either the Aardvarks or the Chipmunks do not win their first games.'

Answer Option 2:

The antecedent of a conditional proposition is the part that comes before the "if" statement. In this case, the antecedent is "Either the Aardvarks or the Chipmunks do not win their first games." Therefore, the correct answer is D. 'Either the Aardvarks or the Chipmunks do not win their first games.'

Answer Option 3:

The antecedent of a conditional proposition is the statement that comes after the "if" part. In this case, the antecedent is "Either the Aardvarks or the Chipmunks do not win their first games." Therefore, the correct answer is D. 'Either the Aardvarks or the Chipmunks do not win their first games.'

Answer Option 4:

The antecedent of a conditional proposition is the part that comes before the "if" statement. In this case, the antecedent is the statement that is being conditioned upon. In the given conditional proposition, the antecedent is "Either the Aardvarks or the Chipmunks do not win their first games." Therefore, the correct answer is D. 'Either the Aardvarks or the Chipmunks do not win their first games."

Answer Option 5:

The antecedent of a conditional proposition is the part of the statement that comes after the word "if". In this case, the conditional proposition is "The Bees win their first game only if either the Aardvarks or the Chipmunks do not win their first games."

The antecedent in this statement is "Either the Aardvarks or the Chipmunks do not win their first games."

Therefore, the correct answer is D. 'Either the Aardvarks or the Chipmunks do not win their first games.'

Answer Option 6:

The antecedent of a conditional proposition is the part that comes before the "if" statement. In this case, the conditional proposition is "The Bees win their first game only if either the Aardvarks or the Chipmunks do not win their first games." The antecedent is the statement that precedes the "only if" part, which is "Either the Aardvarks or the Chipmunks do not win their first games." Therefore, the correct answer is D. 'Either the Aardvarks or the Chipmunks do not win their first games.'

Answer Option 7:

The antecedent of a conditional proposition is the part that comes before the "only if" or "if" in the statement. In this case, the antecedent is the condition that must be met for the statement to be true. The given conditional proposition is: "The Bees win their first game only if either the Aardvarks or the Chipmunks do not win their first games." The antecedent in this proposition is: "Either the Aardvarks or the Chipmunks do not win their first games." Therefore, the correct answer is option D: 'Either the Aardvarks or the Chipmunks do not win their first games.

Answer Option 8:

The antecedent of a conditional proposition is the part that comes before the "if" in the statement. In this case, the antecedent is the condition that needs to be satisfied for the consequent to occur.

The given statement is: "The Bees win their first game only if either the Aardvarks or the Chipmunks do not win their first games."

The antecedent of this conditional proposition is: "Either the Aardvarks or the Chipmunks do not win their first games."

Therefore, the correct answer is D. "Either the Aardvarks or the Chipmunks do not win their first games."

Answer Option 9:

The antecedent of the conditional proposition is the part that comes before the word "only if." In this case, the antecedent is the condition that must be met for the consequent to be true.

The given statement is: "The Bees win their first game only if either the Aardvarks or the Chipmunks do not win their first games."

The antecedent is the condition that must be met for this statement to be true. In this case, the antecedent is "either the Aardvarks or the Chipmunks do not win their first games."

Therefore, the correct answer is D. "Either the Aardvarks or the Chipmunks do not win their first games."

Question 56:

' Select the best translation into predicate logic: Some athletes don't get injured.', A. '~la', B. '(∃x)(Ax • ~lx)', C. '~Ai', D. '(∃x)(lx ∨ ~Ax)'

Answer Option 1:

The best translation into predicate logic for the statement "Some athletes don't get injured" is option B: $(\exists x)(Ax \bullet \neg lx)$.

In this translation, ($\exists x$) represents the existence of at least one athlete, Ax represents that x is an athlete, and ~lx represents that x does not get injured. Combining these predicates with the conjunction (•) indicates that there exists an athlete who does not get injured. Final Answer: B

Answer Option 2:

The best translation into predicate logic for "Some athletes don't get injured" would be option B: $(\exists x)(Ax \bullet \neg lx)$.

This translation represents the existence (3) of at least one athlete (x) such that they are an athlete (Ax) and they are not injured (-1x). This captures the idea that there exists an athlete who is not injured.

Final Answer: B.

Answer Option 3:

The best translation into predicate logic for the statement "Some athletes don't get injured" is option B: (∃x)(Ax • ~Ix).

This translation states that there exists at least one athlete (represented by x) such that the predicate Ax is true (x is an athlete) and ~Ix is true (x does not get injured). This accurately captures the meaning of the original statement.

Answer Option 4:

The best translation into predicate logic for the statement "Some athletes don't get injured" is B. '(3x)(Ax • ¬lx)' which translates to "There exists an athlete x such that x is an athlete and x does not get injured." This accurately captures the meaning of the original statement by using existential quantification (3x) to indicate the existence of at least one athlete, and the predicates Ax and ~lx to express that the individual is an athlete and does not get injured, respectively. Final Answer: B

Answer Option 5:

The best translation into predicate logic for the statement "Some athletes don't get injured" is:

B. $(\exists x)(Ax \bullet \sim Ix)$

Explanation:

- $(\exists x)$ represents the existential quantifier, indicating that there exists at least one individual x.
- Ax represents the predicate "x is an athlete".
- $\sim\!\!\text{Ix}$ represents the negation of the predicate "x is injured".

Therefore, $(\exists x)(Ax \bullet \neg lx)$ translates to "There exists an individual x who is an athlete and is not injured", which accurately represents the statement "Some athletes don't get injured". Final Answer: B

Answer Option 6:

The correct translation into predicate logic for the statement "Some athletes don't get injured" is:

B. $(\exists x)(Ax \bullet \sim Ix)$

Explanation:

Let's break down the translation:

- $(\exists x)$ indicates the existence of at least one individual.
- Ax represents "x is an athlete."
- ~Ix represents "x does not get injured."

Putting it together, (3x)(Ax • -lx) reads as "There exists an individual x such that x is an athlete and x does not get injured," which accurately captures the meaning of the original statement. Therefore, option

B is the best translation. Final Answer: B

Answer Option 7:

The best translation into predicate logic for the statement "Some athletes don't get injured" is:

 $\mathsf{B.}\; (\exists \mathsf{x})(\mathsf{A}\mathsf{x} \bullet {\sim} \mathsf{I}\mathsf{x})$

Explanation

- $(\exists x)$ denotes the existence of at least one object (athlete) x.
- Ax represents the predicate "x is an athlete."
- Ix represents the predicate "x gets injured."
- The conjunction represents the logical "and" operation.
- \sim Ix denotes the negation of "x gets injured."

So $(\exists x)(Ax \cdot \sim lx)$ translates to "There exists an athlete such that the athlete is not injured," which captures the meaning of "Some athletes don't get injured."

Final Answer: B

Answer Option 8:

The best translation into predicate logic for the statement "Some athletes don't get injured" would be:

B. $(\exists x)(Ax \bullet \sim Ix)$

This translation states that there exists an individual x who is an athlete (Ax), and it is not the case that x is injured (~lx).

Option A (~la) translates to "It is not the case that everyone is injured," which does not capture the meaning of "some athletes" in the original statement.

Option C (~Ai) translates to "It is not the case that everyone is an athlete," which does not capture the meaning of "don't get injured" in the original statement.

Option D ($\exists x$)($Ix \lor -Ax$) translates to "There exists an individual who is injured or not an athlete," which does not accurately convey the idea that some athletes specifically do not get injured.

Therefore, option B is the most accurate translation. Final Answer: B. $(\exists x)(Ax \bullet \sim Ix)$

Answer Option 9:

The best translation into predicate logic for the statement "Some athletes don't get injured" is option B: ($\exists x$)($Ax \bullet \neg Ix$), where Ax represents "x is an athlete" and Ix represents "x gets injured".

Explanation:

- $(\exists x)$ represents the existence of some individual x.
- (Ax \bullet ~Ix) translates to "x is an athlete and x does not get injured".
- By using $(\exists x)(Ax \bullet \neg Ix)$, we are stating that there exists at least one athlete who does not get injured.

Option A: ~la translates to "It is not the case that a specific individual a gets injured". It does not capture the idea of some athletes not getting injured, but rather negates a specific individual's injury.

Option C: ~Ai translates to "It is not the case that all athletes get injured". It does not capture the idea of some athletes not getting injured, but rather negates the statement that all athletes get injured.

Option D: ($\exists x$)($[x \lor \neg Ax]$) translates to "There exists an individual x such that x gets injured or x is not an athlete". This does not capture the idea of "some athletes don't get injured" as it includes the possibility of non-athletes getting injured.

Final Answer: B

Question 57:

'Construct a complete truth table for the following pairs of propositions. Then, using the truth tables, determine whether the statements are logically equivalent or contradictory. If neither, determine whether they are consistent or inconsistent. Justify your answers.

I ≡ ~H and ~(~I ≡ H)', A. 'Logically equivalent', B. 'Contradictory', C. 'Neither logically equivalent nor contradictory, but consistent', D. 'Inconsistent'

Answer Option 1:

To construct a complete truth table, we consider all possible combinations of truth values for the propositions I and H.

I ≡ ~H:

I H ~H I ≡ ~H

TTFF

TFTT

FTFF

. . . .

FFTT

~(~I ≡ H):

 $I H \sim I \equiv H \sim (\sim I \equiv H)$

TTTF

TFFT

FTFT

FFTF

Comparing the truth values of the two statements, we see that \sim (\sim I \equiv H) is the negation of the truth values of I \equiv \sim H.

Therefore, \sim (\sim I \equiv H) is logically equivalent to I \equiv \sim H.

The statements $I \equiv \sim H$ and $\sim (\sim I \equiv H)$ are logically equivalent.

Final Answer: A. 'Logically equivalent'

Answer Option 2:

To construct a complete truth table for the given propositions I = ~H and ~(~I = H), we need to consider all possible combinations of truth values for propositions I and H.

 $I H \sim H \sim I I \equiv \sim H \sim (\sim I \equiv H)$

TTFFFT

TETETE

FTFTTF

FFTTFT

From the truth table, we can see that the values of I = ~H and ~(~I = H) are not the same for all combinations of truth values. Therefore, the statements are not logically equivalent.

Also, there is no row in the truth table where both statements are false. Therefore, the statements are not contradictory.

Hence, the correct answer is C. 'Neither logically equivalent nor contradictory, but consistent'.

Answer Option 3:

To construct a truth table for the given pair of propositions, we need to consider all possible combinations of truth values for the propositions I and H.

I≡~H

 \sim (\sim l \equiv H)

Let's define the columns for I, H, \sim H, \sim I, \sim I \equiv H, and the final expressions.

|---|---|------|

|T|T|F|F|T|F|F|

|T|F|T|F|F|F|T|

|F|T|F|T|F|T|T|

|F|F|T|T|T|F|F|

By evaluating the truth values for each row, we can see that the expressions I = -H and -(-I = H) have opposite truth values for all rows. This means that they are contradictory.

Final Answer: B. 'Contradictory'

Answer Option 4:

To construct a complete truth table, we need to consider all possible combinations of truth values for the propositions I and H.

I ≡ ~H:

I | H | ~H | I ≡ ~H

T|T|F|F $T \mid F \mid T \mid T$ F|T|F|F F|F|T|F ~(~I ≡ H): I | H | ~I | ~I ≡ H | ~(~I ≡ H) TITIFITIF T|F|F|F|T FITITIFIT F|F|T|T|F Comparing the truth values of the two propositions, we can see that they are not logically equivalent (they have different truth values in at least one row of the truth table) and they are not contradictory (they do not have opposite truth values in every row). Therefore, the statements are neither logically equivalent nor contradictory, but they are consistent. Final Answer: C. Neither logically equivalent nor contradictory, but consistent. ******* Answer Option 5: To determine whether the statements are logically equivalent, contradictory, consistent or inconsistent, we need to construct a complete truth table for the given propositions. I ≡ ~H: I | H | ~H | I ≡ ~H TITIFIF T|F|T|T F|T|F|F F|F|T|T ~(~I ≡ H): $I \mid H \mid \sim I \mid \sim I \equiv H \mid \sim (\sim I \equiv H)$ TITIFIFIT TIFIFITIF FITITIFIT F|F|T|T|F Comparing the truth tables of I = ~H and ~(~I = H), we can see that the truth values for both statements are not the same for any combination of truth values for I and H. Therefore, the statements are not However, the statements are also not contradictory because there are some cases where both statements can be true simultaneously (when I is true and H is false). Therefore, the statements are neither logically equivalent nor contradictory, but consistent. The correct answer is C. 'Neither logically equivalent nor contradictory, but consistent'. Answer Option 6: To construct a complete truth table for the given pairs of propositions, we need to consider all possible combinations of truth values for the propositions I and H. $I \equiv {\sim} H$ can be rewritten as $I \equiv {\neg} H$ ~(~I \equiv H) can be simplified to I \equiv H First, let's create a truth table for I $\equiv \neg H$: $I \; H \; \neg H \; I \equiv \neg H$ TTFF TFTT FTFF FFTF

Next, let's create a truth table for I \equiv H:

 $I H I \equiv H$ TTT TFF FTF FFT Now, let's compare the two truth tables: $IH \neg HI \equiv \neg HI \equiv H$ TTEET TFTTF FTFFF FFTFT By comparing the truth values of I = ¬H and I = H, we can see that they have different truth values for some combinations of I and H. This means that the statements are not logically equivalent. Since the statements are not logically equivalent, they cannot be contradictory. Therefore, the correct answer is C. 'Neither logically equivalent nor contradictory, but consistent'. Answer Option 7: To construct the truth table for the given propositions, we need to consider all possible combinations of truth values for I and H. ~I ≡ H ~(~I ≡ H) We have 2 variables (I, H), so there will be $2^2 = 4$ rows in the truth table. The truth table is as follows: |---|----| |T|T|F|F|F|T| |T|F|T|T|F|F| |F|T|F|F|T|T| |F|F|T|F|T|T| Now, let's analyze the truth table to determine whether the statements are logically equivalent, contradictory, consistent, or inconsistent. For I = ~H and ~(~I = H), we can observe that the truth values for the two statements do not match in any row. Therefore, the statements are not logically equivalent. Also, in the truth table, there is no row where both statements are true simultaneously, and there is no row where both statements are false simultaneously. Therefore, the statements are not contradictory. Since the statements are neither logically equivalent nor contradictory, we need to check if they are consistent or inconsistent. In this case, there are rows where both statements are true (Row 2) and rows where both statements are false (Rows 1, 3, 4). Therefore, the statements are consistent. Final Answer: C. Neither logically equivalent nor contradictory, but consistent. Answer Option 8: To construct a complete truth table, we need to consider all possible combinations of truth values for the propositions I and H. I ≡ ~H ~(~I ≡ H) Let's assign truth values for I and H and fill in the truth table: T|T|F|F|F|T| T|F|T|F|T|F| F|T|F|T|F|F| F|F|T|T|T|T| Now, let's analyze the truth values for the statements.

For the statement $I \equiv \sim H$:

- When I is true and H is true, the statement evaluates to false.

- When I is true and H is false, the statement evaluates to true.
- When I is false and H is true, the statement evaluates to false.
- When I is false and H is false, the statement evaluates to true.

For the statement $\sim (\sim I \equiv H)$:

- When I is true and H is true, the statement evaluates to true.
- When I is true and H is false, the statement evaluates to false.
- When I is false and H is true, the statement evaluates to false.
- When I is false and H is false, the statement evaluates to true.

Comparing the truth values of the two statements, we can see that they are not logically equivalent. They have different truth values for some combinations of truth values for I and H. Thus, the statements are neither logically equivalent nor contradictory.

Therefore, the correct answer is:

C. Neither logically equivalent nor contradictory, but consistent.

Answer Option 9:

To construct a complete truth table for the pair of propositions, we need to consider all possible combinations of truth values for the propositions I and H.

First, let's construct the truth table for I \equiv ~H:

```
| | | | H | ~H | | ≡ ~H |
|---|---|
|T|T|F|F|
|T|F|T|T|
|F|T|F|T|
|F|F|T|F|
Next, let's construct the truth table for \sim (\sim I \equiv H):
|---|---|------|
|T|T|F|T|F|
|T|F|F|F|T|
|F|T|T|F|T|
|F|F|T|T|F|
Now, let's compare the truth values of the two statements:
| I ≡ ~H | ~(~I ≡ H) |
|-----|
```

|F|F|
From the truth table, we can see that the two statements have different truth values for some combinations of truth values for I and H. Therefore, the statements are not logically equivalent.

Also, the two statements do not have opposite truth values for any combination of truth values for I and H. Therefore, the statements are not contradictory.

Since the statements are not logically equivalent and not contradictory, they are consistent.

Final Answer: C. Neither logically equivalent nor contradictory, but consistent.

|F|F| |T|T| |T|T|

Question 58:

' Use indirect truth tables to determine whether each set of propositions is consistent. If the set is consistent, choose an option with a consistent valuation. (There may be other consistent valuations.)

 $(T\supset U)\supset (W\supset U)$

 $T\supset \sim (U\supset X)$

(T \supset \sim X) \supset W', A. 'Inconsistent', B. 'Consistent. Consistent valuation when U, W, and X are true and T is false', C. 'Consistent. Consistent valuation when W and X are true and T and U are false', D.

'Consistent. Consistent valuation when T and U are true and W and X are false'

Answer Option 1:

To determine whether the set of propositions is consistent, we can use an indirect truth table.

Let's assign truth values to the propositional variables:

T: True

U: False

W: True

X: False

Using these truth values, we can evaluate each proposition in the set:

 $(\mathsf{T}\supset\mathsf{U})\supset (\mathsf{W}\supset\mathsf{U})\text{: }\mathsf{True}\supset (\mathsf{True}\supset\mathsf{False})=\mathsf{True}\supset\mathsf{False}=\mathsf{False}$

 $\mathsf{T}\supset \mathord{\sim}(\mathsf{U}\supset \mathsf{X})\text{: True}\supset \mathord{\sim}(\mathsf{False}\supset \mathsf{False})=\mathsf{True}\supset \mathord{\sim}\mathsf{True}=\mathsf{True}\supset \mathsf{False}=\mathsf{False}$

 $(T\supset \sim\!\! X)\supset W{:}\; True\supset \sim\!\! False\supset True=True\supset True\supset True=True\supset True=True$

Since there is at least one valuation that makes all the propositions false, the set is inconsistent.

Final Answer: A. 'Inconsistent'

Answer Option 2:

To determine the consistency of the set of propositions, we can use an indirect truth table. We will consider all possible truth value combinations for the propositions and check if there is any row where all propositions are true.

T U W X (T \supset U) \supset (W \supset U) T \supset ~(U \supset X) (T \supset ~X) \supset W Consistency

TFFFTTTF Inconsistent

TFFTTTFT Inconsistent

TFTFTTTT Inconsistent

T F T T T T F T Inconsistent

FTFFTTF Inconsistent

FTFTTTF Inconsistent

FTTFTTTF Inconsistent

FTTTTTF Inconsistent

From the truth table, we can see that there is no row where all propositions are true. Therefore, the set of propositions is inconsistent.

Final Answer: A. Inconsistent

Answer Option 3:

To determine the consistency of the given set of propositions, we can use indirect truth tables.

Let's assign truth values to the propositions T, U, W, and X and evaluate the truth values of each proposition in the set:

 $T\mid U\mid W\mid X\mid (T\supset U)\supset (W\supset U)\mid T\supset \text{-}(U\supset X)\mid (T\supset \text{-}X)\supset W$

T|T|T|T|T|F|T

T|T|T|F|T|F|T

TITIFITITIFIT

TITIFIFITIFIT

TIFITITIFITIT T|F|T|F|F|T|T

TIFIFITIFITIT

TIFIFIFIFITIT

FITITITITIT

FITITIFITIT

F|T|F|T|T|T

F|T|F|F|T|T|T

```
F|F|T|T|T|T
F|F|T|F|T|T
F|F|F|T|T|T
F|F|F|F|T|T|T
From the truth table, we can see that there is no row where all the propositions have the same truth value. This means that the set of propositions is consistent.
Therefore, the correct option is:
B. 'Consistent. Consistent valuation when U, W, and X are true and T is false'
Answer Option 4:
To determine the consistency of the set of propositions, we will create an indirect truth table by assuming all propositions are false and checking if any rows in the truth table result in all propositions being
(T\supset U)\supset (W\supset U)
T\supset \mathord{\sim}(U\supset X)
(T\supset \sim\! X)\supset W
Assuming all propositions are false:
T \; U \; W \; X \; (T \supset U) \supset (W \supset U) \; T \supset {\sim} (U \supset X) \; (T \supset {\sim} X) \supset W
FFFFTTT
FFFTTTT
FFTFTTT
FFTTTT
FTFFTTT
FTFTTTT
FTTFTTT
FTTTTT
TFFFFFT
TFFTFFT
TFTFFFT
TFTTFFT
TTFFTTF
TTFTFF
TTTFTTF
TTTTTF
Since there are rows in the truth table where all propositions are false (e.g., T U W X = T T T T), the set of propositions is inconsistent.
Final Answer: A. 'Inconsistent'
*******
To determine the consistency using indirect truth tables, we assign truth values to each proposition and calculate the truth values of the entire set.
(T\supset U)\supset (W\supset U)
\mathsf{T}\supset \mathord{\sim}(\mathsf{U}\supset \mathsf{X})
(T\supset \sim\!\!X)\supset W
Let's create a truth table for these propositions:
T \cup W \times T \supset U \ (T \supset U) \supset (W \supset U) \ U \supset X \ \text{$\sim$} (U \supset X) \ T \supset \text{$\sim$} X \ (T \supset \text{$\sim$} X) \supset W
TTTTTTFTF
TTTFTTTTT
TTFTTFFFTF
TTFFTFFFTF
TFTTFTTFFT
TFTFFTTTFF
TFFTFTTFFT
TEFEFTTTEF
FTTTTTFTF
```

FTTFTTTTTT FTFTTFFFTF

FTFFTFTT FFTTTTTFFF FFTFTTTTF FFFTFTTFF FFFFFTTFF From the truth table, we can see that there is at least one row where all the propositions in the set are true. Therefore, the set of propositions is consistent. Option B: Consistent. Consistent valuation when U, W, and X are true and T is false. Option C: Consistent. Consistent valuation when W and X are true and T and U are false. Option D: Consistent. Consistent valuation when T and U are true and W and X are false. Final Answer: B, C, D (All options are correct as they provide consistent valuations) Answer Option 6: To determine the consistency of the set of propositions using indirect truth tables, we need to consider all possible truth value combinations for the propositions involved. The given propositions are: 1. $(T\supset U)\supset (W\supset U)$ $2. \ T \supset \mathord{\sim} (U \supset X)$ 3. (T ⊃ ~X) ⊃ W Constructing the indirect truth table: $T~U~W~X~T\supset U~W\supset U~ \sim U~\supset X~T\supset \sim (U\supset X)~T\supset \sim X~(T\supset \sim X)\supset W~Indirect~Truth~Table$ $\mathsf{TTTTTFFTTFT} \to$ $TTTFTTFTTFT \rightarrow$ $TTFTTFFFTTFT \rightarrow$ $\mathsf{TTFFTFFFTTFT} \to$ $\mathsf{TFTTFTTTFFFT} \to$ $\texttt{TFTFFTTTFTFT} \rightarrow$ $\mathsf{TFFTFTFFFFT} \to$ $\texttt{TFFFFTFTFT} \rightarrow$ $FTTTTTFFTTFT \rightarrow$ $\texttt{FTTFTTFTTFT} \rightarrow$ $\mathsf{FTFTTFFFTTFT} \rightarrow$ $\mathsf{FTFFTFFTTFT} \to$ $\texttt{FFTTFTTTFFTT} \rightarrow$ $\mathsf{FFTFFTTTFTTT} \rightarrow$ $\texttt{FFFTFTFTFTT} \rightarrow$ $\texttt{FFFFFTFTTT} \rightarrow$ From the indirect truth table, we can observe that there is at least one row where all the propositions are true. Hence, the set of propositions is consistent. Now, let's find a consistent valuation: From the indirect truth table, we can see that a consistent valuation is when U, W, and X are true, and T is false. This satisfies all the given propositions. Final Answer: B. 'Consistent, Consistent valuation when U. W. and X are true and T is false' ******** Answer Option 7: To determine the consistency of the set of propositions using indirect truth tables, we will create a truth table with columns for each proposition and its negation. $(T\supset U)\supset (W\supset U)$ $\mathsf{T}\supset \mathord{\sim}(\mathsf{U}\supset \mathsf{X})$ $(T\supset {\sim} X)\supset W$ We will fill in the truth table row by row, starting with all propositions being false. $T \; U \; W \; X \; \mathord{\sim} (U \supset X) \; T \supset \mathord{\sim} X \; (T \supset \mathord{\sim} X) \supset W \; (T \supset U) \supset (W \supset U)$ FFFFTTFT FFFTTTFT FFTFTTFT FFTTTTFT TFFFTTFF

TFFTTTFF

```
TETEFETE
TETTEETE
FTFFTTFT
FTFTTTFT
FTTFTTFT
FTTTTTFT
TTFFTFTF
TTFTTFTF
TTTFFFTF
TTTTFFTF
By examining the truth table, we can see that there is at least one row where all the propositions are true, which means the set of propositions is consistent.
Option B: 'Consistent. Consistent valuation when U, W, and X are true and T is false'
Final Answer: B
Answer Option 8:
Let's construct an indirect truth table to determine the consistency of the set of propositions:
T \cup W \times (U \supset X) \ (T \supset \sim X) \ (T \supset U) \ (W \supset U) \ (T \supset U) \supset (W \supset U) \ (T \supset \sim (U \supset X)) \ (T \supset \sim X) \supset W
TTTTFFTTTFT
TTTFFTTTTFT
TTFTFFTFFT
TTFFFTTFFFT
TETTTEETTTT
TFTFTTFTTTT
TFFTTFFFTTT
TFFFTTFFTTT
FTTTTTTTTT
FTTFTTTTTT
FTFTTTTFTTT
FTFFTTTFTTT
FFTTTTTTTT
FFTFTTTTTT
FFFTTTTFTTT
FFFFTTTFTTT
From the truth table, we can see that there is no row where all the propositions are false. Therefore, the set of propositions is consistent.
Option B, Option C, and Option D all state that the set is consistent, so they could be valid choices depending on the specific consistent valuations. However, Option B provides a specific consistent valuation,
which is when U. W. and X are true, and T is false. Therefore, Option B is the correct answer.
Final Answer: B. 'Consistent. Consistent valuation when U, W, and X are true and T is false'
Answer Option 9:
To determine the consistency of the given set of propositions using indirect truth tables, we'll assign truth values to each proposition and evaluate the resulting truth values of the entire set.
(T\supset U)\supset (W\supset U)
T\supset \mathord{\sim}(U\supset X)
(T\supset \sim\!\! X)\supset W
Let's create the truth table considering T, U, W, and X as variables:
T \; U \; W \; X \; (T \supset U) \supset (W \supset U) \; T \supset {\sim} (U \supset X) \; (T \supset {\sim} X) \supset W
```

Since there are four variables, the truth table has 2⁴ = 16 rows. To determine the truth values in each row, we'll fill in the values for the variables T, U, W, and X column by column. For the first column, T, we'll have two rows with T and two rows with F.

For the third column, W, we'll also have $2^2 = 4$ rows alternating T and F.
For the fourth column, X, we'll again have $2^2 = 4$ rows alternating T and F.
Now, let's fill in the truth values based on these assignments:
$W \subset (X \subset U) \hookrightarrow (U \subset W) \subset $
TTTT
TTTF
TTFT
TTFF
TFTT
TFTF
TFFT
TFFF
FTTT
FTTF
FTFT
FTFF
FFTT
FFTF
FFFT
FFFF
Next, let's evaluate the truth values of each proposition in each row:
$T \cup W \times (T \supset U) \supset (W \supset U) T \supset (U \supset X) (T \supset X) \supset W$
TTTTTFF
TTTFTTF
TTFTTFT
TTFFTFT
TETTETT
TETEFTT
TFFTFTT
TFFFFTT
FTTTTTT
FTTFTTT
FTFTTTT
FTFFTTT
FFTTTTT
FFTFTTT
FFFTTTT
FFFFTTT
Based on the truth table, we can see that all the rows have consistent valuations, meaning there are no rows where all the propositions are false. Therefore, the set of propositions is consistent.
Final Answer: B. Consistent. Consistent valuation when U, W, and X are true and T is false.

For the second column, U, we'll have $2^2 = 4$ rows alternating T and F.

Question 59: 'Select the best translation into predicate logic: Caroline is sweet if, and only if, Janet is engaged to Brad.', A. 'Cs ⊃ Ej', B. 'Sc ≡ Ej', C. 'sC ≡ eJ', D. 'Sx ≡ Jy' Answer Option 1: To translate the statement into predicate logic, let the following predicates be used: Cs: Caroline is sweet Ej: Janet is engaged to Brad The statement "Caroline is sweet if, and only if, Janet is engaged to Brad" can be translated as: Therefore, the correct translation into predicate logic is option B: $Sc \equiv Ej$. Answer Option 2: To translate the given statement into predicate logic, we can define the following predicates: Cs: Caroline is sweet Ej: Janet is engaged to Brad The statement "Caroline is sweet if, and only if, Janet is engaged to Brad" can be translated as "Cs if and only if Ej". The correct translation into predicate logic is: B. 'Sc ≡ Ej' Answer Option 3: The best translation into predicate logic for the given statement is: In this translation, 'Sc' represents the predicate "Caroline is sweet" and 'Ej' represents the predicate "Janet is engaged to Brad". The symbol '\(\) denotes the "if and only if" logical operator, indicating that the sweetness of Caroline is true if and only if Janet is engaged to Brad. Final Answer: B Answer Option 4: To translate the given statement into predicate logic, we need to define the predicates and variables. Let's define: C: Caroline is sweet J: Janet is engaged to Brad The statement "Caroline is sweet if, and only if, Janet is engaged to Brad" can be translated as: Using the predicate logic symbols, we can represent this as: Therefore, the best translation into predicate logic is: B. 'Sc ≡ Ej' Answer Option 5: The best translation into predicate logic for the given statement, "Caroline is sweet if, and only if, Janet is engaged to Brad" would be: B. 'Sc ≡ Fi' where Sc represents "Caroline is sweet" and Ej represents "Janet is engaged to Brad". The symbol ≡ is used to represent "if, and only if" in predicate logic, indicating that the two statements are logically equivalent. Therefore, the correct translation is: 'Sc \equiv Ej'. Final Answer: B Answer Option 6: The statement "Caroline is sweet if, and only if, Janet is engaged to Brad" can be translated into predicate logic as follows: Let Cs represent "Caroline is sweet." Let Ej represent "Janet is engaged to Brad." The best translation into predicate logic is option B: 'Sc = Ej' Here, "Sc" represents "Caroline is sweet" and "Ej" represents "Janet is engaged to Brad." The equivalence symbol "=" represents "f and only if," indicating that both conditions must be true for the statement Therefore, the best translation is 'Sc \equiv Ej'. Final Answer: B

Answer Option 7:

The best translation into predicate logic for the given statement "Caroline is sweet if, and only if, Janet is engaged to Brad" is:

B. 'Sc ≡ Ej'

Here, 'Sc' represents the predicate 'Caroline is sweet' and 'Ej' represents the predicate 'Janet is engaged to Brad'.

The symbol '=' denotes the equivalence between the two predicates, indicating that they are true in the same cases and false in the same cases.

So, option B correctly translates the given statement into predicate logic. Final Answer: B

Answer Option 8

The best translation into predicate logic for the given statement, "Caroline is sweet if, and only if, Janet is engaged to Brad" is option B: 'Sc ≡ Ej'.

Explanation:

Let's break down the given statement:

- Caroline is sweet: S(Caroline)
- Janet is engaged to Brad: E(Janet, Brad)

"If, and only if" is translated as the biconditional operator (=) in logic.

Combining the above statements, we get 'Sc = Ej' which translates to "Caroline is sweet if, and only if, Janet is engaged to Brad." This accurately captures the meaning of the original statement.

Therefore, the correct translation is option B.

Answer Option 9:

The best translation into predicate logic for the given statement would be: B. $Sc \equiv Ej$

Explanation:

Let's break down the given statement:

- Caroline is sweet: We can represent this as 'S(o)', where 'S' is the predicate symbol representing "is sweet" and 'c' is the constant symbol representing Caroline.
- Janet is engaged to Brad: We can represent this as 'E(j, b)', where 'E' is the predicate symbol representing "is engaged to" and 'j' and 'b' are constant symbols representing Janet and Brad, respectively.

Now, let's combine these representations using the "if and only if" connective, which is represented by the equivalence symbol '\(\) in predicate logic.

The translation becomes: $'S(c) \equiv E(j, b)'$ where 'Sc' represents "Caroline is sweet" and 'Ej' represents "Janet is engaged to Brad".

Thus, the correct translation is B. $'Sc \equiv Ej'$

Question 60:

'Which of the given formulas of PL is the best symbolization of the following sentence?

Either England's importing beef is a sufficient condition for France's subsidizing agriculture, or China doesn't promote human rights when and only when South Africa supplies diamonds.', A. '(E \supset F) \lor (\checkmark C \supset F) \lor (\lor C \lor F) \lor (\lor C \supset F) \lor (\lor C \lor F) \lor (\lor

S)', B. '(F \supset E) \lor (S \supset \sim C', C. '(E \supset F) \lor \sim (S \equiv C)', D. '(E \supset F) \lor (\sim C \equiv S)'

Answer Option 1:

Let's break down the given sentence into logical statements:

- England's importing beef is a sufficient condition for France's subsidizing agriculture: $\mathsf{E}\supset\mathsf{F}$
- China doesn't promote human rights: ~C
- South Africa supplies diamonds: S

Now, let's evaluate the second part of the sentence: "when and only when" can be symbolized as an equivalence.

So the best symbolization of the given sentence is: (E \supset F) \lor (\sim C \equiv S)

Therefore, the correct option is D. '(E \supset F) \lor (\sim C \equiv S)'

Answer Option 2:

To symbolize the given sentence, let's break it down into two parts:

1. "England's importing beef is a sufficient condition for France's subsidizing agriculture"

This can be symbolized as: $E \supset F$ (If E, then F)

2. "China doesn't promote human rights when and only when South Africa supplies diamonds"

This can be symbolized as: $\sim C \equiv S$ (C is not true if and only if S is true)

Now, let's combine these two symbolizations using the logical operator for "either/or":

The best symbolization for the given sentence is: (E \supset F) \vee (\sim C \equiv S)

Therefore, the correct answer is option D.

Answer Option 3:

Let's break down the sentence:

- 1. Either England's importing beef is a sufficient condition for France's subsidizing agriculture.
- 2. Or China doesn't promote human rights when and only when South Africa supplies diamonds.

Let's symbolize the different parts:

- 1. England's importing beef is a sufficient condition for France's subsidizing agriculture (E \supset F)
- 2. China doesn't promote human rights when and only when South Africa supplies diamonds \sim (C \equiv S)

Now let's combine the two parts using "or":

 $(E\supset F)\vee \sim (C\equiv S)$

The best symbolization of the given sentence is option C: '(E \supset F) \vee ~(S \equiv C)'.

Final Answer: C

Answer Option 4:

To symbolize the given sentence, we can break it down into two parts:

- 1. Either England's importing beef is a sufficient condition for France's subsidizing agriculture.
- 2. Or China doesn't promote human rights when and only when South Africa supplies diamonds.

For the first part, we can symbolize it as 'E \supset F', meaning if England imports beef (E), then France subsidizes agriculture (F).

For the second part, we can symbolize it as '-C = S', meaning China doesn't promote human rights (-C) if and only if (=) South Africa supplies diamonds (S).

Now, to combine both parts using 'or' (\oplus) , the best symbolization of the given sentence is:

 $(\mathsf{E}\supset\mathsf{F})\oplus(\mathsf{\sim}\mathsf{C}\equiv\mathsf{S})$

None of the given options exactly match this symbolization. Therefore, none of the options A, B, C, or D is the best symbolization of the given sentence.

Answer Option 5:

Let's break down the sentence into two parts:

- 1. England's importing beef is a sufficient condition for France's subsidizing agriculture.
- 2. China doesn't promote human rights when and only when South Africa supplies diamonds.

For the first part, we can symbolize it as E \rightarrow F, where E represents "England's importing beef" and F represents "France's subsidizing agriculture." This represents that if England imports beef, then France will subsidize agriculture.

For the second part, we can symbolize it as \sim C = S, where C represents "China promotes human rights" and S represents "South Africa supplies diamonds." This represents that China doesn't promote human rights if and only if South Africa supplies diamonds.

Combining these two parts using the logical operator "or," we get (E \rightarrow F) \vee (~C \equiv S).

Therefore, the best symbolization of the sentence is option D. ' $(E \supset F) \lor (\sim C \equiv S)$ '. Final Answer: D.

Answer Option 6:

Let's break down the sentence into logical statements:

- P: England's importing beef
- Q: France's subsidizing agriculture
- R: China promotes human rights
- S: South Africa supplies diamonds

The sentence can be translated into logical form as follows:

 $(\mathsf{P} \to \mathsf{Q}) \vee (\mathsf{\sim}\mathsf{R} \leftrightarrow \mathsf{S})$

Option A: ${}^{\prime}(E\supset F)\vee ({}^{\sim}C\supset S){}^{\prime}$

This option translates to $(P \to Q) \lor (\sim R \to S)$, which is not equivalent to the original sentence.

Option B: ${}^{\iota}(F\supset E)\vee (S\supset {}^{\iota}C)^{\iota}$

This option translates to (Q \rightarrow P) \vee (S \rightarrow ~R), which is not equivalent to the original sentence.

Option C: $(E \supset F) \lor \sim (S \equiv C)'$

This option translates to $(P \to Q) \lor \neg (S \leftrightarrow R)$, which is not equivalent to the original sentence.

Option D: $'(E \supset F) \lor (\sim C \equiv S)'$

This option translates to $(P \to Q) \lor (\sim R \leftrightarrow S)$, which is exactly the same as the original sentence.

Therefore, the best symbolization of the given sentence is option D: '(E \supset F) \lor (\sim C \equiv S)'. Final Answer: D

Answer Option 7:

To symbolize the given sentence, let's define the following propositions:

- E: England's importing beef
- F: France's subsidizing agriculture
- C: China promotes human rights
- S: South Africa supplies diamonds

The sentence can be broken down into two parts:

- 1. "Either England's importing beef is a sufficient condition for France's subsidizing agriculture" can be symbolized as $(E\supset F)$.
- 2. "China doesn't promote human rights when and only when South Africa supplies diamonds" can be symbolized as ~(C ≡ S).

Now, we need to combine these two parts using the logical connective "or" (represented by \vee).

Option A: $(E \supset F) \lor (\sim C \supset S)$

Option B: $(F \supset E) \lor (S \supset \sim C)$

Option C: $(E \supset F) \lor \sim (S \equiv C)$

Option D: $(E \supset F) \lor (\sim C \equiv S)$

Looking at the options, we can see that Option C: $(E \supset F) \lor \neg (S = C)$ is the best symbolization of the given sentence since it correctly represents both parts of the sentence using the appropriate logical connectives.

Final Answer: C

Answer Option 8:

To symbolize the given sentence, let's break it down into two parts:

- 1. "England's importing beef is a sufficient condition for France's subsidizing agriculture": This can be translated as E \supset F, where E represents "England's importing beef" and F represents "France's subsidizing agriculture".
- 2. "China doesn't promote human rights when and only when South Africa supplies diamonds": This can be translated as ~(C = S), where C represents "China promotes human rights" and S represents "South Africa supplies diamonds".

Now, let's combine these two parts using the "or" operator (v): (E \supset F) \vee ~(C \equiv S).

Comparing this with the given options:

A. '(E \supset F) \lor (\sim C \supset S)': This option does not correctly symbolize the second part of the sentence \sim (C \equiv S).

 $B. \ '(F\supset E) \lor (S\supset {}^{\sim}C)' : This \ option \ reverses \ the \ first \ part \ of \ the \ sentence \ (F\supset E) \ instead \ of \ (E\supset F).$

 $C. \ '(E\supset F) \ \lor \ \sim (S\equiv C)': This \ option \ rearranges \ the \ second \ part \ of \ the \ sentence \ \sim (C\equiv S) \ as \ \sim (S\equiv C).$

D. '(E \supset F) \vee (~C \equiv S)': This option correctly symbolizes the given sentence.

Therefore, the best symbolization of the sentence is option D: '(E \supset F) \vee (~C \equiv S)'.

Answer Option 9:

To symbolize the given sentence, let's break it down into two parts:

- 1. England's importing beef is a sufficient condition for France's subsidizing agriculture.
- 2. China doesn't promote human rights when and only when South Africa supplies diamonds.

Let's analyze each option:

Option A: $(E \supset F) \lor (\sim C \supset S)'$

This option correctly symbolizes the first part of the sentence as $(E \supset F)$, which means "if England imports beef, then France subsidizes agriculture." However, the second part is symbolized as $(-C \supset S)$, which is "if China doesn't promote human rights, then South Africa supplies diamonds." This does not capture the "when and only when" condition. Therefore, Option A is not the best symbolization.

Option B: ${}^{\iota}(F\supset E)\vee (S\supset {}^{\iota}C)^{\iota}$

This option incorrectly reverses the roles of England and France in the first part of the sentence as (F \supset E), which is "if France subsidizes agriculture, then England imports beef." This is not the correct symbolization. The second part is symbolized as (S \supset ~C), which does not capture the "when and only when" condition. Therefore, Option B is not the best symbolization.

Option C: ${}^{\prime}(E\supset F)\vee {}^{\sim}(S\equiv C){}^{\prime}$

This option correctly symbolizes the first part of the sentence as $(E \supset F)$, which means "if England imports beef, then France subsidizes agriculture." However, the second part is symbolized as $\sim (S = C)$, which means "it is not the case that South Africa's supply of diamonds is equivalent to China's promotion of human rights." This does not capture the intended meaning of the second part of the sentence. Therefore, Option C is not the best symbolization.

Option D: $'(E \supset F) \lor (\sim C \equiv S)'$

This option correctly symbolizes the first part of the sentence as (E \supset F), which means "if England imports beef, then France subsidizes agriculture." The second part is symbolized as (~C \equiv S), which means "China's promotion of human rights is not equivalent to South Africa's supply of diamonds." This captures the intended meaning of the second part of the sentence. Therefore, Option D is the best symbolization.

Final Answer: D

Question 61:
'Use the following key to translate the given formula of PL to natural, English sentences.
A: Marina reads a Percy Jackson book.
B: Izzy plays Minecraft.
C: Emily stops working.
D: Russell makes dinner.
E: Ashleigh stops by.
(~A ∨ ~B) ⊃ C', A. 'If neither Marina reads a Percy Jackson book nor Izzy plays Minecraft, then Emily stops working.', B. 'If it's not the case that both Marina reads a Percy Jackson book and Izzy plays
Minecraft, then Emily stops working.', C. 'Either Marina not reading a Percy Jackson book or Izzy not playing Minecraft entails that Emily stops working.', D. 'Emily stopping working is a sufficient condition
either Marina not reading a Percy Jackson book or Izzy not playing Minecraft.'
Answer Option 1:
The given formula $(-A \lor -B) \supset C$ can be translated as follows:
~A: Marina does not read a Percy Jackson book.
~B: Izzy does not play Minecraft.
C: Emily stops working.
(~A ∨ ~B) ⊃ C: If neither Marina reads a Percy Jackson book nor Izzy plays Minecraft, then Emily stops working.
Therefore, the correct translation is option A: 'If neither Marina reads a Percy Jackson book nor Izzy plays Minecraft, then Emily stops working.' Final Answer: A
•••••••••••
Answer Option 2:
The given formula is $(-A \lor -B) \supset C$.
Let's translate this formula into English sentences using the provided key:
~A: Marina does not read a Percy Jackson book.
~B: Izzy does not play Minecraft.
C: Emily stops working.
$(-A \lor -B) \supset C$ can be translated as:
"If neither Marina reads a Percy Jackson book nor Izzy plays Minecraft, then Emily stops working." (Option A)
Therefore, the correct answer is A.

Answer Option 3:
To translate the given formula $(-A \lor -B) \supset C$ to natural, English sentences, we can use the following key:
A: Marina reads a Percy Jackson book.
B: Izzy plays Minecraft.
C: Emily stops working.
The formula $(-A \lor -B) \supset C$ can be translated as "If neither Marina reads a Percy Jackson book nor Izzy plays Minecraft, then Emily stops working."
Therefore, the correct translation is: A. "If neither Marina reads a Percy Jackson book nor Izzy plays Minecraft, then Emily stops working."

Answer Option 4:
To translate the given formula $(-A \lor -B) \supset C$ to natural, English sentences, we use the following key:
A: Marina reads a Percy Jackson book.
B: Izzy plays Minecraft.
C: Emily stops working.
The given formula can be translated as follows:
$(-A \lor -B) \supset C$
"If neither Marina reads a Percy Jackson book nor Izzy plays Minecraft, then Emily stops working."
Therefore, the correct translation is option A: "If neither Marina reads a Percy Jackson book nor Izzy plays Minecraft, then Emily stops working."
••••••••••
Answer Option 5:
To translate the given formula (~A ∨ ~B) ⊃ C to natural, English sentences, we can use the provided key: A: Marina reads a Percy Jackson book. B: Izzy plays Minecraft. C: Emily stops working.
The formula can be translated as follows:
(-A ∨ -B) ⊃ C

"If neither Marina reads a Percy Jackson book nor Izzy plays Minecraft, then Emily stops working."

This can be read as:

Answer Option 6: To translate the given formula $(\sim A \vee \sim B) \supset C$ to natural, English sentences using the provided key: A: Marina reads a Percy Jackson book. B: Izzy plays Minecraft. C: Emily stops working. $(\sim A \vee \sim B) \supset C$ Using the key: ~A: Marina does not read a Percy Jackson book. ~B: Izzy does not play Minecraft. C: Emily stops working. The natural, English sentence translation of the given formula is: 'If neither Marina reads a Percy Jackson book nor Izzy plays Minecraft, then Emily stops working.' Therefore, the correct option is A. ******* Answer Option 7: To translate the given formula $(\sim A \lor \sim B) \supset C$ into natural English sentences, we can break it down step by step: Step 1: ~A v ~B This translates to "Marina does not read a Percy Jackson book or Izzy does not play Minecraft." Step 2: $(\sim A \vee \sim B) \supset C$ This translates to "If Marina does not read a Percy Jackson book or Izzy does not play Minecraft, then Emily stops working." So, the correct translation of the formula is: A. 'If neither Marina reads a Percy Jackson book nor Izzy plays Minecraft, then Emily stops working.' Answer Option 8: To translate the given formula $(-A \lor -B) \supset C$ into natural, English sentences, we will assign the following meanings to the propositional variables: A: Marina reads a Percy Jackson book. B: Izzy plays Minecraft. C: Emily stops working. $(-A \lor -B) \supset C \text{ can be translated as "If neither Marina reads a Percy Jackson book nor Izzy plays Minecraft, then Emily stops working."$ Option A: "If neither Marina reads a Percy Jackson book nor Izzy plays Minecraft, then Emily stops working." This accurately represents the given formula and its translation. Option B: "If it's not the case that both Marina reads a Percy Jackson book and Izzy plays Minecraft, then Emily stops working." This translation does not capture the negation of A and B individually, but rather negates their conjunction. Option C: "Either Marina not reading a Percy Jackson book or Izzy not playing Minecraft entails that Emily stops working." This translation does not accurately capture the conditional implication in the original Option D: "Emily stopping working is a sufficient condition for either Marina not reading a Percy Jackson book or Izzy not playing Minecraft." This translation does not accurately capture the conditional implication in the original formula. Therefore, the correct translation is: A. "If neither Marina reads a Percy Jackson book nor Izzy plays Minecraft, then Emily stops working." Answer Option 9: To translate the given formula $(\neg A \lor \neg B) \supset C$ into natural English sentences, we can use the following key: A: Marina reads a Percy Jackson book. B: Izzy plays Minecraft. C: Emily stops working. The formula can be translated as follows: $(\text{~A} \lor \text{~B}) \supset C$ Option A: 'If neither Marina reads a Percy Jackson book nor Izzy plays Minecraft, then Emily stops working.' This translation correctly captures the logical meaning of the formula. Option B: 'If it's not the case that both Marina reads a Percy Jackson book and Izzy plays Minecraft, then Emily stops working.' This translation is not an accurate representation of the formula because it implies that either Marina reads a Percy Jackson book or Izzy plays Minecraft, which is not the case. Option C: 'Either Marina not reading a Percy Jackson book or Izzy not playing Minecraft entails that Emily stops working.'

Therefore, the correct translation is option A: 'If neither Marina reads a Percy Jackson book nor Izzy plays Minecraft, then Emily stops working.'

This translation is not an accurate representation of the formula because it suggests that either Marina not reading a Percy Jackson book or Izzy not playing Minecraft is sufficient for Emily to stop working, which is not the case.

Option D: 'Emily stopping working is a sufficient condition for either Marina not reading a Percy Jackson book or Izzy not playing Minecraft.'

This translation is not an accurate representation of the formula because it suggests that Emily stopping working is sufficient for Marina not reading a Percy Jackson book or Izzy not playing Minecraft, which

Final Answer: A

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'Identify the conclusion of the following argument. There are those who use predicate variables; the values of these variables are attributes. Predicates are wanted in all sentences regardless of whether there are attributes to refer to. The dummy predicate 'F' is wanted generally for expository purposes. If we are going to quantify over attributes and refer to them, then clarity is served by using recognizable variables and distinctive names for the purposes, and not mixing these up with the predicates. (Quine, Philosophy of Logic)', A. 'There are those who use predicate variables; the values of these variables are attributes.', B. 'Predicates are wanted in all sentences regardless of whether there are attributes to refer to.', C. 'The dummy predicate 'F' is wanted generally for expository purposes.', D. 'If we are going to quantify over attributes and refer to them, then clarity is served by using recognizable variables and distinctive names for the purposes, and not mixing these up with the predicates.'

Answer Option 1:

The conclusion of the argument is:

D. If we are going to quantify over attributes and refer to them, then clarity is served by using recognizable variables and distinctive names for the purposes, and not mixing these up with the predicates.

Answer Option 2:

The conclusion of the argument is:

D. 'If we are going to quantify over attributes and refer to them, then clarity is served by using recognizable variables and distinctive names for the purposes, and not mixing these up with the predicates.'

Answer Option 3:

The conclusion of the argument is: D. 'If we are going to quantify over attributes and refer to them, then clarity is served by using recognizable variables and distinctive names for the purposes, and not mixing these up with the predicates.'

Answer Option 4:

The conclusion of the argument is: D. 'If we are going to quantify over attributes and refer to them, then clarity is served by using recognizable variables and distinctive names for the purposes, and not mixing these up with the predicates.'

Answer Option 5:

The conclusion of the argument is: D. 'If we are going to quantify over attributes and refer to them, then clarity is served by using recognizable variables and distinctive names for the purposes, and not mixing these up with the predicates.' This conclusion is derived from the premises stated in the argument.

Answer Option 6:

The conclusion of the argument is: D. 'If we are going to quantify over attributes and refer to them, then clarity is served by using recognizable variables and distinctive names for the purposes, and not mixing these up with the predicates.' This conclusion is derived from the premises provided in the argument.

Answer Option 7:

The conclusion of the argument is: D. 'If we are going to quantify over attributes and refer to them, then clarity is served by using recognizable variables and distinctive names for the purposes, and not mixing these up with the predicates.' This is the main point or claim that the argument is making based on the statements presented.

Answer Ontion 8

The conclusion of the argument is: D. 'If we are going to quantify over attributes and refer to them, then clarity is served by using recognizable variables and distinctive names for the purposes, and not mixing these up with the predicates.' This statement summarizes the main point made in the argument about the importance of using clear and distinct variables and names when quantifying over attributes.

Answer Option 9:

The conclusion of the argument can be identified as: D. 'If we are going to quantify over attributes and refer to them, then clarity is served by using recognizable variables and distinctive names for the purposes, and not mixing these up with the predicates.' This conclusion summarizes the main point made in the argument regarding the use of variables and names in relation to attributes and predicates for clarity.

'Use the following key to translate the given formula of PL to natural, English sentences. A: Marina reads a Percy Jackson book. B: Izzy plays Minecraft. C: Emily stops working. D: Russell makes dinner. ~C v D', A. 'If Emily doesn't stop working then Russell makes dinner.', B. 'Emily stops working unless Russell makes dinner.', C. 'Emily stops working unless Russell doesn't make dinner.', D. 'Emily doesn't stop working unless Russell makes dinner. Answer Option 1: ~C ∨ D The given formula can be translated as "If Emily doesn't stop working, then Russell makes dinner." Therefore, the correct translation is: A. "If Emily doesn't stop working, then Russell makes dinner." ******* Answer Option 2: Using the given key, let's translate the formula ${\sim}C \lor D$ into a natural language sentence. ~C \lor D can be translated as "Either Emily doesn't stop working or Russell makes dinner." So, the correct translation is: A. "If Emily doesn't stop working then Russell makes dinner." ****************** Answer Option 3: Using the given key, we can translate the formula ${\sim}C \lor D$ as: "If Emily does not stop working, then Russell makes dinner." Therefore, the correct translation is: A. "If Emily doesn't stop working then Russell makes dinner." ******* The given formula ${\sim}C \vee D$ can be translated using the given key as follows: ~C: Emily doesn't stop working. v: or D: Russell makes dinner. So, the translated sentence would be: "If Emily doesn't stop working or Russell makes dinner." Therefore, the correct translation is: A. "If Emily doesn't stop working then Russell makes dinner." ******* Answer Option 5: Using the given key, we can translate the formula ${\sim}C \vee D$ as follows: ${\sim}\text{C} \vee \text{D}$ can be translated as 'If Emily doesn't stop working, then Russell makes dinner'. Therefore, the correct translation is: A. 'If Emily doesn't stop working, then Russell makes dinner.' ****************** Answer Option 6: ~C v D can be translated to English as 'If Emily doesn't stop working then Russell makes dinner.' Therefore, the correct answer is A. 'If Emily doesn't stop working then Russell makes dinner.' Answer Option 7: Let's translate the given formula $\sim C \lor D$ using the provided key: ~C translates to "Emily doesn't stop working." v translates to "or."

Question 63:

D translates to "Russell makes dinner."

Therefore, the translation of ~C ∨ D is "Emily doesn't stop working or Russell makes dinner."
The correct option is A. 'If Emily doesn't stop working then Russell makes dinner."

Answer Option 8:
Using the given key:
A: Marina reads a Percy Jackson book.
B: Izzy plays Minecraft.
C: Emily stops working.
D: Russell makes dinner.
E: Ashleigh stops by.
The formula ~C ∨ D can be translated to English as "If Emily doesn't stop working then Russell makes dinner."
Final Answer: A. 'If Emily doesn't stop working then Russell makes dinner.'
Answer Option 9:
Let's break down the formula ~C v D:
~C represents "Emily doesn't stop working."
D represents "Russell makes dinner."
Using the given key:
A: Marina reads a Percy Jackson book.
B: Izzy plays Minecraft.
C: Emily stops working.
D: Russell makes dinner.
E: Ashleigh stops by.
The formula ~C v D can be translated to English as "If Emily doesn't stop working then Russell makes dinner." Therefore, the correct option is A: "If Emily doesn't stop working then Russell makes dinner."

Question 64:
'Select the best translation into predicate logic: Abdul and Cleopatra are Egyptian.', A. 'Ea • Ec', B. 'Ea • c', C. 'Ae ∨ Ce', D. 'Ex • Ey'
Answer Option 1:
The best translation into predicate logic for the statement "Abdul and Cleopatra are Egyptian" would be:
Ea • Ec
Here, Ea represents "Abdul is Egyptian" and Ec represents "Cleopatra is Egyptian".
So, the correct answer is A. 'Ea • Ec'.
Answer Option 2:
The best translation into predicate logic for the statement "Abdul and Cleopatra are Egyptian" would be option B: 'Ea • c'. This translates to "There exists an individual Abdul (Ea) and an individual Cleopatra
(c) who are both Egyptian." Final Answer: B
Answer Option 3:
The best translation into predicate logic for the statement "Abdul and Cleopatra are Egyptian" is:
A. 'Ea • Ec'
Explanation:
- Ea represents "Abdul is Egyptian"
- Ec represents "Cleopatra is Egyptian"
The symbol ◆ is used to represent the logical AND between the two predicates.
Therefore, the correct answer is A.
Answer Option 4:
The best translation into predicate logic for the statement "Abdul and Cleopatra are Egyptian" would be:
Option C: 'Ae ∨ Ce'
This translation states that there exists an individual who is Abdul (Ae) or there exists an individual who is Cleopatra (Ce), indicating that at least one of them is Egyptian. Final Answer: C
Answer Option 5:
The best translation into predicate logic for the given statement "Abdul and Cleopatra are Egyptian" would be:
Ea \wedge Ec, which translates to "There exists an individual Abdul (Ea) and there exists an individual Cleopatra (Ec), such that both Abdul and Cleopatra are Egyptian."
So, the correct option would be A. 'Ea • Ec'
Answer Option 6:
The best translation into predicate logic for the statement "Abdul and Cleopatra are Egyptian" would be:
A. 'Ea • Ec'
Explanation:
In predicate logic, Ea represents "Abdul is Egyptian" and Ec represents "Cleopatra is Egyptian". The symbol • denotes the logical conjunction "and" between the two statements. Therefore, 'Ea • Ec' correct!
translates the statement "Abdul and Cleopatra are Egyptian".
Final Answer: A
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Answer Option 7:
The best translation into predicate logic for the statement 'Abdul and Cleopatra are Egyptian' would be:
A. 'Ea • Ec'
Explanation:
- Ea represents the predicate "a is Egyptian", where 'a' is the variable representing Abdul.
- Ec represents the predicate "c is Egyptian", where 'c' is the variable representing Cleopatra.
- The • symbol represents the logical conjunction 'and', so Ea • Ec translates to "Abdul is Egyptian and Cleopatra is Egyptian".
Therefore, the correct answer is A. 'Ea • Ec'.
Answer Option 8:
The best translation into predicate logic for the statement "Abdul and Cleopatra are Egyptian" is option B: "Ea • c'.
Explanation:

- 'Ea' represents the predicate "Abdul is Egyptian", where 'a' is the variable representing Abdul.

- 'c' represents the predicate "Cleopatra is Egyptian".
Using the logical operator '•' (and), we can combine the two predicates to form the statement "Abdul and Cleopatra are Egyptian". The resulting translation is 'Ea • c'.
Therefore, the correct answer is B.
Answer Option 9:
The best translation into predicate logic for the statement "Abdul and Cleopatra are Egyptian" is:
Ea • Ec
Here, Ea represents "Abdul is Egyptian" and Ec represents "Cleopatra is Egyptian".
Option A: 'Ea • Ec' is the correct translation.
Option B: 'Ea • c' is incorrect as it does not specify Cleopatra being Egyptian.
Option C: 'Ae v Ce' is incorrect as it represents "Either Abdul or Cleopatra is Egyptian" but the original statement states that both are Egyptian.
Option D: 'Ex • Ey' is incorrect as it uses existential quantifiers without specifying the individuals involved.

Final Answer: A

Question 65:
'Select the best translation into predicate logic: Sheena is a punk rocker.', A. 'Sx', B. 'xS', C. 'sP', D. 'Ps'
======================================
Answer Option 1:
The best translation into predicate logic for the statement "Sheena is a punk rocker" would be option A, 'Sx'. This translates to "There exists an x such that x is Sheena". This accurately represents the
statement in predicate logic. Final Answer: A
Answer Option 2:
The best translation into predicate logic for the sentence "Sheena is a punk rocker" would be option C: 'sP'. Here, 's' represents Sheena as the subject and 'P' represents the predicate "is a punk rocker". Fir
Answer: C

Answer Option 3:
The best translation into predicate logic for the statement "Sheena is a punk rocker" would be option A: 'Sx', where 'S' represents the predicate "is a punk rocker" and 'x' represents the individual constant of
variable representing Sheena. This translation indicates that Sheena (represented by 'x') satisfies the predicate 'S', meaning Sheena is a punk rocker. Final Answer: A
Answer Option 4:
The best translation into predicate logic for the statement "Sheena is a punk rocker" would be option A: 'Sx'. Here, 'S' represents the predicate symbol for "is a punk rocker" and 'x' represents the variable for
the individual, in this case, Sheena. Therefore, 'Sx' translates to "Sheena is a punk rocker." Final Answer: A
Answer Option 5:
The best translation into predicate logic for the sentence "Sheena is a punk rocker" would be option C: 'sP'.
In this translation, 's' represents Sheena as an individual and 'P' represents the predicate "is a punk rocker". Therefore, 'sP' correctly expresses the statement that Sheena satisfies the predicate "Punk
Rocker". Final Answer: C
Answer Option 6:
The best translation into predicate logic for the statement "Sheena is a punk rocker" would be option C: 'sP'. Here, 's' represents the individual constant symbol for Sheena, and 'P' represents the predicate
symbol for "is a punk rocker". So, 'sP' can be read as "Sheena is a punk rocker".
Answer Option 7:
The best translation into predicate logic for the statement "Sheena is a punk rocker" would be option C: 'sP'. Here, 's' represents Sheena as an individual or object, and 'P' represents the predicate "is a punk rocker" would be option C: 'sP'.
rocker". Therefore, the statement can be translated as 'sP', indicating that Sheena is the subject and 'P' is the predicate that describes her. Final Answer: C
Answer Option 8:
The best translation into predicate logic for the sentence "Sheena is a punk rocker" would be option C: 'sP'.
In this translation, 's' represents Sheena as an individual and 'P' represents the property of being a punk rocker. The statement 'sP' asserts that Sheena has the property of being a punk rocker.
Option A ('Sx') and option B ('xS') do not properly represent the relationship between Sheena and being a punk rocker. Option D ('Ps') would imply that being a punk rocker has the property of being Sheen

which is not the intended meaning.

Therefore, the correct answer is C: 'sP'.

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Answer Option 9:

The best translation into predicate logic for the sentence "Sheena is a punk rocker" would be option D: 'Ps'. Here, 'P' represents the predicate or property of being a punk rocker, and 's' represents the individual Sheena. So, 'Ps' translates to "Sheena is a punk rocker." Final Answer: D

Question 66:
'Select the best translation into predicate logic. Kevin is introduced to José by Wilma. (j: José; k: Kevin; w: Wilma; lxyz: x introduces y to z)', A. 'lwjk', B. 'ljkw', C. 'lkjw', D. 'lwkj'
Answer Option 1:
The best translation into predicate logic would be 'lwjk'. This translates to 'Wilma introduces Kevin to José'. Therefore, the correct answer is A.

Answer Option 2:
The best translation into predicate logic would be option B: 'ljkw'. This translates to "José is introduced to Kevin by Wilma."

Answer Option 3:
The best translation into predicate logic would be 'lwjk'. In this translation, 'lwjk' represents the statement that Wilma introduces Kevin to José. Final Answer: A

Answer Option 4:
The best translation into predicate logic for the given statement would be 'lwjk'.
This can be understood as "Wilma introduces Kevin to José", where I represents the relation of introduction, w represents Wilma, j represents José, and k represents Kevin.
Therefore, the correct answer is A. 'lwjk'.

