

# Algorithms and Data Structures

## - Lesson 1 -

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# Overview

...of things you should definitely know about if you want a very good grade

- Definition of Algorithms & Data Structures
- Arrays
- Linked Lists
- Stacks
- Queues
- Graphs
- Trees
- Complexity
- Sorting ← Let's talk about this next time.
- ...

# Exact definitions are important!

Algorithms

&

Data Structures

"A finite sequence of well defined instructions which, given an input, will produce in a finite time a required output"

"Structured memory space and a set of operations, or actions, for accessing and manipulating such memory space"

# Exact definitions are important!

## Algorithms

- Finitely many Instructions but also finitely many steps!

"A **finite** sequence of well defined instructions which, given an input, will produce in a **finite time** a required output"

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## Algorithms

- Finitely many Instructions but also finitely many steps
- No ambiguous sequences

"A **finite** sequence of **well defined** instructions which, given an input, will produce in a **finite time** a required output"

# Exact definitions are important!

## Algorithms

- Finitely many Instructions but also finitely many steps

- No ambiguous sequences

- Input and Output

"A **finite** sequence of **well defined** instructions which, given an **input**, will produce in a **finite time** a required **output**"

# Exact definitions are important!

## Algorithms

"A **finite** sequence of **well defined** instructions which, given an **input**, will produce in a **finite time** a **required output**"

- **Finitely many Instructions but also finitely many steps**
- **No ambiguous sequences**
- **Input and Output**
- **Should be correct → does what we want**

# Exact definitions are important!

- Information which is  
located somewhere and  
connected somehow

## Data Structures

"Structured memory space and  
a set of operations,  
or actions, for accessing and  
manipulating such memory  
space"



# Exact definitions are important!

- Information which is located somewhere and connected somehow

- Operations to locate certain information to read / overwrite / delete

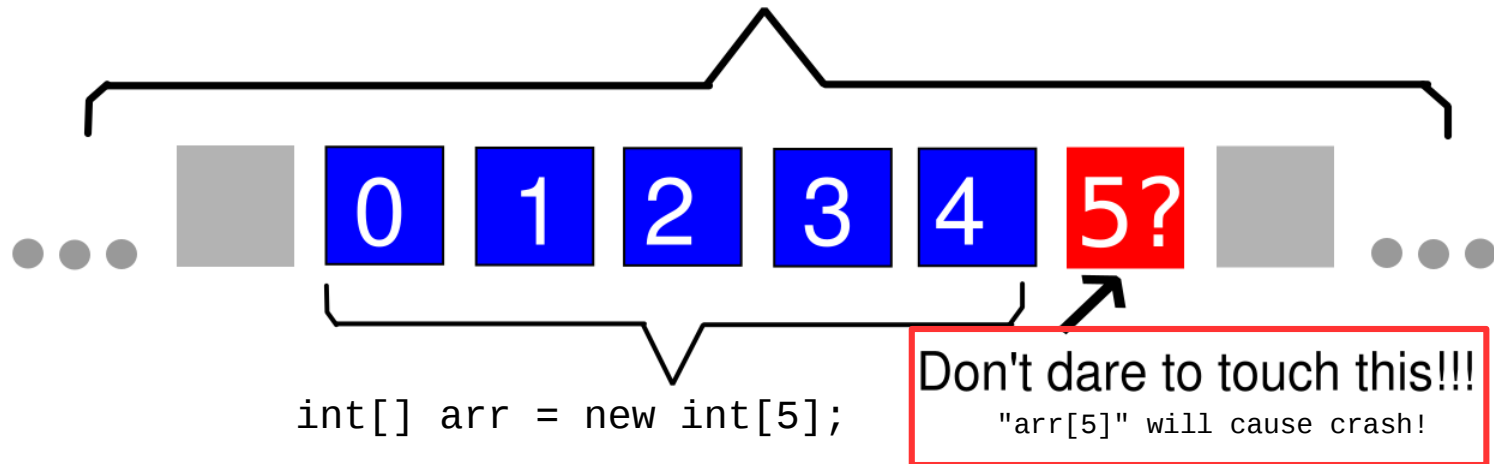
## Data Structures

"Structured memory space and a set of operations, or actions, for accessing and manipulating such memory space"

# Arrays

- Most simple structure but with little freedom
  - Static size once declared
  - Only one datatype

Memory cells in our Computer



# Linked Lists

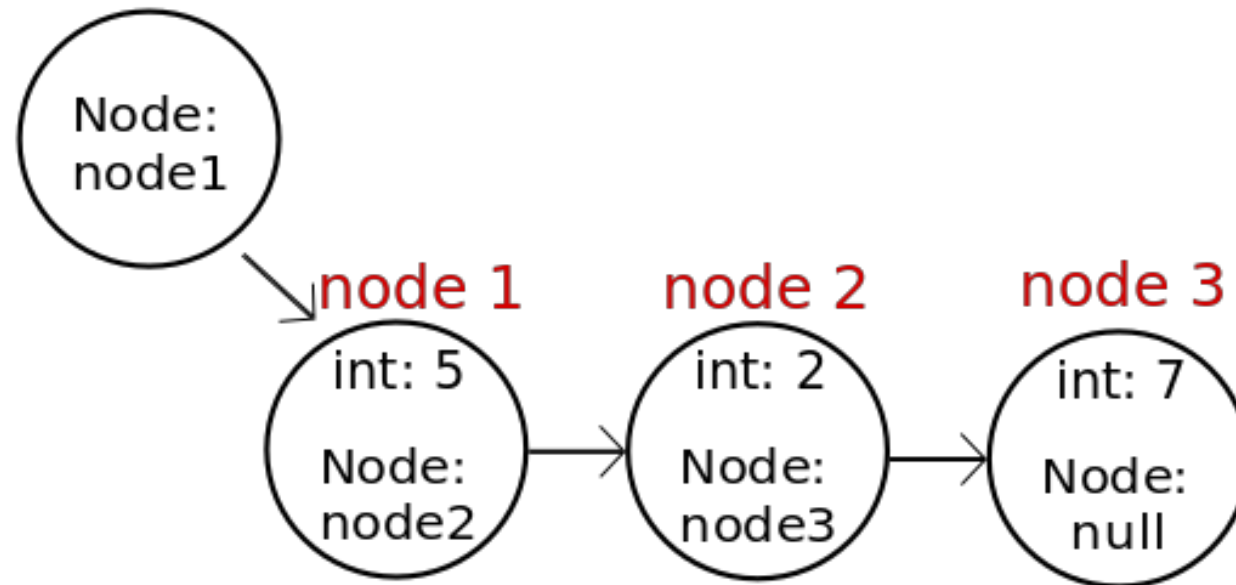
- Consist of "Nodes" which hold data + location of next Node

