# In Search of an Understandable Consensus Algorithm - Raft

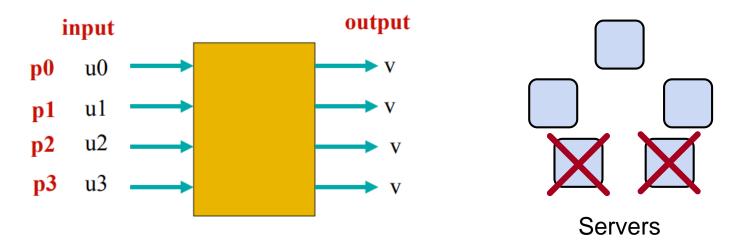
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#### What is Consensus?

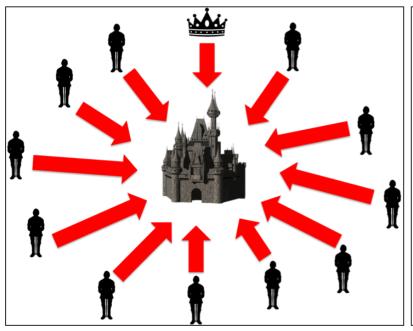
- involves multiple servers agreeing on values
- Recovers from server failures autonomously
  - Minority of servers fail: no problem
  - Majority fail: lose availability, retain consistency

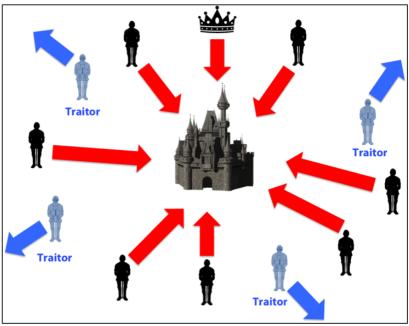


Key to building consistent storage systems

#### Paxos? Too complicated!

- Leslie Lamport, 1989
- Byzantine Generals Problem





**Coordinated Attack Leading to Victory** 

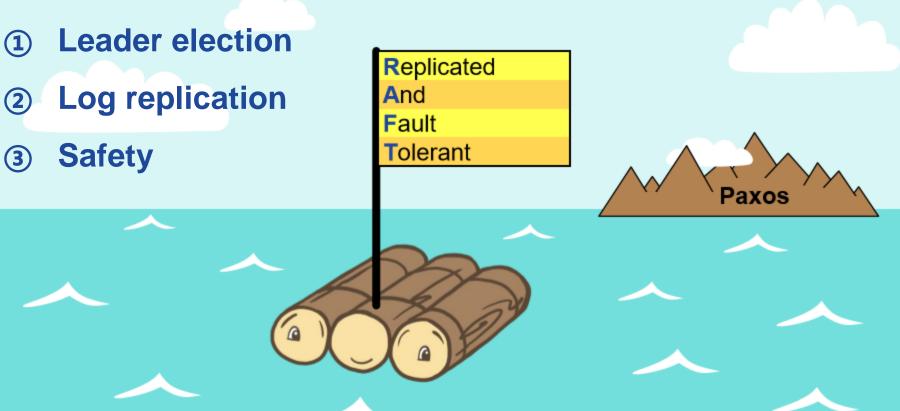
**Uncoordinated Attack Leading to Defeat** 

- Paxos: a consensus algorithm to solve it
  - But... too difficult to be understood & implemented

#### Method: Design a new algorithm

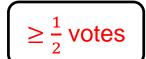
- Named Raft
- Design for Understandability
- Define a better algorithm for building real systems
  - Must be correct, complete, and perform well
  - Must also be understandable
- Be easier to understand than Paxos
  - Fundamentally different decomposition than Paxos
  - Less complexity

