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```

# Basic

# 1.1 vimrc

se nu rnu ai hls et ru ic is sc cul

```
se re=1 ts=4 sts=4 sw=4 ls=2 mouse=a
syntax on
set bg=dark
ino {<CR> {<CR>}<Esc>ko<tab>
gsettings set org.gnome.
    desktop.input-sources xkb-options "['caps:escape']"
1.2 readchar [a419b9]
inline char readchar() {
  static const size_t bufsize = 65536;
  static char buf[bufsize];
  static char *p = buf, *end = buf;
  if (p == end) end = buf +
       fread_unlocked(buf, 1, bufsize, stdin), p = buf;
```

## 1.3 Black Magic [afb343]

```
#include <ext/pb_ds/priority_queue.hpp>
#include <ext/pb_ds/assoc_container.hpp> // rb_tree
#include <ext/rope> // rope
using namespace __gnu_pbds;
using namespace __gnu_cxx; // rope
typedef __gnu_pbds::priority_queue<int> heap;
int main() {
  heap h1, h2; // max heap
  h1.push(1), h1.push(3), h2.push(2), h2.push(4);
  h1.join(h2); // h1 = {1, 2, 3, 4}, h2 = {};
tree<ll, null_type, less<ll>, rb_tree_tag
       , tree_order_statistics_node_update > st;
  tree<ll, ll, less<ll>, rb_tree_tag
        tree_order_statistics_node_update > mp;
  for (int x : {0, 3, 20, 50}) st.insert(x);
  assert(st.
      order_of_key(3) == 1 && st.order_of_key(4) == 2);
  assert(*st.find_by_order
      (2) == 20 && *st.lower_bound(4) == 20);
  rope < char > *root[10]; // nsqrt(n)
  root[0] = new rope < char > ();
  root[1] = new rope < char > (*root[0]);
                              'a');
  // root[1]->insert(pos,
  // root[1]->at(pos); 0-base
  // root[1]->erase(pos, size);
// __int128_t,__float128_t
// for (int i = bs._Find_first
     (); i < bs.size(); i = bs._Find_next(i));
```

#### 1.4 Pragma Optimization [Oca2dd]

```
#pragma GCC optimize("Ofast, no-stack-protector")
#pragma GCC optimize("no-math-errno,unroll-loops")
#pragma GCC target("sse,sse2,sse3,ssse3,sse4")
#pragma GCC target("popcnt,abm,mmx,avx,arch=skylake")
__builtin_ia32_ldmxcsr(__builtin_ia32_stmxcsr()|0x8040)
```

# 2 Graph

#### 2.1 BCC Vertex\* [6763ed]

```
struct BCC { // 0-base
  int n, dft, nbcc;
vi low, dfn, bln, stk, is_ap, cir;
  vvi G, bcc, nG;
  void make_bcc(int u) {
     bcc.eb(1, u);
     for (; stk.back() != u; stk.pop_back())
       bln[stk.back()] = nbcc, bcc[nbcc].pb(stk.back());
     stk.pop_back(), bln[u] = nbcc++;
  void dfs(int u, int p) {
     int child = 0;
    low[u] = dfn[u] = ++dft, stk.pb(u);
     for (int v : G[u])
       if (!dfn[v]) {
         dfs(v, u), ++child;
low[u] = min(low[u], low[v]);
         if (dfn[u] <= low[v]) {</pre>
            is_ap[u] = 1, bln[u] = nbcc;
make_bcc(v), bcc.back().pb(u);
       } else if (dfn[v] < dfn[u] && v != p)</pre>
         low[u] = min(low[u], dfn[v]);
    if (p == -1 && child < 2) is_ap[u] = 0;</pre>
    if (p == -1 && child == 0) make_bcc(u);
  BCC(int _n): n(_n), dft(),
    nbcc(), low(n), dfn(n), bln(n), is_ap(n), G(n) {}
  void add_edge(int u, int v) {
    G[u].pb(v), G[v].pb(u);
  void solve() {
   for (int i = 0; i < n; ++i)</pre>
       if (!dfn[i]) dfs(i, -1);
  void block_cut_tree() {
    cir.resize(nbcc);
     for (int i = 0; i < n; ++i)</pre>
       if (is_ap[i])
         bln[i] = nbcc++;
    cir.resize(nbcc, 1), nG.resize(nbcc);
for (int i = 0; i < nbcc && !cir[i]; ++i)</pre>
       for (int j : bcc[i])
```

if (scc.bln[i] == scc.bln[i + n]) return false;

istrue[i] = scc.bln[i] < scc.bln[i + n];</pre>

```
if (is ap[i])
                                                                   istrue[i + n] = !istrue[i];
          nG[i].pb(bln[j]), nG[bln[j]].pb(i);
 } // up to 2 * n - 2 nodes!! bln[i] for id
                                                                 return true;
                                                               }
                                                            };
2.2 Bridge* [4da29a]
                                                             2.5 MinimumMeanCycle* [3e5d2b]
struct ECC { // 0-base
  int n, dft, ecnt, necc;
                                                             ll road[N][N]; // input here
  vector<int> low, dfn, bln, is_bridge, stk;
                                                             struct MinimumMeanCycle {
  vector<vector<pii>>> G;
                                                               ll dp[N + 5][N], n;
  void dfs(int u, int f) {
                                                               pll solve() {
                                                                 ll a = -1, b = -1, L = n + 1;
for (int i = 2; i <= L; ++i)
    dfn[u] = low[u] = ++dft, stk.pb(u);
    for (auto [v, e] : G[u])
                                                                   for (int k = 0; k < n; ++k)
      if (!dfn[v])
        dfs(v, e), low[u] = min(low[u], low[v]);
                                                                     for (int j = 0; j < n; ++j)</pre>
                                                                       dp[i][j] =
      else if (e != f)
        low[u] = min(low[u], dfn[v]);
                                                                         min(dp[i - 1][k] + road[k][j], dp[i][j]);
                                                                 for (int i = 0; i < n; ++i) {</pre>
    if (low[u] == dfn[u]) {
                                                                   if (dp[L][i] >= INF) continue;
      if (f != -1) is_bridge[f] = 1;
                                                                   ll ta = 0, tb = 1;
      for (; stk.back() != u; stk.pop_back())
        bln[stk.back()] = necc;
                                                                   for (int j = 1; j < n; ++j)</pre>
                                                                     if (dp[j][i] < INF &&</pre>
      bln[u] = necc++, stk.pop_back();
   }
                                                                       ta *
                                                                            (L - j) < (dp[L][i] - dp[j][i]) * tb)
                                                                       ta = dp[L][i] - dp[j][i], tb = L - j;
  ECC(int _n): n(_n), dft()
                                                                   if (ta == 0) continue;
  , ecnt(), necc(), low(n), dfn(n), bln(n), G(n) \{\} void add_edge(int u, int v) \{
                                                                   if (a == -1 || a * tb > ta * b) a = ta, b = tb;
   G[u].pb(pii(v, ecnt)), G[v].pb(pii(u, ecnt++));
                                                                 if (a != -1) {
                                                                   ll g = _
                                                                            _gcd(a, b);
  void solve() {
                                                                   return pll(a / g, b / g);
    is_bridge.resize(ecnt);
    for (int i = 0; i < n; ++i)</pre>
                                                                 return pll(-1LL, -1LL);
      if (!dfn[i]) dfs(i, -1);
                                                               void init(int _n) {
}; // ecc_id(i): bln[i]
                                                                 for (int i = 0; i < n; ++i)</pre>
2.3 SCC* [7ae459]
                                                                   for (int j = 0; j < n; ++j) dp[i + 2][j] = INF;</pre>
struct SCC { // 0-base
                                                            |};
  int n, dft, nscc;
                                                             2.6 Virtual Tree* [1b641b]
  vi low, dfn, bln, instack, stk;
  vvi G:
  void dfs(int u) {
                                                             vector<int> vG[N];
    low[u] = dfn[u] = ++dft;
                                                             int top, st[N];
    instack[u] = 1, stk.pb(u);
    for (int v : G[u])
                                                             void insert(int u) {
      if (!dfn[v])
                                                               if (top == -1) return st[++top] = u, void();
        dfs(v), low[u] = min(low[u], low[v]);
                                                               int p = LCA(st[top], u);
      else if (instack[v] && dfn[v] < dfn[u])
                                                               if (p == st[top]) return st[++top] = u, void();
        low[u] = min(low[u], dfn[v]);
                                                               while (top \Rightarrow= 1 && dep[st[top - 1]] \Rightarrow= dep[p])
    if (low[u] == dfn[u]) {
                                                                 vG[st[top - 1]].pb(st[top]), --top;
                                                               if (st[top] != p)
      for (; stk.back() != u; stk.pop_back())
        bln[stk
                                                                 vG[p].pb(st[top]), --top, st[++top] = p;
      .back()] = nscc, instack[stk.back()] = 0; \\ instack[u] = 0, bln[u] = nscc++, stk.pop_back(); \\
                                                               st[++top] = u;
   }
                                                             void reset(int u) {
 for (int i : vG[u]) reset(i);
                                                               vG[u].clear();
  void add_edge(int u, int v) {
   G[u].pb(v);
                                                             void solve(vector<int> &v) {
  void solve() {
                                                               top = -1;
    for (int i = 0; i < n; ++i)</pre>
                                                               sort(ALL(v),
      if (!dfn[i]) dfs(i);
                                                                [&](int a, int b) { return dfn[a] < dfn[b]; });
                                                               for (int i : v) insert(i);
                                                               while (top > 0) vG[st[top - 1]].pb(st[top]), --top;
}; // scc_id(i): bln[i]
                                                               // do something
2.4 2SAT* [f5630a]
                                                               reset(v[0]);
struct SAT { // 0-base
                                                             2.7 Maximum Clique Dyn* [d50aa9]
  int n;
  vector<bool> istrue;
  SCC scc;
                                                             struct MaxClique { // fast when N <= 100</pre>
  SAT(int _n): n(_n), istrue(n + n), scc(n + n) {}
                                                               bitset<N> G[N], cs[N];
                                                               int ans, sol[N], q, cur[N], d[N], n;
  int rv(int a) {
   return a >= n ? a - n : a + n;
                                                               void init(int _n) {
                                                                 n = _n;
  void add_clause(int a, int b) {
                                                                 for (int i = 0; i < n; ++i) G[i].reset();</pre>
    scc.add_edge(rv(a), b), scc.add_edge(rv(b), a);
                                                               void add_edge(int u, int v) {
  bool solve() {
                                                                 G[u][v] = G[v][u] = 1;
    scc.solve();
    for (int i = 0; i < n; ++i) {</pre>
                                                               void pre dfs(vector<int> &r, int l, bitset<N> mask) {
```

**if** (1 < 4) {

for (int i : r) d[i] = (G[i] & mask).count();

```
sort(ALL(r)
           , [&](int x, int y) { return d[x] > d[y]; });
    vector<int> c(SZ(r));
    int lft = max(ans - q + 1, 1), rgt = 1, tp = 0;
    cs[1].reset(), cs[2].reset();
    for (int p : r) {
      int k = 1;
       while ((cs[k] & G[p]).any()) ++k;
      if (k > rgt) cs[++rgt + 1].reset();
       cs[k][p] = 1;
      if (k < lft) r[tp++] = p;</pre>
    for (int k = lft; k <= rgt; ++k)</pre>
       for (int p = cs[k]._Find_first
           (); p < N; p = cs[k]._Find_next(p))
         r[tp] = p, c[tp] = k, ++tp;
    dfs(r, c, l + 1, mask);
  void dfs(vector<</pre>
       int> &r, vector<int> &c, int l, bitset<N> mask) {
    while (!r.empty()) {
      int p = r.back();
       r.pop_back(), mask[p] = 0;
      if (q + c.back() <= ans) return;</pre>
      cur[q++] = p;
      vector<int> nr;
       for (int i : r) if (G[p][i]) nr.pb(i);
      if (!nr.empty()) pre_dfs(nr, l, mask & G[p]);
      else if (q > ans) ans = q, copy_n(cur, q, sol);
      c.pop_back(), --q;
    }
  int solve() {
    vector<int> r(n);
    ans = q = 0, iota(ALL(r), 0);
    pre_dfs(r, 0, bitset<N>(string(n, '1')));
    return ans;
};
```

#### 2.8 Minimum Steiner Tree\* [62d6fb]

```
struct SteinerTree { // 0-base
  int n, dst[N][N], dp[1 << T][N], tdst[N];
int vcst[N]; // the cost of vertexs</pre>
  void init(int _n) {
    n = _n;
for (int i = 0; i < n; ++i) {</pre>
       fill_n(dst[i], n, INF);
      dst[i][i] = vcst[i] = 0;
    }
  void chmin(int &x, int val) {
    x = min(x, val);
  void add_edge(int ui, int vi, int wi) {
    chmin(dst[ui][vi], wi);
  void shortest_path() {
    for (int k = 0; k < n; ++k)</pre>
      for (int i = 0; i < n; ++i)</pre>
         for (int j = 0; j < n; ++j)</pre>
           chmin(dst[i][j], dst[i][k] + dst[k][j]);
  int solve(const vector<int>& ter) {
    shortest_path();
    int t = SZ(ter), full = (1 << t) - 1;</pre>
    for (int i = 0; i <= full; ++i)</pre>
      fill_n(dp[i], n, INF);
    copy_n(vcst, n, dp[0]);
for (int msk = 1; msk <= full; ++msk) {</pre>
      if (!(msk & (msk - 1))) {
         int who = __lg(msk);
         for (int i = 0; i < n; ++i)</pre>
           dp[msk
                ][i] = vcst[ter[who]] + dst[ter[who]][i];
       for (int i = 0; i < n; ++i)</pre>
         for (int sub = (
             msk - 1) & msk; sub; sub = (sub - 1) & msk)
           chmin(dp[msk][i],
                dp[sub][i] + dp[msk ^ sub][i] - vcst[i]);
       for (int i = 0; i < n; ++i) {</pre>
         tdst[i] = INF;
         for (int j = 0; j < n; ++j)</pre>
           chmin(tdst[i], dp[msk][j] + dst[j][i]);
```

```
copy_n(tdst, n, dp[msk]);
    return *min_element(dp[full], dp[full] + n);
}; // O(V 3^T + V^2 2^T)
```

#### 2.9 Dominator Tree\* [2b8b32]

```
struct dominator_tree { // 1-base
  vector<int> G[N], rG[N];
  int n, pa[N], dfn[N], id[N], T
int semi[N], idom[N], best[N];
                                 Time:
  vector<int> tree[N]; // dominator_tree
  void init(int _n) {
    n = _n;
    for (int i = 1; i <= n; ++i)
      G[i].clear(), rG[i].clear();
  void add_edge(int u, int v) {
    G[u].pb(v), rG[v].pb(u);
  void dfs(int u) {
    id[dfn[u] = ++Time] = u;
    for (auto v : G[u])
      if (!dfn[v]) dfs(v), pa[dfn[v]] = dfn[u];
  int find(int y, int x) {
    if (y <= x) return y;</pre>
    int tmp = find(pa[y], x);
    if (semi[best[y]] > semi[best[pa[y]]])
      best[y] = best[pa[y]];
    return pa[y] = tmp;
  void tarjan(int root) {
    Time = 0;
    for (int i = 1; i <= n; ++i) {</pre>
      dfn[i] = idom[i] = 0;
      tree[i].clear();
      best[i] = semi[i] = i;
    dfs(root);
    for (int i = Time; i > 1; --i) {
      int u = id[i];
      for (auto v : rG[u])
        if (v = dfn[v]) {
          find(v, i);
           semi[i] = min(semi[i], semi[best[v]]);
      tree[semi[i]].pb(i);
      for (auto v : tree[pa[i]]) {
        find(v, pa[i]);
        idom[v] =
           semi[best[v]] == pa[i] ? pa[i] : best[v];
      tree[pa[i]].clear();
    for (int i = 2; i <= Time; ++i) {</pre>
      if (idom[i] != semi[i]) idom[i] = idom[idom[i]];
      tree[id[idom[i]]].pb(id[i]);
  }
};
```

#### 2.10 Minimum Arborescence\* [c7338d]

```
/* TODO
DSU: disjoint set
- DSU(n), .boss(x), .Union(x, y)
min heap <
    T, Info>: min heap for type {T, Info} with lazy tag
 .push({w, i}),
    .top(), .join(heap), .pop(), .empty(), .add_lazy(v)
struct E { int s, t; ll w; }; // O-base
vector<int> dmst(const vector<E> &e, int n, int root) {
  vector<min_heap<ll, int>> h(n * 2);
  for (int i = 0; i < SZ(e); ++i)</pre>
    h[e[i].t].push({e[i].w, i});
  DSU dsu(n * 2);
  vector<int> v(n * 2, -1), pa(n * 2, -1), r(n * 2);
  v[root] = n + 1;
  int pc = n;
  for (int i = 0; i < n; ++i) if (v[i] == -1) {</pre>
    for (int p = i; v[p]
         == -1 \mid \mid v[p] == i; p = dsu.boss(e[r[p]].s)) {
      if (v[p] == i) {
```

```
int q = p; p = pc++;
       do {
         h[q].add_lazy(-h[q].top().X);
         pa[q] = p, dsu.Union(p, q), h[p].join(h[q]);
       } while ((q = dsu.boss(e[r[q]].s)) != p);
     v[p] = i;
     while (!h[p].
         empty() && dsu.boss(e[h[p].top().Y].s) == p)
       h[p].pop();
     if (h[p].empty()) return {}; // no solution
     r[p] = h[p].top().Y;
   }
  vector<int> ans;
 for (int i = pc
       - 1; i >= 0; i--) if (i != root && v[i] != n) {
    for (int f = e[r[i]].t; \sim f \&\& v[f] != n; f = pa[f])
     v[f] = n;
   ans.pb(r[i]);
 }
  return ans; // default minimize, returns edgeid array
```

## 2.11 Vizing's theorem\* [2b5b01]

```
namespace vizing { // returns
     edge coloring in adjacent matrix G. 1 - based
const int N = 105;
int C[N][N], G[N][N], X[N], vst[N], n;
void init(int _n) { n = _n;
  for (int i = 0; i <= n; ++i)
  for (int j = 0; j <= n; ++j)</pre>
      C[i][j] = G[i][j] = 0;
void solve(vector<pii> &E) {
  auto update = [&](int u)
  { for (X[u] = 1; C[u][X[u]]; ++X[u]); };
  auto color = [&](int u, int v, int c) {
    int p = G[u][v];
    G[u][v] = G[v][u] = c;
    C[u][c] = v, C[v][c] = u;
    C[u][p] = C[v][p] = 0;
    if (p) X[u] = X[v] = p;
    else update(u), update(v);
    return p;
  };
  auto flip = [&](int u, int c1, int c2) {
    int p = C[u][c1];
    swap(C[u][c1], C[u][c2]);
if (p) G[u][p] = G[p][u] = c2;
    if (!C[u][c1]) X[u] = c1;
    if (!C[u][c2]) X[u] = c2;
    return p;
  fill_n(X + 1, n, 1);
for (int t = 0; t < SZ(E); ++t) {
    int u = E[t
        ].X, v0 = E[t].Y, v = v0, c0 = X[u], c = c0, d;
    vector<pii> L;
    fill_n(vst + 1, n, 0);
    while (!G[u][v0]) {
      L.emplace_back(v, d = X[v]);
      if (!C[v][c]) for (int a = SZ(
          L) - 1; a >= 0; --a) c = color(u, L[a].X, c);
      else if (!C[u][d]) for (int a = SZ(L))
      ) - 1; a >= 0; --a) color(u, L[a].X, L[a].Y); else if (vst[d]) break;
      else vst[d] = 1, v = C[u][d];
    if (!G[u][v0]) {
      for (; v; v = flip(v, c, d), swap(c, d));
      if (int a; C[u][c0]) {
             a = SZ(L) - 2; a >= 0 && L[a].Y != c; --a);
        for (; a >= 0; --a) color(u, L[a].X, L[a].Y);
 }
} // namespace vizing
```

