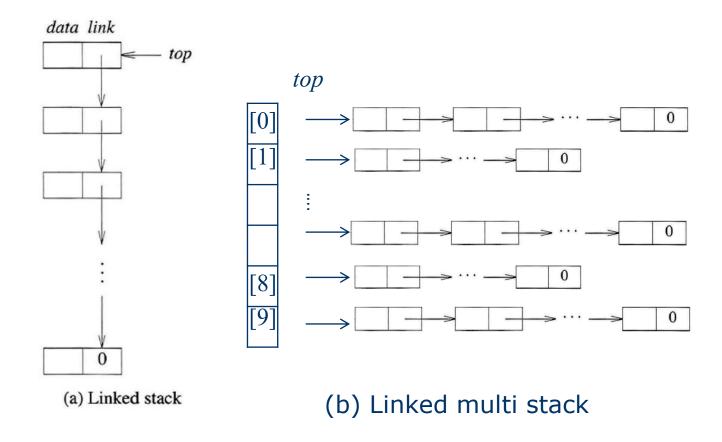
Chap 4. Linked Lists (2)

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- **Chapter 2. Arrays And Structures**
- **Chapter 3. Stacks And Queues**
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Representing n≤MAX_STACKS stacks simultaneously

#define MAX-STACKS 10 /* maximum number of stacks */

typedef struct {
 int key;
 /* other fields */
 } element;

typedef struct stack *stackPointer;

typedef struct stack {
 element data;
 stackPointer link;
 };

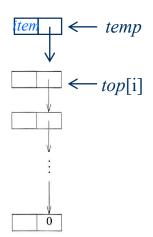
...
};

stackPointer top[MAX-STACKS];

- **♦** top[i] = NULL, 0 ≤ i < MAX STACKS //Initial conditions for the stacks
- * top [i] ==NULL iff the ith stack is empty //Boundary condition
 for the ith stack

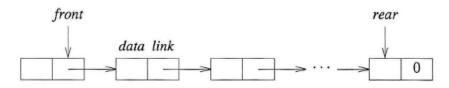
```
void push(int i, element item)
{/* add item to the ith stack */
    stackPointer temp;
    MALLOC(temp, sizeof(*temp));
    temp->tdata = item;
    temp->tlink = top[i];
    top[i] = temp;
}
```

Program 4.5: Add to a linked stack

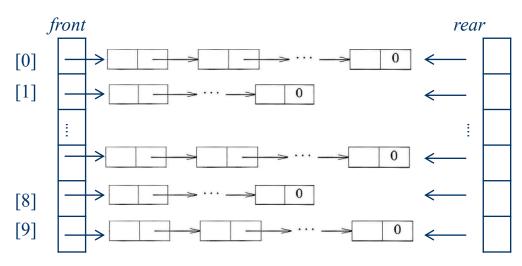


```
element pop(int i)
{/* remove top element from the ith stack */
    element item;
    stackPointer temp= top[i];
    if (!temp)
        return stackEmpty();
    item = temp->data;
    top[i] = temp->link;
    free (temp);
    return item;
```

Program 4.6: Delete from a linked stack



(a) Linked queue



(b) Linked multi queue

 \diamond Representing $n \leq MAX_QUEUES$ queues simultaneously

```
#define MAX-QUEUES 10 // maximum number of queues
typedef struct queue *queuePointer;
typedef struct queue {
    element data;
    queuePointer link;
    };
queuePointer front[MAX-QUEUES], rear[MAX-QUEUES];
```

- * front[i] =NULL, 0 ≤ i <MAX-QUEUES Initial conditions for the queues
 </p>
- * front[i] = NULL iff the ith queue is empty Boundary condition for the ith queue

```
void addq(i, item)
{/* add item to the rear of queue i */
    queuePointer temp;
    MALLOC(temp, sizeof(*temp));
    temp->data = item;
    temp->link = NULL;
    if (front [i])
        rear[i]->link =temp;
    else
        front[i] =temp;
    rear[i] =temp;
}
```

