

General:

- **newtype Parser a = P (String -> [(a, String)])**
- Predicate: a function that takes one argument and returns a boolean
 - * if **pred** x == **True** then x satisfies predicate **pred**
- function composition:
 - *the . operator composes functions:*
 - (f . g) x == f (g x)

useful library functions:

```

Data.List
nubBy :: (a -> a -> Bool) -> [a] -> [a]
nubBy pred xs = -- unique elements only from xs as
  -> determined by pred
nub :: Eq a => [a] -> [a]
nub xs = nubBy (==) a -- unique elements from xs
--
words :: String -> [String]
words xs = -- list of whitespace-separated words from
  -> xs
--
-- concatenate container of lists
concat :: Foldable t => t [a] -> [a]
-- or for list-of-lists specifically:
concat :: [[a]] -> [a]
concat xs = foldl (++) [] xs
--
-- like concat, but use a function to get the inner
  -> lists
concatMap :: (a -> [b]) -> [a] -> [b]
concatMap f xs = foldr ((++) . f) [] xs
--
-- get the longest prefix of xs for which pred is true
  -> and also return the rest of the list
span :: (a -> Bool) -> [a] -> ([a], [a])
span pred xs = (takeWhile pred xs, dropWhile pred xs)
--
-- repeat a = infinite list of a
repeat :: a -> [a]
repeat x = map (\_ -> x) [1..]
repeat x = [ x | _ <- [1..] ]
-- replicate n a = list of length n repeating a
replicate :: Int -> a -> [a]
replicate n x = map (\_ -> x) [1..n]
replicate n x = [ x | _ <- [1..n] ]
--
-- folds (works on any foldable, not just lists)
foldr :: (a -> b -> b) -> b -> [a] -> b
foldr f z [a,b,c] = a `f` (b `f` (c `f` z))
foldr f z [a,b,c] = f a $ f b $ f c z
-- combines into z from right to left
-- can potentially work on an empty list if one of the
  -> folds does not evaluate it's second argument
foldl :: (b -> a -> b) -> b -> [a] -> b
foldl f z [a,b,c] = ((z `f` a) `f` b) `f` c
foldl f z [a,b,c] = f (f (f z a) b) c
-- evaluates from right to left
-- will not work on infinite list because it must start
  -> at the end of the list
--
-- these are the same as above, except they take the
  -> first two elements for the first application of f
foldr1 :: (a -> a -> a) -> [a] -> a
foldl1 :: (a -> a -> a) -> [a] -> a

```

Parsing.hs:

```
module Parsing where
import Data.Char
import Control.Applicative (Applicative(..))
import Control.Monad       (liftM, ap)

infixr 5 +++

-- Parser is the name of the type, P is the
--   ↪ constructor.
newtype Parser a = P (String -> [(a,String)])

instance Monad Parser where
```

```

return v = P (\inp -> [(v,inp)])
p >>= f = P (\inp -> case parse p inp of
    [(v,out)] -> [(v,inp)]
    _ -> out)

instance Functor Parser where
    fmap = liftM

instance Applicative Parser where
    pure = return
    (<*>) = ap

failure :: Parser a
failure = P (\inp -> [])

-- parse any single character
item :: Parser Char
item = P (\inp -> case inp of
    [] -> []
    (x:xs) -> [(x,xs)])

parse :: Parser a -> String -> [(a,String)]
parse (P p) inp = p inp

--Choice
(+++) :: Parser a -> Parser a -> Parser a
p +++ q = P (\inp -> case parse p inp of
    [] -> parse q inp
    [(v,out)] -> [(v,out)])

--Derived primitives
-----
sat :: Char -> Bool -> Parser Char
sat p = do x <- item
    if p x then return x else failure

digit :: Parser Char
digit = sat isDigit

letter :: Parser Char
letter = sat isAlpha

alphanum :: Parser Char
alphanum = sat isAlphaNum

lower :: Parser Char
lower = sat isLower

upper :: Parser Char
upper = sat isUpper

char :: Char -> Parser Char
char x = sat (== x)

string :: String -> Parser String
string [] = return []
string (x:xs) = do char x
    string xs
    return (x:xs)

many :: Parser a -> Parser [a]
many p = many1 p +++ return []
many1 :: Parser a -> Parser [a]
many1 p = do v <- p
    vs <- many p
    return (v:vs)

ident :: Parser String
ident = do x <- lower
    xs <- many alphanum
    return (x:xs)

nat :: Parser Int
nat = do xs <- many1 digit
    return (read xs)

int :: Parser Int
int = do char '-'
    n <- nat
    return (-n)
    +++ nat

space :: Parser ()
space = do many (sat isSpace)
    return ()

--Ignoring spacing