## 2.12 Minimum Clique Cover\* [879472]

```
struct Clique_Cover { // 0-base, 0(n2^n)
int co[1 << N], n, E[N];</pre>
```

```
int dp[1 << N];</pre>
   void init(int _n) {
     n = _n, fill_n(dp, 1 << n, 0);
     fill_n(E, n, 0), fill_n(co, 1 << n, 0);
   void add_edge(int u, int v) {
     E[u] |= 1 << v, E[v] |= 1 << u;
   int solve() {
     for (int i = 0; i < n; ++i)</pre>
      co[1 << i] = E[i] | (1 << i);
     co[0] = (1 << n) - 1;

dp[0] = (n & 1) * 2 - 1;
     for (int i = 1; i < (1 << n); ++i) {</pre>
       int t = i & -i;
       dp[i] = -dp[i ^ t];
       co[i] = co[i ^ t] & co[t];
     for (int i = 0; i < (1 << n); ++i)</pre>
       co[i] = (co[i] & i) == i;
     fwt(co, 1 << n, 1);
     for (int ans = 1; ans < n; ++ans) {</pre>
       int sum = 0; // probabilistic
       for (int i = 0; i < (1 << n); ++i)</pre>
         sum += (dp[i] *= co[i]);
       if (sum) return ans;
     return n;
};
```

## 2.13 NumberofMaximalClique\* [11fa26]

```
struct BronKerbosch { // 1-base
  int n, a[N], g[N][N];
  int S, all[N][N], some[N][N], none[N][N];
  void init(int _n) {
    n = _n;
    for (int i = 1; i <= n; ++i)</pre>
      for (int j = 1; j <= n; ++j) g[i][j] = 0;</pre>
  void add_edge(int u, int v) {
    g[u][v] = g[v][u] = 1;
  void dfs(int d, int an, int sn, int nn) {
    if (S > 1000) return; // pruning
    if (sn == 0 && nn == 0) ++S;
    int u = some[d][0];
    for (int i = 0; i < sn; ++i) {</pre>
      int v = some[d][i];
      if (g[u][v]) continue;
                   tnn = 0;
      copy_n(all[d], an, all[d + 1]);
      all[d + 1][an] = v;
      for (int j = 0; j < sn; ++j)</pre>
        if (g[v][some[d][j]])
          some[d + 1][tsn++] = some[d][j];
      for (int j = 0; j < nn; ++j)</pre>
        if (g[v][none[d][j]])
          none[d + 1][tnn++] = none[d][j];
      dfs(d + 1, an + 1, tsn, tnn);
      some[d][i] = 0, none[d][nn++] = v;
  int solve() {
    iota(some[0], some[0] + n, 1);
    S = 0, dfs(0, 0, n, 0);
    return S;
  }
};
```

# 3 Data Structure

#### 3.1 Discrete Trick

```
vector < int > val;
// build
sort(ALL
      (val)), val.resize(unique(ALL(val)) - val.begin());
// index of x
upper_bound(ALL(val), x) - val.begin();
// max idx <= x
upper_bound(ALL(val), x) - val.begin();
// max idx < x
lower_bound(ALL(val), x) - val.begin();</pre>
```

#### 3.2 zkwseg [aae74d]

```
template < class T, T (*op)(T, T), T (*e)()>
struct segtree {
    int n, sz;
    vector<T> d:
    explicit
          segtree(int n): segtree(vector<T>(n, e())) {}
     explicit segtree(const vector<T> &v) {
        n = v.size(), sz = 1;
while(sz < n) sz <<= 1;</pre>
         d.assign(sz*2, e());
         for(int i=0; i<n; ++i) d[sz+i] = v[i];</pre>
         for(int i=sz-1;
              i \ge 1; --i) d[i] = op(d[i << 1], d[i << 1|1]);
    void set(int p, const T &x) {
         p += sz, d[p] = x;
         while(p>>=1) d[p] = op(d[p<<1], d[p<<1|1]);
    T get(int p) const { return d[sz+p]; }
    T prod(int l, int r) const {
         T sl = e(), sr = e();
         l += sz, r += sz+1;
while(l < r) {
             if(l&1) sl = op(sl, d[l++]);
             if(r&1) sr = op(d[--r], sr);
             l >>= 1, r >>= 1;
         return op(sl. sr):
    }
};
3.3 BIT kth* [e39485]
int bit[N + 1]; // N = 2 ^ k
int query_kth(int k) {
    int res = 0;
```

```
int bit[N + 1]; // N = 2 ^ k
int query_kth(int k) {
   int res = 0;
   for (int i = N >> 1; i >= 1; i >>= 1)
      if (bit[res + i] < k)
            k -= bit[res += i];
   return res + 1;
}</pre>
```

#### 3.4 Interval Container\* [c54d29]

```
/* Add and
     remove intervals from a set of disjoint intervals.
   Will merge the added interval with
      any overlapping intervals in the set when adding.
 * Intervals are [inclusive, exclusive). */
set<pii>::
    iterator addInterval(set<pii>& is, int L, int R) {
  if (L == R) return is.end();
  auto it = is.lower_bound({L, R}), before = it;
  while (it != is.end() && it->X <= R) {</pre>
   R = max(R, it->Y);
    before = it = is.erase(it);
  if (it != is.begin() && (--it)->Y >= L) {
   L = min(L, it->X);
R = max(R, it->Y);
    is.erase(it);
  }
  return is.insert(before, pii(L, R));
void removeInterval(set<pii>& is, int L, int R) {
 if (L == R) return;
  auto it = addInterval(is, L, R);
  auto r2 = it->Y;
  if (it->X == L) is.erase(it);
  else (int&)it->Y = L;
  if (R != r2) is.emplace(R, r2);
```

#### 3.5 Leftist Tree [e91538]

```
struct node {
    ll v, data, sz, sum;
    node *l, *r;
    node(ll k)
        : v(0), data(k), sz(1), l(0), r(0), sum(k) {}
};
ll sz(node *p) { return p ? p->sz : 0; }
ll V(node *p) { return p ? p->sum : 0; }
node *merge(node *a, node *b) {
    if (!a || !b) return a ? a : b;
```

```
if (a->data < b->data) swap(a, b);
a->r = merge(a->r, b);
if (V(a->r) > V(a->l)) swap(a->r, a->l);
a->v = V(a->r) + 1, a->sz = sz(a->l) + sz(a->r) + 1;
a->sum = sum(a->l) + sum(a->r) + a->data;
return a;
}
void pop(node *&o) {
  node *tmp = o;
  o = merge(o->l, o->r);
  delete tmp;
}
```

### 3.6 Heavy light Decomposition\* [b004ae]

```
struct Heavy_light_Decomposition { // 1-base
  int n, ulink[N], deep[N], mxson[N], w[N], pa[N];
  int t, pl[N], data[N], val[N]; // val: vertex data
  vector<int> G[N];
  void init(int _n) {
    n = _n;
    for (int i = 1; i <= n; ++i)
      G[i].clear(), mxson[i] = 0;
  void add_edge(int a, int b) {
    G[a].pb(b), G[b].pb(a);
  void dfs(int u, int f, int d) {
    w[u] = 1, pa[u] = f, deep[u] = d++;
    for (int &i : G[u])
      if (i != f) {
         dfs(i, u, d), w[u] += w[i];
         if (w[mxson[u]] < w[i]) mxson[u] = i;
  void cut(int u, int link) {
  data[pl[u] = ++t] = val[u], ulink[u] = link;
    if (!mxson[u]) return;
    cut(mxson[u], link);
     for (int i : G[u])
      if (i != pa[u] && i != mxson[u])
         cut(i, i);
  void build() { dfs(1, 1, 1), cut(1, 1), /*build*/; }
int query(int a, int b) {
    int ta = ulink[a], tb = ulink[b], res = 0;
    while (ta != tb) {
      if (deep
           [ta] > deep[tb]) swap(ta, tb), swap(a, b);
         ' query(pl[tb], pl[b])
      tb = ulink[b = pa[tb]];
    if (pl[a] > pl[b]) swap(a, b);
    // query(pl[a], pl[b])
  }
};
```

#### 3.7 Centroid Decomposition\* [5a24da]

```
struct Cent_Dec { // 1-base
  vector<pli> G[N];
  pll info[N]; // store info. of itself
pll upinfo[N]; // store info. of climbing up
  int n, pa[N], layer[N], sz[N], done[N];
  ll dis[__lg(N) + 1][N];
  void init(int _n) {
    n = _n, layer[0] = -1;
    fill_n(pa + 1, n, 0), fill_n(done + 1, n, 0);
    for (int i = 1; i <= n; ++i) G[i].clear();</pre>
  void add_edge(int a, int b, int w) {
    G[a].pb(pll(b, w)), G[b].pb(pll(a, w));
  void get_cent(
    int u, int f, int &mx, int &c, int num) {
    int mxsz = 0;
    sz[u] = 1;
    for (pll e : G[u])
      if (!done[e.X] && e.X != f) {
        get_cent(e.X, u, mx, c, num);
        sz[u] += sz[e.X], mxsz = max(mxsz, sz[e.X]);
    if (mx > max(mxsz, num - sz[u]))
      mx = max(mxsz, num - sz[u]), c = u;
  void dfs(int u, int f, ll d, int org) {
    // if required, add self info or climbing info
```

```
dis[layer[org]][u] = d;
    for (pll e : G[u])
      if (!done[e.X] && e.X != f)
        dfs(e.X, u, d + e.Y, org);
  int cut(int u, int f, int num) {
    int mx = 1e9, c = 0, lc;
    get_cent(u, f, mx, c, num);
    done[c] = 1, pa[c] = f, layer[c] = layer[f] + 1;
    for (pll e : G[c])
      if (!done[e.X]) {
        if (sz[e.X] > sz[c])
        lc = cut(e.X, c, num - sz[c]);
else lc = cut(e.X, c, sz[e.X]);
        upinfo[lc] = pll(), dfs(e.X, c, e.Y, c);
    return done[c] = 0, c;
  void build() { cut(1, 0, n); }
  void modify(int u) {
    for (int a = u, ly = layer[a]; a;
         a = pa[a], --ly) {
      info[a].X += dis[ly][u], ++info[a].Y;
      if (pa[a])
        upinfo[a].X += dis[ly - 1][u], ++upinfo[a].Y;
    }
  ll query(int u) {
    ll rt = 0;
    for (int a = u, ly = layer[a]; a;
         a = pa[a], --ly) {
      rt += info[a].X + info[a].Y * dis[ly][u];
      if (pa[a])
          upinfo[a].X + upinfo[a].Y * dis[ly - 1][u];
    return rt;
  }
};
```

#### 3.8 LiChaoST\* [4a4bee]

```
struct L {
  ll m, k, id;
  L(): id(-1) {}
  L(ll a, ll b, ll c): m(a), k(b), id(c) {} ll at(ll x) { return m * x + k; }
class LiChao { // maintain max
private:
  int n; vector<L> nodes;
  void insert(int l, int r, int rt, L ln) {
    int m = (l + r) >> 1;
    if (nodes[rt].id == -1)
      return nodes[rt] = ln, void();
    bool atLeft = nodes[rt].at(l) < ln.at(l);</pre>
    if (nodes[rt].at(m) < ln.at(m))</pre>
      atLeft ^= 1, swap(nodes[rt], ln);
    if (r - l == 1) return;
    if (atLeft) insert(l, m, rt << 1, ln);
else insert(m, r, rt << 1 | 1, ln);</pre>
  ll query(int l, int r, int rt, ll x) {
    int m = (l + r) >> 1; ll ret = -INF;
    if (nodes[rt].id != -1) ret = nodes[rt].at(x);
    if (r - l == 1) return ret;
    if (x
         < m) return max(ret, query(l, m, rt << 1, x));</pre>
    return max(ret, query(m, r, rt << 1 | 1, x));</pre>
public:
 LiChao(int n_) : n(n_), nodes(n * 4) {}
  void insert(L ln) { insert(0, n, 1, ln); }
 ll query(ll x) { return query(0, n, 1, x); }
```

#### 3.9 Link cut tree\* [a35b5d]

```
int dir()
  { return f->ch[0] == this ? 0 : 1; }
  void setCh(Splay *c, int d) {
    ch[d] = c;
    if (c != &nil) c->f = this;
    pull();
  void give_tag(int r) {
    if (r) swap(ch[0], ch[1]), rev ^= 1;
  void push() {
    if (ch[0] != &nil) ch[0]->give_tag(rev);
    if (ch[1] != &nil) ch[1]->give_tag(rev);
 void pull() {
   // take care of the nil!
    size = ch[0] -> size + ch[1] -> size + 1;
    sum = ch[0]->sum ^ ch[1]->sum ^ val;
    if (ch[0] != &nil) ch[0]->f = this;
    if (ch[1] != &nil) ch[1]->f = this;
} Splay::nil;
Splay *nil = &Splay::nil;
void rotate(Splay *x) {
  Splay *p = x - > f;
  int d = x->dir();
  if (!p->isr()) p->f->setCh(x, p->dir());
  else x->f = p->f;
  p->setCh(x->ch[!d], d);
  x->setCh(p, !d);
  p->pull(), x->pull();
void splay(Splay *x) {
  vector < Splay*> splayVec;
  for (Splay *q = x;; q = q->f) {
    splayVec.pb(q);
    if (q->isr()) break;
  reverse(ALL(splayVec));
  for (auto it : splayVec) it->push();
  while (!x->isr()) {
    if (x->f->isr()) rotate(x);
    else if (x->dir() == x->f->dir())
      rotate(x->f), rotate(x);
    else rotate(x), rotate(x);
 }
Splay* access(Splay *x) {
  Splay *q = nil;
  for (; x != nil; x = x->f)
    splay(x), x -> setCh(q, 1), q = x;
void root_path(Splay *x) { access(x), splay(x); }
void chroot(Splay *x){
 root_path(x), x->give_tag(1);
  x->push(), x->pull();
void split(Splay *x, Splay *y) {
  chroot(x), root_path(y);
void link(Splay *x, Splay *y) {
 root_path(x), chroot(y);
  x->setCh(y, 1);
void cut(Splay *x, Splay *y) {
  split(x, y);
  if (y->size != 5) return;
  y->push();
  y - ch[0] = y - ch[0] - f = nil;
Splay* get_root(Splay *x) {
  for (root_path(x); x->ch[0] != nil; x = x->ch[0])
   x->push();
  splay(x);
  return x;
bool conn(Splay *x, Splay *y) {
 return get_root(x) == get_root(y);
Splay* lca(Splay *x, Splay *y) {
 access(x), root_path(y);
  if (y->f == nil) return y;
  return y->f;
```

```
void change(Splay *x, int val) {
  splay(x), x->val = val, x->pull();
int query(Splay *x, Splay *y) {
  split(x, y);
  return y->sum;
3.10 KDTree [375ca2]
namespace kdt {
int root, lc[maxn], rc[maxn], xl[maxn], xr[maxn],
  yl[maxn], yr[maxn];
point p[maxn];
int build(int l, int r, int dep = 0) {
  if (l == r) return -1;
  function < bool(const point &, const point &) > f =
    [dep](const point &a, const point &b) {
       if (dep & 1) return a.x < b.x;</pre>
       else return a.y < b.y;</pre>
    };
  int m = (l + r) >> 1;
  nth_element(p + l, p + m, p + r, f);
  xl[m] = xr[m] = p[m].x;
  yl[m] = yr[m] = p[m].y;
  lc[m] = build(l, m, dep + 1);
  if (~lc[m]) {
    xl[m] = min(xl[m], xl[lc[m]]);
    xr[m] = max(xr[m], xr[lc[m]]);
    yl[m] = min(yl[m], yl[lc[m]]);
yr[m] = max(yr[m], yr[lc[m]]);
  rc[m] = build(m + 1, r, dep + 1);
  if (~rc[m]) {
    xl[m] = min(xl[m], xl[rc[m]]);
    xr[m] = max(xr[m], xr[rc[m]]);
    yl[m] = min(yl[m], yl[rc[m]]);
    yr[m] = max(yr[m], yr[rc[m]]);
  return m;
bool bound(const point &q, int o, long long d) {
  double ds = sqrt(d + 1.0);
  if (q.x < xl[o] - ds || q.x > xr[o] + ds ||
    q.y < yl[o] - ds || q.y > yr[o] + ds)
    return false;
  return true;
long long dist(const point &a, const point &b) {
  return (a.x - b.x) * 1ll * (a.x - b.x) +
      (a.y - b.y) * 1ll * (a.y - b.y);
void dfs(
  const point &q, long long &d, int o, int dep = \theta) {
  if (!bound(q, o, d)) return;
  long long cd = dist(p[o], q);
  if (cd != 0) d = min(d, cd);
  if ((dep & 1) && q.x < p[o].x ||</pre>
    !(dep & 1) && q.y < p[o].y) {
if (~lc[o]) dfs(q, d, lc[o], dep + 1);
if (~rc[o]) dfs(q, d, rc[o], dep + 1);
  } else {
    if (~rc[o]) dfs(q, d, rc[o], dep + 1);
    if (~lc[o]) dfs(q, d, lc[o], dep + 1);
  }
void init(const vector<point> &v) {
  for (int i = 0; i < v.size(); ++i) p[i] = v[i];</pre>
  root = build(0, v.size());
long long nearest(const point &q) {
  long long res = 1e18;
  dfs(q, res, root);
  return res:
} // namespace kdt
4 Flow/Matching
4.1 Bipartite Matching* [784535]
struct Bipartite_Matching { // 0-base
  int mp[N], mq[N], dis[N + 1], cur[N], l, r;
  vector<int> G[N + 1];
```

```
bool dfs(int u) {
  for (int &i = cur[u]; i < SZ(G[u]); ++i) {</pre>
    int e = G[u][i];
```

```
if (mq[e] == l
             || (dis[mq[e]] == dis[u] + 1 && dfs(mq[e])))
         return mp[mq[e] = u] = e, 1;
     return dis[u] = -1, 0;
   bool bfs() {
     queue < int > q;
     fill_n(dis, l + 1, -1);
     for (int i = 0; i < l; ++i)</pre>
       if (!~mp[i])
         q.push(i), dis[i] = 0;
     while (!q.empty()) {
       int u = q.front();
       q.pop();
       for (int e : G[u])
         if (!~dis[mq[e]])
           q.push(mq[e]), dis[mq[e]] = dis[u] + 1;
     return dis[l] != -1;
   int matching() {
     int res = 0;
     fill_n(mp, l, -1), fill_n(mq, r, l);
     while (bfs()) {
       fill_n(cur, l, 0);

for (int i = 0; i < l; ++i)
         res += (!~mp[i] && dfs(i));
     return res; // (i, mp[i] != -1)
   void add_edge(int s, int t) { G[s].pb(t); }
   void init(int _l, int _r) {
     l = _l, r = _r;
for (int i = 0; i <= l; ++i)</pre>
       G[i].clear();
  }
};
4.2 Kuhn Munkres* [4b3863]
```

```
struct KM { // 0-base, maximum matching
  ll w[N][N], hl[N], hr[N], slk[N];
int fl[N], fr[N], pre[N], qu[N], ql, qr, n;
  bool vl[N], vr[N];
  void init(int _n) {
    n = _n;
for (int i = 0; i < n; ++i)</pre>
      fill_n(w[i], n, -INF);
  void add_edge(int a, int b, ll wei) {
    w[a][b] = wei;
  bool Check(int x) {
    if (vl[x] = 1, \sim fl[x])
      return vr[qu[qr++] = fl[x]] = 1;
    while (\sim x) swap(x, fr[fl[x] = pre[x]]);
    return 0;
  void bfs(int s) {
    fill_n(slk
         , n, INF), fill_n(vl, n, 0), fill_n(vr, n, 0);
    ql = qr = 0, qu[qr++] = s, vr[s] = 1;
    for (ll d;;) {
      while (ql < qr)</pre>
        for (int x = 0, y = qu[ql++]; x < n; ++x)
  if (!vl[x] && slk</pre>
               [x] >= (d = hl[x] + hr[y] - w[x][y])) {
             if (pre[x] = y, d) slk[x] = d;
             else if (!Check(x)) return;
      d = INF;
      for (int x = 0; x < n; ++x)
        if (!vl[x] && d > slk[x]) d = slk[x];
      for (int x = 0; x < n; ++x) {</pre>
        if (vl[x]) hl[x] += d;
         else slk[x] -= d;
         if (vr[x]) hr[x] -= d;
      for (int x = 0; x < n; ++x)
        if (!vl[x] && !slk[x] && !Check(x)) return;
    }
  il solve() {
    fill_n(fl
          n, -1), fill_n(fr, n, -1), fill_n(hr, n, 0);
    for (int i = 0; i < n; ++i)</pre>
```

```
\label{eq:hline} \begin{split} &\text{hl[i]} = \text{*max\_element(w[i], w[i] + n);} \\ &\text{for (int } i = 0; \ i < n; \ \text{++i) bfs(i);} \end{split}
        ll res = 0;
         for (int i = 0; i < n; ++i) res += w[i][fl[i]];</pre>
        return res;
   }
};
```

