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2.1.1. The Abstract Data Type

ADT Array is

objects: A set of pairs $\langle index, value \rangle$ where for each value of *index* there is a value from the set *item*. *Index* is a finite ordered set of one or more dimensions, for example, $\{0, \dots, n-1\}$ for one dimension, $\{(0, 0), (0, 1), (0, 2), (1, 0), (1, 1), (1, 2), (2, 0), (2, 1), (2, 2)\}$ for two dimensions, etc.

functions:

for all $A \in \text{Array}, i \in \text{index}, x \in \text{item}, j, \text{size} \in \text{integer}$

Array Create($j, list$) ::= **return** an array of j dimensions where *list* is a j -tuple whose i th element is the size of the i th dimension. *Items* are undefined.

Item Retrieve(A, i) ::= **if** ($i \in \text{index}$) **return** the item associated with index value i in array A
else return error

Array Store(A, i, x) ::= **if** ($i \in \text{index}$)
return an array that is identical to array A except the new pair $\langle i, x \rangle$ has been inserted **else return** error.

end Array

ADT 2.1: Abstract Data Type *Array*

※ Create (2, (3, 4))
3행 4열의 2차원 배열 생성

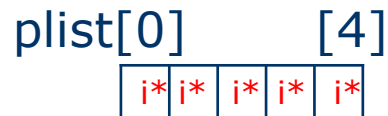
2.1.2 Arrays in C

❖ one-dimensional array

`int list[5];`



`int *plist[5];`



※ i stands for int
i* stands for int pointer

cf) `int (*ary)[5]` , 배열포인터

Variable	Memory address
list[0]	base address = α
list[1]	$\alpha + \text{sizeof}(\text{int})$
list[2]	$\alpha + 2 \cdot \text{sizeof}(\text{int})$
list[3]	$\alpha + 3 \cdot \text{sizeof}(\text{int})$
list[4]	$\alpha + 4 \cdot \text{sizeof}(\text{int})$

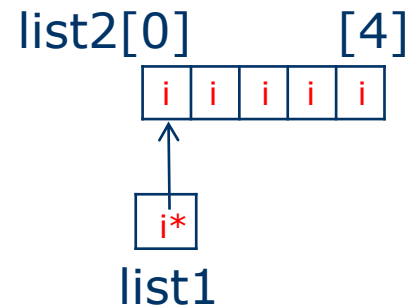
one-dimensional array & pointer

❖ interpretations of pointers: list1, list2

```
int *list1, list2[5];
```

```
list1 = list2; //(variable, constant)
```

```
list1 == &list2[0]  
list1 + i == &list2[i]  
*(list1+i) == list2[i]
```



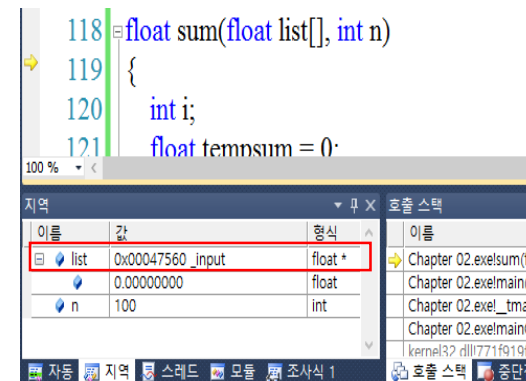
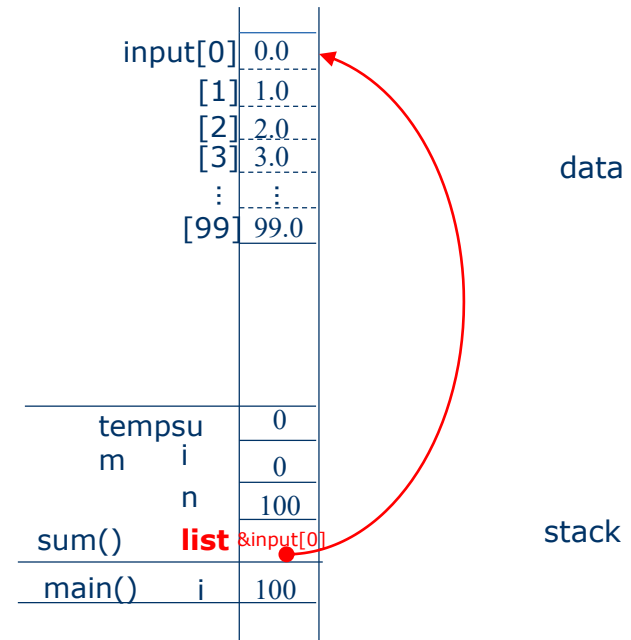
```

#define MAX_SIZE 100
float sum(float [], int);
float input[MAX_SIZE], answer;
void main(void)
{
    int i;
    for (i = 0; i < MAX_SIZE; i++)
        input[i] = i;
    answer = sum(input, MAX_SIZE);
    printf("The sum is: %f\n", answer);
}
float sum(float list[], int n)
{
    int i;
    float tempsum = 0;
    for (i = 0; i < n; i++)
        tempsum += list[i];
    return tempsum;
}

