

# Machine Learning Written Assignment 03

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## Question 01

Focus on and explain the statements and ideas behind of Lemma 3.1 and Lemma 3.2 on the following paper **On the approximation of functions by tanh neural networks**.

## Chapter 3: Uniform approximation of polynomials

本章節首先為雙曲正切函數 ( $\tanh$ ) 神經網路的逼近誤差推導出界限。所以文章針對多項式的均勻逼近誤差，且在 Sobolev 範數下給出界限。固定大小的淺層神經網路，可以在 supremum norm (最大範數) 下以任意精度逼近單項式 (monomials)，(By Pinkus, 1999)。後來 Gühring 和 Raslan (2021) 又將這個結果推廣到 Sobolev 範數。文章繼續提出了一個新的推廣，能對「某一最大次數以下的所有多項式」得到顯式的誤差估計。對於有效率地逼近解析函數 (analytic functions) 非常關鍵。

### Univariate Polynomials

用  $\tanh$  神經網路逼近任意次數的單變數多項式，我們引入  $p$  階中心有限差分算子: The  $p$ -th order central finite difference operator  $\delta_h^p$  for any  $f \in C^{p+2}([a, b])$  and  $p \in \mathbb{N}$ , defined as

$$\delta_h^p[f](x) = \sum_{i=0}^p (-1)^i \binom{p}{i} f\left(x + \left(\frac{p}{2} - i\right)h\right),$$

where  $\binom{p}{i}$  denotes the binomial coefficient (二項式係數). Next, for any  $p \in \mathbb{N}$ ,  $q \in 2\mathbb{N} - 1$  (odd), and  $M > 0$ , we define the following: The monomial

$$f_p : [-M, M] \rightarrow \mathbb{R}, \quad f_p(y) = y^p,$$

and the neural network function

$$\hat{f}_{(q,h)}(y) = \frac{\delta_{hy}^q[\sigma](0)}{\sigma^{(q)}(0) h^q}.$$

### Example: the case $p = 1$

**Before the lemma:** First we introduce two notations that it will appear several times in the report: (以下介紹  $L^p$  空間，Sobolev 空間的定義)

1.  $L^p(\Omega) = \{f \mid \int_{\Omega} |f|^p dx < \infty\}$
2.  $L^\infty(\Omega) = \{f \mid \sup_{\Omega} |f| < \infty\}$
3. Sobolev space is a functional space, denoting by  $W^{k,p}(\Omega) = \{f \in L^p(\Omega) \mid D^\alpha f \in L^p(\Omega), \forall \alpha \in N_0^d, |\alpha| \leq k\}$  (即微分  $\alpha$  次也在  $L^p$  空間)
4.  $W^{k,\infty}(\Omega) = \{f \in L^\infty(\Omega) \mid D^\alpha f \in L^\infty(\Omega), \forall \alpha \in N_0^d, |\alpha| \leq k\}$
5. **Big O notation**, we say  $f(\varepsilon) = O(\varepsilon)$  for small  $\varepsilon$  if  $\lim_{\varepsilon \rightarrow 0} \frac{f(\varepsilon)}{\varepsilon} = c$  (即速度相同，相除極限趨近於某個 constant.)

### Lemma 3.1

Let  $k \in \mathbb{N}_0$  and  $s \in 2\mathbb{N} - 1$ . Then, for every  $\varepsilon > 0$ , there exists a shallow tanh neural network

$$\Psi_{s,\varepsilon} : [-M, M] \rightarrow \mathbb{R}^{(s+1)/2}$$

of width  $(s+1)/2$  such that

$$\max_{\substack{p \leq s \\ p \text{ odd}}} \|f_p - (\Psi_{s,\varepsilon})_{(p+1)/2}\|_{W^{k,\infty}} \leq \varepsilon.$$

In other words, this neural network simultaneously approximates all odd-degree monomials up to degree  $s$  in the  $W^{k,\infty}$  Sobolev norm with accuracy  $\varepsilon$ .

Moreover, the weights of  $\Psi_{s,\varepsilon}$  scale as

$$O\left(\varepsilon^{-s/2} [2(s+2)(2M)^{1/2}]^{s(s+3)}\right),$$

for small  $\varepsilon$  and large  $s$ .

