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# LISP

## Programming Paradigm

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# LISP (Scheme)

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# Scheme as a new LISP-like language

The language was conceived by **John McCarthy**

(Based on his paper "*Recursive Functions of Symbolic Expressions and Their Computation by Machine*" in 1960).

In the mid 70's **Sussman and Steele** (MIT ie [Massachusetts Institute of Technology](#), USA) defined ***Scheme*** as a new LISP-like Language

**Scheme** has mostly been used as a language for teaching Computer programming concepts where as **Common** Lisp is widely used as a practical language

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# The `define` Function

- ❑ Used to define global “variables”;  
e.g.

```
(define f 120)
```

- `define` changes its environment but is not equivalent to an assignment statement. It just gives a name to a value.
- `define` can also be used for other purposes, as we will see.

- ❑ *`setQ` and `setF` are functions that operate more like assignments; also `set!`*

```
(define X 2)
```

```
X
```

```
2
```

```
(* 5 X)
```

```
10
```

```
(define pi 3.14159)
```

```
(define radius 10)
```

```
(* pi (* radius radius))
```

```
314.159
```

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# Expressions

Cambridge prefix notation for *all* Scheme expressions:

`(f x1 x2 ... xn)`

e.g.,

`(+ 2 2) ; evaluates to 4`

`(+ (* 5 4 2) (- 6 2)) ; means 5*4*2+(6-2)`

*Note:* Scheme comments begin with `;`

Cambridge prefix allows operators to have an arbitrary number of arguments.

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# Expression Evaluation

Three steps:

1. replace names of symbols by their current bindings.

2. Evaluate lists as function calls in Cambridge prefix, where a list is a set of elements enclosed in ( );

e.g., (\* a 2)

The first element in the list is always treated as the function unless you specifically say not to.

3. Constants evaluate to themselves.

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# Expressions

cont...

- ❑ An expression can take arbitrary number of arguments

```
(+ 21 35 12 7)  
75
```

```
(* 25 4 12)  
1200
```

- ❑ Nested Expression

```
(+ (* 3 5) (- 10 6))  
19
```

- ❑ There is no limit (in principle) to the depth of such nesting

```
(+ (* 3 (+ (* 2 4) (+ 3 5))) (+ (- 10 7) 6))  
57
```

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# Lists as Function Calls

`(+ 5 2 9 13)` ; evaluates to 29

*but*

`(+ (5 2 9 13))` ; *an error*

Scheme will try to evaluate the second list, interpreting “5” as a function

`(f)` ; **error - f isn't a function**

## Preventing Evaluation

`(define colors (quote (red yellow green)))`

or

`(define colors (` (red yellow green)))`

Quoting tells Scheme/LISP that the following list is not to be evaluated.

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# Quoted Lists

```
(define x f)
```

```
  ; defines x as 120 (value of f)
```

```
(define x `f)
```

```
  ; defines x as the symbol f
```

```
(define color `red)
```

```
  ; color is defined to be red
```

```
(define color red)
```

```
  ; error: no definition of red
```

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# Defining Functions

- ❑ `define` is also used to define functions; according to the following syntax:

*(define name (lambda (arguments) body))*

*or*

*(define (name arguments) body)*

- From the former, one can see that Scheme is an applied lambda calculus.

## Example

*(define (min x y) (if (< x y) x y))*

to return the minimum of x and y

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# Function

The general form of a function definition is

(define (<name> <formal parameters>) <body>)

```
(define (square x) (* x x))
```

Use the procedure *square*

```
(square 21)
```

441

```
(square (+ 2 5))
```

49

```
(square (square 3))
```

81

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# Function

cont...

use **square** as a building block in defining other procedures

```
(define (sum-of-squares x y) (+ (square x) (square y)))
```

```
(sum-of-squares 3 4)
```