→ Can be anywhere in Memory

→ Nodes are easy to add / remove

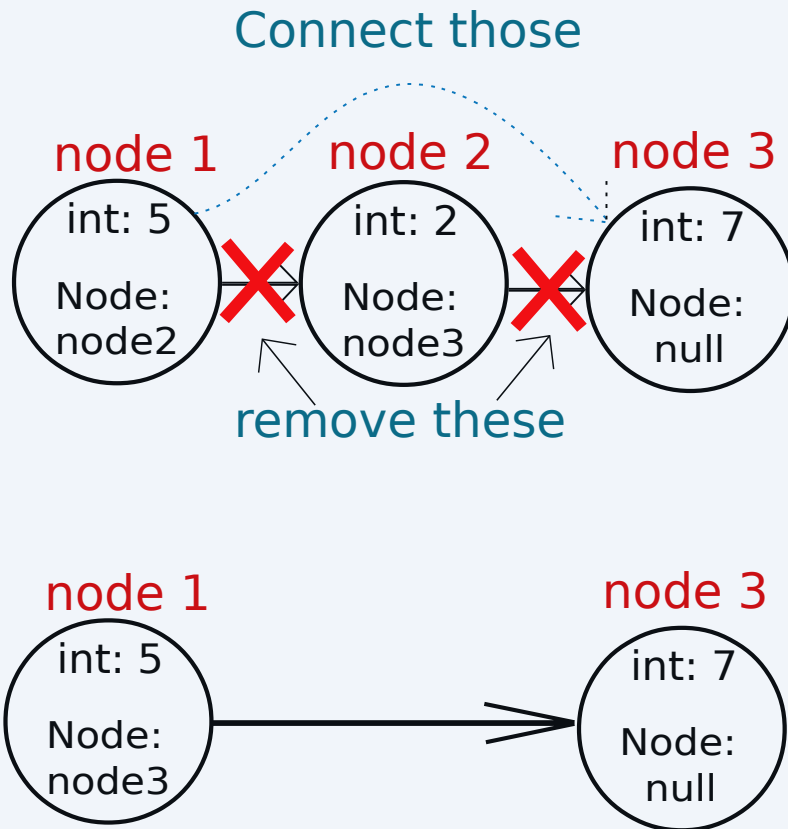
Compare: removing Element in Array → cell of unusable Memory,  
adding an Element in Array → need to reallocate Memory.

## List Object

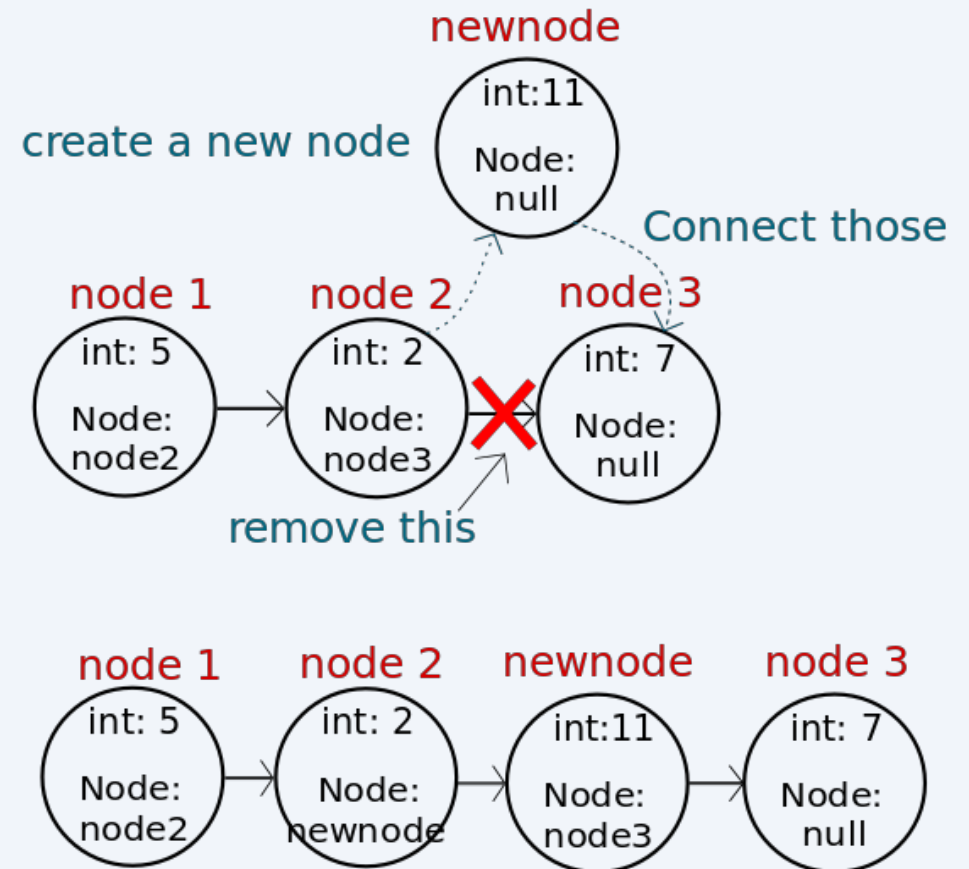


# Linked Lists

## Removing nodes:

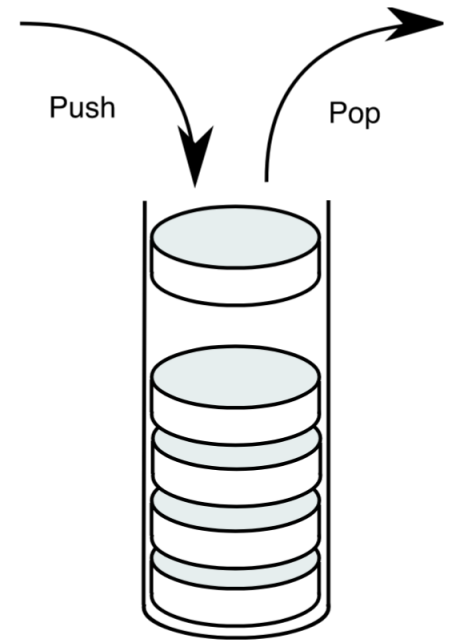


## Adding nodes:



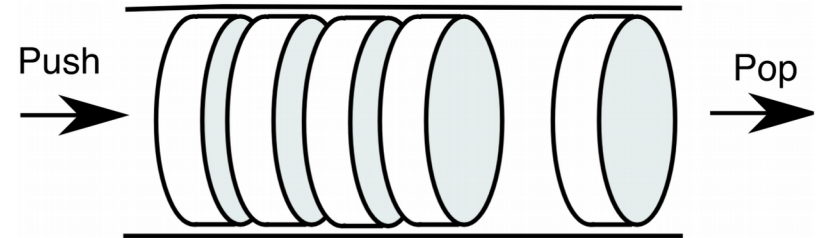
# Stacks

- Last – in – First – out (LiFo)
- Basically a List with constraints:
  - Only first element can be read
  - The same is true for adding and removing nodes
- But what can they be usefull for?



