南方科技大学
SOUTHERN UNIVERSITY OF SCIENCE AND TECHNOLOGY

CS201 DISCRETE MATHEMATICS FOR COMPUTER SCIENCE

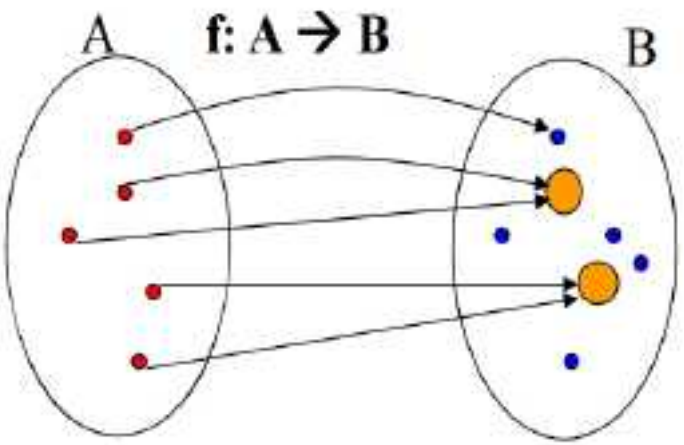
Dr. Qi Wang

Department of Computer Science and Engineering
Office: Room413, CoE South Tower
Email: wangqi@sustech.edu.cn

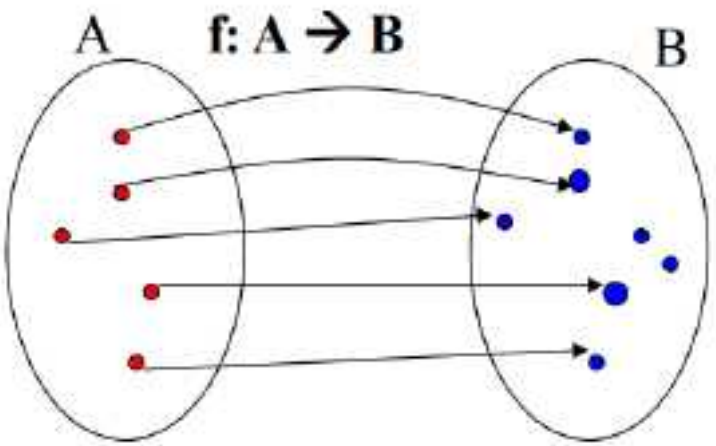
Injective (One-to-One) Function

- A function f is called *one-to-one* or *injective*, if and only if $f(x) = f(y)$ implies $x = y$ for all x, y in the domain of f. In this case, f is called an *injection*.

  Alternatively: A function is *one-to-one* if and only if $f(x) \neq f(y)$ whenever $x \neq y$. (contrapositive!)



Not injective



Injective function

Injective Functions

- **Example 1**:

  $A = \{1, 2, 3\}$, $B = \{a, b, c\}$

  Define f as

  – $1 \mapsto c$

  – $2 \mapsto a$

  – $3 \mapsto c$

  Is f one-to-one?

Injective Functions

- **Example 1**:

  $A = \{1, 2, 3\}$, $B = \{a, b, c\}$

  Define f as

  – $1 \mapsto c$

  – $2 \mapsto a$

  – $3 \mapsto c$

  Is f one-to-one?

- **Example 2**:

  Let $g : \mathbf{Z} \to \mathbf{Z}$, where $g(x) = 2x - 1$

  Is g one-to-one?

Surjective (Onto) Function

- A function f is called *onto* or *surjective*, if and only if for every $b \in B$ there is an element $a \in A$ such that $f(a) = b$. In this case, f is called a *surjection*.

Surjective (Onto) Function

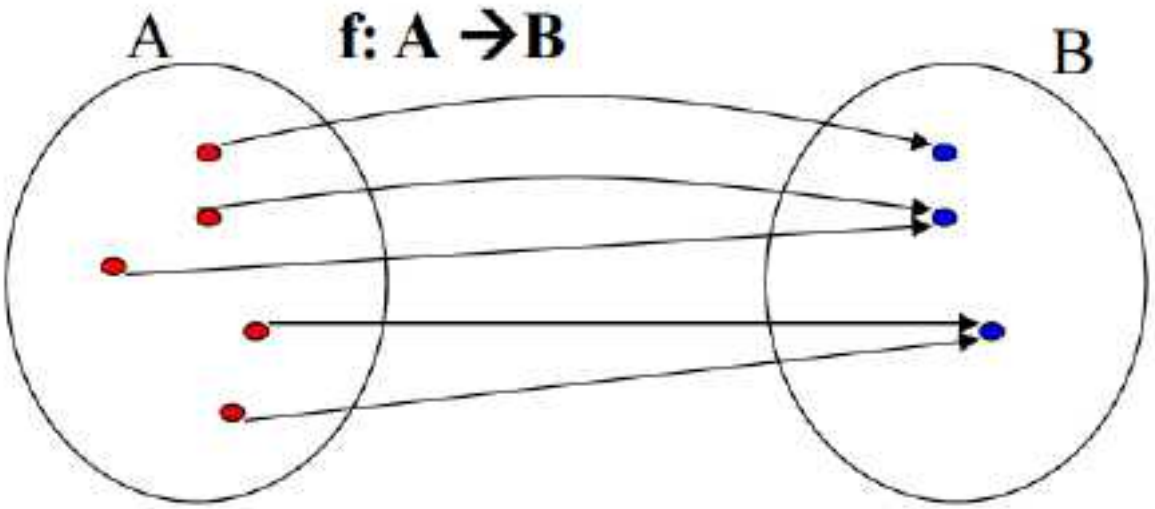
- A function f is called *onto* or *surjective*, if and only if for every $b \in B$ there is an element $a \in A$ such that $f(a) = b$. In this case, f is called a *surjection*.

  Alternatively: A function is *onto* if and only if all codomain elements are covered ($f(A) = B$).

Surjective (Onto) Function

- A function f is called *onto* or *surjective*, if and only if for every $b \in B$ there is an element $a \in A$ such that $f(a) = b$. In this case, f is called a *surjection*.

  Alternatively: A function is *onto* if and only if all codomain elements are covered ($f(A) = B$).



Surjective Functions

- **Example 1**:

  $A = \{1, 2, 3\}$, $B = \{a, b, c\}$

  Define f as

  – $1 \mapsto c$

  – $2 \mapsto a$

  – $3 \mapsto c$

  Is f onto?

Surjective Functions

- **Example 1**:

  $A = \{1, 2, 3\}$, $B = \{a, b, c\}$
  Define f as
  – $1 \mapsto c$
  – $2 \mapsto a$
  – $3 \mapsto c$

  Is f onto?

- **Example 2**:

  Let $A = \{0, 1, \ldots, 9\}$, $B = \{0, 1, 2\}$
  Define $h : A \to B$ as $h(x) = x \bmod 3$.

