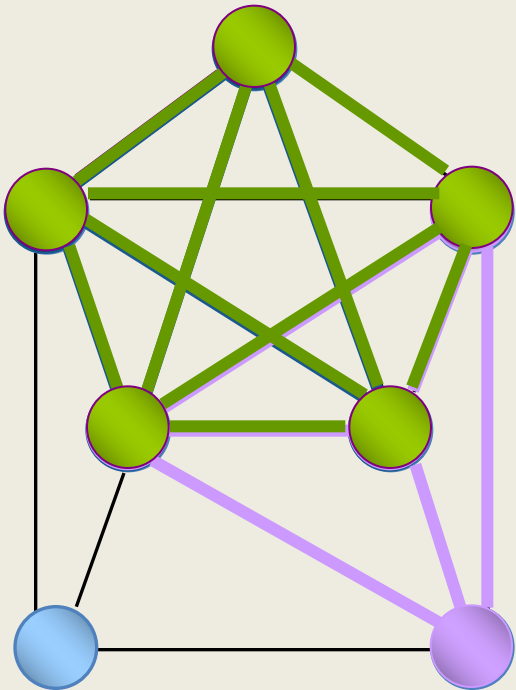




A Heuristic Algorithm for the K- Clique Problem

Problem Outline

$$CLIQUE = \{ \langle G, k \rangle \mid G \text{ has a clique of size } k \}$$



Maximum Clique of Size 5

Clique

Graph $G = (V, E)$, a subset S of the vertices is a clique if there is an edge between every pair of vertices in S , it also means a subset of vertices in V all connected to each other by edges in E

Size of Clique:

number of vertices it contains

Maximal Clique:

a *clique* cannot be enlarged by adding any more vertices

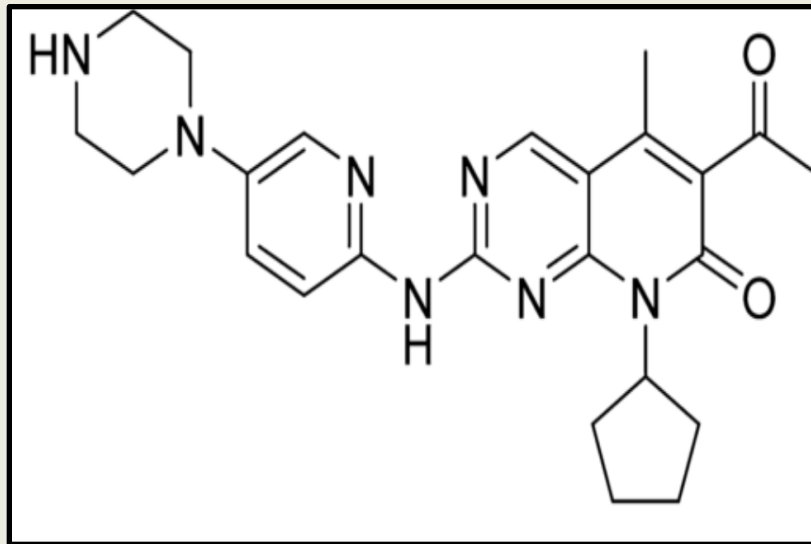
Maximum Clique

the largest *maximal clique* in the graph

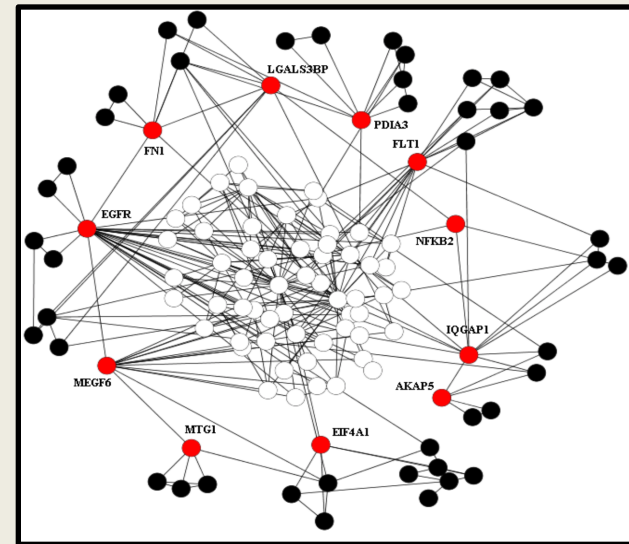
Why is this problem important?



