

# MATH 5312 Advanced Numerical Methods II

## Final Project

Due date: 31 May, Wednesday

Answer all questions with reasoning.

1. Consider the finite difference discretization of the equation

$$\begin{cases} -(a(x)u'(x))' = f(x), & \text{for } x \in (0, 1) \\ u(0) = u(1) = 0. \end{cases}$$

The discretized linear system is  $\mathbf{A}\mathbf{u} = \mathbf{f}$ , where

$$\mathbf{A} = \begin{bmatrix} a_0 + a_1 & -a_1 & & & & & \\ -a_1 & a_1 + a_2 & -a_2 & & & & \\ & -a_2 & a_2 + a_3 & -a_3 & & & \\ & & \ddots & \ddots & \ddots & & \\ & & & -a_{n-2} & a_{n-2} + a_{n-1} & -a_{n-1} & \\ & & & & -a_{n-1} & a_{n-1} + a_n & \end{bmatrix}$$

and  $a_0, a_1, \dots, a_n$  are  $a(x)$  on the grid points. Assume  $C_1 \leq a(x) \leq C_2$  for all  $x \in [0, 1]$ , where  $C_1$  and  $C_2$  are positive constants.

- (a) Prove that  $\mathbf{A}$  is symmetric positive definite.
- (b) Show that both Jacobi and Gauss-Seidel converges for solving  $\mathbf{A}\mathbf{u} = \mathbf{f}$ .
- (c) Prove that

$$4C_1 \sin^2\left(\frac{\pi}{2(n+1)}\right) \leq \lambda_1 \leq \lambda_n \leq 4C_2,$$

where  $\lambda_1$  and  $\lambda_n$  are the smallest and largest eigenvalues of  $\mathbf{A}$  respectively.

- (d) Since  $\mathbf{A}$  is SPD, we may use the preconditioned conjugate gradient (PCG) to solve  $\mathbf{A}\mathbf{u} = \mathbf{f}$ . A candidate preconditioner will be

$$\mathbf{P} = \begin{bmatrix} 2 & -1 & & & \\ -1 & 2 & -1 & & \\ & \ddots & \ddots & \ddots & \\ & & -1 & 2 & -1 \\ & & & -1 & 2 \end{bmatrix}$$

Estimate the number of iterations needed for a solution with  $\epsilon$  precision. Your answer should be as tight as possible. (*This preconditioner will be practically useful for 2D case, because  $\mathbf{P}$  is diagonalizable by discrete sine transform and inverted very efficiently by fast Fourier transform.*)

(a)  $\forall x = (x_1, \dots, x_n)^T \in \mathbb{R}^n$ , we have

$$x^T A x = a_0 x_1^2 + \sum_{i=1}^{n-1} a_i (x_{i+1} - x_i)^2 + a_n x_n^2$$

Note that  $C_1 \leq a(x) \leq C_2$

then  $a_i \geq 0$  for  $i=0, 1, \dots, n \Rightarrow x^T A x \geq 0$

If  $x^T A x = 0$ .

$$\text{Then } x_i = \begin{cases} 0 & i=1, \dots, n \\ x_{i-1} & i=2, \dots, n-1 \end{cases}$$

$\Rightarrow x$  is a zero vector

$\Rightarrow x^T A x = 0$  iff  $x=0$

$\Rightarrow A$  is symmetric positive definite

(b) • Jacobi:

iterative matrix  $G = I - D^{-1}A$

$$\text{Note that } 2D - A = \begin{pmatrix} a_0 + a_1 & a_1 & & & \\ a_1 & a_1 + a_2 & a_2 & & \\ & a_2 & a_2 + a_3 & a_3 & \\ & & \ddots & \ddots & \ddots \\ & & & a_{n-1} & a_{n-1} + a_n \end{pmatrix}$$

$\forall x \in \mathbb{R}^n$ .