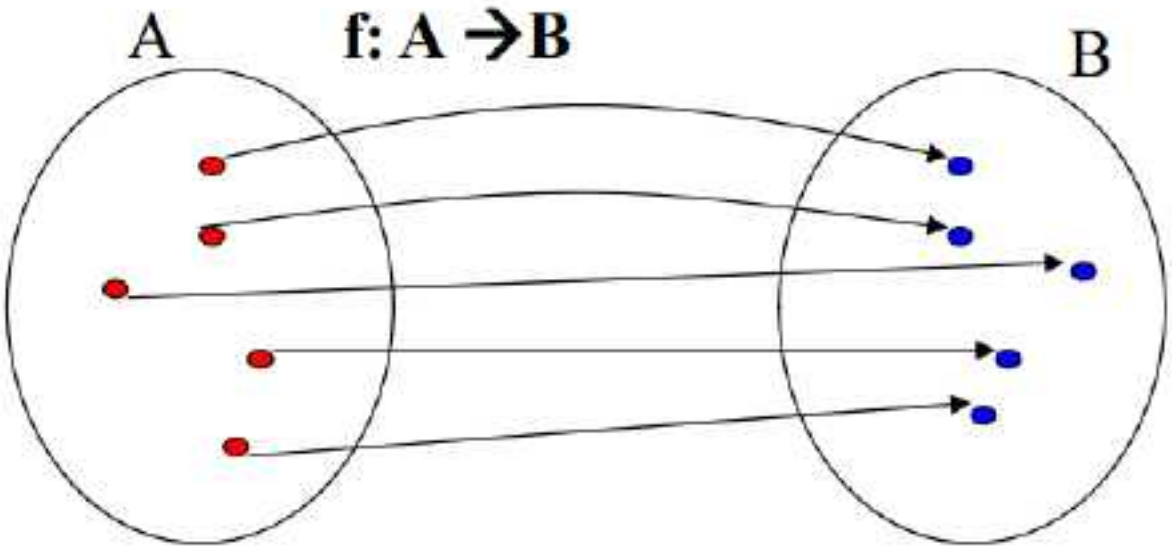
  Is h onto?

Bijective Function (One-to-One Correspondence)

- A function f is called *bijective*, if and only if it is both one-to-one and onto.

Bijective Function (One-to-One Correspondence)

- A function f is called *bijective*, if and only if it is both one-to-one and onto.



Bijective Functions

- **Example 1**:

  $A = \{1, 2, 3\}$, $B = \{a, b, c\}$

  Define f as

  – $1 \mapsto c$

  – $2 \mapsto a$

  – $3 \mapsto b$

  Is f bijective?

Bijective Functions

- **Example 1**:

  $A = \{1, 2, 3\}$, $B = \{a, b, c\}$

  Define f as

  – $1 \mapsto c$

  – $2 \mapsto a$

  – $3 \mapsto b$

  Is f bijective?

- **Example 2**:

  Define $g : \mathbf{N} \to \mathbf{N}$ as $g(x) = \lfloor \frac{x}{2} \rfloor$ (floor function).

  Is g bijective?

Summary

- Suppose that $f : A \to B$.

To show that f is *injective*	Show that if $f(x) = f(y)$ for all $x, y \in A$, then $x = y$
To show that f is not *injective*	Find specific elements $x, y \in A$ such that $x \neq y$ and $f(x) = f(y)$
To show that f is *surjective*	Consider an arbitrary element $y \in B$ and find an element $x \in A$ such that $f(x) = y$
To show that f is not *surjective*	Find a specific element $y \in B$ such that $f(x) \neq y$ for all $x \in A$

Note

- Prove that "for a function $f : A \to B$ with $|A| = |B| = n$, f is one-to-one if and only if f is onto."

Note

- Prove that "for a function $f : A \to B$ with $|A| = |B| = n$, f is one-to-one if and only if f is onto."

Proof.

$\diamond$ only if part: Suppose that f is one-to-one. Let $\{x_1, x_2, \ldots, x_n\}$ be elements of A. Then $f(x_i) \neq f(x_j)$ for $i \neq j$. Therefore, $|f(A)| = |\{f(x_1), \ldots, f(x_n)\}| = n$. But $|B| = n$ and $f(A) \subseteq B$. Therefore, $f(A) = B$.

$\diamond$ if part: Suppose that f is onto. Let $A = \{x_1, x_2, \ldots, x_n\}$ be a listing of the elements of A. Suppose that $f(x_i) = f(x_j)$ for some $i \neq j$. Then, $|\{f(x_1), \ldots, f(x_n)\}| \leq n - 1$. But $|f(A)| = |B| = n$, a contradiction.

Bijective Functions

- "For a function $f : A \to B$ with $|A| = |B| = n$, f is one-to-one if and only if f is onto."

Bijective Functions

- "For a function $f : A \to B$ with $|A| = |B| = n$, f is one-to-one if and only if f is onto."
- "For a function f from A to itself, f is one-to-one if and only if f is onto, where A is infinite."

Bijective Functions

- "For a function $f : A \to B$ with $|A| = |B| = n$, f is one-to-one if and only if f is onto."

- "For a function f from A to itself, f is one-to-one if and only if f is onto, where A is infinite."

  Counterexample:

  $f : \mathbf{Z} \to \mathbf{Z}$, where $f(x) = 2x$.

  f is one-to-one but not onto

  – $1 \mapsto 2$
  – $2 \mapsto 4$
  – $3 \mapsto 6$

  3 has no preimage.

Two Functions on Real Numbers

- Let f_1 and f_2 be functions from A to $\mathbf{R}$. Then $f_1 + f_2$ and $f_1 f_2$ are also functions form A to $\mathbf{R}$ defined for all $x \in A$ by

$$(f_1 + f_2)(x) = f_1(x) + f_2(x)$$
$$(f_1 f_2)(x) = f_1(x) f_2(x)$$

Two Functions on Real Numbers

- Let f_1 and f_2 be functions from A to $\mathbf{R}$. Then $f_1 + f_2$ and $f_1 f_2$ are also functions form A to $\mathbf{R}$ defined for all $x \in A$ by

$$(f_1 + f_2)(x) = f_1(x) + f_2(x)$$
$$(f_1 f_2)(x) = f_1(x) f_2(x)$$

Example:

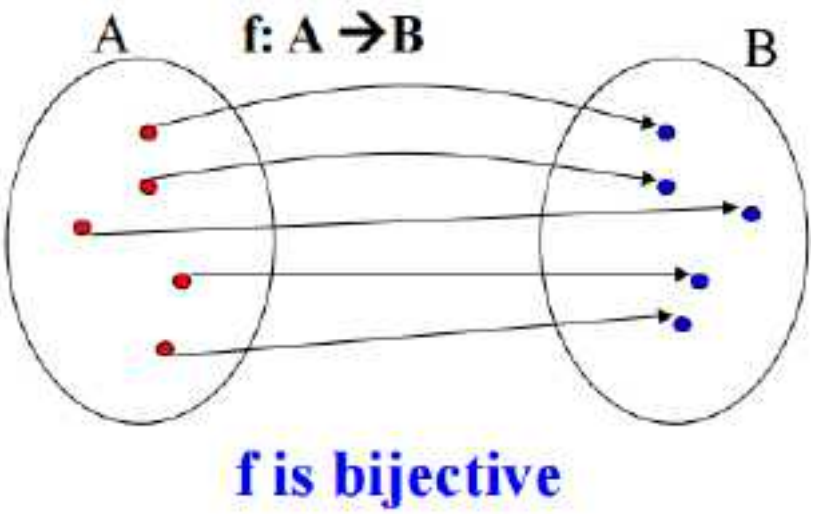
$f_1 = x - 1$ and $f_2 = x^3 + 1$

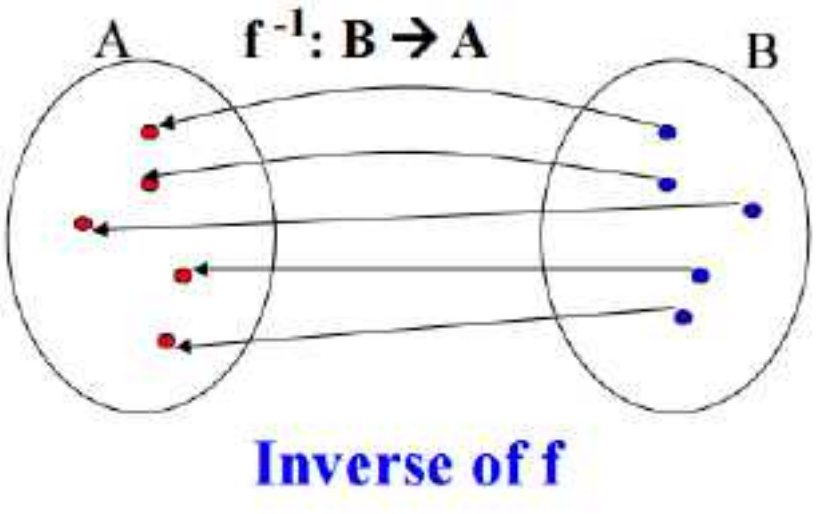
Then

$$(f_1 + f_2)(x) = x^3 + x$$
$$(f_1 f_2)(x) = x^4 - x^3 + x - 1$$

Inverse Functions

- Let $f : A \to B$ be a bijection. The *inverse of f* is the function that assigns to an element b belonging to B the unique element a in A such that $f(a) = b$, denoted by f^{-1}. Hence, $f^{-1}(b) = a$ when $f(a) = b$. In this case, f is called *invertible*.





Inverse Functions

- Note: if f is not a bijection, it is impossible to define the inverse function of f. Why ?

Inverse Functions

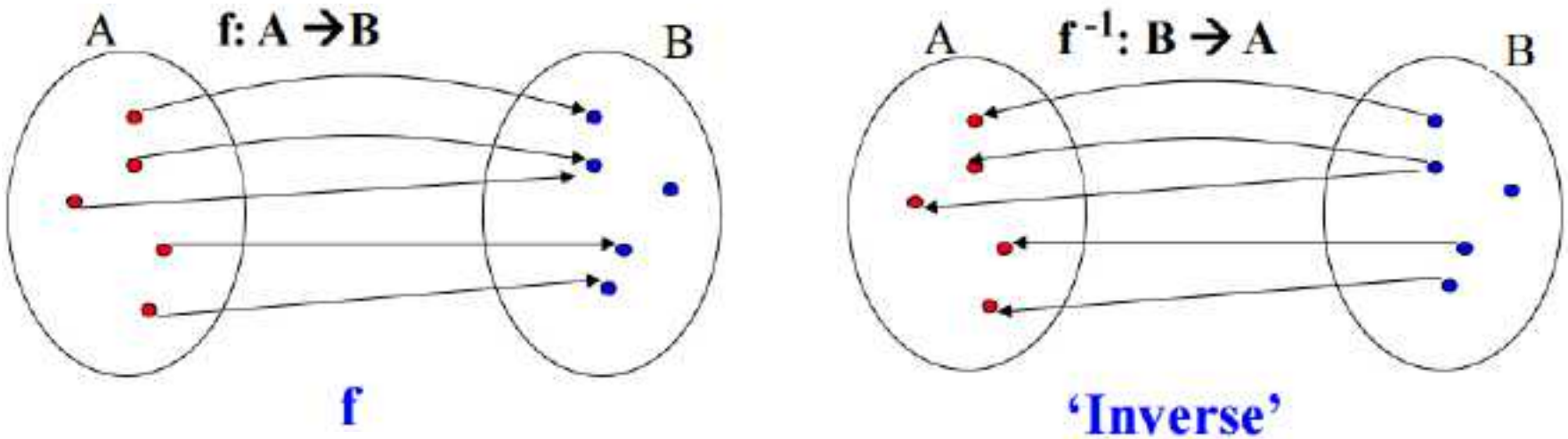
- Note: if f is not a bijection, it is impossible to define the inverse function of f. Why ?

  Assume f is not injective:

Inverse Functions

- Note: if f is not a bijection, it is impossible to define the inverse function of f. Why ?

Assume f is not injective:



The inverse is not a function: one element of B is mapped to two different elements of A

Inverse Functions

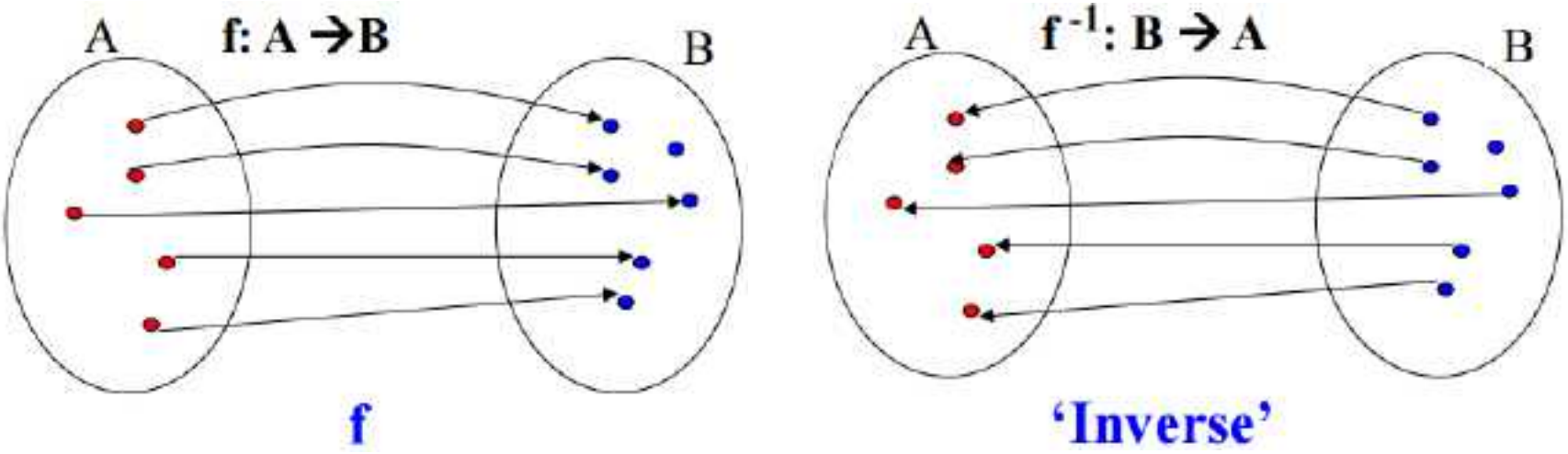
- Note: if f is not a bijection, it is impossible to define the inverse function of f. Why ?

  Assume f is not surjective:

Inverse Functions

- Note: if f is not a bijection, it is impossible to define the inverse function of f. Why ?

Assume f is not surjective:



The inverse is not a function: one element of B is not assigned an element of A

Inverse Functions

- **Example 1**:

  $f : \mathbf{R} \to \mathbf{R}$, where $f(x) = 2x - 1$.

  What is the inverse function f^{-1}?

Inverse Functions

- **Example 1**:

  $f : \mathbf{R} \to \mathbf{R}$, where $f(x) = 2x - 1$.

  What is the inverse function f^{-1}?