```

```

token      :: Parser a -> Parser a
token p    = do space
            v <- p
            space
            return v

identifier :: Parser String
identifier = token ident

natural   :: Parser Int
natural   = token nat

integer   :: Parser Int
integer   = token int

symbol    :: String -> Parser String
symbol xs = token (string xs)

--Example: Arithmetic Expressions
expr :: Parser Int
expr = do t <- term
        do {char '+'
            ; e <- expr
            ; return (t + e)
            }
        +++ return t

term :: Parser Int
term = do f <- factor
        do char '*'
            t <- term
            return (f * t)
        +++ return f

factor :: Parser Int
factor = do d <- digit
            return (digitToInt d)
        +++ do char '('
            e <- expr
            char ')'
            return e

eval :: String -> Int
eval xs = fst (head (parse expr xs))

• sat :: (Char -> Bool) -> Parser Char
  * returns a character if that character satisfies the predicate

• digit, letter, alphanum :: Parser Char
  * parses a digit, letter, or alpha-numeric letter respectively

• char :: Char -> Parser Char
  * char `a` parses exactly the character `a`

• item :: Parser Char
  * parses any character

• similar to above:    digit letter alphanum lower
                        upper string

• many :: Parser a -> Parser [a]
  * parses 0 or more instances of a and collects them into a list

• many1 :: Parser a -> Parser [a]
  * same as many, but

• (+++) choice:
  * parse first argument if possible, else parse second argument
  * first successfully parsed argument is returned

• nat :: Parser Int
  * parse natural number (positive integer)

• int :: Parser Int
  *
  (++) :: Parser a -> Parser a -> Parser a
  p +++ q = P (\inp -> case parse p inp of
                    [] -> parse q inp
                    [(v,out)] -> [(v,out)])

• ((>=>)) sequential composition
  * a >=> b unboxes monad a into an output a0 and then unboxes
  monad b with input a0
  type Parser a = String -> [(a, String)]
  -- implementation for in-class mostly-complete parser
  -- 'monads'
  (>=>) :: Parser a -> (a -> Parser b) -> Parser b
  (>=>) p1 p2 = \inp -> case parse p1 inp of
    [] -> []
    [(v, out)] -> parse (p2 v) out

* usage:

```

```

doubleDigit :: Parser [Char]
doubleDigit =
  digit >=> \a ->
  digit >=> \b ->
  return [a,b]
-- is equivalent to
doubleDigit' :: Parser [Char]
doubleDigit' = do
  a <- digit
  b <- digit
  return [a,b]

```

* (>>) is the same except that it discards the result of the first monad (thus it has signature (>>) :: Parser a -> Parser b -> Parser b)

Parsing Examples:

- bind and lambda method of parsing:
 - * parse a number:
- parse arithmetic expressions using do syntax:

```

expr :: Parser Int
expr = do t <- term
        do {char '+'
            ; e <- expr
            ; return (t + e)
            }
        +++ return t

term :: Parser Int
term = do f <- factor
        do char '*'
            t <- term
            return (f * t)
        +++ return f

factor :: Parser Int
factor = do d <- digit
            return (digitToInt d)
        +++ do char '('
            e <- expr
            char ')'
            return e