$$x^T (2D - A)x = a_0 x_1^2 + \sum_{i=1}^{n-1} a_i (x_i + x_{i+1})^2 + a_n x_n^2 \geq 0$$

$$x^T (2D - A)x = 0 \Leftrightarrow x = 0$$

$\Rightarrow 2D - A$  is SPD

$$\Rightarrow G = I - D^{-1}A, \rho(G) < 1$$

• Gauss-Seidel:

iterative matrix  $G = I - (D - E)^{-1}A$

Assume  $\lambda$  is the eigenvalue of  $G$ , and  $z$  is the corresponding eigenvector

We have  $Gz = \lambda z$

$$\Rightarrow (I - (D - E)^{-1}A)z = \lambda z$$

$$\Rightarrow (D - E)^{-1}Az = (1 - \lambda)z$$

$$\Rightarrow \lambda(D-E)z = E^T z$$

$$\Rightarrow \lambda z^*(D-E)z = z^* E^T z$$

$$\text{Let } z^* E z = \alpha + i\beta, \quad z^* D z = \delta \quad \alpha, \beta, \delta \in \mathbb{R}$$

$$\Rightarrow z^* E z = \alpha - i\beta$$

$$A \text{ is SPD} \Rightarrow D \text{ is SPD} \Rightarrow \delta = z^* D z > 0$$

$$\Rightarrow \lambda(\delta - (\alpha + i\beta)) = \alpha - i\beta$$

$$\Rightarrow \lambda = \frac{\alpha - i\beta}{(\delta - \alpha) - i\beta}$$

$$\Rightarrow |\lambda|^2 = \frac{\alpha^2 + \beta^2}{(\delta - \alpha)^2 + \beta^2}$$

$$A \text{ is SPD} \Rightarrow z^* A z > 0$$

$$z^* A z = z^* (D - E - E^T) z = \delta - 2\alpha$$

$$\Rightarrow \delta > 2\alpha$$

$$\Rightarrow (\delta - \alpha)^2 + \beta^2 > \alpha^2 + \beta^2$$

$$\Rightarrow |\lambda|^2 < 1$$

$$\Rightarrow \rho(G) < 1$$

(c) For a unit vector  $x \in \mathbb{R}^n$ ,  $x = (x_1, x_2, \dots, x_n)^T$ ,  $\|x\|_2 = 1$

$$x^T A x = a_0 x_1^2 + \sum_{i=1}^{n-1} a_i (x_{i+1} - x_i)^2 + a_n x_n^2$$

Note that  $C_1 \leq a_i \leq C_2$  for  $i = 0, 1, \dots, n$

$$C_1 (x_1^2 + \sum_{i=1}^{n-1} (x_{i+1} - x_i)^2 + x_n^2) \leq x^T A x \leq C_2 (x_1^2 + \sum_{i=1}^{n-1} (x_{i+1} - x_i)^2 + x_n^2)$$

$$\text{Denote } L = \begin{pmatrix} 2 & -1 & & \\ -1 & 2 & & \\ & & \ddots & \\ & & & 2 & -1 \\ & & & -1 & 2 \end{pmatrix}$$

$L$  is an 1-D Laplacian matrix

$$\text{Since } \lambda_k(L) = 2(1 - \cos \frac{k}{n+1} \pi) \quad k=1, 2, \dots, n$$

$$2(1 - \cos \frac{\pi}{n+1}) \leq x^T L x = x_1^2 + \sum_{i=1}^{n-1} (x_{i+1} - x_i)^2 + x_n^2 \leq 4$$

$$\text{Also, left hand side } 2(1 - \cos \frac{\pi}{n+1}) = 2 \cdot 2 \sin^2 \frac{\pi}{2(n+1)} = 4 \sin^2 \frac{\pi}{2(n+1)}$$

$$\text{and } C_1 \lambda_{\min}(L) \leq x^T A x \leq C_2 \lambda_{\max}(L)$$

We have

$$4 \sin^2 \frac{\pi}{2(n+1)} C_1 \leq x^T A x \leq 4 C_2$$

$$\Rightarrow 4 C_1 \sin^2 \frac{\pi}{2(n+1)} \leq \lambda_1 \leq \dots \leq \lambda_n \leq 4 C_2$$