❖ Program 4.7: Add to the rear of a linked queue

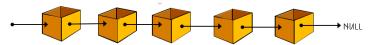


 $front[i] \qquad rear[i]$ $\downarrow \qquad \qquad \downarrow$ $temp \rightarrow item \qquad \rightarrow \qquad \qquad 0$

Circular List Representation

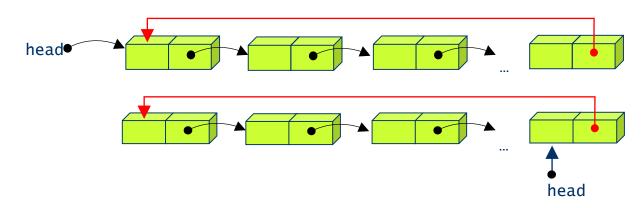
& Chain

A singly linked list in which the last node has a null link



Circular list

 The link filed of the last node points to the first node in the list



Operations For Circularly Linked Lists

```
void insertFront(listPointer *last, listPointer node)
{/* insert node at the front of the circular list whose
     last node is last */
     if (!(*last)) {
    /* list is empty, change last to point to new entry */
          *last = node;
          node->link = node;
    else {
    /* list is not empty, add new entry at front */
          node->link = (*last) ->link;
                                                      node
                                                                                     *last
          (*last) ->link = node;
```

Program 4.18: Inserting at the front of a list

Operations For Circularly Linked Lists

```
void insertLast(listPointer *last, listPointer node)
{/* insert node at the front of the circular list whose
     last node is last */
     if (!(*last)) {
    /* list is empty, change last to point to new entry */
          *last = node;
          node->link = node;
    else {
    /* list is not empty, add new entry at front */
          node->link = (*last) ->link;
          (*last) ->link = node;
                                                 node
                                                                                *last
          *last = node;
```

Program 4.18: Inserting at the front of a list

Operations For Circularly Linked Lists

```
int length(listPointer last)
{/* find the length of the circular list last */
     listPointer temp;
     int count = 0:
     if (last)
         temp = last;
          do {
                                                                               last
          count++;
          temp = temp->link;
                                                                  temp last
          } while (temp !=last);
     return count;
```

Program 4.19: Finding the length of a circular list

4.4 Polynomials

4.4.1 Polynomial Representation

Polynomial

```
A(x) = a_{m-1}x^{e_{m-1}} + \dots + a_0x^{e_0}

• e_{m-1} > e_{m-2} \dots > e_1 > e_0 ≥ 0

• a_i: nonzero coefficients

• e_i: nonnegative integer exponents
```

Representation of Polynomial

```
typedef struct polyNode *polyPointer;
typedef struct polyNode {
   int coef;
   int expon;
   polyPointer link;
   };
polyPointer a,b;
```

4.4 Polynomials

Representation of polynomials

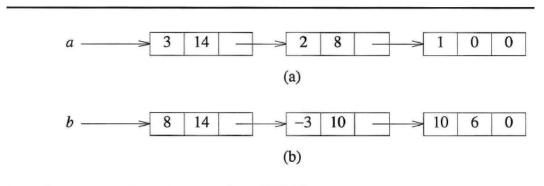


Figure 4.12: Representation of $3x^{14} + 2x^8 + 1$ and $8x^{14} - 3x^{10} + 10x^6$

$$a = 3x^{14} + 2x^{8} + 1$$

$$b = 8x^{14} - 3x^{10} + 10x^{6}$$

$$8 \quad 14 \quad -3 \quad 10 \quad 10 \quad 6 \quad 0$$

$$c \quad 11 \quad 14 \quad 0$$

$$(i) \quad a \rightarrow expon == b \rightarrow expon$$

$$3 \quad 14 \quad -3 \quad 10 \quad 10 \quad 6 \quad 0$$

$$c \quad 11 \quad 14 \quad -3 \quad 10 \quad 10 \quad 6 \quad 0$$

$$c \quad 11 \quad 14 \quad -3 \quad 10 \quad 0$$

$$(ii) \quad a \rightarrow expon < b \rightarrow expon$$

(iii) $a \rightarrow expon > b \rightarrow expon$ Figure 4.13: Generating the first three terms of c = a + b

