

# F11. Design Patterns, Operator-Precedence parser (L6 prep.)

Grammar-driven parser design

L6 Preparation

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# Interpreter

- An **interpreter** is, basically, a program that reads and executes a set of instructions (i.e. another program written in another language), without converting it to machine code.
- For a **compiler** (such as VC++) it is the other way around: instructions are compiled to machine code, but it is not executed.
- The interpreted/compiled language has some form of **grammar**, defining its structures and syntax.
  - Can be simple, such as one-worded instructions (e.g. machine code).
  - Or complex, such as C++.

# Interpreter for logical-comparison expressions

1 && 1 < 2 && 0 || ( 1 == 1 )

- Task: interpret (*evaluate*) expressions containing **logical and comparative operators**
- Logical operators: &&, ||
- Comparison operators: <, >, ==, <=, >=, !=

# Precedence

1 && 1 < 2 && 0 || ( 1 == 1 )

## Language rules

- **Precedence:**

1. Parenthesis ( ... )
2. Comparison op's: <, >, ==, >=, <=, !=
3. And: &&
4. Or: ||

- **Associativity:** operations (with the same operator) are interpreted left-to-right

# Things to note

- The C++ `bool` type implicitly casts to `int`. This allows some weird-looking expressions.
- **Example:**

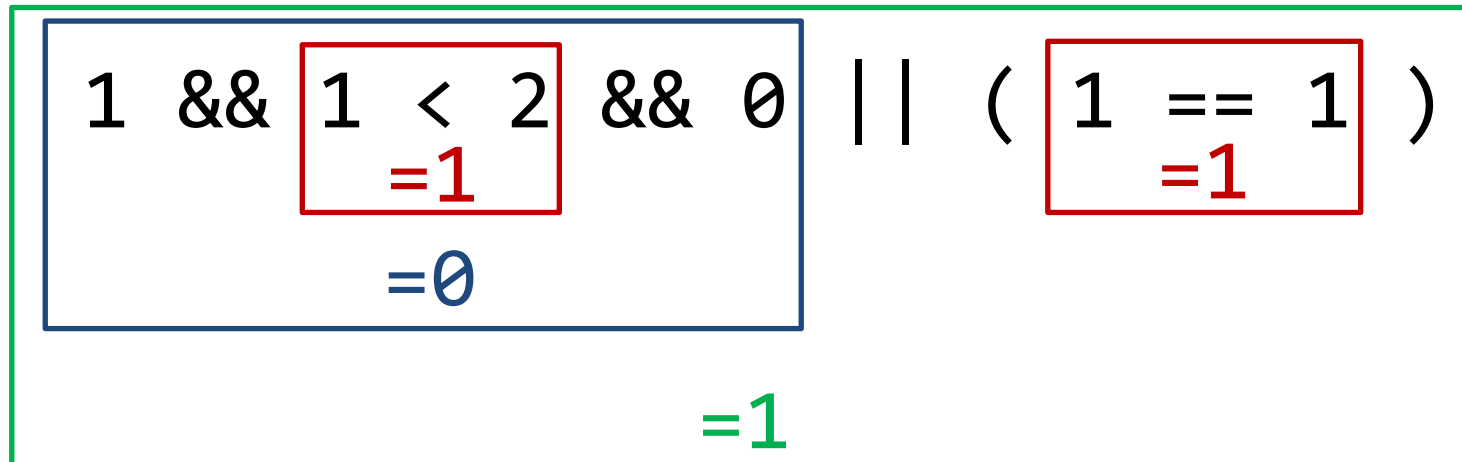
`2 == 2 == 2` equals `false` (!)

  - Same operator so no precedence rules take effect.
  - Left-associativity: `2 == 2` equals `true`
  - Then, `2 == 2 == 2`  $\Leftrightarrow$   
`true == 2`  $\Leftrightarrow$   
`1 == 2`  
equals `false`

implicit cast `bool` to `int`

# Precedence

- We can work out the solution intuitively



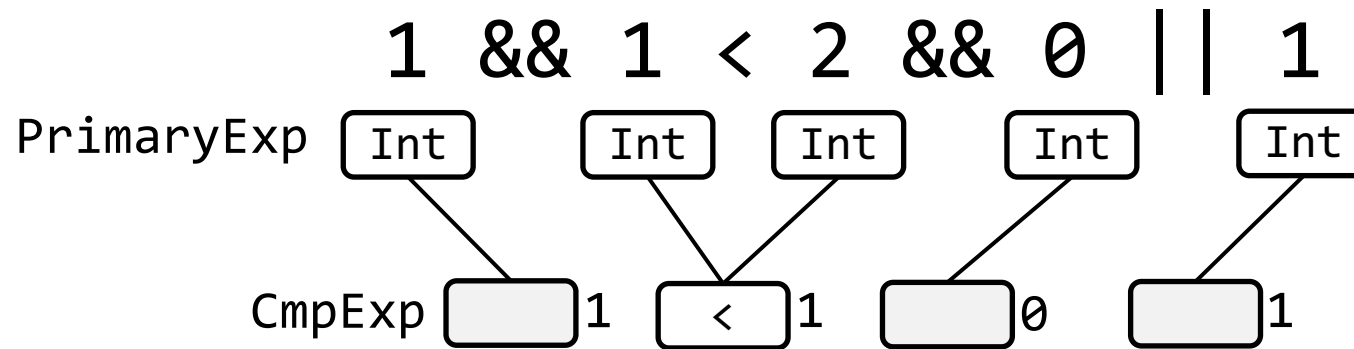
- But how do we design an interpreter to do the job?
  - This figure is actually a good start, but lets remake it as a tree (with a slightly simplified expression to start with)

# Precedence Tree

1 && 1 < 2 && 0 || 1  
PrimaryExp Int Int Int Int Int

- Start by defining **primary expressions**.  
These are “indivisible” parts of the expression.
  - This example: **Numbers** (int)
  - Parenthesis
  - Also, math expressions based on +, −, \*, / → **what you’ll do in L6.**