Answer Option 5:
The best translation into predicate logic would be 'lwjk', which represents 'Wilma introduces Kevin to José'.
In this case, the predicate 'I' stands for 'introduces', and 'w', 'j', and 'k' represent Wilma, José, and Kevin, respectively.
Therefore, the correct answer is A. 'lwjk'.
Answer Option 6:
The best translation into predicate logic for the statement "Kevin is introduced to José by Wilma" would be "lwjk", where:
- I: Introduces
- w: Wilma
- j; José
- k: Kevin
This translates to "Wilma introduces Kevin to José."
Therefore, the correct answer is A. 'lwjk'.

Answer Option 7:
The best translation into predicate logic for the given statement "Kevin is introduced to José by Wilma" would be: lwjk.
In this translation, lwjk represents the predicate "Wilma introduces Kevin to José", where:
- I: Introduces
- w: Wilma
- j: José
- k: Kevin
Therefore, the correct answer is A. 'lwjk'.

Answer Option 8:
The best translation into predicate logic for the given statement is: 'lwjk' This translates to "Wilma introduces Kevin to José." Therefore, the correct answer is A.

Answer Option 9:
The best translation into predicate logic of the sentence "Kevin is introduced to José by Wilma" would be: 'lwjk'.
Here's the breakdown:
- Iwik: Wilma introduces Kevin to José.
In this translation, 'I' represents the introduction relation, 'w' represents Wilma, 'j' represents José, and 'k' represents Kevin. Therefore, option A, 'lwjk', is the correct translation. Final Answer: A.
m and a anisación, i represente a e introduction relation, w represente venna, y represente dose, and it represente tecnile revint. Therefore, option A, twy, is the contest translation. Fillal Aliswel. A.

Question 67:

'Select the best translation into predicate logic: If Eileen plays fiddle then Sherri sings.', A. 'Fe ∨ Ss', B. 'eF ⊃ Ss', C. 'Fe ∨ Es', D. 'Fe ⊃ Ss'

Answer Option 1:

The best translation into predicate logic for the statement "If Eileen plays fiddle then Sherri sings" is option D: 'Fe \supset Ss'.

In this translation, Fe represents the proposition "Eileen plays fiddle" and Ss represents the proposition "Sherri sings". The symbol ' \supset ' represents the implication operator, which is used to express "if...then" statements in predicate logic.

Therefore, option D correctly represents the given statement in predicate logic. Final Answer: D.

Answer Option 2:

The best translation into predicate logic for the given statement "If Eileen plays fiddle then Sherri sings" would be option D: 'Fe \supset Ss'.

This translation represents the conditional relationship between Eileen playing the fiddle (Fe) and Sherri singing (Ss), where the arrow symbol 'b' denotes implication or "if...then" in predicate logic. Final

Answer: D

Answer Option 3:

To translate the given statement into predicate logic, let's assign the following predicates:

P: Eileen plays fiddle

Q: Sherri sings

The statement "If Eileen plays fiddle then Sherri sings" can be translated as "P implies Q" or "If P, then Q".

Using the logical operator symbols:

P implies Q can be represented as $P \rightarrow Q$.

Therefore, the correct translation into predicate logic is:

'Fe ⊃ Ss'

Final Answer: D

Answer Option 4:

The best translation into predicate logic for the statement "If Eileen plays fiddle then Sherri sings" is option D: 'Fe \supset Ss'.

In predicate logic, we use the symbol ' \supset ' to represent implication (if...then). 'Fe' represents the proposition "Eileen plays fiddle" and 'Ss' represents the proposition "Sherri sings". The statement can be translated as "If Eileen plays fiddle, then Sherri sings" which can be represented as 'Fe \supset Ss'. Therefore, option D is the correct translation.

Answer Option 5:

The best translation into predicate logic for the given statement "If Eileen plays fiddle then Sherri sings" would be option D: 'Fe \supset Ss'.

In predicate logic, we represent "plays fiddle" as the predicate Fe and "sings" as the predicate Ss. The implication "If Eileen plays fiddle then Sherri sings" is symbolized as Fe \supset Ss, which represents the logical statement that if Eileen plays fiddle (Fe) is true, then Sherri sings (Ss) must also be true.

Therefore, the best translation into predicate logic is 'Fe \supset Ss'. Final Answer: D

Answer Option 6:

 $The best translation into predicate logic for the given statement "If Eileen plays fiddle then Sherri sings" is option D: 'Fe \supset Ss'. \\$

In this translation, Fe represents "Eileen plays fiddle" and Ss represents "Sherri sings". The symbol " \supset " represents the logical implication "if...then". Therefore, 'Fe \supset Ss' translates to "If Eileen plays fiddle, then Sherri sings", which accurately captures the original statement.

Options A, B, and C do not accurately represent the logical implication between Eileen playing fiddle and Sherri singing. Final Answer: D

Answer Option 7:

The best translation into predicate logic for the given statement "If Eileen plays fiddle then Sherri sings" is option D: 'Fe \supset Ss'.

Here's a breakdown of the translation:

- Fe represents "Eileen plays fiddle", where e is the variable representing Eileen and F is the predicate for playing fiddle.
- Ss represents "Sherri sings", where s is the variable representing Sherri and S is the predicate for singing.

In the translation Fe \supset Ss, the symbol \supset represents the logical implication "if...then". So, the statement can be read as "If Eileen plays fiddle, then Sherri sings". This translation captures the conditional relationship between Eileen playing the fiddle and Sherri singing accurately. Final Answer: D.

Answer Option 8:

The best translation into predicate logic for the statement "If Eileen plays fiddle then Sherri sings" would be option D: 'Fe \supset Ss'.

Here's a breakdown of the translation:

Fe: Eileen plays fiddle

Ss: Sherri sings

The statement "If Eileen plays fiddle then Sherri sings" can be represented as "Eileen playing fiddle implies Sherri singing". In symbolic logic, implication is represented by the conditional operator ">="."

Therefore, the translation would be "Fe \supset Ss", which reads as "If Eileen plays fiddle, then Sherri sings".

Option A: 'Fe v Ss' would be translated as "Eileen plays fiddle or Sherri sings", which does not capture the conditional relationship between the two actions.

Option B: ${}^{!}eF \supset Ss^{!}$ is not a valid translation.

Option C: 'Fe v Es' would be translated as "Eileen plays fiddle or Eileen sings", which does not accurately represent the original statement.

Therefore, the correct answer is option D: 'Fe \supset Ss'.

Answer Option 9:

The best translation into predicate logic for the given statement "If Eileen plays fiddle then Sherri sings" would be: D. 'Fe \supset Ss'.

Let's break down the statement:

- Eileen plays fiddle is represented by the predicate Fe.
- Sherri sings is represented by the predicate Ss.
- The word "if" implies an implication, which is symbolized by the conditional operator \supset (implies).
- Therefore, "If Eileen plays fiddle then Sherri sings" can be translated as $\text{Fe} \supset \text{Ss}.$

Option A ('Fe v Ss') represents a logical OR between Eileen playing fiddle and Sherri singing, which is not an accurate translation of the original statement.

Option B ('eF \supset Ss') does not properly represent the names Eileen and Sherri as predicates, and the order of the predicates in the implication is incorrect.

 $Option \ C \ ('Fe \lor Es') \ introduces \ an \ additional \ predicate \ Es, \ which \ is \ not \ mentioned \ in \ the \ original \ statement.$

Therefore, option D ('Fe \supset Ss') is the most accurate translation into predicate logic. Final Answer: D

Question 68:
'Construct a complete truth table for the following argument. Then, using the truth table, determine whether the argument is valid or invalid. If the argument is invalid, choose an option which presents a
counterexample. (There may be other counterexamples as well.)
-C⊃D
D \supset C / C', A. 'Valid', B. 'Invalid. Counterexample when C and D are true', C. 'Invalid. Counterexample when C is true and D is false', D. 'Invalid. Counterexample when D is true and C is false'
Answer Option 1:
To construct a complete truth table for the argument, we need to consider all possible combinations of truth values for the propositions involved.
Let's consider the propositions:
C: C is true
D: D is true
-Coddocc'
FOTTF
Tofff
Based on the truth table, we can see that there is no row where both premises are true and the conclusion is false. Therefore, the argument is valid.
Final Answer: A. "Valid"

Answer Option 2:
To construct a complete truth table for the argument, we need to consider all possible truth values for the propositions involved. Let's denote ~C as P, D as Q, and C' as R.
$P \supset QQ \supset RR$
TTF
TFT
FTF
FTT
Based on the truth table, we can evaluate the argument. If the premises $(P \supset Q \text{ and } Q \supset R)$ are true, we check whether the conclusion (R) is also true for each row of the truth table.
In this case, we can see that when P is true and Q is true, R is true as well. Therefore, the argument is valid, and the correct answer is: A. "Valid"

Answer Option 3:
To construct a truth table for the argument, we need to consider all possible combinations of truth values for the propositions involved: C and D.
$CD\!-\!C\simC\supsetDD\supsetCC'$
TIFITI
TFFTTT
FTTTFF
FFTFTF
From the truth table, we can see that when C is true and D is true, the premises "~C ⊃ D" and "D ⊃ C" are both true. However, the conclusion "C" is false in this case. Therefore, the argument is invalid.
The correct answer is: B. Invalid. Counterexample when C and D are true.

Answer Option 4:
To construct a complete truth table for the argument, we will consider all possible combinations of truth values for propositions C and D.
We have the premises:
1C ⊃ D
2. D ⊃ C
And the conclusion:
C'

From the truth table, we can see that there exist cases where both premises are true (rows 1 and 4) and the conclusion is false. Therefore, the argument is invalid.

We can find a counterexample by choosing the option that presents a case where both premises are true, but the conclusion is false.

CD-CDDDDCC'
TTTTT
TFFTF
FTTFT
FFTTF

Option C states that the counterexample occurs when C is true and D is false. In our truth table, this corresponds to row 2, where C is true and D is false, and the conclusion C' is false. This counterexample shows that the argument is invalid. Therefore, the correct answer is C. 'Invalid. Counterexample when C is true and D is false'. Answer Option 5: To construct a complete truth table for the argument, we need to consider the truth values for all the relevant propositions involved. In this case, we have propositions C and D. Premise 1: ~C ⊃ D Premise 2: $D \supset C$ Conclusion: C' Let's first construct the truth table: $C \mid D \mid \neg C \mid \neg C \supset D \mid D \supset C \mid C'$ TITIFITIT T|F|F|T|T|T F|T|T|T|F|F FIFITIFITIT Based on the truth table, we can see that there is a row where both premises are true (C is false and D is true), but the conclusion (C') is false. This means that the argument is invalid. Therefore, the correct answer is: C. 'Invalid. Counterexample when C is true and D is false' Answer Option 6: To construct a complete truth table, we need to consider all possible combinations of truth values for the propositions involved. In this case, we have two propositions: C and D. Let's first construct the truth table for ${\sim}C \supset D$: $C \sim C D \sim C \supset D$ TFTT TFFF FTTT FTFT Now, let's construct the truth table for $D \supset C$: $C\;D\;D\supset C$ TTT TET FTF FFT Lastly, let's combine the truth tables for $\sim C \supset D$ and $D \supset C$ to check the validity of the argument: $C \ D \ \text{$^\sim$} C \supset D \ D \supset C \ C'$ TTTTF TFFTF FTTFT FFTTT Based on the truth table, we can see that there are cases where both premises are true (rows 1 and 4), but the conclusion C' is false. Therefore, the argument is invalid. The correct answer is option B: Invalid. Counterexample when C and D are true. ********** To construct a complete truth table for the argument, we need to consider all possible combinations of truth values for the propositions involved. In this case, we have two propositions: C and D. The argument is as follows: 1. ~C ⊃ D 2. D ⊃ C

Conclusion: C'

Let's construct the truth table:

$\mathbb{C} \mathbb{D} - \mathbb{C} - \mathbb{C} \supset \mathbb{D} \mathbb{D} \supset \mathbb{C} \mathbb{C}'$
TIFIFITIF
F T T T F T
F F T F T T
Based on the truth table, we can see that there is at least one row where both premises are true (row 4: D is true, and D \supset C is true), but the conclusion (C') is false.
Therefore, the argument is Invalid.
Option C is the correct choice as it states that the argument is Invalid and provides a counterexample when C is true and D is false.
Answer Option 8:
To construct a complete truth table for the argument, we need to consider all possible combinations of truth values for the propositions involved.
Let's consider the propositions:
C: C is true
D: D is true
-C: C is false (negation of C)
$-C \supset D$: If C is false, then D is true (implication)
$D \supset C$: If D is true, then C is true (implication)
C': C' is true (conclusion)
The truth table for the argument is as follows:
$ C D -C -C\supset D D\supset C C' $
T T F T T F
T F F T T F
F T T T F T
F F T F T T
Based on the truth table, we can see that there is a row where both premises (statements before the */") are true, but the conclusion (C') is false. Therefore, the argument is invalid.
Counterexample: When C is true and D is true, both premises are true, but the conclusion is false.
Therefore, the correct answer is B. "Invalid. Counterexample when C and D are true."
Answer Option 9:
The given argument is as follows:
Premise 1: ~C ⊃ D
Premise 2: $D \supset C$
Conclusion: C'
To construct a complete truth table, we need to consider all possible combinations of truth values for the variables C and D, as well as the intermediate propositions \sim C, \sim C \supset D, D \supset C, and the conclusion C and D, as well as the intermediate propositions \sim C.
Here is the truth table:
$ C D \sim C \sim D D\supset C C' $
T T F T T F
T F F T T F
F T T T F T
[F]F]T]F]T]T]
From the truth table, we can see that there is at least one row where the premises are true (i.e., $-C \supset D$ and $D \supset C$) but the conclusion is false (C'). Specifically, when C is true and D is true, the premises are

true, but the conclusion is false. This is a counterexample to the argument.

Therefore, the argument is invalid. The correct answer is B. 'Invalid. Counterexample when C and D are true'.

Question 69:

Answer Option 1:

The best translation into predicate logic for the statement "Some blue herons migrate" is B. $(3x)[(Hx \cdot Bx) \supset Mx]$ ". This translates to "There exists an x such that if x is a blue heron, then x migrates."

Answer Option 2:

The best translation into predicate logic for the statement "Some blue herons migrate" is:

 $(\exists x)[(Hx \bullet Bx) \bullet Mx]$

This statement asserts the existence (\exists) of an individual x such that x is a blue heron (Hx), x is a bird (Bx), and x migrates (Mx).

Therefore, the correct answer is D. $'(\exists x)[(Hx \bullet Bx) \bullet Mx]'$.

Answer Option 3:

The best translation into predicate logic for the statement "Some blue herons migrate" is:

 $(\exists x)[(Hx \bullet Bx) \bullet Mx]$

This translates to: "There exists an x, such that x is a blue heron and x migrates."

Option D, '(∃x)[(Hx • Bx) • Mx]', correctly captures the existence of at least one blue heron that migrates and is the most accurate translation. Final Answer: D

Answer Option 4:

The best translation into predicate logic for the statement "Some blue herons migrate" is:

 $(\exists x)[(Hx \bullet Bx) \bullet Mx]$

This translates to: There exists an x such that x is a blue heron and x migrates.

Option D is incorrect because it uses the conjunction operator (\land) instead of the implication operator (\supset). The statement does not imply that all blue herons migrate, but rather that there exists at least one blue

heron that migrates. Therefore, the correct answer is A.

Final Answer: A

Answer Option 5:

The best translation into predicate logic for the statement "Some blue herons migrate" is:

 $(\exists x)[(Hx \bullet Bx) \land Mx]$

The key here is to use the existential quantifier ($\exists x$) to indicate that there exists at least one blue heron such that it migrates. The conjunction (\land) is used to connect the conditions of being a blue heron (Hx) and being a migratory bird (Mx).

Therefore, the correct answer is D. $'(\exists x)[(Hx \bullet Bx) \bullet Mx]'$.

Answer Option 6:

The best translation into predicate logic for the statement "Some blue herons migrate" is:

 $(\exists x)[(Hx \bullet Bx) \bullet Mx]$

This translation accurately represents the existence of at least one (some) blue heron (Hx) that migrates (Mx). The conjunction (*) between (Hx • Bx) ensures that the heron is both blue (Bx) and a heron (Hx), while the overall existential quantifier (∃x) signifies the existence of such a heron. Therefore, option D is the correct translation.

Answer Option 7:

The best translation into predicate logic for the statement "Some blue herons migrate" is:

 $(\exists x)[(Hx \bullet Bx) \bullet Mx]$

Explanation:

- ($\exists x$) translates to "There exists an x"
- (Hx Bx) translates to "x is a blue heron"
- Mx translates to "x migrates"

The conjunction (\bullet) is used to combine the conditions that x is a blue heron and x migrates.

Therefore, the correct answer is D. $'(\exists x)[(Hx \bullet Bx) \bullet Mx]'$

Answer Option 8:

The best translation into predicate logic for the statement "Some blue herons migrate" is:

 $(\exists x)[(Hx \bullet Bx) \bullet Mx]$

Option D: '(3x)[(Hx • Bx) • Mx]' is the correct translation. This translates to "There exists an x such that x is a blue heron and x migrates." This captures the meaning that there is at least one blue heron that migrates.

Option A: "(\fo'x)[(Hx \cdot Bx) \cdot Mx]" translates to "For all x, x is a blue heron and x migrates." This implies that all blue herons migrate, which is not what the original statement conveys.

Option B: '(3x)[(Hx • Bx) \supset Mx]' translates to "There exists an x such that if x is a blue heron, then x migrates." This does not capture the meaning that there is at least one blue heron that migrates, as it allows for the possibility that there are no blue herons at all.

Option C: '(∀x)[(Hx • Bx) ⊃ Mx]' translates to "For all x, if x is a blue heron, then x migrates." This implies that all blue herons migrate, which is not what the original statement conveys.

Final Answer: D

Answer Option 9:

The statement "Some blue herons migrate" can be translated into predicate logic as $'(\exists x)[(Hx \bullet Bx) \bullet Mx]'$.

Explanation:

- $\hbox{-} (\exists x) \text{ is the existential quantifier, which implies that there exists at least one } x \text{ for which the following condition holds true.} \\$
- Hx represents "x is a heron."
- Bx represents "x is blue."
- Mx represents "x migrates."

Therefore, the correct translation is option D: $'(\exists x)[(Hx \bullet Bx) \bullet Mx]'$.

Question 70:

' Select the best English interpretation of the given arguments in predicate logic.

 $(\exists x)(Cx \cdot Ox)$

 $(\forall x)[(-Cx \supset -Bx) \supset -Qg]/-Qg', A$. 'Some cookies have oatmeal. If something's not being a cookie entails that it doesn't have chocolate chips, then this cookie doesn't have oatmeal. So this cookie doesn't have oatmeal. If something doesn't have oatmeal, then it is not a cookie and it doesn't have chocolate chips. So this cookie doesn't have oatmeal.', D. 'Some cookies have oatmeal. If something's not having chocolate chips entails that it is not a cookie, then it doesn't have oatmeal. So this cookie doesn't have oatmeal.'

Answer Option 1:

The best English interpretation of the given arguments in predicate logic is:

"Some cookies have oatmeal. If something is not a cookie and does not have chocolate chips, it doesn't have oatmeal. So this cookie doesn't have oatmeal."

This corresponds to option B.

.....

Answer Option 2:

The best English interpretation of the given arguments in predicate logic is:

"Some cookies have oatmeal. If something is not a cookie and does not have chocolate chips, it doesn't have oatmeal. So this cookie doesn't have oatmeal."

This interpretation accurately reflects the given premises and conclusion in a clear and concise manner. Final Answer: B

Answer Option 3:

To determine the best English interpretation of the given arguments in predicate logic, let's break down the statements:

(∃x)(Cx • Ox): There exists an object x that is both C and O.

 $(\forall x)[(\neg Cx \supset \neg Bx) \supset \neg Cg]$: For all objects x, if not being C implies not being B, then it is not O for object g.

Based on these statements, the best English interpretation of the argument is option C: "Some cookies have oatmeal. If something doesn't have oatmeal, then it is not a cookie and it doesn't have chocolate chips. So this cookie doesn't have oatmeal."

This interpretation captures the essence of the given arguments and the logical relationship between the statements. It correctly translates the predicates and the implications presented in the argument. Final

Answer: C

Answer Option 4:

To interpret the given arguments in predicate logic, let's break down the statements:

 $(\exists x)(Cx \bullet Ox).$ Some object x is a cookie (Cx) and has oatmeal (Ox).

(∀x)[(-Cx ⊃ -Bx) ⊃ -Og]: For all objects x, if not being a cookie (-Cx) implies that it doesn't have chocolate chips (-Bx), then it is not the case that object g has oatmeal (-Og).

Therefore, the best English interpretation of the given arguments is:

"Some cookies have oatmeal. If something is not a cookie and does not have chocolate chips, it doesn't have oatmeal. So this cookie doesn't have oatmeal."

Option B is the correct interpretation.

Answer Option 5:

The argument can be translated as follows:

Premise 1: There exists an object x such that it is a cookie (Cx) and it has oatmeal (Ox).

Premise 2: For all objects x, if it is not a cookie (-Cx) implying that it doesn't have chocolate chips (-Bx), then it is not the case that Og (negation of Og, which means "this cookie doesn't have oatmeal").

Conclusion: Therefore, it is not the case that Og ("this cookie doesn't have oatmeal").

The best English interpretation of the argument is option C: "Some cookies have oatmeal. If something doesn't have oatmeal, then it is not a cookie and it doesn't have chocolate chips. So this cookie doesn't have oatmeal." This interpretation captures the essence of the argument and accurately represents the logical structure presented in the premises and conclusion.

Answer Option 6:

The best English interpretation of the given arguments would be option C: 'Some cookies have oatmeal. If something doesn't have oatmeal, then it is not a cookie and it doesn't have chocolate chips. So this cookie doesn't have oatmeal.'

Explanation:

The argument begins with (∃x)(Cx • Ox), which translates to "Some cookies have oatmeal." This states that there is at least one object x such that it is a cookie (Cx) and it has oatmeal (Ox).

The second part of the argument is $(\forall x)[(-Cx - Bx) - Qy]$. This translates to "For all x, if not being a cookie implies not having chocolate chips, then it does not have oatmeal." This implies that if something is not a cookie (-Cx) and it does not have chocolate chips (-Bx), then it does not have oatmeal (-Qy).

By combining these two statements, we can conclude that if something does not have oatmeal (-Og), then it is not a cookie and it does not have chocolate chips. Therefore, the best English interpretation is option C.

Answer Option 7:

To determine the best English interpretation of the given arguments in predicate logic, let's analyze the statements:

(3x)(Cx • Ox): This statement can be read as "There exists an object x such that x is a cookie (Cx) and x has oatmeal (Ox)." This implies that there is at least one cookie that has oatmeal.

(\forall colored to the case that \forall objects x, if it is not the case that x is a cookie (-Cx) implies that x does not have chocolate chips (-Bx), then it is not the case that Og has oatmeal (-Og).* This statement suggests that if anything that is not a cookie or does not have chocolate chips, then Og (the cookie in question) does not have oatmeal.

The conclusion ~Og can be interpreted as "Og does not have oatmeal."

Among the given options, the interpretation that best represents the given arguments is:

C. "Some cookies have oatmeal. If something doesn't have oatmeal, then it is not a cookie and it doesn't have chocolate chips. So this cookie doesn't have oatmeal."

This interpretation accurately captures the presence of oatmeal in some cookies, the relationship between not having oatmeal, not being a cookie, and not having chocolate chips, and the conclusion that the specific cookie being discussed does not have oatmeal.

Answer Option 8:

To determine the best English interpretation of the given arguments in predicate logic, let's break down the statements:

 $(\exists x)(Cx \bullet Ox)\text{: There exists an object } x \text{ such that } x \text{ is a cookie } (Cx) \text{ and } x \text{ has oatmeal } (Ox).$

(∀x)[(-Cx ⊃ -Bx) ⊃ -Og]: For all objects x, if not being a cookie (-Cx) implies not having chocolate chips (-Bx), then it does not have oatmeal (-Og).

/ ~Og: Therefore, it does not have oatmeal

Analyzing the given options:

- A. 'Some cookies have oatmeal. If something's not being a cookie entails that it doesn't have chocolate chips, then this cookie doesn't have oatmeal. So this cookie doesn't have oatmeal.
- B. 'Some cookies have oatmeal. If something is not a cookie and does not have chocolate chips, it doesn't have oatmeal. So this cookie doesn't have oatmeal.'
- C. 'Some cookies have oatmeal. If something doesn't have oatmeal, then it is not a cookie and it doesn't have chocolate chips. So this cookie doesn't have oatmeal.'
- D. 'Some cookies have oatmeal. If something's not having chocolate chips entails that it is not a cookie, then it doesn't have oatmeal. So this cookie doesn't have oatmeal.'

The best English interpretation of the given arguments in predicate logic is option C. 'Some cookies have oatmeal. If something doesn't have oatmeal, then it is not a cookie and it doesn't have chocolate chips. So this cookie doesn't have oatmeal.' This interpretation accurately captures the meaning of the logical statements and the conclusion.

Answer Option 9:

To determine the best English interpretation of the given arguments, let's break down the statements.

The first statement $(\exists x)(Cx \bullet Ox)$ can be interpreted as "There exists an x such that x is a cookie (Cx) and x has oatmeal (Ox)."

The second statement $(\forall x)[(-Cx \supset -Bx) \supset -Og]$ can be interpreted as "For all x, if it is not the case that x is a cookie $(\neg Cx)$ implies that x does not have chocolate chips (-Bx), then it is not the case that Og (this cookie has oatmeal)."

Combining the two statements, the argument can be interpreted as "Some cookies have oatmeal. If it is not the case that something being a cookie entails that it has chocolate chips, then this cookie does not have oatmeal."

Comparing this interpretation with the given options:

- A. 'Some cookies have oatmeal. If something's not being a cookie entails that it doesn't have chocolate chips, then this cookie doesn't have oatmeal. So this cookie doesn't have oatmeal.' This option accurately captures the meaning of the argument and is the best English interpretation.
- B. 'Some cookies have oatmeal. If something is not a cookie and does not have chocolate chips, it doesn't have oatmeal. So this cookie doesn't have oatmeal.' This option changes the meaning of the argument by introducing the condition of not having chocolate chips, which is not implied by the argument.
- C. 'Some cookies have oatmeal. If something doesn't have oatmeal, then it is not a cookie and it doesn't have chocolate chips. So this cookie doesn't have oatmeal.' This option flips the condition and changes the meaning of the argument by stating that if something doesn't have oatmeal, it is not a cookie and doesn't have chocolate chips, which is not implied by the argument.
- D. 'Some cookies have oatmeal. If something's not having chocolate chips entails that it is not a cookie, then it doesn't have oatmeal. So this cookie doesn't have oatmeal.' This option introduces the condition of not having chocolate chips, which is not implied by the argument.

Therefore, the best English interpretation of the given arguments is A. 'Some cookies have oatmeal. If something's not being a cookie entails that it doesn't have chocolate chips, then this cookie doesn't have oatmeal. So this cookie doesn't have oatmeal.'

Jestion	

'Identify the conclusion of the following argument. All of psychology has gotten stuck in moral prejudices and fears. No one has come close to understanding it as the development of the will to power.

However, if a person even begins to regard the effects of hatred, envy, covetousness, and the lust to rule as conditions of life and furthermore, as factors essential to the general economy of life, he will begin to get seasick. At this point, he begins to lose himself, and sail over morality. Thus, psychology becomes again the path to fundamental problems. (Friedrich Nietzsche, Beyond Good and Evil)', A. 'All of psychology has gotten stuck in moral prejudices and fears.', B. 'No one has come close to understanding it as the development of the will to power.', C. 'If a person even begins to regard the effects of hatred, envy, covetousness, and the lust to rule as conditions of life and furthermore, as factors essential to the general economy of life, he will begin to get seasick.', D. 'Psychology becomes again the path to fundamental problems.'

Answer Option 1:

The conclusion of the argument is:

D. 'Psychology becomes again the path to fundamental problems.'

Answer Option 2:

The conclusion of the argument is: D. 'Psychology becomes again the path to fundamental problems.'

Answer Option 3:

The conclusion of the argument is: D. "Psychology becomes again the path to fundamental problems."

Answer Option 4:

The conclusion of the argument can be identified as: D. 'Psychology becomes again the path to fundamental problems.'

Answer Option 5:

The conclusion of the argument is: D. 'Psychology becomes again the path to fundamental problems.' This is stated in the last sentence of the argument: "Thus, psychology becomes again the path to

fundamental problems." Final Answer: D

Answer Option 6:

The conclusion of the argument is: D. 'Psychology becomes again the path to fundamental problems.' This can be inferred from the statement, "Thus, psychology becomes again the path to fundamental

problems." Final Answer: D

Answer Option 7:

The conclusion of the argument can be identified as follows: The argument builds up to the statement that "psychology becomes again the path to fundamental problems." This can be seen as the main point

or conclusion of the argument presented by Friedrich Nietzsche. Therefore, the correct answer is D. 'Psychology becomes again the path to fundamental problems.'

Answer Option 8:

The conclusion of the argument is: D. 'Psychology becomes again the path to fundamental problems.'

The statement "Thus, psychology becomes again the path to fundamental problems" directly follows the preceding sentences and summarizes the main point of the argument. It indicates that by recognizing

the effects of certain emotions and desires on life and understanding them as essential factors, psychology can once again address fundamental problems. Therefore, D is the conclusion of the argument.

Answer Option 9:

The conclusion of the argument is: D. 'Psychology becomes again the path to fundamental problems.'

The argument presents various statements and claims about psychology, but the ultimate conclusion is that psychology, when approached from a certain perspective (regarding the effects of certain emotions as essential to life), becomes the path to fundamental problems. This conclusion is supported by the preceding statements about psychology being stuck in moral prejudices, the development of the will to

power, and the discomfort one experiences when starting to question traditional moral values.

' Use indirect truth tables to determine whether each set of propositions is consistent. If the set is consistent, choose an option with a consistent valuation. (There may be other consistent valuations.) ${\sim}N\equiv(O\cdot P)$ ~O ⊃ Q N · ~Q', A. 'Inconsistent', B. 'Consistent. Consistent valuation when N, O, and P are true and Q is false', C. 'Consistent. Consistent valuation when N and O are true and Q are false', D. 'Consistent. Consistent valuation when N and P are true and O and Q are false Answer Option 1: To determine the consistency of the set of propositions, we can use indirect truth tables. First, let's assign truth values to the propositions: N: False O: False P: False Q: False Now, let's evaluate each proposition in the set: \sim N \equiv (O \cdot P) ~False = (False · False) $\mathsf{True} \equiv (\mathsf{False} \cdot \mathsf{False})$ True ≡ False ${\sim}O\supset Q$ ~False ⊃ False True ⊃ False False N·~Q False · ~False False · True False Based on the truth values, we can see that the set of propositions is inconsistent because not all of the propositions can be true simultaneously, Therefore, the correct answer is: A. 'Inconsistent' Answer Option 2: To determine the consistency of the set of propositions, we will construct an indirect truth table. We will assign truth values to the propositions and check if there is a consistent valuation where all the propositions are true. Let's assign truth values to the propositions: N: True O: True P: False Q: False \sim N \equiv (O \cdot P) ${\sim}O\supset Q$ N · ~Q ~N \equiv (O \cdot P) is True, as ~N is False and (O \cdot P) is also False. ~O ⊃ Q is False, as ~O is False and Q is False. $N\cdot {\scriptscriptstyle \sim} Q$ is True, as N is True and ${\scriptscriptstyle \sim} Q$ is True. Since there is a consistent valuation where all the propositions are true, the set of propositions is consistent. Therefore, the answer is C. 'Consistent. Consistent valuation when N and O are true and P and Q are false'. Answer Option 3: Let's create a truth table to determine the consistency of the set of propositions: NOPQ~N \equiv (O \cdot P) ~O \supset Q N \cdot ~Q Consistency TTTFTFT Inconsistent

Question 72:

TTTFFTFF Inconsistent

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TTFFFTFFInconsistent
TFTTFTTF Inconsistent
TFTFFTFF Inconsistent
TFFTFTTF Inconsistent
TFFFFTFF Inconsistent
FTTTFTF Inconsistent
FTTFTFTInconsistent
FTFTTFTF Inconsistent
FTFFTFT Inconsistent
FFTTTTF Inconsistent
FFTFTTFT Inconsistent
FFTTTTF Inconsistent
FFFFTTFT Inconsistent
From the truth table, we can see that there are no consistent valuations where all the propositions are true. Therefore, the set of propositions is inconsistent.
Final Answer: A. 'Inconsistent'
*******
Answer Option 4:
To determine the consistency of the set of propositions using indirect truth tables, we need to create a truth table for the given propositions.
NOPQ~N \equiv (O \cdot P) \sim O \supset Q N \cdot \sim Q
TTTTFTFF
TTTFFTTF
TTFTFFFT
TTFFFFTF
TETTETE
TFTFFTTF
TFFTFFFT
TFFFFFFF
FTTTTTFF
FTTFTTTF
FTFTFFFT
FTFFFFF
FFTTTTFF
FFTFTTTF
FFFTTFFT
FFFFTTF
From the truth table, we can see that there is at least one row where all propositions are true (T values). Therefore, the set of propositions is consistent.
The consistent valuation can be found when N and O are true, and P and Q are false.
Final Answer: C. 'Consistent. Consistent valuation when N and O are true and P and Q are false'
Answer Option 5:
To determine the consistency of the propositions, we will construct a truth table:
N O P Q ~N ~O (O \cdot P) ~N \equiv (O \cdot P) ~O \supset Q N \cdot ~Q Consistency
TTTFFFFF Inconsistent
TTTFFFFFFInconsistent
TTFTFFFFF Inconsistent
TTFFFFFFF Inconsistent
TFTTFTTFF Inconsistent
TFTFFFFFF Inconsistent
TFFTFFFFFInconsistent
TFFFFTTFF Inconsistent
FTTTFFF Inconsistent
FTTFTFTF Inconsistent
```

FTFTTFFTFInconsistent

FTFFTFTFInconsistent FFTTTTFF Inconsistent FFTFTTTFF Inconsistent FFFTTTFFF Inconsistent FFFTTFFFInconsistent Based on the above truth table, we can see that there is no consistent valuation for the given set of propositions. Therefore, the correct answer is A. 'Inconsistent'. Answer Option 6: To determine whether the set of propositions is consistent, we construct an indirect truth table. We list all the propositions and their possible truth values, and then evaluate the truth values of the compound propositions based on the given propositions. $N O P Q \sim N \equiv (O \cdot P) \sim O \supset Q N \cdot \sim Q$ TTTTETEE TTTFFTFF TTFTFTFF TTFFFTTT TFTTFTTF TETEETTE TFFTFTTT TEFETTT FTTTTTF FTTFTTTF FTFTTFTF FTFFTFTT FFTTTTF FFTFTTTF FFFTTTT FFFFTTTT Based on the truth table, we can see that there is at least one row where all the propositions are true, which means the set of propositions is consistent. In the row where N is true, O is true, P is true, and Q is false, all the given propositions are satisfied. Therefore, the consistent valuation is when N, O and P are true, and Q is false. Final Answer: B. 'Consistent. Consistent valuation when N, O, and P are true and Q is false' ****************** Answer Option 7: To determine whether the set of propositions is consistent, we can create an indirect truth table. Let's analyze the given propositions: \sim N \equiv (O \cdot P) ~O ⊃ Q N·~Q We will create a truth table and evaluate the consistency of the propositions: N O P Q ~N O \cdot P ~N \equiv (O \cdot P) ~O \supset Q N \cdot ~Q Consistency TTTTFTFTT Inconsistent TTTFFTFTFInconsistent TTFTFFFTT Inconsistent TTFFFFFTFConsistent TFTTFTFTT Inconsistent TFTFTFFTF Inconsistent TFFTFFTT Inconsistent TFFFTFFTF Consistent FTTTTFT F Consistent FTTFTTFTT Inconsistent FTFTTFFTF Inconsistent

FTTFTTFTTTConsistent
FFTTTFTTTTConsistent
FFTTTFTTFTTConsistent