## 4.3 MincostMaxflow\* [1c78db]

```
struct MinCostMaxFlow { // 0-base
  struct Edge {
    ll from, to, cap, flow, cost, rev;
  } *past[N];
  vector < Edge > G[N];
  int inq[N], n, s, t;
  ll dis[N], up[N], pot[N];
  bool BellmanFord() {
    fill_n(dis, n, INF), fill_n(inq, n, 0);
    queue<int> q;
    auto relax = [&](int u, ll d, ll cap, Edge *e) {
      if (cap > 0 && dis[u] > d) {
        dis[u] = d, up[u] = cap, past[u] = e;
        if (!inq[u]) inq[u] = 1, q.push(u);
      }
    };
    relax(s, 0, INF, 0);
    while (!q.empty()) {
      int u = q.front();
       q.pop(), inq[u] = 0;
      for (auto &e : G[u]) {
        ll d2 = dis[u] + e.cost + pot[u] - pot[e.to];
         relax
             (e.to, d2, min(up[u], e.cap - e.flow), &e);
      }
    }
    return dis[t] != INF;
  void solve(int
      , int _t, ll &flow, ll &cost, bool neg = true) {
= _s, t = _t, flow = 0. cost = A:
    if (neg) BellmanFord(), copy_n(dis, n, pot);
    for (; BellmanFord(); copy_n(dis, n, pot)) {
      for (int
          i = 0; i < n; ++i) dis[i] += pot[i] - pot[s];
       flow += up[t], cost += up[t] * dis[t];
      for (int i = t; past[i]; i = past[i]->from) {
        auto &e = *past[i];
        e.flow += up[t], G[e.to][e.rev].flow -= up[t];
      }
    }
  void init(int _n) {
    n = _n, fill_n(pot, n, \theta);
    for (int i = 0; i < n; ++i) G[i].clear();</pre>
  void add_edge(ll a, ll b, ll cap, ll cost) {
    G[a].pb(Edge{a, b, cap, 0, cost, SZ(G[b])});
    G[b].pb(Edge{b, a, 0, 0, -cost, SZ(G[a]) - 1});
  }
};
```

#### Maximum Simple Graph Matching\* [0fe1c3]

```
struct Matching { // 0-base
 queue < int > q; int n;
vector < int > fa, s, vis, pre, match;
  vector<vector<int>> G;
  int Find(int u)
  { return u == fa[u] ? u : fa[u] = Find(fa[u]); }
  int LCA(int x, int y) {
    static int tk = 0; tk++; x = Find(x); y = Find(y);
    for (;; swap(x, y)) if (x != n) {
      if (vis[x] == tk) return x;
      vis[x] = tk;
      x = Find(pre[match[x]]);
    }
  void Blossom(int x, int y, int l) {
    for (; Find(x) != l; x = pre[y]) {
      pre[x] = y, y = match[x];
if (s[y] == 1) q.push(y), s[y] = 0;
      for (int z: {x, y}) if (fa[z] == z) fa[z] = l;
    }
  bool Bfs(int r) {
    iota(ALL(fa), 0); fill(ALL(s), -1);
```

```
q = queue < int > (); q.push(r); s[r] = 0;
      for (; !q.empty(); q.pop()) {
         for (int x = q.front(); int u : G[x])
  if (s[u] == -1) {
              if (pre[u] = x, s[u] = 1, match[u] == n) {
                 for (int a = u, b = x, last;
    b != n; a = last, b = pre[a])
                    last =
                         match[b], match[b] = a, match[a] = b;
                 return true;
              q.push(match[u]); s[match[u]] = 0;
           } else if (!s[u] && Find(u) != Find(x)) {
              int l = LCA(u, x);
Blossom(x, u, l); Blossom(u, x, l);
      return false;
   \label{eq:matching} \texttt{Matching(int } \_\texttt{n)} \; : \; \texttt{n(}\_\texttt{n)}, \; \texttt{fa(}\texttt{n} \; + \; \texttt{1)}, \; \texttt{s(}\texttt{n} \; + \; \texttt{1)}, \; \texttt{vis}
   (n + 1), pre(n + 1, n), match(n + 1, n), G(n) {} void add_edge(int u, int v)
   { G[u].pb(v), G[v].pb(u); }
   int solve() {
      int ans = 0;
      for (int x = 0; x < n; ++x)
        if (match[x] == n) ans += Bfs(x);
      return ans;
   } // match[x] == n means not matched
};
```

#### 4.5 Maximum Weight Matching\* [9ffb94]

```
#define REP(i, l, r) for (int i=(l); i<=(r); ++i)</pre>
struct WeightGraph { // 1-based
  struct edge { int u, v, w; }; int n, nx;
  vector<int> lab; vector<vector<edge>> g;
  vector<int> slk, match, st, pa, S, vis;
  vector<vector<int>> flo, flo_from; queue<int> q;
  WeightGraph(int n_{-}): n(n_{-}), nx(n * 2), lab(nx + 1),
    g(nx + 1, vector < edge > (nx + 1)), slk(nx + 1), flo(nx + 1), flo_from(nx + 1, vector(n + 1, 0)) {
    match = st = pa = S = vis = slk;
    REP(u, 1, n) REP(v, 1, n) g[u][v] = \{u, v, \theta\};
  int E(edge e)
  { return lab[e.u] + lab[e.v] - g[e.u][e.v].w * 2; }
  void update_slk(int u, int x, int &s)
  { if (!s || E(g[u][x]) < E(g[s][x])) s = u; }
  void set_slk(int x) {
    slk[x] = 0;
    REP(u, 1, n)
      if (g[u][x].w > 0 && st[u] != x && S[st[u]] == 0)
        update_slk(u, x, slk[x]);
  void q_push(int x) {
    if (x <= n) q.push(x);</pre>
    else for (int y : flo[x]) q_push(y);
  void set_st(int x, int b) {
    st[x] = b:
    if (x > n) for (int y : flo[x]) set_st(y, b);
  vector<int> split_flo(auto &f, int xr) {
    auto it = find(ALL(f), xr);
    if (auto pr = it - f.begin(); pr % 2 == 1)
      reverse(1 + ALL(f)), it = f.end() - pr;
    auto res = vector(f.begin(), it);
    return f.erase(f.begin(), it), res;
  void set_match(int u, int v) {
    match[u] = g[u][v].v;
    if (u <= n) return;</pre>
    int xr = flo_from[u][g[u][v].u];
    auto &f = flo[u], z = split_flo(f, xr);
REP(i, 0, SZ(z) - 1) set_match(z[i], z[i ^ 1]);
    set_match(xr, v); f.insert(f.end(), ALL(z));
  void augment(int u, int v) {
    for (;;) {
      int xnv = st[match[u]]; set_match(u, v);
      if (!xnv) return;
      set_match(v = xnv, u = st[pa[xnv]]);
  int lca(int u, int v) {
    static int t = 0; ++t;
```

```
for (++t; u || v; swap(u, v)) if (u) {
   if (vis[u] == t) return u;
    vis[u] = t, u = st[match[u]];
    if (u) u = st[pa[u]];
 return 0;
void add_blossom(int u, int o, int v) {
 int b = find(n + 1 + ALL(st), 0) - begin(st);
 lab[b] = 0, S[b] = 0, match[b] = match[o];
  vector<int> f = {o};
 for (int t : \{u, v\}) {
    reverse(1 + ALL(f));
    for (int x = t, y; x != o; x = st[pa[y]])
      f.pb(x), f.pb(y = st[match[x]]), q_push(y);
  flo[b] = f; set_st(b, b);
 REP(x, 1, nx) g[b][x].w = g[x][b].w = 0;
  fill(ALL(flo_from[b]), 0);
  for (int xs : flo[b]) {
   REP(x, 1, nx)
      if (g[b][x].w == 0 \mid \mid E(g[xs][x]) < E(g[b][x]))
        g[b][x] = g[xs][x], g[x][b] = g[x][xs];
    REP(x, 1, n)
      if (flo_from[xs][x]) flo_from[b][x] = xs;
  set_slk(b);
void expand_blossom(int b) {
 for (int x : flo[b]) set_st(x, x);
  int xr = flo_from[b][g[b][pa[b]].u], xs = -1;
  for (int x : split_flo(flo[b], xr)) {
    if (xs == -1) { xs = x; continue; }
    pa[xs] = g[x][xs].u, S[xs] = 1, S[x] = 0;
    slk[xs] = 0, set_slk(x), q_push(x), xs = -1;
  for (int x : flo[b])
    if (x == xr) S[x] = 1, pa[x] = pa[b];
    else S[x] = -1, set_slk(x);
 st[b] = 0;
bool on_found_edge(const edge &e) {
 if (int u = st[e.u], v = st[e.v]; S[v] == -1) {
   int nu = st[match[v]]; pa[v] = e.u; S[v] = 1;
   slk[v] = slk[nu] = S[nu] = 0; q_push(nu);
 } else if (S[v] == 0) {
    if (int o = lca(u, v)) add_blossom(u, o, v);
    else return augment(u, v), augment(v, u), true;
 return false:
bool matching() {
 fill(ALL(S), -1), fill(ALL(slk), 0);
  q = queue < int > ();
 REP(x, 1, nx) if (st[x] == x && !match[x])
    pa[x] = S[x] = 0, q_push(x);
  if (q.empty()) return false;
  for (;;) {
    while (SZ(q)) {
      int u = q.front(); q.pop();
      if (S[st[u]] == 1) continue;
      REP(v, 1, n)
        if (g[u][v].w > 0 && st[u] != st[v]) {
          if (E(g[u][v]) != 0)
            update_slk(u, st[v], slk[st[v]]);
          else if
               (on_found_edge(g[u][v])) return true;
    int d = INF;
   REP(b, n + 1, nx) if (st[b] == b && S[b] == 1)
d = min(d, lab[b] / 2);
    REP(x, 1, nx)
      if (int
           s = slk[x]; st[x] == x && s && s[x] <= 0)
        d = min(d, E(g[s][x]) / (S[x] + 2));
    REP(u, 1, n)
      if (S[st[u]] == 1) lab[u] += d;
      else if (S[st[u]] == 0) {
  if (lab[u] <= d) return false;</pre>
        lab[u] -= d;
    REP(b, n + 1, nx) if (st[b] == b && S[b] >= 0)
      lab[b] += d * (2 - 4 * S[b]);
    REP(x, 1, nx)
      if (int s = slk[x]; st[x] == x &&
```

```
s \&\& st[s] != x \&\& E(g[s][x]) == 0)
           if (on_found_edge(g[s][x])) return true;
      REP(b, n + 1, nx)
         if (st[b] == b && S[b] == 1 && lab[b] == 0)
           expand_blossom(b);
    return false:
  pair<ll, int> solve() {
    fill(ALL(match), 0);
    REP(u, 0, n) st[u] = u, flo[u].clear();
    int w_max = 0;
    REP(u, 1, n) REP(v, 1, n) {
  flo_from[u][v] = (u == v ? u : 0);
       w_{max} = max(w_{max}, g[u][v].w);
    fill(ALL(lab), w_max);
    int n_matches = 0; ll tot_weight = 0;
    while (matching()) ++n_matches;
    REP(u, 1, n) if (match[u] && match[u] < u)</pre>
      tot_weight += g[u][match[u]].w;
    return make_pair(tot_weight, n_matches);
  void add_edge(int u, int v, int w)
  { g[u][v].w = g[v][u].w = w; }
};
```

### 4.6 SW-mincut [c705f5]

```
struct SW{ // global min cut, O(V^3)
   #define REP for (int i = 0; i < n; ++i)
   static const int MXN = 514, INF = 2147483647;</pre>
  int vst[MXN], edge[MXN][MXN], wei[MXN];
  void init(int n) {
    REP fill_n(edge[i], n, 0);
  void addEdge(int u, int v, int w){
    edge[u][v] += w; edge[v][u] += w;
  int search(int &s, int &t, int n){
    fill_n(vst, n, 0), fill_n(wei, n, 0);
     s = t = -1;
    int mx, cur;
     for (int j = 0; j < n; ++j) {</pre>
       mx = -1, cur = 0;
       REP if (wei[i] > mx) cur = i, mx = wei[i];
       vst[cur] = 1, wei[cur] = -1;
       s = t; t = cur;
       REP if (!vst[i]) wei[i] += edge[cur][i];
    }
    return mx;
  int solve(int n) {
    int res = INF;
    for (int x, y; n > 1; n--){
  res = min(res, search(x, y, n));
       REP edge[i][x] = (edge[x][i] += edge[y][i]);
       REP {
         edge[y][i] = edge[n - 1][i];
         edge[i][y] = edge[i][n - 1];
       }
    return res;
  }
} sw;
```

#### 4.7 BoundedFlow\*(Dinic\*) [4a793f]

```
struct BoundedFlow { // O-base
    struct edge {
        int to, cap, flow, rev;
    };
    vector<edge> G[N];
    int n, s, t, dis[N], cur[N], cnt[N];
    void init(int _n) {
        n = _n;
        for (int i = 0; i < n + 2; ++i)
            G[i].clear(), cnt[i] = 0;
    }
    void add_edge(int u, int v, int lcap, int rcap) {
        cnt[u] -= lcap, cnt[v] += lcap;
        G[u].pb(edge{v, rcap, lcap, SZ(G[v])});
        G[v].pb(edge{u, 0, 0, SZ(G[u]) - 1});
    }
    void add_edge(int u, int v, int cap) {
        G[u].pb(edge{v, cap, 0, SZ(G[v])});
        G[v].pb(edge{u, 0, 0, SZ(G[u]) - 1});
    }
}</pre>
```

```
int dfs(int u, int cap) {
    if (u == t || !cap) return cap;
    for (int &i = cur[u]; i < SZ(G[u]); ++i) {</pre>
      edge &e = G[u][i];
      if (dis[e.to] == dis[u] + 1 && e.cap != e.flow) {
         int df = dfs(e.to, min(e.cap - e.flow, cap));
        if (df) {
          e.flow += df, G[e.to][e.rev].flow -= df;
          return df;
        }
      }
    dis[u] = -1;
    return 0;
  bool bfs() {
    fill_n(dis, n + 3, -1);
    queue < int > q;
    q.push(s), dis[s] = 0;
    while (!q.empty()) {
      int u = q.front();
      q.pop();
      for (edge &e : G[u])
        if (!~dis[e.to] && e.flow != e.cap)
          q.push(e.to), dis[e.to] = dis[u] + 1;
    return dis[t] != -1;
  int maxflow(int _s, int _t) {
    s = _s, t = _t;
int flow = 0, df;
    while (bfs()) {
      fill_n(cur, n + 3, 0);
      while ((df = dfs(s, INF))) flow += df;
    return flow;
  bool solve() {
    int sum = 0;
    for (int i = 0; i < n; ++i)</pre>
      if (cnt[i] > 0)
        add_edge(n + 1, i, cnt[i]), sum += cnt[i];
      else if (cnt[i] < 0) add_edge(i, n + 2, -cnt[i]);</pre>
    if (sum != maxflow(n + 1, n + 2)) sum = -1;
    for (int i = 0; i < n; ++i)</pre>
      if (cnt[i] > 0)
        G[n + 1].pop_back(), G[i].pop_back();
      else if (cnt[i] < 0)</pre>
        G[i].pop_back(), G[n + 2].pop_back();
    return sum != -1;
  int solve(int _s, int _t) {
    add_edge(_t, _s, INF);
if (!solve()) return -1; // invalid flow
    int x = G[_t].back().flow;
    return G[_t].pop_back(), G[_s].pop_back(), x;
};
```

### 4.8 Gomory Hu tree\* [11be99]

```
MaxFlow Dinic;
int g[MAXN];
void GomoryHu(int n) { // 0-base
  fill_n(g, n, 0);
  for (int i = 1; i < n; ++i) {
    Dinic.reset();
    add_edge(i, g[i], Dinic.maxflow(i, g[i]));
    for (int j = i + 1; j <= n; ++j)
        if (g[j] == g[i] && ~Dinic.dis[j])
        g[j] = i;
  }
}</pre>
```

#### 4.9 Minimum Cost Circulation\* [ba97cf]

```
struct MinCostCirculation { // 0-base
    struct Edge {
        ll from, to, cap, fcap, flow, cost, rev;
    } *past[N];
    vector<Edge> G[N];
    ll dis[N], inq[N], n;
    void BellmanFord(int s) {
        fill_n(dis, n, INF), fill_n(inq, n, 0);
        queue<int> q;
        auto relax = [&](int u, ll d, Edge *e) {
```

```
if (dis[u] > d) {
         dis[u] = d, past[u] = e;
         if (!inq[u]) inq[u] = 1, q.push(u);
     relax(s, 0, 0);
     while (!q.empty()) {
       int u = q.front();
       q.pop(), inq[u] = 0;
       for (auto &e : G[u])
         if (e.cap > e.flow)
           relax(e.to, dis[u] + e.cost, &e);
     }
  void try_edge(Edge &cur) {
     if (cur.cap > cur.flow) return ++cur.cap, void();
     BellmanFord(cur.to);
     if (dis[cur.from] + cur.cost < 0) {</pre>
       ++cur.flow, --G[cur.to][cur.rev].flow;
       for (int
         i = cur.from; past[i]; i = past[i]->from) {
auto &e = *past[i];
         ++e.flow, --G[e.to][e.rev].flow;
     }
     ++cur.cap;
  }
  void solve(int mxlg) {
     for (int b = mxlg; b >= 0; --b) {
  for (int i = 0; i < n; ++i)</pre>
         for (auto &e : G[i])
           e.cap *= 2, e.flow *= 2;
       for (int i = 0; i < n; ++i)</pre>
         for (auto &e : G[i])
           if (e.fcap >> b & 1)
              try_edge(e);
    }
  }
  void init(int _n) { n = _n;
    for (int i = 0; i < n; ++i) G[i].clear();</pre>
   void add_edge(ll a, ll b, ll cap, ll cost) {
     G[a].pb(Edge
          {a, b, 0, cap, 0, cost, SZ(G[b]) + (a == b)});
     G[b].pb(Edge{b, a, 0, 0, 0, -cost, SZ(G[a]) - 1});
} mcmf; // O(VE * ElogC)
```