10

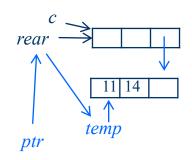
2 | 8

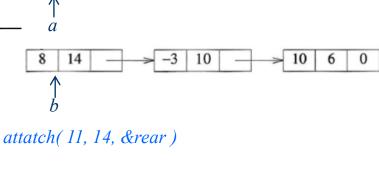
11 14

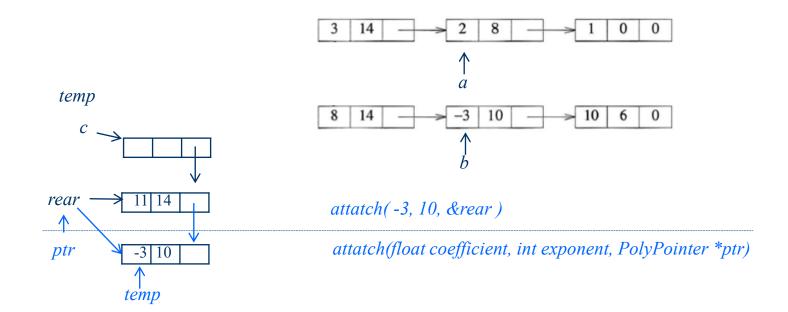
```
polyPointer padd(polyPointer a, polyPointer b)
{/* return a polynomial which is the sum of a and b */
   polyPointer c, rear, temp;
   int sum;
   MALLOC(rear, sizeof(*rear));
   c = rear;
   while (a && b)
      switch (COMPARE(a→expon, b→expon)) {
         case -1: /* a\rightarrowexpon < b\rightarrowexpon */
                attach (b\rightarrowcoef, b\rightarrowexpon, &rear);
                b = b \rightarrow link;
                break;
         case 0: /* a\rightarrowexpon = b\rightarrowexpon */
                sum = a \rightarrow coef + b \rightarrow coef;
                if (sum) attach(sum, a→expon, &rear);
                a = a \rightarrow link; b = b \rightarrow link; break;
         case 1: /* a→expon > b→expon */
                attach (a \rightarrow coef, a \rightarrow expon, &rear);
                a = a \rightarrow link;
   /* copy rest of list a and then list b */
   for (; a; a = a \rightarrow link) attach(a \rightarrow coef, a \rightarrow expon, & rear);
   for (; b; b = b\rightarrowlink) attach(b\rightarrowcoef,b\rightarrowexpon,&rear);
   rear→link = NULL:
   /* delete extra initial node */
   temp = c; c = c \rightarrow link; free(temp);
   return c;
```