  $f^{-1}(x) = (x + 1)/2$

Inverse Functions

- **Example 1**:

  $f : \mathbf{R} \to \mathbf{R}$, where $f(x) = 2x - 1$.

  What is the inverse function f^{-1}?

  $f^{-1}(x) = (x + 1)/2$

- **Example 2**:

  $f : \mathbf{Z} \to \mathbf{Z}$, where $f(x) = 2x - 1$.

  Is f invertible?

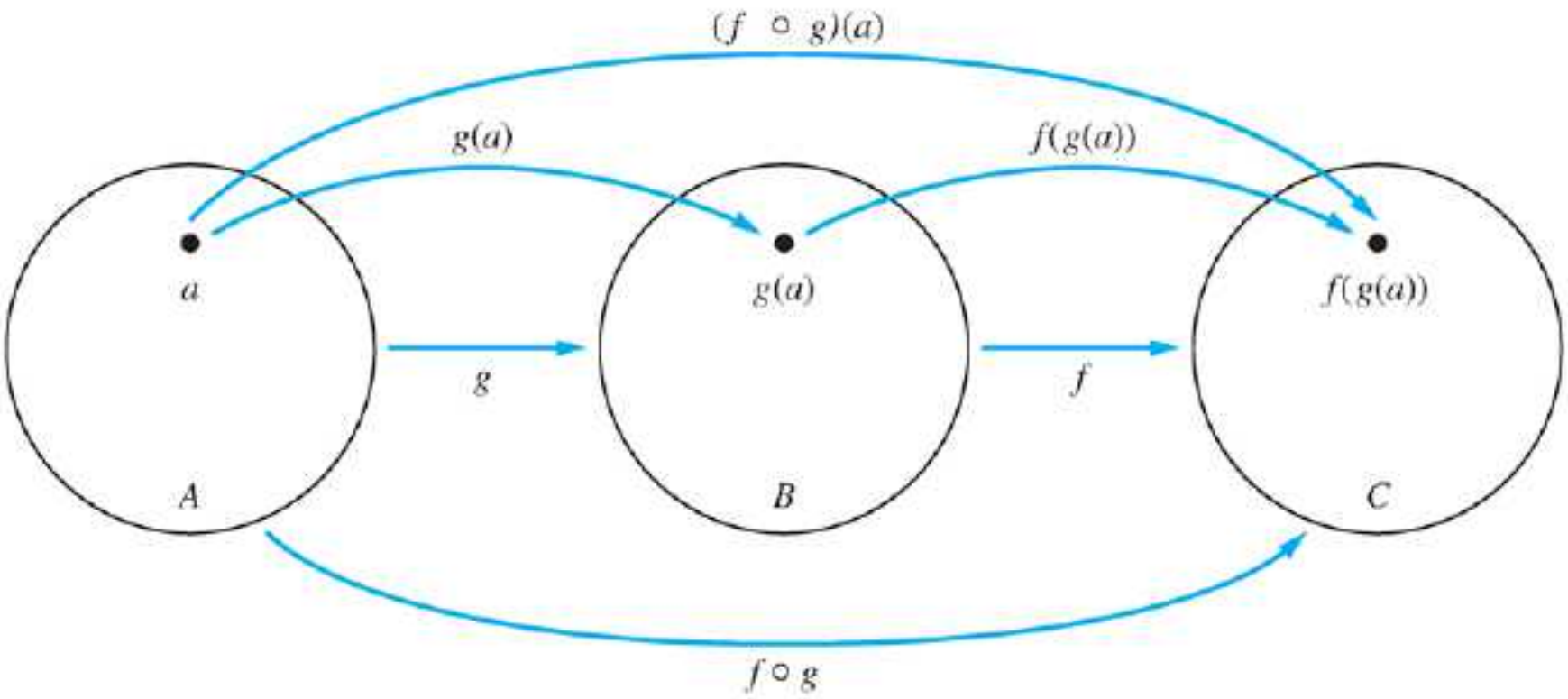
  No, since f is not onto.

Composition of Functions

- Let f be a function from B to C and let g be a function from A to B. The *composition of the functions f and g*, denoted by $f \circ g$, is defined by $(f \circ g)(x) = f(g(x))$.

Composition of Functions

- Let f be a function from B to C and let g be a function from A to B. The *composition of the functions f and g*, denoted by $f \circ g$, is defined by $(f \circ g)(x) = f(g(x))$.



Composition of Functions

- **Example 1**:

  Let $A = \{1, 2, 3\}$ and $B = \{a, b, c, d\}$.

  $g : A \to A$ $\quad$ $f : A \to B$

  $1 \mapsto 3$ $\quad$ $1 \mapsto b$

  $2 \mapsto 1$ $\quad$ $2 \mapsto a$

  $3 \mapsto 2$ $\quad$ $3 \mapsto d$

  What is $f \circ g$?

Composition of Functions

- **Example 1**:

  Let $A = \{1, 2, 3\}$ and $B = \{a, b, c, d\}$.

  $g : A \to A$ $\quad$ $f : A \to B$

  $1 \mapsto 3$ $\quad$ $1 \mapsto b$

  $2 \mapsto 1$ $\quad$ $2 \mapsto a$

  $3 \mapsto 2$ $\quad$ $3 \mapsto d$

  What is $f \circ g$?

  $f \circ g : A \to B$

  $1 \mapsto d$

  $2 \mapsto b$

  $3 \mapsto a$

Composition of Functions

- **Example 2**:

  Let $f : \mathbf{Z} \to \mathbf{Z}$ and $g : \mathbf{Z} \to \mathbf{Z}$, where $f(x) = 2x$ and $g(x) = x^2$.

  What are $g \circ f$ and $f \circ g$?

Composition of Functions

- **Example 2**:

  Let $f : \mathbf{Z} \to \mathbf{Z}$ and $g : \mathbf{Z} \to \mathbf{Z}$, where $f(x) = 2x$ and $g(x) = x^2$.

  What are $g \circ f$ and $f \circ g$?

  $g \circ f : \mathbf{Z} \to \mathbf{Z}$ $\quad g \circ f = 4x^2$

  $f \circ g : \mathbf{Z} \to \mathbf{Z}$ $\quad f \circ g = 2x^2$

Composition of Functions

- **Example 2**:

  Let $f : \mathbf{Z} \to \mathbf{Z}$ and $g : \mathbf{Z} \to \mathbf{Z}$, where $f(x) = 2x$ and $g(x) = x^2$.

  What are $g \circ f$ and $f \circ g$?

  $g \circ f : \mathbf{Z} \to \mathbf{Z}$ $\quad g \circ f = 4x^2$

  $f \circ g : \mathbf{Z} \to \mathbf{Z}$ $\quad f \circ g = 2x^2$

  Note: In general, the order of composition matters.

Composition of Functions

- Suppose that f is a bijection from A to B. Then $f \circ f^{-1} = I_B$ and $f^{-1} \circ f = I_A$, Since

$$(f^{-1} \circ f)(a) = f^{-1}(f(a)) = f^{-1}(b) = a$$
$$(f \circ f^{-1})(b) = f(f^{-1}(b)) = f(a) = b,$$

  where I_A, I_B denote the *identity functions* on the sets A and B, respectively.

Some Important Functions

- The *floor function* assigns a real number x the largest integer that is $\leq x$, denoted by $\lfloor x \rfloor$.
- The *ceiling function* assigns a real number x the smallest integer that is $\geq x$, denoted by $\lceil x \rceil$.