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Now we can use **sum-of-squares** as a building block in constructing further procedures:

```
(define (plus1 x) (+ x 1))
```

```
(define (multby2 x) (* x 2))
```

```
(define (f a) (sum-of-squares (plus1 a) (multby2 a)))
```

or

```
(define (f a) (sum-of-squares (+ a 1) (* a 2)))
```

```
(f 5)
```

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# The Substitution Model for Function Application

```
(define (square x) (* x x))
(define (sum_of_squares x y) (+ (square x) (square y)))
(define (f a) (sum_of_squares (+ a 1) (* a 2)))
(f 5)
```

Trace the evaluation

```
(f 5)
```

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```
(f 5)
(sum_of_square (+ 5 1) (* 5 2))
(+ (square (+ 5 1)) (square (* 5 2)))
(+ (* (+ 5 1) (+ 5 1)) (* (+ 5 2) (+ 5 2)) )
```

.

.

.

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# The Substitution Model for Function Application

## Order of Evaluation

Applicative

(Evaluate the arguments  
and then apply)

```
(f 5)
(sum_of_square (+ 5 1) (* 5 2))
(+ (square 6) (square 10))
(+ (* 6 6) (* 10 10))
(+ 36 100)
136
```

Normal

(Fully expand and then  
evaluate)

```
(f 5)
(sum_of_square (+ 5 1) (* 5 2))
(+ (square (+ 5 1)) (square (* 5 2)))
(+ (* (+ 5 1) (+ 5 1)) (* (+ 5 2) (+ 5 2)))
.
.
136
```

# Control Flow / Conditional statement

## Conditional Statement

➤ `(cond (<p1> <e1>) (<p2> <e2>) . . . . . (<pn> <en>))`

➤ `(cond (<p1> <e1>) (<p2> <e2>) . . . . . (else <en>))`

➤ `(if <p> <e1> <e2>)`

## Consider a function

$$\begin{aligned} |x| &= 0 && \text{if } x=0 \\ &= x && \text{if } x>0 \\ &= -x && \text{if } x<0 \end{aligned}$$

## In Scheme:

➤ `(define (abs x) (cond ((> x 0) x) ((= x 0) 0) ((< x 0) (* -1 x))))`

➤ `(define (abs x) (cond ((< x 0) -x)) (else x))`

➤ `(define (abs x) (if (< x 0) (- x) x))`

# Conditional statement

## An example

```
(define (weather f) (cond ((> f 38) 'too-hot)
                           ((> f 25) 'nice)
                           ((< f 10) 'too-cold)
                           (else 'moderate)))
```

```
(weather 10)
```

```
(weather 23)
```

```
(weather 41)
```



# Logical Operators

```
LOGICAL AND :      (and <e1> ... <en>)  
LOGICAL OR  :      (or <e1> ... <en>)  
LOGICAL NOT :      (not <e>)
```

Consider an expression :

```
(and (> x 5) (< x 10))
```

```
(define (myfun a b) (if (and (> b a) (< b (* a b))) b a))
```

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# Switch/Case

The **case** statement is similar to a Java or C++ **switch** statement:

```
(case month
  (( sep apr jun nov) 30)
  ((feb) 28)
  (else 31)    ; optional
)
```

All cases take an unquoted list of constants, except for the `else`.

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# Recursion

```
Factorial(n) = 1           if n=0  
             = n*Factorial(n-1) if n>0
```

```
(define (factorial n)(if (= n 0)1(* n (factorial (- n 1)))))
```

```
Fibonacci(n) = 0           if n=1  
              = 2           if n=2  
              = Fibonacci(n-1)+Fibonacci(n-2)   if n>2
```