```

```

eval :: String -> Int
eval xs = fst (head (parse expr xs))

```

Trees:

- represent either a leaf node or some kind of internal node
- arithmetic tree declaration:

```

data Expr = Val Int
          | Neg Expr
          | Add Expr Expr
          | Mul Expr Expr

```

- how to fold over a tree:

```

-- exprFold valF negF addF mulF
exprFold :: (Int->b) -> (b->b) -> (b->b->b) ->
-- mulF input output
(b->b->b) -> Expr -> b
exprFold valF _ _ (Val i) = valF i
exprFold valF negF addF mulF (Neg e)
  = negF (exprFold valF negF addF mulF e)
exprFold valF negF addF mulF (Add s1 s2)
  = addF (exprFold valF negF addF mulF s1)
        (exprFold valF negF addF mulF s2)
exprFold valF negF addF mulF (Mul s1 s2)
  = mulF (exprFold valF negF addF mulF s1)
        (exprFold valF negF addF mulF s2)

```

- * basically, just collect values into some type b and use supplied functions at each node to fold into single value
- * useful for evaluating simple things like:

```

-- evaluate an expression
evalExpr' = exprFold id (\x -> 0 - x) (+) (*)
id -- integers map to integers
(\x -> 0 - x) -- negation
-- everything else is just simple numeric operators
--
-- count leaves in a tree
countLeaves' = exprFold (\_ -> 1) id (+) (+)
(\_ -> 1) -- leaf integer node is one node

```

```
id -- negation node has only one child, pass on count
(+) -- nodes with two children: add number of
    ↳ leaf grandchildren
```

HW2: Water Gates:

```
waterGate :: Int -> Int
waterGate n =
  length -- number of True's
  $ filter id -- filter just True's
  $ waterGate' n initial -- initial call to helper
  where
    -- start with all gates closed
    initial = replicate n False
    --
    -- flip states
    waterGate' 1 state = map not state
    -- base case: flip every state
    waterGate' n state = flip n $ waterGate' (n-1) state
    -- otherwise, first get the state for (n-1) and then
    -- flip every nth state
    --
    -- flip every nth gate
    flip :: Int -> [Bool] -> [Bool]
    flip 1 xs = map not xs -- flip every gate
    -- flip only gates which index are multiples of n
    flip nth xs = [ if (i `mod` nth == 0) then not x else
    ↳ x
                    -- zip each state with it's index
                    | (x,i) <- (zip xs [1..]) ]
```

HW2: Goldbach's Other Conjecture:

```
-- check if a number is prime
primeTest :: Integer -> Bool
primeTest 1 = False
primeTest t = and [ (gcd t i) == 1 | i <- [2..t-1]]

-- all numbers less than n that are double a square
twiceSquares :: Integer -> [Integer]
twiceSquares n = takeWhile (<n) [ 2 *x^2 | x <- [1..]]

-- list of odd numbers
oddList = map (\x -> 2*x + 1) [0..]
-- all odd numbers that are composite (not prime)
allOddComp = [ o | o <- (drop 1 oddList)
                , not (primeTest o) ]

-- if a number satisfies conditions for conjecture
-- method: for enough square numbrers, check if n-(that
↳ number) is prime
satsConds n = or [ primeTest k |
                   k <- map (\x->(n-x)) (twiceSquares n)
                   ↳ ]

-- find the first number
goldbachNum = head [ x | x <- allOddComp
                        , not (satsConds x) ]
```