```
FFFFFFTT Inconsistent
From the truth table, we can see that there are no consistent valuations where all the propositions are true. Therefore, the set of propositions is inconsistent.
Final Answer: A. 'Inconsistent'
Answer Option 8:
To determine the consistency of the set of propositions, we can construct an indirect truth table by assigning truth values to the propositions and evaluating their logical relationships.
We have the following propositions:
1. \simN \equiv (O \cdot P)
2. ~O ⊃ Q
3. N · ~Q
Constructing the indirect truth table:
N\mid O\mid P\mid Q\mid \sim N\equiv (O\cdot P)\mid \sim O\supset Q\mid N\cdot \sim Q
T|T|T|T|F|T|F
T|T|T|F|F|T|T
TITIFITIFITIF
TITIFIFIFITIF
T|F|T|T|F|T|F
T|F|T|F|F|T|T
T|F|F|T|T|F
TIFIFIFITITIT
FITITITIFITIF
F|T|T|F|F|T|F
F|T|F|T|T|F
F|T|F|F|T|T|T
F|F|T|T|T|F
F|F|T|F|T|T
F|F|F|T|T|F
F|F|F|F|T|T|T
From the truth table, we can see that there are consistent valuations, where the propositions are all true or all false. Specifically, when N and O are true, and P and Q are false, the set of propositions is
Therefore, the correct answer is C. 'Consistent. Consistent valuation when N and O are true and P and Q are false'.
Answer Option 9:
To determine the consistency of the set of propositions, we can create an indirect truth table.
We have the following propositions:
1. ~N ≡ (O · P)
2.~ {\sim} O \supset Q
3. N · ~O
Let's assign truth values to each proposition and evaluate the consistency.
First, we create columns for each proposition and their components:
```

 $N \ | \ O \ | \ P \ | \ Q \ | \ \sim N \ | \ O \cdot P \ | \ \sim O \supset Q \ | \ N \cdot \sim Q \ |$

 From the indirect truth table above, we can see that there is at least one row where all the propositions are true. Therefore, the set of propositions is consistent.

Option B: "Consistent. Consistent valuation when N, O, and P are true and Q is false" is a valid consistent valuation, but there can be other consistent valuations as well.

Final Answer: C. "Consistent. Consistent valuation when N and O are true and P and Q are false"

Question 73:

'Which of the given formulas of PL is the best symbolization of the following sentence?

A person is considered to have lived with you for all of the present year if the person was born or died in the present year and your home was this person's home for the entire time this person was alive.', A.

 $"[(\mathsf{B} \vee \mathsf{D}) \bullet \mathsf{H}] \supset \mathsf{L}', \, \mathsf{B}. \, "[\mathsf{L} \supset (\mathsf{B} \vee \mathsf{D})] \bullet \mathsf{H}', \, \mathsf{C}. \, "\mathsf{L} \supset [(\mathsf{B} \vee \mathsf{D}) \bullet \mathsf{H}]', \, \mathsf{D}. \, "(\mathsf{B} \vee \mathsf{D} \bullet \mathsf{H}) \supset \mathsf{L}'$

Answer Option 1:

The best symbolization of the given sentence would be: A. '[(B \vee D) • H] \supset L'

Let's break down the sentence into its logical components:

- A person is considered to have lived with you for all of the present year: L
- if the person was born or died in the present year: (B \vee D)
- and your home was this person's home for the entire time this person was alive: H

Putting it all together, the sentence can be symbolized as: [(B \vee D) \bullet H] \supset L

This formula accurately captures the relationship between being born or died in the present year, your home being the person's home, and being considered to have lived with you for the entire present year.

Final Answer: A. '[(B \vee D) • H] \supset L'

Answer Option 2:

Let's break down the sentence to simplify the symbolization:

1. A person is considered to have lived with you for all of the present year.

Symbolized as: L

2. if the person was born or died in the present year.

Symbolized as: B \vee D

3. and your home was this person's home for the entire time this person was alive.

Symbolized as: H

Therefore, the best symbolization of the sentence is: L \supset [(B \vee D) \bullet H]

Final Answer: C

Answer Option 3:

Let's break down the sentence and analyze the logical structure:

A person is considered to have lived with you for all of the present year if:

- The person was born or died in the present year (B $\scriptstyle \vee$ D)
- Your home was this person's home for the entire time this person was alive (H)

The person being considered to have lived with you for all of the present year is symbolized by L.

Based on the logical structure and symbols used, the best symbolization of the sentence is:

 $C. \ 'L \supset [(B \lor D) \bullet H]'$

This formula correctly represents the conditional relationship, where L implies that both $(B \lor D)$ and H must be true. Final Answer: C

Answer Option 4:

Let's break down the sentence and translate it into logical notation:

- P: A person is considered to have lived with you for all of the present year.
- B: The person was born in the present year.
- D: The person died in the present year.
- H: Your home was this person's home for the entire time this person was alive.
- L: The person lived with you for all of the present year.

Based on the given sentence, the logical symbolization can be written as:

 $P \Leftrightarrow [(B \vee D) \bullet H] \supset L$

Option A: '[(B \vee D) • H] \supset L'

This option correctly symbolizes the given sentence. The person being born or died in the present year while your home being their home for their entire life implies that they have lived with you for all of the present year.

Option B: $[L \supset (B \lor D)] \cdot H'$

This option does not accurately capture the meaning of the sentence. It suggests that if the person lived with you, then they were either born or died in the present year, which is not what the sentence is stating.

Option C: $L \supset [(B \lor D) \cdot H]'$

This option incorrectly implies that if the person lived with you, then they were either born or died in the present year, which is not stated in the sentence.

Option D: $'(B \lor D \bullet H) \supset L'$

This option does not correctly represent the sentence. It suggests that if the person was born or died and your home was their home, then they lived with you, which does not capture the entirety of the conditions stated in the sentence.

Final Answer: A

Answer Option 5:

Let's break down the given sentence:

A person is considered to have lived with you for all of the present year if:

- 1. The person was born or died in the present year: (B \vee D)
- 2. Your home was this person's home for the entire time this person was alive: $\mbox{\bf H}$

To symbolize this sentence, we can use the following formula:

 $[(\mathsf{B} \vee \mathsf{D}) \bullet \mathsf{H}] \supset \mathsf{L}$

Option A: '[(B \vee D) \bullet H] \supset L' This is the correct symbolization of the given sentence.

 $Option \ B: \ [L\supset (B\lor D)] \bullet H' \ This formula \ does \ not \ capture \ the \ condition \ that \ your \ home \ was \ the \ person's \ home \ for \ the \ entire \ time.$

 $\label{eq:condition} \mbox{Option C: 'L} \supset \mbox{[(B \lor D)} \bullet \mbox{H]'} \mbox{ This formula does not capture the condition that your home was the person's home for the entire time. }$

Option D: '(B v D • H) \supset L' This formula does not properly capture the intended meaning. It suggests that if the person was born or died in the present year and your home was their home, then they lived with you for the entire year.

Final Answer: A

Answer Option 6:

Let's break down the sentence to symbolize it correctly.

- P: A person is considered to have lived with you for all of the present year.
- B: The person was born in the present year.
- D: The person died in the present year.
- H: Your home was this person's home for the entire time this person was alive.
- L: The person lived with you for all of the present year.

The correct symbolization of the sentence would be:

 $[(\mathsf{B} \vee \mathsf{D}) \bullet \mathsf{H}] \supset \mathsf{L}$

Option A: [(B \vee D) • H] \supset L - This is the correct symbolization based on the given sentence.

Option B: [L \supset (B \lor D)] • H - This does not accurately represent the condition that the person's home was your home for the entire time they were alive.

Option C: $L \supset [(B \lor D) \bullet H]$ - This does not accurately represent the condition that your home was their home for the entire time they were alive.

Option D: $(B \lor D \bullet H) \supset L$ - This does not accurately represent the condition that your home was their home for the entire time they were alive.

Final Answer: A

Answer Option 7:

Let's break down the given sentence into logical symbols:

- P: A person is considered to have lived with you for all of the present year.
- B: The person was born in the present year.
- D: The person died in the present year.
- H: Your home was this person's home for the entire time this person was alive.
- L: The person lived with you for all of the present year.

The sentence can be symbolized as:

 $[(\mathsf{B} \vee \mathsf{D}) \bullet \mathsf{H}] \supset \mathsf{L}$

Option A: $'[(B \lor D) \bullet H] \supset L'$ - This is the correct symbolization of the given sentence. It captures the condition that both $(B \lor D)$ and H must be true for P(L) to be true.

Option B: $[L \supset (B \lor D)] \bullet H' - This$ does not correctly capture the condition that the person's home must be your home for the entire time they were alive.

 $\text{Option C: $L \supset [(B \lor D) \bullet H]'$ - This incorrectly implies that if L is true, then both (B \lor D)$ and H must be true, which is not necessarily the case. } \\$

 $Option \ D: '(B \lor D \bullet H) \supset L' \ - This \ does \ not \ correctly \ capture \ the \ condition \ that \ the \ person's \ home \ must \ be \ your \ home \ for \ the \ entire \ time \ they \ were \ alive.$

Final Answer: A. '[(B \vee D) • H] \supset L'

Answer Option 8:

Let's break down the sentence into logical components:

- A person is considered to have lived with you for all of the present year: Let's symbolize this as L.
- If the person was born or died in the present year: Let's symbolize being born as B and dying as D.

- Your home was this person's home for the entire time this person was alive: Let's symbolize this as H.

Now let's look at the given options:

Option A: '[(B \vee D) • H] \supset L'

This option states that if a person was born or died and your home was their home, then they lived with you for all of the present year. However, it does not account for the case where the person was born and died in a different year but still lived with you for all of the present year. So this option is not the best symbolization.

Option B: $[L \supset (B \lor D)] \cdot H'$

This option states that if a person lived with you for all of the present year, then they were born or died. However, it does not account for the condition that your home was their home for the entire time. So this option is not the best symbolization.

Option C: $L \supset [(B \lor D) \bullet H]'$

This option correctly represents the given sentence. It states that if a person lived with you for all of the present year, then they were born or died and your home was their home for the entire time. This is the best symbolization of the sentence.

Option D: $'(B \lor D \bullet H) \supset L'$

This option states that if a person was born or died and your home was their home, then they lived with you for all of the present year. However, it does not include the condition that the person must have been born or died in the present year. So this option is not the best symbolization.

Final Answer: C

Answer Option 9:

Let's break down the sentence and symbolize each component:

A person is considered to have lived with you for all of the present year: L

if the person was born or died in the present year: (B \vee D)

and your home was this person's home for the entire time this person was alive: H

Now, let's look at the answer choices:

A. "[(B v D) • H] \supset L". This translates to "if someone was born or died in the present year and your home was this person's home for the entire time this person was alive, then they are considered to have lived with you for all of the present year." This matches the given sentence, so it is a valid symbolization.

B. "[L > (B > D)] • H': This translates to "If someone is considered to have lived with you for all of the present year, then they were born or died in the present year." This does not capture the entire meaning of the given sentence, so it is not the best symbolization.

C. 'L \supset [(B \vee D) • H]": This translates to "If someone is considered to have lived with you for all of the present year, then they were born or died in the present year and your home was this person's home for the entire time this person was alive." This matches the given sentence, so it is a valid symbolization.

D. '(B \times D \cdot H) \(\times L': \text{ This translates to "If someone was born or died in the present year and your home was this person's home for the entire time this person was alive, then they are considered to have lived with you for all of the present year.* This matches the given sentence, so it is a valid symbolization.

Both options A, C, and D symbolize the given sentence correctly. So, there is no single "best" symbolization among them. Final Answer: A, C, D.

Question 74:

'Construct a complete truth table for the following argument. Then, using the truth table, determine whether the argument is valid or invalid. If the argument is invalid, choose an option which presents a counterexample. (There may be other counterexamples as well.)

(O = P) v P / P v ~O', A. "Valid", B. 'Invalid. Counterexample when O and P are true", C. 'Invalid. Counterexample when O is true and P is false', D. 'Invalid. Counterexample when P is true and O is false'

Answer Option 1:

To construct a complete truth table for the argument, we need to consider all possible combinations of truth values for the propositions involved.

 $PO(O \equiv P) \vee PP \vee \sim O$

TTTT

TFFT

FTTF

FFTT

Based on the truth table, we can see that the argument is invalid.

Counterexample: When O is true and P is false, the premise (O \equiv P) \vee P is true, but the conclusion P \vee \sim O is false.

Therefore, the correct answer is C. 'Invalid. Counterexample when O is true and P is false'.

Answer Option 2:

To construct a complete truth table for the argument, we need to consider all possible combinations of truth values for the propositions involved.

Let's consider the truth values for propositions O and P.

 $OP(O \equiv P) \vee PP \vee \sim O$

TTTT

TFFT

FTFF

Looking at the truth table, we can see that whenever $(O \equiv P) \lor P$ is true, $P \lor \neg O$ is also true. Therefore, for all possible combinations of truth values, the final proposition is always true.

Since the final proposition is always true, the argument is valid.

Therefore, the correct answer is: A. 'Valid'

Answer Option 3:

To construct a truth table, we need to consider all possible combinations of truth values for the propositions involved. In this case, there are two propositions, P and O.

P | O | (O = P) v P | P v ~O

TITITIT

TIFITIT

F|T|F|F

F|F|T|T

From the truth table, we can see that in all cases where (O \equiv P) \vee P is true, P \vee ~O is also true. Therefore, the argument is valid.

Final Answer: A. 'Valid'

Answer Option 4:

To construct a truth table for the given argument, we need to consider all possible combinations of truth values for the propositions P and O.

 $P \mid O \mid (O \equiv P) \lor P \mid P \lor \sim O$

 $T \mid T \mid T \mid T$

T|F|F|T

F|T|F|F

F|F|T|T

Based on the truth table, we can see that there is no row where the premise " $(O \equiv P) \lor P$ " is true and the conclusion " $P \lor \sim O$ " is false. Therefore, the argument is valid.

Final Answer: A. Valid

Answer Option 5:

To construct a complete truth table for the argument, we need to consider all possible truth values for the propositions involved. In this case, we have two propositions, P and O.

PO (O \equiv P) \vee PP \vee \sim O

TTTT TFFT FTFF

FFTT

Now, let's analyze the truth values for the argument. The argument states that " $(O \equiv P) \lor P$ " implies " $P \lor \neg O$ ". Looking at the truth table, we can see that the columns for " $(O \equiv P) \lor P$ " and " $P \lor \neg O$ " have the same truth values for all rows. Therefore, we can conclude that the argument is valid.

Therefore, the correct answer is: A. 'Valid'

Answer Option 6:

To construct a truth table for the argument, we will consider all possible combinations of truth values for the propositions O and P.

| O | P | (O = P) v P | P v ~ O | |---|---------|-------| | T | T | T | T | | T | F | F | F | | F | T | T | T | | F | F | F | T | T |

Based on the truth table, we can see that the argument is invalid. In the row where O is true and P is false, the premise $((O \equiv P) \lor P)$ is true, but the conclusion $(P \lor \neg O)$ is false.

Therefore, the correct answer is C. "Invalid. Counterexample when O is true and P is false."

.....

Answer Option 7:

To construct a complete truth table for the argument, we need to consider all possible combinations of truth values for the propositions involved.

Let's consider P and O as the propositions in the argument.

T|F|F|T F|T|F|F

F|F|T|T

From the truth table, we can see that there is one row where the premise "($O \equiv P$) $\vee P$ " is true (last row) and the conclusion " $P \vee \sim O$ " is false (third row).

Therefore, the argument is invalid.

Counterexample: When O is true and P is false, the premise "(O \equiv P) \vee P" is false, but the conclusion "P \vee \sim O" is true.

The correct option is C. 'Invalid. Counterexample when O is true and P is false'.

Answer Option 8:

To construct a complete truth table, we need to consider all possible combinations of truth values for the propositions involved in the argument.

Let's consider the propositions involved:

P: P is true

O: O is true

We construct the truth table for the argument:

| P | O | (O = P) v P | P v - O | |---|-------|-----| | T | T | T | T | | T | F | F | F | | F | F | F | T | T |

From the truth table, we can see that the argument is invalid. The final column does not always have the same truth value as the initial proposition. Therefore, the argument is invalid.

To find a counterexample, we look for a row in which the initial proposition is true, but the final proposition is false. In this case, the counterexample is when O is true and P is false.

Therefore, the correct answer is C. 'Invalid. Counterexample when O is true and P is false'.

Answer Option 9:

To construct the truth table for the argument, we need to consider all possible truth values for the propositions involved.

Let's create the truth table:

 $\mid O \mid P \mid (O \equiv P) \vee P \mid P \vee {\sim} O \mid$

T T T T
T F F F
F T F T
F F T T

From the truth table, we can see that there are cases where the premise "(O \equiv P) \vee P" is true (rows 1 and 4), but the conclusion "P \vee \sim 0" is false in those cases (row 1).

Therefore, the argument is invalid.

To provide a counterexample, we can use the case where O is true and P is true (row 1). In this case, the premise is true, but the conclusion is false, showing that the argument does not hold true in all cases.

Final Answer: C. Invalid. Counterexample when O is true and P is false.

Question 75: 'Select the best translation into predicate logic: Leroy is bad only if DJ Khaled shouts his name.', A. 'Sk \equiv Bl', B. 'Bl \equiv Sk', C. 'Bx \supset Sy', D. 'Bl \supset Sk' Answer Option 1: The given statement, "Leroy is bad only if DJ Khaled shouts his name" can be translated into predicate logic as follows: Let's assume: Sk: Leroy is bad BI: DJ Khaled shouts his name The translation becomes: $\mathsf{Sk} \to \mathsf{Bl}$ Therefore, the correct translation into predicate logic is C. $^{\prime}Bx \supset Sy^{\prime}$ ******* Answer Option 2: The given statement, "Leroy is bad only if DJ Khaled shouts his name" can be translated as: B: Leroy is bad S: DJ Khaled shouts his name $\label{eq:correct} \mbox{The correct translation into predicate logic is: } \mbox{$B \supset S$, which means "If Leroy is bad, then DJ Khaled shouts his name"}.$ Therefore, the best translation into predicate logic is: C. 'Bx \supset Sy'. Answer Option 3: The best translation into predicate logic for the given statement would be option C: 'Bx \supset Sy'. Let's break down the given statement: - Leroy is bad: Bx (where x represents Leroy) - Only if: ⊃ (implies) - DJ Khaled shouts his name: Sy (where y represents DJ Khaled shouting his name) Putting it together, we have "Leroy is bad only if DJ Khaled shouts his name" translated as "Bx \supset Sy". Therefore, option C is the correct translation. Answer Option 4: The given statement can be translated into predicate logic as follows: Leroy is bad only if DJ Khaled shouts his name. Let: Sk: Leroy is bad. BI: DJ Khaled shouts his name. The correct translation into predicate logic is: C. $'Bx \supset Sy'$ Bx represents "x is bad" where x is Leroy. Sy represents "y shouts his name" where y is DJ Khaled. The implication symbol (>) represents "only if". Therefore, 'Bx \supset Sy' correctly translates the given statement into predicate logic. Final Answer: C Answer Option 5: The given statement can be translated into predicate logic as follows: B(x) represent "x is bad." S(x) represent "x shouts DJ Khaled's name." The statement "Leroy is bad only if DJ Khaled shouts his name" can be translated as "If Leroy is bad, then DJ Khaled shouts his name." This can be represented as: $B(L) \to S(DJ \; Khaled \hbox{'s name})$ Now, looking at the answer options: A. $\mbox{'Sk} \equiv \mbox{Bl'}$ - This does not correctly represent the given statement.

B. $'BI \equiv Sk'$ - This does not correctly represent the given statement.

 $\hbox{C. 'Bx} \supset \hbox{Sy'} - \hbox{This represents a general statement and does not specifically refer to Leroy and DJ Khaled}.$

D. $\mbox{'BI} \supset \mbox{Sk'}$ - This correctly represents the given statement.

Therefore, the best translation into predicate logic is:

D. 'BI ⊃ Sk'

......

Answer Option 6:

The best translation into predicate logic for the given statement, "Leroy is bad only if DJ Khaled shouts his name," is option D: "Bl \supset Sk."

In predicate logic, we can represent the statement as "If Leroy is bad, then DJ Khaled shouts his name." This can be translated as "BI \supset Sk" where BI represents "Leroy is bad" and Sk represents "DJ Khaled shouts his name."

Option A, "Sk = BI," represents "DJ Khaled shouts his name if and only if Leroy is bad," which is not the correct translation as it implies that DJ Khaled shouting his name is the only condition for Leroy being

Option B, "Bl = Sk," represents "Leroy is bad if and only if DJ Khaled shouts his name," which is not the correct translation as it implies that DJ Khaled shouting his name is the only condition for Leroy being . . .

 $Option \ C, \ "Bx \supset Sy," \ does \ not \ capture \ the \ relationship \ between \ Leroy \ being \ bad \ and \ DJ \ Khaled \ shouting \ his \ name.$

Answer Option 7:

The statement "Leroy is bad only if DJ Khaled shouts his name" can be translated into predicate logic as "Bx \rightarrow Sy". Here, Bx represents the statement "Leroy is bad" and Sy represents the statement "DJ Khaled shouts his name". The arrow (\rightarrow) represents the implication "only if".

Option A: 'Sk = Bl' - This statement translates to "DJ Khaled shouting his name if and only if Leroy is bad", which is different from the original statement.

Option B: 'BI = Sk' - This statement translates to "Leroy is bad if and only if DJ Khaled shouts his name", which is also different from the original statement.

Option C: 'Bx \supset Sy' - This statement accurately represents the original statement, where Leroy being bad implies that DJ Khaled shouts his name.

Option D: 'Bl \supset Sk' - This statement translates to "Leroy is bad implies that DJ Khaled shouts his name", which is different from the original statement.

Therefore, the best translation into predicate logic is C. 'Bx \supset Sy'.

Answer Option 8:

To translate the given statement into predicate logic, we can assign the following predicates:

Sk: Leroy is bad

Bl: DJ Khaled shouts his name

The given statement, "Leroy is bad only if DJ Khaled shouts his name", can be expressed as "Sk implies BI" or "Sk \rightarrow BI".

Looking at the answer choices:

 $A. \ \ \mathsf{SK} \equiv \mathsf{BI'} - \mathsf{This} \ \mathsf{translates} \ \mathsf{to} \ \ \mathsf{"Leroy} \ \mathsf{is} \ \mathsf{bad} \ \mathsf{if} \ \mathsf{and} \ \mathsf{only} \ \mathsf{if} \ \mathsf{DJ} \ \mathsf{Khaled} \ \mathsf{shouts} \ \mathsf{his} \ \mathsf{name} \mathsf{",} \ \mathsf{which} \ \mathsf{is} \ \mathsf{stronger} \ \mathsf{than} \ \mathsf{the} \ \mathsf{given} \ \mathsf{statement}.$

B. 'BI ≡ Sk' - This translates to "DJ Khaled shouts his name if and only if Leroy is bad", which is the converse of the given statement.

C. $Bx \supset Sy'$ - This does not reflect the given statement, as it introduces new variables Bx' and Sy'.

 $D. \ 'Bl \supset Sk' - This \ translates \ to \ "DJ \ Khaled \ shouting \ his \ name \ implies \ that \ Leroy \ is \ bad", \ which \ accurately \ represents \ the \ given \ statement.$

Therefore, the best translation into predicate logic is: D. 'Bl \supset Sk'.

Answer Option 9:

The statement can be translated into predicate logic as follows:

Let B(x) represent "x is bad."

Let S(x) represent "x shouts DJ Khaled's name."

Let L = "Leroy."

The best translation would be "B(L) \rightarrow S(DJ Khaled)." This translates to "Leroy is bad only if DJ Khaled shouts his name."

Comparing this translation to the given options:

 $A.~'Sk \equiv Bl'-This~translation~does~not~capture~the~conditionality~and~the~specific~names~mentioned~in~the~original~statement.$

 $B.\ 'Bl \equiv Sk' - This\ translation\ does\ not\ capture\ the\ conditionality\ and\ the\ specific\ names\ mentioned\ in\ the\ original\ statement.$

 $C.\ 'Bx \supset Sy' \ -\ This\ translation\ uses\ generic\ variables\ and\ does\ not\ capture\ the\ specific\ names\ mentioned\ in\ the\ original\ statement.$

 $D.\ 'BI\supset Sk'\ -\ This\ translation\ captures\ the\ conditionality\ and\ the\ specific\ names\ mentioned\ in\ the\ original\ statement.$

Therefore, the correct translation into predicate logic is option D: 'BI \supset Sk'.

Question 76:

'Which of the given formulas of PL is the best symbolization of the following sentence?

 $If \ either \ George \ enrolls, \ respect to \ Harry \ enrolls, \ then \ Ira \ does \ not \ enroll.', \ A. \ '(G \supset H) \lor I', \ B. \ '(G \lor H) \supset I', \ C. \ '(G \supset H) \lor \neg I', \ D. \ '(G \lor H) \supset \neg I' = \neg$

Answer Option 1:

The best symbolization of the given sentence is option C: $(G \supset H) \lor \sim I'$.

The sentence states that "if either George enrolls or Harry enrolls, then Ira does not enroll."

Symbolizing this, we use:

G: George enrolls

H: Harry enrolls

I: Ira enrolls

The sentence can be translated as "(G \supset H) $\lor \neg$ 1", which means "If George enrolls, then Harry enrolls, or Ira does not enroll." This is the most accurate representation of the original sentence. Final Answer: C

Answer Option 2:

The best symbolization of the given sentence is option C: '(G \supset H) \vee ~I'.

Here's the breakdown of the symbolization:

- George enrolls: G
- Harry enrolls: H
- Ira does not enroll: ~I

The sentence "If either George enrolls or Harry enrolls, then Ira does not enroll" can be represented as " $(G \supset H) \lor \sim I$ ", which means "Either George enrolls and Harry does not enroll, or Ira does not enroll."

So, the best symbolization of the sentence is option C: '(G \supset H) \vee ~I'. Final Answer: C.

Answer Option 3:

The given sentence states, "if either George enrolls or Harry enrolls, then Ira does not enroll." Let's symbolize the statements: G: George enrolls H: Harry enrolls I: Ira enrolls The best symbolization of the given sentence would be: $(G \vee H) \supset I$ Option B: $(G \vee H) \supset I$ means "if either George enrolls or Harry enrolls, then Ira enrolls," which is the opposite of what the original sentence states. Option C: $(G \supset H) \vee I$ means "If George enrolls," which does not accurately represent the original sentence. Option D: $(G \vee H) \supset I$ means "If either George enrolls or Harry enrolls, then Ira does not enroll," which is the correct symbolization of the given sentence. Final Answer: D

Answer Option 4:

The best symbolization of the sentence "If either George enrolls or Harry enrolls, then Ira does not enroll" can be represented as $(G \lor H) \to \neg I$.

Option A: (G \supset H) \lor I: This option does not incorporate the negation of Ira's enrollment.

Option B: $(G \vee H) \supset I$: This option does not incorporate the negation of Ira's enrollment.

Option C: $(G \supset H) \lor \sim I$: This option is incorrect as it allows for the possibility of Ira enrolling in case George enrolls but Harry doesn't.

 $\label{eq:continuous} \text{Option D: } (G \vee H) \supset \neg I\text{: This option is incorrect as it allows for the possibility of Ira not enrolling even when neither George nor Harry enrolls. }$

Therefore, the best symbolization is (G \vee H) $\rightarrow \neg I.$ Final Answer: B

Answer Option 5:

The sentence "If either George enrolls or Harry enrolls, then Ira does not enroll" can be symbolized as " $(G \vee H) \rightarrow \neg I$ ".

Option A: '(G ⊃ H) ∨ I' - This formula does not capture the negation of Ira enrolling when either George or Harry enroll.

 $Option \ B: '(G \lor H) \supset I' - This formula \ does \ not \ capture \ the \ negation \ of \ Ira \ enrolling \ when \ either \ George \ or \ Harry \ enroll.$

 $Option \ C: \ '(G \supset H) \lor \neg I' - This formula \ captures \ the \ possibility \ of \ Ira \ not \ enrolling \ even \ if \ George \ enrolls, \ which \ contradicts \ the \ given \ sentence.$

 $Option \ D: \ '(G \lor H) \supset \neg I' - This \ formula \ correctly \ captures \ the \ negation \ of \ Ira \ enrolling \ when \ either \ George \ or \ Harry \ enroll.$

Therefore, the best symbolization of the given sentence is option D. Final Answer: D.

Answer Option 6:

The given sentence can be symbolized as follows: "If either George enrolls or Harry enrolls, then Ira does not enroll." Let's break it down: G: George enrolls H: Harry enrolls I: Ira enrolls The statement can be written as: $(G \vee H) \rightarrow \neg I$ Option A: $\neg I$ Option A: $\neg I$ Option A: $\neg I$ Option A: $\neg I$ This option is incorrect because it is not representing the negation of Ira enrolling. Option B: $\neg I$ This option is incorrect because it is not representing the negation of Ira enrolling. Option C: $\neg I$ This option is correct because it represents the correct symbolization of the given sentence. It states that either George enrolls implies Harry enrolls, or Ira does not enroll. Option D: $\neg I$ This option is incorrect because it is not representing the correct relationship between George and Harry enrolling and Ira not enrolling. Final Answer: C

Answer Option 7:

The best symbolization of the given sentence is option C: '(G \supset H) \vee ~I'.

The given sentence states that if either George enrolls (G) or Harry enrolls (H), then Ira does not enroll (-1). In logical terms, this can be represented as "($G \supset H$) $\lor \neg$ I", which translates to "if George enrolls, then Ira does not enroll (-1). In logical terms, this can be represented as "($G \supset H$) $\lor \neg$ I", which translates to "if George enrolls, then Ira does not enroll (-1). In logical terms, this can be represented as "($G \supset H$) $\lor \neg$ I", which translates to "if George enrolls, then Ira does not enroll (-1). In logical terms, this can be represented as "($G \supset H$) $\lor \neg$ I", which translates to "if George enrolls, then Ira does not enroll (-1). In logical terms, this can be represented as "($G \supset H$) $\lor \neg$ I", which translates to "if George enrolls, then Ira does not enroll (-1). In logical terms, this can be represented as "($G \supset H$) $\lor \neg$ I", which translates to "if George enrolls, the Ira does not enroll (-1). In logical terms, this can be represented as "($G \supset H$) $\lor \neg$ I", which translates to "if George enrolls, the Ira does not enroll (-1). In logical terms, this can be represented as "($G \supset H$) $\lor \neg$ I", which translates to "if George enrolls, the Ira does not enroll (-1). In logical terms, this can be represented as "($G \supset H$) $\lor \neg$ I", which translates to "if George enrolls, the Ira does not enroll (-1). In logical terms, this can be represented as "($G \supset H$) $\lor \neg$ I", which translates the Ira does not enroll (-1). In logical terms, this can be represented as "($G \supset H$) $\lor \neg$ I", which translates the Ira does not enroll (-1). In logical terms, the Ira does not enroll (-1) in logical terms, the Ira does not enroll (-1) in logical terms, the Ira does not enroll (-1) in logical terms, the Ira does not enroll (-1) in logical terms, the Ira does not enroll (-1) in logical terms, the Ira does not enroll (-1) in logical terms, the Ira does not enroll (-1) in logical terms, the Ira does not enroll (-1) in logical terms, the Ira does not enroll (-1) in logical terms, the Ira does not enroll (-1) in logical terms,

then Harry enrolls, or Ira does not enroll".

 $Option \ A: \ '(G\supset H)\lor I'-This \ option \ is \ incorrect \ because \ it \ does \ not \ capture \ the \ negation \ of \ Ira's \ enrollment.$

Option B: '(G v H) \supset I' - This option is incorrect as it reverses the conditional relationship between George and Harry's enrollment and Ira's enrollment.

Option D: '(G \vee H) \supset ~I' - This option is incorrect as it reverses the negation of Ira's enrollment.

Therefore, the best symbolization is option C: '(G ⊃ H) ∨ ~I'. Final Answer: C

Answer Option 8:

The given sentence states "If either George enrolls or Harry enrolls, then Ira does not enroll." Let's analyze the options:

Option A: $(G \supset H) \lor I$

This option translates to "If George enrolls, then Harry enrolls or Ira enrolls." This is not equivalent to the given sentence.

Option B: $(G \lor H) \supset I$

This option translates to "If George enrolls or Harry enrolls, then Ira enrolls." This is the opposite of what the given sentence states.

Option C: $(G \supset H) \lor \sim I$

This option translates to "If George enrolls, then Harry enrolls or Ira does not enroll." This is not equivalent to the given sentence.

Option D: $(G \lor H) \supset \sim I$

This option translates to "If George enrolls or Harry enrolls, then Ira does not enroll." This accurately represents the given sentence.

Therefore, the best symbolization of the given sentence is option D: '(G \vee H) \supset ~I'. Final Answer: D.

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Answer Option 9:

To symbolize the given sentence, we can break it down into the following components:

George enrolls: G

Harry enrolls: H

Ira does not enroll: ¬I

The sentence states "If either George enrolls or Harry enrolls, then Ira does not enroll."

This can be symbolized as "(G \vee H) \rightarrow ¬I", which means "If George enrolls or Harry enrolls, then Ira does not enroll."

Comparing this with the given options:

A. '(G \supset H) \vee l' - This option does not correctly represent the negation of Ira enrolling.

B. '(G \vee H) \supset I' - This option does not correctly represent the negation of Ira enrolling.

C. '(G \supset H) \lor ~I' - This option correctly represents the disjunction of George enrolling or Harry enrolling, but it does not correctly represent the negation of Ira enrolling.

D. '(G v H) \supset -I' - This option correctly represents the disjunction of George enrolling or Harry enrolling, and it correctly represents the negation of Ira enrolling.

Therefore, the best symbolization of the given sentence is option D: '(G \vee H) \supset ~I'. Final Answer: D.

Question 77: 'Select the best translation into predicate logic: Ms. Jackson lent me some sugar, but neither Janet nor Latoya did.', A. 'Lm v ~(JI v LI)', B. 'Lm • ~(Lj • LI)', C. 'Lm • ~(Lj • LI)', D. 'Lm v ~(Lj • LI)' Answer Option 1: Let's break down the sentence into logical expressions: Lm - Ms. Jackson lent me some sugar JI - Janet lent me some sugar LI - Latoya lent me some sugar The sentence states that Ms. Jackson lent me some sugar, but neither Janet nor Latoya did. This can be translated as: $Lm \land \sim (JI \lor LI)$ The correct translation into predicate logic is: Option A: Lm \vee ~(JI \vee LI) Final Answer: A Answer Option 2: Let's define the predicates: Lm: Ms. Jackson lent me some sugar. Jl: Janet lent me some sugar. LI: Latoya lent me some sugar. The given statement states that Ms. Jackson lent me some sugar, but neither Janet nor Latoya did. This can be translated as "Lm \wedge ~(Jl \vee Ll)", which means Ms. Jackson lent me some sugar and it is not true that either Janet or Latoya lent me some sugar. Therefore, the best translation into predicate logic is A. 'Lm \vee ~(JI \vee LI)'. ******* Answer Option 3: Let the predicate L(x, y) represent "x lent y some sugar", where the domain for x and y is the set of people. The given statement can be translated as: L(Ms. Jackson, me) A ~(L(Janet, me) V L(Latoya, me)) Breaking it down: Ms. Jackson lent me some sugar, which can be translated as L(Ms. Jackson, me). And neither Janet nor Latoya lent me any sugar, which can be translated as ~(L(Janet, me) ~ L(Latoya, Combining these translations, the best option is: C. Lm • ~(Lj \vee Ll) ****************** Answer Option 4: Let the predicates be defined as follows: Lm: Ms. Jackson lent me some sugar. JI: Janet lent me some sugar. The given sentence can be translated as "Ms. Jackson lent me some sugar, but neither Janet nor Latoya did." This can be broken down into two parts: Ms. Jackson lent me some sugar, and neither Janet nor The first part, "Ms. Jackson lent me some sugar," can be translated as Lm. The second part, "neither Janet nor Latoya did," can be translated as ~(Jl v Ll), which means "not (Janet lent me some sugar or Latoya lent me some sugar)." Combining the two parts, the best translation into predicate logic is Lm • ~(JI \vee LI). Therefore, the correct answer is B. 'Lm • ~(Lj • LI)'. Answer Option 5: Let's break down the given sentence into logical propositions: Lm: Ms. Jackson lent me some sugar Jl: Janet lent me some sugar LI: Latoya lent me some sugar

Option B: 'Lm • ~(Lj • Ll)' - This translation is incorrect as it uses a conjunction (^) instead of a disjunction (v) to represent the possibility of Ms. Jackson lending sugar.

Option A: 'Lm v ~(Jl v Ll)' - This translation is incorrect as it uses a disjunction (v) instead of a conjunction (n) to represent the combination of Ms. Jackson lending sugar and neither Janet nor Latoya doing

The sentence "Ms. Jackson lent me some sugar, but neither Janet nor Latoya did" can be translated into predicate logic as:

Lm ∧ ~(Jl ∨ Ll)

Option C: 'Lm • ~(Lj v Ll)' - This translation is incorrect as it uses a conjunction (^) instead of a disjunction (v) to represent the possibility of Ms. Jackson lending sugar.

Option D: 'Lm v ~(Lj • Ll)' - This translation correctly represents the given sentence, where Ms. Jackson lent sugar, and neither Janet nor Latoya did.

Final Answer: D

Answer Option 6:

Let Lm represent "Ms. Jackson lent me some sugar", Jl represent "Janet lent me some sugar", and Ll represent "Latoya lent me some sugar". The statement "Ms. Jackson lent me some sugar, but neither Janet nor Latoya did" can be translated as: Lm \(\times \) (Jl \(\times \) I) In predicate logic, conjunction is represented by the symbol \(\times \) and negation is represented by the symbol \(\times \) Option A: 'Lm \(\times \) (Jl \(\times \) I)' This option does not correctly represent the statement. It translates to "Either Ms. Jackson lent me some sugar or neither Janet nor Latoya did." Option B: 'Lm \(\times \) (Lj \(\times \) LI)' This option does not correctly represent the statement. It translates to "Ms. Jackson lent me some sugar and neither Janet nor Latoya lent me any sugar." Option C: 'Lm \(\times \) (Lj \(\times \) LI)' This option does not correctly represent the statement. It translates to "Either Ms. Jackson lent me some sugar and neither Janet nor Latoya lent me any sugar." Option D: 'Lm \(\times \) (Lj \(\times \) LI)' This option does not correctly represent the statement. It translates to "Either Ms. Jackson lent me some sugar or neither Janet nor Latoya lent me any sugar." Final Answer: C

Answer Option 7:

To translate the given sentence into predicate logic, we can use the following symbols:

Lm: Ms. Jackson lent me some sugar

Lj: Janet lent me some sugar

LI: Latova lent me some sugar

The sentence "Ms. Jackson lent me some sugar, but neither Janet nor Latoya did" can be translated as "Lm \land ~(Lj \lor Ll)".

Option A: 'Lm v ~(JI v LI)'

This translation does not accurately capture the meaning of the original sentence. It states that either Ms. Jackson lent me sugar or Janet or Latoya did not.

Option B: 'Lm • ~(Lj • LI)'

This translation states that Ms. Jackson lent me sugar and neither Janet nor Latoya did. It accurately captures the meaning of the original sentence.

Option C: 'Lm • ~(Lj v Ll)'

This translation states that Ms. Jackson lent me sugar and neither Janet nor Latoya did. It accurately captures the meaning of the original sentence.

Option D: 'Lm v ~(Lj • Ll)'

This translation does not accurately capture the meaning of the original sentence. It states that either Ms. Jackson lent me sugar or Janet and Latoya did not.

Therefore, the best translation into predicate logic is Option B: 'Lm \bullet ~(Lj \bullet Ll)'.

Answer Option 8

Let's break down the given sentence into logical parts:

P: Ms. Jackson lent me some sugar

Q: Janet lent me some sugar

R: Latoya lent me some sugar

The sentence "Ms. Jackson lent me some sugar, but neither Janet nor Latoya did" can be translated as "P and ~(Q or R)".

Now, let's match this translation with the options given:

Option A: Lm $\vee \sim$ (JI \vee LI) - This translation does not include the "and" operator, so it does not capture the meaning of the sentence correctly.

Option B: Lm • ~(Lj • Ll) - This translation uses the wrong variables (Lm, Lj, Ll) instead of P, Q, R. Additionally, it uses the wrong logical operator (• instead of and).

Option C: Lm • ~(Lj v Ll) - This translation uses the wrong variables (Lm, Lj, Ll) instead of P, Q, R. Additionally, it uses the wrong logical operator (• instead of and).

Option D: Lm $\vee \sim$ (Lj * Ll) - This translation uses the wrong variables (Lm, Lj, Ll) instead of P, Q, R. Additionally, it uses the wrong logical operator (* instead of and).

None of the given options correctly translate the sentence into predicate logic.

Final Answer: None of the above.

Answer Option 9:

Let's break down the sentence into logical components.

Lm: Ms. Jackson lent me some sugar.

Jl: Janet lent me some sugar.

LI: Latoya lent me some sugar.

The sentence states that Ms. Jackson lent me some sugar, but neither Janet nor Latoya did.

To represent this in predicate logic, we can use the following translation:

 $Lm \land \sim (JI \lor LI)$

Option A: Lm v ~(JI v LI

This option represents that Ms. Jackson lent me some sugar or neither Janet nor Latoya did. It does not capture the condition that Janet and Latoya did not lend me sugar simultaneously.

Option B: Lm • ~(Lj • Ll)

This option represents that Ms. Jackson lent me some sugar and neither Janet nor Latoya lent me sugar. It captures the condition that neither Janet nor Latoya lent me sugar, but it does not properly represent the fact that Ms. Jackson did lend me sugar.

Option C: Lm • ~(Lj v Ll)

This option represents that Ms. Jackson lent me some sugar and neither Janet nor Latoya lent me sugar. It correctly captures the condition stated in the sentence.

Option D: $Lm \lor \sim (Lj \bullet LI)$

This option represents that either Ms. Jackson lent me some sugar or neither Janet nor Latoya lent me sugar. It does not capture the condition that Ms. Jackson did lend me sugar.

Therefore, the best translation into predicate logic is option C: 'Lm \bullet ~(Lj \vee Ll)'.

Question 78:

'Select the best translation into predicate logic. Some animals are neglected by cruel people. (Ax: x is an animal; Cx: x is cruel; Px: x is a person; Nxy: x neglects y)', A. '(∃x)[(Ax • Cx) • (∃y)(Py • Nyx)]', B.

 $^{\shortmid}(\exists x)[(Cx\bullet Px)\bullet(\exists y)(Ay\bullet Nxy)]^{\shortmid},\ C.\ ^{\shortmid}(\exists x)[(Ax\bullet Px)\bullet(\exists y)(Cy\bullet Nxy)]^{\shortmid},\ D.\ ^{\shortmid}(\exists x)(\exists y)\{[(Ax\bullet Cy)\bullet Py]\bullet Nxy)]\}^{\prime}$

Answer Option 1:

The best translation into predicate logic would be:

A. $'(\exists x)[(Ax \cdot Cx) \cdot (\exists y)(Py \cdot Nyx)]'$

This translates to "There exists an animal x such that x is neglected by some person y, where x is an animal and y is a cruel person."

This accurately captures the meaning of the given statement. Final Answer: A.

Answer Option 2:

The best translation into predicate logic for the given statement "Some animals are neglected by cruel people" would be:

 $(\exists x)(\exists y)[(Ax \cdot Cx) \cdot (Py \cdot Nxy)]$

This translates to: There exist an x and a y such that x is an animal, x is cruel, y is a person, and y neglects x.

Therefore, the correct answer is D. $'(\exists x)(\exists y)\{[(Ax \bullet Cy) \bullet Py] \bullet Nxy)]\}'$

Answer Option 3:

The best translation into predicate logic for the statement *Some animals are neglected by cruel people" would be option D. '(∃x)(∃y){[(Ax • Cy) • Py] • Nxy))'. This translation captures the idea that there exist an animal (x) and a person (y) such that the animal is neglected by the person, and the person is cruel. The use of nested quantifiers is necessary to express the relationship between the two individuals and the properties they possess. Final Answer: D

Answer Option 4:

The best translation into predicate logic for the given sentence is: A. ' $(\exists x)[(Ax \bullet Cx) \bullet (\exists y)(Py \bullet Nyx)]'$.

This translation accurately captures the meaning of the sentence. It states that there exists an animal x such that x is neglected by some person y who is cruel. The first part, (Ax • Cx), represents that x is an animal and x is cruel. The second part, (∃y)(Py • Nyx), represents that there exists a person y who neglects x.

Answer Option 5:

The best translation into predicate logic for the statement "Some animals are neglected by cruel people" is:

 $(\exists x)[(\mathsf{Ax}\bullet\mathsf{Cx})\bullet(\exists y)(\mathsf{Py}\bullet\mathsf{Nxy})]$

Explanation:

- ($\exists x$) represents "There exists an x"
- (Ax Cx) represents "x is an animal and x is cruel"
- (∃y)(Py Nxy) represents "There exists a y such that y is a person and y neglects x"

Therefore, the correct answer is A. ' $(\exists x)[(Ax \bullet Cx) \bullet (\exists y)(Py \bullet Nxy)]'$

Answer Option 6

The best translation into predicate logic for the statement "Some animals are neglected by cruel people" is:

 $(\exists x)(\exists y)[(\mathsf{Ax} \bullet \mathsf{Cx}) \bullet (\mathsf{Py} \bullet \mathsf{Nxy})]$

Explanation:

- ($\exists x$) signifies that there exists at least one object x.
- ($\exists y$) signifies that there exists at least one object y.
- (Ax Cx) represents that x is an animal and x is cruel.
- (Py Nxy) represents that y is a person and y neglects x.

Therefore, the correct answer is A. $(\exists x)[(Ax \bullet Cx) \bullet (\exists y)(Py \bullet Nyx)].$

Answer Option 7:

The best translation into predicate logic for the statement "Some animals are neglected by cruel people" is:

 $(\exists x)[(Ax \cdot Cx) \cdot (\exists y)(Py \cdot Nxy)]$

Explanation:

Ax: x is an animal

Cx: x is cruel

Px: x is a person

Nxy: x neglects y

"Some animals are neglected" can be translated as $(\exists x)(Ax \cdot (\exists y)(Py \cdot Nxy))$. "by cruel people" can be translated as Cx. Combining the two translations, we get: $(\exists x)[(\mathsf{Ax} \bullet \mathsf{Cx}) \bullet (\exists y)(\mathsf{Py} \bullet \mathsf{Nxy})]$ Therefore, the correct answer is A. $(\exists x)[(Ax \bullet Cx) \bullet (\exists y)(Py \bullet Nxy)]$ Answer Option 8: The best translation into predicate logic would be: A. $'(\exists x)[(Ax \bullet Cx) \bullet (\exists y)(Py \bullet Nyx)]'$ Explanation: Let's break down the sentence and translate it step by step: "Some animals are neglected by cruel people." 1. We introduce the variables: - x: represents animals - y: represents people 2. We define the predicates: - Ax: x is an animal - Cx: x is cruel - Px: x is a person - Nxy: x neglects y 3. We translate the sentence: - ($\exists x$): There exists an x such that... $-\left[(\mathsf{Ax}\bullet\mathsf{Cx})\bullet(\exists y)(\mathsf{Py}\bullet\mathsf{Nyx})\right]:...x\text{ is an animal and }x\text{ is cruel, and there exists a }y\text{ such that }y\text{ is a person and }y\text{ neglects }x.$ Option A correctly captures the meaning of the sentence and accurately translates it into predicate logic. Final Answer: A Answer Option 9: The best translation into predicate logic for the given statement, "Some animals are neglected by cruel people" is: $(\exists x)[(Ax \cdot Cx) \cdot (\exists y)(Py \cdot Nxy)]$ Option A: $(\exists x)[(Ax \cdot Cx) \cdot (\exists y)(Py \cdot Nyx)]$ This translation correctly represents that there exists an animal x that is neglected by some cruel person y. Option B: $(\exists x)[(Cx \cdot Px) \cdot (\exists y)(Ay \cdot Nxy)]$ This translation represents that there exists a cruel person x who neglects some animal y, but it doesn't capture the statement that animals are neglected by cruel people. Option C: $(\exists x)[(Ax \cdot Px) \cdot (\exists y)(Cy \cdot Nxy)]$ This translation represents that there exists a person x who is an animal and neglects some cruel person y, but it doesn't capture the statement that animals are neglected by cruel people. Option D: $(\exists x)(\exists y)\{[(Ax \bullet Cy) \bullet Py] \bullet Nxy)]\}$ This translation represents that there exists an animal x and a cruel person y such that a person p neglects them, but it doesn't capture the statement that animals are neglected by cruel people. Final Answer: A

The statement can be broken down as follows:

Question 79: ' Use indirect truth tables to determine whether each set of propositions is consistent. If the set is consistent, choose an option with a consistent valuation. (There may be other consistent valuations.) $(\mathsf{E}\supset\mathsf{F})\supset\mathsf{G}$ $(\sim E \vee F) \equiv (H \vee I)$ $I\supset (J\supset \sim\! H)$ ~G ⊃ ~J', A. 'Inconsistent', B. 'Consistent. Consistent valuation when E is true and F, G, H, I, and J are false', C. 'Consistent. Consistent valuation when E and H are true and F, G, I, and J are false', D. 'Consistent. Consistent valuation when E and I are true and F, G, H, and J are false' Answer Option 1: To determine the consistency of the set of propositions, we need to create a truth table and check if there is any row where all the propositions are true. Let's create the truth table for the given propositions: $\mathsf{E}\;\mathsf{F}\;\mathsf{G}\;\mathsf{H}\;\mathsf{I}\;\mathsf{J}\;(\mathsf{\sim}\mathsf{E}\;\vee\;\mathsf{F})\equiv(\mathsf{H}\;\vee\;\mathsf{I})\;\mathsf{I}\supset(\mathsf{J}\supset\mathsf{\sim}\mathsf{H})\;\mathsf{\sim}\mathsf{G}\supset\mathsf{\sim}\mathsf{J}\;(\mathsf{E}\supset\mathsf{F})\supset\mathsf{G}$ TTTTTTTT TTTTTFTTT TTTTFTTTT TTTTFFTTT TEFFFFFF FFTTTTTT FFTTTFTTT FTTTTTTT FTTTTFFTT Based on the truth table, we can see that there is no row where all the propositions are true. Therefore, the set of propositions is inconsistent. Final Answer: A. 'Inconsistent' ******* Answer Option 2: To determine the consistency of the given set of propositions, we can construct an indirect truth table by assigning truth values to the propositions and checking if there is any row where all propositions are Let's analyze each proposition: 1. $(E \supset F) \supset G$ $\mathsf{E}\;\mathsf{F}\;\mathsf{G}\;(\mathsf{E}\supset\mathsf{F})\;(\mathsf{E}\supset\mathsf{F})\supset\mathsf{G}$ TTTTT TTFTF TFTFT TFFFT FTTTT FTFTF FFTTT FFFTT 2. $(\sim E \vee F) \equiv (H \vee I)$ $\mathsf{E}\;\mathsf{F}\;\mathsf{H}\;\mathsf{I}\;(\mathsf{\sim}\mathsf{E}\;\vee\;\mathsf{F})\;(\mathsf{H}\;\vee\;\mathsf{I})\;(\mathsf{\sim}\mathsf{E}\;\vee\;\mathsf{F})\equiv(\mathsf{H}\;\vee\;\mathsf{I})$ TTTTTTTTFTTFF TFTFTTF TEFETTE FTTTTT FTFTFFT FFTTTT FFFTFTF $3.\ I\supset (J\supset \sim\!\!H)$ $\mathsf{I}\,\mathsf{J}\,\mathsf{H}\,\mathsf{\sim}\mathsf{H}\,\mathsf{J}\supset\mathsf{\sim}\mathsf{H}\,\mathsf{I}\supset(\mathsf{J}\supset\mathsf{\sim}\mathsf{H})$ TTTFFT TTFTTT TFTFTT

TFFTTT

FTTFFT FTFTTT FFTFFT FFFTTT 4. \sim G \supset \sim J $G\ J \sim\!\! G \sim\!\! J \sim\!\! G \supset \sim\!\! J$ TTFFT TFFTT FTTFF FFTTT From the truth table analysis, we can see that there is no row where all propositions are true. Therefore, the set of propositions is inconsistent. Final Answer: A. 'Inconsistent' To determine whether the set of propositions is consistent, we can construct an indirect truth table by assuming all propositions are true and deriving any contradictions. Let's analyze each proposition: $1.\ (E\supset F)\supset G$ $\mathsf{E}\;\mathsf{F}\;\mathsf{G}\;(\mathsf{E}\supset\mathsf{F})\;(\mathsf{E}\supset\mathsf{F})\supset\mathsf{G}$ TTTTT TTFTF TFTFT TFFFT FTTTT FTFTF FFTTT FFFTF From the truth table, there is no contradiction. 2. $(\sim E \vee F) \equiv (H \vee I)$ $E F H I \sim E \vee F H \vee I (\sim E \vee F) \equiv (H \vee I)$ TTTTTT TTFFTTF TFTTTT TFFFTFF FTTTTT FTFFTTF FFTTTT FFFFTTF From the truth table, there is no contradiction. $3.\ I\supset (J\supset \sim\!\!H)$ $I \: J \: H \: {\sim} H \: J \supset {\sim} H \: I \supset (J \supset {\sim} H)$ TTTFTT TTFTFT TFTFTT TFFTTT FTTFTF FTFTFT FFTFTT FFFTTT From the truth table, there is no contradiction. 4. \sim G \supset \sim J $G\ J\ \text{$^\sim$J}\ \text{$^\sim$J}\ \text{$^\sim$J}\ \text{$^\sim$J}$ TTFFT TFFTT

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FTTFF
FFTTT
 From the truth table, there is no contradiction.
 Based on the analysis, the set of propositions is consistent. Hence, the answer is C. 'Consistent. Consistent valuation when E and H are true and F, G, I, and J are false'.
 Answer Option 4:
 To determine the consistency of the given set of propositions, we can use an indirect truth table.
 First, let's assign truth values to the propositions:
E: True, False
F: True, False
G: True, False
H: True, False
I: True, False
J: True, False
We'll evaluate the truth values of each proposition in the set for all possible combinations of truth values for E, F, G, H, I, and J.
For the first proposition, (E\supset F)\supset G:
\mathsf{E} \supset \mathsf{F} evaluates to the following truth table:
\mathsf{E}\,\mathsf{F}\,\mathsf{E}\supset\mathsf{F}
TTT
TFF
FTT
Then, (E \supset F) \supset G evaluates to the following truth table:
EFG(E\supset F)\supset G
\mathsf{T}\,\mathsf{T}\,\mathsf{T}\,\mathsf{T}
 TFTT
FTTT
FFTT
 For the second proposition, (~E \vee F) \equiv (H \vee I):
~E \vee F evaluates to the following truth table:
EF~E~E∨F
TTFT
 TFFF
FTTT
FFTF
\mbox{H} \vee \mbox{I} evaluates to the following truth table:
HIHVI
\mathsf{T}\,\mathsf{T}\,\mathsf{T}
TFT
 FTT
FFF
Then, (\sim E \vee F) \equiv (H \vee I) evaluates to the following truth table:
E F H I (\sim E \vee F) \equiv (H \vee I)
 TTTTT
TFTFF
FTTTT
FFTFT
For the third proposition, I\supset (J\supset \sim\! H) :
J \supset {\sim} H evaluates to the following truth table:
JH~HJ⊃~H
TTFF
TFTT
```

FTFT

```
FFTT
 Then, I\supset (J\supset {\sim}H) evaluates to the following truth table:
I\ J\ H\ I\supset (J\supset \sim\! H)
TTTT
\mathsf{TTFF}
TFTT
 TFFF
FTTT
FTFT
FFTT
 FFFT
For the fourth proposition, {\sim}G \supset {\sim}J:
{\sim}G\supset {\sim}J evaluates to the following truth table:
G\ J\ \text{$^\sim$J}\ \text{$^\sim$J}\ \text{$^\sim$J}
TTFFT
TFFTT
FTTFF
 FFTTT
 Based on the truth tables, we can see that there is no row where all the propositions evaluate to true. Therefore, the set of propositions is inconsistent.
 Final Answer: A. 'Inconsistent'
 *******
 Answer Option 5:
To determine whether the set of propositions is consistent or not, we can construct an indirect truth table by assigning truth values to the propositions involved and evaluating the truth values of the compound
propositions.
Let's analyze the given set of propositions:
1.\; (E\supset F)\supset G
2. (\sim E \vee F) \equiv (H \vee I)
3. I ⊃ (J ⊃ ~H)
To construct the indirect truth table, we assign truth values to the propositions E, F, G, H, I, and J.
\mathsf{E}\;\mathsf{F}\;\mathsf{G}\;\mathsf{H}\;\mathsf{I}\;\mathsf{J}\;(\sim\!\mathsf{E}\;\vee\;\mathsf{F})\equiv (\mathsf{H}\;\vee\;\mathsf{I})\;(\mathsf{E}\supset\mathsf{F})\supset\mathsf{G}\;\mathsf{I}\supset (\mathsf{J}\supset\sim\!\mathsf{H})\;\sim\!\mathsf{G}\supset\sim\!\mathsf{J}
TTTTTTTT
TTTTTFTTTT
TTTTFTTTT
TTTTFFTTTT
TTTFTTTTTT
TTTFTFTTTT
TTTFFTTTTT
TTTFFFTTTT
TFTTTTTTT
TFTTTFFFTT
TETTETTETT
TETTEETETT
TETETTTETT
TFTFTFFTT
TETEETTETT
TFTFFFTFTT
FTTTTTTTT
FTTTTFTTTT
FTTTFTTTT
FTTTFFTTTT
FTTFTTTTT
```

FTTFTFTTTT

FTTFFTTTTT
FTTFFFTTTT
FFTTTTTTT
FFTTTFTTTT
FFTTFTTTTT
FFTTFFTTTT
FFTFTTTTT
FFTFTFTTTT
FFTFFTTTTT
FFTFFFTTTT
Based on the indirect truth table, we can conclude that the set of propositions is consistent.
However, the provided answer choices do not accurately represent the possible consistent valuations. The correct answer should be: 'Consistent. There are multiple consistent valuations for the given set of
propositions.'
Final Answer: 'Consistent. There are multiple consistent valuations for the given set of propositions.'
•••••••
Answer Option 6:
To determine the consistency of the set of propositions, we can construct an indirect truth table by considering all possible truth value combinations for the propositions involved.
Let's assign truth values to the propositions as follows:
E: True or False
F: True or False
G: True or False
H: True or False
l: True or False
J: True or False
Using these truth values, we can evaluate each proposition in the set:
1. (E⊃F)⊃G
We can construct the truth table for this proposition:
$EFG(E\supset F)((E\supset F)\supset G)$
TTTTT
TIFIF
TFTFT
TFFFT
FITTT
FTFTF
FFTTT
FFFTT
2. $(\sim E \vee F) \equiv (H \vee I)$
We can construct the truth table for this proposition:
$EFHI(\simE\veeF)(H\veeI)(\simE\veeF)\equiv(H\veeI)$
TTTTTTT
TTTFTTT
TTFTTTT
TTFFTFF
TFTTTTT
TETEFTE
TFFTTTT
TFFFFFT
FTTTTTT

FTTFTTT FTFTTTT FTFFTFF

```
FFTTTT
FFTFTTT
FFFTTTT
FFFFFF
3.\ I\supset (J\supset \sim\!\!H)
We can construct the truth table for this proposition:  \\
I\ J\ H\ (\sim\!H)\ (J\supset \sim\!H)\ (I\supset (J\supset \sim\!H))
TTTFFT
TTFTTT
TFTFFF
TFFTTT
FTTFFT
FTFTTT
FFTFTT
FFFTTT
4. \simG \supset \simJ
We can construct the truth table for this proposition:
G\ J\ (\sim\!G)\ (\sim\!J)\ (\sim\!G\supset\sim\!J)
TTFFT
TFFTT
FTTFF
FFTTT
Now, let's analyze the results:
From the truth tables, we can see that the set of propositions is consistent. There are consistent valuations for the propositions.
Option C: Consistent. Consistent valuation when E and H are true and F, G, I, and J are false.
Answer Option 7:
To determine the consistency of the set of propositions, let's construct an indirect truth table.
We have 5 propositions: E, F, G, H, and I. We also have the implication and equivalence operators.
Step 1: Assign truth values to the propositions and evaluate the expressions.
E \; F \; G \; H \; I \; J \; (\sim\!E \; \vee \; F) \equiv (H \; \vee \; I) \; I \supset (J \supset \sim\!H) \; \sim\!G \supset \sim\!J \; (E \supset F) \supset G
TTTTTTTT
TTTTTFTFT
TTTTFTTTT
TTTTFFTFFT
TTTFTTTTT
TTTFTFTFT
TTFTTTTTF
TTFTTFTFF
TTFTFTTTTF
TTFTFFTFF
TFTTTTFTTT
TFTTTFFFFT
TFTTFTTTT
TFTTFFFFF
TFTFTTTT
TETETEFET
```

TFFTTTFTTT TFFTTFFFFT TFFTFTTT

TFFTFFFFF TEFETTETTT TFFFTFFFT TEFFETETTT TFFFFFFFT FTTTTTTTT FTTTTFTFFT FTTTFTTTT FTTTFFTFFT FTTFTTTTT FTTFTFFF FTFTTTTTT FTFTTFTFT FTFTFTTTT FTFTFFTFFT FFTTTTTTF FFTTTFTFF FFTTFTTTF FFTTFFTFF FFTFTTTTF FFTFTFFF FFFTTTTTF FFFTTFTFF FFFTFTTTF FFFTFFFF FFFFTTTTF FFFFFFFF FFFFFTTTF FFFFFFFF

Step 2: Check if there is any row where all the expressions are true.

From the truth table, we can see that there is no row where all the expressions are true.

Therefore, the set of propositions is inconsistent.

Final Answer: A. 'Inconsistent'

|T|T|F|F|F|T|T|F|T|F|

Answer Option 8:

To determine the consistency of the set of propositions using indirect truth tables, we will create a truth table for all the propositions involved.

 $\mid E \mid F \mid G \mid H \mid I \mid J \mid (\neg E \vee F) \equiv (H \vee I) \mid I \supset (J \supset \neg H) \mid \neg G \supset \neg J \mid (E \supset F) \supset G \mid$

|T|T|F|F|F|F|T|F|T|F| |T|F|T|T|T|T|T|T|T|T $|\mathsf{T}|\mathsf{F}|\mathsf{T}|\mathsf{T}|\mathsf{T}|\mathsf{F}|\mathsf{T}|\mathsf{T}|\mathsf{T}|\mathsf{T}|$ |T|F|T|T|F|T|F|T|T|T| |T|F|T|T|F|F|F|T|T|T| |T|F|T|F|T|T|T|T|T|T |T|F|T|F|T|F|T|T|T|T |T|F|T|F|F|T|T|T|T|T|T|F|T|F|F|F|T|T|T|T||T|F|F|T|T|T|T|F|T|T||T|F|F|T|T|F|T|F|T|T| |T|F|F|T|F|T|F|F|T|T| |T|F|F|T|F|F|F|F|T|T| $|\mathsf{T}|\mathsf{F}|\mathsf{F}|\mathsf{F}|\mathsf{T}|\mathsf{T}|\mathsf{T}|\mathsf{F}|\mathsf{T}|\mathsf{T}|$ |T|F|F|F|T|F|T|F|T|T||T|F|F|F|F|T|T|F|T|T| |T|F|F|F|F|F|T|F|T|T| |F|T|T|T|T|T|T|T|T|T|T|F|T|T|T|T|F|T|T|T|T|F|T|T|T|F|T|T|T|T|T|F|T|T|T|F|F|T|T|T|T |F|T|T|F|T|T|T|T|T|T |F|T|T|F|T|F|T|T|T|T||F|T|T|F|F|T|T|T|T| $|\mathsf{F}|\mathsf{T}|\mathsf{T}|\mathsf{F}|\mathsf{F}|\mathsf{F}|\mathsf{T}|\mathsf{T}|\mathsf{T}|\mathsf{T}|$ |F|T|F|T|T|T|T|T|T|T |F|T|F|T|T|F|T|T|T|T |F|T|F|T|F|T|T|T|T|T |F|T|F|T|F|F|T|T|T|T |F|T|F|F|T|T|T|T|T|T|F|T|F|F|T|F|T|T|T| |F|T|F|F|F|T|T|T|T| |F|T|F|F|F|F|T|T|T|T| |F|F|T|T|T|T|T|T|T|T|F|F|T|T|T|F|T|T|T|T $|\mathsf{F}|\mathsf{F}|\mathsf{T}|\mathsf{T}|\mathsf{F}|\mathsf{T}|\mathsf{T}|\mathsf{T}|\mathsf{T}|\mathsf{T}|\mathsf{T}|$ |F|F|T|T|F|F|T|T|T|T |F|F|T|F|T|T|T|T|T| |F|F|T|F|T|F|T|T|T|T |F|F|T|F|F|T|T|T|T| |F|F|T|F|F|F|T|T|T|T| |F|F|F|T|T|T|T|F|T| |F|F|F|T|T|F|T|T|F|T| |F|F|F|T|F|T|T|T|F|T| |F|F|F|T|F|F|T|T|F|T| |F|F|F|F|T|T|T|T|T| $|\mathsf{F}|\mathsf{F}|\mathsf{F}|\mathsf{F}|\mathsf{F}|\mathsf{T}|\mathsf{F}|\mathsf{T}|\mathsf{T}|\mathsf{T}|\mathsf{T}|$ |F|F|F|F|F|T|T|T|T| |F|F|F|F|F|F|T|T|T|

From the truth table, we can see that there is no row where all the propositions in the set are false. So, the set of propositions is consistent.

However, the option that correctly represents a consistent valuation is C. Consistent valuation when E and H are true and F, G, I, and J are false.

To determine the consistency of the set of propositions, we can construct an indirect truth table by assigning truth values to each proposition and analyzing the resulting truth values of the entire set.

Let's denote the propositions as follows:

 $A{:}\; (E\supset F)\supset G$

B: $(\sim E \vee F) \equiv (H \vee I)$

 $C{:}\ I\supset (J\supset {\sim} H)$

D: ~G ⊃ ~J

We'll evaluate all possible truth value combinations for propositions A, B, C, and D.

E|F|G|H|I|J|A|B|C|D|

T|T|T|T|T|T|T|T|T|T|T

TITITITIFITITITITI

TITITITIFITITITITI

TITITITIFIFITITITITI

T|T|T|F|T|T|T|T|T|T

TITITIFITIFITITITITI

T|T|T|F|F|T|T|T|T|

T|T|T|F|F|F|T|T|T|T|

TITIFITITIFIFITITI

TITIFITITIFIFIFITITI

T|T|F|T|F|T|F|F|T|T|

T|T|F|T|F|F|F|F|T|T|

T|T|F|F|T|T|F|F|T|T|

T|T|F|F|T|F|F|F|T|T|

TITIFIFIFITIFIFITITI

T|T|F|F|F|F|F|F|T|T|

F|T|T|T|T|T|T|T|T

F|T|T|T|F|T|T|T|

FITITIFITITITITI

F|T|T|T|F|F|T|T|T|T

FITITIFITITITITITI

F|T|T|F|T|F|T|T|T|

F|T|T|F|F|T|T|T|T|

F|T|T|F|F|F|T|T|T|

F|T|F|T|T|T|F|T|T|

F|T|F|T|T|F|T|F|T|T|

F|T|F|T|F|T|T|T|T|

F|T|F|T|F|F|T|T|T|

F|T|F|F|T|T|T|F|T|T|
F|T|F|F|T|F|T|T|

F|T|F|F|F|T|T|T|T|

F|T|F|F|F|F|T|T|T|

F|F|T|T|T|T|F|T|T|

F|F|T|T|T|F|T|F|T|T

FIFITITIFITITITITIT

F|F|T|T|F|F|T|T|T|T

F|F|T|F|T|T|T|F|T|T|

F|F|T|F|T|F|T|F|T|T| F|F|T|F|F|T|T|T|T|

F|F|T|F|F|F|T|T|T|

F|F|F|T|T|T|F|T|T|

F|F|F|T|T|F|T|F|T|T|

F|F|F|T|F|T|T|T|T|

F|F|F|T|F|F|T|T|T|T

Based on the truth table, we can see that there are consistent valuations for all propositions. However, none of the given options match the consistent valuation.

Thus, the correct answer is: 'Consistent, but none of the given options match the consistent valuation.'

Question 80:

 $\text{`Select the best translation into predicate logic: Some students work hard.', A. '(\exists x)(Wx \lor Sx)', B. '(\exists x)(Sx \lor Wx)', C. '(\exists x)(Sx • Wx)', D. '(\exists x)(Sx ⊃ Wx)'}$

Answer Option 1:

The best translation into predicate logic for the statement "Some students work hard" would be option A: $(\exists x)(Wx \vee Sx)$. This translates to "There exists a student x such that x works hard or x is a student."

This captures the meaning that there is at least one student who works hard. Final Answer: A.

Answer Option 2:

The best translation into predicate logic for the statement "Some students work hard" is Option A: $(\exists x)(Wx \lor Sx)$. This translates to "There exists a student x such that x works hard or x is a student." This captures the idea that there is at least one student who works hard. Final Answer: A

captures the idea that there is at least one

Answer Option 3:

The best translation into predicate logic for the statement "Some students work hard" is:

A. $(\exists x)(Wx \lor Sx)$

This translates to "There exists a student such that they work hard or they are a student."

In this translation, $(\exists x)$ represents the existential quantifier "there exists", Wx represents "x works hard", and Sx represents "x is a student."

So, option A is the correct translation, Final Answer: A

Answer Option 4:

The best translation into predicate logic for the statement "Some students work hard" is:

A. $(\exists x)(Wx \vee Sx)$

Here, ($\exists x$) represents the existential quantifier, which means "there exists" and is used to indicate the existence of at least one student. Wx represents the predicate "x works hard" and Sx represents the predicate "x is a student. The disjunction (\lor) is used to indicate that the student either works hard or is a student.

Final Answer: A

Answer Option 5:

The best translation into predicate logic for the statement *Some students work hard" is option A: '(∃x)(Wx ∨ Sx)'. This translates to *There exists some x such that x works hard or x is a student.*

Answer Option 6:

The best translation into predicate logic for the statement "Some students work hard" is: A. $(\exists x)(Wx \lor Sx)$

This translates to: There exists a student \boldsymbol{x} such that \boldsymbol{x} works hard or \boldsymbol{x} is a student.

This statement captures the idea that there is at least one student who works hard.

Answer Option 7:

The best translation into predicate logic for the statement "Some students work hard" is:

A. $(\exists x)(Wx \lor Sx)$

This translates to: There exists an \boldsymbol{x} such that \boldsymbol{x} works hard or \boldsymbol{x} is a student.

Option A correctly captures the statement that there exists at least one student who works hard. Final Answer: A

Answer Option 8:

The best translation into predicate logic for the statement "Some students work hard" is:

A. $(\exists x)(Wx \lor Sx)$

This translates to "There exists at least one student x such that x works hard or x is a student." It accurately captures the idea that there is at least one student who works hard.

Option B, $(\exists x)(Sx \lor Wx)$, translates to "There exists at least one student x such that x is a student or x works hard." This statement allows for the possibility that there could be a student who is not working . . .

Option C, ($\exists x$)(Sx • Wx), translates to "There exists at least one student x such that x is a student and x works hard." This statement requires that all students both work hard and are students, which may not have the coop.

Option D, ($\exists x$)($Sx \supset Wx$), translates to "There exists at least one student x such that if x is a student, then x works hard." This statement allows for the possibility that there are students who do not work hard. Therefore, the best translation is A. ($\exists x$)($Wx \lor Sx$).

Answer Option 9:

The best translation into predicate logic for the statement "Some students work hard" is option A: ' $(\exists x)(Wx \vee Sx)$ '.

This translation represents the existence (\exists) of at least one student (x) who either works hard (Wx) or is a student (Sx).

Option B '($\exists x$)($Sx \lor Wx$)' represents the existence of at least one student who is either a student (Sx) or works hard (Wx), which does not capture the intended meaning of the original statement.

Option C '($\exists x$)(Sx • Wx)' represents the existence of at least one student who is both a student (Sx) and works hard (Wx), which is stricter than the original statement.

Option D '($\exists x$)(Sx ⊃ Wx)' represents the existence of at least one student who, if they are a student (Sx), then they work hard (Wx), which is a stronger condition than the original statement.

Final Answer: A

Question 81:

'Which of the given formulas of PL is the best symbolization of the following sentence?

Either England's importing beef is not a necessary condition for France's subsidizing agriculture, or China promotes human rights if South Africa supplies diamonds.', A. $'(-F \supset E) \lor (C \equiv -S)'$, B. $'(-E \supset F) \lor (C \supset$

 $\supset S)',\,C.\;{}^{\prime}(\sim\!\!F\supset E)\vee(C\equiv S)',\,D.\;{}^{\prime}\!\!\sim\!\!(F\supset E)\vee(S\supset C)'$

Answer Option 1:

Let's break down the sentence to identify the propositions:

- England's importing beef is not a necessary condition for France's subsidizing agriculture: \sim F \supset E
- China promotes human rights if South Africa supplies diamonds: $\mathsf{S}\supset\mathsf{C}$

The sentence states that either the first proposition is true or the second proposition is true. The best symbolization of this sentence would be the disjunction (OR) of the two propositions:

 $(\sim F \supset E) \lor (S \supset C)$

None of the given options exactly match this symbolization. Therefore, none of the provided options are the best symbolization of the sentence.

......

Answer Option 2:

To symbolize the given sentence, let's break it down into two parts:

Part 1: "England's importing beef is not a necessary condition for France's subsidizing agriculture"

Let's symbolize England's importing beef as E, France's subsidizing agriculture as F, and "not a necessary condition" as α .

The symbolization for Part 1 is: ~E ⊄ F.

Part 2: "China promotes human rights if South Africa supplies diamonds"

Let's symbolize China promotes human rights as C and South Africa supplies diamonds as S.

The symbolization for Part 2 is: $S \supset C$.

Now, let's combine both parts using the logical operator "or" (v).

The best symbolization for the given sentence is:

(~E ⊄ F) ∨ (S ⊃ C)

Final Answer: D

Answer Option 3:

To symbolize the given sentence, let's break it down:

- 1. "England's importing beef is not a necessary condition for France's subsidizing agriculture" can be symbolized as " \sim E \supset F".
- 2. "China promotes human rights if South Africa supplies diamonds" can be symbolized as "S \supset C".

Now, let's put these symbolizations together:

"Either England's importing beef is not a necessary condition for France's subsidizing agriculture, or China promotes human rights if South Africa supplies diamonds" can be symbolized as "(~E ⊃ F) ∨ (S ⊃ C)"

Comparing this with the options:

A. '(~F \supset E) \vee (C \equiv ~S)' is not an accurate symbolization of the given sentence.

B. '(~E \supset F) \vee (C \supset S)' is not an accurate symbolization of the given sentence.

C. '(~F \supset E) \vee (C \equiv S)' is not an accurate symbolization of the given sentence.

D. '~(F \supset E) \vee (S \supset C)' is not an accurate symbolization of the given sentence.

Thus, the best symbolization of the given sentence is: (A) '(~E \supset F) \vee (S \supset C)'.

Answer Option 4:

To symbolize the given sentence, let's break it down into two parts:

- 1. England's importing beef is not a necessary condition for France's subsidizing agriculture.
- 2. China promotes human rights if South Africa supplies diamonds.

Let's analyze each option:

A. $(\sim F \supset E) \lor (C \equiv \sim S)'$

This option symbolizes the first part correctly as $(-F\supset E)$, but it symbolizes the second part incorrectly as $(C\equiv \sim S)$. It should be $(C\supset S)$ instead. Therefore, this option is not the best symbolization.

B. $(\sim E \supset F) \lor (C \supset S)'$

This option symbolizes the first part incorrectly as (-E \supset F). It should be (E \supset -F) instead. The second part, (C \supset S), is correctly symbolized. Therefore, this option is not the best symbolization.

C. $'(\sim F \supset E) \lor (C \equiv S)'$

This option symbolizes the first part correctly as $(-F \supset E)$. However, it symbolizes the second part incorrectly as $(C \equiv S)$. It should be $(C \supset S)$ instead. Therefore, this option is not the best symbolization.

D. '~(F \supset E) \lor (S \supset C)'

This option symbolizes the first part correctly as \neg (F \supset E). It symbolizes the second part correctly as (S \supset C). Therefore, this option is the best symbolization.

Final Answer: D

Answer Option 5:

To determine the best symbolization of the given sentence, let's break down the sentence into its components:

- 1. England's importing beef is not a necessary condition for France's subsidizing agriculture.
- 2. China promotes human rights if South Africa supplies diamonds.

Let's analyze each answer choice:

A. $'(\sim F \supset E) \lor (C \equiv \sim S)'$

This formula represents the first component correctly (England's importing beef is not a necessary condition for France's subsidizing agriculture), but the second component is not accurately represented. It states that China promotes human rights if South Africa does not supply diamonds, which is different from the original sentence.

 $\mathsf{B.} \ '(\mathsf{\sim}\mathsf{E}\supset\mathsf{F}) \lor (\mathsf{C}\supset\mathsf{S})'$

This formula does not accurately represent the first component. It states that if England does not import beef, France subsidizes agriculture, which is the opposite of the original sentence. The second component is represented correctly.

C. $'(\sim F \supset E) \lor (C \equiv S)'$

This formula represents both components accurately. It states that England's importing beef is not a necessary condition for France's subsidizing agriculture, and China promotes human rights if and only if South Africa supplies diamonds.

D. ' \sim (F \supset E) \vee (S \supset C)'

This formula does not accurately represent the first component. It states the negation of England's importing beef being a necessary condition for France's subsidizing agriculture. The second component is represented correctly.

Based on the above analysis, the best symbolization of the given sentence is C. '(\sim F \supset E) \lor (C \equiv S)'. Final Answer: C.

Answer Option 6:

To symbolize the given sentence, let's break it down.

- 1. England's importing beef is not a necessary condition for France's subsidizing agriculture: This can be symbolized as ~E ⊃ F, where ~E represents "not England's importing beef" and F represents "France's subsidizing agriculture".
- 2. China promotes human rights if South Africa supplies diamonds: This can be symbolized as S ⊃ C, where S represents "South Africa supplies diamonds" and C represents "China promotes human rights".
 Now, let's combine the two statements using the logical operator "or" (∨) to symbolize the overall sentence.

Option A: $(\sim F \supset E) \lor (C \equiv \sim S)'$

This option does not accurately symbolize the given sentence. It introduces additional conditions (-F and ≡) that are not present in the original sentence.

Option B: $(\sim E \supset F) \lor (C \supset S)'$

This option accurately symbolizes the given sentence. It represents the two conditions mentioned above and combines them with the logical operator "or" (v).

Option C: $(\sim F \supset E) \lor (C \equiv S)'$

This option does not accurately symbolize the given sentence. It introduces an additional condition (=) that is not present in the original sentence.

Option D: '~ $(F\supset E)\lor (S\supset C)'$

This option does not accurately symbolize the given sentence. It introduces a negation (~) that is not present in the original sentence.

Based on the analysis, the best symbolization of the given sentence is Option B: '(\sim E \supset F) \lor (C \supset S)'. Final Answer: B

Answer Option 7:

Let's break down the sentence:

- 1. "England's importing beef is not a necessary condition for France's subsidizing agriculture" can be symbolized as $\sim F \supset E$.
- 2. "China promotes human rights if South Africa supplies diamonds" can be symbolized as $S \supset C$.

Now let's look at the options:

A. '(-F \supset E) \lor (C \equiv ~S)': This option includes the logical equivalence operator (\equiv), which is not necessary for symbolizing the given sentence. Additionally, it does not capture the structure of the sentence correctly. Therefore, option A is not the best choice.

B. '(-E \supset F) \lor (C \supset S)': This option incorrectly reverses the necessary condition relationship and does not accurately represent the given sentence. Therefore, option B is not the best choice.

C. '(-F > E) v (C = S)': This option correctly represents the first part of the sentence, -F > E, but the second part, C = S, does not accurately represent the relationship between China promoting human rights and South Africa supplying diamonds. Therefore, option C is not the best choice.

D. '~(F ⊃ E) ∨ (S ⊃ C)': This option correctly represents the first part of the sentence, ~(F ⊃ E), which is equivalent to ~F ∨ E. The second part, S ⊃ C, accurately represents the relationship between China promoting human rights and South Africa supplying diamonds. Therefore, option D is the best choice.

Final Answer: D

Answer Option 8:

To symbolize the given sentence, we can break it down into two parts:

- 1. England's importing beef is not a necessary condition for France's subsidizing agriculture.
- 2. China promotes human rights if South Africa supplies diamonds.

Let's analyze each option to determine the best symbolization:

Option A: $(\sim F \supset E) \lor (C \equiv \sim S)'$

This option symbolizes the first part correctly by representing it as (-F ⊃ E), meaning "If France does not subsidize agriculture, then England does not import beef." However, the second part (C = -S), meaning "China promotes human rights if and only if South Africa does not supply diamonds," does not accurately capture the second part of the sentence.

Option B: $(\sim E \supset F) \lor (C \supset S)$

This option symbolizes the first part as (-E \supset F), meaning "If England does not import beef, then France subsidizes agriculture," which is the opposite of what the sentence states. Additionally, the second part (C \supset S), meaning "If China promotes human rights, then South Africa supplies diamonds," also does not capture the intended meaning.

Option C: $(\sim F \supset E) \lor (C \equiv S)'$

This option symbolizes the first part correctly as $(-F \supset E)$, but the second part (C = S), meaning "China promotes human rights if and only if South Africa supplies diamonds," is not accurate based on the original sentence.

Option D: $'\sim$ (F \supset E) \vee (S \supset C)'

This option represents the first part as –(F \supset E), meaning "England's importing beef is not a necessary condition for France's subsidizing agriculture," which aligns with the original sentence. The second part (S \supset C), meaning "If South Africa supplies diamonds, then China promotes human rights," also correctly represents the second part of the sentence.

Based on the analysis, the best symbolization of the given sentence is Option D: '~ $(F\supset E)\lor (S\supset C)$ '. Final Answer: D

Answer Option 9:

To symbolize the given sentence, let's break it down into two parts:

- 1. England's importing beef is not a necessary condition for France's subsidizing agriculture
- 2. China promotes human rights if South Africa supplies diamonds

Let's analyze each option:

Option A: $(\sim F \supset E) \lor (C \equiv \sim S)'$

This option symbolizes the first part correctly as $(-F \supset E)$ which means "if France does not subsidize agriculture, then England does import beef." However, the second part $(C \equiv -S)$ symbolizes "China promotes human rights if South Africa does not supply diamonds," which is the opposite of what is stated in the original sentence.

Option B: $'(\sim E \supset F) \lor (C \supset S)'$

This option symbolizes the first part as (-E > F) which means "if England does not import beef, then France subsidizes agriculture." This is the opposite of what is stated in the original sentence. The second part (C > S) symbolizes "China promotes human rights if South Africa supplies diamonds," which is correctly symbolized.

Option C: $(\sim F \supset E) \lor (C \equiv S)$

This option symbolizes the first part correctly as (-F \supset E) which means "if France does not subsidize agriculture, then England imports beef." However, the second part (C = S) symbolizes "China promotes human rights if South Africa supplies diamonds," which is the opposite of what is stated in the original sentence.

Option D: $'\sim$ (F \supset E) \lor (S \supset C)'

This option symbolizes the first part as $\neg(F \supset E)$ which means "France subsidizing agriculture is not a necessary condition for England to import beef." This is the opposite of what is stated in the original sentence. The second part $(S \supset C)$ symbolizes "South Africa supplying diamonds is a condition for China to promote human rights," which is the opposite of what is stated in the original sentence.

Based on the analysis, none of the options correctly symbolize the given sentence. Therefore, none of the options are the best symbolization.

Question 82:

'Construct a complete truth table for the following argument. Then, using the truth table, determine whether the argument is valid or invalid. If the argument is invalid, choose an option which presents a counterexample. (There may be other counterexamples as well.)

~A · ~B

(A v C) v B / C', A. "Valid', B. "Invalid. Counterexample when A and B are true and C is false', C. "Invalid. Counterexample when A is true and B and C are false', D. "Invalid. Counterexample when B is true and

A and C are false!

Answer Option 1:

To construct a truth table for the argument, we need to consider all possible combinations of truth values for the variables A, B, and C.

A B C ~A ~B ~A · ~B (A ∨ C) ∨ B C

TTTEFETT

TTEFFFTF

TFTFTFTT

.

TFFFTFTF

FTTTFFTT

FTFTFFF

FFTTTTT

FEFTTTEE

Using the truth table, we can evaluate the argument. The argument states that $\sim\!A\cdot\sim\!B$ and $(A\vee C)\vee B$ together implies C.

Looking at the truth table, we can see that whenever $\neg A \cdot \neg B$ is true and $(A \lor C) \lor B$ is true, C is always true. There are no counterexamples where both $\neg A \cdot \neg B$ and $(A \lor C) \lor B$ are true and C is false.

Therefore, the argument is valid.

Final Answer: A. Valid

Answer Option 2:

To construct the truth table, we consider all possible combinations of truth values for the propositions A, B, and C.

A B C ~A ~B ~A \cdot ~B (A \vee C) \vee B C

TTTFFFTTT

TTFFFFTFF

TFTFTTTT

TFFFTFF

FTTTFFTTT

FTFTFFFF

FFTTTFTTT

FFFTTFFF

From the truth table, we can observe that there are no rows where the premises $(\sim A \cdot \sim B)$ and $((A \lor C) \lor B)$ are true and the conclusion C is false. Therefore, the argument is valid.

Final Answer: A. Valid

Answer Option 3:

To construct a complete truth table for the argument, we will consider all possible combinations of truth values for the variables A, B, and C.

A | B | C | ~A | ~B | ~A \cdot ~B | (A \vee C) \vee B | C

T|T|T|F|F|F|T|T

T|T|F|F|F|F|T|F

TIFITIFITIFITIT

T|F|F|F|T|F|T|F

F|T|T|T|F|F|T|T F|T|F|T|F|F|T|F

F|F|T|T|T|T|T

F|F|F|T|T|F|F

Based on the truth table, we can see that whenever the premises (-A · -B) and ((A ∨ C) ∨ B) are true, the conclusion C is also true. Therefore, the argument is valid.

Therefore, the correct answer is A. 'Valid'.

Answer Option 4:

To construct a complete truth table for the argument, we need to consider all possible combinations of truth values for the variables A, B, and C.

A B C ~A ~B ~A \cdot ~B (A \vee C) \vee B C

TTTEFETT TTEFEETE TETETETT

TFFFTFTF FTTTFFTT

FTFTFFTF

FFTTTTT

FFFTTTFF

From the truth table, we can see that for each row where the premises $(-A \cdot -B)$ and $((A \lor C) \lor B)$ are both true, the conclusion C is also true. Therefore, the argument is valid.

Final Answer: A. 'Valid'

Answer Option 5:

To construct a complete truth table for the argument, we consider all possible combinations of truth values for the propositions A, B, and C.

Let's first define the propositions in the argument:

P1: ~A · ~B

P2: (A v C) v B

Conclusion: C

Now, let's construct the truth table:

 $\mid A \mid B \mid C \mid \neg A \cdot \neg B \mid (A \lor C) \lor B \mid C \mid$

|---|---|---|---|

|T|T|T|F|T|T|

|T|T|F|F|T|F|

|T|F|T|F|T|T|

|T|F|F|F|T|F|

|F|T|T|F|T|T|

|F|T|F|F|T|F|

|F|F|T|T|T|T|

|F|F|F|T|F|F|

Based on the truth table, we can see that in all cases where the premises (~A · ~B and (A ∨ C) ∨ B) are true, the conclusion (C) is also true. Hence, the argument is valid.

Therefore, the answer is A. "Valid".

Answer Option 6:

 $To \ construct \ a \ complete \ truth \ table \ for \ the \ argument, \ we \ need \ to \ consider \ all \ possible \ truth \ values \ for \ the \ propositions \ A, \ B, \ and \ C.$

A | B | C | ~A | ~B | ~A \cdot ~B | (A \vee C) \vee B | C

T|T|T|F|F|F|T|T

TITIFIFIFIFITIF

TIFITIFITIFITIT

T|F|F|F|T|F|T|F

F|T|T|F|F|T|T

F|T|F|T|F|F|T|F

F|F|T|T|T|T|T F|F|F|T|T|T|F|F

Based on the truth table, we can see that there is no row where both the premises (-A·-B) and ((A v C) v B) are true and the conclusion (C) is false. Therefore, the argument is valid.

Final Answer: A. 'Valid'

Answer Option 7:

To construct a complete truth table for the argument, we need to consider all possible combinations of truth values for the variables A, B, and C.

 $A\mid B\mid C\mid \neg A\mid \neg B\mid \neg A\cdot \neg B\mid (A\vee C)\vee B\mid C$

T|T|T|F|F|F|T|T T|T|F|F|F|F|T|F T|F|T|F|T|F|T|T T|F|F|F|T|F|T|F FITITITIFIFITIT F|T|F|T|F|F|T|F F|F|T|T|T|T|T|T

Looking at the truth table, we can see that the argument is invalid because there are cases where the premise (-A · -B) is true, but the conclusion C is false. Specifically, in the case where A is true, B is true, and C is false, the premise is true, but the conclusion is false.

Therefore, the correct answer is B. 'Invalid. Counterexample when A and B are true and C is false'.

F|F|F|T|T|T|F|F

Answer Option 8:

In order to construct a complete truth table for the argument, we need to consider all possible truth value combinations for the variables A, B, and C.

Let's break down the argument into its components:

1. ~A · ~B

2. (A v C) v B

3. C

We will examine all possible combinations of truth values for A, B, and C and evaluate the truth value of each component of the argument.

A | B | C | \sim A | \sim B | \sim A $\cdot \sim$ B | (A \vee C) \vee B | C

T|T|T|F|F|F|T|T

TITIFIFIFIFITIF

T|F|T|F|T|F|T|T

T|F|F|F|T|F|T|F

F|T|T|T|F|F|T|T F|T|F|T|F|F|F|F

F|F|T|T|T|T|T F|F|F|T|T|T|F|F

Looking at the truth table, we can see that whenever all of the premises are true, the conclusion is also true. Therefore, the argument is valid.

Final Answer: A. Valid

Answer Option 9:

To construct a complete truth table for the given argument, we need to consider all possible combinations of truth values for the propositions A, B, and C.

Premise 1: ~A · ~B

Premise 2: $(A \lor C) \lor B$

We need to evaluate the conclusion, which is C.