$$(d) \quad P = \alpha \begin{pmatrix} 2 & -1 & & & \\ -1 & 2 & -1 & & \\ & \ddots & \ddots & \ddots & \\ & & -1 & 2 & -1 \\ & & & -1 & 2 \end{pmatrix}$$

$$\lambda_{\max}(P^T A) \leq \lambda_{\max}(P^T) \cdot \lambda_{\max}(A) \leq \frac{C_2}{\alpha} \frac{\lambda_{\max}(L)}{\lambda_{\min}(L)}$$

$$x^T A x = x^T P P^T A x$$

Denote  $\lambda_{\min}(P^T A)$  as the smallest eigenvalue of  $P^T A$ , and  $x$  is the corresponding eigenvector.

$$\lambda_{\min}(A) \leq x^T A x = x^T P \lambda_{\min}(P^T A) \leq \lambda_{\max}(P) \lambda_{\min}(P^T A)$$

$$\lambda_{\min}(P^T A) \geq \frac{\lambda_{\min}(A)}{\lambda_{\max}(P)} \geq \frac{C_1}{\alpha} \frac{\lambda_{\min}(L)}{\lambda_{\max}(L)}$$

$$\gamma \triangleq \frac{\lambda_{\max}(P^T A)}{\lambda_{\min}(P^T A)} \leq \frac{C_2 \lambda_{\max}^2(L)}{C_1 \lambda_{\min}^2(L)} = \frac{C_2}{C_1 \tan^2 \frac{\pi}{2(n+1)}}$$

When  $k < n$

$$\|x_k - x_*\|_A \leq 2 \left( \frac{\sqrt{\gamma} - 1}{\sqrt{\gamma} + 1} \right)^k \|x_0 - x_*\|_A$$

$\varepsilon$  precision:

$$2 \left( \frac{\sqrt{\gamma} - 1}{\sqrt{\gamma} + 1} \right)^k \|x_0 - x_*\|_A < \varepsilon$$

$$\Rightarrow k \geq O \left( \frac{\log \frac{1}{\varepsilon} - \log \frac{1}{2 \|x_0 - x_*\|_A}}{\log \frac{\sqrt{\gamma} - 1}{\sqrt{\gamma} + 1}} \right)$$

Suppose  $\varepsilon$  and  $\|x_0 - x_*\|_A$  are constants.

$$\frac{\sqrt{\gamma} - 1}{\sqrt{\gamma} + 1} \leq \frac{\sqrt{C_2} - \sqrt{C_1} \tan \frac{\pi}{2(n+1)}}{\sqrt{C_2} + \sqrt{C_1} \tan \frac{\pi}{2(n+1)}} \leq 1 - \sqrt{\frac{C_1}{C_2}} \tan \frac{\pi}{2(n+1)}$$

$$\log \frac{\sqrt{\gamma} - 1}{\sqrt{\gamma} + 1} \sim -\frac{1}{n} \Rightarrow k \geq O(n)$$

When  $k \geq n$

$$\|x_k - x_*\|_A = 0$$

2. The singular value decomposition (SVD) is a fundamental decomposition with numerous applications. In this question, we derive the SVD by the eigenvalue decomposition, and develop an algorithm for it. Let  $\mathbf{A} \in \mathbb{R}^{m \times n}$  with  $m \geq n$ . Since  $\mathbf{A}^T \mathbf{A} \in \mathbb{R}^{n \times n}$  is square and symmetric positive semi-definite (SPSD), there exists an eigenvalue decomposition

$$\mathbf{A}^T \mathbf{A} = \mathbf{V} \mathbf{\Lambda} \mathbf{V}^T,$$

where  $\mathbf{\Lambda} = \text{diag}(\lambda_1, \dots, \lambda_n)$  are eigenvalues, and  $\mathbf{V} = [\mathbf{v}_1, \dots, \mathbf{v}_n] \in \mathbb{R}^{n \times n}$  are corresponding eigenvectors.