# Queues

- First – in – First – out (FiFo)
- Basically a List with constraints:
  - Only first element can be read and removed
  - Adding nodes only possible at the very end\*
- And what can they be usefull for?



\*Just in theory as you had to manipulate the last element to connect it to a new one

# Graphs

- Def: Ordered pair of two sets:

$$(V, E)$$

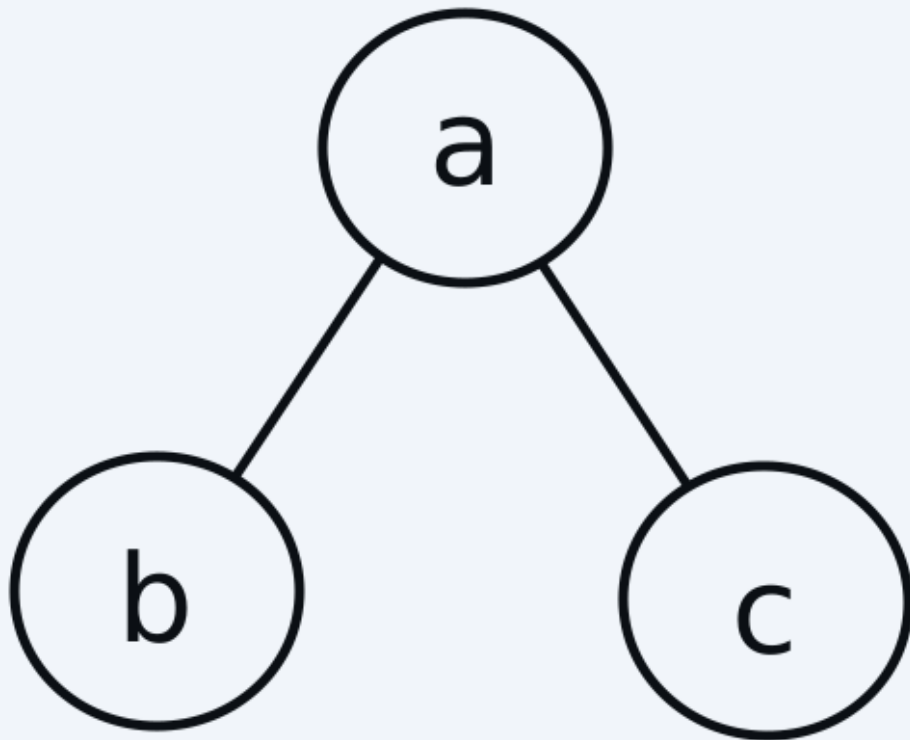
Where  $V$  is the set of Vertices while  $E$  denotes the set of Edges connecting two Vertices (In general).

- So what does such a Graph look like?

# Graphs

For example:  $G = ( \{a,b,c\} , \{(a,b) , (a,c)\} )$

Visual representation:



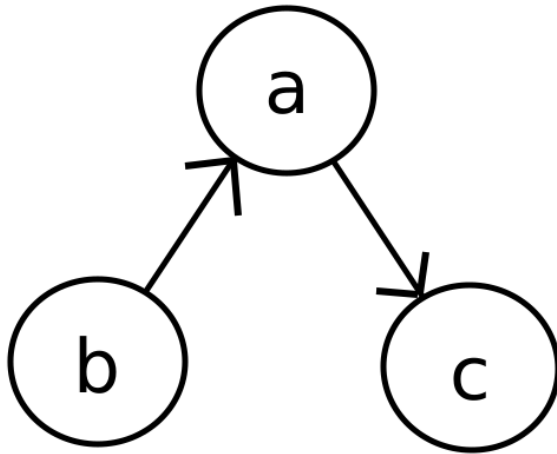
Adjacency Matrix:

	a	b	c
a	0	1	1
b	1	0	0
c	1	0	0



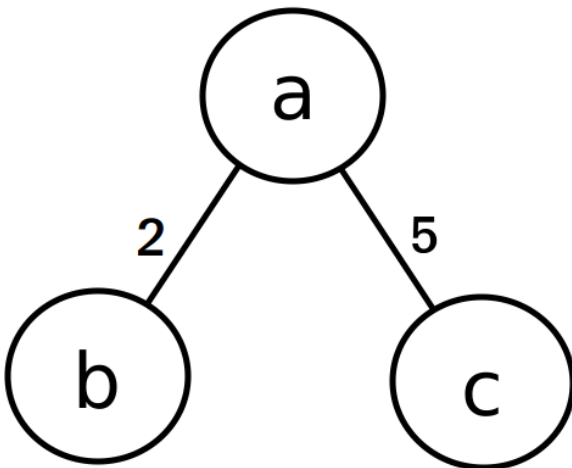
# Graphs

Directed Graph:



	a	b	c
a	0	0	1
b	1	0	0
c	0	0	0

Weighted Graph:

