



# Data-Center Network Management

Network Management

Prof. Dr. Panagiotis Papadimitriou



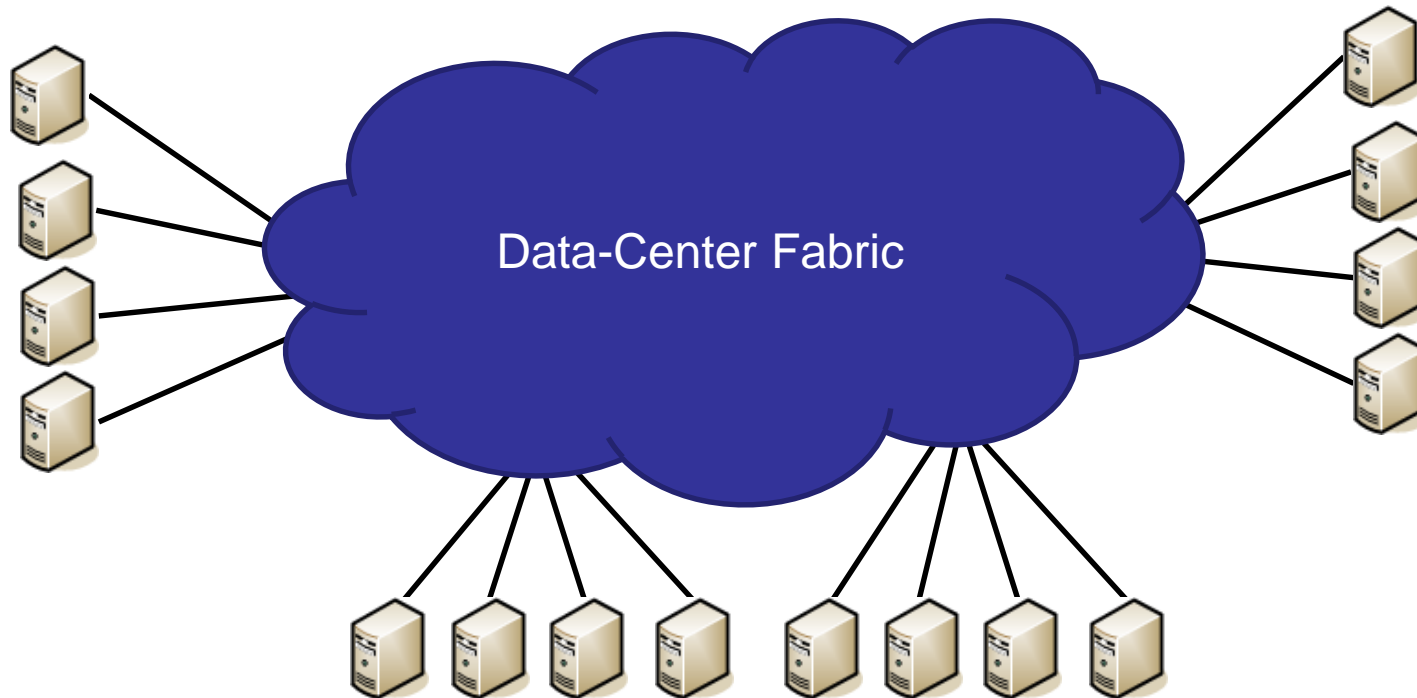
# Data-Center Networks

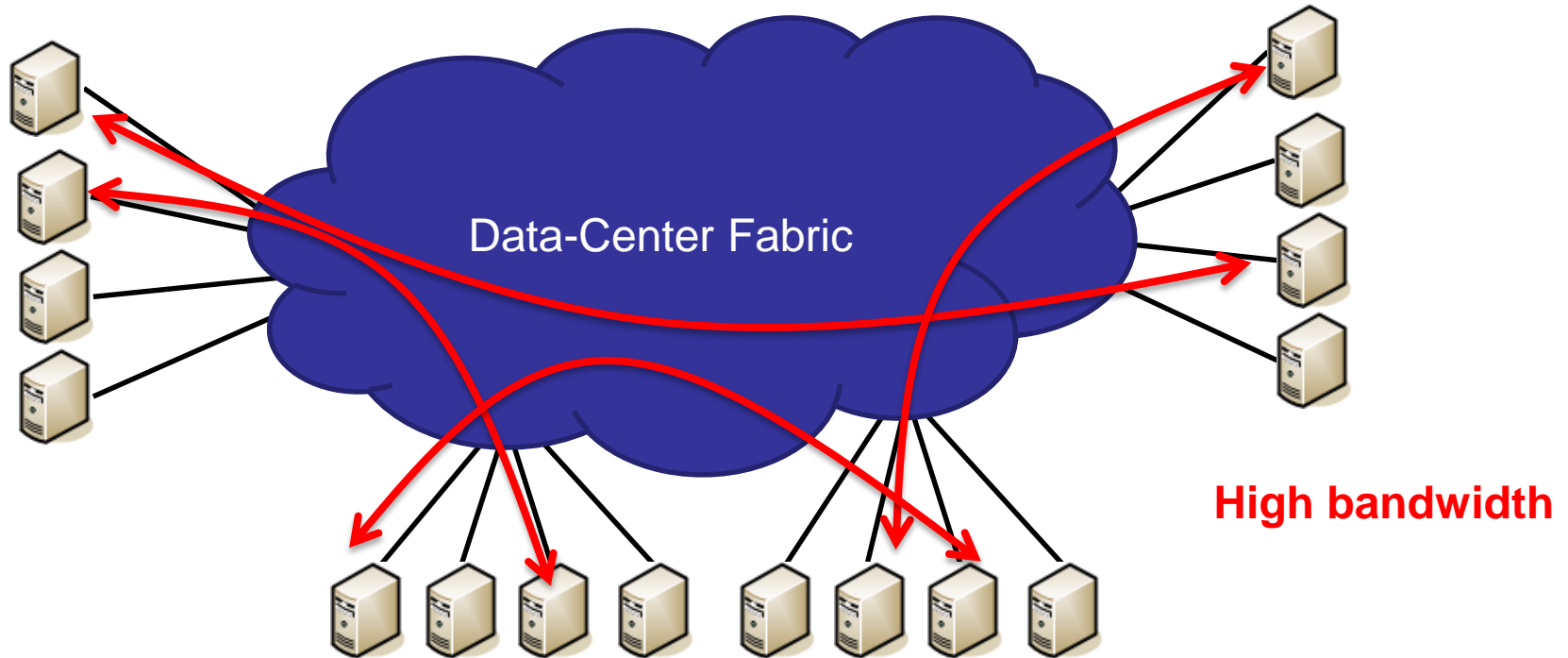


- Major players build and operate data centers:
  - Amazon, Google, HP, Microsoft, Facebook, etc.
- Features:
  - Massive scale:
    - Usually tens of thousands servers (up to hundreds of thousands servers)
  - Commoditization:
    - Wide use of commodity (inexpensive) hardware (i.e., servers and switches)
  - Server virtualization:
    - Widespread adoption of server virtualization to maximize resource utilization
    - A large number of virtual machines may be hosted on a single server using technologies such as VMWare, Xen, etc.



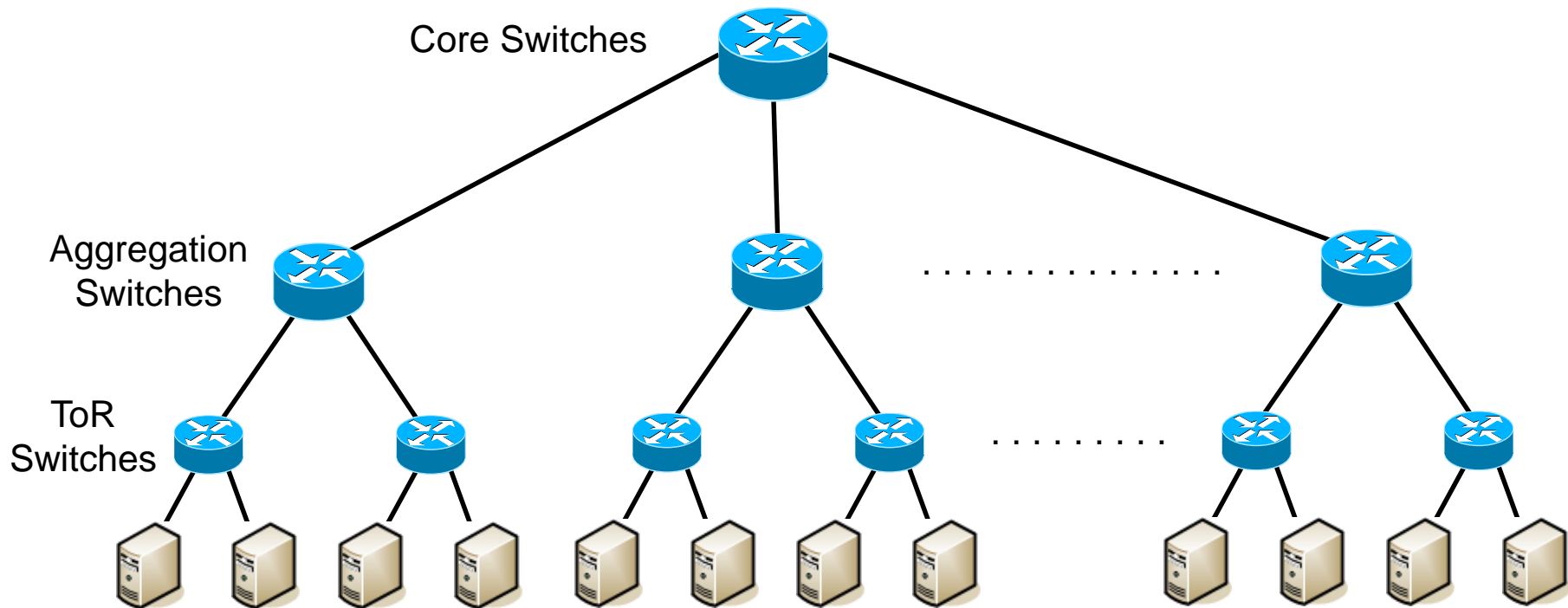
- Main Goals:
  - Performance:
    - Fast execution of applications (e.g., Map-Reduce jobs)
  - Energy efficiency:
    - Increased consolidation (e.g., shut down redundant servers) to achieve energy savings
  - Minimal configuration overhead:
    - Host/service discovery
    - Mobility support (e.g., VM migration)
  - Scalability:
    - Scaling data-center network to tens or hundreds of servers





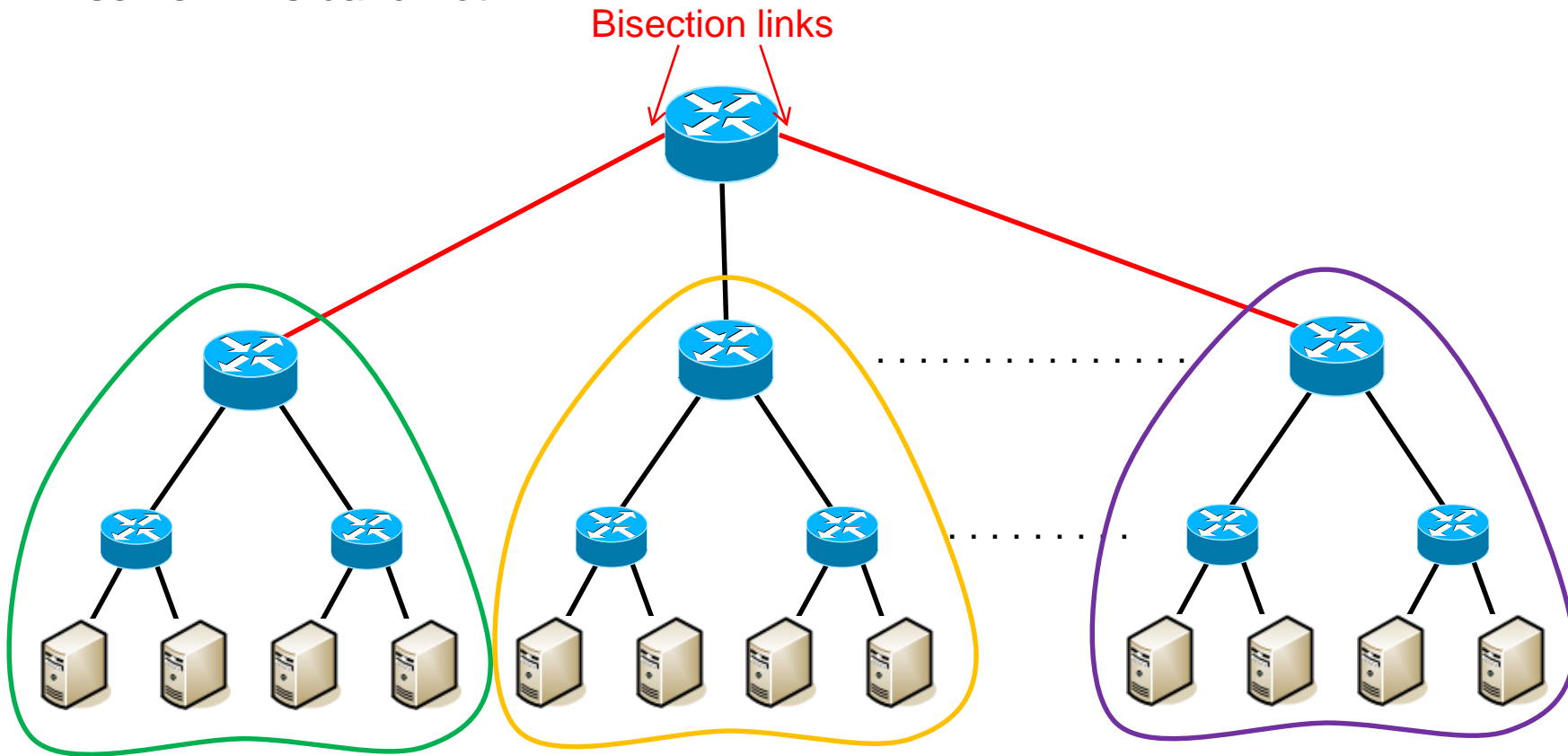


- Tree-based topologies:
  - Suitable either for flows that enter/leave the DC or for data transfers within the same rack
  - Poor performance with data transfers among different racks
  - Not fault-tolerant





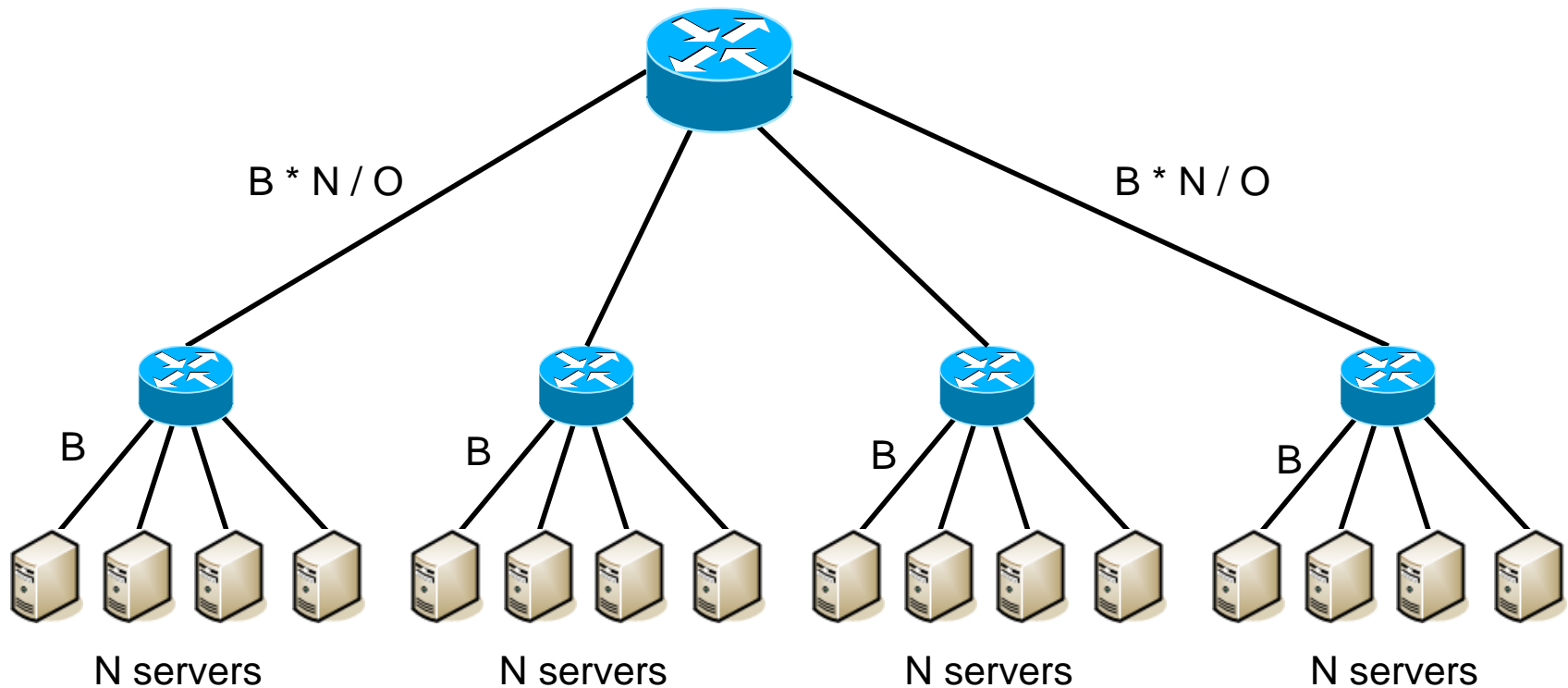
- Bisection bandwidth is the sum of the core link bandwidths
- Full bisection bandwidth:
  - The total bandwidth of the core links should be equal to the sum of all server links bandwidth





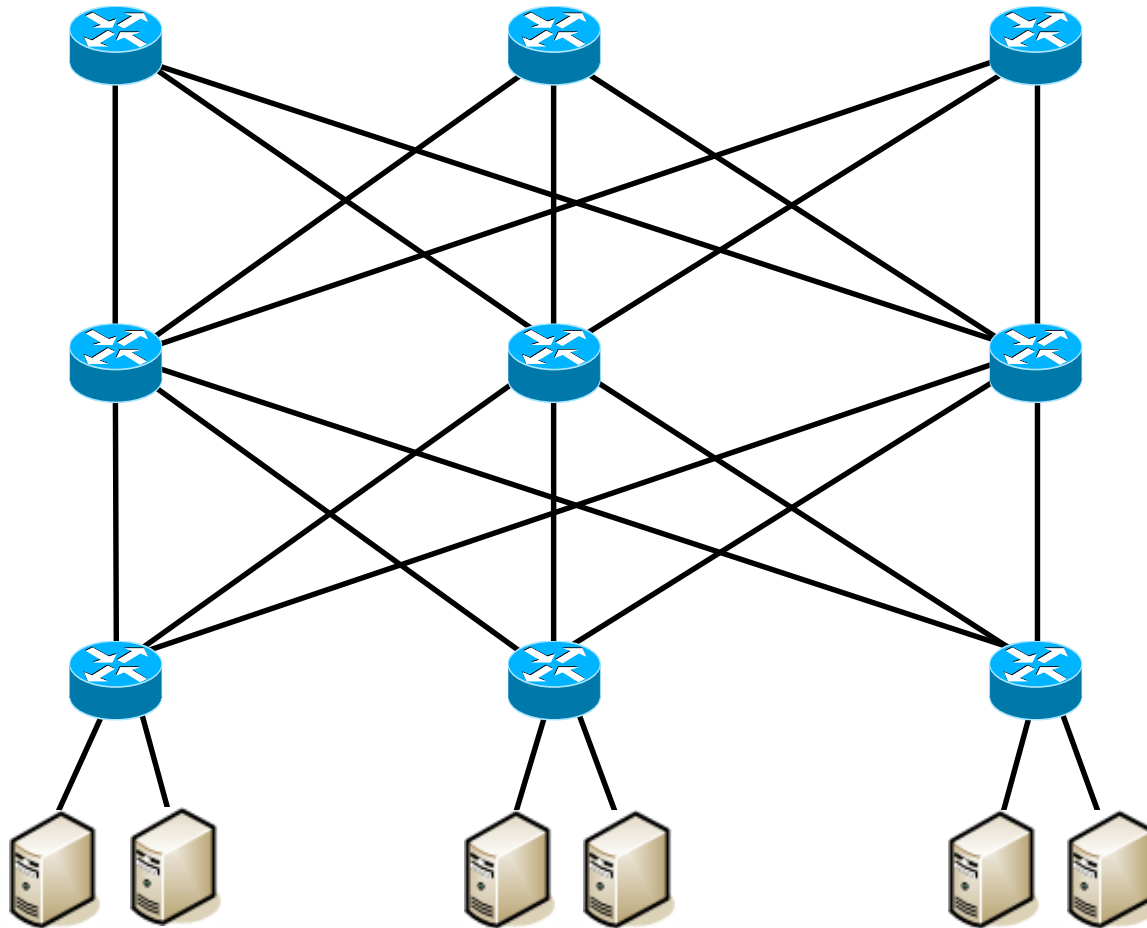


- Oversubscription factor (O):
  - is adjusted taking into account the locality of data transfers and the traffic rate/pattern



