Graph - 3Wednesday December 2, 2020 8:48 AM Graph Connectivity Vertex womedning, K(G) G= Kn, a complète graph of n vertices. $K(Kn) = (n-1) \rightarrow no$ of vertices to be removed to produce a graph with single vertex In general, $0 \le K(G) \le (n-1)$, if G has n vertices. ~ K(G)=B, iff G is disconnected, - Larger the value of K(G), more connected the graph. Gy: cut vertex:b,c,e. Vertex cut = {b} or {c} or {e} $K(G_1) = 1.$ Cut vertex: C. Vertex ut = {C}

Quick Notes Page 1

Gis: out vertex =
$$\phi$$

Vertex out = $\{b, f, g\}$
 $\mathcal{L}(G_5) = 3$

Edge connectivity

$$\lambda(G)=0$$
, if G is not tomested and consists of one vertex.

Foragraph with n Vertices,

cut edge =
$$\phi$$

edge cut = $S(a,b)$, (a,q) , $S(b,c)$, (g,f) ...
Edge connectivity, $\gamma(G_3) = 2$.

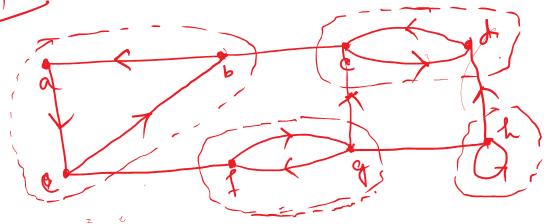
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edge cut = { (n, c), (a, f), (f, 7). $\gamma(G_3) = 3.$

For $G_1 = (V, E)$ to be a non-complete connected graph with [V] >3, the minimum Regree of a vertex of G is the upper bound for both the vertex connectivity and edge Connectivity of Gr.

i.e. $K(G) \leq min_{v \in V} deg(v)$ and

Y(G) & min ver deg(o)



v 4 SCCs.

~ For a graph to have SSCs, it must have cycles.

Both G & H have simple arounts of length three,

four, and five.

All four isomorphic invariants agree.

GDH are potentially isomorphic.