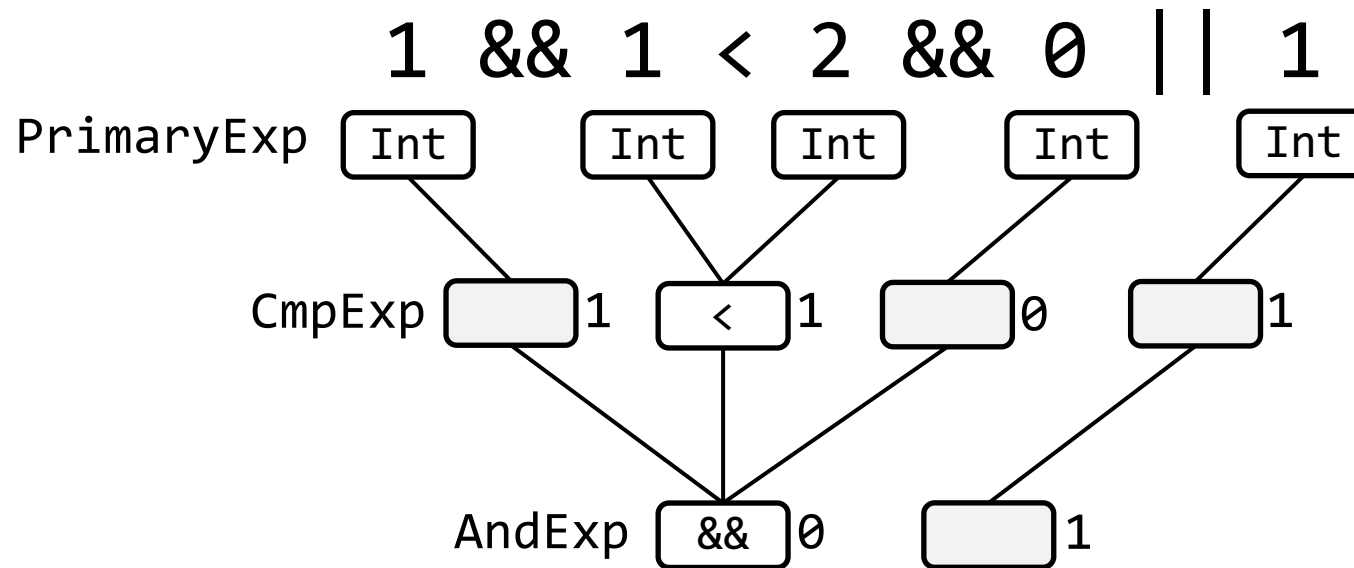
# Precedence Tree



- Go on with the top priority operations– **comparisons**.
  - Primaries (`PrimaryExp`) **separated by a comparison operator** are grouped in **comparison expressions** (`CmpExp`) and evaluated.
  - Primaries without a comparison operator also form `CmpExp`'s – but without a right-hand-side operator (left-hand-side is then simply forwarded).