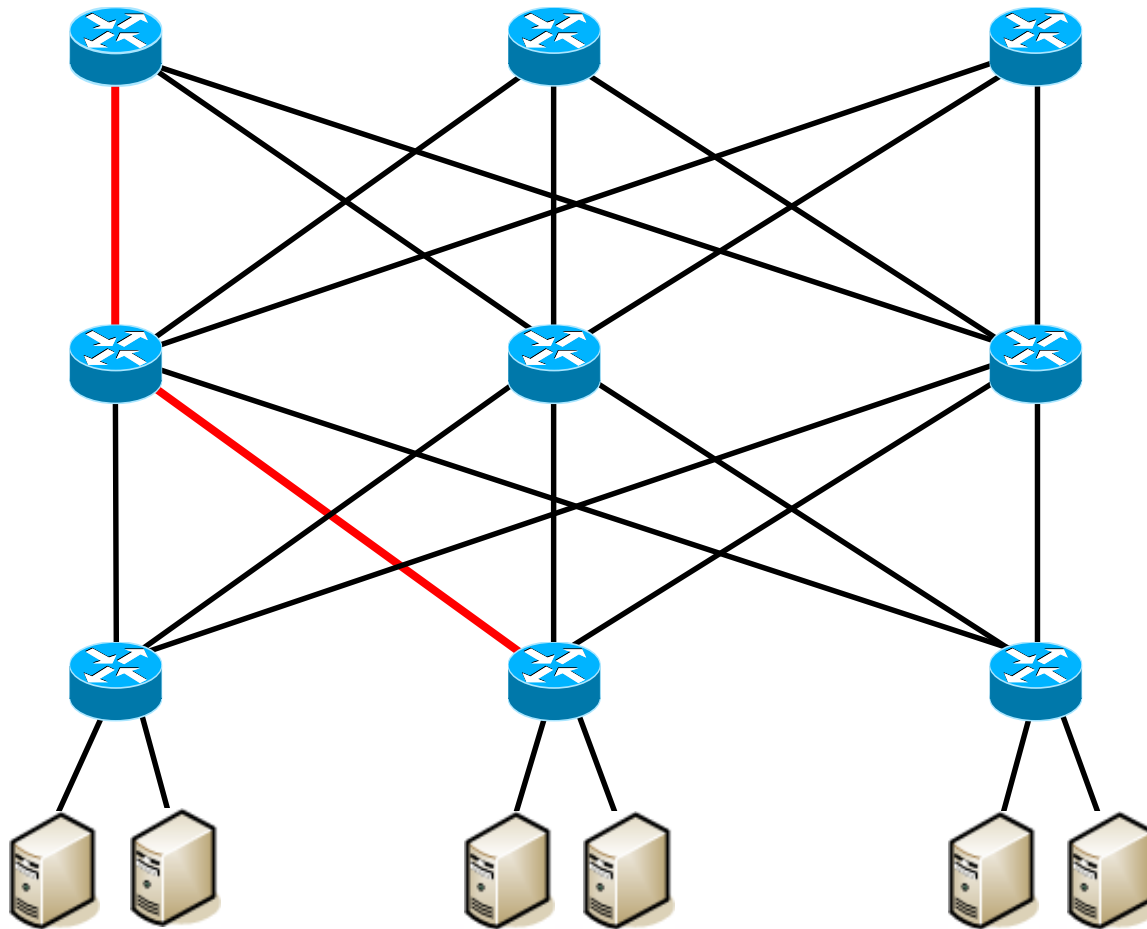


- Multiple Equal Cost Paths (ECMP)
- Full bisection bandwidth



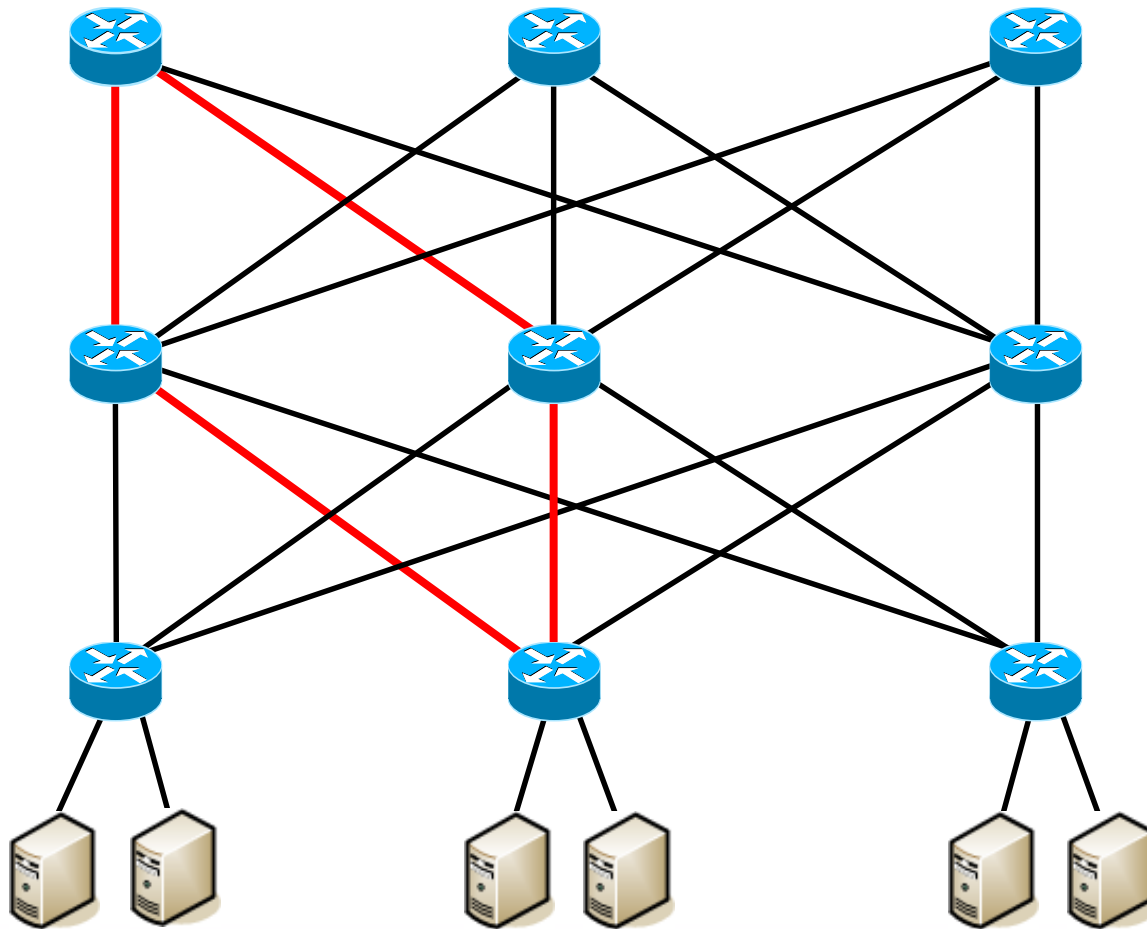


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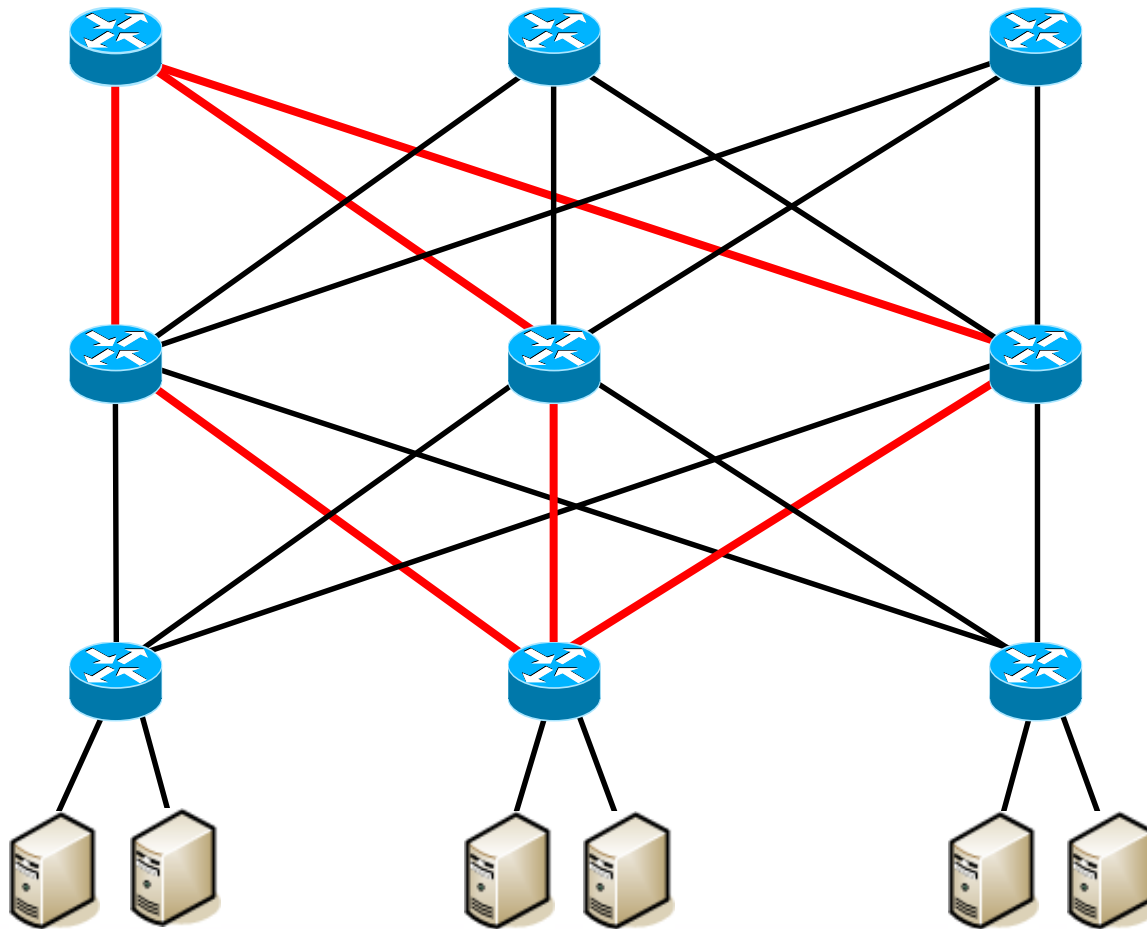


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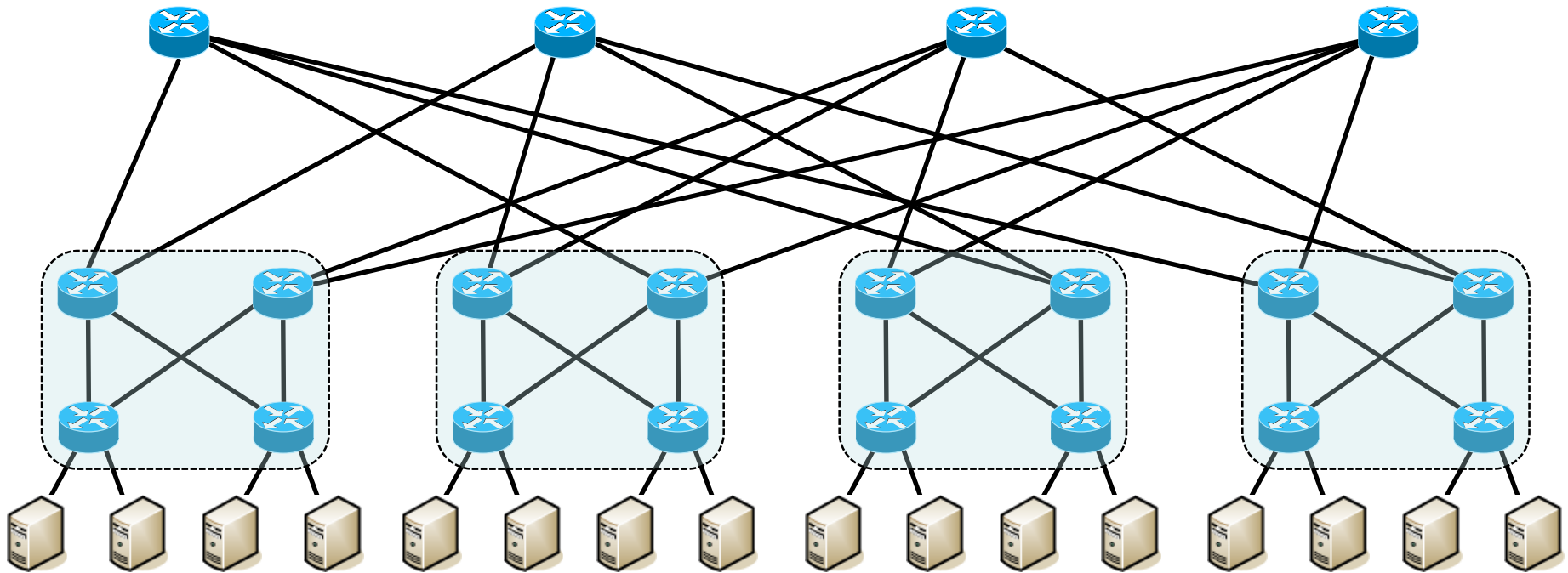


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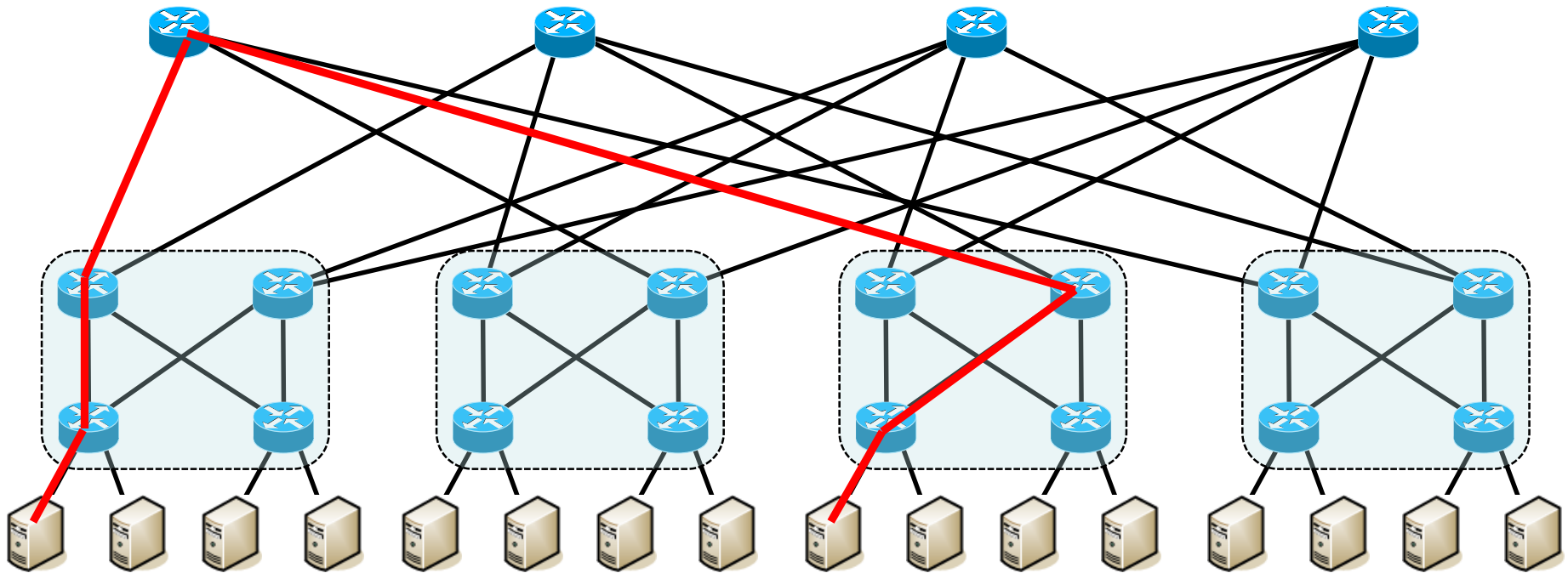


- Special case of a Clos topology:
  - $(k/2)^2$   $k$ -port core switches,  $k$  pods each one with two layers of  $k/2$  switches and  $k^2/4$  servers,  $k^3/4$  servers in total



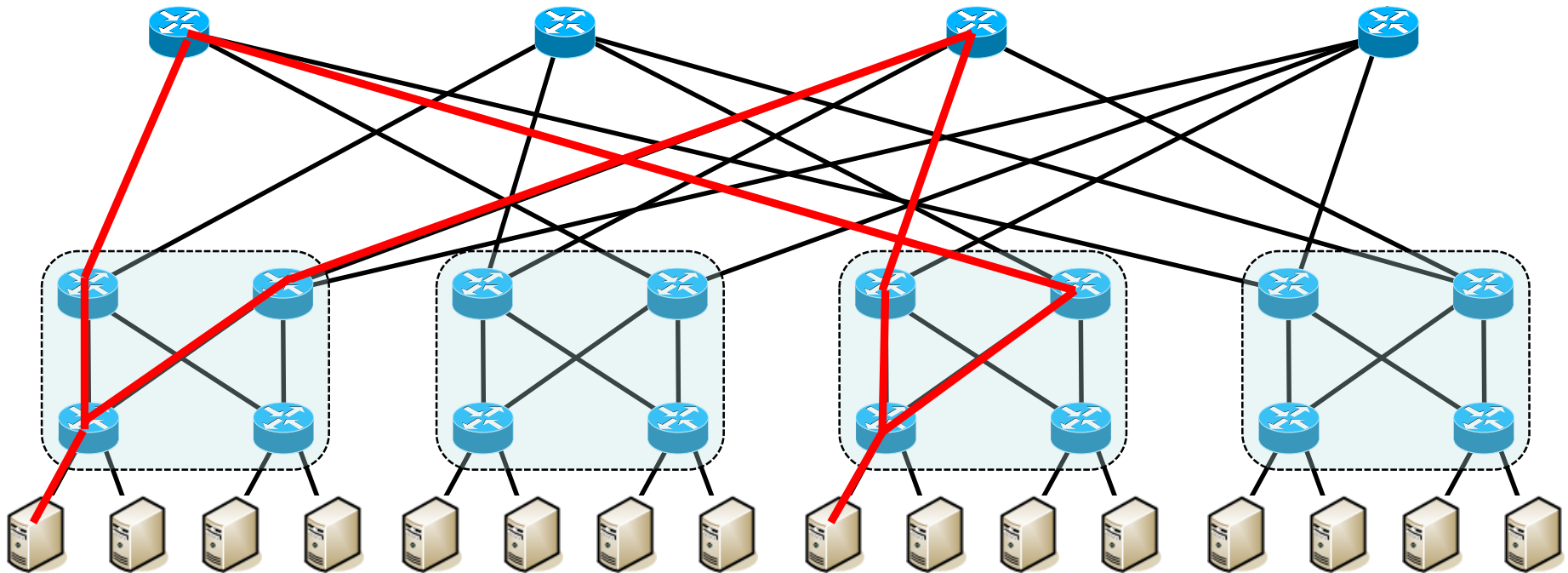


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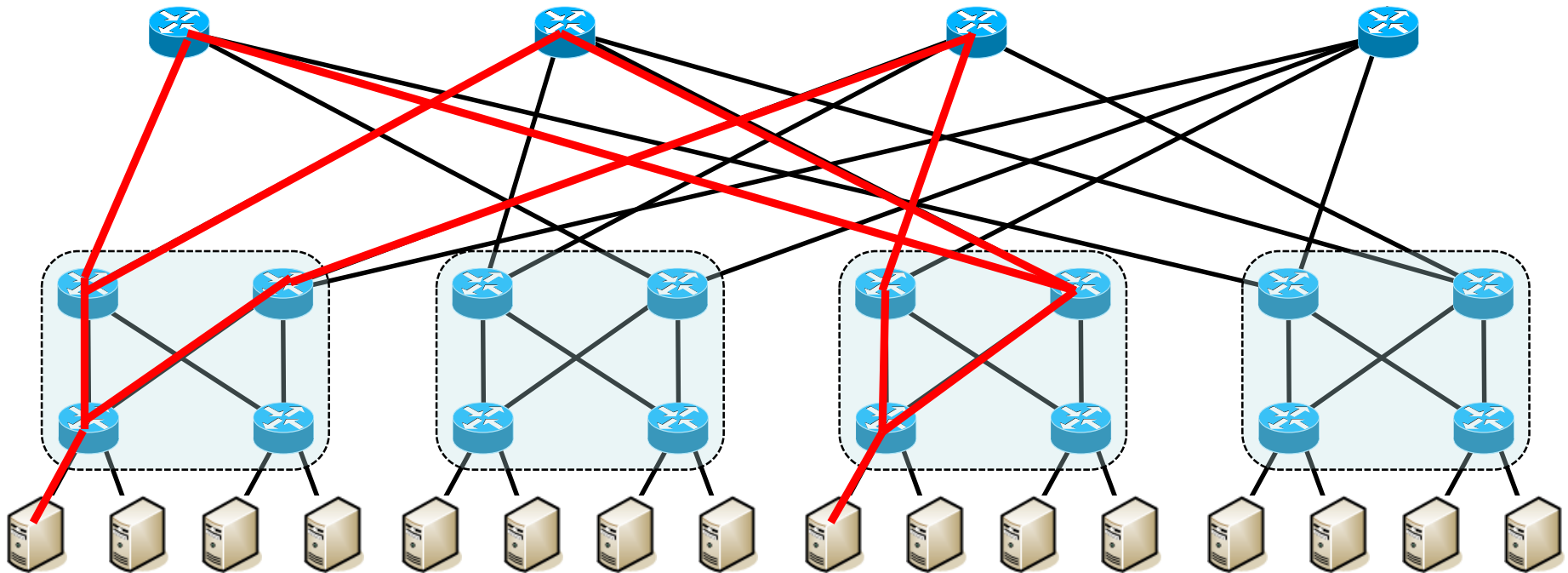
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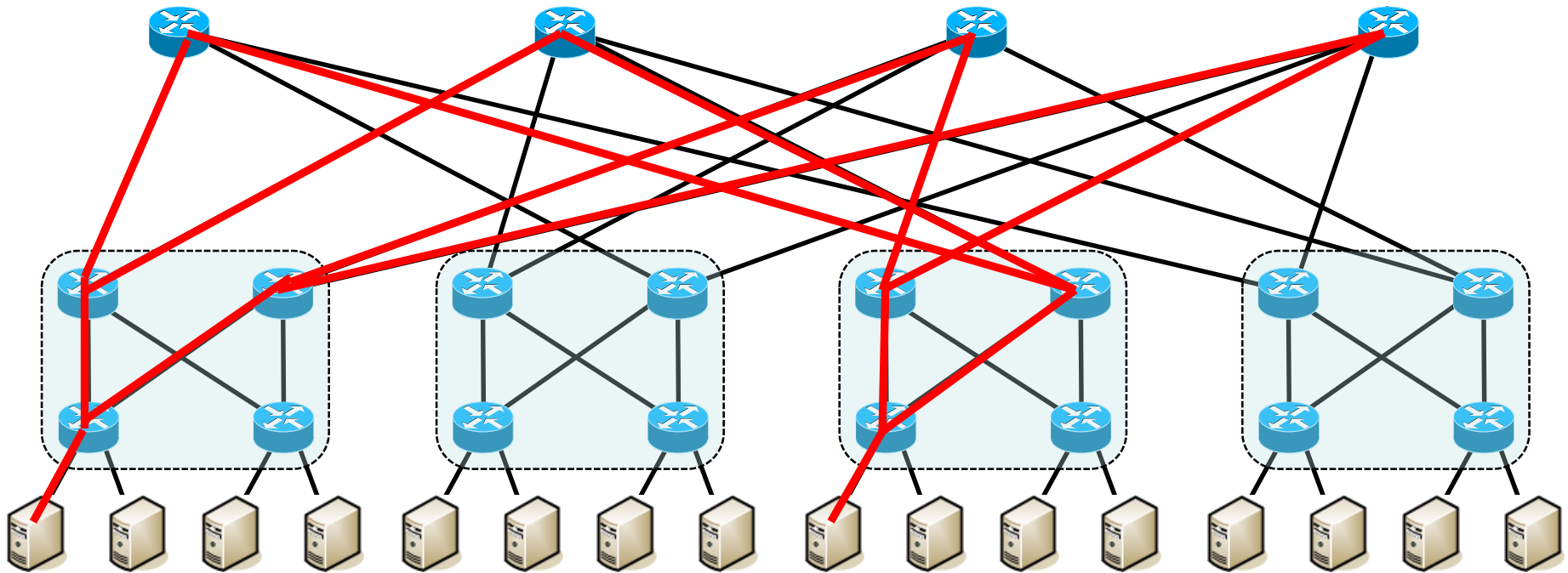