- (a) Prove that  $\mathbf{A} \mathbf{A}^T \in \mathbb{R}^{m \times m}$  has at most  $n$  nonzero eigenvalues, which are also  $\lambda_1, \dots, \lambda_n$ .  
 (b) Therefore,  $\mathbf{A} \mathbf{A}^T$  has an eigenvalue decomposition

$$\mathbf{A} \mathbf{A}^T = \mathbf{U} \mathbf{\Lambda} \mathbf{U}^T,$$

where  $\mathbf{U} = [\mathbf{u}_1, \dots, \mathbf{u}_n] \in \mathbb{R}^{m \times n}$  are eigenvectors of  $\mathbf{A} \mathbf{A}^T$  corresponding to  $\lambda_1, \dots, \lambda_n$  respectively. Assume all eigenvalues  $\lambda_1, \dots, \lambda_n$  are all simple (though this assumption can be removed). Prove that there exists  $\sigma_i \geq 0$ ,  $i = 1, \dots, n$ , such that

$$\begin{cases} \mathbf{A} \mathbf{v}_i = \sigma_i \mathbf{u}_i, \\ \mathbf{A}^T \mathbf{u}_i = \sigma_i \mathbf{v}_i, \\ \sigma_i^2 = \lambda_i, \end{cases} \quad i = 1, \dots, n.$$

- (c) Define  $\sigma_i = \sqrt{\lambda_i}$ ,  $i = 1, \dots, n$ . Prove that  $\mathbf{A}$  has a decomposition

$$\mathbf{A} = \mathbf{U} \mathbf{\Sigma} \mathbf{V}^T,$$

where  $\mathbf{\Sigma} = \text{diag}(\sigma_1, \dots, \sigma_n)$ . This decomposition is SVD, and  $(\sigma_i, \mathbf{u}_i, \mathbf{v}_i)$  are called singular value, left and right singular vectors of  $\mathbf{A}$  respectively.

- (d) Similar to eigenvalues of symmetric matrices, singular values also have many nice variational properties. Prove

$$\sigma_1 = \max_{\|\mathbf{u}\|_2=1, \|\mathbf{v}\|_2=1} \mathbf{u}^T \mathbf{A} \mathbf{v},$$

where  $\sigma_1$  is the largest singular value of  $\mathbf{A}$ . (There are other identities similar to the min-max theorem of eigenvalues.)

- (e) Use (or not use) (d) to prove

$$\sigma_1 \mathbf{u}_1 \mathbf{v}_1^T = \arg \min_{\text{rank}(\mathbf{B})=1} \|\mathbf{A} - \mathbf{B}\|_F^2.$$

(That is, SVD gives the best rank-1 approximation. This can be extended to any best rank- $r$  approximation, and this makes SVD a fundamental tool in many applications.)

- (f) Propose a power iteration to compute the leading left and right singular vectors of  $\mathbf{A}$ . Your algorithm should use fewest possible matrix-vector products in each iteration. (All eigenvalue algorithms can be extended to SVD.)

(a)  $A^T A = V \Lambda V^T$ .  $\Lambda = \text{diag}(\lambda_1, \dots, \lambda_n)$ .  $A \in \mathbb{R}^{m \times n}$ ,  $m \geq n$

Consider matrix:  $\begin{pmatrix} 0 & A \\ A^T & 0 \end{pmatrix} \in \mathbb{R}^{(m+n) \times (m+n)}$

