



CS201 DISCRETE MATHEMATICS FOR COMPUTER SCIENCE

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Generating Functions

- We may use *generating functions* to **characterize** sequences.

◇ The sequence $\{a_k\}$ with $a_k = 3$

$$\sum_{k=0}^{\infty} 3x^k$$

◇ The sequence $\{a_k\}$ with $a_k = 2^k$

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Definition The *generating function* for the sequence $a_0, a_1, \dots, a_k, \dots$ of real numbers is the infinite series

$$G(x) = a_0 + a_1x + \dots + a_kx^k$$



Counting and Generating Functions

- **Problem 2** Find the number of solutions of

$$x_1 + x_2 + x_3 = 17,$$

where x_1, x_2, x_3 are **nonnegative** integers with $2 \leq x_1 \leq 5$,
 $3 \leq x_2 \leq 6$, $4 \leq x_3 \leq 7$.



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Using *generating functions*, the number is the **coefficient** of x^{17} in the expansion of

$$(x^2 + x^3 + x^4 + x^5)(x^3 + x^4 + x^5 + x^6)(x^4 + x^5 + x^6 + x^7)$$



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$$C(n + r - 1, r) = C(19, 17) = C(19, 2)$$



r -Combinations from a Set

- **Definition** An r -combination with repetition allowed, or a *multiset of size r* , chosen from a set of n elements, is an unordered selection of elements with repetition allowed.



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Read more on pp. 537-548.



Cartesian Product

- Let $A = \{a_1, a_2, \dots, a_m\}$ and $B = \{b_1, b_2, \dots, b_n\}$, the *Cartesian product* $A \times B$ is the set of pairs

$$\{(a_1, b_1), (a_2, b_2), \dots, (a_1, b_n), \dots, (a_m, b_n)\}$$



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Cartesian product defines a set of all **ordered** arrangements of elements in the two sets.



Binary Relation

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Example: Let $A = \{a, b, c\}$ and $B = \{1, 2, 3\}$

- ◇ Is $R = \{(a, 1), (b, 2), (c, 2)\}$ a relation from A to B ?
- ◇ Is $Q = \{(1, a), (2, b)\}$ a relation from A to B ?
- ◇ Is $P = \{(a, a), (b, c), (b, a)\}$ a relation from A to A ?



Representing Binary Relations

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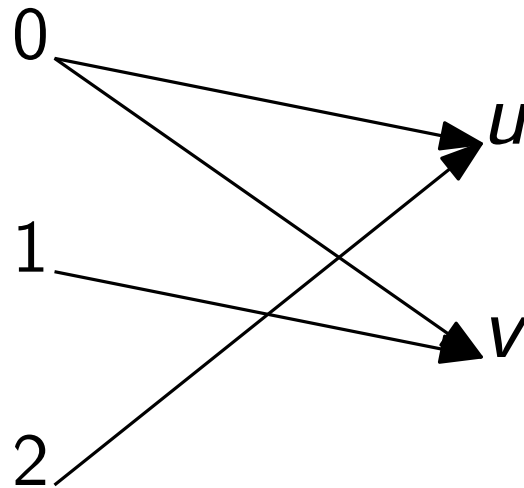
Example: Let $A = \{0, 1, 2\}$ and $B = \{u, v\}$, and
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|-----|-----|-----|
| 0 | × | × |
| 1 | × | |
| 2 | | × |



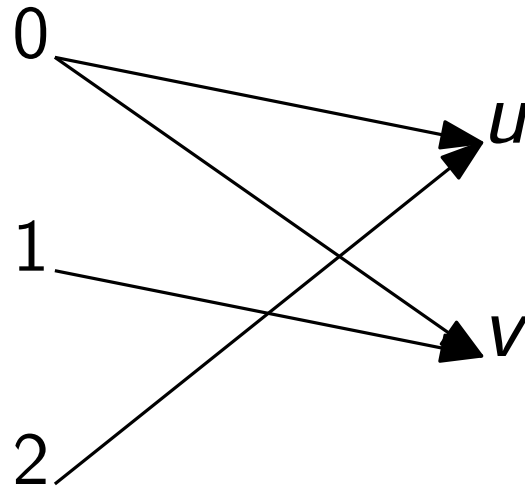
Relations and Functions

- Relations represent **one to many relationships** between elements in A and B .



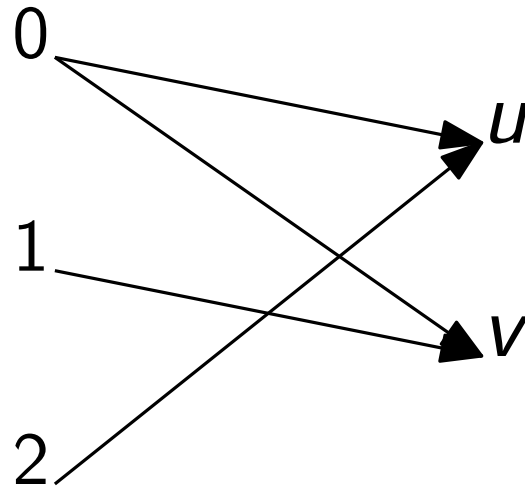
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What is the **difference** between a **relation** and a **function** from A to B ?