Lemma 3.1 的目標是建立一個具體的淺層 tanh 神經網路，這個網路寬度大約是  $(s+1)/2$ 。它能夠在 Sobolev 範數下，也就是說它不是只逼近「一個」單項式，而是同時逼近所有次數  $\leq s$  的奇數次單項式  $f_p(y) = y^p$ 。同時誤差量度使用的是  $W^{k,\infty}$  norm（即函數及其導數到  $k$  階為止的最大誤差）。此 Lemma 保證：對所有奇數次  $p \leq s$ ，都可以把誤差壓到小於  $\varepsilon$ 。因此可以以任意精度  $\varepsilon$  逼近所有小於等於某個奇數次  $s$  的單項式，並且給出權重大小隨  $s, \varepsilon$  的成長率。（給出權重大小的量級估計，且不只是函數值的誤差（sup-norm），而是包括導數的誤差都被控制。）

**Proof.** Let  $p \leq s$  be odd and let  $0 < h < \frac{2}{pM}$ . Fix  $0 \leq m \leq \min\{k, p+1\}$ .

（條件：取一個奇數次多項式次數  $p \leq s$ ，步長  $0 < h < \frac{2}{pM}$ ，並令導數階數  $m = \min\{k, p+1\}$ ）。

利用有限差分運算子  $\delta_h^p$  和泰勒定理（Taylor's theorem）可得到：For each  $x$  there exists  $\xi_{x,i}$  such that

$$\frac{d^m}{dx^m} \delta_h^p = \sum_{i=0}^p (-1)^i \binom{p}{i} \left(\frac{p}{2} - i\right)^m h^m \sigma^{(m)}\left(\left(\frac{p}{2} - i\right)hx\right).$$

Expanding  $\sigma^{(m)}$  in Taylor series at 0 gives

$$\begin{aligned} \frac{d^m}{dx^m} \delta_h^p &= \sum_{i=0}^p (-1)^i \binom{p}{i} \left(\frac{p}{2} - i\right)^m h^m \left[ \sum_{l=m}^{p+1} \frac{\sigma^{(l)}(0)}{(l-m)!} \left(\frac{p}{2} - i\right)^{l-m} (hx)^{l-m} \right] \\ &\quad + \sum_{i=0}^p (-1)^i \binom{p}{i} \left(\frac{p}{2} - i\right)^m h^m \left[ \frac{\sigma^{(p+2)}(\xi_{x,i})}{(p+2-m)!} \left(\frac{p}{2} - i\right)^{p+2-m} (hx)^{p+2-m} \right]. \end{aligned}$$

（上述式子分兩項，第一部分是「有限展開到  $p+1$  階」的泰勒級數。第二部分是誤差項，取在某個中間點  $\xi_{x,i}$ ）

From Katsuura in 2009, we know that

$$\sum_{i=0}^p (-1)^i \binom{p}{i} \left(\frac{p}{2} - i\right)^l = \begin{cases} p!, & l = p, \\ 0, & l \neq p, \quad l = 0, \dots, p. \end{cases} \quad (1)$$

Moreover, the above equality remains valid also for  $l = p+1$ , since all summands change sign when  $i$  is replaced by  $p-i$ .

((1) 式的意思是，有限差分結構會「過濾掉」除了  $l = p$  的所有項，保留單一最高次項。而且 (1) 對於  $l = p+1$  仍然成立)

Using this, we can rewrite the first term above as

$$\begin{aligned} &\sum_{i=0}^p (-1)^i \binom{p}{i} \left(\frac{p}{2} - i\right)^m h^m \left[ \sum_{l=m}^{p+1} \frac{\sigma^{(l)}(0)}{(l-m)!} \left(\frac{p}{2} - i\right)^{l-m} (hx)^{l-m} \right] \\ &= h^m \left[ \sum_{l=m}^{p+1} \frac{\sigma^{(l)}(0)}{(l-m)!} (hx)^{l-m} \right] \sum_{i=0}^p (-1)^i \binom{p}{i} \left(\frac{p}{2} - i\right)^l. \end{aligned}$$

By the above result, only the term  $l = p$  survives, so for  $0 \leq m \leq p$  we obtain

$$h^m \frac{\sigma^{(p)}(0)}{(p-m)!} (hx)^{p-m} p!, \text{ and for } m = p+1, \text{ the expression vanishes.}$$