When  $\lambda \neq 0$

$$\begin{aligned} \left| \lambda I - \begin{pmatrix} 0 & A \\ A^T & 0 \end{pmatrix} \right| &= \begin{vmatrix} \lambda I_m & -A \\ -A^T & \lambda I_n \end{vmatrix} \\ &= \frac{1}{\lambda^n} \begin{vmatrix} \lambda I_m & -A \\ -\lambda A^T & \lambda^2 I_n \end{vmatrix} \\ &= \frac{1}{\lambda^n} \begin{vmatrix} \lambda I_m & -A \\ 0 & \lambda^2 I_n - A^T A \end{vmatrix} \\ &= \lambda^{m-n} |\lambda^2 I_n - A^T A| \end{aligned}$$

$$\begin{aligned} \text{Also, } \begin{vmatrix} \lambda I_m & -A \\ -A^T & \lambda I_n \end{vmatrix} &= \frac{1}{\lambda^m} \begin{vmatrix} \lambda^2 I_m & -\lambda A \\ -A^T & \lambda I_n \end{vmatrix} \\ &= \frac{1}{\lambda^m} \begin{vmatrix} \lambda^2 I_m - A A^T & 0 \\ -A^T & \lambda I_n \end{vmatrix} \\ &= \lambda^{n-m} |\lambda^2 I_m - A A^T| \end{aligned}$$

$$\Rightarrow |\lambda^2 I_m - A A^T| = \lambda^{2(m-n)} |\lambda^2 I_n - A^T A|$$

When  $\lambda = 0$

$$|\lambda^2 I_m - A A^T| = \lambda^{2(m-n)} |\lambda^2 I_n - A^T A| \text{ also holds.}$$

Consider nonzero eigenvalues of  $A A^T$

$$\text{If } |\lambda^2 I_m - A A^T| = 0 \text{ and } \lambda \neq 0$$

$$\text{Then } |\lambda^2 I_n - A^T A| = 0$$

$$\Rightarrow \lambda \in \{\lambda_1, \dots, \lambda_n\}$$

$\Rightarrow A A^T$  has at most  $n$  nonzero eigenvalues which are also  $\lambda_1, \dots, \lambda_n$

(b)  $A^T A = V \Lambda V^T$ .  $V \in \mathbb{R}^{n \times n}$

Define  $w_1 = A v_1$ ,  $w_2 = A v_2$ , ...,  $w_n = A v_n$

Let  $W = [w_1, \dots, w_n] \in \mathbb{R}^{m \times n}$

Note that  $v^T A^T A v = \lambda \Rightarrow W^T W = \Lambda$

Define  $u_i = \frac{w_i}{\sqrt{\lambda_i}}$  ( $\lambda_i \neq 0$ )

For  $\lambda_i = 0$ . select orthonormal  $u = [u_1, \dots, u_n]$

$$\Rightarrow Av_i = \sqrt{\lambda_i} u_i$$

Denote  $\sigma_i = \sqrt{\lambda_i}$ .

$$\Sigma = \text{diag} \{ \sigma_1, \dots, \sigma_n \}$$

$$\Rightarrow AV = U\Sigma$$

$$\Rightarrow U^T A = \Sigma V^T$$

$$\Rightarrow A^T U = V \Sigma$$

$$\Rightarrow A^T u_i = \sigma_i v_i$$

In conclusion,

$$\begin{cases} Av_i = \sigma_i u_i \\ A^T u_i = \sigma_i v_i \\ \sigma_i^2 = \lambda_i \end{cases} \quad i = 1, 2, \dots, n$$

(c) From (b) we have  $AV = U\Sigma$

$$\text{Then, } A = U\Sigma V^T$$

(d) For  $U \in \mathbb{R}^{m \times n}$ ,  $\Sigma \in \mathbb{R}^{n \times n}$

Define  $\tilde{U} \in \mathbb{R}^{m \times m}$  as the completion of  $U$ .

$$\tilde{U} [1:n] = U.$$

$$\tilde{U}^T \tilde{U} = \tilde{U} \tilde{U}^T = I$$

$$\text{Define } \tilde{\Sigma} \in \mathbb{R}^{m \times n}, \quad \tilde{\Sigma} = \begin{pmatrix} \Sigma \\ 0 \end{pmatrix}$$

$$\text{Then } A = U\Sigma V^T = \tilde{U} \tilde{\Sigma} V^T$$

$$\max_{\substack{\|u\|_2=1 \\ \|v\|_2=1}} u^T A v = \max_{\substack{\|u\|_2=1 \\ \|v\|_2=1}} u^T \tilde{U} \tilde{\Sigma} V^T$$