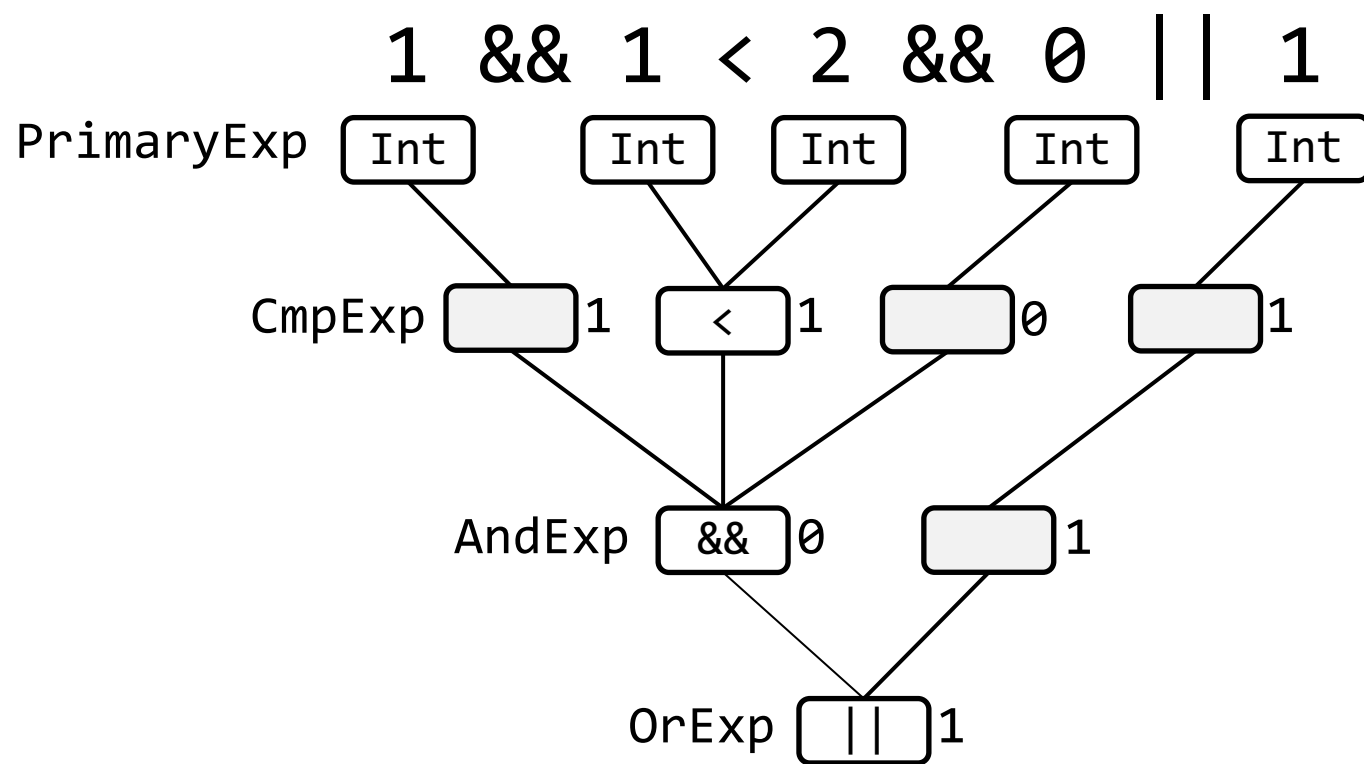


# Precedence Tree



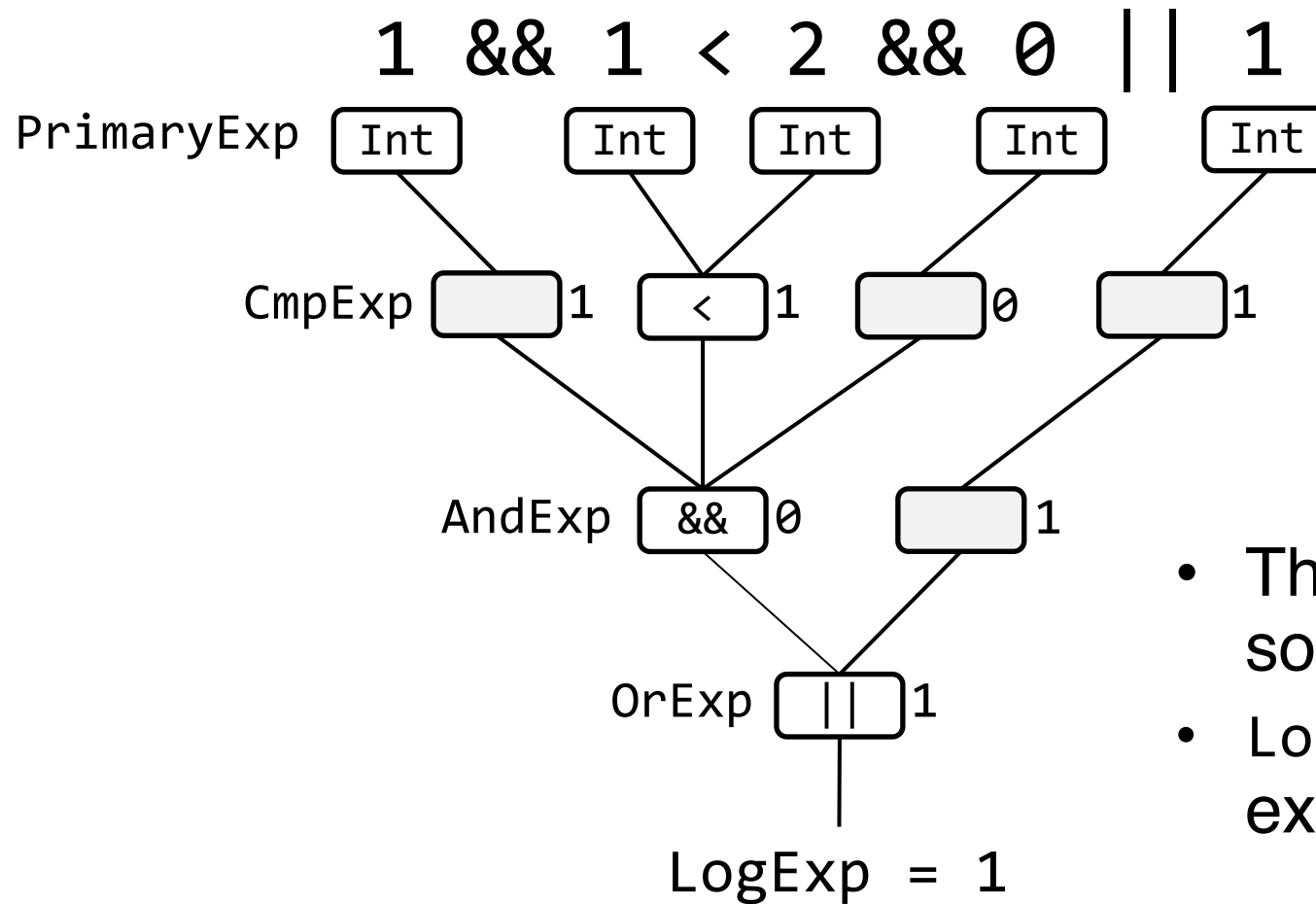
- Same procedure with the second highest-priority operator: **And (&&)**
- Note that some operators (like &&) may have **more than two operands**. In this case, the expression is evaluated left-to-right

# Precedence Tree



- Finally, the operator with lowest priority: **Or (||)**

# Precedence Tree



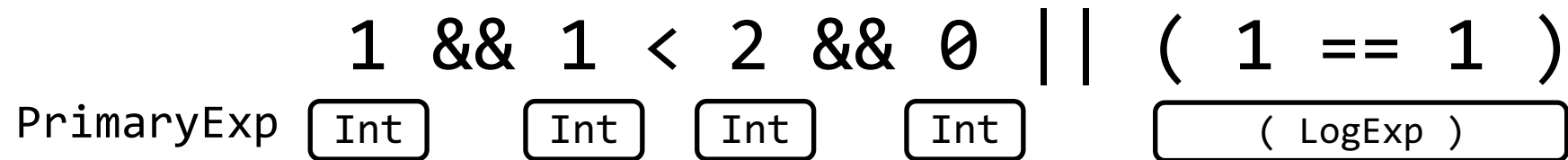
- The evaluation of `OrExp` is the solution to the expression
- `LogExp` captures the entire expression.

