

Proof Complexity and Solving LAB

Winter Semester 2020

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<https://github.com/JoshuaBlinkhorn/SAT-LAB>

Goals

- Implementation of SAT solving algorithms
 - (a) 2-SAT (polynomial time)
 - (b) DPLL (decision tree)
 - (c) CDCL
 - clause learning
 - watched literals
 - decision heuristics
 - restart strategy
 - (d) QBF expansion..
- Practical programming experience
 - use your favourite language (Python, C, C++, Java, ..)
 - recommended: **Python**

Conjunctive Normal Form (CNF)

- a literal is a variable or its negation
- a clause is a disjunction of literals
- a CNF is a conjunction of clauses

$$(x_1 \vee x_2 \vee x_3) \wedge (\neg x_1 \vee \neg x_2) \wedge (\neg x_3) \wedge (\neg x_1 \vee x_3)$$

The DIMACS CNF Encoding

- machine readable encoding
- variables are natural numbers: $x_1 \mapsto 1$, $x_2 \mapsto 2$ etc.
- negation represented by minus: $\neg x_1 \mapsto -1$, $\neg x_2 \mapsto -2$ etc.

$$(x_1 \vee x_2 \vee x_3) \wedge (\neg x_1 \vee \neg x_2) \wedge (\neg x_3) \wedge (\neg x_1 \vee x_3)$$

```
p cnf 3 4
1 2 3 0
-1 -2 0
-3 0
-1 3 0
```

The DIMACS CNF Encoding

$$(x_1 \vee x_2 \vee x_3) \wedge (\neg x_1 \vee \neg x_2) \wedge (\neg x_3) \wedge (\neg x_1 \vee x_3)$$

```
p cnf 3 4
1 2 3 0
-1 -2 0
-3 0
-1 3 0
```

- begins with the problem line

p cnf *[number of variables]* *[number of clauses]*

- other lines represent clauses as zero-terminated literal sets

First Task – Random CNF Generator

- A random (n, c, k) -CNF is:
 - a random CNF with n variables and c clauses
 - all clauses of width k
 - no tautological clauses
 - no repeated clauses
 - no repeated literals in clauses
 - clauses selected uniformly at random
- Implement a tool:

`random-cnf t n c k`

which outputs t random (n, c, k) -CNF formulas in DIMACS

Second Task – Testing

- Install a SAT solver

<http://fmv.jku.at/lingeling/>

- Test your random CNFs with the solver
- Can you find out anything interesting by varying the parameters n , c and k ?