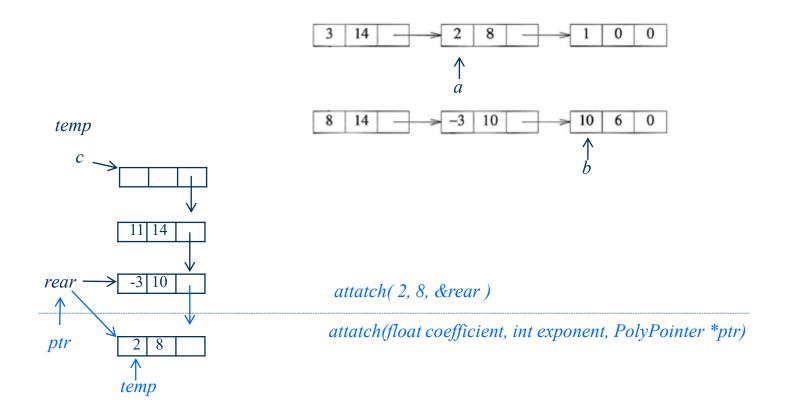
Program 4.9: Add two polynomials

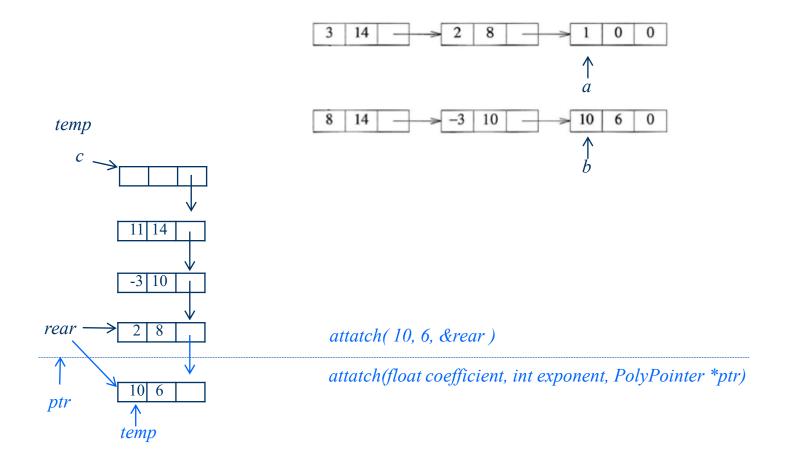
Program 4.10: Attach a node to the end of a list

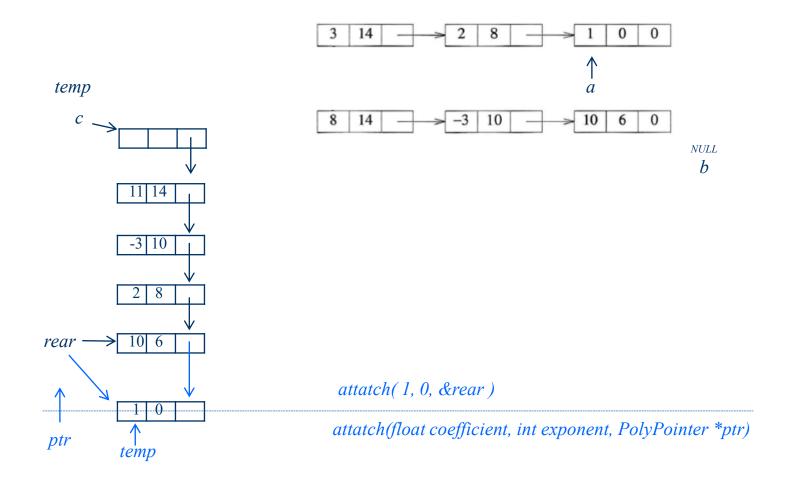


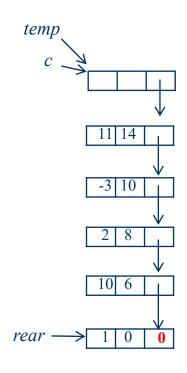


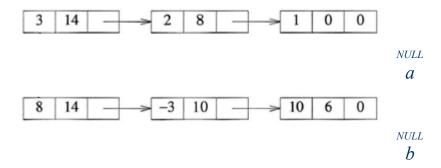












Analysis of padd

Three cost measures for this algorithm

(1)Coefficient additions

$$A(x) = a_{m-1}x^{e_{m-1}} + \dots + a_0x^{e_0}$$

$$B(x) = b_{n-1}x^{f_{n-1}} + \dots + b_0x^{f_0}$$
 where a_i , $b_i \neq 0$ and $e_{m-1} > e_{m-2} \dots > e_1 > e_0 \geq 0$,
$$f_{n-1} > f_{n-2} \dots > f_1 > f_0 \geq 0$$

 $0 \le \text{number of coefficient additions} \le \min\{m, n\}$

(2) Exponent comparisons

- One comparison on each iteration of the while loop
- The number of iterations is bounded by m + n
 ex) m+n-1 iterations, for example m == n

$$e_{m-1} > f_{n-1} > e_{m-2} > f_{n-2} \ldots > e_1 > f_1 > e_0 > f_0$$
 다항식 a의 지수 다항식 b의 지수