# Precedence Tree

1 && 1 < 2 && 0 || ( 1 == 1 )

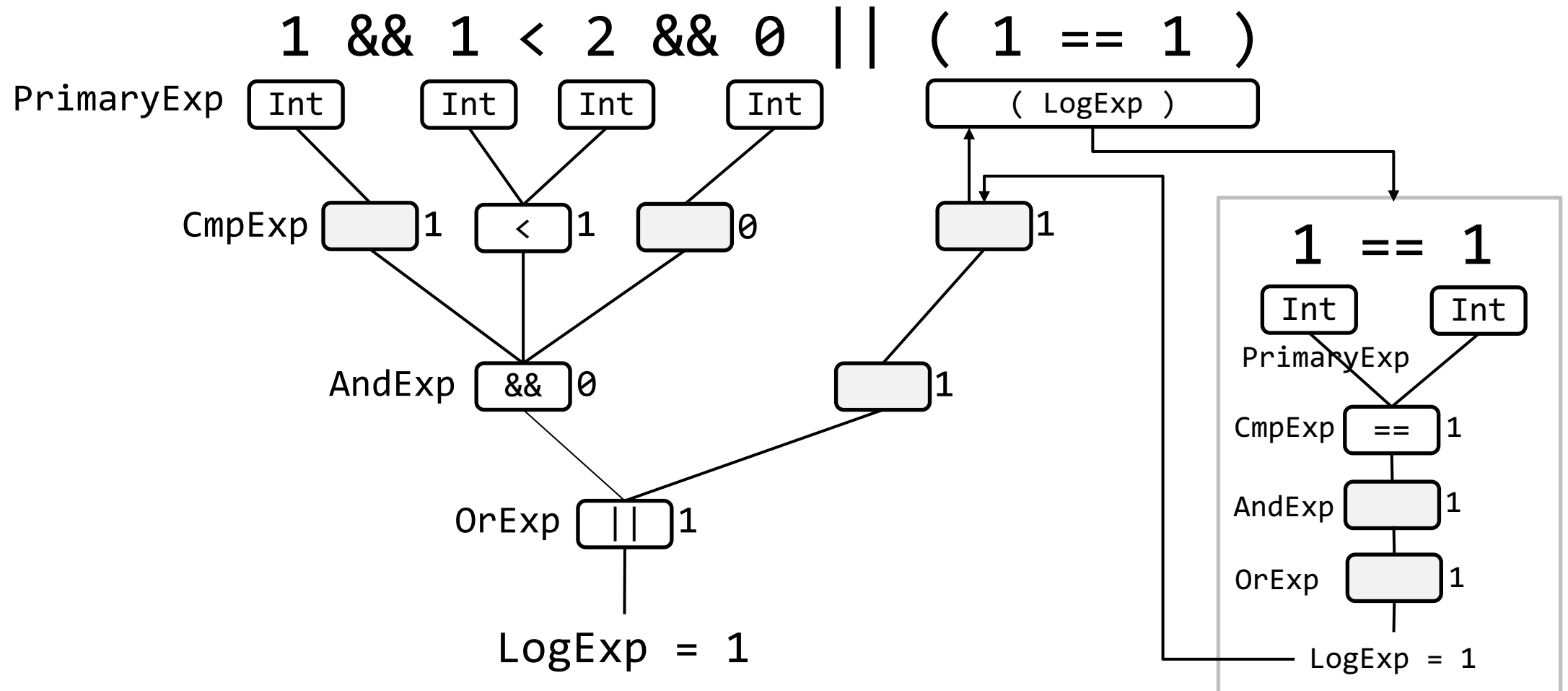
- So what about parentheses?

# Precedence Tree



- A parenthesis is another form of **Primary expression**
  - “(“ **LogExp** “)”
- Parenthesis grants localized priority to a sub-expression with any combination of primaries and operators – i.e. a **sub-LogExp**.
- The inside of a parenthesis is thus evaluated as a **separate LogExp-expression**.