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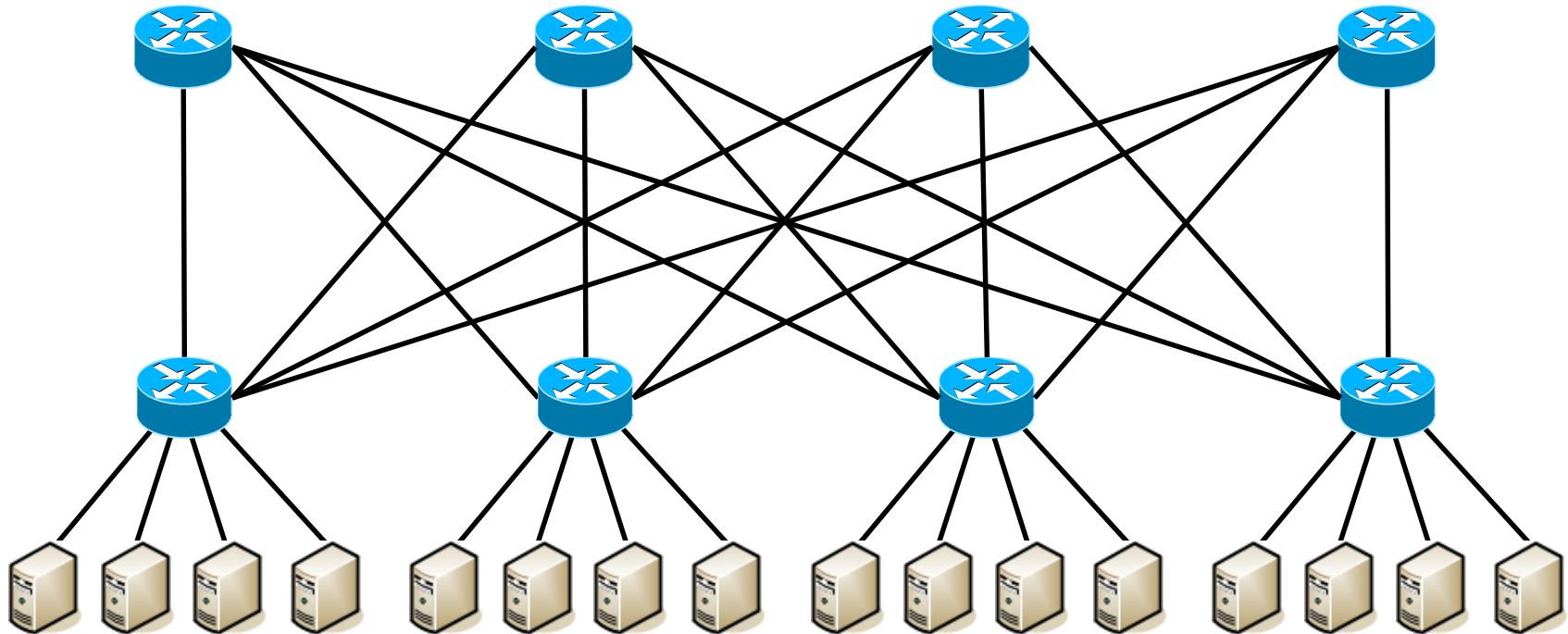


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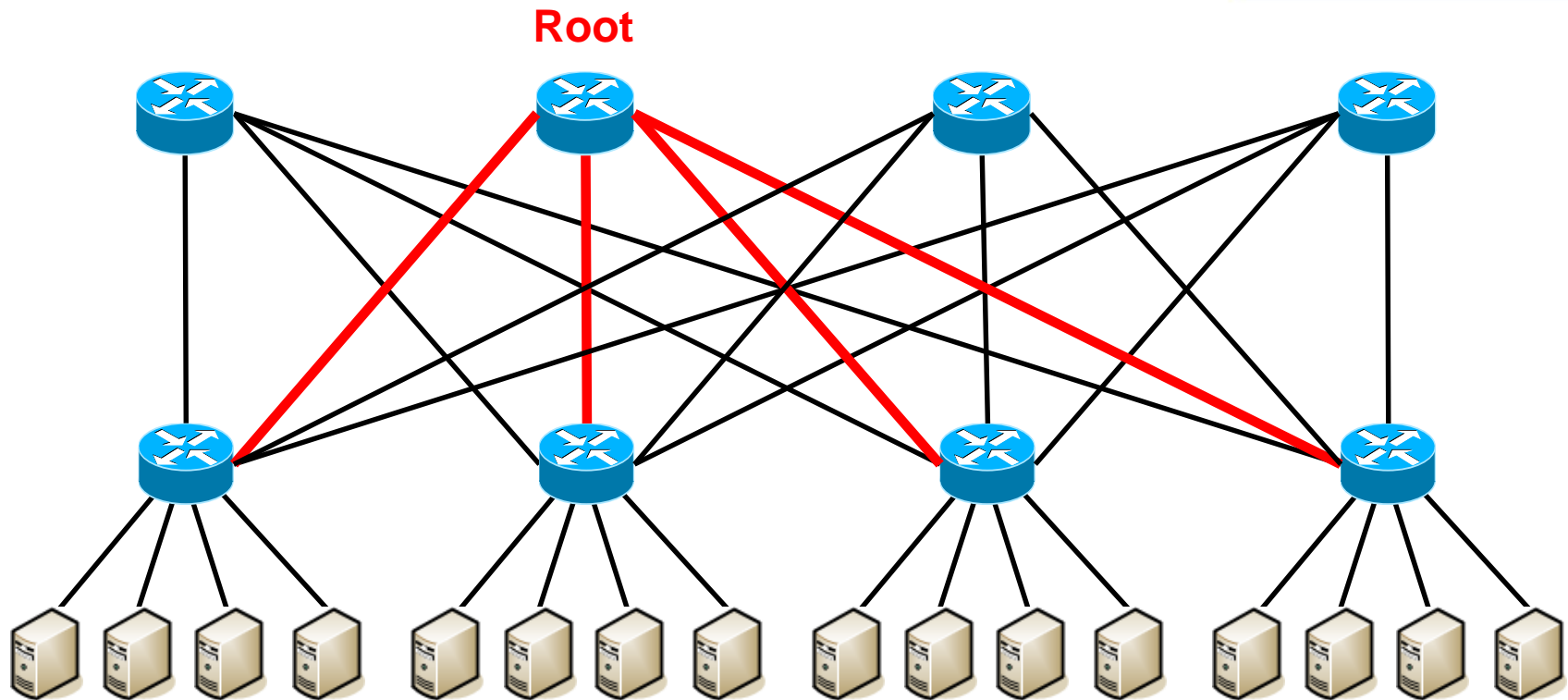




- Network-wide flooding:
  - Large number of control messages while disseminating host information
- Large forwarding tables:
  - Switches have to maintain large forwarding tables
    - Forwarding table size is proportional to the number of hosts due to flat addressing
- Broadcast traffic:
  - ARP and DHCP broadcasts consume bandwidth and processing resources at hosts and switches
- Spanning tree:
  - Waste of bandwidth at unused links
  - Load imbalance



- Spanning-tree inefficiencies:
  - Limited bandwidth utilization, since many links are not used
  - Requires very expensive core switches with large switching capacity



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- Configuration overhead:
  - Administrators should assign manually prefixes to subnets
  - Host IP address assignment via DHCP should be consistent with prefix assignment
  - Network topology changes require manual reconfiguration
- Limited mobility support:
  - Host mobility is restricted within a subnet where the host can maintain its IP address
  - Mobility across subnets requires reconfiguration (e.g., IP address reassignment) and may cause service disruption

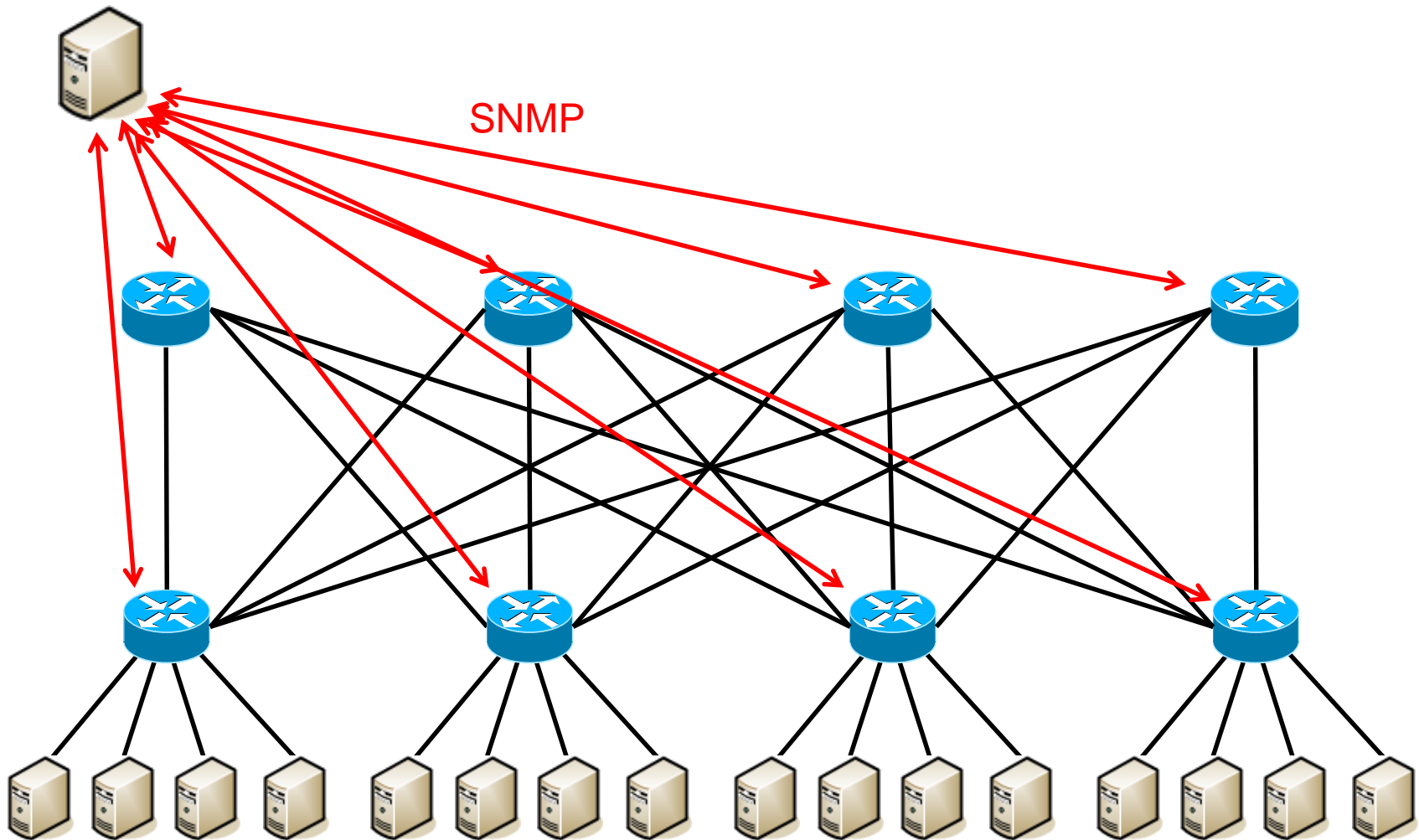


# Smart Path Assignment in Networks (SPAIN)



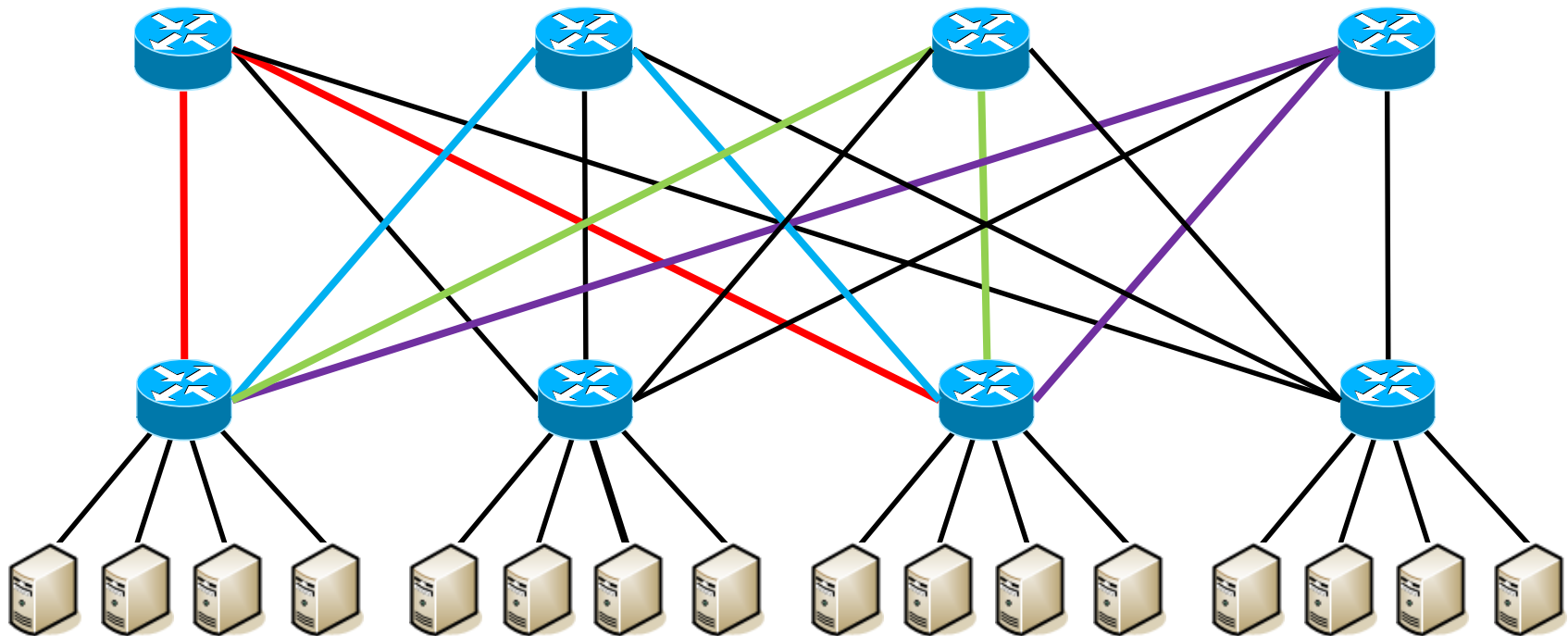
- Goal:
  - Provide multi-path routing using VLANs
  - Should work on arbitrary DC network topologies
  
- Offline computation of the network:
  - Topology discovery
  - Path computation
  - Assignment of paths to VLANs





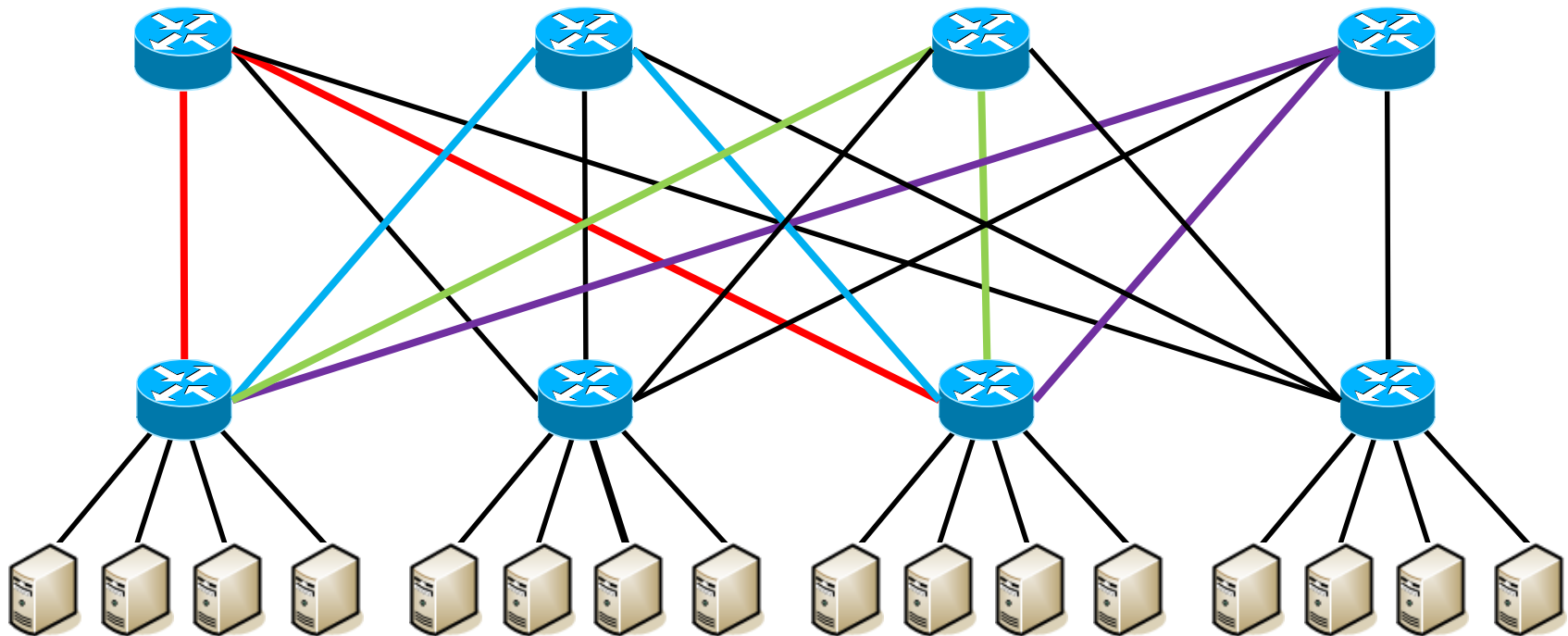


- Compute the smallest set of paths that exploit all redundancy
- Consider only paths between edge switches

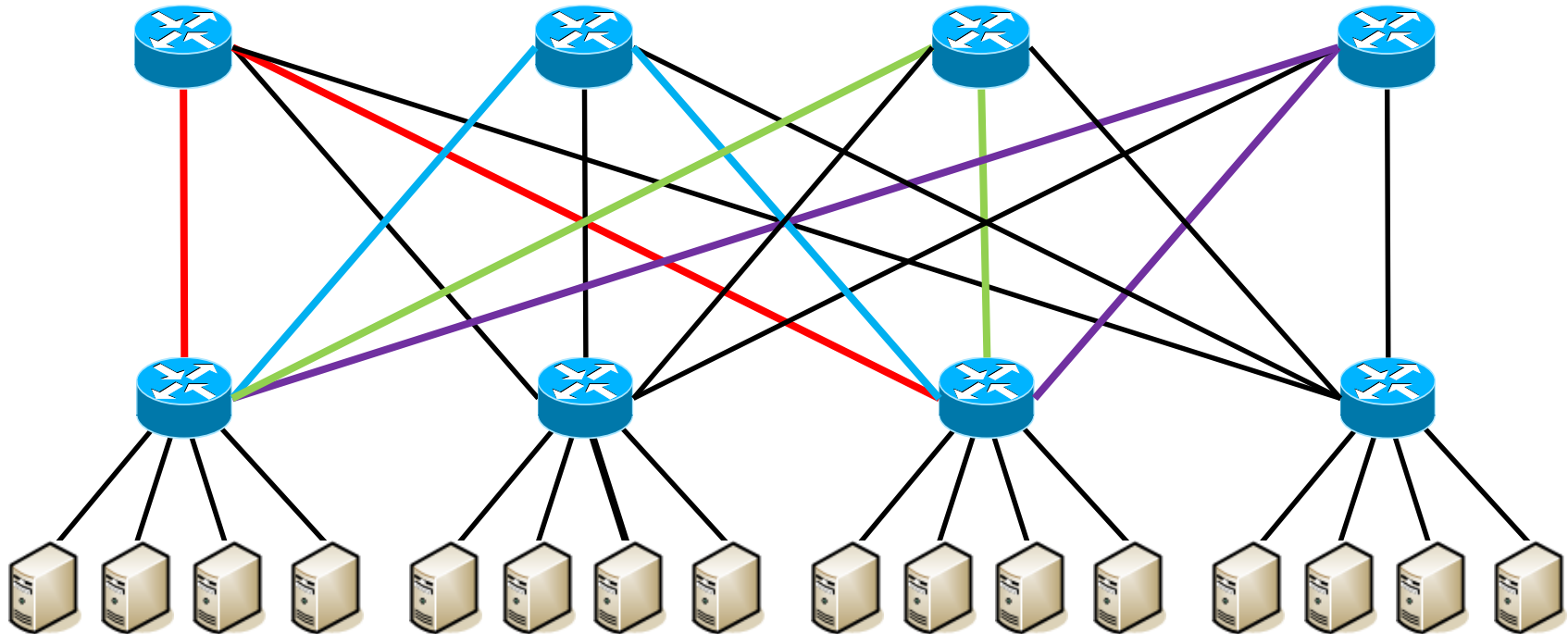




- Simple assignment:
  - Each path is assigned to a separate VLAN
  - Limited by the maximum number of VLANs (4096)
  - Scales only to a small number of switches



- Assignment proposed by SPAIN:
  - 1 VLAN is used for a set of paths
  - Greedy VLAN packing algorithm for optimizing path assignment





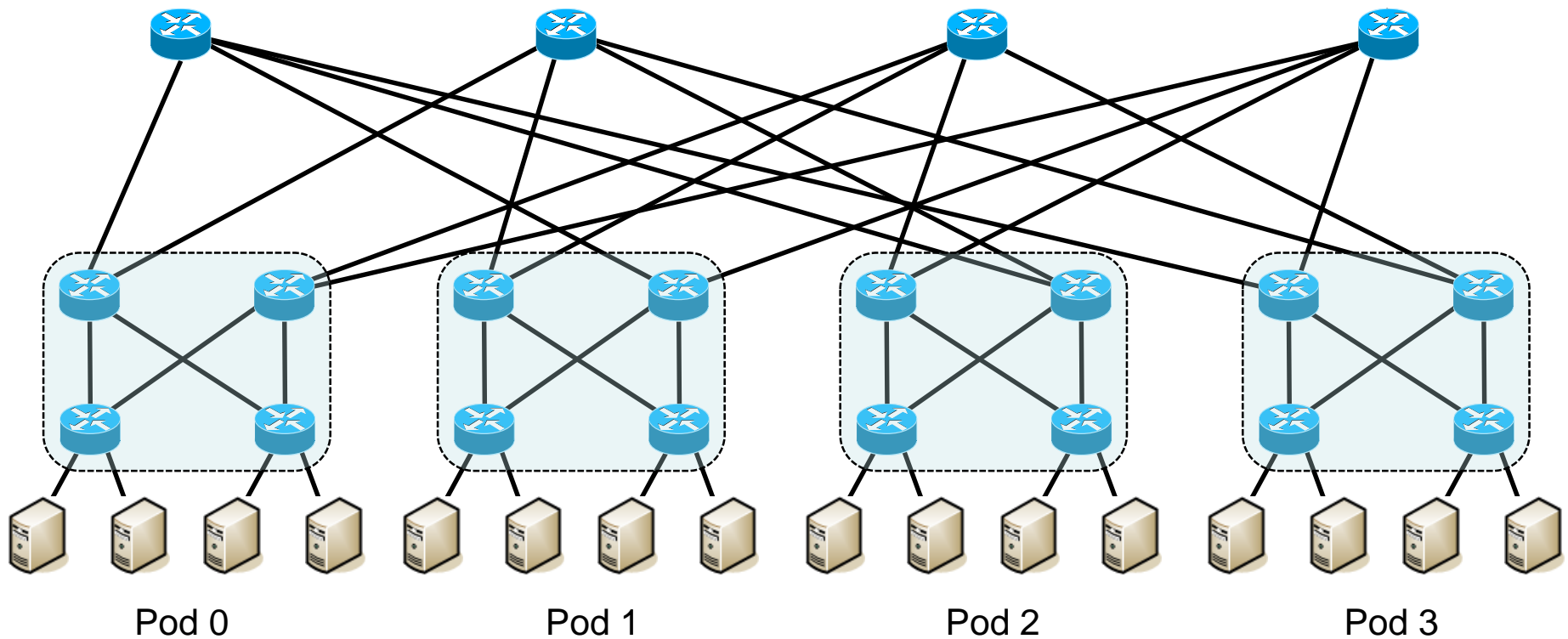
# PortLand



- PortLand is a single logical layer-2 data center network fabric that scales to millions of (virtual) end-points
- PortLand internally separates host identity from host location:
  - uses IP address as host identifier
  - introduces “Pseudo MAC” (PMAC) addresses internally to encode endpoint location
- PortLand runs on commodity switch hardware with unmodified hosts