## 4.10 Flow Models

- Maximum/Minimum flow with lower bound / Circulation problem
  - 1. Construct super source S and sink T.
  - 2. For each edge (x,y,l,u), connect  $x \rightarrow y$  with capacity u-l.
  - 3. For each vertex v, denote by in(v) the difference between the sum of incoming lower bounds and the sum of outgoing lower bounds.
  - 4. If in(v)>0, connect  $S\to v$  with capacity in(v), otherwise, connect  $v\to T$  with capacity -in(v).
    - To maximize, connect  $t \to s$  with capacity  $\infty$  (skip this in circulation problem), and let f be the maximum flow from S to T. If  $f \neq \sum_{v \in V, in(v) > 0} in(v)$ , there's no solution. Otherwise, the maximum flow from s to t is the answer.
    - To minimize, let f be the maximum flow from S to T. Connect  $t \to s$  with capacity  $\infty$  and let the flow from S to T be f'. If  $f+f' \neq \sum_{v \in V, in(v)>0} in(v)$ , there's no solution. Otherwise, f' is the answer.
  - 5. The solution of each edge e is  $l_e+f_e$ , where  $f_e$  corresponds to the flow of edge e on the graph.
- Construct minimum vertex cover from maximum matching M on bipartite graph (X,Y)
  - 1. Redirect every edge:  $y \rightarrow x$  if  $(x,y) \in M$ ,  $x \rightarrow y$  otherwise.
  - 2. DFS from unmatched vertices in X.
  - 3.  $x \in X$  is chosen iff x is unvisited.
  - 4.  $y \in Y$  is chosen iff y is visited.
- · Minimum cost cyclic flow
  - 1. Consruct super source S and sink T
  - 2. For each edge (x,y,c), connect  $x \to y$  with (cost,cap) = (c,1) if c > 0, otherwise connect  $y \to x$  with (cost,cap) = (-c,1)
  - 3. For each edge with  $c\,{<}\,0$  , sum these cost as K , then increase d(y) by 1, decrease d(x) by 1
  - 4. For each vertex v with d(v)>0, connect  $S\to v$  with (cost,cap)=(0,d(v))
  - 5. For each vertex v with d(v)<0, connect  $v\to T$  with (cost,cap)=(0,-d(v))
  - 6. Flow from S to T, the answer is the cost of the flow C+K
- Maximum density induced subgraph
  - 1. Binary search on answer, suppose we're checking answer  ${\cal T}$

- 2. Construct a max flow model, let K be the sum of all weights
- 3. Connect source  $s \rightarrow v$ ,  $v \in G$  with capacity K
- 4. For each edge (u,v,w) in G, connect  $u \to v$  and  $v \to u$  with capacity w
- 5. For  $v\in G$ , connect it with sink  $v\to t$  with capacity  $K+2T-(\sum_{e\in E(v)}w(e))-2w(v)$
- 6. T is a valid answer if the maximum flow f < K|V|
- · Minimum weight edge cover
  - 1. For each  $v \in V$  create a copy v', and connect  $u' \to v'$  with weight w(u,v).
  - 2. Connect  $v \to v'$  with weight  $2\mu(v)$ , where  $\mu(v)$  is the cost of the cheapest edge incident to v.
  - 3. Find the minimum weight perfect matching on G'.
- Project selection problem
  - 1. If  $p_v > 0$ , create edge (s,v) with capacity  $p_v$ ; otherwise, create edge (v,t) with capacity  $-p_v$ .
  - 2. Create edge (u,v) with capacity w with w being the cost of choosing u without choosing v.
- 3. The mincut is equivalent to the maximum profit of a subset of projects.
- · Dual of minimum cost maximum flow
  - 1. Capacity  $c_{uv}$  , Flow  $f_{uv}$  , Cost  $w_{uv}$  , Required Flow difference for vertex  $b_u$  .
  - 2. If all  $w_{uv}$  are integers, then optimal solution can happen when all  $p_u$  are integers.

$$\begin{aligned} & \min \sum_{uv} w_{uv} f_{uv} \\ & -f_{uv} \geq -c_{uv} \Leftrightarrow \min \sum_{u} b_{u} p_{u} + \sum_{uv} c_{uv} \max(0, p_{v} - p_{u} - w_{uv}) \\ & \sum_{v} f_{vu} - \sum_{v} f_{uv} = -b_{u} \end{aligned}$$

## 5 String

#### 5.1 KMP [db6286]

```
vector<int> fail_func(const string &s) {
    int n = s.size();
    vector<int> f(n);
    for(int i=1; i<n; ++i) {
        int j = f[i-1];
        while(j > 0 && s[i] != s[j]) j = f[j-1];
        f[i] = j + (s[i] == s[j]);
    }
    return f;
}
int kmp_count(const string &s, const string &p) {
    int n = s.size(), m = p.size();
    vector<int> f = fail_func(p+"$");
    int cnt = 0, j = 0;
    for(int i=0; i<n; ++i) {
        while(j > 0 && s[i] != p[j]) j = f[j-1];
        j += (s[i] == p[j]);
        if(j == m) cnt++;
    }
    return cnt;
}
```

#### 5.2 Z-value\* [e4e6b3]

```
vi z_value(string s) {
   int n = s.size(), b = 0;
   vi z(n);
   for(int i=1; i<n; ++i) {
      if(z[b]+b < i) z[i] = 0;
      else z[i] = min(z[b]+b-i, z[i-b]);
      while(s[z[i]] == s[z[i]+i]) z[i]++;
      if(z[i]+i > z[b]+b) b = i;
   }
   return z;
}
```

#### 5.3 Manacher\* [1ad8ef]

## 5.4 SAIS\* [6f26bc]

```
auto sais(const auto &s) {
   const int n = SZ(s), z = ranges::max(s) + 1;
   if (n == 1) return vector{0};
  vector<int> c(z); for (int x : s) ++c[x];
  partial_sum(ALL(c), begin(c));
  vector < int > sa(n); auto I = views::iota(0, n);
  vector < bool > t(n, true);
for (int i = n - 2; i >= 0; --i)
     t[i] = (
          s[i] == s[i + 1] ? t[i + 1] : s[i] < s[i + 1]);
   auto is_lms = views::filter([&t](int x) {
     return x && t[x] && !t[x - 1];
  });
   auto induce = [&] {
     for (auto x = c; int y : sa)
       if (y--) if (!t[y]) sa[x[s[y] - 1]++] = y;
     for (auto x = c; int y : sa | views::reverse)
       if (y--) if (t[y]) sa[--x[s[y]]] = y;
  };
  vector < int > lms, q(n); lms.reserve(n);
for (auto x = c; int i : I | is_lms)
    q[i] = SZ(lms), lms.pb(sa[--x[s[i]]] = i);
   induce(); vector < int > ns(SZ(lms));
for (int j = -1, nz = 0; int i : sa | is_lms) {
     if (j >= 0) {
       int len = min({n - i, n - j, lms[q[i] + 1] - i});
       begin(s) + i, begin(s) + i + len);
     j = i;
  fill(ALL(sa), 0); auto nsa = sais(ns);
for (auto x = c; int y : nsa | views::reverse)
     y = lms[y], sa[--x[s[y]]] = y;
  return induce(), sa;
// sa[i]: sa[i]-th suffix
       is the i-th lexicographically smallest suffix.
// hi[i]: LCP of suffix sa[i] and suffix sa[i - 1].
struct Suffix {
  int n; vector<int> sa, hi, ra;
  Suffix
     (const auto &_s, int _n) : n(_n), hi(n), ra(n) {
vector < int > s(n + 1); // s[n] = 0;
copy_n(_s, n, begin(s)); // _s shouldn't contain 0
     sa = sais(s); sa.erase(sa.begin());
     for (int i = 0; i < n; ++i) ra[sa[i]] = i;
for (int i = 0, h = 0; i < n; ++i) {</pre>
        if (!ra[i]) { h = 0; continue; }
        for (int j = sa[ra[i] - 1]; max
             (i, j) + h < n & s[i + h] == s[j + h];) ++h;
       hi[ra[i]] = h ? h-- : 0;
  }
};
```

## 5.5 Aho-Corasick Automatan\* [794a77]

```
struct AC_Automatan {
  int nx[len][sigma], fl[len], cnt[len], ord[len], top;
  int rnx[len][sigma]; // node actually be reached
  int newnode() {
    fill_n(nx[top], sigma, -1);
    return top++:
  void init() { top = 1, newnode(); }
  int input(string &s) {
    int X = 1;
    for (char c : s) {
   if (!~nx[X][c - 'A']) nx[X][c - 'A'] = newnode();
   X = nx[X][c - 'A'];
    return X; // return the end node of string
  void make_fl() {
    queue < int > q;
    q.push(1), fl[1] = 0;
    for (int t = 0; !q.empty(); ) {
      int R = q.front();
      q.pop(), ord[t++] = R;
for (int i = 0; i < sigma; ++i)</pre>
         if (~nx[R][i]) {
           int X = rnx[R][i] = nx[R][i], Z = fl[R];
           for (; Z && !~nx[Z][i]; ) Z = fl[Z];
```

#### 5.6 Smallest Rotation [4f469f]

```
string mcp(string s) {
  int n = SZ(s), i = 0, j = 1;
  s += s;
  while (i < n && j < n) {
    int k = 0;
    while (k < n && s[i + k] == s[j + k]) ++k;
    if (s[i + k] <= s[j + k]) j += k + 1;
    else i += k + 1;
    if (i == j) ++j;
  }
  int ans = i < n ? i : j;
  return s.substr(ans, n);
}</pre>
```

## 5.7 De Bruijn sequence\* [a09470]

```
constexpr int MAXC = 10, MAXN = 1e5 + 10;
struct DBSeq {
  int C, N, K, L, buf[MAXC * MAXN]; // K <= C^N</pre>
  void dfs(int *out, int t, int p, int &ptr) {
     if (ptr >= L) return;
     if (t > N) {
       if (N % p) return;
       for (int i = 1; i <= p && ptr < L; ++i)</pre>
         out[ptr++] = buf[i];
     } else
       buf[t] = buf[t - p], dfs(out, t + 1, p, ptr);
       for (int j = buf[t - p] + 1; j < C; ++j)</pre>
         buf[t] = j, dfs(out, t + 1, t, ptr);
    }
  }
  void solve(int _c, int _n, int _k, int *out) {
     int p = 0;
    C = _c, N = _n, K = _k, L = N + K - 1;
dfs(out, 1, 1, p);
if (p < L) fill(out + p, out + L, 0);</pre>
} dbs;
```

#### 5.8 Extended SAM\* [64c3b7]

```
int len[N * 2], link[N * 2]; // maxlength, suflink
int next[N * 2][CNUM], tot; // [0, tot), root = 0
int lenSorted[N * 2]; // topo. order
int cnt[N * 2]; // occurence
int newnode() {
  fill_n(next[tot], CNUM, 0);
  len[tot] = cnt[tot] = link[tot] = 0;
  return tot++;
void init() { tot = 0, newnode(), link[0] = -1; }
int insertSAM(int last, int c) {
  int cur = next[last][c];
  len[cur] = len[last] + 1;
  int p = link[last];
  while (p != -1 && !next[p][c])
  next[p][c] = cur, p = link[p];
if (p == -1) return link[cur] = 0, cur;
  int q = next[p][c];
  if (len
       [p] + 1 == len[q]) return link[cur] = q, cur;
  int clone = newnode();
  for (int i = 0; i < CNUM; ++i)</pre>
    next[
         clone][i] = len[next[q][i]] ? next[q][i] : 0;
  len[clone] = len[p] + 1;
  while (p != -1 && next[p][c] == q)
    next[p][c] = clone, p = link[p];
  link[link[cur] = clone] = link[q];
  link[q] = clone;
  return cur;
void insert(const string &s) {
```

```
int cur = 0:
     for (auto ch : s) {
       int &nxt = next[cur][int(ch - 'a')];
       if (!nxt) nxt = newnode();
       cnt[cur = nxt] += 1;
    }
  }
   void build() {
     queue < int > q;
     q.push(0);
     while (!q.empty()) {
      int cur = q.front();
       q.pop();
       for (int i = 0; i < CNUM; ++i)
         if (next[cur][i])
           q.push(insertSAM(cur, i));
     vector<int> lc(tot);
     for (int i = 1; i < tot; ++i) ++lc[len[i]];</pre>
     partial_sum(ALL(lc), lc.begin());
     for (int i
         = 1; i < tot; ++i) lenSorted[--lc[len[i]]] = i;
   void solve() {
   for (int i = tot - 2; i >= 0; --i)
       cnt[link[lenSorted[i]]] += cnt[lenSorted[i]];
};
```

### 5.9 PalTree\* [d7d2cf]

```
struct palindromic_tree {
  struct node
    int next[26], fail, len;
    int cnt, num; // cnt: appear times, num: number of
                   // pal. suf.
    node(int l = 0) : fail(0), len(l), cnt(0), num(0) {
      for (int i = 0; i < 26; ++i) next[i] = 0;</pre>
  };
  vector<node> St:
  vector < char > s;
  int last, n;
  palindromic_tree() : St(2), last(1), n(0) {
    St[0].fail = 1, St[1].len = -1, s.pb(-1);
  inline void clear() {
    St.clear(), s.clear(), last = 1, n = 0;
    St.pb(0), St.pb(-1);
    St[0].fail = 1, s.pb(-1);
  inline int get_fail(int x) {
    while (s[n - St[x].len - 1] != s[n])
     x = St[x].fail;
    return x;
  inline void add(int c) {
  s.push_back(c -= 'a'), ++n;
    int cur = get_fail(last);
    if (!St[cur].next[c]) {
      int now = SZ(St);
      St.pb(St[cur].len + 2);
      St[now].fail
        St[get_fail(St[cur].fail)].next[c];
      St[cur].next[c] = now;
      St[now].num = St[St[now].fail].num + 1;
    last = St[cur].next[c], ++St[last].cnt;
  inline void count() { // counting cnt
    auto i = St.rbegin();
    for (; i != St.rend(); ++i) {
      St[i->fail].cnt += i->cnt;
  inline int size() { // The number of diff. pal.
    return SZ(St) - 2;
```

#### 5.10 Main Lorentz [dflae3]

```
vector<pair<int, int>> rep[kN]; // 0-base [l, r]
void main_lorentz(const string &s, int sft = 0) {
  const int n = s.size();
  if (n == 1) return;
  const int nu = n / 2, nv = n - nu;
```

```
const string u = s.substr(0, nu), v = s.substr(nu),
        ru(u.rbegin
             (), u.rend()), rv(v.rbegin(), v.rend());
  main_lorentz(u, sft), main_lorentz(v, sft + nu);
  const
      auto z1 = z_value(ru), z2 = z_value(v + '#' + u),
             z3 = z_value
                  (ru + '#' + rv), z4 = z_value(v);
  auto get_z = [](const vector<int> &z, int i) {
    return
         (0 <= i and i < (int)z.size()) ? z[i] : 0; };
  auto add rep
       = [&](bool left, int c, int l, int k1, int k2) {
         int L = max(1, l - k2), R = min(l - left, k1);
    if (L > R) return;
    if (left)
         rep[l].emplace_back(sft + c - R, sft + c - L);
    else rep[l].emplace_back
         (sft + c - R - l + 1, sft + c - L - l + 1);
  for (int cntr = 0; cntr < n; cntr++) {</pre>
    int l, k1, k2;
    if (cntr < nu) {</pre>
      l = nu - cntr;
      k1 = get_z(z1, nu - cntr);
k2 = get_z(z2, nv + 1 + cntr);
    } else {
      l = cntr - nu + 1;
      k1 = get_z(z3, nu + 1 + nv - 1 - (cntr - nu));
      k2 = get_z(z4, (cntr - nu) + 1);
    if (k1 + k2 >= l)
      add_rep(cntr < nu, cntr, l, k1, k2);</pre>
  } // p \in [l, r] => s[p, p + i) = s[p + i, p + 2i)
```

## 6 Math

### 6.1 ax+by=gcd(only exgcd \*) [7b833d]

```
pll exgcd(ll a, ll b) {
   if (b == 0) return pll(1, 0);
   ll p = a / b;
   pll q = exgcd(b, a % b);
   return pll(q.Y, q.X - q.Y * p);
}
/* ax+by=res, let x be minimum non-negative
g, p = gcd(a, b), exgcd(a, b) * res / g
if p.X < 0: t = (abs(p.X) + b / g - 1) / (b / g)
else: t = -(p.X / (b / g))
p += (b / g, -a / g) * t */</pre>
```

## 6.2 Floor and Ceil [692c04]

```
int floor(int a, int b)
{ return a / b - (a % b && (a < 0) ^ (b < 0)); }
int ceil(int a, int b)
{ return a / b + (a % b && (a < 0) ^ (b > 0)); }
```

#### 6.3 Floor Enumeration [7cbcdf]

```
// enumerating x = floor(n / i), [l, r]
for (int l = 1, r; l <= n; l = r + 1) {
   int x = n / l;
   r = n / x;
}</pre>
```

#### 6.4 Mod Min [9118e1]

```
// min{k | l <= ((ak) mod m) <= r}, no solution -> -1
ll mod_min(ll a, ll m, ll l, ll r) {
  if (a == 0) return l ? -1 : 0;
  if (ll k = (l + a - 1) / a; k * a <= r)
    return k;
  ll b = m / a, c = m % a;
  if (ll y = mod_min(c, a, a - r % a, a - l % a))
    return (l + y * c + a - 1) / a + y * b;
  return -1;
}</pre>
```

#### 6.5 Gaussian integer gcd [0e7740]

```
cpx gaussian_gcd(cpx a, cpx b) {
#define rnd
    (a, b) ((a >= 0 ? a * 2 + b : a * 2 - b) / (b * 2))
    ll c = a.real() * b.real() + a.imag() * b.imag();
    ll d = a.imag() * b.real() - a.real() * b.imag();
```

### 6.6 Miller Rabin\* [06308c]

### 6.7 Simultaneous Equations [a231be]

```
struct matrix { //m variables, n equations
   int n, m;
   fraction M[MAXN][MAXN + 1], sol[MAXN];
   int solve() { //-1: inconsistent, >= 0: rank
     for (int i = 0; i < n; ++i) {</pre>
       int piv = 0;
       while (piv < m && !M[i][piv].n) ++piv;</pre>
       if (piv == m) continue;
        for (int j = 0; j < n; ++j) {</pre>
          if (i == j) continue;
fraction tmp = -M[j][piv] / M[i][piv];
          for (int k = 0; k <=
                m; ++k) M[j][k] = tmp * M[i][k] + M[j][k];
       }
     int rank = 0;
     for (int i = 0; i < n; ++i) {</pre>
       int piv = 0;
       while (piv < m && !M[i][piv].n) ++piv;</pre>
       if (piv == m && M[i][m].n) return -1;
else if (piv
             < m) ++rank, sol[piv] = M[i][m] / M[i][piv];</pre>
     return rank;
  }
};
```

