Relation - 3 Monday, October 12, 2020 10:35 AM

Closure of a relation $R = \{(a,b) | a(b) | U \} (a,a) | a \in Z\}$ $R = \{(a,b) | a(b) | a(b) \}$

Inverse of a relation

If R is a relation from A to B (RCAXB) then
the relation RT from B to A can be defined
by interchanging the elements of All the ordered
pains of R.

RT = { (y, x) EBXA (n,y) ER}.

Prito: R = { (a, b) | a (Db).

 $|RUR^{-1} = \{(a,b) | a > b\} \cup \{(b,a) | a < b\}$

=> RURT = { (1, b) | a = b}.

Theorem 1:

(i) (a,b) ERⁿ -> path of length n from a tob.

(ii) Path of length n from a tob -> (a,b) ERⁿ

(ii) Path of length on from a tob -> (a,b) ER" Proof(i): Let (a, b) ER" is true from for n= 1,2,3.... It follows that (a, b) ER implies. (a, c) ER and (c, b) ERn-1 for some CE A day the definion of composition of (a, b) ER'n. Because (a,c) ER means there is a josth of length I from a to C, and (c,b) ER "> means there is a path of length (n-i) from C to b. It follows that there is a path of length in from a to b. l'ruf (ii): We will use mathematical industion to prove

: By definition, for n=1, a path exists b) w a and b. Therefore, (a, b) ER (basis step)

Assume that the theorem is true for the printive integer n i.e. (a,b) E Rh (inductive hypothesis)

ve To complete the proof we have to show (a,b) fr" there is a path of length (n+i) from a to b

Iff there is an element CEAS.t.

there is a path of length one from a to C i.e. (a, () & R (basis step), and a path of length n from c to b i.e. (c, b) & R (inductive hypothesis). Consequently, by using basis step and inductive hypothesis, there is a path of length (n+1) from a to b iff (a,b) & R. This completes our proof.

 $A = \{1, 2, 3, 4\}$ 2 + 3 2 + 3 2 + 3 3 + 4 4 + 4

 $R = \left\{ (1, 2), (2, 3), (3, 4), (1, 4) \right\}$ $R^{2} = \left\{ (1, 3), (2, 4) \right\}$ $R = \left\{ (1, 3), (2, 4) \right\}$

R = {1,4}

(i) $R^{n} \subseteq R \rightarrow R$ is transitive (ii) Ris transitive -> Rn CR. Profin: Let Rn CR is true for n=1,2,3,.... : It follows that RER torhich implies that if (a, b) ER and (b, c) ER, then by the definition of composition (a,c) ER. Because, Rt C'R, this means (A,c) ER. Hence, Ris transitive. Prof(i): We will use mathematical induction to prove if Ristransitive, then RMCR. By definonaleton, for n=1, R C R (basis) Assume that this theorem is true for the positive integer n i.e. Rn C R (inductive nypothesis) To complete the inductive step, we must show that RM+1 CR. Let (a, b) E Rn+1 Then as R = ROR, there is an element C with CEA, s.t. (a,c) ER and (c,b) ER. Consequently,

Using inductive hypothesis, with R' C C then (c, b) ER: Furthermore, as Rio transitive, and (a, c) ER and (c, DER, it follows that (a, b) ER. This shows that Rnt1 CR, Completing the front. Krost-2. To prove that R is the transitive chase of a relation R we need to prove: (i) R* Contains R (R* 2R) ~ (ii) R* is transitive R* = UR R= n=1 (III) R* is contained in every relation S that contains R. (i) By definition of connectivity relation,

R* 2 R (R* Contains R) (1i) If (a, b) ER and (b, c) ER, then there are paths from a to b and from b to c. We obtain a path from a to c by starting from a and reach b, then from Starting from b to reach C. Hence, (a, c) FR. Therefore, R* is transitive.

(111) Duppose S is a transitive relation Containing R. Because, S is transitive, Sn is Also transitive and Sn C S (from theorem - 2). Furthermore, be once (S*= US* and SKCS, it follows that S C S. Now, it is to be noted that if RCS, then R*CS*, because consequently, R*CS*CS. Thus, any retransitive relation that Contains R must also contain R*.

Therefore, R* is the transitive closure of @ R= { (a, b) | a has niet b} R' - consists of All pairs (a, b) such that there are people 21, 72, -2 2n-1 such that a has met 21, 21 has met 22,, and une has met b.

R* -> Contains (a, b) if there is a segnence
of people, Starting with a and
enting with b, S. t. each person in the
segnence has met the next term in the
Segnence.

Problem.

 $R = \{(a, b) | a \text{ shones border with b} \}$

Rn - Assessists of the pains (a, b) when it is possible to go from state on to state by crossing exactly a state borders.

R* -> Consists of pairs (a, b) where it is possible to go from state a to b by Consing as many state borders as necessary.