Let's construct the truth table:

 $\mid A \mid B \mid C \mid \neg A \mid \neg B \mid \neg A \cdot \neg B \mid (A \vee C) \vee B \mid C \mid$

|T|T|T|F|F|F|T|T|

|T|T|F|F|F|F|T|F|

|T|F|T|F|T|F|T|T|

|T|F|F|F|T|F|T|F|

|F|T|T|T|F|F|T|T|

|F|T|F|T|F|F|T|F|

|F|F|T|T|T|T|T|T|

|F|F|F|T|T|T|F|F|

Based on the truth table, we can see that whenever the premises are true, the conclusion is also true. There is no row in the truth table where the premises are true, but the conclusion is false. Therefore, the argument is valid.

Final Answer: A. 'Valid'

Question 83:
' Select the best translation into predicate logic. All birds live in some nest. (Bx: x is a bird; Nx: x is a nest; Lxy: x lives in y)', A. ' $(\forall x)(\forall y)[(Bx \bullet Ny) \supset Lxy]'$, B. ' $(\forall x)[Bx \supset (\exists y)(Ny \bullet Lxy)]'$, C. ' $(\forall x)[(\exists y)(Ny \bullet Bx) \supset (\exists y)(Ny \bullet Cxy)]'$
Lxy]', D. ' $(\forall x)(\exists y)[(Bx \bullet Ny) \supset Lxy]'$
Answer Option 1:
The best translation into predicate logic for the given statement would be: B. $'(\forall x)[Bx \supset (\exists y)(Ny \bullet Lxy)]'$ The statement "All birds live in some nest" can be translated as "For all x, if x is a bird, then there exists
a y such that y is a nest and x lives in y." This can be represented symbolically as $(\forall x)[Bx \supset (\exists y)(Ny \bullet Lxy)]$. Final Answer: B
Answer Option 2:
The best translation into predicate logic for the statement "All birds live in some nest" is:
Bx: x is a bird
Nx: x is a nest
Lxy: x lives in y
The correct translation is:
$(\forall x)[Bx\supset (\exists y)(Ny \cdot Lxy)]$
Option B is the correct translation into predicate logic. Final Answer: B

Answer Option 3:
The best translation into predicate logic for the given statement is:
B. $(\forall x)[Bx \supset (\exists y)(Ny \cdot Lxy)]$
Explanation:
- Bx: x is a bird
- Nx: x is a nest
- Lxy: x lives in y
The statement "All birds live in some nest" can be translated as "For all x, if x is a bird, then there exists a y such that y is a nest and x lives in y."
Symbolically, this can be represented as $(\forall x)[Bx \supset (\exists y)(Ny \bullet Lxy)]$. Therefore, option B is the best translation.
Answer Option 4:
The best translation into predicate logic for the statement "All birds live in some nest" is:
B. $(\forall x)[Bx \supset (\exists y)(Ny \cdot Lxy)]$
Explanation:
- $(\forall x)$ indicates that the following statement holds for all x .
- Bx represents "x is a bird."
- (∃y) indicates that there exists some y.
- Ny represents "y is a nest."
- Lxy represents "x lives in y."
By combining these predicates, the translation $(\forall x)[Bx \supset (\exists y)(Ny \bullet Lxy)]$ accurately represents the statement that all birds live in some nest. Final Answer: B
Answer Option 5:
The best translation into predicate logic for the statement "All birds live in some nest" would be: B. $(\forall x)[Bx \supset (\exists y)(Ny \bullet Lxy)]'$
Breaking it down:
Bx: x is a bird.
Nx: x is a nest.
Lxy: x lives in y.
The statement "All birds live in some nest" can be translated as: For every x (every bird), if x is a bird (Bx), then there exists a y (some nest) such that y is a nest (Ny) and x lives in y (Lxy).
Therefore, the correct translation is: $'(\forall x)[Bx \supset (\exists y)(Ny \bullet Lxy)]'$. Final Answer: B.
Answer Option 6:
The best translation into predicate logic for the statement "All birds live in some nest" is:
B. $'(\forall x)[Bx \supset (\exists y)(Ny \cdot Lxy)]'$
Explanation:
- Bx: x is a bird

- Nx: x is a nest

- Lxy: x lives in y

The statement "All birds live in some nest" can be translated as:

"For all x, if x is a bird, then there exists a y such that y is a nest and x lives in y."

In predicate logic notation, this can be represented as:

 $(\forall x)[\mathsf{B} x \supset (\exists y)(\mathsf{N} y \bullet \mathsf{L} xy)]$

This translation accurately captures the meaning of the original statement. Final Answer: B.

Answer Option 7:

The best translation into predicate logic for the statement "All birds live in some nest" is option B: $'(\forall x)[Bx \supset (\exists y)(Ny \bullet Lxy)]'$.

Let's break it down:

Bx: x is a bird

Nx: x is a nest

Lxy: x lives in y

The statement says that for all x (x represents any object), if x is a bird (Bx), then there exists some y (y represents any object) such that y is a nest (Ny) and x lives in y (Lxy).

This can be translated as '(\forall x)[Bx \to (\forall y)(Ny \cdot Lxy)]', where (\forall x') represents "for all x", Bx represents "x is a bird", (\forall y) represents "there exists some y", Ny represents "y is a nest", and Lxy represents "x lives

Therefore, option B is the best translation into predicate logic. Final Answer: B

Answer Option 8:

The best translation into predicate logic for the given statement is:

 $\mathsf{B.}\ '(\forall x)[\mathsf{Bx} \supset (\exists y)(\mathsf{Ny} \bullet \mathsf{Lxy})]'$

Explanation

The statement "All birds live in some nest" can be translated as "For all birds x, if x is a bird, then there exists a nest y such that y is a nest and x lives in y".

In predicate logic, this can be represented as:

 $(\forall x)[Bx\supset (\exists y)(Ny \bullet Lxy)]$

where Bx represents "x is a bird", Ny represents "y is a nest", and Lxy represents "x lives in y".

Option A: '(∀x)(∀y)[(Bx • Ny) ⊃ Lxy]' is incorrect because it states that all birds and all nests are related by the relation "x lives in y", which is not what the original statement implies.

Option C: '(∀x)[(∃y)(Ny • Bx) ⊃ Lxy]' is incorrect because it says that if there exists a nest y and x is a bird, then x lives in y, which is not the same as saying that all birds live in some nest.

Option D: $(\forall x)(\exists y)[(Bx \bullet Ny) \supset Lxy]'$ is incorrect because it says that for all birds x, there exists some nest y such that if x is a bird and y is a nest, then x lives in y. This is a stronger statement than the original statement, which only says that all birds live in some nest, not necessarily a specific nest.

Final Answer: B

Answer Option 9:

The best translation into predicate logic for the given sentence is:

 $\mathsf{B.}\ '(\forall x)[\mathsf{Bx}\supset (\exists y)(\mathsf{Ny}\bullet\mathsf{Lxy})]'$

Explanation:

Let's break down the sentence:

"All birds live in some nest."

- We can start by representing the statement "x is a bird" as Bx.
- Similarly, we can represent "x is a nest" as Nx.
- Finally, we can represent "x lives in y" as Lxy.

Now, let's analyze the options:

Option A: $(\forall x)(\forall y)[(Bx \cdot Ny) \supset Lxy]'$

This option states that for every x and y, if x is a bird and y is a nest, then x lives in y. However, this does not capture the idea that all birds live in some nest. It states that all birds live in all nests, which is not what the original sentence is saying.

Option B: $'(\forall x)[Bx \supset (\exists y)(Ny \bullet Lxy)]'$

This option states that for every x, if x is a bird, then there exists a y such that y is a nest and x lives in y. This accurately captures the idea that all birds live in some nest.

Option C: $'(\forall x)[(\exists y)(Ny \bullet Bx) \supset Lxy]'$

This option states that for every x, if there exists a y such that y is a nest and x is a bird, then x lives in y. This does not accurately capture the idea that all birds live in some nest. It allows for the possibility that there are nests where no birds live.

Option D: $'(\forall x)(\exists y)[(Bx \cdot Ny) \supset Lxy]'$

This option states that for every x, there exists a y such that if x is a bird and y is a nest, then x lives in y. This allows for the possibility that there are nests where no birds live, which is not what the original sentence is saying.

Based on the analysis above, the best translation into predicate logic is Option B: $(\forall x)[Bx \supset (\exists y)(Ny \bullet Lxy)]'$. Final Answer: B......

Ougation	04

'Identify the conclusion of the following argument. Utilitarian theory is embarrassed by the possibility of utility monsters. A utility monster gets enormous gains in utility, greater than the utility lost by the sacrifices of others. The sacrifices of even large numbers of other people are outweighed by the gains of the utility monster. The theory seems to require that we all be sacrificed in the monste's maw in order to increase total utility. (Robert Nozick, Anarchy, State, and Utopia)', A. 'Utilitarian theory is embarrassed by the possibility of utility monsters.', B. 'A utility monster gets enormous gains in utility.', C. 'A utility monster gets gains in utility greater than the utility lost by the sacrifices of others.', D. The sacrifices of even large numbers of other people are outweighed by the gains of the utility monster.'

Answer Option 1:

The conclusion of the argument is: D. 'The sacrifices of even large numbers of other people are outweighed by the gains of the utility monster.'

Answer Option 2:

The conclusion of the argument is: D. 'The sacrifices of even large numbers of other people are outweighed by the gains of the utility monster.'

Answer Option 3:

The conclusion of the argument is: D. 'The sacrifices of even large numbers of other people are outweighed by the gains of the utility monster.'

Answer Option 4:

The conclusion of the argument is: D. 'The sacrifices of even large numbers of other people are outweighed by the gains of the utility monster.'

Answer Option 5:

The conclusion of the argument is: D. 'The sacrifices of even large numbers of other people are outweighed by the gains of the utility monster.'

Answer Option 6:

The conclusion of the argument is: D. 'The sacrifices of even large numbers of other people are outweighed by the gains of the utility monster.' This is the main point or claim that the argument is making based on the information provided.

.....

Answer Option 7:

The conclusion of the argument is: D. The sacrifices of even large numbers of other people are outweighed by the gains of the utility monster.' This is the main point made in the argument, highlighting the issue with the utilitarian theory when faced with the concept of utility monsters.

Answer Option 8:

The conclusion of the argument is: D. 'The sacrifices of even large numbers of other people are outweighed by the gains of the utility monster.' This statement summarizes the main point made in the argument regarding the impact of utility monsters on utilitarian theory.

Answer Option 9:

The conclusion of the argument is: 'The theory seems to require that we all be sacrificed in the monster's maw in order to increase total utility.' Thus, the correct option is D. 'The sacrifices of even large numbers of other people are outweighed by the gains of the utility monster.'

Question 85:
Use the following key to translate the given formula of PL to natural, English sentences.
A: Marina reads a Percy Jackson book.
B: Izzy plays Minecraft.
C: Emily stops working.
D: Russell makes dinner.
E: Ashleigh stops by.
(A ∨ ~E) ≡ (D • ~B)', A. 'Either Marina reading a Percy Jackson book or Ashleigh not stopping by is a necessary condition for both Russell making dinner and Izzy not playing Minecraft.', B. 'Either Marina
reading a Percy Jackson book or Ashleigh not stopping by is a necessary and sufficient condition for Russell making dinner and Izzy playing Minecraft.', C. 'Marina reads a Percy Jackson book unless
Ashleigh stops by if, and only if, both Russell makes dinner and Izzy doesn't play Minecraft.', D. 'Either Marina reads a Percy Jackson book or Ashleigh does not stop by just in case both Russell makes dinner
and Izzy doesn't play Minecraft.'
Answer Option 1:
Using the given key:
A: Marina reads a Percy Jackson book.
B: Izzy plays Minecraft.
C: Emily stops working.
D: Russell makes dinner.
E: Ashleigh stops by.
The formula $(A \lor -E) \equiv (D \bullet -B)$ can be translated to natural English sentences as:
"Either Marina reading a Percy Jackson book or Ashleigh not stopping by is a necessary and sufficient condition for both Russell making dinner and Izzy not playing Minecraft."
Final Answer: B

Answer Option 2:
Let's break down the given formula and translate it using the key:
$(A \vee \sim E) \equiv (D \bullet \sim B)$
Translation:
Either Marina reads a Percy Jackson book or Ashleigh does not stop by if and only if both Russell makes dinner and Izzy does not play Minecraft.
The correct option that matches the translation is D: 'Either Marina reads a Percy Jackson book or Ashleigh does not stop by just in case both Russell makes dinner and Izzy doesn't play Minecraft.'

Answer Option 3:
The given formula is $(A \lor \neg E) \equiv (D \bullet \neg B)$.
Using the key:
A: Marina reads a Percy Jackson book.
B: Izzy plays Minecraft.
C: Emily stops working.
D: Russell makes dinner.
E: Ashleigh stops by.
We can translate the formula as:
$(A \vee {\sim} E) \equiv (D \bullet {\sim} B)$
Translation:
Either Marina reads a Percy Jackson book or Ashleigh does not stop by if and only if both Russell makes dinner and Izzy does not play Minecraft.
The correct option that represents the translation is:
D. 'Either Marina reads a Percy Jackson book or Ashleigh does not stop by just in case both Russell makes dinner and Izzy doesn't play Minecraft.'
Answer Option 4:
To translate the given formula $(A \lor \sim E) \equiv (D \bullet \neg B)$ into natural, English sentences, we can use the following key:
A: Marina reads a Percy Jackson book.
B: Izzy plays Minecraft.
C: Emily stops working.
D: Russell makes dinner.
E: Ashleigh stops by.
$(A \vee \neg E) \equiv (D \cdot \neg B)$ can be translated as:

"Either Marina reads a Percy Jackson book or Ashleigh does not stop by if and only if both Russell makes dinner and Izzy does not play Minecraft." The correct option that reflects this translation is: C. 'Marina reads a Percy Jackson book unless Ashleigh stops by if, and only if, both Russell makes dinner and Izzy doesn't play Minecraft.' Answer Option 5: To translate the given formula $(A \lor \sim E) \equiv (D \bullet \sim B)$ to natural, English sentences, we can use the following key: A: Marina reads a Percy Jackson book. B: Izzy plays Minecraft. C: Emily stops working. D: Russell makes dinner. E: Ashleigh stops by. $(\mathsf{A} \vee \mathsf{\sim} \mathsf{E}) \equiv (\mathsf{D} \bullet \mathsf{\sim} \mathsf{B})$ Translation: "Either Marina reads a Percy Jackson book or Ashleigh does not stop by" is equivalent to "Russell makes dinner and Izzy does not play Minecraft." The correct translation is: D. "Either Marina reads a Percy Jackson book or Ashleigh does not stop by just in case both Russell makes dinner and Izzy doesn't play Minecraft." Answer Option 6: Given PL formula: $(A \lor \sim E) \equiv (D \bullet \sim B)$ Translation using the key: A: Marina reads a Percy Jackson book. B: Izzy plays Minecraft. C: Emily stops working. D: Russell makes dinner. E: Ashleigh stops by. Translation of $(A \lor \sim E) \equiv (D \cdot \sim B)$: "Either Marina reads a Percy Jackson book or Ashleigh does not stop by, if and only if, both Russell makes dinner and Izzy does not play Minecraft." Therefore, the correct option is D. "Either Marina reads a Percy Jackson book or Ashleigh does not stop by just in case both Russell makes dinner and Izzy doesn't play Minecraft." Answer Option 7: The given formula is $(A \lor \sim E) \equiv (D \bullet \sim B)$. Let's break it down step by step: (A v ~E): Marina reads a Percy Jackson book or Ashleigh does not stop by. (D • ~B): Russell makes dinner and Izzy does not play Minecraft. Now, let's translate each part into English sentences: (A \vee ~E): Either Marina reads a Percy Jackson book or Ashleigh does not stop by. (D • ~B): Both Russell makes dinner and Izzy does not play Minecraft. Lastly, let's put them together and simplify the formula: (A v ~E) = (D • ~B): Either Marina reading a Percy Jackson book or Ashleigh not stopping by is a necessary and sufficient condition for both Russell making dinner and Izzy not playing Minecraft. So, the correct translation is: B. 'Either Marina reading a Percy Jackson book or Ashleigh not stopping by is a necessary and sufficient condition for Russell making dinner and Izzy playing Minecraft.' Answer Option 8: In order to translate the given formula (A $\lor \lor$ E) \equiv (D $\bullet \lor$ AB) into natural English sentences, we can break it down into smaller parts. A: Marina reads a Percy Jackson book. B: Izzy plays Minecraft. C: Emily stops working. D: Russell makes dinner E: Ashleigh stops by. $(\mathsf{A} \vee \mathsf{\sim} \mathsf{E}) \equiv (\mathsf{D} \bullet \mathsf{\sim} \mathsf{B})$ Let's break it down: (A \vee ~E): Either Marina reads a Percy Jackson book or Ashleigh does not stop by. (D • ~B): Both Russell makes dinner and Izzy does not play Minecraft.

Now let's combine the two parts using the equivalence operator \equiv :

(Either Marina reads a Percy Jackson book or Ashleigh does not stop by) is true if and only if (Both Russell makes dinner and Izzy does not play Minecraft) is true.

So the correct translation is:

'Either Marina reads a Percy Jackson book or Ashleigh does not stop by is a necessary and sufficient condition for Russell making dinner and Izzy not playing Minecraft.'

Therefore, the correct answer is B. 'Either Marina reading a Percy Jackson book or Ashleigh not stopping by is a necessary and sufficient condition for Russell making dinner and Izzy playing Minecraft.'

Answer Option 9:

To translate the given formula $(A \lor \sim E) \equiv (D \bullet \sim B)$ using the given key:

A: Marina reads a Percy Jackson book.

B: Izzy plays Minecraft.

C: Emily stops working.

D: Russell makes dinner.

E: Ashleigh stops by.

The formula $(A \lor \sim E) \equiv (D \bullet \sim B)$ can be translated as follows:

Option A: Either Marina reading a Percy Jackson book or Ashleigh not stopping by is a necessary condition for both Russell making dinner and Izzy not playing Minecraft.

Option B: 'Either Marina reading a Percy Jackson book or Ashleigh not stopping by is a necessary and sufficient condition for Russell making dinner and Izzy playing Minecraft.'

Option C: 'Marina reads a Percy Jackson book unless Ashleigh stops by if, and only if, both Russell makes dinner and Izzy doesn't play Minecraft.'

Option D: 'Either Marina reads a Percy Jackson book or Ashleigh does not stop by just in case both Russell makes dinner and Izzy doesn't play Minecraft.'

The correct translation based on the given formula is Option D: 'Either Marina reads a Percy Jackson book or Ashleigh does not stop by just in case both Russell makes dinner and Izzy doesn't play

Minecraft.'

Question 86:

'Which of the given formulas of PL is the best symbolization of the following sentence?

Brittany likes fish and lizards, but not cats.', A. 'F • L • ~C', B. '(F • L) • ~C', C. '(F • L) \lor ~C', D. 'C • (F • L)'

Answer Option 1:

The best symbolization of the sentence "Brittany likes fish and lizards, but not cats" would be option B, '(F • L) • ~C'. This formula accurately represents the statement that Brittany likes both fish and lizards (F

• L), and also represents that she does not like cats (~C), by using the logical operator for conjunction (•). This option correctly captures the meaning of the original sentence. Final Answer: B

Answer Option 2:

The best symbolization for the given sentence, "Brittany likes fish and lizards, but not cats," would be option B: '(F • L) • ~C'.

This symbolization accurately represents that Brittany likes both fish and lizards (F • L) and also that she does not like cats (-C), and combines these two conditions with the conjunction operator (•).

Therefore, option B correctly captures the intended meaning of the sentence. Final Answer: B

Answer Option 3:

The best symbolization of the sentence "Brittany likes fish and lizards, but not cats" would be B. '(F • L) • ~C'.

This is because the sentence states that Brittany likes both fish and lizards (F • L) and also explicitly mentions that she does not like cats (~C). The conjunction operator (•) is used to represent the "and" between fish and lizards, and the negation (~) is used to represent the "not" in relation to cats.

Therefore, option B is the most accurate symbolization of the given sentence. Final Answer: B. '(F • L) • ~C'

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Answer Option 4:

The best symbolization of the sentence "Brittany likes fish and lizards, but not cats" would be option B: '(F • L) • ~C'.

This is because the sentence states that Brittany likes fish and lizards, which can be represented by 'F • L'. The '•' symbol represents the logical operator for conjunction, which means "and". The sentence also states that Brittany does not like cats, which can be represented by '-C'. The '-' symbol represents the logical operator for negation, which means "not".

Therefore, combining 'F • L' and '~C' with another conjunction operator '•' gives us '(F • L) • ~C', which accurately symbolizes the sentence "Brittany likes fish and lizards, but not cats".

Final Answer: B

Answer Option 5:

The best symbolization of the given sentence, "Brittany likes fish and lizards, but not cats", is option B: '(F • L) • ~C'.

Explanation

Let's break down the sentence into its components:

- Brittany likes fish: F
- Brittany likes lizards: L
- Brittany does not like cats: ~C

The sentence states that Brittany likes both fish and lizards, which can be represented as (F • L). It also states that she does not like cats, which can be represented as ~C.

In order for the sentence to be true, both conditions must hold. Therefore, the best representation would be '(F • L) • ~C', which means "Brittany likes fish and lizards, but not cats". Final Answer: B

Answer Option 6:

The best symbolization of the given sentence would be B. '(F • L) • ~C'.

In the sentence, we have two statements "Brittany likes fish" and "Brittany likes lizards", which are represented as F and L respectively. The statement "but not cats" is represented by the negation of liking cats, which is ~C.

Using the logical connectives, the conjunction operator (•) is used to connect the statements "Brittany likes fish" and "Brittany likes lizards" together, as she likes both fish and lizards. Then, the conjunction operator (•) is used again to connect the overall conjunction of fish and lizards with the negation of liking cats, as she does not like cats.

Thus, the best symbolization is '(F • L) • ~C'. Final Answer: B.

Answer Option 7:

The best symbolization of the sentence 'Brittany likes fish and lizards, but not cats' would be option B, '(F \bullet L) \bullet \sim C'.

Explanation:

- "F" represents the statement "Brittany likes fish".
- "L" represents the statement "Brittany likes lizards".
- "C" represents the statement "Brittany likes cats".

The conjunction symbol "•" is used to represent "and". Therefore, '(F • L)' represents the statement "Brittany likes fish and lizards".

The negation symbol "~" is used to represent "not". Therefore, '~C' represents the statement "Brittany does not like cats".

Finally, the overall statement '(F \bullet L) \bullet ~C' represents "Brittany likes fish and lizards, but not cats".

Answer Option 8:

The best symbolization of the given sentence, "Brittany likes fish and lizards, but not cats" would be option B: (F • L) • ~C.

This symbolization correctly represents the conjunction of "Brittany likes fish" (F) and "Brittany likes lizards" (L), and then further conjoins it with the negation of "Brittany likes cats" (~C), which captures the meaning of "but not cats"

Option A: F • L • ~C would incorrectly imply that Brittany likes fish, lizards, and does not like cats simultaneously, which is not accurate.

Option C: (F • L) \vee ~C would mean that Brittany either likes both fish and lizards or does not like cats, which is not the intended meaning of the sentence.

 $Option \ D: \ C \bullet (F \bullet L) \ would \ mean \ that \ Brittany \ likes \ cats, \ fish, \ and \ lizards \ simultaneously, \ which \ is \ also \ not \ accurate.$

Final Answer: B

Answer Option 9:

The best symbolization of the sentence "Brittany likes fish and lizards, but not cats" would be option B: '(F \bullet L) \bullet ~C'.

This symbolization accurately represents that Brittany likes both fish and lizards (F • L), and also states that she does not like cats (-C). The conjunction operator (•) is used to represent "and" in this context, and the negation (-) is used to indicate "not".

Option A: 'F • L • ~C' suggests that liking fish, lizards, and not liking cats are all separate conditions, which is not what the sentence conveys.

Option C: '(F • L) v ~ C' uses the disjunction operator (v), which represents "or", indicating that Brittany either likes fish and lizards or does not like cats, which is not an accurate representation of the sentence.

Option D: 'C • (F • L)' suggests that liking cats is a condition that must be met in order for Brittany to like fish and lizards, which is not what the sentence conveys.

Therefore, option B is the best symbolization of the given sentence. Final Answer: B.

Question 87: 'Which of the given formulas of PL is the best symbolization of the following sentence? If either Harry enrolls or George does not enroll, then Jim enrolls.', A. 'H \vee (\sim G \supset J)', B. ' \sim (H \vee G) \supset J', C. '(H \vee \sim G) \supset J', D. '(H \supset \sim G) \vee J' = 1.5 (H \vee \sim G) \sim G) \vee J' = 1.5 (H \vee \sim G) \sim Answer Option 1: The given sentence can be symbolized as follows: "Either Harry enrolls or George does not enroll" can be expressed as H \vee ~G. "Jim enrolls" can be expressed as J. Putting it all together, the best symbolization of the sentence would be: $(H \lor \sim G) \rightarrow J.$ Therefore, the correct answer is option C: $'(H \vee {\sim} G) \supset J'$. ******* Answer Option 2: The best symbolization of the given sentence "If either Harry enrolls or George does not enroll, then Jim enrolls" is option C: "(H $\vee \sim$ G) \supset J". Here is the breakdown: - "Either Harry enrolls or George does not enroll" can be represented as "H v ~G" where H represents "Harry enrolls" and ~G represents "George does not enroll". - "Then Jim enrolls" can be represented as "J". - Combining the two parts with an implication, we get "(H \vee ~G) \supset J". Therefore, the correct answer is C: "(H \vee ~G) \supset J". Answer Option 3: The given sentence can be symbolized as: "If either Harry enrolls or George does not enroll, then Jim enrolls." Let's break it down: - Harry enrolls: H - George does not enroll: ~G - Jim enrolls: J Now, let's look at the options: Option A: $'H \lor (\sim G \supset J)'$ This option states that if George does not enroll, then Jim enrolls. However, it does not cover the case where Harry enrolls. This option states that if neither Harry enrolls nor George enrolls, then Jim enrolls. It covers the case where George does not enroll but does not cover the case where Harry enrolls. Option C: '(H \vee ~G) \supset J' This option states that if either Harry enrolls or George does not enroll, then Jim enrolls. It accurately symbolizes the given sentence. Option D: $'(H \supset \sim G) \lor J'$ This option states that if Harry enrolls implies that George does not enroll or Jim enrolls. It does not cover the case where George does not enroll without Harry enrolling. Therefore, the best symbolization of the given sentence is option C: '(H \vee ~G) \supset J'. Final Answer: C Answer Option 4: To symbolize the given sentence, we can break it down into the following parts: Harry enrolls: H George does not enroll: ~G Jim enrolls: J

The sentence "If either Harry enrolls or George does not enroll, then Jim enrolls" can be symbolized as "(H \vee ~G) \rightarrow J".

Comparing this symbolization with the given options:

A. 'H \vee (\sim G \supset J)': This option is not correct as it does not capture the logical structure of the original sentence.

 $B. \ \ ^{\prime} \sim (H \lor G) \supset J' : This \ option \ is \ not \ correct as \ it \ does \ not \ correctly \ represent the \ condition "either Harry enrolls or George does \ not \ enroll".$

C. '(H \vee ~G) \supset J': This option is correct as it accurately symbolizes the given sentence.

 $D.\ '(H\supset \sim\! G)\lor J':\ This\ option\ is\ not\ correct\ as\ it\ does\ not\ correctly\ represent\ the\ "or"\ condition\ in\ the\ original\ sentence.$

Therefore, the best symbolization of the given sentence is option C: $'(H \lor \sim G) \supset J'$. Final Answer: C

Answer Option 5:

To symbolize the given sentence, let's break it down into parts:

- 1. Harry enrolls: H
- 2. George does not enroll: ~G

3. Jim enrolls: J The sentence "If either Harry enrolls or George does not enroll, then Jim enrolls" can be symbolized as " $(H \lor \sim G) \to J$ ". Comparing this symbolization with the given options: A. 'H v (~G > J)': This option does not represent the given sentence accurately because it says "either Harry enrolls or if George does not enroll, then Jim enrolls". $B. \ \ '-(H\vee G)\supset J': This \ option \ says \ "If \ neither \ Harry \ enrolls \ nor \ George \ enrolls, \ then \ Jim \ enrolls" \ which is not the \ same \ as the \ given \ sentence.$ C. '(H ∨ ~G) ⊃ J': This is the correct symbolization of the given sentence as explained above. D. $'(H \supset \neg G) \lor J'$: This option says "If Harry enrolls implies George does not enroll or Jim enrolls" which is not the same as the given sentence. Therefore, the best symbolization of the given sentence is option C: '(H \vee ~G) \supset J'. ****************** Answer Option 6: Let's break down the sentence into logical symbols: H: Harry enrolls G: George enrolls The sentence, "If either Harry enrolls or George does not enroll, then Jim enrolls," can be symbolized as "(H \vee ~G) \rightarrow J." Comparing this with the given options: A. 'H v (~G \supset J)': This option does not correctly capture the condition "either Harry enrolls or George does not enroll." It only specifies the case where George enrolls. B. '~(H v G) \supset J': This option incorrectly negates the condition "either Harry enrolls or George does not enroll." It states the negation of "Harry enrolls or George enrolls" instead of "either Harry enrolls or George does not enroll." C. '(H v ~G) \supset J': This option correctly captures the condition "either Harry enrolls or George does not enroll." It implies that if either Harry enrolls or George does not enroll, then Jim enrolls. D. '(H $\supset \sim$ G) \vee J': This option does not correctly capture the condition "either Harry enrolls or George does not enroll." It implies that if Harry enrolls, then George does not enroll. Therefore, the best symbolization of the given sentence is option C: '(H \vee ~G) \supset J'. ******* Answer Option 7: To symbolize the given sentence, we can break it down into its components: 1. Harry enrolls: H 2. George does not enroll: ~G 3. Jim enrolls: J The given sentence states "If either Harry enrolls or George does not enroll, then Jim enrolls." This can be symbolized as " $(H \lor \sim G) \to J$." Let's compare this symbolization with the options given: Option A: 'H v (-G \supset J)' - This option does not correctly capture the condition that either Harry enrolls or George does not enroll. It only captures the case when George does not enroll and Jim enrolls. Option B: '~(H v G) \supset J' - This option does not include the condition that either Harry enrolls or George does not enroll. It only captures the case when neither Harry nor George enrolls and Jim enrolls. Incorrect. Option C: '(H v ~G) \supset J' - This option correctly symbolizes the given sentence. It captures the condition that either Harry enrolls or George does not enroll, and then Jim enrolls. This is the best symbolization Option D: '(H \supset ~G) \vee J' - This option does not correctly capture the condition that either Harry enrolls or George does not enroll. It only captures the case when Harry enrolls and George does not enroll, or Jim enrolls. Incorrect Final Answer: C ******** Answer Option 8: The given sentence can be symbolized as follows: If either Harry enrolls or George does not enroll, then Jim enrolls. Let's break down the sentence into its logical components: Harry enrolls: H George does not enroll: ~G Jim enrolls: J

This symbolization means "Harry enrolls or if George does not enroll, then Jim enrolls." It does not capture the condition "if either Harry enrolls or George does not enroll." Therefore, it is not the best

Now, let's analyze the answer choices:

Option A: $H \lor (\sim G \supset J)$

Option B: \sim (H \vee G) \supset J

This symbolization means "If it is not the case that either Harry or George enrolls, then Jim enrolls." It does not accurately capture the condition "if either Harry enrolls or George does not enroll." Therefore, it is not the best symbolization.

Option C: $(H \lor \sim G) \supset J$

This symbolization means "If either Harry enrolls or George does not enroll, then Jim enrolls." It accurately captures the condition given in the sentence. Therefore, it is a good symbolization.

Option D: $(H \supset \sim G) \lor J$

This symbolization means "If Harry enrolls, then George does not enroll, or Jim enrolls." It does not accurately capture the condition "if either Harry enrolls or George does not enroll." Therefore, it is not the

Based on the analysis, the best symbolization of the given sentence is Option C: (H \vee ~G) \supset J. Final Answer: C.

Answer Option 9:

The statement "If either Harry enrolls or George does not enroll, then Jim enrolls" can be symbolized as follows:

Let H represent "Harry enrolls", G represent "George enrolls", and J represent "Jim enrolls".

The statement can be broken down into two parts:

- 1. Either Harry enrolls or George does not enroll: $H \vee {\scriptstyle \sim} G$
- 2. If the above condition is true, then Jim enrolls: (H \vee ~G) \to J

Comparing the given options:

Option A: $H \lor (\sim G \supset J)'$

This option does not accurately represent the condition "either Harry enrolls or George does not enroll". The symbolization (-G \supset J) implies that if George does not enroll, then Jim enrolls, which is not what the sentence states.

Option B: '~(H \vee G) \supset J'

This option accurately represents the condition "either Harry enrolls or George does not enroll" as ~(H ∨ G), which means "neither Harry nor George enrolls". The symbolization (~H ∨ ~G) ⊃ J represents the correct implication that if neither Harry nor George enrolls, then Jim enrolls.

Option C: '(H \vee ~G) \supset J'

This option accurately represents the condition "either Harry enrolls or George does not enroll" as (H v ~G). However, it incorrectly implies that if the given condition is true, then Jim enrolls, which is not what the sentence states.

Option D: $'(H\supset \sim G)\vee J'$

This option does not accurately represent the condition "either Harry enrolls or George does not enroll". The symbolization (H \supset ~G) implies that if Harry enrolls, then George does not enroll, which is not what the sentence states.

Therefore, the best symbolization of the given sentence is Option B: '~(H \vee G) \supset J'.

	nn 88	

'Construct a complete truth table for the following argument. Then, using the truth table, determine whether the argument is valid or invalid. If the argument is invalid, choose an option which presents a counterexample. (There may be other counterexamples as well.)

 $\sim X \supset Y$

 $Y \supset Z\,$

~Z / ~X', A. "Valid', B. 'Invalid. Counterexample when X, Y, and Z are true', C. 'Invalid. Counterexample when Y and Z are true and Z is false', D. 'Invalid. Counterexample when X is true and Y and Z are false'

Answer Option 1:

To construct a complete truth table for the argument, we need to consider all possible combinations of truth values for the propositions X, Y, and Z.

 $X | Y | Z | \sim X \supset Y | Y \supset Z | \sim Z | \sim X |$

T|T|T|T|F|F|

TITIFITIFITIFI

T|F|T|F|T|F|F|

TIFIFIFITITIFI

F|T|T|T|T|F|T|

FITIFITIFITITI

F|F|T|T|T|F|T|

F|F|F|T|T|T|T|

The given argument is: ~X \supset Y, Y \supset Z, ~Z / ~X

We can see that in all the rows where $\sim X$ is true, $\sim X \supset Y$ is true. In all the rows where Y is true. And in all the rows where Z is false, $\sim Z$ is true.

Therefore, in every row where ~X is true, Y is true, and ~Z is true, ~X is true as well. Hence, the argument is valid.

Final Answer: A. Valid

Answer Option 2:

To construct the truth table, we consider all possible combinations of truth values for X, Y, and Z.

 $X\mid Y\mid Z\mid \sim X\mid \sim X\supset Y\mid Y\supset Z\mid \sim Z\mid \sim X$

T|T|T|F|T|T|F|F

T|T|F|F|T|F|T|F

T|F|T|F|F|T|F|F

T|F|F|F|F|T|T|F

F|T|T|T|T|F|T

F|T|F|T|T|F|T|T

F|F|T|T|T|F|T

F|F|F|T|T|T|T

The given argument is:

 ${\sim} X \supset Y$

 $Y \supset Z$

....×

From the truth table, we can see that whenever all the premises are true, the conclusion is also true. Therefore, the argument is valid.

Final Answer: A. 'Valid'

Answer Option 3:

To construct a complete truth table for the argument, we need to consider all possible truth values for the propositions involved: ~X, Y, Z.

~X ⊃ Y

 $Y\supset Z$

~Z

We have four columns for ~X, Y, Z, and ~X. We will evaluate the truth values for each row by substituting different combinations of truth values for ~X, Y, Z.

~X | Y | Z | ~X

TITITIF TITIFIF TIFITIF $\mathsf{T} | \mathsf{F} | \mathsf{F} | \mathsf{F}$ FITITIT F|T|F|T F|F|T|T F|F|F|T Now, we evaluate the truth values for ~X \supset Y, Y \supset Z, and ~Z. $\sim X \supset Y \mid Y \supset Z \mid \sim Z \mid \sim X$ T|T|F|F TITIFIF TITIFIF TITIFIF TITITIT TIFITIT $\mathsf{T} \, | \, \mathsf{T} \, | \, \mathsf{T} \, | \, \mathsf{T}$ TITITIT From the truth table, we can see that the conclusion ~X is not always true when the premises ~X ⊃ Y, Y ⊃ Z, and ~Z are all true. In the first four rows, when ~X is true, the conclusion ~X is false. This means that there is a counterexample where the premises are true, but the conclusion is false. Thus, the argument is invalid. Therefore, the correct answer is B. 'Invalid. Counterexample when X, Y, and Z are true'. ********** Answer Option 4: To construct a truth table for the argument, we need to consider all possible combinations of truth values for the propositions X, Y, and Z. $X \mid Y \mid Z \mid \sim X \mid \sim X \supset Y \mid Y \supset Z \mid \sim Z$ TITITIFITIF TITIFIFITIFIT T|F|T|F|F|T|F T|F|F|F|F|T|T FITITITITIF F|T|F|T|T|F|T F|F|T|T|T|F F|F|F|T|T|T Based on the truth table, we can determine the truth values of \sim X (not X), \sim X \supset Y (if not X, then Y), Y \supset Z (if Y, then Z), and \sim Z (not Z) for each row. In the last column, we see that there is a row where ~X is true and ~X \supset Y, Y \supset Z, and ~Z are all true. This means that the premises of the argument are true in that row, but the conclusion ~X is false. Therefore, the argument is invalid. Option B is correct: Invalid. Counterexample when X, Y, and Z are true. Answer Option 5: To construct a complete truth table for the argument, we need to consider all possible truth values for the propositions X, Y, and Z. Given premises: $1.~{\sim}X\supset Y$ $2.\ Y\supset Z$ We will evaluate the conclusion ~X based on these premises. Truth table: $| X | Y | Z | \sim X \supset Y | Y \supset Z | \sim Z | \sim X |$ |---|---|-----|-----| |T|T|T|T|F|F|

|T|T|F|T|F|T|F|

|T|F|T|F|T|F|F| |T|F|F|T|T|T|F| |F|T|T|T|T|F|T| |F|T|F|T|F|T|T| |F|F|T|T|T|F|T| |F|F|F|T|T|T|T| From the truth table, we can see that whenever ~X is false (F), the conclusion ~X is false (F) as well. Therefore, the argument is valid. Final Answer: A. Valid ****************** Answer Option 6: To construct a complete truth table for the argument, we need to consider all possible combinations of truth values for the variables X, Y, and Z. X | Y | Z | ~X | ~X \supset Y | Y \supset Z | ~Z | ~X (Given) | Conclusion T|T|T|F|T|T|F|F|Invalid T|T|F|F|T|F|T|F|Invalid T|F|T|F|F|T|F|F|Invalid T|F|F|F|F|T|T|F|Invalid F|T|T|T|T|F|T|Invalid F|T|F|T|T|F|T|T|Invalid F|F|T|T|T|F|T|Invalid F|F|F|T|T|T|T|T|ValidFrom the truth table, we can see that there is no row where all the premises are true and the conclusion is false. Therefore, the argument is invalid. To find a counterexample, we look for a row where all the premises are true and the conclusion is false. From the truth table, we can see that when X is true, Y is false, and Z is false, all the premises are true, but the conclusion is false. Counterexample: X is true, Y is false, and Z is false. (Option D) Final Answer: B. Invalid. Counterexample when X, Y, and Z are true Answer Option 7: To construct a complete truth table for the argument, we need to consider all possible truth values for the propositions involved. Let's start by filling in the truth values for -X, Y, and Z: $\mid X\mid Y\mid Z\mid \neg X\mid \neg X\supset Y\mid Y\supset Z\mid \neg Z\mid \neg X \text{ (Conclusion)}\mid$ |----|----|----|-----|-----|-----|-----|T|T|T|F|T|T|F|F| |T|T|F|F|T|F|T|F| |T|F|T|F|F|T|F|F| |T|F|F|F|F|T|T|F| |F|T|T|T|T|T|F|T||F|T|F|T|F|F|T|T| |F|F|T|T|T|T|F|T| |F|F|F|T|T|T|T|T Now, let's analyze the truth table to determine whether the argument is valid or invalid. For the argument to be valid, every row where all the premises are true should also have the conclusion as true. Looking at the truth table, we can see that there is a row where the premises $(\neg X \supset Y \text{ and } Y \supset Z)$ are true, but the conclusion $(\neg X)$ is false. Specifically, when X is false, Y is true, and Z is true, the premises are true, but the conclusion is false. Therefore, the argument is invalid. The correct answer is option D: "Invalid. Counterexample when X is true and Y and Z are false". Answer Option 8: To construct a complete truth table for the argument, we need to consider all possible truth values for the propositions involved: X, Y, and Z. Let's analyze the argument statement by statement: 1. $\sim X \supset Y$ 2. Y ⊃ Z 3. ~Z

Conclusion: ~X
Here's the truth table:

The argument is valid if the conclusion is true in every row where all the premises are true. In this case, the conclusion ~X is true in all rows where ~X $\supset Y$, $Y \supset Z$, and ~Z are true.

Looking at the truth table, we can see that there is no row where $\sim X \supset Y, Y \supset Z$, and $\sim Z$ are all true, but the conclusion $\sim X$ is false. Therefore, the argument is invalid.

Counterexample: Option D - Invalid. Counterexample when X is true and Y and Z are false. In the row where X is true, Y is false, Z is false, and -X, -X \(\supset\) Y, Y \(\supset\) Z, and -Z are all true, the conclusion -X is false.

Final Answer: D. Invalid. Counterexample when X is true and Y and Z are false.

Answer Option 9:

To construct a complete truth table for the argument, we need to list all possible combinations of truth values for the propositions involved. In this case, we have three propositions: X, Y, and Z.

The truth table would look as follows:

In the last column (~X), we have the negation of X.

Now, let's analyze the truth table to determine whether the argument is valid or invalid.

We can see that whenever -X is true (-X = T), $-X \supset Y$ is always true $(-X \supset Y = T)$. Similarly, whenever Y is true (Y = T), $Y \supset Z$ is always true $(Y \supset Z = T)$.

However, we can also see that in the rows where ~Z is true (~Z = T), ~X is false (~X = F). This means that ~X ⊃ Y is false (~X ⊃ Y = F), which contradicts the first proposition.

Therefore, the argument is invalid.

To find a counterexample, we need to find a row where all premises are true but the conclusion is false. In this case, the counterexample is when X is true, Y is true, and Z is false. In this row, ~X is false, ~X ⊃

 $Y \ \text{is false, } Y \supset Z \ \text{is true, but } \sim Z \ \text{is true. This satisfies all premises but contradicts the conclusion } \sim X.$

Thus, the correct answer is: C. 'Invalid. Counterexample when Y and Z are true and Z is false'.

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'S Select the best translation into predicate logic. Some kind people help some endangered species. (Ex: x is endangered; Kx: x is kind; Px: x is a person; Sx: x is a species; Hxy: x helps y)', A. '(∃x){(Kx • Px)}

 $\bullet \ (\exists y)[[(Ey \bullet Sy) \bullet Hxy]]', \ B. \ (\exists x)[\exists y)[[(Kx \bullet Sx) \bullet (Ey \bullet Py)] \bullet Hxy]', \ C. \ (\exists x)(\exists y)[[(Kx \bullet Sx) \bullet (Ey \bullet Py)] \bullet Hxy]', \ D. \ (\exists x)[\exists y)[[(Kx \bullet Sx) \bullet (Ey \bullet Py)] \bullet Hxy]', \ D. \ (\exists x)[\exists y)[[(Kx \bullet Sx) \bullet (Ey \bullet Py)] \bullet Hxy]', \ D. \ (\exists x)[\exists y)[[(Kx \bullet Sx) \bullet (Ey \bullet Py)] \bullet Hxy]', \ D. \ (\exists x)[\exists y)[[(Kx \bullet Sx) \bullet (Ey \bullet Py)] \bullet Hxy]', \ D. \ (\exists x)[\exists y)[[(Kx \bullet Sx) \bullet (Ey \bullet Py)] \bullet Hxy]', \ D. \ (\exists x)[(Ey \bullet Px) \bullet (Ey \bullet Px)] \bullet Hxy]', \ D. \ (\exists x)[(Ey \bullet Px) \bullet (Ey \bullet Px)] \bullet Hxy]', \ D. \ (\exists x)[(Ey \bullet Px) \bullet (Ey \bullet Px)] \bullet Hxy]', \ D. \ (\exists x)[(Ey \bullet Px) \bullet (Ey \bullet Px)] \bullet Hxy]', \ D. \ (\exists x)[(Ey \bullet Px) \bullet (Ey \bullet Px)] \bullet Hxy]', \ D. \ (\exists x)[(Ey \bullet Px) \bullet (Ey \bullet Px)] \bullet Hxy]', \ D. \ (\exists x)[(Ey \bullet Px) \bullet (Ey \bullet Px)] \bullet Hxy]', \ D. \ (\exists x)[(Ey \bullet Px) \bullet (Ey \bullet Px)] \bullet Hxy]', \ D. \ (\exists x)[(Ey \bullet Px) \bullet (Ey \bullet Px)] \bullet Hxy]', \ D. \ (\exists x)[(Ey \bullet Px) \bullet (Ey \bullet Px)] \bullet Hxy]', \ D. \ (\exists x)[(Ey \bullet Px) \bullet (Ey \bullet Px)] \bullet Hxy]', \ D. \ (\exists x)[(Ey \bullet Px) \bullet (Ey \bullet Px)] \bullet Hxy]', \ D. \ (\exists x)[(Ey \bullet Px) \bullet (Ey \bullet Px)] \bullet Hxy]', \ D. \ (\exists x)[(Ey \bullet Px) \bullet (Ey \bullet Px)] \bullet Hxy]', \ D. \ (\exists x)[(Ey \bullet Px) \bullet (Ey \bullet Px)] \bullet Hxy]', \ D. \ (\exists x)[(Ey \bullet Px) \bullet (Ey \bullet Px)] \bullet Hxy]', \ D. \ (\exists x)[(Ey \bullet Px) \bullet (Ey \bullet Px)] \bullet Hxy]', \ D. \ (\exists x)[(Ey \bullet Px) \bullet (Ey \bullet Px)] \bullet Hxy]', \ D. \ (\exists x)[(Ey \bullet Px) \bullet (Ey \bullet Px)] \bullet Hxy]', \ D. \ (\exists x)[(Ey \bullet Px) \bullet (Ey \bullet Px)] \bullet Hxy]', \ D. \ (\exists x)[(Ey \bullet Px) \bullet (Ey \bullet Px)] \bullet Hxy]', \ D. \ (\exists x)[(Ey \bullet Px) \bullet (Ey \bullet Px)] \bullet Hxy]', \ D. \ (\exists x)[(Ey \bullet Px) \bullet (Ey \bullet Px)] \bullet Hxy]', \ D. \ (\exists x)[(Ey \bullet Px) \bullet (Ey \bullet Px)] \bullet Hxy]', \ D. \ (\exists x)[(Ey \bullet Px) \bullet (Ey \bullet Px)] \bullet Hxy]', \ D. \ (\exists x)[(Ey \bullet Px) \bullet (Ey \bullet Px)] \bullet Hxy]', \ D. \ (\exists x)[(Ey \bullet Px) \bullet (Ey \bullet Px)] \bullet Hxy]', \ D. \ (\exists x)[(Ey \bullet Px) \bullet (Ey \bullet Px)] \bullet Hxy]', \ D. \ (\exists x)[(Ey \bullet Px) \bullet (Ey \bullet Px)] \bullet Hxy]', \ D. \ (\exists x)[(Ey \bullet Px) \bullet (Ey \bullet Px)] \bullet Hxy]', \ D. \ (\exists x)[(Ey \bullet Px) \bullet (Ey \bullet Px)] \bullet Hxy]', \ D. \ (\exists x)[(Ey \bullet Px) \bullet (Ey \bullet Px)] \bullet Hxy]', \ D. \ (\exists x)[(Ey \bullet Px) \bullet (Ey \bullet Px)] \bullet Hxy]', \ D. \ (\exists x)[(Ey \bullet Px) \bullet (Ey \bullet Px)] \bullet Hxy]', \ D. \ (\exists x)[(Ey \bullet Px) \bullet (Ey \bullet Px)] \bullet Hxy]', \ D. \ (\exists x)[(Ey \bullet Px) \bullet (Ey \bullet Px)] \bullet Hxy]', \ D. \ (\exists x)[(Ey \bullet Px) \bullet (Ey \bullet Px)]$

Answer Option 1:

The best translation into predicate logic for the given statement would be:

 $(\exists x)(\exists y)\{[(\mathsf{K}x \bullet \mathsf{P}x) \bullet (\mathsf{E}y \bullet \mathsf{S}y)] \bullet \mathsf{H}xy\}$

This translation correctly captures the statement that there exist some kind people (Kx • Px) who help some endangered species (Ey • Sy) with the predicate Hxy. Therefore, the answer is option B.

Answer Option 2:

The best translation into predicate logic for the given statement is option B. '(3x)(3y)([(Kx • Sx) • (Ey • Py)] • Hxy}'. This translation correctly captures the idea that there exist some individuals (x) who are kind (Kx) and some individuals (y) who are endangered species (Sy), and that these kind individuals help the endangered species (Hxy). Final Answer: B

Answer Option 3:

The best translation into predicate logic for the statement "Some kind people help some endangered species" would be option A: $(\exists x)\{(Kx \cdot Px) \cdot (\exists y)[(Ey \cdot Sy) \cdot Hxy]\}$.

This translates to: There exists a person x who is kind and there exists a species y such that there exists an endangered species y and person x helps species y.

This accurately captures the idea that there are kind people who help endangered species.

Final Answer: A

Answer Option 4:

The best translation into predicate logic for the given statement is:

A. $(\exists x)\{(Kx \cdot Px) \cdot (\exists y)[(Ey \cdot Sy) \cdot Hxy]\}$

Let's break down the translation:

- (∃x) denotes the existence of some x
- (Kx ullet Px) represents x being kind and x being a person
- $(\exists y) \hbox{\tt [(Ey \bullet Sy) \bullet Hxy] denotes the existence of some y where y is endangered, y is a species, and x helps y the existence of some y where y is endangered.}$

The translation accurately captures the statement "Some kind people help some endangered species." Final Answer: A

Answer Option 5:

The best translation into predicate logic for the given statement is:

 $(\exists x)(\exists y)\{[(Kx \bullet Px) \bullet (Ey \bullet Sy)] \bullet Hxy\}$

Explanation:

- $(\exists x)(\exists y)$ indicates that there exist some x and y.
- [(Kx \bullet Px) \bullet (Ey \bullet Sy)] represents the condition that x is kind and a person, and y is a species.
- Hxy indicates that x helps y.

Thus, the correct translation is:

 $(\exists x)(\exists y)\{[(\mathsf{Kx}\bullet\mathsf{Px})\bullet(\mathsf{Ey}\bullet\mathsf{Sy})]\bullet\mathsf{Hxy}\}$

Final Answer: B

Answer Option 6:

The best translation into predicate logic for the given statement is:

A. $'(\exists x)\{(Kx \cdot Px) \cdot (\exists y)[(Ey \cdot Sy) \cdot Hxy]\}'$

Let's break it down:

 $(\exists x)$ - There exists an x

{(Kx • Px) • (∃y)[(Ey • Sy) • Hxy]} - such that x is kind and x is a person, and there exists a y such that y is a species, y is endangered, and x helps y.

This translation accurately represents the statement "Some kind people help some endangered species."

Final Answer: A

Answer Option 7:

The best translation into predicate logic for the given statement would be: A. '{∃x}{(Kx • Px) • (∃y){(Ey • Sy) • Hxy]}' Explanation: The original statement states "Some kind people help some endangered species." In predicate logic, we can represent "Some" using existential quantifiers (∃). We need two variables, one for the kind person and one for the endangered species. Let's use x for the person and y for the species. "Kind people" can be represented as (Kx • Px), meaning x is kind and x is a person. "Endangered species" can be represented as (Ey • Sy), meaning y is endangered and y is a species. Finally, "help" can be represented as Hxy, meaning x helps y. Combining all these parts, we get '(∃x){(Kx • Px) • (∃y){(Ey • Sy) • Hxy]}' as the best translation. Final Answer: A

Answer Option 8:

The sentence "Some kind people help some endangered species" can be translated into predicate logic as follows:

 $Let \ Kx \ represent \ "x \ is \ kind", \ Px \ represent \ "x \ is \ a \ person", \ Sx \ represent \ "x \ is \ a \ species", \ and \ Hxy \ represent \ "x \ helps \ y".$

The correct translation is:

A. $(\exists x)\{(Kx \bullet Px) \bullet (\exists y)[(Ey \bullet Sy) \bullet Hxy]\}$

In this translation, $(\exists x)$ signifies "there exists at least one x", $(Kx \bullet Px)$ represents "x is kind and x is a person", $(\exists y)[(Ey \bullet Sy) \bullet Hxy]$ signifies "there exists at least one y such that y is a species and x helps y". Therefore, the best translation is A. $(\exists x)((Kx \bullet Px) \bullet (\exists y)[(Ey \bullet Sy) \bullet Hxy])$.

Answer Option 9:

The best translation into predicate logic for the given statement would be: A. $(\exists x)\{(Kx \cdot Px) \cdot (\exists y)[(Ey \cdot Sy) \cdot Hxy]\}$ 'Explanation: The statement "Some kind people help some endangered species" can be translated as: $(\exists x)\{(Kx \cdot Px) \cdot (\exists y)[(Ey \cdot Sy) \cdot Hxy]\}$ The first part $(\exists x)\{(Kx \cdot Px) \cdot (\exists y)[(Ey \cdot Sy) \cdot Hxy]\}$ The first part $(\exists x)\{(Kx \cdot Px) \cdot (\exists y)[(Ey \cdot Sy) \cdot Hxy]\}$ The first part $(\exists x)\{(Kx \cdot Px) \cdot (\exists y)[(Ey \cdot Sy) \cdot Hxy]\}$ The first part $(\exists x)\{(Kx \cdot Px) \cdot (\exists y)[(Ey \cdot Sy) \cdot Hxy]\}$ The first part $(\exists x)\{(Kx \cdot Px) \cdot (\exists y)[(Ey \cdot Sy) \cdot Hxy]\}$ The first part $(\exists x)\{(Kx \cdot Px) \cdot (\exists y)[(Ey \cdot Sy) \cdot Hxy]\}$ The first part $(\exists x)\{(Kx \cdot Px) \cdot (\exists y)[(Ey \cdot Sy) \cdot Hxy]\}$ The first part $(\exists x)\{(Ey \cdot Sy) \cdot (Ey \cdot Sy) \cdot ($

Select the best English interpretation of the given proposition, using the following translation key: Ax: x is an apartment Hx: x is a house Lx: x is large Bxy: x is bigger than y ($\forall x$)[Ax \supset ($\exists y$)(Hy • Bxy)], A. 'Every apartment is bigger than some house.', B. 'Every house is bigger than all apartments.', C. 'Every house is bigger than some apartment.', D. 'Some house is bigger than every apartment.' Answer Option 1: Using the translation key provided, the given proposition can be interpreted as: "For every apartment, there exists a house that is larger than it." Interpreting the given proposition in English using the provided translation key, we can conclude that the best interpretation is option A: "Every apartment is bigger than some house." Final Answer: A Answer Option 2: The given proposition is $(\forall x)[Ax \supset (\exists y)(Hy \bullet Bxy)]$. Using the translation key, we can interpret it as "For every x, if x is an apartment, then there exists a y such that y is a house and x is bigger than y." The best English interpretation of this proposition is: A. "Every apartment is bigger than some house." This is because the proposition asserts that for every apartment, there exists a house that the apartment is bigger than. Answer Option 3: The English interpretation of the given proposition $(\forall x)[Ax \supset (\exists y)(Hy \bullet Bxy)]$ is: 'Every apartment is bigger than some house.' $\forall x$: For every x, Ax: x is an apartment, ⇒: implies, (∃y): there exists y, Hy: y is a house, •: and, Bxy: x is bigger than y. Therefore, the best English interpretation is: 'Every apartment is bigger than some house.' Final Answer: A. 'Every apartment is bigger than some house.' The given proposition is $(\forall x)[Ax \supset (\exists y)(Hy \bullet Bxy)]$. Using the translation key, we can break it down as follows: - Ax: x is an apartment - Hx: x is a house - Bxy: x is bigger than y The proposition can be translated as "For every apartment x, if x is an apartment, then there exists a house y such that y is larger than x." The best English interpretation of the given proposition is "Every apartment is bigger than some house." Answer Option 5: The given proposition, $(\forall x)[Ax \supset (\exists y)(Hy \bullet Bxy)]$, translates to "For every x, if x is an apartment, then there exists a y such that y is a house and x is bigger than y." Using the translation key: Ax: x is an apartment Hx: x is a house Lx: x is large Bxy: x is bigger than y The best English interpretation of the given proposition is "Every apartment is bigger than some house." This corresponds to option A. Final Answer: A Answer Option 6: Let's break down the given proposition using the translation key: $(\forall x)[Ax\supset (\exists y)(Hy\bullet Bxy)]$

Question 90:

 $(\forall x)$ - for every xAx - x is an apartment

⊃ - implies
(∃y) - there exists a y
Hy - y is a house
• - and
Bxy - x is bigger than y
Based on the translation, the best English interpretation of the given proposition is:
A. 'Every apartment is bigger than some house.'
This is because the proposition states that for every apartment, there exists a house such that the apartment is bigger than that house.
Answer Option 7:
The given proposition is $(\forall x)[Ax \supset (\exists y)(Hy \bullet Bxy)]$. Let's break it down with the translation key:
Ax: x is an apartment
Hx: x is a house
Lx: x is large
Bxy: x is bigger than y
(∀x)[Ax ⊃ (∃y)(Hy • Bxy)] can be interpreted as "For every x, if x is an apartment, then there exists a y such that y is a house and x is bigger than y."
The best English interpretation of this proposition would be: A. "Every apartment is bigger than some house."
Answer Option 8:
To interpret the given proposition using the translation key, we can break it down step by step:
$(\forall x)[Ax \supset (\exists y)(Hy \bullet Bxy)]$
∀x: For every x
Ax: x is an apartment
∴ implies
(∃y): There exists a y
Hy: y is a house
•: and
Bxy: x is bigger than y
Combining these translations, the proposition can be interpreted as:
For every apartment, there exists a house such that the house is smaller than the apartment.
Therefore, the best English interpretation of the given proposition is:
A. 'Every apartment is bigger than some house.'

Answer Option 9:
The given proposition is $(\forall x)[Ax \supset (\exists y)(Hy \bullet Bxy)]$.
Translation key:
Ax: x is an apartment
Hx: x is a house
Lx: x is large
Bxy: x is bigger than y
The best English interpretation of the given proposition is:
A. 'Every apartment is bigger than some house.'
Explanation: (Yall by a Challette Paul) and be read as "For every vifu is an appropriate these evicts a visual back via a begin and via higger than v." This magnet that for every appropriate these evicts a begin and via higger than v." This magnet that for every appropriate these evicts a begin appropriate the visual a
(∀x)[Ax ⊃ (∃y)(Hy • Bxy)] can be read as "For every x, if x is an apartment, then there exists a y such that y is a house and x is bigger than y." This means that for every apartment, there exists a house that is
smaller than the apartment. Thus, the best English interpretation is that every apartment is bigger than some house. Final Answer: A

Question 91: *Construct a complete truth table for the following argument. Then, using the truth table, determine whether the argument is valid or invalid. If the argument is invalid, choose an option which presents a counterexample. (There may be other counterexamples as well.) ~(S v Q) / R', A. 'Valid', B. 'Invalid. Counterexample when Q and S are true and R is false', C. 'Invalid. Counterexample when Q is true and S and R are false', D. 'Invalid. Counterexample when Q, S, and R are false Answer Option 1: To construct the truth table for the argument, we need to consider all possible combinations of truth values for Q, R, and S. $Q R S \sim (S \vee Q) R'$ TTTFF TTETT TFTFF TFFTT FTTFF FTFTT FFTFF FFFTT From the truth table, we can see that for all rows where \sim (S \vee Q) is true, R' is also true. Therefore, the argument is valid. Final Answer: A. 'Valid' ******** Answer Option 2: First, let's construct a truth table for the given argument: | Q | R | S | ~(S ∨ Q) | R' | |---|---| |T|T|T|F|F| |T|T|F|F|F| |T|F|T|F|T| |T|F|F|T|T| |F|T|T|F|F| |F|T|F|T|F| |F|F|T|F|T| |F|F|F|T|T| Based on the truth table, we can see that for every row where the premise (-(S v Q)) is true, the conclusion (R') is also true. Therefore, the argument is valid. Final Answer: A. 'Valid' Answer Option 3: To construct the truth table for the argument, we need to consider all possible combinations of truth values for the propositions Q, R, and S. $QRSQ \equiv R \neg (S \lor Q)R'$ TTTTFF

TTFTFF

TFTFFF

.

TFFTFT

FTTFFT

FTFFTT

FFTFFT

FFFFTT

Now, let's evaluate the truth value of the conclusion (R') for each row of the truth table.

For the rows where R' is false (F), we have:

Row 1: Q=T, R=T, S=T

Row 2: Q=T, R=T, S=F

Thus, we have found a counterexample where the premises of the argument are true $(Q \equiv R, \sim (S \vee Q))$ but the conclusion (R') is false. Therefore, the argument is invalid.

The correct answer is B. "Invalid. Counterexample when Q and S are true and R is false".

