Write an assembler that reads the source code of an assembly program (for an imaginary machine) and displays the machine language for that program.

The following assembler instructions must be supported.

Instruction	opcode	Length	Format	Example	
add	0	2	0	add R1, R2, R3	R3 = R1 + R2
sub	1	2	0	sub R1, R2, R3	R3 = R1 - R2
mult	2	2	0	mult R1, R2, R3	R3 = R1 * R2
div	3	2	0	div R1, R2, R3	R3 = R1 / R2
load	4	4	1	load R1,addr[R7]	R1 ← memory value
laddr	5	4	1	laddr R1,addr[R7]	R1 ← address
store	6	4	1	store R1,addr[R7]	R1 → memory
call	7	4	2	call addr	push return addr on stack
					and jump to addr
rtn	8	1	4	rtn	pop address from stack
					and jump to there
jump	9	4	2	jump addr[R7]	jump to address
jz	A	4	1	jz R1,addr[R7]	jump to address only if
					R1 is zero
jneg	В	4	1	jn R1,addr[R7]	jump to address only if
					R1 is negative
push	C	1	3	push R1	push R1 on stack
pop	D	1	3	pop R1	pop stack into R1
lpsw	Е	1	3	lpsw R1	load PSW into R1
spsw	F	1	3	spsw R1	store R1 into PSW
data		4	5	data 123	created a data word

Machine language instructions have the following formats:

# 0: arithmetic

opcode	Reg 1	Reg 2	Reg 3
4 bits	4 bits	4 bits	4 bits

# 1: Load, store, conditional jump

opcode	Reg	index	addr
4 bits	4 bits	4 bits	20 bits

2: call, jump

opcode	unused	index	addr
4 bits	4 bits	4 bits	20 bits

3: push, pop, lpsw, spsw

opcode	Reg	
4 bits	4 bits	

#### 4: rtn

opcode	unused	
4 bits	4 bits	

#### 5: data

J. data	
	integer data value
	32 bits

A comment in the assembly program is indicated by an asterisk, "\*", in the first column of the source code. Labels must start in the first column. Your assembler does not have to detect every possible error in the assembler program, but it should produce a reasonable error message if an undefined label is used.

For a maximum score of 95 points, your program can ignore all index registers. The presence of an [index register] may or may not be considered an error. The index field of all instructions must be set to zero if index registers are not implemented. For a maximum score of 100, all memory addressing instructions must support an optional index register.

### **Example output:**

```
address machine
                 source
             * Example assembler program for COMP360
 0
    58000034
                   laddr R8, cat
                   load R1, dog[R8]
 4
    41800030
 8
    42800038
                   load R2, goat[R8]
                         R1, R2, R3
         123 more
                   add
 С
                         R3, R2, R4
        2324
                   mult
 е
10
   b400000c
                   jη
                         R4, more
14
        1415
                         R4, R1, R5
                   sub
16
    a300001c
                   jΖ
                         R3, nodiv
        3536
                   div
                         R5, R3, R6
1a
    66000025 nodiv store R6, addr
1c
20
          С6
                   push
                         R6
21
    7000002d
                   call
                         mthd
25
          fa addr spsw
                         R10
26
                         R10, R2, R7
         a27
                   add
28
                   lpsw
          e7
                         R7
29
    900000c
                   jump
                         more
                example method
2d
         128 mthd
                   add
                         R1, R2, R8
2f
          80
                   rtn
                data
30
           1 dog
                   data 1
34
           0 cat
                   data 0
38
         801 goat data 2049
```

### **Hints:**

You may wish to create an object to hold the symbol table entries. This would contain:

- opcode, address or other value of the symbol
- length of instruction in bytes
- format of the machine instruction

A hash table works well to hold the symbol table. Symbol table entries can be keyed by the mnemonic or address label.

The assembler can be written to use one or two passes. After the first pass, a two pass assembler closes the input file and then reads the input again. The goal of the first pass is to create a symbol table containing the address of all symbols. The second pass creates the output. It uses the symbol table to get the value of all names used in the program.

General outline for pass 1.

- Initialize the current address to zero
- Read a line of assembler
- Split the line into tokens
- If label, put label and current address in symbol table
- Get mnemonic and find in symbol table
- Add instruction length to current address
- repeat until end of file

# General outline for pass 2

- Initialize the current address to zero
- Read a line of assembler
- Split the line into tokens
- Get mnemonic information from symbol table
- Get additional fields and create the machine language
- Display the results

The following Java classes and methods may be useful in writing your assembler. Similar C++ methods exist.

java.util.HashMap java.util.StringTokenizer Integer.toUnsignedString(intNum, 16)