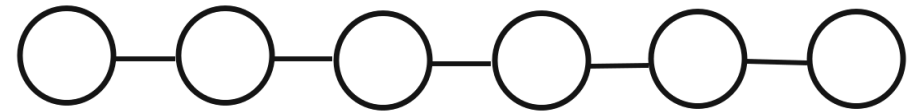


	a	b	c
a	0	2	0
b	2	0	5
c	0	5	0

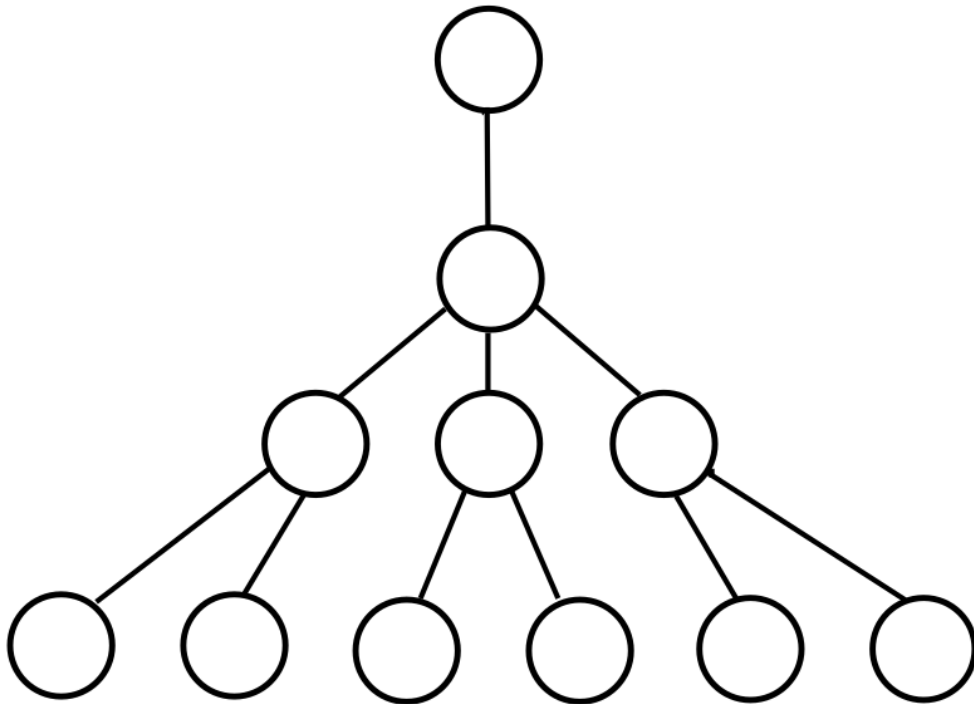
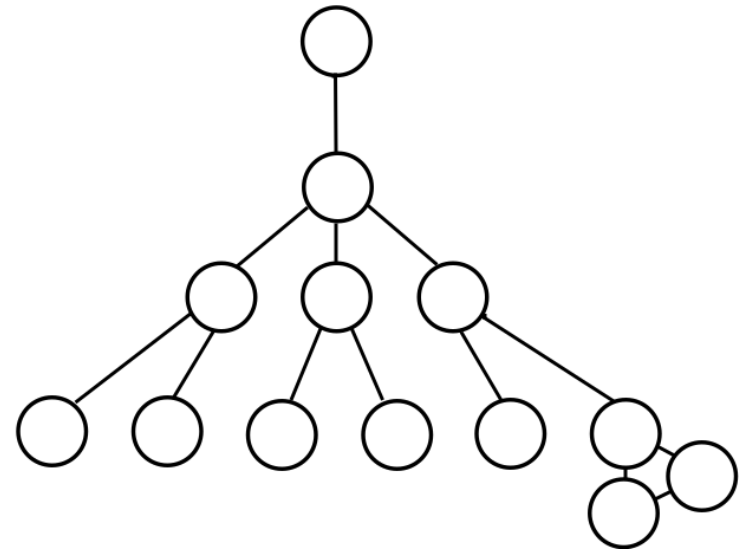
# Trees

Remember: Graphs without Circles

Still a tree by definition:

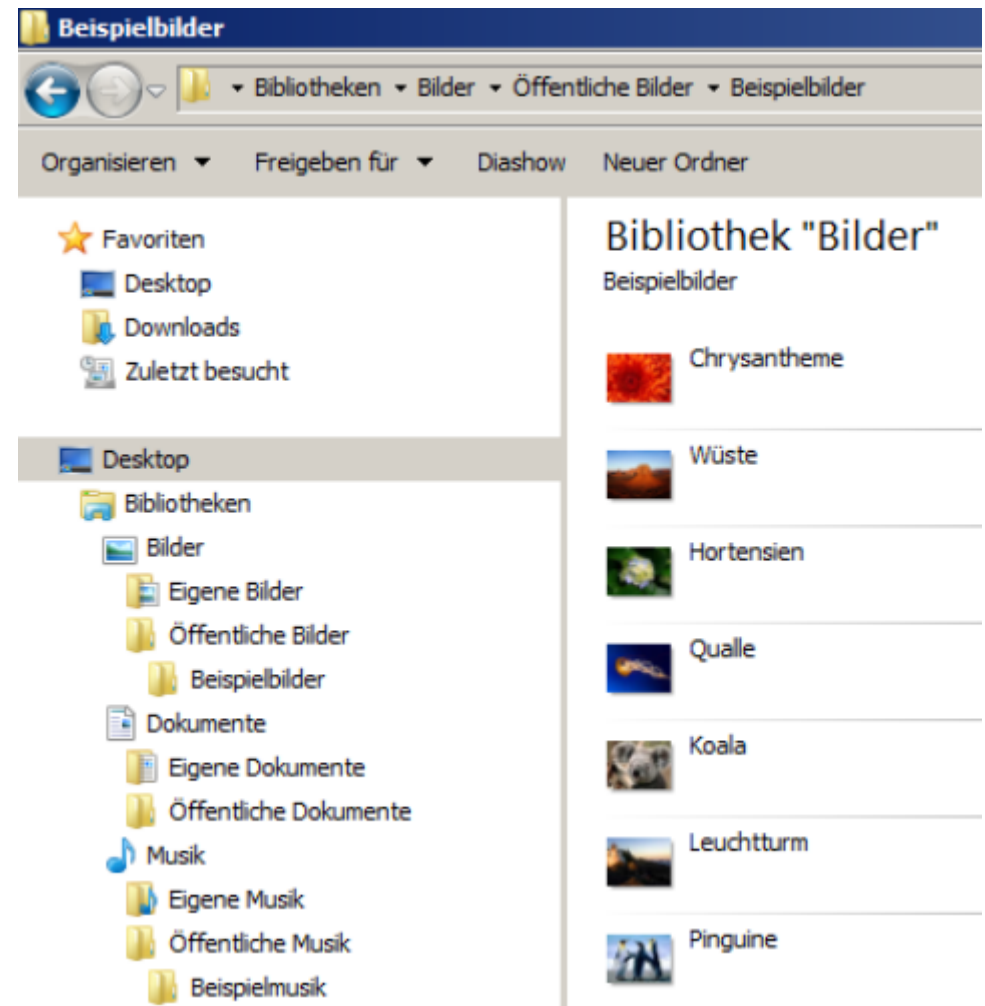
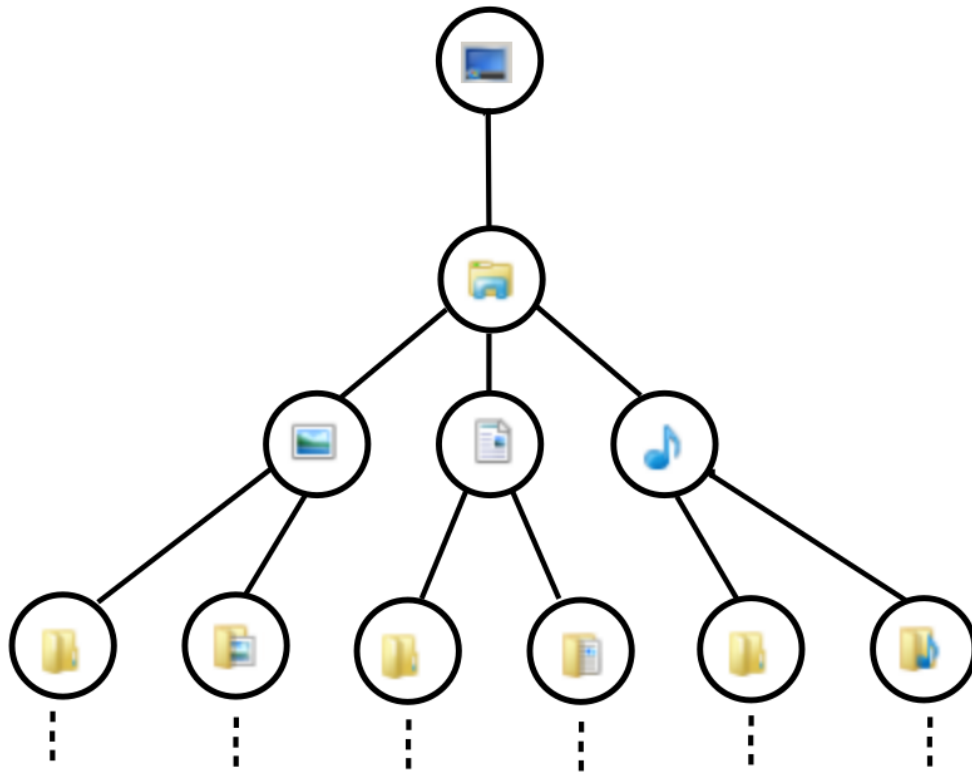