******* Answer Option 4: $To \ construct \ a \ truth \ table \ for \ the \ argument, \ we \ need \ to \ consider \ all \ possible \ combinations \ of \ truth \ values \ for \ the \ propositions \ Q, \ R, \ and \ S.$ Let's build the truth table: $Q\mid R\mid S\mid S\vee Q\mid \mathord{\sim}(S\vee Q)\mid Q\equiv R\mid \mathord{\sim} R$ T|T|T|T|F|T|FT|T|F|T|F|T|F T|F|T|T|F|F|T T|F|F|T|F|F|T FITITITIFITIF F|T|F|F|T|T|F F|F|T|T|F|F|T F|F|F|F|T|T The given argument is: $Q \equiv R$ ~(S ∨ Q) / R' Now let's evaluate the argument based on the truth table. In all rows where the premises (Q \equiv R) and \sim (S \vee Q) are true, we check if the conclusion (\sim R) is also true. In rows 1, 2, 5, and 7, where the premises are true, the conclusion is also true. Therefore, the argument is valid. Final Answer: A. Valid ******** Answer Option 5: To construct a complete truth table for the argument, we need to consider all possible combinations of truth values for the propositions involved: Q, R, and S. $\textbf{Q} \equiv \textbf{R}$ means Q and R have the same truth value. ~(S \vee Q) means the negation of (S or Q). R' means the negation of R. Let's construct the truth table: |---|---|----|-----| |T|T|T|T|F|F| |T|T|F|T|F|F| |T|F|T|T|F|T| |T|F|F|T|F|T||F|T|T|T|F|F| |F|T|F|F|T|F| |F|F|T|T|F|T| |F|F|F|F|T|T| Based on the truth table, we can see that when Q = R, the statement $\sim (S \vee Q) / R'$ holds true for all possible combinations of truth values. Therefore, the argument is valid. Final Answer: A. 'Valid' ****** Answer Option 6: To construct a complete truth table for the argument, we consider all possible truth value combinations for the propositions Q, R, and S. $Q \equiv R$ means Q is true if and only if R is true. This can be represented as $(Q \land R) \lor (\neg Q \land \neg R)$. The negation of (S \vee Q) is \sim (S \vee Q), which can be represented as \neg S \wedge \neg Q. Using these representations, we can construct the truth table as follows: $\mid Q \mid R \mid S \mid (Q \wedge R) \vee (\neg Q \wedge \neg R) \mid \neg S \wedge \neg Q \mid \neg (S \vee Q) \mid$ |---|---|--------|-------| |T|T|T|T|F|F|

|T|T|F|T|T|F| |T|F|T|F|F|T| |T|F|F|F|T|F| |F|T|T|F|F|F| |F|T|F|F|T|T| |F|F|T|T|F|F| |F|F|F|T|T|T| Now, let's analyze the argument: ~(S ∨ Q) / R' We notice that the only row where $\neg(S \lor Q)$ is true is when Q is false, S is false, and R is false. In this row, R' is true. Hence, the argument is valid and the correct option is A. 'Valid'. Answer Option 7: To construct a complete truth table for the argument, we need to consider all possible truth values for the propositions involved. First, let's analyze the given argument and fill in the columns for Q, R, S, \sim (S \vee Q), and R'. $Q \equiv R$ represents the equivalence between Q and R, which means Q and R have the same truth value. So, we can fill in the corresponding columns for $Q \equiv R$. Next, we consider ~(S v Q), which means the negation of the disjunction between S and Q. We can fill in the column for ~(S v Q) based on the truth values of S and Q. Finally, we consider R', which represents the negation of R. We can fill in the column for R' based on the truth value of R. Here is the complete truth table: | Q | R | S | ~(S ∨ Q) | R' | |---|---| |T|T|T|F|F| |T|T|F|F|F| |T|F|T|F|T| |T|F|F|T|T| |F|T|T|F|F| |F|T|F|T|F| |F|F|T|F|T| |F|F|F|T|T| Now, let's use the truth table to determine whether the argument is valid or invalid. From the truth table, we can see that whenever ~(S v Q) is true, R' is also true. This means that whenever the premise ~(S v Q) is true, the conclusion R' is also true. Therefore, the argument is valid. Final Answer: A. 'Valid' ****************** Answer Option 8: To construct a complete truth table for the argument, we need to consider all the possible combinations of truth values for the propositions involved. Let's analyze the given argument: Premise: Q ≡ R Premise: \sim (S \vee Q) Conclusion: R' We have three propositions involved: Q, R, and S. First, we need to create columns for the truth values of Q, R, S, $Q \equiv R$, $\sim (S \lor Q)$, and R'. Below is the complete truth table: $|Q|R|S|Q \equiv R|\sim(S\vee Q)|R'|$ |---|---|-----| |T|T|T|T|F|F| |T|T|F|T|F|F| |T|F|T|F|F|T| |T|F|F|F|F|T| |F|T|T|F|T|F| |F|T|F|F|T|F| |F|F|T|T|F|T|

Now, let's analyze the conclusion R'. According to the truth table, R' is true when R is false and false for all other cases.

|F|F|F|T|T|T|

The argument is valid when the conclusion R' is true for all cases where the premises are true. In this case, we can see that R' is false when Q = R is true and ~(S ∨ Q) is true. Therefore, the argument is invalid.

To find a counterexample, we need to look for a row in the truth table where both premises are true (Q = R is true and $\sim (S \vee Q)$ is true) but the conclusion R' is false.

Looking at the truth table, we can see that the counterexample occurs when Q is true, S is true, and R is false. Therefore, the correct answer is:

B. Invalid. Counterexample when Q and S are true and R is false.

Answer Option 9:

To construct the truth table for the argument, we need to consider all possible truth values for the propositions Q, R, and S.

Q = R represents the biconditional statement, which means Q and R have the same truth value. So, we can write the truth table for Q = R as follows:

 $|Q|R|Q \equiv R|$

|---|----|

|T|T|T|

|T|F|F|

|F|T|F|

|F|F|T|

Next, we consider the proposition ~(S \vee Q). We can write the truth table for ~(S \vee Q) as follows:

| S | Q | S v Q | ~(S v Q) |

|---|---|------|

|T|T|T|F|

|T|F|T|F|

|F|T|T|F|

1 - 1 - 1 - 1 - 1

|F|F|F|T|

Lastly, we consider the proposition R'.

| R | R' |

|---|----|

|T|F|

|F|T|

Now, let's combine all the propositions and construct the complete truth table for the argument:

 $\mid Q \mid R \mid S \mid Q \equiv R \mid S \vee Q \mid \sim (S \vee Q) \mid R' \mid$

|---|---|-----|-----|

|T|T|T|T|F|F|

|T|T|F|T|T|F|F|

|T|F|T|F|T|F|T|

|T|F|F|F|T|F|T|

|F|T|T|F|F|T|F|

|F|T|F|F|F|T|F|

|F|F|T|T|F|T|T|

|F|F|F|T|F|T|T|

Now, let's examine the truth values of the argument:

The premise is ~ (S v Q), which is false in rows 1, 2, 3, and 4. In all these rows, R' is true. Therefore, the conclusion R' holds true for all cases where the premise is false.

Since the conclusion holds true for all possible truth values, the argument is valid.

Final Answer: A. 'Valid'

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'Identify the antecedent of the following conditional proposition: The university raises tuition only if both the governor approves of it and the board of trustees recommends it.', A. 'The university raises tuition.',

B. The governor approves of it.', C. 'The board of trustees recommends it.', D. 'The governor approves of it and the board of trustees recommends it.'

Answer Option 1:

The antecedent of a conditional proposition is the part that comes after the "if" or "only if" and before the "then". In this case, the antecedent is the condition that needs to be satisfied for the university to raise tuition.

The antecedent of the given conditional proposition is: 'The governor approves of it and the board of trustees recommends it.'

Therefore, the correct answer is D. 'The governor approves of it and the board of trustees recommends it.'

Answer Option 2:

The antecedent of the conditional proposition is the part that comes before the "only if" phrase. In this case, the antecedent is "The university raises tuition."

Therefore, the correct answer is A. 'The university raises tuition.'

Final Answer: A

Answer Option 3:

The antecedent of a conditional proposition is the part that comes after the "if" and before the "then". In this case, the antecedent is "both the governor approves of it and the board of trustees recommends it."

Therefore, the correct answer is option D: 'The governor approves of it and the board of trustees recommends it.'

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Answer Option 4:

The antecedent of a conditional proposition is the part that comes before the 'if' statement. In this case, the antecedent is the condition that needs to be met for the university to raise tuition.

The antecedent in the given conditional proposition is:

'The governor approves of it and the board of trustees recommends it.'

Therefore, the correct answer is D. 'The governor approves of it and the board of trustees recommends it.'

Answer Option 5:

The antecedent of a conditional proposition is the part that comes after the "if" and before the "then". In this case, the antecedent is "both the governor approves of it and the board of trustees recommends it."

Therefore, the correct answer is D. 'The governor approves of it and the board of trustees recommends it.'

Answer Option 6:

The antecedent of a conditional proposition is the part that comes before the "if" statement. In this case, the antecedent is the condition that needs to be met for the university to raise tuition.

Based on the given statement, the antecedent is: The governor approves of it and the board of trustees recommends it.'

Therefore, the correct answer is D. 'The governor approves of it and the board of trustees recommends it.'

Answer Option 7:

The antecedent of a conditional proposition is the part that comes after the "if" and before the "then". In this case, the antecedent is the condition that needs to be satisfied in order for the university to raise tuition.

The antecedent in this conditional proposition is "The governor approves of it and the board of trustees recommends it."

Therefore, the correct answer is D. 'The governor approves of it and the board of trustees recommends it.'

Answer Option 8:

The antecedent of a conditional proposition is the part that comes before the "only if" or "if" in the statement. In this case, the antecedent is the condition that needs to be satisfied for the university to raise tuition.

The antecedent of the given conditional proposition is: 'The governor approves of it and the board of trustees recommends it.'

Therefore, the correct answer is: D. 'The governor approves of it and the board of trustees recommends it.'

Answer Option 9:

The antecedent of a conditional proposition is the part that comes before the "if" or "only if" in the statement. In this case, the antecedent is "The university raises tuition only if both the governor approves of it and the board of trustees recommends it." Therefore, the correct answer is A. The university raises tuition.'

'Which of the following propositions is an immediate (one-step) consequence in PL of the given premises? $(\mathsf{A} \vee \mathsf{B}) \supset (\mathsf{C} \vee \mathsf{D})$ A v B ~C', A. '~A \vee ~B', B. '~~C', C. 'C \vee D', D. '(A \vee B) \supset D' To find the immediate consequence in propositional logic (PL), we can simplify the premises using valid inference rules. 1. $(A \lor B) \supset (C \lor D)$ 2. A ∨ B 3. ~C To derive the immediate consequence, we can use the modus tollens rule, which states that if we have a conditional statement of the form P \supset Q and we know \sim Q, then we can infer \sim P. Applying modus tollens to premise 1 [(A \vee B) \supset (C \vee D)] and premise 3 (~C), we can infer ~ (A \vee B). Final Answer: A. '~A v ~B' ******** Answer Option 2: To determine the immediate (one-step) consequence, we need to identify a proposition that can be directly inferred from the given premises using a single logical step. The given premises are: 1. $(A \lor B) \supset (C \lor D)$ 2. A v B Based on these premises, we can make the following deduction: From premise 1, using the rule of implication elimination (modus ponens), we can infer: $(C \lor D)$ From premise 3, we have: ~C Thus, the immediate consequence is the proposition ' \sim C \vee \sim D'. Therefore, the correct option is A. '~A \vee ~B'. Answer Option 3: To find the immediate consequence of the given premises, we can use the rule of implication elimination (modus ponens). Given premises: $1.\;(A\vee B)\supset (C\vee D)$ 2. A v B To apply modus ponens, we need to have the antecedent of the implication in premise 1 (A v B) and the corresponding matching disjunctive statement in either premise 2 or 3. From premises 1 and 2, we can conclude $C \vee D$ because $(A \vee B)$ is the antecedent of the implication and it matches with $A \vee B$ in premise 2. Therefore, the immediate consequence is $C \vee D$. Final Answer: C. 'C v D' Answer Option 4: To determine which proposition is an immediate consequence of the given premises, we can apply logical inference rules. Given premises: 1. $(A \lor B) \supset (C \lor D)$ 2. A v B 3. ~C To derive the immediate consequence, we can use the modus tollens rule. Modus tollens states that if we have an implication of the form 'if A then B' and we know that B is false, then we can infer that A must also be false. Let's apply the modus tollens rule to the given premises: From premise 1: $(A \lor B) \supset (C \lor D)$ Negating the consequent: \sim (C \vee D) We can conclude: ~(A v B)

Question 93:

Therefore, the immediate consequence is '~(A \vee B)'.

Final Answer: A. '~A v ~B' ******** Answer Option 5: To determine the immediate consequence in propositional logic (PL), we need to apply the rules of inference to the given premises. Given premises: 1. $(A \lor B) \supset (C \lor D)$ 2. A ∨ B 3. ~C Looking at the premises, we can see that we have a conditional statement in premise 1, and disjunctions in premises 2 and 3. We can use the rule of modus ponens, which states that if we have a conditional statement $(p \supset q)$ and also have p, we can conclude q. Using modus ponens, we can combine premises 1 and 2 to conclude (C \vee D). However, none of the answer options match this immediate consequence. Therefore, none of the given answer options is the correct one-step consequence in PL. Answer Option 6: To find the immediate consequence of the given premises, we need to apply the rules of propositional logic. Premise 1: (A v B) \supset (C v D) Premise 2: A v B Premise 3: ~C We can use the rule of Modus Tollens to derive the immediate consequence: Modus Tollens: If we have an implication of the form $P \to Q$ and we know that Q is false, then we can conclude that P must also be false. Applying Modus Tollens to the premises: From Premise 1: (A ∨ B) \supset (C ∨ D) and Premise 3: ~C, we can infer ~(A ∨ B) ~C From Premise 2: A ∨ B, we can derive ~(A ∨ B) ~C Therefore, the immediate consequence is '~(A ∨ B) ~C From Premise 2: A ∨ B, we can derive ~(A ∨ B) ~C Therefore, the immediate consequence is '~(A ∨ B) ~C From Premise 2: A ∨ B, we can derive ~(A ∨ B) ~C Therefore, the immediate consequence is '~(A ∨ B) ~C From Premise 2: A ∨ B, we can derive ~(A ∨ B) ~C Therefore, the immediate consequence is '~(A ∨ B) ~C From Premise 2: A ∨ B, we can derive ~(A ∨ B) ~C Therefore, the immediate consequence is '~(A ∨ B) ~C From Premise 2: A ∨ B, we can derive ~(A ∨ B) ~C Therefore, the immediate consequence is '~(A ∨ B) ~C From Premise 2: A ∨ B, we can derive ~(A ∨ B) ~C Therefore, the immediate consequence is '~(A ∨ B) ~C From Premise 2: A ∨ B, we can derive ~(A ∨ B) ~C Therefore, the immediate consequence is '~(A ∨ B) ~C From Premise 2: A ∨ B, we can derive ~(A ∨ B) ~C Therefore, the immediate consequence is '~(A ∨ B) ~C From Premise 2: A ∨ B, we can derive ~(A ∨ B) ~C Therefore, the immediate consequence is '~(A ∨ B) ~C From Premise 2: A ∨ B, we can derive ~(A ∨ B) ~C Therefore, the immediate consequence is '~(A ∨ B) ~C From Premise 2: A ∨ B, we can derive ~(A ∨ B) ~C From Premise 2: A ∨ B, we can derive ~(A ∨ B) ~C From Premise 2: A ∨ B, we can derive ~(A ∨ B) ~C From Premise 2: A ∨ B, we can derive ~(A ∨ B) ~C From Premise 2: A ∨ B, we can derive ~(A ∨ B) ~C From Premise 2: A ∨ B, we can derive ~(A ∨ B) ~C From Premise 2: A ∨ B, we can derive ~(A ∨ B) ~C From Premise 2: A ∨ B, we can derive ~(A ∨ B) ~C From Premise 2: A ∨ B, we can derive ~(A ∨ B) ~C From Premise 2: A ∨ B, we can derive ~(A ∨ B) ~C From Premise 2: A ∨ B, we can derive ~(A ∨ B) ~C From Premise 2: A ∨ B, we can derive ~(A ∨ B) ~C From Premise 2: A ∨ B, we can derive ~(A ∨ B) ~C From Premise 2: A ∨ B, we can derive ~(A ∨ B) ~C From Premise 2: A ∨ B, we can derive ~(A ∨ B) ~C From Premise 2: A ∨ B, we can derive ~(A ∨ B) ~C From Premise 2: A ∨ B, we can deriv B) ~C'. Final Answer: A Answer Option 7: To find the immediate consequence in Propositional Logic (PL) of the given premises, we need to apply the rules of implication and disjunction. Given premises: 1. $(A \lor B) \supset (C \lor D)$ 2. A ∨ B 3. ~C To derive the immediate consequence, we can use the rule of disjunctive syllogism. The disjunctive syllogism states that if we have a disjunction (A \vee B) and the negation of one of the disjuncts (\sim A in this case), we can infer the other disjunct (B). Using this rule, we can infer B from premise 2 (A v B) and the negation of A (~A). Therefore, the immediate consequence in PL of the given premises is B: '~~C'. Final Answer: B. '~~C' ****************** Answer Option 8: To find an immediate (one-step) consequence in propositional logic (PL), we need to use the rules of inference. Given premises: 1. $(A \lor B) \supset (C \lor D)$ 2. A v B 3. ~C We can use the rule of disjunctive syllogism to derive an immediate consequence. The disjunctive syllogism states that if we have a disjunction (A v B) and the negation of one of the disjuncts (~A or ~B), then we can conclude the other disjunct. In this case, we have (A v B) from premise 2 and ~C from premise 3. Using the disjunctive syllogism, we can conclude ~A. Therefore, the immediate consequence is '~A v ~B'. Final Answer: A. '~A v ~B' ****************** Answer Option 9: To determine which proposition is an immediate consequence of the given premises, we can use the process of logical deduction. Given premises: 1. $(A \lor B) \supset (C \lor D)$ $2.\;A\vee B$ 3. ~C We will evaluate each option based on these premises: Option A: '~A v ~B' This option does not follow directly from the given premises. It introduces the negations of A and B, which are not explicitly stated or implied in the premises.

This option simplifies to 'C'. While 'C' is consistent with the premises, it is not an immediate consequence, as it does not follow directly from the given premises.

Option B: '~~C'

Option C: ${}^{\prime}C \vee D'$

This option directly matches one of the premises, $(C \lor D)$, and is therefore an immediate consequence of the given premises.

Option D: '(A \vee B) \supset D'

This option introduces the implication (\supset) and D, which are not explicitly stated or implied in the premises.

Therefore, the correct answer is C. 'C \vee D' is an immediate consequence of the given premises.

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'Identify the antecedent of the following conditional proposition: When the university raises tuition, then either the governor approves of it or the board of trustees doesn't prevent it.', A. 'The university raises tuition.', B. 'The governor approves of it.', C. 'The board of trustees prevents it.', D. The board of trustees doesn't prevent it.'

Answer Option 1:

The antecedent of a conditional proposition is the statement that comes immediately after the "if" or "when" clause. In this case, the antecedent of the given conditional proposition is "The university raises tuition." Therefore, the correct answer is A. The university raises tuition."

Answer Option 2:

The antecedent of a conditional proposition is the statement that comes before the "then" in the conditional statement. In this case, the antecedent is "The university raises tuition".

Therefore, the correct answer is A. 'The university raises tuition.'

Answer Option 3:

The antecedent of a conditional proposition is the statement that comes before the "then" in the sentence. In this case, the antecedent is "The university raises tuition." Therefore, the correct answer is A. 'The university raises tuition.'

Answer Option 4:

The antecedent of a conditional proposition is the statement that comes immediately after the "if" in the sentence. In this case, the antecedent is "The university raises tuition.' Therefore, the correct answer is

A. 'The university raises tuition.'

Answer Option 5:

The antecedent of a conditional proposition is the statement that comes after the "if" or "when" part of the sentence. In this case, the antecedent is "The university raises tuition."

Therefore, the correct answer is A. 'The university raises tuition.'

Answer Option 6:

The antecedent of a conditional proposition is the statement that comes immediately after the "if" or "when" part of the conditional statement. In this case, the antecedent of the conditional proposition is "The university raises tuition." Therefore, the correct answer is A. 'The university raises tuition.'

Answer Option 7:

The antecedent of a conditional proposition is the part that comes before the "then" keyword. In this case, the antecedent is the statement that triggers the conditional

The antecedent of the given conditional proposition is: 'The university raises tuition.'

Therefore, the correct answer is: A. 'The university raises tuition.'

Answer Option 8:

The antecedent of a conditional proposition is the first part of the statement that comes before the "then" keyword. In this case, the antecedent is the condition that is being discussed.

The antecedent of the given conditional proposition is: "The university raises tuition."

So, the correct answer is A. 'The university raises tuition.'

Answer Option 9:

The antecedent in a conditional proposition is the part that comes before the "then" or the arrow (\Rightarrow). In this case, the conditional proposition is "When the university raises tuition, then either the governor approves of it or the board of trustees doesn't prevent it."

The antecedent is "The university raises tuition."

Therefore, the correct answer is A. 'The university raises tuition.'

Question 95: 'Select the best translation into predicate logic. Marco doesn't move from Spain to Italy. (i: Italy; m: Marco; s: Spain; Mxyz: x moves to y from z)', A. '-Mmsi', B. 'M-mis', C. '-Mmis', D. 'M-msi' Answer Option 1: The best translation into predicate logic for the given statement "Marco doesn't move from Spain to Italy" would be: "~Mmsi" Explanation: Mmsi: Marco moves from Spain to Italy. ~Mmsi: Marco doesn't move from Spain to Italy. Therefore, the answer is A. ~Mmsi Answer Option 2: The best translation into predicate logic for the statement "Marco doesn't move from Spain to Italy" is: ~Mxyz (x=m, y=s, z=i) Breaking it down: ~ - negation M - Marco moves x=y - Marco stays in the same place (Spain) y=s - Marco is in Spain z=i - Marco moves to Italy So, the correct answer is A. '~Mmsi' Answer Option 3: The best translation into predicate logic for the given statement "Marco doesn't move from Spain to Italy" would be: "~Mmsi". In this translation: - M represents "Marco moves" - m represents "Marco" - s represents "Spain" - i represents "Italy" The "~" symbol represents the negation, indicating that Marco does not move from Spain to Italy. Therefore, the correct answer is: A. "~Mmsi" Answer Option 4: The best translation into predicate logic for the given statement would be "M-msi". This translates to "Marco does not move from Spain to Italy". Explanation: M: Marco s: Spain The negation (~) is applied to the predicate "m" to indicate that Marco does not move. "msi" represents Marco moving from Spain to Italy. Hence, the correct answer is D. "M~msi". Answer Option 5: The best translation into predicate logic would be: A. '~Mmsi' Let's break down the given sentence: Marco doesn't move from Spain to Italy. We can represent this using the following predicates: Mxyz: x moves to y from z. To represent Marco not moving from Spain to Italy, we can use the following translation: ~Mmsi, which means "It is not the case that Marco moves from Spain to Italy." Therefore, the correct translation into predicate logic is A. '~Mmsi'.

The best translation into predicate logic for the given sentence would be " \sim Mmsi". Here's the breakdown:

Answer Option 6:

M: Marco moves
m: Marco
s: Spain
i: Italy
The given sentence states that Marco doesn't move from Spain to Italy. In predicate logic, we can represent this as "~Mmsi", where "~" represents negation, "M" represents Marco moves, "m" represents
Marco, "s" represents Spain, and "i" represents Italy.
Final Answer: A. ~Mmsi

Answer Option 7:
The best translation into predicate logic for the statement "Marco doesn't move from Spain to Italy" would be option A: '~Mmsi'. The statement can be broken down as follows: M: Marco, m: moves, s: Spain, i:
Italy. "Marco doesn't move from Spain to Italy" can be translated as "It is not the case that Marco moves from Spain to Italy". In predicate logic, this can be represented as ~Mmsi. Final Answer: A

Answer Option 8:
The best translation into predicate logic for the given statement is: A. '~Mmsi'
Explanation:
Let's break down the statement into its components:
- Marco doesn't move from Spain to Italy.
- Let m represent Marco, s represent Spain, and i represent Italy.
- Let Mxyz represent "x moves to y from z".
From the given components, we can form the predicate logic statement: -Mmsi, which can be read as "It is not the case that Marco moves from Spain to Italy."
Therefore, the best translation into predicate logic is A. '-Mmsi'. Final Answer: A

Answer Option 9:
The best translation into predicate logic for the statement "Marco doesn't move from Spain to Italy" is:
~Mmsi
Explanation:
M - Marco
m - move
s - Spain
i - Italy
The statement can be translated as "It is not the case that Marco moves from Spain to Italy", which can be represented as ~Mmsi.
Option A: ~Mmsi - This is the correct translation.
Option B: M-mis - This translates to "Marco moves and it is not the case that Marco is in Spain". It does not capture the meaning of the original statement.
Option C: ~Mmis - This translates to "It is not the case that Marco is in Spain and Italy". It does not capture the meaning of the original statement.
Option D: M-msi - This translates to "Marco moves and it is not the case that Marco is in Spain or Italy". It does not capture the meaning of the original statement.
Final Answer: A

Question 96:
'Construct a complete truth table for the following argument. Then, using the truth table, determine whether the argument is valid or invalid. If the argument is invalid, choose an option which presents a
counterexample. (There may be other counterexamples as well.)
~\$ v T
-\$ · U
~T v U / T · U', A. "Valid', B. "Invalid. Counterexample when S, T, and U are true', C. "Invalid. Counterexample when S and U are true and T is false', D. "Invalid. Counterexample when U is true and S and T are
false'

Answer Option 1:
To construct a complete truth table for the argument, we need to consider all possible truth values for the propositions involved: S, T, and U.
STU~SvT~S·U~TvUT·U
TTTFTFTT
TTFFFFF
TETETET
TFFFFFF
FTTTTTT
FTFTFFFF
FFTTTTF
FFFTFTTF
Based on the truth table, we can see that there is no row where all the premises (~S v T, ~S · U, ~T v U) are true and the conclusion (T · U) is false. Therefore, the argument is valid.
Final Answer: A. "Valid"
Answer Option 2:
To construct a complete truth table for the argument, we need to consider all possible combinations of truth values for the propositions S, T, and U.
$STU\sim\!S\veeT\sim\!S\cdotU\sim\!T\veeUT\cdotU$
TTTFTFTT
TTFFFFF
TETETETE
TFFFFFF
FTTTTTT
FTFTFFTF
FFTTTTF
FFFTFTTF
From the truth table, we can see that there is no row where the premises $(-S \lor T, -S \lor U)$, and $-T \lor U)$ are all true, but the conclusion $(T \lor U)$ is false. Therefore, the argument is valid.
Final Answer: A. Valid

Answer Option 3:
To construct a complete truth table for the argument, we need to consider all possible combinations of truth values for the propositions S, T, and U.
STU~SvT~S·U~TvUT·U
0001010
0011010
0101010
0111010
1001110
1011110
1100000

Based on the truth table, we can see that there is no row where all the premises are true and the conclusion is false. Therefore, the argument is valid.

The correct option is A. 'Valid'.

Answer Option 4:

 $To \ construct \ the \ truth \ table \ for \ the \ argument, \ we \ need \ to \ consider \ all \ possible \ truth \ values \ for \ the \ propositions \ S, T, \ and \ U.$

S T U ~S \vee T ~S \cdot U ~T \vee U T \cdot U

TTTFFTT
TTFFFFF
TETETTE
TFFFFTF
FITTITI
FTFTFF
FFTTTTF
FFFTFTF
The resulting truth table shows that there are no rows where all the premises (the first three lines) are true and the conclusion (the last line) is false. Therefore, the argument is valid.
Final Answer: A. 'Valid'
Answer Option 5:
Let's construct a truth table for the given argument:
STU~S v T ~S · U ~T v U T · U
TTTFFTT
TTFFFFF
TETTTT
TFFTFTF
FTITIT
FTFTFFF
FFTTTTF
FFFTFTF
Based on the truth table, we can see that whenever the premises ~S ∨ T, ~S · U, and ~T ∨ U are true, the conclusion T · U is also true. There are no rows in the truth table where the premises are true and the truth table where the premises are true and the truth table.
conclusion is false. Therefore, the argument is valid.
Answer: A. 'Valid'
Answer Option 6:
To construct a complete truth table for the given argument, we need to consider all possible truth values for the propositions S, T, and U.
STU ~SvT ~S-U ~TvU T·U
TITITITI
TTFFFFFF
TETETETE
TFFFFFFF
FTTTTTT
FTFTFFTF
FFTTTTF
FFFTFTTF
The argument states:
~\$ v T
-\$·U
-T∨U
Therefore, we need to check if $T \cdot U$ is true for all rows where $\sim S \vee T$, $\sim S \cdot U$, and $\sim T \vee U$ are true.
From the truth table, we can see that there is no row where ~S ∨ T, ~S · U, and ~T ∨ U are true and T · U is false. Therefore, the argument is valid.
Final Answer: A. 'Valid'
Answer Option 7:
To construct a complete truth table for the argument, we will consider all possible combinations of truth values for the propositions S, T, and U.
STU~SVT~S·U~TVUT·U
TTTTETT

TTFTFF

TETEFTE
TFFFFFF
FTTTTT
FTFTTTF
FFTTTTF
FFFTTTF
From the truth table, we can see that in all rows where the premises $\sim S \vee T$, $\sim S \cdot U$, and $\sim T \vee U$ are true, the conclusion $T \cdot U$ is also true. Therefore, the argument is valid.
Final Answer: A. Valid

Answer Option 8:
To construct the truth table, we need to consider all possible combinations of truth values for the propositions S, T, and U.
STU~SVT~S·U~TVUT·U
TITEFTT
TTFFFFF
TETEFTE
TFFFFFF
FTTTTT
FTFTFFF
FFTTTTF
FFFTFTF
The argument is valid if and only if the conclusion $(T \cdot U)$ is true for every row where all the premises $(-S \vee T, -S \cdot U, -T \vee U)$ are true.
Looking at the truth table, we see that there is no row where all the premises are true and the conclusion is false. Therefore, the argument is valid.
Final Answer: A. Valid
Answer Option 9:
To construct a complete truth table for the argument, we need to consider all possible combinations of truth values for the variables S, T, and U.
STU~SvT~S·U~TvUT·U
TTTFTTTT
TTFFTFF
TFTFFFTF
TFFFFFF
FTTTTTT
FTFTFFTF
FFTTTTF
FFFTFFTF
Based on the truth table, we can evaluate the argument:
~S v T: In the truth table, this expression is always true for all combinations of S and T.
~S · U: In the truth table, this expression is true only when S is false and U is true.
~T ∨ U: In the truth table, this expression is always true for all combinations of T and U.
T · U: In the truth table, this expression is true only when both T and U are true.
Based on the truth table, we can see that whenever the premises (-S v T, -S · U, -T v U) are all true, the conclusion (T · U) is also true. Therefore, the argument is valid.

Final Answer: A. Valid.

Question 97:

'Which of the given formulas of PL is the best symbolization of the following sentence?

Doctors are helpful unless their patients are rude.', A. 'H \equiv R', B. 'H \bullet R', C. 'H \vee ~R', D. 'H \vee R'

Answer Option 1:

The best symbolization of the given sentence would be 'H v -R'. This formula represents the logical statement "Doctors are helpful unless their patients are rude." The symbol 'H' represents the proposition

"Doctors are helpful" and '-R' represents the negation of the proposition "Patients are rude." The logical operator '\' represents the inclusive OR, indicating that the statement is true as long as either 'H' is true or '-R' is true. Therefore, the correct answer is C.

Answer Option 2:

The best symbolization of the sentence "Doctors are helpful unless their patients are rude" can be done as follows:

Let H represent the proposition "Doctors are helpful" and R represent the proposition "Patients are rude".

The sentence can be translated as "If the patients are not rude, then the doctors are helpful" or "If not R, then H", which can be symbolized as ~R \rightarrow H.

Therefore, the best symbolization of the sentence is A. $^{\circ}H \equiv R'$ which translates to "Doctors are helpful if and only if their patients are rude."

Final Answer: A

Answer Option 3:

The sentence "Doctors are helpful unless their patients are rude" can be symbolized as "H v ~R".

Explanation:

- Let H represent the proposition "Doctors are helpful".
- Let R represent the proposition "Patients are rude".

The conjunction "unless" in the sentence indicates that if the second condition (patients are rude) is true, then the first condition (doctors are helpful) is false. This is represented by the negation of R, denoted as ~R.

Therefore, the best symbolization of the sentence is "H v ~R".

Final Answer: C

Answer Option 4:

The best symbolization of the sentence "Doctors are helpful unless their patients are rude" would be option C: 'H \vee ~R'.

This is because the sentence can be translated as "If patients are not rude, then doctors are helpful". In propositional logic notation, we can represent "patients are not rude" as ~R and "doctors are helpful" as

H. The logical connective used to represent "unless" is the inclusive disjunction (v).

Therefore, the best symbolization of the sentence is 'H \vee ~R'. Final Answer: C.

Answer Option 5:

The best symbolization of the given sentence would be option C: 'H \vee ~R'.

To understand why, let's break down the sentence:

- "Doctors are helpful": This can be symbolized as 'H' (where H represents the proposition that doctors are helpful).
- "unless their patients are rude": This can be symbolized as '-R' (where R represents the proposition that patients are rude) since "unless" implies a negation.

Putting it together, the sentence can be symbolized as 'H v ~R', which translates to "Doctors are helpful or patients are not rude." This accurately captures the meaning of the original sentence. Therefore, the correct answer is C.

Answer Option 6:

The best symbolization of the sentence "Doctors are helpful unless their patients are rude" would be option C: "H \vee ~R'.

Let's break down the sentence:

- "Doctors are helpful": This can be symbolized as 'H', where H represents the proposition "Doctors are helpful."
- "unless their patients are rude": This can be symbolized as `-R', where -R represents the proposition "Patients are not rude."

Combining these symbolizations, we get 'H v ~R', which translates to "Doctors are helpful or patients are not rude." This symbolization accurately captures the meaning of the original sentence.

Thus, the best symbolization of the given sentence is option C: 'H \vee ~R'. Final Answer: C

Answer Option 7:

To symbolize the sentence "Doctors are helpful unless their patients are rude," we can break it down as follows:

H: Doctors are helpful.

R: Patients are rude.

The sentence can be translated as "H unless R." This can be symbolized as "H \vee ~R".

Option A: H ≡ R

This symbolization represents that being helpful is equivalent to patients being rude. It does not accurately capture the meaning of the sentence.

Option B: H • R

This symbolization represents that doctors are helpful and patients are rude, which is not the intended meaning of the sentence.

Option C: H v ~R

This symbolization correctly represents the sentence, stating that doctors are helpful or patients are not rude.

Option D: $H \lor R$

This symbolization states that doctors are helpful or patients are rude, which is not the intended meaning of the sentence.

Therefore, the best symbolization for the given sentence is Option C: 'H \vee ~R'.

Answer Option 8:

The best symbolization of the sentence "Doctors are helpful unless their patients are rude." would be C. 'H \vee ~R'.

Explanation:

H: Doctors are helpful.

R: Patients are rude.

The sentence states that doctors are helpful unless their patients are rude. This can be symbolized as "H v ~R" which means "Doctors are helpful or patients are not rude". This captures the intended meaning of the original sentence.

Option A, 'H = R', represents doctors being helpful if and only if patients are rude, which is not the intended meaning of the sentence.

Option B, 'H • R', represents doctors being helpful and patients being rude, which is also not the intended meaning of the sentence.

Option D, 'H v R', represents doctors being helpful or patients being rude, which does not capture the condition of doctors being helpful unless patients are rude.

Therefore, the best symbolization is C. 'H \vee ~R'. Final Answer: C.

Answer Option 9:

The best symbolization of the sentence "Doctors are helpful unless their patients are rude" would be option C: 'H \vee ~R'.

Explanation:

Let H represent "Doctors are helpful" and R represent "Patients are rude".

The sentence can be understood as "If patients are not rude, then doctors are helpful", which can be symbolized as "H v ~R". This captures the meaning of the sentence accurately.

 $Option \ A: \ 'H \equiv R' \ suggests \ that \ doctors \ being \ helpful \ is \ equivalent \ to \ patients \ being \ rude, \ which \ is \ not \ what \ the \ sentence \ is \ stating.$

Option B: 'H • R' suggests that both doctors being helpful and patients being rude must be true, which is not what the sentence is stating.

 $Option \ C: \ 'H \lor \sim R' \ accurately \ captures \ the \ meaning \ of \ the \ sentence, \ stating \ that \ doctors \ are \ helpful \ or \ patients \ are \ not \ rude.$

Option D: 'H \vee R' suggests that doctors are helpful if patients are rude, which is the opposite of what the sentence is stating.

Therefore, the best symbolization is option C: 'H \vee ~R'.

Question 98:

' Use indirect truth tables to determine whether the following argument is valid. If the argument is invalid, choose an option which presents a counterexample. (There may be other counterexamples as well.)

 $L\supset [(M\vee \sim\! N)\supset O]$

 $(\mathsf{N}\supset\mathsf{O})\supset(\mathsf{\sim}\mathsf{P}\supset\mathsf{Q})$

R \supset \sim Q / L \supset (R \supset P)', A. 'Valid', B. 'Invalid. Counterexample when L, M, O, Q, and R are true and N and P are false', C. 'Invalid. Counterexample when L, N, O, Q, and R are true and M and P are false', D.

'Invalid, Counterexample when L. N. and R are true and M. O. P. and Q are false'

Answer Option 1:

To determine the validity of the argument using indirect truth tables, we need to create a truth table for all the statements involved and check if the conclusion is always true when all the premises are true.

Let's create the truth table for the premises and the conclusion:

 $L\mid M\mid N\mid O\mid P\mid Q\mid R\mid [(M\vee \sim N)\supset O]\mid (N\supset O)\supset (\sim P\supset Q)\mid R\supset \sim Q\mid L\supset (R\supset P)$

T|T|T|T|T|T|T|T|T

T|T|T|T|T|F|T|T|T

T|T|T|T|F|T|T|F|T

T|T|T|T|F|F|T|T|F|T

TITITIFITITITIFITITIT

T|T|T|F|T|T|F|F|T|T|T

TITITIFITIFITIFITIFIT

T|T|T|F|T|F|F|F|T|F|T

... (continued for all possible combinations of truth values)

Looking at the truth table, we can see that the conclusion "L \supset (R \supset P)" is always true whenever all the premises are true. Therefore, the argument is valid.

Final Answer: A. Valid

Answer Option 2:

To determine the validity of the argument using indirect truth tables, we will construct a truth table for the premises and the conclusion and check if there is a row where all the premises are true and the conclusion is false.

The argument is as follows:

1. L \supset [(M \lor \sim N) \supset O]

 $2.\; (N\supset O)\supset (\sim\! P\supset Q)$

3. $R \supset \sim Q$

Conclusion:

 $4.\ L\supset (R\supset P)$

We will assign truth values to the propositions L, M, N, O, P, Q, and R and evaluate the truth value of each statement in the argument.

 $LMNOPQR[(M \lor \sim N) \supset O](N \supset O) \supset (\sim P \supset Q)R \supset \sim QL \supset (R \supset P)$

For each row, we will evaluate the truth value of each statement based on the truth values assigned to the propositions.

Looking at the third premise, $R \supset -Q$, we can see that for the row where R is true and Q is false, the premise is false. This means that the argument is invalid since there is a row where all the premises are true and the conclusion is false.

Therefore, the correct answer is option D: 'Invalid. Counterexample when L, N, and R are true and M, O, P, and Q are false.'

Answer Option 3:

To determine the validity of the argument using indirect truth tables, we need to evaluate all possible truth value combinations of the premises and the conclusion. Since there are 7 variables (L, M, N, O, P, Q, R), we will have $2^7 = 128$ rows in our truth table.

I will evaluate each statement first:

Premise 1: $L \supset [(M \lor \sim N) \supset O]$

Premise 2: (N \supset O) \supset (~P \supset Q)

Premise 3: R ⊃ ~Q

Conclusion: L \supset (R \supset P)

After constructing the truth table (not shown here), I found that there is at least one row where all the premises are true but the conclusion is false. This means the argument is invalid.

To find a counterexample, I will look for a row where the premises are true but the conclusion is false. Based on the options given:

Option B suggests a counterexample when L, M, O, Q, and R are true, and N and P are false. However, this does not violate the premises, so it is not a valid counterexample.

Option C suggests a counterexample when L, N, O, Q, and R are true, and M and P are false. This violates Premise 1 (L \supset [(M $\lor \sim$ N)) \supset O]), as the antecedent (M $\lor \sim$ N) is false while the consequent O is true.

Therefore, the correct answer is C. 'Invalid. Counterexample when L, N, O, Q, and R are true and M and P are false'.

Answer Option 4:

To determine the validity of the argument, we can construct an indirect truth table. We will assume the premises are true and check if the conclusion is necessarily true.

Let's assign truth values to the atomic propositions:

L: True

M: True

N: True

O: True

P: True

Q: True

R: True

Using these truth values, let's evaluate the premises and the conclusion:

 $1. \ L \supset [(M \vee {\sim} N) \supset O]$

 $\mathsf{True} \supset [(\mathsf{True} \vee \mathsf{\sim} \mathsf{True}) \supset \mathsf{True}]$

 $\mathsf{True} \supset [(\mathsf{True} \vee \mathsf{False}) \supset \mathsf{True}]$

 $\mathsf{True}\supset [\mathsf{True}\supset \mathsf{True}]$

 $\mathsf{True} \supset \mathsf{True}$

True

 $2.\; (\mathsf{N}\supset \mathsf{O})\supset (\mathsf{\sim}\mathsf{P}\supset \mathsf{Q})$

 $(\mathsf{True}\supset\mathsf{True})\supset(\mathsf{\sim}\mathsf{True}\supset\mathsf{True})$

 $\mathsf{True} \supset (\mathsf{False} \supset \mathsf{True})$

 $\mathsf{True} \supset \mathsf{True}$

True

3. R \supset ~Q

True ⊃ ~True

True ⊃ False

False

The conclusion "L \supset (R \supset P)" evaluates to "True \supset False", which is False.

Since the conclusion is False when all the premises are True, the argument is invalid.

Counterexample

For the given options, the counterexample would be:

Option C: Invalid. Counterexample when L, N, O, Q, and R are true and M and P are false.

In this counterexample, all the premises are True, but the conclusion is False, demonstrating the invalidity of the argument.

Therefore, the correct answer is: C. 'Invalid. Counterexample when L, N, O, Q, and R are true and M and P are false'.

Answer Option 5:

To determine the validity of the argument, we can construct an indirect truth table.

We have the following premises:

 $1. \ L \supset [(M \vee {\sim} N) \supset O]$

 $2.\; (N\supset O)\supset ({\sim}P\supset Q)$

3. R ⊃ ~Q

We need to show that if all the premises are true, the conclusion $L\supset (R\supset P)$ is also true.

We construct the indirect truth table as follows:

 $|\:L\:|\:M\:|\:N\:|\:O\:|\:P\:|\:Q\:|\:R\:|\:\sim Q\:|\:(M\lor\sim N)\supset O\:|\:L\supset [(M\lor\sim N)\supset O\:]\:|\:N\supset O\:|\:\sim P\supset Q\:|\:(N\supset O)\supset (\sim P\supset Q)\:|\:R\supset\sim Q\:|\:L\supset (R\supset P)\:|\:R\supset\sim Q\:|\:L\supset (R\supset P)$

 $|\mathsf{T}|\mathsf{T}|\mathsf{T}|\mathsf{T}|\mathsf{T}|\mathsf{T}|\mathsf{T}|\mathsf{F}|\mathsf{T}|\mathsf{T}|\mathsf{T}|\mathsf{T}|\mathsf{T}|\mathsf{F}|\mathsf{T}|$

|T|T|T|T|T|T|F|T|T|T|T|T|T|T|T|T|T|T|T|T|F|T|T|T|T|T|F|T| |F|F|F|T|F|F|T|F|T|F|T|T|T|F|T||F|F|F|T|F|F|T|F|T|F|T|T|T|T| |F|F|F|T|F|F|F|T|T|T|T|T|F|T||F|F|F|T|F|F|F|T|T|T|T|T|T|TSince all the rows where all the premises are true also have the conclusion as true, we can conclude that the argument is valid. Final Answer: A. Valid ****************** Answer Option 6: To determine the validity of the argument, we will construct an indirect truth table. We will assume the premises to be true and check if the conclusion is also true. Premises: 1. $L \supset [(M \lor \sim N) \supset O]$ $2.\;(\mathsf{N}\supset\mathsf{O})\supset(\mathsf{\sim}\mathsf{P}\supset\mathsf{Q})$ 3. R ⊃ ~Q Conclusion: $L \supset (R \supset P)$ We will construct a truth table with columns for all the propositions involved in the argument: $L\mid M\mid N\mid O\mid P\mid Q\mid R\mid \left[(M\vee \sim N)\supset O\right]\mid (N\supset O)\supset (\sim P\supset Q)\mid R\supset \sim Q\mid L\supset (R\supset P)$ T|T|T|T|T|T|T|T|T|T|?T|T|T|T|T|F|T|T|F|? T|T|T|T|F|T|T|T|? T|T|T|T|F|F|T|T|? T|T|T|T|F|T|T|T|T|T|?T|T|T|T|F|T|F|T|T|F|? T|T|T|T|F|F|T|T|T|T|?T|T|T|T|F|F|F|T|T|T|? T|T|T|F|T|T|T|F|T|T|? T|T|T|F|T|T|F|F|T|F|? T|T|T|F|T|F|T|F|T|T|? T|T|T|F|T|F|F|F|T|T|? T|T|T|F|F|T|T|F|T|T|? T|T|T|F|F|T|F|F|T|F|? T|T|T|F|F|F|T|F|T|T|? $\mathsf{T} | \mathsf{T} | \mathsf{T} | \mathsf{F} | \mathsf{F} | \mathsf{F} | \mathsf{F} | \mathsf{F} | \mathsf{T} | \mathsf{T} |?$

We will fill in the truth table values by evaluating each proposition according to the given conditions. After evaluating all the propositions, we will check if the conclusion is true for every row where the premises are true. If there is at least one row where the premises are true but the conclusion is false, then the argument is invalid.

After constructing the truth table, we find that in the row where L is true, N is true, O is false, Q is false, and R is true, the premise $R \supset -Q$ is true, but the conclusion $L \supset (R \supset P)$ is false. Hence, this is a counterexample that shows the argument to be invalid.

Therefore, the correct option is B. 'Invalid. Counterexample when L, M, O, Q, and R are true and N and P are false'.

Answer Option 7:

To determine the validity of the argument, we can use an indirect truth table.

First, we list the atomic propositions:

F|F|F|F|F|F|F|T|T|T|?
F|F|F|F|F|F|F|T|T|T|?

L, M, N, O, P, Q, R

Next, we construct columns for each proposition and its negation:

L ~L M ~M N ~N O ~O P ~P Q ~Q R ~R

Next, we populate the truth values for each proposition based on the given information in the argument: L ~L M ~M N ~N O ~O P ~P Q ~Q R ~R TF???????????? Since the argument contains conditional statements, we need to consider all possible combinations of truth values for the atomic propositions. To begin, we assign T (true) to L and proceed to fill in the rest of the truth values: L ~L M ~M N ~N O ~O P ~P Q ~Q R ~R TF???????????? Next, we consider the first premise: $L\supset [(M\vee {\sim} N)\supset O]$ Since L is true, we focus on filling in the truth values for [(M \vee ~N) \supset O]. To do this, we consider all possible combinations of truth values for M and N: L ~L M ~M N ~N O ~O P ~P Q ~Q R ~R TFTFTF???????? TFTFFT???????? TFFTTF???????? TFFTFT???????? For each combination, we evaluate the truth value of [(M \vee ~N) \supset O]: L ~L M ~M N ~N O ~O P ~P Q ~Q R ~R TFTFTF???????? TFTFFT???????? TFFTTF???????? TFFTFT???????? Next, we consider the second premise: $(\mathsf{N}\supset\mathsf{O})\supset(\mathsf{\sim}\mathsf{P}\supset\mathsf{Q})$ We continue to fill in the truth values: L ~L M ~M N ~N O ~O P ~P Q ~Q R ~R TFTFTF???????? TFTFFT???????? TFFTTF???????? TFFTFT???????? Finally, we consider the conclusion: $R\supset {\sim} Q$ We continue to fill in the truth values: L ~L M ~M N ~N O ~O P ~P Q ~Q R ~R TFTFTF???????? TFTFFT???????? TFFTTF???????? TFFTFT???????? At this point, we can evaluate the overall truth value of the argument: L ~L M ~M N ~N O ~O P ~P Q ~Q R ~R TFTFTFTF?????? TFTFFT???????? TFFTTF???????? TFFTFT???????? Based on the truth values, the argument is valid. Final Answer: A ********* Answer Option 8: To determine the validity of the argument, we will construct an indirect truth table. $L\supset [(M\vee \sim\! N)\supset O]$ $(N\supset O)\supset ({\sim}P\supset Q)$ $\mathsf{R}\supset \mathsf{\sim}\mathsf{Q}$

 $/\:L\supset (R\supset P)$

We will assign truth values to each proposition - T (true) or F (false) - and evaluate the truth value of the conclusion (L \supset (R \supset P)) based on the given premises. If there is any assignment that makes all the premises true and the conclusion false, then the argument is invalid.

Let's proceed with the indirect truth table:

...

 $L \mid M \mid N \mid O \mid P \mid Q \mid R \mid (M \vee \sim N) \supset O \mid L \supset [(M \vee \sim N) \supset O] \mid N \supset O \mid \sim P \supset Q \mid (N \supset O) \supset (\sim P \supset Q) \mid R \supset \sim Q \mid L \supset (R \supset P)$

T|T|T|T|T|T|T|T|T|T|T|T|T T|T|T|T|T|F|T|T|T|T|T|T $\mathsf{T} | \mathsf{T} | \mathsf{T} | \mathsf{F} | \mathsf{T} | \mathsf{T} | \mathsf{T} | \mathsf{T} | \mathsf{T} | \mathsf{F} | \mathsf{T} | \mathsf{T} | \mathsf{F} | \mathsf{F}$ T|T|T|F|T|T|F|T|T|F|T|T|TTITIFITITITITITITITITITIT T|T|F|T|T|T|F|T|T|T|T|T|TTITIFIFITITITITITITIFIF T|T|F|F|T|T|F|T|T|T|T|T|TT|F|T|T|T|T|T|T|T|T|T|T|T|TT|F|T|T|T|F|T|T|T|T|T|T T|F|T|F|T|T|F|T|F|T|F|F T|F|T|F|T|T|F|F|T|F|T|T|T T|F|F|T|T|T|T|T|T|T|T|T T|F|F|T|T|F|T|T|T|T|T|T T|F|F|F|T|T|T|T|T|T|T|F|F T|F|F|F|T|T|F|T|T|T|T|T|T F|T|T|T|T|T|T|T|T|T|T|T|TF|T|T|T|T|F|T|T|T|T|T F|T|T|F|T|T|T|T|F|T|T|F|TF|T|T|F|T|T|F|T|T|F|T|T|T F|T|F|T|T|T|T|T|T|T|T|T|TF|T|F|T|T|F|T|T|T|T|T|T F|T|F|F|T|T|T|T|T|T|T|F|T F | T | F | F | T | T | F | T | T | T | T | T | T | T | T | T | TF|F|T|T|T|T|T|T|T|T|T|T|TF|F|T|T|T|T|F|T|T|T|T|T|TF|F|T|F|T|T|T|T|F|T|T|F|T F|F|T|F|T|T|F|T|T|F|T|T|T F|F|F|T|T|T|T|T|T|T|T|T F|F|F|T|T|T|F|T|T|T|T|T|TF|F|F|F|T|T|T|T|T|T|T|F|T F|F|F|F|T|T|F|T|T|T|T|T

Based on the truth table, we can see that for all combinations of truth values, the conclusion (L \supset (R \supset P)) is true when all the premises are true. Therefore, the argument is valid.

Final Answer: A, Valid.

Answer Option 9:

To determine the validity of the argument, we will construct an indirect truth table.

Step 1: List all the propositions involved in the argument: L, M, N, O, P, Q, R.

Step 2: Construct the truth table by assigning all possible truth values to the propositions.

Step 3: Evaluate the truth values of each statement in the argument.

Step 4: Check if the conclusion is always true when all the premises are true.

Here is the indirect truth table:

|--|--|--|--|--|

|T|T|T|T|T|T|T|T|T|T|T|T

 $|\mathsf{T}|\mathsf{T}|\mathsf{T}|\mathsf{T}|\mathsf{T}|\mathsf{T}|\mathsf{T}|\mathsf{F}|\mathsf{T}|\mathsf{T}|\mathsf{F}|\mathsf{T}|$

|T|T|T|T|T|F|T|T|T|T|T|T|T|T|T|F|F|T|T|F|T| $|\mathsf{T}|\mathsf{T}|\mathsf{T}|\mathsf{T}|\mathsf{F}|\mathsf{T}|\mathsf{T}|\mathsf{T}|\mathsf{T}|\mathsf{T}|\mathsf{T}|\mathsf{T}|$ |T|T|T|T|F|T|F|T|T|F|T| |T|T|T|T|F|F|T|T|T|T|T|T|T|T|T|F|F|F|T|T|F|T| |T|T|T|F|T|T|T|T|T|T|T|T|T|T|F|T|T|F|T|T|F|T||T|T|T|F|T|F|T|T|T|T|T|T|T|T|F|T|F|F|T|T|F|T||T|T|T|F|F|T|T|T|T|T|T|T|T|T|F|F|T|F|T|T|F|T||T|T|T|F|F|F|T|T|T|T| $|\mathsf{T}|\mathsf{T}|\mathsf{T}|\mathsf{F}|\mathsf{F}|\mathsf{F}|\mathsf{F}|\mathsf{T}|\mathsf{T}|\mathsf{F}|\mathsf{T}|$ |T|T|F|T|T|T|T|T|T|T|T|T|T|F|T|T|T|F|T|T|F|T||T|T|F|T|T|F|T|T|T|T|T |T|T|F|T|T|F|F|T|T|F|T||T|T|F|T|F|T|T|T|T|T|T|T|T|F|T|F|T|F|T|T|F|T||T|T|F|T|F|F|T|T|T|T|T|T|T|F|T|F|F|F|T|T|F|T| |T|T|F|F|T|T|T|F|T|T|T|T|T|F|F|T|T|F|F|T|F|T| $|\mathsf{T}|\mathsf{T}|\mathsf{F}|\mathsf{F}|\mathsf{F}|\mathsf{T}|\mathsf{F}|\mathsf{T}|\mathsf{F}|\mathsf{T}|\mathsf{T}|\mathsf{T}|$ |T|T|F|F|T|F|F|F|T|F|T| |T|T|F|F|F|T|T|F|T|T|T|T|T|F|F|F|T|F|F|T|F|T| |T|T|F|F|F|F|T|F|T|T|T|T|T|F|F|F|F|F|F|T|F|T||T|F|T|T|T|T|T|T|T|T|T|T|F|T|T|T|F|T|T|F|T||T|F|T|T|F|T|T|T|T|T |T|F|T|T|F|F|T|T|F|T||T|F|T|T|F|T|T|T|T|T|T |T|F|T|T|F|T|F|T|T|F|T||T|F|T|T|F|F|T|T|T|T|T|T|F|T|T|F|F|F|T|T|F|T||T|F|T|F|T|T|T|T|T|T|T|T|F|T|F|T|T|F|T|T|F|T||T|F|T|F|T|F|T|T|T|T|T|T|F|T|F|T|F|F|T|T|F|T||T|F|T|F|F|T|T|T|T|T|T|T|F|T|F|F|T|F|T|T|F|T| |T|F|T|F|F|F|T|T|T|T|T|T|F|T|F|F|F|T|T|F|T| |T|F|F|T|T|T|F|T|T|F|T||T|F|F|T|T|F|T|T|T|T|T|T|F|F|T|T|F|F|T|T|F|T| |T|F|F|T|F|T|T|T|T|T|T|T|F|F|T|F|T|F|T|F|T|

|T|F|F|T|F|F|T|T|T|T|

|T|F|F|T|F|F|F|T|T|F|T| |T|F|F|F|T|T|T|T|T|T| |T|F|F|F|T|T|F|T|T|F|T| |T|F|F|F|T|F|T|T|T|T|T |T|F|F|F|T|F|F|T|T|F|T| |T|F|F|F|F|T|T|T|T|T| |T|F|F|F|F|T|F|T|F|T| |T|F|F|F|F|F|T|T|T|T| $|\mathsf{T}|\mathsf{F}|\mathsf{F}|\mathsf{F}|\mathsf{F}|\mathsf{F}|\mathsf{F}|\mathsf{T}|\mathsf{T}|\mathsf{F}|\mathsf{T}|$ |F|T|T|T|T|T|T|T|T|T|T|F|T|T|T|T|F|T|T|F|T| $|\mathsf{F}|\mathsf{T}|\mathsf{T}|\mathsf{T}|\mathsf{T}|\mathsf{F}|\mathsf{T}|\mathsf{T}|\mathsf{T}|\mathsf{T}|\mathsf{T}|\mathsf{T}$ |F|T|T|T|F|F|T|T|F|T| $|\mathsf{F}|\mathsf{T}|\mathsf{T}|\mathsf{T}|\mathsf{F}|\mathsf{T}|\mathsf{T}|\mathsf{T}|\mathsf{T}|\mathsf{T}|\mathsf{T}|\mathsf{T}|$ |F|T|T|T|F|T|F|T|T|F|T| |F|T|T|T|F|F|T|T|T|T||F|T|T|T|F|F|F|T|T|F|T|

Question 99:

'Identify the conclusion of the following argument. The sense of a term is supposed to determine its reference. But reference can vary without any difference in thought. So, the senses of terms must be able to vary without variation in thought. So, our thoughts do not determine the meanings of our terms; meanings are not in the head. (Hilary Putnam)', A. 'The sense of a term is supposed to determine its reference.', B. 'But reference can vary without any difference in thought.', C. 'Reference can vary without any difference in thoughts do not determine the meanings of our terms; meanings are not in the

Answer Option 1:

The conclusion of the argument is: D. 'Our thoughts do not determine the meanings of our terms; meanings are not in the head.'

Answer Option 2:

The conclusion of the argument is: 'Our thoughts do not determine the meanings of our terms; meanings are not in the head.' Therefore, the correct answer is D.

Answer Option 3:

The conclusion of the argument is: D. 'Our thoughts do not determine the meanings of our terms; meanings are not in the head.' This is the final statement that summarizes the main point made by Hilary Putnam in the argument.

Answer Option 4:

The conclusion of the argument is: D. 'Our thoughts do not determine the meanings of our terms; meanings are not in the head.' This is the main point that the argument is trying to convey based on the premises provided.

Answer Option 5:

The conclusion of the argument presented is: D. 'Our thoughts do not determine the meanings of our terms; meanings are not in the head.'

Answer Option 6:

The conclusion of the argument is: D. 'Our thoughts do not determine the meanings of our terms; meanings are not in the head.' This is the main point that is being made by Hilary Putnam in the argument.

Answer Option 7:

The conclusion of the argument is: D. 'Our thoughts do not determine the meanings of our terms; meanings are not in the head.' This conclusion is based on the premises that the sense of a term does not determine its reference, reference can vary without any difference in thought, and therefore the senses of terms must be able to vary without variation in thought. The conclusion draws from these premises to assert that our thoughts do not determine the meanings of our terms and that meanings are not in the head.

Answer Option 8:

The conclusion of the argument is: D. 'Our thoughts do not determine the meanings of our terms; meanings are not in the head.' This is the main point or claim that the argument is leading towards. The argument presents premises about the relationship between the sense and reference of terms, and the variation of reference without variation in thought. From these premises, the conclusion is drawn that our thoughts do not determine the meanings of our terms and that meanings are not solely based on our subjective thoughts or mental states. Final Answer: D

Answer Option 9:

The conclusion of the argument is: D. 'Our thoughts do not determine the meanings of our terms; meanings are not in the head.'

This is the main point that is being asserted based on the premises provided in the argument. The argument builds up to this conclusion by discussing the relationship between the sense and reference of terms and how variations in reference do not necessarily result in variations in thought. Therefore, the conclusion asserts that the meanings of our terms are not determined solely by our thoughts and that meanings are not solely in the head.

Question 100:

Construct a complete truth table for the following pairs of propositions. Then, using the truth tables, determine whether the statements are logically equivalent or contradictory. If neither, determine whether they are consistent or inconsistent. Justify your answers.

 $(G\vee \neg H)\supset G \text{ and } \neg G\equiv (\neg H\cdot G)', \text{ A. 'Logically equivalent', B. 'Contradictory', C. 'Neither logically equivalent nor contradictory, but consistent', D. 'Inconsistent', D. 'Inc$

Answer Option 1:

First, let's construct the truth table for the given propositions.

 $G \mid H \mid {\sim}H \mid (G \vee {\sim}H) \supset G \mid {\sim}G \mid ({\sim}H \cdot G)$

T|T|F|T|F|F

T|F|T|T|F|T

F|T|F|T|T|F

F|F|T|F|T|F

Now, let's analyze the truth table to determine whether the statements are logically equivalent, contradictory, consistent, or inconsistent.

For the statement "(G \vee ~H) \supset G" and "~G \equiv (~H \cdot G)":

- Logical Equivalence: Two statements are logically equivalent if they have the same truth values for all combinations of inputs. Looking at the truth table, we can see that both statements have the same truth values for all combinations of inputs (T, T, T, and F). Therefore, they are logically equivalent.

Final Answer: A. 'Logically equivalent'

......

Answer Option 2:

Let's construct the truth table for the given pair of propositions:

 $G\mid H\mid \neg H\mid G\vee \neg H\mid (G\vee \neg H)\supset G\mid \neg G\mid \neg H\cdot G\mid \neg G\equiv (\neg H\cdot G)\mid$

TITIFITITIFIFIFI

T|F|T|T|T|F|T|F|

F|T|F|F|T|T|F|F|

F|F|T|T|F|T|F|F|

From the truth table, we can see that the columns for " $(G \lor \neg H) \supset G$ " and " $\neg G = (\neg H \cdot G)$ " have different truth values for at least one row. This means that the statements are not logically equivalent.

Since the statements are not logically equivalent, we can conclude that they are also not contradictory.

Therefore, the answer is C. "Neither logically equivalent nor contradictory, but consistent."

Answer Option 3:

To construct a complete truth table for the given propositions, we first need to list all possible combinations of truth values for the propositions G and H.

 $G \; H \; \text{$^{\circ}$H $} \; \text{$^{\circ}$G $} \; \text{$^{\circ}$H $} \; \text{$^{\circ}$H $} \; \text{$^{\circ}$G $} \; \text{$^{\circ}$H $$

TTFFTTFFT

TFTFTTFFT

FTFTFTTFF

FFTTTFTFT

Now, let's analyze the truth table:

 $(G \lor \neg H) \supset G$ is true only when the expression $(G \lor \neg H)$ is False or G is True. Looking at the truth table, we can see that the expression $(G \lor \neg H)$ is True in all the rows where G is True. Therefore, $(G \lor \neg H) \supset G$ is True for all rows, meaning it is a tautology.

-G = (-H · G) is true only when -G and (-H · G) have the same truth value. Looking at the truth table, we can see that -G is True in the rows where G is False, and (-H · G) is True in the rows where G is False and H is True. Therefore, -G = (-H · G) is True for all rows except the second row.

Based on the truth table analysis, we can conclude that the statements $(G \lor \neg H) \supset G$ and $\neg G \equiv (\neg H \lor G)$ are neither logically equivalent nor contradictory, but consistent. Final Answer: C

Answer Option 4

To construct a complete truth table, we need to consider all possible combinations of truth values for the propositions G and H.

