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- Extends Thread Class
- Implements Runnable Interface
- Thread Properties and Methods

- States of a Java Thread
- Useful Methods for Java Threads
- Life Cycle of a thread

- Synchronising in Java
 - Example of using Multiple Locks
- Deadlocks

④ Communicating Between threads

Why use threads

- ## Threading in the JVM

- ## The Methods of Creating Threads

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Information on the thread Class

- ### Example

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Information on the Runnable Interface

- ```
– public abstract void run();
```

## Why use a Runnable Interface over a Thread Abstract Class

- This provides more opportunity for polymorphism

## Thread Methods

```
Thread.sleep(long n)
```

- Makes the current thread pause for n milliseconds
- During this time, the thread enters a "sleeping" state
- Must be in a try-catch block as it throws InterruptedException

```
1 try {
2 Thread.sleep(1000); // Sleep for 1 second
3 } catch (InterruptedException e) {
4 e.printStackTrace();
5 }
```

## Thread Priorities

- Threads have priorities from 1 (lowest) to 10 (highest)
- Set using `setPriority(int priority)`
- Constants available: `MIN_PRIORITY`, `NORM_PRIORITY`, `MAX_PRIORITY`
- Higher priority threads are generally executed first

Thread.yield()

- Temporarily pauses current thread to let other threads execute
- Thread moves from "running" state to "ready" state
- No guarantee which thread will execute next
- Primarily used for thread scheduling optimization
  - Helps prevent any single thread from monopolizing CPU time
  - Particularly useful in busy-wait scenarios
  - Can improve overall system responsiveness
- Important considerations:
  - Behavior is highly platform-dependent
  - Only a scheduling hint - may be ignored by the JVM
  - Not recommended for synchronization purposes
- Example usage:

```
1 if(threadNeedsToWait) {
2 Thread.yield(); // Give other threads a chance
3 }
```

## Interrupting Threads

- Use `interrupt()` to interrupt a thread
- Interrupted threads throw `InterruptedException`
- Check if interrupted using `isInterrupted()`
- Useful for stopping long-running or sleeping threads

## What are the States of a Thread

- ## Setters and Getters for Thread Names

- ## Daemon Threads

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```
isAlive()
```

- join()

- The `join()` method is used to make one thread wait for another thread's completion
- When `threadA.join()` is called from `threadB`:
  - `threadB` will pause its execution
  - `threadB` will wait until `threadA` completes
  - This ensures sequential execution when needed
- Has three variants:
  - `join()` - waits indefinitely
  - `join(long millis)` - waits for specified milliseconds
  - `join(long millis, int nanos)` - adds nanosecond precision
- Common use cases:
  - Waiting for background tasks to complete
  - Ensuring proper order of operations
  - Coordinating dependent thread activities



## Thread States and Transitions

- New State
  - Thread is created but not yet started
  - Transitions to Runnable when `start()` is called
- Runnable State
  - Thread is executing or ready to execute
  - Can transition to:
    - \* Running - when selected by scheduler
    - \* Blocked/Waiting - when requesting resources
    - \* Terminated - when execution completes
- Blocked/Waiting States
  - Thread is temporarily inactive
  - Caused by:
    - \* `sleep()` calls
    - \* I/O operations
    - \* Lock acquisition
    - \* `join()` calls
- Terminated State
  - Thread has completed execution
  - Cannot be restarted

# Synchronisation

## The need for Synchronisation

- All threads in the JVM **share the heap**
  - Any fields in an object can be accessible to a number of threads at the same time
- After a thread has upated a field with a value, it's possible another thread comes along and updates the same field with a different value
- When one thread is trying to read the value from a field, another thread may be trying to change that value
- This can lead to a lot of funkiness to do with value change order
- Even simple operations like incrementing are not atomic:
  - A statement like `z = z + 1` compiles to multiple bytecode instructions
  - These include loading from heap, incrementing on stack, and storing back
  - Threads can be interrupted between any of these steps
  - This creates race conditions even for seemingly atomic operations

## Incrementation Example

- See below, the final output for `z` should be 4