```

array parameter
float * list → pointer parameter

Program 2.1: Example array program



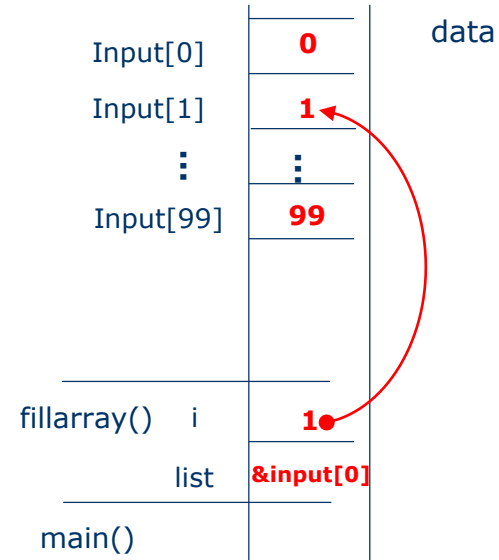
```

1  #include <stdio.h>
2  #define MAX_SIZE 100
3  float sum(float[], int);
4  void fillarray(float[], int);
5  float input[MAX_SIZE], answer;
6  void main(void)
7  {
8      fillarray(input, MAX_SIZE);
9      answer = sum(input, MAX_SIZE);
10     printf("The sum is: %f\n", answer);
11 }
12 void fillarray(float list[], int n)
13 {
14     int i;
15     for (i = 0; i < n; i++)
16         list[i] = i;
17 }
18 float sum(float list[], int n)
19 {
20     int i;
21     float tempsum = 0;
22     for (i = 0; i < n; i++)
23         tempsum += list[i];
24     return tempsum;
25 }

```

■ left-side of equal sign

- the value produced on the right-hand side is stored in the location (*list+i*)



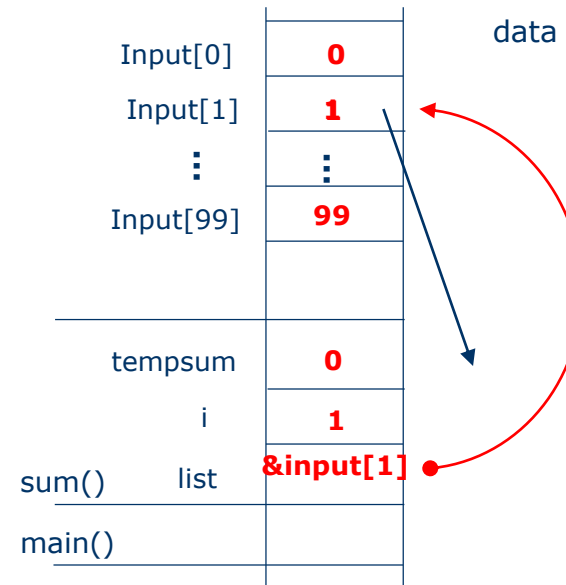
```

1  #include <stdio.h>
2  #define MAX_SIZE 100
3  float sum(float[], int);
4  void fillarray(float[], int);
5  float input[MAX_SIZE], answer;
6  void main(void)
7  {
8      fillarray(input, MAX_SIZE);
9      answer = sum(input, MAX_SIZE);
10     printf("The sum is: %f\n", answer);
11 }
12 void fillarray(float list[], int n)
13 {
14     int i;
15     for (i = 0; i < n; i++)
16         list[i] = i;
17 }
18 float sum(float list[], int n)
19 {
20     int i;
21     float tempsum = 0;
22     for (i = 0; i < n; i++)
23         tempsum += list[i];
24     return tempsum;
25 }

```

■ right-side of
equal sign

- the value pointed at
by $(list+i)$ is returned




```

1  #include <stdio.h>
2  #define MAX_SIZE 100
3  float sum(float[], int);
4  void fillarray(float[], int);
5  float input[MAX_SIZE], answer;
6  void main(void)
7  {
8      fillarray(input, MAX_SIZE);
9      answer = sum(input, MAX_SIZE);
10     printf("The sum is: %f\n", answer);
11 }
12 void fillarray(float list[], int n)
13 {
14     int i;
15     for (i = 0; i < n; i++)
16         list[i] = i;
17 }
18 float sum(float list[], int n)
19 {
20     int i;
21     float tempsum = 0;
22     for (i = 0; i < n; i++)
23         tempsum += list[i];
24     return tempsum;
25 }

```

In C, array parameters have their values altered, despite the fact that the parameter passing is done using *call-by-value*.

