Complete the interpreter (code provided) to support procedures parametrized by expressions under call-by-name (lazy) semantics. The syntax, typing rules and semantic equations for this extension are found in Schmidt. You will need to extend  $meaning\_of\_declaration$  to properly handle procedure declarations with lazily evaluated  $\tau$ exp parameters that will be evaluated lazily.

The equations to implement are on page  $80^1$ .

```
[\![\pi \vdash \mathsf{proc}\ I_1(I_2 : \tau exp) = C : \{I_1 : \tau exp \to comm\} dec]\!] \ es = (\{I_1 : p\}, s)
where pfs' = [\![\pi \mapsto \{I_2 : \tau exp\} \vdash C : comm]\!] (e \mapsto \{I_2 : f\})s'
```

This rule becomes two rules when we specialize  $\tau$  to be int or bool as follows:

For the first rule, you update the proc\_int environment and for the second you must update the proc\_bool environment.

Note that for type assignments,  $\pi_1 \mapsto \pi_2$  is implemented by the following ML code: (bar\_union\_ta pi1 pi2) and for environments  $e_1 \mapsto e_2$  is implemented by the call (union\_env e1 e2).

For completeness I include the semantics of the call here as well (though it is implemented in the code provided).

$$[\![\pi \vdash \mathsf{call}\ I_1(E) : comm]\!] e s = p f s$$
  
where  $\langle I_1, p \rangle \in e$  and  $f s' = [\![\pi \vdash E : \tau exp]\!] e s$ 

$$p \in \llbracket \tau \exp \rrbracket \to \llbracket comm \rrbracket$$

The inclusion of the  $Env_{\pi}$  argument to p is incorrect. You can see this in the definition of the type env in the code.

<sup>&</sup>lt;sup>1</sup>Note that the line above the semantic equations on page 80 should say