### 6.8 Pollard Rho\* [cfe72f]

```
map<ll, int> cnt;
void PollardRho(ll n) {
  if (n == 1) return;
  if (prime(n)) return ++cnt[n], void();
  if (n % 2
      == 0) return PollardRho(n / 2), ++cnt[2], void();
  ll x = 2, y = 2, d = 1, p = 1;
  #define f(x, n, p) ((mul(x, x, n) + p) % n)
  while (true) {
    if (d != n && d != 1) {
      PollardRho(n / d);
      PollardRho(d);
      return:
    if (d == n) ++p;
    x = f(x, n, p), y = f(f(y, n, p), n, p);
d = gcd(abs(x - y), n);
  }
}
```

#### 6.9 Simplex Algorithm [6b4566]

```
const int MAXN = 11000, MAXM = 405;
const double eps = 1E-10;
double a[MAXN][MAXM], b[MAXN], c[MAXM];
double d[MAXN][MAXM], x[MAXM];
int ix[MAXN + MAXM]; // !!! array all indexed from 0
// max{cx} subject to {Ax<=b,x>=0}
// n: constraints, m: vars !!!
// x[] is the optimal solution vector
// usage :
```

```
// value = simplex(a, b, c, N, M);
double simplex(int n, int m){
  fill_n(d[n], m + 1, 0);
  fill_n(d[n + 1], m + 1, 0);
  iota(ix, ix + n + m, \theta);
  int r = n, s = m - 1;
  for (int i = 0; i < n; ++i) {</pre>
    for (int j = 0; j < m - 1; ++j) d[i][j] = -a[i][j];</pre>
    d[i][m - 1] = 1;
    d[i][m] = b[i];
    if (d[r][m] > d[i][m]) r = i;
  copy_n(c, m - 1, d[n]);
d[n + 1][m - 1] = -1;
  for (double dd;; ) {
    if (r < n) {
       swap(ix[s], ix[r + m]);
       d[r][s] = 1.0 / d[r][s];
       for (int j = 0; j <= m; ++j)</pre>
         if (j != s) d[r][j] *= -d[r][s];
       for (int i = 0; i <= n + 1; ++i) if (i != r) {
  for (int j = 0; j <= m; ++j) if (j != s)</pre>
           d[i][j] += d[r][j] * d[i][s];
         d[i][s] *= d[r][s];
      }
    }
    r = s = -1;
    for (int j = 0; j < m; ++j)
  if (s < 0 || ix[s] > ix[j]) {
         if (d[n + 1][j] > eps ||
              (d[n + 1][j] > -eps && d[n][j] > eps))
    if (s < 0) break;</pre>
    for (int i = 0; i < n; ++i) if (d[i][s] < -eps) {</pre>
      if (r < 0 ||
           (dd = d[r][m]
                 / d[r][s] - d[i][m] / d[i][s]) < -eps ||
           (dd < eps && ix[r + m] > ix[i + m]))
         r = i:
    if (r < 0) return -1; // not bounded
  if (d[n + 1][m] < -eps) return -1; // not executable</pre>
  double ans = 0;
  fill_n(x, m, 0);
  for (int i = m; i <
    n + m; ++i) { // the missing enumerated x[i] = 0
if (ix[i] < m - 1){
ans += d[i - m][m] * c[ix[i]];
       x[ix[i]] = d[i-m][m];
    }
  return ans;
```

## 6.9.1 Construction

Primal	Dual
Maximize $c^{T}x$ s.t. $Ax \leq b$ , $x \geq 0$	Minimize $b^{\intercal}y$ s.t. $A^{\intercal}y \ge c$ , $y \ge 0$
Maximize $c^{T}x$ s.t. $Ax \leq b$	Minimize $b^{T}y$ s.t. $A^{T}y = c$ , $y \ge 0$
Maximize $c^{T}x$ s.t. $Ax = b$ . $x > 0$	Minimize $b^{T} y$ s.t. $A^{T} y > c$

 $ar{\mathbf{x}}$  and  $ar{\mathbf{y}}$  are optimal if and only if for all  $i\in[1,n]$ , either  $\bar{x}_i=0$  or  $\sum_{j=1}^m A_{ji}\bar{y}_j=c_i$  holds and for all  $i\in[1,m]$  either  $\bar{y}_i=0$  or  $\sum_{j=1}^n A_{ij}\bar{x}_j=b_j$  holds.

- 1. In case of minimization, let  $c_i' = -c_i$
- 2.  $\sum_{1 \leq i \leq n} A_{ji} x_i \geq b_j \rightarrow \sum_{1 \leq i \leq n} -A_{ji} x_i \leq -b_j$
- $3. \sum_{1 \le i \le n} A_{ji} x_i = b_j$ 
  - $\sum_{1 \le i \le n} A_{ji} x_i \le b_j$
  - $\sum_{1 \leq i \leq n}^{-} A_{ji} x_i \geq b_j$
- 4. If  $x_i$  has no lower bound, replace  $x_i$  with  $x_i x_i'$

#### 6.10 chineseRemainder [a53b6d]

```
ll solve(ll x1, ll m1, ll x2, ll m2) {
    ll g = gcd(m1, m2);
    if ((x2 - x1) % g) return -1; // no sol
    m1 /= g; m2 /= g;
    pll p = exgcd(m1, m2);
    ll lcm = m1 * m2 * g;
    ll res = p.first * (x2 - x1) * m1 + x1;
    // be careful with overflow
    return (res % lcm + lcm) % lcm;
}
```

### 6.11 Factorial without prime factor\* [c324f3]

```
// O(p^k + log^2 n), pk = p^k
ll prod[MAXP];
ll fac_no_p(ll n, ll p, ll pk) {
  prod[0] = 1;
  for (int i = 1; i <= pk; ++i)
      if (i % p) prod[i] = prod[i - 1] * i % pk;
      else prod[i] = prod[i - 1];
ll rt = 1;
for (; n; n /= p) {
    rt = rt * mpow(prod[pk], n / pk, pk) % pk;
    rt = rt * prod[n % pk] % pk;
}
  return rt;
} // (n! without factor p) % p^k</pre>
```

#### 6.12 QuadraticResidue\* [e0bf30]

```
int Jacobi(int a, int m) {
   int s = 1;
   for (; m > 1; ) {
   a %= m;
     if (a == 0) return 0;
     const int r = __builtin_ctz(a);
if ((r & 1) && ((m + 2) & 4)) s = -s;
     if (a & m & 2) s = -s;
     swap(a, m);
   return s;
int QuadraticResidue(int a, int p) {
   if (p == 2) return a & 1;
   const int jc = Jacobi(a, p);
   if (jc == 0) return 0;
   if (jc == -1) return -1;
   int b, d;
   for (; ; ) {
     b = rand() % p;
     d = (1LL * b * b + p - a) % p;
     if (Jacobi(d, p) == -1) break;
   int f0 = b, f1 = 1, g0 = 1, g1 = 0, tmp;
for (int e = (1LL + p) >> 1; e; e >>= 1) {
     if (e & 1) {
       tmp = (1LL *
           g0 * f0 + 1LL * d * (1LL * g1 * f1 % p)) % p;
       g1 = (1LL * g0 * f1 + 1LL * g1 * f0) % p;
       g0 = tmp;
     tmp = (1LL     * f0 * f0 + 1LL * d * (1LL * f1 * f1 % p)) % p;
     f1 = (2LL * f0 * f1) % p;
     f0 = tmp;
   }
   return g0;
}
```

#### 6.13 PiCount\* [cad6d4]

```
ll PrimeCount(ll n) { // n \sim 10^13 => < 2s
  if (n <= 1) return 0;</pre>
  int v = sqrt(n), s = (v + 1) / 2, pc = 0;
  vector<int> smalls(v + 1), skip(v + 1), roughs(s);
  vector<ll> larges(s);
  for (int i = 2; i <= v; ++i) smalls[i] = (i + 1) / 2;</pre>
  for (int i = 0; i < s; ++i) {</pre>
    roughs[i] = 2 * i + 1;
    larges[i] = (n / (2 * i + 1) + 1) / 2;
  for (int p = 3; p <= v; ++p) {</pre>
    if (smalls[p] > smalls[p - 1]) {
      int q = p * p;
      ++pc;
      if (1LL * q * q > n) break;
      skip[p] = 1;
      for (int i = q; i <= v; i += 2 * p) skip[i] = 1;</pre>
      int ns = 0;
      for (int k = 0; k < s; ++k) {
        int i = roughs[k];
        if (skip[i]) continue;
        ll d = 1LL * i * p;
        larges[ns] = larges[k] - (d <= v ? larges</pre>
             [smalls[d] - pc] : smalls[n / d]) + pc;
        roughs[ns++] = i;
```

```
}
     s = ns;
     for (int j = v / p; j >= p; --j) {
              smalls[j] - pc, e = min(j * p + p, v + 1);
       for (int i = j * p; i < e; ++i) smalls[i] -= c;</pre>
  }
for (int k = 1; k < s; ++k) {</pre>
  const ll m = n / roughs[k];
  ll t = larges[k] - (pc + k - 1);
for (int l = 1; l < k; ++l) {</pre>
     int p = roughs[l];
    if (1LL * p * p > m) break;
t -= smalls[m / p] - (pc + l - 1);
  larges[0] -= t;
return larges[0];
```

### 6.14 Discrete Log\* [da27bf]

```
int DiscreteLog(int s, int x, int y, int m) {
  constexpr int kStep = 32000;
  unordered_map<int, int> p;
  int b = 1;
  for (int i = 0; i < kStep; ++i) {</pre>
    p[y] = i;
y = 1LL * y * x % m;
    b = 1LL * b * x % m;
  for (int i = 0; i < m + 10; i += kStep) {
    s = 1LL * s * b % m;</pre>
    if (p.find(s) != p.end()) return i + kStep - p[s];
  return -1;
int DiscreteLog(int x, int y, int m) {
  if (m == 1) return 0;
  int s = 1;
  for (int i = 0; i < 100; ++i) {</pre>
    if (s == y) return i;
s = 1LL * s * x % m;
  if (s == y) return 100;
  int p = 100 + DiscreteLog(s, x, y, m);
  if (fpow(x, p, m) != y) return -1;
  return p;
```

#### 6.15 Berlekamp Massey [3eb6fa]

```
template <typename T>
vector<T> BerlekampMassey(const vector<T> &output) {
   vector<T> d(SZ(output) + 1), me, he;
   for (int f = 0, i = 1; i <= SZ(output); ++i) {
  for (int j = 0; j < SZ(me); ++j)
    d[i] += output[i - j - 2] * me[j];
  if ((d[i] -= output[i - 1]) == 0) continue;</pre>
      if (me.empty()) {
         me.resize(f = i);
         continue;
      vector<T> o(i - f - 1);
      T k = -d[i] / d[f]; o.pb(-k);
      for (T x : he) o.pb(x * k);
      o.resize(max(SZ(o), SZ(me)));
      for (int j = 0; j < SZ(me); ++j) o[j] += me[j];
if (i - f + SZ(he) >= SZ(me)) he = me, f = i;
      me = o;
   }
   return me;
}
```

#### 6.16 Primes

```
/* 12721 13331 14341 75577 123457 222557
     556679 999983 1097774749 1076767633 100102021
    999997771 1001010013 1000512343 987654361 999991231
     999888733 98789101 987777733 999991921 1010101333
     1010102101 1000000000039 100000000000037
     2305843009213693951 4611686018427387847
     9223372036854775783 18446744073709551557 */
```

#### 6.17 Theorem

· Cramer's rule

· Vandermonde's Identity

$$C(n+m,k) = \sum_{i=0}^{k} C(n,i)C(m,k-i)$$

• Kirchhoff's Theorem

Denote L be a  $n \times n$  matrix as the Laplacian matrix of graph G, where  $L_{ii} = d(i)$ ,  $L_{ij} = -c$  where c is the number of edge (i,j) in G.

- The number of undirected spanning in G is  $|\det(\tilde{L}_{11})|$ .
- The number of directed spanning tree rooted at r in G is  $|\det(\tilde{L}_{rr})|$ .

Let D be a n imes n matrix, where  $d_{ij} = x_{ij}$  ( $x_{ij}$  is chosen uniformly at random) if i < j and  $(i,j) \in E$ , otherwise  $d_{ij} = -d_{ji}.$   $\frac{rank(D)}{2}$  is the maximum matching on G.

- · Cayley's Formula
  - Given a degree sequence  $d_1, d_2, ..., d_n$  for each labeled vertices, there are  $\frac{(n-2)!}{(d_1-1)!(d_2-1)!\cdots(d_n-1)!}$  spanning trees.
  - Let  $T_{n,k}$  be the number of *labeled* forests on n vertices with k components, such that vertex  $1,2,\ldots,k$  belong to different components. Then  $T_{n,k} = kn^{n-k-1}$ .
- Erdős–Gallai theorem

A sequence of nonnegative integers  $d_1 \ge \cdots \ge d_n$  can be represented as the degree sequence of a finite simple graph on n vertices if and only if

$$d_1+\dots+d_n \text{ is even and } \sum_{i=1}^k d_i \leq k(k-1)+\sum_{i=k+1}^n \min(d_i,k) \text{ holds for every } 1\leq k\leq n.$$

· Gale-Ryser theorem

A pair of sequences of nonnegative integers  $a_1 \ge \cdots \ge a_n$  and  $b_1, \dots, b_n$ is bigraphic if and only if  $\sum_{i=1}^n a_i = \sum_{i=1}^n b_i$  and  $\sum_{i=1}^k a_i \leq \sum_{i=1}^n \min(b_i,k)$  holds for

every  $1 \le k \le n$ .

Fulkerson-Chen-Anstee theorem

A sequence  $(a_1,\ b_1),\ ...\ ,\ (a_n,\ b_n)$  of nonnegative integer pairs with  $a_1 \geq \cdots \geq a_n$  is digraphic if and only if  $\sum_{i=1}^n a_i = \sum_{i=1}^n b_i$  and

$$\sum_{i=1}^k a_i \leq \sum_{i=1}^k \min(b_i,k-1) + \sum_{i=k+1}^n \min(b_i,k) \text{ holds for every } 1 \leq k \leq n.$$

· Pick's theorem

For simple polygon, when points are all integer, we have  $A = \#\{\text{lattice points in the interior}\} + \frac{\#\{\text{lattice points on the boundary}\}}{2} - 1.$ 

- · Möbius inversion formula
  - $f(n) = \sum_{d|n} g(d) \Leftrightarrow g(n) = \sum_{d|n} \mu(d) f(\frac{n}{d})$
  - $f(n) = \sum_{n|d} g(d) \Leftrightarrow g(n) = \sum_{n|d} \mu(\frac{d}{n}) f(d)$
- Spherical cap
  - A portion of a sphere cut off by a plane.
  - r: sphere radius, a: radius of the base of the cap, h: height of the cap,  $\theta$ : arcsin(a/r).
  - Volume =  $\pi h^2 (3r h)/3 = \pi h (3a^2 + h^2)/6 = \pi r^3 (2 + \cos \theta)(1 \theta)$  $\cos\theta)^2/3$ .
  - Area  $= 2\pi rh = \pi(a^2 + h^2) = 2\pi r^2(1 \cos\theta)$ .
- · Lagrange multiplier
  - Optimize  $f(x_1,...,x_n)$  when k constraints  $g_i(x_1,...,x_n) = 0$ .
  - Lagrangian function  $\mathcal{L}(x_1,\,...\,,\,x_n,\,\lambda_1,\,...\,,\,\lambda_k) = f(x_1,\,...\,,\,x_n)$   $\sum_{i=1}^k \lambda_i g_i(x_1,...,x_n)$ .
  - The solution corresponding to the original constrained optimization is always a saddle point of the Lagrangian function.
- · Nearest points of two skew lines
  - Line 1:  $v_1 = p_1 + t_1 d_1$
  - Line 2:  $v_2 = p_2 + t_2 d_2$
  - $\boldsymbol{n} = \boldsymbol{d}_1 \times \boldsymbol{d}_2$
  - $n_1 = d_1 \times n$
  - $n_2 = d_2 \times n$ -  $c_1 = p_1 + \frac{(p_2 - p_1) \cdot n_2}{d_1 \cdot n_2} d_1$
  - $c_2 = p_2 + \frac{(p_1 p_2) \cdot n_1}{d_2 \cdot n_1} d_2$

· Derivatives/Integrals

Integration by parts: 
$$\int_a^b f(x)g(x)dx = [F(x)g(x)]_a^b - \int_a^b F(x)g'(x)dx \\ \left| \frac{d}{dx}\sin^{-1}x = \frac{1}{\sqrt{1-x^2}} \right| \frac{d}{dx}\cos^{-1}x = -\frac{1}{\sqrt{1-x^2}} \left| \frac{d}{dx}\tan^{-1}x = \frac{1}{1+x^2} \right| \\ \frac{d}{dx}\tan x = 1 + \tan^2x \quad \int \tan ax = -\frac{\ln|\cos ax|}{a} \\ \int e^{-x^2} = \frac{\sqrt{\pi}}{2}\operatorname{erf}(x) \left| \int xe^{ax}dx = \frac{e^{ax}}{a^2}(ax-1) \right| \\ \int \sqrt{a^2 + x^2} = \frac{1}{2} \left( x\sqrt{a^2 + x^2} + a^2 \operatorname{asinh}(x/a) \right)$$

Spherical Coordinate

$$(x,y,z)\!=\!(r\!\sin\!\theta\!\cos\!\phi,\!r\!\sin\!\theta\!\sin\!\phi,\!r\!\cos\!\theta)$$

$$(r,\!\theta,\!\phi)\!=\!(\sqrt{x^2\!+\!y^2\!+\!z^2},\!\mathsf{acos}(z/\sqrt{x^2\!+\!y^2\!+\!z^2}),\!\mathsf{atan2}(y,\!x))$$

Rotation Matrix

$$M(\theta)\!=\!\begin{bmatrix}\!\cos\!\theta & -\!\sin\!\theta \\ \!\sin\!\theta & \!\cos\!\theta \end{bmatrix}\!,\!R_x(\theta_x)\!=\!\begin{bmatrix} 1 & 0 & 0 \\ 0 & \!\cos\!\theta_x & -\!\sin\!\theta_x \\ 0 & \!\sin\!\theta & \!\cos\!\theta \end{bmatrix}$$