2.1.2 Arrays in C

```
int one[] = {0, 1, 2, 3, 4};  
print1(&one[0], 5);
```

```
void print1(int *ptr, int rows)  
{/* print out a one-dimensional array using a pointer */  
    int i;  
    printf("Address Contents\n");  
    for (i = 0; i < rows; i++)  
        printf("%8u%5d\n", ptr + i, *(ptr + i));  
    printf("\n");  
}
```

Program 2.2: One-dimensional array accessed by address

Address	Contents
12244868	0
12344872	1
12344876	2
12344880	3
12344884	4

2.2 Dynamically Allocated Arrays

2.2.1 One-dimensional Arrays

```
#define MALLOC(p,s) \
    if ( ! ( (p) = malloc ( s ) ) ) { \
        fprintf(stderr, "Insufficient memory"); \
        exit(EXIT_FAILURE);\
    }

int i,n,*list;
printf("Enter the number of numbers to generate: ");
scanf_s ( "%d", &n);
if( n < 1 ) {
    fprintf(stderr, "Improper value of n\n");
    exit(EXIT_FAILURE);
    MALLOC(list, n * sizeof(int));
}
```

2.2.2 Two-Dimensional Arrays

❖ A multidimensional array in C

- *Array-of-arrays* representation

`int x[3][5];`

`x[i]` : a pointer to zeroth element of row *i* of the array

`x[i][j]` : an element accessed by the address,
`x[i]+j*sizeof(int)`

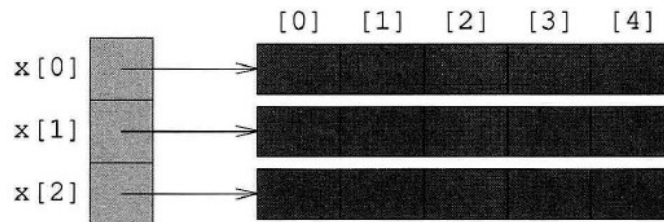
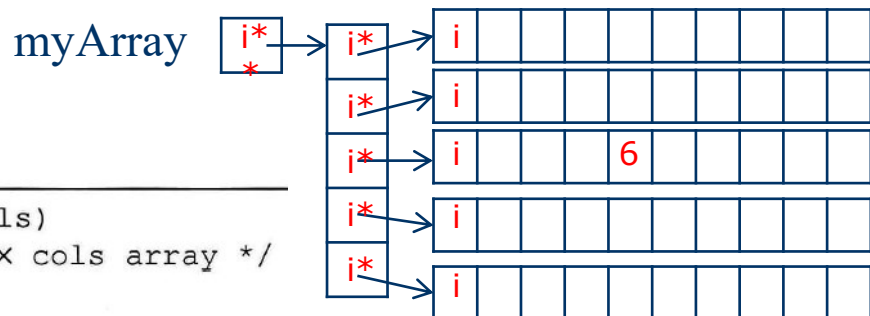


Figure 2.2: Array-of-arrays representation

```
int **myArray;
myArray = make2dArray(5,10);
myArray[2][4] = 6;
```



```
int** make2dArray(int rows, int cols)
{/* create a two dimensional rows X cols array */
    int **x, i;

    /* get memory for row pointers */
    MALLOC(x, rows * sizeof (*x));

    /* get memory for each row */
    for (i = 0; i < rows; i++)
        MALLOC(x[i], cols * sizeof (**x));
    return x;
}
```

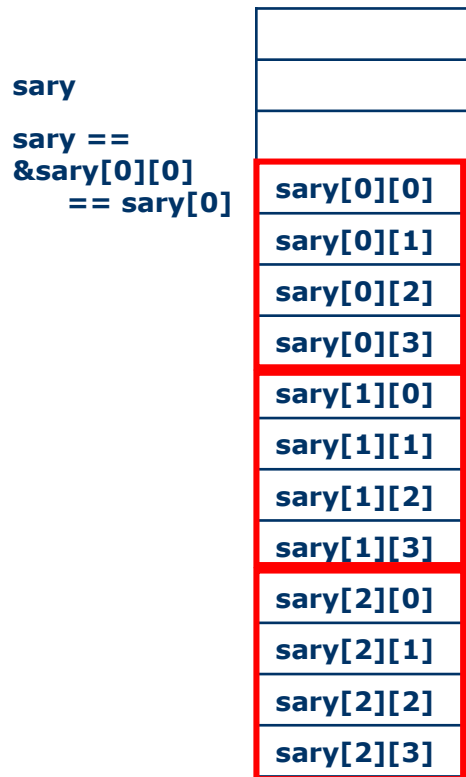
포인터배열

int *arr2[10] : 포인터 배열, arr2는 상수
int (*arr2)[10] : 배열 포인터, arr2는 [10]을 가
르키는 포인터 변수