Since  $\|\tilde{U} u^T\| = 1$ ,  $\|V V^T\| = 1$ , we can write RHS as  $\max_{\substack{\|u\|_2=1 \\ \|v\|_2=1}} u^T \tilde{\Sigma} V^T$

$$\text{i.e. } \max_{\substack{\|u\|_2=1 \\ \|v\|_2=1}} u^T A v = \max_{\substack{\|u\|_2=1 \\ \|v\|_2=1}} u^T \tilde{\Sigma} V^T$$

$$(u^T \tilde{\Sigma} V^T)^2 = \left( \sum_{i=1}^n \sigma_i u_i v_i \right)^2 \leq \left( \sum_{i=1}^n u_i^2 \right) \left( \sum_{i=1}^n \sigma_i^2 v_i^2 \right) = \sigma_1^2$$

$$\Rightarrow u^T \tilde{\Sigma} V^T \leq \sigma_1$$

$$\Rightarrow \max_{\substack{\|u\|_2=1 \\ \|v\|_2=1}} u^T A v \leq \sigma_1$$

Select  $u = u_1$ ,  $v = v_1$ .

$$u_1^T A v_1 = \sigma_1 \Rightarrow \sigma_1 \text{ could be achieved}$$

$$\Rightarrow \max_{\substack{\|u\|_2=1 \\ \|v\|_2=1}} u^T A v = \sigma_1$$

$$(e) \text{ rank}(B) = 1$$

$$\Rightarrow \exists c \in \mathbb{R}, u \in \mathbb{R}^{n \times 1}, v \in \mathbb{R}^n, \|u\|_2 = \|v\|_2 = 1 \text{ s.t.}$$

$$B = c u v^T.$$

$$\begin{aligned} \|A - B\|_F^2 &= \|A - c u v^T\|_F^2 \\ &= \text{Tr}((A - c u v^T)^T (A - c u v^T)) \\ &= \text{Tr}(A^T A) - 2c \text{Tr}(v u^T A) + c^2 \end{aligned}$$

For fixed  $u, v$ ,

$$\|A - B\|_F^2 \text{ achieves its minimum} \Leftrightarrow c = \text{Tr}(v u^T A)$$

$$\Rightarrow \|A - B\|_F^2 \geq \text{Tr}(A^T A) - \text{Tr}(v u^T A)^2$$

$$\text{Note that } \text{Tr}(v u^T A) = u^T A v, \text{Tr}(A^T A) = \sum_{i=1}^n \sigma_i^2$$

$$\text{From (d), we have } \min_{\substack{\|u\|_2=1 \\ \|v\|_2=1}} u^T A v = \sigma_1$$

$$\Rightarrow \|A - B\|_F^2 \geq \sum_{i=1}^n \sigma_i^2 - \sigma_1^2 = \sum_{i=2}^n \sigma_i^2$$

$$\Rightarrow \min_{\text{rank}(B)=1} \|A - B\|_F^2 \geq \sum_{i=2}^n \sigma_i^2$$

$$\text{When } B = \sigma_1 u_1 v_1^T,$$

$$\|A - B\|_F^2 = \text{Tr}(A^T A) - \sigma_1^2 = \sum_{i=2}^n \sigma_i^2$$

$$\Rightarrow \sigma_1 u_1 v_1^T \in \arg \min_{\text{rank}(B)=1} \|A - B\|_F^2$$

$$(f) \text{ Let } X = A^T A \in \mathbb{R}^{n \times n}$$

The algorithm for computing leading right singular vector of  $A$  is:

$$\text{Select } x^{(0)} \in \mathbb{R}^n \text{ s.t. } \|x^{(0)}\|_2 = 1$$

For  $k = 1, 2, \dots$

$$z^{(k)} = X x^{(k-1)}$$

$$x^{(k)} = z^{(k)} / \|z^{(k)}\|_2$$

End

$$\text{Denote } v = x^{(\text{max.iter})}$$



Then  $\tilde{u} = A^T v$

$$u = \tilde{u} / \|\tilde{u}\|$$

The leading right singular vector of  $A$  is  $v$ .