# Precedence Tree



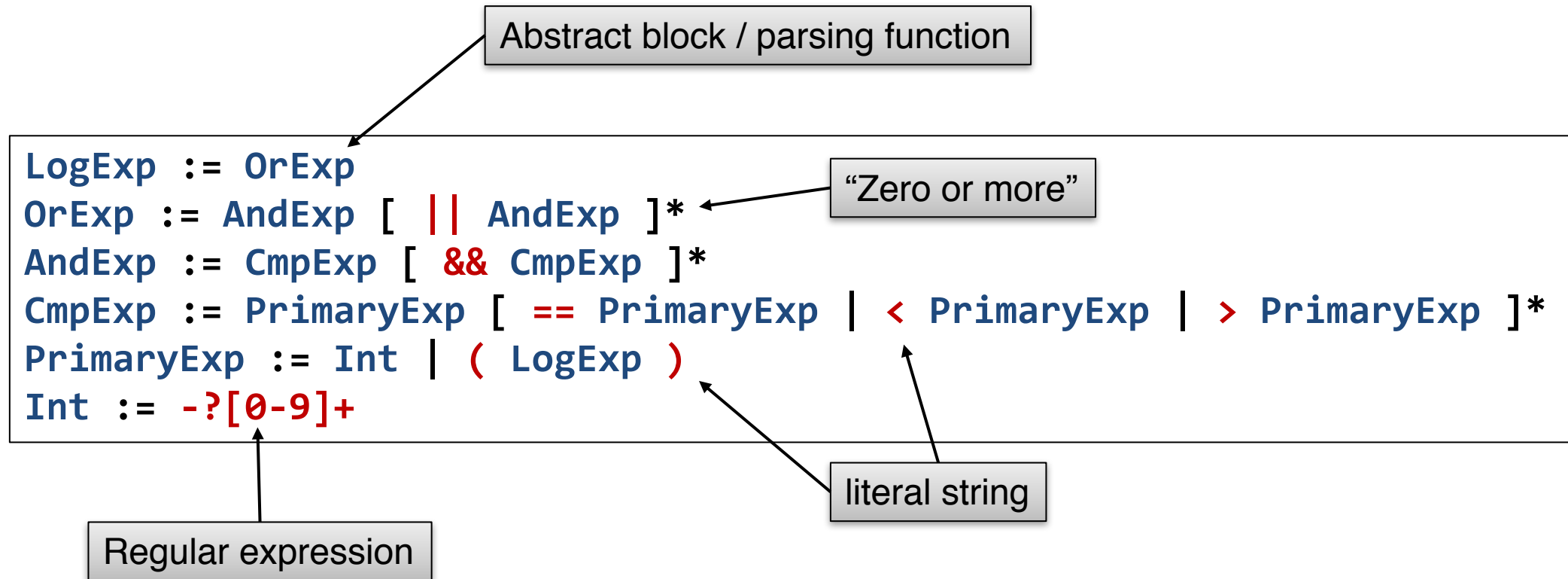
# Abstract Grammar

- **Abstract grammar** for logical-comparison expressions

```
LogExp := OrExp
OrExp  := AndExp [ || AndExp ]*
AndExp := CmpExp [ && CmpExp ]*
CmpExp := PrimaryExp [ == PrimaryExp | < PrimaryExp | > PrimaryExp ]*
PrimaryExp := Int | ( LogExp )
Int := -?[0-9]+
```

# Abstract Grammar

- **Abstract grammar** for logical-comparison expressions





# Operator-Precedence parser

The grammar gives important hints on how we can **design the parser**.

- Let each block/operator-expression be a separate **parsing function**
  - LogExp, OrExp, AndExp, CmpExp, PrimaryExp
- Parsing functions ...
  - **Consume** (step through) tokens of the expression.  
*Token* = sequence of characters separated by whitespace.
  - **Obtain operands by calling other parsing functions**. Usually the one with next-level (higher) priority.
  - **Evaluates operands**, if matched by its role (e.g. “+” for AndExp)  
*Example*: OrExp obtains its operands (one or more) by “asking” AndExp.  
AndExp then consumes & evaluates tokens, and return the result.

# **EXAMPLE CODE**

## **OPERATOR-PRECEDENCE PARSER**

# Initialization

- Assume we are implementing an Interpreter for a logical-comparison grammar.
- The class has the following basic content:

```
// Tokenized expression
std::vector<std::string> tokens =
    { "1", "&&", "1", "<", "2", "&&", "0", "||", "1", "(", "1", "==" , "1", ")" };

// Current position
int position = 0;

// Reserved end of expression string symbol (ETX = end-of-text)
const std::string ETX = "\u0003";
```

# Auxiliaries: peek

- **peek** – get current or future token

```
// Return current token
std::string peek()
{
    return peek(0);
}

// Return token @steps ahead
std::string peek(int steps)
{
    if (position+steps >= tokens.size()) return ETX;

    return tokens[position+steps];
}
```

# Auxiliaries: consume

- **consume** – step one token forward

```
// Advance to the next token.  
// @token is a safeguard to make sure the caller knows what is being consumed.  
void consume(const std::string& token)  
{  
    std::string next_token = peek();  
    if (next_token == ETX)  
        throw std::runtime_error("Consumed past last token\n");  
    if (next_token != token)  
        throw std::runtime_error("Could not consume token " + token + "\n");  
  
    ++position;  
}
```

# evaluate

- **evaluate** – have the interpreter evaluate an expression
  - Immediately queries LogExp (private member function), which queries OrExp, the parsing function with lowest priority

```
bool evaluate()
{
    return parse_LogExp();
}
```

```
LogExp := OrExp
OrExp := AndExp [ || AndExp ]*
AndExp := CmpExp [ && CmpExp ]*
CmpExp := PrimaryExp [ == PrimaryExp |
                        < PrimaryExp |
                        > PrimaryExp ]*
PrimaryExp := Int | ( LogExp )
Int := -?[0-9]+
```

```
bool parse_LogExp()
{
    return parse_OrExp();
}
```

# OrExp

```
LogExp := OrExp
OrExp  := AndExp [ || AndExp ]*
AndExp := CmpExp [ && CmpExp ]*
CmpExp := PrimaryExp [ == PrimaryExp |
                        < PrimaryExp |
                        > PrimaryExp ]*
PrimaryExp := Int | ( LogExp )
Int := -?[0-9]+
```

```
bool parse_OrExp()
{
    // Parse the left-hand-side block
    bool result = parse_AndExp();

    // Parse right-hand-side blocks
    std::string next_token = peek();
    while (1)
    {
        if (next_token == "||")
        {
            consume("||");
            result = (result || parse_AndExp());
        }
        else
            break;

        next_token = peek();
    }

    return result;
}
```

# AndExp

```
LogExp := OrExp
OrExp := AndExp [ || AndExp ]*
AndExp := CmpExp [ && CmpExp ]*
CmpExp := PrimaryExp [ == PrimaryExp |
                        < PrimaryExp |
                        > PrimaryExp ]*
PrimaryExp := Int | ( LogExp )
Int := -?[0-9]+
```

```
bool parse_AndExp()
{
    // Parse the left-hand-side block
    bool result = parse_CmpExp();

    // Parse right-hand-side blocks
    std::string next_token = peek();
    while (1)
    {
        if (next_token == "&&")
        {
            consume("&&");
            result = result && parse_CmpExp();
        }
        else
            break;

        next_token = peek();
    }

    return result;
}
```



# CmpExp

```
LogExp := OrExp
OrExp := AndExp [ || AndExp ]*
AndExp := CmpExp [ && CmpExp ]*
CmpExp := PrimaryExp [ == PrimaryExp |
                        < PrimaryExp |
                        > PrimaryExp ]*
PrimaryExp := Int | ( LogExp )
Int := -?[0-9]+
```

```
bool parse_CmpExp()
{
    // Parse the left-hand-side block
    int result = parse_PrimaryExp();

    // Parse right-hand-side blocks
    std::string next_token = peek();
    while (1)
    {
        if (next_token == "==")
        {
            consume("==");
            result = (result == parse_PrimaryExp());
        }
        else if (next_token == "<")
        {
            consume("<");
            result = (result < parse_PrimaryExp());
        }
        else if (next_token == ">")
        {
            consume(">");
            result = (result > parse_PrimaryExp());
        }
        else
            break;

        next_token = peek();
    }

    return (bool)result;
}
```