Some Important Functions

- The *floor function* assigns a real number x the largest integer that is $\leq x$, denoted by $\lfloor x \rfloor$.
- The *ceiling function* assigns a real number x the smallest integer that is $\geq x$, denoted by $\lceil x \rceil$.

TABLE 1 Useful Properties of the Floor and Ceiling Functions.
(n is an integer, x is a real number)

(1a)	$\lfloor x \rfloor = n$ if and only if $n \leq x < n+1$
(1b)	$\lceil x \rceil = n$ if and only if $n-1 < x \leq n$
(1c)	$\lfloor x \rfloor = n$ if and only if $x-1 < n \leq x$
(1d)	$\lceil x \rceil = n$ if and only if $x \leq n < x+1$
(2)	$x-1 < \lfloor x \rfloor \leq x \leq \lceil x \rceil < x+1$
(3a)	$\lfloor -x \rfloor = -\lceil x \rceil$
(3b)	$\lceil -x \rceil = -\lfloor x \rfloor$
(4a)	$\lfloor x+n \rfloor = \lfloor x \rfloor + n$
(4b)	$\lceil x+n \rceil = \lceil x \rceil + n$

Some Important Functions

Ex. 1: Prove or disprove that if x is a real number, then $\lfloor 2x \rfloor = \lfloor x \rfloor + \lfloor x + \frac{1}{2} \rfloor$.

Ex. 2: Prove or disprove that $\lceil x + y \rceil = \lceil x \rceil + \lceil y \rceil$ for all real numbers x and y.

Some Important Functions

Ex. 1: Prove or disprove that if x is a real number, then $\lfloor 2x \rfloor = \lfloor x \rfloor + \lfloor x + \frac{1}{2} \rfloor$.

Ex. 2: Prove or disprove that $\lceil x + y \rceil = \lceil x \rceil + \lceil y \rceil$ for all real numbers x and y.

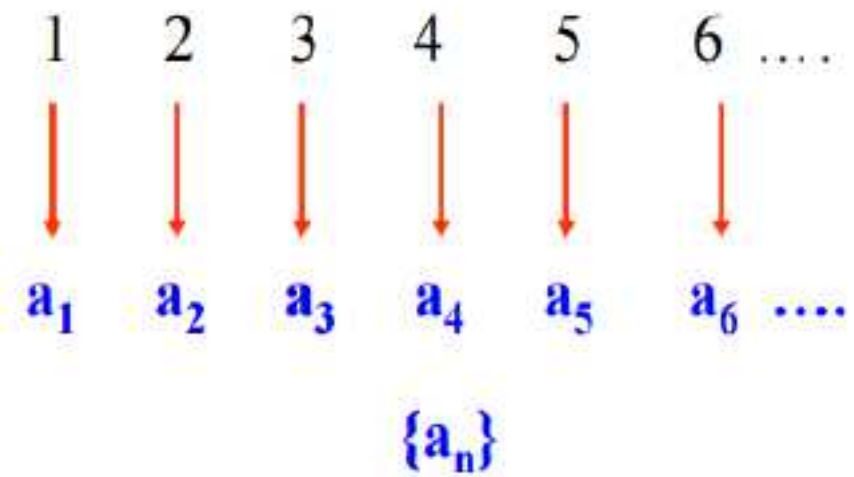
- The factorial function $f : \mathbf{N} \to \mathbf{Z}^+$ is the product of the first n positive integers when n is a nonnegative integer, denoted by $f(n) = n!$.

Sequences

- A *sequence* is a function from a subset of the set of integers (typically the set $\{0, 1, 2, \ldots\}$ or $\{1, 2, 3, \ldots\}$ to a set S. We use the notation a_n to denote the image of the integer n. ($\{a_n\}$ represents the ordered list $a_1, a_2, a_3, \ldots$)

Sequences

- A *sequence* is a function from a subset of the set of integers (typically the set $\{0, 1, 2, \ldots\}$ or $\{1, 2, 3, \ldots\}$ to a set S. We use the notation a_n to denote the image of the integer n. ($\{a_n\}$ represents the ordered list $a_1, a_2, a_3, \ldots$)



Sequences

1.1 Basic Concepts and Notation

In general, a *sequence* is an ordered list of elements from a set S. Formally, a *finite sequence* with elements over S is a function from the index set $\{0, 1, \ldots, N-1\}$ to S for some integer $N \geq 0$, and N is called the *length* of the sequence. An *infinite sequence* with elements over S is a function from the integer group $\mathbf{Z}$ to S, and a *semi-infinite* sequence with elements over S is a function from the semi-group $\{0, 1, \ldots\}$ to S. If the set S is a finite field $\mathbb{F}_q$ with q elements, we say that the sequence is a *q-ary sequence* over $\mathbb{F}_q$. In particular, if $S = \mathrm{GF}(2)$, the sequence is called a *binary* sequence.

For a sequence $\mathbf{s} = (s_i)_{i \geq 0}$, if there exist integers $r > 0$ and $u \geq 0$ such that

$$s_{i+r} = s_i \quad \text{for all } i \geq u, \tag{1.1}$$

the sequence is said to be *ultimately periodic* with parameters (r, u), and r is called a *period* of the sequence $\mathbf{s}$. The smallest number r satisfying (1.1) is called the *least period*

Sequences

- **Examples**:
  - $a_n = n^2$, where $n = 1, 2, 3, \ldots$
  - $a_n = (-1)^n$, where $n = 0, 1, 2, \ldots$
  - $a_n = 2^n$, where $n = 0, 1, 2, \ldots$

Sequences

- **Examples**:
  - $a_n = n^2$, where $n = 1, 2, 3, \ldots$
  - $a_n = (-1)^n$, where $n = 0, 1, 2, \ldots$
  - $a_n = 2^n$, where $n = 0, 1, 2, \ldots$

- An *arithmetic progression* is a sequence of the form $a, a+d, a+2d, a+3d, \ldots, a+nd, \ldots$, where the *initial term* a and *common difference* d are real numbers.

Sequences

- **Examples**:
  - $a_n = n^2$, where $n = 1, 2, 3, \ldots$
  - $a_n = (-1)^n$, where $n = 0, 1, 2, \ldots$
  - $a_n = 2^n$, where $n = 0, 1, 2, \ldots$

- An *arithmetic progression* is a sequence of the form $a, a+d, a+2d, a+3d, \ldots, a+nd, \ldots$, where the *initial term* a and *common difference* d are real numbers.

  Example:
  - $a_n = -1 + 4n$, where $n = 0, 1, 2, 3, \ldots$

Geometric Progression

- A geometric progression is a sequence of the form $a, ar, ar^2, \ldots, ar^n, \ldots$, where the *initial term* a and the *common ratio* r are real numbers.

Geometric Progression

- A geometric progression is a sequence of the form $a, ar, ar^2, \ldots, ar^n, \ldots$, where the *initial term* a and the *common ratio* r are real numbers.

  Example:

  - $a_n = (1/2)^n$, where $n = 0, 1, 2, 3, \ldots$

Geometric Progression

- A geometric progression is a sequence of the form $a, ar, ar^2, \ldots, ar^n, \ldots$, where the *initial term* a and the *common ratio* r are real numbers.

  Example:

  - $a_n = (1/2)^n$, where $n = 0, 1, 2, 3, \ldots$

  Question:

  Given a sequence, how to find a rule for generating the sequence?