Program 2.3: Dynamically create a two-dimensional array

정적 Array

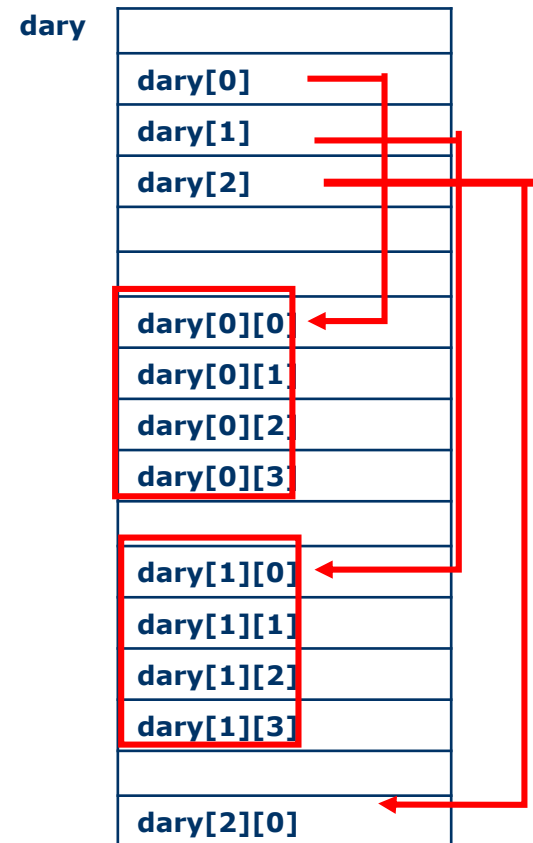
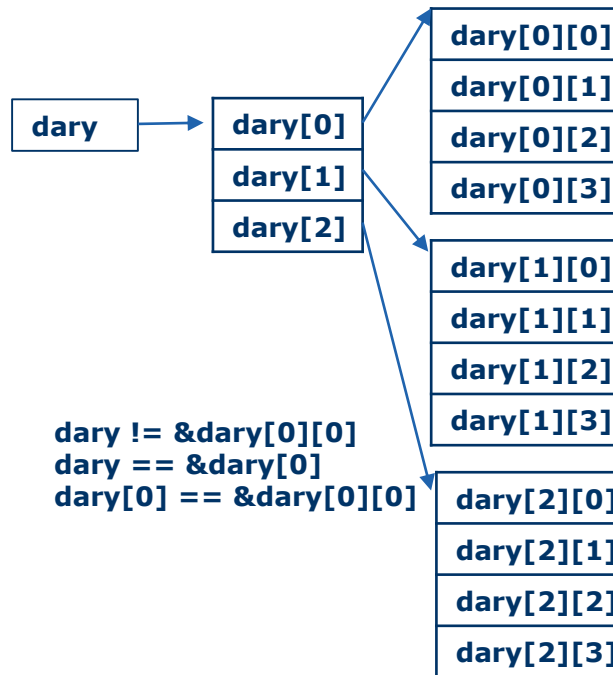
`int sary[3][4]`



동적 Array

`int **dary`

`dary`는 동적으로 할당 받은 2차원 배열[3][4]의 시작 주소를 가지고있는 포인터임



calloc

```
int *x, n;  
x = (int *) calloc(n, sizeof(int));  
/* allocated bits are set to 0*/
```

Write clean robust programs

```
#define CALLOC(p, n, s)\  
if (!(p) = calloc(n,s)) {\  
    fprintf(stderr, "Insufficient memory"); \  
    exit(EXIT_FAILURE);\  
}
```



realloc

```
/* changes the size of memory block pointed by x to  
s*sizeof(int) */  
int *old, *x, s;  
old = x;  
if ( (x = (int *)realloc(x, s*sizeof(int))) == NULL ) {  
    free(old);  
    exit(EXIT_FAILURE);  
}
```

...

free(x);



재 할당이 적을 경우: 메모리 반납



재 할당이 클 경우: 메모리 추가할당

realloc

```
#define REALLOC(p,s)\
if (! ( (p) = realloc (p, s))) {\
    fprintf(stderr, "Insufficient memory");\
    exit(EXIT_FAILURE);\
}
```