# PrimaryExp

```
LogExp := OrExp
OrExp := AndExp [ || AndExp ]*
AndExp := CmpExp [ && CmpExp ]*
CmpExp := PrimaryExp [ == PrimaryExp |
                        < PrimaryExp |
                        > PrimaryExp ]*
PrimaryExp := Int | ( LogExp )
Int := -?[0-9]+
```

```
int parse_PrimaryExp()
{
    int value;
    std::string next_token = peek();

    // Number
    if (is_int(next_token))
    {
        value = std::stoi(next_token);
        consume(next_token);
    }
    // Parenthesis expression: ( LogExp )
    else if (next_token == "(")
    {
        consume("(");
        value = parse_LogExp();
        if (peek() == ")")
            consume(")");
        else
            throw std::runtime_error("Expected: )\n");
    }
    // No valid PrimaryExp found, which is an error
    else
        throw std::runtime_error("expected int or ( )");

    return value;
}
```

# Leftovers


- Regular expression matching
  - Either: via `std::regex` and `regex_match` (`#include <regex>`)
    - Flexible
  - Or: check string manually char-by-char
    - Feasible if pattern is simple
  - Only do the actual cast once the pattern has been matched
    - `int value = std::stoi(token);`

```
bool is_int(const std::string& token)
{
    // ...
}
```

```
LogExp := OrExp
OrExp  := AndExp [ || AndExp ]*
AndExp := CmpExp [ && CmpExp ]*
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```