由於 (1) 式的結果，除了  $l = p$  以外的項都會消失，因此，這代表著有限差分能精確抓出第  $p$  階導數項，這就是為什麼它能近似多項式。同時，上式等價於下面式子：

$$\frac{d^m}{dx^m} \delta_h^p = h^p \sigma^{(p)}(0) f_p^{(m)}(x).$$

整理來自泰勒展開式的剩餘項後，可以得到近似誤差

$$\hat{f}_{p,h}^{(m)}(x) - f_p^{(m)}(x) = \sum_{i=0}^p (-1)^i \binom{p}{i} \frac{1}{(p+2-m)!} \frac{\sigma^{(p+2)}(\xi_{x,i})}{\sigma^{(p)}(0)} \left(\frac{p}{2} - i\right)^{p+2} h^2 x^{p+2-m}.$$

可以注意到誤差中的  $h^2$ ，所以有限差分對於 monomial 的近似是二階精確的。

文章中的 Lemmas A.1 and A.4 給出了  $\sigma = \tanh$  的各階導數的上下界估計，將其帶入上式，可得：For  $m \leq \min\{k, p+1\}$ :

$$\|f_p - \hat{f}_{p,h}\|_{W^{m,\infty}} \leq \sum_{i=0}^p \binom{p}{i} \frac{|\sigma^{(p+2)}(\xi_{x,i})|}{|\sigma^{(p)}(0)|} \left|\frac{p}{2} - i\right|^{p+2} h^2 M^{p+2}.$$

Estimating the binomial sums gives

$$\|f_p - \hat{f}_{p,h}\|_{W^{m,\infty}} \leq 2^p (2(p+2))^{p+3} \left(\frac{p}{2}\right)^{p+2} h^2 M^{p+2} \leq (2(p+2)pM)^{p+3} h^2.$$

Hence, if  $k \leq p+1$ ,

$$\|f_p - \hat{f}_{p,h}\|_{W^{k,\infty}} \leq (2(p+2)pM)^{p+3} h^2.$$

Now suppose  $k > p+1$ , and consider  $m \geq p+2$ . In this case  $f_p^{(m)} = 0$ , so it suffices to bound  $\hat{f}_{p,h}^{(m)}$ . For  $0 < h < 1$ ,

$$\begin{aligned} |\hat{f}_{p,h}^{(m)}(x)| &= \left| \frac{1}{h^p \sigma^{(p)}(0)} \sum_{i=0}^p (-1)^i \binom{p}{i} \left(\frac{p}{2} - i\right)^m h^m \sigma^{(m)}\left(\left(\frac{p}{2} - i\right)hx\right) \right| \\ &\leq 2 \sum_{i=0}^p \binom{p}{i} \left|\frac{p}{2} - i\right|^m h^2 (2m)^{m+1}. \end{aligned}$$

(如果  $k \leq p+1$ ，直接得到這個估計；但如果  $k > p+1$ ，因為  $f_p^{(m)} = 0$ ，我們只需界定近似函數的高階導數即可，最後仍可以得到誤差  $\sim h^2$ )

Bounding the sum gives

$$|\hat{f}_{p,h}^{(m)}(x)| \leq 2^{p+1} \left(\frac{p}{2}\right)^k (2k)^{k+1} h^2 \leq (2pk)^{k+1} h^2.$$

We thus obtain, for arbitrary  $k \in \mathbb{N}$ ,

$$\|f_p - \hat{f}_{p,h}\|_{W^{k,\infty}} \leq \left( (2(p+2)pM)^{p+3} + (2pk)^{k+1} \right) h^2 = \varepsilon.$$

(即對於任意 Sobolev 階數  $k$ ，真實 monomial  $f_p(x) = x^p$  與近似函數  $\hat{f}_{p,h}$  的差距，可以用一個與  $h^2$  成正比的誤差上界來控制，而常數則由  $p, k, M$  決定。

Furthermore, observe that the weights scale as

$$O\left(\max_i \binom{p}{i} h^{-p}\right).$$

For  $\varepsilon \rightarrow 0$  and large  $p$ , it holds that

$$O(h^{-p}) = O\left(\varepsilon^{-p/2} ((p+2)(2M)^{1/2})^p (p+3)\right),$$

where the implied constant depends on  $k$ .

