

PROPOSITION is

Tautology \longrightarrow Always True

Contradiction \longrightarrow Always False

Contingency \longrightarrow Sometimes True Sometimes False

✓ Logical equivalence

Each statements in ~~both~~ Truth Table has same truth value for two different expressions.

✓ inverse

$$p \rightarrow q \xrightarrow{\text{inverse}} \sim p \rightarrow \sim q$$

✓ converse

$$p \rightarrow q \xrightarrow{\text{converse}} q \rightarrow p$$

✓ contra-positive

$$p \rightarrow q \xrightarrow{\text{contra-positive}} \sim q \rightarrow \sim p$$

✓ ATOMIC PROPOSITION

COMPOUND PROPOSITION

\downarrow Well Formed Formulas (wffs)

✓ Representation using BOOLEAN ALGEBRA
LOGIC GATE

Duality Principle

Dual of $(A \cap B) \cup C$ is $(A \cup B) \cap C$

- interchange union into intersection & vice versa
- interchange Null set with universal set & vice versa

Deductions

If $\boxed{\text{Vijay eats his vegetable}}$ then $\boxed{\text{he can have cookie.}}$

Vijay ate his vegetable.

Therefore, he gets a cookie.

$$\frac{P \rightarrow Q \quad P}{\therefore Q}$$

P	Q	$P \rightarrow Q$
T	T	T
T	F	F
F	T	T
F	F	T

$$P \rightarrow Q \quad \sim P \vee Q$$

$$\frac{P \rightarrow Q}{P \vee Q}$$

$(P \rightarrow Q) \text{ and } P$

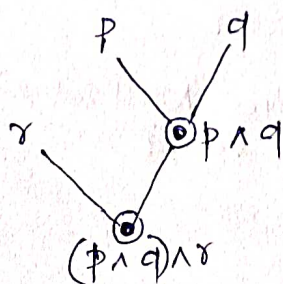
$(P \rightarrow Q) \text{ and } P = T$

T	T	T
F	T	F
T	F	T
F	F	T

When $P \rightarrow Q = T$, $P = T$ then $Q = T$
(~~1st~~ ^{1st} row of the truth table)

This is a valid deduction rule.

Expression tree



Normal Forms

- Disjunctive Normal Form (DNF)

$$(P \wedge Q) \vee (\neg P \wedge \neg Q)$$

- Conjunctive Normal Form (CNF)

$$(P \vee Q) \wedge (\neg P \vee \neg Q)$$

$$P \cdot Q + \bar{P} \bar{Q} \quad \underline{\text{DNF}}$$

$$(P + Q) \cdot (\bar{P} + \bar{Q}) \quad \underline{\text{CNF}}$$

Conjunction $P \wedge Q$
Disjunction $P \vee Q$

Arguments

$P_1, P_2, P_3, \dots, P_n \vdash Q$ is valid if

Q is true whenever $P_1, P_2, P_3, \dots, P_n$ are true.

If an argument is not valid then it is called fallacy.

$$\frac{P \rightarrow Q \quad P}{\therefore Q}$$

$$\longrightarrow P \rightarrow Q, P \vdash Q$$

Valuation

Each row of truth table corresponds to valuations

p	q	F
T	T	T
T	F	F
F	T	T
F	F	F

$$V(x)$$

Valuation is a function that takes proposition x & produces truth value $\{T, F\}$

Interpretation

An interpretation function takes a proposition formula Φ and a valuation V , and returns truth value of the formula.

$$\Phi = p \rightarrow (p \vee q)$$

$$\text{Valuation } V(p) = T \text{ and } V(q) = F$$

Then interpretation function

$$\Phi^V = V(p) \rightarrow (V(p) \vee V(q))$$

$$= T \rightarrow (T \vee F)$$

$$= T \rightarrow T$$

$$= T$$

Satisfiability & Validity

A proposition that has at least one ^{true} interpretation is called satisfiable.

A proposition that is false in all interpretations is called unsatisfiable.

True in all interpretation — valid
False in at least one interpretation — invalid