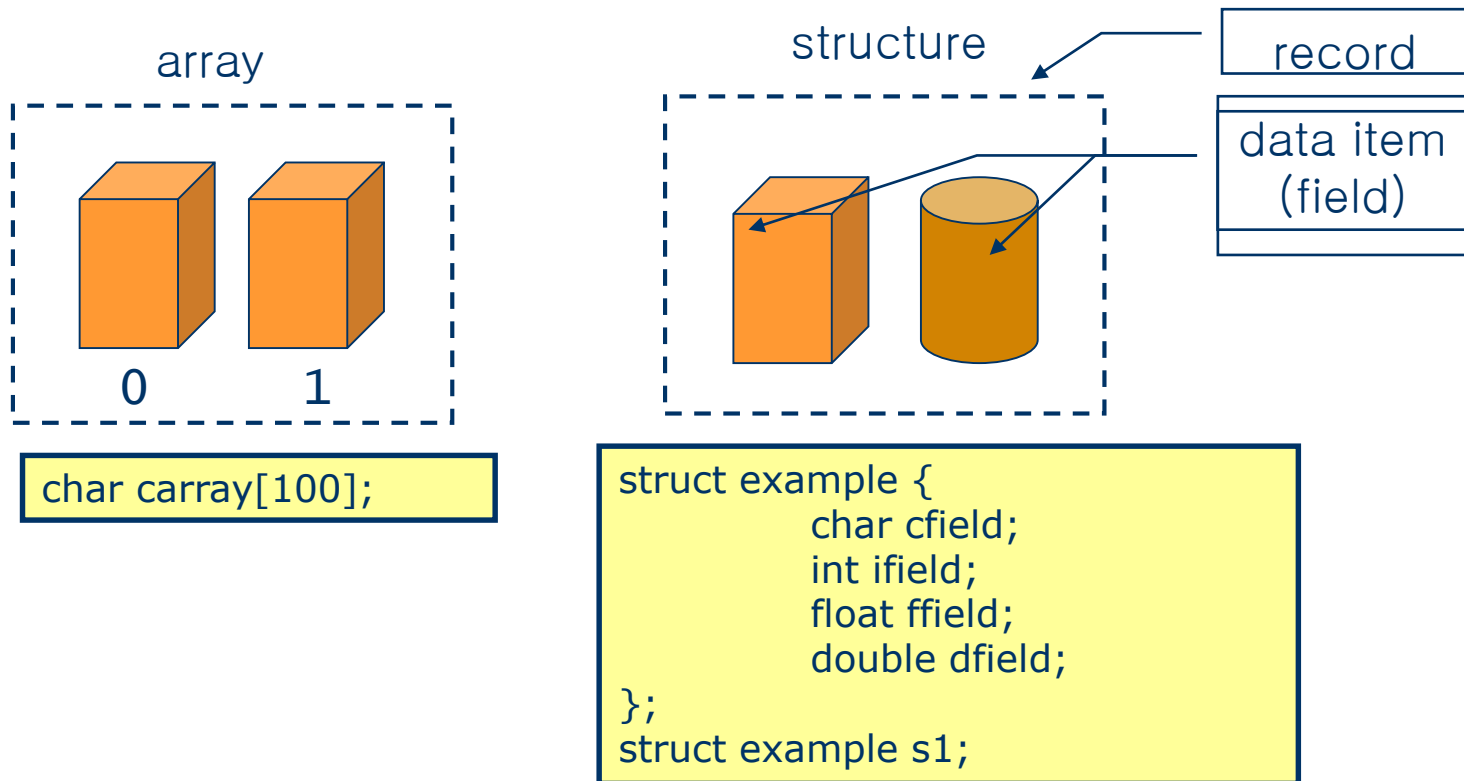
❖ Correct error.

```
#define REALLOC(o, p, s) \
    if (!((p) = realloc(o, s))) {\
        free(o);\
        fprintf(stderr, "Insufficient memory"); \
        exit(EXIT_FAILURE);\
    }

...
REALLOC(old, x, s*sizeof(int)); // x = realloc(x,
s*sizeof(int))
```

2.3 Structures and Union

2.3.1 Structures



2.3 Structures and Union

2.3.1 Structures

❖ Called a *record*

❖ Collection of data items

- Each item is identified as to its type and name

```
struct person{  
    char name[10];  
    int age;  
    float salary;  
} man;
```

- man is a variable
 - strcpy_s(man.name, sizeof(name), "kim");
 - man.age = 10;
- Structure member operator : dot(.)
 - struct person a, b;
 - a and b are variables.

2.3 Structures and Union

❖ Using the typedef statement

```
typedef structure {  
    char name[10]  
    int age;  
    float salary;  
} human;
```

```
human person1, person2; // Declaration of variables  
if(strcmp(person1.name, person2.name))  
    printf("The two people do not have the same name\n");  
else  
    printf("The two people have the same name\n");  
•
```

2.3 Structures and Union

❖ Structure assignment :

person1=person2;

- in ANSI C, OK!
 - However, don't use the assignment operation when the structure has a pointer to a memory space. Why?

❖ Check of equality or inequality :

if(person1==person2)

- cannot be checked directly

2.3 Structures and Union

❖ **strcpy_s format**

- `errno_t strcpy_s(char *dest, rsize_t dest_size, const char *src);`
 - Return Value
 - Zero if successful; otherwise, an error.

```
#include <stdio.h>    // for printf
#include <string.h>    // for strcpy_s

int main(void)
{
    char stringBuffer[80];
    strcpy_s(stringBuffer, sizeof(stringBuffer), "Hello world from ");
    printf("stringBuffer = %s\n", stringBuffer);
}
```

2.3 Structures and Union

❖ strcmp format

- `int strcmp (const char* str1, const char* str2);`
 - return 0 : str1 and str2 are equal
 - return >0 : str1 is greater than str2
 - return <0 : str2 is lower than str2

```
#include <stdio.h>
#include <string.h>
int main() {
    char str1[] = "abcd", str2[] = "abCd", str3[] = "abcd";
    int result;
    // comparing strings str1 and str2
    result = strcmp(str1, str2);
    printf("strcmp(str1, str2) = %d\n", result);
    // comparing strings str1 and str3
    result = strcmp(str1, str3);
    printf("strcmp(str1, str3) = %d\n", result);
    return 0;
}
```

2.3 Structures and Union

❖ Check of equality or inequality(cont')

```
int humansEqual(humanBeing person1, humanBeing
person2)
{/* return TRUE if person1 and person2 are the same
human
being otherwise return FALSE */
if (strcmp(person1.name, person2.name))
    return FALSE;
if (person1.age != person2.age)
    return FALSE;
if (person1.salary != person2.salary)
    return FALSE;
return TRUE;
}
```


2.3 Structures and Union

❖ A structure within a structure

```
typedef struct {  
    int month;  
    int day;  
    int year;  
} date;  
typedef struct {  
    char name [10];  
    int age;  
    float salary;  
    date dob;  
} humanBeing;
```

```
humanBeing person1, person2;  
person1.dob.month = 2;  
person1.dob.day = 11;  
person1.dob.year = 1944;
```

2.3.2 Unions

The fields share their memory space
Only one field is “active” at any given time.