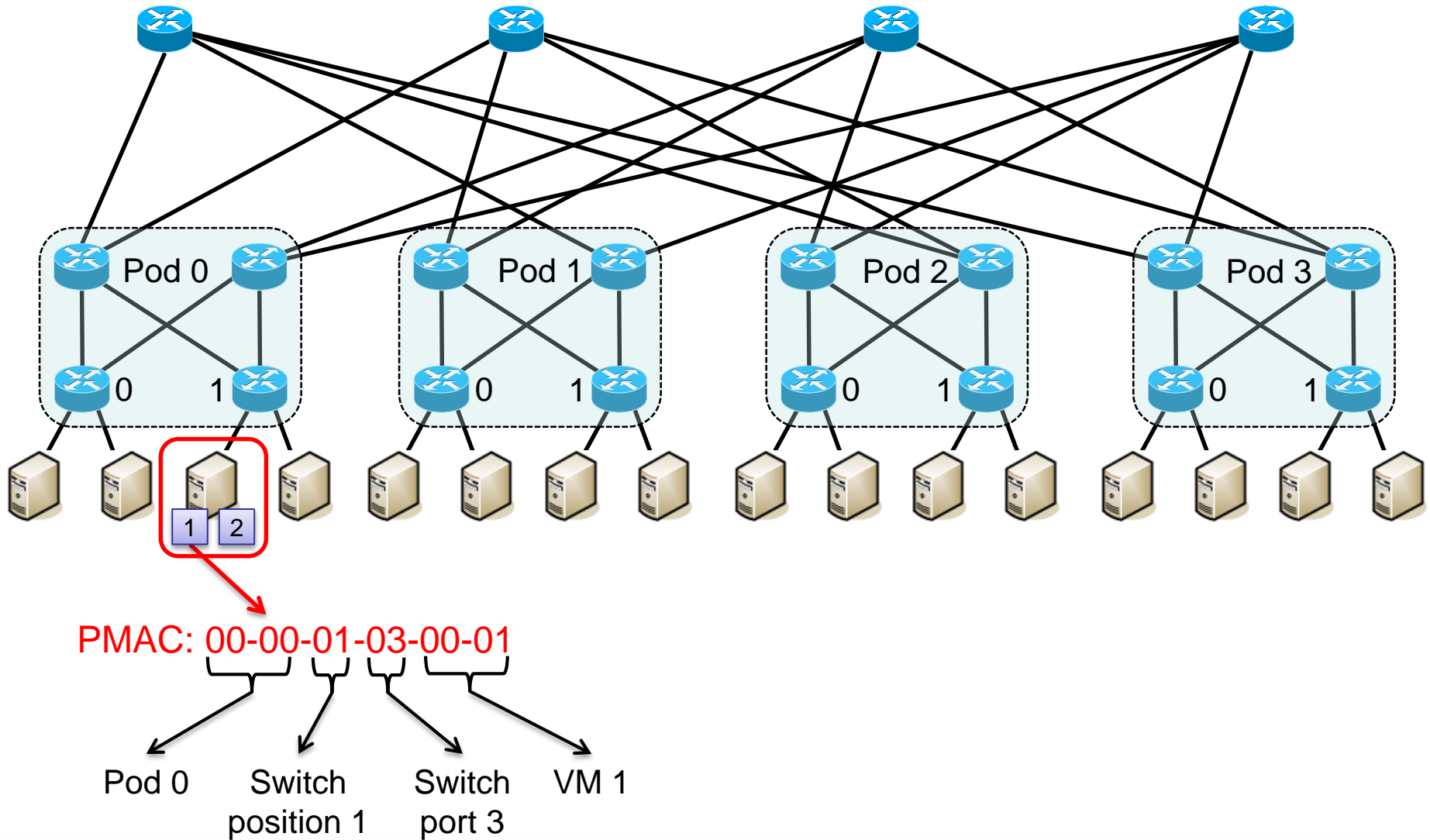
- PortLand assumes hierarchical structure of data center networks:
  - e.g., fat-tree topology (multi-rooted tree)

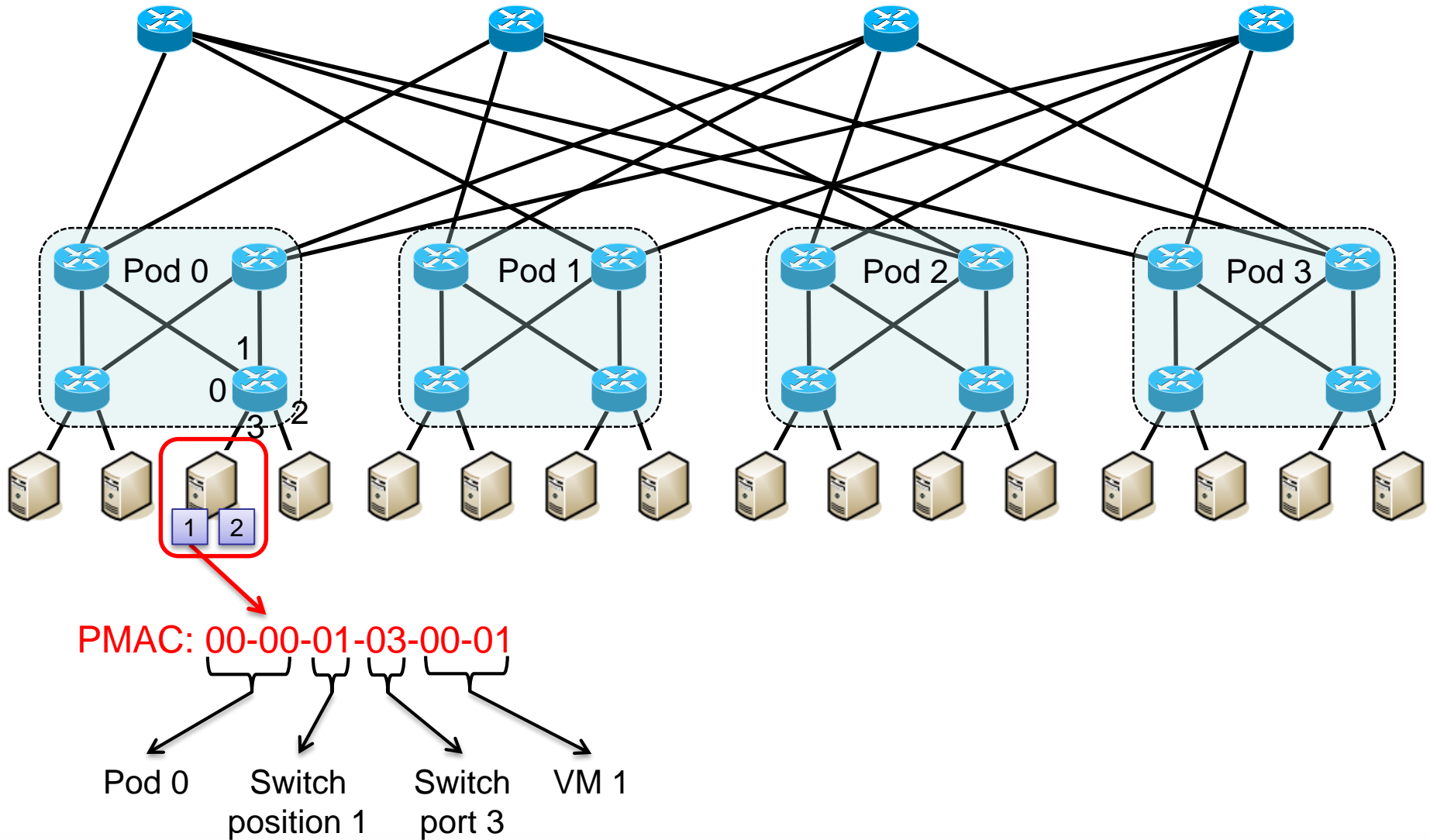




- PMAC addressing for packet forwarding and routing:
  - Besides IP and MAC, each end-host is assigned with a unique PMAC address
  - PMAC encodes the location of each end-host
- PMAC address format: `pod.position.port.vmid`
  - `pod`: pod number of the edge switch (16 bits)
  - `position`: position of the edge switch within the pod (8 bits)
  - `port`: switch port number (8 bits)
  - `vmid`: virtual machine ID for demultiplexing (16 bits)

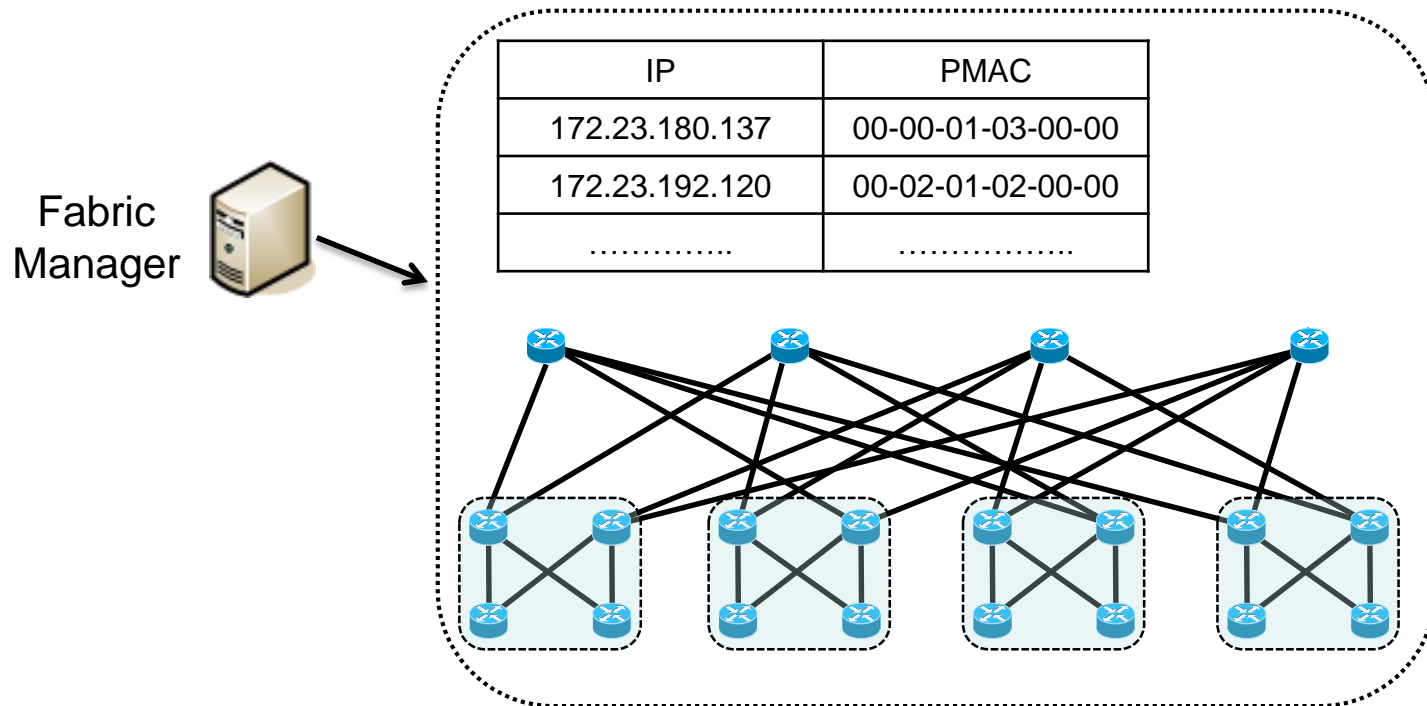






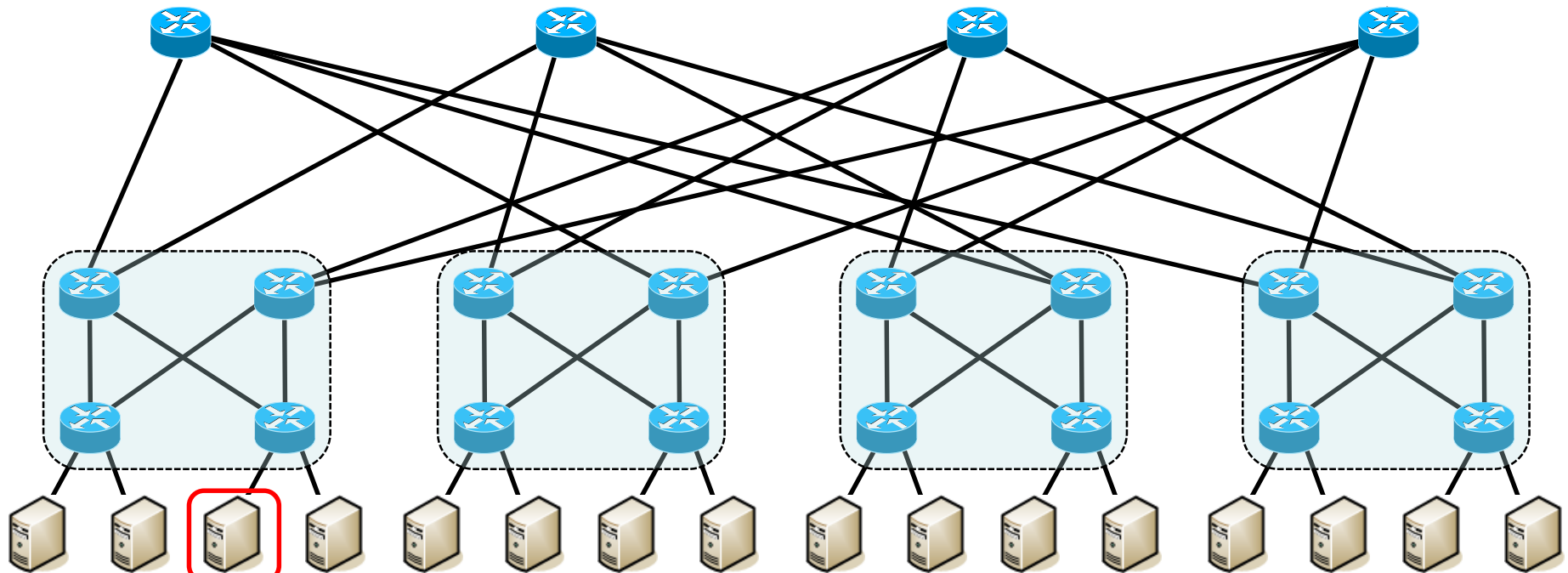


- Centralized Fabric Manager:
  - maintains (IP, PMAC) bindings, assisting ARP resolution
  - maintains the switch-level topology, facilitating fault-tolerant routing
  - maintains soft state, eliminating the need for manual configuration





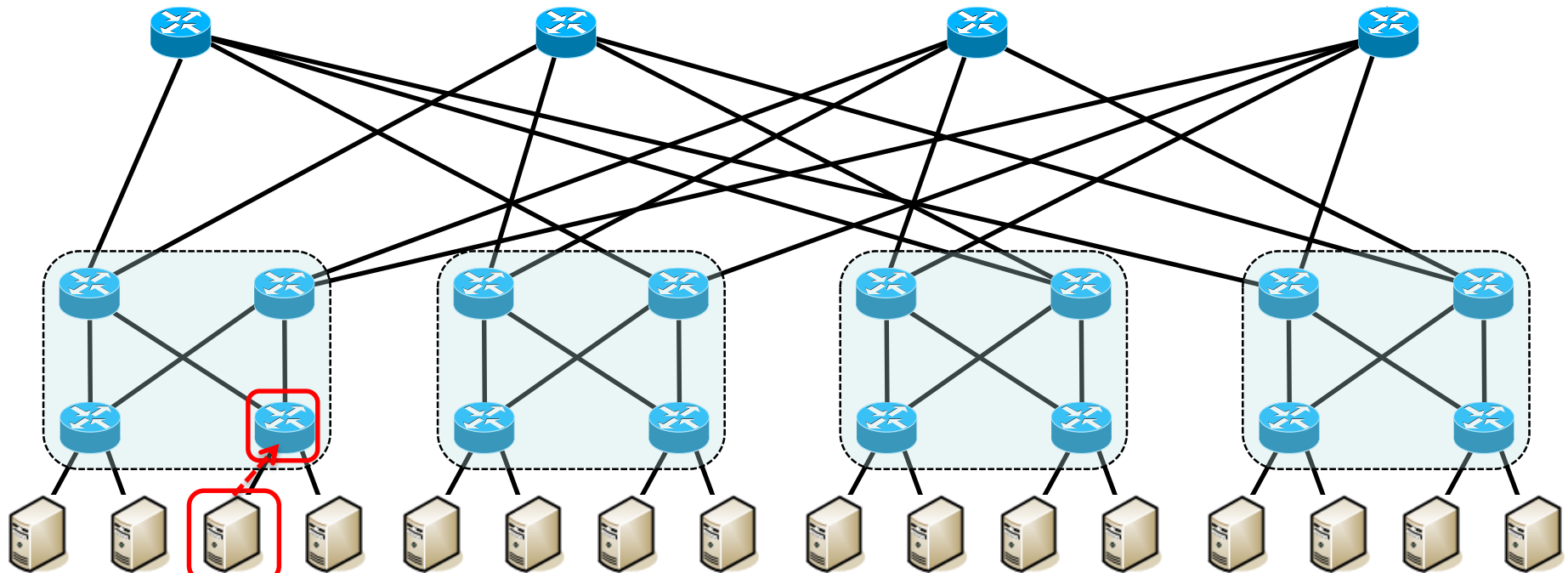
Fabric  
Manager



IP	MAC
172.23.180.137	84-2B-2B-A5-A5-77



Fabric  
Manager

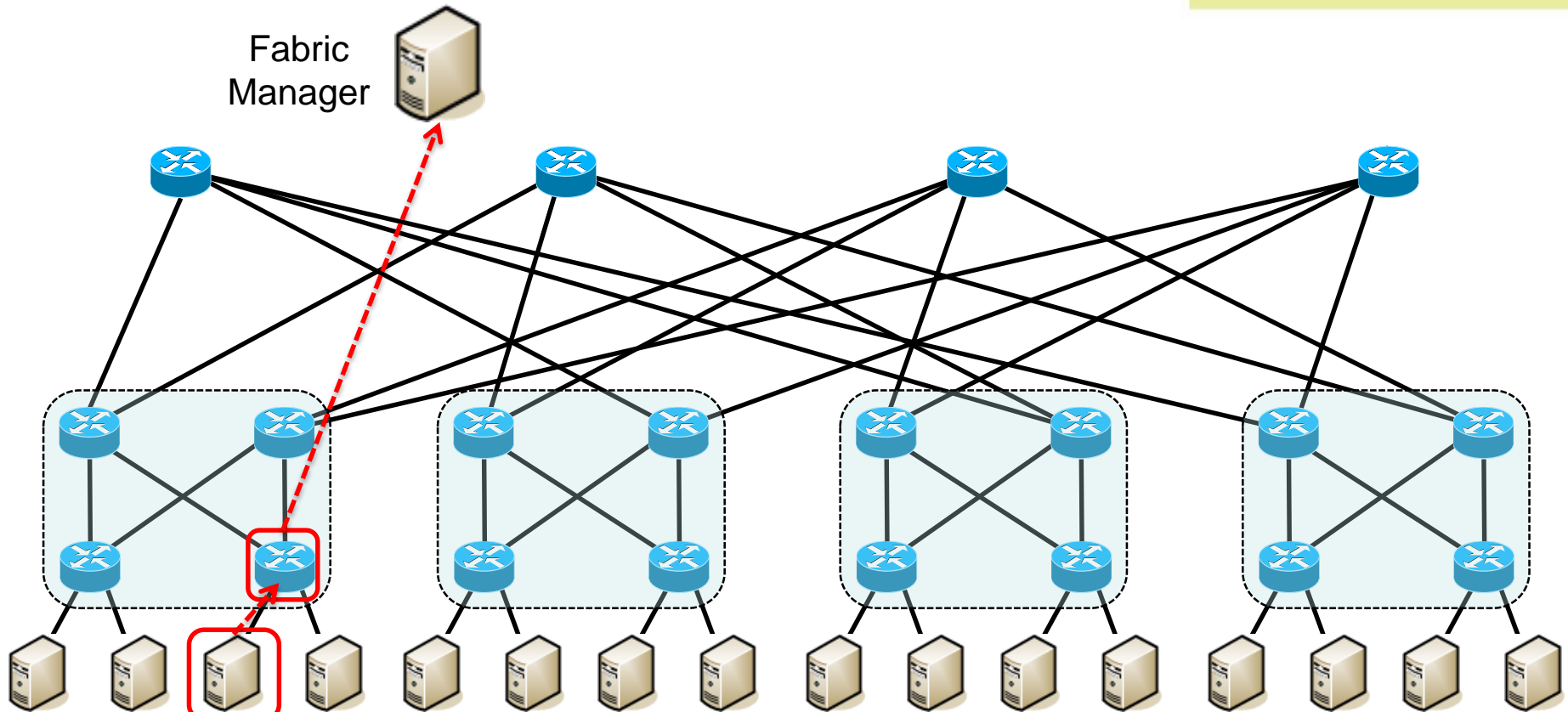


Edge switch generates  
PMAC 00-00-01-03-00-00

IP	MAC
172.23.180.137	84-2B-2B-A5-A5-77



Fabric  
Manager



Edge switch generates  
PMAC 00-00-01-03-00-00

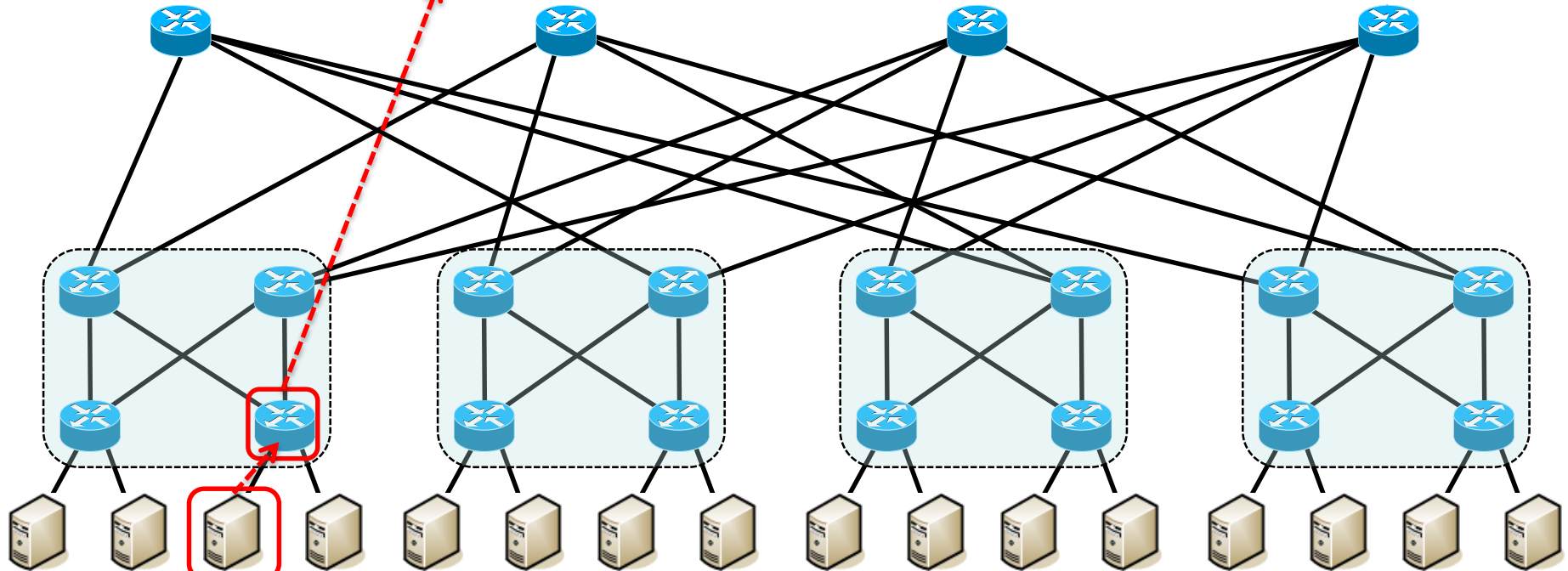
IP	MAC
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Fabric  
Manager



