

## Laboratory practice No. 4: Greedy algorithms

**Juan David Echeverri Villada**  
Universidad Eafit  
Medellín, Colombia  
jdecheverv@eafit.edu.co

**Juan Sebastian Guerra Hernandez**  
Universidad Eafit  
Medellín, Colombia  
jsguerrah@eafit.edu.co

### 3) Practice for final project defense presentation

- 3.1
- 3.2
- 3.3
- 3.4
- 3.5
- 3.6

### 4) Practice for midterms

4.1  $i = j$

4.2  $min > adjacencyMatrix[element][i]$

4.3

4.3.1

Step	A	B	C	D	E	F	G	H
1	A	20, A	inf	80, A	inf	inf	90, A	inf
2	B	20, A	inf	80, A	inf	30, B	90, A	inf
3	F	20, A	40, F	70, F	inf	30, B	90, A	inf
4	C	20, A	40, F	50, C	inf	30, B	90, A	60, C
5	D	20, A	40, F	50, C	inf	30, B	70, D	60, C
6	H	20, A	40, F	50, C	inf	30, B	70, D	60, C
7	G	20, A	40, F	50, C	inf	30, B	70, D	60, C
8	E	20, A	40, F	50, C	inf	30, B	70, D	60, C

4.3.2 The best route to go from A to G is: A – B – F – D – G.

4.4

4.4.1 Line 10:  $temp/2$

4.4.2 Line 11:  $temp + minimo$

4.4.3 B

4.5

4.5.1 D

**PhD. Mauricio Toro Bermúdez**

Professor | School of Engineering | Informatics and Systems

Email: mtorobe@eafit.edu.co | Office: Building 19 – 627

Phone: (+57) (4) 261 95 00 Ext. 9473

**ESTRUCTURA DE DATOS 2**  
**Código ST0247**

4.5.2 *By making mergeSort to the set of  $n$  numbers, we can ensure that the smallest are in the first positions. The sum of the first  $k$  numbers, arranged from least to greatest, will always be the minimum sum with  $k$  numbers. The complexity of making mergeSort is  $O(n \log n)$ , and the access to the  $k$  numbers is  $O(k)$ , but by multiplicity the final complexity of the algorithm is  $O(n \log n)$ .*

**4.6 A**

4.6.1  $i+1$

4.6.2  $res+1$

4.6.3  $last = i$

4.6.4 2

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