

Module 2: Application Layer (Lecture – 4)

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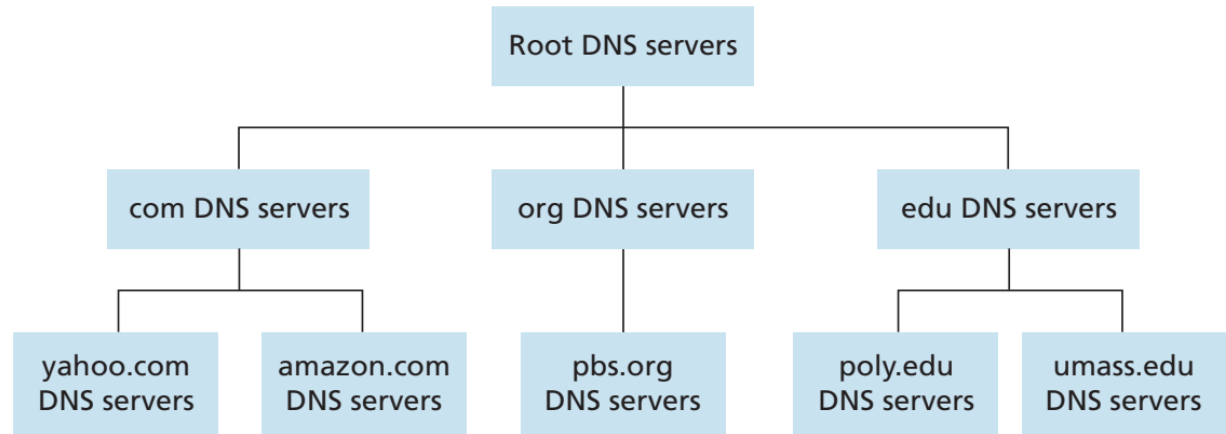
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Domain Name System (DNS)

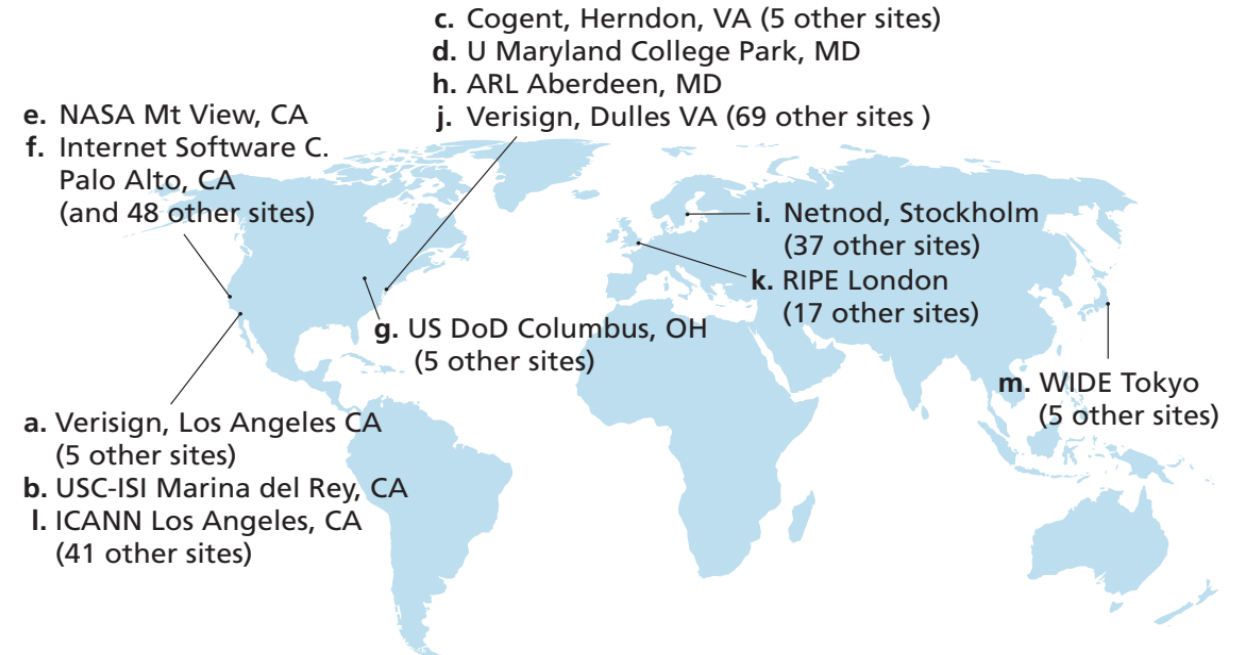
- DNS services other than **hostname-to-IP address translation**
 - **Host aliasing**
 - A host with **complicated hostname (canonical hostname)** can have **one or more alias names**
 - **Alias name: typically more mnemonic than canonical hostnames**
 - Example: **relay1.west-coast.enterprise.com** (canonical name) can have **enterprise.com** or **www.enterprise.com** as alias names
 - DNS can be invoked by an application to obtain **canonical hostname** for a supplied **alias hostname** as well as the **IP address** of the host
 - **Load distribution**
 - Performs **load distribution** among **replicated servers** (e.g., replicated Web servers) running on **different end system** and each having **different IP address**
 - DNS database contains the **list of replicated servers**
 - In response to typical **name-to-address translation query**, DNS server responds with the **entire set of IP addresses** – changes the **ordering of the addresses** with each reply
 - Client typically **sends HTTP request message** to the **IP address listed first in the set** – **distributes traffic among replicated servers**
- **DNS: both service and protocol**
 - **Service: a simple straight-forward name-to-address translation service**
 - **Complex in nature**, consists of a **large number of DNS servers** distributed across the globe/Internet
 - **Protocol: an application-layer protocol** that specifies **how** the **DNS servers** and **querying hosts** communicate
 - **Why is DNS distributed by design?**
 - **Centralized DNS server: does not scale**
 - **A single point of failure:** if the **DNS server** crashes, so does the **entire Internet**
 - **Traffic volume:** has to handle **all DNS queries** across the Internet
 - **Distant centralized database:** cannot be “close” to all querying clients - **significant delays** for clients who **far away**
 - **Maintenance:** have to maintain a **huge database** for all Internet hosts – need to **update frequently** to account for **every new join and leave**

