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## Security goals

### (a) Possible violated security goals

- Confidentiality: The attacker may have read private documents, letters, photos or looked through her computer.
- Integrity: Items could have been modified or damaged, like something in documents.
- Availability: Objects might have been stolen or left unusable, like the attacker formatted the hard disk of her laptop.

### (b) Security mechanisms from different classes

- prevention: security door, bars on windows, a safe for valuable items.
- detection: Alarm system, surveillance cameras.
- Analysis: evaluating the damage after the break in would help to understand what happened and help with recovery.

## Simple combinatorics

### (a) ROT13

ROT13 uses a fixed rotation of 13 positions. So the key space consists of exactly one element.

### (b) Vigenère cipher with known key length $n$

If the alphabet consists of 26 letters. Each of the  $n$  key positions can be chosen independently from 26 possibilities so it would be  $26^n$ .

### (c) AES with a 256 bit key

An AES 256 key has length of 256 bits. Each bit can be 0 or 1, independently, so it would be  $2^{256}$

### (d) Monoalphabetic substitution cipher with $k$ letters

A monoalphabetic substitution over an alphabet of size  $k$  is a permutation of these  $k$  letters. so it would be  $k!$  and if the alphabet consists of 26 letters it would be  $26!$ .

## XOR

Alice sends two ciphertexts using the same XOR key  $k$ , so in other words:

$$c_0 = m_0 \oplus k, \quad c_1 = m_1 \oplus k,$$

and Eve captures  $c_0$  and  $c_1$ . Eve knows that one of the plaintexts  $m$  is either  $m_0$  or  $m_1$ .

## How can Eve recover the other plaintext message?

First we look at:

$$c_0 \oplus c_1 = (m_0 \oplus k) \oplus (m_1 \oplus k) = m_0 \oplus m_1,$$

and that is because  $k \oplus k = 0$ .

If Eve correctly guesses one plaintext (let's say  $m_0$ ), she can compute the other as:

$$m_1 = m_0 \oplus (m_0 \oplus m_1) = m_0 \oplus (c_0 \oplus c_1).$$

The same goes for  $m_1$ . So once one plaintext is known, the other plaintext is easy to derive without knowing the key  $k$ .

## Can Eve also recover the exact key $k$ used by Alice?

If she knows that  $c_0$  encrypts  $m_0$ , then she can also do:

$$k = c_0 \oplus m_0.$$

So:

- With only  $c_0$  and  $c_1$  but she doesn't know which plaintext was sent, then she cannot know the correct key.
- As soon as she knows one correct plaintext/ciphertext pair (maybe from guessing), she can compute  $k$  and then decrypt any other ciphertext encrypted with the same key.