Social Network

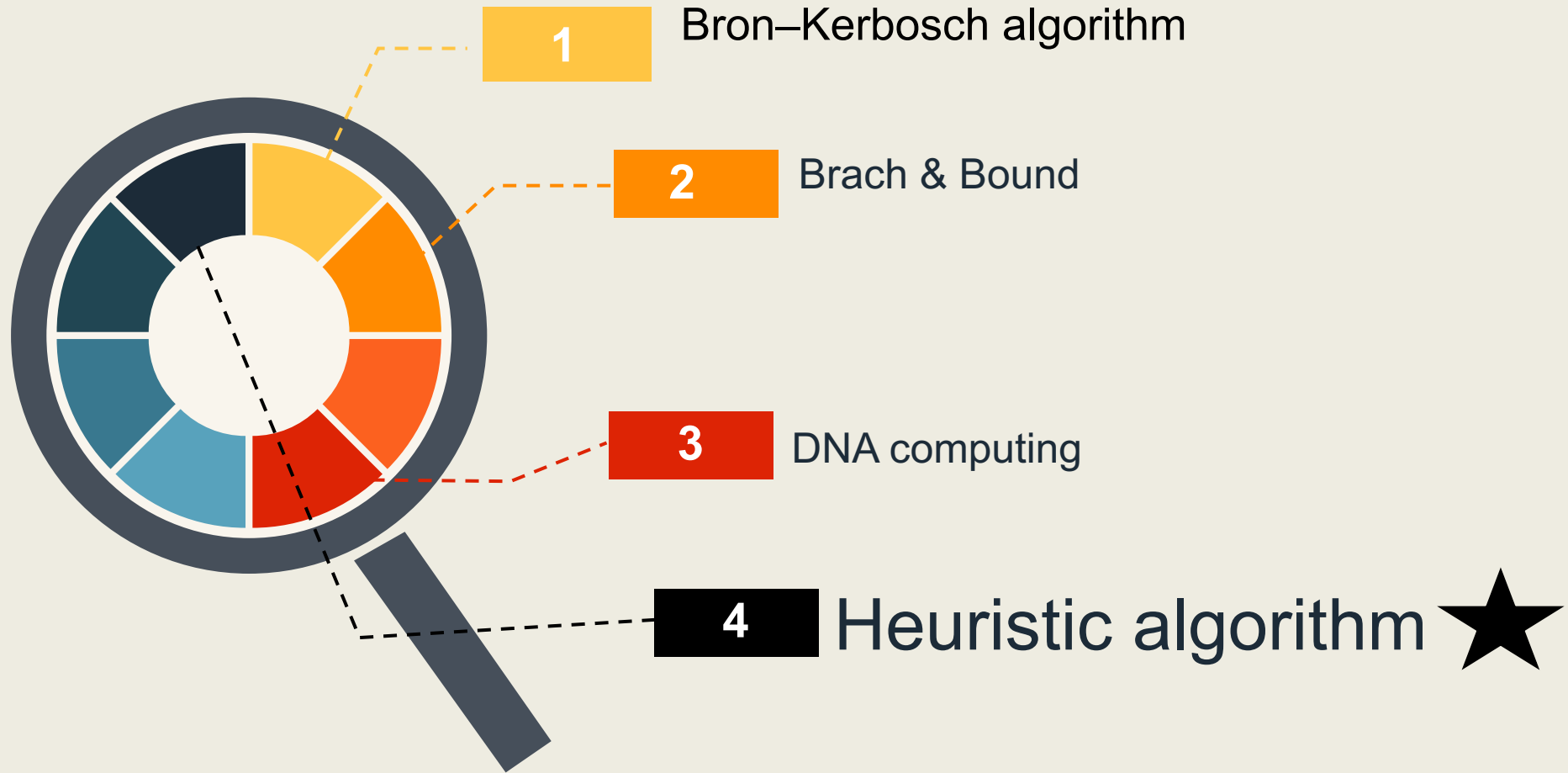


Computational chemistry



Bio-information

Solution Overview



What is a Heuristic Search?