Not a tree anymore:



# Trees

## Practical example: Data in Operating Systems



# Trees

Based on our example:

What if you search for certain documents?

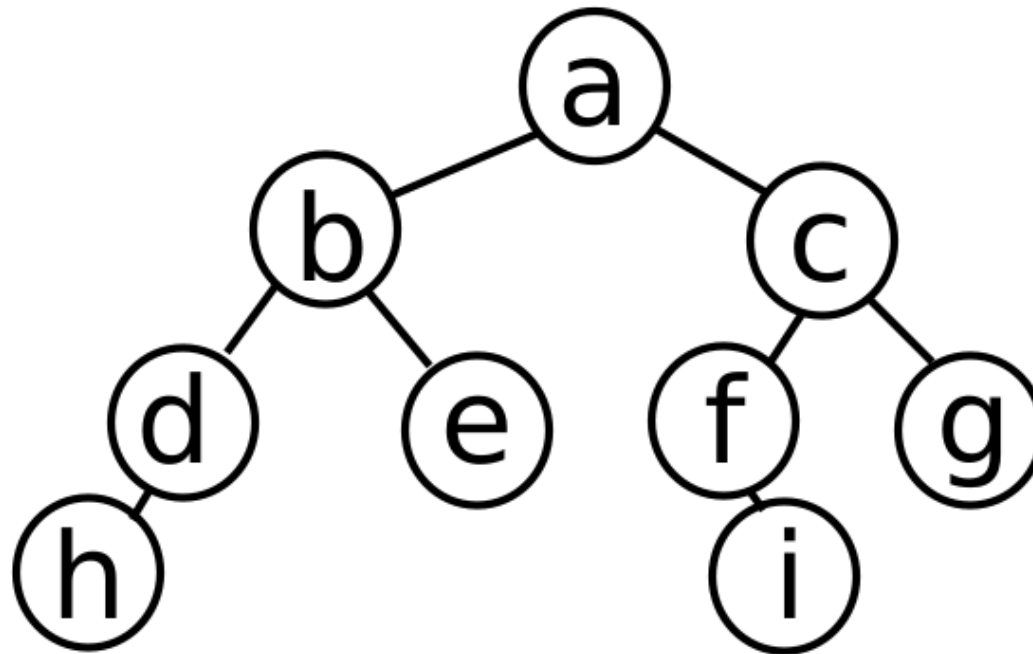
→ You need to *traverse* the tree:

- Pre – Order
- Post – Order
- In – Order
- Level –  
Order

Very good explanation if you need to repeat:

[https://en.wikipedia.org/wiki/Tree\\_traversal](https://en.wikipedia.org/wiki/Tree_traversal)

# Traversal of Trees



- Pre – Order      a, b, d, h, e, c, f, i, g
- Post – Order    h, d, e, b, i, f, g, c, a
- In – Order      h, d, b, e, a, f, i, c, g
- Level – Order   a, b, c, d, e, f, g, h, i

# Some use cases

List: Representing and adding polynomials

Stack: Recognizing palindromes

Queue: Orders from a fast Device to a slow one like a printer. → Data has to „wait“

Just think about some more!

# Assignment

## Huffman Code

- Lossless Data Compression
- Still relevant in later courses like Coding Theory
- Uses Trees
- Well documented on the web

# Assignment

By the way, how would you encode

Lossless? Data Compression

Still relevant in later courses like Coding Theory

Uses Trees

Well documented on the web



# Assignment

By the way, how would you encode  
Lossless?