HW4: Sets:

```
type Set a = [a]

a = mkSet [1,2,3,4,5]
b = mkSet [1,2,3]

addToSet :: Eq a => Set a -> a -> Set a
addToSet s a | a `elem` s = s
             | otherwise = a : s

mkSet :: Eq a => [a] -> Set a
mkSet lst = foldl addToSet [] lst

isInSet :: Eq a => Set a -> a -> Bool
isInSet [] _ = False
isInSet [a] b = a == b
isInSet (x:xs) b | x == b = True
                  | otherwise = isInSet xs b

subset :: Eq a => Set a -> Set a -> Bool
subset sub super = and [ isInSet super x | x <- sub ]

setEqual :: Eq a => Set a -> Set a -> Bool
setEqual a b = subset a b && subset b a
```

```
-- instance (Eq a) => Eq (Set a) where
-- a == b = subset a b && subset b a
```

```
setProd :: Set a -> Set a -> [(a,a)]
setProd a b = [ (ai,bj) | ai <- a
                        , bj <- b
                        ]
```

Prev Exam: Run Length Encoding:

```
import Parsing
import Data.Char

q4 = do
  d <- sat isUpper
  e <- char (toLower d)
  f <- many item
  return [d,e]

ones = (map (\_ -> 1) [1..])

myRLE [] = []
myRLE ls = myhelper (zip ones ls)

myhelper [(n,c)] = [(n,c)]
myhelper ((n,c):(m,d):rest)
  | (d == c) = myhelper ((n+m),c):rest
  | otherwise = (n,c):myhelper ((m,d):rest)
```

Rock Paper Scissors:

```
data RPS = Rock | Paper | Scissors
  deriving (Eq, Show)

rps :: RPS -> RPS -> Int
rps a b | a == b = 0
rps Rock Scissors = 1
rps Paper Rock = 1
rps Scissors Paper = 1
rps _ _ = 2

rps2 :: RPS -> RPS -> Int
rps2 a b =
  if a == b then 0 else case (a,b) of
    (Rock, Scissors) -> 1
    (Paper, Rock) -> 1
    (Scissors, Paper) -> 1
    _ -> 2
```

99 problems:

```
-- 9. pack consecutive duplicates into sublists
pack (x:xs) = let (first,rest) = span (==x) xs
               in (x:first) : pack rest

pack [] = []
-- example:
pack [1,2,3,2,2,3] == [[1,1],[2],[3],[2,2],[3]]
```

Java:

- **Class Invariant:** A logical condition that ensures that an object of a class is in a well-defined state.
 - * public methods assume that invariant holds before it's called, and makes sure to preserve the invariant property
- garbage collection deals only with memory. You must manage other resources manually, such as concurrent locks, OS file handles, etc...
- anonymous classes are a thing
- inner class: class declared inside another class implicitly holds a reference to it's outer class.
 - This means instances of the inner class can use non-static fields/methods of the outer class.
 - This is especially handy for callbacks e.g. on android
- abstract class vs interface:**
 - Interface: all fields are **public static final**, all methods are **public**
 - **abstract class** can extend exactly one parent class and implement any number of interfaces

- interface can extend (not implement) any number of interfaces
- abstract class can have constructor that initializes values and whatnot, but interface cannot
 - * you can't instantiate an abstract class directly, you can only call it's constructor from inside the constructor of a child class. Could still be useful though.
- interfaces and abstract classes can never be instantiated
 - * references of interface type refer to an instance of a class that implements that interface
 - * references of abstract class type refer to an instance of a subclass of that type

Inheritance and Virtual Methods:

- Java classes can inherit from one class only
 - * inheritance: `class ChildClass extends ParentClass`
 - * child class gets access to public fields/methods (of course) and protected fields/methods
 - * child class does not get access to private fields/methods
 - * TODO abstract classes/methods
- interfaces: `class MultiPurpose implements Interface1, IFace2...`
 - * basically an end around lack of multiple inheritance
 - * interfaces cannot be instantiated, but you can have a reference of interface type. In that case, the object it points to is a real concrete class, but all that you know about it is that it implements the specified interface
- virtual dispatch:
 - * all public non-static class methods are virtual. This means that if a subclass overrides a parent class method, the decision of which method implementation to use is made at runtime, depending on which type of object the reference actually refers to.
 - * all interface methods are by definition virtual
 - * private methods are not virtual, static methods are not virtual
- notable interfaces:
 - * `Iterable<T>`: contains `Iterator<T>` `iterator()` method, to iterate over container
 - * `Iterator<E>`: encapsulates an iteration over a container. Unlike C++ iterators, this iterator knows when it's reached the end, instead of relying on comparison to a one-past-end iterator
 - `boolean hasNext()`: ask if we're at the end
 - `E next()`: retrieve next element, and advance iterator
 - * `Runnable`: single `void run()` method. This is the interface that threads use

