Module 7: Processing two lists or numbers

Readings: HtDP, section 17.

Processing two lists simultaneously

We now move to using recursion in a complicated situation, namely writing functions that consume two lists (or two data types, each of which has a recursive definition).

Following Section 17 of the textbook, we will distinguish three different cases, and look at them in order of complexity.

One goal of this module is to learn how to choose among cases.

Case 1: a list going along for the ride

my-append consumes two lists list1 and list2, and produces the list with empty in list1 replaced by list2.

The function append is a built-in function in Scheme that consumes two or more lists. Be careful not to use (append (list 3) my-list) instead of (cons 3 my-list). Use append only when the first list has length greater than one or when there are more than two lists to join together.

Do we need to process the individual elements of list1?

What about the individual elements of list2?

```
(define (my-list-fun alist)
  (cond
    [(empty? alist) . . . ]
    [else ... (first alist) ...
           ... (my-list-fun (rest alist)) ...]))
(define (my-alongforride-fun list1 list2)
  (cond
    [(empty? list1) . . . ]
    [else ... (first list1) ...
           ... (my-alongforride-fun (rest list1) list2) ...]))
```

The function my-append

```
;; my-append: (listof any) (listof any) \rightarrow (listof any)
;; Produces list formed by replacing empty in list1 with list2.
;; Examples: (my-append empty (list 1 2)) \Rightarrow (list 1 2)
;; (my-append (cons 1 empty) (cons 2 empty)) \Rightarrow
;; (cons 1 (cons 2 empty))
(define (my-append list1 list2)
  (cond
    [(empty? list1) list2]
    [else (cons (first list1) (my-append (rest list1) list2))]))
```

Condensed trace of my-append

```
(\mathsf{my}\text{-append }(\mathsf{cons}\ 1\ (\mathsf{cons}\ 2\ \mathsf{empty})))\\ \Rightarrow (\mathsf{cons}\ 1\ (\mathsf{my}\text{-append }(\mathsf{cons}\ 2\ \mathsf{empty}))\\ \qquad \qquad (\mathsf{cons}\ 3\ (\mathsf{cons}\ 4\ \mathsf{empty}))))\\ \Rightarrow (\mathsf{cons}\ 1\ (\mathsf{cons}\ 2\ (\mathsf{my}\text{-append }\mathsf{empty}\\ \qquad \qquad \qquad (\mathsf{cons}\ 3\ (\mathsf{cons}\ 4\ \mathsf{empty})))))\\ \Rightarrow (\mathsf{cons}\ 1\ (\mathsf{cons}\ 2\ (\mathsf{cons}\ 3\ (\mathsf{cons}\ 4\ \mathsf{empty})))))
```

Case 2: processing in lockstep

total-value determines the total value of items sold, given prices of items and numbers of each item.

Example: the total of prices $(1\ 2\ 3)$ and numbers $(4\ 5\ 6)$ is $1\cdot 4 + 2\cdot 5 + 3\cdot 6 = 4 + 10 + 18 = 32$.

We can store the items (and numbers) in lists, so $(1\ 2\ 3)$ becomes (list 1 2 3).

Key observation: the lists have the same length and items in a given position correspond to each other.

The function total-value

```
;; total-value: (listof num[>0]) (listof nat) \rightarrow num[\geq0] ;; Produces the total value of items with prices in pricelist ;; and numbers in numlist, where the list lengths are equal. ;; Examples: (total-value empty empty) \Rightarrow 0 ;; (total-value (list 2) (list 3)) \Rightarrow 6 ;; (total-value (list 2 3) (list 4 5)) \Rightarrow 23 (define (total-value pricelist numlist) . . . )
```

Developing a template

pricelist is either empty or a cons, and the same is true of numlist, yielding four options of empty/non-empty lists.

However, because the two lists must be the same length, (empty? pricelist) is true if and only if (empty? numlist) is true.

Out of the four possibilities, two are invalid for proper data.

The template is thus simpler than in the general case.

Extracting the lockstep template

Completing total-value

Condensed trace of total-value

```
(total-value (cons 2 (cons 3 empty))
              (cons 4 (cons 5 empty)))
\Rightarrow (+ (* 2 4) (total-value (cons 3 empty) (cons 5 empty)))
\Rightarrow (+ 8 (total-value (cons 3 empty) (cons 5 empty)))
\Rightarrow (+ 8 (+ (* 3 5) (total-value empty empty)))
\Rightarrow (+ 8 (+ 15 (total-value empty empty)))
\Rightarrow (+ 8 (+ 15 0))
\Rightarrow (+ 8 15)
\Rightarrow 23
```

Case 3: processing at different rates

If the lists being consumed, list1 and list2, are of different lengths, all four possibilities for their being empty/nonempty are possible:

```
(and (empty? list1) (empty? list2))
(and (empty? list1) (cons? list2))
(and (cons? list1) (empty? list2))
(and (cons? list1) (cons? list2))
```

Exactly one of these is true, but all must be tested in the template.