$L \rightarrow 00, O \rightarrow 01, S \rightarrow 10, E \rightarrow 11$

00 01 10 10 00 11 10 10

16 Bit.

# Assignment

You can do it this way but you will

$L \rightarrow 00, O \rightarrow 01, S \rightarrow 10, E \rightarrow 11$

more space in Memory than necessary

# Assignment

Frequently represented letter gets shorter Code

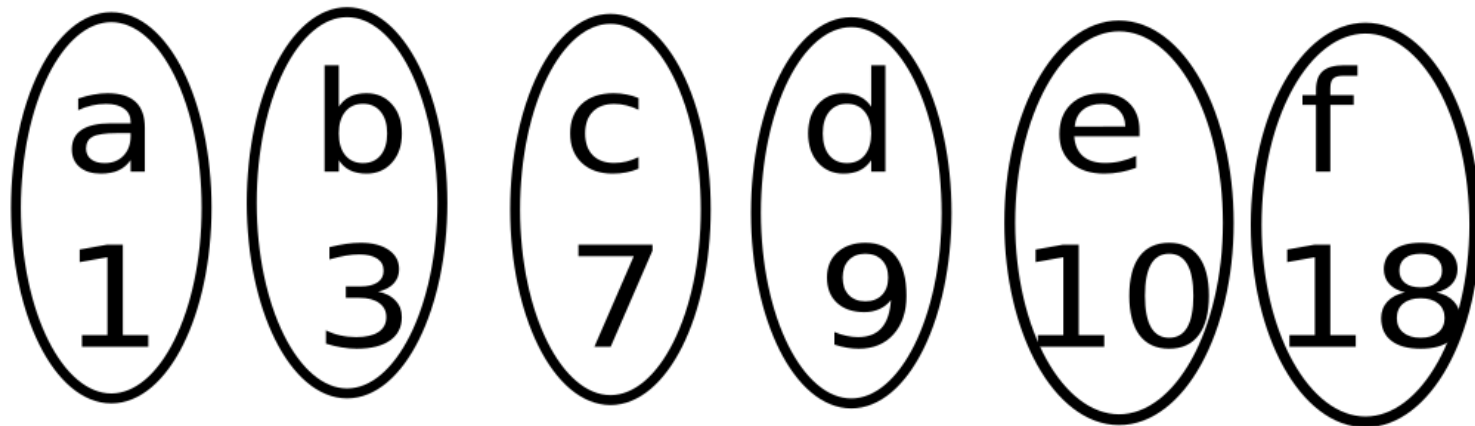
Letter	Frequency	Code
s	4	1
l	2	01
o	1	001
e	1	000

→ 0100 1110 1000 11 – only 14 Bit!

# Assignment

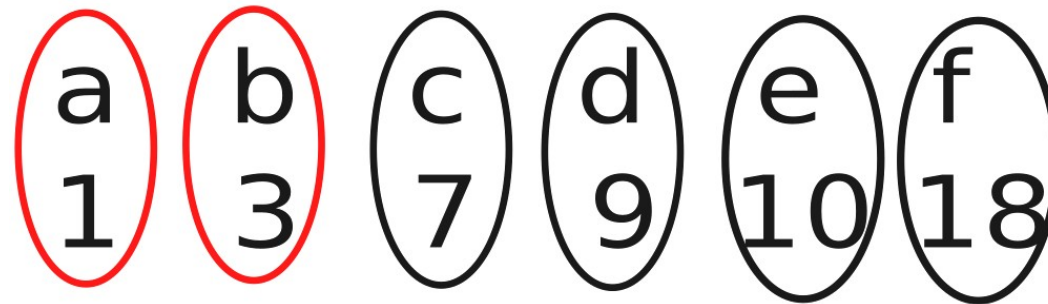
How do we achieve this Coding scheme?

- Input:
  - Letter
  - Corresponding Frequency
- Treat every node as a Subtree:

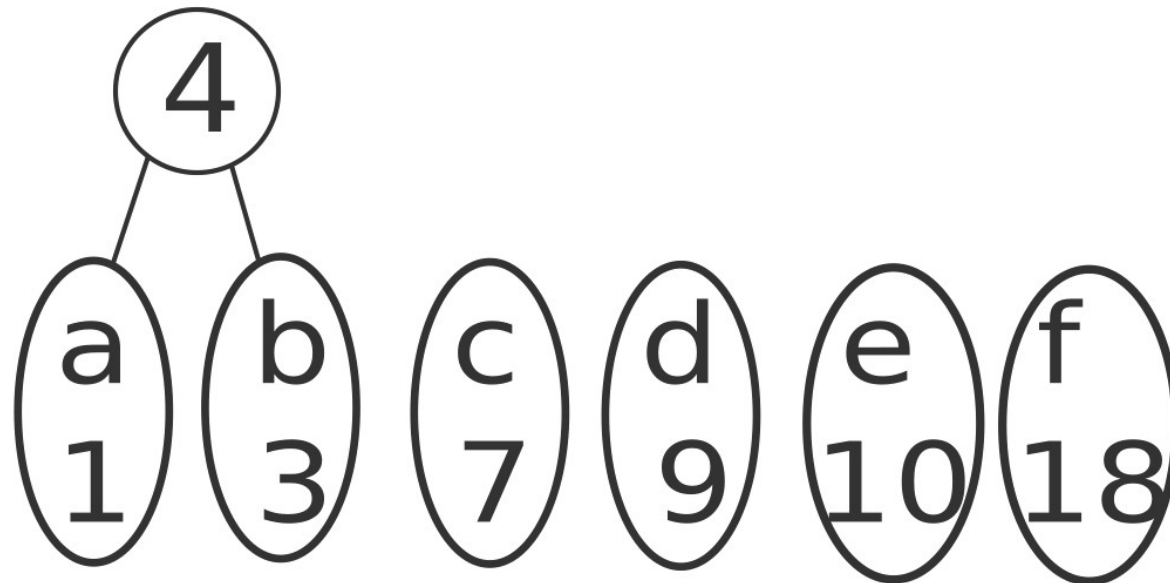


# Assignment

Search the two smallest roots

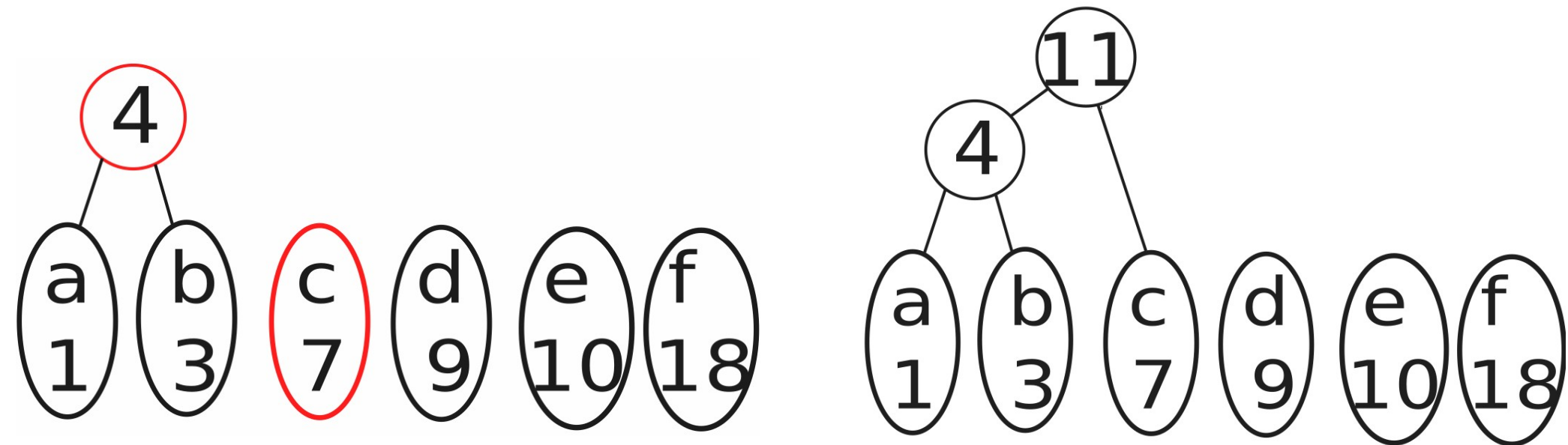


Connect to a new Tree and add frequencies



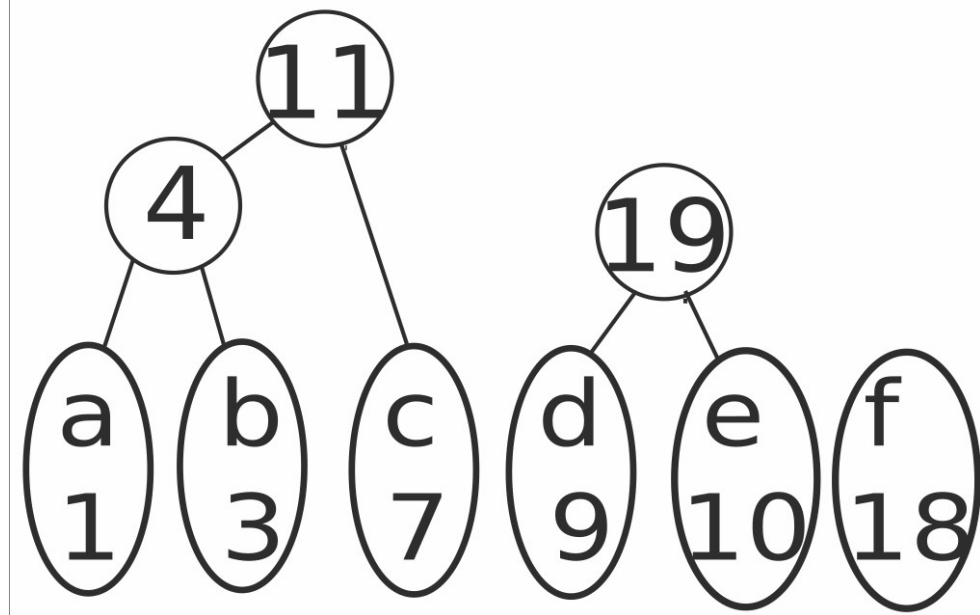
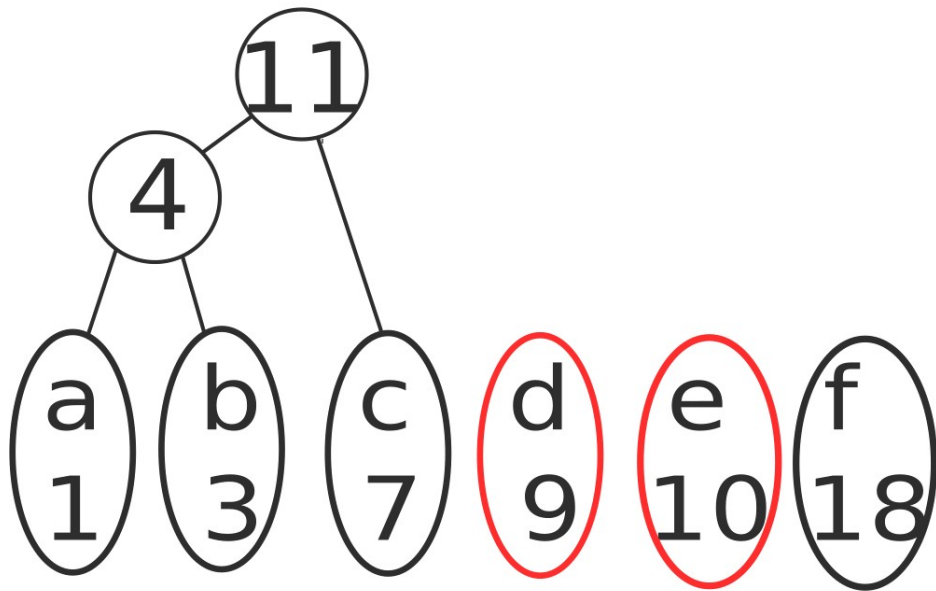
# Assignment

Do it rekursively on the new Data



# Assignment

And so on...