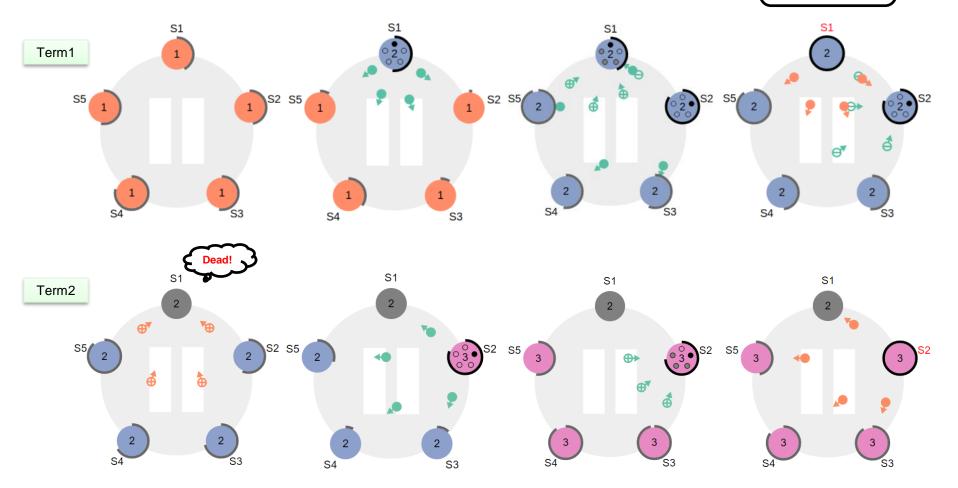
# 3 main parts:



#### 1. Leader election

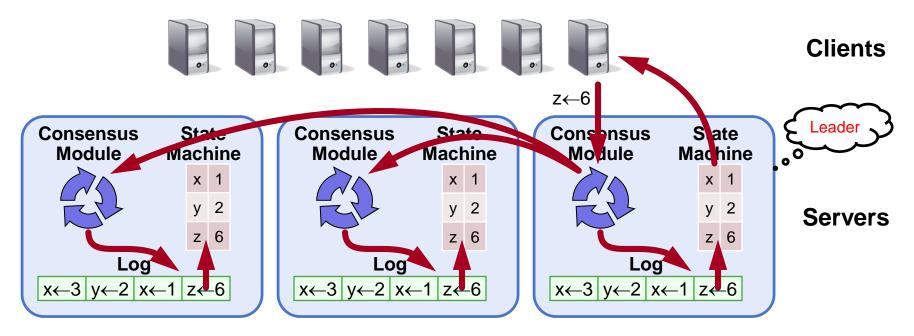
- Select one of the servers to act as cluster leader
- If the leader is crashed, choose a new one





#### 2. Log replication

Replicated State Machines

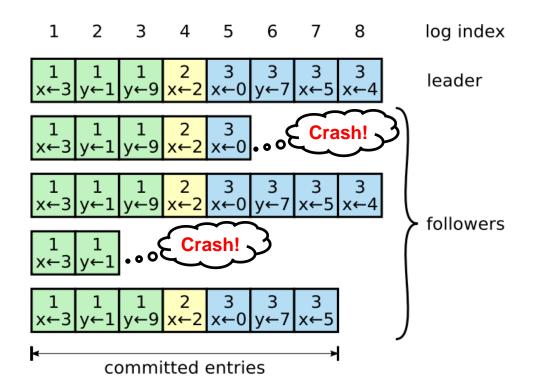


- Replicated log ⇒ replicated state machine
  - All servers execute same commands in same order
- Consensus module ensures proper log replication
- Node crashes / Leader changes ⇒ replication may be wrong!

## 2. Log replication

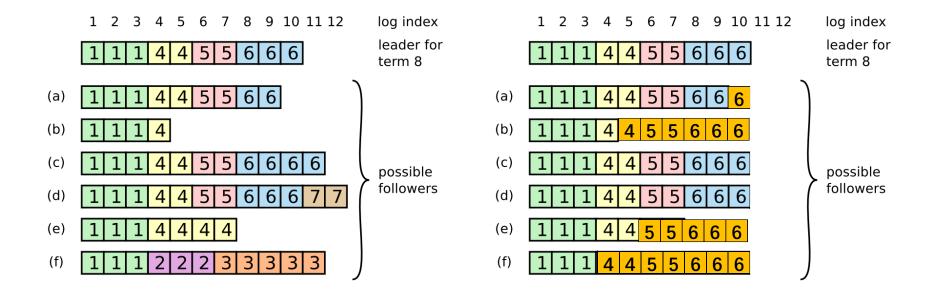
#### Node crashes

- Safely replication
  - More than half of nodes have replicated logs
  - then leader commits



#### 2. Log replication

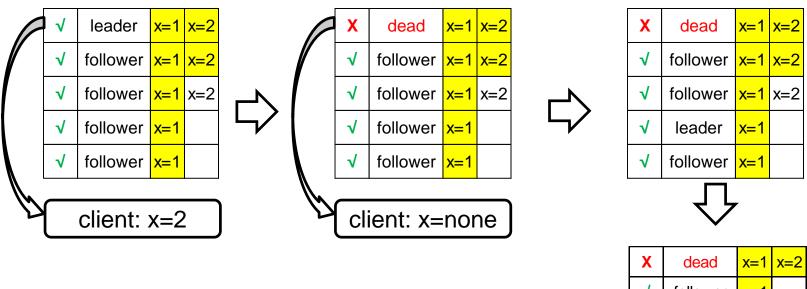
- Leader changes
- Find the last common log from the end to the front
- Force followers to replicate different logs!