DNS: A distributed, hierarchical data

- DNS: uses a **large number of databases**; organized in a **hierarchical fashion**; **distributed** around the **world**
- **No single DNS server** has **all** of the **mappings** for all hosts in the Internet
- **Three classes** of DNS servers are **organized hierarchically** (see fig.):
 - Root DNS servers
 - Top-level domain (TLD) DNS servers
 - Authoritative DNS servers
- **Root DNS Servers**
 - 13 root DNS servers exist in the Internet (as of 2012)
 - Each server is actually a network of replicated servers for both security and reliability



Portion of the Hierarchy of DNS Servers



DNS: A distributed, hierarchical database

• Top-level Domain (TLD) DNS Servers

- Responsible for top level domains – **com, org, net, edu, gov**, and all **country top-level domains** (uk, fr, ca, jp, in, etc.)
- **Verisign Global Registry**: maintains the TLD servers for the “**com**” top-level domain
- **Educause**: maintains the TLD servers for the “**edu**” top-level domain

• Authoritative DNS Servers

- Contain **organization-specific DNS records**
- **Map** the **names** of **publicly accessible hosts** (such as Web servers and mail servers) on the **Internet** to **IP addresses**
- Most universities and large companies **implement and maintain** their own **primary and secondary (backup) authoritative DNS servers**
- Alternatively, the organizations can pay to have these records stored in the authoritative DNS server of some service provider