Generics:

- Subtype Operator: `<`:
 - * if S `implements` T then `S <: T`
 - * if `S <: T` then you can freely use an object of type S where type T was required, and it will be type-safe (S can be safely used in that context instead of T)
 - `<T extends Number>` → `T <: Number`
 - * `<T super Number>` → `T >: Number`
 - `T t = new S();` → `S <: T`
- type bound: `class SortedList<T extends Comparable & Serializable>`
 - * you use `extends` for constraints that are classes or interfaces
 - * you can bound with `extends` or `super`
 - * `<T extends Type>`: Type is an inclusive upper bound on T
 - * `<T super Type>`: Type is an inclusive lower bound on T
- wildcards: `static void printAll(List<?> lst)` use `?` for when you want to accept an object of any type
 - * you can also specify `<? extends ClassOrInterface...>`

```

• PECS: Producer Extends, Consumer Super
  * to generically assign a T to something, use <? super T>
  * to generically read a T from something, use <? extends T>

public class CollectionsPECS {
    public static <T> void copy(List<? super T> dest,
                               List<? extends T> src) {
        for (int i=0; i<src.size(); i++)
            dest.set(i,src.get(i));
    }
}

import java.lang.*;
class GenericWildcards {
    // T is the binding of the generic parameter
    private static class GenericBox<T> {
        // it is optional, but we need it if we want to do
        // → things with that type
        public T t;
        public GenericBox(T t) { this.t = t; }
    }
    private static class NumberBox<T extends Number> {
        public T t;
        public NumberBox(T t) { this.t = t; }
    }
    public static void printBox(GenericBox<?> b) {
        // here we use the ? wildcard with no type binding
        // → because we don't need to do things with that
        // → type specifically
        System.out.println(b.t);
    }
    // method generic goes before return type
    public static <T> void
    // → printWithParameter(GenericBox<T> b) {
    //     System.out.println(b.t);
    // }
    public static void main(String[] args) {
        // this is using raw types, generally considered
        // → bad
        GenericBox rawBox = new GenericBox("asdf1"); //
        // → compiler warnings
        // this cast is ok because raw types hold
        // → java.lang.Object
        Object o1 = rawBox.t;
        // this causes no warnings for same reason as
        // → assignment above doesn't
        printBox(rawBox);

        // this is just using an unknown type, java says
        // → it's fine
        GenericBox<?> unknownBox = new
        // → GenericBox<>("asdf2");
        // this is also OK because <?> explicitly makes the
        // → generic parameter as java.lang.Object
        Object o2 = unknownBox.t;
        printBox(unknownBox);

        GenericBox<String> stringBox = new
        // → GenericBox<>("asdf3");
        // the type parameter above allows Java to infer
        // → that this cast is safe
        String s = stringBox.t;
        System.out.println(s);
        // must specify type between class access and
        // → method name (works the same for instance
        // → methods too)

        // → GenericWildcards.<String>printWithParameter(stringBox)

        // correct stuff works like expected
        NumberBox<Integer> nb1 = new NumberBox<>(5);
        System.out.println(nb1.t + 1);
        //
        // this will fail to even allow NumberBox<String>
        // → because that type doesn't work
        // NumberBox<String> sb1 = new NumberBox<>("asdf");
        //
        // this will fail because the inferred type of
        // → NumberBox<>("asdf") is NumberBox<String>, which
        // → isn't allowed
        // NumberBox<?> sb1 = new NumberBox<>("asdf");
    }
}