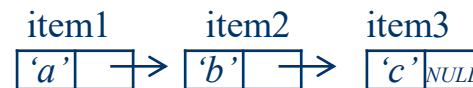
```
typedef struct {  
    enum tagField {female, male} sex;  
    union {  
        int children;  
        int beard;  
    } u;  
} sexType;
```

```
typedef struct {  
    char name[10];  
    int age;  
    float salary;  
    date dob;  
    sexType sexinfo;  
} humanBeing;  
humanBeing person1, person2;  
person1.sexinfo.sex = male;  
person1.sexinfo.u.beard = FALSE;  
person2.sexinfo.sex = female;  
person2.sexinfo.u.children = 4;
```

2.3.4 Self-Referential Structures

- ❖ A structure in which one or more of its components is a pointer to itself.

```
typedef struct list {  
    char data;  
    list *link ;  
} ;  
list item1, item2, item3;  
item1.data 'a';  
item2 .data 'b';  
item3.data 'c';  
item1.link = &item2;  
item2.link = &item3;  
item3.link = NULL;
```



2.4 Polynomials

2.4.1 The Abstract Data Type

❖ ***Ordered list or linear list***

- an ordered set of data items

❖ ex) Days-of-week

(Sun, Mon, Tue, Wed, Thu, Fri, Sat) : *list*
1st 2nd 3rd 4th 5th 6th 7th : *order*

- denote as ($item_0, item_1, \dots, item_{n-1}$)
- empty list : ()

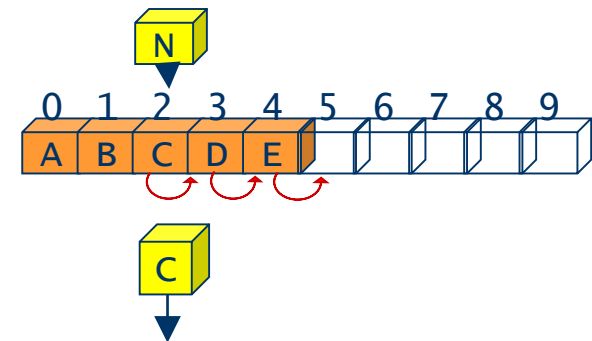
Implementation of Ordered List

❖ Array

- associate the list element, $item_i$, with the array index i
- *sequential mapping*
- retrieve, replace an item, or find the length of a list, in constant time
- **problems in insertion and deletion**
 - sequential mapping forces us to move items

❖ Linked List

- *Non-sequential mapping*
- Chapter 4



2.4 Polynomials

❖ Manipulation of symbolic polynomials

- $A(x) = 3x^{20} + 2x^5 + 4$, $B(x) = x^4 + 10x^3 + 3x^2 + 1$

❖ degree : the largest exponent of a polynomial

❖ assumption: unique exponents arranged in decreasing order

- When $A(x) = \sum a_i x^i$ and $B(x) = \sum b_i x^i$,
- $A(x) + B(x) = \sum (a_i + b_i) x^i$
- $A(x) B(x) = \sum (a_i x^i \sum (b_j x^j))$

2.4 Polynomials

❖ operations on ordered list

- find the length n
- read the items in a list from right to left (or left to right)
- retrieve i th item, $0 \leq i < n$
- replace i th item's value, $0 \leq i < n$
- insert i th position, $0 \leq i < n$
- delete i th item, $0 \leq i < n$

2.4 Polynomials

ADT Polynomial is

objects: $p(x) = a_1x^{e_1} + \dots + a_nx^{e_n}$; a set of ordered pairs of $\langle e_i, a_i \rangle$ where a_i in *Coefficients* and e_i in *Exponents*, e_i are integers ≥ 0

functions:

for all $poly, poly1, poly2 \in \text{Polynomial}$, $coef \in \text{Coefficients}$, $expon \in \text{Exponents}$

<i>Polynomial</i> Zero()	::=	return the polynomial, $p(x) = 0$
<i>Boolean</i> IsZero(<i>poly</i>)	::=	if (<i>poly</i>) return FALSE else return TRUE
<i>Coefficient</i> Coef(<i>poly</i> , <i>expon</i>)	::=	if (<i>expon</i> \in <i>poly</i>) return its coefficient else return zero
<i>Exponent</i> LeadExp(<i>poly</i>)	::=	return the largest exponent in <i>poly</i>
<i>Polynomial</i> Attach(<i>poly</i> , <i>coef</i> , <i>expon</i>)	::=	if (<i>expon</i> \in <i>poly</i>) return error else return the polynomial <i>poly</i> with the term $\langle coef, expon \rangle$ inserted
<i>Polynomial</i> Remove(<i>poly</i> , <i>expon</i>)	::=	if (<i>expon</i> \in <i>poly</i>) return the polynomial <i>poly</i> with the term whose exponent is <i>expon</i> deleted else return error
<i>Polynomial</i> SingleMult(<i>poly</i> , <i>coef</i> , <i>expon</i>)	::=	return the polynomial $poly \cdot coef \cdot x^{expon}$
<i>Polynomial</i> Add(<i>poly1</i> , <i>poly2</i>)	::=	return the polynomial $poly1 + poly2$
<i>Polynomial</i> Mult(<i>poly1</i> , <i>poly2</i>)	::=	return the polynomial $poly1 \cdot poly2$

end Polynomial

ADT 2.2: Abstract data type *Polynomial*

Representation of polynomials in C

```
#define MAX-DEGREE 101
/*Max degree of polynomial+1*/
typedef struct {
    int degree;
    float coef[MAX-DEGREE];
} polynomial;
```