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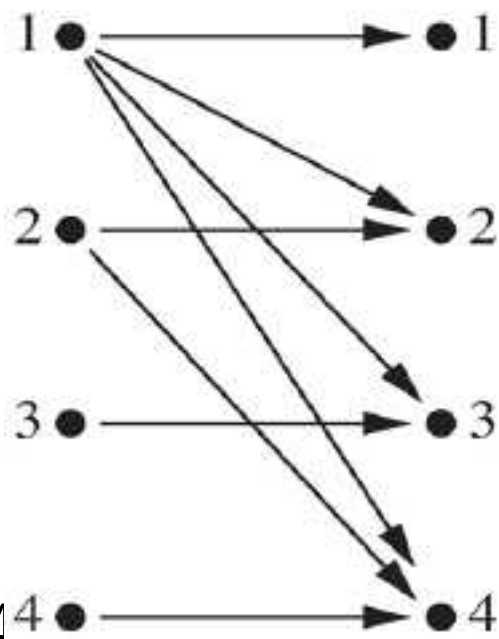


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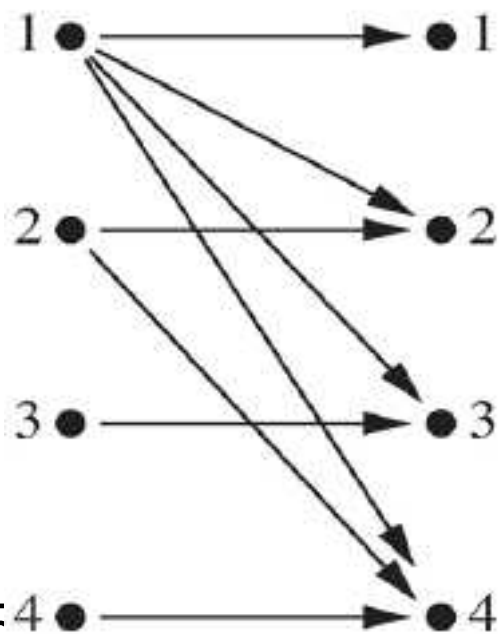


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| R | 1 | 2 | 3 | 4 |
|-----|---|---|---|---|
| 1 | × | × | × | × |
| 2 | | × | | × |
| 3 | | | × | |
| 4 | | | | × |



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The number of subsets of a set with k elements is 2^k



Properties of Relations

- **Reflexive Relation:** A relation R on a set A is called *reflexive* if $(a, a) \in R$ for **every** element $a \in A$.



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Yes. $(1, 1), (2, 2), (3, 3), (4, 4) \in R_{div}$



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$$MR_{div} = \begin{matrix} & \begin{matrix} 1 & 2 & 3 & 4 \end{matrix} \\ \begin{matrix} 1 \\ 2 \\ 3 \\ 4 \end{matrix} & \begin{bmatrix} 1 & 1 & 1 & 1 \\ 0 & 1 & 0 & 1 \\ 0 & 0 & 1 & 0 \\ 0 & 0 & 0 & 1 \end{bmatrix} \end{matrix}$$



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A relation R is reflexive if and only if MR has 1 in every position on its main diagonal.



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No. $(1, 1) \notin R$



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A relation R is antisymmetric if and only if $m_{ij} = 1$ implies $m_{ji} = 0$ for $i \neq j$.



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No. $(1, 2), (2, 1) \in R_{\neq}$ but $(1, 1) \notin R_{\neq}$.



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Set operations: **union, intersection, difference, etc.**



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We may also combine relations by **matrix operations**.



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“only if” part: by induction.



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How many subsets on $n(n-1)$ elements are there?