(權重的數值大小與二項式係數的最大值、以及  $h^{-p}$  成正比。即當我們要更精準近似  $\varepsilon = 0$  時，權重會急遽放大，並且放大的速率大約是  $\varepsilon^{-p/2}$ ，還要乘上一個依賴於  $p, M$  的因子。)

Next, using Stirling's approximation, we find that for  $0 \leq i \leq p$  it holds that

$$\binom{p}{i} \leq \binom{p}{\lfloor (p-1)/2 \rfloor} \leq \frac{e p^{(p+1)/2}}{2\pi ((p-1)/2)^{p/2} ((p+1)/2)^{p/2+1}} = O\left(\frac{2^p}{\sqrt{p}}\right).$$

這裡使用 Stirling's 近似來估計二項式係數最大值，可知最大的二項式係數大約和  $\frac{2^p}{\sqrt{p}}$  同階，因此，權重的最終尺度是：

$$O\left(\varepsilon^{-p/2} (2(p+2)(2M)^{1/2})^p (p+3)\right).$$

Regarding the network architecture, note that the neurons needed for all  $\hat{f}_{p,h}$  are already available in the network  $\hat{f}_{s,h}$ . This allows us to define the shallow tanh neural network  $\Psi_{s,\varepsilon}$  by

$$(\Psi_{s,\varepsilon})_p = \hat{f}_{p,h},$$

such that it only has  $(s+1)/2$  neurons in its hidden layer. The width follows directly from its definition and from the fact that  $\sigma$  is an odd function.

(要構造  $\hat{f}_{p,h}$  所需的神經元，已經包含在  $\hat{f}_{s,h}$  的網路裡。要同時近似所有次數  $\leq s$  的單項式，其實只需要  $\frac{s+1}{2}$  個神經元。)  $\square$

We now extend the previous result to monomials with even degree. To this end, we rely on the observation that for  $n \in \mathbb{N}$  and  $\alpha > 0$ , it holds that

$$y^{2n} = \frac{1}{2\alpha(2n+1)} \left( (y+\alpha)^{2n+1} - (y-\alpha)^{2n+1} - 2 \sum_{k=0}^{n-1} \binom{2n+1}{2k} \alpha^{2(n-k)+1} y^{2k} \right). \quad (2)$$

This formula allows us to construct recursively defined tanh neural network approximations of even powers of  $y$ . The following lemma quantifies the uniform approximation accuracy of these networks in the Sobolev norm.

(接著我們處理偶數次 monomial (形如  $y^{2n}$ )，其可以用「相鄰兩個奇次幂」以及更低階偶次幂的線性組合 (即 (2) 式) 來表示。這也提供了一個可行的遞迴構造方式：先近似奇次 monomial，再利用這個公式去近似偶次 monomial。)

### Lemma 3.2

Let  $k \in \mathbb{N}_0$ ,  $s \in 2\mathbb{N} - 1$  and  $M > 0$ . For every  $\varepsilon > 0$ , there exists a shallow tanh neural network

$$\psi_{s,\varepsilon} : [-M, M] \rightarrow \mathbb{R}^s$$

of width  $\frac{3}{2}(s+1)$  such that

$$\max_{p \leq s} \|f_p - (\psi_{s,\varepsilon})_p\|_{W^{k,\infty}} \leq \varepsilon.$$

Furthermore, the weights scale as

$$O\left(\varepsilon^{-s/2} (\sqrt{M}(s+2))^{\frac{3s(s+3)}{2}}\right)$$

for small  $\varepsilon$  and large  $s$ .

(給定  $k \in \mathbb{N}_0$  (非負整數)、 $s \in 2\mathbb{N} - 1$  (奇數)、以及  $M > 0$ 。對於每個  $\varepsilon > 0$ ，都存在一個淺層 tanh 神經網路，隱藏層寬度為  $\frac{3}{2}(s+1)$ ，並且這個神經網路可以同時近似所有次方單項式  $f_p(y) = y^p$  ( $p \leq s$ )，而且誤差在 Sobolev 範數下小於  $\varepsilon$ 。另外，這個網路的權重大小隨著上面給的參數縮放。

因此，Lemma 3.2 是 Lemma 3.1 的推廣，從「奇數次幂」擴展到「偶數次幂」，最後保證整個多項式空間 (最高次數  $s$ ) 都能用一個淺層 tanh 網路統一逼近。)

**Proof.** For  $h, M > 0$  and  $p \leq s$ , we define  $\hat{f}_{p,h} : [-M-1, M+1] \rightarrow \mathbb{R}$  by

$$f_p(y) = y_p, \quad \hat{f}_{q,h}(y) = \frac{\delta_h^q y}{\sigma^{(q)}(0) h^q}.$$

