

# Dynamically Discovering Likely Program Invariants to Support Program Evolution

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# Motivation

- Programs evolve, and they need to sustain some properties over changes
- By explicitly stating the specification inside program code, the one who modifies the code can understand the properties that must be preserved at the modification
- An invariant is a form of program specification that asserts that a certain predicate at a program location always evaluates to true when it gets executed
  - Invariants are embedded in the target program code (e.g., assert statement)
  - Invariants can be used to protect the program from adversarial changes that violate assumptions that the program correctly behaves

# Ex. Invariants of CircularFifoQueue

```
private transient E[] elements;
private transient int start = 0;
private transient int end = 0;
private transient boolean full = false;
private final int maxElements;

01 public boolean add(final E element) {
02     if (null == element)
03         throw new NullPointerException
04         ( "Attempted to add null object to queue");
05     if (isAtFullCapacity())
06         remove();
07     elements[end++] = element;
08     if (end >= maxElements)
09         end = 0;
10     if (end == start)
11         full = true ;
12     return true ;
13 }
```

- `elements.length <= maxElements`
- `full == true -> start == end`

- `start >= 0`
- `start < maxElements`
- `end >= 0`
- `end < maxElements`

- `element != null`
- `isAtFullCapacity() == false`
- `full == false`

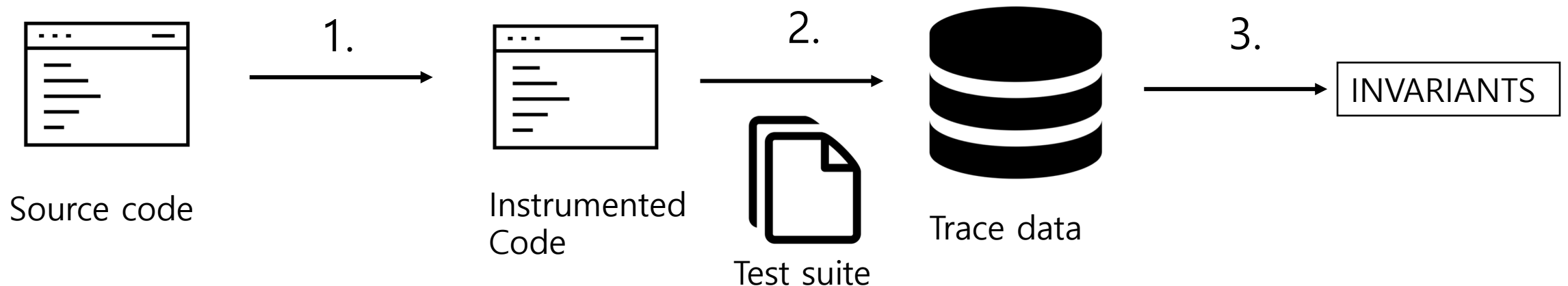
- `full -> full`
- `element ∈ elements`

# Daikon

- Daikon is a tool that discovers invariants of a program by observing the target program executions
  - Input : a set of target Java classes and their test cases
  - Output
    1. Preprocessing (DynComp): decls file which is a list of grouped variables of the same abstract data type
    2. Invariant extractor (Chicory): a list of invariants of given target classes
      - class invariants,
      - method pre-/post-condition
      - loop invariant

# Workflow

1. Instrument source program to trace variables of interest
2. Run instrumented program over test suite
3. Check properties (infer invariants) over instrument variables and derived variables



# Ex. Likely Invariants Inferred by Daikon

```
01 public boolean add(final E element) {
02     if (null == element)
03         throw new NullPointerException
04         ( "Attempted to add null object to queue");
05     if (isAtFullCapacity())
06         remove();
07     elements[end++] = element;
08     if (end >= maxElements)
09         end = 0;
10     if (end == start)
11         full = true ;
12     return true ;
13 }
```

- True positive

```
this.start == orig(this.start)
return == true
orig(arg0) == this.elements[orig(this.end)]
```

- False positive

```
this.start <= orig(this.end)
this.end one of { 0, 1, 2 }
arg0 != null
```

- c.f. False negative

```
this.full == false
this.maxelements < 100
```

# Tutorial

- <http://bit.ly/2LsC04Z>

# Invariant Inference Process Overview

1. Obtain a program trace which is a list of a variable name and its value at an execution point (method entry/exist points)
2. Create the initial set of likely invariant candidates by combining programs variables according to the invariant templates.
3. Exam whether a candidate is consistent with the given program trace
  - throw away violating invariants from the candidate set
4. Exam whether a candidate is confident according to the given program trace
5. Return the consistent & confident invariants as the result



# Step 1. Trace Data

- Trace the values of all variables in the scope at every program point along an execution
- Program points include every loop head, every procedure entry/exit point

# Ex. CircularFifoQueue fields and add method

```
private transient E[] elements;
private transient int start = 0;
private transient int end = 0;
private transient boolean full = false;
private final int maxElements;

01 public boolean add(final E element) { ← P1
02     if (null == element)
03     throw new NullPointerException
04     ( "Attempted to add null object to queue");
05     if (isAtFullCapacity())
06         remove();
07     elements[end++] = element;
08     if (end >= maxElements)
09         end = 0;
10     if (end == start)
11         full = true ;
12     return true ; ← P2
13 }
```

Point	Assignment (variable-value map)
...	...
P1	this.elements = [1, 2, 3, ], this.start = 0, this.end = 3, this.full = false, this.maxElement = 10 element = 4
P2	this.elements = [1, 2, 3, ], this.start = 0, this.end = 3, this.full = false, this.maxElement = 10 element = 4, \result = true
...	...
P1	this.elements = [1, 2, 3 ], this.start = 0, this.end = 0, this.full = true, this.maxElement = 3, element = 4
P2	this.elements = [4, 2, 3, ], this.start = 1, this.end = 1, this.full = true, this.maxElement = 3, element = 4, \result = true
...	...