```
(define (fib n)(cond ((= n 1) 0)  
                    ((= n 2) 1)  
                    (else (+ (fib (- n 1))(fib (- n 2))))))
```

# Arithmetic expression

- **(exp x)** - which returns the value of  $e^x$
- **(log x)** - which returns the value of the natural logarithm of  $x$
- **(sin x)** - which returns the value of the sine of  $x$       **(cos x)(tan x)**
- **(sqrt x)** - which returns the principle square root of  $x$
- **(max  $x_1 x_2 \dots$ )** - which returns the largest number from the list of given numbers      **(min  $x_1 x_2 \dots$ )**
- **(quotient  $x_1 x_2$ )** - which returns the quotient of  $x_1 x_2$ .
- **(remainder  $x_1 x_2$ )** - which returns the integer remainder of  $x_1 x_2$
- **(modulo  $x_1 x_2$ )** - returns  $x_1$  modulo  $x_2$
- **(gcd num1 num2 ...)** - which returns the greatest common divider from the list of given numbers
- **(lcm num1 num2 ...)** - which returns the least common multiple from the list of given numbers
- **(expt base power)** - which returns the value of base raised to power

# Lists

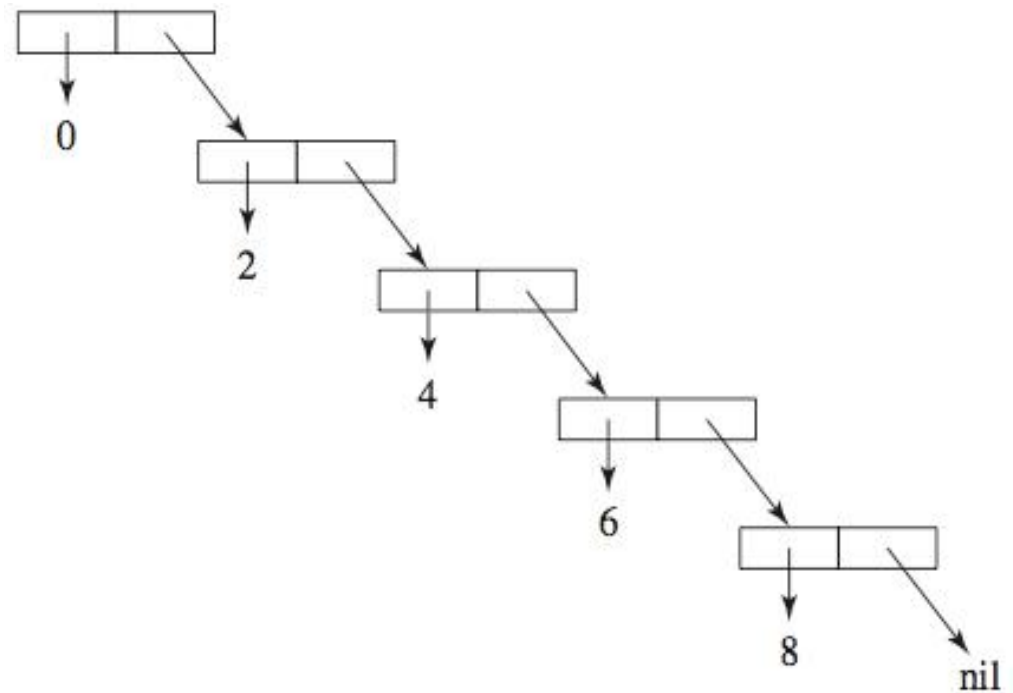
- ❑ A *list* is a series of expressions enclosed in parentheses.