# Parser flow

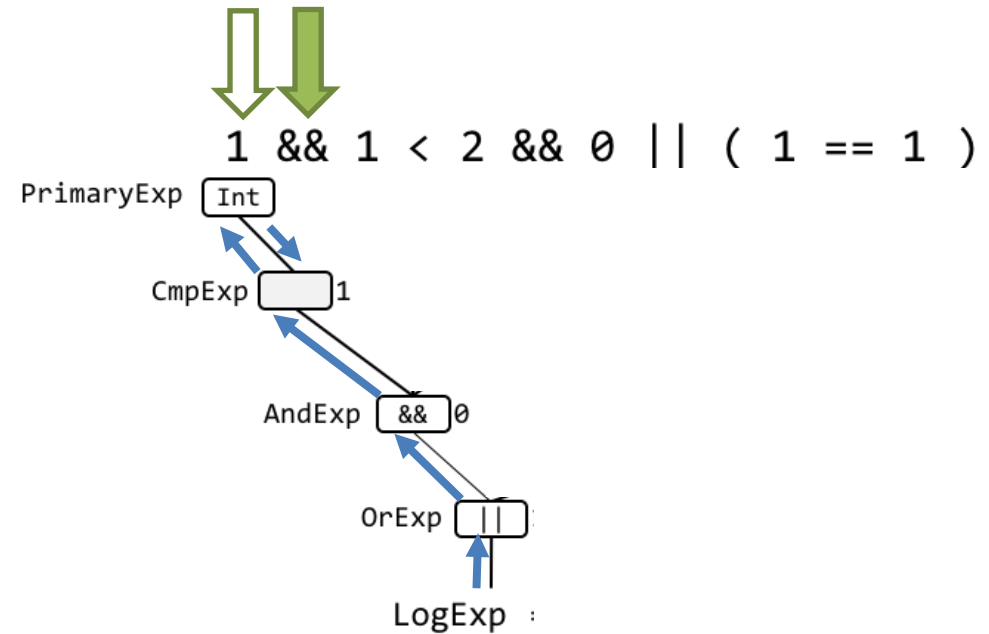
- Parsing starts by:
  1. Setting a step-variable to the first token
  2. Calling the LogExp parsing function

  
1 && 1 < 2 && 0 || ( 1 == 1 )

 LogExp

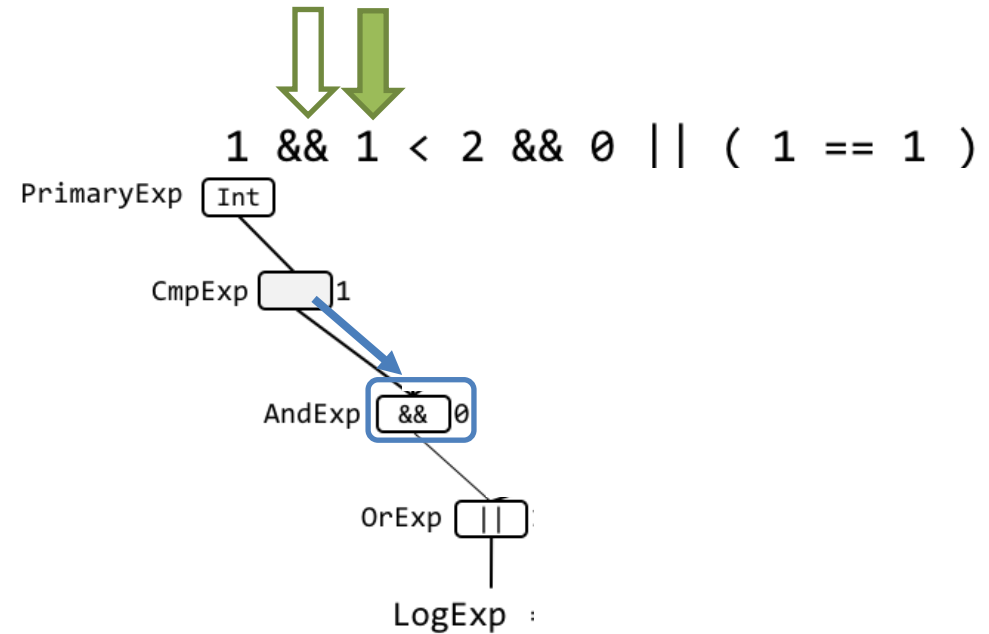
# Parser flow

3. LogExp immediately calls OrExp. Which calls AndExp and so on, until a PrimaryExp is reached.
4. PrimaryExp consumes the token and returns its value ("1") to CmpExp.



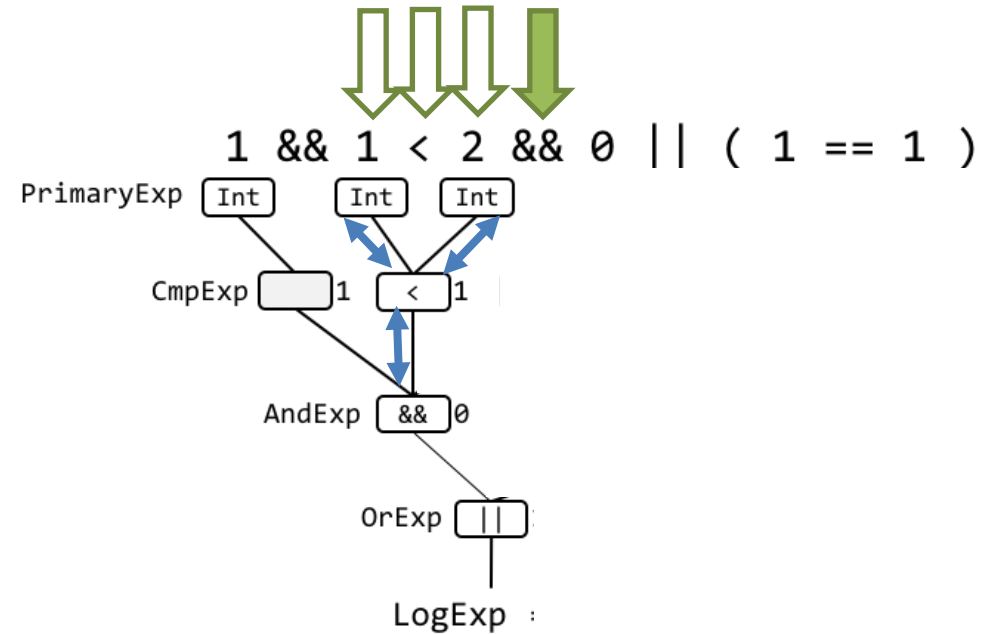
# Parser flow

5. CmpExp does not find a comparison operator, so the value is sent back to AndExp.
6. AndExp finds its operator, “&&”, and consumes it.



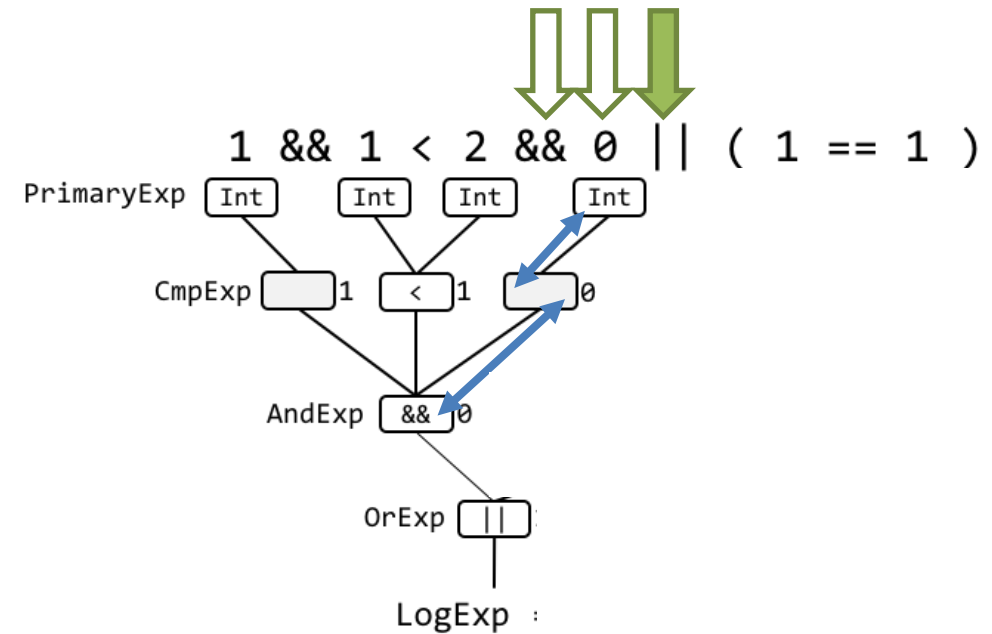
# Parser flow

7. AndExp then queries CmpExp again, which finds and evaluates two operands (consuming three tokens)
8. AndExp now has two operands, which are valued.



# Parser flow

9. Since there is another “&&”, AndExp consumes it and obtains a third operand the same way as before.



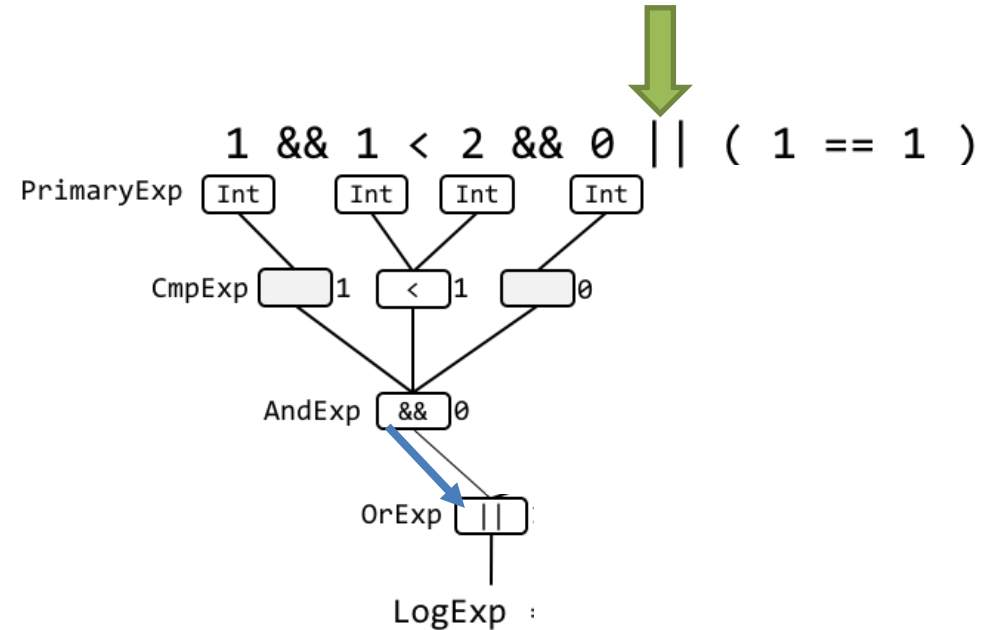


# Parser flow

10. AndExp re-evaluates based on the last operand and returns its result to OrExp.

11. OrExp now has its first operand, and an impending “||”.

And so on.

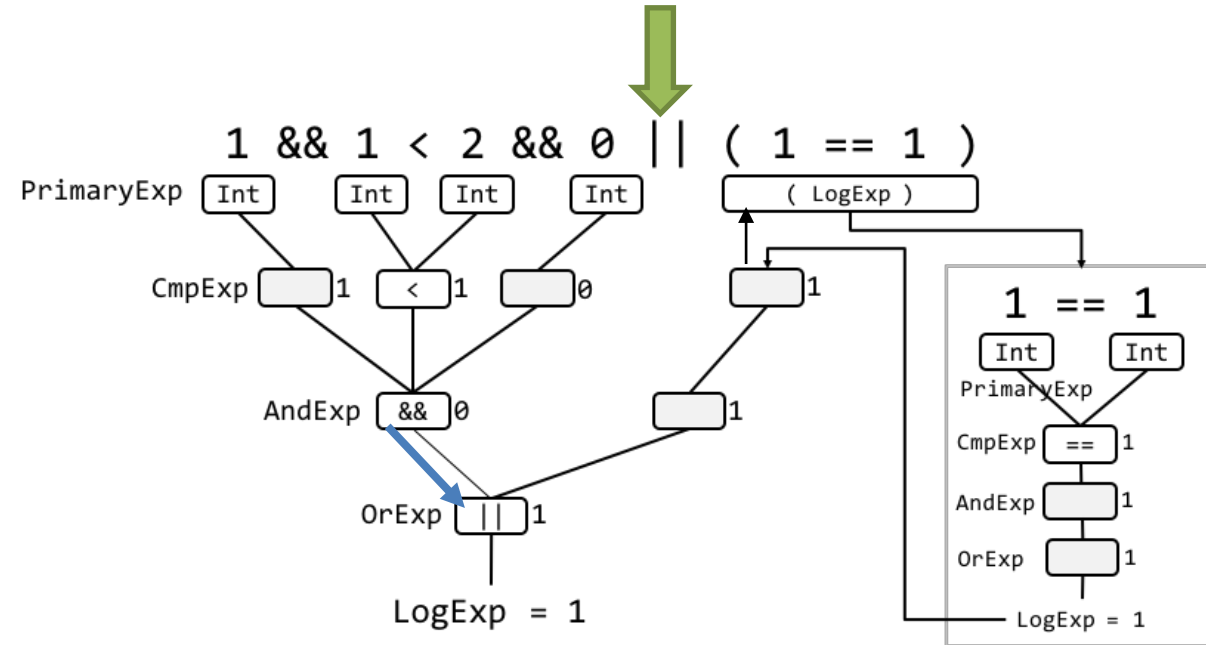


# Parser flow

10. AndExp re-evaluates based on the last operand and returns its result to OrExp.

11. OrExp now has its first operand, and an impending “||”.

And so on.