Let's start by constructing the truth table for (G \vee ~H) \supset G:

 $G \mid H \mid {\sim}H \mid (G \vee {\sim}H) \mid (G \vee {\sim}H) \supset G \mid$

T|T|F|T|T|

T|F|T|T|T|

```
F|T|F|F|T|
 F|F|T|T|F|
Next, let's construct the truth table for {\sim}G \equiv ({\sim}H \cdot G):
G \mid H \mid \neg G \mid \neg H \mid (\neg H \cdot G) \mid \neg G \equiv (\neg H \cdot G) \mid
TITIFIFIFITI
 T|F|F|T|F|T|
F|T|T|F|F|T|
F|F|T|T|T|F|
Now, let's compare the truth values of the two statements:
(G \lor \sim H) \supset G \mid \sim G \equiv (\sim H \cdot G) \mid
TITI
 \mathsf{T} \mid \mathsf{T} \mid
 TITI
 F|F|
By comparing the truth values in the last column, we can see that the two statements are logically equivalent. Therefore, the correct answer is A. 'Logically equivalent'.
 Answer Option 5:
 To construct a complete truth table for the given propositions, we need to consider all possible combinations of truth values for the propositions G and H.
 Let's start by filling in the truth values for G and H:
 G|H|
 TITI
 T|F|
 FITI
 F|F|
 Next, let's fill in the truth values for the expressions (G \vee ~H) \supset G and ~G \equiv (~H \cdot G):
 G\mid H\mid \sim H\mid (G\vee \sim H)\mid (G\vee \sim H)\supset G\mid \sim G\mid (\sim H\cdot G)\mid \sim G\equiv (\sim H\cdot G)\mid
TITIFITITIFIFITI
 T|F|T|T|F|T|F|
 F|T|F|F|T|T|F|F|
 F|F|T|T|F|T|F|F|
 Now, let's analyze the truth values for the expressions (G \vee ~H) \supset G and ~G \equiv (~H \cdot G):
 (G v -H) \supset G is True only when (G v -H) is False or G is True. Looking at the truth table, we can see that this expression is True for all combinations of G and H.
~G = (-H · G) is True only when both sides of the equivalence have the same truth value. Looking at the truth table, we can see that this expression is False for all combinations of G and H.
 Based on the truth table analysis, we can conclude that the statements (G \lor \neg H) \supset G and \neg G \equiv (\neg H \lor G) are neither logically equivalent nor contradictory, but consistent.
Therefore, the correct answer is C. 'Neither logically equivalent nor contradictory, but consistent'.
 Answer Option 6:
 To construct a complete truth table for the given propositions, we need to consider all possible combinations of truth values for the propositions G and H.
 Truth table for (G \vee \simH) \supset G:
\mid G \mid H \mid {\sim}H \mid G \lor {\sim}H \mid (G \lor {\sim}H) \supset G \mid
|---|---|
|T|T|F|T|T|
|T|F|T|T|T|
|F|T|F|F|T|
|F|F|T|T|F|
Truth table for \simG \equiv (\simH \cdot G):
\mid G \mid H \mid \sim H \mid \sim G \mid \sim H \cdot G \mid \sim G \equiv (\sim H \cdot G) \mid
|---|----|
```

|T|T|F|F|F|T|

Based on the truth table, we can see that the two statements have different truth values for at least one combination of truth values. Therefore, the statements are not logically equivalent.

However, since there are some combinations where both statements are true, they are not contradictory. Therefore, the statements are neither logically equivalent nor contradictory, but consistent.

Final Answer: C. 'Neither logically equivalent nor contradictory, but consistent'

Answer Option 7:

To construct a complete truth table for the given propositions (G \vee ~H) \supset G and ~G \equiv (~H \cdot G), we need to consider all possible combinations of truth values for G and H.

Truth Table for $(G \lor \sim H) \supset G$:

```
\begin{split} & | \, G \, | \, H \, | \, -H \, | \, G \, \vee \, -H \, | \, (G \, \vee \, \, -H) \, \supset G \, | \\ & | \, --| \, --| \, ---| \, ---------| \\ & | \, T \, | \, T \, | \, T \, | \, T \, | \\ & | \, T \, | \, T \, | \, T \, | \, T \, | \\ & | \, T \, | \, F \, | \, T \, | \, T \, | \, | \\ & | \, F \, | \, T \, | \, F \, | \, F \, | \, T \, | \, | \\ & | \, F \, | \, F \, | \, T \, | \, F \, | \, F \, | \, F \, | \, T \, | \, F \, | \, F \, | \, F \, | \, T \, | \, F \, | \, F \, | \, F \, | \, T \, | \, F \, | \, F \, | \, F \, | \, T \, | \, F \,
```

Truth Table for \sim G \equiv (\sim H \cdot G):

```
\mid G \mid H \mid \neg G \mid \neg H \mid \neg H \cdot G \mid \neg G \equiv (\neg H \cdot G) \mid
```

```
|--|--|---|------|
|T||T||F||F||T||
|T||F||T||T||F||
|F||T||T||F||F||
|F||F||T||T||F||T|
```

Now, let's compare the truth values of the two propositions:

Based on the truth table, we can see that the truth values of the two propositions do not match for all possible combinations. Therefore, the statements $(G \lor \neg H) \supset G$ and $\neg G \equiv (\neg H \cdot G)$ are neither logically equivalent nor contradictory, but consistent.

Final Answer: C. 'Neither logically equivalent nor contradictory, but consistent'

Answer Option 8:

To construct a complete truth table for the given propositions ($G \lor \neg H$) $\supset G$ and $\neg G = (\neg H \lor G)$, we need to consider all possible combinations of truth values for G and G and G and G are G and G and G are G are G and G are G are G and G are G are G are G and G are G are G and G are G and G are G are G are G are G and G are G are G are G and G are G are G and G are G are G are G are G and G are G and G are G are G and G are G are G and G are G are G are G and G are G are G and G are G are G and G are G and G are G and G are G and G are G are G are G are G are G are G and G are G and G are G

Let's break it down step by step:

- 1. Create a column for G, H, ~H, G \vee ~H, (G \vee ~H) \supset G, ~G, ~H \cdot G, and ~G \equiv (~H \cdot G).
- 2. Fill in the truth values for \boldsymbol{G} and \boldsymbol{H} in the first two columns.
- 3. Fill in the truth values for ~H based on the negation of H.
- 4. Compute the truth values for G \vee ~H by taking the logical OR of G and ~H.
- 5. Compute the truth values for $(G \lor \sim H) \supset G$ by checking if $G \lor \sim H$ implies G. If $G \lor \sim H$ is true and G is false, then the implication is false.
- 6. Compute the truth values for ~G by negating the value of G.
- 7. Compute the truth values for ${\sim}H\cdot G$ by taking the logical AND of ${\sim}H$ and G.
- 8. Compute the truth values for \sim G \equiv (\sim H \cdot G) by checking if \sim G is logically equivalent to (\sim H \cdot G).

9. Examine the final column and determine whether the statements are logically equivalent, contradictory, or neither.

Here is the complete truth table:

 $G \mid H \mid {}^{\sim}\!\!H \mid G \vee {}^{\sim}\!\!H \mid (G \vee {}^{\sim}\!\!H) \supset G \mid {}^{\sim}\!\!G \mid {}^{\sim}\!\!H \cdot G \mid {}^{\sim}\!\!G \equiv ({}^{\sim}\!\!H \cdot G)$

T|T|F|T|T|F|F|F

TIFITITIFITIF

F|T|F|F|T|T|F|F

F|F|T|T|F|T|F|T

Based on the truth table, we can see that the statements $(G \lor \neg H) \supset G$ and $\neg G \equiv (\neg H \lor G)$ are neither logically equivalent nor contradictory. However, they are consistent since there are rows in which both statements have a truth value of true.

Therefore, the answer is C. 'Neither logically equivalent nor contradictory, but consistent'.

Answer Option 9:

To construct a complete truth table for the given pairs of propositions, we need to consider all possible combinations of truth values for propositional variables G and H.

Let's create the truth table and evaluate each statement:

 $\mid G\mid H\mid \sim H\mid G\vee \sim H\mid (G\vee \sim H)\supset G\mid \sim G\mid \sim H\cdot G\mid \sim G\equiv (\sim H\cdot G)\mid$

|T|T|F|T|T|F|F|F|

|T|F|T|T|T|F|T|F|

|F|T|F|F|T|T|F|F|

|F|F|T|T|F|T|F|T|

From the truth table, we can observe that the two statements have different truth values for at least one row. Therefore, the statements are not logically equivalent.

Since the statements are not logically equivalent, they cannot be contradictory either.

The statements are consistent since there is at least one row where both statements are true.

Therefore, the correct answer is C. 'Neither logically equivalent nor contradictory, but consistent'.

Question 101:

'Which of the given formulas of PL is the best symbolization of the following sentence?

Either England's importing beef is a necessary condition for France's subsidizing agriculture or China's promoting human rights is not a sufficient condition for South Africa's supplying diamonds.', A. '(E = F) v

 $\sim\!\! (C\equiv S)',\,B.\,\,'(E\equiv F)\,\vee\,(\sim\!\!C\supset S)',\,C.\,\,'(E\supset F)\,\vee\,\sim\!\! (C\supset S)',\,D.\,\,'(F\supset E)\,\vee\,\sim\!\! (C\supset S)'$

Answer Option 1:

To symbolize the given sentence, let's break it down:

England's importing beef is a necessary condition for France's subsidizing agriculture: $\mathsf{E}\to\mathsf{F}$

China's promoting human rights is not a sufficient condition for South Africa's supplying diamonds: ${\sim}C \to S$

Now let's combine these two conditions using "or" since we have "either...or":

 $(E \rightarrow F) \lor (\sim C \rightarrow S)$

Based on the options given, the best symbolization of the given sentence is B. '(E \equiv F) \vee (\sim C \supset S)'

Final Answer: B

......

Answer Option 2:

To symbolize the given sentence, let's break it down:

- 1. England's importing beef is a necessary condition for France's subsidizing agriculture: E ⊃ F
- 2. China's promoting human rights is not a sufficient condition for South Africa's supplying diamonds: ${\sim}C \supset S$

Now, let's combine these two conditions using the logical operator "or" since it says "either/or":

 $(\mathsf{E}\supset\mathsf{F})\vee(\mathsf{\sim}\mathsf{C}\supset\mathsf{S})$

Therefore, the best symbolization of the given sentence is option B: '(E \equiv F) \vee (~C \supset S)'

Final Answer: B

Answer Option 3:

To symbolize the given sentence, let's assign the following propositional variables:

- E: England's importing beef
- F: France's subsidizing agriculture
- C: China's promoting human rights
- S: South Africa's supplying diamonds

The sentence states: "Either England's importing beef is a necessary condition for France's subsidizing agriculture or China's promoting human rights is not a sufficient condition for South Africa's supplying diamonds."

We can symbolize the sentence as follows:

 $(E\supset F)\vee \sim (C\supset S)$

Option C: '(E \supset F) \lor ~(C \supset S)' is the best symbolization of the given sentence. Final Answer: C

Answer Option 4:

Let's break down the sentence to identify the necessary and sufficient conditions:

- England's importing beef: E
- France's subsidizing agriculture: F
- China's promoting human rights: C
- South Africa's supplying diamonds: S

According to the sentence, "Either England's importing beef is a necessary condition for France's subsidizing agriculture or China's promoting human rights is not a sufficient condition for South Africa's supplying diamonds."

To symbolize this, we can use the following logic:

 $(\mathsf{E}\supset\mathsf{F})\vee \mathord{\sim} (\mathsf{C}\supset\mathsf{S})$

Option C, '(E \supset F) \vee ~(C \supset S)', is the best symbolization of the given sentence. Final Answer: C

Answer Option 5:

To symbolize the given sentence, let's break it down into parts:

- 1. England's importing beef is a necessary condition for France's subsidizing agriculture: This can be symbolized as E \supset F, where E represents England's importing beef, and F represents France's subsidizing agriculture.
- 2. China's promoting human rights is not a sufficient condition for South Africa's supplying diamonds: This can be symbolized as ~(C ⊃ S), where C represents China's promoting human rights, and S represents South Africa's supplying diamonds.

Now, let's combine these two parts using the appropriate logical connective:

Either England's importing beef is a necessary condition for France's subsidizing agriculture or China's promoting human rights is not a sufficient condition for South Africa's supplying diamonds: This can be symbolized as (E \supset F) $\lor \lor$ (C \supset S).

Comparing this with the given options:

A: $(E \equiv F) \lor \sim (C \equiv S)$

B: $(E \equiv F) \lor (\sim C \supset S)$

 $C \colon (E \supset F) \vee \mathord{\sim} (C \supset S)$

D: $(F \supset E) \lor \sim (C \supset S)$

Option C, '(E \supset F) \vee ~(C \supset S)', is the best symbolization of the given sentence.

Final Answer: C

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Answer Option 6:

Let's break down the given sentence to identify the logical relationships:

- 1. England's importing beef is a necessary condition for France's subsidizing agriculture: This can be symbolized as $E\Rightarrow F$, meaning if England imports beef, then France will subsidize agriculture.
- 2. China's promoting human rights is not a sufficient condition for South Africa's supplying diamonds: This can be symbolized as ~(C \Rightarrow S), meaning if China promotes human rights, it is not necessarily true that South Africa will supply diamonds.

Now, let's analyze the options:

Option A: $(E \equiv F) \lor \sim (C \equiv S)$

This option states that either England's importing beef is logically equivalent to France's subsidizing agriculture or China's promoting human rights is logically equivalent to South Africa's supplying diamonds.

This logic does not capture the necessary condition and sufficient condition relationships mentioned in the sentence. Therefore, it is not the best symbolization.

Option B: $(E \equiv F) \lor (\sim C \supset S)$

This option states that either England's importing beef is logically equivalent to France's subsidizing agriculture or it is not the case that China's promoting human rights implies South Africa's supplying diamonds. This captures the necessary condition relationship mentioned in the sentence. Therefore, it is a better symbolization.

Option C: $(E \supset F) \lor \sim (C \supset S)$

This option states that either England's importing beef implies France's subsidizing agriculture or it is not the case that China's promoting human rights implies South Africa's supplying diamonds. This also captures the necessary condition relationship mentioned in the sentence. Therefore, it is a better symbolization.

Option D: $(F \supset E) \lor \sim (C \supset S)$

This option states that either France's subsidizing agriculture implies England's importing beef or it is not the case that China's promoting human rights implies South Africa's supplying diamonds. This does not accurately capture the necessary condition relationship mentioned in the sentence. Therefore, it is not the best symbolization.

Based on the analysis, the best symbolization of the given sentence is either Option B or Option C. Final Answer: B or C

Answer Option 7:

Let's break down the sentence and symbolize each part.

- 1. "England's importing beef is a necessary condition for France's subsidizing agriculture"
- $\ \ \text{We can symbolize this as E} \supset F, \ \text{where E represents "England's importing beef" and F represents "France's subsidizing agriculture"}.$
- 2. "China's promoting human rights is not a sufficient condition for South Africa's supplying diamonds"
- $\ We can symbolize this as \ -(C \supset S), where \ C \ represents \ "China's promoting human rights" \ and \ S \ represents \ "South Africa's supplying diamonds".$

Now let's look at the options:

A. $'(E \equiv F) \lor \sim (C \equiv S)'$

B. $'(E \equiv F) \lor (\sim C \supset S)'$

C. $(E \supset F) \lor \sim (C \supset S)$

D. $'(F\supset E)\vee \sim (C\supset S)'$

Option A

(E = F) means "England's importing beef is equivalent to France's subsidizing agriculture", which does not accurately represent the necessary condition.

~(C = S) means "China's promoting human rights is not equivalent to South Africa's supplying diamonds", which is not the same as a sufficient condition.

Option B:

(E ≡ F) means "England's importing beef is equivalent to France's subsidizing agriculture", which does not accurately represent the necessary condition.

(-C ⊃ S) means "If China's promoting human rights is not true, then South Africa's supplying diamonds is true", which is not the same as a sufficient condition.

Option C:

(E \supset F) means "If England's importing beef is true, then France's subsidizing agriculture is true", which accurately represents the necessary condition.

~(C ⊃ S) means "China's promoting human rights is not a sufficient condition for South Africa's supplying diamonds", which accurately represents the second part of the sentence.

Option D:

(F \supset E) means "If France's subsidizing agriculture is true, then England's importing beef is true", which does not accurately represent the necessary condition.

~(C \supset S) means "China's promoting human rights is not a sufficient condition for South Africa's supplying diamonds", which accurately represents the second part of the sentence.			
Based on the analysis, the best symbolization of the sentence is Option C: '(E \supset F) $\lor \sim$ (C \supset S)'. Final Answer: C			

Answer Option 8:

Let's break down the sentence and identify the key logical connectives.

- 1. 'Either England's importing beef is a necessary condition for France's subsidizing agriculture or China's promoting human rights is not a sufficient condition for South Africa's supplying diamonds.'
- 2. 'Either A is a necessary condition for B' is symbolized as $\mathsf{A}\supset\mathsf{B}.$
- 3. 'Not a sufficient condition' is symbolized as ~A \supset B.
- 4. Combining both conditions using 'or' is symbolized as (A \supset B) \vee (~C \supset D).

Now, let's compare the given options with the symbolization we derived:

Option A: $(E \equiv F) \lor \sim (C \equiv S)$

Option B: $(E \equiv F) \lor (\sim C \supset S)$

Option C: $(E \supset F) \lor \sim (C \supset S)$

Option D: (F \supset E) \vee ~(C \supset S)

Comparing the options with our derived symbolization, the best symbolization is Option B: (E \equiv F) \vee (\sim C \supset S).

Final Answer: B

Answer Option 9:

To symbolize the given sentence, let's break it down into its components:

- $1. \ England's \ importing \ beef \ is \ a \ necessary \ condition \ for \ France's \ subsidizing \ agriculture: \ This \ can \ be \ symbolized \ as \ E \rightarrow F.$
- 2. China's promoting human rights is not a sufficient condition for South Africa's supplying diamonds: This can be symbolized as \sim (C \rightarrow S).

Combining these two components with the logical operator "or" gives us the symbolization: $(E \to F) \lor \neg (C \to S)$.

Looking at the options:

 $A.\ (E=F)\lor -(C=S): This option uses the biconditional operator\ (=), but the given sentence does not indicate an equivalence between the conditions. Therefore, option A is not the best symbolization.$

B. (E = F) \vee (-C \supset S): This option uses the conditional operator (\supset) for the second condition, which is incorrect. The correct operator should be the negation of the conditional (-) to symbolize "not a sufficient condition." Therefore, option B is not the best symbolization.

 $C. \ (E\supset F) \lor \neg (C\supset S) \\ : This option correctly symbolizes the two conditions with the correct operators. This is the best symbolization of the given sentence.$

 $D. \ (F \supset E) \lor \neg (C \supset S) \\ : This \ option \ interchanges \ the \ positions \ of \ E \ and \ F \ in \ the \ first \ condition. \ Therefore, \ option \ D \ is \ not \ the \ best \ symbolization.$

Final Answer: 0

Question 102:

' Select the best English interpretation of the given arguments in predicate logic.

 $Wn \lor Wm$

 $(\forall x)[Lx \supset (Dx \supset \sim Wx)]$

Ln • Dn / ~ (\fo x) – Wx', A. 'Either Nancy or Marvin are at work.' B. 'Either Nancy or Marvin are at work.' All lawyers are not at work.' B. 'Either Nancy or Marvin are at work. All lawyers are out to dinner if they are not at work. Nancy is a lawyer and out to dinner. So not everything is not at work.' C. 'Either Nancy or Marvin are at work. All lawyers are out to dinner if they are not at work. Nancy is a lawyer and out to dinner. So not everything is at work.' D. 'Either Nancy or Marvin are at work. All lawyers are not at work if they are out to dinner. Nancy is a lawyer and out to dinner. So not everything is at work.'

Answer Option 1:

Let's break down the argument using the given predicates:

- 1. Wn \lor Wm This statement says that either Nancy (n) or Marvin (m) is at work.
- 2. $(\forall x)[Lx\supset (Dx\supset -Wx)]$ This statement says that for all x (everyone), if they are a lawyer (Lx), then if they are out to dinner (Dx), they are not at work (-Wx).
- 3. Ln Dn This statement says that Nancy is a lawyer (Ln) and she is out to dinner (Dn).

Using these statements, we can conclude that Nancy is not at work. However, we cannot conclude that "not everything is not at work" from the given premises.

Therefore, the best English interpretation of the argument is:

'Either Nancy or Marvin are at work. All lawyers are out to dinner if they are not at work. Nancy is a lawyer and out to dinner. So Nancy is not at work.'

The correct answer is B.

Answer Option 2:

Let's break down the given arguments in predicate logic:

- 1. Wn v Wm: This statement can be interpreted as "Either Nancy (n) or Marvin (m) is at work."
- 2. $(\forall x)[Lx\supset (Dx\supset \neg Wx)]$: This statement can be interpreted as "For all individuals (x), if they are a lawyer (Lx), then if they are out to dinner (Dx), they are not at work $(\neg Wx)$."
- 3. Ln Dn: This statement can be interpreted as "Nancy is a lawyer (L) and Nancy is out to dinner (D)."
- 4. \sim (\forall x) \sim Wx: This statement can be interpreted as "It is not the case that for all individuals (x), they are not at work (\sim Wx)."

Based on these statements, the best English interpretation of the given argument is:

"Either Nancy or Marvin is at work. All lawyers are not at work if they are out to dinner. Nancy is a lawyer and out to dinner. So not everything is at work."

Therefore, the correct option is A

Answer Option 3:

Let's break down the argument step by step:

- 1. Wn \vee Wm: This statement states that either Nancy (N) is at work (W) or Marvin (M) is at work.
- 2. $(\forall x)[Lx\supset (Dx\supset \neg Wx)]$: This statement is a universally quantified statement that says for all x, if x is a lawyer (L), then if x is out to dinner (D), then x is not at work (W).
- 3. Ln Dn: This statement states that Nancy is a lawyer (L) and Nancy is out to dinner (D).

 $Conclusion: \ \, \sim (\forall x) \sim Wx: \ \, \text{This statement says that not everything is not at work, which can be simplified as "something is at work"}.$

From the given premises and conclusion, the best English interpretation of the argument is:

C. "Either Nancy or Marvin are at work. All lawyers are out to dinner if they are not at work. Nancy is a lawyer and out to dinner. So not everything is at work."

Note that options A, B, and D do not accurately capture the logic of the argument and the implications of the premises and conclusion.

Answer Option 4:

Let's analyze the given arguments step by step:

- 1. Wn \vee Wm: This statement represents that either Nancy (N) or Marvin (M) is at work.
- 2. (\forall x)[Lx \(\to (Dx \) -\forall x)]: This statement represents that for all x, if x is a lawyer (Lx), then if x is out to dinner (Dx), then x is not at work (-\forall x). In other words, if someone is a lawyer and they are out to dinner, they are not at work.
- 3. Ln Dn: This statement represents that Nancy (N) is a lawyer (Ln) and she is out to dinner (Dn).

Now, let's draw the logical deductions based on the given statements:

From 1 and 3, we can deduce that Nancy is at work: Wn.

From 2 and 3, we can deduce that Nancy is not at work: ~Wn.

Since we have reached a contradiction (Wn and ~Wn), we can conclude that the argument is invalid.

None of the given answer options accurately interpret the given argument. Therefore, none of the options A, B, C, or D are correct interpretations of the argument.

Answer Option 5:

To interpret the given arguments in predicate logic, let's analyze the statements one by one

- 1. Wn \lor Wm: This statement states that either Nancy (n) or Marvin (m) is at work.
- 2. (∀x)[Lx ⊃ (Dx ⊃ ~Wx)]: This statement expresses that for all individuals x, if x is a lawyer (Lx), then if x is out to dinner (Dx), x is not at work (Wx).
- 3. Ln Dn: This statement states that Nancy (n) is a lawyer (Ln) and she is out to dinner (Dn).

Now, let's deduce the conclusion from the premises:

Given that either Nancy or Marvin is at work (Wn \vee Wm) and the statement ($\forall x$)[Lx \supset (Dx \supset \sim Wx)] implies that if someone is a lawyer and out to dinner, they are not at work, we can conclude that Nancy,

being a lawyer and out to dinner (Ln • Dn), is not at work (~Wn).

So, the best English interpretation of the given arguments is option A: "Either Nancy or Marvin is at work. All lawyers are not at work if they are out to dinner. Nancy is a lawyer and out to dinner. So not everything is not at work."

Answer Option 6:

Let's analyze the given arguments in predicate logic step by step:

- 1. Wn \lor Wm This statement states that either Nancy (n) or Marvin (m) is at work.
- 2. (∀x)[Lx ⊃ (Dx ⊃ ~Wx)] This statement is a universal statement that says for all x, if x is a lawyer (Lx), then if x is out to dinner (Dx), then x is not at work (~Wx). This means that all lawyers are not at work if
- 3. Ln Dn This statement states that Nancy (n) is a lawyer (Ln) and is out to dinner (Dn).

Based on these premises, we can conclude ~(\forall x)-Wx. This statement can be understood as "It is not the case that everything is not at work." In other words, not all things/people are not at work, which can be logically interpreted as "Some things/people are at work."

Therefore, the best English interpretation of the given arguments is:

'Either Nancy or Marvin are at work. All lawyers are not at work if they are out to dinner. Nancy is a lawyer and out to dinner. So not everything is not at work.'

Hence, the correct answer is A.

Answer Option 7:

To analyze the argument and select the best English interpretation, let's break down the given statements:

Wn v Wm: Either Nancy (N) or Marvin (M) is at work.

 $(\forall x)[Lx\supset (Dx\supset \sim Wx)] : \text{For all } x, \text{ if } x \text{ is a lawyer (L), then if } x \text{ is out to dinner (D), } x \text{ is not at work (W)}.$

Ln • Dn: Nancy is a lawyer (L) and out to dinner (D).

The conclusion we want to derive is $\sim (\forall x) \sim Wx$, which is equivalent to saying that not everything is not at work.

Now, let's analyze the premises and conclusion together:

- 1. Either Nancy or Marvin is at work. (Wn v Wm)
- 2. For all x, if x is a lawyer, then if x is out to dinner, x is not at work. $(\forall x)[Lx \supset (Dx \supset \neg Wx)]$
- 3. Nancy is a lawyer and out to dinner. (Ln Dn)

From premise 1, we know that either Nancy or Marvin is at work.

From premise 2, we know that for all lawyers, if they are out to dinner, they are not at work.

From premise 3, we know that Nancy is a lawyer and out to dinner.

Based on the given premises, we can conclude that Nancy is not at work.

Therefore, the best English interpretation of the argument is:

'Either Nancy or Marvin are at work. All lawyers are not at work if they are out to dinner. Nancy is a lawyer and out to dinner. So not everything is not at work.'

The correct answer is A.

Answer Option 8:

To determine the best English interpretation of the given arguments in predicate logic, let's break down the statements:

- $1.\ Wn\lor Wm:\ This\ statement\ means\ that\ either\ Nancy\ (N)\ is\ at\ work\ (Wn)\ or\ Marvin\ (M)\ is\ at\ work\ (Wm).$
- 2. $(\forall x)[Lx\supset (Dx\supset \neg Wx)]$: This statement can be interpreted as "For all individuals x, if x is a lawyer (Lx), then if x is not at work $(\neg Wx)$, then x is out to dinner (Dx)."
- 3. Ln Dn: This statement indicates that Nancy (N) is a lawyer (Ln) and she is out to dinner (Dn).
- $4. \ {\sim} (\forall x) {\sim} Wx: This \ statement \ can be \ interpreted \ as \ "It \ is \ not \ the \ case \ that \ for \ all \ individuals \ x, \ x \ is \ not \ at \ work."$

Now, let's analyze the given arguments:

The given argument can be expressed symbolically as:

 $(\mathsf{Wn} \vee \mathsf{Wm}) \wedge (\forall x)[\mathsf{Lx} \supset (\mathsf{Dx} \supset \mathsf{\sim} \mathsf{Wx})] \wedge (\mathsf{Ln} \bullet \mathsf{Dn}) \, / \, \mathsf{\sim} (\forall x) \mathsf{\sim} \mathsf{Wx}$

Applying the logical rules of inference, we can conclude that "It is not the case that for all individuals x, x is not at work." This means that at least one individual is at work.

The best English interpretation of the given arguments is:

"Either Nancy or Marvin is at work. All lawyers are out to dinner if they are not at work. Nancy is a lawyer and out to dinner. Therefore, it is not the case that everyone is not at work."

So, the correct answer is A. "Either Nancy or Marvin is at work. All lawyers are not at work if they are out to dinner. Nancy is a lawyer and out to dinner. So not everything is not at work."

Answer Option 9:

Let's break down the argument step by step:

- $1.\ Wn \lor Wm This\ statement\ asserts\ that\ either\ Nancy\ (represented\ by\ Wn)\ or\ Marvin\ (represented\ by\ Wm)\ is\ at\ work.$
- 2. $(\forall x)[Lx \supset (Dx \supset \neg Wx)]$ This statement is a universal quantification $(\forall x)$ that applies to all individuals represented by the variable x. In this case, it states that for all x, if x is a lawyer (represented by Lx), then if x is out to dinner (represented by Dx), then x is not at work (represented by $\neg Wx$).
- 3. Ln Dn This statement asserts that Nancy (represented by Ln) is a lawyer (represented by Ln) and is out to dinner (represented by Dn).
- 4. The conclusion is $-(\forall x)$ -Wx This statement is the negation of the universal quantification $(\forall x)$ -Wx, which asserts that it is not the case that everything is not at work.

Now, let's analyze the given answer choices:

Option A: 'Either Nancy or Marvin are at work. All lawyers are not at work if they are out to dinner. Nancy is a lawyer and out to dinner. So not everything is not at work.' - This interpretation does not accurately represent the given argument.

Option B: 'Either Nancy or Marvin are at work. All lawyers are out to dinner if they are not at work. Nancy is a lawyer and out to dinner. So not everything is not at work.' - This interpretation does not accurately represent the given argument.

Option C: 'Either Nancy or Marvin are at work. All lawyers are out to dinner if they are not at work. Nancy is a lawyer and out to dinner. So not everything is at work.' - This interpretation does not accurately represent the given argument.

Option D: 'Either Nancy or Marvin are at work. All lawyers are not at work if they are out to dinner. Nancy is a lawyer and out to dinner. So not everything is at work.' - This interpretation accurately represents the given argument.

Therefore, the correct interpretation is Option D: 'Either Nancy or Marvin are at work. All lawyers are not at work if they are out to dinner. Nancy is a lawyer and out to dinner. So not everything is at work.'

Question 103:

'Identify the conclusion of the following argument. Remember to remove any conclusion indicators. It is wrong for society to kill a murderer. This follows for the reason that if a murderer is wrong in killing his victim, then society is also wrong in killing the murderer. And a murderer is wrong in killing his victim, then society to kill a murderer.', B. This follows for the reason that if a murderer is wrong in killing his victim, then society is also wrong in killing the murderer.', C. 'If a murderer is wrong in killing his victim, then society is also wrong in killing the murderer.', D. 'And a murderer is wrong in killing his victim.'

The conclusion of the argument is: A. 'It is wrong for society to kill a murderer.' ********* Answer Option 2: The conclusion of the argument is: A. 'It is wrong for society to kill a murderer.' Answer Option 3: The conclusion of the argument is: A. 'It is wrong for society to kill a murderer.' Answer Option 4: The conclusion of the argument is: A. 'It is wrong for society to kill a murderer.' Answer Option 5: The conclusion of the argument is: A. 'It is wrong for society to kill a murderer.' ******* Answer Option 6: The conclusion of the argument is: C. 'If a murderer is wrong in killing his victim, then society is also wrong in killing the murderer.' ********* Answer Option 7: The conclusion of the argument is: 'It is wrong for society to kill a murderer.' This conclusion is derived from the premises that state that if a murderer is wrong in killing his victim, then society is also wrong in killing the murderer, and that a murderer is wrong in killing his victim. Therefore, the correct answer is option A: 'It is wrong for society to kill a murderer.' Answer Option 8: The conclusion of the argument is: "It is wrong for society to kill a murderer." This conclusion is derived from the premises given in the argument, which state that if a murderer is wrong in killing his victim, then society is also wrong in killing the murderer, and that a murderer is wrong in killing his victim. Therefore, the correct answer is A. 'It is wrong for society to kill a murderer.' ********* The conclusion of the argument is: "It is wrong for society to kill a murderer." This conclusion is derived from the reasons provided in the argument, which state that if a murderer is wrong in killing his victim, then society is also wrong in killing the murderer. The final reason given in the argument is that a murderer is wrong in killing his victim. Therefore, the conclusion is that it is wrong for society to kill a murderer. Final Answer: A. 'It is wrong for society to kill a murderer.'

Question 104: 'Which of the following propositions is not an immediate (one-step) consequence in PL of the given premises? $I\supset (J\equiv K)$ $(J \equiv K) \supset (L \lor M)$ $I\supset {\sim}L',\ A.\ '{\sim}L',\ B.\ 'J\equiv K',\ C.\ 'I\supset (L\vee M)',\ D.\ '{\sim}L\supset (J\equiv K)'$ To determine which proposition is not an immediate consequence of the given premises, we can analyze each option based on the rules of propositional logic. Premises: 1. $I \supset (J \equiv K)$ 2. $(J \equiv K) \supset (L \lor M)$ 3. I \supset ~L Option A: '~L' We can see that '~L' can be derived as an immediate consequence from premise 3, as 'I \supset ~L' is one of the premises. Option B: $'J \equiv K'$ We can derive 'J \equiv K' as an immediate consequence by applying modus ponens using premises 1 and 2. Option C: 'L > (L v M)' We can derive 'I \supset (L \lor M)' as an immediate consequence by applying modus ponens using premises 1 and 2. Option D: $^{\prime}\sim L\supset (J\equiv K)^{\prime}$ We cannot directly derive $' \sim L \supset (J \equiv K)'$ as an immediate consequence from the given premises. This option does not follow logically from the premises. Therefore, the proposition that is not an immediate consequence of the given premises is option D: $-L \supset (J \equiv K)$. Final Answer: D ******* Answer Option 2: To determine which proposition is not an immediate consequence of the given premises, we can utilize the rules of propositional logic. The given premises are: $1.\ I\supset (J\equiv K)$ 2. $(J \equiv K) \supset (L \lor M)$ We can analyze each option: Option A: '~L' In the given premises, there is no direct implication or logical connection involving ~L. Therefore, '~L' is not an immediate consequence. Option B: $'J \equiv K'$ This proposition is already included in the premises as ' $(J \equiv K)$ '. Therefore, ' $J \equiv K$ ' is an immediate consequence. Option C: $I \supset (L \lor M)'$ Using the premises, we can infer the following chain: $I \supset (J \equiv K)$ (premise 1) $(J \equiv K) \supset (L \lor M)$ (premise 2) Therefore, $I \supset (L \vee M)$ is an immediate consequence. Option D: $^{\prime}\sim L\supset (J\equiv K)^{\prime}$ This proposition is not an immediate consequence because it is not directly supported by the given premises. Although we have $I \supset \sim L$, there is no direct connection between $\sim L$ and $(J \equiv K)$. Based on the analysis, the proposition that is not an immediate consequence of the given premises is D. Therefore, the answer is D. Answer Option 3: To determine which proposition is not an immediate consequence in propositional logic (PL) of the given premises, we need to apply the rules of inference and logical equivalences to deduce the possible conclusions. Given premises: 1. $I \supset (J \equiv K)$ 2. $(J \equiv K) \supset (L \lor M)$

To deduce this proposition from the premises, we can use Modus Tollens (MT) on premise 3: I \supset -L. However, this proposition is an immediate consequence.

3. I ⊃ ~L

Let's examine the options:
Option A: '~L'

Option B: $'J \equiv K'$

This proposition is an immediate consequence since it is given in the premises.

Option C: $^{\prime}I\supset (L\vee M)^{\prime}$

Using Disjunctive Syllogism (DS) on premises 1 and 2, we can deduce 'I \supset (L \vee M)'. Therefore, this proposition is an immediate consequence.

Option D: $^{\prime}\sim L\supset (J\equiv K)^{\prime}$

To deduce this proposition, we need to use rules of inference and logical equivalences. However, it is not an immediate consequence of the given premises.

Therefore, the proposition that is NOT an immediate (one-step) consequence in PL of the given premises is option D: $-L \supset (J \equiv K)$. Final Answer: D.

Answer Option 4:

To determine which proposition is not an immediate consequence of the given premises, we need to examine each option and see if it can be derived directly from the premises using one step.

Given premises:

 $1.\ I\supset (J\equiv K)$

2. $(J \equiv K) \supset (L \lor M)$

3. I \supset ~L

Option A: '~L'

To derive '-L' from the premises, we can use Modus Tollens on premise 3, which gives us -I. However, we cannot directly derive '-L' from the given premises alone. Therefore, option A is not an immediate consequence.

Option B: $'J \equiv K'$

This proposition is already provided as one of the premises, so it is an immediate consequence.

Option C: $^{\prime}I\supset (L\vee M)^{\prime}$

To derive $'I\supset (L\lor M)'$ from the premises, we can use Hypothetical Syllogism, combining premise 1 and premise 2. Therefore, option C is an immediate consequence.

Option D: $'\sim L \supset (J \equiv K)'$

To derive '-L \supset (J \equiv K)' from the premises, we can use Contrapositive on premise 3, which gives us L \supset ~I. However, we cannot directly derive '-L \supset (J \equiv K)' from the given premises alone. Therefore, option D

is not an immediate consequence.

Final Answer: A

Answer Option 5:

To determine which proposition is not an immediate consequence of the given premises, we can use logical deduction rules to evaluate each option.

Given premises

 $1. \ I \supset (J \equiv K)$

 $2.~(J\equiv K)\supset (L\vee M)$

3. I ⊃ ~L

Let's evaluate each option:

Option A: '~L'

We can derive this proposition using premise 3. Since I ⊃ ~L and I is given, we can deduce ~L. Therefore, '~L' is an immediate consequence.

Option B: 'J ≡ K'

We can derive this proposition using premise 1. Since $I \supset (J \equiv K)$ and I is given, we can deduce $J \equiv K$. Therefore, $J \equiv K'$ is an immediate consequence.

Option C: $I \supset (L \lor M)'$

We can derive this proposition using premises 1 and 2. By applying the transitive rule of implication, we have I \supset (L \lor M) as an immediate consequence.

Option D: ${}^{\backprime}\sim L\supset (J\equiv K)^{\backprime}$

To determine if this proposition is an immediate consequence, we need to evaluate it based on the given premises. However, we cannot derive this proposition directly using the given premises. Therefore,

 $\ ^{\shortmid }{\sim }L\supset (J\equiv K)^{\shortmid }$ is not an immediate consequence.

Final Answer: D. ' \sim L \supset (J \equiv K)'

Answer Option 6:

To determine which proposition is not an immediate consequence of the given premises, we need to check the truth values of each proposition based on the truth values of the premises.

Given premises:

1. I ⊃ (J ≡ K)

2. $(J \equiv K) \supset (L \lor M)$

3. I ⊃ ~L

Let's evaluate each option:

Option A: '~L'

Since -L is not mentioned in any of the premises, we cannot determine its truth value based on the given information. Therefore, '-L' is not an immediate consequence of the premises.

Option B: $'J \equiv K'$

This proposition is directly stated in the first premise. Therefore, $J \equiv K'$ is an immediate consequence of the premises.

Option C: $^{\prime}I\supset (L\vee M)^{\prime}$

This proposition is not directly stated in the premises. It involves a combination of I, L, and M with the logical operator \supset . Therefore, "I \supset (L \lor M)" is not an immediate consequence of the premises.

Option D: $^{\prime}\sim L\supset (J\equiv K)^{\prime}$

This proposition is not directly stated in the premises. It involves a combination of ~L, J, and K with the logical operator \supset . Therefore, '~L \supset ($J \equiv K$)' is not an immediate consequence of the premises.

Final Answer: A

Answer Option 7:

To determine which proposition is not an immediate consequence of the given premises, we can check if the proposition can be derived directly from the premises using one step in propositional logic (PL).

Given premises:

 $1.\ I\supset (J\equiv K)$

2. $(J \equiv K) \supset (L \lor M)$

3. I \supset ~L

Let's consider each option:

Option A: '~L'

We can see that '~L' is not directly present in any of the premises, so it cannot be an immediate consequence.

Ontion B: 'I = K'

This proposition is directly present in the first premise $(I \supset (J \equiv K))$, so it is an immediate consequence.

Option C: $'I \supset (L \lor M)'$

This proposition is not directly present in any of the premises, so it cannot be an immediate consequence.

Option D: $'\sim L \supset (J \equiv K)'$

We can see that '-L' is present in the third premise (I \supset -L), and '(J \equiv K)' is present in the first premise (I \supset (J \equiv K)'. Therefore, by combining these premises, we can derive the proposition '-L \supset (J \equiv K)'.

Hence, it is an immediate consequence.

Final Answer: Option C, 'I \supset (L \vee M)', is not an immediate consequence of the given premises.

Answer Option 8:

To determine which proposition is not an immediate consequence in PL (Propositional Logic) of the given premises, we need to check if we can derive the proposition directly from the premises using one step using standard inference rules.

Given premises:

1. I ⊃ (J ≡ K)

2. $(J \equiv K) \supset (L \lor M)$

3. I ⊃ ~L

Let's evaluate each option:

Option A: '~L'

We can derive this proposition directly from the third premise by using Modus Tollens, which states that if we have a conditional statement P \supset Q and its negation \sim Q, then we can infer the negation of the antecedent \sim P. Therefore, ' \sim L' is an immediate consequence of the given premises.

Option B: 'J ≡ K'

We can derive this proposition directly from the first premise as it is stated explicitly. Therefore, $"J \equiv K"$ is an immediate consequence of the given premises.

Option C: $I \supset (L \lor M)'$

To derive this proposition, we need to combine the second premise and the third premise using the hypothetical syllogism. However, this would require an additional step and is not an immediate consequence of the given premises. Therefore, 1 \supset (L \vee M) is not an immediate consequence.

Option D: $'\sim L \supset (J \equiv K)'$

To derive this proposition, we would need to combine the first premise and the third premise using the hypothetical syllogism. However, this would require an additional step and is not an immediate consequence of the given premises. Therefore, $-L \supset (J = K)$ is not an immediate consequence.

Final Answer: C ('I \supset (L \vee M)') is not an immediate consequence in PL of the given premises.

Answer Option 9:

To determine which proposition is not an immediate consequence of the given premises, we can evaluate each option using the rules of propositional logic.

Premises:

 $1. \ I \supset (J \equiv K)$

2. $(J \equiv K) \supset (L \vee M)$

3. I ⊃ ~L

Option A: '~L'

To obtain '-L' as an immediate consequence, we need 'L' to be false. However, the premises do not provide any information about the truth value of 'L'. Therefore, '-L' is not an immediate consequence.

Option B: $'J \equiv K'$

This proposition can be derived as an immediate consequence by applying the transitive property of equivalence. From premise 1, $I\supset (J\equiv K)$, and from premise 2, $(J\equiv K)\supset (L\vee M)$, we can conclude $I\supset (L\vee M)$ using the transitive property. Then, combining this with premise 3, $I\supset \sim L$, we can derive ' $J\equiv K$ ' as an immediate consequence.

Option C: $I \supset (L \lor M)'$

This proposition can also be derived as an immediate consequence by applying the transitive property of implication. From premise 1, $I \supset (J \equiv K)$, and from premise 2, $(J \equiv K) \supset (L \lor M)$, we can directly conclude $I \supset (L \lor M)$ using the transitive property.

Option D: $'\sim L\supset (J\equiv K)'$

To obtain $\ '-L \supset (J \equiv K)'$ as an immediate consequence, we need to show that when $\ 'L'$ is false, $\ 'J \equiv K'$ must hold. However, the premises do not provide any direct relationship between $\ '-L'$ and $\ '(J \equiv K)'$. Therefore, $\ '-L \supset (J \equiv K)'$ is not an immediate consequence.

Final Answer: A. '~L'

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'Construct a complete truth table for the following argument. Then, using the truth table, determine whether the argument is valid or invalid. If the argument is invalid, choose an option which presents a counterexample. (There may be other counterexamples as well.)

I ≡ ~J

I v J / I', A. 'Valid', B. 'Invalid. Counterexample when I and J are true', C. 'Invalid. Counterexample when I is true and J is false', D. 'Invalid. Counterexample when J is true and I is false'

Answer Option 1:

To construct a complete truth table for the argument, we need to consider all possible combinations of truth values for the variables I and J.

 $IJ \sim JI \equiv \sim JI \vee JI'$

TTEETT

TFTFTT

FTFTTF

FFTTFF

From the truth table, we can see that the argument is invalid because there is a row where the premises (I v J) is true but the conclusion (I') is false. Specifically, when I is false and J is true, the premise is true (F v T = T) but the conclusion is false (F').

Therefore, the correct answer is D. "Invalid. Counterexample when J is true and I is false."

Answer Option 2:

 $I \equiv \sim J$ represents that I is logically equivalent to the negation of J.

 $\mbox{I} \vee \mbox{J}$ represents the logical OR of I and J.

I' represents the negation of I.

To construct a complete truth table, we need to consider all possible combinations of truth values for I and J.

 $IJ \sim JI \equiv \sim JI \vee JI'$

....

TTFFTF

TFTFTF

FTFTTT

FFTTFT

From the truth table, we can see that there are cases where the argument is invalid. In particular, when I is true and J is false, the premise (I v J) is true, but the conclusion (I') is false. Therefore, the argument is invalid.

The correct answer is D. 'Invalid. Counterexample when J is true and I is false'.

Answer Option 3:

 $I \equiv {\sim}J$ means I is equivalent to the negation of J.

The argument is as follows:

Premise: I v J

Conclusion: I'

To construct a complete truth table, we need to consider all possible combinations of truth values for I and J.

| J ~J | ≡ ~J | ∨ J |

TTFFTF

TFTTTF

FTFTTT

FFTFFT

In the truth table, we can see that when the premise $I \vee J$ is true, the conclusion I' is false. This means that the argument is invalid.

To find a counterexample, we need to find a row in the truth table where the premise I v J is true, but the conclusion I' is false. Looking at the truth table, we can see that when I is true and J is false, the premise I v J is true, but the conclusion I' is false. Therefore, the counterexample is when I is true and J is false.

Final Answer: C. Invalid. Counterexample when I is true and J is false.

Answer Option 4:

To construct a complete truth table, we need to consider all possible combinations of truth values for the propositions I and J, as well as evaluate the given argument.

 $I \equiv {\sim}J$ translates to I is equivalent to not J.

The argument I \vee J / I' can be interpreted as "If I is true or J is true, then I' (not I) is true."

Let's create the truth table: $|\; I\; |\; J\; |\; {\sim} J\; |\; I\; \equiv\; {\sim} J\; |\; I\; \vee\; J\; |\; I'\; |\;$ |---|----|-----| |T|T|F|F|T|F| |T|F|T|T|T|F| |F|T|F|T|T|T| |F|F|T|F|F|T| Based on the truth table, we can see that there is no row where the premise (I v J) is true but the conclusion (I') is false. Therefore, the argument is valid. The correct answer is A. 'Valid'. Answer Option 5: To construct a complete truth table, we need to consider all possible truth value combinations for the variables I and J. I ≡ ~J $I\vee J/I'$ Let's create the truth table: | | | | | ~ | | | = ~ | | | | \land | | | | T|T|F|F|T|TTIFITITIT FITIFITIT F|F|T|F|F|F Based on the truth table, we can see that the argument is invalid because there are cases where the premises are true (I = ~J and I ∨ J) but the conclusion (I') is false. A counterexample is when I is true and J is false: $I \equiv \sim J$ is true ($T \equiv \sim F$ is true) $I \vee J \text{ is true (T} \vee F \text{ is true)}$ I' is false (F is false) Therefore, the correct answer is: C. Invalid. Counterexample when I is true and J is false. Answer Option 6: To construct a complete truth table, we need to consider all possible combinations of truth values for the propositions involved. In this case, we have propositions I and J. $I \equiv \sim J$ represents the equivalence between I and the negation of J. $I \vee J$ represents the disjunction (logical OR) between I and J. Based on these propositions, we can construct the truth table as follows: |---|----|-----| |T|T|F|F|T|F| |T|F|T|T|T|F| |F|T|F|T|T|T| |F|F|T|F|F|T| In the truth table, I' represents the negation of I. Now, let's analyze the truth table to determine the validity of the argument. From the truth table, we can see that whenever I v J is true, I' is also true. This means that the argument is valid since the conclusion I' is always true whenever the premise I v J is true. Therefore, the correct answer is A. 'Valid'. ******** Answer Option 7: To construct the truth table for the argument, we need to consider all possible combinations of truth values for the variables I and J. $I \;|\; J \;|\; {\sim} J \;|\; I \equiv {\sim} J \;|\; I \;\vee\; J \;|\; I'$ T|T|F|F|T|F TIFITIFITIF

FITIFITIT

F|F|T|T|F|T

Now, we can analyze the truth table to determine the validity of the argument. In this argument, the premise is I v J and the conclusion is I'. For the argument to be valid, whenever the premise is true, the conclusion must also be true.

Looking at the truth table, we can see that there is a row where the premise (I v J) is true (T), but the conclusion (I') is false (F). Specifically, when I is false (F) and J is true (T), the premise is true, but the conclusion is false.

Therefore, the argument is invalid.

The correct option is D. 'Invalid. Counterexample when J is true and I is false'.

Answer Option 8:

To construct a complete truth table, we need to consider all possible combinations of truth values for the variables I and J and evaluate the argument.

 $I \equiv \sim J$ represents the equivalence between I and the negation of J.

Let's start by constructing the truth table for I \equiv ~J:

|||J||-J||=-J| |--|--|----| |T|T|F|F| |T|F|T|T| |F|T|F|F| |F|F|T|T|

Next, let's evaluate the argument I \vee J / I':

Based on the truth table, we see that the argument is invalid because there exists a row where the premise (I v J) is true, but the conclusion (I') is false. Specifically, when I is false and J is true, the premise is true, but the conclusion is false.

The correct answer is D. 'Invalid. Counterexample when J is true and I is false'.

Answer Option 9:

To construct a truth table for the argument, we need to consider all possible combinations of truth values for the propositions I and J.

From the truth table, we can see that whenever I v J is true, I' is also true. This means that the argument is valid because the conclusion I' always follows from the premise I v J.

Therefore, the correct answer is A. 'Valid'.

Question 106:

'Select the best translation into predicate logic: Some folks are lucky.', A. '(∃x)(Fx • Lx)', B. 'FI', C. 'Lf', D. '(∀x)(Fx • Lx)'

Answer Option 1:

The best translation into predicate logic for the statement "Some folks are lucky" would be: A. $(\exists x)(\mathsf{Fx} \bullet \mathsf{Lx})$

This translates to "There exists an x such that x is a folk and x is lucky", indicating that there is at least one person who is both a folk and lucky. Final Answer: A

Answer Option 2:

The best translation into predicate logic for the statement "Some folks are lucky" would be option A: $(\exists x)(\mathsf{Fx} \bullet \mathsf{Lx})$. This translates to "There exists an x such that x is a folk and x is lucky." This accurately captures the meaning of the statement by using the existential quantifier (3x) to indicate that there is at least one individual who is both a folk (Fx) and lucky (Lx). Final Answer: A

Answer Option 3:

The best translation into predicate logic for the statement "Some folks are lucky" would be: A. (3x)(Fx • Lx) This translates to "There exists an individual x such that x is fortunate (Fx) and x is lucky (Lx)." This accurately captures the meaning of the original statement. Final Answer: A

Answer Option 4:

The best translation into predicate logic for the statement "Some folks are lucky" would be option A: (3x)(Fx • Lx).

This translates to "There exists an individual x such that they are lucky (Lx) and also a folk (Fx)."

This accurately represents the statement that there is at least one person who is lucky.

Answer Option 5:

The best translation into predicate logic for the statement "Some folks are lucky" is option A: ($\exists x$)(Fx • Lx), which translates to "There exists an x such that x is lucky and x is a folk." This accurately captures the idea that there exists at least one person who is both lucky and a folk. Final Answer: A. $(\exists x)(\mathsf{Fx} \bullet \mathsf{Lx})$

Answer Option 6:

The best translation of the statement "Some folks are lucky" into predicate logic is A. '(∃x)(Fx • Lx)'. This can be read as "There exists an x such that x is a folk (F(x)) and x is lucky (L(x))". It correctly captures the idea that there exists at least one person who is both a folk and lucky. Option B, 'Fl', translates to "There exists an x such that x is lucky", which does not capture the notion of being a folk. Option C, 'Lf', translates to "For all x, x is lucky", which does not capture the idea that only some people are lucky. Option D. '(\forall x)(\forall x)(\forall x)', translates to "For all x, x is a folk and x is lucky", which would mean that every single person is lucky, rather than just some people.

Answer Option 7:

The best translation into predicate logic for the statement "Some folks are lucky" would be option A: ($\exists x$)(Fx • Lx), which translates to "There exists an individual x such that x is lucky (Lx) and x is a folk (Fx)." This correctly represents the idea that there is at least one lucky person among the group of folks. Option B: FI, translates to "There exists an individual x such that x is lucky (Lx)." This doesn't capture the idea of folks or a group of people. Option C: Lf, translates to "There exists an individual x such that x is a folk (Fx)." This doesn't capture the idea of luck. Option D: (∀x)(Fx • Lx), translates to "For every individual x, x is lucky (Lx) and x is a folk (Fx)." This represents the idea that every folk is lucky, which contradicts the original statement. Final Answer: A

Answer Option 8:

The best translation into predicate logic for the statement "Some folks are lucky" is option A: '(3x)(Fx • Lx)'. This translates to "There exists an x such that x is a folk (F(x)) and x is lucky (L(x))". This accurately captures the idea that there is at least one person who is both a folk and lucky. Option B 'FI' translates to "There is an x such that x is lucky", which does not capture the notion of being a folk. Option C 'Lf' translates to "There is an x such that x is a folk", which does not capture the notion of being lucky. Option D '(∀x)(Fx • Lx)' translates to "For all x, x is a folk and x is lucky", which implies that every person is both a folk and lucky, which is not the same as "Some folks are lucky". Therefore, option A is the best translation into predicate logic. Final Answer: A

Answer Option 9:

The best translation into predicate logic for the statement "Some folks are lucky" is A. '(∃x)(Fx • Lx)'.

Explanation:

- (3x) represents the existential quantifier, which means "there exists" or "some".
- Fx represents the predicate "x is a folk" or "x is a person".
- Lx represents the predicate "x is lucky".

So, $(\exists x)(\mathsf{Fx} \bullet \mathsf{Lx})$ translates to "There exists a person x such that x is a folk and x is lucky", which accurately captures the meaning of the statement "Some folks are lucky".

Option B 'FI' would translate to "There exists a person x such that x is lucky", which does not capture the idea that the person is a folk.

Option C 'Lf' would translate to "There exists a person x such that x is a folk", which does not capture the idea that the person is lucky.

Option D '(\forall x)(Fx • Lx)' translates to "For all x, x is a folk and x is lucky", which is not the same as saying that some folks are lucky.

Question 107:
'Select the best translation into predicate logic. Alexis is greeted by Ben. (a: Alexis; b: Ben; Gxy: x greets y)', A. 'Gab', B. '~Gba', C. 'aGb', D. 'Gba'
Answer Option 1:
The best translation into predicate logic for the statement "Alexis is greeted by Ben" would be option D: 'Gba'. This translates to "Ben greets Alexis" using the predicate Gxy, where Gba represents "Ben greets"
Alexis".
Answer Option 2:
The best translation into predicate logic for the statement "Alexis is greeted by Ben" is D. "Gba". Here, Gxy represents "x greets y" and substituting a for Alexis and b for Ben would give us Gba, which means
Ben greets Alexis. Final Answer: D.

Answer Option 3:
The best translation into predicate logic for the statement "Alexis is greeted by Ben" would be option C: "aGb". This translates to "Alexis greets Ben" in predicate logic, where 'a' represents Alexis, 'G'
represents the greeting relation, and 'b' represents Ben. Final Answer: C.

Answer Option 4:
The best translation into predicate logic for the statement "Alexis is greeted by Ben" is option C: 'aGb'. In this translation, 'a' represents Alexis, 'G' represents the relation "is greeted by", and 'b' represents
Ben. Therefore, 'aGb' translates to "Alexis is greeted by Ben". Final Answer: C

Answer Option 5:
The best translation into predicate logic for the statement "Alexis is greeted by Ben" would be option D: 'Gba'.
In predicate logic, 'Gxy' represents the statement "x greets y".
Here, 'Gba' represents the statement "Ben greets Alexis", which accurately translates the given statement.
Therefore, the correct answer is D.

Answer Option 6:
The best translation into predicate logic for the statement "Alexis is greeted by Ben" is D. 'Gba'.
In predicate logic, we represent the relationship between the subject and object using a predicate symbol. Here, we can use the predicate symbol Gxy to represent "x greets y".
So, Gba represents "Ben greets Alexis", which accurately captures the statement "Alexis is greeted by Ben".
Answer Option 7:
The best translation into predicate logic for the statement "Alexis is greeted by Ben" is D. "Gba".
In the given options:
A. "Gab" - This translates to "Ben is greeted by Alexis", which is the opposite of the given statement.
B. "-Gba" - This translates to "Ben is not greeted by Alexis", which does not accurately represent the given statement.
C. "aGb" - This translates to "Alexis greets Ben", which is the opposite of the given statement.
Therefore, the correct translation is D. "Gba", which represents "Ben greets Alexis", accurately reflecting the original statement.
Answer Option 8:
The best translation into predicate logic for the statement "Alexis is greeted by Ben" would be:
G(ab)
Here, G represents the predicate "greets," and (ab) represents the individuals Alexis and Ben.
Option A: 'Gab' is not the correct translation as it does not include the predicate "greets."
Option B: '~Gba' translates to "Ben is not greeted by Alexis," which is the opposite of the given statement.
Option C: 'aGb' is not the correct translation as it reverses the order of the individuals.
Option D: 'Gba' translates to "Ben is greeted by Alexis," which is not the same as the given statement.
Final Answer: A
••••••
Answer Option 9:

The best translation into predicate logic for the statement "Alexis is greeted by Ben" would be Option D: "Gba".

In predicate logic, Gxy represents "x greets y". In this case, we can translate the statement as "Ben greets Alexis", which can be represented as Gba.

Option A, B, and C do not accurately represent the statement. Option A, "Gab", would mean "Alexis greets Ben" instead of Ben greeting Alexis. Option B, "-Gba", represents the negation of "Ben greets Alexis", which is not what the statement is saying. Option C, "aGb", represents "Alexis greets Ben" and does not capture the idea that Ben is greeting Alexis. Therefore, Option D is the correct translation.

Question 108:

'Select the best translation into predicate logic: Not all lobsters are red.', A. '~(∀x)(Lx ⊃ Rx)', B. '~(∃x)(Lx ⊃ Rx)', C. '~(∃x)(Lx • Rx)', D. '~(∀x)(Lx • Rx)'

Answer Option 1:

The best translation into predicate logic for the statement "Not all lobsters are red" is option B. '~(∃x)(Lx ⊃ Rx)'. This translates to "There does not exist an x such that if x is a lobster, then x is red", which accurately captures the meaning of the original statement. Final Answer: B.

Answer Option 2:

The best translation into predicate logic for the statement "Not all lobsters are red" is option B: $\[-(\exists x)(Lx \supset Rx)'. \]$ This translates to "There does not exist an x such that if x is a lobster, then x is red." This accurately captures the meaning of the original statement. Final Answer: B

Answer Option 3:

The best translation into predicate logic for the statement 'Not all lobsters are red' is:

B. $'\sim(\exists x)(Lx\supset Rx)$

This translates to "There does not exist an x such that if x is a lobster, then x is red." which accurately represents the statement 'Not all lobsters are red.'

Final Answer: B

Answer Option 4:

The best translation into predicate logic for the statement "Not all lobsters are red" is option B. '- $(\exists x)(Lx \supset Rx)$ '. This translates to "There does not exist an x such that if x is a lobster, then x is red." This accurately captures the meaning of the original statement. Final Answer: B

Answer Option 5:

The best translation into predicate logic for the statement "Not all lobsters are red." is option B. $(-3x)(Lx \supset Rx)$ ". This translates to "There does not exist an x such that if x is a lobster, then x is red." which accurately captures the meaning of the original statement. Final Answer: B

Answer Option 6:

The best translation into predicate logic for the statement "Not all lobsters are red" is option B: '~ $(\exists x)(Lx\supset Rx)$ '.

Explanation:

Let Lx represent "x is a lobster" and Rx represent "x is red".

The statement "Not all lobsters are red" can be translated as "It is not the case that there exists an x such that if x is a lobster, then x is red." This can be represented as \sim (3x)(Lx \supset Rx).

Answer Option 7:

The best translation into predicate logic for the statement "Not all lobsters are red" is option B. ' \sim ($\exists x$)($Lx \supset Rx$)'.

Explanation:

To translate "Not all lobsters are red" into predicate logic, we can break it down into two parts:

- 1. "All lobsters are red" is translated as $'(\forall x)(Lx \supset Rx)'$, which means for all x (lobsters), if x is a lobster (Lx), then x is red (Rx).
- 2. "Not all" is translated as the negation of the statement, which is '~($\forall x$)(Lx \supset Rx)'.

Therefore, the best translation is option B. '~($\exists x$)(Lx \supset Rx)'.

Answer Option 8:

The best translation into predicate logic for the statement "Not all lobsters are red" is option B: $'\sim(\exists x)(Lx\supset Rx)'$.

Explanation:

Let's break down the statement. "Not all lobsters are red" implies that there exists at least one lobster that is not red.

In predicate logic, we can represent "There exists" using the existential quantifier (\exists) , and we can represent "for all" using the universal quantifier (\forall) .

Let's define the predicates:

Lx: x is a lobster

Rx: x is red

The statement "Not all lobsters are red" can be translated as "It is not the case that for all lobsters, they are red." This can be symbolized as:

 $-(\forall x)(Lx \supset Rx)$

This option correctly captures the meaning of the statement. Therefore, the answer is option B: $'\sim (\exists x)(Lx\supset Rx)'$.

Answer Option 9:

The best translation into predicate logic for the statement "Not all lobsters are red" is option B: '~ $(\exists x)(Lx\supset Rx)$ '.

Explanation:

- The statement "Not all lobsters are red" can be rephrased as "There does not exist a lobster that implies it is red".
- In predicate logic, the universal quantifier $(\forall x)$ is used to denote "for all" or "all", and the existential quantifier $(\exists x)$ is used to denote "there exists" or "some".
- The symbol '>' represents implication, which means "implies".
- The negation symbol '~' is used to indicate "not".

Therefore, the best translation for the statement "Not all lobsters are red" is $'\sim(\exists x)(Lx\supset Rx)'$, which can be read as "There does not exist a lobster that implies it is red". Final Answer: B.

Question 109:

"Use indirect truth tables to determine whether the following argument is valid. If the argument is invalid, choose an option which presents a counterexample. (There may be other counterexamples as well.)

 $H \equiv (I \cdot \sim J)$

K≡~H

~(K \supset J) / I \equiv J', A. 'Valid', B. 'Invalid. Counterexample when I, H, and K are true and J is false', C. 'Invalid. Counterexample when I and H are true and J and K are false', D. 'Invalid. Counterexample when K is

true and I, H, and J are false'

Answer Option 1:

To determine the validity of the argument, we can construct an indirect truth table. We will assign truth values to the premises and the conclusion and check if the premises are true while the conclusion is

Let's assign truth values as follows:

 $H \equiv (I \cdot \sim J)$

K = ~H

 $\mathord{\sim} (\mathsf{K} \supset \mathsf{J})$

 $I \equiv J$

We will create a truth table for these propositions:

 $I \; H \; J \; K \; \mathord{\sim} (K \supset J) \; I \equiv J$

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In all the rows where the premises are true (T), the conclusion is also true (T). Hence, the argument is valid.

Final Answer: A. 'Valid'

Answer Option 2:

To determine the validity of the argument, we can construct an indirect truth table by assuming the negation of the conclusion and checking if the premises are still true. If the premises are true and the conclusion's negation is true, then the argument is invalid.

Let's construct the indirect truth table:

 $|\; I\; |\; J\; |\; H\; |\; K\; |\; K\supset J\; |\; \sim (K\supset J)\; |\; I\equiv J\; |\; \sim I\equiv \sim J\; |\;$

|---|---|--------|-------|

|T|T|T|F|T|F|T|T|

|T|F|F|T|F|T|F|T|

 $|\mathsf{F}|\mathsf{T}|\mathsf{F}|\mathsf{T}|\mathsf{T}|\mathsf{F}|\mathsf{F}|\mathsf{T}|$

|F|F|T|F|T|F|F|T|

Since we have found a row where all the premises are true and the conclusion's negation is true (row 2), the argument is invalid.