Geometric Progression

- A geometric progression is a sequence of the form $a, ar, ar^2, \ldots, ar^n, \ldots$, where the *initial term* a and the *common ratio* r are real numbers.

Example:

◇ $a_n = (1/2)^n$, where $n = 0, 1, 2, 3, \ldots$

Question:

Given a sequence, how to find a rule for generating the sequence?

$$8, 42, 226, 1232, 6646, 35362, 185868, \ldots$$

Recursively Defined Sequences

- The n-th element of the sequence $\{a_n\}$ is defined recursively in terms of the previous elements of the sequence and the initial elements of the sequence.

Recursively Defined Sequences

- The n-th element of the sequence $\{a_n\}$ is defined recursively in terms of the previous elements of the sequence and the initial elements of the sequence.

Examples:

◇ $a_n = a_{n-1} + 2$ assuming $a_0 = 1$, for $n \geq 1$

◇ $f_n = f_{n-1} + f_{n-2}$ for $n = 2, 3, 4, \ldots$ (*Fibonacci sequence*)

Summations

- The *summation of the terms of a sequence* is

$$\sum_{j=m}^{n} a_j = a_m + a_{m+1} + \cdots + a_n$$

The variable j is referred to as *the index of summation* and the choice of the letter j is arbitrary.

  - ◇ m is the *lower limit*
  - ◇ n is the *upper limit* of the summation

Summations

- The *summation of the terms of a sequence* is

$$\sum_{j=m}^{n} a_j = a_m + a_{m+1} + \cdots + a_n$$

The variable j is referred to as *the index of summation* and the choice of the letter j is arbitrary.

  - ◇ m is the *lower limit*
  - ◇ n is the *upper limit* of the summation

$$\sum_{j=1}^{n}(ax_j + by_j) = a\sum_{j=1}^{n} x_j + b\sum_{j=1}^{n} y_j$$

$$\sum_{i=1}^{m}\sum_{j=1}^{n} a_i b_j = \sum_{i=1}^{m} a_i \sum_{j=1}^{n} b_j$$

Summations

- The sum of the first n terms of the arithmetic progression $a, a+d, a+2d, \ldots, a+nd$ is

$$S = \sum_{j=0}^{n} (a + jd) = (n+1)a + d\sum_{j=0}^{n} j = (n+1)a + d\frac{n(n+1)}{2}$$

- The sum of the first n terms of the geometric progression $a, ar, ar^2, \ldots, ar^k$ is

$$S = \sum_{j=0}^{n} (ar^j) = a\sum_{j=0}^{n} r^j = a\frac{r^{n+1} - 1}{r - 1}$$

Examples

- **Examples**:
  - $S = \sum_{j=1}^{5}(2+3j)$
  - $S = \sum_{j=3}^{5}(2+3j)$
  - $S = \sum_{i=1}^{4}\sum_{j=1}^{2}(2i-j)$
  - $S = \sum_{j=0}^{3}2(5)^{j}$
  - $S = \sum_{i=1}^{4}\sum_{j=1}^{3}ij$

Examples

- **Examples**:
  - $S = \sum_{j=1}^{5}(2+3j)$ 55
  - $S = \sum_{j=3}^{5}(2+3j)$ 42
  - $S = \sum_{i=1}^{4}\sum_{j=1}^{2}(2i-j)$ 28
  - $S = \sum_{j=0}^{3} 2(5)^j$ 312
  - $S = \sum_{i=1}^{4}\sum_{j=1}^{3} ij$ 60

Infinite Series

- Infinite geometric series can be computed in the closed form for $|x| < 1$.

Infinite Series

- Infinite geometric series can be computed in the closed form for $|x| < 1$.

$$\sum_{k=0}^{\infty} x^k = \lim_{n\to\infty} \sum_{k=0}^{n} x^k = \lim_{n\to\infty} \frac{x^{n+1} - 1}{x - 1} = \frac{1}{1 - x}$$

Infinite Series

- Infinite geometric series can be computed in the closed form for $|x| < 1$.

$$\sum_{k=0}^{\infty} x^k = \lim_{n\to\infty} \sum_{k=0}^{n} x^k = \lim_{n\to\infty} \frac{x^{n+1}-1}{x-1} = \frac{1}{1-x}$$

$$\sum_{k=0}^{\infty} k x^{k-1} = \frac{1}{(1-x)^2}$$

Some Useful Summation Formulas

TABLE 2 Some Useful Summation Formulae.

Sum	*Closed Form*
$\sum_{k=0}^{n} ar^k \ (r \neq 0)$	$\frac{ar^{n+1} - a}{r - 1}, r \neq 1$
$\sum_{k=1}^{n} k$	$\frac{n(n+1)}{2}$
$\sum_{k=1}^{n} k^2$	$\frac{n(n+1)(2n+1)}{6}$
$\sum_{k=1}^{n} k^3$	$\frac{n^2(n+1)^2}{4}$
$\sum_{k=0}^{\infty} x^k, \lvert x \rvert < 1$	$\frac{1}{1-x}$
$\sum_{k=1}^{\infty} kx^{k-1}, \lvert x \rvert < 1$	$\frac{1}{(1-x)^2}$

Cardinality of Sets

- Recall: the cardinality of a finite set is defined by the number of the elements in the set.

Cardinality of Sets

- Recall: the cardinality of a finite set is defined by the number of the elements in the set.

- The sets A and B have *the same cardinality* if there is a one-to-one correspondence between elements in A and B.

Cardinality of Sets

- Recall: the cardinality of a finite set is defined by the number of the elements in the set.

- The sets A and B have *the same cardinality* if there is a one-to-one correspondence between elements in A and B.

- If there is a one-to-one function from A to B, the cardinality of A is less than or the same as the cardinality of B, denoted by $|A| \leq |B|$. Moreover, when $|A| \leq |B|$ and A and B have different cardinalities, we say that the cardinality of A is less than the cardinality of B, denoted by $|A| < |B|$.

Countable Sets

- A set that is **either finite** or **has the same cardinality as the set of positive integers $\mathbf{Z}^+$** is called *countable*.
  A set that is **not countable** is called *uncountable*.

Countable Sets

- A set that is **either finite** or **has the same cardinality as the set of positive integers $\mathbf{Z}^+$** is called *countable*.
  A set that is **not countable** is called *uncountable*.

  Why are these called **countable**?

Countable Sets

- A set that is **either finite** or **has the same cardinality as the set of positive integers $\mathbf{Z}^+$** is called *countable*.
  A set that is **not countable** is called *uncountable*.

  Why are these called **countable**?

  ◇ The elements of the set can be **enumerated and listed**.

Hilbert's Grand Hotel

- The Grand Hotel has **countably infinite number of rooms**, each occupied by a guest. We can always accommodate a new guest at this hotel. How is this possible?

Hilbert's Grand Hotel

- The Grand Hotel has **countably infinite number of rooms**, each occupied by a guest. We can always accommodate a new guest at this hotel. How is this possible?