IP	PMAC
172.23.180.137	00-00-01-03-00-00



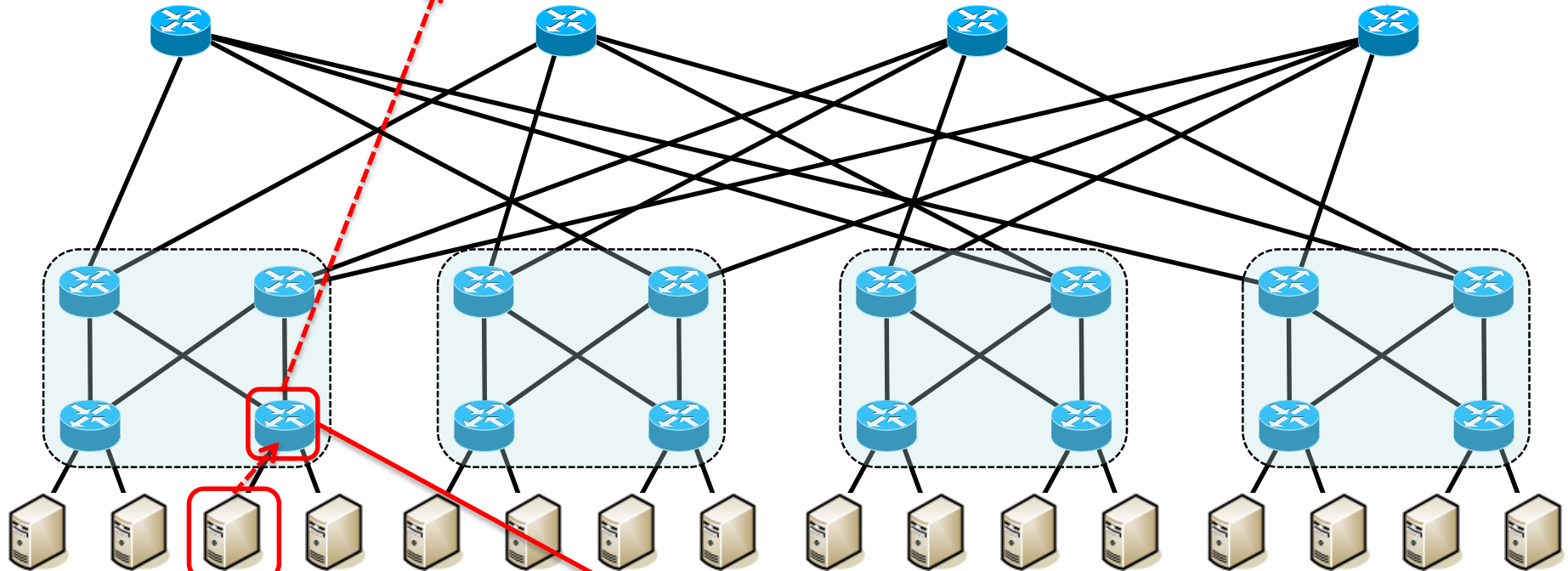
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Fabric  
Manager



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172.23.180.137	00-00-01-03-00-00



IP	MAC
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IP	MAC	PMAC
172.23.180.137	84-2B-2B-A5-A5-77	00-00-01-03-00-00



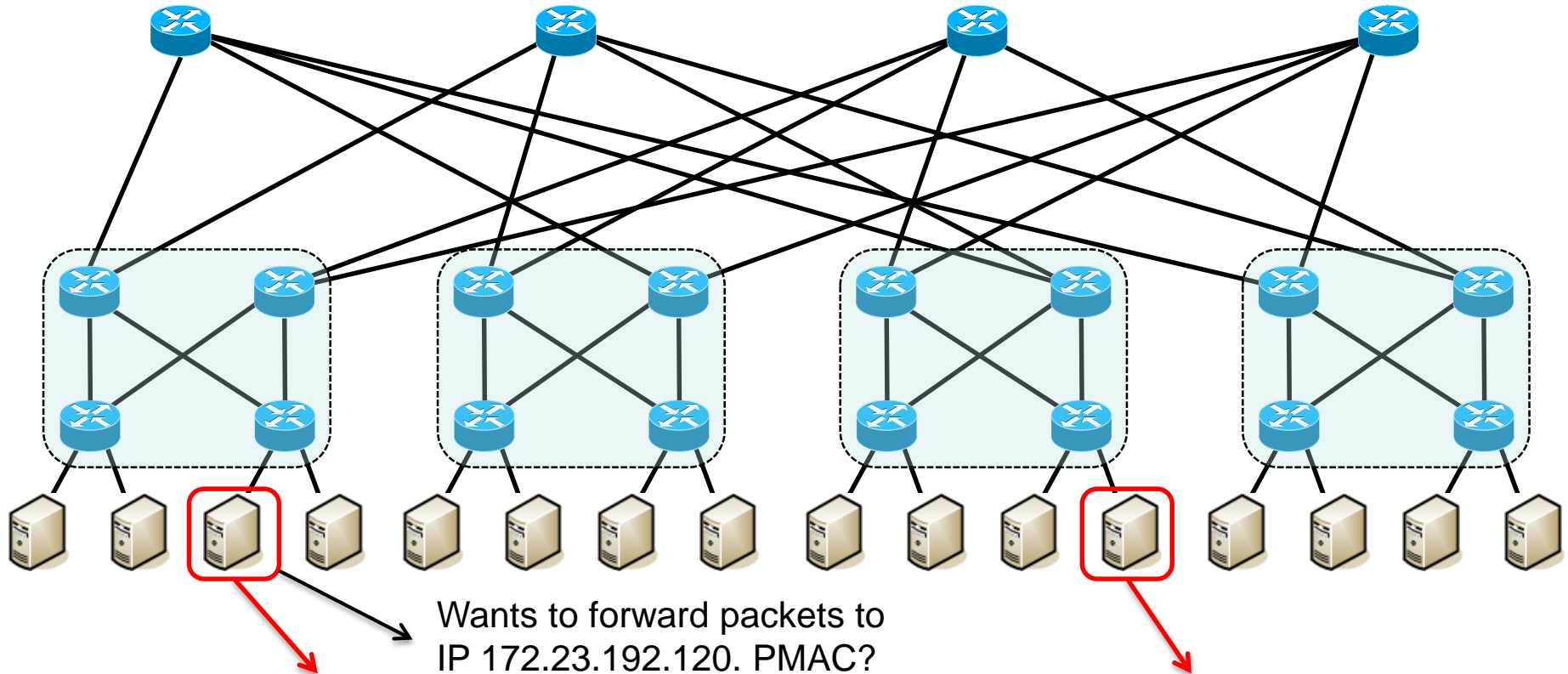
# Proxy-based Address Resolution



Fabric  
Manager



IP	PMAC
172.23.180.137	00-00-01-03-00-00
172.23.192.120	00-02-01-02-00-00



IP	MAC
172.23.180.137	84-2B-2B-A5-A5-77

IP	MAC
172.23.192.120	A3-B4-21-87-D4-12

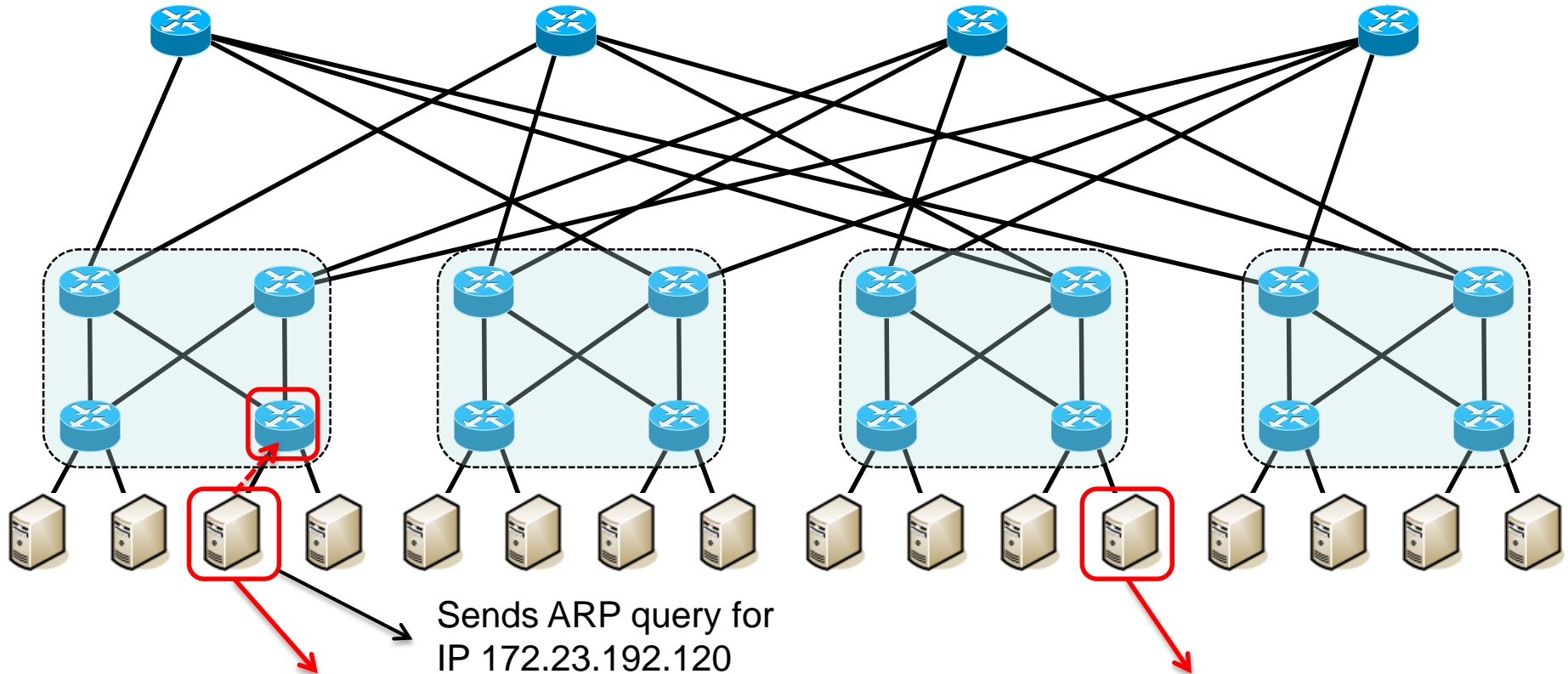
# Proxy-based Address Resolution



Fabric  
Manager



IP	PMAC
172.23.180.137	00-00-01-03-00-00
172.23.192.120	00-02-01-02-00-00



IP	MAC
172.23.180.137	84-2B-2B-A5-A5-77

IP	MAC
172.23.192.120	A3-B4-21-87-D4-12

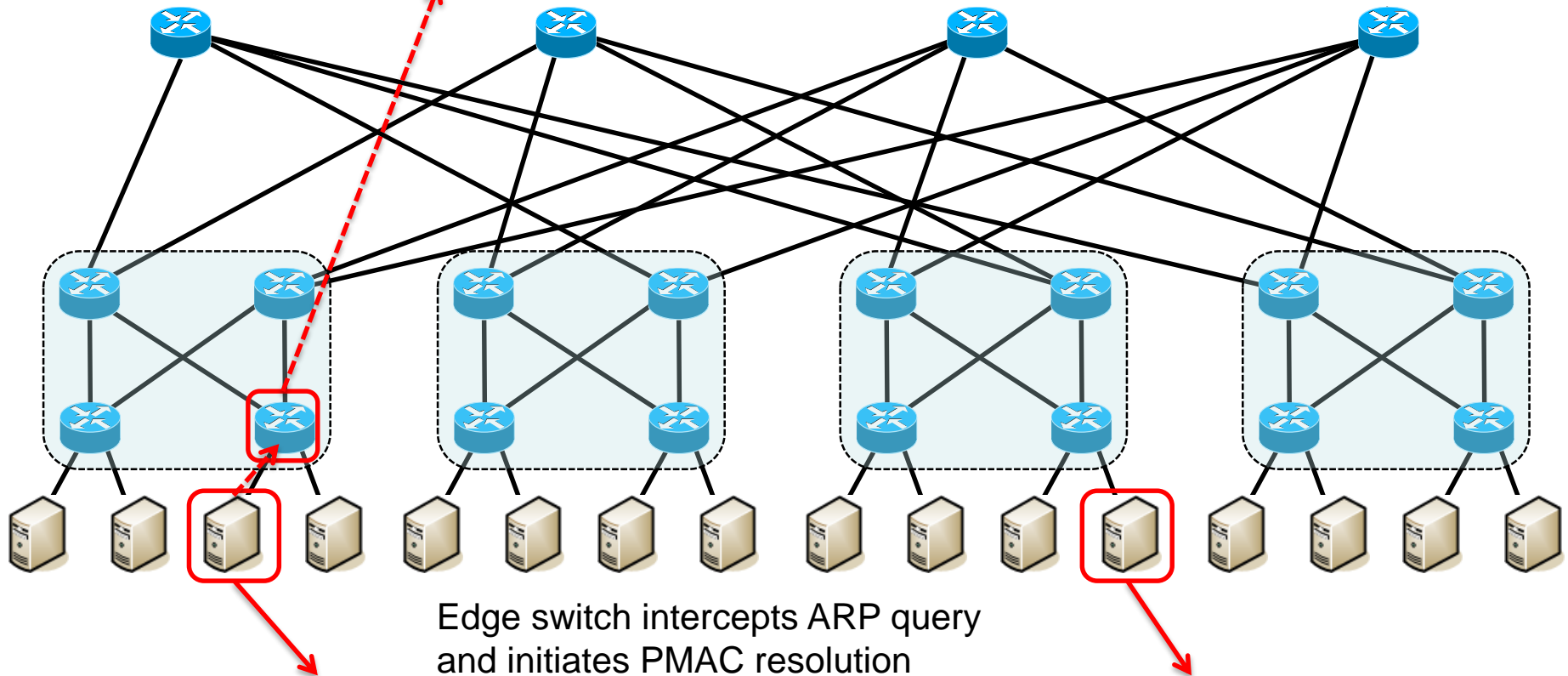
# Proxy-based Address Resolution



Fabric  
Manager



IP	PMAC
172.23.180.137	00-00-01-03-00-00
172.23.192.120	00-02-01-02-00-00



IP	MAC
172.23.180.137	84-2B-2B-A5-A5-77

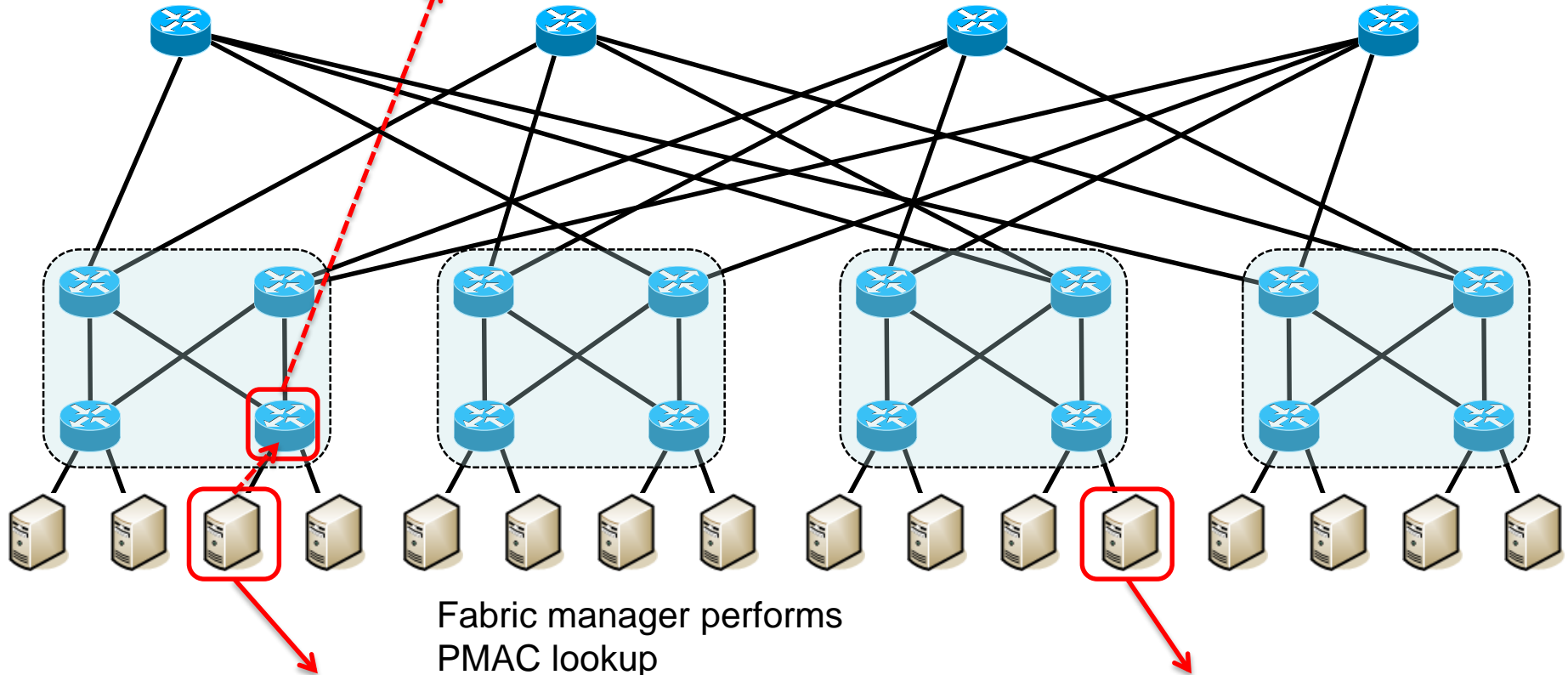
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172.23.192.120	A3-B4-21-87-D4-12

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Fabric  
Manager

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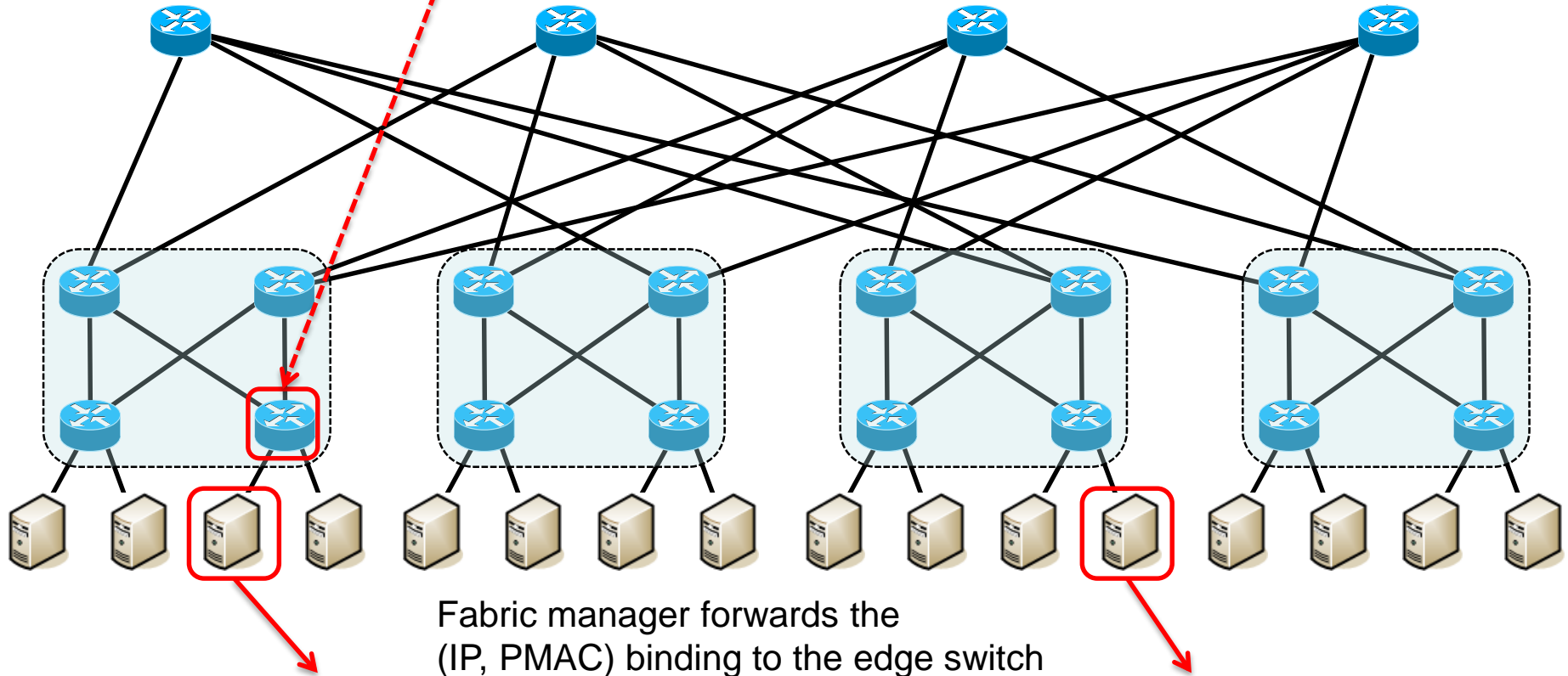
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172.23.192.120	A3-B4-21-87-D4-12

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Fabric  
Manager

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172.23.180.137	00-00-01-03-00-00
172.23.192.120	00-02-01-02-00-00