- A Search Technique that employs a rule that increases the likelihood of finding a solution.
- Domain specific knowledge must be added to improve efficiency.
- **In Our Case:**
 - Node Degree
 - Position relative to other nodes



Solution Details

Test Case Generation

- Generate a random graph of N vertices with a density D .
- Set target clique size k

Preparation

- Remove all vertices with degree $< k$
- Keep doing this while you can still delete vertices.

Initialization

- For v in V
- $v.location = i, i++$

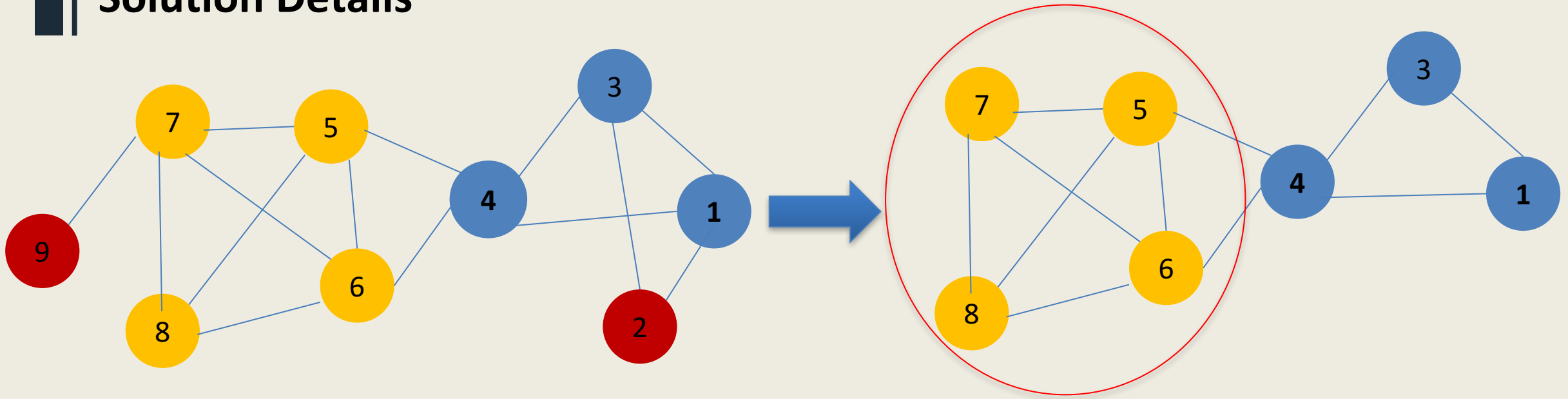
Movement

- For each edge (u,v)
- Move the vertices to the midpoint between them.

Search

- For v in V
- Draw a circle of “unit” distance from v
- Enough points in v ?
- All connected?
- Return true
- Return false.