For  $\varepsilon > 0$ , choose  $h > 0$  small enough depending on  $\varepsilon$ , let  $\alpha \leq 1$ , and let  $y \in [-M, M]$ . We then define

$$(\psi_{s,\varepsilon}(y))_p = \hat{f}_{p,h}(y), \quad \text{for } p \text{ odd,}$$

and for  $p = 2n$  even, we set  $(\psi_{s,\varepsilon}(y))_{2n}$  recursively by

$$(\psi_{s,\varepsilon}(y))_0 = 1,$$

and

$$(\psi_{s,\varepsilon}(y))_{2n} = \frac{1}{2\alpha(2n+1)} \left( \hat{f}_{2n+1,h}(y+\alpha) - \hat{f}_{2n+1,h}(y-\alpha) - 2 \sum_{k=0}^{n-1} \binom{2n+1}{2k} \alpha^{2(n-k)+1} (\psi_{s,\varepsilon}(y))_{2k} \right). \quad (3)$$

(對奇數次單項式: 我們直接用  $\hat{f}_{p,h}$  就能逼近, Lemma 3.1 已經保證誤差  $\leq \varepsilon$ 。對於偶數次單項式, 我們使用上面的遞迴公式 (3) 來構造: 將偶數次幂  $y^{2n}$  展開為「高次幂差」減去「低次幂的線性組合」。遞迴的 idea 是: 如果低次偶數次方已經能被近似, 則更高次的偶數次方也能近似。)

We introduce the notation

$$E_p = \|f_p - (\psi_{s,\varepsilon})_p\|_{W^{k,\infty}}.$$

We claim that for all  $\varepsilon > 0$ , there exists  $h > 0$  such that for all  $p \leq s$ ,

$$E_p \leq E_p^* = \frac{2^{p/2}(1+\alpha)^{(p^2+p)/2}}{\alpha^{p/2}} \varepsilon. \quad (4)$$

**Step 1: Odd  $p$ .** Choosing  $h$  as in Lemma 3.1 implies

$$\max_{\substack{p \leq s \\ p \text{ odd}}} E_p \leq \varepsilon, \quad (5)$$

which proves (4) for odd  $p$ , since  $(1+\alpha)/\alpha \geq 1$ .

**Step 2: Even  $p$ , base case.** For  $p = 2$ , we find

$$E_2 \leq \frac{2\varepsilon}{6\alpha} \leq E_2^*,$$

which establishes the base step.

**Step 3: Even  $p$ , induction step.** Let  $n \in \mathbb{N}$  with  $2n+1 \leq s$  and  $n > 1$ . Assume by the induction hypothesis that

$$E_{2k} \leq E_{2k}^*, \quad \forall k < n.$$

From (3) it follows that

$$E_{2n} \leq \frac{1}{2\alpha(2n+1)} \left( E_{2n+1} + E_{2n+1} + 2 \sum_{k=1}^{n-1} \binom{2n+1}{2k} \alpha^{2(n-k)+1} E_{2k} \right). \quad (6)$$

By the induction hypothesis and the monotonicity of  $E_{2k}^*$  in  $k$ , we have

$$E_{2k} \leq E_{2k}^* \leq E_{2(n-1)}^*.$$

Using also (5) and the fact that  $\varepsilon \leq E_{2(n-1)}^*$ , we can estimate (6) as

$$E_{2n} \leq \frac{1}{\alpha(2n+1)} \left( \max_{\substack{p \leq s \\ p \text{ odd}}} E_p + \sum_{k=1}^{n-1} \binom{2n+1}{2k} \alpha^{2(n-k)+1} E_{2(n-1)}^* \right).$$

This yields

$$E_{2n} \leq \frac{1}{\alpha} \left( E_{2(n-1)}^* + (1+\alpha)^{2n+1} E_{2(n-1)}^* \right) \leq \frac{2}{\alpha} (1+\alpha)^{2n+1} E_{2(n-1)}^*.$$

Recalling the definition of  $E_{2(n-1)}^*$ , we obtain

$$E_{2n} \leq \left( \frac{2}{\alpha} (1+\alpha)^{2n+1} \right)^n \varepsilon = E_{2n}^*.$$

(此處進行誤差分析  $E_p$ , 對奇數情況: Lemma 3.1 直接給  $E_p \leq \varepsilon \leq E_p^*$ . 而對偶數情況: 使用數學歸納法。Base case ( $p = 2$ ): 驗算成立。歸納步: 假設所有  $2k < 2n$  都成立, 利用展開式與估計不等式, 把

高次誤差  $E_{2n}$  控制在  $E_{2n}^*$  之內。綜合奇數與偶數情況，得到所有  $p \leq s$  的誤差都被控制在  $\varepsilon$  以內，完成證明。）