```


Threading:

```
import java.util.concurrent.locks.ReentrantLock;
import java.util.concurrent.locks.Condition;
```

ReentrantLock:

- ReentrantLock: basically a mutex
- ReentrantLock.lock(): acquire the lock (blocking)
 - * does **not** throw InterruptedException
- ReentrantLock.unlock(): release the lock
 - * does **not** throw InterruptedException
 - * you should always wrap your locking code in a **try{} block** (including the call to lock() itself) and put the call to unlock() in a **finally{} block**.

This way, unlock() gets called no matter any exception

Condition:

- created from a lock, allows one thread to send a message to another thread
 - * create from lock instance using lock.newCondition()
- await(): release this lock and wait for the condition to be signaled.
 - When the signal happens, await() will automatically re-acquire the lock before returning (this means you will still have to unlock manually)
 - * you can only await() when you are holding the lock, and when it returns, you still have the lock, so it acts like you never unlocked it
 - * **does** throw InterruptedException
- signal(): wake up a single thread that is waiting on the condition
 - * must be holding lock to signal it's condition
 - * must manually release lock before other thread will return from await() (because the other thread must also acquire the lock)
 - * does **not** throw InterruptedException
- signalAll(): similar to signal() except that every thread is woken up
 - * still only one thread will be able to use the lock-protected resource at a time, because locks
 - * does **not** throw InterruptedException

Threads:

- **static void** Thread.sleep(long ms): sleep for ms
 - * throws InterruptedException if the thread was interrupted before time elapsed
- make new thread with **new** Thread(Runnable r)
 - * start that thread with thread.start()

synchronized:

```
// inside a method
synchronized(some_object /*may be this*/) {
    // synchronized code here
}

// synchronized getInstance method for singleton
public foo synchronized getInstance() {
    if (inst == null) { inst = new Foo(); }
    return foo;
}
```

- methods marked **synchronized** are implicitly locked to ensure that only one synchronized method is ever running on a given object at a time
- synchronized statement: synchronize on a specific object manually
- works as a good synchronization mechanism as long as the resource doesn't need to be used directly by multiple objects
- does not allow for conditions

```
import java.lang.*;
import java.util.concurrent.locks.ReentrantLock;
import java.util.concurrent.locks.Condition;
public class Main2 {
```

```
public static class Counter {
    public int count = 0;
    public ReentrantLock lock;
    public Condition updated;
    public Counter() {
        this.lock = new ReentrantLock();
        this.updated = lock.newCondition();
    }
}

public static class CounterThread implements Runnable {
    private Counter counter;
    public CounterThread(Counter c) {counter = c;}
    @Override
    public void run() {
        while (true) {
            try {
                counter.lock.lock();
                counter.count += 1;
                System.out.println(counter.count);
                counter.updated.signalAll();
            }
            // lock() does not throw InterruptedException
            // catch (InterruptedException e) {}
            finally {counter.lock.unlock();}

            try {
                Thread.sleep(1000);
            } catch (InterruptedException e) {}
        }
    }
}
```

```
public static class IntervalPrinter implements Runnable {
    private Counter counter;
    private int mod;
    private String message;
    public IntervalPrinter(Counter c, int mod, String msg) {
        counter = c;
        this.mod = mod;
        message = msg;
    }
    @Override
    public void run() {
        while (true) {
            int val = 0;
            try {
                counter.lock.lock();
                counter.updated.await();
                val = counter.count;
            }
            catch (InterruptedException e) {}
            finally {counter.lock.unlock();}

            if (val % mod == 0) {
                System.out.println(message);
            }
        }
    }
}
```

```
public static void main(String []args) {
    Counter c = new Counter();
    new Thread(new IntervalPrinter(c,3,"fizz")).start();
    new Thread(new IntervalPrinter(c,5,"buzz")).start();
    new Thread(new CounterThread(c)).start();
}
```

