Primal-Dual Formulation

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1 Simplified Problem

Let us take a look at a simplified problem of each residence scheduling only a single load of unit rating for one interval anywhere within the time window. This means that each residence has to schedule their load for only a single interval within the time window of length T. The residence n aims to obtain the optimal load schedule $\{p_n^t\}_{t=1}^T \in \{0,1\}$. The power schedules are coupled across space through the voltage constraints (1d). The overall optimization problem is shown below. For every interval t, the load is either switched on with $p_n^t = 1$ or it is turned off $p_n^t = 0$. (1c) ensures that the load at each residence n is switched on for only a single interval.

$$\min \quad \sum_{n=1}^{N} \sum_{t=1}^{T} c^t p_i^t \tag{1a}$$

over
$$\{\{p_i^t\}_{t=1}^T\}_{n=1}^N \in \{0,1\}$$
 (1b)

s.to.
$$\sum_{t=1}^{w} p_n^t = 1 \quad \forall n = 1, 2, \dots, N$$
 (1c)

$$\alpha \le -\sum_{j=1}^{N} R_{nj} p_j^t \le \beta \quad \forall n = 1, \dots, N \quad \forall t = 1, \dots, T$$
 (1d)

The above problem can be solved for the optimal solution in a centralized manner. However, in reality, the information regarding all residences need not be available to the central entity. This requires a distributed approach of finding the optimal solution. We use a primal-dual approach to address this problem [1, 2]. Let $\mathbf{P} \in \{0,1\}^{N \times T}$ be the power schedule where the entry at the n^{th} row and t^{th} column denotes the

Let $\mathbf{P} \in \{0,1\}^{N \times T}$ be the power schedule where the entry at the n^{th} row and t^{th} column denotes the power schedule for the time interval t at residence n. Let the N Lagrangian multipliers corresponding to the equality constraints be stacked in the vector $\mathbf{\lambda} \in \mathbb{R}^N$. Let $\mathbf{U} \in \mathbb{R}^{N \times T}$ be the Lagrangian multipliers corresponding to the inequality constraints. Therefore, we can write the Lagrangian function as

$$\mathcal{L}(\mathbf{P}, \mathbf{U}, \lambda) = \mathbf{1}^{T} \mathbf{P} \mathbf{c} + \lambda (\mathbf{1} - \mathbf{P} \mathbf{1}) + \mathsf{Tr} \left[\mathbf{U}^{T} \left(\alpha \mathbf{1} \mathbf{1}^{T} + \mathbf{R} \mathbf{P} \right) \right]$$
(2)

Next steps

- Compute the partial derivative of the Lagrangian function \mathcal{L} with respect to each row of \mathbf{P} , each row of \mathbf{U} and each entry of $\boldsymbol{\lambda}$: $\frac{\partial \mathcal{L}}{\partial \mathbf{p}_n}$, $\frac{\partial \mathcal{L}}{\partial \mathbf{u}_n}$, $\frac{\partial \mathcal{L}}{\partial \boldsymbol{\lambda}_n}$
- Update the primal and dual variables over successive iterations

$$\mathbf{p}_{n}^{(k+1)} = \mathbf{p}_{n}^{(k)} + \gamma_{\mathbf{p}} \frac{\partial \mathcal{L}}{\partial \mathbf{p}_{n}}$$
 (3a)

$$\mathbf{u}_{n}^{(k+1)} = \mathbf{u}_{n}^{(k)} + \gamma_{\mathbf{u}} \frac{\partial \mathcal{L}}{\partial \mathbf{u}_{n}}$$
(3b)

$$\lambda_n^{(k+1)} = \lambda_n^{(k)} + \gamma_\lambda \frac{\partial \mathcal{L}}{\partial \lambda_n}$$
 (3c)

(3d)

References

- [1] Mark Eisen, Clark Zhang, Luiz F. O. Chamon, Daniel D. Lee, and Alejandro Ribeiro. Learning Optimal Resource Allocations in Wireless Systems. *IEEE Transactions on Signal Processing*, 67(10):2775–2790, May 2019.
- [2] Sarthak Gupta, Vassilis Kekatos, and Walid Saad. Optimal real-time coordination of energy storage units as a voltage-constrained game, 2018.