IP	MAC
172.23.180.137	84-2B-2B-A5-A5-77

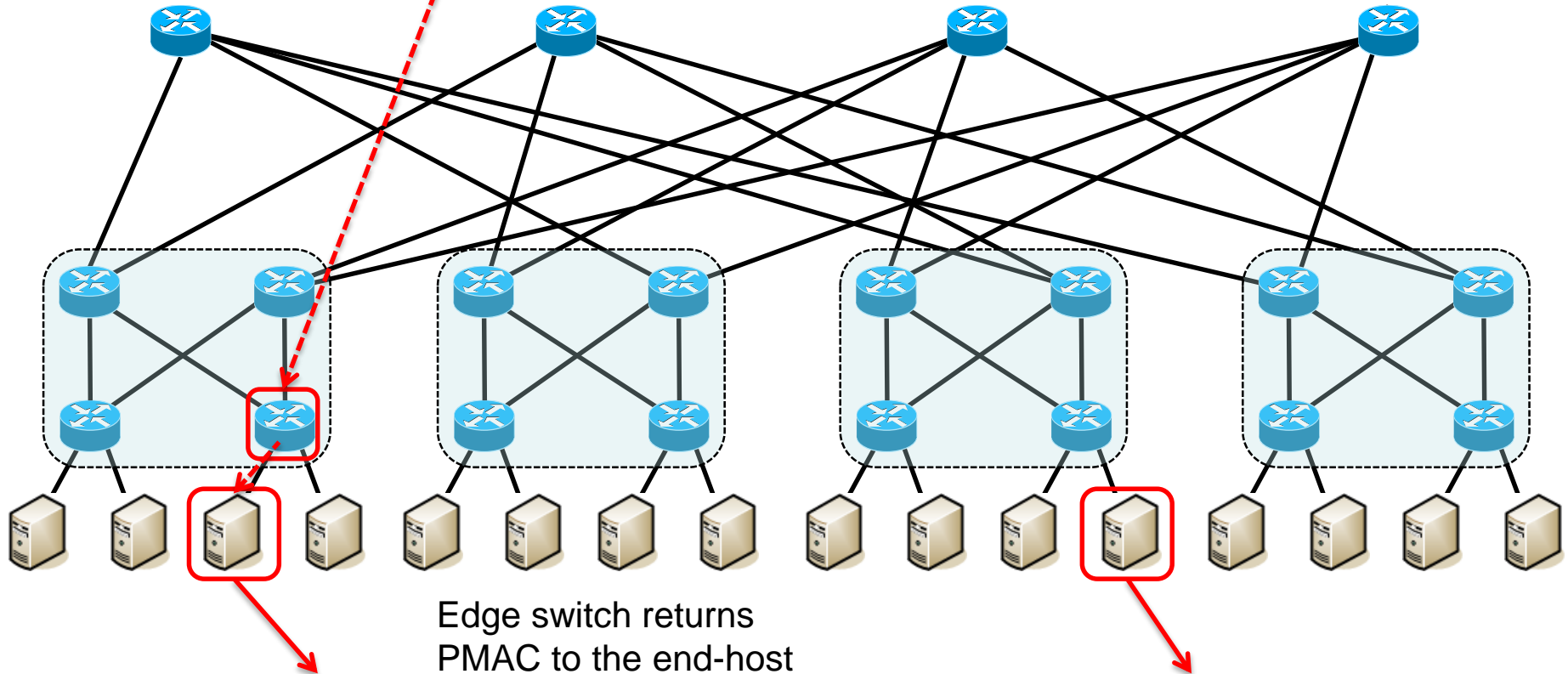
IP	MAC
172.23.192.120	A3-B4-21-87-D4-12

# Proxy-based Address Resolution



Fabric  
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IP	PMAC
172.23.180.137	00-00-01-03-00-00
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IP	MAC
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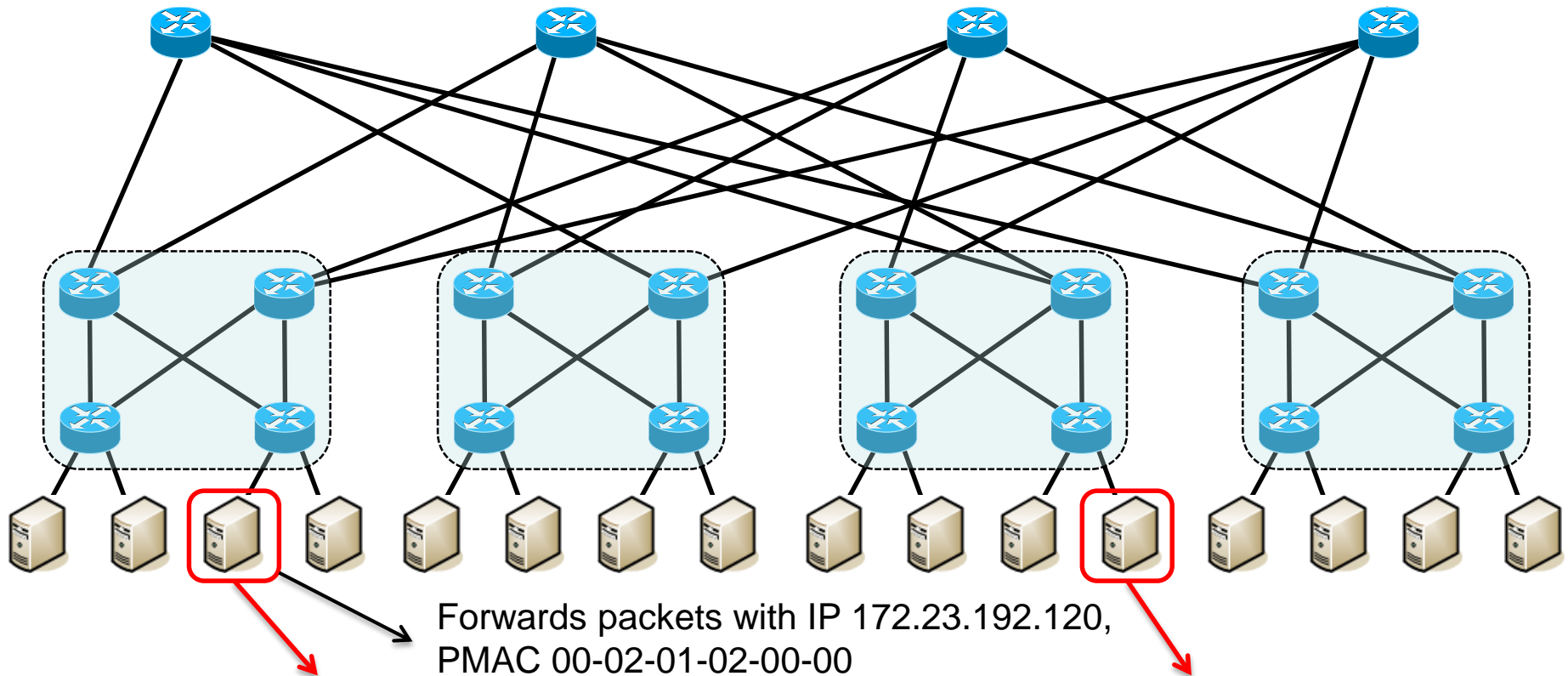
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172.23.192.120	A3-B4-21-87-D4-12



Fabric  
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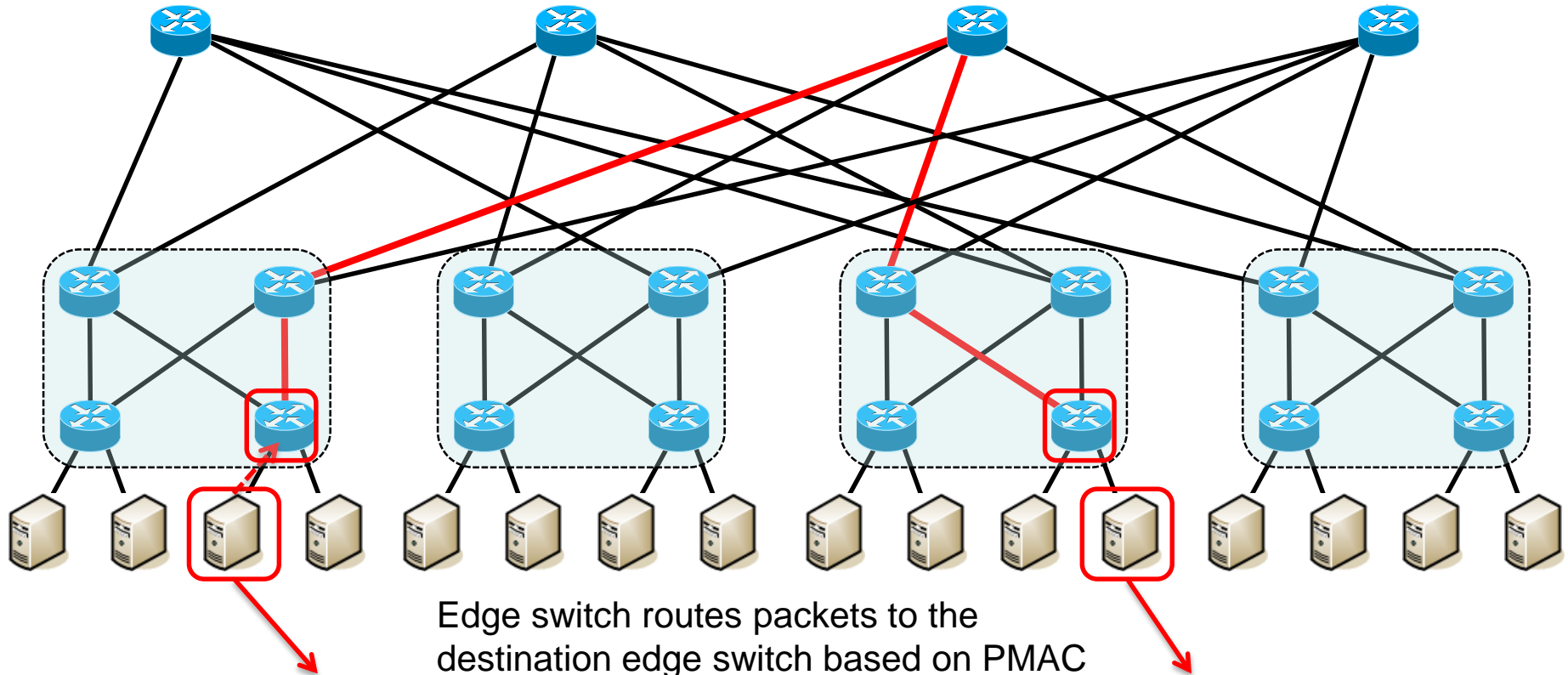
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172.23.192.120	A3-B4-21-87-D4-12



Fabric  
Manager



IP	PMAC
172.23.180.137	00-00-01-03-00-00
172.23.192.120	00-02-01-02-00-00



Edge switch routes packets to the  
destination edge switch based on PMAC

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172.23.180.137	84-2B-2B-A5-A5-77

IP	MAC
172.23.192.120	A3-B4-21-87-D4-12

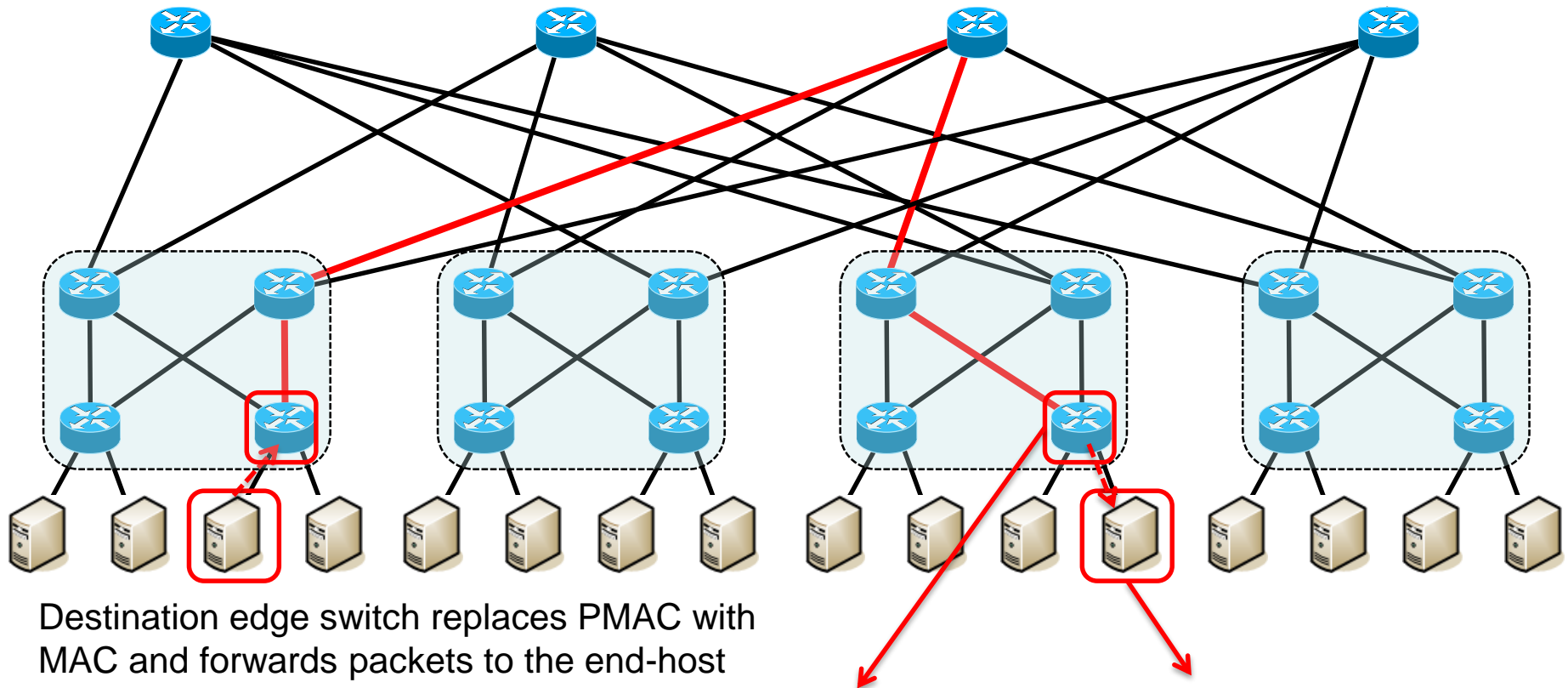




Fabric  
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172.23.180.137	00-00-01-03-00-00
172.23.192.120	00-02-01-02-00-00



IP	MAC	PMAC
172.23.192.120	A3-B4-21-87-D4-12	00-02-01-02-00-00

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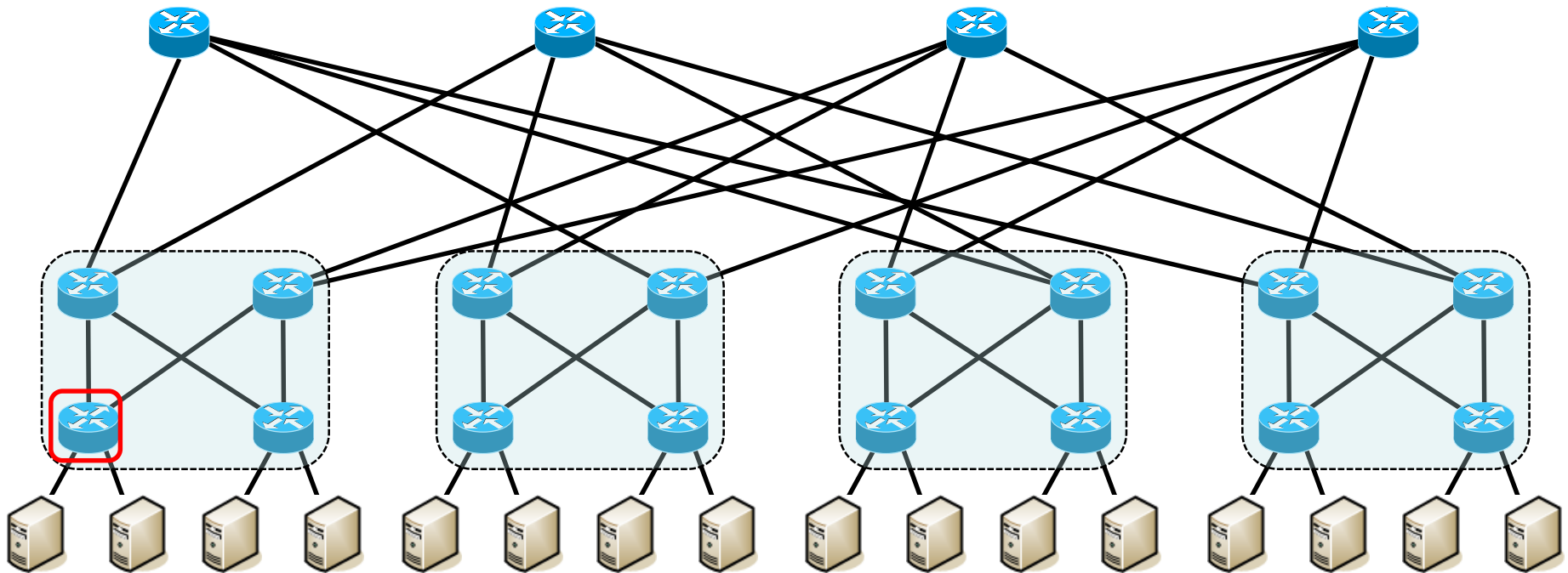
- Location Discovery Messages (LDMs) are exchanged between neighboring switches
- Switches self-discover the following location information on boot:
  - Tree level / role:
    - Based on neighbor identity
  - Position number:
    - Aggregation switches assist edge switches in choosing a unique position number
  - Pod number:
    - Fabric manager assists switches in choosing pod number



- LDMs include the following information:
  - Switch identifier:
    - Globally unique ID for each switch (e.g., the lowest MAC address of all local ports)
  - Pod number:
    - Unique pod number shared by all switches in the same pod
  - Position number:
    - Unique number for each switch in the same pod
  - Tree level:
    - Number that indicates whether a switch is an edge (0), aggregation (1) or core switch (2)
  - Up/down:
    - A bit that indicates whether a switch port is facing downward or upward in the fat tree



