Introduction to Operating Systems

Chapter 10: Multiple processors systems

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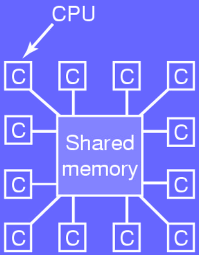
Fall 2017

Outline

1. Multiprocessor setup
2. Multiprocessor scheduling
3. Distributed systems

Shared memory model

- CPU communicate through the shared memory
- Every CPU has equal access to entire physical memory
- Access time: 2-10 ns



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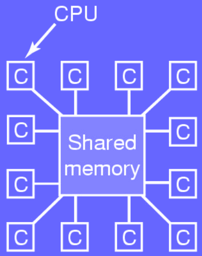

Three main appraoches:

- Each CPU has its own OS: no sharing, all independent
- Master-slave multiprocessors: one CPU handles all the requests
- Symmetric MultiProcessor: solution used in practice

SMP approach

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- Add a mutex when entering those regions
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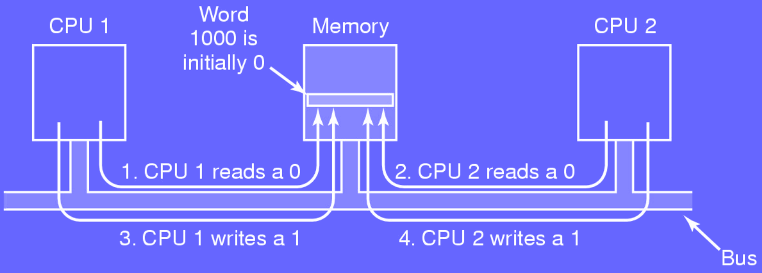
Challenges:

- How to divide up the OS
- Easy to run into deadlock with the shared tables
- On the long run, hard to keep consistency between programmers

Synchronisation strategy with a single CPU:

Synchronisation strategy with a single CPU: `TSL` instruction



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1. Lock the bus by asserting a special line on the bus
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New multiprocessor issue: slows down the whole system

Solution: use a local cache such as not to block the bus to often

Cache with multiprocessors

Multiprocessors cache implementation:

- Requesting CPU reads the lock and gets a copy in its cache
- Polling is done using the value in cache
- When the lock is remove the cache protocol invalidates all the remote copies
- Updated value is to be fetched

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Problem:

- Mutex is 1 bit, but a whole block is copied
- TSL require write access
- Cache block that is modified is invalidated
- Cache need to be recopied

Read and TSL

New approach:

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Problem: what if two CPUs see the lock being free and apply the `TSL` instruction? Does it lead to a race condition?

- The value returned by the read is only a hint
- Only one CPU gets the lock
- The `TSL` instruction prevent any race condition

More solutions

Ethernet binary exponential backoff algorithm:

- Do not poll at regular intervals
- Add a loop where waiting time is doubled at each iteration
- Setup a maximum waiting time

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Set a mutex for each CPU requesting the lock:

- When a CPU request a lock it attaches itself at the end of the list of CPUs requesting the lock
- When the initial lock is released it frees the lock of the first CPU on the list
- The first CPU enters the critical region, does its work and releases the lock
- Next CPU on the list can start its work

Spinning or switching?

New perspective:

- On a uniprocessor: the time spent on waiting is wasted
- On a multiprocessor: one CPU is waiting while another works

Dilemma: switching is expensive but looping is a waste

Best choice:

- Only know which solution was best after
- Impossible to have an always accurate optimal decision

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Solution: mix of waiting and switching with a (variable) threshold

Outline

Setup

Single threaded process: only need to schedule the process

Uniprocessor: which thread to run?

Multiprocessor:

- Which thread to run?
- Which CPU to run it on?

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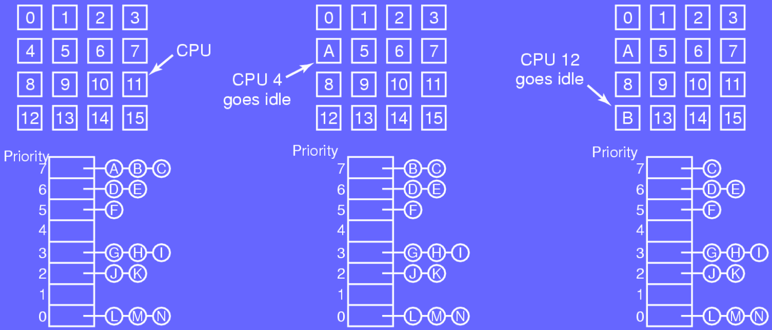
Multiprocessor:

- Which thread to run?
- Which CPU to run it on?
- Threads of a process run on the same CPU $\rightarrow$ no need to reload the whole process
- Threads of a process run in parallel $\rightarrow$ threads can cooperate more efficiently