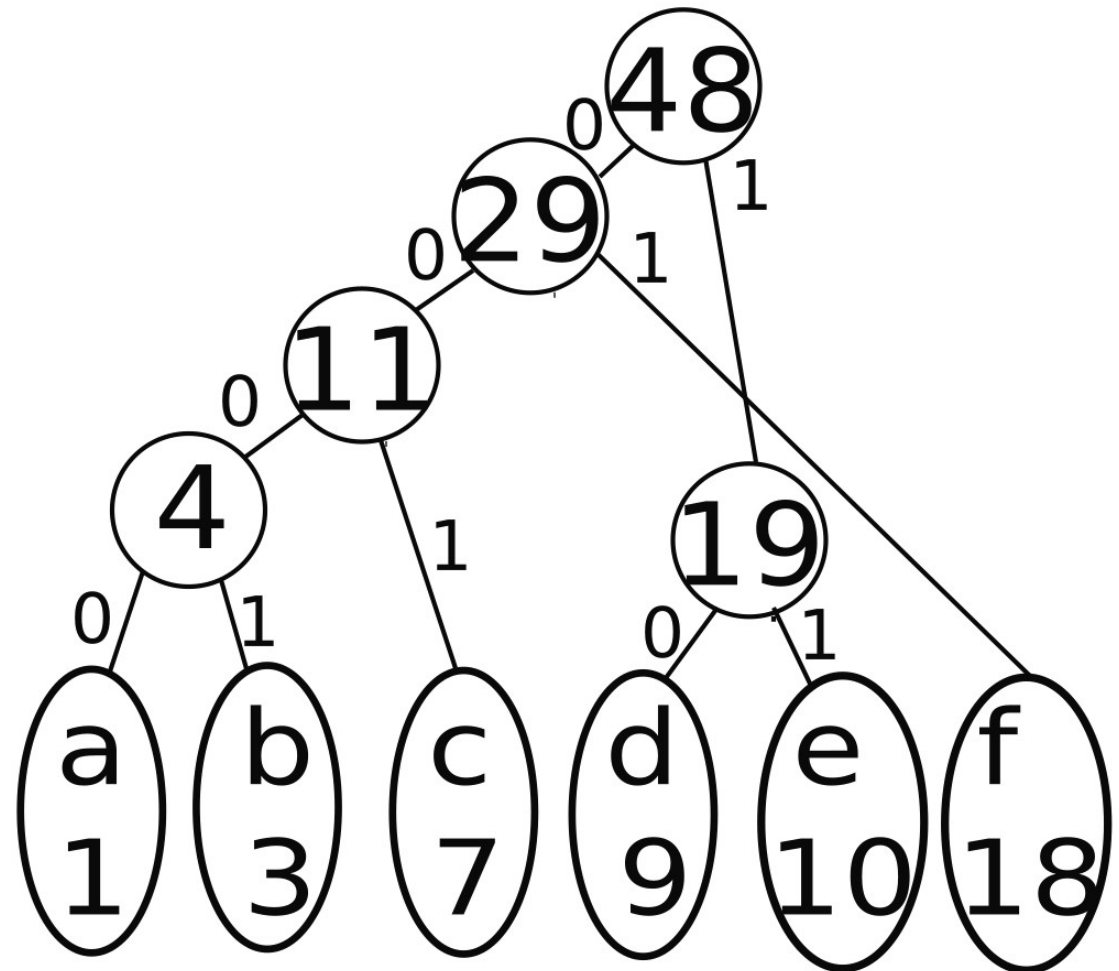
# Assignment

You stop if there is only one node left.

For every node which is not a Leaf:

- Right edge: 1
- Left edge: 0

Huffman -  
tree

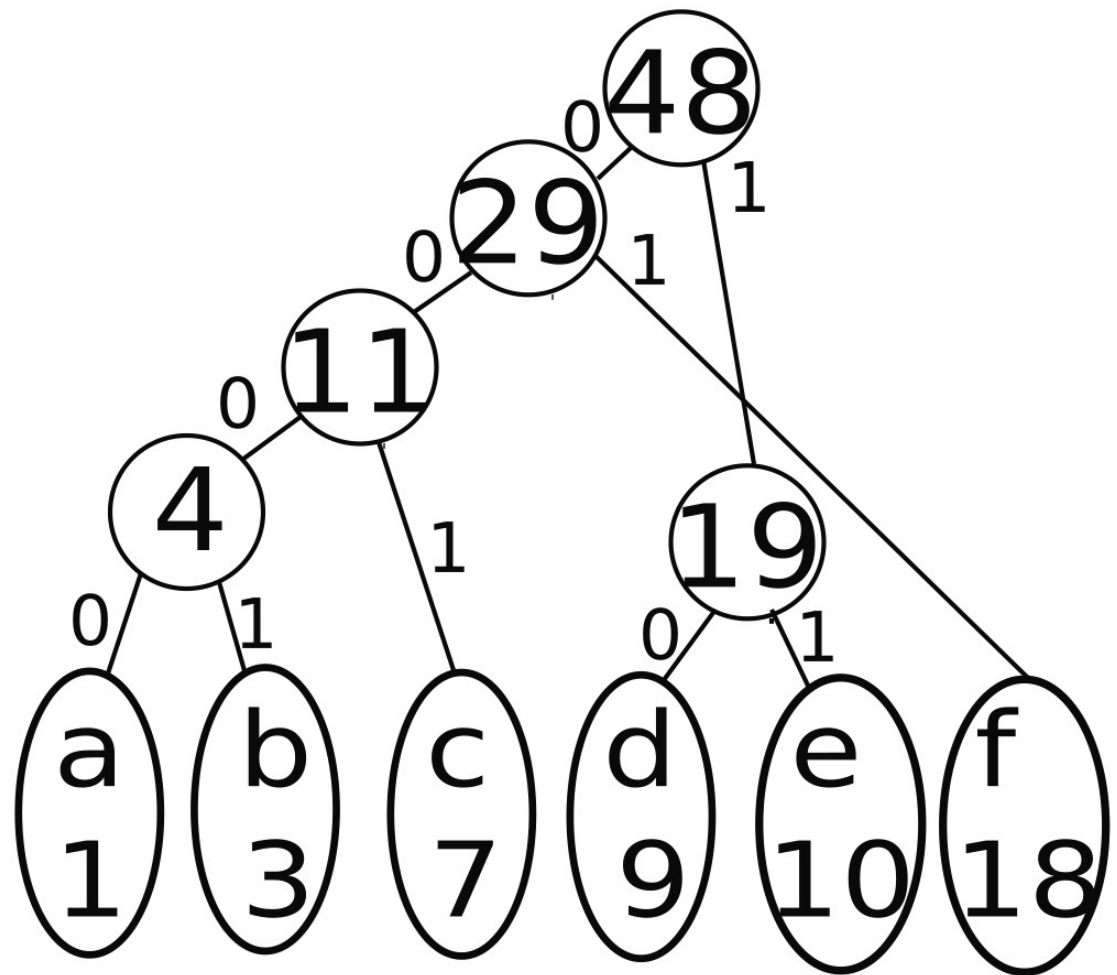




# Assignment

Following the Edges from Root to Leaf you can just read the coding scheme:

Letter	Code
a	0000
b	0001
c	001
d	10
e	11
f	01



# Assignment

Your job:

Write a program which takes the following frequencies as an input and gives out a coding scheme according to Huffman!

(Either take this one or the one shown in the Exercise. They are equally difficult)

Letter	a	b	c	d	e	f	g	h
Frequency	14	20	7	10	23	5	16	3

# Assignment

It will be alright to hard – code the Input data.  
Most important that the algorithm works like  
the shown way.

Output can look like this (for Example):

a - 11

b - 101

c - 100

d - 01

e - 001

f - 0001

g - 00001

h - 000000|

Don't need to represent the tree  
graphically. It's handled  
internally

# Assignment Requirements

- Use Java, C or C++
- Compilable source code (gcc or javac)
- No IDEs
- Readable source code
- Self-contained submission
- Sufficient comments (rule of thumb: another Developers can get program logic alone by reading the comments)

# Assignment Requirements

- 5 Assignments
- In each!  
Data: Name, matriculation number, SS2018
- max 6 point for each assignment
- 30% exercise and 70% final exam

# Assignment Requirements

- I suggest to write a „node“ - Class to build the tree.
- Again: Do not copy code. Not from the Internet. Not from your partner. We will find out.
- Deadline: 15. May 2018 23:59
- Mailto: [michael.schwarzkopf@uni-weimar.de](mailto:michael.schwarzkopf@uni-weimar.de)

Good luck!