Solution Details



In a clique of size k , each node maintains degree $\geq k-1$
Nodes with degree $< k-1$ will not be included in the maximum clique

In order to find a clique > 3 , remove all nodes with degree $\leq 3-1=2$

Remove nodes 2 and 9

Remove nodes 1 and 3

Remove node 4

Empirical Results

Algorithm

Inputs

```
package project4;

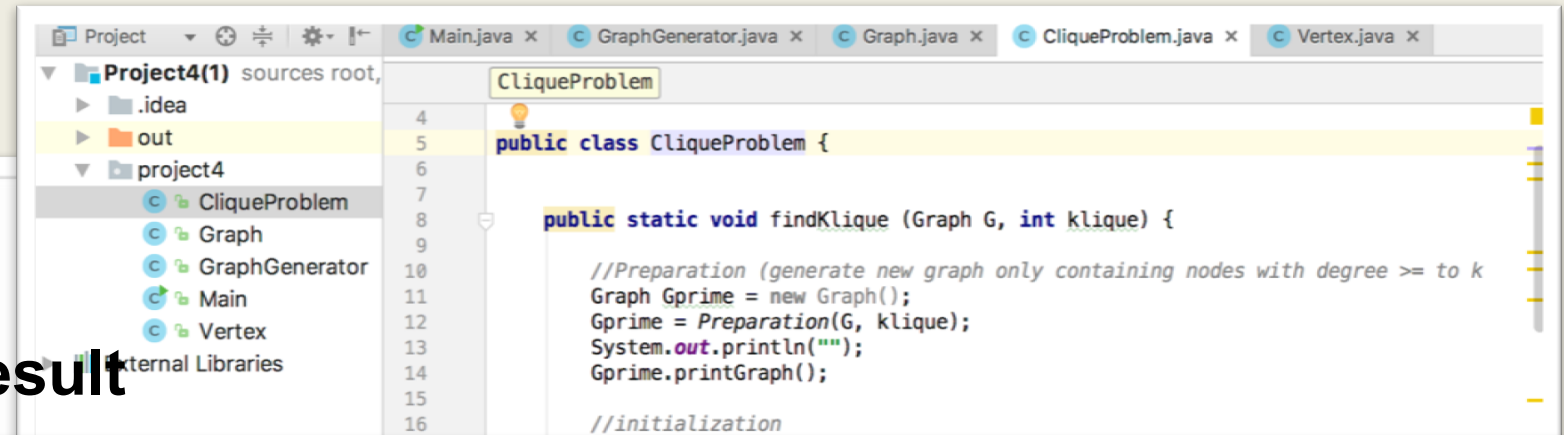
public class Main {

    public static void main(String[] args) {
        // TODO Auto-generated method stub

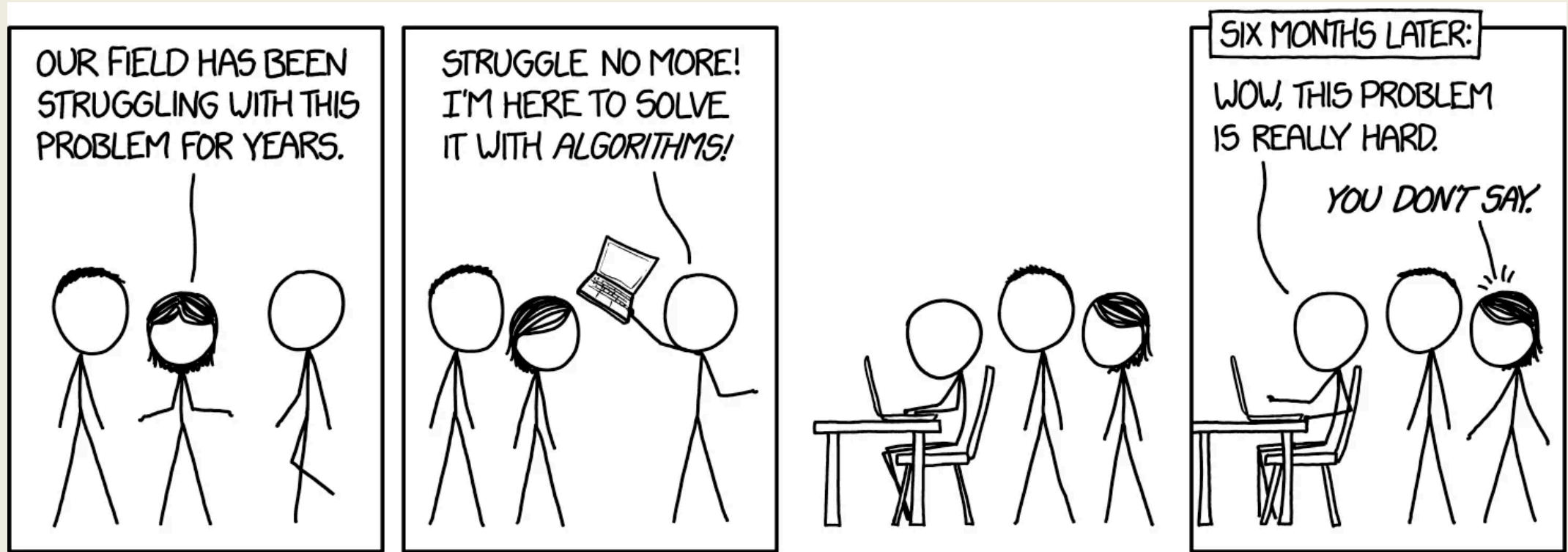
        int n = 1000;
        int k = 100;
        int density = 100000;
        Graph G = new Graph();
        G = GraphGenerator.testCaseGen(n);
        G.printGraph();
    }
}
```