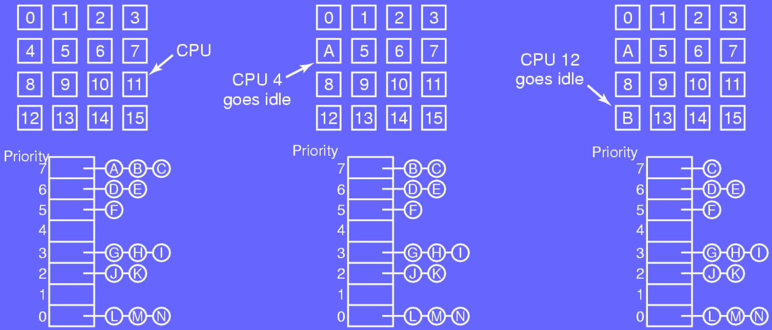
Timesharing



- Single data structure for ready processes
- Simple and efficient implementation

Limitation: a thread holds a spin lock but reaches the end of its quantum

Timesharing



- Single data structure for ready processes
- Simple and efficient implementation

Limitation: a thread holds a spin lock but reaches the end of its quantum → other CPUs spin waiting for lock to be released

Improvements on timesharing

Smart scheduling:

- A thread holding a spin lock sets a flag
- Scheduler lets such thread run after the end of the quantum
- Clear the flag when lock is released

Affinity scheduling:

- Thread is assigned a CPU when it is created
- Try as much as possible to run a thread on the same CPU
- Cache affinity is maximized

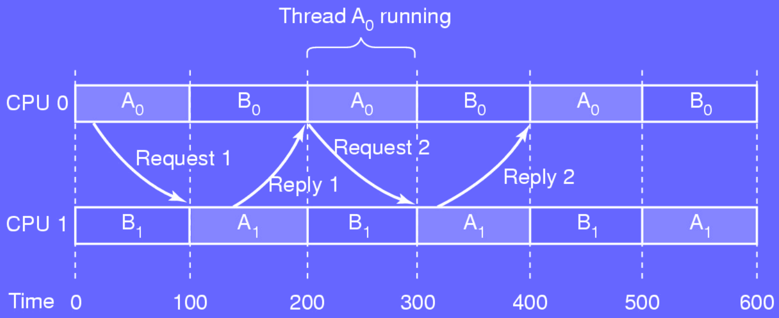
Spacesharing

When a process is created the scheduler checks if the number of free CPUs is larger then the number of threads:

- Yes: run one thread per CPU until completion
- No: wait for more free CPUs
- Keep the CPU even during I/O

Optimization: try to apply shortest job first $\rightarrow$ in practice FIFO better

The communication problem



Gang scheduling

Simple idea: schedule processes by group

- Group related threads into a gang
- All gang members run simultaneously on different CPUs
- All members start and end at the same time
- No intermediary scheduling decision is taken

Limitations: what if one gang member issue an I/O request or finishes earlier than others?

Outline

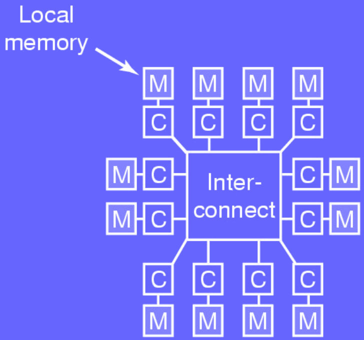
Main characteristics

Characteristics of distributed systems:

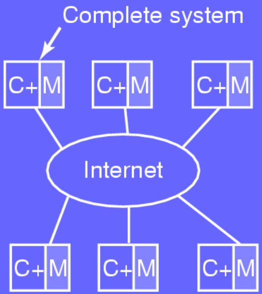
- Composed of autonomous entities, each one with its own memory
- Communication is done over a network using message passing
- The system must tolerate node failures
- All the nodes perform the same task

Cluster: set of connected computers working together

Structure



Multicomputer



Distributed system

Setup and challenges

General idea of how a cluster works:

- Computing nodes connected over a LAN
- A clustering middleware sits on the top of the node
- Users view a large computer

Example.
A single master node handles the management of the scheduling and slave nodes

Main challenges:

- Scheduling: where should a job be scheduled?
- Load balancing: should a job be rescheduled on a different node?

Example

Apache Hadoop:

- Opensource framework for distributed file system
- Written in Java
- Very large files stored across multiple nodes
- Used and enhanced by Yahoo!, Facebook, Amazon, Microsoft, Google, IBM...

More on distributed computing

Advances in network technologies lead to the development of:

- Volunteer computing: volunteer offer part of their computational power to some project

- Grids: collection of computer resources from multiple locations to reach a common goal

- Jungle computing: network not necessarily composed of "regular computers"

Key points

- What is the main difference between a distributed system and a multiprocessor system?
- Explain the SMP approach
- How is `TSL` working for multiprocessors?
- Cite three scheduling strategies targeting multiprocessors
- What is a middleware?

Thank you!
Enjoy the winter break. . .