Therefore, the correct answer is:

B. Invalid. Counterexample when I, H, and K are true and J is false.

Answer Option 3:

To determine the validity of the argument, we will create an indirect truth table. We will assign truth values to the statements and check if the conclusion always follows from the premises.

 $H \equiv (I \cdot {\sim} J) \mid K \equiv {\sim} H \mid {\sim} (K \supset J) \mid I \equiv J'$

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From the truth table, we can see that there are cases where the premises are true but the conclusion is false. Specifically, when I and H are true, J is false, and K is false, the premises are true but the
conclusion is false. Therefore, the argument is invalid.
Counterexample: I and H are true, J is false, and K is false.
Final Answer: C. Invalid. Counterexample when I and H are true and J and K are false.

Answer Option 4:
To determine the validity of the argument using indirect truth tables, we need to consider all possible truth value combinations for the propositions involved.
Let's create a truth table for the argument:
$IJKH \equiv (I \cdot \sim J)K \equiv \sim H \sim (K \supset J)I \equiv J$
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TTFFTF
TFTTFF
TFTFFF
TFFTTT
TFFFTF
FTTTFF
FTTFFF
FTFTFT
FTFFFF
FFTTTT
FFTFTT
FFFTTT
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In the truth table, the argument is invalid because there is at least one row where the premises are true (denoted by 'T') and the conclusion is false (denoted by 'F').
The counterexample can be found in the row where I is true, H is true, K is false, and J is false.
Therefore, the correct answer is option C: Invalid. Counterexample when I and H are true and J and K are false.

Answer Option 5:
To determine the validity of the argument, we can use an indirect truth table:
Step 1: Create a truth table for all the propositions involved in the argument:
P Q R S T U
Step 2: Fill in the truth values for the premises and the conclusion:
P Q R S T U

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Step 3: Determine the truth values of the compound propositions in each row:
P Q R S T U
Step 4: Determine the truth value of the conclusion in each row:
P Q R S T U
Step 5: Determine if there is a row in the truth table where all the premises are true and the conclusion is false. If there is such a row, then the argument is invalid. If not, the argument is valid.
After completing the indirect truth table, we find that there is no row where all the premises are true and the conclusion is false. Therefore, the argument is valid.
Final Answer: A. 'Valid'

Answer Option 6:
To determine the validity of the argument, we can use indirect truth tables.
First, we create a truth table to evaluate the premises and the conclusion:
First, we create a truth table to evaluate the premises and the conclusion:
 J H K -(K⊂J) ≣J
T T J H K -(K⊃J) ≅
 J
J
 J
J
I J H K -(K>J) Ы T T T F F T
J
J
I J H K -(K⊃J) ≡J T T F T F F Next, we fill in the truth values for the given statements: I J H K -(K⊃J) ≡J T T T F
i J H K -⟨K⊃J) ≡ i T i F F F Next, we fill in the truth values for the given statements: i J H K -⟨K⊃J) ≡ i T i F i F

Now, let's consider the conclusion, I=J. We can fill in the truth values for that statement: $I \mid J \mid H \mid K \mid {\sim}(K {\supset} J) \mid I {\equiv} J \mid$ TITIIIFITI TIFIIITIFI FITIIITIFI F|F|||T|T| Looking at the truth table, we see that there is at least one row (second row) in which all the premises are true, but the conclusion is false. Therefore, the argument is invalid. The correct answer is option B: 'Invalid. Counterexample when I, H, and K are true and J is false', which provides a specific counterexample where the premises are true and the conclusion is false. Answer Option 7: To determine the validity of the argument, we need to construct a truth table for all the propositions involved. $P\mid Q\mid R\mid S\mid (P\Rightarrow Q)\mid (R\Rightarrow S)\mid {\sim}((P\Rightarrow Q)\supset (R\Rightarrow S))\mid (Q\equiv S)$ T|T|T|T|T|F|T TITITIFITIFITIF TITIFITITITIFIT T|T|F|F|T|T|F|F TIFITITIFITITIF T|F|T|F|F|F|T|T T|F|F|T|F|T|F|F T|F|F|F|F|T|F|T F|T|T|T|T|T|F|TF|T|T|F|T|F|T|F F|T|F|T|T|T|F|T F|T|F|F|T|T|F|F F|F|T|T|T|T|F|T F|F|T|F|T|F|T|F F|F|F|T|T|T|F|T F|F|F|F|T|T|F|F From the truth table, we can see that the argument is invalid. There is a row where the premises are true (I, H, and K are true and J is false) but the conclusion (I = J) is false. Therefore, the correct answer is B. 'Invalid. Counterexample when I, H, and K are true and J is false' Answer Option 8: To determine the validity of the argument, we can construct an indirect truth table. We will assume the premises are true and check if the conclusion must also be true. $H \equiv (1 \cdot \sim J)$ K ≡ ~H ~(K ⊃ J) $I \equiv J$ We will consider all possible truth value combinations for the propositions I, J, H, and K. Starting with I, J, H, and K as false: I | J | H | K | K ⊃ J | ~(K ⊃ J) | I ≡ J F|F|F|T|T|F|T In this case, the conclusion $I\equiv J$ is true. Next, let's consider the case where I, J, H, and K are all true: $I \mid J \mid H \mid K \mid K \supset J \mid \sim (K \supset J) \mid I \equiv J$

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In this case, the conclusion $I\equiv J$ is also true. Now, let's examine the options for counterexamples: Option A: Valid - The argument is valid. This option does not provide any counterexample. Option B: Invalid. Counterexample when I, H, and K are true and J is false - We have already considered this case in our table, and the conclusion was true. So this is not a valid counterexample. Option C: Invalid. Counterexample when I and H are true and J and K are false - Let's consider this case in our table: $| | J | H | K | K \supset J | \sim (K \supset J) | I \equiv J$ T|F|T|F|T|F|F In this case, the conclusion $I \equiv J$ is false, which means this option provides a valid counterexample. Option D: Invalid. Counterexample when K is true and I, H, and J are false - Let's consider this case in our table: $| | J | H | K | K \supset J | \sim (K \supset J) | I \equiv J$ F|F|F|T|F|T|T In this case, the conclusion $I \equiv J$ is true, so this option does not provide a valid counterexample. Therefore, the correct answer is: C. Invalid. Counterexample when I and H are true and J and K are false Answer Option 9: To determine the validity of the argument, we will use indirect truth tables. Let's create a truth table to check the validity of the argument: $I\mid J\mid H\mid K\mid K\supset J\mid {\sim}(K\supset J)\mid I\equiv J$ T|T|||| T|F|||| F|T|||| F|F|||| Based on the given information, we can fill in the values for H and K: $I | J | H | K | K \supset J | \sim (K \supset J) | I \equiv J$ T|T|||T T|F|||| F|T|||| F|F|||| Now, let's fill in the values for H and K using the given equivalences: I | J | H | K | K ⊃ J | ~(K ⊃ J) | I ≡ J T|T|||T T|F|||| F|T|||| F|F|||| From the given equivalences, we can deduce that: $H\equiv (I\cdot \sim J)$ K = ~H Now, let's fill in the values for H and K: $I \mid J \mid H \mid K \mid K \supset J \mid \sim (K \supset J) \mid I \equiv J$ $\mathsf{T} \, | \, \mathsf{T} \, | \, | \, | \, | \, | \, |$ T|F|||| F|T|||| F|F|||| Using the equivalences $H \equiv (I \cdot \sim J)$ and $K \equiv \sim H$, we can calculate the values for H and K:

 $I \mid J \mid H \mid K \mid K \supset J \mid \sim (K \supset J) \mid I \equiv J$

TITITIFIII TIFIFITIII F|T|F|T||| F|F|T|F||| Now, let's fill in the values for $K\supset J$ and $\sim (K\supset J)$: $| | J | H | K | K \supset J | \sim (K \supset J) | | \equiv J$ T|T|T|F|T|F| TIFIFITITIFI F|T|F|T|T|F| F|F|T|F|T|F| Finally, let's fill in the values for I \equiv J: $I | J | H | K | K \supset J | \sim (K \supset J) | I \equiv J$ TITITIFITIFIT TIFIFITITIFIF F|T|F|T|T|F|F F|F|T|F|T|F|T Now, let's check if the argument is valid or invalid: The argument states ~(K \supset J) / I \equiv J Looking at the truth table, we can see that when the premise $\neg (K \supset J)$ is true (denoted by F), the conclusion $I \equiv J$ is also true (denoted by T). Therefore, the argument is valid. Final Answer: A

Question 110:

Construct a complete truth table for the following pairs of propositions. Then, using the truth tables, determine whether the statements are logically equivalent or contradictory. If neither, determine whether they are consistent or inconsistent. Justify your answers.

 $\sim (J \vee K) \cdot L \text{ and } (L \supset J) \cdot K', \text{ A. 'Logically equivalent', B. 'Contradictory', C. 'Neither logically equivalent nor contradictory, but consistent', D. 'Inconsistent', D. 'Inc$

Answer Option 1:

To construct a complete truth table for the given propositions, we consider all possible combinations of truth values for the propositions J, K, and L.

 $J \ K \ L \ {\sim} (J \lor K) \cdot L \ (L \supset J) \cdot K$

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From the truth table, we can see that $\neg(J \lor K) \cdot L$ and $(L \supset J) \cdot K$ have different truth values for some combinations of J, K, and L. Therefore, the statements are not logically equivalent.

Also, there is no combination of J, K, and L for which \sim (J \vee K) \cdot L and (L \supset J) \cdot K are both true or both false. Therefore, the statements are not contradictory.

Hence, the statements are neither logically equivalent nor contradictory, but consistent.

Final Answer: C. 'Neither logically equivalent nor contradictory, but consistent'

Answer Option 2:

To determine whether the statements "-(J \vee K) \cdot L" and "(L \supset J) \cdot K" are logically equivalent, contradictory, consistent, or inconsistent, we need to construct a truth table for both statements and compare the resulting truth values.

Let's construct the truth table for both statements:

 $J \ K \ L \ {\sim} (J \lor K) \cdot L \ (L \supset J) \cdot K$

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From the truth table, we can see that the truth values of the two statements are not always the same. Therefore, the statements " $-(J \lor K) \cdot L$ " and " $(L \supset J) \cdot K$ " are not logically equivalent.

Furthermore, there is no row in the truth table where both statements have a truth value of true. Therefore, the statements are not consistent.

The statements " $(J \vee K) \cdot L$ " and " $(L \supset J) \cdot K$ " are neither logically equivalent nor contradictory, but they are inconsistent. Thus, the correct answer is C. 'Neither logically equivalent nor contradictory, but consistent.'

Answer Option 3:

To construct a complete truth table for the given propositions, we will consider all possible combinations of truth values for J, K, and L.

 $J \mathrel{K} L \mathrel{\sim} (J \mathrel{\vee} K) \cdot L \mathrel{(L \mathrel{\supset} J)} \cdot K$

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FFFTT

By comparing the truth values for \sim (J \vee K) \cdot L and (L \supset J) \cdot K, we can see that they are not equivalent. The truth values for \sim (J \vee K) \cdot L and (L \supset J) \cdot K differ in rows 2, 3, 5, and 6. Therefore, the statements \sim (J \vee K) \cdot L and (L \supset J) \cdot K are not logically equivalent.

Final Answer: C. 'Neither logically equivalent nor contradictory, but consistent' Answer Option 4: Let's construct a truth table for the propositions: $|\:J\:|\:K\:|\:L\:|\: \hbox{$\scriptstyle \sim$} (J \lor K) \cdot L\:|\: (L \supset J) \cdot K\:|$ |---|---| |T|T|T|F|T||T|T|F|F|F| |T|F|T|F|F| |T|F|F|F|F| |F|T|T|F|T| |F|T|F|T|F| |F|F|T|T|T| |F|F|F|T|T| Looking at the truth table, we can see that \sim (J \vee K) \cdot L and (L \supset J) \cdot K do not have the same outputs for all combinations of truth values for J, K, and L. Therefore, the statements are not logically equivalent. However, there is no combination of truth values for J, K, and L that makes both propositions true or both propositions false. Therefore, the statements are not contradictory. Hence, the answer is C. 'Neither logically equivalent nor contradictory, but consistent'. Answer Option 5: To construct a complete truth table for the given propositions, let's analyze the individual propositions first: ~(J ∨ K) · L: $J \mathrel{K} \sim (J \mathrel{\vee} K) \sim (J \mathrel{\vee} K) \cdot L$ TTFF TEEE FTFT FFTF $(L\supset J)\cdot K$: $\mathsf{J}\,\mathsf{K}\,\mathsf{L}\,(\mathsf{L}\supset\mathsf{J})\,(\mathsf{L}\supset\mathsf{J})\cdot\mathsf{K}$ TTTTT TTFTF TETET TFFTF FTTTT FTFTF FFTTT FFFTF Now let's combine the two propositions and create a truth table for their conjunction: $J \mathrel{K} \mathrel{L} \mathord{\sim} (J \mathrel{\vee} \mathsf{K}) \cdot \mathrel{L} (L \mathrel{\supset} J) \cdot \mathrel{K} \mathord{\sim} (J \mathrel{\vee} \mathsf{K}) \cdot \mathrel{L} \wedge (L \mathrel{\supset} J) \cdot \mathsf{K}$ TTTFTF TTFFTF TFTTFF TEFETE FTTFTF FTFFTF

However, they are also not contradictory, as there are rows where both statements evaluate to true (rows 1, 4, and 7). Therefore, the statements –(J \vee K) \vee L and (L \supset J) \vee K are consistent.

By examining the last column of the truth table, we can see that \sim (J \vee K) \cdot L and (L \supset J) \cdot K are not logically equivalent since they have different truth values for some combinations of J, K, and L.

Since they are not logically equivalent, we can conclude that they are neither contradictory nor consistent. Therefore, the correct answer is:

C. "Neither logically equivalent nor contradictory, but consistent" $\,$

Answer Option 6:

FFTTTT FFFTTF

To construct the truth table, we need to consider all possible combinations of truth values for the propositions J, K, and L.

```
J \mathrel{K} L \mathrel{\sim} (J \mathrel{\vee} K) \cdot L \mathrel{(L \mathrel{\supset} J)} \cdot K
TTTFT
TTFFT
TFTFF
TEFEE
FTTFT
FTFFT
FFTTT
FFFTT
Now let's compare the truth values for \sim (J \vee K) \cdot L and (L \supset J) \cdot K:
J \mathrel{K} L \mathrel{\sim} (J \mathrel{\vee} K) \cdot L \mathrel{(L \mathrel{\supset} J)} \cdot K
TTTFT
TTFFT
TFTFF
TEFEE
 FTTFT
FTFFT
 FFTTT
FFFTT
 From the truth table, we can see that \sim (J \vee K) \cdot L and (L \supset J) \cdot K have the same truth values for all possible combinations of truth values for J, K, and L. Therefore, \sim (J \vee K) \cdot L and (L \supset J) \cdot K have the same truth values for all possible combinations of truth values for J, K, and L. Therefore, \sim (J \vee K) \cdot L and (L \supset J) \cdot K have the same truth values for all possible combinations of truth values for J, K, and L. Therefore, \sim (J \vee K) \cdot L and (L \supset J) \cdot K have the same truth values for J.
 equivalent.
Final Answer: A. 'Logically equivalent'
 *******
Answer Option 7:
To construct a complete truth table for the given pair of propositions, we need to consider all possible truth values for the variables J, K, and L.
Let's break it down step by step:
 1. \sim(J \vee K) \cdot L
J \ K \ L \ \text{$^{\sim}$} (J \ \lor \ K) \ \text{$^{\sim}$} (J \ \lor \ K) \cdot L
TTTFF
TTFFF
TFTFT
TEEFE
FTTFF
FTFTF
FFTTT
FFFTF
2. (L \supset J) \cdot K
J\:K\:L\:L\supset J\:(L\supset J)\cdot K
TTTTT
TTFTT
TETTE
 TFFTF
FTTTT
FTFTT
FFTTF
 FFFTT
 Now, let's compare the two truth tables to determine whether the statements are logically equivalent, contradictory, consistent, or inconsistent.
Comparing the two truth tables, we can see that for every combination of truth values for J, K, and L, the values of \sim (J \vee K) \vee L and (L \supset J) \vee K are the same.
 Therefore, the statements {\sim}(J\vee K)\cdot L and (L\supset J)\cdot K are logically equivalent.
 Final Answer: A. 'Logically equivalent'
```

Answer Option 8: First, let's construct the truth table for the given pairs of propositions: $J \mathrel{K} L \mathrel{\sim} (J \mathrel{\vee} K) \cdot L \mathrel{(L \mathrel{\supset} J)} \cdot K$ TTTFT TTFFF TETEE TFFFF FTTFF FTFFF FFTTT FFFTF Now, let's compare the truth values of ${\sim}(J\vee K)\cdot L$ and $(L\supset J)\cdot K$ in each row. In rows 1, 2, 3, 4, 5, and 6, \sim (J \vee K) \cdot L evaluates to F and (L \supset J) \cdot K evaluates to F. Therefore, \sim (J \vee K) \cdot L and (L \supset J) \cdot K are logically equivalent in these rows. In rows 7 and 8, \sim (J \vee K) \cdot L evaluates to T and (L \supset J) \cdot K evaluates to T. Therefore, \sim (J \vee K) \cdot L and (L \supset J) \cdot K are logically equivalent in these rows as well. Since \sim (J \vee K) \cdot L and (L \supset J) \cdot K have the same truth values in all rows, they are logically equivalent. Therefore, the correct answer is: A. 'Logically equivalent'. Answer Option 9: To construct a complete truth table for the given expressions, we need to consider all possible combinations of truth values for the propositional variables J, K, and L. Let's start by creating a truth table for $\text{--}(J \vee K) \cdot L$: $J \mid K \mid L \mid J \vee K \mid {\sim} (J \vee K) \mid {\sim} (J \vee K) \cdot L$ TITITIFIE TITIFITIFIF T|F|T|T|F|F TIFIFITIFIF F|T|T|T|F|F FITIFITIFIF F|F|T|F|T|T F|F|F|F|T|F Now, let's create a truth table for (L \supset J) \cdot K: $J \,|\, K \,|\, L \,|\, L \supset J \,|\, (L \supset J) \cdot K$ $T \mid T \mid T \mid T \mid T$ TITIFIFIF T|F|T|T|F TIFIFIFIF F|T|T|T|TFITIFITIT F|F|T|T|F F|F|F|T|F Now, let's compare the columns of the two truth tables to determine if the statements are logically equivalent, contradictory, consistent, or inconsistent. For each row, we compare the values of ${\sim}(J\vee K)\cdot L$ and $(L\supset J)\cdot K$: $J \mid K \mid L \mid \sim (J \lor K) \cdot L \mid (L \supset J) \cdot K$ TITITIFIT T|T|F|F|F TIFITIFIF T|F|F|F|F FITITIFIT

F|T|F|F|T F|F|T|T|F F|F|F|F|F Based on the truth table, we can see that $\sim (J \vee K) \cdot L$ and $(L \supset J) \cdot K$ have different truth values for some rows (e.g., the second and third rows). Therefore, the two statements are not logically equivalent.

However, they are also not contradictory since there are some rows where they have the same truth value (e.g., the first and eighth rows).

Therefore, the two statements are neither logically equivalent nor contradictory, but consistent.

Final Answer: C. Neither logically equivalent nor contradictory, but consistent.

Question 111:

 $\label{eq:controller} \ ^{\prime} \ Select \ the \ best \ translation \ into \ predicate \ logic: \ All \ first-year \ students \ take \ a \ swim \ test.', A. \ '(\forall x)(Sx\supset Fx)', B. \ '(\forall x)(Fx\supset Sx)', C. \ '(\forall x)(Fx\supset Sx)', D. \ '(\forall x)(Fx\subseteq Sx)', D$

Answer Option 1:

The best translation into predicate logic for the statement "All first-year students take a swim test" would be:

A. $(\forall x)(Sx \supset Fx)$

In this translation, Sx represents "x takes a swim test" and Fx represents "x is a first-year student". The symbol \supset represents implication, indicating that if someone takes a swim test (Sx), then they are a first-year student (Fx). The universal quantifier (\forall x) signifies that this statement applies to all individuals x.

Final Answer: A

Answer Option 2:

The best translation into predicate logic for the statement "All first-year students take a swim test" would be:

A. $(\forall x)(Sx \supset Fx)$

This translation states that for all individuals 'x', if 'x' takes a swim test (Sx), then 'x' is a first-year student (Fx).

The other options do not accurately capture the meaning of the original statement. Option B states that if someone is a first-year student, then they take a swim test, which is not the same as saying all first-year students take a swim test. Option C states that all individuals are either first-year students or take a swim test, which also does not capture the intended meaning. Option D states that being a first-year student is equivalent to taking a swim test, which is not accurate in this context.

Final Answer: A

Answer Option 3:

The best translation into predicate logic for the statement "All first-year students take a swim test" would be:

A. $(\forall x)(Sx \supset Fx)$

In this translation, " $\forall x$ " quantifies over all individuals x, "Sx" represents "x takes a swim test", and "Fx" represents "x is a first-year student". The proposition $Sx \supset Fx$ captures the idea that if someone takes a swim test (Sx), then they must be a first-year student (Fx). This translation accurately represents the statement that all first-year students take a swim test. Final Answer: A

Answer Option 4:

The best translation into predicate logic for the statement "All first-year students take a swim test" is:

$A.~(\forall x)(Sx\supset Fx)$

In this translation, Sx represents "x takes a swim test" and Fx represents "x is a first-year student". The arrow (\supset) represents implication, indicating that if someone takes a swim test (Sx), then they must be a first-year student (Fx). The universal quantifier (\forall x) indicates that this statement applies to all individuals.

Therefore, the best translation is A. $(\forall x)(Sx\supset Fx)$

Answer Option 5:

The best translation into predicate logic for the statement "All first-year students take a swim test" would be:

A. $(\forall x)(Sx\supset Fx)$

Here, Sx represents "x takes a swim test" and Fx represents "x is a first-year student". The symbol \supset represents implication, so $Sx \supset Fx$ translates to "If x takes a swim test, then x is a first-year student". The universal quantifier $(\forall x)$ indicates that the statement applies to all individuals x.

Final Answer: A

Answer Option 6:

The best translation into predicate logic for the statement "All first-year students take a swim test" would be:

A. $(\forall x)(Sx \supset Fx)$

In this translation, $(\forall x)$ represents "for all x" or "all," Sx represents "x takes a swim test," and Fx represents "x is a first-year student." The symbol \supset represents the implication or "if...then" relationship. Therefore, the translation $(\forall x)(Sx \supset Fx)$ accurately represents the statement that all first-year students take a swim test.

Final Answer: A

Answer Option 7:

The best translation into predicate logic for the statement "All first-year students take a swim test" is:

A. $(\forall x)(Sx \supset Fx)$

Here's how the translation works:

- Let Sx represent "x is a swim test"
- Let Fx represent "x is a first-year student"

The statement "All first-year students take a swim test" can be translated as "For every x, if x is a swim test (Sx), then x is a first-year student (Fx)."

Answer Option 8:
The best translation into predicate logic for the statement "All first-year students take a swim test" is: A. $(\forall x)(Sx \supset Fx)'$
Explanation:
Let Sx represent "x is a student" and Fx represent "x is a first-year student" and Tx represent "x takes a swim test."
The statement "All first-year students take a swim test" can be translated as "For all x, if x is a first-year student, then x takes a swim test."
This can be represented in predicate logic as $'(\forall x)(Sx \supset Fx)'$, where $(\forall x)$ denotes "for all x" and \supset denotes "implies".
Therefore, the correct translation is A. $'(\forall x)(Sx \supset Fx)'$.
••••••
Answer Option 9:
The best translation into predicate logic for the statement "All first-year students take a swim test" is:
$A. \ (\forall x)(Sx\supset Fx)$
Explanation:
- Let Sx represent "x takes a swim test"
- Let Fx represent "x is a first-year student"
- $(\forall x)$ represents "for all x" or "for every x"
The translation $(\forall x)(Sx \supset Fx)$ states that for every x, if x takes a swim test (Sx) , then x is a first-year student (Fx) . This accurately represents the statement "All first-year students take a swim test".
Option B: $(\forall x)(Fx\supset Sx)$ translates to "For every x, if x is a first-year student (Fx), then x takes a swim test (Sx)". This does not capture the meaning of the original statement.
Option C: $(\forall x)(Fx \lor Sx)$ translates to "For every x, x is a first-year student (Fx) or x takes a swim test (Sx)". This allows for the possibility that some first-year students may not take a swim test, which is not
what the original statement implies.
Option D: $(\forall x)(Fx \equiv Sx)$ translates to "For every x, x is a first-year student (Fx) if and only if x takes a swim test (Sx)". This implies that taking a swim test is a necessary and sufficient condition for being a

Therefore, the correct answer is A. $(\forall x)(Sx\supset Fx).$

first-year student, which is not necessarily true.

Final Answer: A

Question 112:
'Select the best translation into predicate logic: Johnny came late.', A. 'jL', B. 'Lj', C. 'Cj v Lj', D. 'Lx'
Answer Option 1:
The best translation into predicate logic for the statement "Johnny came late" would be option D: "Lx".
This represents that there exists an individual x (which can be interpreted as Johnny) such that x came late (Lx).
••••••••••
Answer Option 2:
The best translation into predicate logic for the statement "Johnny came late" is option B, which is 'Lj'. This statement represents that Johnny (j) came late (L). Therefore, 'Lj' accurately translates the given
statement. Final Answer: B
•••••••
Answer Option 3:
The best translation into predicate logic for the sentence "Johnny came late" would be option A: 'jL'. Here, 'j' represents Johnny and 'L' represents came late. The statement 'jL' translates to "Johnny came
late." Final Answer: A
Answer Option 4:
The best translation into predicate logic for the sentence "Johnny came late" would be option B: "Lj".
Here, "L' represents the predicate "came late", and "j' represents the individual constant "Johnny". Therefore, "Lj' translates to "Johnny came late". Final Answer: B
Answer Option 5:
The best translation into predicate logic for the statement "Johnny came late" would be A. 'jL'.
Here, j represents Johnny and L represents the predicate "came late". The statement can be translated as "Johnny came late", which can be represented as jL in predicate logic. Final Answer: A.
Answer Option 6:
The best translation into predicate logic for the statement "Johnny came late" would be option A: 'jL'. This translates to "There exists an individual j such that j came late." Here, j represents Johnny and L
represents the condition of coming late. This translation captures the essence of the statement accurately. Final Answer: A
Answer Option 7:
The best translation into predicate logic for the statement "Johnny came late" would be: C. $C_j \vee L_j$
Explanation:
- 'Cj' represents the predicate "Johnny came".
- 'Lj' represents the predicate "Johnny is late".
Therefore, the logical expression 'Cj v Lj' correctly translates the statement "Johnny came late" into predicate logic. Final Answer: C

Answer Option 8:
The best translation into predicate logic for the statement "Johnny came late" would be:
Cj, where Cj represents "Johnny came".
Option A: 'jL' would translate to "Johnny is late", which is not the same as "Johnny came late".
Option B: 'Lj' would translate to "Late is Johnny", which does not make sense in this context.
Option C: 'Cj v Lj' would translate to "Johnny came or late is Johnny", which is not an accurate representation of the original statement.
Option D: 'Lx' would translate to "Someone is late", which is not specific to Johnny.
Therefore, the correct answer is A. 'Cj'

Answer Option 9:
The best translation into predicate logic for the statement "Johnny came late" would be: Lj.
This translates to "There exists an individual, Johnny, such that Johnny is late."
Option A, "jL", translates to "There exists an individual, j, such that j is late," which does not accurately represent the statement.
Option B, 'Lj', translates to "Johnny is late," which accurately represents the statement.
Option C, 'Cj v Lj', translates to "There exists an individual, Cj, such that Cj or Lj," which introduces an unnecessary predicate Cj.
Option D, 'Lx', translates to "There exists an individual, x, such that x is late," which is too general and does not specifically refer to Johnny.
Therefore, the best translation into predicate logic is option B, 'Li'.

Question 113:

'Select the best translation into predicate logic: No senior graduates without a major.', A. ' $(\forall x)(Sx \supset -Gx)'$, B. ' $(\neg (\forall x)(Sx \supset Gx)', C$. ' $(\forall x)(\neg Sx \supset \neg Gx)', C$. ' $(x)(\neg Sx \supset \neg$

Answer Option 1:

The best translation into predicate logic for the statement "No senior graduates without a major" is option C: $(\forall x)(-Sx \supset -Gx)$. This translation states that for all individuals x, if x is not a senior (-Sx), then x does not graduate without a major (-Gx). This captures the meaning of the original statement accurately. Final Answer: C.

Answer Option 2:

The best translation into predicate logic for the statement "No senior graduates without a major" is:

C. $(\forall x)(\sim Sx \supset \sim Gx)$

This translates to "For all x, if x is not a senior (\neg Sx), then x does not graduate without a major (\neg Gx)".

Answer Option 3:

The best translation into predicate logic for the statement "No senior graduates without a major" is:

C. $(\forall x)(\sim Sx \supset \sim Gx)$

This translation states that for every individual x, if x is not a senior (-Sx), then x does not graduate without a major (-Gx). In other words, it captures the idea that there are no seniors who graduate without a major. Final Answer: C

Answer Option 4:

The best translation into predicate logic for the statement "No senior graduates without a major" is:

A. $(\forall x)(Sx \supset \sim Gx)$

This translation states that for all individuals x, if x is a senior (Sx), then it is not the case that x graduates without a major (-Gx). This accurately captures the meaning of the original statement.

Therefore, the correct answer is A.

Answer Option 5:

To translate the given statement into predicate logic, we can break it down as follows: "No senior graduates without a major" Let's define our predicates: S(x): x is a senior, G(x): x graduates, M(x): x has a major. The logical structure of the statement is: $(\forall x)$ ($S(x) \rightarrow (G(x) \land M(x))$) Breaking it down: $(\forall x)$: For all x (seniors) $S(x) \rightarrow (G(x) \land M(x))$: If x is a senior, then x graduates and has a major. Therefore, the best translation into predicate logic is: A: $(\forall x)$ ($S(x) \rightarrow G(x)$) Final Answer: A

Answer Option 6:

The best translation into predicate logic for the statement "No senior graduates without a major" is option A: $(\forall x)(Sx \supset \sim Gx)$.

Let's break down the statement:

- "No senior graduates" translates to $(\forall x)$ ~ $(Sx \to Gx)$, which means for all x, it is not the case that if x is a senior, then x graduates.
- "without a major" translates to ~Gx, meaning x does not graduate.

So, combining the two translations, we get $(\forall x)(Sx \supset \neg Gx)$, which means for all x, if x is a senior, then x does not graduate. This is the best translation of the given statement.

Final Answer: A

Answer Option 7:

The best translation into predicate logic for the statement "No senior graduates without a major" is:

A. $(\forall x)(Sx \supset \sim Gx)$

This translates to: For all x, if x is a senior, then it is not the case that x graduates without a major.

This accurately captures the meaning of the original statement. Option B is incorrect because it translates to "It is not the case that for all x, if x is a senior, then x graduates without a major," which is not the same as the original statement. Option C is incorrect because it translates to "For all x, if x is not a senior, then it is not the case that x graduates without a major," which is not the same as the original statement. Option D is incorrect because it translates to "For all x, it is not the case that x graduates without a major," which does not capture the restriction to seniors in the original statement. Final Answer: A

Answer Option 8:

The given statement, "No senior graduates without a major" can be translated into predicate logic as:

 $(\forall x)(Sx\supset Gx), \text{ where } Sx \text{ represents "}x \text{ is a senior" and } Gx \text{ represents "}x \text{ graduates"}.$

Option A: $(\forall x)(Sx \supset \sim Gx)$

This translates to "For all x, if x is a senior, then x does not graduate", which does not accurately represent the given statement.

Option B: \sim ($\forall x$)($Sx \supset Gx$)

This translates to "It is not the case that for all x, if x is a senior, then x graduates", which is not the correct translation.

Option C: $(\forall x)(\sim Sx \supset \sim Gx)$

This translates to "For all x, if x is not a senior, then x does not graduate", which does not accurately represent the given statement.

Option D: $(\forall x) \sim (Sx \supset \sim Gx)$

This translates to "For all x, it is not the case that if x is a senior, then x does not graduate", which accurately represents the given statement.

Therefore, the correct translation into predicate logic is option D: $(\forall x) \sim (Sx \supset \sim Gx)$.

Answer Option 9:

The statement "No senior graduates without a major" can be translated into predicate logic as follows:

Let Sx represent "x is a senior" and Gx represent "x graduates" and Mx represent "x has a major".

The best translation would be:

 $(\forall x)(Sx \to (Gx \land Mx))$

This states that for all individuals x, if x is a senior, then x graduates and has a major.

Among the given options:

 $A. \ (\forall x) (Sx \supset \neg Gx) \ - \ This \ translates \ to \ "For \ all \ individuals \ x, \ if \ x \ is \ a \ senior, \ then \ x \ does \ not \ graduate", \ which \ is \ not \ the \ intended \ meaning.$

 $B. \sim (\forall x)(Sx \supset Gx) - \text{This translates to "It is not the case that for all individuals } x, \text{ if } x \text{ is a senior, then } x \text{ graduates", which is not the intended meaning.}$

 $C. \ (\forall x) (\neg Sx \supset \neg Gx) \ - \ This \ translates \ to \ "For \ all \ individuals \ x, \ if \ x \ is \ not \ a \ senior, \ then \ x \ does \ not \ graduate", \ which \ is \ not \ the \ intended \ meaning.$

D. $(\forall x)$ ~ $(Sx \supset \neg Gx)$ - This translates to "For all individuals x, it is not the case that if x is a senior, then x does not graduate", which is not the intended meaning.

Therefore, the correct translation is A. $(\forall x)(Sx \rightarrow (Gx \land Mx))$.

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'Identify the conclusion of the following argument. From a historical point of view, genetics is a young science. The principles that govern inheritance traits were described less than 150 years ago.

Furthermore, around the turn of the twentieth century, the laws of inheritance were rediscovered. But even then, the importance of DNA was not really understood until the 1950s.', A. 'But even then, the importance of DNA was not really understood until the 1950s.', C. 'From a historical point of view, genetics is a young science.', D. 'The principles that govern inheritance traits were described less than 150 years ago.'

Answer Option 1:

The conclusion of the argument is:

C. 'From a historical point of view, genetics is a young science.'

Answer Option 2:

The conclusion of the argument is: C. 'From a historical point of view, genetics is a young science.' The argument provides historical evidence and facts to support the conclusion that genetics is a young science.

Answer Option 3:

The conclusion of the argument is the main point or claim that the author is trying to make. In this case, the conclusion can be identified as:

B. 'The importance of DNA was not really understood until the 1950s.'

Answer Option 4:

The conclusion of the argument is the main point or claim being made based on the information provided. In this case, the conclusion can be identified as: B. The importance of DNA was not really understood until the 1950s. This statement summarizes the main idea conveyed in the argument that despite the principles of inheritance being described and rediscovered, it was not until the 1950s that the significance of DNA was truly grasped. Final Answer: B

Answer Option 5:

The conclusion of the argument is: 'The importance of DNA was not really understood until the 1950s.'

This is supported by the statement: "But even then, the importance of DNA was not really understood until the 1950s."

Therefore, the correct answer is B. 'The importance of DNA was not really understood until the 1950s.'

Answer Option 6:

The conclusion of the argument is the main point or the final statement that is being made. In this case, the conclusion is: 'From a historical point of view, genetics is a young science.' This statement summarizes the main idea of the argument and is supported by the information provided about the history and development of genetics. Therefore, the correct answer is option C: 'From a historical point of view, genetics is a young science.'

Answer Option 7:

The conclusion of the argument can be identified as the main point or the final statement that the author is trying to convey. In this case, the conclusion can be found at the end of the passage: "But even then, the importance of DNA was not really understood until the 1950s." Therefore, the correct answer is A. 'But even then, the importance of DNA was not really understood until the 1950s.'

. . . .

Answer Option 8

The conclusion of the argument is the main point or the final statement that the argument is trying to make. In this case, the conclusion of the argument is: C. 'From a historical point of view, genetics is a young science.' This is because the argument presents various points and evidence to support the idea that genetics is a relatively new field, and this is the main point that is being emphasized throughout the argument. Final Answer: C

Answer Option 9:

The conclusion of the argument is the main point or the ultimate claim that the argument is trying to make. In this case, the conclusion can be identified as:

B. 'The importance of DNA was not really understood until the 1950s.'

This is because the argument mentions the historical development of genetics and how the understanding of DNA's importance came much later in the 1950s. Therefore, the conclusion can be inferred to be that the understanding of DNA's importance did not come until the 1950s. Final Answer: B. The importance of DNA was not really understood until the 1950s.'

Question 115:

'Which of the given formulas of PL is the best symbolization of the following sentence?

Either England's importing beef is a necessary condition for France's subsidizing agriculture or China promotes human rights given that South Africa doesn't supply diamonds.', A. '(E \supset F) \lor (C \equiv \sim S)', B. '(F \supset

E) \lor (\sim S \equiv C)', C. '(E \supset F) \lor (S $\supset\sim$ C)', D. '(F \supset E) \lor (\sim S \supset C)'

Answer Option 1:

Let's break down the given sentence:

"Either England's importing beef is a necessary condition for France's subsidizing agriculture or China promotes human rights given that South Africa doesn't supply diamonds."

Let's symbolize the different parts of the sentence:

England's importing beef is a necessary condition for France's subsidizing agriculture: $\mathsf{E} \to \mathsf{F}$

China promotes human rights given that South Africa doesn't supply diamonds: $\sim S \to C$

To represent the statement "Either A or B", we use the logical operator ">" (OR).

Combining the symbolized parts, the best representation of the given sentence is:

 $(\mathsf{E} \to \mathsf{F}) \lor (\mathsf{\sim}\mathsf{S} \to \mathsf{C})$

Therefore, the correct answer is: A. '(E \supset F) \vee (C \equiv ~S)'

Answer Option 2:

To symbolize the given sentence, let's break it down into two parts:

- England's importing beef is a necessary condition for France's subsidizing agriculture: This can be symbolized as E → F, where E represents England's importing beef and F represents France's subsidizing agriculture.
- 2. China promotes human rights given that South Africa doesn't supply diamonds: This can be symbolized as ¬S → C, where ¬S represents the negation of South Africa supplying diamonds and C represents

China promoting human rights.

Combining these two parts with the "or" condition, the best symbolization of the given sentence is:

 $(E \rightarrow F) \lor (\sim S \rightarrow C)$

Therefore, the correct answer is A. '(E \supset F) \vee (C \equiv ~S)'.

Answer Option 3:

Let's break down the given sentence:

- 1. England's importing beef is a necessary condition for France's subsidizing agriculture: $\mathsf{E} \supset \mathsf{F}$
- 2. China promotes human rights: C
- 3. South Africa doesn't supply diamonds: ~S

Now, let's consider the second part of the sentence "China promotes human rights given that South Africa doesn't supply diamonds". This can be symbolized as: ~S ⊃ C.

Combining both parts of the sentence with an "or" connector, we have:

 $(\mathsf{E}\supset\mathsf{F})\vee(\mathsf{\sim}\mathsf{S}\supset\mathsf{C})$

So, the best symbolization of the given sentence is Option D: '(F \supset E) \lor (\sim S \supset C)'. Final Answer: D

Answer Option 4:

To symbolize the given sentence, we need to break it down into its logical components.

Let:

E = England's importing beef

F = France's subsidizing agriculture

C = China promotes human rights

S = South Africa supplies diamonds

The sentence can be broken down into two parts:

1. "England's importing beef is a necessary condition for France's subsidizing agriculture"

This can be symbolized as $\mathsf{E}\supset\mathsf{F}$

2. "China promotes human rights given that South Africa doesn't supply diamonds"

This can be symbolized as ~S \supset C

Now, we need to connect these two parts using an "or" operator since they represent two different conditions. The correct symbolization of the sentence is:

 $(E\supset F)\lor (\sim S\supset C)$

Option D: '(F \supset E) \lor (\sim S \supset C)' is the best symbolization of the given sentence.

Answer Option 5:

To symbolize the given sentence, let's break it down into two parts:

- 1. England's importing beef is a necessary condition for France's subsidizing agriculture.
- 2. China promotes human rights given that South Africa doesn't supply diamonds.

For the first part, we can symbolize it as $(E\supset F)$, where E represents "England's importing beef" and F represents "France's subsidizing agriculture."

For the second part, we can symbolize it as (-S \supset C), where S represents "South Africa supplies diamonds" and C represents "China promotes human rights."

Now let's combine the two parts using the logical operator "or" since it says "either or":

 $(E \supset F) \lor (\sim S \supset C)$

Comparing this with the given options:

A. $(E\supset F)\lor (C\equiv \sim S)$

 $\mathsf{B.}\; (\mathsf{F}\supset \mathsf{E}) \vee (\mathsf{\sim}\mathsf{S} \equiv \mathsf{C})$

C. $(E \supset F) \lor (S \supset \sim C)$

 $\mathsf{D.}\; (\mathsf{F}\supset \mathsf{E}) \vee (\mathsf{\sim}\mathsf{S}\supset \mathsf{C})$

We can see that option A: $(E \supset F) \lor (C \equiv \sim S)$ is the best symbolization for the given sentence. Final Answer: A.

Answer Option 6:

Let's break down the given sentence into its components:

- 1. England's importing beef is a necessary condition for France's subsidizing agriculture: $\mathsf{E} \supset \mathsf{F}$
- 2. China promotes human rights: C
- 3. South Africa doesn't supply diamonds: ~S

The sentence states that either 1 or 2 is true.

Therefore, the correct symbolization would be (E \supset F) \vee C.

Comparing this with the given options:

Option A: (E \supset F) \lor (C \equiv \sim S) - This option also includes the condition of South Africa not supplying diamonds, which is not mentioned in the original sentence. Not the best symbolization.

Option B: $(F \supset E) \lor (-S \equiv C)$ - This option rearranges the conditions and includes the equivalence of South Africa not supplying diamonds and China promoting human rights, which is not mentioned in the original sentence. Not the best symbolization.

Option C: (E \supset F) \lor (S \supset C) - This option includes the condition of South Africa supplying diamonds, which is the opposite of what is mentioned in the original sentence. Not the best symbolization.

Option D: (F \supset E) \vee (\neg S \supset C) - This option rearranges the conditions and includes the condition of South Africa not supplying diamonds, which is not mentioned in the original sentence. Not the best symbolization.

Therefore, the best symbolization of the given sentence is option A: $(E\supset F)\lor C$. Final Answer: A.

Answer Option 7:

To symbolize the given sentence, we can break it down into two parts:

- 1. England's importing beef is a necessary condition for France's subsidizing agriculture.
- 2. China promotes human rights given that South Africa doesn't supply diamonds.

Let's analyze the options:

Option A: $'(E\supset F)\lor (C\equiv \sim S)'$

This option correctly symbolizes the first part of the sentence as $(E\supset F)$. However, the second part is symbolized as $(C\equiv -S)$, which does not capture the given condition.

Option B: $(F \supset E) \lor (\sim S \equiv C)$

This option incorrectly symbolizes the first part of the sentence as (F \supset E). The second part is symbolized as (~S = C), which is not the same as "given that South Africa doesn't supply diamonds."

Option C: $'(E\supset F)\lor (S\supset \sim C)'$

This option correctly symbolizes the first part of the sentence as (E \supset F). However, the second part is symbolized as (S \supset -C), which is not the same as "given that South Africa doesn't supply diamonds."

Option D: ${}^{\iota}(F\supset E)\vee ({}^{\star}S\supset C){}^{\iota}$

This option correctly symbolizes the first part of the sentence as (F \supset E). The second part is symbolized as (-S \supset C), which correctly represents "given that South Africa doesn't supply diamonds."

Therefore, the best symbolization of the sentence is Option D: '(F \supset E) \lor (\sim S \supset C)'. Final Answer: D

Answer Option 8:

Let's break down the given sentence:

- 1. Either England's importing beef is a necessary condition for France's subsidizing agriculture
- 2. or China promotes human rights given that South Africa doesn't supply diamonds.

Let's symbolize the individual parts:

- 1. England's importing beef is a necessary condition for France's subsidizing agriculture: $\mathsf{E} \supset \mathsf{F}$
- 2. China promotes human rights given that South Africa doesn't supply diamonds: $\sim S \supset C$

Now let's combine these symbolizations using the logical operator for "either/or", which is the inclusive OR operator (\vee).

Option A: $(E \supset F) \lor (C \equiv \sim S)$

The symbolization for the first part is correct, but the symbolization for the second part, (C = ~S), is not accurate. It should be ~S ⊃ C. So this option is not the best symbolization.

Option B: $(F \supset E) \lor (\sim S \equiv C)$

The symbolization for the first part, $(F \supset E)$, is not accurate. It should be $(E \supset F)$. The symbolization for the second part, $(-S \equiv C)$, is also not accurate. It should be $-S \supset C$. So this option is not the best symbolization.

Option C: $(E \supset F) \lor (S \supset \sim C)$

The symbolization for the first part is correct, but the symbolization for the second part, $(S \supset C)$, is not accurate. It should be $\neg S \supset C$. So this option is not the best symbolization.

Option D: $(F \supset E) \lor (\sim S \supset C)$

The symbolization for both parts is accurate. The first part, (F \supset E), corresponds to "England's importing beef is a necessary condition for France's subsidizing agriculture", and the second part, (-S \supset C), corresponds to "China promotes human rights given that South Africa doesn't supply diamonds". This option is the best symbolization of the given sentence.

Final Answer: D

.....

Answer Option 9:

To symbolize the sentence accurately, we need to break it down into its constituent parts and represent them using propositional logic symbols.

Let's analyze the sentence:

- 1. "Either England's importing beef is a necessary condition for France's subsidizing agriculture" can be symbolized as (E \supset F).
- This means that if England imports beef (E), then France will subsidize agriculture (F).
- 2. "China promotes human rights given that South Africa doesn't supply diamonds" can be symbolized as (S \supset ~C).
- This means that if South Africa does not supply diamonds (S), then China promotes human rights (~C).

Now, combining the two parts using the logical connective "or" (\vee), we get:

 $(E \supset F) \lor (S \supset \sim C)$

So, the best symbolization of the given sentence is option C: $'(E\supset F)\lor (S\supset \sim C)'.$

Question 116:

'Which of the given formulas of PL is the best symbolization of the following sentence?

Either England's importing beef is a sufficient condition for France's subsidizing agriculture or China's promoting human rights is a necessary condition for South Africa's supplying diamonds.', A. '(E > F) v (S

 $\supset C)', \ B. \ '(F\supset E) \lor (C\supset S)', \ C. \ '(E\supset F) \lor (C\supset S)', \ D. \ '(F\supset E) \lor (S\supset C)'$

Answer Option 1:

To symbolize the given sentence, we can break it down into two parts:

- 1. England's importing beef is a sufficient condition for France's subsidizing agriculture.
- 2. China's promoting human rights is a necessary condition for South Africa's supplying diamonds.

Let's analyze each option:

A. $(E \supset F) \lor (S \supset C)$

This option symbolizes the first part correctly but not the second part.

B. $(F \supset E) \lor (C \supset S)$

This option symbolizes the second part correctly but not the first part.

C. $(E \supset F) \lor (C \supset S)$

This option symbolizes both parts correctly, making it a potential correct answer.

D. $(F \supset E) \lor (S \supset C)$

This option symbolizes neither part correctly.

Based on the analysis, the best symbolization of the given sentence is option C: '(E \supset F) \lor (C \supset S)'. Final Answer: C

Answer Option 2:

To symbolize the given sentence, we can break it down into two parts:

- 1. England's importing beef is a sufficient condition for France's subsidizing agriculture
- 2. China's promoting human rights is a necessary condition for South Africa's supplying diamonds

For the first part, the symbolization should be in the form $(A\supset B)$ where A represents England's importing beef and B represents France's subsidizing agriculture.

For the second part, the symbolization should be in the form (C \supset D) where C represents China's promoting human rights and D represents South Africa's supplying diamonds.

Now, let's analyze the options:

 $\mathsf{A.}\; (\mathsf{E}\supset\mathsf{F}) \vee (\mathsf{S}\supset\mathsf{C})$

This option combines the two parts correctly, where E represents England's importing beef, F represents France's subsidizing agriculture, S represents South Africa's supplying diamonds, and C represents

China's promoting human rights. This is the best symbolization of the given sentence.

 $\mathsf{B.}\; (\mathsf{F}\supset\mathsf{E}) \vee (\mathsf{C}\supset\mathsf{S})$

This option interchanges the conditions for the first and second parts, which is not logically equivalent to the original sentence.

C. $(E \supset F) \lor (C \supset S)$

This option correctly symbolizes the first part, but it interchanges the conditions for the second part, which is not logically equivalent to the original sentence.

 $\mathsf{D}.\; (\mathsf{F}\supset \mathsf{E}) \vee (\mathsf{S}\supset \mathsf{C})$

This option interchanges the conditions for the first and second parts, which is not logically equivalent to the original sentence.

Therefore, the best symbolization of the given sentence is: $(E \supset F) \lor (S \supset C)$ which corresponds to option A. Final Answer: A

Answer Option 3:

In the given sentence, we can symbolize the following statements:

- E: England's importing beef
- F: France's subsidizing agriculture
- C: China's promoting human rights
- S: South Africa's supplying diamonds

The sentence states, "Either England's importing beef is a sufficient condition for France's subsidizing agriculture or China's promoting human rights is a necessary condition for South Africa's supplying diamonds."

To symbolize this, we can use the following form:

 $(E\supset F)\lor (C\supset S)$

Option A: $'(E\supset F)\lor (S\supset C)'$

Option B: ${}^{\prime}(F\supset E)\vee (C\supset S){}^{\prime}$

Option C: $'(E\supset F)\lor (C\supset S)'$

Option D: '(F \supset E) \lor (S \supset C)'

The best symbolization of the sentence is Option C: '(E \supset F) \lor (C \supset S)'. Final Answer: C

Answer Option 4:
To determine the best symbolization of the given sentence, let's break down the sentence into its logical components:
England's importing beef is a sufficient condition for France's subsidizing agriculture.
Symbolized as: E ⊃ F
2. China's promoting human rights is a necessary condition for South Africa's supplying diamonds.
Symbolized as: C⊃S
Now, when combining the two conditions using "or" (eitheror), we get:
$(E\supset F)\lor (C\supset S)$
Looking at the options:
A. $'(E\supset F)\lor (S\supset C)'$
This option has the conditions in the wrong order compared to the given sentence.
$B.\ '(F\supsetE)\lor (C\supsetS)'$
This option has the conditions swapped compared to the given sentence.
C. $'(E\supset F)\lor (C\supset S)'$
This option correctly represents the given sentence.
D. $(F \supset E) \lor (S \supset C)$
This option has the conditions swapped compared to the given sentence.
Therefore, the best symbolization of the given sentence is option C: $(E \supset F) \lor (C \supset S)$.
Answer Option 5:
Let's break down the sentence:
1. England's importing beef is a sufficient condition for France's subsidizing agriculture: This can be symbolized as E \supset F, where E represents England's importing beef and F represents France's subsidizing
agriculture.
2. China's promoting human rights is a necessary condition for South Africa's supplying diamonds: This can be symbolized as C \supset S, where C represents China's promoting human rights and S represents
South Africa's supplying diamonds.
Now, let's combine these two conditions using the logical operator "or" since the original sentence states "eitheror":
$(E\supset F)\lor (C\supset S)$
Comparing this with the given options, we can see that the best symbolization is option C: $(E \supset F) \lor (C \supset S)'$
Final Answer: C

Answer Option 6:
To symbolize the given sentence, we need to identify the main components and their logical relationships. Let:
E: England's importing beef
F: France's subsidizing agriculture
C: China's promoting human rights
S: South Africa's supplying diamonds
The sentence states:
(Either England's importing beef is a sufficient condition for France's subsidizing agriculture) OR (China's promoting human rights is a necessary condition for South Africa's supplying diamonds)
The best symbolization of this sentence is option C: $(E \supset F) \lor (C \supset S)'$
Explanation:
. (E ⊃ F) represents the statement "England's importing beef is a sufficient condition for France's subsidizing agriculture."
(C ⊃ S) represents the statement "China's promoting human rights is a necessary condition for South Africa's supplying diamonds."
The sentence states that either one of these conditions is true, hence the use of the logical OR (v) operator.
Final Answer: C
Answer Option 7:
Let's break down the sentence given:
"Either England's importing beef is a sufficient condition for France's subsidizing agriculture or China's promoting human rights is a necessary condition for South Africa's supplying diamonds."
Let's symbolize the propositions:
E: England's importing beef
F: France's subsidizing agriculture
C: China's promoting human rights

S: South Africa's supplying diamonds

To symbolize "England's importing beef is a sufficient condition for France's subsidizing agriculture," we use the implication operator (\supset) as follows:

 $(E \supset F)$

To symbolize "China's promoting human rights is a necessary condition for South Africa's supplying diamonds," we can use the implication operator (\supset) again, but with the propositions switched:

 $(C\supset S)$

To symbolize the entire sentence using the logical OR operator (v) for the word "either," we can combine the two symbolic expressions:

 $(E \supset F) \lor (C \supset S)$

Therefore, the best symbolization of the sentence is option A: $'(E\supset F)\vee (S\supset C)'$.

Answer Option 8:

To symbolize the given sentence, we need to identify the correct propositions and their relationships.

Let's assign the following propositions:

- E: England's importing beef
- F: France's subsidizing agriculture
- C: China's promoting human rights
- S: South Africa's supplying diamonds

The sentence states: Either England's importing beef is a sufficient condition for France's subsidizing agriculture or China's promoting human rights is a necessary condition for South Africa's supplying diamonds

The correct symbolization of this sentence is: (E \supset F) \lor (C \supset S)

Option A: '(E \supset F) \lor (S \supset C)'

This option swaps the propositions F and S, which is incorrect.

Option B: $'(F \supset E) \lor (C \supset S)'$

This option swaps the propositions E and F, which is incorrect.

Option C: $'(E \supset F) \lor (C \supset S)'$

This option correctly symbolizes the given sentence based on the assigned propositions.

Option D: ${}^{\prime}(F\supset E)\vee(S\supset C){}^{\prime}$

This option swaps the propositions ${\sf E}$ and ${\sf F},$ which is incorrect.

Therefore, the best symbolization of the given sentence is '(E \supset F) \vee (C \supset S)'.

Final Answer: 0

Answer Option 9:

To symbolize the given sentence, let's assign the following propositions:

- E: England's importing beef
- F: France's subsidizing agriculture
- C: China's promoting human rights
- S: South Africa's supplying diamonds

The sentence states that either England's importing beef is a sufficient condition for France's subsidizing agriculture or China's promoting human rights is a necessary condition for South Africa's supplying

The symbolization of this sentence can be represented as: $(E \supset F) \lor (C \supset S)$

Option A: (E \supset F) \vee (S \supset C) - This option does not retain the correct conditions for importing beef, subsidizing agriculture, promoting human rights, and supplying diamonds. It is not the best symbolization.

 $Option \ B: \ (F\supset E)\lor (C\supset S) \ -\ This \ option \ swaps \ the \ conditions \ for \ importing \ beef \ and \ subsidizing \ agriculture. \ It \ is \ not \ the \ best \ symbolization.$

Option C: $(E \supset F) \lor (C \supset S)$ - This option correctly symbolizes the given sentence. It retains the correct conditions for importing beef, subsidizing agriculture, promoting human rights, and supplying diamonds.

It is the best symbolization.

Option D: (F \supset E) \lor (S \supset C) - This option swaps the conditions for importing beef and subsidizing agriculture as well as promoting human rights and supplying diamonds. It is not the best symbolization.

Final Answer: C

Question 117:

'Which of the given formulas of PL is the best symbolization of the following sentence?

Dylan likes neither math nor science.', A. 'M \vee S', B. '~(M \bullet S)', C. '~(M \vee S)', D. '~M \vee ~S'

Answer Option 1:

The best symbolization of the sentence "Dylan likes neither math nor science" would be option C: '-(M v S)'. This translates to "It is not the case that Dylan likes math or science." This captures the meaning of Dylan not liking both math and science. Final Answer: C

Answer Option 2:

The sentence "Dylan likes neither math nor science" can be symbolized as "¬(M ∨ S)" because it means Dylan does not like either math or science. Therefore, the best symbolization of the sentence is option

C: '~(M v S)'. Final Answer: C.

Answer Option 3:

The statement "Dylan likes neither math nor science" can be symbolized as follows: Let M represent "Dylan likes math" and S represent "Dylan likes science". The word "neither" implies a negation.

Therefore, the correct symbolization would be: ~M ^ ~S. Option C, '~(M \ S)', is the best symbolization of the given sentence. Final Answer: C

Answer Option 4:

The sentence "Dylan likes neither math nor science" can be symbolized as "¬(M v S)" where M represents math and S represents science.

Option A: 'M v S' symbolizes that Dylan likes either math or science, which is not the correct interpretation of the given sentence.

Option B: '~(M • S)' symbolizes that Dylan does not like both math and science, which is not the correct interpretation of the given sentence.

Option C: '~(M v S)' is the correct symbolization of the sentence "Dylan likes neither math nor science". This option represents that Dylan does not like either math or science.

Option D: '~M ∨ ~S' symbolizes that Dylan dislikes either math or science, which is not the correct interpretation of the given sentence.

Therefore, the best symbolization of the sentence is option C: '~(M \vee S)'. Final Answer: C.

Answer Option 5:

The sentence "Dylan likes neither math nor science" can be symbolized as '~(M v S)', where M represents liking math and S represents liking science. This formula states that it is not the case that Dylan likes

Option A: 'M v S' represents the statement that Dylan likes math or science. This is not the best symbolization because it does not capture the meaning of the original sentence that Dylan likes neither math

Option B: '-(M • S)' represents the statement that it is not the case that Dylan likes both math and science. This is not the best symbolization because it does not capture the meaning of the original sentence that Dylan likes neither math nor science.

Option C: '~(M ∨ S)' is the best symbolization because it correctly represents the original sentence "Dylan likes neither math nor science."

Option D: '-M v -S' represents the statement that Dylan does not like math or Dylan does not like science. This is not the best symbolization because it allows for the possibility that Dylan may like one of the subjects but not the other, whereas the original sentence states that Dylan likes neither math nor science.

Final Answer: C

Answer Option 6:

The sentence "Dylan likes neither math nor science" can be symbolized as "¬(M ∨ S)", which translates to "It is not the case that Dylan likes math or science." Option C, "¬(M ∨ S)", correctly symbolizes this sentence. Final Answer: C

Answer Option 7:

The sentence "Dylan likes neither math nor science" can be symbolized as "¬(M ∨ S)", where M represents "Dylan likes math" and S represents "Dylan likes science". This symbolization captures the idea that Dylan does not like math or science, which is what the original sentence is expressing. Therefore, the correct answer is C. '¬(M ∨ S)'.

Answer Option 8:

The best symbolization for the sentence "Dylan likes neither math nor science" would be option D. '~M v ~S'.

Let's break it down:

M represents "Dylan likes math" and S represents "Dylan likes science".

The sentence "Dylan likes neither math nor science" can be translated to "It is not the case that Dylan likes math or Dylan likes science".

In logical notation, this can be represented as '~M v ~S', which means "Dylan does not like math or Dylan does not like science".

Therefore, the best symbolization is option D. ' \sim M \vee \sim S'.

Answer Option 9:

The best symbolization of the given sentence "Dylan likes neither math nor science" is "~(M v S)", which corresponds to option C. This is because '~(M v S)' translates to "It is not the case that Dylan likes either math or science", which accurately captures the meaning of the original sentence. Option A ('M v S') translates to "Dylan likes math or science", which is the opposite of the intended meaning. Option B ('~(M • S)') translates to "It is not the case that Dylan likes both math and science", which is also different from the original sentence. Option D ('~M v ~S') translates to "Dylan does not like math or Dylan does not like science", which is a slightly different interpretation. Therefore, option C is the best symbolization. Final Answer: C

Question 118: 'Select the best translation into predicate logic: All robots are artifacts.', A. '($\forall x$)(Ax \vee Rx)', B. '($\forall x$)(Ax \vee Rx)', C. '($\forall x$)(Rx \vee Ax)', D. '($\forall x$)(Rx \vee Ax)' Answer Option 1: The best translation into predicate logic for the statement "All robots are artifacts" would be: C. $((\forall x)(Rx \supset Ax))'$ This translates to: For all x, if x is a robot, then x is an artifact. Answer Option 2: The best translation into predicate logic for the statement "All robots are artifacts" is: $(\forall x)(Rx \rightarrow Ax)$ Option C: $'(\forall x)(Rx\supset Ax)'$ is the correct translation. Final Answer: C. The best translation into predicate logic for the statement "All robots are artifacts" is option C. $(\forall x)(Rx \supset Ax)$. This translates to "For all x, if x is a robot, then x is an artifact." Final Answer: C Answer Option 4: The best translation into predicate logic for the statement "All robots are artifacts" would be: C. $(\forall x)(Rx \supset Ax)$ This translation captures the meaning that for every object x, if x is a robot (Rx), then x is an artifact (Ax). Answer Option 5: The best translation into predicate logic for the statement "All robots are artifacts" is: $(\forall x)(Rx \rightarrow Ax)$ Rx represents "x is a robot" Ax represents "x is an artifact" This translation states that for all x, if x is a robot, then x is an artifact. Therefore, the correct answer is C. $'(\forall x)(Rx\supset Ax)'$ Answer Option 6: The best translation into predicate logic for the statement "All robots are artifacts" is: $(\forall x)(Rx \supset Ax)$ Explanation: Rx represents "x is a robot" The statement can be translated as "For all x, if x is a robot, then x is an artifact." This is represented by $(\forall x)(Rx\supset Ax)$. Option C: $(\forall x)(Rx\supset Ax)$ is the correct translation into predicate logic. Answer Option 7: The best translation into predicate logic for the statement "All robots are artifacts" is option C. $(\forall x)(Rx\supset Ax)'.$ - ($\forall x$) represents "for all x" or "for every x", indicating that the statement applies to all objects x. - Rx represents "x is a robot." - Ax represents "x is an artifact." The translation '(∀x)(Rx ⊃ Ax)' can be read as "For all x, if x is a robot, then x is an artifact." This accurately captures the meaning of the original statement. Final Answer: C Answer Option 8: The best translation into predicate logic for the statement "All robots are artifacts" is: C. ' $(\forall x)(Rx \supset Ax)$ '. Explanation: Let's break down the statement:

- "All robots" can be translated to $(\forall x)(Rx)$, which means "For all x, x is a robot."

- "are artifacts" can be translated to (Rx \supset Ax), which means "If x is a robot, then x is an artifact."

Combining these translations, we get $(\forall x)(Rx \supset Ax)'$, which represents the statement "For all x, if x is a robot, then x is an artifact." This accurately captures the meaning of the original statement "For all x, if x is a robot, then x is an artifact."	nent

Answer Option 9:

The best translation into predicate logic for the statement "All robots are artifacts" is option C: $(\forall x)(Rx\supset Ax)'$.

Explanation:

Let's break down the statement:

- "All robots": This can be represented as $(\forall x)(Rx),$ which reads as "For all x, x is a robot".
- "are artifacts": This can be represented as (Rx \supset Ax), which reads as "If x is a robot, then x is an artifact".

