

Worksheet: Complexity Analysis

1

```
def search(L, target):
    for x in L:
        if x == target: #line A
            return True
    return False
```

assume this is asked unless otherwise stated

Let $f(N)$ be the number of times line A executes, with $N = \text{len}(L)$. What is $f(N)$ in each case?

Worst Case (target is at end of list): $f(N) = \underline{\hspace{2cm}}$

Best Case (target is at beginning of list): $f(N) = \underline{\hspace{2cm}}$

Average Case (target in middle of list): $f(N) = \underline{\hspace{2cm}}$

A **step** is any unit of work with bounded execution time (it doesn't keep getting slower with growing input size).

We classify algorithm complexity by classifying the **order of growth** of a function $f(N)$, where f gives the number of steps the algorithm must perform for a given input size.

Big O definition: if $f(N) \leq C * g(N)$ for large N values and some fixed constant C , then $f(N) \in O(g(N))$

2

Let $f(N) = 2N^2 + N + 12$

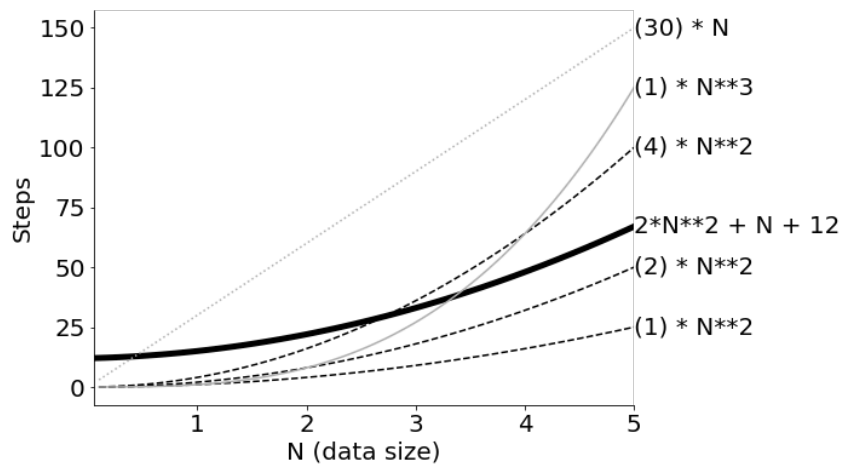
If we want to show $f(N) \in O(N^3)$, what is a good lower bound on N ? Let's have $C=1$.

To show $f(N) \in O(N^2)$, do we pick 1, 2, or 4 for the C ? After picking C , what should we choose for N 's lower bound?

What is more informative to show?

$f(N) \in O(N^3)$ or $f(N) \in O(N^2)$?

Somebody claims $f(N) \in O(N)$, offering $C=30$ and $N>0$. Suggest an N value to counter their claim.



3

```
nums = [...]
first100sum = 0
for x in nums[:100]:
    first100sum += x
print(first100sum)
```

If we increase the size of `nums` from 20 items to 100 items, the code will probably take times longer to run.

If we increase the size of `nums` from 100 to 1000, will the code take longer? Yes / No

The complexity of the code is $O(\underline{\hspace{2cm}})$, with $N = \text{len}(\text{nums})$.

4

Each of the following list operations are either $O(1)$ or $O(N)$, where N is $\text{len}(L)$. Circle those you think are $O(N)$.

`L.insert(0, x)` `L.pop(0)` `x = L[0]` `x = max(L)` `x = len(L)`
`L.append(x)` `L.pop(-1)` `L2.extend(L)` `x = sum(L)` `found = x in L`

5

```
L = [...]
for x in L:
    avg = sum(L) / len(L)
    if x > 2*avg:
        print("outlier", x)
```

What is the big O complexity?

Is there a way to optimize the code?

6

```
A = [...]
B = [...]
```

```
for x in A:
    for y in B:
        print(x*y)
```

how would you define the variable(s) to describe the size of the input data?

The complexity of code is

$O(\rule{1cm}{0.4pt})$

7

assume L is already sorted, $N = \text{len}(L)$

```
def binary_search(L, target):
```

```
    left_idx = 0 # inclusive
```

```
    right_idx = len(L) # exclusive
```

```
    while right_idx - left_idx > 1:
```

```
        mid_idx = (right_idx + left_idx) // 2
```

```
        mid = L[mid_idx]
```

```
        if target >= mid:
```

```
            left_idx = mid_idx
```

```
        else:
```

```
            right_idx = mid_idx
```

```
    return right_idx > left_idx and L[left_idx] == target
```

how many times does this step run
when $N = 1$? $N = 2$? $N = 4$? $N = 8$?

If $f(N)$ is the number of times this step runs, then $f(N) = \rule{1cm}{0.4pt}$

The complexity of binary search is

$O(\rule{1cm}{0.4pt})$

8

```
s1 = tuple("...") # could be any string
```

```
s2 = tuple("...")
```

version A

```
import itertools
```

```
matches = False
```

```
for p in itertools.permutations(s1):
```

```
    if p == s2:
```

```
        matches = True
```

version B

```
s1 = sorted(s1)
```

```
s2 = sorted(s2)
```

```
matches = (s1 == s2)
```

assumed sorted is $O(N \log N)$

what is the complexity of version A? $O(\rule{1cm}{0.4pt})$

what is the complexity of version B? $O(\rule{1cm}{0.4pt})$

9

```
def selection_sort(L):
```

```
    for i in range(len(L)):
```

```
        idx_min = i
```

```
        for j in range(i, len(L)):
```

```
            if L[j] < L[idx_min]:
```

```
                idx_min = j
```

```
        # swap values at i and idx_min
```

```
        L[idx_min], L[i] = L[i], L[idx_min]
```

```
nums = [2, 4, 3, 1]
```

```
selection_sort(nums)
```

```
print(nums)
```

if this runs $f(N)$ times, where $N = \text{len}(L)$,

then $f(N) = \rule{1cm}{0.4pt}$

The complexity of selection sort is

$O(\rule{1cm}{0.4pt})$