Surname

Student No.

08 - 09 - 2Course Name Principles of Compiling Examination Term Score Software Examination Form Close test Test Duration 150 Mins Related Major Engineering

There are 9 problems in this paper. You can write the answers in English or Chinese on the attached paper sheets.

- 1. Please construct context-free grammars with ε-free productions for the following language (10%).
  - $\{\omega | \omega \in (a,b,c,d)^* \text{ and the numbers of a's and b's and c's occurred in } \}$  $\omega$  are even, and  $\omega$  starts with b or c, ends with a}
- 2. Please construct a DFA with minimum states for the following regular expression. (10%)  $(a|(a|(a|b^*))^*)^*(a|b^*)$
- 3. Please eliminate the left recursions (if there are) and extract maximum common left factors (if there are) from the following context free grammar, and then decide the resulted grammar is whether a LL(1) grammar by constructing the related LL(1) parsing table.(15%)

 $P \rightarrow b S d$ 

$$S \rightarrow S ; A|A$$

$$A \rightarrow B|C$$

$$B \rightarrow a$$

$$C \rightarrow D|D e A$$

$$D \rightarrow E B$$

$$E \rightarrow i F t$$

$$F \rightarrow F \circ G|G$$

$$G \rightarrow b$$

4. Please show that the following operator grammar is whether an operator precedence grammar by **constructing the related parsing table.** (10%)

$$E \rightarrow E \ a \ F|F$$

$$F \rightarrow F \circ T \mid T$$

$$T\rightarrow (E)|n T|b$$

5. Please construct a LR(1) parsing table for the following ambiguous grammar with the additional conditions that \*,  $\otimes$  and  $\oplus$  have the properties of left associative law, and \* has higher precedence than  $\otimes$ ,  $\otimes$  has higher precedence than

$$\oplus$$
.(15%)

## $E \rightarrow E \oplus E | E \otimes E | E^* | (E) | a | b$

Please construct an annotated parse tree for the input string 123.123 where the syntax-directed definition is as following (10%):

**Productions Semantic Rules**  $S \rightarrow L^{(1)}.L^{(2)}$  $S.val = L^{(1)}.val + L^{(2)}.val / 4^{L(2).len}$  $S \rightarrow L$ S.val=L.val  $L \rightarrow L^{(1)}B$ L.val= $L^{(1)}$ .val\*4+B.val, L.len= $L^{(1)}$ .len+1 L.val=B.val, L.len=1  $L\rightarrow B$  $B\rightarrow 0$ B.val=0  $B\rightarrow 1$ B.val=1 $B\rightarrow 2$ B.val=2 $B\rightarrow 3$ B.val=3

7. We assume that the storage organization and the form of activation record used in C language program run-time stack storage allocation are as following. Please construct the run-time stack map when it gets the maximum size at the third time for the following C program (10%).

共 6 页 第 3 页

## Storage Organization of C Language

```
Activation Record of the function called by Main function
Activation Record of Main function
Global Variable Data Area
```

The Form of Activation Record of any a function in C

Unit for Returned Value of function

Internal Temporary Work Units

Local Variable Data Storage Units

Formal parameter Storage Units

Number of Formal Parameters

Returned Address

Caller's SP (Start Address of caller's activation record)

SP

## The C program is as the following:

```
#include <stdio.h>
int x,y;
int main()
   x=6;
   y=f(x);
 }
 int f(int n)
   if (n<=1)
        return 1;
   else if(n==2)
         return 2;
      else
         int t1,t2,t3,t;
         t1=f(n-1);
         t2=f(n-2);
                                 共 6 页
                                               第 4 页
```

```
t3=f(n-3);
         t=t1+t2;
         t=t+t3;
         return t
       }
}
```

**Notes:** 1) Here we assume that the caller's sp of Main function is the start address of global variable data area, and the returned address in the activation record of a function (including Main function) is filled by the operating system automatically, you might not care it.

- 2) The initial value of variable X is 6, the start address of stack used in the program is K.
- 3) The stack map may get its maximum size for several times, here we ask you draw the stack map at maximum size for the third time.
- 8. Please translate the following program fragment into quadruple sequence. (10%)

```
i=1:
m=0;
loop=0;
n=0;
while (loop==0 && i<=10) {
  i=1;
  while (loop ==0 && j <= i)
    if (a[i,j] != a[j,i])
       {
          loop=1;
          m=i;
          n=j;
     else j=j+1;
  if (loop==0) i=i+1;
                                       第 5 页
```

Notes: Here we assume that the declaration of array A is array [1..10,1..10], each data element of array A would **use 4 storage unit,** and the start address of array A's storage area is addrA.

9. Please **construct the DAG** for the following basic block. We assume that only **variable P be used later**, please optimize the block and **rewrite the block** in optimized code form.(10%)

E=A+B

F=E-C

G=F\*D

H=A+B

I=H-C

L=I+G

M=I\*I

M = 2\*M

N=L+M

P=N+M