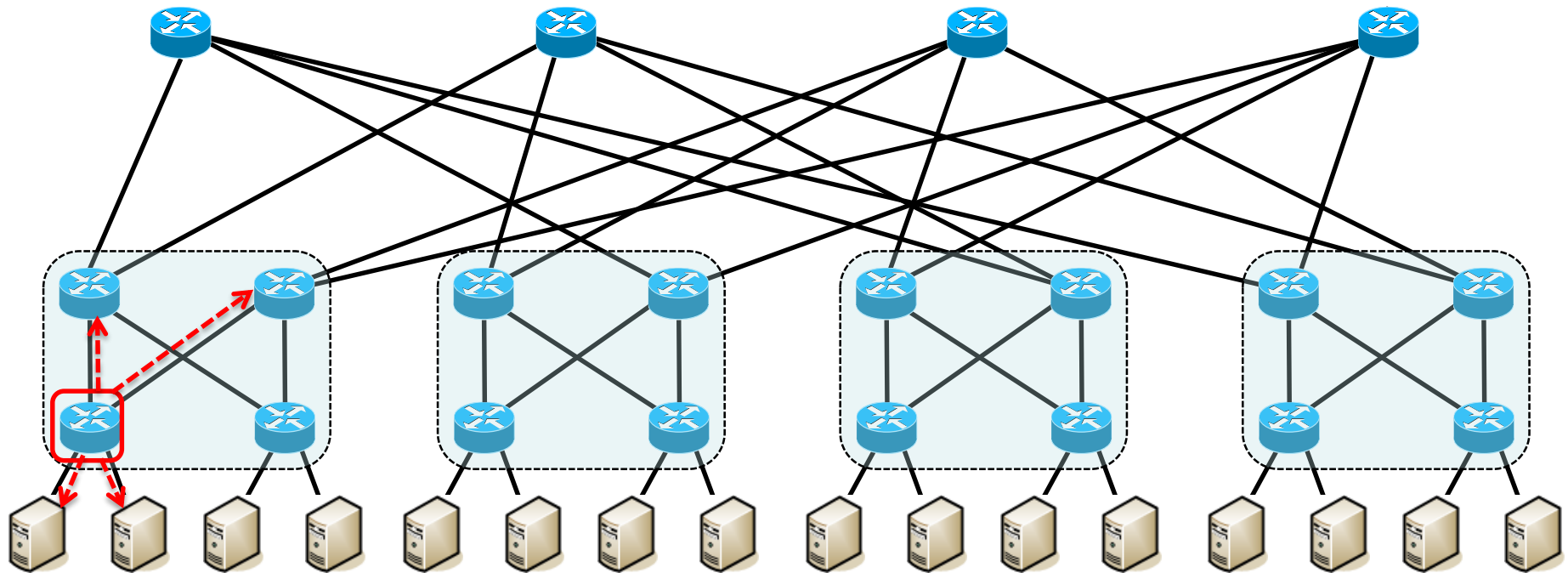
Fabric  
Manager



Switch ID	Pod Number	Position	Tree Level
A0-2B-FB-23-34-01			



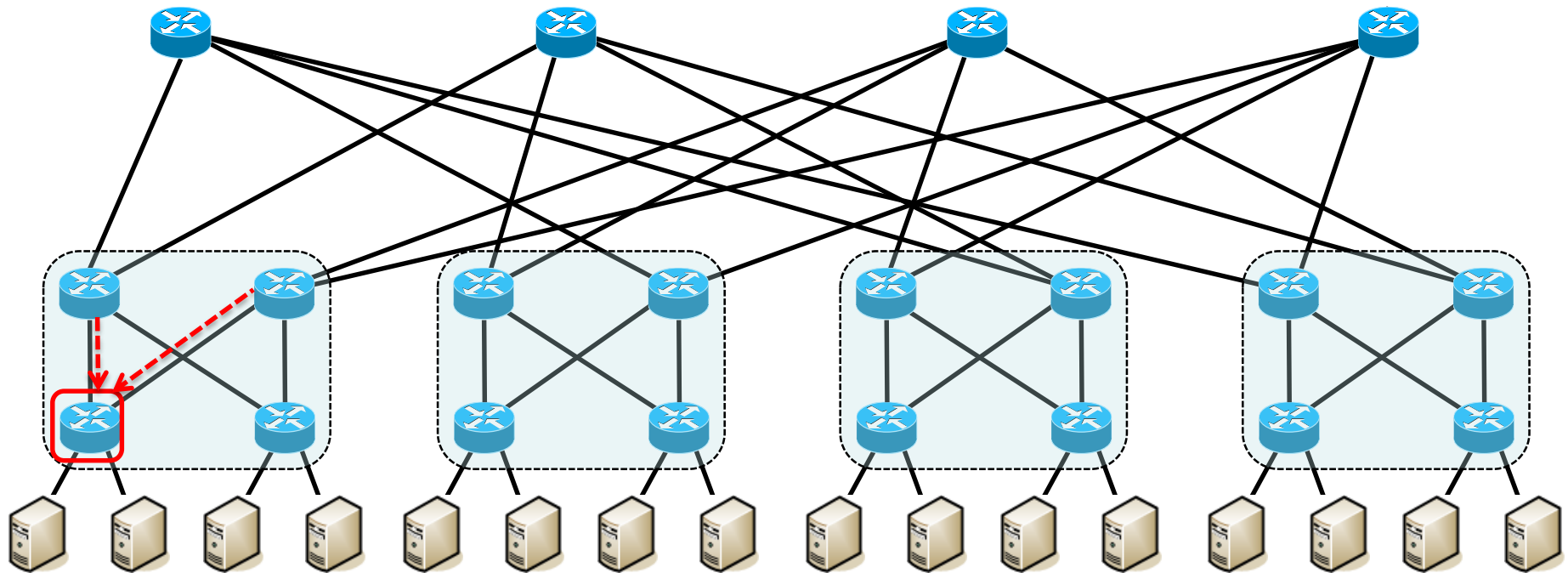
Fabric  
Manager



Switch ID	Pod Number	Position	Tree Level
A0-2B-FB-23-34-01			



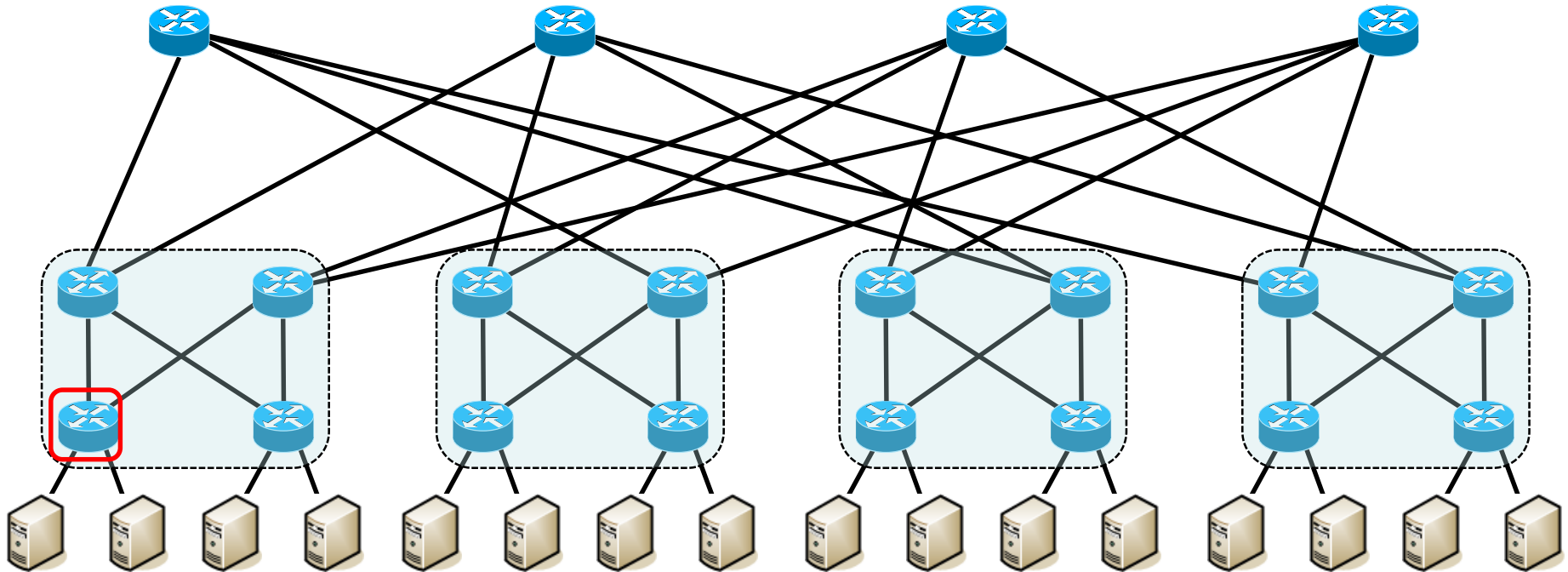
Fabric  
Manager



Switch ID	Pod Number	Position	Tree Level
A0-2B-FB-23-34-01			



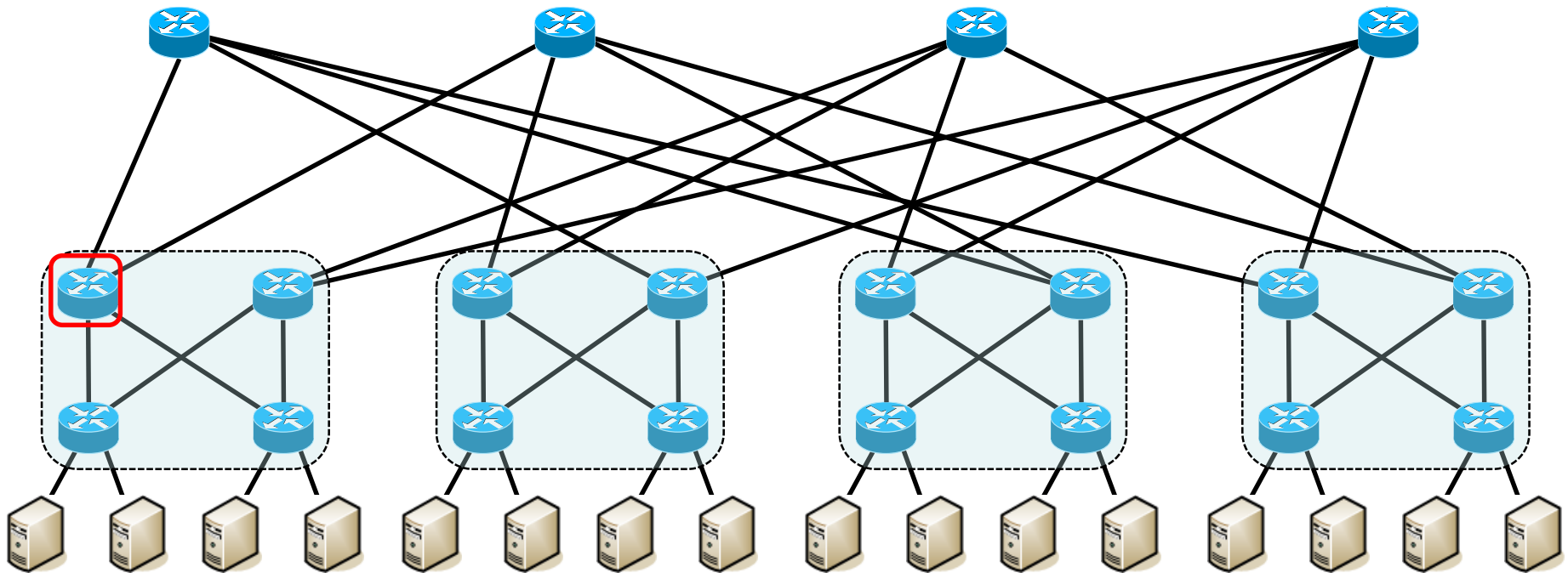
Fabric  
Manager



Switch ID	Pod Number	Position	Tree Level
A0-2B-FB-23-34-01			0



Fabric  
Manager

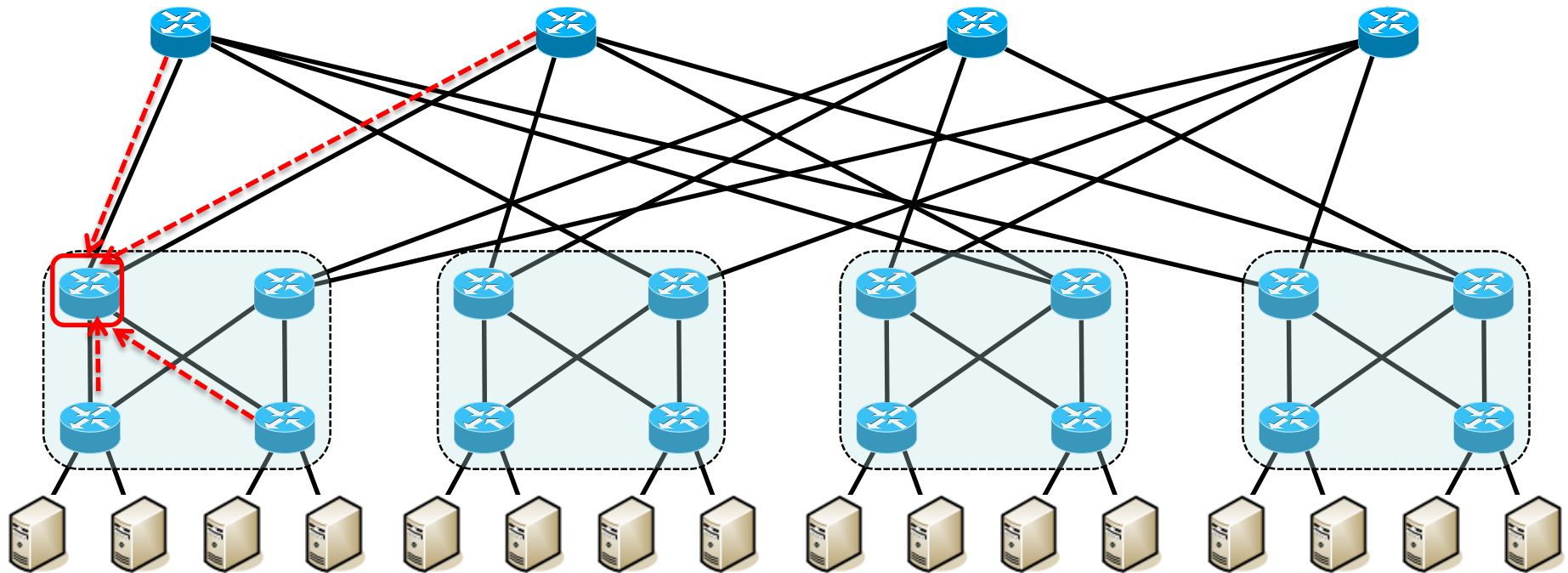


Switch ID	Pod Number	Position	Tree Level
A1-25-EB-23-2A-10			





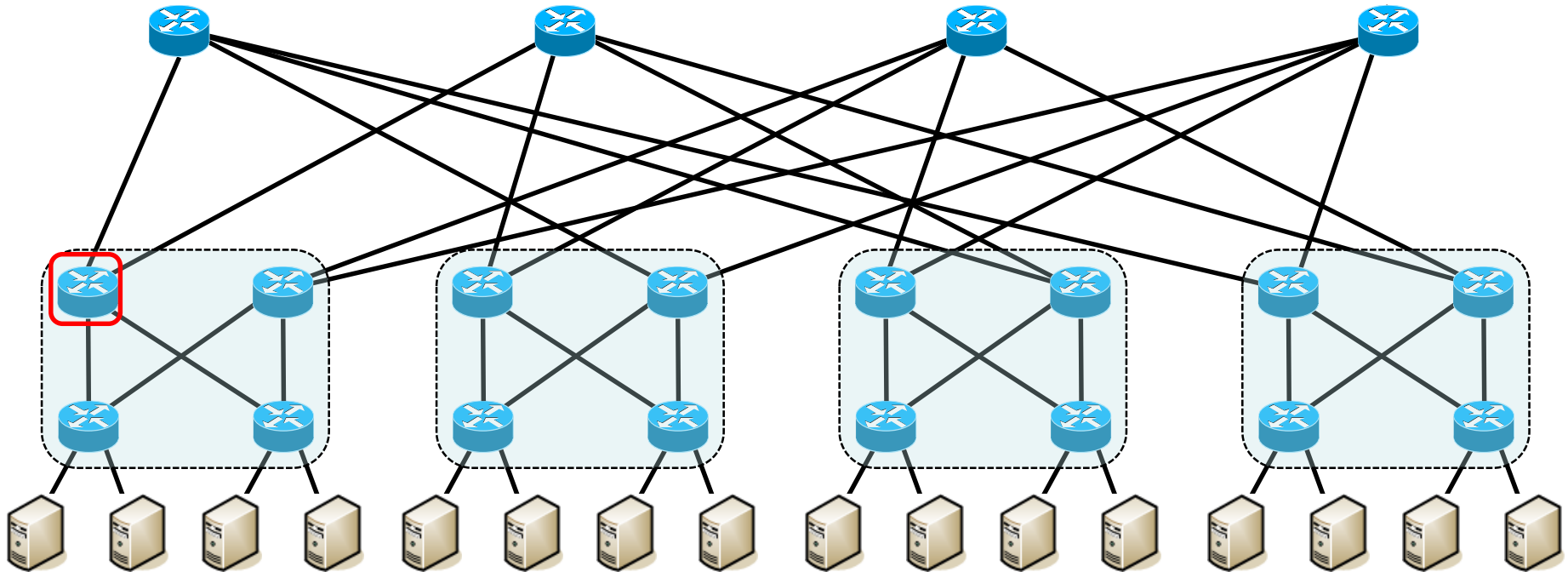
Fabric  
Manager



Switch ID	Pod Number	Position	Tree Level
A1-25-EB-23-2A-10			



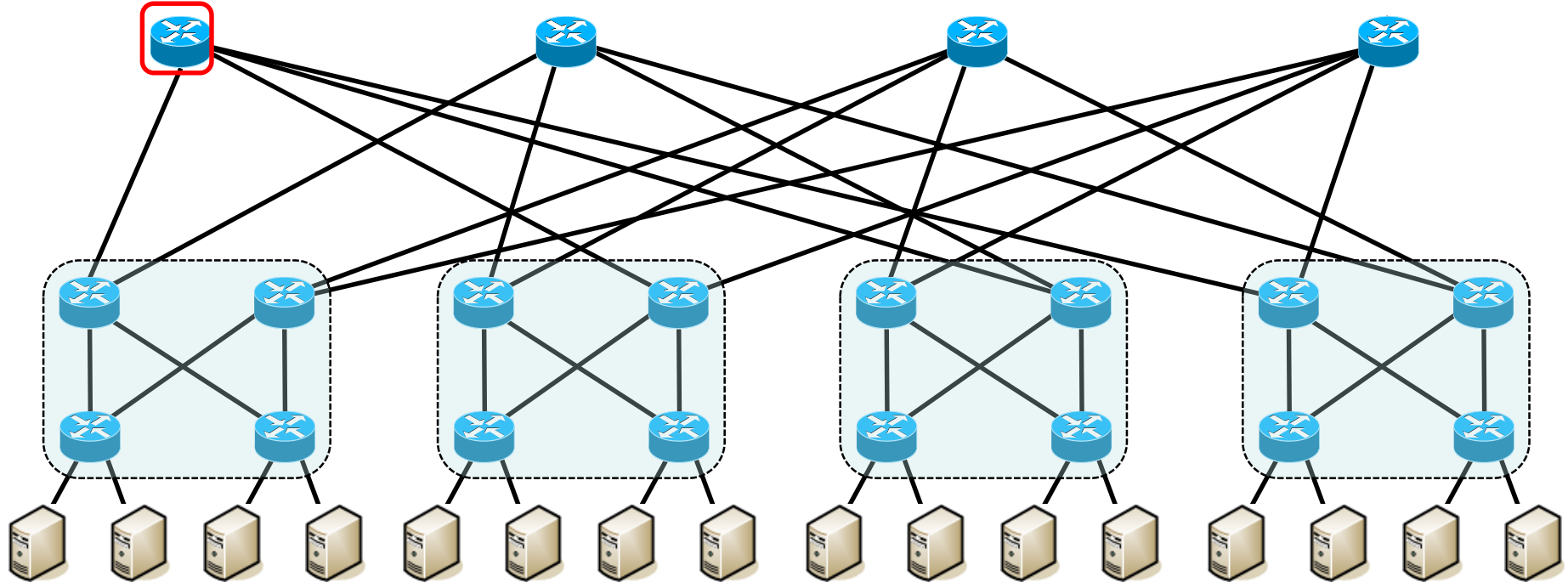
Fabric  
Manager



Switch ID	Pod Number	Position	Tree Level
A1-25-EB-23-2A-10			1



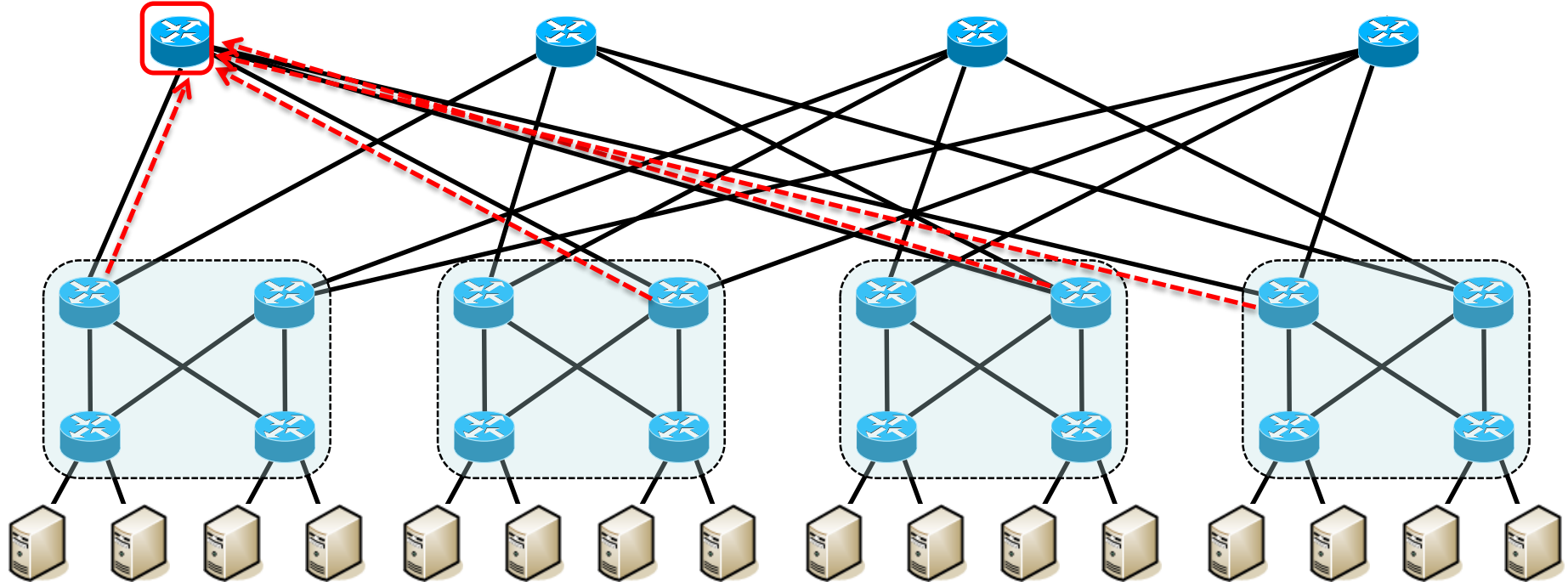
Fabric  
Manager



Switch ID	Pod Number	Position	Tree Level
B1-25-34-13-2A-10			



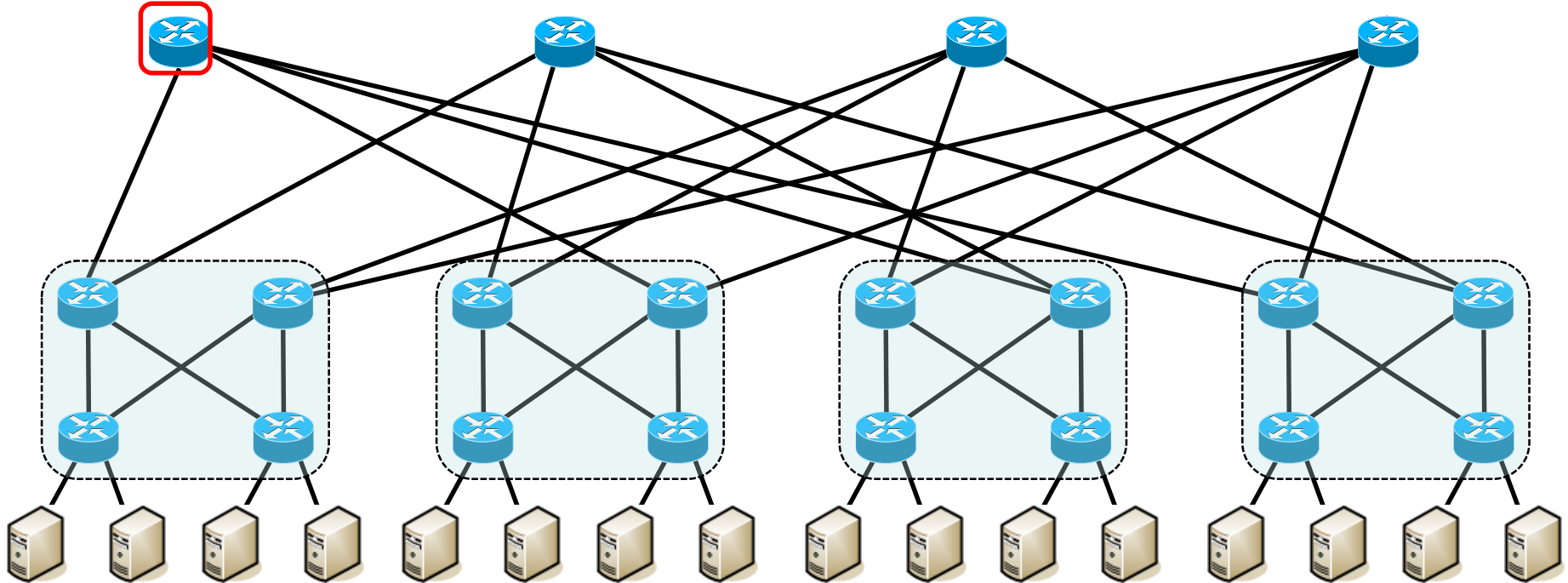
Fabric  
Manager



Switch ID	Pod Number	Position	Tree Level
B1-25-34-13-2A-10			



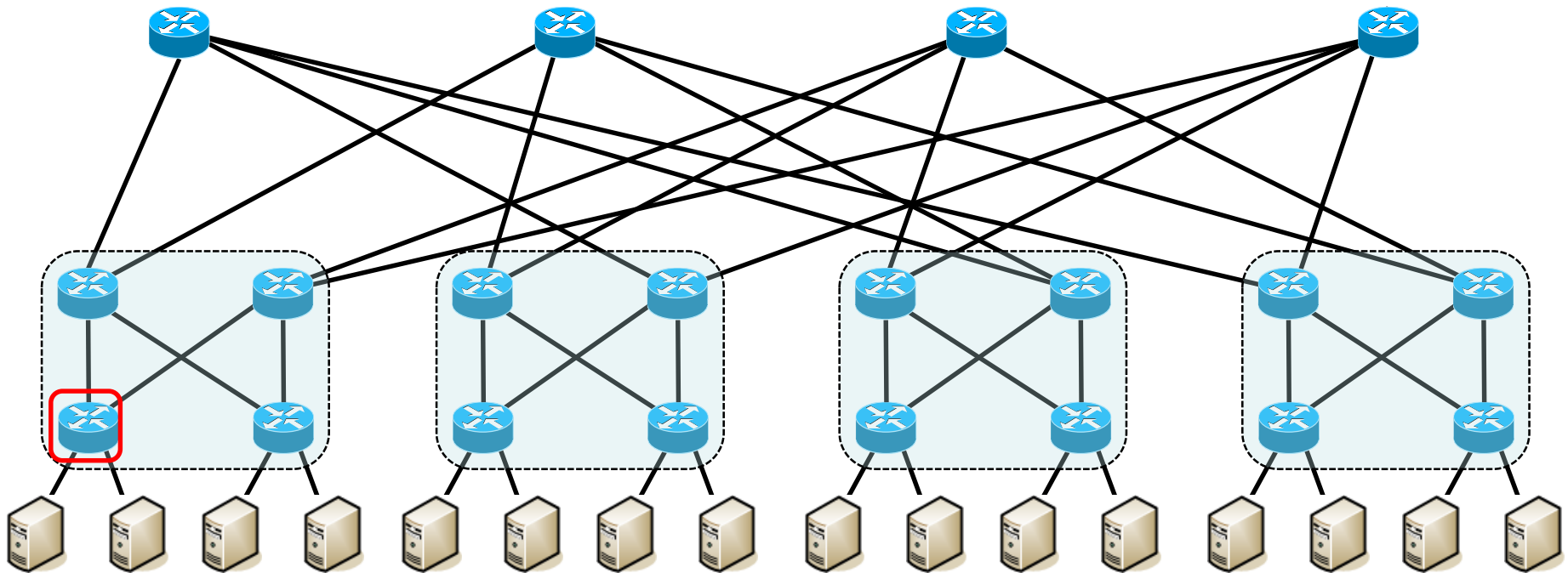
Fabric  
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Switch ID	Pod Number	Position	Tree Level
B1-25-34-13-2A-10			2