❖ $A(x) = \sum_{i=0}^n a_i x^i$ would be represented as :

- $a.\text{degree} = n$
- $a.\text{coef}[i] = a_{n-i}, \quad n < \text{MAX_DEGREE}$

❖ ✖ $a.\text{coef}[i]$ is the coefficient of x^{n-i}

Representation of polynomials in C

$$A(x) = 2x^{1000} + 1 \text{ and } B(x) = x^4 + 10x^3 + 3x^2 + 1$$

1000

2	0	0	0	0	0	0	0	0	0	0	0	0	1
---	---	---	---	---	---	-------	---	---	---	---	---	---	---	---

4

1	1 0	3	0	1
---	--------	---	---	---

Initial version of *padd* function (cont')

$$D(x) = 0$$

$$A(x) = 2x^{1000} + 2x^3$$

$$B(x) = x^4 + 10x^3 + 3x^2 + 1$$

(step1)

$$D(x) = 2x^{1000}$$

$$A(x) = 2x^3$$

$$B(x) = x^4 + 10x^3 + 3x^2 + 1$$

(step2)

$$D(x) = 2x^{1000} + x^4$$

$$A(x) = 2x^3$$

$$B(x) = 10x^3 + 3x^2 + 1$$

(step3)

$$D(x) = 2x^{1000} + x^4 + 12x^3$$

$$A(x) = 0$$

$$B(x) = 3x^2 + 1$$

(step4)

$$D(x) = 2x^{1000} + x^4 + 12x^3 + 3x^2 + 1$$

$$A(x) = 0$$

$$B(x) = 0$$

2.4.2 Polynomial Representation

```
#define COMPARE(x, y) ( ((x) < (y)) ? -1 : ((x) == (y)) ? 0: 1 ) ※ p.12
```

```
/* d = a + b, where a, b, and d are polynomials */
d = Zero()
while (! IsZero(a) && ! IsZero(b)) do {
    switch COMPARE(LeadExp(a), LeadExp(b)) {
        case -1: d =
            Attach(d, Coef(b, LeadExp(b)), LeadExp(b));
            b = Remove(b, LeadExp(b));
            break;
        case 0: sum = Coef(a, LeadExp(a))
            + Coef(b, LeadExp(b));
            if (sum) {
                Attach(d, sum, LeadExp(a));
                a = Remove(a, LeadExp(a));
                b = Remove(b, LeadExp(b));
            }
            break;
        case 1: d =
            Attach(d, Coef(a, LeadExp(a)), LeadExp(a));
            a = Remove(a, LeadExp(a));
    }
    insert any remaining terms of a or b into d
}
```

Program 2.5: Initial version of *padd* function

Representation of polynomials in C(cont')

```
MAX-TERMS 100 /*size of terms array*/  
typedef struct {  
    float coef;  
    int expon;  
} polynomial;  
polynomial terms[MAX-TERMS];  
int avail = 0;
```

❖ $A(x) = 2x^{1000} + 1$ and $B(x) = x^4 + 10x^3 + 3x^2 + 1$

<i>coef</i>	2	1
<i>exp</i>	1000	0
	0	1

Representation of polynomials in C(cont')

Using one array to represent two polynomials

$$A(x) = 2x^{1000} + 1 \text{ and } B(x) = x^4 + 10x^3 + 3x^2 + 1$$

	<i>startA</i>	<i>finishA</i>	<i>startB</i>		<i>finishB</i>	<i>avail</i>
	↓	↓	↓		↓	↓
<i>coef</i>	2	1	1	10	3	1
<i>exp</i>	1000	0	4	3	2	0
	0	1	2	3	4	5
						6