```
1 public class UnsafeCounter {
2 private int z = 2; // Initial value is 2
3 void incr() {
4 int x = z; // Thread 1: reads z=2
5 // Thread 2: reads z=2 (before Thread 1 updates z)
6 x = x + 1; // Thread 1: x becomes 3
7 // Thread 2: x becomes 3 (both threads working with original z value)
8 z = x; // Thread 1: sets z=3
9 // Thread 2: sets z=3 (overwrites Thread 1's update)
10 } // Final z=3, even though we wanted z=4
11 }
12 public class ThreadDemo {
13 public static void main(String[] args) {
14 UnsafeCounter counter = new UnsafeCounter();
15 Thread t1 = new Thread(() -> counter.incr());
16 Thread t2 = new Thread(() -> counter.incr());
17 t1.start();
18 t2.start();
19 // After both threads finish, z might still be 3 instead of 4!
20 }
21 }
```

## Locks

- ## Synchronised Blocks

## Synchronised Methods

## The problem with Synchronisation

## 'this' lock and multiple locks

# Synchronising in Java - Example of using Multiple Locks

## Synced Bank account using Sync Blocks

```
1 public class BankAccount {
2 private int savingsBalance = 0;
3 private int checkingBalance = 0;
4
5 // Separate lock objects for each account type
6 private final Object savingsLock = new Object();
7 private final Object checkingLock = new Object();
8
9 // Deposit to savings account
10 public void depositToSavings(int amount) {
11 synchronized (savingsLock) {
12 savingsBalance += amount;
13 }
14 }
15
16 // Deposit to checking account
17 public void depositToChecking(int amount) {
18 synchronized (checkingLock) {
19 checkingBalance += amount;
20 }
21 }
22
23 // Get savings balance
24 public int getSavingsBalance() {
25 synchronized (savingsLock) {
26 return savingsBalance;
27 }
28 }
29
30 // Get checking balance
31 public int getCheckingBalance() {
32 synchronized (checkingLock) {
33 return checkingBalance;
34 }
35 }
36 }
```

- This example allows threads to safely update or read the savings and checking balances independently.
- A thread updating the savings account does not block another thread updating the checking account, increasing concurrency.

# Deadlocks

## What is a Deadlock?

- A deadlock occurs when two or more threads are waiting for each other to release resources
- Each thread holds a resource that another thread needs
- None can proceed because they're all waiting
- Real-life example:
  - Two cars meet on a narrow bridge from opposite directions
  - Neither can proceed because both need the other to back up
  - Both are waiting for the other to move first

## Code Example of Deadlock

```
1 public class DeadlockExample {
2 private Object lock1 = new Object();
3 private Object lock2 = new Object();
4
5 public void method1() {
6 synchronized(lock1) {
7 System.out.println("Method 1 has lock1");
8 try { Thread.sleep(100); } catch (InterruptedException e) {}
9 synchronized(lock2) {
10 System.out.println("Method 1 has lock2");
11 }
12 }
13 }
14
15 public void method2() {
16 synchronized(lock2) {
17 System.out.println("Method 2 has lock2");
18 try { Thread.sleep(100); } catch (InterruptedException e) {}
19 synchronized(lock1) {
20 System.out.println("Method 2 has lock1");
21 }
22 }
23 }
24 }
```

## Preventing Deadlocks

- Always acquire locks in the same order
- Use timeouts when acquiring locks
- Avoid nested locks where possible
- Use `tryLock()` methods from `java.util.concurrent.locks`
- Implement deadlock detection mechanisms

- Sometimes threads need to communicate or coordinate
  - For example, a consumer thread may need to **wait** for a producer to put data in the shared buffer
  - Java has methods that allow threads to pause execution until a certain criterion is met
    - \* Threads can also **notify** other threads when such condition has changed

- A monitor is a mechanism that ensures:
  - Only one thread can execute a synchronized method/block at a time
  - Other threads must wait for the active thread to finish
  - Provides a way for threads to coordinate through wait/notify
- When a thread enters a synchronized block:
  - It acquires the object's monitor
  - Other threads cannot acquire the same monitor
  - Monitor is released when thread exits the synchronized block

- The `wait()` method causes **the thread it is called in** to release the lock it holds and **wait** until another thread calls `notify()` or `notifyAll()`
- Must be called inside a **synchronized block** or **method**
- Can also include a timeout by adding numbers in milliseconds to the constructor
- The `wait()` method is called on the same object as the lock in the synchronised block

- The `notify()` method wakes up a **single** thread that has previously been **waiting** on the object's **monitor**
- The awakened thread only proceed if the synchronised block is **unlocked** - it will still wait until this is true before proceeding
- The `notify()` method is called on the same object as the lock in the synchronised block

- The `notifyAll()` method has the same core functionality as `notify()` but wakes up **all** threads waiting on the object's **monitor**
- Only one thread will acquire the lock and proceed; others will continue waiting for the lock.
- The `notifyAll()` method is called on the same object as the lock in the synchronised block