Lists represent both functions and data.

The empty list is written `()`.

e.g., `(0 2 4 6 8)` is a list of even numbers.

Here's how it's stored:



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# List Node Structure

- ❑ Each list node is a record with two fields: the *car* and the *cdr*: *car* is (a pointer to) the first field, *cdr* is (a pointer to) the second.
  - ❑ Note that in the previous example the *cdr* of the last node ( |8|nil| ) is *nil*.
    - This is equivalent to the null pointer in C/C++
  - ❑ The nil value can be represented as ( ), which is also the representation for an empty list.
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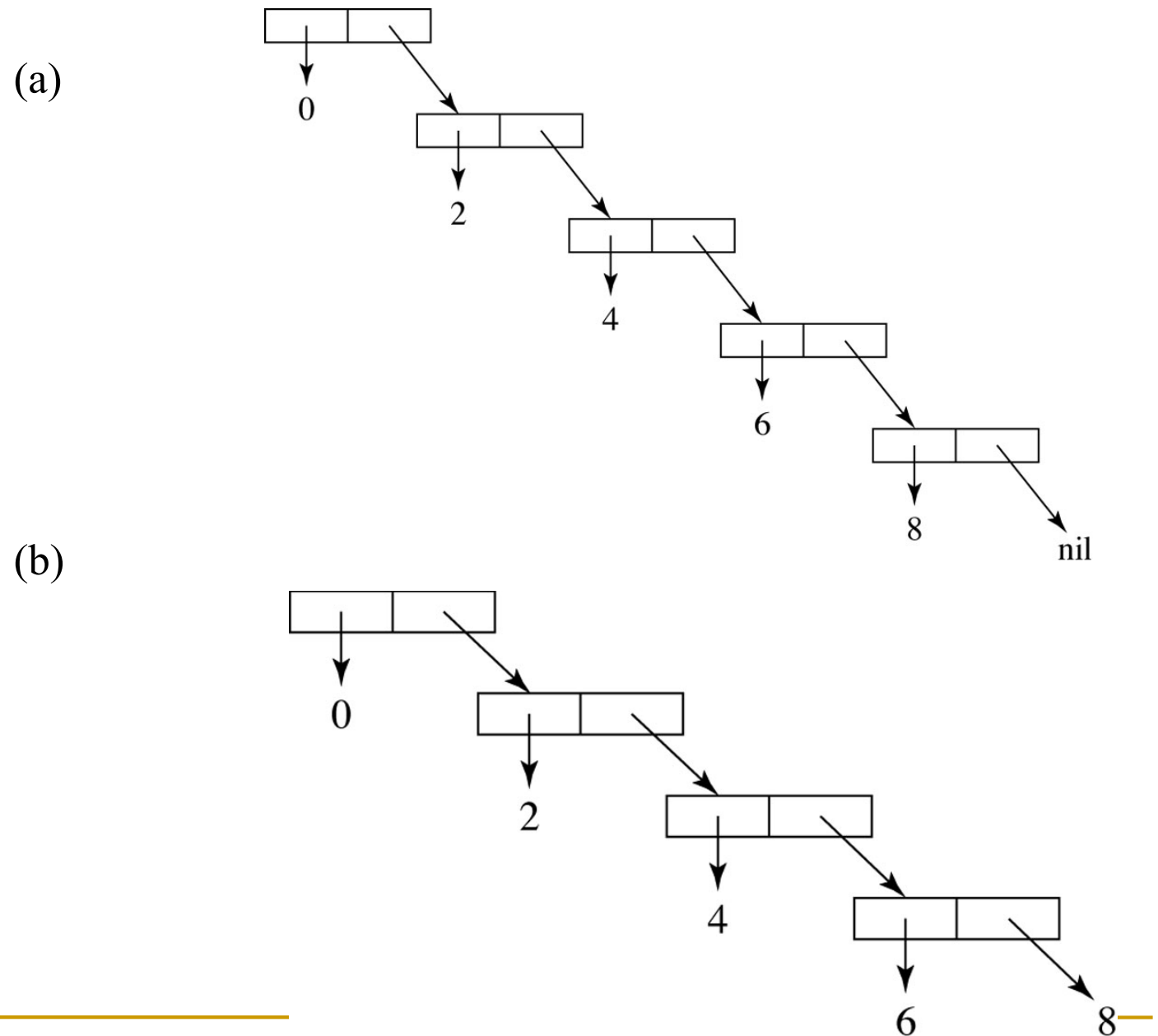
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# “Proper” Lists & Dotted Lists

- ❑ Proper lists are assumed to end with the value ( ) which is implemented by null reference
  - ❑ In a dotted list the last cons has some value other than nil as the value of the cdr field.
    - “Dotted” lists are written (0 2 4 6 . 8) The last node in the previous list would have been |6|8|