 $Combining \ these \ two \ parts, \ we \ get \ '(\forall x)(Rx \supset Ax)', \ which \ translates \ to \ "For \ all \ x, \ if \ x \ is \ a \ robot, \ then \ x \ is \ an \ artifact".$

 $Therefore, option \ C, \ '(\forall x)(Rx \supset Ax)', is the best translation into predicate logic for the given statement. \ Final Answer: \ C.$

Question 119:

'Which of the given formulas of PL is the best symbolization of the following sentence?

Since the university raised tuition, the president approved of it and the board didn't prevent it.', A. '($P \cdot A = D \cup U'$, B. '($U \supset P \cdot B'$, C. ' $U \supset (P \cdot A = D) \cup U'$, D. ' $U \supset (P \cdot A = D) \cup U'$, B.' ($U \supset (P \cdot A = D) \cup U'$)

Answer Option 1:

Let's break down the sentence into its constituent parts:

- "The university raised tuition" can be symbolized as U.
- "The president approved of it" can be symbolized as P.
- "The board didn't prevent it" can be symbolized as ~B (negation of B).

Combining these symbols, the best symbolization of the given sentence would be:

 $U \supset (P \bullet \sim B)$

Option C, ${}^{\prime}U\supset (P \ {}^{\bullet} \ {}^{\backprime}B)',$ is the correct symbolization of the given sentence. Final Answer: C

Answer Option 2:

The best symbolization for the given sentence would be option C: 'U \supset (P \bullet ~B)'.

Let's break down the sentence:

- 1. "Since the university raised tuition" This can be represented as 'U'.
- 2. "the president approved of it" This can be represented as 'P'
- 3. "the board didn't prevent it" This can be represented as '~B', where '~' denotes negation.

Combining these representations, we get "U \supset (P • ~B)" which translates to "If the university raised tuition, then the president approved of it and the board didn't prevent it".

Therefore, the best symbolization for the given sentence is option ${\sf C}.$

Answer Option 3:

To symbolize the given sentence, let's break it down into its constituent parts:

- P: The university raised tuition.
- U: The president approved of it.
- B: The board didn't prevent it.

The sentence can be represented as "Since P, U and ${\sim}B.$ "

To symbolize "Since P," we use P as a premise.

To symbolize "U and ~B," we use U as the antecedent and ~(B) as the consequent.

Therefore, the best symbolization of the sentence is: 'P \supset (U \bullet ~B)'.

Comparing this with the given options:

A. $'(P \bullet \sim B) \supset U'$ - Incorrect. It does not capture the "Since" relationship.

B. ${}^{\shortmid}(U\supset P)\bullet B^{\prime}$ - Incorrect. It does not capture the "Since" relationship.

C. ${}^{\prime}U \supset (P \bullet {}^{\prime} B)^{\prime}$ - Correct. It correctly symbolizes the given sentence.

D. $'U \supset (P \vee {\sim} B)'$ - Incorrect. It represents a different logical relationship.

Final Answer: C

Answer Option 4:

Let's break down the sentence to identify the appropriate symbols:

- "Since the university raised tuition" can be represented as U (university raised tuition).
- "the president approved of it" can be represented as P (president approved).
- "the board didn't prevent it" can be represented as ~B (board didn't prevent).

Now let's evaluate the options:

Option A: $'(P \bullet {\scriptstyle \sim} B) \supset U'$

This option translates to "If the president approves and the board doesn't prevent, then the university raised tuition." This does not accurately represent the original sentence, so it is not the best symbolization.

Option B: '(U ⊃ P) • B'

This option translates to "If the university raised tuition, then the president approves, and the board prevents." This does not accurately represent the original sentence, so it is not the best symbolization.

Option C: $'U \supset (P \cdot \sim B)'$

This option translates to "If the university raised tuition, then both the president approves and the board doesn't prevent." This accurately represents the original sentence, making it the best symbolization.

Option D: $'U \supset (P \lor \sim B)'$

This option translates to "If the university raised tuition, then either the president approves or the board doesn't prevent." This does not accurately represent the original sentence, so it is not the best symbolization.

Final Answer: C

Answer Option 5:

To symbolize the sentence "Since the university raised tuition, the president approved of it and the board didn't prevent it", we can break it down into three parts:

- 1. The university raised tuition: Let's symbolize this as U.
- 2. The president approved of it: Let's symbolize this as P.
- 3. The board didn't prevent it: Let's symbolize this as ~B.

Now, we need to connect these three parts using logical operators to represent the relationship between them.

Looking at the options:

A. '(P • ~B) \supset U': This option represents that if the president approved and the board didn't prevent, then the university raised tuition. However, we need to represent the opposite relationship, where the university raising tuition leads to the president approving and the board not preventing. So, this option is not the best symbolization.

B. '(U > P) • B': This option represents that if the university raised tuition, then the president approves, and the board prevents it. However, we need the opposite relationship, where the university raising tuition leads to the president approving and the board not preventing. So, this option is not the best symbolization.

C. U \supset (P • -B): This option represents that if the university raised tuition, then the president approves and the board doesn't prevent it. This accurately represents the relationship between the three parts of the sentence. So, this option is a good symbolization.

D. 'U \supset (P \lor ~B)': This option represents that if the university raised tuition, then either the president approves or the board doesn't prevent it. This doesn't accurately represent the original sentence, as it allows for the possibility of the president approving and the board preventing. So, this option is not the best symbolization.

Based on the above analysis, the best symbolization of the given sentence is option C: $"U \supset (P \cdot {\sim} B)"$. Final Answer: C.

Answer Option 6:

Let's analyze the sentence to find the best symbolization:

- 1. Since the university raised tuition: This can be symbolized as 'U'.
- 2. the president approved of it: This can be symbolized as 'P'
- 3. the board didn't prevent it: This can be symbolized as '~B' (not B).

Now, let's examine the given options:

Option A: $'(P \bullet \sim B) \supset U'$

This option translates to "If the president approves of it and the board doesn't prevent it, then the university raised tuition." This doesn't accurately capture the original sentence, as it doesn't include the notion of causality or the fact that the university raising tuition is the reason for the other events.

Option B: $'(U \supset P) \bullet B'$

This option translates to "If the university raised tuition, then the president approves of it, and the board prevents it." This not only changes the meaning of the sentence but also introduces a contradiction by stating that the board prevents the tuition increase while also stating that the university raised tuition.

Option C: $'U \supset (P \cdot \sim B)'$

This option translates to "If the university raised tuition, then the president approves of it and the board doesn't prevent it." This accurately captures the original sentence, as it represents the causal relationship between the university raising tuition and the subsequent actions of the president and the board.

Option D: 'U ⊃ (P ∨ ~B)'

This option translates to "If the university raised tuition, then either the president approves of it or the board doesn't prevent it." This does not accurately capture the original sentence, as it allows for the possibility that both the president approves and the board doesn't prevent the tuition increase.

Based on the analysis above, the best symbolization of the given sentence is Option C: 'U \supset (P \bullet ~B)'. The final answer is C.

Answer Option 7:

To symbolize the sentence "Since the university raised tuition, the president approved of it and the board didn't prevent it," we need to break it down into its constituent parts.

Let's assign the following propositions:

- U: The university raised tuition.
- P: The president approved of it.
- B: The board prevented it.

The sentence can be translated as follows:

"U ⊃ (P • ~B)"

Option A: $'(P \bullet \sim B) \supset U'$

This option translates to "The president approved of it and the board didn't prevent it, implies that the university raised tuition." This is not an accurate representation of the original sentence.

Option B: $'(U \supset P) \bullet B'$

This option translates to "The university raised tuition implies that the president approved of it, and the board prevented it." This is not an accurate representation of the original sentence.

Option C: $'U \supset (P \bullet \sim B)'$

This option correctly symbolizes the sentence "Since the university raised tuition, the president approved of it and the board didn't prevent it." It accurately represents the logical relationship between the propositions

Option D: $'U \supset (P \lor \sim B)'$

This option translates to "If the university raised tuition, then either the president approved of it or the board prevented it." This is not an accurate representation of the original sentence.

Therefore, the best symbolization of the given sentence is option C: 'U ⊃ (P • ~B)'.

Answer Option 8:

To symbolize the given sentence, we can break it down into three parts:

- 1. The university raised tuition: U
- 2. The president approved of it: P
- 3. The board didn't prevent it: ~B

Now, let's evaluate the options:

 $A.\ '(P \bullet {\scriptstyle \sim} B) \supset U'$

This option states that if the president approved of it and the board didn't prevent it, then the university raised tuition. This does not accurately represent the given sentence because it does not capture the causality between the university raising tuition and the president's approval or the board's prevention.

B. $'(U \supset P) \bullet B'$

This option states that if the university raised tuition, then the president approved of it, and the board prevented it. This does not accurately represent the given sentence because it states that the board prevented the university from raising tuition, which contradicts the original statement.

 $C. \ 'U \supset (P \bullet {\scriptstyle \sim} B)'$

This option states that if the university raised tuition, then both the president approved of it and the board didn't prevent it. This accurately represents the given sentence, as it captures the causality and the conditions mentioned.

 $\mathsf{D.} \; \mathsf{'U} \supset (\mathsf{P} \vee \mathsf{\sim} \mathsf{B})\mathsf{'}$

This option states that if the university raised tuition, then either the president approved of it or the board didn't prevent it. This does not accurately represent the given sentence because it allows for the possibility that the president approved and the board prevented.

Therefore, the best symbolization of the given sentence is option C: $"U\supset (P \bullet \sim B)".$

Answer Option 9:

Let's analyze the given sentence: "Since the university raised tuition, the president approved of it and the board didn't prevent it."

Let's symbolize the propositions:

- U: The university raised tuition.
- P: The president approved of it.
- B: The board prevented it.

The sentence can be broken down into three parts:

- 1. The university raised tuition.
- 2. The president approved of it.
- 3. The board didn't prevent it.

Now let's examine the options:

Option A: (P • ~B) ⊃ U

This option translates to "If the president approved of it and the board didn't prevent it, then the university raised tuition." However, this doesn't capture the idea of causality in the original sentence. It only states a conditional relationship.

Option B: $(U \supset P) \bullet B$

This option translates to "If the university raised tuition, then the president approved of it, and the board prevented it." This doesn't capture the original sentence accurately.

Option C: $U \supset (P \bullet \sim B)$

This option translates to "If the university raised tuition, then the president approved of it, and the board didn't prevent it." This captures the original sentence accurately, as it states that the president approves and the board doesn't prevent when the university raises tuition.

Option D: U \supset (P \lor \sim B)

This option translates to "If the university raised tuition, then either the president approves or the board didn't prevent it." This doesn't capture the original sentence accurately, as it allows for either the president's approval or the board's non-prevention.

Based on the analysis, the best symbolization of the given sentence is Option C: $"U \supset (P \cdot \sim B)"$. Final Answer: C

Q	ue	tic	1	21

'Identify the conclusion of the following argument. We shouldn't abolish capital punishment. If we do, prisons will become more crowded. If prisons become more crowded, then we will have to build more prisons.', A. 'We shouldn't abolish capital punishment.', B. 'If we do, prisons will become more crowded.', C. 'If prisons become more crowded, then we will have to build more prisons.', D. 'We will have to build more prisons.'

Answer Option 1:

The conclusion of the argument is: D. 'We will have to build more prisons.'

This is the final statement that follows from the previous premises. It is the result of the reasoning presented in the argument.

Answer Option 2:

The conclusion of the argument is: 'We don't want to build more prisons.'

This is evident from the statement "And we don't want to build more prisons." This conclusion is drawn based on the reasoning provided in the argument.

Final Answer: D. 'We will have to build more prisons.'

Answer Option 3:

The conclusion of the argument is: D. 'We will have to build more prisons.'

This is the final statement that follows from the premises given in the argument. The argument presents a chain of reasoning, starting with the statement that abolishing capital punishment will lead to overcrowded prisons, which in turn will require the construction of more prisons. The conclusion is the final outcome of this reasoning, which is that we will have to build more prisons.

Answer Option 4:

The conclusion of the argument is: D. 'We will have to build more prisons.'

Answer Option 5:

The conclusion of the argument is: D. 'We will have to build more prisons.'

Answer Option 6:

The conclusion of the argument is: D. 'We will have to build more prisons.'

Answer Option 7:

The conclusion of the argument is: D. 'We will have to build more prisons.'

Answer Option 8:

The conclusion of the argument is: D. We will have to build more prisons. This is the final statement that is directly supported by the premises given in the argument.

Answer Option 9:

The conclusion of the argument is: D. We will have to build more prisons.' The argument presents a chain of reasoning that leads to the conclusion that if capital punishment is abolished, more prisons will

need to be built.

Question 121:

'Which of the given formulas of PL is the best symbolization of the following sentence?

Either England's not importing beef is a necessary condition for France's subsidizing agriculture, or China promotes human rights if and only if South Africa doesn't supply diamonds.', A. '(F \(\to - E\) \(\to (C = \-S)', \)

B. $'(\sim E\supset F)\lor (C\supset \sim S)', C. '\sim (F\supset E)\lor (\sim S\supset C)', D. '(F\supset E)\lor (C\equiv \sim S)'$

Answer Option 1:

To symbolize the given sentence, let's break it down into two parts:

- 1. "England's not importing beef is a necessary condition for France's subsidizing agriculture"
- 2. "China promotes human rights if and only if South Africa doesn't supply diamonds"

We can represent the first part as "(~E → F)", where E represents England's importing beef and F represents France's subsidizing agriculture. Here, we use the implication symbol (→) to represent "is a necessary condition for."

Similarly, we can represent the second part as " $(C \leftrightarrow \sim S)$ ", where C represents China promoting human rights and S represents South Africa supplying diamonds. Here, we use the biconditional symbol (\leftrightarrow) to represent "if and only if."

Now, we need to combine both parts using the logical operator for disjunction (OR).

Option A: " $(F \supset \sim E) \lor (C \equiv \sim S)$ "

This option does not accurately symbolize the given sentence. It represents the first part correctly, but it incorrectly represents the second part using the biconditional (=) instead of the implication (-).

Option B: "(\sim E \supset F) \vee (C \supset \sim S)"

This option does not accurately symbolize the given sentence. It represents the first part incorrectly by placing the negation (-) outside of the implication (→).

Option C: " \sim (F \supset E) \vee (\sim S \supset C)"

This option does not accurately symbolize the given sentence. It incorrectly represents both parts by placing the negation (~) outside of the implications (→) and using incorrect order of propositions.

Option D: " $(F \supset E) \lor (C \equiv \sim S)$ "

This option accurately symbolizes the given sentence by representing both parts correctly using the implication (\rightarrow) and biconditional (\leftrightarrow) symbols and combining them with the disjunction (OR) operator.

Final Answer: D

Answer Option 2:

To symbolize the given sentence, let's break it down into two parts:

1. Either England's not importing beef is a necessary condition for France's subsidizing agriculture

Symbolization: (F \supset ~E)

2. China promotes human rights if and only if South Africa doesn't supply diamonds

Symbolization: (C \equiv ~S)

Now, let's combine these two parts using the logical operator for "or" (v):

(F ⊃ ~E) v (C ≡ ~S)

Comparing this with the given options, we can see that the best symbolization is:

A. $'(F\supset \sim E) \lor (C\equiv \sim S)'$

Final Answer: A

Answer Option 3:

Let's break down the sentence to its logical components:

- England's not importing beef is a necessary condition for France's subsidizing agriculture: This can be symbolized as E $\supset \sim$ F.
- China promotes human rights if and only if South Africa doesn't supply diamonds: This can be symbolized as $C \equiv \sim S$.

Now, let's consider the sentence as a whole:

"Either England's not importing beef is a necessary condition for France's subsidizing agriculture, or China promotes human rights if and only if South Africa doesn't supply diamonds."

This can be symbolized as (E \supset ~F) \lor (C \equiv ~S).

Therefore, the correct symbolization of the given sentence is A. '(F $\supset \sim$ E) \vee (C $\equiv \sim$ S)'.

Answer Option 4:

To symbolize the given sentence, let's break it down into two parts:

- 1. Either England's not importing beef is a necessary condition for France's subsidizing agriculture
- 2. Or China promotes human rights if and only if South Africa doesn't supply diamonds

Let's analyze each part individually:

- 1. England's not importing beef is a necessary condition for France's subsidizing agriculture can be symbolized as: E \rightarrow ~F
- 2. China promotes human rights if and only if South Africa doesn't supply diamonds can be symbolized as: C \leftrightarrow ~S

Now, let's combine these two parts using the logical connective 'or':

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1.~(E \rightarrow {\sim}F) \lor (C \leftrightarrow {\sim}S)
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Comparing this with the given options:

A. $(F\supset \sim E) \lor (C\equiv \sim S)$

 $\mathsf{B.}\; (\mathsf{\sim}\mathsf{E}\supset\mathsf{F}) \lor (\mathsf{C}\supset\mathsf{\sim}\mathsf{S})$

 $C. \ \text{$\sim$}(\mathsf{F}\supset\mathsf{E}) \lor (\text{\sim}\mathsf{S}\supset\mathsf{C})$

D. $(F \supset E) \lor (C \equiv \sim S)$

Option A, B, and D do not have the correct symbolization for the first part of the sentence.

Option C has the correct symbolization for the first part, but it has an incorrect symbolization for the second part.

Therefore, the best symbolization of the given sentence is option C: '~(F \supset E) \lor (~S \supset C)'.

Answer Option 5:

Let's break down the sentence to determine the best symbolization:

- England's not importing beef is a necessary condition for France's subsidizing agriculture: This can be symbolized as ~E → F, where ~E represents "England's not importing beef" and F represents "France's subsidizing agriculture".
- 2. China promotes human rights if and only if South Africa doesn't supply diamonds: This can be symbolized as C = ~S, where C represents "China promotes human rights" and ~S represents "South Africa doesn't supply diamonds".

Now, let's evaluate the given options:

Option A: '(F \supset ~E) \lor (C \equiv ~S)': This option incorrectly asserts that France's subsidizing agriculture is a necessary condition for England's not importing beef.

Option B: '(-E > F) v (C > -S)': This option correctly symbolizes the first condition but incorrectly states that China promoting human rights is a necessary condition for South Africa not supplying diamonds.

Option C: '~(F \supset E) \lor (~S \supset C)': This option incorrectly negates the first condition and asserts that France's subsidizing agriculture is a necessary condition for England's importing beef.

 $\text{Option D: '}(F\supset E) \lor (C\equiv \sim S)\text{': This option correctly symbolizes both conditions and is the best symbolization of the given sentence. } \\$

Therefore, the correct answer is D. '(F \supset E) \lor (C \equiv ~S)'.

Answer Option 6:

To symbolize the given sentence, we can break it down into two main parts:

1. 'Either England's not importing beef is a necessary condition for France's subsidizing agriculture'

This can be symbolized as (\sim E \rightarrow F), where E represents England's importing beef and F represents France's subsidizing agriculture.

2. 'China promotes human rights if and only if South Africa doesn't supply diamonds'

This can be symbolized as ($C \equiv \sim S$), where C represents China promoting human rights and S represents South Africa supplying diamonds.

Combining these two parts using the logical operator 'or' (v), we get the symbolization: (\sim E \rightarrow F) \vee (C \equiv \sim S).

Comparing the given options:

A. $(F \supset \sim E) \lor (C \equiv \sim S)$

B. $(\sim E \supset F) \lor (C \supset \sim S)$

 $C. \ {\sim} (\mathsf{F} \supset \mathsf{E}) \lor ({\sim} \mathsf{S} \supset C)$

D. $(F \supset E) \lor (C \equiv \sim S)$

Option A does not have the correct order of ~E \rightarrow F.

Option B does not have the correct order of $\sim E \rightarrow F$ and $C \equiv \sim S$.

Option C is not the best symbolization as it has negations in front of the implications.

Option D has the correct order and is the best symbolization of the given sentence.

Therefore, the correct answer is D. '(F \supset E) \lor (C \equiv ~S)'.

Answer Option 7:

To symbolize the given sentence, we can break it down into two parts:

1) "England's not importing beef is a necessary condition for France's subsidizing agriculture"

 $This \ can \ be \ symbolized \ as \ \neg E \rightarrow F, \ where \ E \ represents \ England's \ importing \ beef \ and \ F \ represents \ France's \ subsidizing \ agriculture.$

2) "China promotes human rights if and only if South Africa doesn't supply diamonds"

 $This \ can \ be \ symbolized \ as \ C \equiv -S, \ where \ C \ represents \ China \ promotes \ human \ rights \ and \ S \ represents \ South \ Africa \ supplying \ diamonds.$

Now, we need to combine these two parts using logical connectives. The correct symbolization for the given sentence is the disjunction (OR) between the two parts.

Option A: $(F \supset \sim E) \lor (C \equiv \sim S)'$

This option correctly symbolizes the two parts, $\sim E \rightarrow F$ and $C \equiv \sim S$, and combines them with a disjunction.

Option B: $(\sim E \supset F) \lor (C \supset \sim S)'$

This option incorrectly changes the order of ~E and F in the first part, which results in a different meaning.

Option C: '~(F \supset E) \lor (~S \supset C)'

This option incorrectly adds negations to the first part, which changes its meaning.

Option D: $'(F \supset E) \lor (C \equiv \sim S)'$

This option incorrectly changes the direction of the implication in the first part, which results in a different meaning.

Therefore, the best symbolization for the given sentence is Option A: $(F \supset \sim E) \lor (C \equiv \sim S)'$.

Final Answer: A

Answer Option 8

To symbolize the sentence, let's break it down into two parts:

1. "England's not importing beef is a necessary condition for France's subsidizing agriculture."

This can be symbolized as ~E → F, where E represents England's importing beef and F represents France's subsidizing agriculture.

2. "China promotes human rights if and only if South Africa doesn't supply diamonds."

 $This \ can be \ symbolized \ as \ C \leftrightarrow \sim S, \ where \ C \ represents \ China \ promoting \ human \ rights \ and \ S \ represents \ South \ Africa \ supplying \ diamonds.$

Now, we need to combine these two parts with an "or" condition. The correct symbolization would be:

'Either England's not importing beef is a necessary condition for France's subsidizing agriculture, or China promotes human rights if and only if South Africa doesn't supply diamonds.

Symbolization: $(\sim E \rightarrow F) \lor (C \leftrightarrow \sim S)$

Comparing the options:

A. $'(F\supset \sim E)\lor (C\equiv \sim S)'$: This option is not the best symbolization as it has the wrong direction in the first part ($\sim E\to F$ instead of $\sim E\to F$).

 $B. \ '(\sim\!E\supset\!F) \lor (C\supset\sim\!S)'; \ This \ option \ is \ also \ not \ the \ best \ symbolization \ as \ it \ has \ the \ wrong \ direction \ in \ the \ first \ part \ (\sim\!E\to\!F) \ instead \ of \ \sim\!E\to\!F).$

 $C. \ \ \text{`$-(F\supset E)$} \lor (\text{$-S$}) \lor (\text{$-S$})'. \ \text{This option is not the best symbolization as it has the wrong direction in the second part ($C\leftrightarrow -S$)}.$

 $D. \ '(F \supset E) \lor (C \equiv \sim S)': This \ option \ is \ not \ the \ best \ symbolization \ as \ it \ has \ the \ wrong \ direction \ in \ the \ first \ part \ (\sim E \to F \ instead \ of \ \sim E \to F).$

Therefore, the best symbolization is option B: $'(\sim E \supset F) \lor (C \supset \sim S)'$.

Answer Option 9:

To symbolize the given sentence, let's break it down into two parts:

- 1. England's not importing beef is a necessary condition for France's subsidizing agriculture.
- 2. China promotes human rights if and only if South Africa doesn't supply diamonds.

Now, let's analyze each option:

Option A: $'(F\supset \sim E) \lor (C\equiv \sim S)'$

This option correctly symbolizes the first part of the sentence as (F \supset ~E), which means "France's subsidizing agriculture is a necessary condition for England not importing beef." However, it does not correctly symbolize the second part of the sentence.

Option B: $(\sim E \supset F) \lor (C \supset \sim S)'$

This option incorrectly symbolizes the first part of the sentence as (-E \supset F), which means "England not importing beef is a necessary condition for France's subsidizing agriculture." This is the opposite of what was stated in the sentence. It also correctly symbolizes the second part of the sentence.

This option incorrectly symbolizes the first part of the sentence as ¬(F⊃E), which means "It is not the case that France's subsidizing agriculture is a necessary condition for England importing beef." This is the opposite of what was stated in the sentence. It also incorrectly symbolizes the second part of the sentence.

Option D: $'(F \supset E) \lor (C \equiv \sim S)'$

This option incorrectly symbolizes the first part of the sentence as (F \supset E), which means "France's subsidizing agriculture is a necessary condition for England importing beef." This is the opposite of what was stated in the sentence. It correctly symbolizes the second part of the sentence.

Therefore, the best symbolization of the given sentence is Option B: '(~E \supset F) \lor (C \supset ~S)'. Final Answer: B

'Use indirect truth tables to determine whether the following argument is valid. If the argument is invalid, choose an option which presents a counterexample. (There may be other counterexamples as well.)
$(K \cdot L) \supset M$
N⊃~M
N v K / L ⊃ K', A. "Valid', B. "Invalid. Counterexample when L, M, and N are true and K is false', C. "Invalid. Counterexample when L and M are true and K and N are false', D. "Invalid. Counterexample when L
and N are true and K and M are false'
Answer Option 1:
To determine the validity of the argument, we can use an indirect truth table.
Assume the following truth values for the premises:
$(K \cdot L) \supset M$: true
N ⊃ ~M: true
N ∨ K: true
We need to determine the truth value of $L\supset K$ based on these assumptions.
To do this, we construct a truth table considering L and K as variables.
$L[K]N[(K \cdot L) \supset M]N \supset -M]N \lor K[L \supset K$
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In this case, we can see that for the assumption $N \vee K$ to be true, it is not necessary for $L \supset K$ to be true.
Therefore, the argument is invalid.
Counterexample: When L and N are true, K and M are false, the premises are true but the conclusion (L ⊃ K) is false.
Final Answer: C. 'Invalid. Counterexample when L and M are true and K and N are false'
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Answer Option 2:
To determine the validity of the argument, we can use an indirect truth table. We will analyze all possible combinations of truth values for the given propositions and check if the conclusion is always true when
the premises are true.
P1: (K⋅L) ⊃ M
P2: N ¬ ~M
P3: N v K
Conclusion: L > K
We construct the truth table as follows:
KLMN(K·L) > MN > ~MN v KL > K
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Question 122:

TFFTFFTT TEFFETE FTTTTFTT FTTFTTTT FTFTTFTT FTFFTTFT FFTTTFFT FFTFTTFT FFFTTFTT FFFFTTFT From the truth table, we can see that there is a row where all the premises are true (T) but the conclusion is false (F). This means the argument is invalid. Counterexample: Option C provides a counterexample when L and M are true, and K and N are false. In this case, all the premises are true, but the conclusion (L \supset K) is false. Final Answer: B. Invalid. Counterexample when L and M are true and K and N are false. Answer Option 3: To determine the validity of the argument, we will use an indirect truth table. $P{:}\; (K\cdot L)\supset M$ Q: $N \supset \sim M$ R: N v K $S{:}\ L\supset K$ We construct a truth table for the premises and conclusion: $\mathsf{K} \mid \mathsf{L} \mid \mathsf{M} \mid \mathsf{N} \mid (\mathsf{K} \cdot \mathsf{L}) \supset \mathsf{M} \mid \mathsf{N} \supset \mathsf{\sim} \mathsf{M} \mid \mathsf{N} \vee \mathsf{K} \mid \mathsf{L} \supset \mathsf{K} \mid \mathsf{L} \supset \mathsf{K}'$ T|T|T|T|F|T|T|TTITITIFITITITIT T|T|F|T|F|F|T|T|TTITIFIFIFITITITIT T|F|T|T|T|F|T|T|TT|F|T|F|T|T|T|TT|F|F|T|F|F|T|T|T T|F|F|F|F|T|F|T|T FITITITITITIT F|T|T|F|T|T|F|T|T FITIFITITIFITIT F|T|F|F|T|T|F|T|T F|F|T|T|T|F|T|T|TF|F|T|F|T|T|F|T|TF|F|F|T|T|F|T|T F|F|F|F|T|T|F|T|T As we can see from the truth table, the conclusion $L\supset K'$ is always true whenever the premises are true. Therefore, the argument is valid. Final Answer: A. Valid ****************** Answer Option 4: To determine the validity of the argument, we can create an indirect truth table. We have the premises: 1. $(K \cdot L) \supset M$ 2. $N \supset \sim M$ And the conclusion:

We will assign truth values to the variables K, L, M, and N and evaluate the truth value of the conclusion.

 $K \mid L \mid M \mid N \mid N \vee K \mid L \supset K' \mid (N \vee K) \supset L \supset K'$

TITITITITIT T|T|T|F|T|T|TTITIFITITIT T|T|F|F|T|T|TTIFITITIFIF TIFITIFITIT T|F|F|T|T|F|FTIFIFIFITITIT F|T|T|T|T|T|TFITITIFIFITIT FITIFITITIT F|T|F|F|F|T|T F|F|T|T|T|F|F F|F|T|F|F|F|T F|F|F|T|T|F|F F|F|F|F|F|T As we can see, there is at least one row where the premises are true and the conclusion is false, specifically when L and N are true, and K and M are false. Therefore, the argument is invalid. The correct answer is D. 'Invalid. Counterexample when L and N are true and K and M are false' Answer Option 5: To determine the validity of the argument, we can use an indirect truth table. We will assume the premises to be true and check if the conclusion is also true in all cases. Assumptions: $1.\; (K\cdot L)\supset M$ 2. N ⊃ ~M $3.\;N\vee K$ We will create a truth table with columns for K, L, M, N, $(K \cdot L) \supset M$, $N \supset \sim M$, $N \vee K$, and $L \supset K'$. $\mathsf{K} \mid \mathsf{L} \mid \mathsf{M} \mid \mathsf{N} \mid (\mathsf{K} \cdot \mathsf{L}) \supset \mathsf{M} \mid \mathsf{N} \supset \mathsf{\sim} \mathsf{M} \mid \mathsf{N} \vee \mathsf{K} \mid \mathsf{L} \supset \mathsf{K}'$ TITITITIFITIT T|T|T|F|T|T|T TITIFITIFIFITIT TITIFIFIFITITIT TIFITITITITIT T|F|T|F|T|T|T|TT|F|F|T|F|F|T|TT|F|F|F|F|T|F|F

FITITITITITIT FITITIFITITIT F|T|F|T|T|F|T|TF|T|F|F|T|T|F|F F|F|T|T|T|F|F|T F|F|T|F|T|T|F|T F|F|F|T|T|F|T|T F|F|F|F|T|T|F|T

From the truth table, we can see that there is at least one row where the premises are true (T) and the conclusion $L\supset K'$ is false (F). Specifically, when L and M are true, and K and N are false, the conclusion

Therefore, the argument is invalid.

The correct answer is C. 'Invalid. Counterexample when L and M are true and K and N are false'.

Answer Option 6:

To determine the validity of the argument, we can use an indirect truth table. We will assign truth values to the premises and check if the conclusion is always true when the premises are true.

Let's construct the truth table: $\mathsf{K} \mid \mathsf{L} \mid \mathsf{M} \mid \mathsf{N} \mid (\mathsf{K} \cdot \mathsf{L}) \supset \mathsf{M} \mid \mathsf{N} \supset \mathsf{\sim} \mathsf{M} \mid \mathsf{N} \vee \mathsf{K} \mid \mathsf{L} \supset \mathsf{K}$ TITITITIFITIT $\mathsf{T} | \mathsf{T} | \mathsf{T} | \mathsf{F} | \mathsf{T} | \mathsf{T} | \mathsf{T} | \mathsf{T} | \mathsf{T}$ T|T|F|T|F|F|T|T T|T|F|F|F|T|F|T T|F|T|T|T|F|T|T TIFITIFITITIT T|F|F|T|F|F|T|TT|F|F|F|F|T|F|T F|T|T|T|T|F|T|TF|T|T|F|T|T|F|T F|T|F|T|T|F|T|T F|T|F|F|T|T|F|T F|F|T|T|T|F|T|TF|F|T|F|T|T|F|T F|F|F|T|T|F|T|T F|F|F|F|T|T|F|T We can see from the truth table that there are no rows where all the premises are true and the conclusion is false. Therefore, the argument is valid. Final Answer: A. 'Valid' Answer Option 7: To determine the validity of the argument, we can create an indirect truth table. $P{:}\; (K\cdot L)\supset M$ Q: N ⊃ ~M $R \colon N \vee K$ $S{:}\ L\supset K'$ We need to check if the conclusion S is always true when the premises P and Q are true. PQRS $\mathsf{T}\,\mathsf{T}\,\mathsf{T}\,\mathsf{T}$ TTTF TTFT TTFF TFTT TFTT TFFT TFFT FTTT FTTT FTFT FTFT FFTT FFTT FFFT FFFT In all the cases where P and Q are true, S is also true. Therefore, the argument is valid. Final Answer: A. Valid ******* Answer Option 8: To determine the validity of the argument, we can use an indirect truth table approach.

To do so, we list all the variables involved in the argument and assign truth values to all possible combinations:

 $\mathsf{K} \mathrel{L} \mathsf{M} \mathrel{N} \mathsf{L} \supset \mathsf{K} \; (\mathsf{K} \cdot \mathsf{L}) \supset \mathsf{M} \; \mathsf{N} \supset \mathsf{\sim} \mathsf{M} \; \mathsf{N} \vee \mathsf{K} \; \mathsf{L} \supset \mathsf{K}'$

TTTTT TTTTTTTTTT ${\tt TTTTT}$ TETTTE TFTTTF TFTTTF TFTTTF FTTTTT FTTTTT FTTTT FTTTTT FFTTFT FFTTFT FFTTFT FFTTFT In the last column, $L\supset K'$, we can see that it is always true. There is no row where the premise $(N\vee K)$ is true and the conclusion $L\supset K'$ is false. Therefore, the argument is 'Valid'. Final Answer: A ********* Answer Option 9: To determine the validity of the argument, we can use an indirect truth table. We have the premises: 1. $(K \cdot L) \supset M$ 2. $N \supset \sim M$ 3. N v K And we want to determine the conclusion: $\mathsf{L}\supset\mathsf{K}$ We construct a truth table with all the relevant propositions: K, L, M, and N. We also include the intermediate propositions from the premises. $K \mid L \mid M \mid N \mid (K \cdot L) \supset M \mid N \supset {} \sim \! M \mid N \vee K \mid L \supset K$ TITITITIFITIT TITITIFITITITIT T|T|F|T|F|F|T|TT|T|F|F|F|T|T|T $\mathsf{T} | \mathsf{F} | \mathsf{T} | \mathsf{T} | \mathsf{T} | \mathsf{F} | \mathsf{T} | \mathsf{F}$ T|F|T|F|T|T|F T|F|F|T|F|F|T|F T|F|F|F|F|T|T|F F|T|T|T|T|F|F|TFITITIFITITIT F|T|F|T|T|F|T|TF|T|F|F|T|T|F|TF|F|T|T|T|F|T|T F|F|T|F|T|T|F|TF|F|F|T|T|F|T|T F|F|F|F|T|T|F|T From the truth table, we can see that there is no row where all the premises are true and the conclusion is false. Therefore, the argument is valid. Final Answer: A

Question 123:

* Construct a complete truth table for the following argument. Then, using the truth table, determine whether the argument is valid or invalid. If the argument is invalid, choose an option which presents a counterexample. (There may be other counterexamples as well.)

(G ≡ H) · ~I

~G v (~H v I) / G', A. 'Valid', B. 'Invalid. Counterexample when H and I are true and G is false', C. 'Invalid. Counterexample when H is true and I and G are false', D. 'Invalid. Counterexample when H, I, and G are false

Answer Option 1:

To construct a complete truth table, we need to consider all possible combinations of truth values for the variables G, H, and I.

 $G H I (G \equiv H) \cdot \sim I \sim G \vee (\sim H \vee I) G'$

TTTFFTT

TTEETTT

TETEETT

TFFFTTT

FTTFFTF

FTFFTTF FFTTTTF

FFFFTTF

Based on the truth table, we can see that there are two rows where the premises (~G v (~H v I)) are true, but the conclusion G' is false. These rows are:

- 1. When G is true, H is true, and I is false.
- 2. When G is false, H is false, and I is false.

Therefore, the argument is invalid. The correct answer is C. 'Invalid. Counterexample when H is true and I and G are false.'

Answer Option 2:

To construct the truth table, we evaluate the given argument by considering all possible combinations of truth values for the propositions G, H, and I.

 $G H I (G \equiv H) \cdot \sim I \sim G \vee (\sim H \vee I) G'$

TTTTFTT

TTFFTTT

TETEETT

TEEFETT

FTTFFTF

FTFFFTF

FFTFFTF

FFFTTTF

The truth table shows that in all cases where the premises ((G = H) · -I) and (-G ∨ (-H ∨ I)) are true, the conclusion G' is also true. Therefore, the argument is valid.

Final Answer: A. Valid.

Answer Option 3:

To construct a truth table for the argument, we need to consider all possible combinations of truth values for the propositions G, H, and I.

 $G H I \sim I (G \equiv H) (\sim H \vee I) \sim G \vee (\sim H \vee I) G'$

TTTETTTT

TTFTTTT

TETEETTT

TEETETTT

FTTFFTTF

FTFTTFTF

FFTFTTTF

FFFTTTTT

In the last column, G' is false in all cases where G is false. Therefore, the argument is invalid since there exists at least one row where the premises are true but the conclusion is false.

To show a counterexample, we can choose the row where G is false, H is true, and I is false.

Therefore, the correct answer is: C. 'Invalid. Counterexample when H is true and I and G are false'

Answer Option 4:

 $G H I \sim I (G \equiv H) (G \equiv H) \cdot \sim I \sim G \sim H \sim G \vee (\sim H \vee I) G'$ TTTETEETT TTFTTFFFF TFTFFFFTTT TEETEETTT FTTFFFTFTF FTETEETETE FFTFTFTTFF FFFTTTTFF The argument is valid if and only if the conclusion (G') is true for every row where the premises ((G \equiv H) \cdot ~I) and (~G \vee (~H \vee I)) are true. In the truth table, we can see that there is no row where both premises are true and the conclusion is false. Therefore, the argument is valid. Final Answer: A. 'Valid' Answer Option 5: We will construct a truth table for the given argument: $G \; H \; I \; \neg I \; G \equiv H \; (G \equiv H) \; \cdot \; \neg I \; \neg G \; \neg H \; \neg H \; \lor \; I \; \neg G \; \lor \; (\neg H \; \lor \; I) \; G'$ TTTFTFFFFTT TTFTTFFFTTF TETEFFETTT TFFTFFFTTF FTTFFTFTT FTFTFTTTF FFTFTTTTFTT FFFTTTTTFTF From the truth table, we can see that the argument is invalid. There are rows where the premises are true $(-G \vee (-H \vee I))$ is true) but the conclusion (G') is false. Counterexample: When H is true, I is false, and G is false, the premises $(-G \lor (-H \lor I))$ are true but the conclusion (G') is false. B. Invalid. Counterexample when H and I are true and G is false. ****************** Answer Option 6: To construct a complete truth table for the argument, we need to consider all possible combinations of truth values for the propositions G, H, and I. $G H I (G \equiv H) \cdot \sim I \sim G \vee (\sim H \vee I) G'$ TTTTFTF TTFFFTF TETETTE TFFFTTF FTTFFTT FTFFTTT FETTETT FFFTTTT Now, let's evaluate the argument based on the truth table: When G is true (T), the conclusion G' is false (F) in all cases. When G is false (F), the conclusion G' is true (T) in all cases. Therefore, the argument is invalid because there are cases where the premises are true, but the conclusion is false. To find a counterexample, we need to find a row in the truth table where the premises are true (T) and the conclusion is false (F). In this case, the counterexample is when H is true (T), I is false (F), and G is false (F). Therefore, the correct answer is: C. Invalid. Counterexample when H is true and I and G are false. Answer Option 7:

To construct a complete truth table for the argument, we need to consider all possible combinations of truth values for the variables G, H, and I.

To construct a truth table for the argument, we need to consider all possible truth values for the propositions G, H, and I.

```
G H I (G \equiv H) \cdot \sim I \sim G \vee (\sim H \vee I) G'
TTTFFTT
 TTFFFTT
 TETEETT
TFFFFTT
FTTFFTF
FTFFTTF
 FFTFTTF
FFFTTTT
In the truth table, we evaluate the statements (G \equiv H) \cdot ~I, ~G \vee (~H \vee I), and G'.
 (G \equiv H) \cdot \sim I evaluates to F in rows 1, 2, 3, and 4, and T in rows 5, 6, 7, and 8.
 ~G \vee (~H \vee I) evaluates to T in all rows except row 5, where it evaluates to F.
 G' evaluates to T in rows 2, 3, 4, 5, 6, 7, and 8, and F in row 1.
 Based on the truth table, we can see that there is a row (row 5) where the premises (-G v (-H v I)) do not imply the conclusion (G'). Therefore, the argument is invalid.
The correct answer is: C. 'Invalid. Counterexample when H is true and I and G are false'.
 *******
 Answer Option 8:
 To determine the validity of the argument, we need to construct a complete truth table for the premises and the conclusion and check if there is any row where all the premises are true and the conclusion is
 false.
Let's break down the argument:
 Premises:
 1. (G ≡ H) · ~I
2. \simG \vee (\simH \vee I)
 Conclusion:
Let's construct the truth table:
\mid G \mid H \mid I \mid (G \equiv H) \cdot \sim I \mid \sim G \vee (\sim H \vee I) \mid G' \mid
|---|---|
|T|T|T|F|T|F|
|T|T|F|F|T|F|
|T|F|T|F|T|F|
|T|F|F|T|T|F|
|F|T|T|F|T|T|
|F|T|F|F|F|T|
|F|F|T|T|T|T|
|F|F|F|F|T|T|
From the truth table, we can see that there is a row where all the premises are true (row 3, where G = F, H = F, I = T) but the conclusion is false (G' = T). Therefore, the argument is invalid.
 The correct answer is option C: Invalid. Counterexample when H is true and I and G are false.
 To construct a truth table for the argument, we need to consider all possible combinations of truth values for the propositions G, H, and I.
 The given premises are:
 1. (G \equiv H) \cdot \sim I
2. ~G v (~H v I)
We also need to determine the truth value of the conclusion G.
 Here is the complete truth table for the argument:
\mid G \mid H \mid I \mid (G \equiv H) \cdot \sim I \mid \sim G \vee (\sim H \vee I) \mid G' \mid
|T|T|T|F|T|F|
|T|T|F|T|T|F|
|T|F|T|F|T|F|
```

|T|F|F|T|T|F| |F|T|T|F|T|F|

|F|T|F|T|T|F| |F|F|T|T|T|F| |F|F|F|F|T|T|

Now, let's analyze the truth table to determine whether the argument is valid or invalid.

The argument is valid if and only if the conclusion (G') is true for all rows where both premises are true.

In the truth table, we can see that there are two rows where both premises are true (rows 6 and 7), but the conclusion G' is false in both of those rows.

Therefore, the argument is invalid, and the counterexample is when H is true and I and G are false.

Final Answer: C. Invalid. Counterexample when H is true and I and G are false.

Question 124:
'Construct a complete truth table for the following argument. Then, using the truth table, determine whether the argument is valid or invalid. If the argument is invalid, choose an option which presents a
counterexample. (There may be other counterexamples as well.)
\sim (M = \sim N)
M / N', A. 'Valid', B. 'Invalid. Counterexample when M and N are true', C. 'Invalid. Counterexample when M is true and N is false', D. 'Invalid. Counterexample when N is true and M is false'
Answer Option 1:
To construct a complete truth table for the argument, we need to consider all possible combinations of truth values for M and N.
$M/N/\sim (M = \sim N)$
T/T/F
T/F/T
F/T/T
F/F/F
The argument is valid if the conclusion ~(M = ~N) is true for every row where the premises M and N are true. Looking at the truth table, we see that the conclusion is true for every row where the premises are
true. Therefore, the argument is valid.
Final Answer: A. Valid
Answer Option 2:
To construct a complete truth table for the argument, we consider all possible combinations of truth values for M and N and evaluate the expression $\neg (M \equiv \neg N)$.
M N \sim N \sim \sim N
FITIFIFIT
F F T T F
Since the expression -(M = -N) evaluates to false in all rows, the argument is valid.
Final Answer: A. Valid
Answer Option 3:
To construct a truth table, we need to consider all possible combinations of truth values for the propositions M and N.
$M/N \mid \sim (M = \sim N)$
TT F
TF T
FT T
FF F
From the truth table, we can see that there are cases where the premise \sim (M \equiv \sim N) is true and the conclusion is false. Specifically, when M is true and N is false, the premise is true but the conclusion is false.
Therefore, the argument is invalid. The correct answer is option C: 'Invalid. Counterexample when M is true and N is false'.
Answer Option 4:
To construct a complete truth table for the argument, let's consider all possible combinations of truth values for M and N:
$M \mid N \mid \sim (M = \sim N)$
TITIF
TIFIT
F T T
F F F
From the truth table, we can see that there are cases where \sim (M \equiv \sim N) is true and cases where it is false. Therefore, the argument is invalid.
To find a counterexample, we can look for a row in the truth table where the premises (M and N) are true, but the conclusion is false. In this case, the counterexample is when M is true and N is false.
Therefore, the correct answer is: C. 'Invalid. Counterexample when M is true and N is false'

Answer Option 5:

To construct a complete truth table for the argument, we need to consider all possible combinations of truth values for M and N. $M/N \sim (M \equiv \sim N)$ T/TF T/FT F/TT F/FF From the truth table, we can see that there are cases where the statement ~(M = ~N) is false, indicating that the argument is invalid. To find the counterexample, we need to identify a row in the truth table where the statement \sim (M \equiv \sim N) is false. Looking at the truth table, when M is true and N is false, the statement ~(M ≡ ~N) evaluates to true. Therefore, the correct answer is C. 'Invalid. Counterexample when M is true and N is false'. Answer Option 6: $To\ construct\ a\ complete\ truth\ table\ for\ the\ argument,\ we\ need\ to\ consider\ all\ possible\ combinations\ of\ truth\ values\ for\ M\ and\ N.$ $M/N \mid \sim (M \equiv \sim N)$ TTI TFI FT FFI To determine whether the argument is valid or invalid, we need to check if the statement ~(M = ~N) is always true for every row in the truth table. ~(M = ~N) evaluates to: \sim (T \equiv \sim T) : F ~(T ≡ ~F) : F \sim (F \equiv \sim T) : T ~(F ≡ ~F) : F Since there is at least one row (M = F, N = T) where the statement \sim (M \equiv \sim N) is true, the argument is invalid. Thus, the correct answer is D. 'Invalid, Counterexample when N is true and M is false'. Answer Option 7: To construct a complete truth table for the argument and determine its validity, we need to consider all possible combinations of truth values for M and N. $M/N \mid \sim (M \equiv \sim N)$ TTI TFI FTI FFI To evaluate \sim (M \equiv \sim N), we consider the truth value of the equivalence statement (M \equiv \sim N) and negate it. $M / N \mid M \equiv \sim N \mid \sim (M \equiv \sim N)$ TT|F|T TF|T|F FT|T|F FF|F|T Based on the truth table, we can see that the argument is invalid because there are instances where the premise (~(M = ~N)) is true, but the conclusion is false. The counterexample is when M is true and N is false. In this case, ~(M = ~N) is true, but the argument states that the conclusion is false, which is not the case. Therefore, the correct option is C. 'Invalid. Counterexample when M is true and N is false'.

Answer Option 8:

To construct a complete truth table, we need to consider all possible combinations of truth values for the variables M and N.

 $M / N / M \equiv \sim N / \sim (M \equiv \sim N)$

T/T/F/F

T/F/T/T

F/T/T/T

F/F/F/F

The truth table above represents all possible combinations of truth values for M and N, as well as the resulting truth value for the proposition \sim (M \equiv \sim N).

From the truth table, we can see that there are cases where $-(M \equiv -N)$ is false (F) even when M and N have different truth values. This means that the argument is invalid.

To find a counterexample, we can look for a row in the truth table where the premise \sim ($M \equiv \sim$ N) is true (T) and the conclusion is false (F). In this case, the row where M is true and N is false satisfies this condition.

Thus, the argument is invalid. Counterexample: M is true and N is false.

Final Answer: C. Invalid. Counterexample when M is true and N is false

Answer Option 9:

To construct a complete truth table for the argument, we need to consider all possible combinations of truth values for M and N.

 $M / N \mid \sim (M \equiv \sim N) \mid$

TT|T|

TF|F|

FT|F|

FFIFI

The truth table above shows all possible combinations of truth values for M and N, and the corresponding truth value of the proposition \sim (M \equiv \sim N).

Now, let's analyze the truth values of the proposition \sim (M \equiv \sim N) in each row of the truth table:

In the first row where M and N are both true, ${\sim}(M\equiv{\sim}N)$ is true.

In the second row where M is true and N is false, \sim (M \equiv \sim N) is false.

In the third row where M is false and N is true, \sim (M \equiv \sim N) is false.

In the fourth row where M and N are both false, ${\sim}(M\equiv{\sim}N)$ is false.

Since there exists at least one row (the second row) in which the proposition \sim (M \equiv \sim N) is false, the argument is invalid.

 $Counterexample: M \ is \ true \ and \ N \ is \ false. \ In \ this \ case, \ {\scriptstyle \sim}(M\equiv {\scriptstyle \sim}N) \ is \ false, \ which \ contradicts \ the \ conclusion \ that \ the \ argument \ is \ valid.$

Therefore, the correct answer is C. 'Invalid. Counterexample when M is true and N is false'.

Question 125:

'Select the best translation into predicate logic. Some CEO is wealthier than everyone. (Cx: x is a CEO; Px: x is a person; Wxy: x is wealthier than y)', A. '(∃x)[Cx • (∀y)(Py • Wxy)]', B. '(∃x)[Cx ⊃ (∀y)(Py ⊃

 $\mathsf{Wxy)}]',\,\mathsf{C}.\,\,'(\exists x)[\mathsf{C}x\bullet(\forall y)(\mathsf{P}y\supset\mathsf{W}xy)]',\,\mathsf{D}.\,\,'(\forall x)[\mathsf{P}x\supset(\exists y)(\mathsf{C}y\supset\mathsf{W}xy)]'$

Answer Option 1:

The best translation into predicate logic for the statement "Some CEO is wealthier than everyone" is option A: $'(\exists x)[Cx \cdot (\forall y)(Py \cdot Wxy)]'$.

This translation reads as: "There exists an x such that x is a CEO and for all y, if y is a person, then x is wealthier than y."

This accurately captures the idea that there is at least one CEO who is wealthier than every person. Final Answer: A

Answer Option 2:

The best translation into predicate logic for the given statement "Some CEO is wealthier than everyone" is:

 $(\exists x)[Cx \bullet (\forall y)(Py \supset Wxy)]$

Explanation:

- ($\exists x$) denotes existence, meaning there exists some x.
- Cx represents x is a CEO.
- $(\forall y)(Py\supset Wxy)$ represents for every y, if y is a person, then x is wealthier than y.

So, the correct answer is A. $(\exists x)[Cx \bullet (\forall y)(Py \supset Wxy)].$

Answer Option 3:

The best translation into predicate logic for the statement "Some CEO is wealthier than everyone" is: $(\exists x)[Cx \cdot (\forall y)(Py \cdot Wxy)]'$.

Explanation:

- (3x) indicates the existence of at least one x (CEO) that satisfies the following conditions.
- Cx represents x being a CEO.
- (∀y)(Py Wxy) states that for all y (persons), they satisfy the conditions of being a person (Py) and x (CEO) being wealthier than y (Wxy).

Therefore, option A is the correct translation. Final Answer: A.

Answer Option 4:

The best translation into predicate logic for the given statement "Some CEO is wealthier than everyone" would be:

 $(\exists x)[Cx \bullet (\forall y)(Py \supset Wxy)]$

Explanation:

- $(\exists x)$ signifies the existence of at least one x (some) that satisfies the following conditions.
- Cx represents x being a CEO.
- (∀y) signifies for all y (everyone).
- (Py \supset Wxy) represents the condition that if y is a person, then x is wealthier than y.

Therefore, the correct translation is option C: ' $(\exists x)[Cx \bullet (\forall y)(Py \supset Wxy)]'$.

Answer Option 5:

The best translation into predicate logic for the statement "Some CEO is wealthier than everyone" is option A: $'(\exists x)[Cx \bullet (\forall y)(Py \bullet Wxy)]'$.

Let's break down the translation:

Cx: x is a CEO

Py: y is a person

Wxy: x is wealthier than y

The statement "Some CEO is wealthier than everyone" can be translated as: There exists an x such that x is a CEO and for all y, if y is a person, then x is wealthier than y.

This translation captures the idea that there is at least one CEO who is wealthier than every person.

Final Answer: A. $(\exists x)[Cx \bullet (\forall y)(Py \bullet Wxy)]'$

Answer Option 6:

The correct translation into predicate logic for the given statement is:

" $(\exists x)[Cx \bullet (\forall y)(Py \bullet Wxy)]$ "

Explanation:

- $(\exists x)$ denotes the existence of at least one CEO.
- Cx represents x being a CEO.
- ($\forall y$) denotes for all individuals y.

- Wxy represents x being wealthier than y.
The statement translates to "There exists some individual x who is a CEO such that for all individuals y, if y is a person, then x is wealthier than y."
Therefore, the correct answer is A. " $(\exists x)[Cx \bullet (\forall y)(Py \bullet Wxy)]$ ".
Answer Option 7:
The best translation into predicate logic for the statement "Some CEO is wealthier than everyone" would be option A: '(∃x)[Cx • (∀y)(Py • Wxy)]'.
Explanation:
- (∃x) represents "There exists".
- Cx represents "x is a CEO".
- $(\forall y)$ represents "For all y ".
- Py represents "y is a person".
- Wxy represents "x is wealthier than y".
Putting it all together:
(3x)[Cx • (∀y)(Py • Wxy)] translates to "There exists a person x who is a CEO and for all persons y, x is wealthier than y". This accurately captures the meaning of the original statement.
Final Answer: A

Answer Option 8:
The best translation into predicate logic for the statement "Some CEO is wealthier than everyone" would be:
$(\exists x)[Cx \cdot (\forall y)(Py \cdot Wxy)]$
Explanation:
- (∃x) indicates the existence of an individual x.
- Cx represents "x is a CEO".
$- (\forall y) \ indicates \ that \ the \ following \ statement \ applies \ to \ all \ individuals \ y.$
- Py represents "y is a person".
- Wxy represents "x is wealthier than y".
The statement can be broken down as follows:
$- (\exists x) [Cx \bullet (\forall y) (Py \bullet Wxy)] : \text{There exists an individual } x \text{ who is a CEO such that for all individuals } y \text{ who are people, } x \text{ is wealthier than } y.$
This translation correctly captures the statement that there is at least one CEO who is wealthier than everyone else.
Therefore, the correct answer is A. $'(\exists x)[Cx \cdot (\forall y)(Py \cdot Wxy)]'$.
Answer Option 9:
The best translation into predicate logic for the statement "Some CEO is wealthier than everyone" is:
$(\exists x)[Cx \bullet (\forall y)(Py \supset Wxy)]$
Explanation:
- (3x) indicates the existence of at least one x that satisfies the following conditions.
- Cx: x is a CEO.
- $(\forall y)(Py \supset Wxy)$: for all y , if y is a person (Py) , then x is wealthier than y (Wxy) .
Option A: (∃x)[Cx • (∀y)(Py • Wxy)]
This option uses the conjunction operator (\bullet) incorrectly. It should be $(\forall y)(Py \supset Wxy)$ instead of $(\forall y)(Py \bullet Wxy)$.
Option B: $(\exists x)[Cx \supset (\forall y)(Py \supset Wxy)]$
This option uses the implication operator (\supset) incorrectly. It should be ($\forall y$)($Py \supset Wxy$) instead of ($Cx \supset (\forall y)(Py \supset Wxy)$).
Option C: $(\exists x)[Cx \bullet (\forall y)(Py \supset Wxy)]$
This option is the correct translation and matches the original statement.
Option D: $(\forall x)[Px \supset (\exists y)(Cy \supset Wxy)]$
This option incorrectly uses the universal quantifier $(\forall x)$ instead of the existential quantifier $(\exists x)$ to indicate the existence of at least one CEO.
Final Answer: A

- Py represents y being a person.

Question 126: Select the best translation into predicate logic. All children go to some school. (Cx: x is a child; Sx: x is a school; Gxy: x goes to y)', A. '(∀x)(∃y)([Cx • Sy) • Gxy)', B. '(∀x)(Sx ⊃ (∃y)(Cy • Gxy))', C. '(∀x)(Cx ⊃ (∃y)(Cy ⊃ (∃x)(Cx ⊃ (∃y)(Cy ⊃ (∃x)(Cx ⊃ $(\exists y)(Sy \bullet Gxy)]',\, D.\ '(\forall x)[Sx \bullet (\exists y)(Cy \bullet Gxy)]'$ Answer Option 1: The best translation into predicate logic for the given statement would be: $(\forall x)[Cx \supset (\exists y)(Sy \cdot Gxy)]'$, which is Option C. Let's break down the translation: Cx: x is a child Sx: x is a school Gxy: x goes to y The statement "All children go to some school" can be translated as "For all x, if x is a child, then there exists a y such that y is a school and x goes to y". This can be represented as: $(\forall x)[Cx\supset (\exists y)(Sy\bullet Gxy)]$ Therefore, the correct answer is Option C. Answer Option 2: The best translation into predicate logic for the statement "All children go to some school" is: $(\forall x)[Cx\supset (\exists y)(Sy\bullet Gxy)]$ Explanation: - ($\forall x$) translates to "for all x" - Cx represents "x is a child" - Sy represents "y is a school" - Gxy represents "x goes to y" So, $(\forall x)[Cx \supset (\exists y)(Sy \bullet Gxy)]$ can be read as "For all x, if x is a child, then there exists a y such that y is a school and x goes to y." This accurately captures the idea that all children go to some school. Therefore, the correct answer is C. Answer Option 3: The best translation into predicate logic for the statement "All children go to some school" would be: $\mathsf{B.}\ '(\forall x)[\mathsf{S}x\supset (\exists y)(\mathsf{C}y\bullet\mathsf{G}xy)]'$ Let's break down the statement: - "All children" is represented as $(\forall x)Cx,$ which means for all $x,\,x$ is a child. - "go to some school" is represented as (∃y)(Sy • Gxy), which means there exists a y such that y is a school and x goes to y. Combining these parts using the implication operator (⊃), we get: $(\forall x)[Sx\supset (\exists y)(Cy\bullet Gxy)]$ This translation accurately represents the original statement "All children go to some school". Final Answer: B.

Answer Option 4:

The best translation into predicate logic is option C: $\label{eq:continuous} \text{(} \forall x)[Cx \supset (\exists y)(Sy \bullet Gxy)]\text{'}.$

Explanation:

- $(\forall x)$ is used to express "for all x" or "for every x".
- Cx represents "x is a child".
- \supset is the symbol for implication, meaning "implies" or "if...then".
- ($\exists y$) is used to express "there exists a y" or "for some y".
- Sy represents "y is a school".
- Gxy represents "x goes to y".

Putting it all together, the translation $[\forall x][Cx \supset (\exists y)[Sy \bullet Gxy)]$ can be read as "For every x, if x is a child, then there exists a y such that y is a school and x goes to y." This accurately captures the statement

"All children go to some school" in predicate logic. Final Answer: C

Answer Option 5:

The best translation into predicate logic for the statement "All children go to some school" would be option B: $(\forall x)[Sx \supset (\exists y)(Cy \bullet Gxy)]^t$.

Let's break down the translation:

- ($\forall x$) signifies "for all x", meaning the statement applies to every individual.
- Sx represents "x is a school".

- $(\exists y)$ signifies "there exists a y", meaning there is at least one y that satisfies the condition. - Cy represents "y is a child". - Gxy represents "x goes to y". Putting it all together, the translation reads as: "For every x, if x is a school, then there exists a y such that y is a child and x goes to y". This accurately captures the statement "All children go to some school". So, option B is the correct translation. Final Answer: B. Answer Option 6: The best translation into predicate logic for the statement "All children go to some school" is option C: $(\forall x)[Cx \supset (\exists y)(Sy \bullet Gxy)]'.$ Explanation: - $(\forall x)$ indicates that the statement applies to all individuals x.- Cx represents "x is a child". - ($\exists y$) represents "there exists a y". - Sy represents "y is a school". - Gxy represents "x goes to y". Thus, when we translate the statement into predicate logic, we express that for every child x, if x is a child (Cx), then there exists a school y such that y is a school (Sy) and x goes to y (Gxy). This accurately represents the original statement "All children go to some school". Final Answer: C. '($\forall x$)[Cx \supset ($\exists y$)(Sy • Gxy)]' Answer Option 7: The best translation into predicate logic for the given statement "All children go to some school" is: $(\forall x)(\exists y)[(Cx \bullet Sy) \bullet Gxy]$ Explanation: - $(\forall x)$ represents "For all x" or "For every x". - ($\exists y$) represents "There exists a y" or "There is a y". - Cx represents "x is a child". - Sy represents "y is a school". - Gxy represents "x goes to y". The translation $(\forall x)(\exists y)[(Cx \cdot Sy) \cdot Gxy]$ states that for every child x, there exists a school y such that x is a child, y is a school, and x goes to y. This accurately captures the statement *All children go to some Therefore, the correct answer is A. $(\forall x)(\exists y)[(Cx \bullet Sy) \bullet Gxy]$. ****************** Answer Option 8: The best translation into predicate logic for the statement "All children go to some school" would be option B: $(\forall x)[Sx \supset (\exists y)(Cy \bullet Gxy)]^t$. Explanation: Cx: x is a child Sx: x is a school Gxy: x goes to y The statement "All children go to some school" can be translated as "For every x, if x is a child, then there exists a y such that y is a school and x goes to y." This can be represented using the following predicate logic formula: $(\forall x)[Cx\supset (\exists y)(Sy\bullet Gxy)]$ Here, $(\forall x)$ represents "For every x," $Cx \supset (\exists y)(Sy \bullet Gxy)$ represents "if x is a child, then there exists a y such that y is a school and x goes to y." Therefore, option B is the correct translation into predicate logic. Final Answer: B Answer Option 9: The best translation into predicate logic for the statement "All children go to some school" is: $(\forall x)(\exists y)[(Cx \bullet Sy) \bullet Gxy]$ This translation states that for all x (individuals), there exists a y (school) such that x is a child and y is a school, and x goes to y. Option A: $(\forall x)(\exists y)[(Cx \bullet Sy) \bullet Gxy]$ - This is the correct translation. Option B: $(\forall x)[Sx \supset (\exists y)(Cy \bullet Gxy)]$ - This translation states that if x is a school, then there exists a y such that y is a child and x goes to y. This does not capture the meaning of the original statement. Option C: $(\forall x)[Cx\supset (\exists y)(Sy \bullet Gxy)]$ - This translation states that if x is a child, then there exists a y such that y is a school and x goes to y. This does not capture the meaning of the original statement.

Option D: (∀x)[Sx • (∃y)(Cy • Gxy)] - This translation states that for all x (individuals), if x is a school, then there exists a y such that y is a child and x goes to y. This does not capture the meaning of the

original statement Final Answer: A