The leading left singular vector of  $A$  is  $u$

3. We have used QR decomposition for solving the least squares (LS) problems in GMRES method

$$\min_{x \in \mathbb{R}^n} \|Ax - b\|_2^2$$

where  $A \in \mathbb{R}^{m \times n}$  is rank  $r$  with  $m \geq n$ . Actually, LS problems arise in many other applications in applied mathematics and engineering, and there are other solvers for them. We consider to use the singular value decomposition (SVD) to solve LS problems.

(a) If we take only the non-zero singular values, then we obtain the compact SVD of  $A$

$$A = U \Sigma V^T,$$

where  $U \in \mathbb{R}^{m \times r}$  satisfies  $U^T U = I$ ,  $V \in \mathbb{R}^{n \times r}$  satisfies  $V^T V = I$ , and  $\Sigma = \text{diag}(\sigma_1, \sigma_2, \dots, \sigma_r)$  with  $\sigma_1 \geq \sigma_2 \geq \dots \geq \sigma_r > 0$ . Prove that

$$\text{Ran}(A) = \text{Ran}(U), \quad \text{Ker}(A) = \text{Ran}(V)^\perp,$$

where  $(\cdot)^\perp$  stands for the orthogonal complementary.

(b) Assume  $r = n$ . Prove that the solution of LS is unique and is given by

$$x = V \Sigma^{-1} U^T b.$$

(c) Continuing (b): Let  $\tilde{x}$  be the solution of the LS when the input  $b$  is perturbed to  $\tilde{b}$ . Prove that

$$\frac{\|x - \tilde{x}\|_2}{\|x\|_2} \leq C \frac{\|b - \tilde{b}\|_2}{\|b\|_2},$$

where  $C = \frac{\sigma_1}{\sigma_n}$  if the LS is solved by formula in (b) and  $C = \frac{\sigma_1^2}{\sigma_n^2}$  if the LS solution is obtained by solving the normal equation.

(d) Assume  $r < n$ . Prove that all solutions of LS are given by

$$x = V \Sigma^{-1} U^T b + y, \quad y \in \text{Ker}(A),$$

and  $x_0 := V \Sigma^{-1} U^T b$  is the solution of LS with the minimum 2-norm among all solutions.

(a) SVD decomposition of  $A$ :

$$A = U \Sigma V^T$$

$$\Rightarrow \text{Ran}(A) \subseteq \text{Ran}(U)$$

$$\text{Also, since } U = A V \Sigma^{-1}$$

$$\Rightarrow \text{Ran}(U) \subseteq \text{Ran}(A)$$

$$\text{Hence } \text{Ran}(A) = \text{Ran}(U)$$

$$A = U \Sigma V^T \Rightarrow \text{Ran}(V)^\perp = \text{Ker}(V^T) \subseteq \text{Ker}(A)$$

$$V^T = \Sigma^{-1} U^T A \Rightarrow \text{Ker}(A) \subseteq \text{Ker}(V^T) = \text{Ran}(V)^\perp$$

$$\Rightarrow \text{Ker}(A) = \text{Ran}(V)^\perp$$

$$(b) \quad \|Ax - b\|_2^2 = \|U \Sigma V^T x - b\|_2^2 \\ = \|\Sigma V^T x - U^T b\|_2^2$$

$$(\text{Let } y = V^T x) \quad = \|\Sigma y - U^T b\|_2^2$$

Since  $\min_y \|\Sigma y - U^T b\|_2^2$  has a unique solution  $y = \Sigma^{-1} U^T b$

