Proof Complexity and Solving LAB

Clause Learning

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Goals

- Implemenatation of SAT solving algorithms
 - (a) 2-SAT (polynomial time)
 - (b) DPLL
 - (c) CDCL
 - watched literals
 - clause learning
 - decision heuristics
 - restart strategy
 - (d) QBF expansion..
- Practical programming experience
 - use your favourite language (Python, C, C++, Java, ..)
 - recommended: Python

CDCL Pseudocode

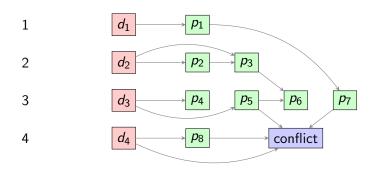
```
function CDCL-solver(\Phi)
                                                               #assuming \Phi is preprocessed
decision-level \leftarrow 0
while there are unassigned variables
   decision-level++
   decide()
                                                                   #adds assignment to trail
   C_{\text{conflict}} \leftarrow \text{propagate}()
                                                            #returns conflict clause or null
   while C<sub>conflict</sub> is not null
      if decision-level = 0 return UNSAT
      C_{\text{learned}} \leftarrow \text{analyse-conflict}(C_{\text{conflict}})
      if C_{\text{conflict}} is unit
         backtrack(0)
         assign unit literal
      else
         backtrack(asserting-level(C_{learned}))
                                                                       #changes trail and DL
         \Phi \leftarrow \Phi \wedge C_{\text{learned}}
      C_{\text{conflict}} \leftarrow \text{propagate()}
   apply-restart-policy()
return SAT
```

Resolution

$$(C \lor x), (D \lor \neg x) \vdash (C \lor D)$$

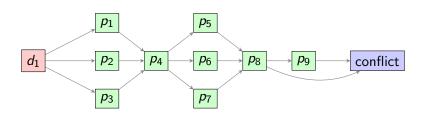
- $(C \lor D)$ is the resolvent
- x is the pivot
- logically correct rule: premises imply resolvent
- sound and complete proof system for UNSAT

Clause Learning - Cutting the Implication Graph



Unique Implication Points

- A unique implication point (UIP) is:
 - a node at the highest decision level
 - every path from highest decision to conflict passes through it



Common Implementation

- advantage: clauses learned from UIPs are always asserting
- asserting means 'becomes unit at a previous decision level'
- the asserting level is second highest decision level in learned clause
- hence: backtrack to asserting level ('backjumping')
- remember: we don't learn unit clauses we make assignments instead
- easy implementation:
- resolve conflict clause with reason clauses until there is exactly one variable at highest decision level
- This is the 1UIP learning scheme

Thoughts

- Don't worry too much about why 1UIP schemes work well
- My opinion: the community has not agreed on a clear reason
- Theoreticians: SAT solver are automated decision procedures
- Practitioners: SAT solvers perform local search

Clause Learning Task

- implement a 1UIP clause learning scheme in your CDCL solver
- check correctness

compare the solving time to DPLL-style learning