#### 6.18 Estimation

### 6.19 Euclidean Algorithms

- $m = \lfloor \frac{an+b}{c} \rfloor$
- Time complexity:  $O(\log n)$

$$\begin{split} f(a,b,c,n) &= \sum_{i=0}^n \lfloor \frac{ai+b}{c} \rfloor \\ &= \begin{cases} \lfloor \frac{a}{c} \rfloor \cdot \frac{n(n+1)}{2} + \lfloor \frac{b}{c} \rfloor \cdot (n+1) \\ +f(a \bmod c,b \bmod c,c,n), & a \geq c \vee b \geq c \\ 0, & n < 0 \vee a = 0 \\ nm - f(c,c - b - 1,a,m - 1), & \text{otherwise} \end{cases} \\ g(a,b,c,n) &= \sum_{i=0}^n i \lfloor \frac{ai+b}{c} \rfloor \\ &= \begin{cases} \lfloor \frac{a}{c} \rfloor \cdot \frac{n(n+1)(2n+1)}{6} + \lfloor \frac{b}{c} \rfloor \cdot \frac{n(n+1)}{2} \\ +g(a \bmod c,b \bmod c,c,n), & a \geq c \vee b \geq c \\ 0, & n < 0 \vee a = 0 \end{cases} \end{split}$$

$$\begin{split} h(a,b,c,n) &= \sum_{i=0}^n \lfloor \frac{ai+b}{c} \rfloor^2 \\ &= \begin{cases} \lfloor \frac{a}{c} \rfloor^2 \cdot \frac{n(n+1)(2n+1)}{6} + \lfloor \frac{b}{c} \rfloor^2 \cdot (n+1) \\ &+ \lfloor \frac{a}{c} \rfloor \cdot \lfloor \frac{b}{c} \rfloor \cdot n(n+1) \\ &+ h(a \bmod c, b \bmod c, c, n) \\ &+ 2 \lfloor \frac{a}{c} \rfloor \cdot g(a \bmod c, b \bmod c, c, n) \\ &+ 2 \lfloor \frac{b}{c} \rfloor \cdot f(a \bmod c, b \bmod c, c, n), & a \geq c \lor b \geq c \\ 0, & n < 0 \lor a = 0 \\ nm(m+1) - 2g(c, c - b - 1, a, m - 1) \\ &- 2f(c, c - b - 1, a, m - 1) - f(a, b, c, n), & \text{otherwise} \end{cases} \end{split}$$

 $\frac{1}{2} \cdot (n(n+1)m - f(c,c-b-1,a,m-1) \\ -h(c,c-b-1,a,m-1)),$ 

otherwise

#### 6.20 General Purpose Numbers

Bernoulli numbers

$$\begin{split} &B_0 - 1, B_1^{\pm} = \pm \frac{1}{2}, B_2 = \frac{1}{6}, B_3 = 0 \\ &\sum_{j=0}^m \binom{m+1}{j} B_j = 0, \text{EGF is } B(x) = \frac{x}{e^x - 1} = \sum_{n=0}^\infty B_n \frac{x^n}{n!}. \\ &S_m(n) = \sum_{k=1}^n k^m = \frac{1}{m+1} \sum_{k=0}^m \binom{m+1}{k} B_k^+ n^{m+1-k} \end{split}$$

- Stirling numbers of the second kind Partitions of n distinct elements into exactly k groups.

$$\begin{split} S(n,k) &= S(n-1,k-1) + kS(n-1,k), S(n,1) = S(n,n) = 1 \\ S(n,k) &= \frac{1}{k!} \sum_{i=0}^k (-1)^{k-i} {k \choose i} i^n \\ x^n &= \sum_{i=0}^n S(n,i)(x)_i \end{split}$$

Pentagonal number theorem

$$\prod_{n=1}^{\infty} (1-x^n) = 1 + \sum_{k=1}^{\infty} (-1)^k \left( x^{k(3k+1)/2} + x^{k(3k-1)/2} \right)$$

• Catalan numbers  $C_n^{(k)} = \frac{1}{(k-1)n+1} \binom{kn}{n}$   $C^{(k)}(x) = 1 + x[C^{(k)}(x)]^k$ 

• Eulerian numbers

Number of permutations  $\pi \in S_n$  in which exactly k elements are greater than the previous element. k j:s s.t.  $\pi(j) > \pi(j+1)$ , k+1 j:s s.t.  $\pi(j) \ge j$ , k j:s s.t.  $\pi(j) \ge j$ :s s.t.  $\pi(j) \ge j$ :s

### 6.21 Tips for Generating Functions

- Ordinary Generating Function  $A(x) = \sum_{i \ge 0} a_i x^i$ 
  - $A(rx) \Rightarrow r^n a_n$ -  $A(x) + B(x) \Rightarrow a_n + b_n$ -  $A(x)B(x) \Rightarrow \sum_{i=0}^n a_i b_{n-i}$
  - $A(x)B(x) \Rightarrow \sum_{i=0} a_i o_{n-i}$ -  $A(x)^k \Rightarrow \sum_{i_1+i_2+\dots+i_k=n} a_{i_1} a_{i_2} \dots a_{i_k}$
  - $xA(x)' \Rightarrow na_n$
  - $-\frac{A(x)}{1-x} \Rightarrow \sum_{i=0}^{n} a_i$
- Exponential Generating Function  $A(x) = \sum_{i>0} \frac{a_i}{i!} x_i$ 
  - $A(x)+B(x) \Rightarrow a_n+b_n$
  - $A^{(k)}(x) \Rightarrow a_{n+k}$
  - $A(x)B(x) \Rightarrow \sum_{i=0}^{n} {n \choose i} a_i b_{n-i}$
  - $A(x)^k \Rightarrow \sum_{i_1+i_2+\dots+i_k=n} \binom{n}{i_1,i_2,\dots,i_k} a_{i_1} a_{i_2} \dots a_{i_k}$
  - $xA(x) \Rightarrow na_n$
- · Special Generating Function
  - $(1+x)^n = \sum_{i\geq 0} \binom{n}{i} x^i$
  - $-\frac{1}{(1-x)^n} = \sum_{i\geq 0} \binom{i}{n-1} x^i$

## 7 Polynomial

## 7.1 Fast Fourier Transform [56bdd7]

```
template < int MAXN >
struct FFT {
    using val_t = complex < double >;
    const double PI = acos(-1);
    val_t w[MAXN];
    FFT() {
        for (int i = 0; i < MAXN; ++i) {
            double arg = 2 * PI * i / MAXN;
            w[i] = val_t(cos(arg), sin(arg));
        }
    void bitrev(val_t *a, int n); // see NTT
    void trans
        (val_t *a, int n, bool inv = false); // see NTT;
    // remember to replace LL with val_t
};</pre>
```

## 7.2 Number Theory Transform\* [f68103]

```
//(2^16)+1, 65537, 3
//7*17*(2^23)+1, 998244353, 3
//1255*(2^20)+1, 1315962881, 3
//51*(2^25)+1, 1711276033, 29
template < int MAXN, ll P, ll RT > // MAXN must be 2^k
struct NTT {
  ll w[MAXN];
  ll mpow(ll a, ll n);
  ll minv(ll a) { return mpow(a, P - 2); }
  NTT() {
    ll dw = mpow(RT, (P - 1) / MAXN);
    w[0] = 1;
    for (int
         i = 1; i < MAXN; ++i) w[i] = w[i - 1] * dw % P;
  void bitrev(ll *a, int n) {
    int i = 0;
    for (int j = 1; j < n - 1; ++j) {</pre>
      for (int k = n >> 1; (i ^= k) < k; k >>= 1);
      if (j < i) swap(a[i], a[j]);</pre>
  void operator()(
       ll *a, int n, bool inv = false) { //0 <= a[i] < P
    bitrev(a, n);
for (int L = 2; L <= n; L <<= 1) {
       int dx = MAXN / L, dl = L >> 1;
       for (int i = 0; i < n; i += L) {</pre>
         for (int
           j = i, x = 0; j < i + dl; ++j, x += dx) { ll tmp = a[j + dl] * w[x] % P;
```

#### 7.3 Fast Walsh Transform\* [c9cdb6]

```
/* x: a[j], y: a[j + (L >> 1)]
or: (y += x * op), and: (x += y * op)
xor: (x, y = (x + y) * op, (x - y) * op)

invop: or, and, xor = -1, -1, 1/2 */
void fwt(int *a, int n, int op) { //or
  for (int L = 2; L <= n; L <<= 1)</pre>
    for (int i = 0; i < n; i += L)</pre>
       for (int j = i; j < i + (L >> 1); ++j)
         a[j + (L >> 1)] += a[j] * op;
const int N = 21;
int f[
    N_{1}^{-}[1 << N], g[N][1 << N], h[N][1 << N], ct[1 << N];
    subset_convolution(int *a, int *b, int *c, int L) {
  // c_k = \sum_{i | j = k, i & j = 0} a_i * b_j
  int n = 1 << L;
  for (int i = 1; i < n; ++i)</pre>
    ct[i] = ct[i & (i - 1)] + 1;
  for (int i = 0; i < n; ++i)</pre>
    f[ct[i]][i] = a[i], g[ct[i]][i] = b[i];
  for (int i = 0; i <= L; ++i)
  fwt(f[i], n, 1), fwt(g[i], n, 1);</pre>
  for (int i = 0; i <= L; ++i)</pre>
    for (int j = 0; j <= i; ++j)</pre>
       for (int x = 0; x < n; ++x)
         h[i][x] += f[j][x] * g[i - j][x];
  for (int i = 0; i <= L; ++i)</pre>
  fwt(h[i], n, -1);
for (int i = 0; i < n; ++i)</pre>
    c[i] = h[ct[i]][i];
```

#### 7.4 Polynomial Operation [105808]

```
fi(s, n) for (int i = (int)(s); i < (int)(n); ++i)
template < int MAXN, ll P, ll RT> // MAXN = 2^k
struct Poly : vector<ll> { // coefficients in [0, P)
 using vector<ll>::vector;
  static NTT < MAXN, P, RT > ntt;
  int n() const { return (int)size(); } // n() >= 1
 Poly(const Poly &p, int m) : vector<ll>(m) {
    copy_n(p.data(), min(p.n(), m), data());
 Poly& irev()
      { return reverse(data(), data() + n()), *this; }
  Poly& isz(int m) { return resize(m), *this; }
 Poly& iadd(const Poly &rhs) { // n() == rhs.n()
    fi(0, n()) if
         (((*this)[i] += rhs[i]) >= P) (*this)[i] -= P;
    return *this;
 Poly& imul(ll k) {
    fi(0, n()) (*this)[i] = (*this)[i] * k % P;
    return *this;
  Poly Mul(const Poly &rhs) const {
    int m = 1;
    while (m < n() + rhs.n() - 1) m <<= 1;</pre>
    Poly X(*this, m), Y(rhs, m);
    ntt(X.data(), m), ntt(Y.data(), m);
    fi(0, m) X[i] = X[i] * Y[i] % P;
    ntt(X.data(), m, true);
    return X.isz(n() + rhs.n() - 1);
  Poly Inv() const { // (*this)[0] != 0, 1e5/95ms
   if (n() == 1) return {ntt.minv((*this)[0])};
    int m = 1;
    while (m < n() * 2) m <<= 1;</pre>
```

```
Poly Xi = Poly(*this, (n() + 1) / 2).Inv().isz(m);
  Poly Y(*this, m);
  ntt(Xi.data(), m), ntt(Y.data(), m);
  fi(0, m) {
    Xi[i] *= (2 - Xi[i] * Y[i]) % P;
    if ((Xi[i] %= P) < 0) Xi[i] += P;</pre>
  ntt(Xi.data(), m, true);
  return Xi.isz(n());
Poly Sqrt()
     const { // Jacobi((*this)[0], P) = 1, 1e5/235ms
  if (n()
      == 1) return {QuadraticResidue((*this)[0], P)};
  Poly
      X = Poly(*this, (n() + 1) / 2).Sqrt().isz(n());
  return
       X.iadd(Mul(X.Inv()).isz(n())).imul(P / 2 + 1);
pair < Poly, Poly > DivMod
    (const Poly &rhs) const { // (rhs.)back() != 0
  if (n() < rhs.n()) return {{0}, *this};
const int m = n() - rhs.n() + 1;</pre>
  Poly X(rhs); X.irev().isz(m);
  Poly Y(*this); Y.irev().isz(m);
  Poly Q = Y.Mul(X.Inv()).isz(m).irev();
  X = rhs.Mul(Q), Y = *this;
  fi(0, n()) if ((Y[i] -= X[i]) < 0) Y[i] += P;
  return {Q, Y.isz(max(1, rhs.n() - 1))};
Poly Dx() const {
  Poly ret(n() - 1);
      ret.n()) ret[i] = (i + 1) * (*this)[i + 1] % P;
  return ret.isz(max(1, ret.n()));
Poly Sx() const {
  Poly ret(n() + 1);
  fi(0, n())
       ret[i + 1] = ntt.minv(i + 1) * (*this)[i] % P;
  return ret;
Poly _tmul(int nn, const Poly &rhs) const {
  Poly Y = Mul(rhs).isz(n() + nn - 1);
  return Poly(Y.data() + n() - 1, Y.data() + Y.n());
vector<ll> _eval(const
     vector<ll> &x, const vector<Poly> &up) const {
  const int m = (int)x.size();
  if (!m) return {};
  vector<Poly> down(m * 2);
  // down[1] = DivMod(up[1]).second;
  // fi(2, m *
        2) down[i] = down[i / 2].DivMod(up[i]).second;
  down[1] = Poly(up[1])
      .irev().isz(n()).Inv().irev()._tmul(m, *this);
  fi(2, m * 2) down[i]
       = up[i ^ 1]._tmul(up[i].n() - 1, down[i / 2]);
  vector<ll> y(m);
  fi(0, m) y[i] = down[m + i][0];
  return y;
static vector<Poly> _tree1(const vector<ll> &x) {
  const int m = (int)x.size();
  vector<Poly> up(m * 2);
  fi(0, m) up[m + i] = \{(x[i] ? P - x[i] : 0), 1\};
  for (int i = m - 1; i
      > 0; --i) up[i] = up[i * 2].Mul(up[i * 2 + 1]);
  return up;
}
vector
    <ll> Eval(const vector<ll> &x) const { // 1e5, 1s
  auto up = _tree1(x); return _eval(x, up);
static Poly Interpolate(const vector
    <ll> &x, const vector<ll> &y) { // 1e5, 1.4s
  const int m = (int)x.size();
  vector<Poly> up = _tree1(x), down(m * 2);
  vector < ll > z = up[1].Dx()._eval(x, up);
fi(0, m) z[i] = y[i] * ntt.minv(z[i]) % P;
  fi(0, m) down[m + i] = {z[i]};
  for (int i = m -
       1; i > 0; --i) down[i] = down[i * 2].Mul(up[i
      * 2 + 1]).iadd(down[i * 2 + 1].Mul(up[i * 2]));
  return down[1];
```

```
Poly Ln() const { // (*this)[0] == 1, 1e5/170ms
    return Dx().Mul(Inv()).Sx().isz(n());
  Poly Exp() const { // (*this)[0] == 0, 1e5/360ms
    if (n() == 1) return {1};
Poly X = Poly(*this, (n() + 1) / 2).Exp().isz(n());
    Poly Y = X.Ln(); Y[0] = P - 1;
    fi(0, n())
         if ((Y[i] = (*this)[i] - Y[i]) < 0) Y[i] += P;</pre>
    return X.Mul(Y).isz(n());
  }
  ^{\prime}// M := P(P - 1). If k >= M, k := k \% M + M.
  Poly Pow(ll k) const {
    int nz = 0;
    while (nz < n() && !(*this)[nz]) ++nz;</pre>
    if (nz * min(k, (ll)n()) >= n()) return Poly(n());
    if (!k) return Poly(Poly {1}, n());
    Poly X(data() + nz, data() + nz + n() - nz * k);
    const ll c = ntt.mpow(X[0], k % (P - 1));
    return X.Ln().imul
        (k % P).Exp().imul(c).irev().isz(n()).irev();
  static ll
      LinearRecursion(const vector<ll> &a, const vector
    <ll> &coef, ll n) { // a_n = |sum c_j| a_n (n-j)
const int k = (int)a.size();
    assert((int)coef.size() == k + 1);
    Poly C(k + 1), W(Poly \{1\}, k), M = \{0, 1\};
    fi(1, k + 1) C[k - i] = coef[i] ? P - coef[i] : 0;
    C[k] = 1;
    while (n) {
      if (n % 2) W = W.Mul(M).DivMod(C).second;
      n /= 2, M = M.Mul(M).DivMod(C).second;
    fi(0, k) ret = (ret + W[i] * a[i]) % P;
    return ret;
 }
};
#undef fi
using Poly_t = Poly<131072 * 2, 998244353, 3>;
template <> decltype(Poly_t::ntt) Poly_t::ntt = {};
```

## 7.5 Value Polynomial [96cde9]

```
struct Poly {
  mint base; // f(x) = poly[x - base]
   vector<mint> poly;
  Poly(mint b = 0, mint x = 0): base(b), poly(1, x) {}
   mint get_val(const mint &x) {
     if (x >= base && x < base + SZ(poly))
       return poly[x - base];
     mint rt = 0;
     vector<mint> lmul(SZ(poly), 1), rmul(SZ(poly), 1);
     for (int i = 1; i < SZ(poly); ++i)</pre>
       lmul[i] = lmul[i - 1] * (x - (base + i - 1));
     for (int i = SZ(poly) - 2; i >= 0; --i)
  rmul[i] = rmul[i + 1] * (x - (base + i + 1));
     for (int i = 0; i < SZ(poly); ++i)
  rt += poly[i] * ifac[i] * inegfac</pre>
           [SZ(poly) - 1 - i] * lmul[i] * rmul[i];
     return rt;
   void raise() { // g(x) = sigma\{base:x\} f(x)
     if (SZ(poly) == 1 && poly[0] == 0)
       return:
     mint nw = get_val(base + SZ(poly));
     poly.pb(nw);
     for (int i = 1; i < SZ(poly); ++i)</pre>
       poly[i] += poly[i - 1];
  }
};
```

### 7.6 Newton's Method

Given F(x) where

$$F(x) = \sum_{i=0}^{\infty} \alpha_i (x - \beta)^i$$

for  $\beta$  being some constant. Polynomial P such that F(P)=0 can be found iteratively. Denote by  $Q_k$  the polynomial such that  $F(Q_k)=0$  (mod  $x^{2^k}$ ), then

$$Q_{k+1}\!=\!Q_k\!-\!\frac{F(Q_k)}{F'(Q_k)}\pmod{x^{2^{k+1}}}$$

## 8 Geometry

## 8.1 Default Code [ca77e7]

```
using ld = ll;
using pdd = pair<ld, ld>;
using Line = pair<pdd, pdd>;
struct Cir{ pdd 0; ld R; };
#define X first
#define Y second
// const ld eps = 1e-7;
pdd operator+(pdd a, pdd b)
{ return {a.X + b.X, a.Y + b.Y}; }
pdd operator - (pdd a, pdd b)
{ return {a.X - b.X, a.Y - b.Y}; }
pdd operator*(ld i, pdd v)
{ return {i * v.X, i * v.Y};
pdd operator*(pdd v, ld i)
{ return {i * v.X, i * v.Y}; }
pdd operator/(pdd v, ld i)
{ return {v.X / i, v.Y / i}; }
ld dot(pdd a, pdd b)
{ return a.X * b.X + a.Y * b.Y; }
ld cross(pdd a, pdd b)
{ return a.X * b.Y - a.Y * b.X; }
ld abs2(pdd v)
{ return v.X * v.X + v.Y * v.Y; };
ld abs(pdd v)
{ return sqrt(abs2(v)); };
int sgn(ld v)
{ return v > 0 ? 1 : (v < 0 ? -1 : 0); }
// int sgn(
     ld v){ return v > eps ? 1 : ( v < -eps ? -1 : 0); }</pre>
int ori(pdd a, pdd b, pdd c)
{ return sgn(cross(b - a, c - a)); }
bool collinearity(pdd a, pdd b, pdd c)
{ return ori(a, b, c) == 0; }
bool btw(pdd p, pdd a, pdd b)
{ return collinearity
     (p, a, b) && sgn(dot(a - p, b - p)) <= 0; }
bool seg_intersect(pdd p1, pdd p2, pdd p3, pdd p4){
   if(btw(p1, p3, p4) || btw(p2
     , p3, p4) || btw(p3, p1, p2) || btw(p4, p1, p2))
return true;
   return ori(p1, p2, p3) * ori(p1, p2, p4) < 0 &&</pre>
     ori(p3, p4, p1) * ori(p3, p4, p2) < 0;
pdd intersect(pdd p1, pdd p2, pdd p3, pdd p4){
  ld a123 = cross(p2 - p1, p3 - p1);
  ld a124 = cross(p2 - p1, p4 - p1);
  return (p4 * a123 - p3 * a124) / (a123 - a124);
pdd perp(pdd p1)
{ return pdd(-p1.Y, p1.X); }
pdd projection(pdd p1, pdd p2, pdd p3)
{ return p1 + (
     p2 - p1) * dot(p3 - p1, p2 - p1) / abs2(p2 - p1); }
pdd reflection(pdd p1, pdd p2, pdd p3)
{ return p3 + perp(p2 - p1
     ) * cross(p3 - p1, p2 - p1) / abs2(p2 - p1) * 2; }
pdd linearTransformation
     (pdd p0, pdd p1, pdd q0, pdd q1, pdd r) {
   pdd dp = p1 - p0
, dq = q1 - q0, num(cross(dp, dq), dot(dp, dq));
   return q0 + pdd(
        cross(r - p0, num), dot(r - p0, num)) / abs2(dp);
\} // from line p0--p1 to q0--q1, apply to r
```