## Step 2. Inferring Invariants

- Detect invariant at a specific program execution point
- Create set of potential invariants over variables in the code and derived variables that are not explicitly stated in the code using given templates
  - The sets will be created for unary, binary, and ternary tuple of variables

# Templates for inferring invariants (1/2)

- Invariants over variable
  - Any Variable
    - Constant value, uninitialized variable, small value set
  - Single numeric
    - Range limits, nonzero, modulus, non modulus
  - Two numeric
    - Linear relationship, ordering comparison, functions, invariants over  $x+y$ , invariants over  $x-y$
  - Three numeric
    - Linear relationship, functions
  - Single sequence
    - Range (min, max, ordered), element ordering, invariants over all sequence elements
  - Two Sequence
    - Linear relationship, comparison, subsequence relationship, reversal
  - A Sequence and a numeric
    - membership

# Ex. CircularFifoQueue Inferring Invariants over variables

- Unary :
  - `maxElements == 10`
  - $0 \leq \text{element} \leq 5$
  - `return == true`
- Binary:
  - `orig(maxElements) == maxElements`
  - `orig(elements) != elements`
- Ternary:
  - `orig(end) == 3*end + 2*maxElements + 1`

# Templates for inferring invariants(1/2)

- Invariants over derived variables
  - Derived from any sequence
    - length (size)
    - Extremal elements (  $s[0]$ ,  $s[1]$ ,  $s[\text{size}(s)-1]$ , ...)
  - Derived from any numeric sequence
    - Sum
    - Minimum element
    - Maximum element
  - Derived from any sequence  $s$  and numeric variable  $i$ 
    - Element at index
    - Subsequences ( $s[0 \dots i]$ )
  - Derived from function invocations
    - Number of calls
- Invariants over derived variables are introduced in stages because computed invariants can derive other variables

## Ex. CircularFifoQueue Inferring Invariants over derived variables

- Unary :
  - $\text{sum}(\text{elements}) < 20$
  - $\text{min}(\text{elements}) == 0$
- Binary:
  - $0 \leq \text{size}(\text{elements}) < \text{maxElements}$
  - $\text{size}(\text{orig}(\text{elements})) < \text{size}(\text{elements})$

# Step 3. Filter Valid Invariants

- Checks unary, binary, and ternary invariants for tuple of variables by examining each sample
  - Each set of possible invariants is tested against various combination of traced variables
- When sample not satisfying invariant is met, the invariant is thrown away and will not be checked for the rest of the samples



# Ex. CircularFifoQueue fields and add method

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06         remove();
07     elements[end++] = element;
08     if (end >= maxElements)
09         end = 0;
10     if (end == start)
11         full = true ;
12     return true ; ← P2
13 }
```

Point	Assignment (variable-value map)
...	...
P1	this.elements = [1, 2, 3, ], this.start = 0, this.end = 3, this.full = false, this.maxElement = 10 element = 4
P2	this.elements = [1, 2, 3, ], this.start = 0, this.end = 3, this.full = false, this.maxElement = 10 element = 4, \result = true
...	...
P1	this.elements = [1, 2, 3 ], this.start = 0, this.end = 0, this.full = true, this.maxElement = 3, element = 4
P2	this.elements = [4, 2, 3, ], this.start = 1, this.end = 1, this.full = true, this.maxElement = 3, element = 4, \result = true
...	...

# Ex. CircularFifoQueue Inferring Invariants over variables

- Unary :
  - `maxElements == 10`
    - Falsified in the second sample when `maxElements == 3`
  - `0 ≤ element ≤ 5`
    - Passed for both samples
  - `return == true`
    - Passed for both samples
- Binary:
  - `start < end`
    - Falsified for the second sample when `start == end == 1`
  - `orig(elements) != elements`
    - Passed for both samples
- Ternary:
  - `orig(end) == 3*end + 2*maxElements + 1`
    - Falsified in the first sample when `end == 3, maxElements == 10`

## Ex. CircularFifoQueue Inferring Invariants over derived variables

- Unary :
  - $\text{sum}(\text{elements}) < 20$ 
    - Passed by all samples
  - $\text{min}(\text{elements}) == 0$ 
    - Falsified in the first sample when  $\text{min}(\text{element}) == 1$
- Binary:
  - $0 \leq \text{size}(\text{elements}) < \text{maxElements}$ 
    - Passed by all samples
  - $\text{size}(\text{orig}(\text{elements})) < \text{size}(\text{elements})$ 
    - Falsified in the second sample when  $\text{size}(\text{orig}(\text{elements})) == \text{size}(\text{elements}) == 3$

# Step 4. Spurious Invariants

- Not all of unfalsified invariants are important / correct
- Correctness of generated invariants depend largely on the test suite
- Solution
  - Use of larger, complete test suites
  - Method for computing invariant confidence
  - Prune invariants logically implied by other invariants

# Invariant Confidence

- Plausibility of an invariant
  - Inadequate number of samples of particular variable means patterns observed over the variable may be a coincidence
- Compute probability that such property appear by change in random input
- If probability is smaller than user specified confidence parameter, it is considered non coincidental and is reported

## Ex. Computing Confidence for CircularFifoQueue

- Suppose that Daikon infers an invariant `maxElements != 0`
- Hypothesis: Daikon conjectures that the values of a non-negative integer `maxElements` are uniformly distributed over range between 0 and  $R$
- P-value: the probability of missing the `maxElements` being 0 case in  $N$  samples:  
$$\left(1 - \frac{1}{R}\right)^N$$
- Report an inferred invariant only if the p-value does not exceed the given confidence level (e.g., 0.01)