Result

```
179, 2, 013, 010, 017, 019, 021, 022, 024, 020, 030, 037, 003, 040, 007, 000, 032, 033, 010, 039, 000, 001, 022, 070,
Clique of size 100 found at Vertex #79
[187, 0, 4, 6, 804, 819, 821, 822, 824, 826, 830, 837, 605, 848, 607, 608, 852, 853, 616, 859, 860, 861, 622, 870, 875,
Clique of size 100 found at Vertex #187
[82, 3, 7, 816, 817, 819, 821, 822, 824, 826, 830, 837, 605, 848, 607, 608, 852, 853, 616, 859, 860, 861, 622, 870, 875,
Clique of size 100 found at Vertex #82
[88, 0, 4, 6, 804, 819, 821, 822, 824, 826, 830, 837, 605, 848, 607, 608, 852, 853, 616, 859, 860, 861, 622, 870, 875,
Clique of size 100 found at Vertex #88
[192, 0, 4, 6, 804, 819, 821, 822, 824, 826, 830, 837, 605, 848, 607, 608, 852, 853, 616, 859, 860, 861, 622, 870, 875,
Clique of size 100 found at Vertex #192
[198, 3, 7, 816, 817, 819, 821, 822, 824, 826, 830, 837, 605, 848, 607, 608, 852, 853, 616, 859, 860, 861, 622, 870, 875,
Clique of size 100 found at Vertex #198
[93, 8, 804, 819, 821, 822, 824, 826, 830, 837, 605, 848, 607, 608, 852, 853, 616, 859, 860, 861, 622, 870, 875, 880,
Clique of size 100 found at Vertex #93
[97, 6, 804, 819, 821, 822, 824, 826, 830, 837, 605, 848, 607, 608, 852, 853, 616, 859, 860, 861, 622, 870, 875, 880,
Clique of size 100 found at Vertex #97
[99, 8, 804, 819, 821, 822, 824, 826, 830, 837, 605, 848, 607, 608, 852, 853, 616, 859, 860, 861, 622, 870, 875, 880,
Clique of size 100 found at Vertex #99
Finished
```



Limitations



References/Citations

1

https://en.wikipedia.org/wiki/Bron-Kerbosch_algorithm

2

https://en.wikipedia.org/wiki/Clique_problem

3

Rossi, R. A., Gleich, D. F., Gebremedhin, A. H., Patwary, M. M. A., & Ali, M. (2013). A fast parallel maximum clique algorithm for large sparse graphs and temporal strong components. *CoRR*, *abs/1302.6256*.

4

Tomita, E., Tanaka, A., & Takahashi, H. (2006). The worst-case time complexity for generating all maximal cliques and computational experiments. *Theoretical Computer Science*, *363*(1), 28-42.