#### 8.2 PointSegDist\* [57b6de]

### 8.3 Heart [468aeb]

```
pair < pdd
    , ld > circenter(pdd a, pdd b, pdd c){ // SCOPE HASH
    pdd m1 = (a+b)/2, m2 = (b+c)/2;
    pdd cent = intersect
        (m1, m1 + perp(b-a), m2, m2 + perp(c-b));
```

```
return {cent, abs(a-cent)};
pair<pdd, ld> incenter
    (pdd p1, pdd p2, pdd p3) { // radius = area / s * 2
 abs(p2 - p3), b = abs(p1 - p3), c = abs(p1 - p2);
pdd cent = (a * p1 + b * p2 + c * p3) / (a + b + c);
  return {cent, abs(a-cent)};
pdd masscenter(pdd p1, pdd p2, pdd p3)
{ return (p1 + p2 + p3) / 3; }
pdd orthcenter(pdd p1, pdd p2, pdd p3)
{ return masscenter
    (p1, p2, p3) * 3 - circenter(p1, p2, p3) * 2; }
```

### 8.4 point in circle [02a7ca]

```
// return q'
    s relation with circumcircle of tri(p[0],p[1],p[2])
bool inCC(const array<pll, 3> &p, pll q) {
    _int128 det = 0;
  for (int i = 0; i < 3; ++i)</pre>
    det += __int128(abs2(p[i]) - abs2(q)) *
         cross(p[(i + 1) % 3] - q, p[(i + 2) % 3] - q);
  return det > 0; // in: >0, on: =0, out: <0
}
```

#### 8.5 Convex hull\* [ca62f8]

```
vector<int> getConvexHull(vector<pdd>& pts){
  vector<int> id(SZ(pts));
  iota(iter(id), 0);
  sort(iter(id
      ), [&](int x, int y){ return pts[x] < pts[y]; });
  vector<int> hull;
  for(int tt = 0; tt < 2; tt++){</pre>
    int sz = SZ(hull);
    for(int j : id){
      pdd p = pts[j];
      while(SZ(hull) - sz >= 2 &&
          cross(pts
              [hull.back()] - pts[hull[SZ(hull) - 2]],
            p - pts[hull[SZ(hull) - 2]]) <= 0)</pre>
        hull.pop_back();
      hull.pb(j);
    hull.pop_back();
    reverse(iter(id));
  return hull;
}
```

### 8.6 PointInConvex\* [82b81e]

```
bool PointInConvex
    (const vector<pll> &C, pll p, bool strict = true) {
  int a = 1, b = SZ(C) - 1, r = !strict;
  if (SZ(C) == 0) return false;
  if (SZ(C) < 3) return r && btw(p, C[0], C.back());</pre>
  if (ori(C[0], C[a], C[b]) > 0) swap(a, b);
  if (ori
      (C[0], C[a], p) >= r \mid\mid ori(C[0], C[b], p) <= -r)
    return false;
  while (abs(a - b) > 1) {
    int c = (a + b) / 2;
    (ori(C[0], C[c], p) > 0 ? b : a) = c;
  return ori(C[a], C[b], p) < r;</pre>
```

#### 8.7 TangentPointToHull\* [523bc1]

```
/* The point should be strictly out of hull
  return arbitrary point on the tangent line */
pii get_tangent(vector<pll> &C, pll p) {
 auto gao = [&](int s) {
    return cyc_tsearch(SZ(C), [&](int x, int y)
    { return ori(p, C[x], C[y]) == s; });
  return pii(gao(1), gao(-1));
\} // return (a, b), ori(p, C[a], C[b]) >= 0
```

## 8.8 Intersection of line and convex [157258]

```
int TangentDir(vector<pll> &C, pll dir) {
 return cyc_tsearch(SZ(C), [&](int a, int b) {
   return cross(dir, C[a]) > cross(dir, C[b]);
 });
```

```
#define cmpL(i) sign(cross(C[i] - a, b - a))
pii lineHull(pil a, pll b, vector<pil> &C) {
  int A = TangentDir(C, a - b);
  int B = TangentDir(C, b - a);
  int n = SZ(C);
  if (cmpL(A) < 0 \mid \mid cmpL(B) > 0)
  return pii(-1, -1); // no collision
auto gao = [&](int l, int r) {
    for (int t = l; (l + 1) % n != r; ) {
       int m = ((l + r + (l < r? 0 : n)) / 2) % n;
       (cmpL(m) == cmpL(t) ? l : r) = m;
     return (l + !cmpL(r)) % n;
  pii res = pii(gao(B, A), gao(A, B)); // (i, j)
if (res.X == res.Y) // touching the corner i
    return pii(res.X, -1);
  if (!
       cmpL(res.X) && !cmpL(res.Y)) // along side i, i+1
     switch ((res.X - res.Y + n + 1) % n) {
       case 0: return pii(res.X, res.X);
       case 2: return pii(res.Y, res.Y);
  /* crossing sides (i, i+1) and (j, j+1)
  crossing corner i is treated as side (i, i+1)
  returned
        in the same order as the line hits the convex */
  return res;
} // convex cut: (r, l]
```

## minMaxEnclosingRectangle\* [180fb8]

```
const double INF = 1e18, qi = acos(-1) / 2 * 3;
pdd solve(vector<pll> &dots) {
#define diff(u, v) (dots[u] - dots[v])
#define vec(v) (dots[v] - dots[i])
  hull(dots);
  double Max = 0, Min = INF, deg;
  int n = SZ(dots);
  dots.pb(dots[0]);
  for (int i = 0, u = 1, r = 1, l = 1; i < n; ++i) {
    pll nw = vec(i + 1);
    while (cross(nw, vec(u + 1)) > cross(nw, vec(u)))
      u = (u + 1) \% n;
    while (dot(nw, vec(r + 1)) > dot(nw, vec(r)))
      r = (r + 1) \% n;
    if (!i) l = (r + 1) \% n;
    while (dot(nw, vec(l + 1)) < dot(nw, vec(l)))</pre>
      l = (l + 1) \% n;
    Min = min(Min, (double)(dot(nw, vec(r)) - dot
        (nw, vec(l))) * cross(nw, vec(u)) / abs2(nw));
    deg = acos(dot(diff(r
        , l), vec(u)) / abs(diff(r, l)) / abs(vec(u)));
    deg = (qi - deg) / 2;
    Max = max(Max, abs(diff))
        (r, l)) * abs(vec(u)) * sin(deg) * sin(deg));
  return pdd(Min, Max);
```

## 8.10 VectorInPoly\* [c6d0fa]

```
b, c) >= 0, valid: "strict" angle from a-b to a-c
bool btwangle(pll a, pll b, pll c, pll p, int strict) {
      ori(a, b, p) >= strict && ori(a, p, c) >= strict;
// whether vector
   {cur, p} in counter-clockwise order prv, cur, nxt
bool inside
    (pll prv, pll cur, pll nxt, pll p, int strict) {
  if (ori(cur, nxt, prv) >= 0)
    return btwangle(cur, nxt, prv, p, strict);
  return !btwangle(cur, prv, nxt, p, !strict);
```

#### 8.11 PolyUnion\* [434d93]

```
ld rat(pll a, pll b) {
  return sgn(b.X) ? (ld)a.X / b.X : (ld)a.Y / b.Y;
  // all poly. should be ccw
ld polyUnion(vector<vector<pll>>> &poly) {
  ld res = 0;
  for (auto &p : poly)
    for (int a = 0; a < SZ(p); ++a) {</pre>
```

```
pll A = p[a], B = p[(a + 1) % SZ(p)];
       vector<pair<ld, int >> segs = \{\{0, 0\}, \{1, 0\}\};
       for (auto &q : poly) {
         if (&p == &q) continue;
         for (int b = 0; b < SZ(q); ++b) {</pre>
           pll C = q[b], D = q[(b + 1) \% SZ(q)];
           int sc = ori(A, B, C), sd = ori(A, B, D);
           if (sc != sd && min(sc, sd) < 0) {</pre>
             ld sa = cross(D
                   - C, A - C), sb = cross(D - C, B - C);
             segs.eb(sa / (sa - sb), sgn(sc - sd));
           if (!sc && !sd &&
                &q < &p && sgn(dot(B - A, D - C)) > 0) {
             segs.eb(rat(C - A, B - A), 1);
             segs.eb(rat(D - A, B - A), -1);
        }
      sort(all(segs));
      for (auto &s : segs) s.X = clamp(s.X, 0.0, 1.0);
      ld sum = 0;
      int cnt = segs[0].second;
      for (int j = 1; j < SZ(segs); ++j) {
  if (!cnt) sum += segs[j].X - segs[j - 1].X;</pre>
         cnt += segs[j].Y;
      res += cross(A, B) * sum;
  return res / 2;
}
```

#### 8.12 Trapezoidalization [4d3bca]

```
template < class T>
struct SweepLine {
  struct cmp {
    cmp(const SweepLine &_swp): swp(_swp) {}
    bool operator()(int a, int b) const {
  if (abs(swp.get_y(a) - swp.get_y(b)) <= swp.eps)</pre>
        return swp.slope_cmp(a, b);
      return swp.get_y(a) + swp.eps < swp.get_y(b);</pre>
    const SweepLine &swp;
  } _cmp;
  T curTime, eps, curQ;
  vector<Line> base;
  multiset < int , cmp > sweep;
  multiset<pair<T, int>> event;
  vector<typename multiset<int, cmp>::iterator> its;
      <typename multiset<pair<T, int>>::iterator> eits;
  bool slope_cmp(int a, int b) const {
    assert(a != -1);
    if (b == -1) return 0;
    return sign(cross(base
        [a].Y - base[a].X, base[b].Y - base[b].X)) < 0;
  T get_y(int idx) const {
    if (idx == -1) return curQ;
    Line l = base[idx];
    if (l.X.X == l.Y.X) return l.Y.Y;
    return ((curTime - l.X.X) * l.Y.Y
        + (l.Y.X - curTime) * l.X.Y) / (l.Y.X - l.X.X);
  void insert(int idx) {
    its[idx] = sweep.insert(idx);
    if (its[idx] != sweep.begin())
      update_event(*prev(its[idx]));
    update_event(idx);
    event.emplace(base[idx].Y.X, idx + 2 * SZ(base));
  void erase(int idx) {
    assert(eits[idx] == event.end());
    auto p = sweep.erase(its[idx]);
    its[idx] = sweep.end();
    if (p != sweep.begin())
      update_event(*prev(p));
  void update_event(int idx) {
    if (eits[idx] != event.end())
      event.erase(eits[idx]);
    eits[idx] = event.end();
    auto nxt = next(its[idx]);
    if (nxt ==
         sweep.end() || !slope_cmp(idx, *nxt)) return;
```

```
auto t = intersect(base[idx].
        X, base[idx].Y, base[*nxt].X, base[*nxt].Y).X;
    if (t + eps < curTime || t</pre>
         >= min(base[idx].Y.X, base[*nxt].Y.X)) return;
    eits[idx] = event.emplace(t, idx + SZ(base));
  void swp(int idx) {
    assert(eits[idx] != event.end());
    eits[idx] = event.end();
    int nxt = *next(its[idx]);
    swap((int&)*its[idx], (int&)*its[nxt]);
    swap(its[idx], its[nxt]);
    if (its[nxt] != sweep.begin())
      update_event(*prev(its[nxt]));
    update_event(idx);
  // only expected to call the functions below
  SweepLine(T t, T e, vector
      <Line> vec): _cmp(*this), curTime(t), eps(e)
       curQ(), base(vec), sweep(_cmp), event(), its(SZ
    (vec), sweep.end()), eits(SZ(vec), event.end()) {
for (int i = 0; i < SZ(base); ++i) {</pre>
      auto &[p, q] = base[i];
      if (p > q) swap(p, q);
      if (p.X <= curTime && curTime <= q.X)</pre>
        insert(i);
      else if (curTime < p.X)</pre>
        event.emplace(p.X, i);
  void setTime(T t, bool ers = false) {
    assert(t >= curTime);
    while (!event.empty() && event.begin()->X <= t) {</pre>
      auto [et, idx] = *event.begin();
      int s = idx / SZ(base);
      idx %= SZ(base);
      if (abs(et - t) <= eps && s == 2 && !ers) break;</pre>
      curTime = et;
      event.erase(event.begin());
      if (s == 2) erase(idx);
      else if (s == 1) swp(idx);
      else insert(idx);
    curTime = t;
  T nextEvent() {
    if (event.empty()) return INF;
    return event.begin()->X;
  int lower_bound(T y) {
    curQ = y;
    auto p = sweep.lower_bound(-1);
    if (p == sweep.end()) return -1;
    return *p;
  }
};
8.13 Polar Angle Sort* [b20533]
  -1: a // b (if same), 0/1: a < b
int cmp(pll a, pll b, bool same = true) {
#define is_neg(k) (
    sign(k.Y) < 0 \mid \mid (sign(k.Y) == 0 \&\& sign(k.X) < 0))
  int A = is_neg(a), B = is_neg(b);
  if (A != B) return A < B;</pre>
  if (sign(cross(a, b)) == 0)
    return same ? abs2(a) < abs2(b) : -1;
  return sign(cross(a, b)) > 0;
8.14 Half plane intersection* [e13bb6]
pll area_pair(Line a, Line b)
{ return pll(cross(a.Y
      - a.X, b.X - a.X), cross(a.Y - a.X, b.Y - a.X)); }
bool isin(Line l0, Line l1, Line l2) {
  // Check inter(l1, l2) strictly in l0
  auto [a02X, a02Y] = area_pair(l0, l2);
  auto [a12X, a12Y] = area_pair(l1, l2);
  if (a12X - a12Y < 0) a12X *= -1, a12Y *= -1;
  return (__int128)
a02Y * a12X - (__int128) a02X * a12Y > 0; // C^4
```

/\* Having solution, check size > 2 \*/ /\* --^-- Line.X --^-- Line.Y --^-- \*/

vector<Line> halfPlaneInter(vector<Line> arr) {

sort(all(arr), [&](Line a, Line b) -> int {

```
if (cmp(a.Y - a.X, b.Y - b.X, 0) != -1)
return cmp(a.Y - a.X, b.Y - b.X, 0);
    return ori(a.X, a.Y, b.Y) < 0;</pre>
  });
  deque<Line> dq(1, arr[0]);
  for (auto p : arr) {
    if (cmp(
         dq.back().Y - dq.back().X, p.Y - p.X, 0) == -1)
       continue;
    while (SZ(dq
        ) >= 2 && !isin(p, dq[SZ(dq) - 2], dq.back()))
       dq.pop_back();
    while (SZ(dq) >= 2 && !isin(p, dq[0], dq[1]))
       dq.pop_front();
    dq.pb(p);
  while (SZ(dq)
        >= 3 && !isin(dq[0], dq[SZ(dq) - 2], dq.back()))
    dq.pop_back();
  while (SZ(dq) >= 3 \&\& !isin(dq.back(), dq[0], dq[1]))
    dq.pop_front();
  return vector<Line>(all(dq));
}
```

### 8.15 RotatingSweepLine [37eaad]

```
struct Event {
  pll d; int u, v;
  bool operator < (const Event &b) const {</pre>
     int ret = cmp(d, b.d, false);
     return ret == -1 ? false : ret; } // no tie-break
void rotatingSweepLine(const vector<pll> &p) {
  const int n = SZ(p);
  vector<Event> e; e.reserve(n * (n - 1));
  for (int i = 0; i < n; i++)</pre>
    for (int j = 0; j < n
          j++) // pos[i] < pos[j] when the event occurs
       if (i != j) e.pb(p[j] - p[i], i, j);
  sort(all(e));
  vector<int> ord(n), pos(n);
  iota(all(ord), 0);
  sort(all(ord), [&](int i, int j) { // initial order
       return p[i].Y != p[
  j].Y ? p[i].Y < p[j].Y : p[i].X < p[j].X; });
for (int i = 0; i < n; i++) pos[ord[i]] = i;</pre>
   // initialize
  for (int i = 0, j = 0; i < SZ(e); i = j) {</pre>
    // do something
     vector<pii> tmp;
    for (; j < SZ(e) && !(e[i] < e[j]); j++)</pre>
       tmp.pb(pii(e[j].u, e[j].v));
     sort(all(tmp), [&](pii x, pii y){
         return pii(pos[x.fi],
             pos[x.se]) < pii(pos[y.fi], pos[y.se]); });</pre>
    for (auto [x, y] : tmp) // pos[x] + 1 == pos[y]
       tie(ord[pos[x]], ord[pos[y]], pos[x], pos[y]) =
         make tuple
             (ord[pos[y]], ord[pos[x]], pos[y], pos[x]);
  }
}
```

#### 8.16 Minimum Enclosing Circle\* [2c7ccd]

```
using ld = long double;
mt19937 rng(950223); // O(n) expected time
pair<pdd, ld> minimumEnclosingCircle(vector<pdd> &pts){
  shuffle(all(pts), rng);
  pdd c = pts[0];
  ld r = 0;
  for(int i = 1; i < SZ(pts); i++){
  if(abs(pts[i] - c) <= r) continue;</pre>
    c = pts[i]; r = 0;
    for(int j = 0; j < i; j++){</pre>
      if(abs(pts[j] - c) <= r) continue;</pre>
      c = (pts[i] + pts[j]) / 2;
       r = abs(pts[i] - c);
       for(int k = 0; k < j; k++){</pre>
         if(abs(pts[k] - c) > r)
           auto [c
                , r] = circenter(pts[i], pts[j], pts[k]);
   }
  return {c, r};
```

