

Rules of Inference - 2

Monday, September 7, 2020 10:37 AM

p: "You send me an email message" —

q: "I will finish writing the program" —

r: "I will go to sleep early" —

s: "I will wake up feeling refreshed" —

Premises.

$$\textcircled{1} \quad p \rightarrow q$$

$$\textcircled{3} \quad r \rightarrow s$$

$$\textcircled{2} \quad \neg p \rightarrow r$$

Conclusion :

$$\neg q \rightarrow s$$

Steps

$$1. \quad p \rightarrow q$$

$$2. \quad \neg q \rightarrow \neg p$$

$$3. \quad \neg p \rightarrow r$$

$$4. \quad \neg q \rightarrow r$$

$$5. \quad r \rightarrow s$$

$$6. \quad \neg q \rightarrow s$$

Reason

Premise

Contrapositive of (1).

Premise

Hypothetical Syllogism on (2) & (3)

Premise

Hypothetical Syllogism on (4) & (5)

Prob-4

$\neg R_K :=$ I was reading newspaper in the kitchen.

$G_K :=$ My glasses are on the kitchen table.

$S_B :=$ I saw them at breakfast.

$R_L :=$ I was reading the newspaper in the living room.

$G_C :=$ My glasses are on the coffee table.

Steps

1. $\neg R_K \rightarrow G_K$
2. $G_K \rightarrow S_B$
3. $\neg R_K \rightarrow S_B$
4. $\neg S_B$
5. $\neg \neg R_K$
6. $R_L \vee \neg R_K$
7. R_L
8. $R_L \rightarrow G_C$
9. G_C

Reasons

Premise (a)

Premise (b)

Hypothetical syllogism on
(1) and (2)

Premise (c)

Modus tollens on (3) &
(4)

Premise (d)

Disjunctive syllogism on
(5) & (6)

Premise (e)

Modus ponens on (7) &
(8)

\rightarrow My glasses are on the coffee table.

Prob-5

$p :=$ "It rained"

$q :=$ "It is foggy"

$p :=$ "It rained"
 $f :=$ "It is foggy"
 $s :=$ "The Sailing race will be held"
 $d :=$ "The life saving demo will go on"
 $t :=$ "The trophy will be awarded".

Premises:

- ① $(\neg r \vee \neg f) \rightarrow (s \wedge d)$
- ② $\neg s \rightarrow t$
- ③ $\neg t$

Conclusion: r

Steps

- ① $(\neg r \vee \neg f) \rightarrow (s \wedge d) =$
- ② $\neg s \rightarrow t$
- ③ $\neg t$
- ④ $\neg s$
- ⑤ $\neg(s \wedge d) \rightarrow \neg(\neg r \vee \neg f)$

$$\underline{\neg s \vee \neg d} \rightarrow r \wedge f$$

$\therefore r$

Reasons.

Premise

Premise

Premise

$$p \rightarrow q \equiv \neg p \rightarrow \neg q$$

Modus tollens on ② & ⑤

Equivivalence $p \rightarrow q \equiv \neg p \rightarrow \neg q$
on ①

Applying De Morgan's
laws on ⑤

Addition on ④

-
- (7) $\neg S \vee \neg d$
 (8) $r \wedge f$
 (9) $\neg r$

laws $\neg\neg$
 Addition on (4)
 Modus ponens on (6) &
 (7)
 Simplification on (8)

Prob6
 $D(x) := x \text{ is in this discrete mathematics class.}$
 $C(x) := x \text{ has taken a course in computer sc.}$

Premises:

- 1 $\forall x (D(x) \rightarrow C(x))$
2 $D(\text{Mark})$

Conclusion:

$$\underline{C(\text{Mark})}$$

Steps

- (1) $\forall x (D(x) \rightarrow C(x))$
 (2) $D(\text{Mark}) \rightarrow C(\text{Mark})$
 (3) $D(\text{Mark})$
 (4) $C(\text{Mark})$

Reason

Premise

Universal instantiation
on (1)

Premise

Modus ponens on (2) & (3).

④ $C(x)$

Modus ponens on (2) & (3).

Prob-7
Premises:
 $C(x) = "x \text{ is } \in \text{ in this class}"$

$B(x) = "x \text{ has read the book}"$

$P(x) = "x \text{ has passed the first exam}"$

Premises:
① $\exists x (C(x) \wedge \neg B(x))$

② $\forall x (C(x) \rightarrow P(x))$

Conclusion: $\exists x (P(x) \wedge \neg B(x)) \models$

$P(x)$ Δ $\neg B(x)$

Steps

1. $\exists x (C(x) \wedge \neg B(x))$

2. $C(a) \wedge \neg B(a)$

3. $C(a)$

4. $\forall x (C(x) \rightarrow P(x))$

5. $C(a) \rightarrow P(a)$

6. $P(a)$

Reason

Premise

Existential instantiation
on (1)

Simplification on (2)

Premise

Universal instantiation
on (4)

Modus ponens on (3) &
(5)

6. $P(a)$, Modus ponens on (3)
 (5)
 7. $\neg B(a)$ Simplification on (2)
 8. $P(a) \wedge \neg B(a)$ Conjunction on (6)
 (7)
 9. $\exists a (P(a) \wedge \neg B(a))$ Existential Generalization
 on (8)

~~Combining Rules of Inference for Propositions and Quantified Statements.~~

① Universal Modus tollens: Combines modus tollens and universal instantiation.

$$\forall a (P(a) \rightarrow Q(a))$$

$\neg Q(a)$, where a is a particular element in the domain

$$\therefore \neg P(a)$$

② Universal Modus ponens: Combines modus ponens and universal instantiation.

$$\forall a (P(a) \rightarrow Q(a))$$

$P(a)$, where a is a particular element in the domain

$$\therefore Q(a)$$