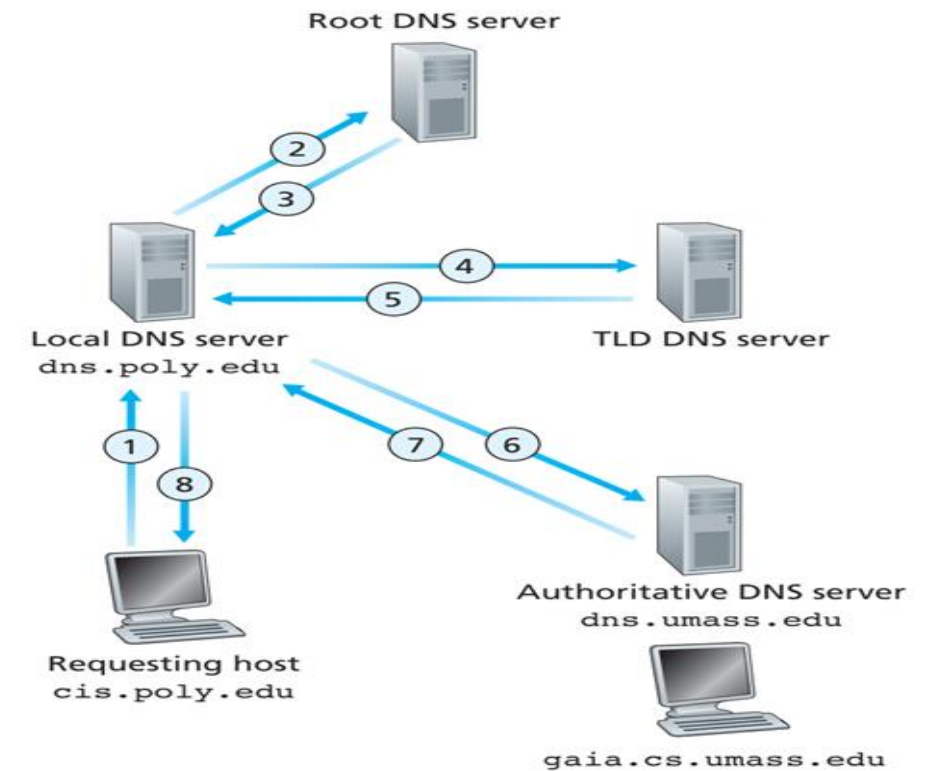


Fig.1: Interaction of various DNS Servers (Iterative)

• Local DNS Server

- Does not strictly belong to the hierarchy of servers, but is nevertheless central to the DNS architecture
- Each ISP (e.g., a university, company, etc.) has a **local DNS server** (also called default name server)
- When a host connects an ISP, the ISP provides the **IP address of one or more of its local DNS servers** (typically through DHCP)
- **Local DNS server: close to the host**
- Acts as a **proxy** to a **host's DNS query** and **forwards it into the DNS server hierarchy**

DNS: A distributed, hierarchical database

• Local DNS Server

- Makes use of both **recursive** and **iterative** queries
- In Fig. 1, query sent from the requesting host (cis.poly.edu) to the local DNS server (dns.poly.edu) is **recursive** - query asks the **local DNS server** to obtain **mapping** on behalf of the **host**
 - The rest of the queries are **iterative** as all the replies are **directly returned** to the **local DNS server**
- Fig. 2 shows a **DNS query chain** for which all of the queries are **recursive**
- In practice, the queries typically follow the pattern in Fig. 1
 - Query from the **requesting host** to the **local DNS server** is **recursive**, and the **remaining queries** are **iterative**

• DNS Caching

- Exploits caching in order to **improve the delay performance** and to **reduce the number of DNS messages** in the Internet
- **Local DNS server** can cache a **received mapping** (from hostname to IP address) in its **memory**
- If a query arrives for the name hostname, the DNS server can provide the **desired IP address** from its **local cache** (even if the DNS server is not authoritative for the hostname)
- **Discards cache information** after **period of time** as **mapping between hostname and IP address** is not permanent
- Local DNS server can also cache **IP addresses of TLD servers** – by **passes visit to the root DNS servers** in the query chain

