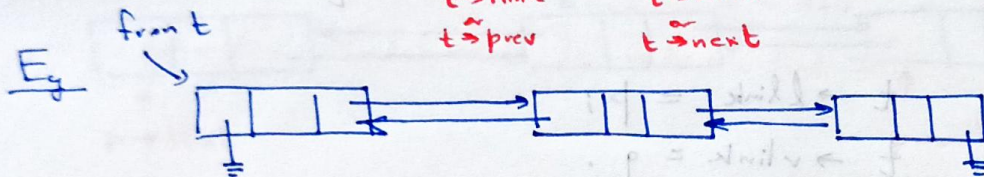
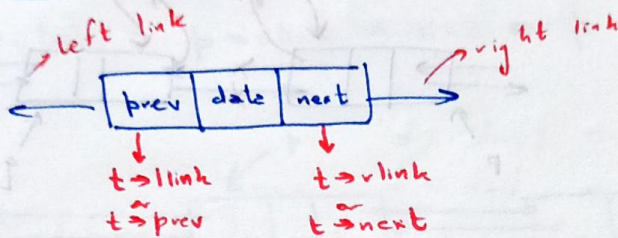


# Double Linked List

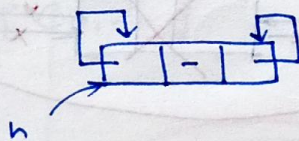


## Circular Double Linked List

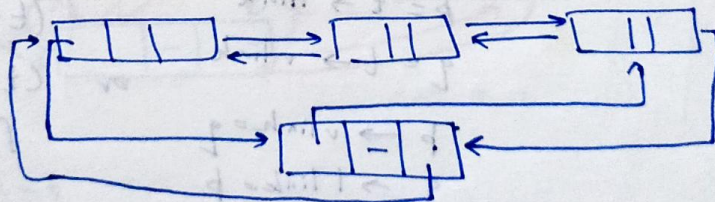
→ we use another node, which is called "head node", it's data is "don't care"



Empty list :



not empty list :

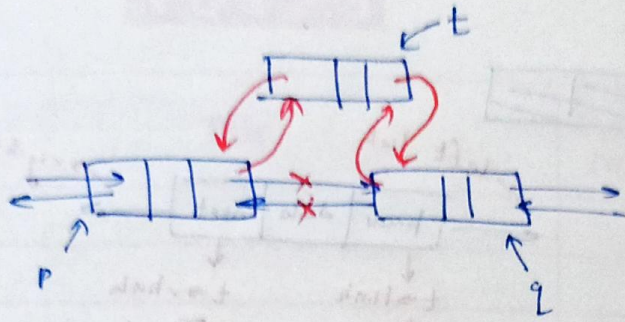


→ head node are used to simplify algorithms



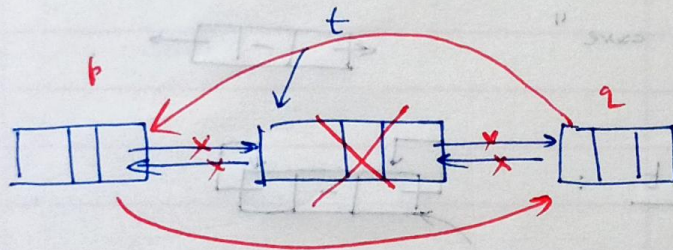
## Insertion

Eg



$t \rightarrow llink = p;$   
 $t \rightarrow rlink = q;$   
 $p \rightarrow rlink = t;$   
 $q \rightarrow llink = t;$

## Delete



$p = t \rightarrow llink$   
 $q = t \rightarrow rlink$

$(t \rightarrow llink) \rightarrow rlink = (t \rightarrow rlink)$   
 $(t \rightarrow rlink) \rightarrow llink = t \rightarrow llink$   
 $free(t)$

$free(t)$

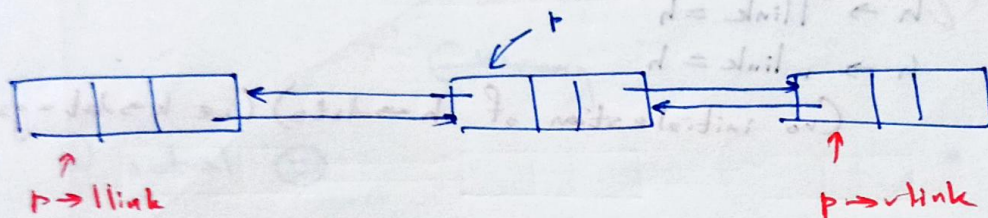
while deleting, it may be case that the remaining node ~~is already~~ becomes empty, i.e. only head node remains., you can't delete head node (use if to check first)