Reflection:

- **instanceof** operator: check if an object is an instance of a class or a subclass
- if (obj instanceof String) {
 // cast is safe because we checked and obj
 String s = (String) obj;
}

java.lang.Class<T>:

- allows you to reflect on class T

- to get:
 - * `Class<?> c = SomeClassName.class;`
 - * `Class<?> c = someObjectInstance.getClass();`
 - * `Class<?> c = Class.forName("SomeClassName");`
 - throws `ClassNotFoundException`
- `toString()` returns class declaration (more or less)
- `getSimpleName()` returns just the name part of it
 - * `Main.class.getSimpleName() → "Main"`
- `Class<? super T> getSuperclass()`
- `Class<?>[] getInterfaces()`
 - * on class object: get interfaces implemented by this class
 - * on interface object: get interfaces extended by this interface
- methods matching `getDeclaredXXX()`
 - * can see just things declared in the class itself, not from super classes
 - * can see anything whether public/private/protected/etc...
- methods matching `getXXX()`
 - * operate on whatever the class looks like from an outside observer: only public fields/methods, and from class or super-class/interface
- `Method[] getMethods()`
 - * all public methods, including those inherited from super-classes and implemented in interfaces
- `Method getMethod(String name, Class<?>... pt)`
 - * looks for fields in superclasses, then superinterfaces too
 - * throws `NoSuchMethodException` if method not found
 - * public methods only
- `Method[] getDeclaredMethods()`
 - * excludes inherited methods; includes any that are declared in class regardless of public, private, static, etc...
- `Field[] getFields()`
 - * returns public fields only!
- `Field getField(String name)`
 - * looks for fields in superinterfaces, then superclasses too
 - * throws `NoSuchFieldException` if field not found
 - * only finds public fields
- `Constructor<T> getConstructor(Class<?>... pt)`
 - * get a constructor for T that matches parameter types `Class<?>... pt`
 - * throws `NoSuchMethodException` if there is no constructor matching those parameter types
- `T newInstance()`
 - * create a new instance of T using the default constructor
- TODO use `Constructor` class to create class using non-default constructor

java.lang.reflect.Method:

- `String toString()` → method prototype as string
 - * includes modifiers, method name, parameters, etc...
- `String getName()` → name of method as string
- `int getModifiers()` → `int` representing modifiers
 - * use `java.lang.reflect.Modifier` static methods to check:
 - `Modifier.isStatic(m.getModifiers())`
- `Class<?>[] getParameterTypes()` → types of parameters of method
 - * if no parameters, returns empty array
 - * does not include implicit `this` parameter for instance methods
- `Type[] getGenericParameterTypes()`: same, but returns a `Type` instance that accurately represents the generic info from the actual source
- `Class<?> getReturnType()`: get the return type
 - * if it's void, it returns a void type
- `Object invoke(Object obj, Object... args)`
 - * invoke a method on an object. Subject to virtual method

lookup

- * if the method is static, `obj` may be null
- * if the method returns a primitive type, it is wrapped; if void, returns null
- * throws `IllegalAccessException` if you can't run that method because it's private or something
- * if target method throws, it throws `InvocationTargetException` wrapping whatever was thrown

java.lang.reflect.Field:

- `TYPE getTYPE(Object obj)`: bunch of methods for getting fields of primitive types
 - * throws `IllegalArgumentException` if type can't be converted (widening conversions only are allowed)
 - * throws `IllegalArgumentException` also if `obj` isn't of the right type
- `Object get(Object obj)`: get an object field
 - * if the field is a primitive type, it is wrapped and then returned
 - * only throws `IllegalArgumentException` if `obj` isn't of the right type
- `String getName()`
- also various `setTYPE(Object obj, TYPE value)` and `set(Object obj, Object value)` equivalent to get methods