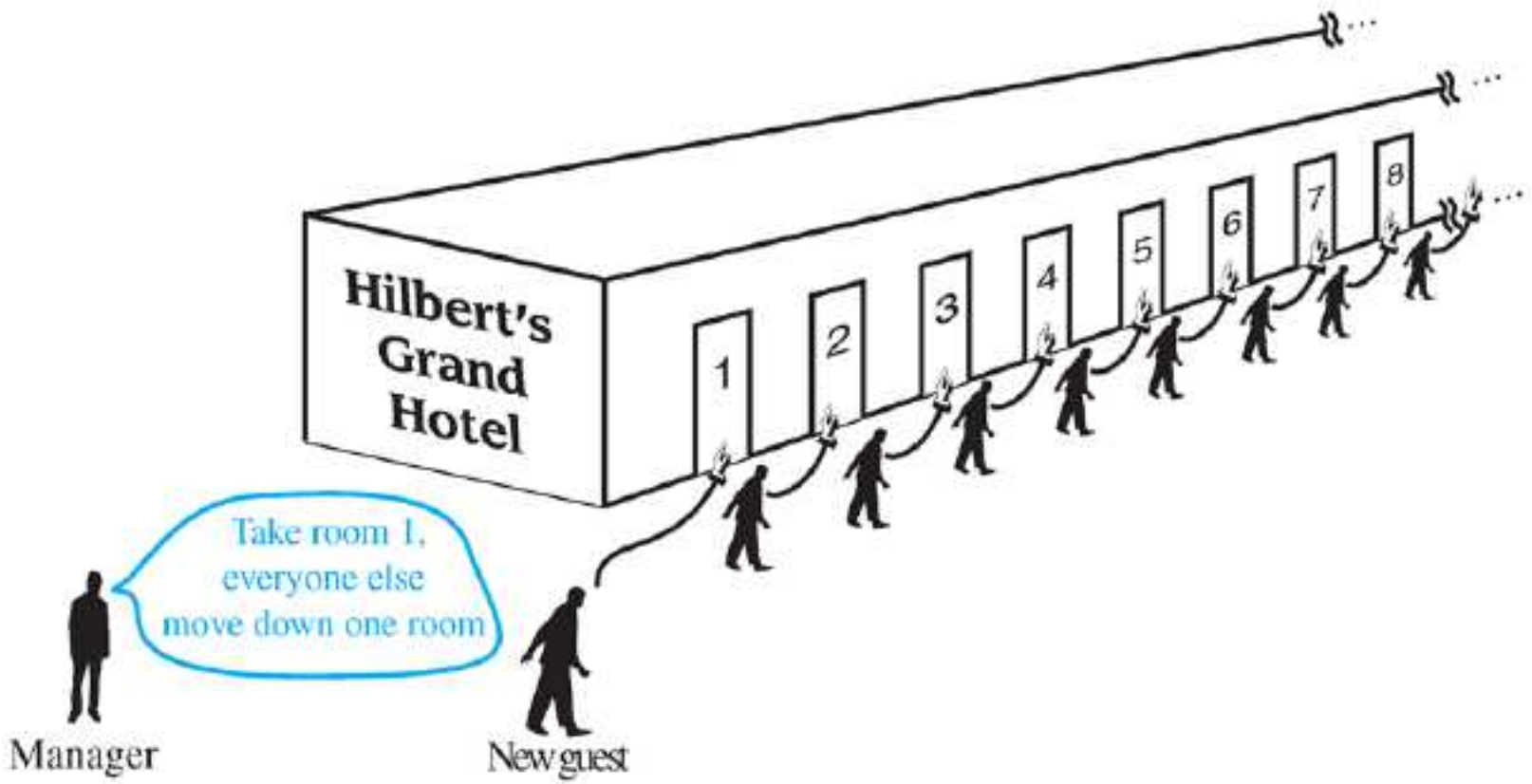


FIGURE 2 A New Guest Arrives at Hilbert's Grand Hotel.

Countable Sets

- **Example 1**

  $A = \{0, 2, 4, 6, \ldots\}$ – set of even numbers. Is it countable?

Countable Sets

- **Example 1**

  $A = \{0, 2, 4, 6, \ldots\}$ – set of even numbers. Is it countable?

  Using the definition: Is there a **bijection** $f : \mathbf{Z}^+ \to A$?

Countable Sets

- **Example 1**

  $A = \{0, 2, 4, 6, \ldots\}$ – set of even numbers. Is it countable?

  Using the definition: Is there a **bijection** $f : \mathbf{Z}^+ \to A$?

  Define a function $f : x \mapsto 2x - 2$. This is a bijection!

  one-to-one Why?

  onto Why?

Countable Sets

- **Example 1**

  $A = \{0, 2, 4, 6, \ldots\}$ – set of even numbers. Is it countable?

  Using the definition: Is there a **bijection** $f : \mathbf{Z}^+ \to A$?

  Define a function $f : x \mapsto 2x - 2$. This is a bijection!

  one-to-one Why?

  if $2x - 2 = 2y - 2$, then $x = y$

  onto Why?

  $\forall x \in A$, $(x + 2)/2$ is the preimage in $\mathbf{Z}^+$

Countable Sets

- **Example 2 (Theorem)**

  The set of integers $\mathbf{Z}$ is countable.

Countable Sets

- **Example 2 (Theorem)**

  The set of integers $\mathbf{Z}$ is countable.

  Solution:

  We can list a sequence:

  $$0, 1, -1, 2, -2, 3, -3, \ldots$$

  or define a bijection from $\mathbf{Z}^+$ to $\mathbf{Z}$:
  – when n is even: $f(n) = n/2$
  – when n is odd: $f(n) = -(n-1)/2$

Countable Sets

- **Example 3 (Theorem)**

  The set of (positive) rational numbers is countable.

Countable Sets

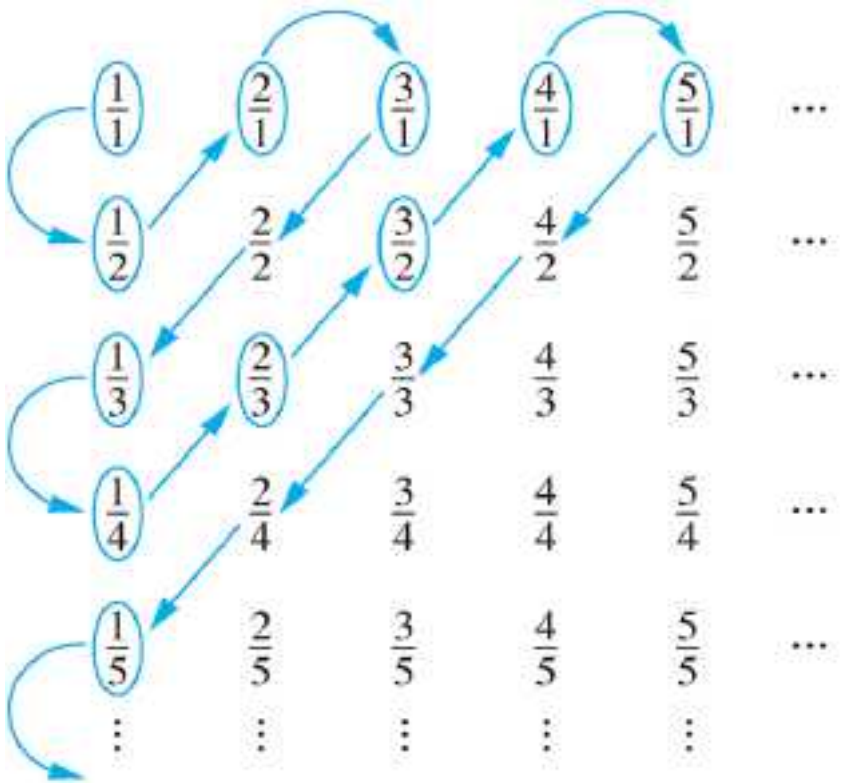
- **Example 3 (Theorem)**

  The set of (positive) rational numbers is countable.