### 8.17 Intersection of two circles\* [f7a2fe]

```
bool CCinter(Cir &a, Cir &b, pdd &p1, pdd &p2) {
  pdd o1 = a.0, o2 = b.0;
  double r1 =
        a.R, r2 = b.R, d2 = abs2(o1 - o2), d = sqrt(d2);
  if(d < max
        (r1, r2) - min(r1, r2) || d > r1 + r2) return 0;
  pdd u = (o1 + o2) * 0.5
        + (o1 - o2) * ((r2 * r2 - r1 * r1) / (2 * d2));
  double A = sqrt((r1 + r2 + d) *
        (r1 - r2 + d) * (r1 + r2 - d) * (-r1 + r2 + d));
  pdd v
        = pdd(o1.Y - o2.Y, -o1.X + o2.X) * A / (2 * d2);
  p1 = u + v, p2 = u - v;
  return 1;
}
```

### 8.18 Intersection of polygon and circle\* [d4d295]

```
// Divides into multiple triangle, and sum up
const double PI=acos(-1);
double _area(pdd pa, pdd pb, double r){
  if(abs(pa)<abs(pb)) swap(pa, pb);</pre>
  if(abs(pb)<eps) return 0;</pre>
  double S, h, theta;
  double a=abs(pb),b=abs(pa),c=abs(pb-pa);
  double cosB = dot(pb,pb-pa) / a / c, B = acos(cosB);
  double cosC = dot(pa,pb) / a / b, C = acos(cosC);
  if(a > r){
    S = (C/2)*r*r;
    h = a*b*sin(C)/c;
    if (h < r && B
         < PI/2) S -= (acos(h/r)*r*r - h*sqrt(r*r-h*h));
  else if(b > r){
    theta = PI - B - asin(sin(B)/r*a);
    S = .5*a*r*sin(theta) + (C-theta)/2*r*r;
  else S = .5*sin(C)*a*b;
  return S;
double area_poly_circle(const
     vector<pdd> poly,const pdd &0,const double r){
  double S=0:
  for(int i=0;i<SZ(poly);++i)</pre>
    S+=_area(poly[i]-0,poly[(i+1)%SZ(poly
        )]-0,r)*ori(0,poly[i],poly[(i+1)%SZ(poly)]);
  return fabs(S);
}
```

#### 8.19 Intersection of line and circle\* [76e533]

#### 8.20 Tangent line of two circles [5ac5a5]

```
vector < Line
     > CCtan(const Cir& c1 , const Cir& c2 , int sign1){
   // sign1 = 1 for outer tang, -1 for inter tang
   vector<Line> ret;
   double d sq = abs2(c1.0 - c2.0);
   if (sgn(d_sq) == 0) return ret;
   double d = sqrt(d_sq);
   pdd v = (c2.0 - c1.0) / d;
   double c = (c1.R - sign1 * c2.R) / d;
   if (c * c > 1) return ret;
   double h = sqrt(max(0.0, 1.0 - c * c));
   for (int sign2 = 1; sign2 >= -1; sign2 -= 2) {
  pdd n = pdd(v.X * c - sign2 * h * v.Y,
                 v.Y * c + sign2 * h * v.X);
     pdd p1 = c1.0 + n * c1.R;
     pdd p2 = c2.0 + n * (c2.R * sign1);
     if (!sgn(p1.X - p2.X) && !sgn(p1.Y - p2.Y))
       p2 = p1 + perp(c2.0 - c1.0);
     ret.pb(Line(p1, p2));
   return ret;
}
```

#### 8.21 Minkowski Sum\* [399d43]

```
void reorder_poly(vector<pdd>& pnts){
   int mn = 0;
   for(int i = 1; i < SZ(pnts); i++)</pre>
     if(pnts[i].Y < pnts[mn].Y || (pnts[i].Y ==</pre>
           pnts[mn].Y && pnts[i].X < pnts[mn].X)) mn = i;</pre>
   rotate(pnts.begin(), pnts.begin() + mn, pnts.end());
vector<pdd> minkowski(vector<pdd> P, vector<pdd> Q){
   reorder_poly(P); reorder_poly(Q);
   int psz = SZ(P), qsz = SZ(Q);
  P.pb(P[0]); P.pb(P[1]); Q.pb(Q[0]); Q.pb(Q[1]);
vector<pdd> ans; int i = 0, j = 0;
  while(i < psz || j < qsz){
  ans.pb(P[i] + Q[j]);</pre>
     int t = sgn(cross(P[i + 1]-P[i], Q[j + 1]-Q[j]));
     if(t >= 0) i++; if(t <= 0) j++;</pre>
   return ans;
}
```

#### 9 Else

#### 9.1 Cyclic Ternary Search\* [9017cc]

```
* bool pred(int a, int b);
f(\theta) \sim f(n-1) is a cyclic-shift U-function
return idx s.t. pred(x, idx) is false forall x*/
int cyc_tsearch(int n, auto pred) {
  if (n == 1) return 0;
  int l = 0, r = n; bool rv = pred(1, 0);
while (r - l > 1) {
  int m = (l + r) / 2;
    if (pred(0, m) ? rv: pred(m, (m + 1) % n)) r = m;
    else l = m;
  return pred(l, r % n) ? l : r % n;
```

### Mo's Algorithm(With modification) [f05c5b]

```
Mo's Algorithm With modification
Block: N^{2/3}, Complexity: N^{5/3}
struct Query {
  int L, R, LBid, RBid, T;
  Query(int l, int r, int t):
    L(l), R(r), LBid(l / blk), RBid(r / blk), T(t) {}
  bool operator<(const Query &q) const {</pre>
     if (LBid != q.LBid) return LBid < q.LBid;</pre>
     if (RBid != q.RBid) return RBid < q.RBid;</pre>
     return T < b.T;</pre>
  }
void solve(vector<Query> query) {
  sort(ALL(query));
  int L=0, R=0, T=-1;
for (auto q : query) {
     while (T < q.T) addTime(L, R, ++T); // TODO
while (T > q.T) subTime(L, R, T--); // TODO
     while (R < q.R) add(arr[++R]); // TODO</pre>
    while (L > q.L) add(arr[--L]); // TODO
while (R > q.R) sub(arr[R--]); // TODO
     while (L < q.L) sub(arr[L++]); // TODO</pre>
     // answer query
}
```

#### Mo's Algorithm On Tree [8331c2] 9.3

```
Mo's Algorithm On Tree
Preprocess:
1) LCA
2) dfs with in[u] = dft++, out[u] = dft++
3) ord[in[u]] = ord[out[u]] = u
4) bitset<MAXN> inset
struct Query {
  int L, R, LBid, lca;
  Query(int u, int v) {
    int c = LCA(u, v);
    if (c == u || c == v)
      q.lca = -1, q.L = out[c ^ u ^ v], q.R = out[c];
    else if (out[u] < in[v])</pre>
      q.lca = c, q.L = out[u], q.R = in[v];
```

```
else
       q.lca = c, q.L = out[v], q.R = in[u];
     q.Lid = q.L / blk;
   bool operator < (const Query &q) const {</pre>
     if (LBid != q.LBid) return LBid < q.LBid;</pre>
     return R < q.R;</pre>
  }
void flip(int x) {
     if (inset[x]) sub(arr[x]); // TODO
else add(arr[x]); // TODO
     inset[x] = ~inset[x];
void solve(vector<Query> query) {
  sort(ALL(query));
   int L = 0, R = 0;
  for (auto q : query) {
     while (R < q.R) flip(ord[++R]);</pre>
     while (L > q.L) flip(ord[--L]);
     while (R > q.R) flip(ord[R--]);
     while (L < q.L) flip(ord[L++]);</pre>
     if (~q.lca) add(arr[q.lca]);
     // answer query
     if (~q.lca) sub(arr[q.lca]);
  }
}
```

#### 9.4 Additional Mo's Algorithm Trick

- · Mo's Algorithm With Addition Only
- Sort querys same as the normal Mo's algorithm.
  - For each query [l,r]:
  - If l/blk = r/blk, brute-force.
  - If  $l/blk \neq curL/blk$ , initialize  $curL := (l/blk+1) \cdot blk$ , curR := curL-1
  - If r > curR, increase curR
- decrease curL to fit l, and then undo after answering
- Mo's Algorithm With Offline Second Time
  - Require: Changing answer  $\equiv$  adding f([l,r],r+1).
  - Require: f([l,r],r+1) = f([1,r],r+1) f([1,l),r+1).
  - Part1: Answer all f([1,r],r+1) first.
  - Part2: Store  $curR \rightarrow R$  for curL (reduce the space to O(N)), and then answer them by the second offline algorithm.
- Note: You must do the above symmetrically for the left boundaries. **Hilbert Curve** [1274a3]

```
ll hilbert(int n, int x, int y) {
  ll res = 0;
  for (int s = n / 2; s; s >>= 1) {
    int rx = (x \& s) > 0;
    int ry = (y & s) > 0;
res += s * 1ll * s * ((3 * rx) ^ ry);
    if (ry == 0) {
      if (rx == 1) x = s - 1 - x, y = s - 1 - y;
       swap(x, y);
    }
  }
  return res;
 // n = 2^k
```

#### 9.6 DynamicConvexTrick\* [673ffd]

```
// only works for integer coordinates!! maintain max
struct Line {
  mutable ll a, b, p;
  bool operator
       <(const Line &rhs) const { return a < rhs.a; }
  bool operator<(ll x) const { return p < x; }</pre>
};
struct DynamicHull : multiset<Line, less<>>> {
  static const ll kInf = 1e18;
  ll Div(ll a,
       ll b) { return a / b - ((a ^ b) < 0 && a % b); }
  bool isect(iterator x, iterator y) {
    if (y == end()) { x->p = kInf; return 0; }
    if (x
    ->a == y->a) x->p = x->b > y->b ? kInf : -kInf; else x->p = Div(y->b - x->b, x->a - y->a);
    return x->p >= y->p;
  void addline(ll a, ll b) {
    auto z = insert({a, b, 0}), y = z++, x = y;
    while (isect(y, z)) z = erase(z);
    if (x != begin
         () && isect(--x, y)) isect(x, y = erase(y));
    while ((y = x) != begin
         () && (--x)->p >= y->p) isect(x, erase(y));
```

```
ll auerv(ll x) {
     auto l = *lower_bound(x);
     return l.a * x + l.b;
};
9.7
      All LCS* [78a378]
void all_lcs(string s, string t) { // 0-base
  vector<int> h(SZ(t));
  iota(ALL(h), 0);
  for (int a = \theta; a < SZ(s); ++a) {
    int v = -1;
     for (int c = 0; c < SZ(t); ++c)</pre>
       if (s[a] == t[c] || h[c] < v)</pre>
         swap(h[c], v);
     // LCS(s[0, a], t[b, c]) =
// c - b + 1 - sum([h[i] >= b] | i <= c)
    // h[i] might become -1 !!
}
9.8 DLX* [fbcf6c]
#define TRAV(i, link, start)
      for (int i = link[start]; i != start; i = link[i])
template <
    bool E> // E: Exact, NN: num of 1s, RR: num of rows
struct DLX {
  int lt[NN], rg[NN], up[NN], dn[NN
       ], rw[NN], cl[NN], bt[NN], s[NN], head, sz, ans;
  int rows, columns;
  bool vis[NN];
  bitset<RR> sol, cur; // not sure
  void remove(int c) {
     if (E) lt[rg[c]] = lt[c], rg[lt[c]] = rg[c];
    TRAV(i, dn, c) {
       if (E) {
        TRAV(j, rg, i)
           up[dn[j]]
                = up[j], dn[up[j]] = dn[j], --s[cl[j]];
      } else {
         lt[rg[i]] = lt[i], rg[lt[i]] = rg[i];
    }
  }
  void restore(int c) {
    TRAV(i, up, c) {
      if (E) {
         TRAV(j, lt, i)
           ++s[cl[j]], up[dn[j]] = j, dn[up[j]] = j;
      } else {
         lt[rg[i]] = rg[lt[i]] = i;
       }
     if (E) lt[rg[c]] = c, rg[lt[c]] = c;
  void init(int c) {
     rows = 0, columns = c;
     for (int i = 0; i < c; ++i) {</pre>
      up[i] = dn[i] = bt[i] = i;
       lt[i] = i == 0 ? c : i - 1;
       rg[i] = i == c - 1 ? c : i + 1;
      s[i] = 0;
    rg[c] = 0, lt[c] = c - 1;
    up[c] = dn[c] = -1;
    head = c, sz = c + 1;
  void insert(const vector<int> &col) {
    if (col.empty()) return;
     int f = sz;
     for (int i = 0; i < (int)col.size(); ++i) {</pre>
       int c = col[i], v = sz++;
       dn[bt[c]] = v;
      up[v] = bt[c], bt[c] = v;
rg[v] = (i + 1 == (int)col.size() ? f : v + 1);
       rw[v] = rows, cl[v] = c;
       ++s[c];
      if (i > 0) lt[v] = v - 1;
    }
    ++rows, lt[f] = sz - 1;
  int h() {
    int ret = 0;
     fill_n(vis, sz, false);
    TRAV(x, rg, head) {
```

```
if (vis[x]) continue;
       vis[x] = true,
      TRAV(i, dn, x) TRAV(j, rg, i) vis[cl[j]] = true;
     return ret;
  void dfs(int dep) {
     if (dep + (E ? 0 : h()) >= ans) return;
     if (rg[head
         ] == head) return sol = cur, ans = dep, void();
     if (dn[rg[head]] == rg[head]) return;
     int w = rg[head];
     TRAV(x, rg, head) if (s[x] < s[w]) w = x;
     if (E) remove(w);
     TRAV(i, dn, w) {
      if (!E) remove(i);
      TRAV(j, rg, i) remove(E ? cl[j] : j);
       cur.set(rw[i]), dfs(dep + 1), cur.reset(rw[i]);
       TRAV(j, lt, i) restore(E ? cl[j] : j);
       if (!E) restore(i);
     if (E) restore(w);
   int solve() {
     for (int i = 0; i < columns; ++i)</pre>
      dn[bt[i]] = i, up[i] = bt[i];
     ans = 1e9, sol.reset(), dfs(0);
     return ans;
};
```

#### 9.9 Matroid Intersection

Start from  $S = \emptyset$ . In each iteration, let

•  $Y_1 = \{x \notin S \mid S \cup \{x\} \in I_1\}$ •  $Y_2 = \{x \notin S \mid S \cup \{x\} \in I_2\}$ 

If there exists  $x \in Y_1 \cap Y_2$ , insert x into S. Otherwise for each  $x \in S, y \notin S$ , create edges

•  $x \to y \text{ if } S - \{x\} \cup \{y\} \in I_1.$ 

•  $y \to x$  if  $S - \{x\} \cup \{y\} \in I_2$ .

Find a *shortest* path (with BFS) starting from a vertex in  $Y_1$  and ending at a vertex in  $Y_2$  which doesn't pass through any other vertices in  $Y_2$ , and alternate the path. The size of  $\hat{S}$  will be incremented by 1 in each iteration. For the weighted case, assign weight w(x) to vertex x if  $x \in S$  and -w(x) if  $x \notin S$ . Find the path with the minimum number of edges among all minimum

# length paths and alternate it. 9.10 AdaptiveSimpson\* [4074b3]

```
template < typename Func, typename d = double >
struct Simpson {
  using pdd = pair<d, d>;
  Func f;
  pdd mix(pdd l, pdd r, optional<d> fm = {}) {
   d h = (r.X - l.X) / 2, v = fm.value_or(f(l.X + h));
   return {v, h / 3 * (l.Y + 4 * v + r.Y)};
  d eval(pdd l, pdd r, d fm, d eps) {
     pdd m((l.X + r.X) / 2, fm);
     d s = mix(l, r, fm).second;
     auto [flm, sl] = mix(l, m);
     auto [fmr, sr] = mix(m, r);
     d delta = sl + sr - s;
     if (abs(delta
          ) <= 15 * eps) return sl + sr + delta / 15;
     return eval(l, m, flm, eps / 2) +
       eval(m, r, fmr, eps / 2);
  d eval(d l, d r, d eps) {
     return eval
          ({l, f(l)}, {r, f(r)}, f((l + r) / 2), eps);
  d = val2(d l, d r, d eps, int k = 997) {
     d h = (r - l) / k, s = 0;
     for (int i = 0; i < k; ++i, l += h)
s += eval(l, l + h, eps / k);</pre>
     return s;
  }
};
template < typename Func >
\label{eq:simpson} \mbox{Simpson(Func f) } \{ \mbox{ $return $\{f\}$; } \}
9.11 Simulated Annealing [de78c6]
```

```
double factor = 100000;
const int base = 1e9; // remember to run ~ 10 times
for (int it = 1; it <= 1000000; ++it) {</pre>
    // ans:
        answer, nw: current value, rnd(): mt19937 rnd()
```

```
National Taiwan University 5 Hours Codebook
    if (exp(-(nw - ans
                                                              print(Decimal(input()) * Decimal(input()))
         ) / factor) >= (double)(rnd() % base) / base)
                                                              from fractions import Fraction
         ans = nw;
                                                              Fraction
    factor *= 0.99995;
                                                                   ('3.14159').limit_denominator(10).numerator # 22
}
9.12 Tree Hash* [34aae5]
ull seed;
ull shift(ull x) {
  x ^= x << 13;
  x ^= x >> 7;
  x ^= x << 17;
  return x;
ull dfs(int u, int f) {
  ull sum = seed;
  for (int i : G[u])
    if (i != f)
      sum += shift(dfs(i, u));
  return sum;
9.13 Binary Search On Fraction [765c5a]
struct Q {
  ll p, q;
  Q go(Q b, ll d) { return \{p + b.p*d, q + b.q*d\}; \}
bool pred(Q);
// returns smallest p/q in [lo, hi] such that
  pred(p/q) is true, and 0 <= p,q <= N
Q frac_bs(ll N) {
  Q lo{0, 1}, hi{1, 0};
  if (pred(lo)) return lo;
  assert(pred(hi));
  bool dir = 1, L = 1, H = 1;
  for (; L || H; dir = !dir) {
    ll len = 0, step = 1;
    for (int t = 0; t < 2 && (t ? step/=2 : step*=2);)</pre>
      if (Q mid = hi.go(lo, len + step);
          mid.p > N || mid.q > N || dir ^ pred(mid))
      else len += step;
    swap(lo, hi = hi.go(lo, len));
    (dir ? L : H) = !!len;
  return dir ? hi : lo;
}
9.14 Min Plus Convolution* [09b5c3]
// a is convex a[i+1]-a[i] <= a[i+2]-a[i+1]
vector<int> min_plus_convolution
    (vector<int> &a, vector<int> &b) {
  int n = SZ(a), m = SZ(b);
vector < int > c(n + m - 1, INF);
auto dc = [&](auto Y, int l, int r, int jl, int jr) {
    if (l > r) return;
    int mid = (l + r) / 2, from = -1, &best = c[mid];
    for (int j = jl; j <= jr; ++j)
  if (int i = mid - j; i >= 0 && i < n)</pre>
        if (best > a[i] + b[j])
          best = a[i] + b[j], from = j;
    Y(Y, l,
        mid - 1, jl, from), Y(Y, mid + 1, r, from, jr);
  };
  return dc(dc, 0, n - 1 + m - 1, 0, m - 1), c;
9.15 Bitset LCS [330ab1]
cin >> n >> m;
for (int i = 1, x; i <= n; ++i)</pre>
  cin >> x, p[x].set(i);
for (int i = 1, x; i <= m; i++) {</pre>
  cin >> x, (g = f) |= p[x];
  f.shiftLeftByOne(), f.set(0);
  ((f = g - f) ^= g) &= g;
```

## 10 Python 10.1 Misc

cout << f.count() << '\n';</pre>