$\min_x \|Ax - b\|_2^2$  also has a unique solution  $V^T x = \Sigma^{-1} U^T b$

$$\Rightarrow x = V \Sigma^{-1} U^T b$$

(c) •  $\tilde{x}$  is the solution of LS by the formula in (b)

$$\tilde{x} = V \Sigma^{-1} U^T \tilde{b}$$

$$\|x - \tilde{x}\|_2 = \|V \Sigma^{-1} U^T (b - \tilde{b})\|_2 \\ \leq \|V \Sigma^{-1} U^T\|_2 \|b - \tilde{b}\|_2 \\ = \|\Sigma^{-1}\|_2 \|b - \tilde{b}\|_2 \\ = \frac{1}{\sigma_n} \|b - \tilde{b}\|_2$$

Note that  $x = V \Sigma^{-1} U^T b$

$$\Rightarrow b = U \Sigma V^T x$$

Thus  $\|b\|_2 = \|U \Sigma V^T x\|_2$

$$\leq \|\Sigma\|_2 \|x\|_2 \\ = \sigma_1 \|x\|_2$$

$$\Rightarrow \frac{1}{\|x\|_2} \leq \frac{\sigma_1}{\|b\|_2}$$

Therefore,

$$\frac{\|x - \tilde{x}\|_2}{\|x\|_2} \leq \frac{\sigma_1}{\sigma_n} \frac{\|b - \tilde{b}\|_2}{\|b\|_2}$$

•  $\tilde{x}$  is the solution obtained by solving the normal equation

$$A^T A x = A^T b \Rightarrow x = (A^T A)^{-1} A^T b$$

$$A^T A \tilde{x} = A^T \tilde{b} \Rightarrow \tilde{x} = (A^T A)^{-1} A^T \tilde{b}$$

$$\begin{aligned} \|x - \tilde{x}\|_2 &\leq \|(A^T A)^{-1} A^T\|_2 \|b - \tilde{b}\|_2 \\ &\leq \|(A^T A)^{-1}\|_2 \|A^T\|_2 \|b - \tilde{b}\|_2 \\ &= \frac{\sigma_1}{\sigma_n^2} \|b - \tilde{b}\|_2 \end{aligned}$$

$$\text{Since } \frac{1}{\|x\|_2} \leq \frac{\sigma_1}{\|b\|_2}$$

$$\Rightarrow \frac{\|x - \tilde{x}\|_2}{\|x\|_2} \leq \frac{\sigma_1^2}{\sigma_n^2} \frac{\|b - \tilde{b}\|_2}{\|b\|_2}$$

$$\begin{aligned} (a) \quad \|Ax - b\|^2 &= \|U \Sigma V^T x - b\|^2 \\ &= \|\tilde{U} \tilde{\Sigma} \tilde{V}^T x - b\|^2 \end{aligned}$$

where  $\tilde{U}, \tilde{\Sigma}, \tilde{V}$  are full vectors and singular values with  $U$ .

$$\text{Let } P = \tilde{V}^T x$$

$$\text{Assume } P = \begin{bmatrix} p_r \\ p_{n-r} \end{bmatrix}$$

$$\|Ax - b\|^2 = \|\tilde{U} \tilde{\Sigma} P - b\|^2$$

$$\text{Solution: } p_r = \tilde{\Sigma}^{-1} \tilde{U}^T b$$

$$\begin{aligned} \Rightarrow x = \tilde{V} P &= [V_r \ V_{n-r}] \begin{bmatrix} p_r \\ p_{n-r} \end{bmatrix} \\ &= V \tilde{\Sigma}^{-1} \tilde{U}^T b + V_{n-r} p_{n-r} \\ &= V \Sigma^{-1} U^T b + V_{n-r} p_{n-r} \end{aligned}$$

$$\text{Let } y = V_{n-r} p_{n-r} \in \text{Ker}(A)$$

$$\Rightarrow x = V \Sigma^{-1} U^T b + y.$$