Next, we optimize inequality (4) by choosing the optimal value of  $\alpha$ . Lemma A.2 shows that the optimal choice is  $\alpha = 1/s$ . We therefore conclude that, for any  $\varepsilon > 0$ , there exists a shallow tanh neural network

$$\psi_{s,\varepsilon},$$

with width independent of  $\varepsilon$ , such that

$$\max_{p \leq s} \|f_p - (\psi_{s,\varepsilon})_p\|_{W^{k,\infty}} \leq \sqrt{e} (2es)^{s/2} \varepsilon.$$

(接下來對式子 (4) 進行最佳化，方法是選擇一個最佳的參數  $\alpha$ 。根據 Lemma A.2 的結果，最佳的  $\alpha$  值就是  $\alpha = 1/s$ 。這裡  $\alpha$  是一個控制平移間距的參數，最佳化  $\alpha$  的目的是讓近似誤差最小。)

(因此，我們得到結論：對任何誤差精度  $\varepsilon > 0$ ，都可以構造一個淺層的  $\tanh$  神經網路，其寬度不依賴於  $\varepsilon$ ，滿足上面的誤差上界：這表示：我們能找到一個固定寬度的淺層網路，使得對所有  $p \leq s$ ，它對多項式的近似誤差在 Sobolev norm  $W^{k,\infty}$  意義下，被一個常數倍的  $\varepsilon$  控制住。)

Replacing  $\varepsilon \mapsto \frac{\varepsilon}{\sqrt{e}(2es)^{s/2}}$  recovers the claimed error bound in the statement of the lemma.

(此處調整  $\varepsilon$  的尺度，來對齊 Lemma 3.2 開頭所要求的誤差形式。)

To quantify the size of the weights, we observe from equation (27) that the weight bound from Lemma 3.1 needs to be multiplied by the factor

$$\max_k \binom{s}{2k} \left(\frac{1}{s}\right)^{s-2k} \leq s \sum_{j=0}^s \binom{s}{j} \left(\frac{1}{s}\right)^{s-j} \leq s \left(1 + \frac{1}{s}\right)^s = O(s),$$

where we used the binomial theorem.

(估計網路的權重大小，我們利用式子(3)。這顯示在 Lemma 3.1 的權重上限之外，我們需要再乘上一個因子：利用二項式展開，可以證明這個最大值被上面式子控制：並利用二項式展開，可以證明這個最大值權重增長到約  $O(s)$  的等級，不會比指數更糟。)

Therefore, the final weight bound can be seen to satisfy

$$O\left(\varepsilon^{-s/2} s \left(2(2es)^{1/4} [2(M+1)]^{1/2} (s+2)\right)^{s(s+3)}\right) = O\left(\varepsilon^{-s/2} (\sqrt{M}(s+2))^{\frac{3s(s+3)}{2}}\right),$$

for small  $\varepsilon$  and large  $s$ . This proves the weight bound as stated in the lemma.(證明了 Lemma 3.2 中所聲明的權重上界。)

Finally, we note that the constructed approximations indeed correspond to a shallow tanh neural network of the claimed size. In fact, by observing that  $\sigma$  is an odd function, and using

$$\delta_h^p[f](x) = \sum_{i=0}^p (-1)^i \binom{p}{i} f\left(x + \left(\frac{p}{2} - i\right) h\right),$$

together with Lemma 3.1 and equation (3), we conclude that a shallow tanh neural network suffices. Specifically, the values of the  $\frac{3}{2}(s+1)$  neurons in the hidden layer are given by

$$\sigma\left(\left(\frac{s}{2} - i\right)h(y + \beta)\right), \quad i = 0, 1, \dots, \frac{s-1}{2}, \quad \beta \in \{-\alpha, 0, \alpha\}.\square$$

(由上面可以確認我們成功構造出一個隱藏層大小受控的淺層  $\tanh$  網路來近似多項式，並且誤差與權重都有清楚的上界。)

## Question 02—Unanswered questions

1. In the paper we discussed in the class (**On the approximation of functions by tanh neural networks**), we drew the neural network. If a network has more layer and more neuron, will it be better, that is, has better accuracy and less loss?

## References

- [1] On the approximation of functions by tanh neural networks, Tim DeRyck, Samuel Lanthaler, Siddhartha Mishra, Seminar for Applied Mathematics, ETH Zürich, Rämistrasse 101, 8092 Zürich, Switzerland.