Solution:

Constructing the list: first list p/q with $p+q=2$, next list p/q with $p+q=3$, and so on.

$1, 1/2, 2, 3, 1/3, 1/4, 2/3, \ldots$



Countable Sets

- **Example 4 (Theorem)**

  The set of finite strings S over a finite alphabet A is countably infinite. (Assume an alphabetical ordering of symbols in A)

Countable Sets

- **Example 4 (Theorem)**

  The set of finite strings S over a finite alphabet A is countably infinite. (Assume an alphabetical ordering of symbols in A)

  Solution:

  We show that the strings can be listed in a sequence. First list

  (i) all the strings of length 0 in alphabetical order.

  (ii) then all the strings of length 1 in lexicographic order.

  (iii) and so on.

  This implies a bijection from $\mathbf{Z}^+$ to S.

Countable Sets

- **Example 5**

  The set of all Java programs is countable.

Countable Sets

- **Example 5**

  The set of all Java programs is countable.

  Solution:

  Let S be the set of strings constructed from the characters which may appear in a Java program. Use the ordering from the previous example. Take each string in turn

  – feed the string into a Java compiler

  – if the complier says YES, this is a syntactically correct Java program, we add this program to the list

  – we move on to the next string

  In this way, we construct a bijection from $\mathbf{Z}^+$ to the set of Java programs.

Uncountable Sets

- **Theorem**

  The set of real numbers $\mathbf{R}$ is uncountable.

Uncountable Sets

- **Theorem**

  The set of real numbers **R** is uncountable.

Proof by contradiction:

Assume that **R** is countable. Then every subset of **R** is countable (why?), in particular, the interval from 0 to 1 is countable. This implies that the elements of this set can be listed as $r_1, r_2, r_3, \ldots$, where

- $r_1 = 0.d_{11}d_{12}d_{13}d_{14}\cdots$
- $r_2 = 0.d_{21}d_{22}d_{23}d_{24}\cdots$
- $r_3 = 0.d_{31}d_{32}d_{33}d_{34}\cdots$

all $d_{ij} \in \{0, 1, 2, \ldots, 9\}$.

Uncountable Sets

- **Theorem**

  The set of real numbers **R** is uncountable.

Proof by contradiction:

We want to show that not all real numbers in the interval between 0 and 1 are in this list.

Form a new number called $r = 0.d_1 d_2 d_3 d_4 \cdots$, where $d_i = 2$ if $d_{ii} \neq 2$, and $d_i = 3$ if $d_{ii} = 2$.

Uncountable Sets

- **Theorem**

  The set of real numbers **R** is uncountable.

Proof by contradiction:

We want to show that not all real numbers in the interval between 0 and 1 are in this list.

Form a new number called $r = 0.d_1 d_2 d_3 d_4 \cdots$, where $d_i = 2$ if $d_{ii} \neq 2$, and $d_i = 3$ if $d_{ii} = 2$.

Example: suppose	r1 = 0.75243...	d1 = 2
	r2 = 0.524310...	d2 = 3
	r3 = 0.131257...	d3 = 2
	r4 = 0.9363633...	d4 = 2
	...	...
	rt = 0.23222222...	dt = 3

Uncountable Sets

- **Theorem**

  The set of real numbers **R** is uncountable.

Proof by contradiction:

We claim that r is different from each number in the list.

Each expansion is unique, if we exclude an infinite string of 9's. r and r_i differ in the i-th decimal place for all i.

Uncountable Sets

- **Theorem**

  The set of real numbers **R** is uncountable.

Proof by contradiction:

We claim that r is different from each number in the list.

Each expansion is unique, if we exclude an infinite string of 9's. r and r_i differ in the i-th decimal place for all i.

This is called *Cantor diagonalization argument*.

Uncountable Sets

- **Theorem**

  The set $\mathcal{P}(\mathbb{N})$ is uncountable.

Uncountable Sets

- **Theorem**

  The set $\mathcal{P}(\mathbb{N})$ is uncountable.

Proof by contradiction:

Assume that $\mathcal{P}(\mathbb{N})$ is countable. This implies that the elements of this set can be listed as $S_0, S_1, S_2, \ldots$, where $S_i \subseteq \mathbb{N}$, and each S_i can be represented uniquely by the bit string $b_{i0}\, b_{i1}\, b_{i2} \ldots$, where $b_{ij} = 1$ if $j \in S_i$ and $b_{ij} = 0$ if $j \notin S_i$

- $S_0 = b_{00}\, b_{01}\, b_{02}\, b_{03} \cdots$
- $S_1 = b_{10}\, b_{11}\, b_{12}\, b_{13} \cdots$
- $S_2 = b_{20}\, b_{21}\, b_{22}\, b_{23} \cdots$

$\vdots$

all $b_{ij} \in \{0, 1\}$.

Uncountable Sets

- **Theorem**

  The set $\mathcal{P}(\mathbb{N})$ is uncountable.

  Proof by contradiction:

  Form a new set called $R = b_0 b_1 b_2 b_3 \cdots$, where $b_i = 0$ if $b_{ii} = 1$, and $b_i = 1$ if $b_{ii} = 0$.

Uncountable Sets

- **Theorem**

  The set $\mathcal{P}(\mathbb{N})$ is uncountable.

  Proof by contradiction:

  Form a new set called $R = b_0 b_1 b_2 b_3 \cdots$, where $b_i = 0$ if $b_{ii} = 1$, and $b_i = 1$ if $b_{ii} = 0$.

  We claim that R is different from each set in the list.

  Each bit string is unique, and R and S_i differ in the i-th bit for all i.

Schröder-Bernstein Theorem

- **Theorem**

  If A and B are sets with $|A| \leq |B|$ and $|B| \leq |A|$, then $|A| = |B|$. In other words, if there are one-to-one functions f from A to B and g from B to A, then there is a one-to-one correspondence between A and B.

Schröder-Bernstein Theorem

- **Theorem**

  If A and B are sets with $|A| \leq |B|$ and $|B| \leq |A|$, then $|A| = |B|$. In other words, if there are one-to-one functions f from A to B and g from B to A, then there is a one-to-one correspondence between A and B.

  Example

  Show that $|(0, 1)| = |(0, 1]|$.

Schröder-Bernstein Theorem

- **Theorem**

  If A and B are sets with $|A| \leq |B|$ and $|B| \leq |A|$, then $|A| = |B|$. In other words, if there are one-to-one functions f from A to B and g from B to A, then there is a one-to-one correspondence between A and B.

  Example

  Show that $|(0, 1)| = |(0, 1]|$.

  $f(x) = x$; $g(x) = x/2$

Computable vs Uncomputable

- **Definition**

  We say that a function is *computable* if there is a computer program in some programming language that finds the values of this function. If a function is not computable, we say it is *uncomputable*.

Computable vs Uncomputable

- **Definition**

  We say that a function is *computable* if there is a computer program in some programming language that finds the values of this function. If a function is not computable, we say it is *uncomputable*.

  Theorem*

  There are functions that are not computable.

Computable vs Uncomputable

- **Definition**

  We say that a function is *computable* if there is a computer program in some programming language that finds the values of this function. If a function is not computable, we say it is *uncomputable*.

  Theorem*

  There are functions that are not computable.

  Cantor's theorem*

  If S is a set, then $|S| < |\mathcal{P}(S)|$.

Next Lecture

- complexity ...