## Module 2: More Functional Abstraction

#### Topics:

- Making abstract list functions
- •Making single-use, unnamed functions
- •Making functions that produce functions

Readings: HtDP 21, 22.1, 22.2, Intermezzo 4

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## Recall: Implementation of map

# Design Recipe: What is the contract for map?

- What is the type of a function?
  - Its contract
- For example:

```
;; sqr: num -> num[>=0]
So type of sqr is (num->num[>=0])
;; string-length: string -> nat
So type of string-length is (string -> nat)
```

• But, map must work for many types of functions

## Using map

Example	Map consumes
(map sqr 1st)	<pre>(num -&gt; num[&gt;=0]) and (listof num)</pre>
(map not 1st)	<pre>(boolean-&gt;boolean) and (listof boolean)</pre>
<pre>(map string-length los)</pre>	<pre>(string -&gt; nat) and (listof string)</pre>

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## Using variables in contracts

Contract for **map** must show relationships among the parameters:

- Let X be the type of data in the consumed list
- Let Y be the type of data in the produced list

```
;; map: (X -> Y) (listof X)
;; -> (listof Y)

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```

**Exercises** 

Complete the contracts for

- filter
- foldr

## Another abstract list function: build-list

build-list can be used for tasks such as:

- Creating a list with the first 10 even numbers, starting from 0.
- Creating a list of the first 25 perfect squares, starting from 0.
- Creating

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```
(list (make-posn 0 0) (make-posn 1 1) ... (make-posn 999 999))
```

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## Solving without build-list

- First, solve a related, simpler problem:
   create (list 0 1 2 ... N-1) for any nat N
- Recall: N is a natural number if
  - -N=0, or

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- N=K+1, where K is a natural number
- Recall the "count-up" template from CS115

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## Creating a "count-up" list

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# Use **count-up** to produce list of first n evens

- Approach 1:
  - Modify count-up directly to create an even number 2\*i instead of just i

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### A solution for first-n-evens

# Use **count-up** to produce list of first evens

- Approach 2:
  - Use count-up as a helper function and map a function from i to 2\*i.

```
(define (first-n-evens n)
  (local
    [(define (double k) (* 2 k))]
    (map double (count-up n))))
```

## Using build-list directly

```
(define (first-n-evens n)
    (local
        [(define (double k) (* 2 k))]
        (build-list n double)))

Contract:
;; build-list: nat (nat -> X)
;;        -> (listof X)

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```

#### Other abstract list functions

 See section 21.2 (Figure 57) for summaries of Scheme's abstract list processing functions

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# **lambda** and unnamed, single-use functions

Consider the function double in first-n-evens:

- It only exists inside first-n-evens
- It has a name only so that it can be passed to build-list
- → New language feature: lambda
- → New language level: Intermediate
  Student with lambda

#### lambda

- creates a function value
- function can have any number of parameters

```
(lambda (x) (+ 1 (sqr x)))
```

⇒A function value that consumes a single parameter and produces its square plus 1 lts value is the function itself.

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## Using lambda

```
(define (first-n-evens n)
  (build-list n
        (lambda (k) (* k 2))))
```

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## Defining functions with lambda

```
(define double
  (lambda (k) (* k 2)))
```

is equivalent to

```
(define (double k)
   (* k 2))
```

## Evaluating a lambda expression

What is the value of

```
((lambda (x) (* x 2)) 4) ?
```

 Since the first part is a function value, apply function simplification rules

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### Evaluating another **lambda** expression

What is the value of

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#### Our lambda Rule

Use lambda when the function is

- Single use
- Reasonably short (2-3 lines)

Complete **multiples**-of which consumes a nat **n** and a list of integers, and produces a list of only those which are multiples of n.

### Solution

### Format of lambda

```
(lambda
  (list-of-parameters)
  body-of-function)
```

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### Functions as values

We can use function values anywhere:

- As arguments to other functions
- In lists
- As the value produced by another function

## Functions as parameters: another example

```
;; transformed?: (X -> Y) X Y -> boolean
;; Produces true if the result of
;; applying f to s equals t, else false.
;; Examples:
;; (transformed? sqr 5 25) => true
;; (transformed? (lambda (s)
;; substring s 0 1))"ABC" "a") => false
(define (transformed? f s t)
    (equal? (f s) t))
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```

#### Function values in lists

### Functions that make functions

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## Writing Examples and Tests for adder-maker

- adder-maker produces a function
- We can't test two functions for equality
- Instead, we write examples and tests for what the produced function does, e.g. show that (adder-maker 5) is a function that adds 5 to its argument.

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#### More details for adder-maker

### Goals of Module 2

- Understand use of build-list
- Implement functions that consume or produce functions
- Understand lambda and how and when it should be used
- Understand that functions are like other values in Scheme