Example: merging two sorted lists

We wish to design a function merge that consumes two lists, each of distinct elements sorted in ascending order, and produces one such list with all elements in it.

As an example:

```
(merge (list 1 8 10) (list 2 4 6 12)) \Rightarrow (list 1 2 4 6 8 10 12)
```

We need more examples to be confident of how to proceed.

```
(merge\ empty\ empty) \Rightarrow empty
(merge empty
        (cons 2 empty)) \Rightarrow (cons 2 empty)
(merge (cons 1 (cons 3 empty))
        empty) \Rightarrow (cons 1 (cons 3 empty))
(merge (cons 1 (cons 4 empty))
        (cons 2 empty)) \Rightarrow (cons 1 (cons 2 (cons 4 empty)))
(merge (cons 3 (cons 4 empty))
        (cons 2 empty)) \Rightarrow (cons 2 (cons 3 (cons 4 empty)))
```

The outline of merge

```
(define (merge list1 list2)
  (cond
    [(and (empty? list1) (empty? list2)) ....]
    [(and (empty? list1) (cons? list2)) ....]
    [(and (cons? list1) (empty? list2)) ....]
    [(and (cons? list1) (cons? list2)) ....]))
```

Refining the outline

```
(define (merge list1 list2)
  (cond
    [(and (empty? list1) (empty? list2)) ...]
    [(and (empty? list1) (cons? list2)) ... (first list2) ... (rest list2) ...]
    [(and (cons? list1) (empty? list2)) ... (first list1) ... (rest list1) ...]
    [(and (cons? list1) (cons? list2)) ???]))
```

Further refinements

There are many different possible natural recursions for the last cond answer ??? :

```
... (first list2) ... (merge list1 (rest list2)) ...
... (first list1) ... (merge (rest list1) list2) ...
... (first list1) ... (first list2) ... (merge (rest list1) (rest list2)) ...
```

We write all of these down, realizing that not all will be used, and eliminate unnecessary ones in reasoning about any particular function.

```
(define (my-twolist-fun lst1 lst2)
  (cond
    [(and (empty? lst1) (empty? lst2)) ...]
    [(and (empty? lst1) (cons? lst2))
     ... (first lst2) ... (my-twolist-fun empty (rest lst2)) ...]
    [(and (cons? lst1) (empty? lst2))
     ... (first lst1) ... (my-twolist-fun (rest lst1) empty)...]
    [(and (cons? lst1) (cons? lst2))
     ... (first lst2) ... (my-twolist-fun lst1 (rest lst2)) ...
     ... (first lst1) ... (my-twolist-fun (rest lst1) lst2) ...
     ... (first lst1) ... (first lst2)... (my-twolist-fun (rest lst1) (rest lst2)).
```

```
(define (merge list1 list2)
  (cond
    [(and (empty? list1) (empty? list2)) empty]
    [(and (empty? list1) (cons? list2)) list2]
    [(and (cons? list1) (empty? list2)) list1]
    [(and (cons? list1) (cons? list2))
       (cond
          [(< (first list1) (first list2))
              (cons (first list1) (merge (rest list1) list2))]
          [ else
              (cons (first list2) (merge list1 (rest list2)))])))
```

Condensed trace of merge

```
(merge (cons 3 (cons 4 empty))
        (cons 2 (cons 5 (cons 6 empty))))
\Rightarrow (cons 2 (merge (cons 3 (cons 4 empty))
            (cons 5 (cons 6 empty))))
\Rightarrow (cons 2 (cons 3 (merge (cons 4 empty))
                     (cons 5 (cons 6 empty)))))
⇒ (cons 2 (cons 3 (cons 4 (merge empty
                             (cons 5 (cons 6 empty)))))
\Rightarrow (cons 2 (cons 3 (cons 4 (cons 5 (cons 6 empty)))))
```

Consuming a list and a number

In Module 6, we saw how to use structural recursion on natural numbers as well as lists.

We can extend our idea for computing on two lists to computing on a list and a number, or on two numbers.

Example: computing the sublist starting at the nth occurrence of some symbol.

The sublist of (list 'd 'z 'z 'h 'z 'd) starting at the 3rd occurrence of 'z is (list 'z 'd).

The function nth-occur-suffix

```
;; nth-occur-suffix: symbol nat[>=1] (listof symbol) \rightarrow (listof symbol) ;; Produces sublist starting at nth occurrence of sym in alist. ;; Examples: (nth-occur-suffix 'a 1 empty) \Rightarrow empty ;; (nth-occur-suffix 'a 1 (list 'a 'c)) \Rightarrow (list 'a 'c) ;; (nth-occur-suffix 'a 2 empty) \Rightarrow empty ;; (nth-occur-suffix 'a 2 (list 'a 'b 'a 'c)) \Rightarrow (list 'a 'c) (define (nth-occur-suffix sym n alist) ...)
```

Developing the template

The recursion will involve the parameters n and alist, once again giving four possibilities:

```
(and (= n 1) (empty? alist))
(and (= n 1) (cons? alist))
(and (> n 1) (empty? alist))
(and (> n 1) (cons? alist))
```