(3) Creations of new nodes for c

• The maximum number of terms in c is m + n

From
$$(1)\sim(3)$$
,

• the total time complexity is O(m + n)

4.4.3 Erasing Polynomials

```
void erase(polyPointer *ptr)
{/* erase the polynomial pointed to by ptr */
    polyPointer temp;
    while (*ptr) {
        temp = *ptr;
        *ptr = (*ptr)->link;
        free (temp);
                            erase(&a)
                            erase(polyPointer *ptr)
```

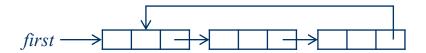
Available space list

- A chain of nodes that have been "freed"
- Use getNode and retNode, instead of malloc
 & free

$$avail \longrightarrow \bigcirc \bigcirc \bigcirc$$

When maintaining it,

 we can obtain an efficient erase algorithm for circular list.



```
polyPointer getNode(void)
{/* provide a node for use */
     polyPointer node;
     if (avail) {
                                                                                   0
          node = avail:
          avail = avail->link;
                                                                       node = getNode()
                                                     avail
                                                                          getNode(void)
     else
                                                  node
          MALLOC(node, sizeof(*node));
     return node;
}
Program 4.12: getNode function
void retNode(polyPointer node)
{/* return a node to the available list */
                                                                       retNode(node)
                                                    avail <
                                                               avail
     node->link = avail:
                                                                retNode(polyPointer node)
                                                 node
     avail = node;
Program 4.13: retNode function
```

```
void cerase(polyPointer *ptr)
{/* erase the circular list pointed to by ptr */
     polyPointer temp;
     if ( *ptr) {
          temp = (*ptr)->link;
          (*ptr)->link = avail;
                                                        avail
          avail = temp;
          *ptr = NULL;
                                                            A
Program 4.14: Erasing a circular list
                                                        temp
                                                                                     ptr
                                                                                     NULL
                                                cerase(&first)
```

A header node is added to easily handle polynomial addition in a circular

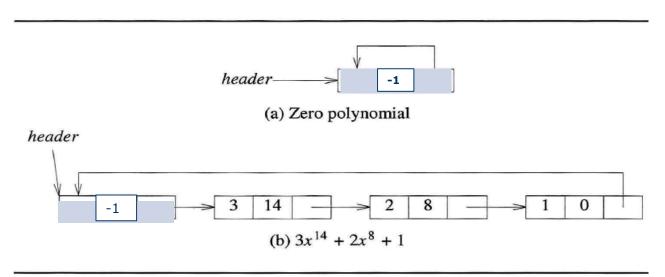


Figure 4.15: Example polynomials with header nodes

```
polyPointer cpadd(polyPointer a, polyPointer b)
{/* polynomials a and b are singly linked circular lists
    with a header node. Return a polynomial which is
    the sum of a and b */
   polyPointer startA, c, lastC;
   int sum, done = FALSE;
   startA = a;
                         /* record start of a */
                         /* skip header node for a and b*/
   a = a \rightarrow link;
   b = b \rightarrow link;
                           /* get a header node for sum */
   c = getNode();
   c \rightarrow expon = -1; lastC = c;
   do {
      switch (COMPARE(a→expon, b→expon)) {
         case -1: /* a→expon < b→expon */
                attach(b→coef, b→expon, &lastC);
               b = b \rightarrow link;
                break;
         case 0: /* a\rightarrowexpon = b\rightarrowexpon */
                if (startA == a) done = TRUE;
                else {
                   sum = a \rightarrow coef + b \rightarrow coef;
                   if (sum) attach(sum, a→expon, &lastC);
                   a = a \rightarrow link; b = b \rightarrow link;
                break:
         case 1: /* a→expon > b→expon */
                attach(a \rightarrow coef, a \rightarrow expon, &lastC);
                a = a \rightarrow link;
   } while (!done);
   lastC \rightarrow link = c;
   return c;
```

Program 4.15: Adding two polynomials represented as circular lists with header nodes