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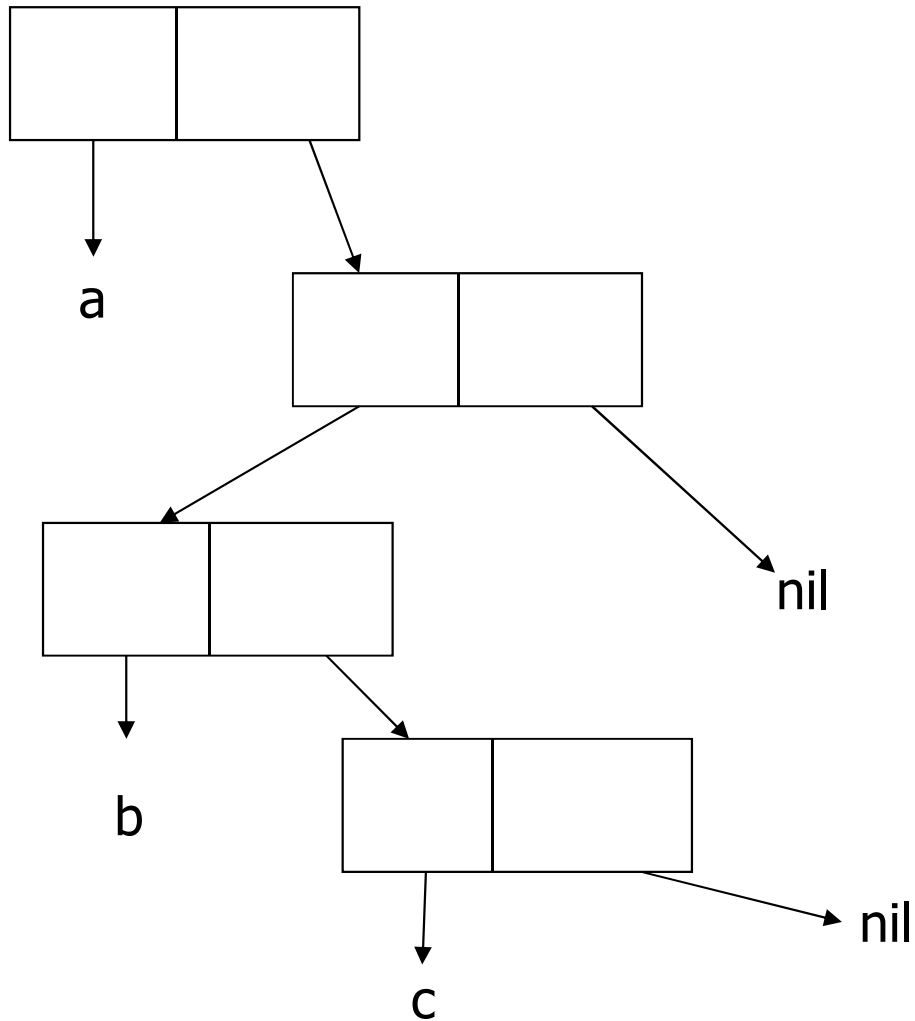
# Structure of a List in Scheme

Notice the difference between the list with nil as the last entry, and the "dotted" list shown in part b.





# List Elements Can Be Lists Themselves



This represents  
the list (a (b c))

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# List Functions

*car:* returns the first element ("head") of a list

*cdr:* returns the tail of the list, which is itself a list

*cons:* used to build lists

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# List Transforming Functions

Suppose we write

```
(define evens '(0 2 4 6 8)).
```

Then:

```
(car evens)      ; gives 0
(cdr evens)      ; gives (2 4 6 8)
( (null? '())    ; gives #t, or true
(cdr(cdr evens)); (4 6 8)
(car `(6 8))    ; 6 (quoted to stop eval)
```

The *cons* requires two arguments: an element and a list; e.g.,

```
(cons 8 ( ))      ; gives the 1-element list (8)
(cons 6 (cons 8 ( ))) ; gives the list (6 8)
(cons 6 `(8))     ; also gives the list (6 8)
(cons 4 (cons 8 9)) ; gives the dotted list
                    ; (4 8 . 9) since 9 is not a list
```