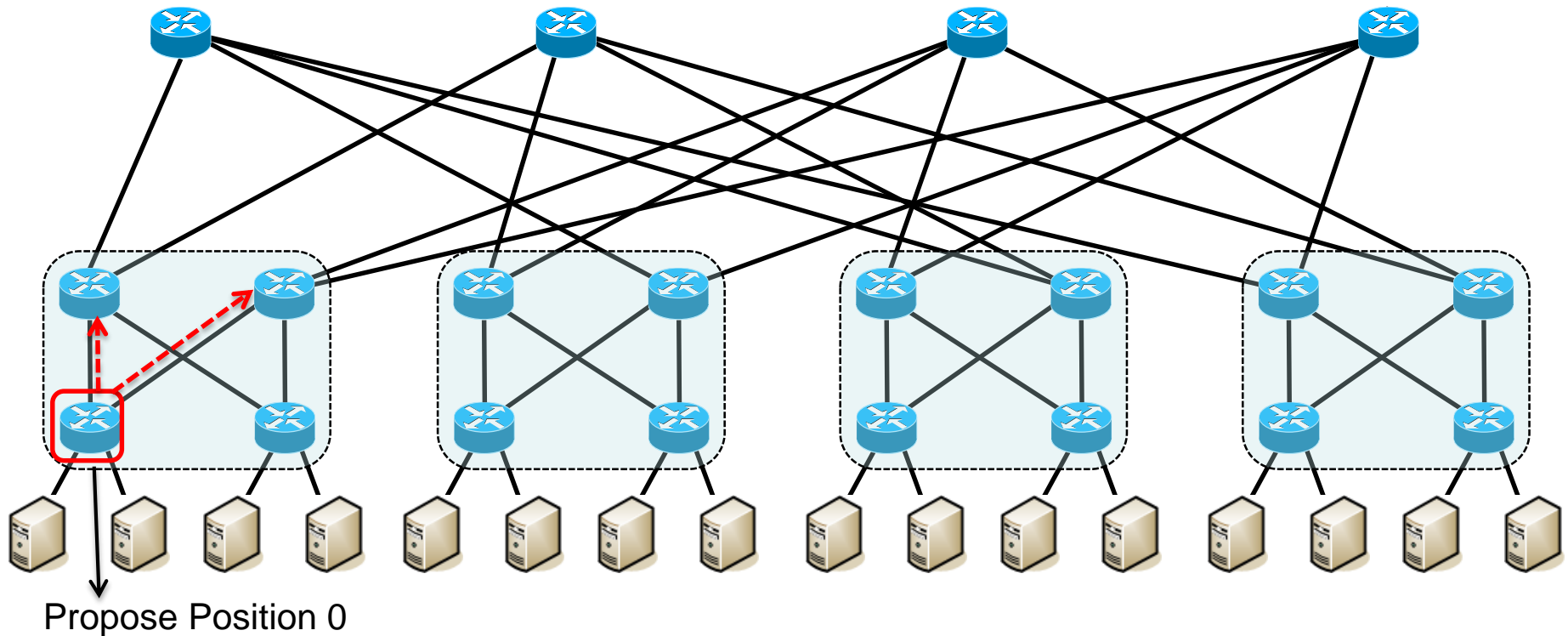
Fabric  
Manager



Switch ID	Pod Number	Position	Tree Level
A0-2B-FB-23-34-01			0



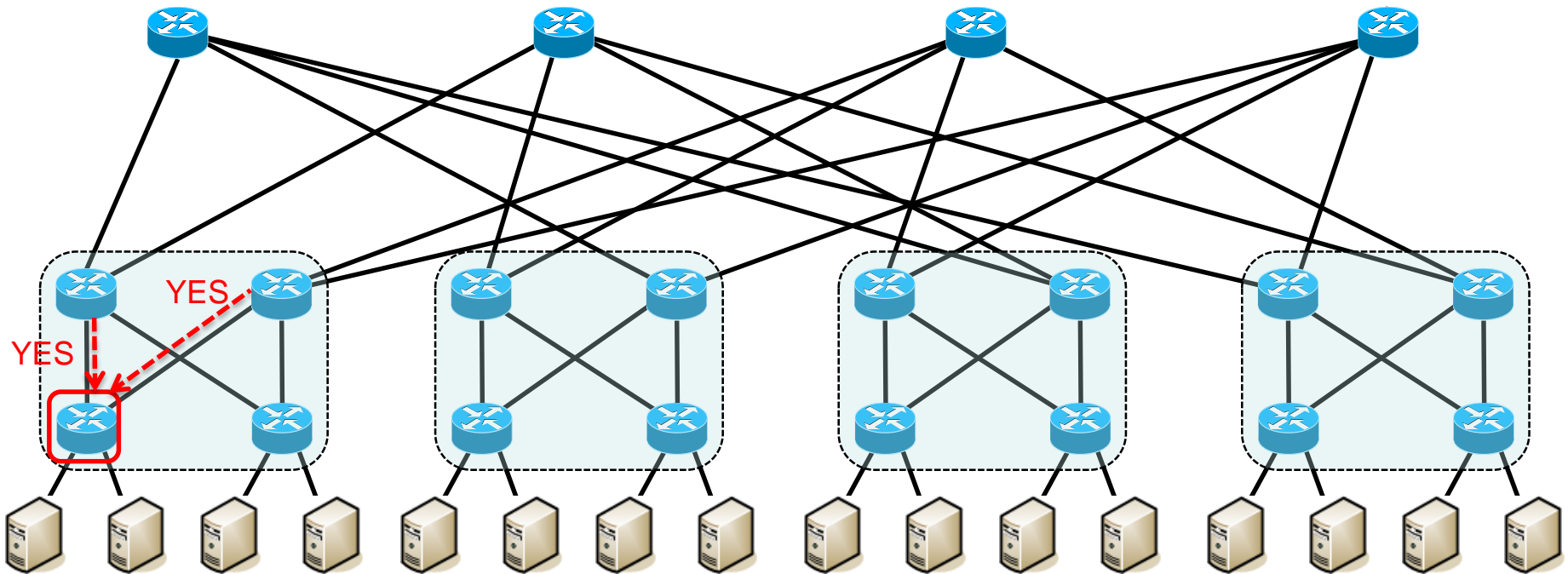
Fabric  
Manager



Switch ID	Pod Number	Position	Tree Level
A0-2B-FB-23-34-01			0



Fabric  
Manager

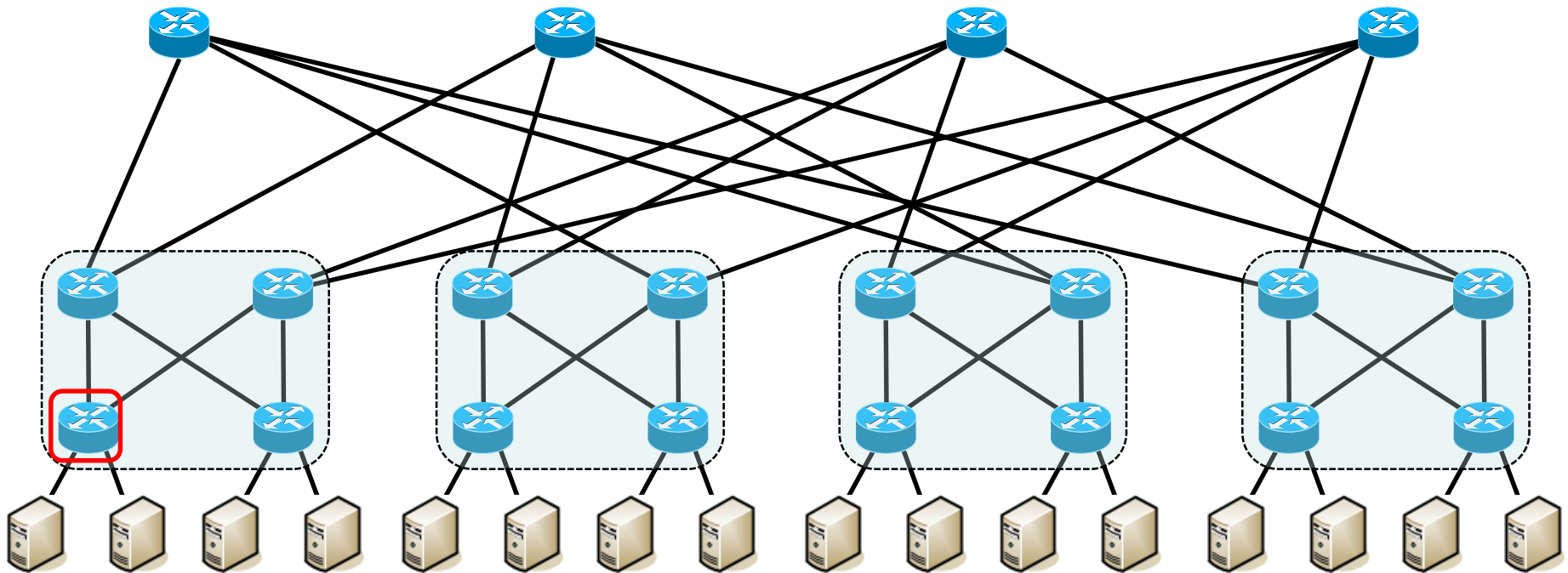


Switch ID	Pod Number	Position	Tree Level
A0-2B-FB-23-34-01			0





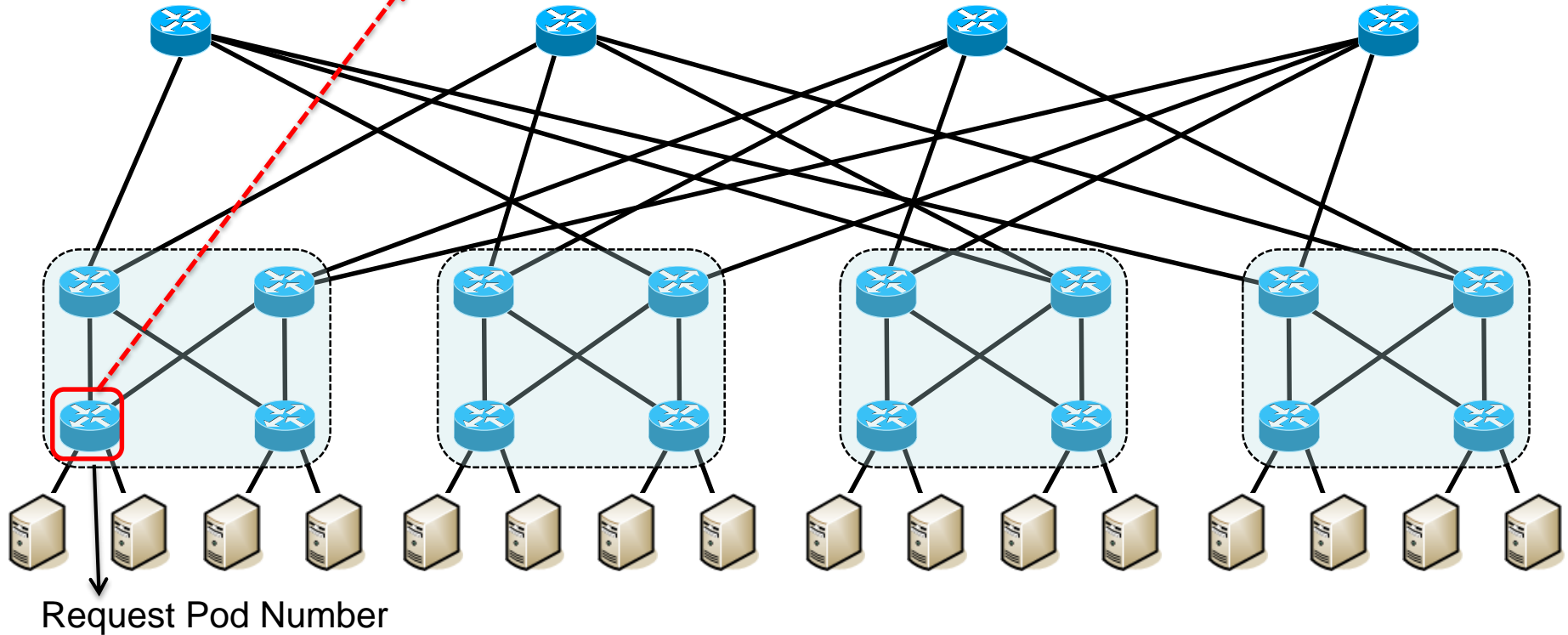
Fabric  
Manager



Switch ID	Pod Number	Position	Tree Level
A0-2B-FB-23-34-01		0	0



Fabric Manager



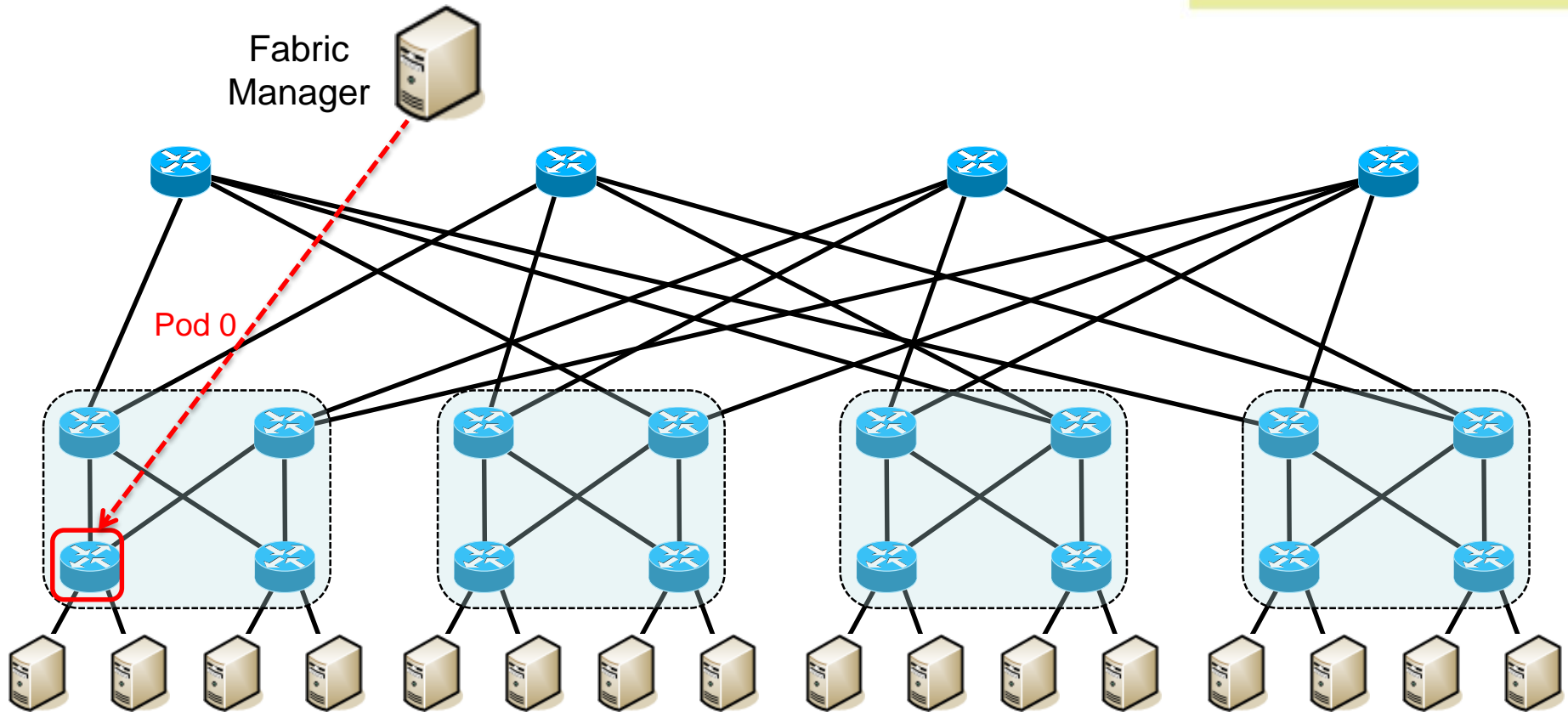
Switch ID	Pod Number	Position	Tree Level
A0-2B-FB-23-34-01		0	0



Fabric Manager



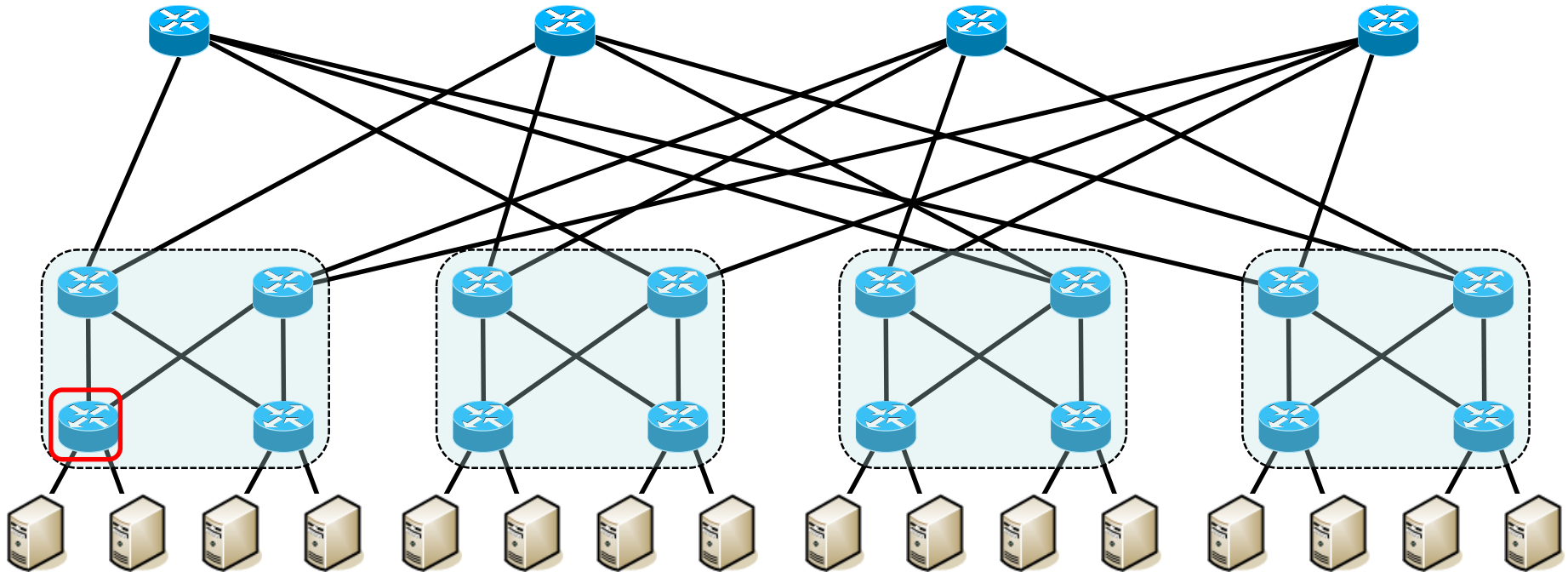
Pod 0



Switch ID	Pod Number	Position	Tree Level
A0-2B-FB-23-34-01		0	0



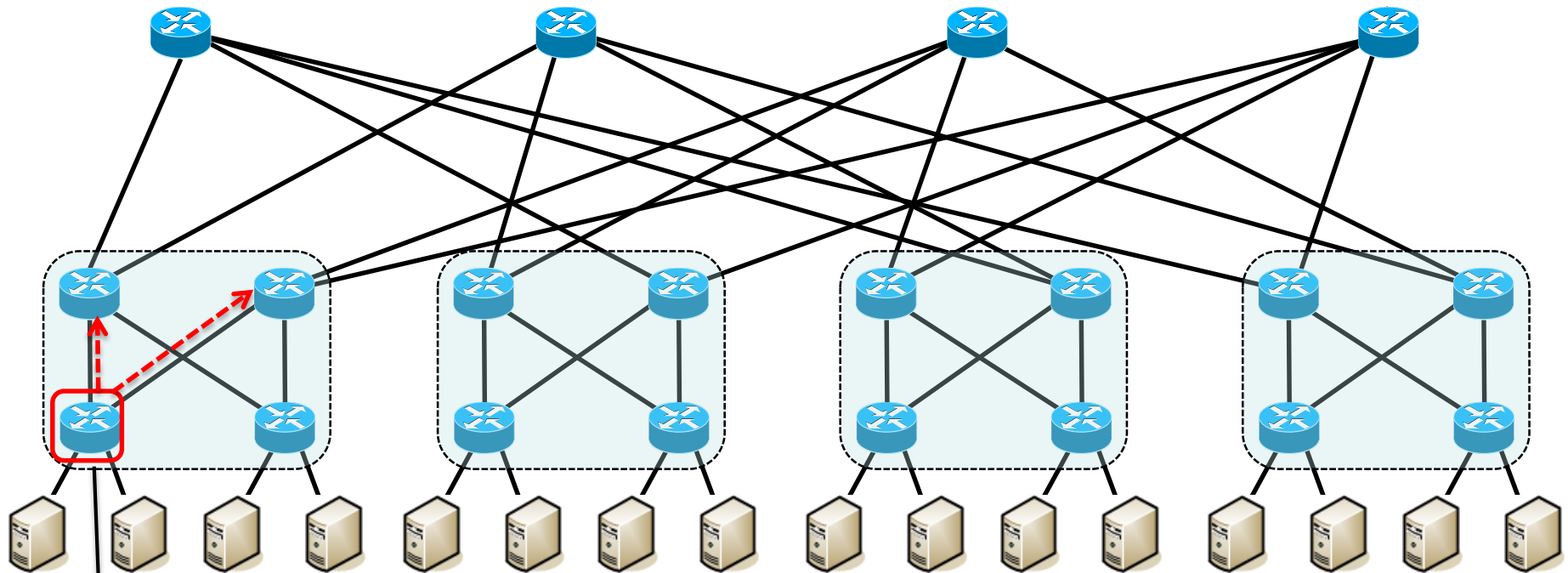
Fabric  
Manager



Switch ID	Pod Number	Position	Tree Level
A0-2B-FB-23-34-01	0	0	0



Fabric  
Manager

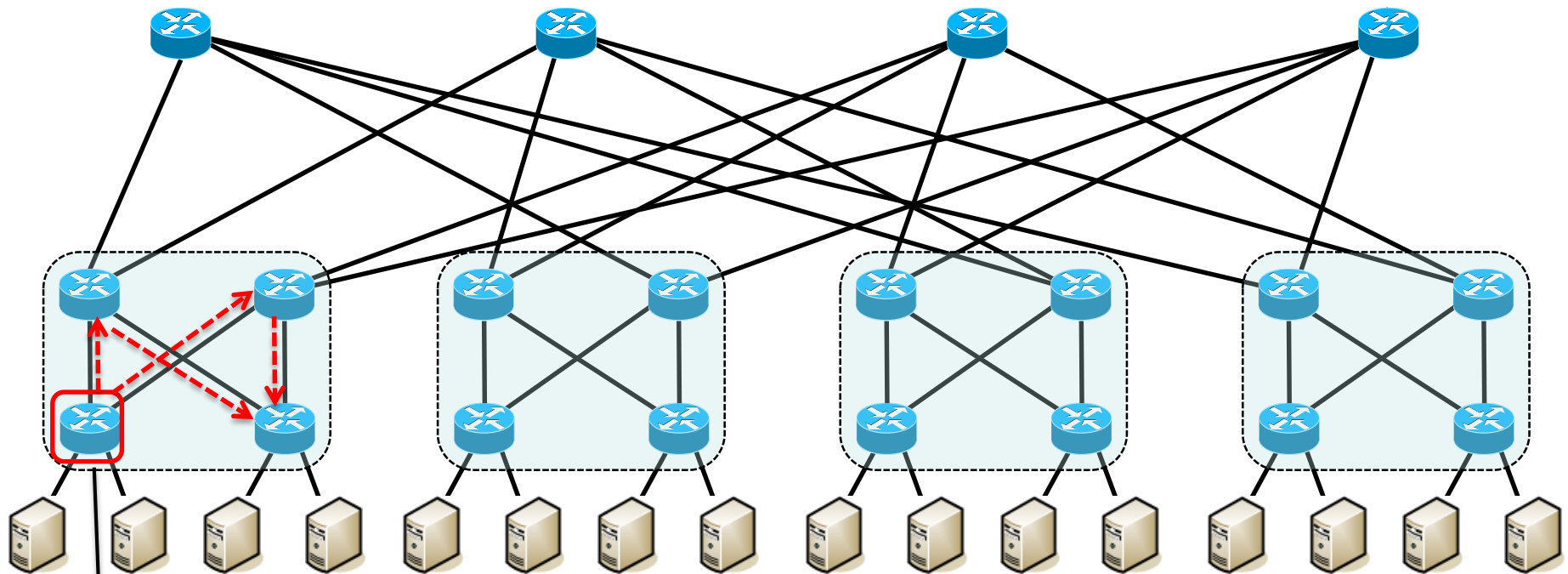


Distribute Pod Number

Switch ID	Pod Number	Position	Tree Level
A0-2B-FB-23-34-01	0	0	0



Fabric  
Manager



Distribute Pod Number

Switch ID	Pod Number	Position	Tree Level
A0-2B-FB-23-34-01	0	0	0



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