iterations
 $\leq m+n-1$

$\leq m$

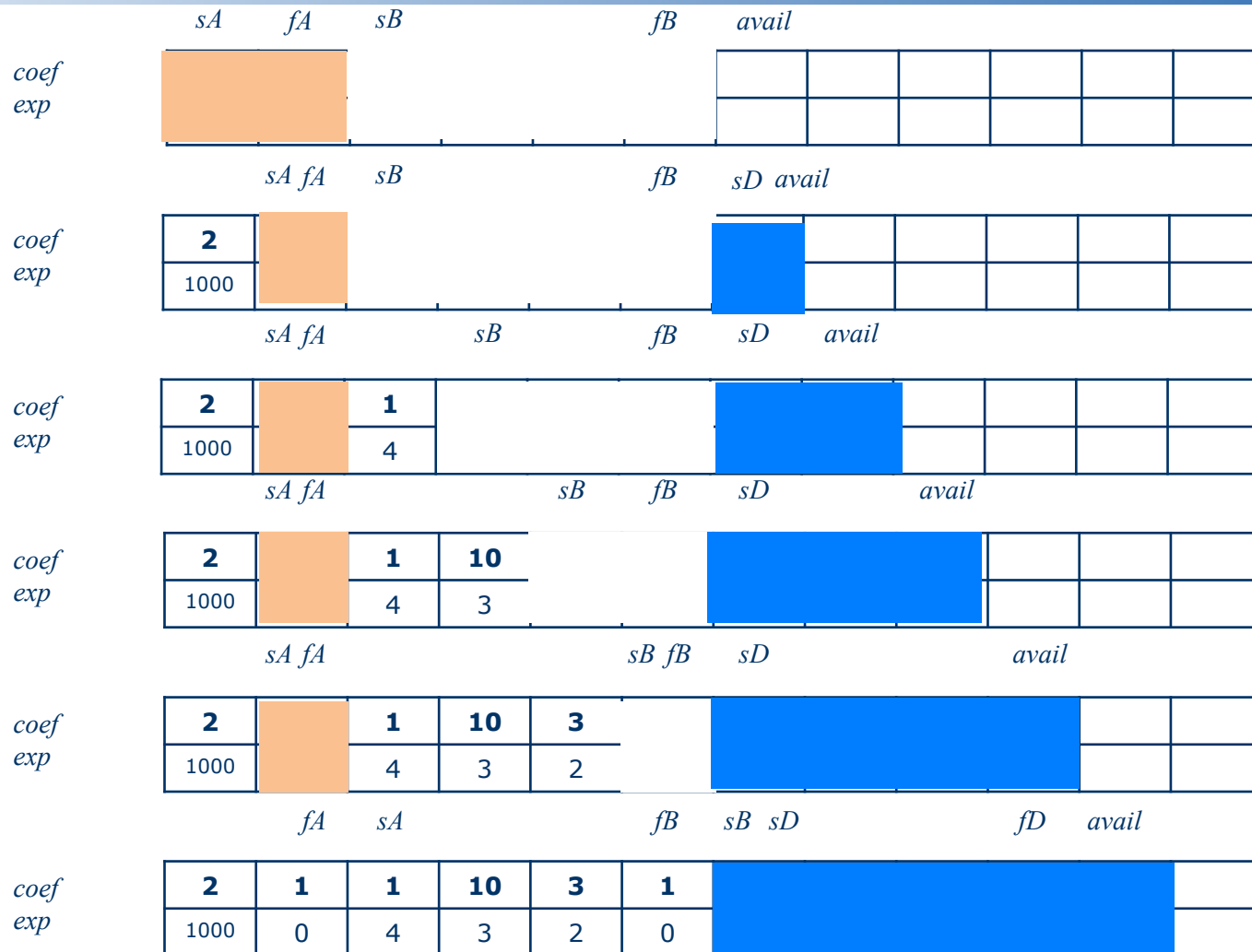
$\leq n$

```
void padd(int startA,int finishA,int startB, int finishB,
          int *startD,int *finishD)
{/* add A(x) and B(x) to obtain D(x) */
  float coefficient;
  *startD = avail;
  while (startA <= finishA && startB <= finishB)
    switch(COMPARE(terms[startA].expon,
                  terms[startB].expon)) {
      case -1: /* a expon < b expon */
        attach(terms[startB].coef,terms[startB].expon);
        startB++;
        break;
      case 0: /* equal exponents */
        coefficient = terms[startA].coef +
                      terms[startB].coef;
        if (coefficient)
          attach(coefficient,terms[startA].expon);
        startA++;
        startB++;
        break;
      case 1: /* a expon > b expon */
        attach(terms[startA].coef,terms[startA].expon);
        startA++;
    }
  /* add in remaining terms of A(x) */
  for(; startA <= finishA; startA++)
    attach(terms[startA].coef,terms[startA].expon);
  /* add in remaining terms of B(x) */
  for( ; startB <= finishB; startB++)
    attach(terms[startB].coef, terms[startB].expon);
  *finishD = avail-1;
}
```

Program 2.6: Function to add two polynomials

```
void attach(float coefficient, int exponent)
{
    /* add a new term to the polynomial */
    if (avail >= MAX_TERMS) {
        fprintf(stderr, "Too many terms in the polynomial\n");
        exit(EXIT_FAILURE);
    }
    terms[avail].coef = coefficient;
    terms[avail++].expon = exponent;
}
```

Program 2.7: Function to add a new term



2.4.3 Polynomial Addition

❖ Analysis of *padd*

- Let m and n be the number of nonzero terms in A and B, respectively.
- ① If $m > 0$ and $n > 0$, **while loop**
 - each iteration : $O(1)$
 - The iteration terminates when either *startA* or *startB* exceeds *finishA* or *finishB*, respectively
 - The number of iterations is bounded by $m+n-1$
 - the worst case : ex) $a(x) = x^6 + x^4 + x^2 + x^0$, $b(x) = x^7 + x^5 + x^3 + x^1$
- ② The remaining **two for loops** $\rightarrow O(m+n)$
- ①&② $\rightarrow O(m+n)$