Once again, exactly one of these four possibilities is true.

```
(define (my-list-and-num-fun n alist)
  (cond
     [(and (= n 1) (empty? alist)) . . . ]
     [(and (= n 1) (cons? alist))]
      ... (first alist) ... (my-list-and-num-fun 1 (rest alist)) ...]
     [(and (> n 1) (empty? alist))]
      ... (my-list-and-num-fun (sub1 n) empty) ... ]
     [(and (> n 1) (cons? alist))]
      ... (first alist) ... (my-list-and-num-fun n (rest alist)) ...
      ... (my-list-and-num-fun (sub1 n) alist) ...
      ... (first alist) ... (my-list-and-num-fun (sub1 n) (rest alist)) ...]))
```

```
(define (nth-occur-suffix sym n alist)
  (cond
    [(and (= n 1) (empty? alist)) empty]
    [(and (= n 1) (cons? alist))]
     (cond [(equal? sym (first alist)) alist]
            [else (nth-occur-suffix sym n (rest alist))])]
    [(and (> n 1) (empty? alist)) empty]
    [(and (> n 1) (cons? alist))]
     (cond [(equal? sym (first alist))
            (nth-occur-suffix sym (sub1 n) (rest alist))]
            [else (nth-occur-suffix sym n (rest alist))])])
```

```
(define (nth-occur-suffix sym n alist)
  (cond [(empty? alist) empty]
        [(= n 1); alist is nonempty]
         (cond [(equal? sym (first alist)) alist]
                [else (nth-occur-suffix sym n (rest alist))])]
         [else; alist is nonempty, n > 1
          (cond [(equal? sym (first alist))
                  (nth-occur-suffix sym (sub1 n) (rest alist))]
                 [else (nth-occur-suffix sym n (rest alist))])])
```

Condensed trace of nth-occur-suffix

```
(nth-occur-suffix 'z 3
             (cons 'd (cons 'z (cons 'z
                    (cons 'h (cons 'z (cons 'd empty))))))
\Rightarrow (nth-occur-suffix 'z 3
          (cons 'z (cons 'z (cons 'h (cons 'z (cons 'd empty)))))
⇒ (nth-occur-suffix 'z 2 (cons 'z (cons 'h (cons 'z (cons 'd empty)))))
⇒ (nth-occur-suffix 'z 1 (cons 'h (cons 'z (cons 'd empty)))))
⇒ (nth-occur-suffix 'z 1 (cons 'z (cons 'd empty)))
\Rightarrow (cons 'z (cons 'd empty))
```

Midpoints of pairs of posns

Suppose we have two lists of posns and wish to find a list of midpoints between each pair of corresponding posns.

Which template should we use?

Testing list equality

We can apply the templates we have created to the question of deciding whether or not two lists of numbers are equal.

```
;; list=?: (listof num) (listof num) \rightarrow boolean
;; Produces true if list1 and list2 are equal, false otherwise.
```

Which template is most appropriate?

Applying the two list template

```
(define (list=? list1 list2)
  (cond
    [(and (empty? list1) (empty? list2)) ...]
    [(and (empty? list1) (cons? list2))
     ... (first list2) ... (list=? empty (rest list2)) ...]
    [(and (cons? list1) (empty? list2))
     ... (first list1) ... (list=? (rest list1) empty) ...]
    [(and (cons? list1) (cons? list2)) ??? ]))
```

Reasoning about list equality

Two empty lists are equal.

If one list is empty and the other is not, the lists are not equal.

If two nonempty lists are equal, then their first elements are equal, and their rests are equal.

The natural recursion in this case is

```
(list=? (rest list1) (rest list2))
```

The function list=?

```
(define (list=? list1 list2)
  (cond
    [(and (empty? list1) (empty? list2)) true]
    [(and (empty? list1) (cons? list2)) false]
    [(and (cons? list1) (empty? list2)) false]
    [(and (cons? list1) (cons? list2))
       (and (= (first list1) (first list2))
             (list=? (rest list1) (rest list2)))]))
```

Another approach to the problem

Another way of viewing the problem comes from the observation that if the lists are equal, they will have the same length.

We can then use the structure of one list in our function, checking that the structure of the other list matches.

This implies the use of the lockstep template.

Here is the result of applying the lockstep template.

Reasoning about list equality again

If the first list is empty, the answer depends on whether or not the second list is empty.

If the first list is not empty, then the following should all be true:

- the second list must be nonempty
- the first elements must be equal
- the rests must be equal

Notice the importance of order of checking.

Built-in list equality

As you know, Scheme provides the predicate equal? which tests structural equivalence. It can compare two atomic values, two structures, or two lists. Each of the nonatomic objects being compared can contain other lists or structures.

At this point, you can see how you might write equal? if it were not already built in. It would involve testing the type of data supplied, and doing the appropriate comparison, recursively if necessary.

Goals of this module

You should understand the three approaches to designing functions that consume two lists (or a list and a number, or two numbers) and know which one is (or ones are) suitable in a given situation.