## properties

for a node  $p$ :

$$p \rightarrow llink \rightarrow rlink = p \rightarrow rlink \rightarrow llink = p$$

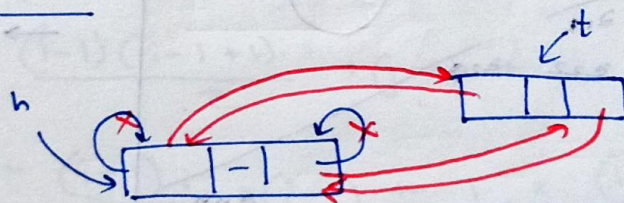


## Summary

→ a node should have three things.

- i) llink
- ii) rlink
- iii) Data

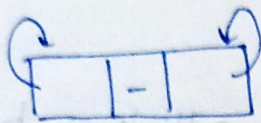
## Insertion in empty list [Circular]



$$\begin{aligned} t \rightarrow llink &= h \\ t \rightarrow rlink &= h \\ h \rightarrow rlink &= t \\ h \rightarrow llink &= t \end{aligned}$$



## initial config of double linked list



$h \rightarrow \text{link} = h$

$h \rightarrow \text{rlink} = h$

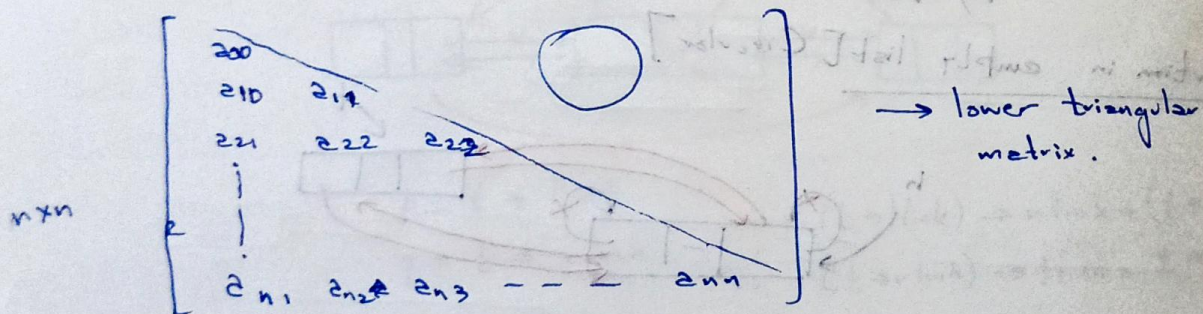
(No initialization of  $h \rightarrow \text{data}$ ) (i.e.  $h \rightarrow \text{data} = \text{garbage}$ )

## Sparse Matrix in Linked Form

Revision.  
Reviews of Sparse Matrix representation in array form

I)

If we have a matrix like this.



we can simplify the algo for these special matrix.

→ no. of elements in the array =  $1 + 2 + 3 + \dots + n = \frac{n(n+1)}{2}$

→ we store element in format

$a_{11}, a_{12}, a_{13}, \dots, a_{1n}, a_{21}, a_{22}, a_{23}, \dots, a_{2n}, \dots, a_{n1}, a_{n2}, \dots, a_{nn}$

so we don't even need position array.



→ finding pos. Q)  $B[L] = a_{ij}$ , find L

i) possible ~~algo~~:

$$L = i(i+1+2+3+\dots+i) + j$$

$$= \frac{i(i+1)}{2} + j \quad (\text{wrong})$$

Wrong

(i-1) wrong due to zero index matter

ii) actual ☺

$$B[1] = a_{11}$$

$$B[2] = a_{12}$$

$$B[3] = a_{22}$$

$$B[4] = a_{31}$$

→ assuming array is 1 indexed

stop of  $a_{ij}$

→ left of  $a_{ij}$

$$L = (1+2+\dots+(i-1)) + (j-1) + 1$$

index of actual  $a_{ij}$

$$= \frac{(i-1)(i-1+1)}{2} + (j-1) + 1$$

$$L = \frac{(i-1)i}{2} + (j-1) + 1 = \frac{(i-1)i}{2} + j$$

address  
index of

$$a_{ij} = \alpha + \left[ \frac{i(i-1)}{2} + j - 1 \right] \times W$$

size of element

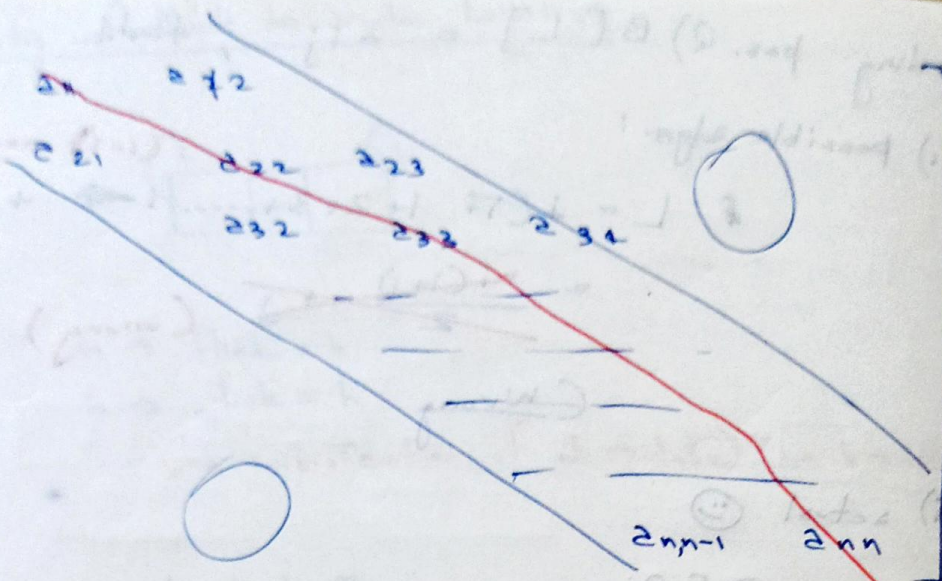
index of  $a_{ij}$

$(i-1) \times i = i^2 - i$   
 $1 + i - i = 1$  for the whole formula



II)

$n \times n$



→ main diagonal, its upward and lower diagonal are non zero  
all other zero

→ called Tridiagonal matrix.

$$B[1] = a_{11}$$

$$B[2] = a_{12}$$

$$B[3] = a_{21} + (1-i)(1-i) =$$

⋮

$$a) B[L] = a_{ij} \quad L = ?$$

$$L = \left[ \overset{\text{rows}}{2 + \underset{\substack{\uparrow \\ \text{first row}}}{3(i-2)}} \right] + \overset{\text{columns}}{\underset{\substack{\downarrow \\ \text{next rows}}}{(j-1) \times 3 + 1}} \overset{\text{itself}}{+ 1}$$

$$\text{element above } a_{ij} = 2 + 3(i-2)$$

$$\text{element } \text{left of } a_{ij} = j - i + 1$$



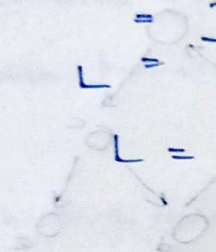
∴ index of  $a_{ij}$  in B

$$L = [3(i-2) + 2] + [(j-i+1)] + 1$$

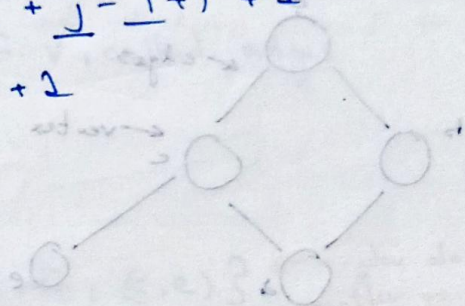
$$= 3i - 6 + 2 + j - i + 1 + 1$$

$$L = 2i + j - 3 + 1$$

$$L = 2i + j - 2$$



digraph



digraph

undirected graph  $n_i \leftarrow$

$$a_i = d_i$$

directed graph  $n_i \leftarrow$

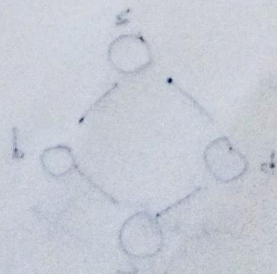
$$a_i \neq d_i$$

Connected graph

→ any vertex will have at least one path  
→ a path between every pair of vertices exist

Disconnected graph

→ There exist a vertex which is not connected to any other vertices



graph