# RegEx Primer

Variable := [a-zA-z][a-zA-z0-9]*	First char is any letter a-z (upper or lower case), then zero or more letters or digits. Examples: "x", "XY", "x123", "abc123def"
Float := -?[0-9]+\.[0-9]+f	Optional (zero or one) minus sign '-', then one or more digits, then dot (written "."), then one or more digits. Ends with 'f'. Examples: "1.0f", "-00.11f"
Int := -?[0-9]+	Optional (zero or one) minus sign '-', then one or more digits. Examples: "1", "234", "-567", "0089"

Från cppreference.com ☺

*Some people, when confronted with a problem, think "I know, I'll use regular expressions". Now they have two problems.*

# L6

# L6: Interpreter

- Implement an interpreter for the *M@* language
  - **Math expressions**, and some additional statements
    - Structure similar to the interpreter in this lecture
  - Support for **variables**: use of / assignment to
  - Read code from a file
  - Tokenize, parse etc
  - Print results to an output stream
- See handout for more details.

# Grammar for L6

M@ Grammar for L6

Stmt := ConfigStmt | AssgStmt | PrintStmt

ConfigStmt := **config** [ **dec** | **hex** | **bin** ]

AssgStmt := Variable = MathExp

PrintStmt := **print** MathExp

MathExp := SumExp

SumExp := ProductExp [ **+** ProductExp | **-** ProductExp ]\*

ProductExp := PrimaryExp [ **\*** PrimaryExp | **/** PrimaryExp ]\*

PrimaryExp := Int | Variable | ( MathExp )

Variable := [a-zA-z][a-zA-z0-9]\*

Int := -?[0-9]+

# DEMO