

Arrays and Linked Lists

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1 Arrays

1.1 Fixed Length

1.1.1 Properties

- Contiguous chunk of memory
 - All the elements of the array are stored in adjacent memory locations, one right after the other, without any gaps in between
- Has an access time of $\Theta(1)$
- Very efficient use of memory
 - Often provides best performance

1.1.2 Disadvantages

- Fixed length - Can be variable at extra time cost
- Insertion and deletion to/from the middle have $\Theta(n)$ time complexity

1.2 Variable Length

- Most `add(elem)` operations are $\Theta(1)$
- When a chunk of memory for an array is full, we need to copy all elements to a new larger chunk in order to add more elements
 - This resizing operation has a time complexity of $\Theta(n)$, where n is the number of elements in the array

1.3 General Time Analysis of Adding Elements

- To perform N adds with an initial capacity of n
- We must perform m copies where:

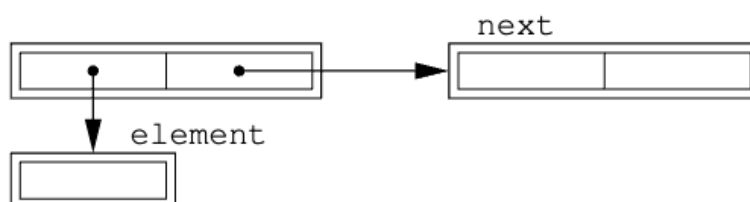
$$n \cdot 2^{m-1} < N \leq n \cdot 2^m$$

- Total Elements copied

$$\sum_{i=1}^m 2^{i-1} \cdot n = n(2^m - 1)$$

2 Linked Lists

- Non-Contiguous data use pointers to reference their units of data
- Removes the disadvantage of Contiguous data structures of adding and removing data from the middle being expensive
- A linked list is built up of nodes
- Each node contains a value and a reference to the next node in the list
- The final value in the list has a null reference



2.1 Singly Linked List

- A linked list that simply is a string of nodes to eachother in **one direction**
- Has a single pointer to the **next node only**
- Differs to a doubly linked lists which has pointers to the *next and previous* nodes

2.1.1 Java Implementation

```
1  /**
2   * A simple implementation of a singly linked list in Java.
3   * This class demonstrates the basic operations of a linked list.
4   */
5  public class MyLinkedList<E> {
6      // The head node of the linked list
7      private Node<E> head;
8      // The number of elements in the linked list
9      private int noElements;
10
11     /**
12      * A static nested class representing a node in the linked list.
13      * Each node contains an element and a reference to the next node.
14      */
15     private static class Node<T> {
16         private T element; // The data stored in the node
17         private Node<T> next; // Reference to the next node in the list
18     }
19
20     /**
21      * Constructor to create an empty linked list.
22      * Initially, the head is null and the number of elements is zero.
23      */
24     public MyLinkedList() {
25         head = null;
26         noElements = 0;
27     }
28
29     /**
30      * Returns the number of elements in the linked list.
31      * @return the size of the linked list
32      */
33     public int size() {
34         return noElements;
35     }
36
37     /**
38      * Checks if the linked list is empty.
39      * @return true if the linked list is empty, false otherwise
40      */
41     public boolean isEmpty() {
```

```

42     return head == null;
43 }
44
45 /**
46  * Adds a new element to the front of the linked list.
47  * @param element the element to be added
48  * @return true if the element is added successfully
49  */
50 public boolean add(E element) {
51     // Create a new node with the given element
52     Node<E> newNode = new Node<E>();
53     newNode.element = element; // Set the element of the new node
54     newNode.next = head; // Set the next reference of the new node to the
55     ↪ current head
56     head = newNode; // Update the head to the new node
57     noElements++; // Increment the number of elements
58     return true;
59 }
60
61 /**
62  * Removes the head (first element) of the linked list.
63  * @return true if the head is removed successfully, false if the list is empty
64  */
65 public boolean remove_head() {
66     if (!isEmpty()) {
67         head = head.next; // Update the head to the next node
68         noElements--; // Decrement the number of elements
69         return true;
70     }
71     return false; // Return false if the list is empty
72 }
73
74 /**
75  * Checks if the linked list contains a specific element.
76  * @param obj the element to check for
77  * @return true if the element is found, false otherwise
78  */
79 public boolean contains(E obj) {
80     for (Node<E> current = head; current != null; current = current.next) {
81         if (obj.equals(current.element)) {
82             return true;
83         }
84     }
85     return false;
86 }

```

3 Stack Implementation with Linked Lists

```
1 public class LinkedListStack<E>
2 {
3     private MyLinkedList<E> list = new MyLinkedList<E>();
4
5     boolean push(E obj) { // New element becomes the head
6         return list.add(obj);
7     }
8
9     E peek() { // Only return the head of the linked list
10        return list.get_head();
11    }
12
13    E pop() { // Removes the head of the linked list
14        if (isEmpty()) throw EmptyStackException;
15        T elem=list.get_head();
16        list.remove_head();
17        return elem;
18    }
19    boolean isEmpty() {
20        return list.isEmpty();
21    }
22 }
```

- Stack operations of linked list take constant time - $\Theta(1)$
- There is a hidden cost of creating and destroying `Node` objects
- Memory requirement is $\Theta(n)$
- Array implementation of stacks are better in practice - constant time with no hidden cost

4 Queue Implementation with Linked List

- For Queues, we add at one end and remove from the other.
- Use a **Head and Tail**
 - Enqueue at the back
 - Dequeue at the head

5 Java Linked Lists

- Allows add and remove at **both** ends of the list
- Uses a doubly linked list with pointers to the next and previous nodes
- Uses a new node that has a null value but **next** and **previous** reference the head and tail. Last element references this node.

6 Skip Lists

6.1 Introduction

- Skip lists are a data structure that allows fast search within an ordered sequence of elements.
- They are a probabilistic alternative to balanced trees.
- Skip lists use multiple levels of linked lists to achieve logarithmic time complexity for search, insertion, and deletion operations.

6.2 Structure

- A skip list consists of multiple layers, where each layer is a sorted linked list.
- The bottom layer contains all the elements, and each higher layer acts as an "express lane" with fewer elements.
- Each element in a layer has a reference to the next element in the same layer and possibly to an element in the layer below.

6.3 Operations

6.3.1 Search

- Start from the top-left element.
- Move right until the next element is greater than or equal to the target.
- Move down one level and repeat until the target is found or the bottom level is reached.
- Average time complexity: $O(\log n)$.

6.3.2 Insertion

- Perform a search to find the position where the new element should be inserted.
- Insert the element in the bottom layer.
- Promote the element to higher layers with a certain probability (e.g., 50% chance).
- Average time complexity: $O(\log n)$.

6.3.3 Deletion

- Perform a search to find the element to be deleted.
- Remove the element from all layers where it appears.
- Average time complexity: $O(\log n)$.

6.4 Advantages and Disadvantages

- **Advantages:**
 - Simple to implement.
 - Provides fast search, insertion, and deletion operations.
 - Probabilistic balancing without the need for complex rotations.
- **Disadvantages:**
 - Performance can degrade if the randomization does not work well.
 - Requires additional memory for multiple levels of pointers.