## 3. Safety

#### Hidden error of execute

each node executes the same commands in the same order?



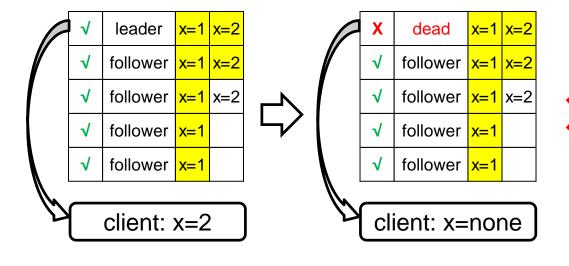
What is wrong?

X	dead	x=1	x=2
>	follower	x=1	
<b>&gt;</b>	follower	x=1	
<b>&gt;</b>	leader	x=1	
<b>√</b>	follower	x=1	

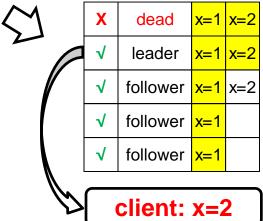
**client:** x=1???

## 3. Safety

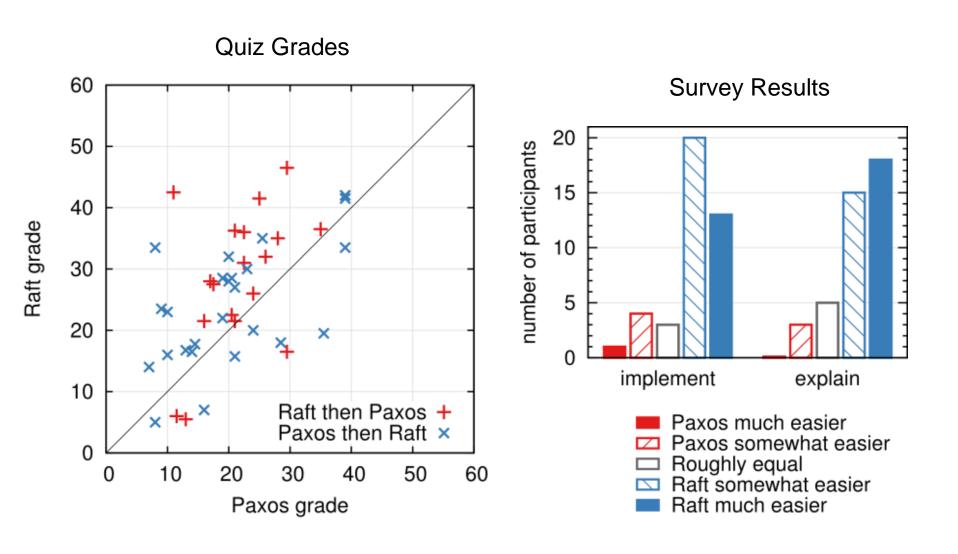
Node with <u>smaller committed log number</u> cannot be elected as leader





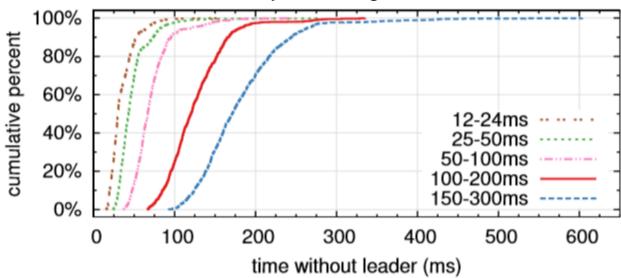


## **Result: Raft User Study**



#### **Result: Performance**

- The time to detect and replace a crashed leader.
- "x-y ms" means election timeout.
  - Downtime can be reduced by reducing the election timeout.



#### Recommendation

- Use a conservative election timeout such as 150–300ms
- Avoid unnecessary leader changes

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## **Questions**