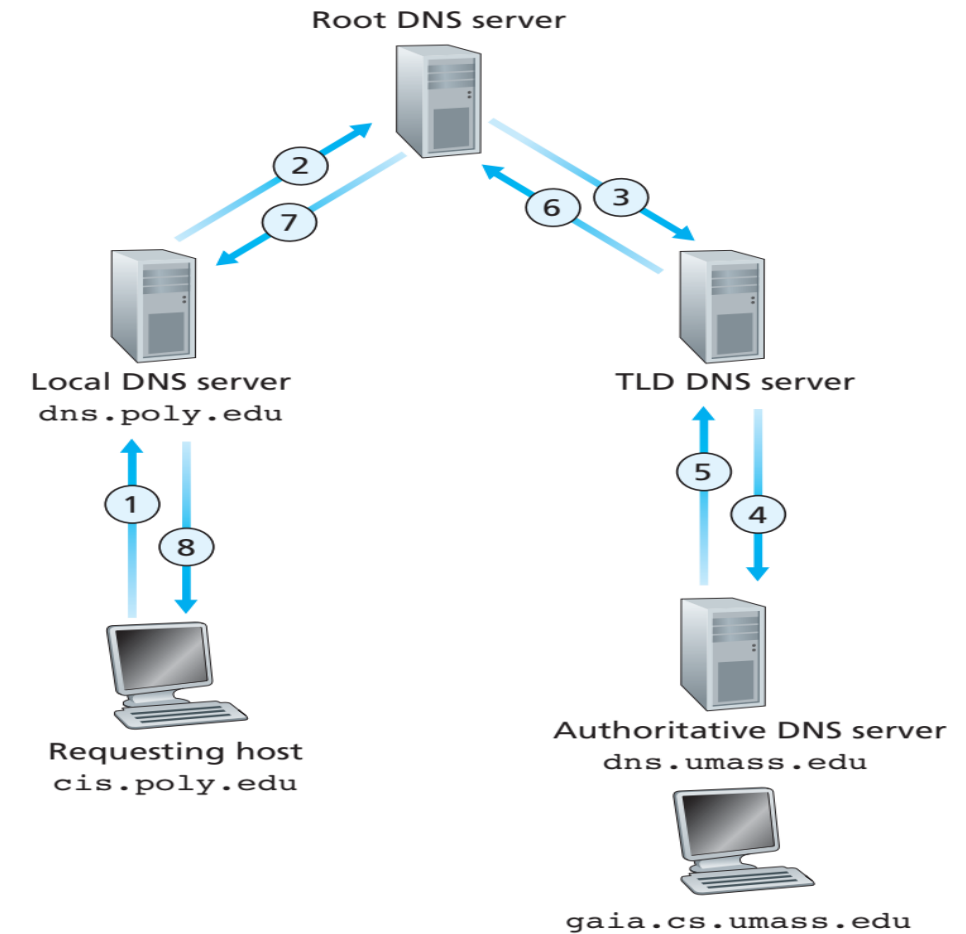


Fig.2: Recursive Queries in DNS

DNS Resource Records

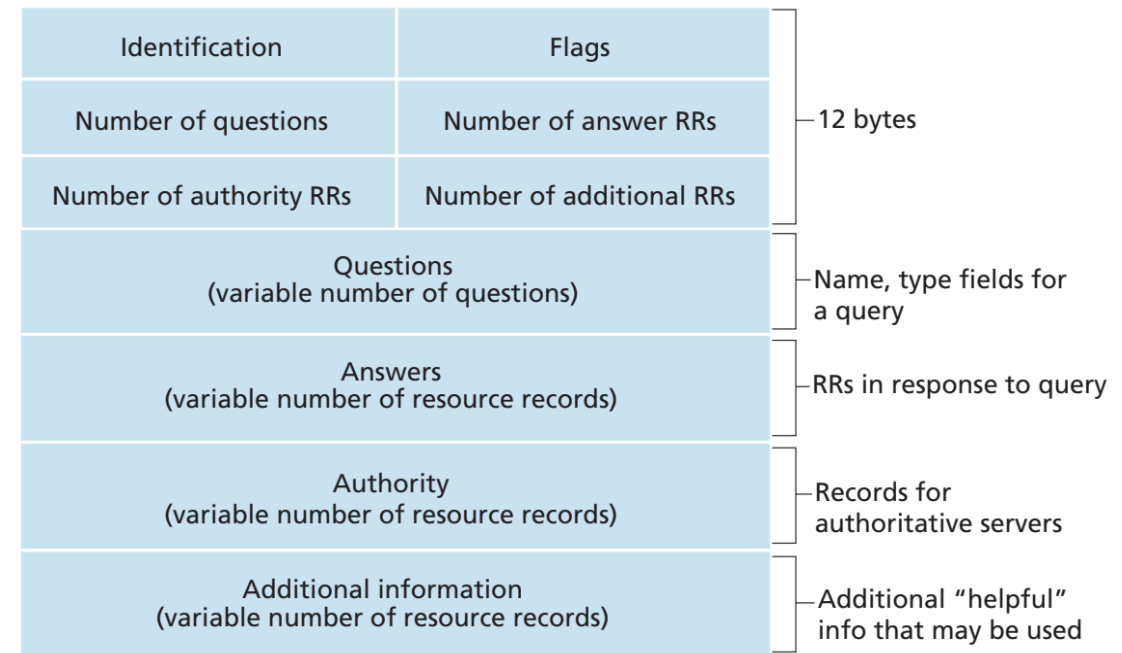
- DNS distributed databases store **resource records (RR)** including RRs that provide **hostname-to-IP address mappings**
- Each **DNS reply message** carries **one or more** resource records
- Resource record: a **four-tuple** containing the following fields: (*Name, Value, Type, TTL*)
- TTL: time-to-live of the resource record – determines when a resource record should be removed from the cache
- Different types of **resource records** that are replied back through **DNS messages** are tabularized below:

Type	Name	Value	Response Message (Example)
A	Hostname	IP address	<i>relay1.bar.foo.com, 145.37.93.126, A</i>
NS	Domain (foo.com)	Hostname of an authoritative DNS server that knows how to obtain the IP addresses for the hosts in the domain	<i>foo.com, dns.foo.com, NS</i>
CNAME	Alias hostname	Canonical hostname	<i>foo.com, relay1.bar.foo.com, CNAME</i>
MX	Alias hostname of a mail server	Canonical hostname of the mail server	<i>foo.com, mail.bar.foo.com, MX</i>

DNS Messages

- **Header section (12-bytes)**

- **Identification (16-bit)**: identifies the query – copied into the reply message to a query for client's convenience
- **Flags**: consists of a number of flags; some of them are:
 - A **1-bit query/reply flag**: query (0) or reply (1)
 - A **1-bit authoritative flag**: set in a reply message if a DNS server is an authoritative server for a queried name
 - A **1-bit recursive-desired flag**: set when a client (host or DNS server) desires that the DNS server performs recursion if it doesn't have the record
 - A **1-bit recursion-available field**: set in a reply if the DNS server supports recursion
- **Four “number-of” fields**: number of occurrences of the four types of data sections that follow the header
- **Question section**: contains **information** about the **query** that is being made
 - **Name field** (contains the name being queried)
 - **Type field** (type of question being asked about the name)



DNS Message Format

- **Answer section**: contains the resource records for the name that was originally queried
 - **Resource record – Type, Value, TTL**
 - A reply can have **multiple RRs** in the answer since a **hostname** can have **multiple IP addresses** (e.g., replicated Web server)
- **Authority section**: contains **information** of other authoritative servers
- **Additional section**: contains other helpful records
 - **Reply to MX query** – providing the **canonical hostname** of a mail server
 - **Type A record** – **IP address** of the **canonical hostname** of the mail server

Inserting Records into the DNS Database

- For any **new startup company**, the first step is to **register its domain name**
- **Registrar: commercial entity that verifies the uniqueness of the domain name**
 - Accredited by the **Internet Corporation for Assigned Names and Numbers (ICANN)**
 - Responsible **for entering the domain name into the DNS database** at the cost of fees
- Organization needs to provide the registrar with the **names** and **IP addresses** of its **primary and secondary authoritative DNS servers**
- Registrar makes sure that a **Type NS (name) record** and a **Type A (IP address) record** are entered into the **corresponding TLD servers (e.g., “com”)**
- Example: for a startup domain ***somecompany.com***, following RRs are entered -
 - ***(somecompany.com, dns1. somecompany.com, NS)***
 - ***(dns1. somecompany.com, 212.212.212.1, A)***
 - ***(dns2.somecompany.com, 212.212.212.2, A)***
- **Type A RR for the Web server (*www.somecompany.com*)** and the **Type MX RR for the mail server (*mail.somecompany.com*)** are entered into the **authoritative DNS servers**
 - Configured statically by the system manager