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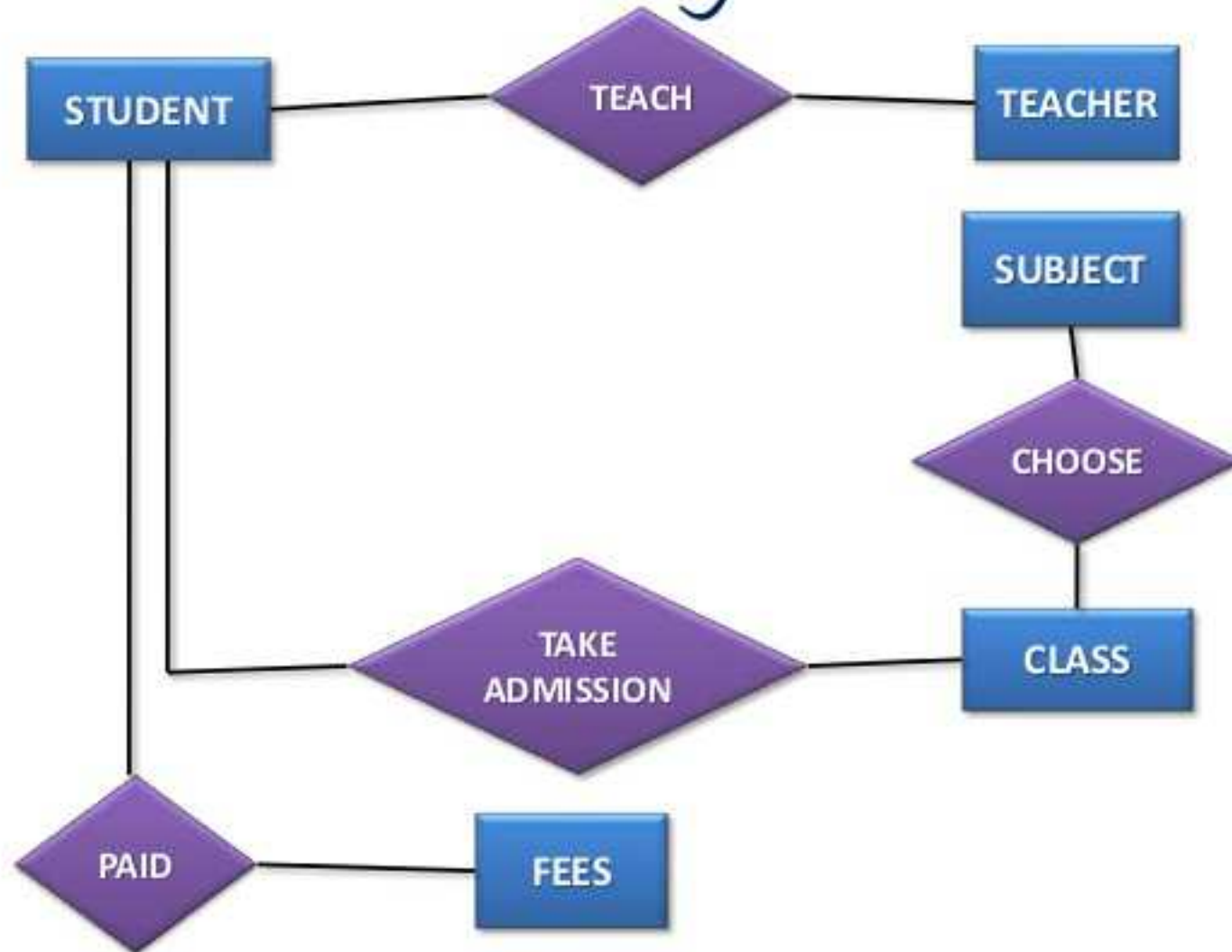
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- A *relational database* is essentially an *n*-ary relation R .
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- A *composite key* for the database is a set of domains $\{A_i, A_j, \dots\}$ such that R contains **at most 1 *n*-tuple** $(\dots, a_i, \dots, a_j, \dots)$ for each composite value $(a_i, a_j, \dots) \in A_i \times A_j \times \dots$.



Relational Databases

E-R Diagram



Selection Operators

- Let A be any *n -ary domain* $A = A_1 \times \cdots \times A_n$, and let $C : A \rightarrow \{T, F\}$ be any *condition* (predicate) on elements (n -tuples) of A .



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$$- \forall R \subseteq A,$$

$$\begin{aligned} s_C(R) &= R \cap \{a \in A \mid s_C(a) = T\} \\ &= \{a \in R \mid s_C(a) = T\}. \end{aligned}$$



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- Then, $\textit{SUpperLevel}$ is the selection operator that takes any relation R on A (database of students) and produces a relation consisting of just the upper-level classes (juniors and seniors).



Projection Operators

- Let $A = A_1 \times \cdots \times A_n$ be any n -ary domain, and let $\{i_k\} = (i_1, \dots, i_m)$ be a sequence of indices all falling in the range 1 to n .
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- Then the *projection operator* on n -tuples

$$P_{\{i_k\}} : A \rightarrow A_{i_1} \times \cdots \times A_{i_m}$$

is defined by

$$P_{\{i_k\}}(a_1, \dots, a_n) = (a_{i_1}, \dots, a_{i_m})$$



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- This operator can be usefully applied to a whole relation $R \subseteq Cars$ (database of cars) to obtain a list of *model/color* combinations available.



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- A, B, C can also be sequences of elements rather than single elements.



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- Suppose that R_2 is a room assignment table relating *Courses* to *Rooms* and *Times*.
- Then $J(R_1, R_2)$ is like your **class schedule**, listing *(professor, course, room, time)*.



Next Lecture

- relation II ...

